# **54**

# POSIX SHARED MEMORY

In previous chapters, we looked at two techniques that allow unrelated processes to share memory regions in order to perform IPC: System V shared memory (Chapter 48) and shared file mappings (Section 49.4.2). Both of these techniques have potential drawbacks:

- The System V shared memory model, which uses keys and identifiers, is not consistent with the standard UNIX I/O model, which uses filenames and descriptors. This difference means that we require an entirely new set of system calls and commands for working with System V shared memory segments.
- Using a shared file mapping for IPC requires the creation of a disk file, even if we are not interested in having a persistent backing store for the shared region. Aside from the inconvenience of needing to create the file, this technique incurs some file I/O overhead.

Because of these drawbacks, POSIX.1b defined a new shared memory API: POSIX shared memory, which is the subject of this chapter.

POSIX talks about shared memory *objects*, while System V talks about shared memory *segments*. These differences in terminology are historical—both terms are used for referring to regions of memory shared between processes.

### 54.1 **Overview**

POSIX shared memory allows to us to share a mapped region between unrelated processes without needing to create a corresponding mapped file. POSIX shared memory is supported on Linux since kernel 2.4.

SUSv3 doesn't specify any of the details of how POSIX shared memory is to be implemented. In particular, there is no requirement for the use of a (real or virtual) file system to identify shared memory objects, although many UNIX implementations do employ a file system for this purpose. Some UNIX implementations create the names for shared memory objects as files in a special location in the standard file system. Linux uses a dedicated *tmpfs* file system (Section 14.10) mounted under the directory /dev/shm. This file system has kernel persistence—the shared memory objects that it contains will persist even if no process currently has them open, but they will be lost if the system is shut down.

> The total amount of memory in all POSIX shared memory regions on the system is limited by the size of the underlying tmpfs file system. This file system is typically mounted at boot time with some default size (e.g., 256 MB). If necessary, the superuser can change the size of the file system by remounting it using the command mount -o remount, size=<num-bytes>.

To use a POSIX shared memory object, we perform two steps:

- Use the *shm\_open()* function to open an object with a specified name. (We described the rules governing the naming of POSIX shared memory objects in Section 51.1.) The  $shm\_open()$  function is analogous to the open() system call. It either creates a new shared memory object or opens an existing object. As its function result, *shm\_open()* returns a file descriptor referring to the object.
- 2. Pass the file descriptor obtained in the previous step in a call to mmap() that specifies MAP SHARED in the flags argument. This maps the shared memory object into the process's virtual address space. As with other uses of mmap(), once we have mapped the object, we can close the file descriptor without affecting the mapping. However, we may need to keep the file descriptor open for subsequent use in calls to *fstat()* and *ftruncate()* (see Section 54.2).

The relationship between  $shm\_open()$  and mmap() for POSIX shared memory is analogous to that between *shmget()* and *shmat()* for System V shared memory. The origin of the two-step process (shm\_open() plus mmap()) for using POSIX shared memory objects instead of the use of a single function that performs both tasks is historical. When the POSIX committee added this feature, the mmap() call already existed ([Stevens, 1999]). In effect, all that we are doing is replacing calls to open() with calls to  $shm_open()$ , with the difference that using *shm\_open()* doesn't require the creation of a file in a disk-based file system.

Since a shared memory object is referred to using a file descriptor, we can usefully employ various file descriptor system calls already defined in the UNIX system (e.g., ftruncate()), rather than needing new special-purpose system calls (as is required for System V shared memory).

### 54.2 **Creating Shared Memory Objects**

The shm\_open() function creates and opens a new shared memory object or opens an existing object. The arguments to  $shm\_open()$  are analogous to those for open().

```
/* Defines 0 * constants */
#include <fcntl.h>
                             /* Defines mode constants */
#include <sys/stat.h>
#include <sys/mman.h>
int shm_open(const char *name, int oflag, mode t mode);
                               Returns file descriptor on success, or -1 on error
```

The *name* argument identifies the shared memory object to be created or opened. The oflag argument is a mask of bits that modify the behavior of the call. The values that can be included in this mask are summarized in Table 54-1.

**Table 54-1:** Bit values for the *shm\_open()* of lag argument

Flag	Description
O_CREAT	Create object if it doesn't already exist
0_EXCL	With 0_CREAT, create object exclusively
O_RDONLY	Open for read-only access
O_RDWR	Open for read-write access
O_TRUNC	Truncate object to zero length

One of the purposes of the *oflag* argument is to determine whether we are opening an existing shared memory object or creating and opening a new object. If oflag doesn't include O\_CREAT, we are opening an existing object. If O\_CREAT is specified, then the object is created if it doesn't already exist. Specifying 0 EXCL in conjunction with 0 CREAT is a request to ensure that the caller is the creator of the object; if the object already exists, an error results (EEXIST).

