

Data Mining Techniques

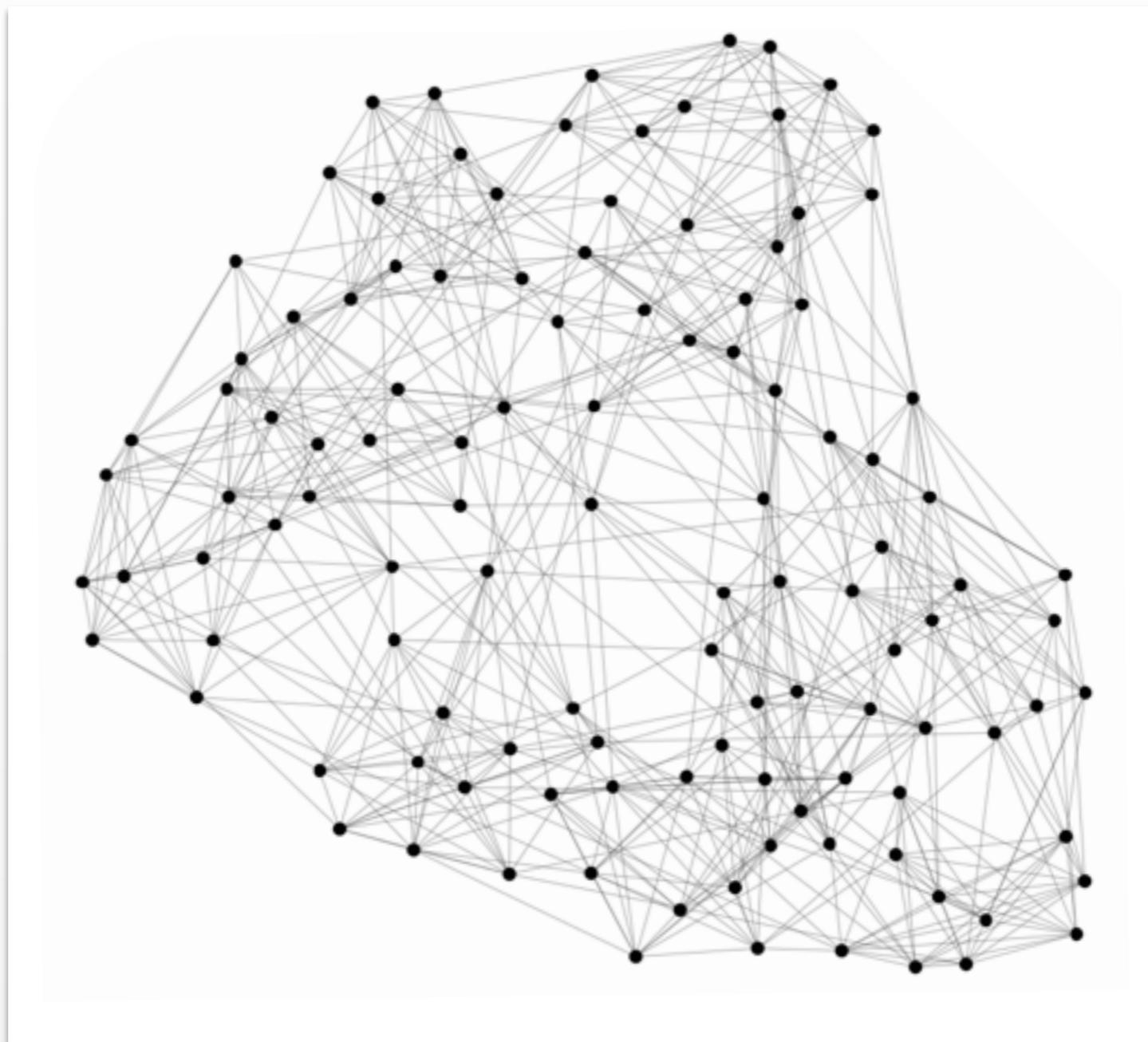
CS 6220 - Section 3 - Fall 2016

Lecture 19: Social Networks

Jan-Willem van de Meent
(*credit*: Leskovec et al Chapter 10,
Aggarwal Chapter 19)

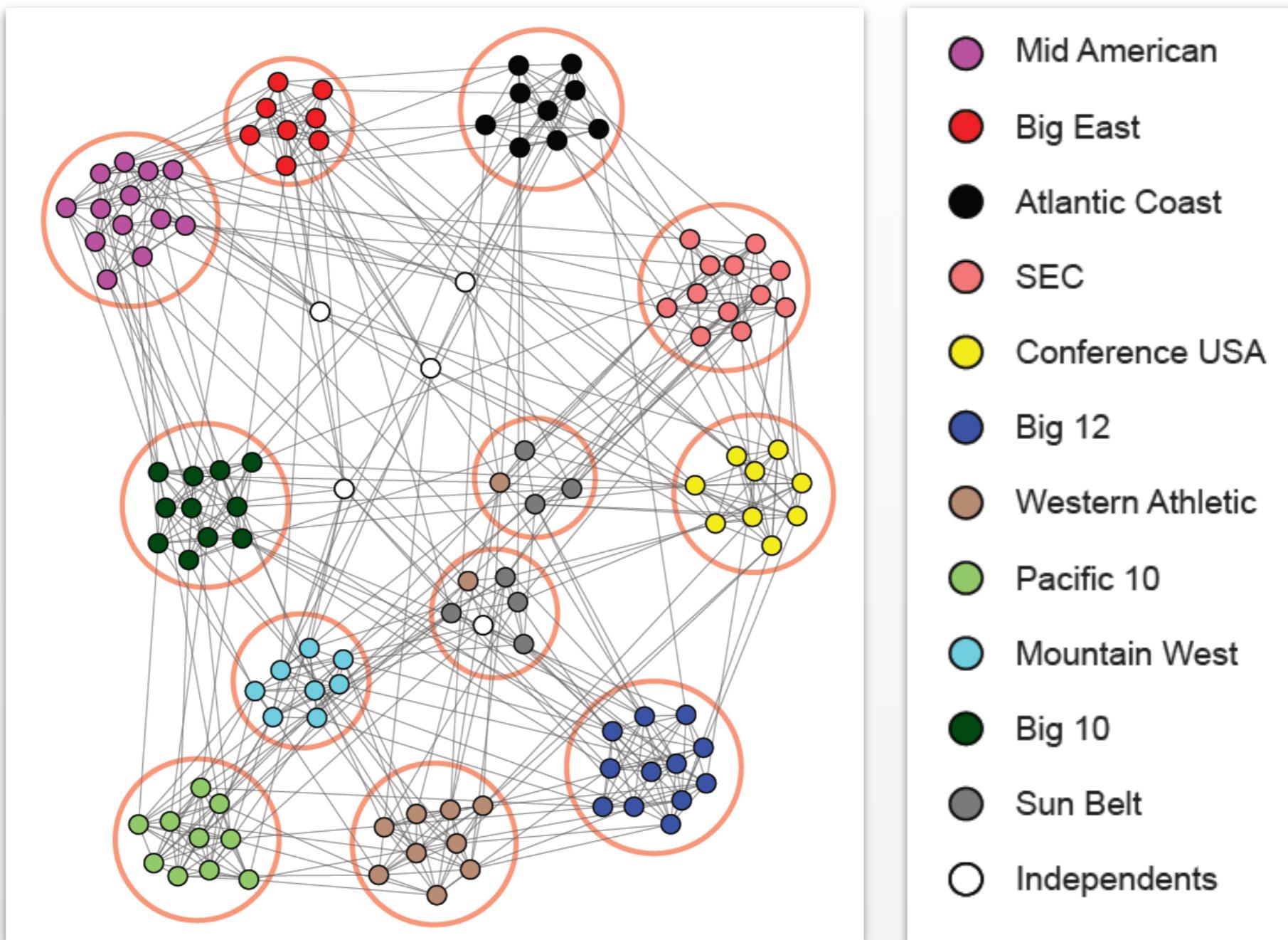


Community Detection



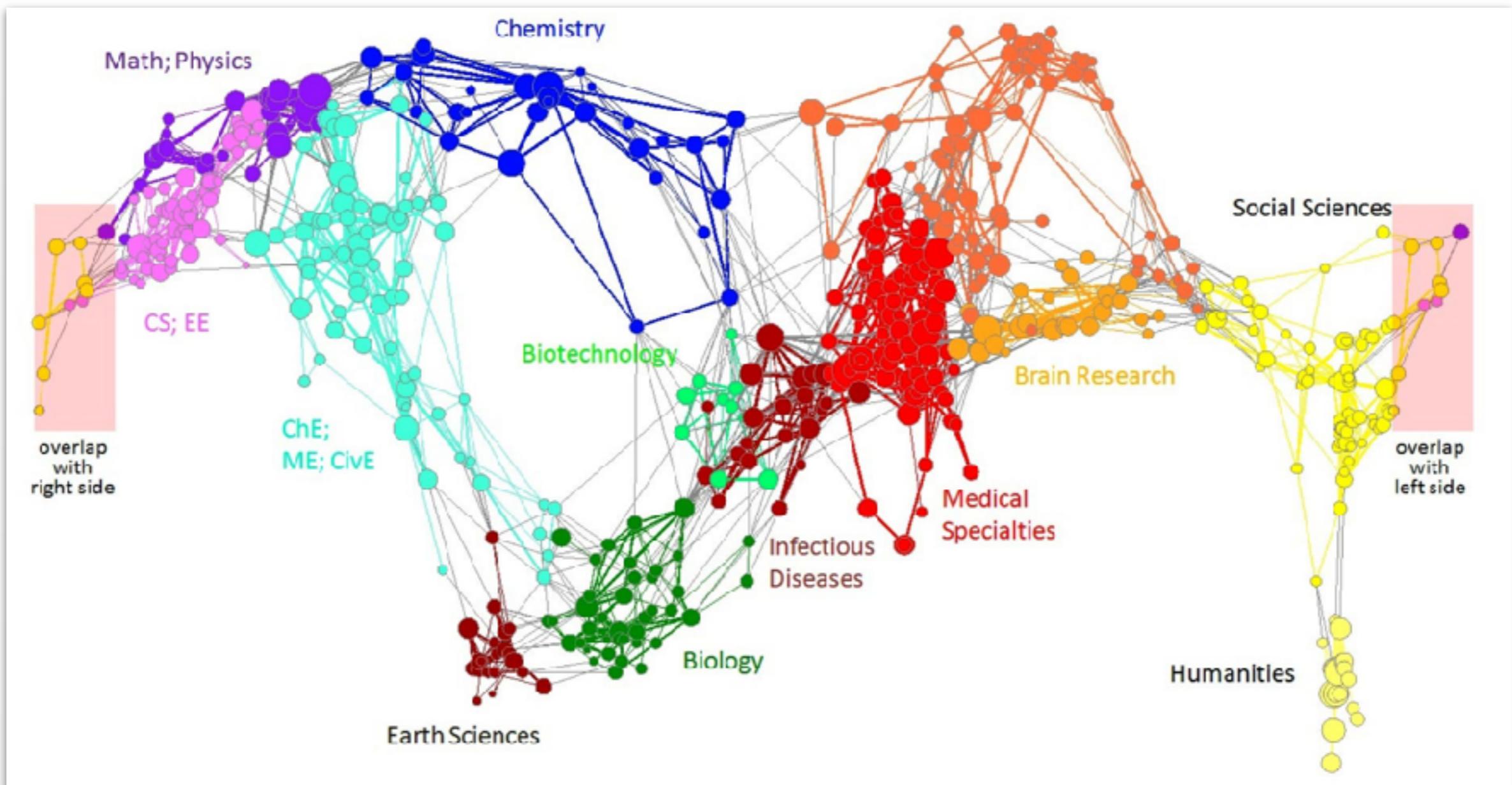
Problem: Can we identify groups of densely connected nodes?

Communities: Football Conferences



*Nodes: Football Teams, Edges: Matches,
Communities: Conferences*

Communities: Academic Citations

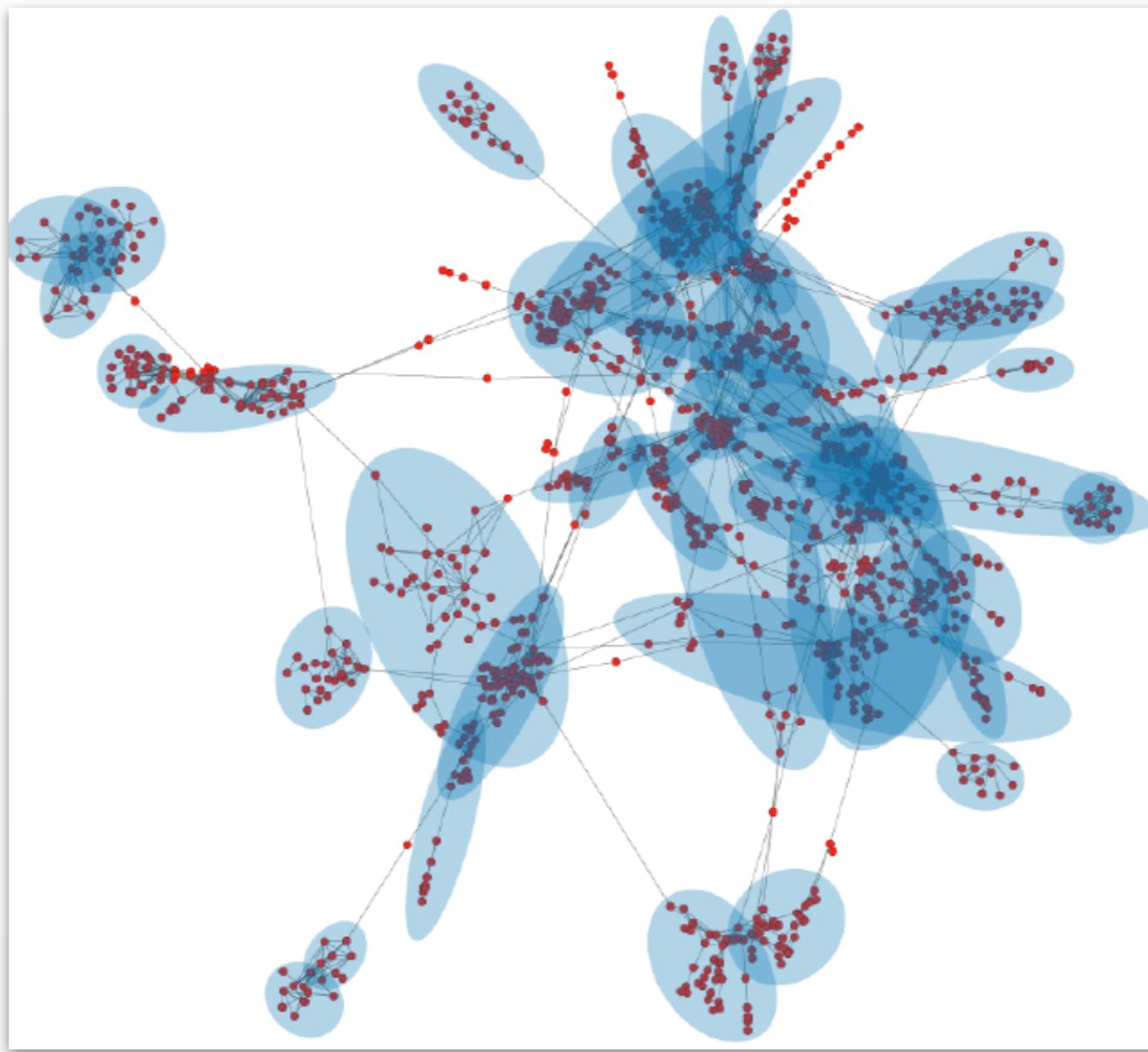


Source: Citation networks and Maps of science [Börner et al., 2012]

Nodes: Journals, Edges: Citations,
Communities: Academic Disciplines

(Adapted from: Mining of Massive Datasets, <http://www.mmds.org>)

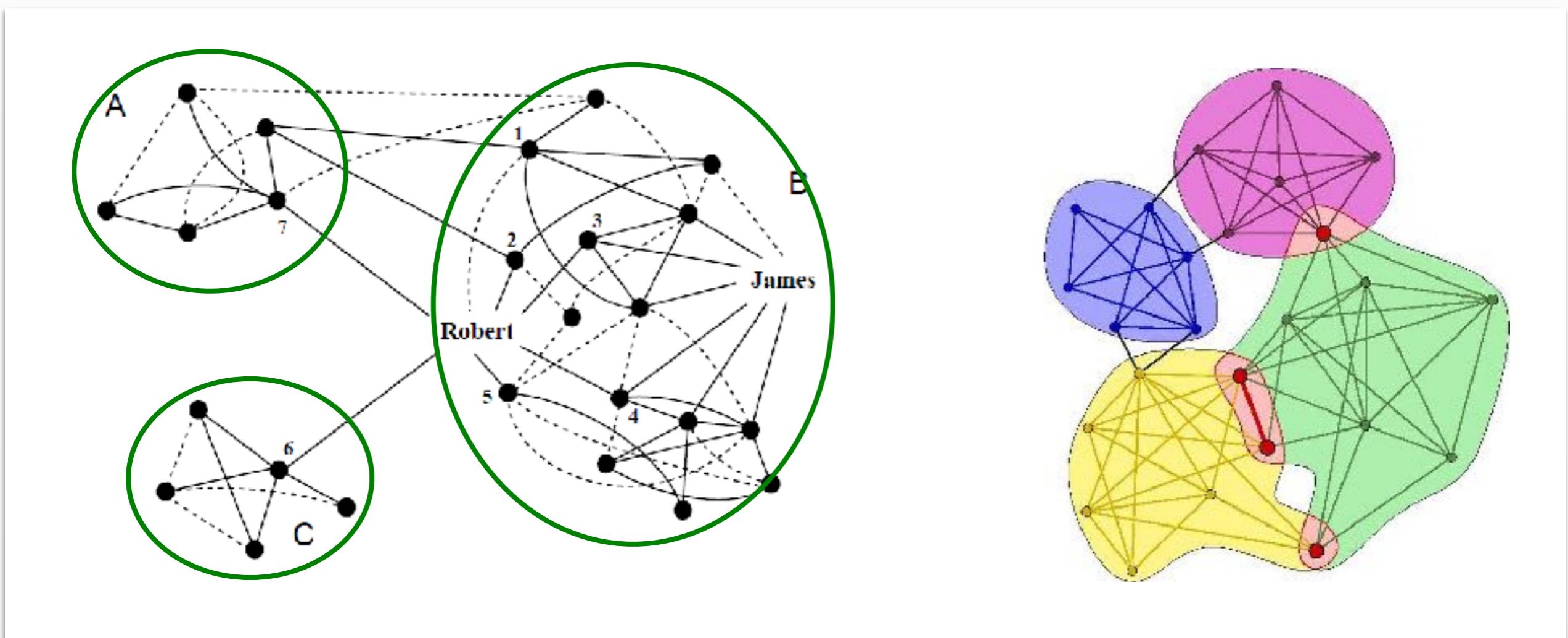
Communities: Protein-Protein Interactions



