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### **Department of Computer Engineering**

Batch: A1 Roll No.: 16010123012

Experiment / assignment / tutorial No. 7

Grade: AA / AB / BB / BC / CC / CD /DD

Signature of the Staff In-charge with date

Title: Indexing - Create Index and observe the evaluation statistics of query execution with and without Index

**Objective:** Implement indexing to improve query execution plans

# **Expected Outcome of Experiment:**

CO 4: Analyse Advanced Database Concepts like indexing, hashing, query processing, query optimization, normalization.

### **Books/ Journals/ Websites referred:**

- 1. Dr. P.S. Deshpande, SQL and PL/SQL for Oracle 10g.Black book, Dreamtech Press
- 2. www.db-book.com
- 3. Korth, Slberchatz, Sudarshan : "Database Systems Concept", 5<sup>th</sup> Edition , McGraw Hill
- 4. Elmasri and Navathe,"Fundamentals of database Systems", 4<sup>th</sup> Edition,PEARSON Education.

Resources used: PostgreSQL

### **Theory:**

A database index is a data structure that improves the speed of operations in a table. Indexes can be created using one or more columns, providing the basis for both rapid random lookups and efficient ordering of access to records.

While creating index, it should be taken into consideration which all columns will be used to make SQL queries and create one or more indexes on those columns.

To add an index for a column or a set of columns, you use the CREATE INDEX statement as follows:

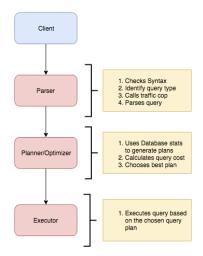
CREATE INDEX index\_name ON table\_name (column\_list)

Query life cycle





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### **Planner and Executor:**

The planner receives a query tree from the rewriter and generates a (query) plan tree that can be processed by the executor most effectively.

The planner in Database is based on pure cost-based optimization -

### **EXPLAIN command:**

This command displays the execution plan that the PostgreSQL/MySQL planner generates for the supplied statement. The execution plan shows how the table(s) referenced by the statement will be scanned — by plain sequential scan, index scan, etc. — and if multiple tables are referenced, what join algorithms will be used to bring together the required rows from each input table.

As in the other RDBMS, the EXPLAIN command in Database displays the plan tree itself. A specific example is shown below:-

Database: testdb=#

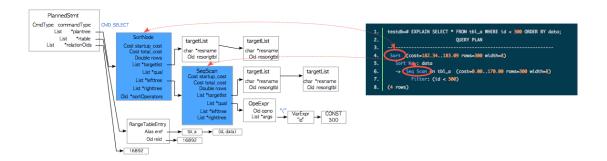
- 1. EXPLAIN SELECT \* FROM tbl a WHERE id < 300 ORDER BY data;
- 2. QUERY PLAN
- 3. -----
- 4. Sort (cost=182.34..183.09 rows=300 width=8)
- 5. Sort Key: data
- 6. -> Seq Scan on tbl\_a (cost=0.00..170.00 rows=300 width=8)
- 7. Filter: (id < 300)
- 8. (4 rows)

A simple plan tree and the relationship between the plan tree and the result of the EXPLAIN command in PostgreSQL.

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### **Nodes**

The first thing to understand is that each indented block with a preceeding "->" (along with the top line) is called a node. A node is a logical unit of work (a "step" if you will) with an associated cost and execution time. The costs and times presented at each node are cumulative and roll up all child nodes.

### Cost:

It is not the time but a concept designed to estimate the cost of an operation. The first number is start-up cost (cost to retrieve first record) and the second number is the cost incurred to process entire node (total cost from start to finish).

Cost is a combination of 5 work components used to estimate the work required: sequential fetch, non-sequential (random) fetch, processing of row, processing operator (function), and processing index entry.

**Rows** are the approximate number of rows returned when a specified operation is performed.

(In the case of select with where clause rows returned is

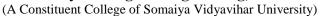
Rows = cardinality of relation \* selectivity )

Width is an average size of one row in bytes.

### **Explain Analyze command:**

The EXPLAIN ANALYZE option causes the statement to be actually executed, not only planned. Then actual run time statistics are added to the display, including the total elapsed time expended within each plan node (in milliseconds) and the total number of rows it actually returned. This is useful for seeing whether the planner's estimates are close to reality.

