

## ITC Tutorial 2

Q.1 Encode the message "CAFFEINE" using Shannon-Fano

Q.2 Compress the message "HEARTBEAT" using 1) Huffman encoding  
2) Shannon Fano coding  
Calculate and compare the code efficiency in both cases

Q.3 Define Kraft's inequality theorem. Check whether it is satisfying in questions 1 and 2

Q.1  $n_i$   $P(n_i)$   
word - "CAFFEINE"

$$\text{New } P(C) = \frac{1}{8}$$

$$P(A) = \frac{1}{8}$$

$$P(F) = \frac{1}{4}$$

$$P(E) = \frac{1}{4}$$

$$P(I) = \frac{1}{8}$$

$$P(N) = \frac{1}{8}$$

$x_i$	$P_{x_i}$	col1	col2	col3	code	codeword length
E	0.25	0	0		00	2
F	0.25	0	1		01	2
A	0.125	1	0	0	100	3
G	0.125	1	0	1	101	3
N	0.125	1	1	0	110	3
I	0.125	1	1	1	111	3

$$L = \sum P x_i$$

$$= 0.25 \times 2 + 0.25 \times 2 + 0.125 \times 3 + 0.125 \times 3 + 0.125 \times 3 + 0.125 \times 3$$

$$0.5 + 0.5 + 1.5$$

$$L = 2.5 \text{ bits / symbol}$$

$$H(x) = 0.25 \log(4) + 0.25 \log 4 + 0.125 \left( \log \left( \frac{1}{0.125} \right) \right) \times 4$$

$$H(x) = 1 \times 2.5$$

$$= 2.5 \text{ bits symbol}$$

$$\eta = \frac{H(x)}{L}$$

$$= \frac{2.5}{2.5}$$

$$\eta = 1$$

$$\eta \% = 100\%$$

$\therefore$  The efficiency is 100%.



Q.2 b) word = "HEARTBEAT" - By method shannon fano coding

$$P(H) = \frac{1}{9}$$

$$P(R) = \frac{1}{9}$$

$$P(E) = \frac{2}{9}$$

$$P(B) = \frac{1}{9}$$

$$P(A) = \frac{2}{9}$$

$$P(T) = \frac{2}{9}$$

$x_i$	$P_{x_i}$	col 1	col 2	col 3	code	codeword length
A	0.22	0	0	0	00	2
E	0.22	0	1	1	01	2
T	0.22	1	0	0	100	3
B	0.11	1	0	1	101	3
H	0.11	1	1	0	110	3
R	0.11	1	1	1	111	3

$$\bar{L} = \sum P x_i$$

$$0.22 \times 2 + 0.22 \times 2 + 0.22 \times 3 + 0.11 \times 3 + 0.11 \times 3 + 0.11 \times 3$$

$$= 0.44 + 0.44 + 0.66 + 0.33 + 0.33 + 0.33$$

$$\bar{L} = 2.53 \sim 2.55$$

$$H = - \sum_{i=1}^6 P_{x_i} \log \left( \frac{1}{P_{x_i}} \right)$$

$$0.22 \log \left( \frac{1}{0.22} \right) + 0.22 \log \left( \frac{1}{0.22} \right) + 0.22 \log \left( \frac{1}{0.22} \right)$$

$$+ 0.11 \log \left( \frac{1}{0.11} \right) + 0.11 \log \left( \frac{1}{0.11} \right) + 0.11 \log \left( \frac{1}{0.11} \right)$$