*Nodes: Proteins, Edges: Physical interactions,
Communities: Functional Modules*

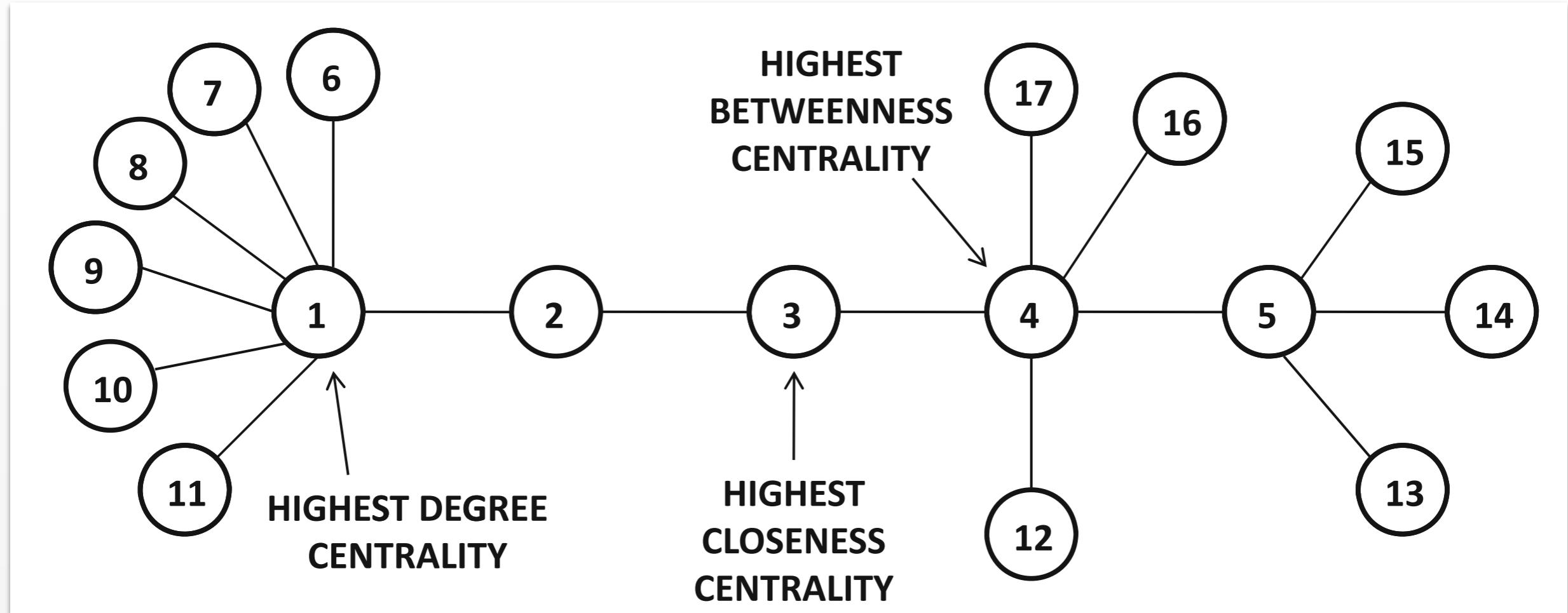
Community Detection

Graph Partitioning



We will work with **undirected** (unweighted) networks

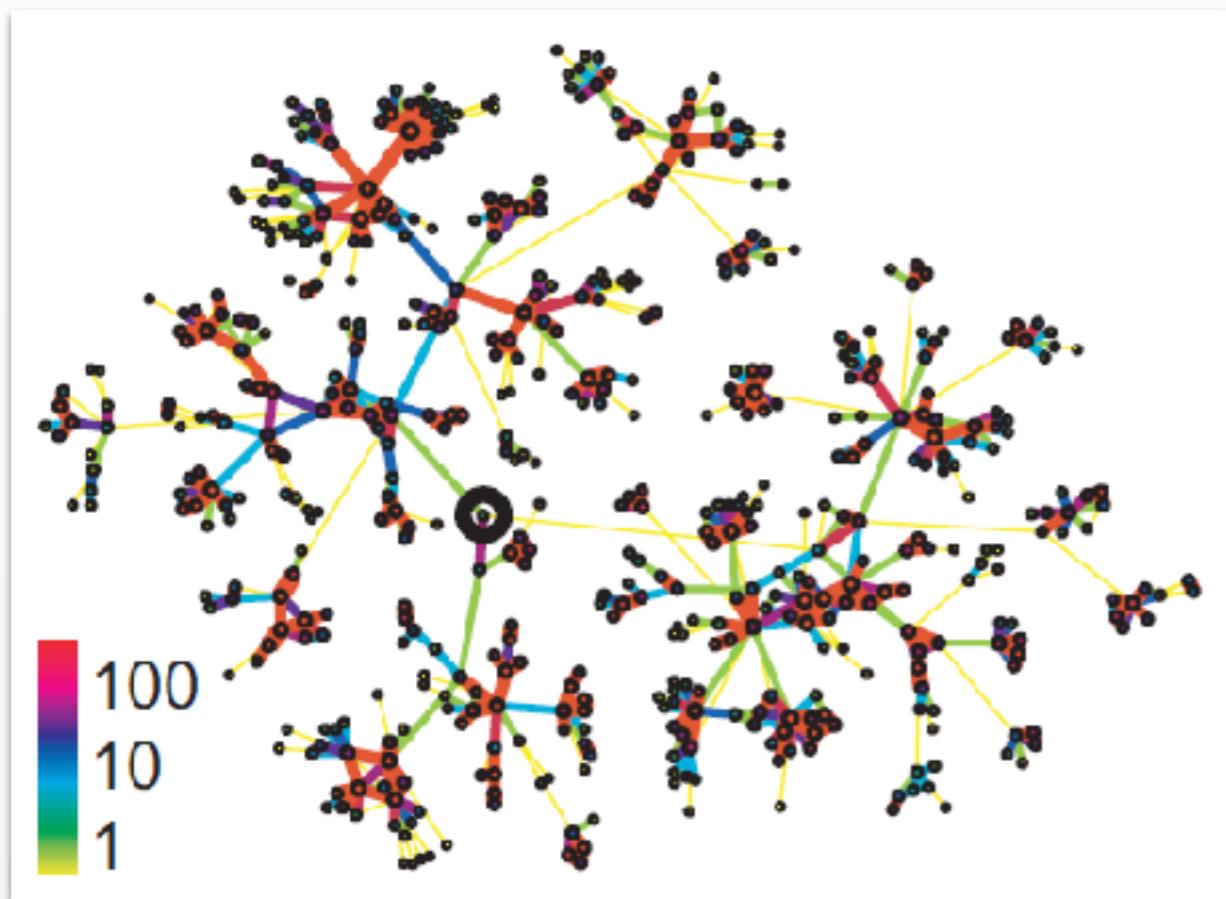
Centrality Measures



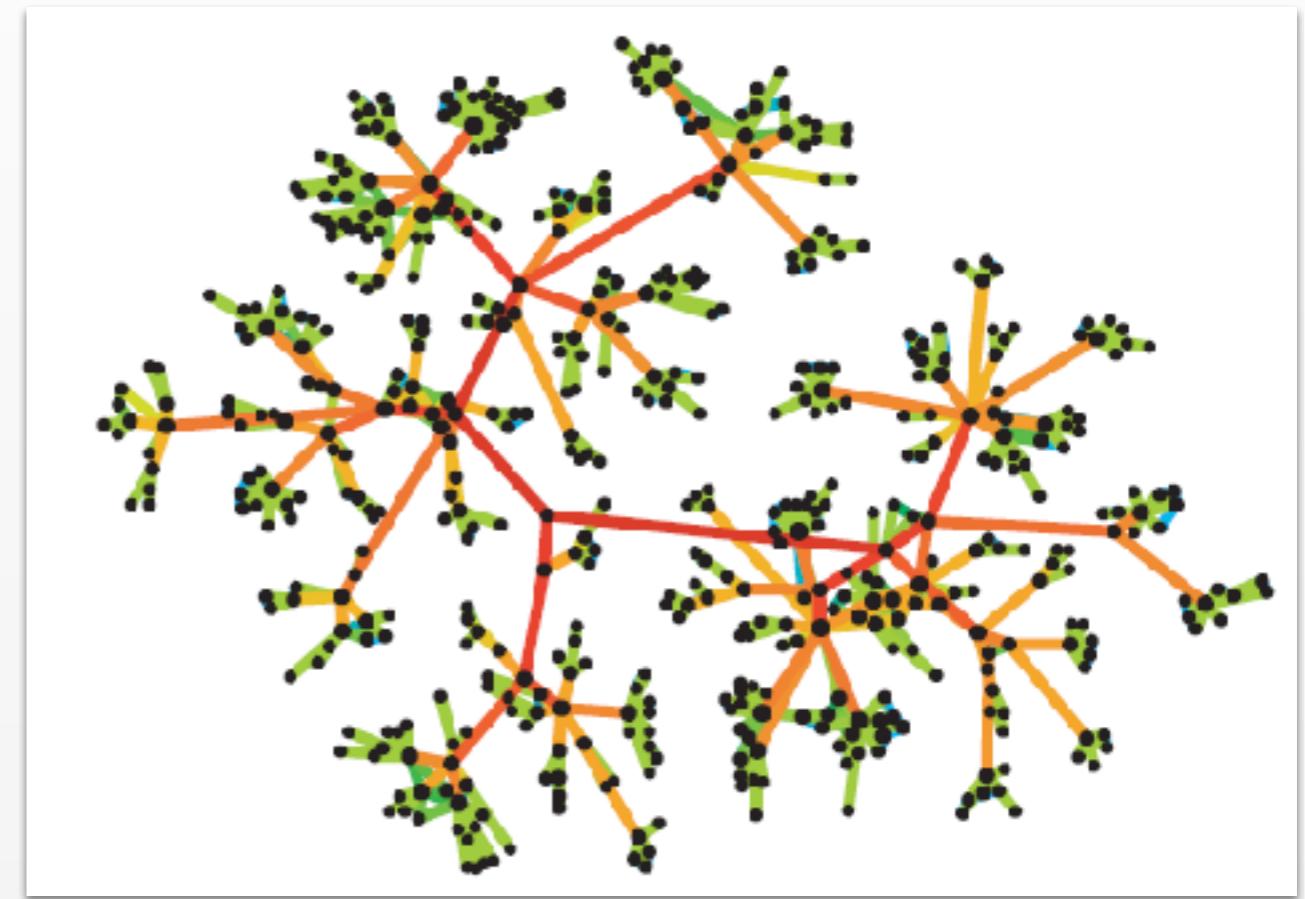
- *Betweenness*: Number of shortest paths
- *Closeness*: Average distance to other nodes
- *Degree*: Number of connections to other nodes

Betweenness

Edge Strength (call volume)

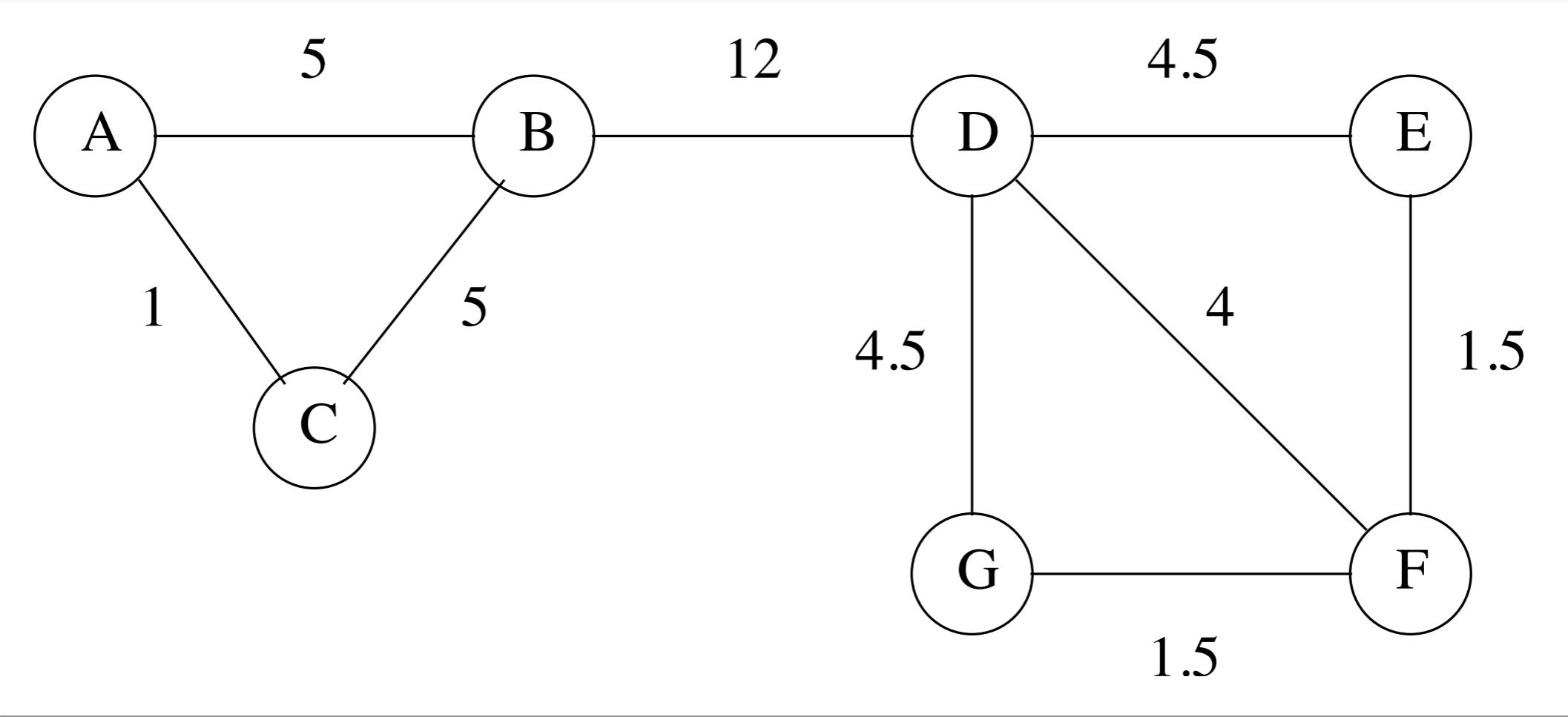


Edge Betweenness



- *Betweenness*: Number of shortest paths passing through a node or edge

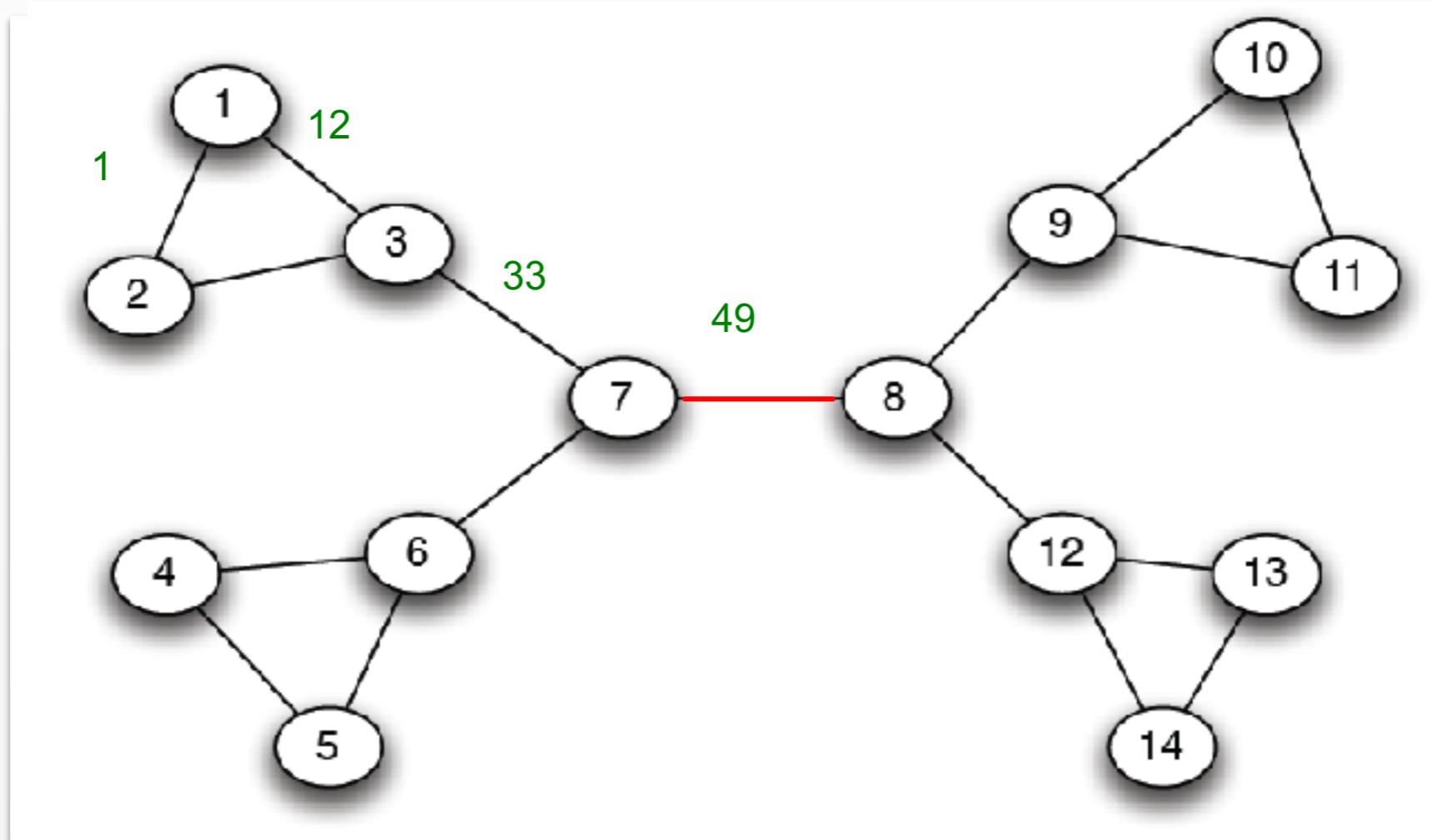
Edge Betweenness



- Count number of shortest paths passing through each edge (*can be done with weighted edges*)
- If there are multiple paths of equal length, then split counts

