

# Languages and Compilers

## **Semantic Analysis:**

## **Expressions and More**

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# Today's Plan...

- Analysing Expressions.
- Other Types of Identifiers.

# Analysing Expressions.

- Handling Variables and Literal Constants.
  - Type Checking.
- Statements & Expressions.
  - Assignment Example.
  - Expression Example.

# Analysing Variables and Literals.

- If the RD parser drives the semantic analysis:
  - (as we saw last week – not the only way)
- To semantically analyse the use of a variable:
  - ✓ check the name has been declared, and
  - ✓ check for appropriate type compatibility.
- To semantically analyse the use of a **literal**:  
(constant values, e.g. 42, 3.141, “hello world”)
  - ✓ check for appropriate type compatibility.

# General Strategy.

- Recognisers return type information.
- Let's have some Semantics methods:
  - *CheckId(token)* – recognise a variable identifier, and return its type (if your BNF contains Variable, you might not even need the helper!)
  - *CheckTypesSame(type1, type2)* – check whether two types are compatible (in a simple language, whether they're the same type or not)

```
private void recDecl () {  
    int varType;  
    if (have("integer")) {  
        mustBe("integer");  
        varType = LanguageType.Integer;  
    } ...  
    ...DeclareId (scanner.CurrentToken,  
                 varType);  
    mustBe (Token.IdentifierToken);  
}
```

```
private void recLet () {  
    mustBe("let");  
    int varType =  
        ...CheckId (scanner.CurrentToken);  
    mustBe (Token.IdentifierToken);  
    mustBe (":=");  
    int exprType = recExpr();  
    ...CheckTypesSame(varType, exprType);  
}
```

# The *CheckId* Method ...

- Supply with the required token.

```
public int CheckId (IToken id) {
```

- (remember `int` is how Ardkit represents type expressions)

- If not an identifier then return a dummy type.

```
    if (!id.Is (Token.IdentifierToken))  
        return LanguageType.Undefined;
```

## ... *CheckId* continued.

- If the variable name is not defined in the current scope (see last week!) then report a semantic error, and return a dummy type ...

```
if (!Scope.CurrentScope.IsDefined (id.TokenValue)) {  
    semanticError (new NotDeclaredError (id));  
    return LanguageType.Undefined;  
}
```

- ... else look up and return the type

```
return Scope.CurrentScope.Get (id.TokenValue).Type;  
} // end CheckId method.
```



# The *CheckLiteral* method.

- We could have a very similar helper method that takes a token representing a literal, and returns its type...

```
int CheckLiteral (IToken lit) {  
    if (lit.Is (Token.IntegerToken))  
        return LanguageType.Integer;  
    else if (lit.Is (Token.RealToken))  
        return LanguageType.Real;  
    else  
        return LanguageType.Undefined;  
}
```

- Or you could do this in `recLiteral`, if your BNF required you to write that

# *CheckTypesSame...*

- The arguments are the two types:

```
public void CheckTypesSame (?????,  
                           int oldType, int newType) {
```

- Complain if they don't match:

```
    if (oldType != newType) {  
        semanticError (new TypeConflictError  
                      (?????, newType, oldType));  
    }  
} // end checkType method.
```

# *CheckTypesSame* completed

- The arguments are the two types:

```
public void CheckTypesSame (IToken token,  
                           int oldType, int newType) {
```

- Complain if they don't match:

```
    if (oldType != newType) {  
        semanticError (new TypeConflictError  
                      (token, newType, oldType));  
    }  
} // end checkType method.
```

- We also need the token to give a source location in the error message!

# *CheckTypesSame* in practice

- This is a very simple kind of type checking! No implicit type conversion...
- Many real languages would need more complex logic in typechecking methods
  - identifying when one type can be automatically converted to another
  - handling user-defined types
  - etc. etc.

# Dealing with Expressions.

- Consider the EBNF extracts

```
<Expr> ::= <Factor> ( (+|-) <Factor>)* ;  
<Factor> ::= Ident | Integer | "(" <Expr> ")" ;
```

```
private void recExpr () {  
    recFactor ();  
    while (have("+") || have("-")) {  
        if (have("+")) mustBe ("+");  
        else           mustBe ("-");  
        recFactor ();  
    }  
}
```

- How do we add semantics to this?
- Make each production return the type!

# Flashback to the Type Inference lecture...

$$\begin{array}{c} (\text{height} / \text{num}) * (\text{i} + 1) \\ \text{float} \quad \text{int} \quad \text{int} \quad \text{int} \\ \text{float} \quad \text{int} \\ \text{float} \end{array}$$

- Start with the simplest parts of the expression (the **terminal symbols**), and work outwards
- 1 is a constant, of type `int`
- ...
- So the type of this expression is `float`

# Dealing with Expressions.

```
private int recFactor () {  
    int ty;  
    if (have (Token.IdentifierToken)) {  
        ty = semantics.CheckId (scanner.CurrentToken);  
        mustBe (Token.IdentifierToken);  
    }  
    else if (have (Token.IntegerToken)) {  
        ty = LanguageType.Integer;  
        mustBe (Token.IntegerToken);  
    }  
    else {  
        mustBe ("(");  
        ty = recExpr();  
        mustBe (")");  
    }  
    return ty;  
}
```

# Dealing with Expressions.

```
private int recExpr () {  
    int lTy = recFactor ();  
    while (have("+") || have("-")) {  
        if (have("+")) mustBe ("+");  
        else           mustBe ("-");  
        int rTy = recFactor ();  
  
        // a+b and a-b require a and b to  
        // have the same type.  
        semantics.CheckTypesSame(scanner.CurrentToken,  
                                   lTy, rTy);  
    }  
    return lTy;  
}
```

(What would we need to do to handle  $a/b$  ?)



