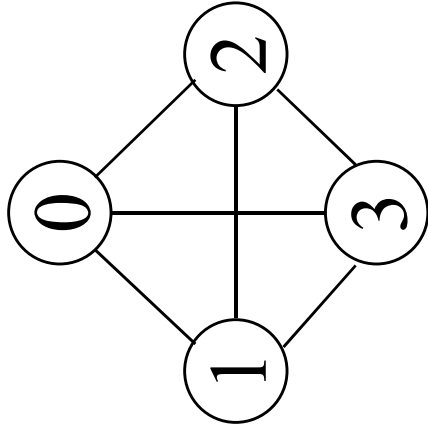


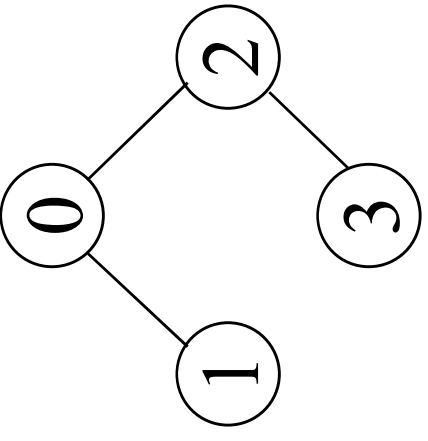
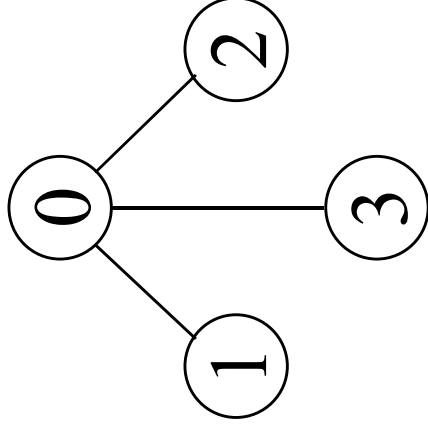
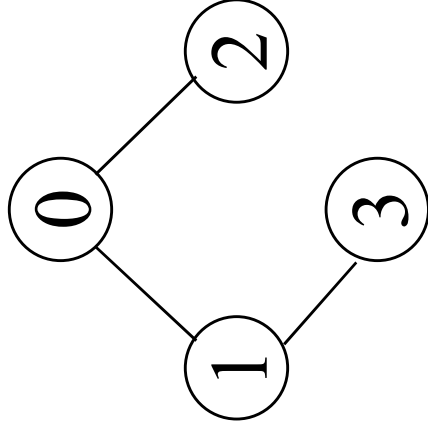
# Spanning Trees

- When graph  $G$  is connected, a depth first or breadth first search starting at any vertex will visit all vertices in  $G$
- A **spanning tree** is any tree that consists solely of edges in  $G$  and that includes all the vertices
- $E(G)$ :  $T$  (**tree edges**) +  $N$  (**nontree edges**)  
where  $T$ : set of edges used during search  
 $N$ : set of remaining edges

# Examples of Spanning Tree



$G_1$



Possible spanning trees

# Spanning Trees

- Either dfs or bfs can be used to create a spanning tree
  - When dfs is used, the resulting spanning tree is known as a **depth first spanning tree**
  - When bfs is used, the resulting spanning tree is known as a **breadth first spanning tree**
- While adding a nontree edge into any spanning tree, this will create a cycle

# Minimum Cost Spanning Tree

- The cost of a spanning tree of a weighted undirected graph is the sum of the costs of the edges in the spanning tree
- A minimum cost spanning tree is a spanning tree of least cost
- Three different algorithms can be used
  - Kruskal
  - Prim
  - Sollin

Select  $n-1$  edges from a weighted graph of  $n$  vertices with minimum cost.

# Greedy Strategy

- An optimal solution is constructed in stages
- At each stage, the best decision is made at this time
- Since this decision cannot be changed later, we make sure that the decision will result in a feasible solution
- Typically, the selection of an item at each stage is based on a least cost or a highest profit criterion

# Kruskal's Idea

- Build a minimum cost spanning tree  $T$  by adding edges to  $T$  one at a time
- Select the edges for inclusion in  $T$  in nondecreasing order of the cost
- An edge is added to  $T$  if it does not form a cycle
- Since  $G$  is connected and has  $n > 0$  vertices, exactly  $n-1$  edges will be selected

