

Batch: A1 Roll No.: 16010123012

Experiment / assignment / tutorial No.: 08

Grade: AA / AB / BB / BC / CC / CD / DD

Signature of the Staff In-charge with date

Experiment No.:8

TITLE: Study and configure RIP protocol using Cisco Packet tracer

AIM: To study and configure RIP protocol using Cisco Packet tracer

Expected Outcome of Experiment:
CO:

Books/ Journals/ Websites referred:

1. A. S. Tanenbaum, "Computer Networks", Pearson Education, Fourth Edition
2. B. A. Forouzan, "Data Communications and Networking", TMH, Fourth Edition

Pre Lab/ Prior Concepts:

IPv4 Addressing, Subnetting, Distance Vector Protocol, Router configuration Commands.

New Concepts to be learned: RIP Protocol and its configuration.

RIP (Routing Information Protocol)

RIP is a standardized Distance Vector protocol, designed for use on smaller networks. RIP was one of the first true Distance Vector routing protocols and is supported on a wide variety of systems.

RIP adheres to the following Distance Vector characteristics:

- RIP sends out periodic routing updates (every 30 seconds)
- RIP sends out the full routing table every periodic update.
- RIP uses a form of distance as its metric (in this case, hop count).
- RIP uses the Bellman-Ford Distance Vector algorithm to determine the best "path" to a particular destination

Other characteristics of RIP include:

- RIP supports IP and IPX routing.
- RIP utilizes UDP port 520
- RIP routes have an administrative distance of 120.
- RIP has a maximum hop count of 15 hops.

RIP Versions

RIP has two versions, Version 1 (RIPv1) and Version 2 (RIPv2).

RIPv1 (RFC 1058) is *classful*, and thus does not include the subnet mask with its routing table updates. Because of this, RIPv1 does not support **Variable Length Subnet Masks (VLSMs)**. When using RIPv1, networks must be contiguous, and subnets of a major network must be configured with identical subnet masks. Otherwise, route table inconsistencies (or worse) will occur. RIPv1 sends updates as broadcasts to address 255.255.255.255.

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RIPv2 (RFC 2543) is *classless*, and thus does include the subnet mask with its routing table updates. RIPv2 fully supports VLSMs, allowing discontinuous networks and varying subnet masks to exist.

Other enhancements offered by RIPv2 include:

- Routing updates are sent via multicast, using address 224.0.0.9
- Encrypted authentication can be configured between RIPv2 routers
- Route tagging is supported.

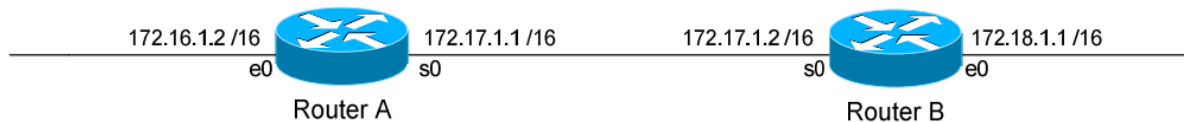
RIPv2 can interoperate with RIPv1. By default:

- RIPv1 routers will send only Version 1 packets
- RIPv1 routers will receive both Version 1 and 2 updates
- RIPv2 routers will both send and receive only Version 2 updates

We can control the version of RIP a particular interface will “send” or “receive.”

Unless RIPv2 is manually specified, a Cisco will default to RIPv1 when configuring RIP.

RIPv1 Basic Configuration



Routing protocol configuration occurs in Global Configuration mode. On Router A, to configure RIP, we would type:

```
Router(config)# router rip
Router(config-router)# network 172.16.0.0
Router(config-router)# network 172.17.0.0
```

The first command, router rip, enables the RIP process.

The network statements tell RIP which networks you wish to advertise to other RIP routers. We simply list the networks that are directly connected to our router. Notice that we specify the networks at their classful boundaries, and we do not specify a subnet mask.