$$3 = 2.49$$

$$\eta = \frac{H}{L}$$

$$\eta = \frac{2.49}{2.55}$$

$$= 0.976$$

$$\eta \% = 97.6\% \text{ efficiency}$$



Q.2 a) By haffmaan encoding

$x_i$        $P_{x_i}$       code      codeword length

E      0.22

A      0.22

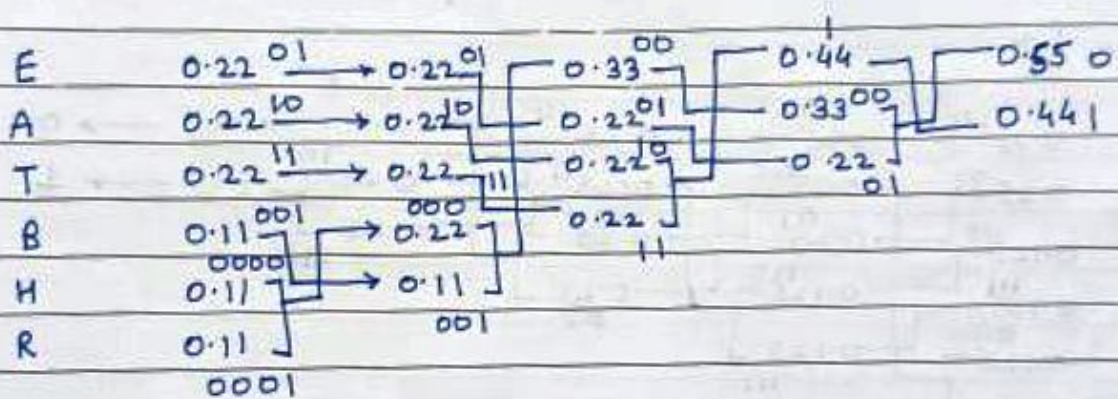
T      0.22

B      0.11

H      0.11

R      0.11

$x_i$        $P_{x_i}$



$x_i$        $P_{x_i}$       code      codeword length

E      0.22      01      2

A      0.22      10      2

T      0.22      11      2

B      0.11      001      3

H      0.11      0000      4

R      0.11      0001      4

$$H(x) = \sum_{i=1}^6 P_{x_i} \log_2 \left( \frac{1}{P_{x_i}} \right)$$

$$0.22 \times 2 + 0.22 \times 2 + 0.22$$

$$0.22 \log \left( \frac{1}{0.22} \right) + 0.22 \log \left( \frac{1}{0.22} \right) + 0.22 \log \left( \frac{1}{0.22} \right)$$

$$+ 0.11 \log \left( \frac{1}{0.11} \right) + 0.11 \log \left( \frac{1}{0.11} \right) + 0.11 \log \left( \frac{1}{0.11} \right)$$

$$= 2.49$$

$$\bar{L} = \frac{0.22 \times 2 + 0.22 \times 2 + 0.22 \times 2 + 0.11 \times 3 + 0.11 \times 4}{0.11 \times 4}$$

$$\frac{0.44 \times 3 + 0.33 + 0.44 + 0.44}{1.32 + 0.77 + 0.44}$$

$$1.32 + 0.77 + 0.44$$

$$\bar{L} = 2.53$$

$$\therefore \eta = \frac{2.49}{2.53}$$

$$0.984$$

$$\therefore \eta = 98.4\%$$

$\therefore$  The efficiency is 98.4%.



### 3) Kraft's inequality condition

Kraft's inequality is a necessary and sufficient condition to prove existence of prefix code (uniquely decodable code) given  $n$  symbols and  $l_i$  are no of bits used to represent a symbol for all  $i=1$  to  $n$

$$L = \sum_{i=1}^n 2^{-l_i} \leq 1$$

$l_i$  = codeword length

For Q1 - (Question 1)

$x_i$	$P_{x_i}$	codeword length
E	0.25	2
F	0.25	2
A	0.125	3
C	0.125	3
N	0.125	3
I	0.125	3

$$\text{formula} = \sum_{i=1}^n 2^{-l_i} \leq 1$$

$$\therefore \frac{1}{2^{12}} + \frac{1}{2^2} + \frac{1}{2^3} + \frac{1}{2^3} + \frac{1}{2^3} + \frac{1}{2^3}$$

$$= \frac{1}{4} + \frac{1}{4} + \frac{1}{8} + \frac{1}{8} + \frac{1}{8} + \frac{1}{8}$$

$$= \frac{1}{2} + \frac{1}{2}$$

$$= 1$$

$$\therefore \sum_{i=1}^n 2^{-l_i} = 1$$

$\therefore$  condition satisfied

$\therefore$  Q.1 follows Kraft's inequality condition

for Q.2a)

$n_i$	$P_{n_i}$	codeword length
E	0.22	2
A	0.22	2
T	0.22	2
B	0.11	3
H	0.11	4
R	0.11	4

$$\text{formula} = \sum_{i=1}^n 2^{-l_i} \leq 1$$

$$\frac{1}{2^2} + \frac{1}{2^2} + \frac{1}{2^2} + \frac{1}{2^3} + \frac{1}{2^4} + \frac{1}{2^4}$$

$$= \frac{1}{4} + \frac{1}{4} + \frac{1}{4} + \frac{1}{8} + \frac{1}{16} + \frac{1}{16}$$

$$= \frac{3}{4} + \frac{1}{8} + \frac{1}{8}$$

$$= \frac{3}{4} + \frac{1}{4}$$

$$= 1$$

$$\therefore \sum_{i=1}^n 2^{-l_i} = 1$$

$\therefore$  Q.2a) follows Kraft's inequality condition



For Q.2 b)

$x_i$	$P_{x_i}$	codeword Length
A	0.22	2
E	0.22	2
T	0.22	3
B	0.11	3
H	0.11	3
R	0.11	3

$$\therefore \text{formula} = \sum_{i=1}^n 2^{-l_i} \leq 1$$

$$\therefore \frac{1}{4} + \frac{1}{4} + \frac{1}{8} + \frac{1}{8} + \frac{1}{8} + \frac{1}{8}$$

$$= \frac{1}{2} + \frac{1}{4} + \frac{1}{4}$$

$$= \frac{1}{2} + \frac{1}{2}$$

$$= 1$$

$$\therefore \sum_{i=1}^n 2^{-l_i} = 1$$

$\therefore$  Q.2 b) follows Kraft's inequality condition