### Storage Systems

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## Storage hierarchy

Primary (Cache, Memory)

Secondary (Hard disks, SSD)

**Tertiary** (Tape, Optical Disk) Cost, Energy Consumption

#### Access times

LEVEL	ACCESS TIME	TYPICAL SIZE
Registers	"instantaneous"	under 1KB
Level 1 Cache	1-3 ns	64KB per core
Level 2 Cache	3-10 ns	256KB per core
Level 3 Cache	10-20 ns	2-20 MB per chip
Main Memory	30-60 ns	4-32 GB per system
Hard Disk	3,000,000-10,000,000 ns	over 1TB

Ssd: 25,000 ns

Resource: <a href="https://arstechnica.com/information-technology/2012/06/inside-the-ssd-revolution-how-solid-state-disks-really-work/">https://arstechnica.com/information-technology/2012/06/inside-the-ssd-revolution-how-solid-state-disks-really-work/</a>

#### **CRUD**

Basic functions of a storage device

#### • CRUD:

- (C)reate/write
- (R)ead
- (U)pdate/overwrite
- (D)elete

#### Characterizing a storage device

- Capacity (bytes)
  - How much data it can hold
- Cost (\$\$\$)
  - Price per byte of storage
- Bandwidth (bytes/sec)
  - Number of bytes that can be transferred per second
  - Note that read and write bandwidth may be different
- Latency (secs)
  - Time between initiating a request and an action
  - In the case of storage, action is to deliver 1<sup>st</sup> Byte

#### Time to complete an operation

- Time to complete an operation depends on both bandwidth and latency
  - CompletionTime = Latency + Size/Bandwidth
- The time of a workload will depend on
  - Technology, e.g., hard drive/ssd
  - Operation type, e.g., read/write
  - Number of operations in the workload
  - Access pattern (random vs. sequential)

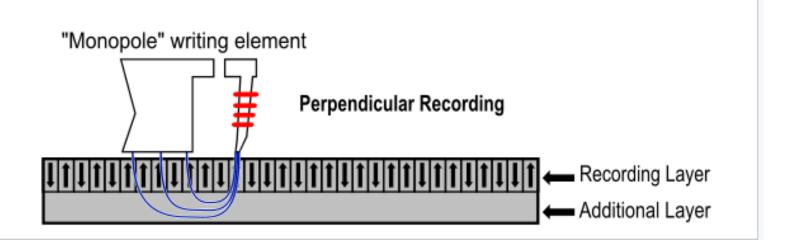
#### Access pattern

- Sequential
  - Data to be accessed are located next to each other or sequentially on the device

- Random
  - Access data located randomly on storage device

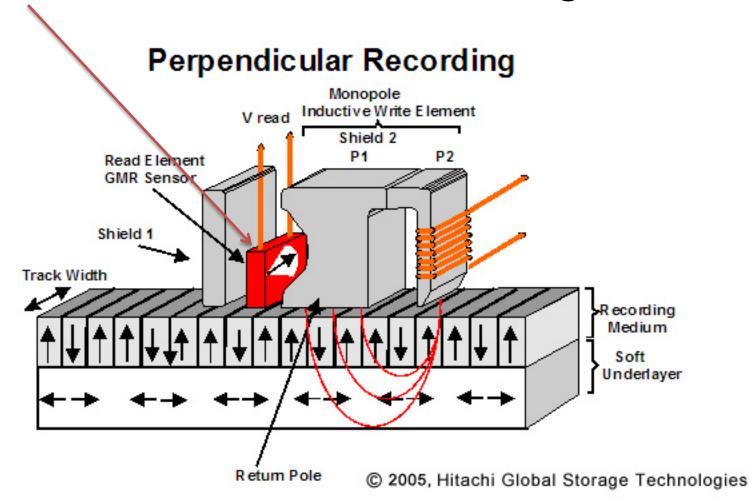
## Magnetic recording

- Write head
  - Applies electrical current to write head
  - Changes direction of magnetic field under head



## Reading

Read head senses direction of magnetic field



## Road map

Tapes



Hard disk

Solid state drive

#### Linear tape

 Data recorded on parallel tracks that span the length of the tape



#### **Tapes**

- Current technology is LTO
  - Linear Tape-Open (an open standard)
- Characteristics
  - Capacity up to 6.25 TB per tape (LTO-7)
  - Drive cost ~ \$2500
  - Tape cost ~ \$45 for 2.5TB tape
- Tape access time (~ minute)
  - Time to mount the tape
  - Time to wind the tape to correct position
- Data transmission rates ~ 250MB/sec



\$200 300MB/s LTO-7

#### Performance characteristics

- High latency/low cost makes tape most appropriate for "archival" storage
  - Low frequency of reads
  - Very large data objects

- Random access will be slow due to latency
  - Sequential reads will be fast

### Linear tape file system

- Two partitions on tape
  - First contains metadata and directories. Tape reader can find and load this very quickly
  - Second contains blocks for data