The oflag argument also indicates the kind of access that the calling process will make to the shared memory object, by specifying exactly one of the values 0 RDONLY or O RDWR.

The remaining flag value, 0 TRUNC, causes a successful open of an existing shared memory object to truncate the object to a length of zero.

> On Linux, truncation occurs even on a read-only open. However, SUSv3 says that results of using 0\_TRUNC with a read-only open is undefined, so we can't portably rely on a specific behavior in this case.

When a new shared memory object is created, its ownership and group ownership are taken from the effective user and group IDs of the process calling shm\_open(), and the object permissions are set according to the value supplied in the *mode* bitmask argument. The bit values for *mode* are the same as for files (Table 15-4, on page 295). As with the open() system call, the permissions mask in mode is masked against the process umask (Section 15.4.6). Unlike open(), the mode argument is always required for a call to shm\_open(); if we are not creating a new object, this argument should be specified as 0.

The close-on-exec flag (FD CLOEXEC, Section 27.4) is set on the file descriptor returned by *shm\_open()*, so that the file descriptor is automatically closed if the process performs an exec(). (This is consistent with the fact that mappings are unmapped when an *exec()* is performed.)

When a new shared memory object is created, it initially has zero length. This means that, after creating a new shared memory object, we normally call ftruncate() (Section 5.8) to set the size of the object before calling mmap(). Following the mmap() call, we may also use ftruncate() to expand or shrink the shared memory object as desired, bearing in mind the points discussed in Section 49.4.3.

When a shared memory object is extended, the newly added bytes are automatically initialized to 0.

At any point, we can apply *fstat()* (Section 15.1) to the file descriptor returned by shm\_open() in order to obtain a stat structure whose fields contain information about the shared memory object, including its size (st\_size), permissions (st\_mode), owner (st\_uid), and group (st\_gid). (These are the only fields that SUSv3 requires fstat() to set in the stat structure, although Linux also returns meaningful information in the time fields, as well as various other less useful information in the remaining fields.)

The permissions and ownership of a shared memory object can be changed using *fchmod()* and *fchown()*, respectively.

# **Example program**

Listing 54-1 provides a simple example of the use of shm\_open(), ftruncate(), and mmap(). This program creates a shared memory object whose size is specified by a command-line argument, and maps the object into the process's virtual address space. (The mapping step is redundant, since we don't actually do anything with the shared memory, but it serves to demonstrate the use of mmap().) The program permits the use of command-line options to select flags (O\_CREAT and O\_EXCL) for the  $shm\_open()$  call.

In the following example, we use this program to create a 10,000-byte shared memory object, and then use ls to show this object in /dev/shm:

```
$ ./pshm create -c /demo shm 10000
$ 1s -1 /dev/shm
total 0
-rw----
             1 mtk
                                    10000 Jun 20 11:31 demo_shm
                        users
```

Listing 54-1: Creating a POSIX shared memory object

pshm/pshm create.c

```
#include <sys/stat.h>
#include <fcntl.h>
#include <sys/mman.h>
#include "tlpi hdr.h"
```

```
static void
usageError(const char *progName)
    fprintf(stderr, "Usage: %s [-cx] name size [octal-perms]\n", progName);
    fprintf(stderr, " -c Create shared memory (O_CREAT)\n");
fprintf(stderr, " -x Create exclusively (O EXCL)\n"):
    exit(EXIT FAILURE);
}
int
main(int argc, char *argv[])
    int flags, opt, fd;
    mode t perms;
    size t size;
    void *addr;
    flags = 0 RDWR;
    while ((opt = getopt(argc, argv, "cx")) != -1) {
        switch (opt) {
        case 'c': flags |= 0_CREAT;
                                                  break;
        case 'x': flags |= 0 EXCL;
                                                  break;
        default:
                    usageError(argv[0]);
        }
    }
    if (optind + 1 >= argc)
        usageError(argv[0]);
    size = getLong(argv[optind + 1], GN ANY BASE, "size");
    perms = (argc <= optind + 2) ? (S_IRUSR | S_IWUSR) :</pre>
                getLong(argv[optind + 2], GN_BASE_8, "octal-perms");
    /* Create shared memory object and set its size */
    fd = shm_open(argv[optind], flags, perms);
    if (fd == -1)
        errExit("shm open");
    if (ftruncate(fd, size) == -1)
        errExit("ftruncate");
    /* Map shared memory object */
    addr = mmap(NULL, size, PROT READ | PROT WRITE, MAP SHARED, fd, 0);
    if (addr == MAP FAILED)
        errExit("mmap");
    exit(EXIT_SUCCESS);
}
                                                                    — pshm/pshm_create.c
```

### 54.3 **Using Shared Memory Objects**

Listing 54-2 and Listing 54-3 demonstrate the use of a shared memory object to transfer data from one process to another. The program in Listing 54-2 copies the string contained in its second command-line argument into the existing shared memory object named in its first command-line argument. Before mapping the object and performing the copy, the program uses ftruncate() to resize the shared memory object to be the same length as the string that is to be copied.

