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Batch: B2

ITC Tut 3.

(Q1) Arithmetic coding "PROOF".

(Q2) LZW coding "banana-badadaa".
↳ efficiency

(Q3) Run Length (RLR) 1111...000...001111
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15 20 4

Q1) Probability of  $P = \frac{1}{5} = 0.2$

Formula for d

= upper limit (ul)

Probability of  $R = \frac{1}{5} = 0.2$

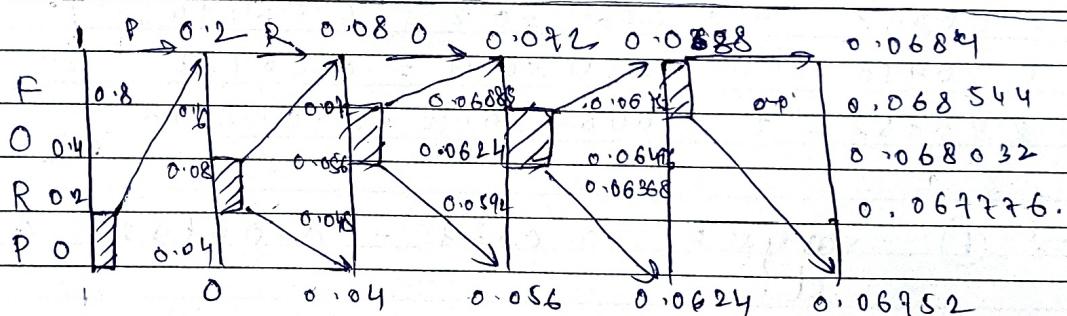
- lower limit (ll)

formula for Range

Probability of  $O = \frac{2}{5} = 0.4$

$R = ll + \alpha \times P(n)$

Probability of  $F = \frac{1}{5} = 0.2$



For (P) range =  $0.2 - 0 = 0.2$

$$0 + (0.2 - 0) \times 0.2 = 0.04$$

$$0.04 + (0.2) \times 0.2 = 0.08$$

$$0.08 + (0.2) \times (0.4) = 0.16$$

$$0.16 + (0.2) \times (0.2) = 0.2$$

For (R) range =  $0.08 - 0.04 = 0.04$

$$0.04 + 0.04 \times 0.2 = 0.048$$

$$0.048 + 0.04 \times 0.2 = 0.056$$

$$0.056 + 0.04 \times 0.4 = 0.072$$

$$0.072 + 0.04 \times 0.2 = 0.08$$

For (O) range  $\times$   $0.092 - 0.056 = 0.016$

$$0.056 + 0.016 \times 0.2 = 0.0592$$

$$0.0592 + 0.016 \times 0.2 = 0.06088$$

$$0.06088 + 0.016 \times 0.2 = 0.0624$$

$$0.0624 + 0.016 \times 0.4 = 0.0688$$

For (O) (2nd time) range  $\times$   $0.0688 - 0.0624 = 0.0064$

$$R = Ql + dxP(n)$$

$$0.0624 + 0.0064 \times 0.2 = 0.06368$$

$$0.06368 + 0.0064 \times 0.2 = 0.06496$$

$$0.06496 + 0.0064 \times 0.4 = 0.06752$$

$$0.06752 + 0.0064 \times 0.2 = 0.0688$$

For (F) range  $\times$   $0.0688 - 0.06752 = 0.00128$

$$0.06752 + 0.00128 \times 0.2 = 0.06776$$

$$0.06776 + 0.00128 \times 0.2 = 0.068032$$

$$0.068032 + 0.00128 \times 0.4 = 0.068544$$

$$0.068544 + 0.00128 \times 0.2 = 0.0688$$

Final Ans is lower value of last alphabet.

For (O) range {first time}  $= 0.092 - 0.056 = 0.036$

$$0.056 + 0.016 \times 0.2 = 0.0592$$

$$0.0592 + 0.016 \times 0.2 = 0.0624$$

$$0.0624 + 0.016 \times 0.4 = 0.0688$$

$$0.0688 + 0.016 \times 0.2 = 0.072$$

Ans 3)  $s = \underbrace{1111111111111}_{(15)} \underbrace{00000000000000000}_{(20)} \underbrace{001111}_{(4)}$

It can be coded using Run-length coding as.

$(11111)$ ,  $(10011, 0)$ ,  $(100, 1)$   
 (5) (6) (4)

compressing ratio is the ratio between the no. of bits required after encoding to the no. of bits required before encoding  
 i.e. for the original message

Here compression ratio =  $\frac{15}{89}$  ~~approx~~

$$\frac{15}{89} = 0.3846$$

Ans 2) "banana - badada"

Let dictionary be assigned as:-

b: 0

a: 1

n: 2

-: 3

d: 4

| Input           | Current Str. | Seen | Encoded                        | New dict |
|-----------------|--------------|------|--------------------------------|----------|
| b               | b            | Yes  | -                              |          |
| b a             | ba           | No   | 0                              | ba 5     |
| ba n            | an           | No   | 0,1                            | an 6     |
| ban a           | na           | No   | 0,1,2                          | na 7     |
| banana          | an           | Yes  | 0,1,2,6                        | -        |
| banana a        | na           | Yes  | 0,1,2,6,7                      | -        |
| banana -        | a-           | No   | 0,1,2,6,7,1                    | a- 8     |
| banana -b       | -b           | No   | 0,1,2,6,7,1,3                  | -b 9     |
| banana-b a      | ba           | Yes  | 0,1,2,6,7,1,3,5                | -        |
| banana-ba d     | ad           | No   | 0,1,2,6,7,1,3,5,1              | ad 10    |
| banana-bad a    | da           | No   | 0,1,2,6,7,1,3,5,1,4            | da 11    |
| banana-bada d   | ad           | Yes  | 0,1,2,6,7,1,3,5,1,4,10         | -        |
| banana-badad a  | da           | Yes  | 0,1,2,6,7,1,3,5,1,4,10,11      | -        |
| banana-badada a | aa           | No   | 0,1,2,6,7,1,3,5,1,4,10,11,1    | aa 12    |
| banana-badadaa  | -            | Yes  | 0,1,2,6,7,1,3,5,1,4,10,11,1,12 |          |

Ans - LZW of "banana-badadaa" is.

[0,1,2,6,7,1,3,5,1,4,10,11,1,12]

$$\text{Efficiency} = \frac{\text{Available}}{\text{Encoded}} = \frac{12}{14} = 0.8571 \\ = 85.71\%$$