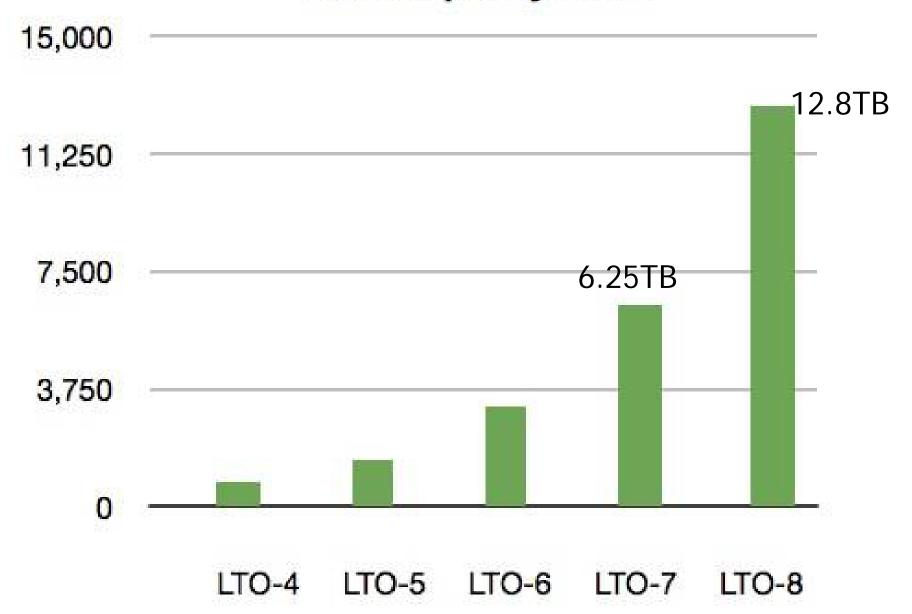
- Directory structure coded in XML
  - Self describing file format...



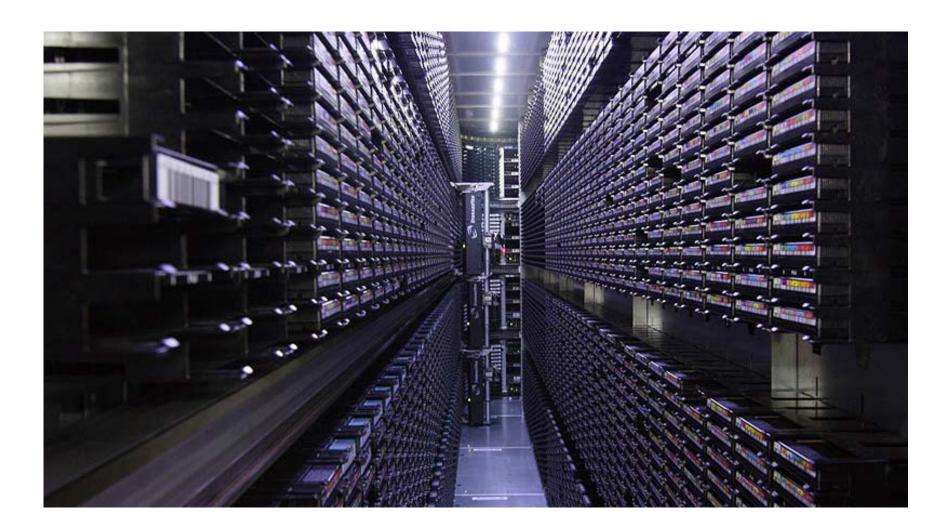
# Tape Cartridge



#### Raw Capacity in GB



# A tape library



## Inside a robotic tape library

https://www.youtube.com/watch?v=nYfTtvpQ
 778



## Road map

Tapes

Hard disk



Solid state drive

#### Hard disk drives

- Perhaps the most pervasive form of storage
- Basic Idea:
  - One or more spinning magnetic platters
    - Typically two surfaces per platter
  - Disk arm positions over the radial position (tracks) where data are stored
    - It swings across tracks (but do not extend/shrink)
  - Data is read/written by a read/write head as platter spins



WESTERN
DIGITAL
WD30EZRS
Caviar Green 3TB
64MB cache SATA
3.0Gb/s 3.5
internal hard
drive (Bare Drive)