Listing 54-2: Copying data into a POSIX shared memory object

```
pshm/pshm_write.c
#include <fcntl.h>
#include <sys/mman.h>
#include "tlpi_hdr.h"
main(int argc, char *argv[])
{
    int fd;
                                /* Size of shared memory object */
    size_t len;
   char *addr;
    if (argc != 3 || strcmp(argv[1], "--help") == 0)
        usageErr("%s shm-name string\n", argv[0]);
    fd = shm_open(argv[1], 0_RDWR, 0);
                                           /* Open existing object */
    if (fd == -1)
       errExit("shm_open");
   len = strlen(argv[2]);
    if (ftruncate(fd, len) == -1)
                                            /* Resize object to hold string */
        errExit("ftruncate");
   printf("Resized to %ld bytes\n", (long) len);
    addr = mmap(NULL, len, PROT READ | PROT WRITE, MAP SHARED, fd, 0);
    if (addr == MAP FAILED)
        errExit("mmap");
    if (close(fd) == -1)
                                            /* 'fd' is no longer needed */
        errExit("close");
   printf("copying %ld bytes\n", (long) len);
                                            /* Copy string to shared memory */
   memcpy(addr, argv[2], len);
   exit(EXIT_SUCCESS);
}
                                                                  pshm/pshm_write.c
```

The program in Listing 54-3 displays the string in the existing shared memory object named in its command-line argument on standard output. After calling shm\_open(), the program uses fstat() to determine the size of the shared memory and uses that size in the call to mmap() that maps the object and in the write() call that prints the string.

```
pshm/pshm_read.c
```

```
#include <fcntl.h>
#include <sys/mman.h>
#include <sys/stat.h>
#include "tlpi hdr.h"
int
main(int argc, char *argv[])
    int fd;
    char *addr;
    struct stat sb;
    if (argc != 2 || strcmp(argv[1], "--help") == 0)
        usageErr("%s shm-name\n", argv[0]);
    fd = shm_open(argv[1], 0_RDONLY, 0);  /* Open existing object */
    if (fd == -1)
        errExit("shm_open");
    /* Use shared memory object size as length argument for mmap()
       and as number of bytes to write() */
    if (fstat(fd, &sb) == -1)
        errExit("fstat");
    addr = mmap(NULL, sb.st_size, PROT_READ, MAP_SHARED, fd, 0);
    if (addr == MAP FAILED)
        errExit("mmap");
    if (close(fd) == -1);
                                            /* 'fd' is no longer needed */
        errExit("close");
    write(STDOUT_FILENO, addr, sb.st_size);
    printf("\n");
    exit(EXIT_SUCCESS);
}
```

pshm/pshm read.c

The following shell session demonstrates the use of the programs in Listing 54-2 and Listing 54-3. We first create a zero-length shared memory object using the program in Listing 54-1.

```
$ ./pshm create -c /demo shm 0
$ 1s -1 /dev/shm
                                         Check the size of object
total 4
                                 O Jun 21 13:33 demo_shm
-rw----
              1 mtk
                       users
```

We then use the program in Listing 54-2 to copy a string into the shared memory object:

```
$ ./pshm_write /demo_shm 'hello'
                                         Check that object has changed in size
$ 1s -1 /dev/shm
total 4
                                 5 Jun 21 13:33 demo_shm
-rw----
              1 mtk
                       users
```

From the output, we can see that the program resized the shared memory object so that it is large enough to hold the specified string.

Finally, we use the program in Listing 54-3 to display the string in the shared memory object:

```
$ ./pshm_read /demo_shm
hello
```

Applications must typically use some synchronization technique to allow processes to coordinate their access to shared memory. In the example shell session shown here, the coordination was provided by the user running the programs one after the other. Typically, applications would instead use a synchronization primitive (e.g., semaphores) to coordinate access to a shared memory object.

