CSE 632: Analysis of Algorithms II: Randomized Algorithms

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Lecture 6: Concentration Inequalities

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1 Markov's Inequality

Lemma 1 (Markov's Inequality). Let X be a non-negative random variable with mean $\mu = \mathbb{E}X$. Then we have

$$\forall a \geq 0: \ \Pr\left(X \geq a\mu\right) \leq \frac{1}{a};$$
 Equivalently,
$$\forall a \geq 0: \ \Pr\left(X \geq a\right) \leq \frac{\mu}{a}.$$

Proof. We have

$$\begin{split} \mathbb{E}X &= \Pr(X < a) \, \mathbb{E}[X \mid X < a] + \Pr(X \ge a) \, \mathbb{E}[X \mid X \ge a] \\ &\ge \Pr(X \ge a) \, \mathbb{E}[X \mid X \ge a] \\ &\ge \Pr(X \ge a) \cdot a. \end{split}$$

The lemma then follows.

Example 2. Suppose we toss a fair coin n times. What is the probability that we get at least 0.6n Heads? How fast does the probability diminish as n grows?

Let $X_i = 1$ if the *i*th toss is Head and $X_i = 0$ otherwise. So $\mathbb{E}X_i = 1/2$. Let $X = \sum_{i=1}^n X_i$ be the number of Heads in *n* tosses. So $\mathbb{E}X = n/2$. By Markov's Inequality, we have

$$\Pr(X \ge 0.6n) \le \frac{0.5n}{0.6n} = \frac{5}{6}.$$

2 Chebyshev's Inequality

Lemma 3 (Chebyshev's Inequality). Let X be a random variable with mean $\mu = \mathbb{E}X$ and variance $\sigma^2 = \text{Var}(X)$ where $\sigma \geq 0$ is the standard deviation. Then we have

$$\forall a \ge 0: \Pr(|X - \mu| \ge a\sigma) \le \frac{1}{a^2};$$

Equivalently, $\forall a \ge 0: \Pr(|X - \mu| \ge a) \le \frac{\sigma^2}{a^2}.$

Proof. Let $Y = (X - \mu)^2$. Note that $Y \ge 0$ and $\mathbb{E}Y = \mathbb{E}[(X - \mu)^2] = \sigma^2$. We then have

$$\Pr(|X - \mu| \ge a) = \Pr(Y \ge a^2)$$

$$\le \frac{\mathbb{E}Y}{a^2}$$

$$= \frac{\sigma^2}{a^2},$$
(Markov's Inequality)

as wanted.

Example 4. As in Example 2, each X_i has variance 1/4 and hence Var(X) = n/4. By Chebyshev's Inequality, we have

$$\Pr(X \ge 0.6n) = \Pr(X - \mathbb{E}X \ge 0.1n)$$
$$= \frac{1}{2} \Pr(|X - \mathbb{E}X| \ge 0.1n)$$
$$\le \frac{1}{2} \frac{n/4}{(0.1n)^2} = \frac{12.5}{n} \xrightarrow{n \to \infty} 0.$$

3 Chernoff Bounds

Lemma 5 (Chernoff Bounds). Let X_1, \ldots, X_n be independent random variables taking values in [0,1]. Let $X = \sum_{i=1}^n X_i$ be their sum with mean $\mu = \mathbb{E}X = \sum_{i=1}^n \mathbb{E}X_i$. Then we have

$$\forall \delta \ge 0: \Pr(X \ge (1+\delta)\mu) \le \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{\mu} \tag{1}$$

and
$$\Pr\left(X \le (1-\delta)\mu\right) \le \left(\frac{e^{-\delta}}{(1-\delta)^{1-\delta}}\right)^{\mu}$$
. (2)

A looser but simpler version is

$$\forall \delta \in (0,1): \Pr\left(X \ge (1+\delta)\mu\right) \le e^{-\frac{\delta^2 \mu}{3}} \tag{3}$$

and
$$\Pr\left(X \le (1 - \delta)\mu\right) \le e^{-\frac{\delta^2 \mu}{2}}$$
. (4)

Proof. Our plan is to pick a "nice" auxiliary function φ and apply Markov's inequality to $Y = \varphi(X)$. Namely, we expect to have

$$\Pr\left(X \geq (1+\delta)\mu\right) \overset{\text{(i)}}{\leq} \Pr\left(\varphi(X) \geq \varphi((1+\delta)\mu)\right)$$

$$\overset{\text{(ii)}}{\leq} \frac{\mathbb{E}[\varphi(X)]}{\varphi((1+\delta)\mu)}$$
(Markov's Inequality)
$$\overset{\text{(iii)}}{\leq} \text{ nice bound,}$$

where (i) requires φ is monotone increasing, (ii) requires $\varphi \geq 0$ since Markov's inequality applies only to non-negative random variables, and (iii) requires we can upper bound $\mathbb{E}[\varphi(X)]$ nicely.