About \$120 at Amazon

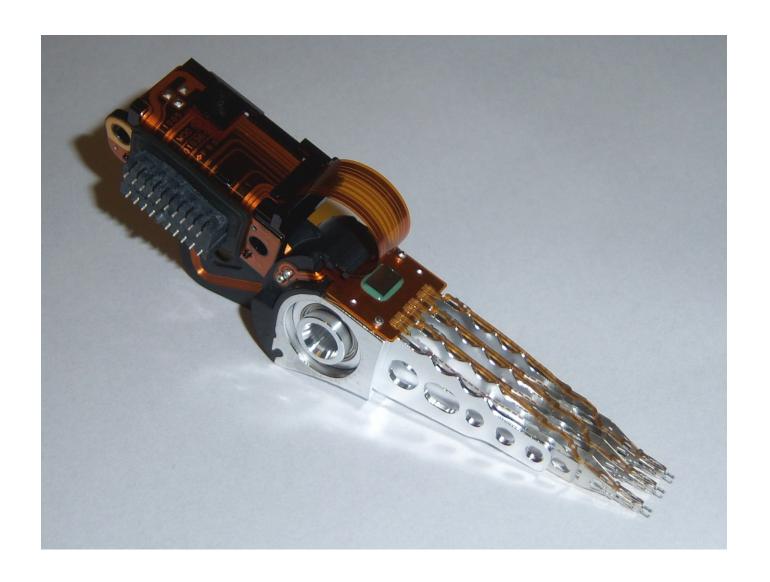
#### Internal of hard disk



# Disk arm and platter

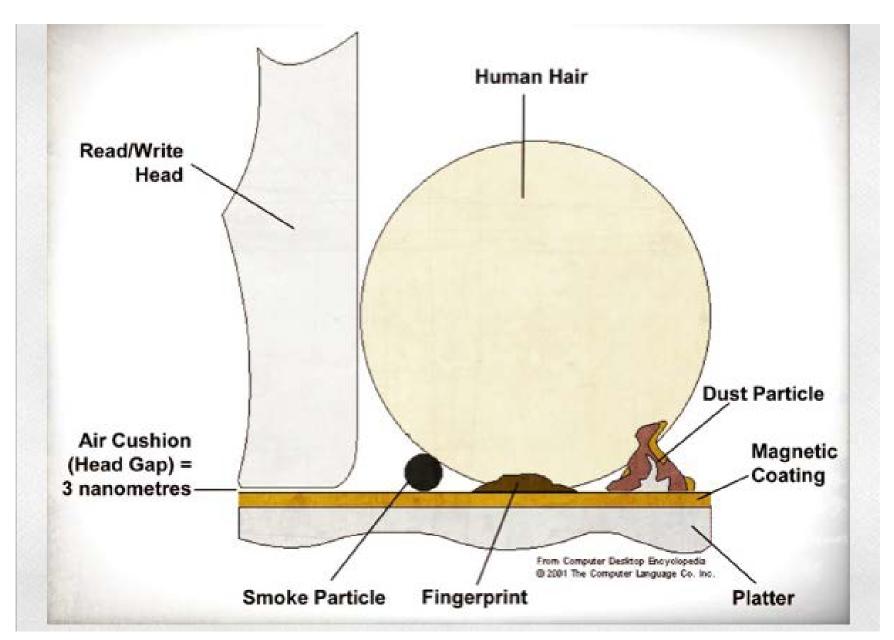


## Disk arm & head close-up



# Disk head close-ups





#### Disk head movement

- Hard disk head movement while copying files between two folders (e.g., partition c to d)
  - https://www.youtube.com/watch?v=BIB49F6ExkQ



# 2GB Storage in 1980s (\$250,000!)



## Physical characteristics

- 3.5" (diameter, common in desktops)
- 2.5" (common in laptops)
- Rotational speed
  - -4,800 RPM
  - 5,400 RPM
  - 7,200 RPM
  - 10,000 RPM
- Between 5-7 platters
- Current capacity up to 10TB (Western Digital)

#### Disk organization

Each platter consists of a number of tracks

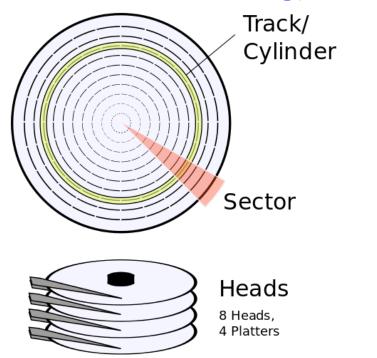
- Each track is divided into N fixed size sectors
  - Typical sector size is 512 bytes (old) or 4KB (new)
  - Sectors can be numbered from 0 to N-1
  - Entire sector is written "atomically"
    - All or nothing

## CHS (cylinder-head-sector)

Early way to address a sector

Now LBA (<u>Logical Block Addressing</u>) more

common



https://en.wikipedia.org/wiki/Cylinder-head-sector

## Example

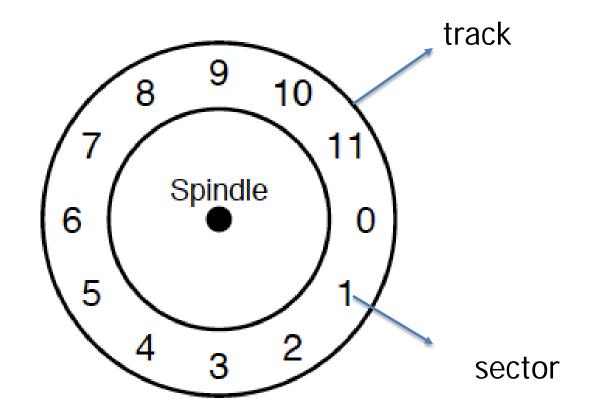
• # cylinders: 256

# heads: 16 (i.e., 8 platters, 2 heads/platter)

• # sectors/track: 64

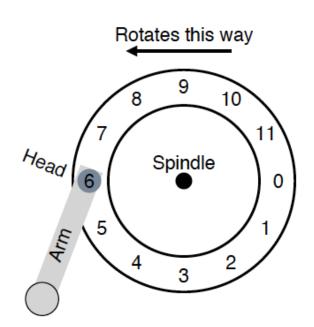
Sector size = 4KB

## A simple disk drive (one track only)



### Rotational latency

- Waiting for the right sector to rotate under the head
  - On average: about ½ of time of a full rotation
  - Worst case?
  - Best case?

