CS420 - Lecture 9

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Spring 2023



• Always true: local operations appear to local thread to execute in program order

- Always true: local operations appear to local thread to execute in program order
- Sequential consistency: Operations appear to all threads to execute in same order (operations from different threads can interleave arbitrarily).

```
thread 0
---
x = 1
print(y)
```

```
thread 1
---
y = 1
print(x)
```

```
thread 0
---
x = 1
print(y, z)
```

```
thread 1
---
y = 1
print(x, z)
```

```
thread 2
---
z = 1
print(x, y)
```

```
thread 0
---
x = 1
print(y, z)
```

```
thread 1
---
y = 1
print(x, z)
```

```
thread 2
---
z = 1
print(x, y)
```

• 001011

```
thread 0 thread 1 thread 2 --- x = 1 y = 1 z = 1 print(y, z) print(x, y)
```

- 001011
- 001011 : t0, t0, t1, t1, t2, t2
- 101011

```
thread 0 thread 1 thread 2 --- x = 1 y = 1 z = 1 print(y, z) print(x, y)
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thread 0 thread 1 thread 2 --- x = 1 y = 1 z = 1 print(y, z) print(x, z) print(x, y)
```

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• 001011 : t0, t0, t1, t1, t2, t2

• 101011

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• 000011

Compilers

```
1 X = 0

2 for i in range(100):

3 X = 1

4 print X
```

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- Sequential consistency: Operations appear to all threads to execute in same order (operations from different threads can interleave arbitrarily).
- Most processors ARE NOT sequentially consistent
- User should write only race-free code: If two threads perform conflicting accesses to a shared variable, then the accesses must be explicitly ordered by an OpenMP synchronization operation
- The OpenMP compiler and runtime will make sure that properly synchronized accesses appear to occur in the right order

```
. . .
int i=1, j=3, k=4, m=5, n=6;
                                                       i=1; j=3; k=4; m=5; n=6
omp_set_num_threads(2);
#pragma omp parallel
                                                                                   fork
if(omp_get_thread_num()==0) {
                                                       i=7; m=j; k=8
i=7; m=j; k=8;
#pragma omp barrier
                                                                                   barrier
else {
                                                                    n=i; j=9; k=10
#pragma omp barrier
n=i; j=9; k=10;
                                                                                    join
                                                       printf
printf("i=%d, j=%d, k=%d, m=%d, n=%d \n", i, j, l
```

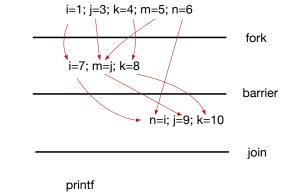
i=7; m=j; k=8

barrier

join

printf

output is: i=7, j=9, k=10, m=3, n=7



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Synchronization

- Ordering constructs: barrier, fork, join
- Mutual Exclusion constructs: atomic, critical, lock