Ex: EXPLAIN (ANALYZE) SELECT \* FROM foo;





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# QUERY PLAN — Seq Scan on foo (cost=0.00..18334.10 rows=1000010 width=37) (actual time=0.012..61.524 rows=1000010 loops=1) Total runtime: 90.944 ms (2 rows)

The command displays the following additional parameters:

- **actual time** is the actual time in milliseconds spent to get the first row and all rows, respectively.
- rows is the actual number of rows received with Seq Scan.
- **loops** is the number of times the Seq Scan operation had to be performed.
- **Total runtime** is the total time of query execution.

Query plans for select with where clause can be sequential scan, Index Scan, Index only Scan, Bitmap Index Scan etc.

Query plans for joins are Nested loop join, Hash join, Merge join etc.

Indexing: CREATE INDEX constructs an index on the specified column(s) of the specified relation, which can be a table or a materialized view. Indexes are primarily used to enhance database performance (though inappropriate use can result in slower performance).

**Syntax** 

```
CREATE [ UNIQUE ] INDEX [ CONCURRENTLY ] [ [ IF NOT EXISTS ] name ] ON [
ONLY ] table_name [ USING method ]

( { column_name | (expression ) } [ COLLATE collation ] [ opclass [ (opclass_parameter = value [, ...] ) ] ] [ ASC | DESC ] [ NULLS { FIRST | LAST } ] [, ...] )

[ INCLUDE (column_name [, ...] ) ]

[ NULLS [ NOT ] DISTINCT ]

[ WITH (storage_parameter [= value] [, ...] ) ]

[ TABLESPACE tablespace_name ]

[ WHERE predicate ]
```

example



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```
explain Analyse select * from std2 where branch ='ext'
create index dept on std2(Branch)
select * from std2
explain Analyse select * from std2 where branch ='ext'
drop index dept
```

### **Implementation Screenshots:**

Comprehend how indexes improves the performance of query applied for your database . Demonstrate for the following types of query on your database

- a. Simple select query
- b. Select query with where clause
- c. Select query with order by query
- d. Select query with JOIN
- e. Select query with aggregation

```
CREATE TABLE users (
    id SERIAL PRIMARY KEY,
    name TEXT,
    email TEXT
);

INSERT INTO users (name, email) VALUES
('Alice', 'alice@example.com'),
('Bob', 'bob@example.com'),
('Charlie', 'charlie@example.com');

EXPLAIN ANALYZE SELECT * FROM users WHERE name = 'Alice';

CREATE INDEX idx_users_name ON users (name);