Girvan-Newman Algorithm

(hierarchical divisive clustering according to betweenness)



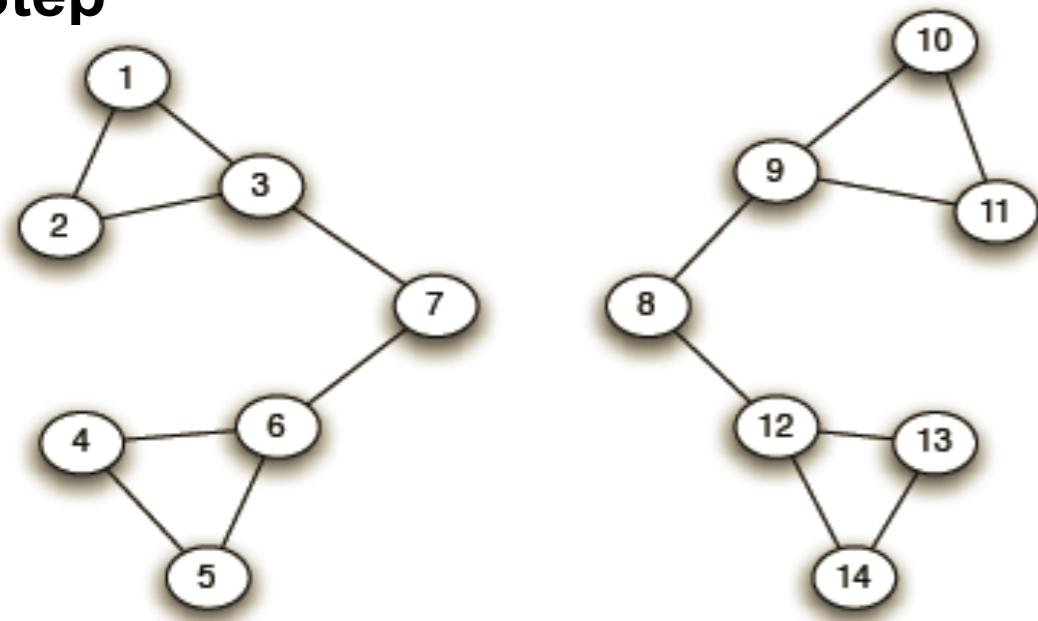
Repeat until k clusters found

1. Calculate betweenness
2. Remove edge(s) with highest betweenness

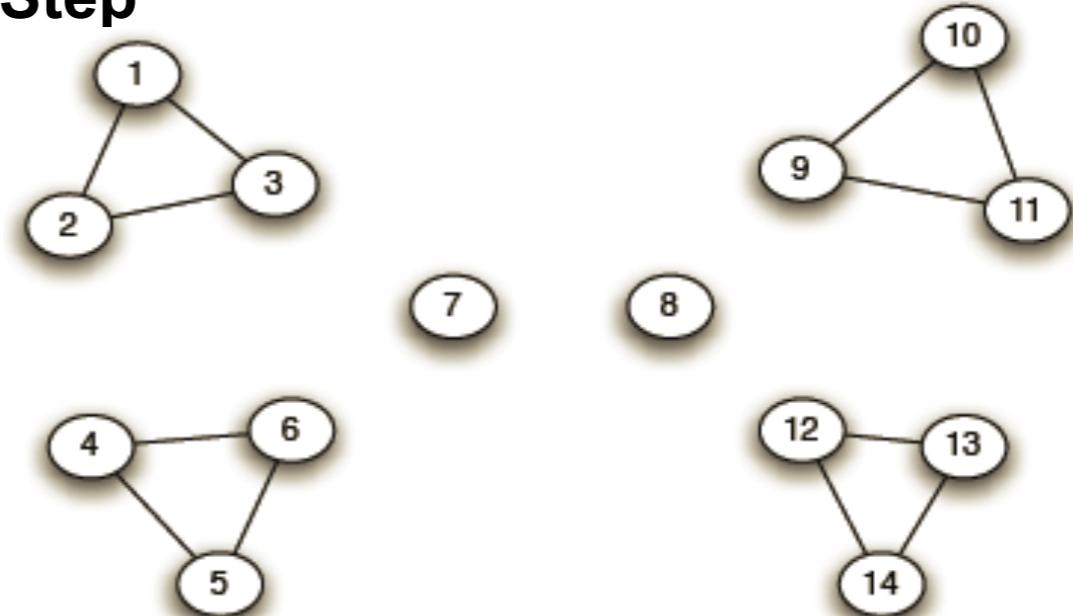
Girvan-Newman Algorithm

(hierarchical divisive clustering according to betweenness)

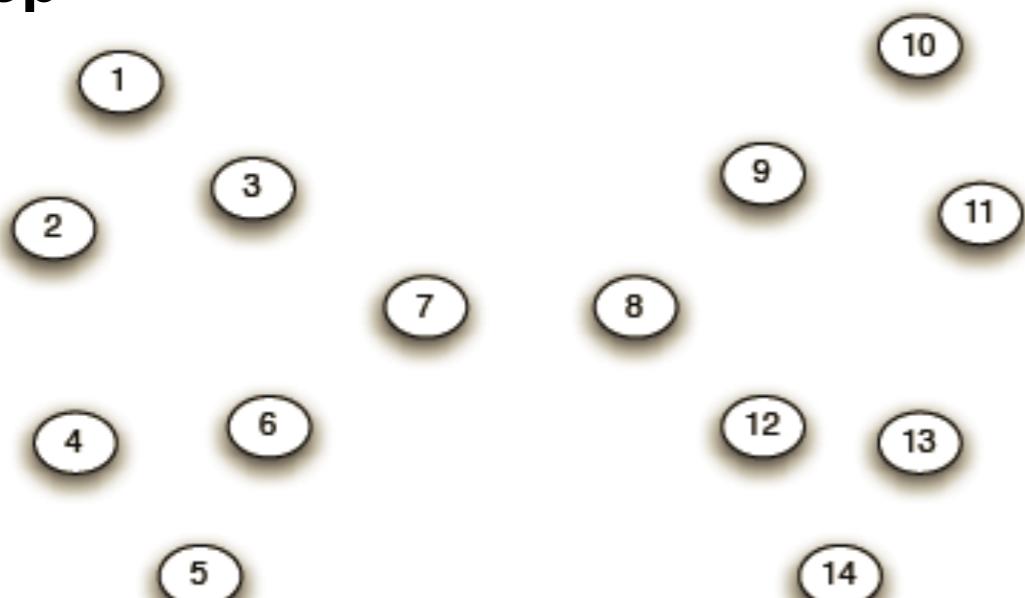
Step



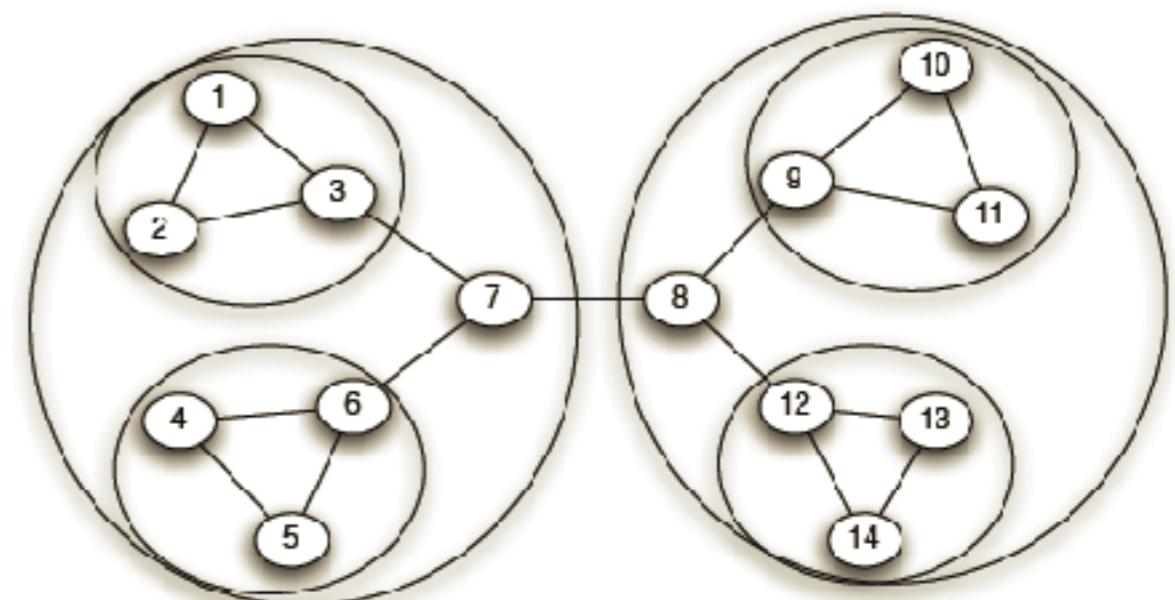
Step



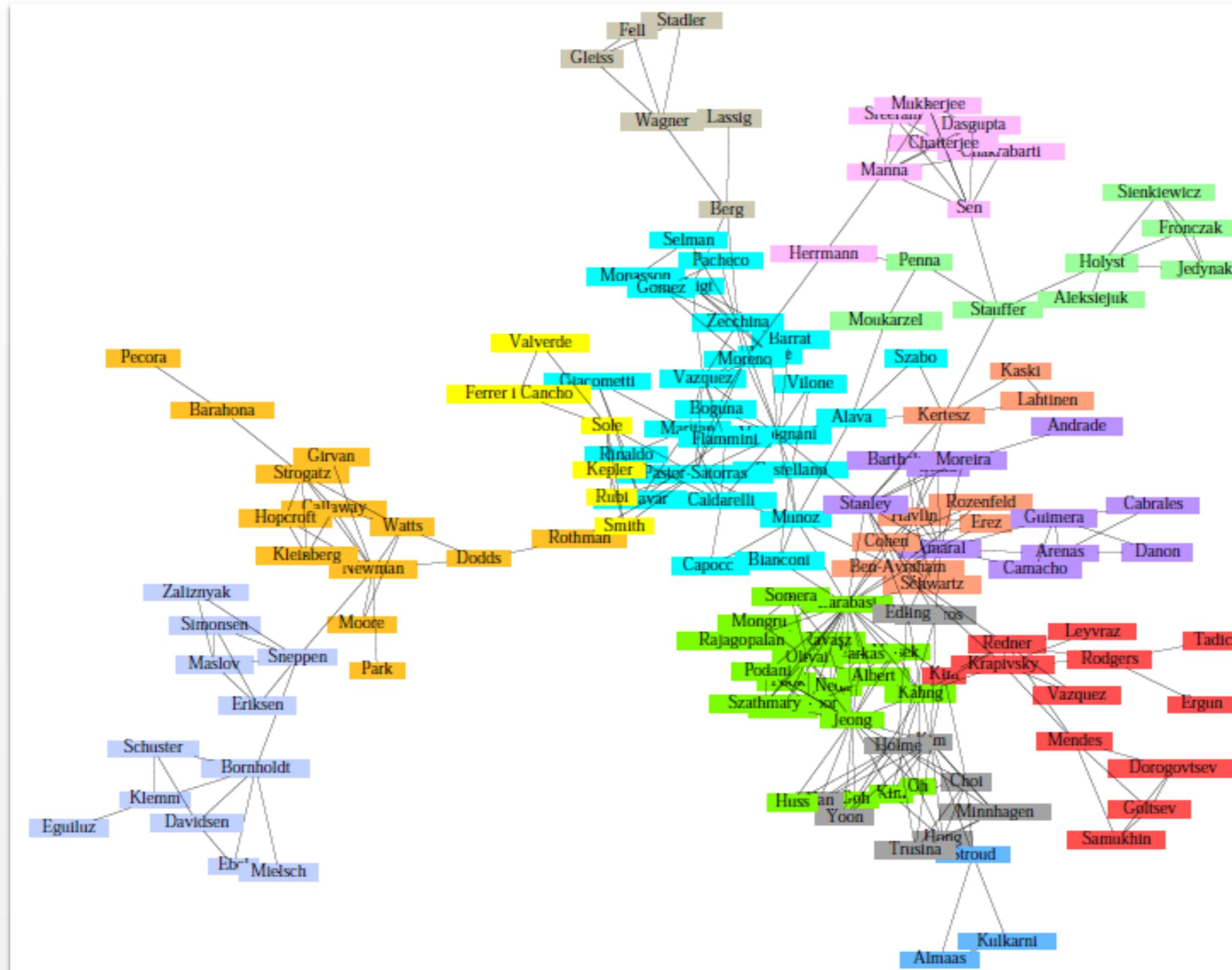
Step



Hierarchical network



Girvan-Newman: Physics Citations



(Adapted from: Mining of Massive Datasets, <http://www.mmds.org>)

Girvan-Newman

Two problems

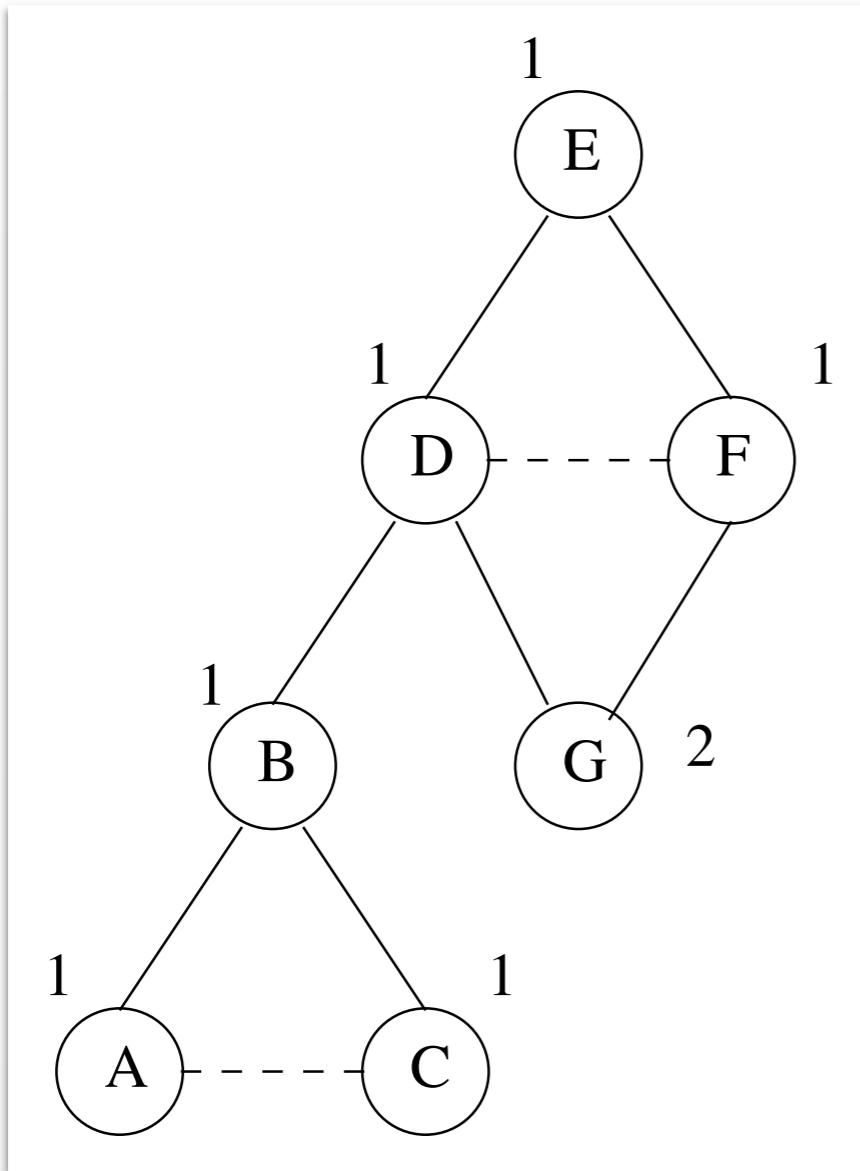
1. How can we compute the betweenness for all edges?
2. How can we choose the number of components k?