To configure Router B:

```
Router(config)# router rip
Router(config-router)# network 172.17.0.0
Router(config-router)# network 172.18.0.0
```

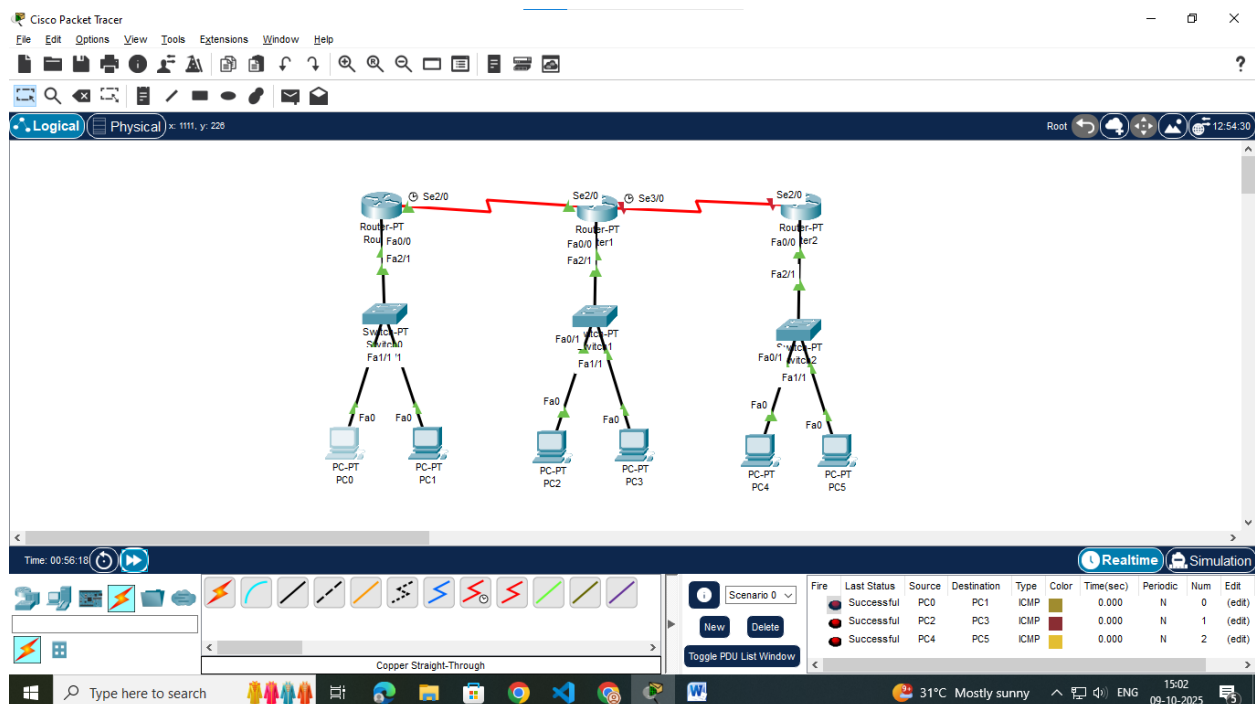
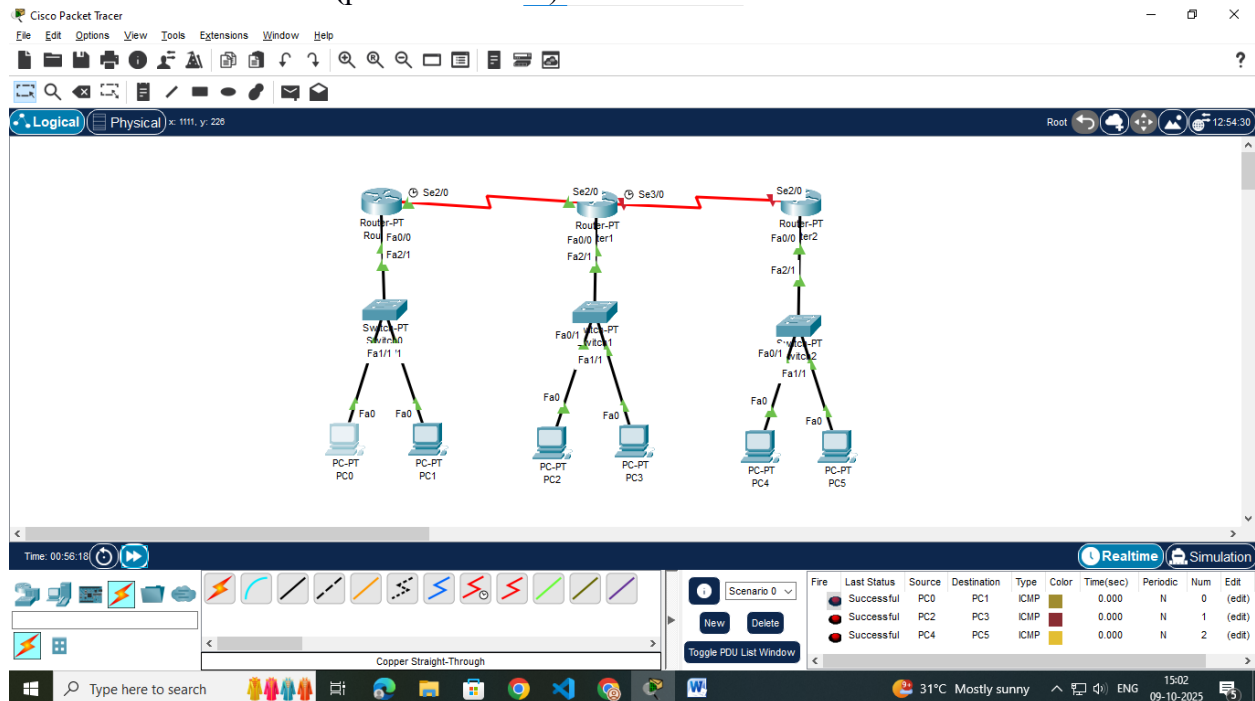
The routing table on Router A will look like:

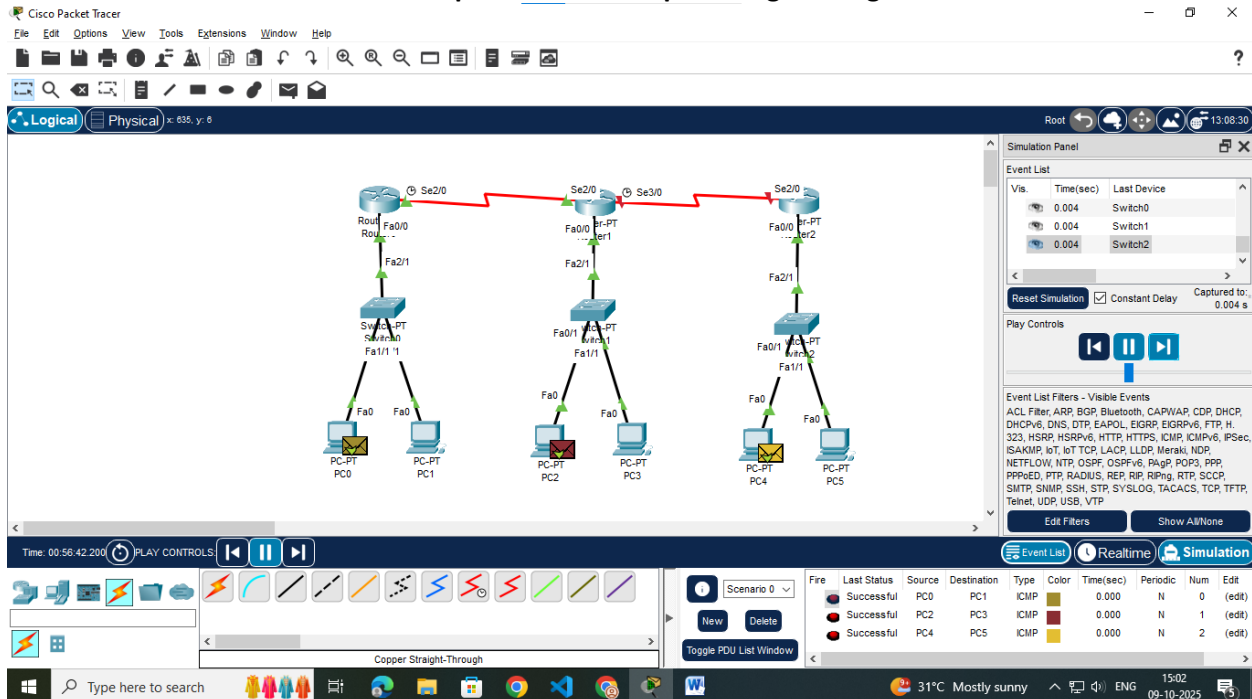
```
RouterA# show ip route
Gateway of last resort is not set
C    172.16.0.0 is directly connected, Ethernet0
C    172.17.0.0 is directly connected, Serial0
R    172.18.0.0 [120/1] via 172.17.1.2, 00:00:00, Serial0
```