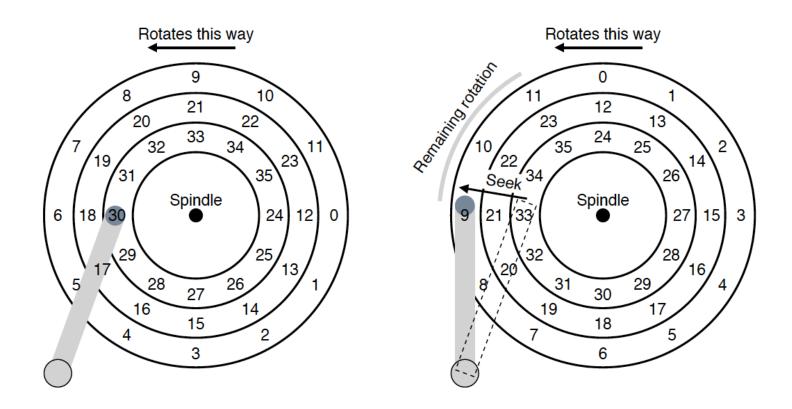


#### Rotation time

Assume 10,000 RPM (rotations per minute)

$$\frac{60000 \, ms}{1000 \, rotations} = \frac{6 \, ms}{rotation}$$

#### Multiple tracks: add seek times



Average seek time is about 1/3 max seek time (see reading: Chapter 37, page 9 for more details)

#### Transfer time

Assume that we transfer 512KB

Assume 128 MB/sec transmission bandwidth

• Transfer time:

512KB/128MB \* 1000ms = 4ms

#### Completion time

- $T = T_{seek} + T_{rotation} + T_{transfer}$ 
  - T<sub>seek</sub>: Time to get the disk head on right track
  - T<sub>rotation</sub>: Time to wait for the right sector to rotate under the head
  - T<sub>transfer</sub>: Time to actually transfer the data

#### Example

- Capacity 4TB
- # platters: 4
- # heads: 8
- Bytes per sector: 4096
- Transmission bandwidth: 100MB/sec
- Maximum seek time: 12ms
- RPM: 10,000

#### Time to transfer a file

- The file occupies 100 sectors (sequentially)
- Avg. seek time =?
- Avg. rotational latency =?
- Transfer time = ?

#### Sector vs. block

Block has 1 or more sectors

Disk typically transfers one block at a time

- We will assume one block = one sector
  - Unless stated otherwise

#### Sequential operations

- May assume all sectors involved are on same track
  - We may need to seek to the right track
  - And rotate to the first sector

But no rotation/seeking needed afterward

#### Actual bandwidth

- Consider a workload w
  - E.g., w = sequential access of 100 blocks of data
  - Denote size (# of bytes) of data in w as |w|
  - E.g., w = 400KB (100 blocks, 4KB/block)

Suppose completion time for w = t

Actual bandwidth (with respect to w) = |w|/t

#### Sequential vs. random

- Consider disk with 7ms avg seek, 10,000 RPM platter speed and 50 MB/sec transfer rate, 4KB/block
- Sequential access of 10 MB
  - Completion time = 7 + 3 + 10/50\*1000 = 210ms
  - Actual bandwidth = 10MB/210ms = 47.62 MB/s
- Random access of 10 MB (2,500 blocks)
  - Completion time = 2500 \* (7 + 3 + 4/50) = 25.2s
  - Actual bandwidth = 10MB / 25.2s = .397 MB/s

# Road map

Tapes

Hard disk

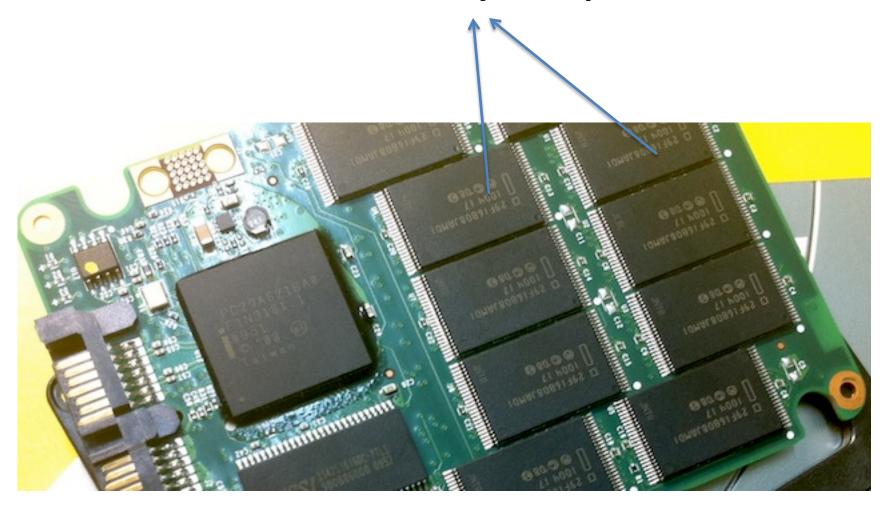
Solid state drive



#### **Solid State Drive**



# Memory chips



#### **Solid State Drives**

- All electronic, made from flash memory
- Lower energy consumption than hard drive
- Significantly more expensive, less capacity
  - About a factor of 5-10 more expensive
- Limited lifetime, can only write a limited number of times.
  - E.g., 100K program/write cycles for SLC (single-level cell) memory