Calculating Betweenness

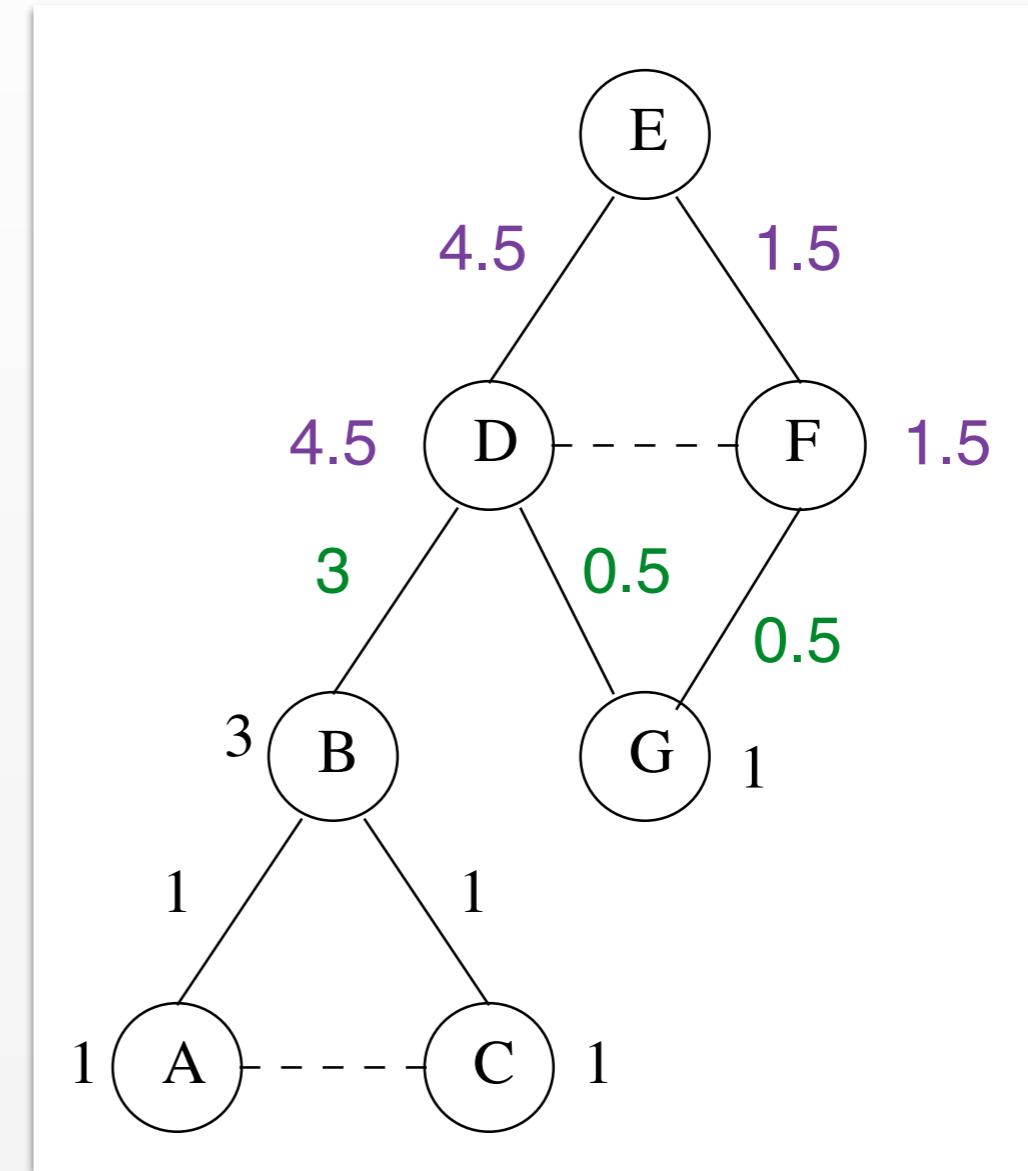
How can we count all shortest paths?

- Loop over nodes in graph
 - Perform breadth-first search to find shortest paths to other nodes
 - Increment counts for edges traversed by shorts paths
- Divide final betweenness by 2
(since all paths counted twice)

Counting Shortest Paths



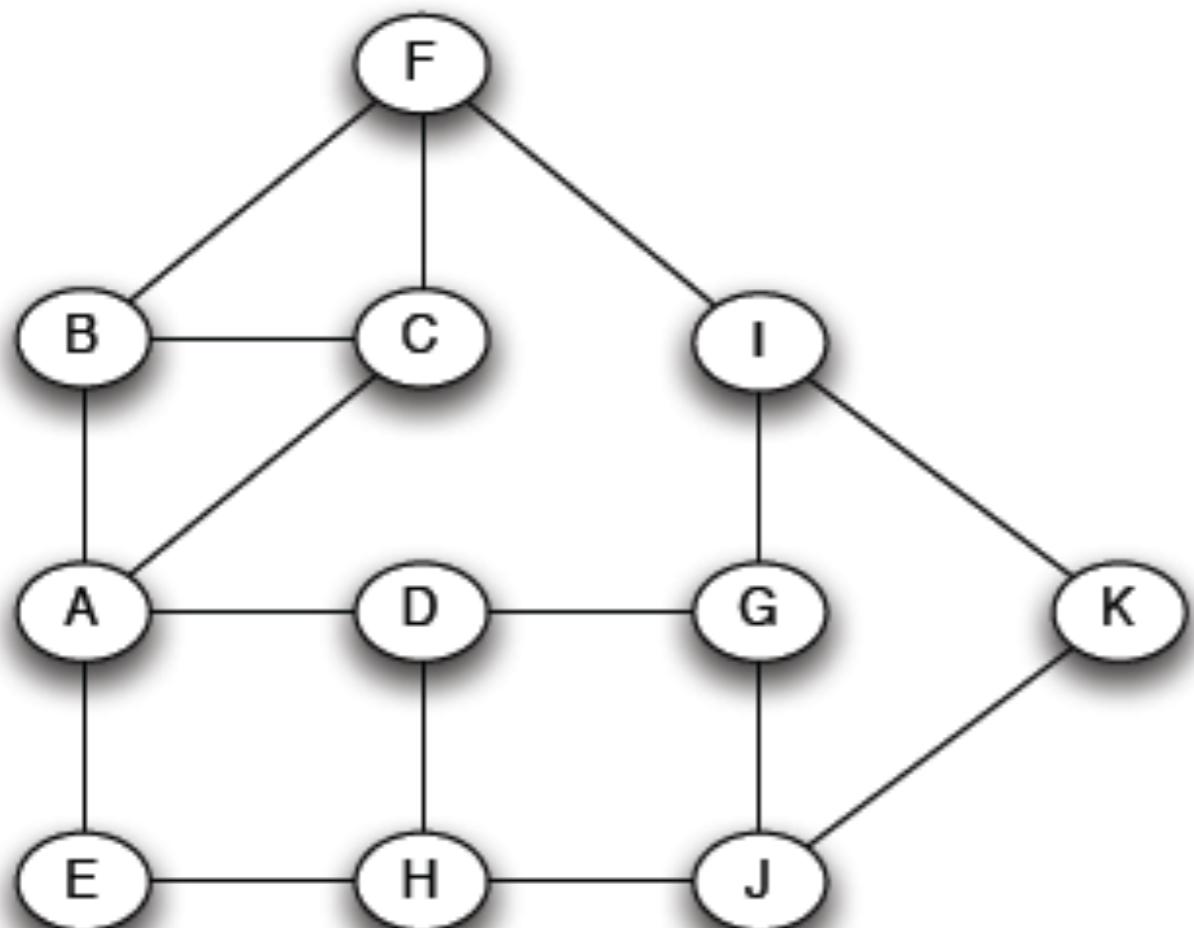
Count number of
shortest paths from
(E) to each node



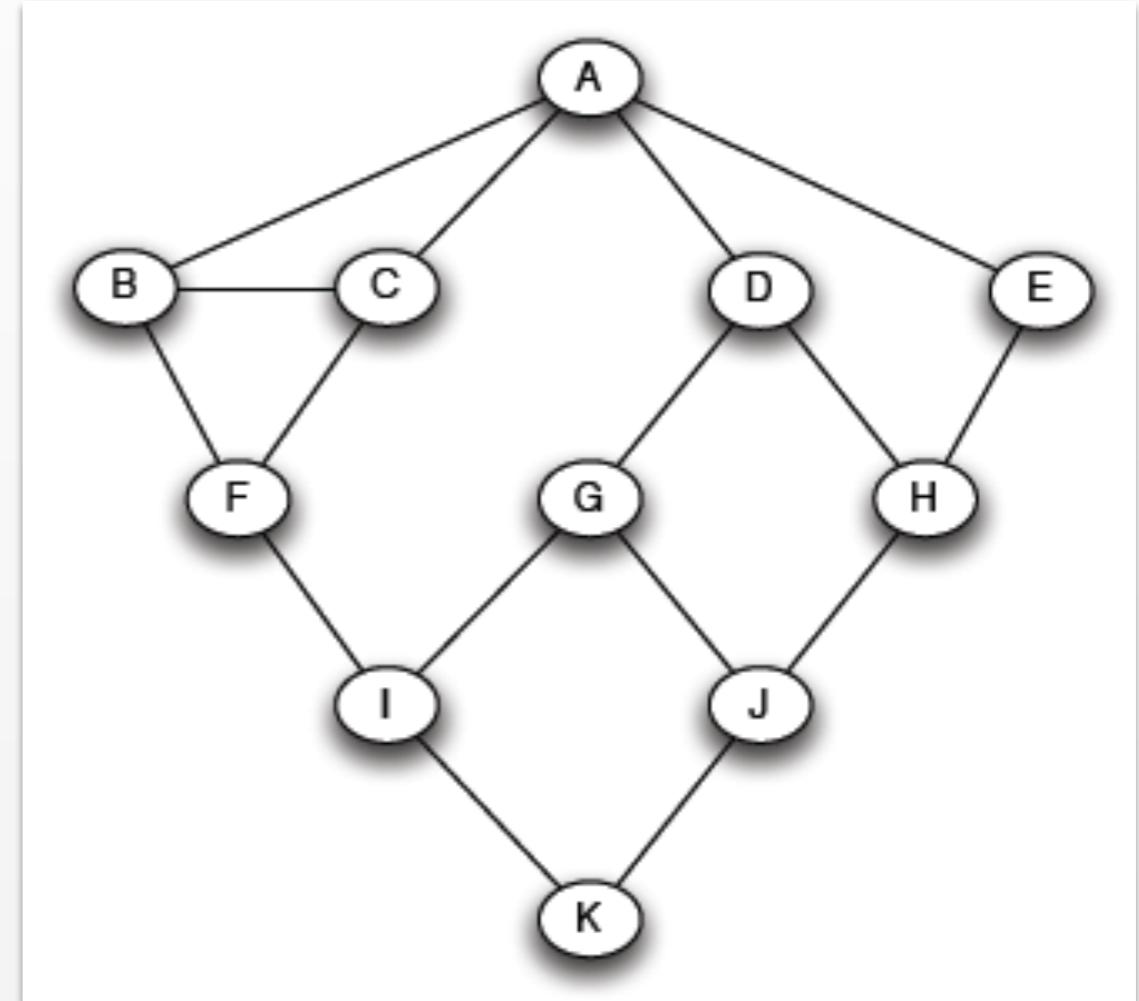
Accumulate credit
upwards, dividing
across shortest paths