The routing table on Router B will look like:

```
RouterB# show ip route
Gateway of last resort is not set
C    172.17.0.0 is directly connected, Serial0
C    172.18.0.0 is directly connected, Ethernet0
R    172.16.0.0 [120/1] via 172.17.1.1, 00:00:00, Serial0
```

IMPLEMENTATION: (printout of code)



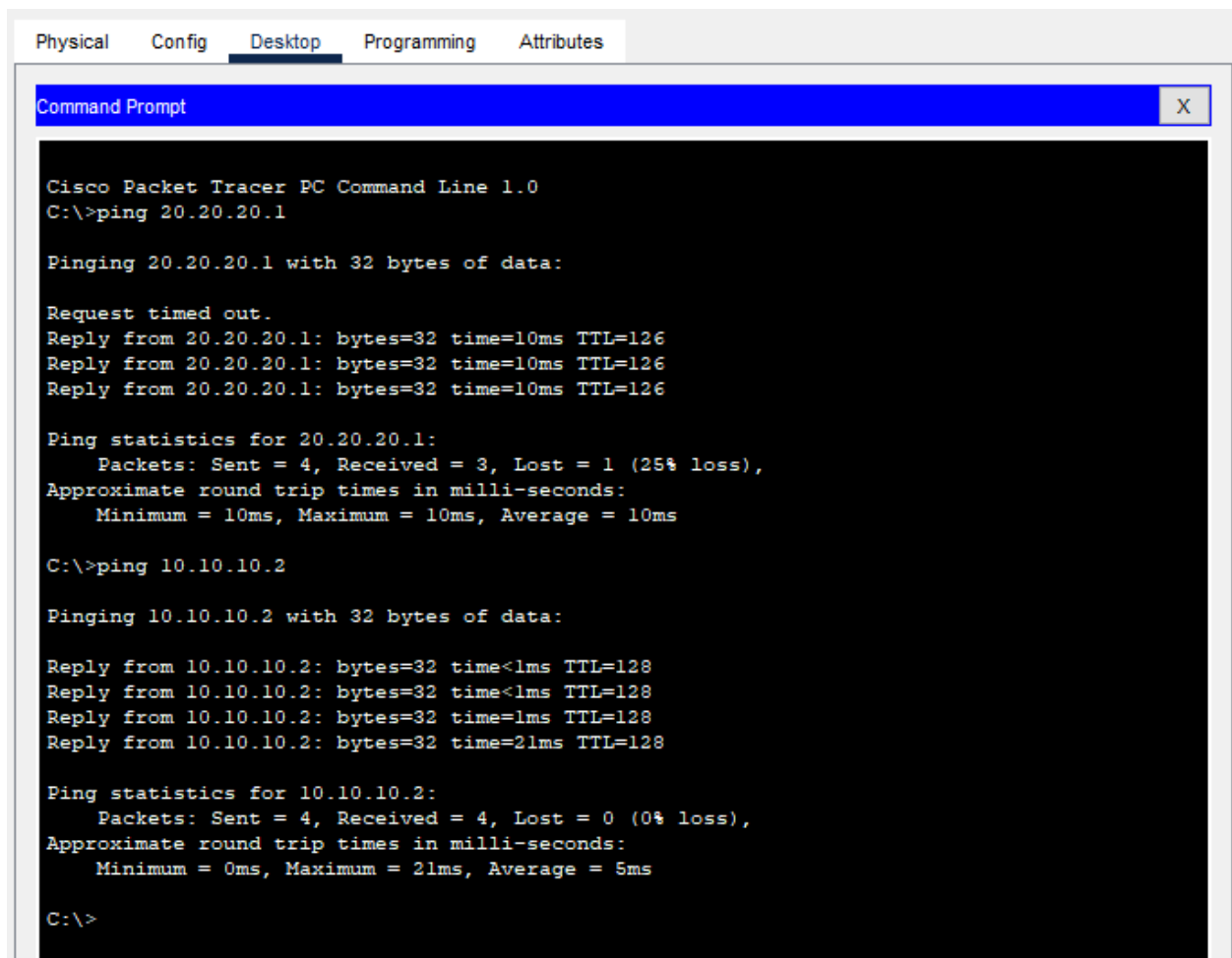


The image shows a Cisco Packet Tracer network diagram and its simulation panel. The network consists of three routers (R1, R2, R3) connected in a line. R1 is connected to R2 via S2/0/0 and S2/0/1. R2 is connected to R3 via S2/0/0 and S2/0/1. Each router has two Fast Ethernet interfaces (Fa0/0, Fa0/1) connected to two PCs (PC0, PC1 for R1; PC2, PC3 for R2; PC4, PC5 for R3). The simulation panel on the right shows the Event List with a table of events:

Vis	Time(sec)	Last Device
0.004	0.004	Switch0
0.004	0.004	Switch1
0.004	0.004	Switch2

The simulation panel also includes a Play Controls section with buttons for Play, Pause, and Stop. Below the Event List, there is a table showing the status of the simulation:

Fire	Last Status	Source	Destination	Type	Color	Time(sec)	Periodic	Num	Edit
Successful	Successful	PC0	PC1	ICMP	Yellow	0.000	N	0	(edit)
Successful	Successful	PC2	PC3	ICMP	Yellow	0.000	N	1	(edit)
Successful	Successful	PC4	PC5	ICMP	Yellow	0.000	N	2	(edit)