# 54.4 Removing Shared Memory Objects

SUSv3 requires that POSIX shared memory objects have at least kernel persistence; that is, they continue to exist until they are explicitly removed or the system is rebooted. When a shared memory object is no longer required, it should be removed using  $shm\_unlink()$ .

```
#include <sys/mman.h>
int shm_unlink(const char *name);
                                            Returns 0 on success, or -1 on error
```

The *shm\_unlink()* function removes the shared memory object specified by *name*. Removing a shared memory object doesn't affect existing mappings of the object (which will remain in effect until the corresponding processes call munmap() or terminate), but prevents further shm\_open() calls from opening the object. Once all processes have unmapped the object, the object is removed, and its contents are lost.

The program in Listing 54-4 uses shm\_unlink() to remove the shared memory object specified in the program's command-line argument.

**Listing 54-4:** Using *shm\_unlink()* to unlink a POSIX shared memory object

```
pshm/pshm_unlink.c
#include <fcntl.h>
#include <sys/mman.h>
#include "tlpi_hdr.h"
main(int argc, char *argv[])
    if (argc != 2 || strcmp(argv[1], "--help") == 0)
       usageErr("%s shm-name\n", argv[0]);
    if (shm_unlink(argv[1]) == -1)
        errExit("shm_unlink");
    exit(EXIT_SUCCESS);
}
                                                                   pshm/pshm_unlink.c
```

### 54.5 **Comparisons Between Shared Memory APIs**

By now, we have considered a number of different techniques for sharing memory regions between unrelated processes:

- System V shared memory (Chapter 48);
- shared file mappings (Section 49.4.2); and
- POSIX shared memory objects (the subject of this chapter).

Many of the points that we make in this section are also relevant for shared anonymous mappings (Section 49.7), which are used for sharing memory between processes that are related via *fork()*.

A number of points apply to all of these techniques:

- They provide fast IPC, and applications typically must use a semaphore (or other synchronization primitive) to synchronize access to the shared region.
- Once the shared memory region has been mapped into the process's virtual address space, it looks just like any other part of the process's memory space.
- The system places the shared memory regions within the process virtual address space in a similar manner. We outlined this placement while describing System V shared memory in Section 48.5. The Linux-specific /proc/PID/maps file lists information about all types of shared memory regions.
- Assuming that we don't attempt to map a shared memory region at a fixed address, we should ensure that all references to locations in the region are calculated as offsets (rather than pointers), since the region may be located at different virtual addresses within different processes (Section 48.6).
- The functions described in Chapter 50 that operate on regions of virtual memory can be applied to shared memory regions created using any of these techniques.

There are also a few notable differences between the techniques for shared memory:

- The fact that the contents of a shared file mapping are synchronized with the underlying mapped file means that the data stored in a shared memory region can persist across system restarts.
- System V and POSIX shared memory use different mechanisms to identify and refer to a shared memory object. System V uses its own scheme of keys and identifiers, which doesn't fit with the standard UNIX I/O model and requires separate system calls (e.g., shmctl()) and commands (ipcs and ipcrm). By contrast, POSIX shared memory employs names and file descriptors, and consequently shared memory objects can be examined and manipulated using a variety of existing UNIX system calls (e.g., fstat() and fchmod()).
- The size of a System V shared memory segment is fixed at the time of creation (via *shmget*()). By contrast, for a mapping backed by a file or by a POSIX shared memory object, we can use ftruncate() to adjust the size of the underlying object, and then re-create the mapping using munmap() and mmap() (or the Linux-specific mremap()).

Historically, System V shared memory was more widely available than *mmap()* and POSIX shared memory, although most UNIX implementations now provide all of these techniques.

With the exception of the final point regarding portability, the differences listed above are advantages in favor of shared file mappings and POSIX shared memory objects. Thus, in new applications, one of these interfaces may be preferable to System V shared memory. Which one we choose depends on whether or not we require a persistent backing store. Shared file mappings provide such a store; POSIX shared memory objects allow us to avoid the overhead of using a disk file when a backing store is not required.

## 54.6 Summary

A POSIX shared memory object is used to share a region of memory between unrelated processes without creating an underlying disk file. To do this, we replace the call to open() that normally precedes mmap() with a call to shm\_open(). The shm\_open() call creates a file in a memory-based file system, and we can employ traditional file descriptor system calls to perform various operations on this virtual file. In particular, ftruncate() must be used to set the size of the shared memory object, since initially it has a length of zero.

We have now described three techniques for sharing memory regions between unrelated processes: System V shared memory, shared file mappings, and POSIX shared memory objects. There are several similarities between the three techniques. There are also some important differences, and, except for the issue of portability, these differences favor shared file mappings and POSIX shared memory objects.

### 54.7 **Exercise**

Rewrite the programs in Listing 48-2 (svshm xfr writer.c) and Listing 48-3 (svshm xfr reader.c) to use POSIX shared memory objects instead of System V shared memory.