# SSD vs. Hard Drive (price)

#### MLC (multi-level cell)



Samsung 850 EVO 1TB 2.5-Inch SATA III Internal SSD (MZ-75E1T0B/AM)

by Samsung

\$309<sup>99</sup> *Prime* 

Get it by Thursday, Jan 19

More Buying Choices \$309.99 new (80 offers) \$273.37 used (6 offers)



Hard Disk Size: 1 TB

Hardware Interface: sata 3 0 gb



Seagate 1TB Barracuda SATA 6GB/s 128MB Cache 2.5-Inch 7mm Internal Bare/OEM Hard Drive (ST1000LM048)

by Seagate

\$59<sup>22</sup> \$59.99 **\*Prime** | FREE Same-Day

Get it by TODAY, Jan 17

More Buying Choices \$58.24 new (27 offers)



Product Features

The thinnest and highest-capacity 2.5-Inch Hard Drive available

Electronics: See all 7,236 items

#### **Solid State Drives**

- Same form-factor and control interface as magnetic disks
- Significantly better latency
  - No seek or rotational delay
- Consistent bandwidth for sequential & random:
  - Benefits from improved latency
  - However, writes take significantly longer then reads

#### Writing to SSD is complicated

- Can not overwrite a page
  - Need to erase its block (at a certain point) instead

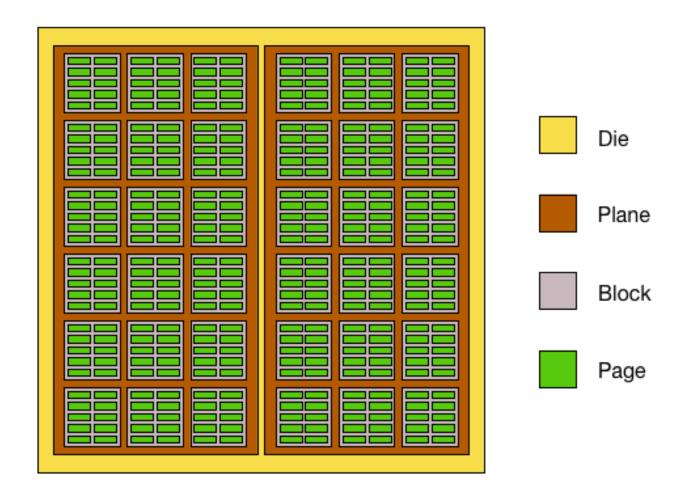
SSD controllers take care of all these details

#### SSD

- Contains a number of flash memory chips
  - Chip -> dies -> planes -> blocks -> pages (rows) -> cells
  - Cells are made of floating-gate transistors

- Page is the smallest unit of data transfer between SSD and main memory
  - Much like a block in hard disk

# Die Layout



#### Dies, planes, block, and pages

- Typically, a chip may have 1, 2, or 4 dies
- A die may have 1 or 2 planes
- A plane has a number of blocks
  - Block is the smallest unit that can be erased
- A block has a number of pages
  - Page is the smallest unit that can be read, programmed/written

#### Typical page and block sizes

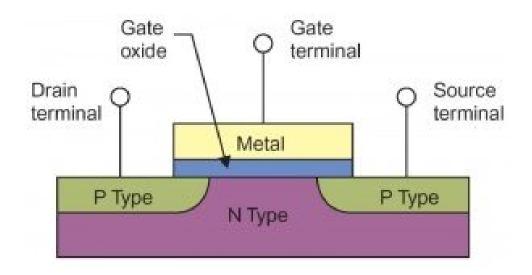
Common page sizes: 2K, 4K, 8K, and 16K

A block typically has 128 to 256 pages

=> Block size: 256KB to 4MB

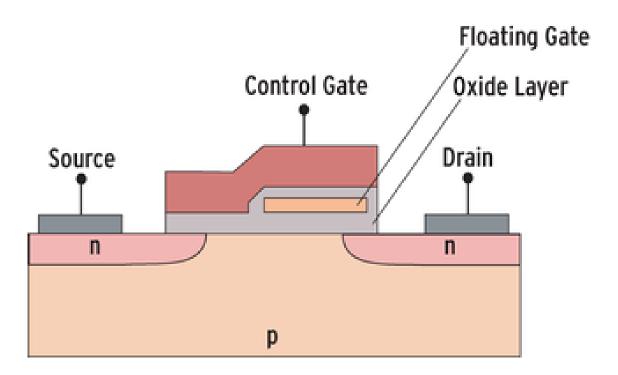
# Normal transistor (MOSFET)