Counting Paths: Larger Example

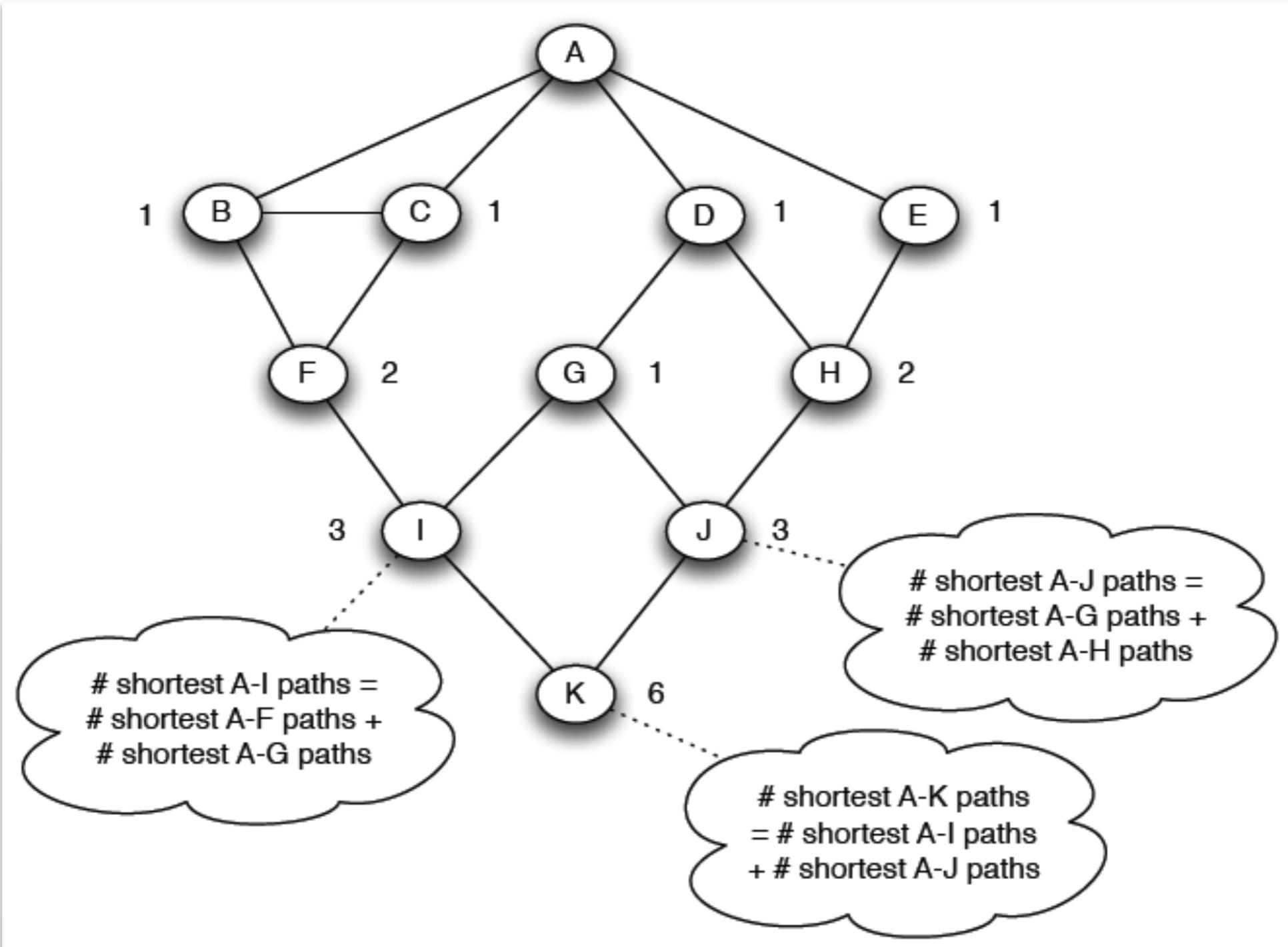
Original Graph



Breadth-first Ordering from A

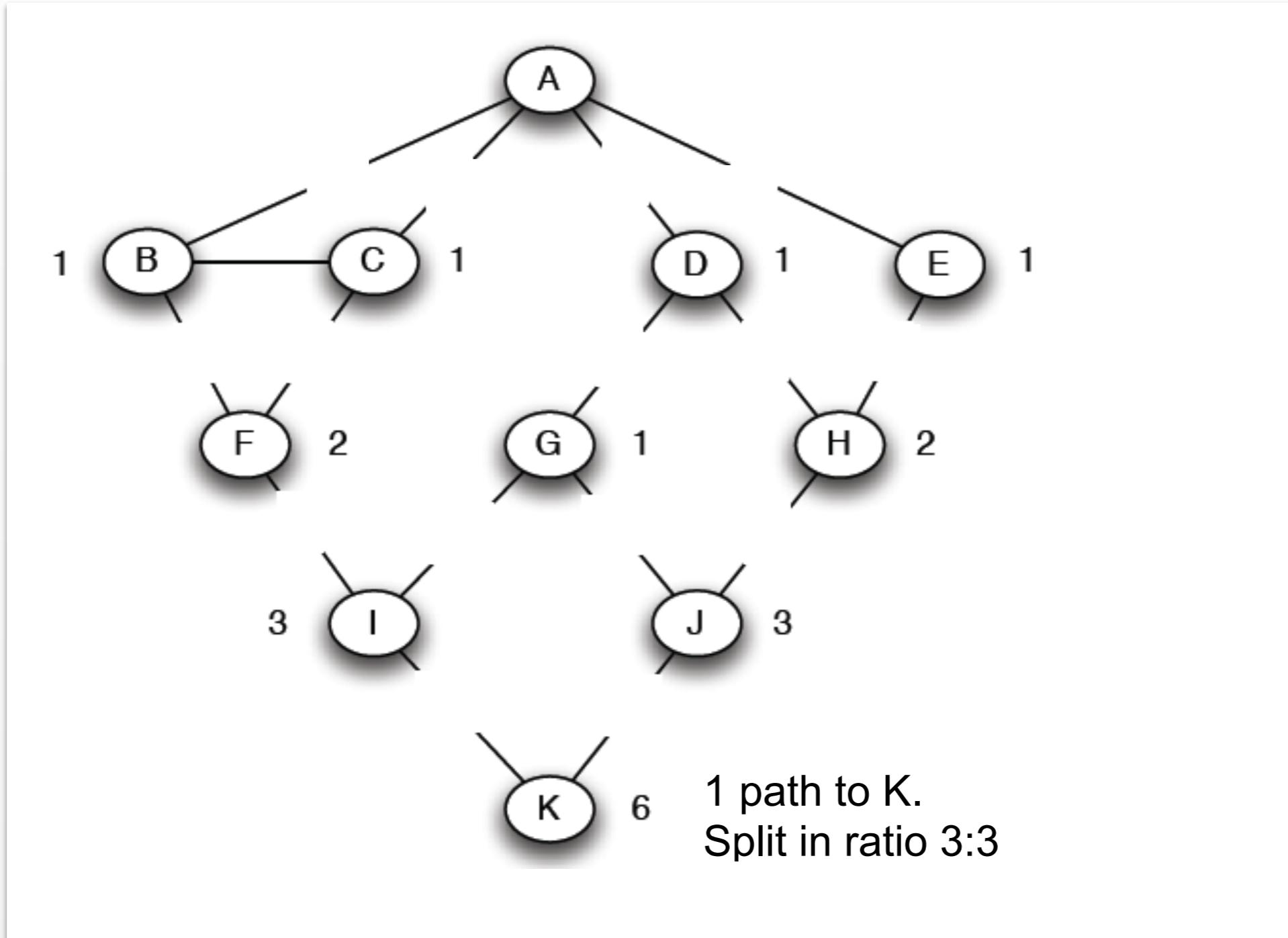


Counting Paths: Larger Example



Step 1. Count number of shortest paths from to each node

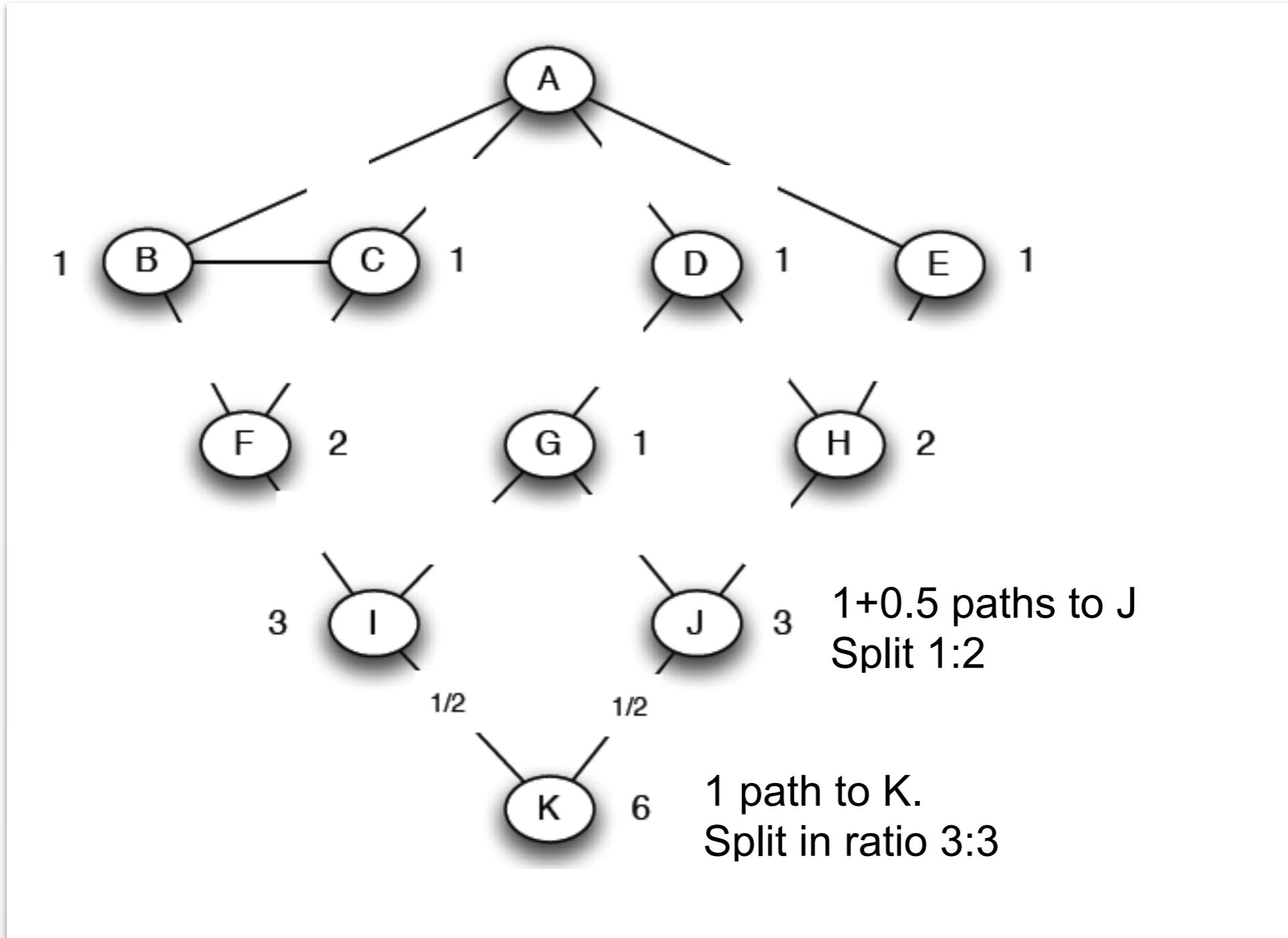
Counting Paths: Larger Example



Step 2. Propagate credit upwards, splitting according to number of paths to parents

(Adapted from: Mining of Massive Datasets, <http://www.mmds.org>)

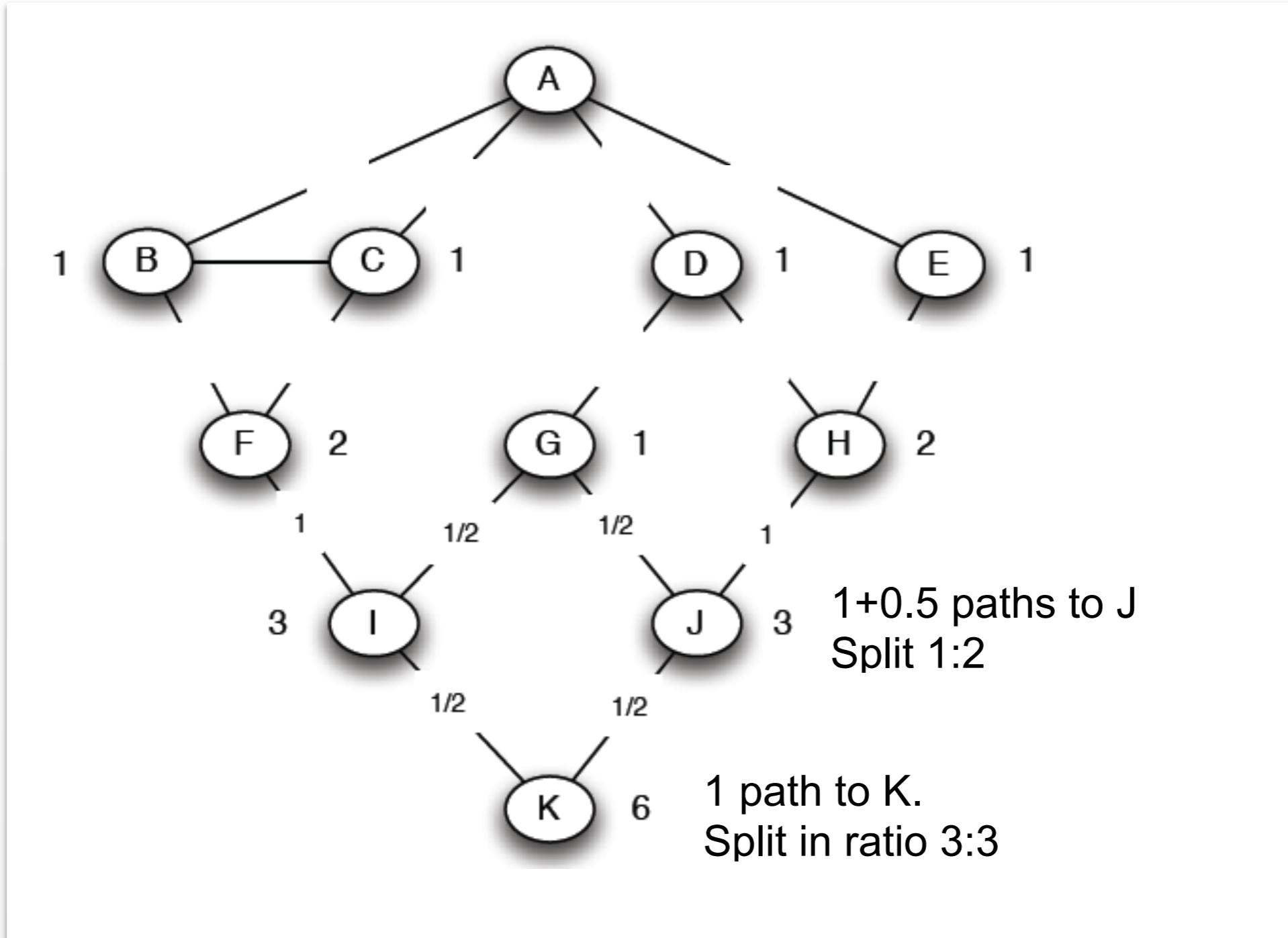
Counting Paths: Larger Example



Step 2. Propagate credit upwards, splitting according to number of paths to parents

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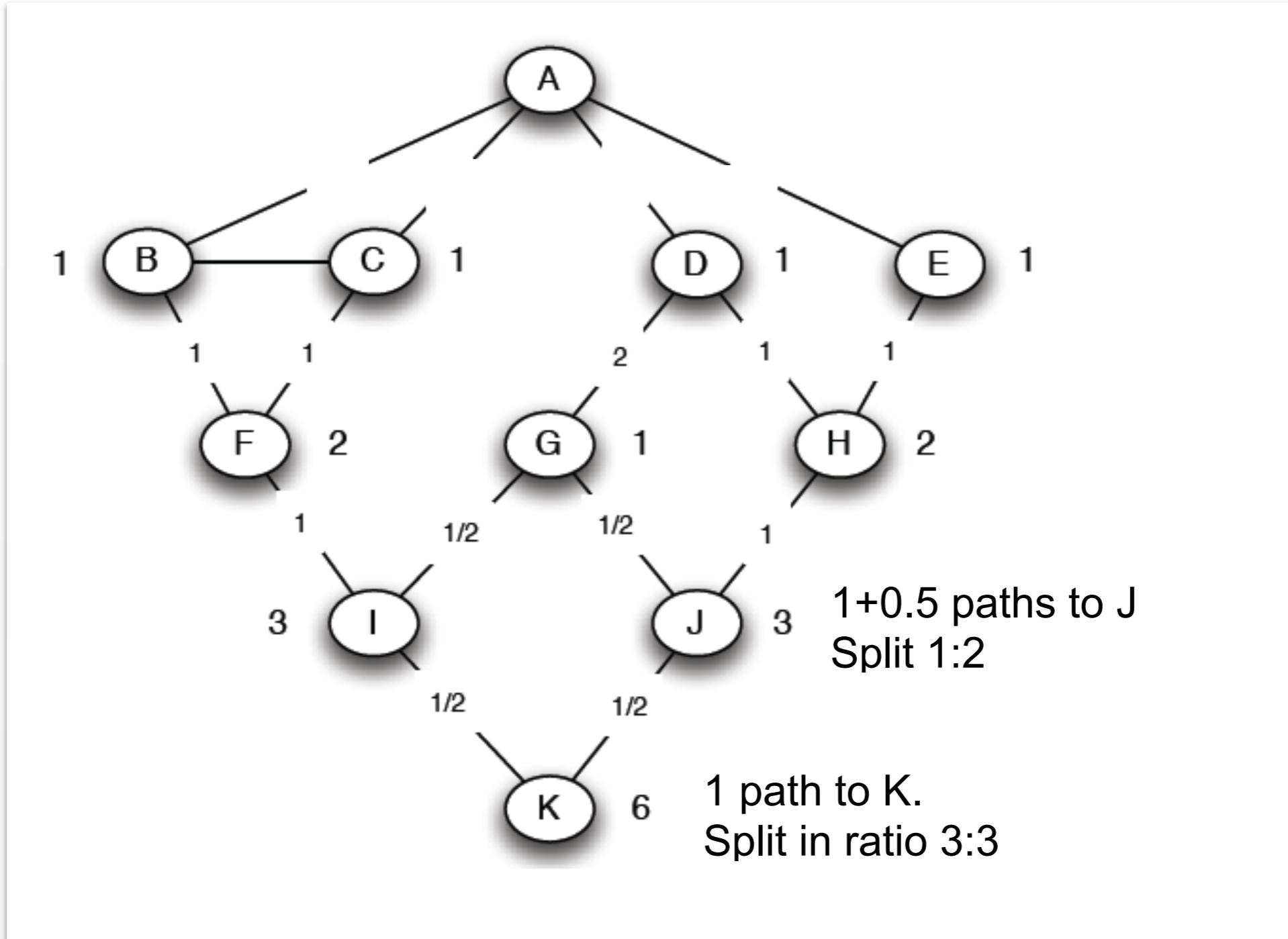
Counting Paths: Larger Example



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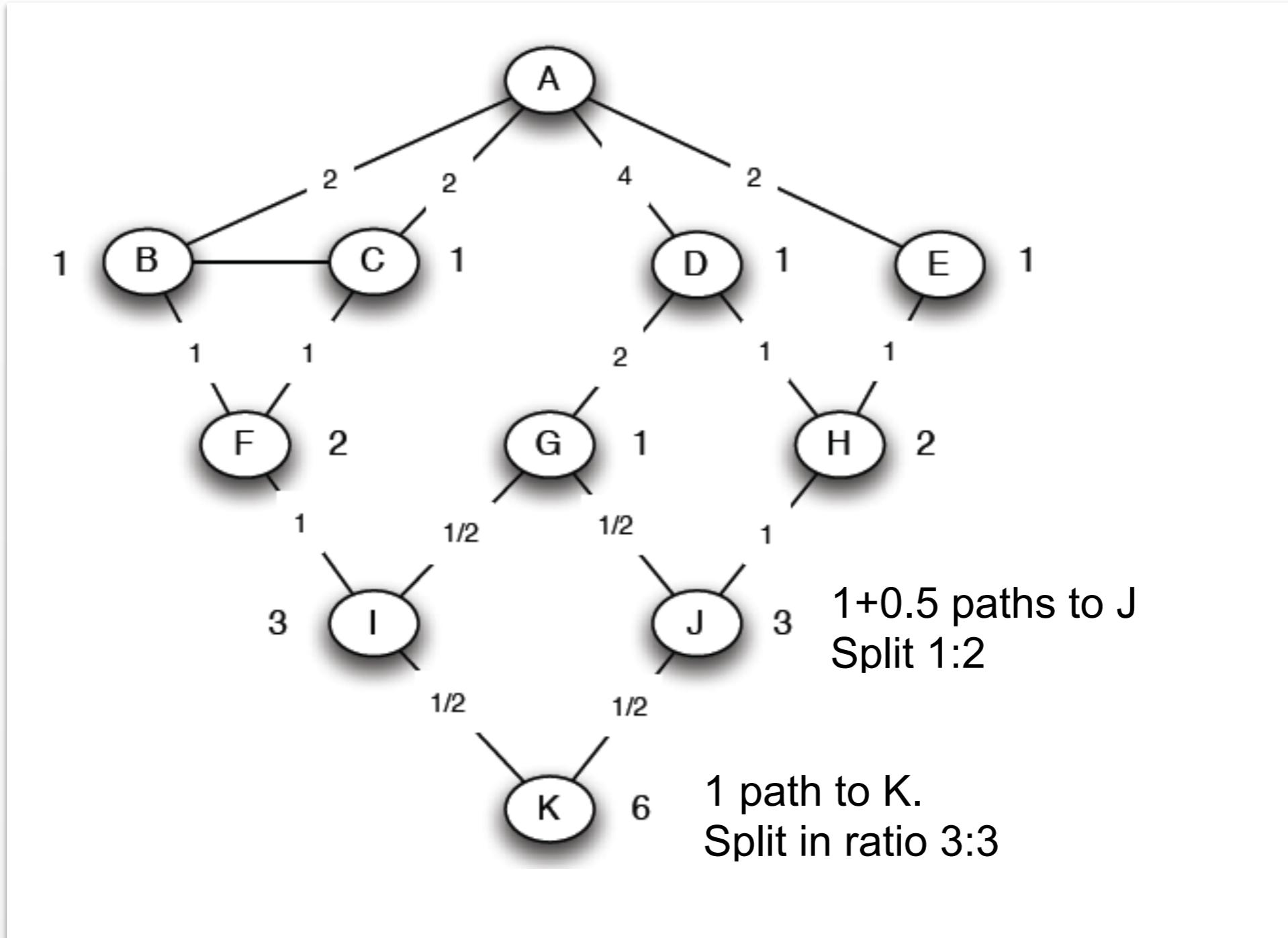
Counting Paths: Larger Example



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Counting Paths: Larger Example