Our choice of φ is the exponential function $\varphi(x) = e^{tx}$ where $t \ge 0$ is a parameter to be decided. Note that

$$\mathbb{E}[\varphi(X)] = \mathbb{E}\left[e^{tX}\right] = \mathbb{E}\left[e^{t\sum_{i=1}^{n} X_i}\right] = \mathbb{E}\left[\prod_{i=1}^{n} e^{tX_i}\right] = \prod_{i=1}^{n} \mathbb{E}\left[e^{tX_i}\right],$$

where the last equality is because X_1, \ldots, X_n are independent.

Fact 6. Suppose Z is a random variable taking values in [0,1]. Then we have

$$\mathbb{E}\left[e^{tZ}\right] \le 1 + (e^t - 1)\mathbb{E}Z.$$

Proof. For all $x \in [0,1]$, it holds $e^{tx} \le 1 + (e^t - 1)x$. Therefore,

$$\mathbb{E}\left[e^{tZ}\right] \le \mathbb{E}\left[1 + (e^t - 1)Z\right] = 1 + (e^t - 1)\mathbb{E}Z$$

as claimed. \Box

By Fact 6, we deduce that

$$\mathbb{E}\left[e^{tX}\right] = \prod_{i=1}^{n} \mathbb{E}\left[e^{tX_{i}}\right] \leq \prod_{i=1}^{n} \left(1 + (e^{t} - 1)\mathbb{E}X_{i}\right)$$

$$\leq \prod_{i=1}^{n} e^{(e^{t} - 1)\mathbb{E}X_{i}}$$

$$= e^{\sum_{i=1}^{n} (e^{t} - 1)\mathbb{E}X_{i}}$$

$$= e^{(e^{t} - 1)\mu}.$$
(5)

Therefore,

$$\Pr(X \ge (1+\delta)\mu) \le \Pr\left(e^{tX} \ge e^{t(1+\delta)\mu}\right) \qquad \text{(since } t \ge 0)$$

$$\le \frac{\mathbb{E}\left[e^{tX}\right]}{e^{(1+\delta)t\mu}} \qquad \text{(Markov's Inequality)}$$

$$\le \frac{e^{(e^t-1)\mu}}{e^{(1+\delta)t\mu}} \qquad \text{(by Eq. (5))}$$

$$= e^{\left(e^t-1-(1+\delta)t\right)\mu}. \qquad (6)$$

Notice that Eq. (6) holds for all $t \ge 0$. We would pick a best $t \ge 0$ that minimizes the exponent $e^t - 1 - (1 + \delta)t$. By Calculus our choice is $t = \log(1 + \delta) \ge 0$, and for this $t \to 0$ becomes

$$\Pr\left(X \ge (1+\delta)\mu\right) \le \left(\frac{e^{\delta}}{(1+\delta)^{1+\delta}}\right)^{\mu},$$

which establishes Eq. (1).

The proof of Eq. (2) is very similar. Let $t \leq 0$ to be decided, and we have

$$\Pr(X \le (1 - \delta)\mu) \le \Pr\left(e^{tX} \ge e^{t(1 - \delta)\mu}\right) \qquad \text{(since } t \le 0)$$

$$\le \frac{\mathbb{E}\left[e^{tX}\right]}{e^{(1 - \delta)t\mu}} \qquad \text{(Markov's Inequality)}$$

$$\le \frac{e^{(e^t - 1)\mu}}{e^{(1 - \delta)t\mu}} \qquad \text{(by Eq. (5))}$$

$$= e^{\left(e^t - 1 - (1 - \delta)t\right)\mu}. \qquad (7)$$

We choose $t = \log(1 - \delta) \le 0$ which minimizes the exponent $e^t - 1 - (1 - \delta)t$, and then Eq. (7) becomes

$$\Pr\left(X \le (1 - \delta)\mu\right) \le \left(\frac{e^{-\delta}}{(1 - \delta)^{1 - \delta}}\right)^{\mu},$$

establishing Eq. (2).

The bounds Eqs. (3) and (4) are obtained by upper bounding the right-hand sides of Eqs. (1) and (2) with simpler functions. \Box

Example 7. As in Example 2, by Chernoff Bound we have

$$\Pr\left(X \ge 0.6n\right) = \Pr\left(X \ge (1 + 0.2) \,\mathbb{E}X\right) \le \left(\frac{e^{0.2}}{1.2^{1.2}}\right)^{0.5n} \le 0.991^n \stackrel{n \to \infty}{\longrightarrow} 0.$$

In particular, the probability decays to zero exponentially fast.