- When voltage applied to gate terminal
  - semi-conducting region (purple) becomes conductive



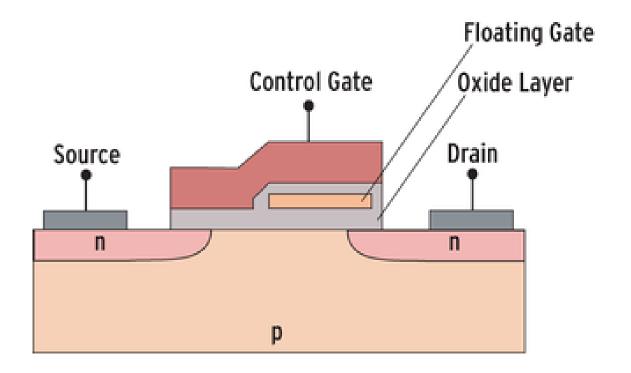
# Floating gate transistor

- Contain an additional gate: floating gate
  - Floating, since isolated by oxide layer
  - (thus not connected to other components)



# Floating gate transistor

 By applying high positive/negative voltage to control gate, electrons can be attracted to or repelled from floating gate



# Floating gate transistor

• State = 1, if no electrons in the floating gate

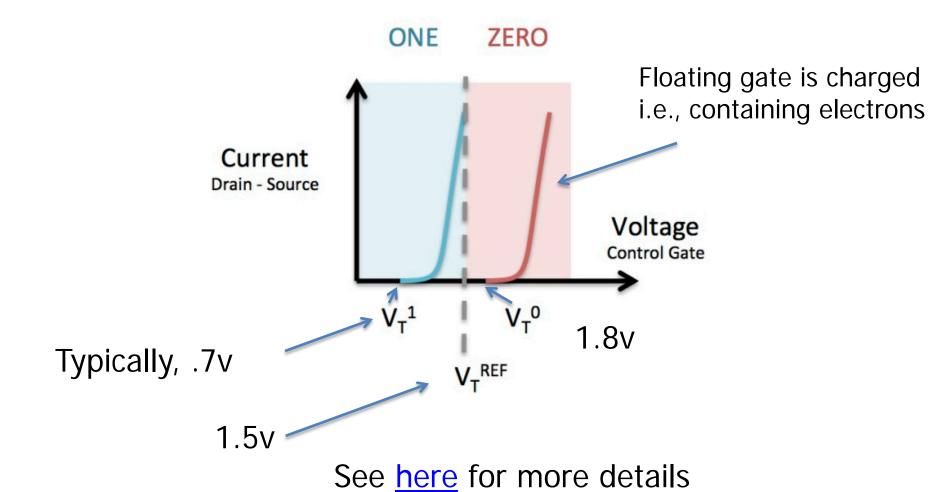
- State = 0, if there are electrons (negative charges)
  - Electrons stuck there even when power is off
  - So state is retained

#### Read operations

 Electrons on the floating gate affect the threshold voltage for the floating gate transistor to conduct

 Higher voltage needed when gate has electrons

#### Read operations



# Read operations

Apply V<sub>int</sub> (intermediate voltage)