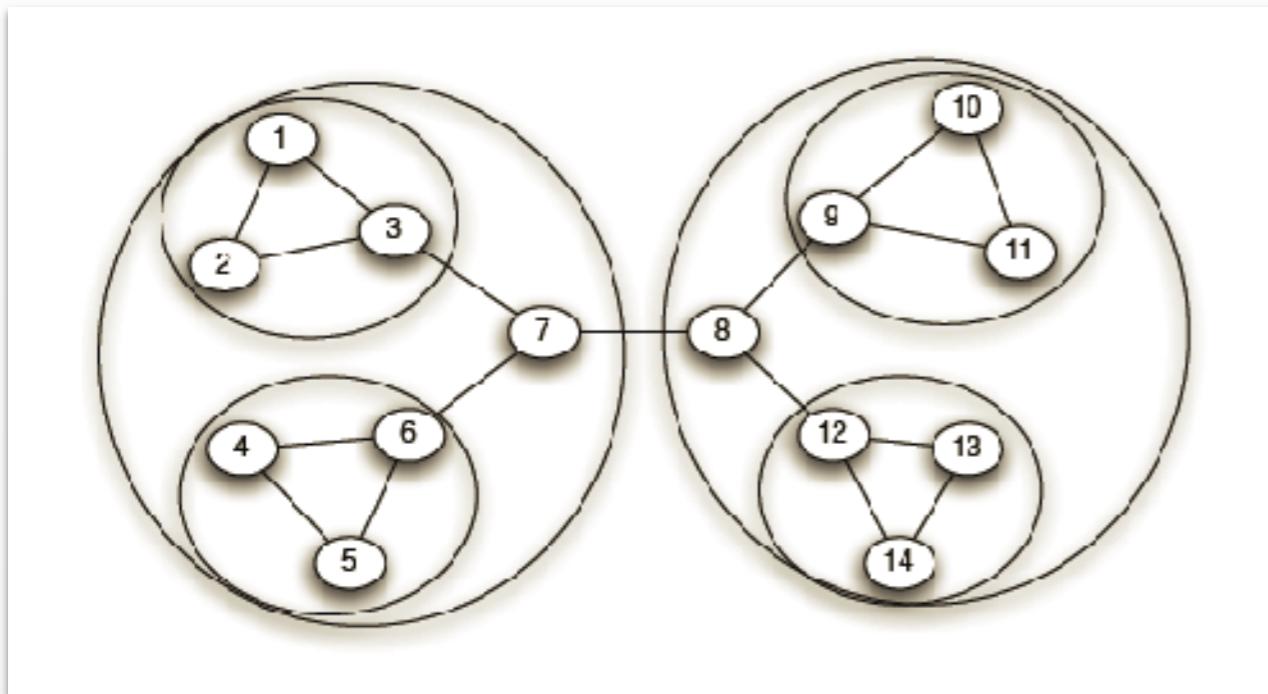


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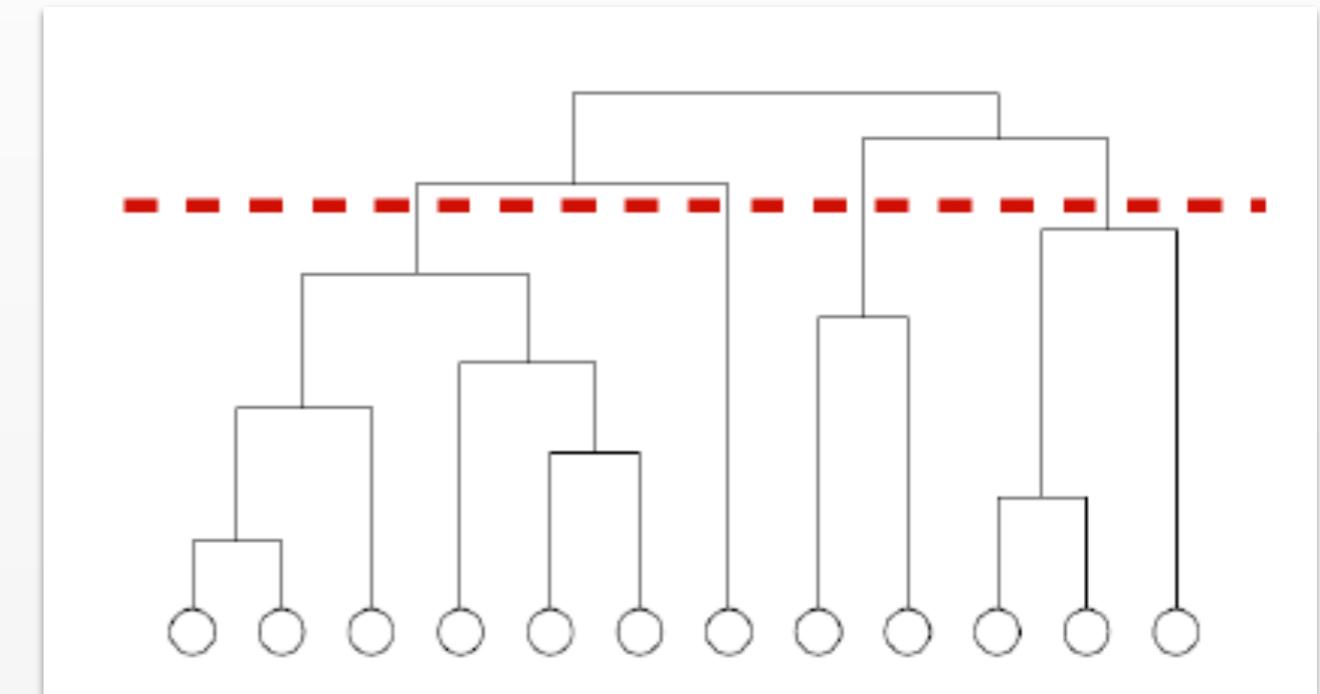
(Adapted from: Mining of Massive Datasets, <http://www.mmds.org>)

Determining the Number of Communities

Hierarchical decomposition



Choosing a cut-off



Analogous problem to deciding on number
of clusters in hierarchical clustering

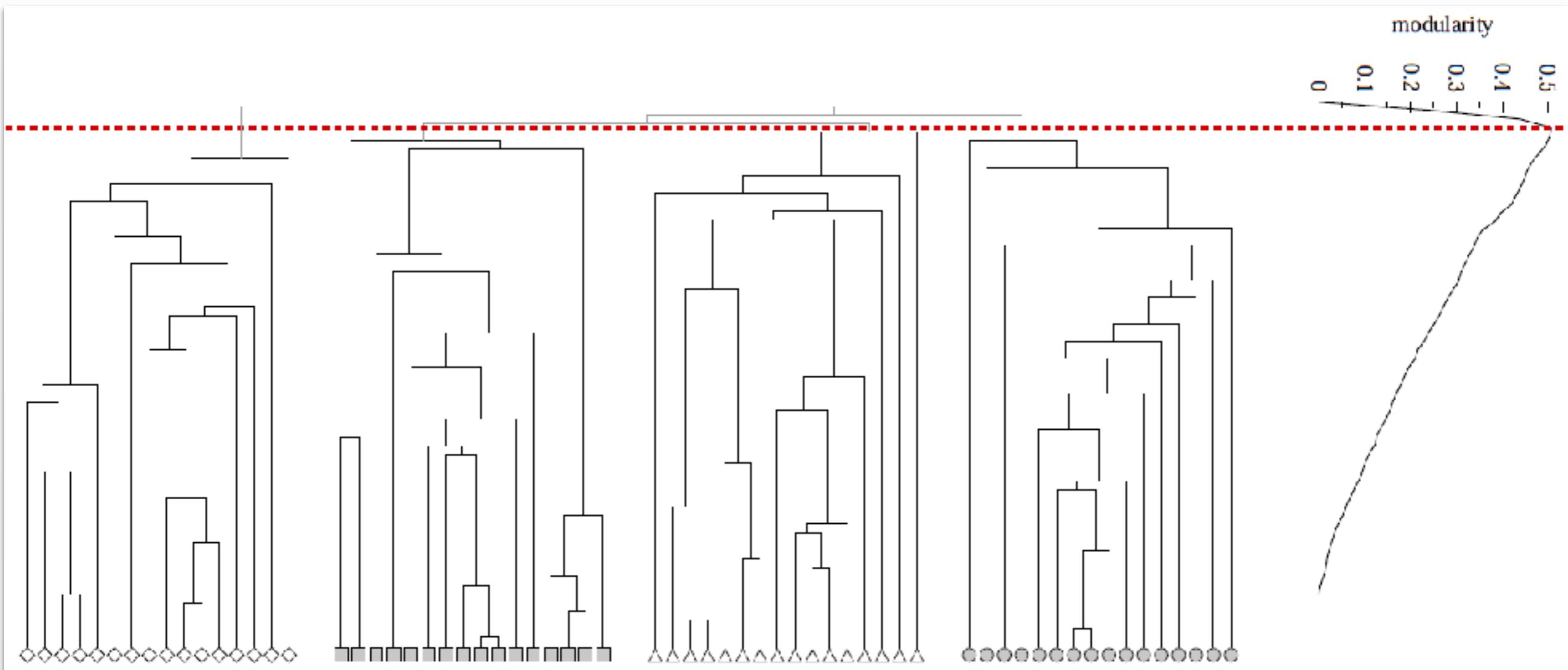
Modularity

Idea: Compare fraction of edges within module to fraction that would be observed for random connections

$$Q = \frac{1}{2m} \sum_{uv} \left[A_{vw} - \frac{k_v k_w}{2m} \right] \delta(c_u, c_v)$$

- m : Number of edges in graph
- A_{uv} : Adjacency matrix (1 if edge exists 0 otherwise)
- k_u : Degree of node u
- c_u : Cluster assignment for node u

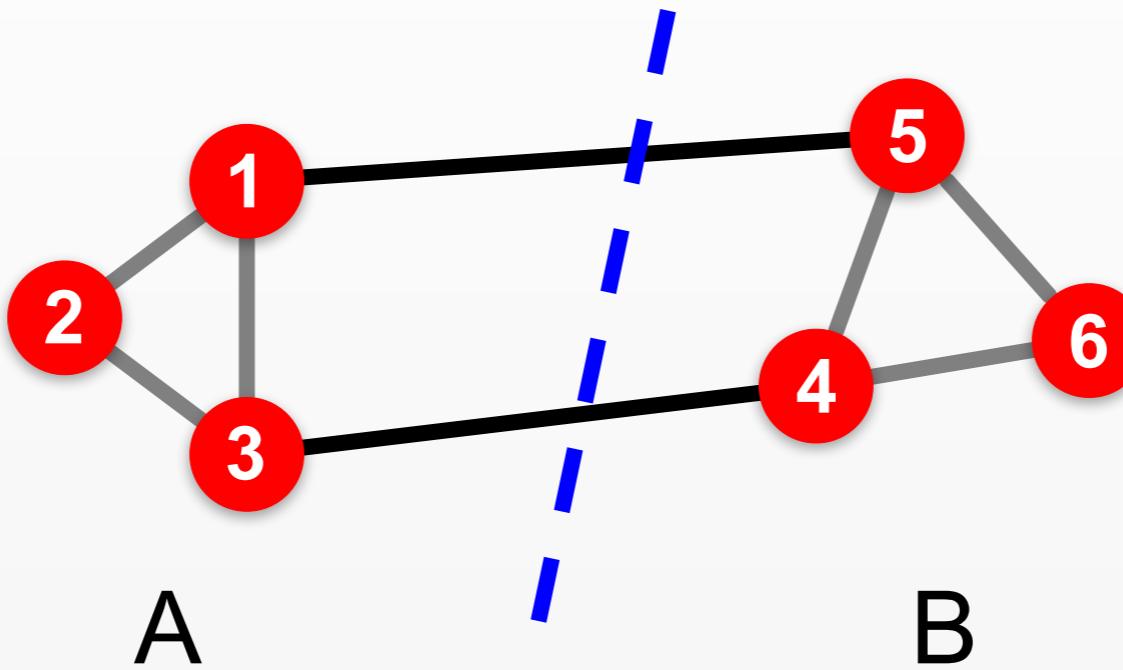
Modularity



Use modularity to optimize connectivity within modules

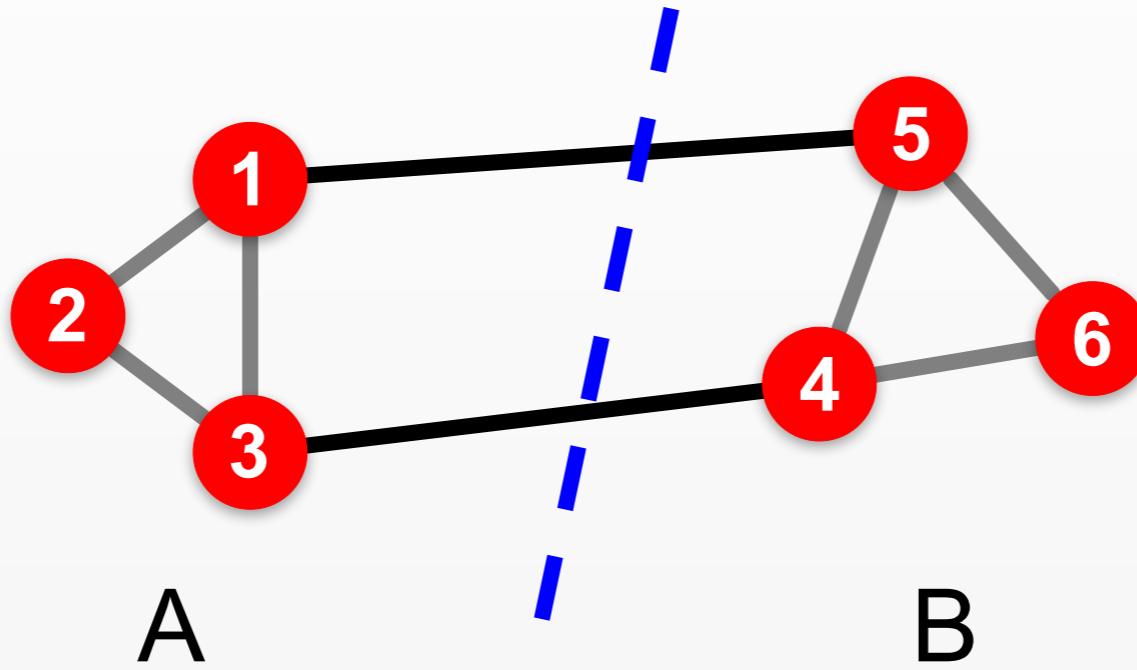
Spectral Clustering

Graph Partitioning



- What makes a good partition?
 - Maximize the within-group connections
 - Minimize the between-group connections

Graph Cuts



Degree

$$d_i = \sum_j A_{ij}$$

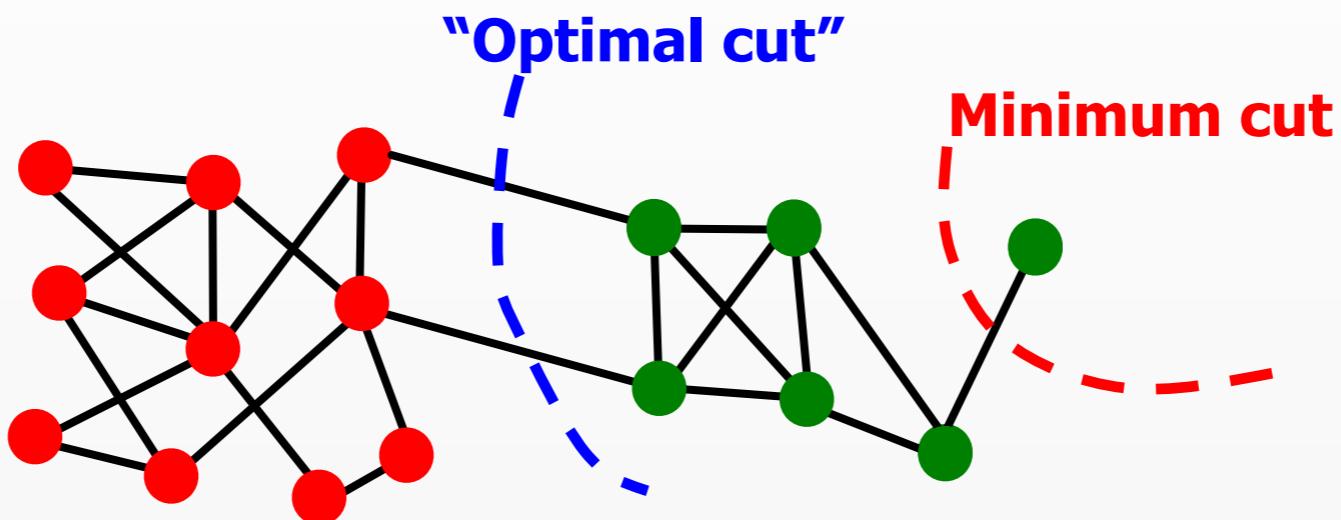
Volume

$$\text{vol}(A) = \sum_j d_i$$

Cut

$$\text{cut}(A, B) = \sum_{i \in A, j \in B} A_{ij}$$

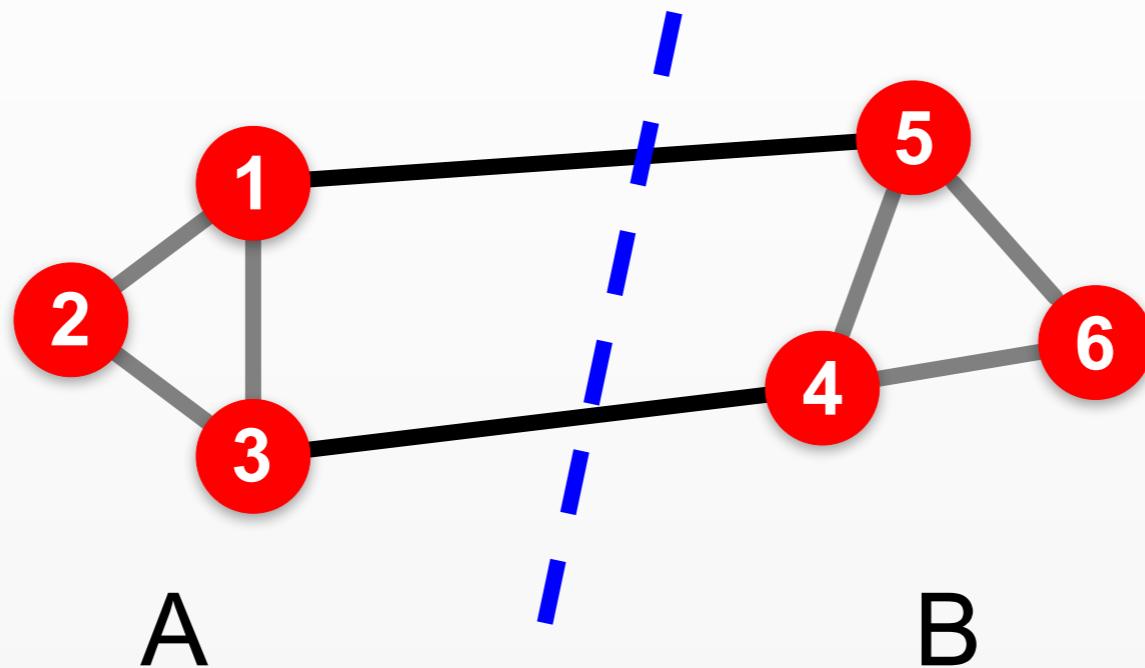
Minimal Cuts



$$\arg \min_{A,B} \text{cut}(A,B)$$

Problem: minimal cut is not necessarily a good splitting criterion

Normalized Cuts



$$\text{ncut}(A, B) = \frac{\text{cut}(A, B)}{\text{vol}(A)} + \frac{\text{cut}(A, B)}{\text{vol}(B)}$$

Degree

$$d_i = \sum_j A_{ij}$$

Volume

$$\text{vol}(A) = \sum_j d_i$$

Cut

$$\text{cut}(A, B) = \sum_{i \in A, j \in B} A_{ij}$$

Find Optimal Cut [Fiedler'73]