# Assignment Example.

- Consider the EBNF extracts

```
<Statement> ::= <Assign> | ... ;  
<Assign> ::= Ident = <Expr> ;
```

```
private void recAssign () {  
    int varTy = semantics.CheckId(scanner.CurrentToken);  
    mustBe (Token.IdentifierToken);  
    mustBe ("=");  
    IToken expToken = scanner.CurrentToken;  
    int expTy = recExpr();  
    checkTypesSame(expToken, varTy, expTy);  
}
```

# Assignment Example.

- Cons

Still void – there's no type information to return from an assignment (in this language!)

<Statement> ::= <Assign> | ... ;  
<Assign> ::= Ident = <Expr>

Save the token at this point, so we report the error at the right place in the source code...

```
private void recAssign () {  
    int varTy = semantics.CheckType(Token.IdentifierToken),  
    mustBe (Token.IdentifierToken),  
    mustBe ("=");  
    IToken expToken = scanner.CurrentToken;  
    int expTy = recExpr();  
    checkTypesSame(expToken, varTy, expTy);  
}
```

# Wider Use Of Expressions.

- Expressions are also used in other kinds of statements – e.g. `if/while` conditions.
- The semantic methods defined here can also be used in those contexts.
  - e.g. for a `while` loop's condition, we could parse the expression, and check that its type matches `LanguageType.Boolean`

# Other Uses Of Identifiers.

- User-Defined Names.
- Constants, arrays and Subroutines.
- Symbol Class Hierarchy.
- User-Defined Types.

# User-Defined Names.

- One of the main roles of semantic analysis is checking the use of user-defined names (or identifiers).
- To date we have looked at simple variables but there may be other uses of names
  - constants,
  - arrays,
  - functions,
  - user-defined types
  - etc. etc.

# Extending The *Symbol* class.

- The *Symbol* class must be extended to handle the different uses of identifiers.
  - Responsibility of this class is to store an identifier name with its attributes.
  - These different uses of names require modified attributes.
  - See the *Ardkit* toolkit provision of the subclasses to represent variables, constants, arrays and functions.

# Arrays.

- A language may allow an identifier to denote an array.
- An array is a fixed-size collection of locations, (usually) all of the same type.
- Arrays supported by the Ardkit *ArraySymbol* class may have multiple dimensions with specified bounds but must be regular.
- Languages might not need to treat arrays specially – functional languages generally don't

# Functions.

- A language might have functions, procedures or methods as components; we will use the generic term “function” for these.
- Functions are (usually) named blocks of code, with optional parameters, that may return a value.
- Must distinguish between requirements for
  - the function declaration (formal parameters)
  - the function call (actual parameters)



# Function Call.

- Assume the following function declaration in some language:

```
int sum (int x, int y)  
begin ... .. end
```

- What attributes are required to analyse the call?

```
i = sum (a, 5)
```

- Looking up **sum** in the symbol table, the following attributes will allow the call to be analysed.
  - type returned by the function,
  - number of parameters, and
  - type of each parameter
- Note that the **names** of the formal parameters are not required for the analysis of the call.

# Function Declaration.

- Consider the following function declaration in some language:

```
int sum (int x, int y)  
begin ... .. end
```

- What attributes are required to analyse the function body?
- Within the function body the code will be analysed in the context of the function name and the collection of formal parameters.
- We create a new scope for the function body, and put the formal parameters into that scope's symbol table.

```
begin ... .. end
```

# Ardkit Support for Functions.

- Refer to the Ardkit reference page for the *FunctionSymbol* class.
- Review the source code for the symbol classes' implementations.
- Note that the API is targetted at processing
  - the declaration of a function signature (setting the symbol's attributes),
  - the declaration of the function body (getting the formal parameter attributes), and
  - the calling of a function (getting the attributes).

# User-Defined Types.

- A language may allow user-defined types to be declared.
- Type-checking code needs to take these new types into account.
- In Ardkit, we'd need to create new symbol classes for the kinds of types that users can declare
  - e.g. a product/struct type would need to list the fields it contains with types and names...

Summary.

# We Have Seen How ...

- ... variables and literal constants need to be semantically analysed,
- ... expressions can be type-checked by passing type information between recogniser methods,
- ... identifiers can be used with different meanings with different semantic analysis requirements, and
- ... a symbol class hierarchy can represent these different meanings.