If the current is detected, gate has no electrons

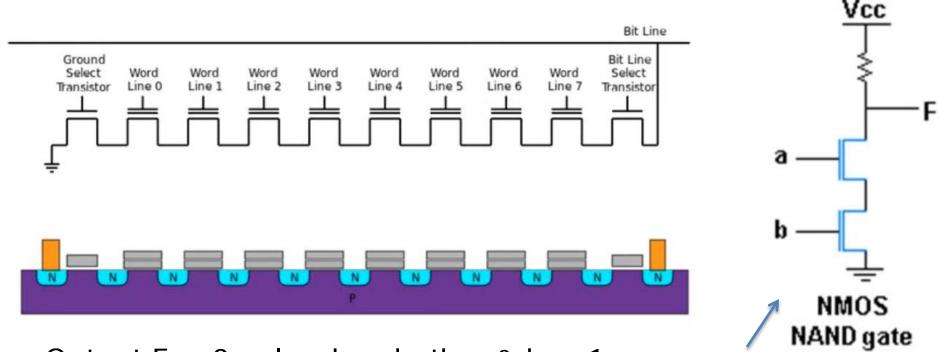
$$=>$$
 bit  $=$  1

If no current, gate must have electrons

$$=>$$
 bit  $=$  0

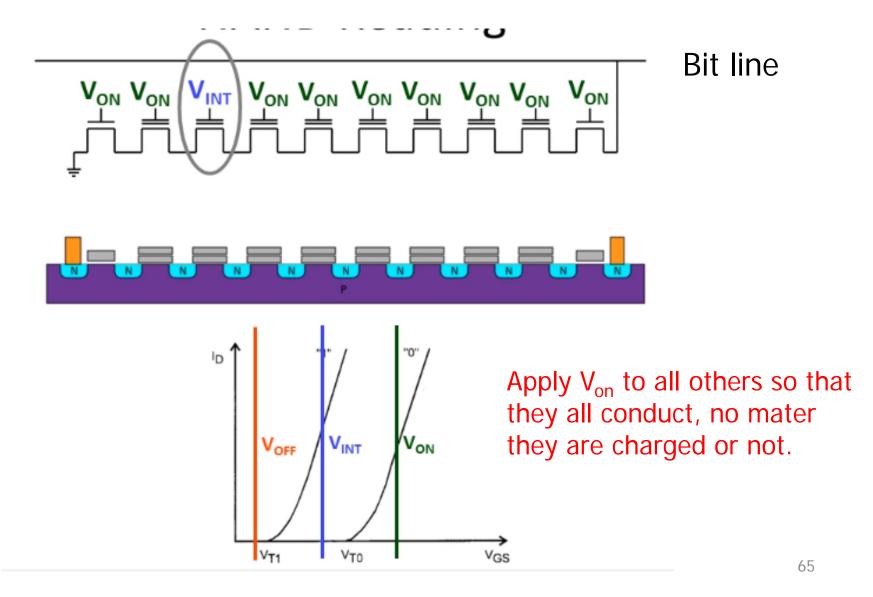
# NAND flash layout

- Transistors are strung together in a series
  - Similar to the transistors in an NAND gate



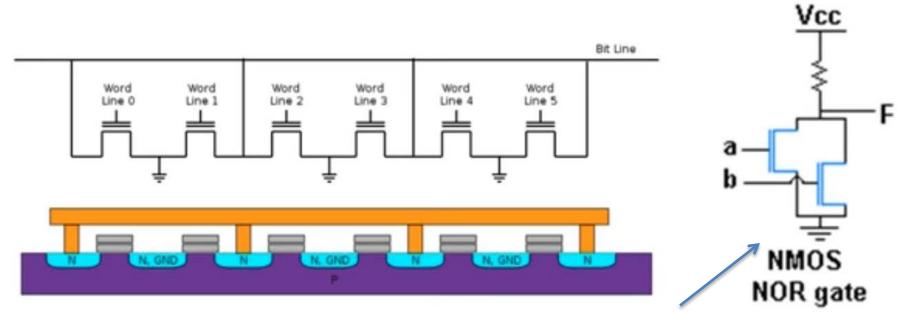
Output F = 0 only when both a & b = 1I.e., F will be grounded only when both a and b conduct

# NAND reading



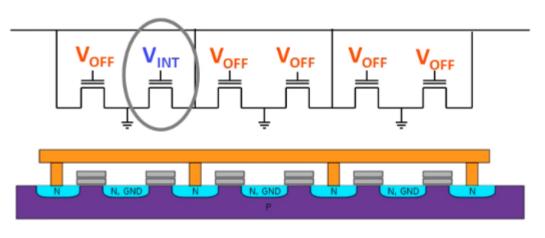
# NOR flash layout

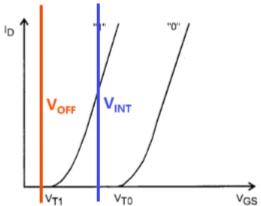
- Each transistor is directly connected to bit line (also ground)
  - Similar to transistors in a NOR gate (to output F)



F will be grounded as long as one of a & b conducts

#### NOR reading





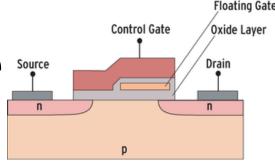
Apply  $V_{\text{OFF}}$  to all others so that they do NOT conduct, charged or not.

#### NAND vs NOR

- NOR has individual bit lines and ground connections
  - circuit more complex

- NAND does not have individual bit lines, nor separate ground connections
  - → save space, larger capacity

#### Write and erase source



- Write: 1 => 0
  - Apply high POSITVE voltage (>> voltage for read)
     to the control gate
  - Attract electrons from channel to floating gate (through quantum tunneling)
- Erase: 0 => 1
  - Need to apply much higher NEGATIVE voltage
  - Get rid of electrons from floating gate
  - May stress surrounding cells
  - So dangerous to do on individual pages

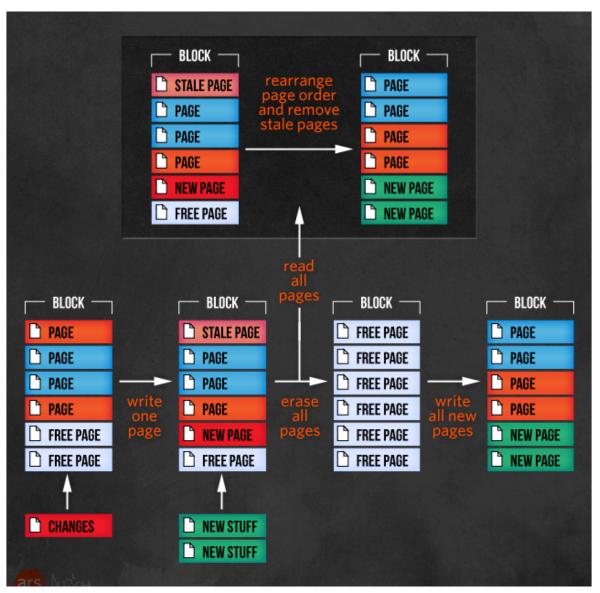
# Read/write units

- Page is the smallest unit for read and write (write is also called program, 1->0)
- Block is the smallest unit for erase (0->1)
  - i.e., make cells "empty" (i.e., no electrons)