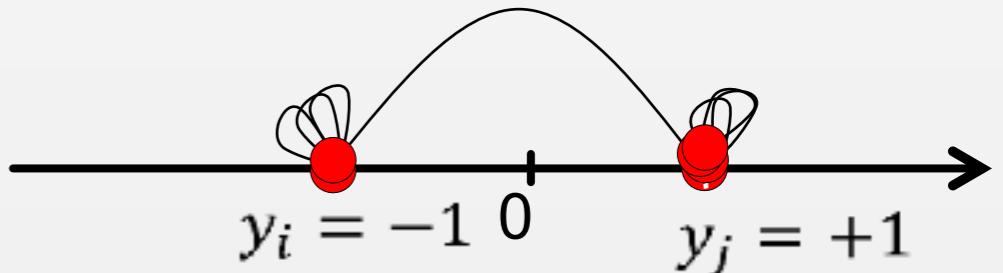
- Back to finding the optimal cut
- Express partition (A, B) as a vector

$$y_i = \begin{cases} +1 & \text{if } i \in A \\ -1 & \text{if } i \in B \end{cases}$$

- We can minimize the cut of the partition by finding a non-trivial vector x that **minimizes**:

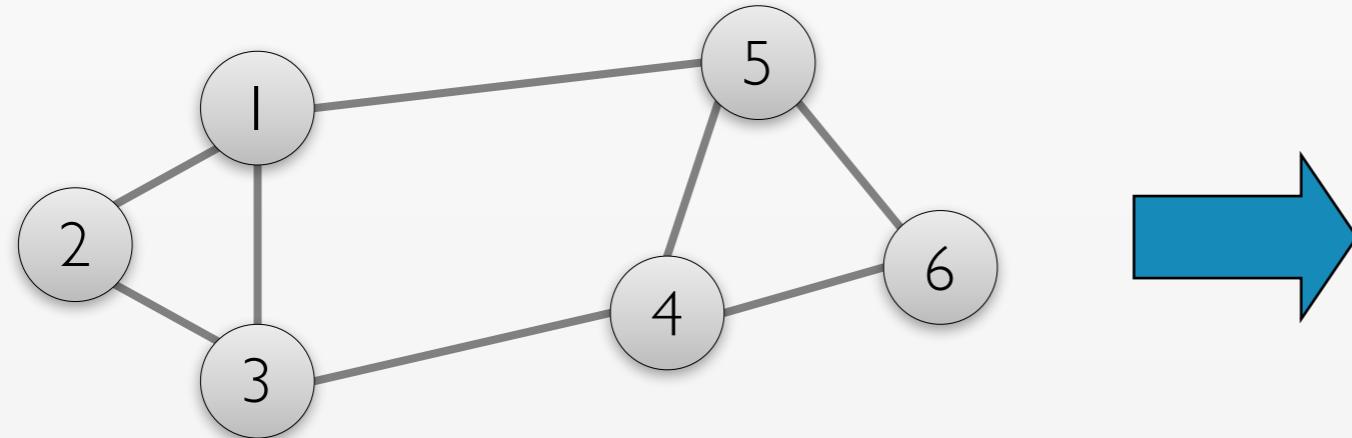
$$y^* = \operatorname{argmin}_{y \in \{-1, 1\}^n} \sum_{(i,j) \in E} (y_i - y_j)^2$$

Can't solve exactly. Let's relax y and allow it to take any real value.



Matrix Representations

- Adjacency matrix (A):
 - $n \times n$ matrix
 - $A = [a_{ij}]$, $a_{ij} = 1$ if edge between node i and j

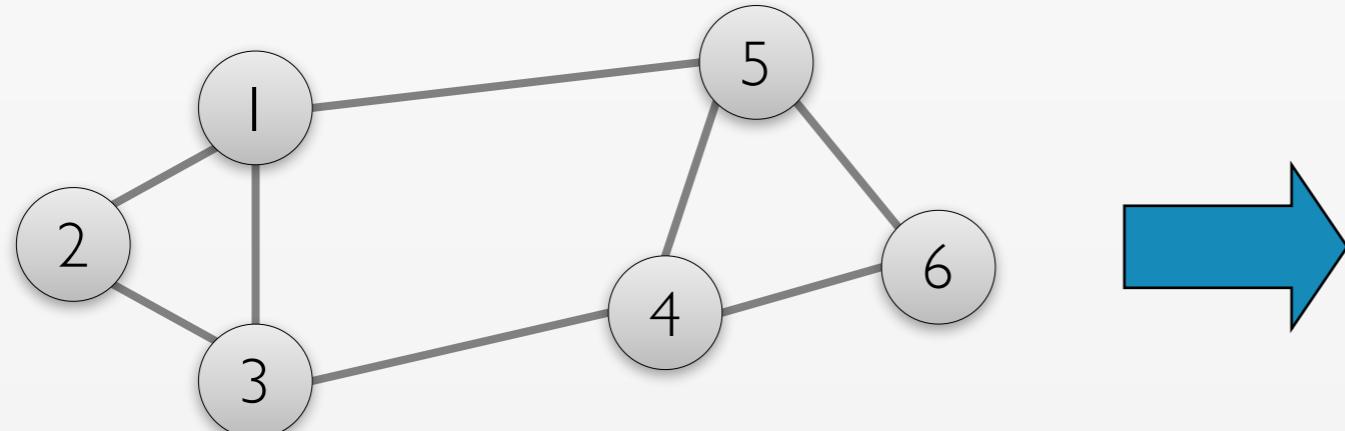


	1	2	3	4	5	6
1	0	1	1	0	1	0
2	1	0	1	0	0	0
3	1	1	0	1	0	0
4	0	0	1	0	1	1
5	1	0	0	1	0	1
6	0	0	0	1	1	0

- Important properties:
 - Symmetric matrix
 - Eigenvectors are real and orthogonal

Matrix Representations

- Degree matrix (D):
 - $n \times n$ diagonal matrix
 - $D = [d_{ii}]$, d_{ii} = degree of node i

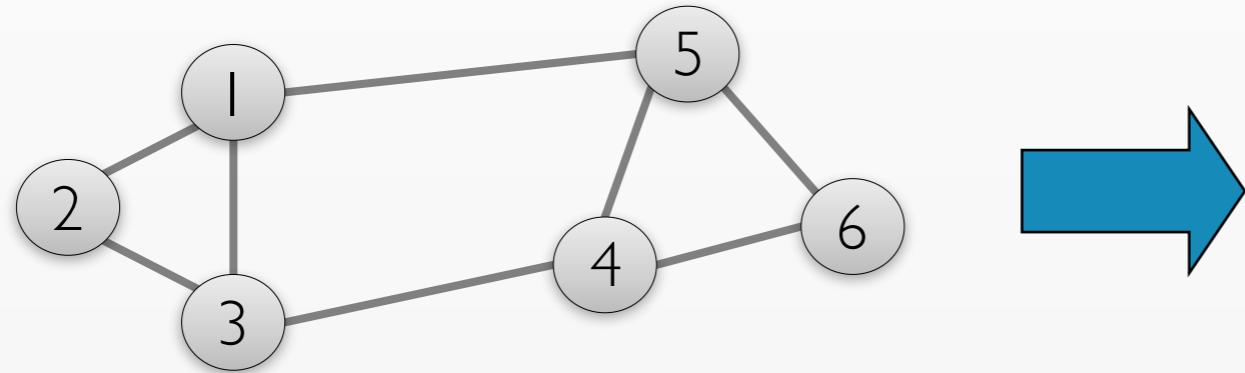


	1	2	3	4	5	6
1	3	0	0	0	0	0
2	0	2	0	0	0	0
3	0	0	3	0	0	0
4	0	0	0	3	0	0
5	0	0	0	0	3	0
6	0	0	0	0	0	2

Matrix Representations

■ Laplacian matrix (L):

- $n \times n$ symmetric matrix



	1	2	3	4	5	6
1	3	-1	-1	0	-1	0
2	-1	2	-1	0	0	0
3	-1	-1	3	-1	0	0
4	0	0	-1	3	-1	-1
5	-1	0	0	-1	3	-1
6	0	0	0	-1	-1	2

■ What is trivial eigenpair?

- $x = (1, \dots, 1)$ then $L \cdot x = \mathbf{0}$ and so $\lambda = \lambda_1 = 0$

■ Important properties:

- **Eigenvalues** are non-negative real numbers
- **Eigenvectors** are real and orthogonal

Second Eigenvalue

- Fact: For symmetric matrix M :

$$\lambda_2 = \min_x \frac{x^T M x}{x^T x}$$

- What is the meaning of $\min x^T L x$ on G ?

- $x^T L x = \sum_{i,j=1}^n L_{ij} x_i x_j = \sum_{i,j=1}^n (D_{ij} - A_{ij}) x_i x_j$
- $= \sum_i D_{ii} x_i^2 - \sum_{(i,j) \in E} 2x_i x_j$
- $= \sum_{(i,j) \in E} \underbrace{(x_i^2 + x_j^2 - 2x_i x_j)}_{\text{Node } i \text{ has degree } d_i. \text{ So, value } x_i^2 \text{ needs to be summed up } d_i \text{ times.}} = \sum_{(i,j) \in E} (\mathbf{x}_i - \mathbf{x}_j)^2$

Node i has degree d_i . So, value x_i^2 needs to be summed up d_i times.
But each edge (i,j) has two endpoints so we need $x_i^2 + x_j^2$

Second Eigenvector of Laplacian

■ What else do we know about x ?

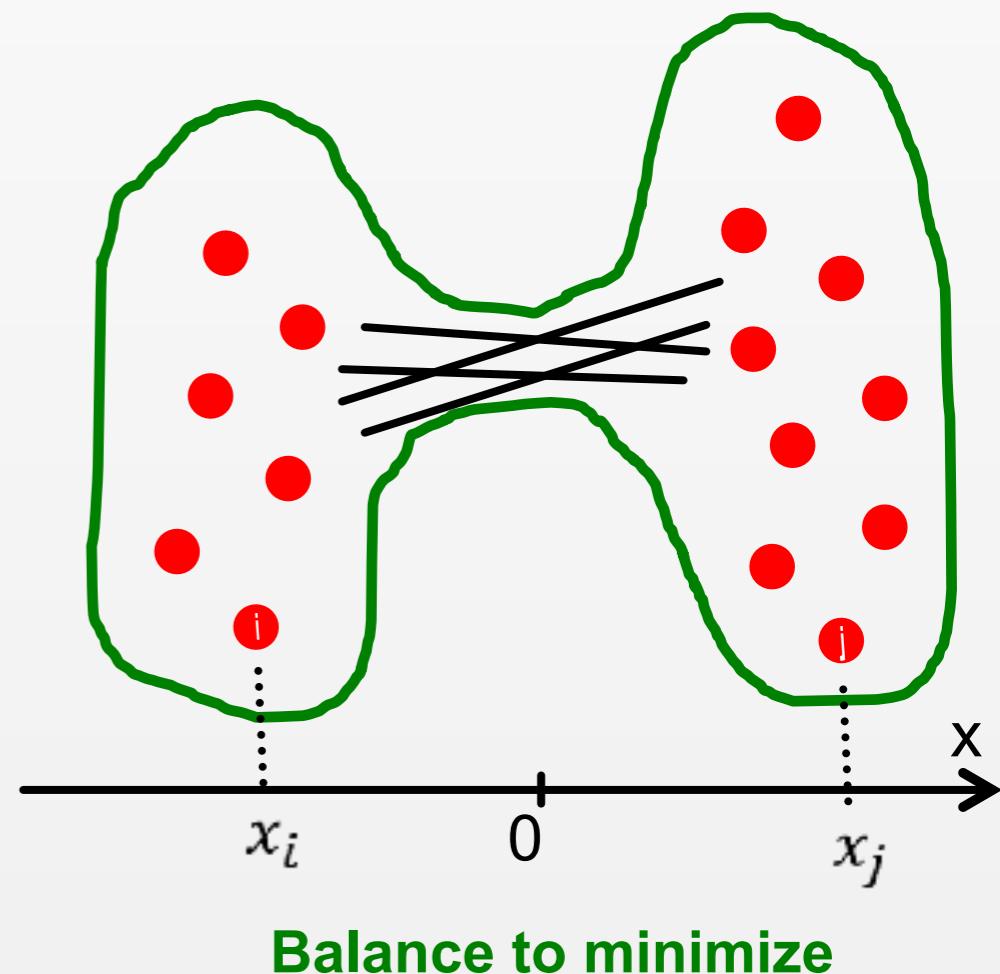
- x is unit vector: $\sum_i x_i^2 = 1$
- x is orthogonal to 1st eigenvector $(1, \dots, 1)$ thus:
 $\sum_i x_i \cdot 1 = \sum_i x_i = 0$

■ Remember:

$$\lambda_2 = \min_{\substack{\text{All labelings} \\ \text{of nodes } i \text{ so} \\ \text{that } \sum x_i = 0}} \frac{\sum_{(i,j) \in E} (x_i - x_j)^2}{\sum_i x_i^2}$$

We want to assign values x_i to nodes i such that few edges cross 0.

(we want x_i and x_j to subtract each other)



Rayleigh Theorem

$$\min_{y \in \mathbb{R}^n} f(y) = \sum_{(i,j) \in E} (y_i - y_j)^2 = y^T L y$$

- $\lambda_2 = \min_y f(y)$: The minimum value of $f(y)$ is given by the 2nd smallest eigenvalue λ_2 of the Laplacian matrix L
- $x = \arg \min_y f(y)$: The optimal solution for y is given by the corresponding eigenvector x , referred as the **Fiedler vector**

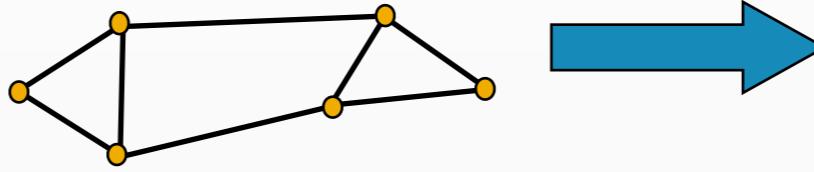
Spectral Clustering Algorithms

- Three basic stages:
 - 1) Pre-processing
 - Construct a matrix representation of the graph
 - More generally, construct similarity matrix
 - 2) Decomposition
 - Compute eigenvalues and eigenvectors of the matrix
 - Map each point to a lower-dimensional representation based on one or more eigenvectors
 - 3) Grouping
 - Assign points to two or more clusters, based on the new representation

Spectral Partitioning Algorithm

- 1) Pre-processing:

- Build Laplacian matrix \mathbf{L} of the graph



	1	2	3	4	5	6
1	3	-1	-1	0	-1	0
2	-1	2	-1	0	0	0
3	-1	-1	3	-1	0	0
4	0	0	-1	3	-1	-1
5	-1	0	0	-1	3	-1
6	0	0	0	-1	-1	2

- 2) Decomposition:

- Find eigenvalues λ and eigenvectors \mathbf{x} of the matrix \mathbf{L}

$$\lambda = \begin{matrix} 0.0 \\ 1.0 \\ 3.0 \\ 3.0 \\ 4.0 \\ 5.0 \end{matrix}$$

$$\mathbf{x} = \begin{matrix} 0.4 & 0.3 & -0.5 & -0.2 & -0.4 & -0.5 \\ 0.4 & 0.6 & 0.4 & -0.4 & 0.4 & 0.0 \\ 0.4 & 0.3 & 0.1 & 0.6 & -0.4 & 0.5 \\ 0.4 & -0.3 & 0.1 & 0.6 & 0.4 & -0.5 \\ 0.4 & -0.3 & -0.5 & -0.2 & 0.4 & 0.5 \\ 0.4 & -0.6 & 0.4 & -0.4 & -0.4 & 0.0 \end{matrix}$$

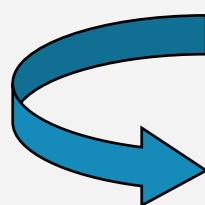
- Map vertices to corresponding components of λ_2

1	0.3
2	0.6
3	0.3
4	-0.3
5	-0.3
6	-0.6

How do we now find the clusters?