The image shows the Cisco Packet Tracer PC Command Line Interface (CLI) for a PC. The CLI is titled "Cisco Packet Tracer PC Command Line 1.0". The user has entered the command "ping 20.20.20.1". The output shows that the ping request timed out. The user then enters the command "ping 10.10.10.2". The output shows that the ping request was successful. The CLI also displays ping statistics for both destinations.

```
Cisco Packet Tracer PC Command Line 1.0
C:\>ping 20.20.20.1

Pinging 20.20.20.1 with 32 bytes of data:

Request timed out.
Reply from 20.20.20.1: bytes=32 time=10ms TTL=126
Reply from 20.20.20.1: bytes=32 time=10ms TTL=126
Reply from 20.20.20.1: bytes=32 time=10ms TTL=126

Ping statistics for 20.20.20.1:
    Packets: Sent = 4, Received = 3, Lost = 1 (25% loss),
    Approximate round trip times in milli-seconds:
        Minimum = 10ms, Maximum = 10ms, Average = 10ms

C:\>ping 10.10.10.2

Pinging 10.10.10.2 with 32 bytes of data:

Reply from 10.10.10.2: bytes=32 time<1ms TTL=128
Reply from 10.10.10.2: bytes=32 time<1ms TTL=128
Reply from 10.10.10.2: bytes=32 time=1ms TTL=128
Reply from 10.10.10.2: bytes=32 time=21ms TTL=128

Ping statistics for 10.10.10.2:
    Packets: Sent = 4, Received = 4, Lost = 0 (0% loss),
    Approximate round trip times in milli-seconds:
        Minimum = 0ms, Maximum = 21ms, Average = 5ms

C:\>
```

PC0

Physical Config **Desktop** Programming Attributes

IP Configuration X

Interface: FastEthernet0

IP Configuration

☐ DHCP ☒ Static

IPv4 Address: 10.10.10.1

Subnet Mask: 255.0.0.0

Default Gateway: 10.10.10.0

DNS Server: 0.0.0.0

IPv6 Configuration

☐ Automatic ☒ Static

IPv6 Address: /

Link Local Address: FE80::2E0:A3FF:FE8D:19C6

Default Gateway:

DNS Server:

802.1X

☐ Use 802.1X Security

Authentication: MD5

Username:

Password:

Router0

Physical **Config** CLI Attributes

GLOBAL

Settings

Algorithm Settings

ROUTING

Static

RIP

INTERFACE

FastEthernet0/0

FastEthernet1/0

Serial2/0

Serial3/0

FastEthernet4/0

FastEthernet5/0

RIP Routing

Network: Add

Network Address

10.0.0.0

40.0.0.0

Remove

Equivalent IOS Commands

Press RETURN to get started!

Press RETURN to get started!

Router>enable

Router#

Router#configure terminal

Enter configuration commands, one per line. End with CNTL/Z.

Router(config)#router rip

Router(config-router)#

Router0

Physical

Config

CLI

Attributes

GLOBAL

Settings

Algorithm Settings

ROUTING

Static

RIP

INTERFACE

FastEthernet0/0

FastEthernet1/0

Serial2/0

Serial3/0

FastEthernet4/0

FastEthernet5/0

FastEthernet0/0

Port Status

Bandwidth

Duplex

MAC Address

IP Configuration

IPv4 Address

Subnet Mask

Tx Ring Limit

Equivalent IOS Commands

```

Router>enable
Router#
Router#configure terminal
Enter configuration commands, one per line. End with CNTL/Z.
Router(config)#router rip
Router(config-router)#
Router(config-router)#
Router(config-router)#end
Router#configure terminal
Enter configuration commands, one per line. End with CNTL/Z.
Router(config)#
Router(config)#
Router(config)#interface FastEthernet0/0
Router(config-if)#

```

Router0

Physical

Config

CLI

Attributes

GLOBAL

Settings

Algorithm Settings

ROUTING

Static

RIP

INTERFACE

FastEthernet0/0

FastEthernet1/0

Serial2/0

Serial3/0

FastEthernet4/0

FastEthernet5/0

Serial2/0

Port Status

Duplex

Clock Rate

IP Configuration

IPv4 Address

Subnet Mask

Tx Ring Limit

Equivalent IOS Commands

```

Router(config)#router rip
Router(config-router)#
Router(config-router)#
Router(config-router)#end
Router#configure terminal
Enter configuration commands, one per line. End with CNTL/Z.
Router(config)#
Router(config)#
Router(config)#interface FastEthernet0/0
Router(config-if)#
Router(config-if)#exit
Router(config)#interface Serial2/0
Router(config-if)#

```

CONCLUSION:

Post Lab Questions

1. are two popular examples of distance vector routing protocols.

A. OSPF and RIP

B. RIP and BGP

C. BGP and OSPF

D. BGP and SPF

2. A routing table contains information entered manually.

A. static

B. dynamic

C. hierarchical

D. non static

3. A routing table is updated periodically using one of the dynamic routing protocols.

A. static

B. dynamic

C. hierarchical

D. non static

4. Which of the following is not the category of dynamic routing algorithm.

A. Distance vector protocols

B. Link state protocols

C. Hybrid protocols

D. Automatic state protocols

5. In forwarding, the mask and destination addresses are both 0.0.0.0 in routing table.

A. next-hop

B. network-specific

C. host-specific

D. default

6. Differentiate between Distance Vector Routing and Link State Routing.

Distance Vector Routing	Link State Routing
Bandwidth required is less due to local sharing, small packets and no flooding.	Bandwidth required is more due to flooding and sending of large link state packets.
Based on local knowledge, since it updates table based on information from neighbours.	Based on global knowledge, it have knowledge about entire network.
Make use of Bellman Ford Algorithm.	Make use of Dijkstra's algorithm.
Traffic is less.	Traffic is more.



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Department of Computer Engineering



Converges slowly i.e, good news spread fast and bad news spread slowly.	Converges faster.
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Date: 09/10/2025

Signature of faculty in-charge