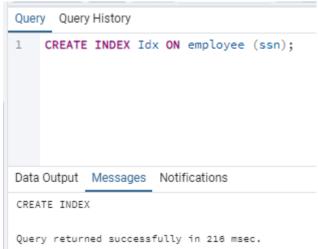
EXPLAIN ANALYZE SELECT * FROM users WHERE name = 'Alice';
```



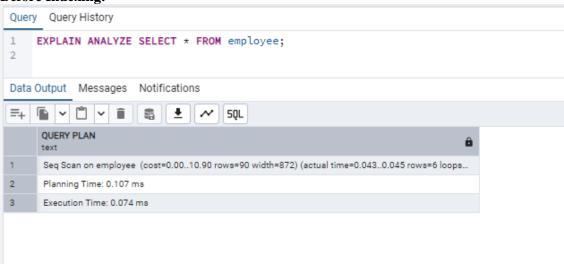
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# Index:



### **Before Indexing:**



### **After Indexing:**

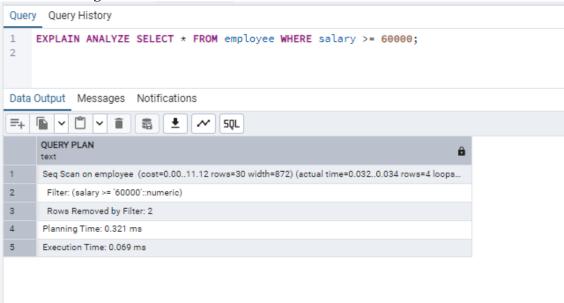




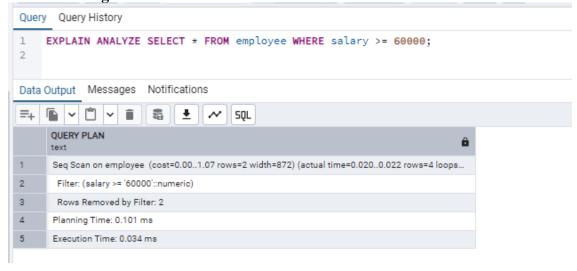
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### **Before Indexing:**



### **After Indexing:**

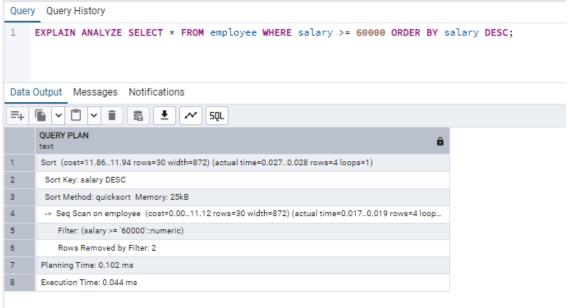




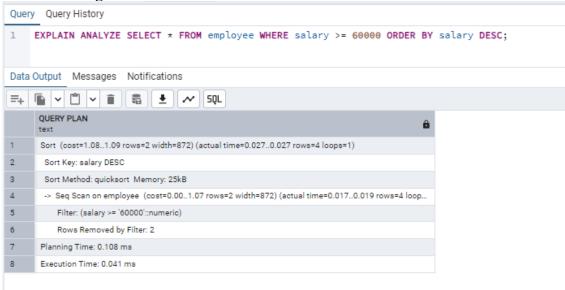
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### **After Indexing:**

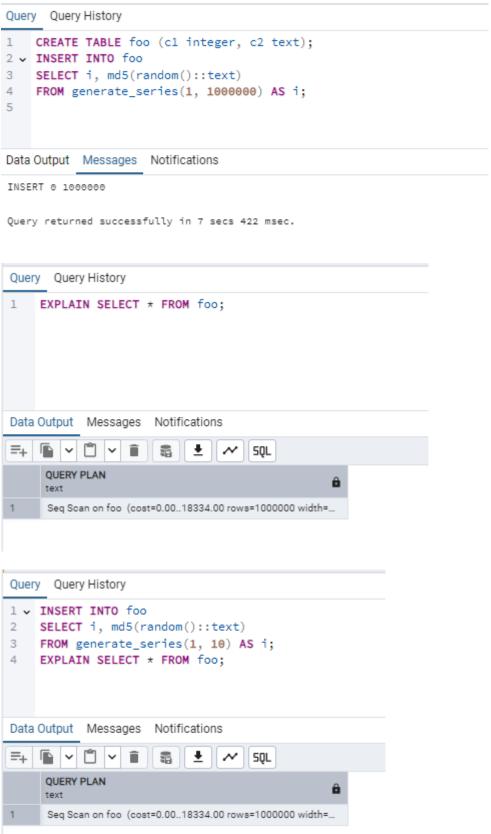




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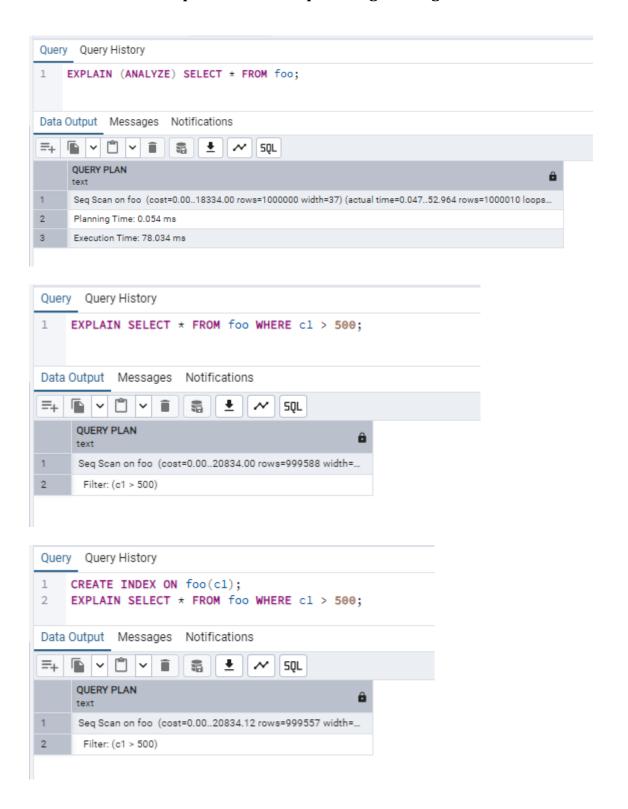
### FOO-





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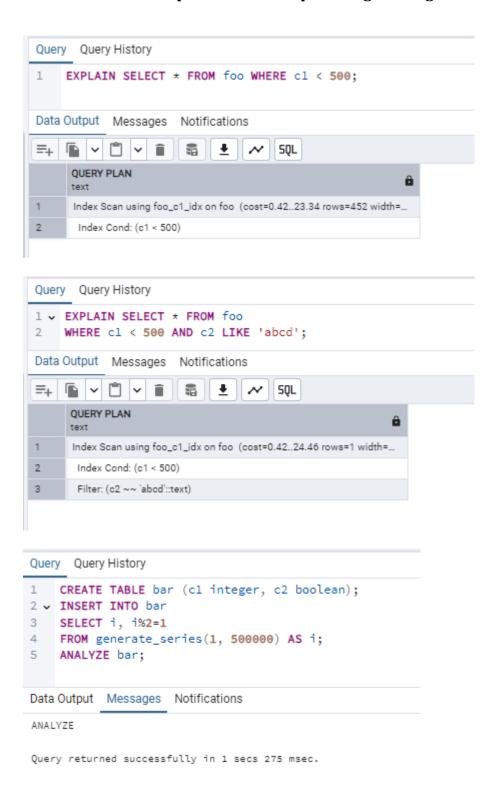
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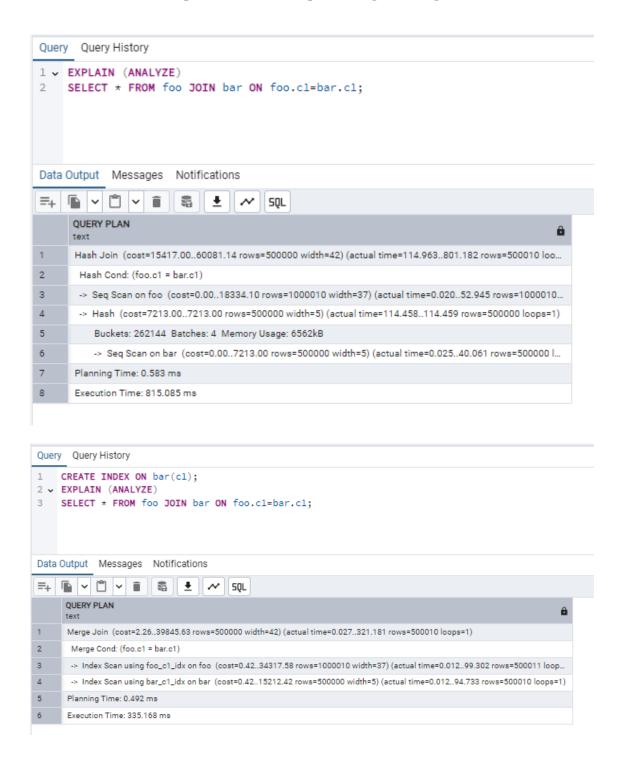
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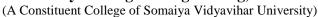
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### **Post Lab Question:**

1. How does a B-tree differ from a B+-tree? Why is a B+-tree usually preferred as an access structure to a data file?

A B-tree and a B+-tree are both balanced tree structures used in databases and file systems for efficient data retrieval. In a B-tree, both internal and leaf nodes store keys and data, whereas in a B+-tree, only the leaf nodes hold the actual data, and





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internal nodes contain only keys used for navigation. Additionally, leaf nodes in a B+-tree are linked, enabling efficient sequential and range queries, unlike in a B-tree where leaf nodes are not connected. Search operations in a B-tree can end at any node, while in a B+-tree, they always reach the leaf level, making the search path uniform. Though B+-trees may be slightly taller, their smaller internal nodes fit more keys per disk block, reducing disk I/O. This structure also simplifies insertions and deletions, as changes mainly affect leaf nodes. Overall, a B+-tree is usually preferred as an access structure to a data file because it supports faster range queries, better disk access patterns, and more predictable performance, making it ideal for indexing large datasets.

### **Conclusion:**

I have successfully implemented indexing techniques and analyzed their impact on query performance using SQL commands such as EXPLAIN ANALYZE. This experiment demonstrated how indexing optimizes query execution by minimizing data scan time and improving efficiency. By comparing query plans before and after indexing, I observed significant performance gains, especially in operations involving filtering and sorting.