Spectral Partitioning

- 3) Grouping:
 - Sort components of reduced 1-dimensional vector
 - Identify clusters by splitting the sorted vector in two
- How to choose a splitting point?
 - Naïve approaches:
 - Split at 0 or median value
 - More expensive approaches:
 - Attempt to minimize normalized cut in 1-dimension (sweep over ordering of nodes induced by the eigenvector)



1	0.3
2	0.6
3	0.3
4	-0.3
5	-0.3
6	-0.6

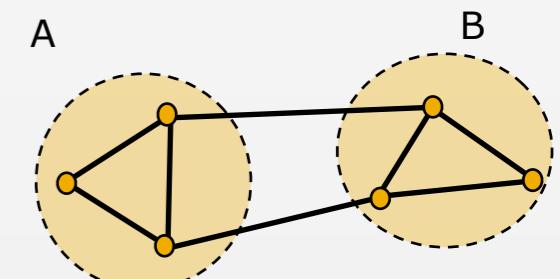


Split at 0:

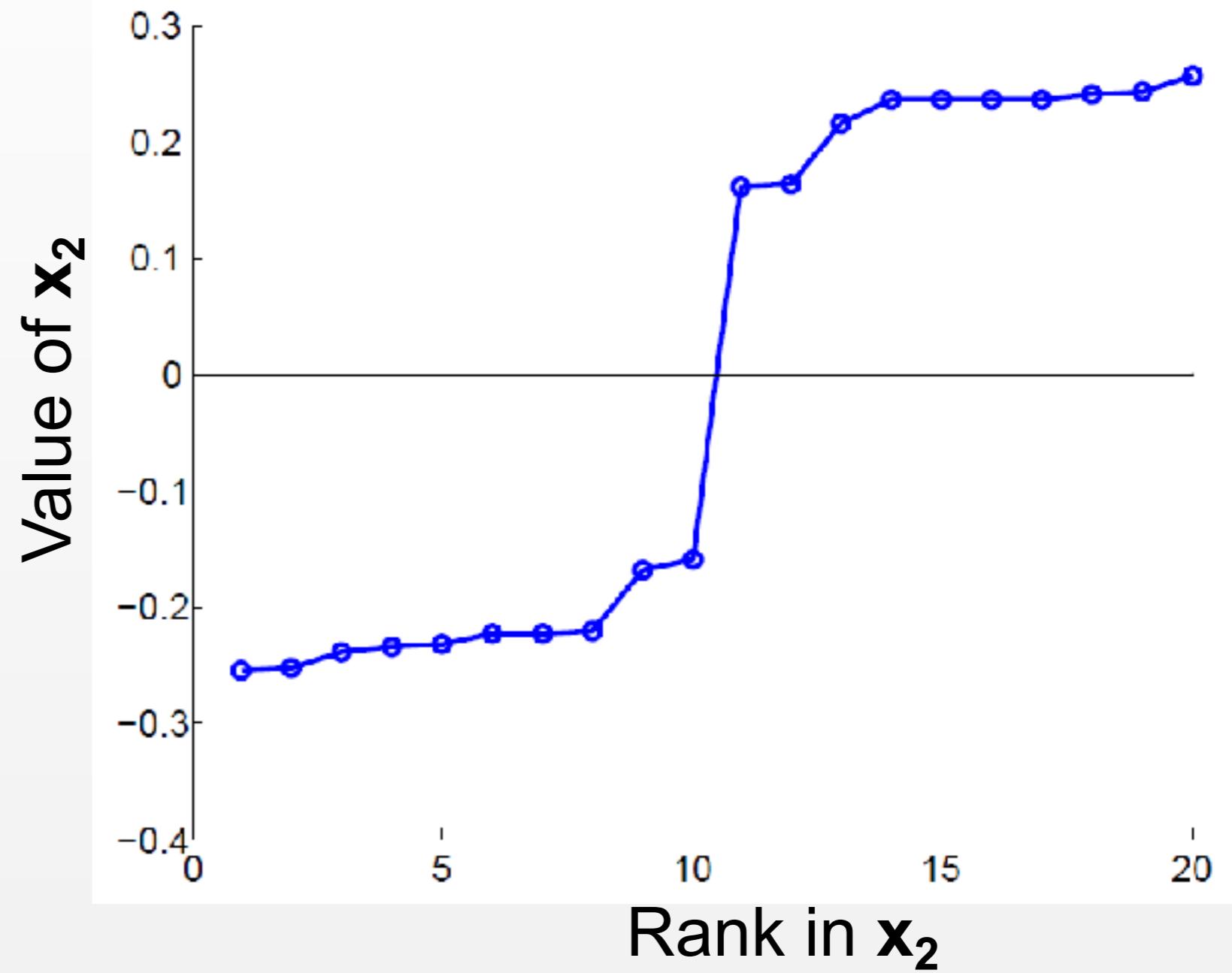
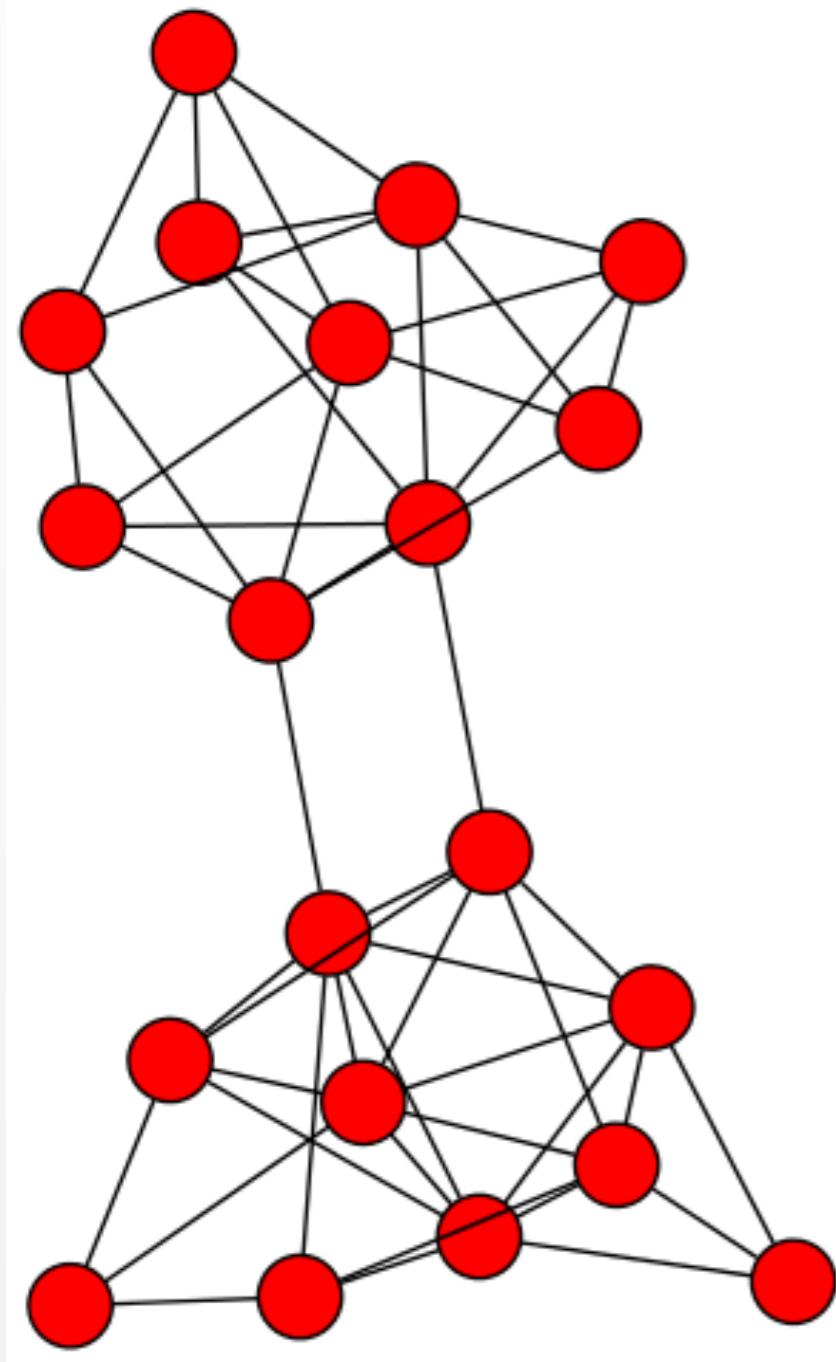
Cluster A: Positive points
Cluster B: Negative points

1	0.3
2	0.6
3	0.3

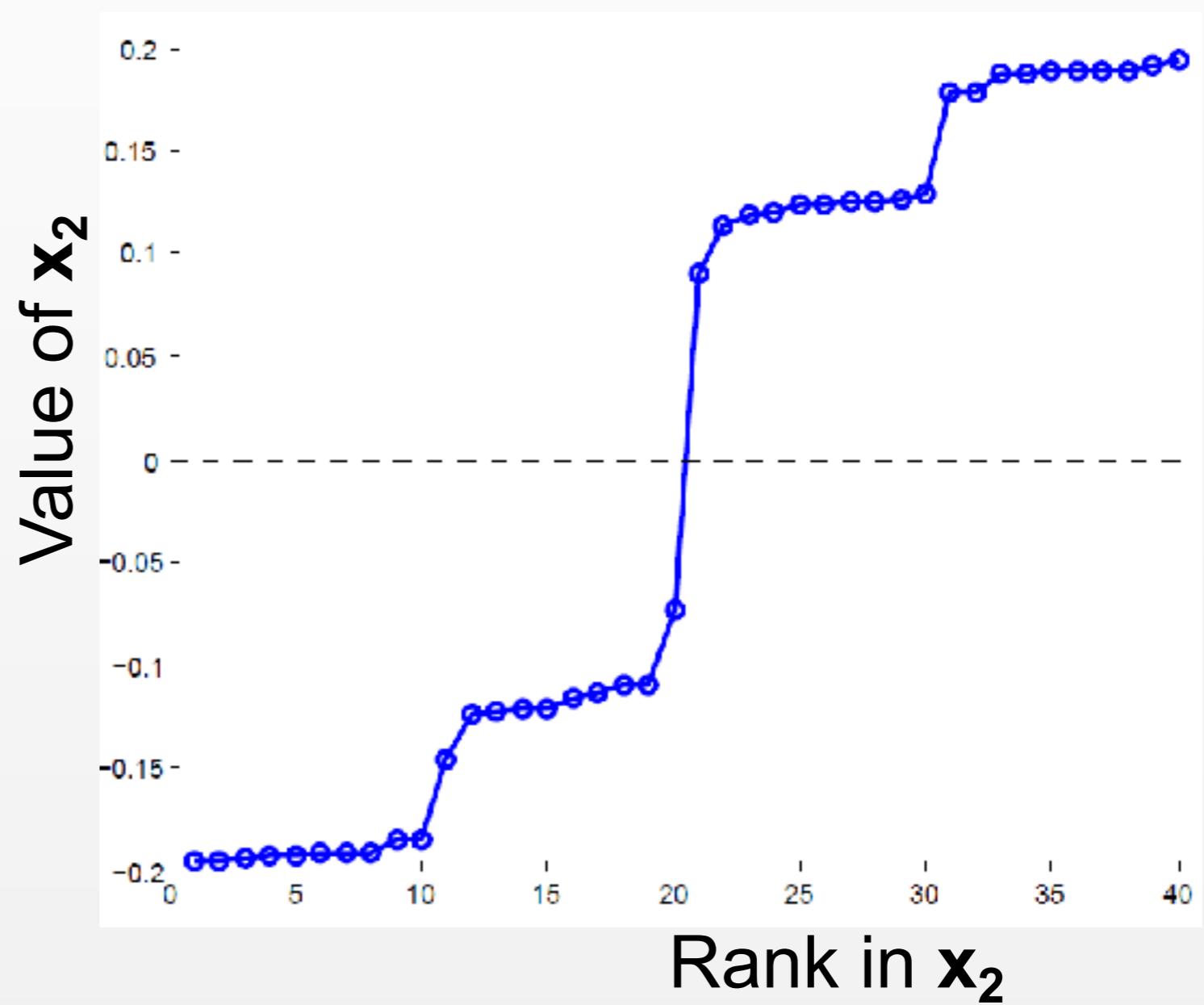
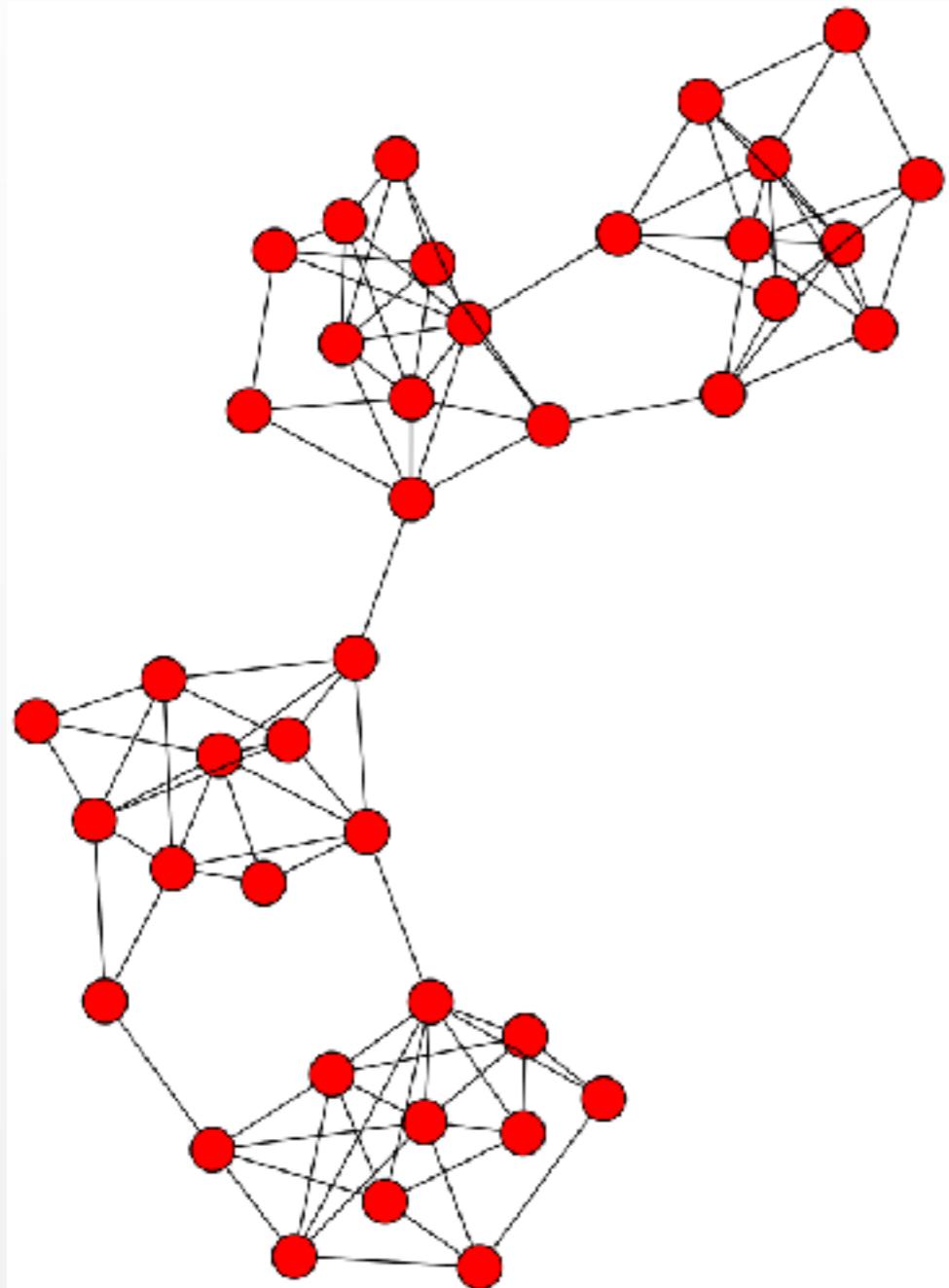
4	-0.3
5	-0.3
6	-0.6



Example: Spectral Partitioning



Example: Spectral Partitioning

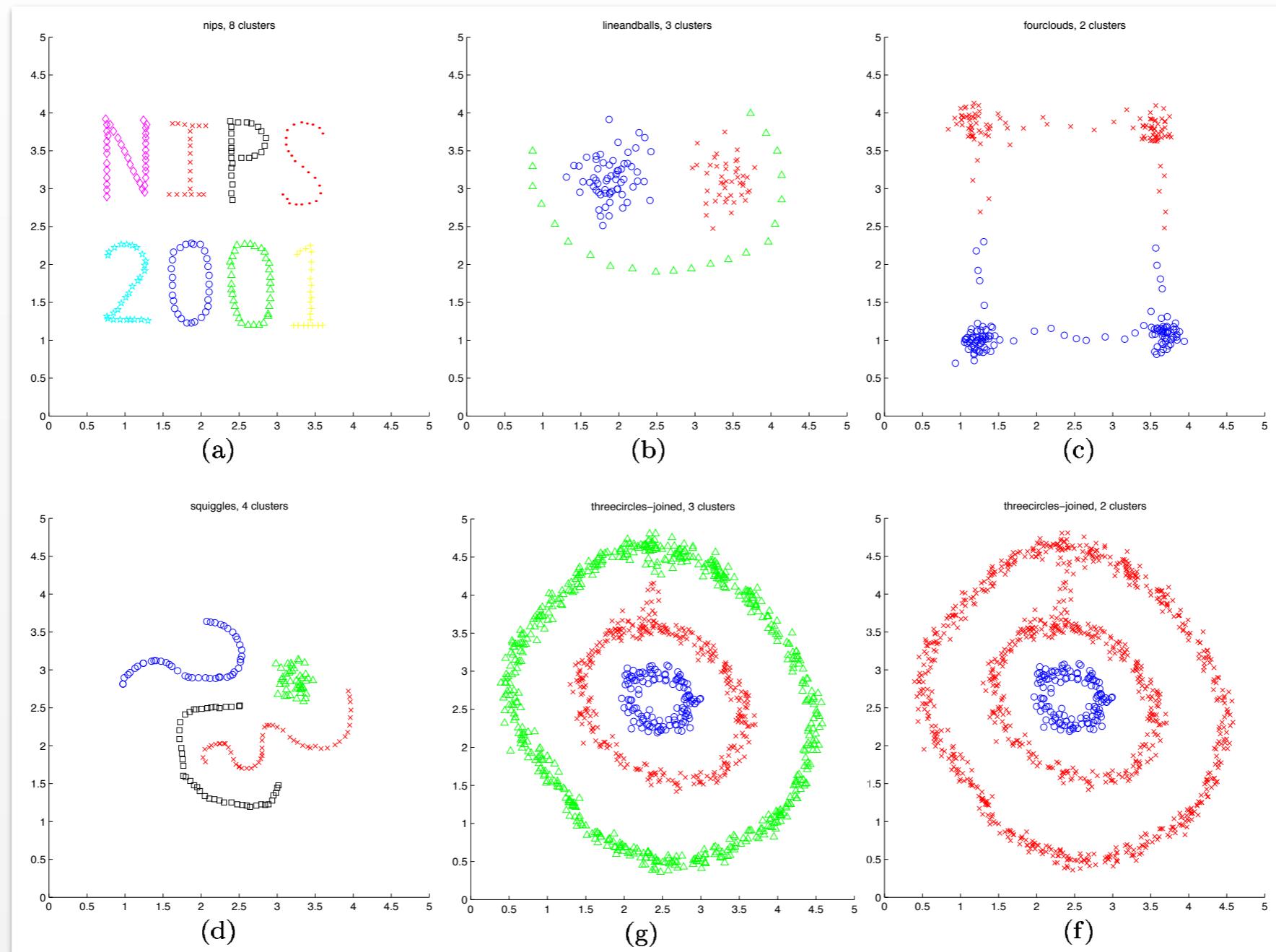


(Adapted from: Mining of Massive Datasets, <http://www.mmds.org>)

k -Way Spectral Clustering

- How do we partition a graph into k clusters?
- Two basic approaches:
 - Recursive bi-partitioning [Hagen et al., '92]
 - Recursively apply bi-partitioning algorithm in a hierarchical divisive manner
 - Disadvantages: Inefficient, unstable
 - Cluster multiple eigenvectors [Shi-Malik, '00]
 - Build a reduced space from multiple eigenvectors
 - Commonly used in recent papers

Spectral Clustering as General-purpose Method



source: Ng, Jordan and Weiss, NIPS 2001

Define “edge weight” W using some similarity metric
(e.g. a kernel function)