Operation	Area
Read	Page
Program (Write)	Page
Erase	Block

# Example



# Latencies: read, write, and erase

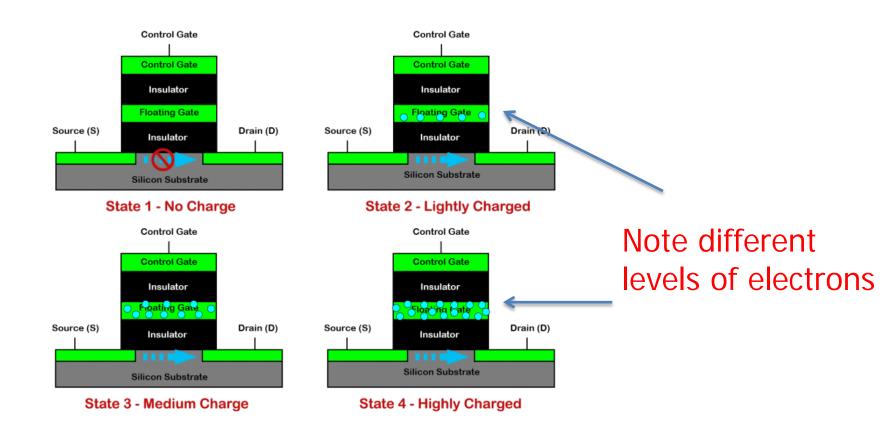
	SLC	MLC	TLC	HDD	RAM	
P/E cycles	100k	10k	5k	*	*	
Bits per cell	1	2	3	*	*	
Seek latency (µs)	*	*	*	9000	*	
Read latency (µs)	25	50	100	2000-7000	0.04-0.1	
Write latency (µs)	250	900	1500	2000-7000	0.04-0.1	
Erase latency (μs)	1500	3000	5000	*	*	
Notes	* metric is not applicable for that type of memory					
Sources	Sources  P/E cycles [20]  SLC/MLC latencies [1]  TLC latencies [23]  Hard disk drive latencies [18, 19, 25]  RAM latencies [30, 52]  L1 and L2 cache latencies [52]					

# P/E cycle

- P: program/write; E: erase
- Every write & erase damages oxide layer surrounding the floating-gate to some extent

- P/E cycle:
  - Data are written to cells (P): cell value from 1 -> 0
  - Then erased (E): 0 -> 1

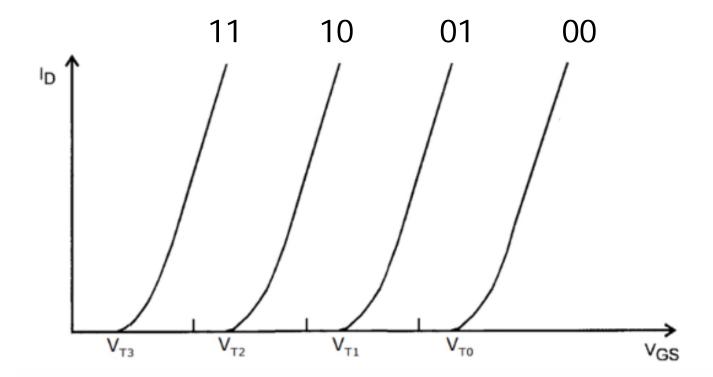
# Multi-level cell (MLC)



http://www.cactus-tech.com/resources/blog/details/solidstate-drive-primer-2-slc-mlc-and-tlc-nand-flash

# MLC reading

• 2 bits, 3 intermediate voltages



# SLC compared with MLC

#### • SLC:

- Less complex
- Faster
- More reliable
- Less storage
- More costly

#### Read more

Solid-state revolution: in-depth on how SSDs really work

- How do SSDs work?
  - http://www.extremetech.com/extreme/210492extremetech-explains-how-do-ssds-work

- How Flash Memory Works
  - https://www.youtube.com/watch?v=msi5GDz9JIw
- Floating Gate Basics
  - http://www.cse.scu.edu/~tschwarz/coen180/LN/fl ash.html
- Friend of Flash
  - http://www.nnc3.com/mags/LM10/Magazine/Arc hive/2008/86/040-041 logfs/article.html

- Understanding Flash: Floating Gates and Wear
  - https://flashdba.com/2015/01/09/understandingflash-floating-gates-and-wear/

- From Transistors to Functions
  - http://www.cs.bu.edu/~best/courses/modules/Tr ansistors2Gates/

- Solid State Drive Primer
  - https://www.cactustech.com/resources/blog/details/solid-state-driveprimer-1-the-basic-nand-flash-cell

- How Does a Transistor Work?
  - https://www.youtube.com/watch?v=IcrBqCFLHIY&feature=youtu.be

- How Flash Memory Works
  - https://www.youtube.com/watch?v=s7JLXs5es7I