# Examples for Kruskal's Algorithm

0 ~~10~~ 5

2 ~~12~~ 3

1 ~~14~~ 6

1 ~~16~~ 2

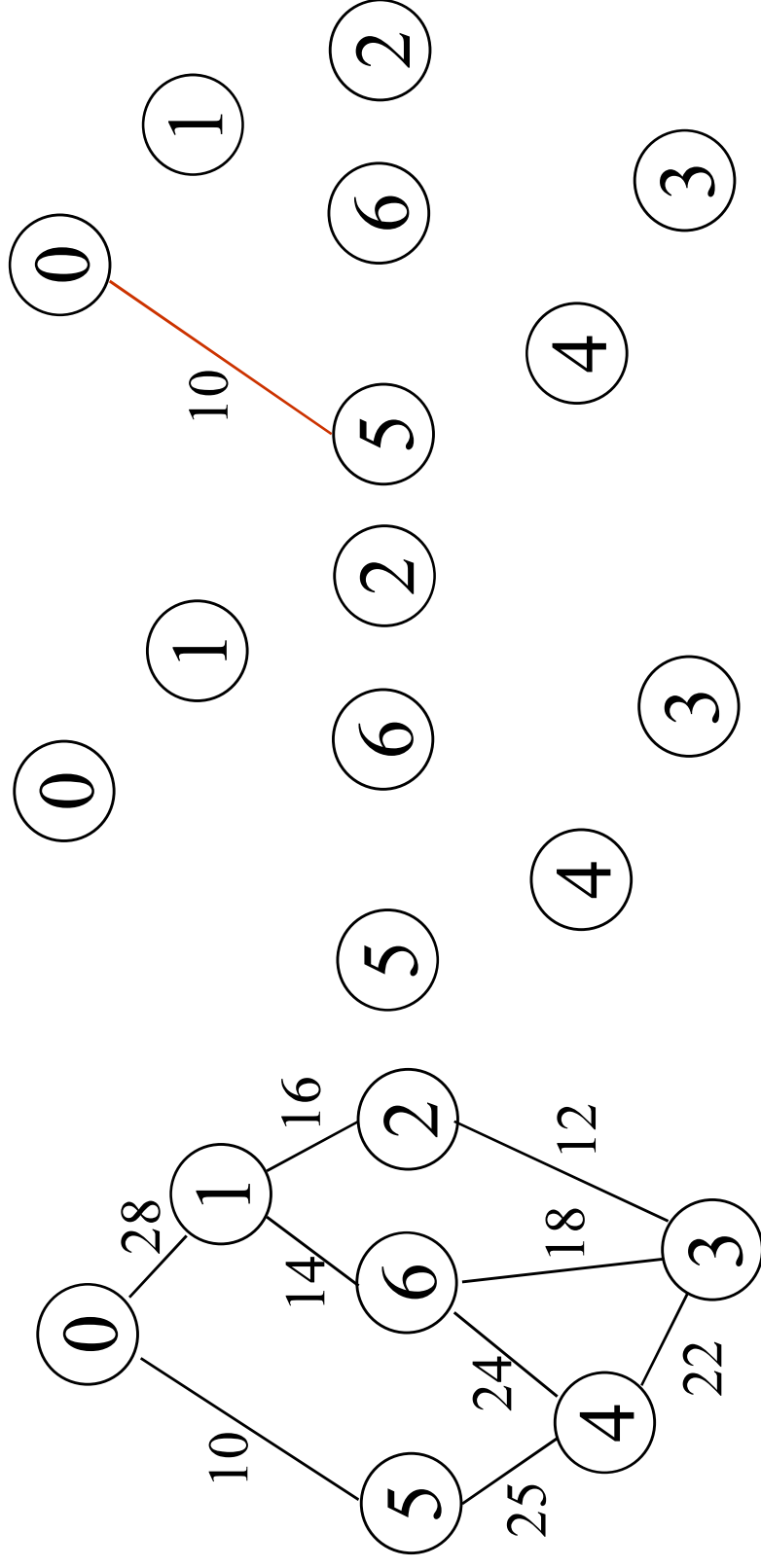
3 ~~18~~ 6

3 ~~22~~ 4

4 ~~24~~ 6

4 ~~25~~ 5

0 ~~28~~ 1



0-~~10~~-5

2-~~12~~-3

1-~~14~~-6

1-~~16~~-2

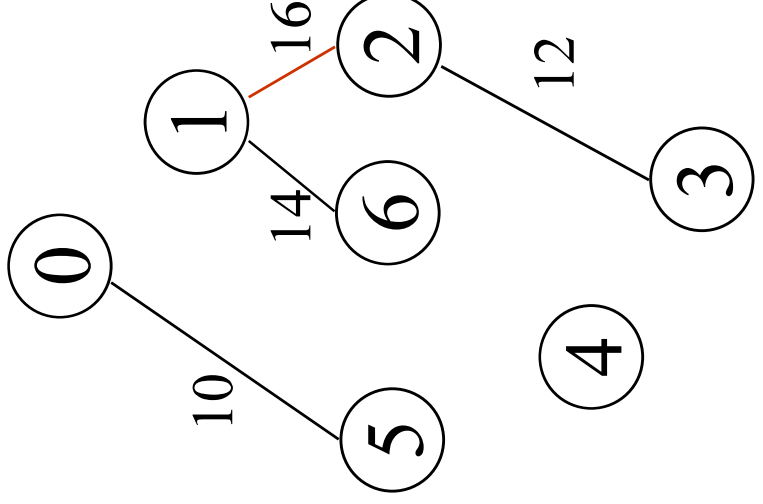
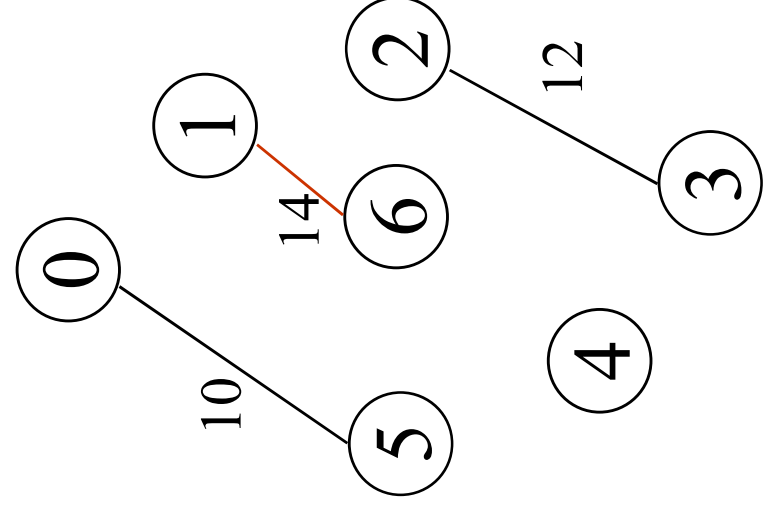
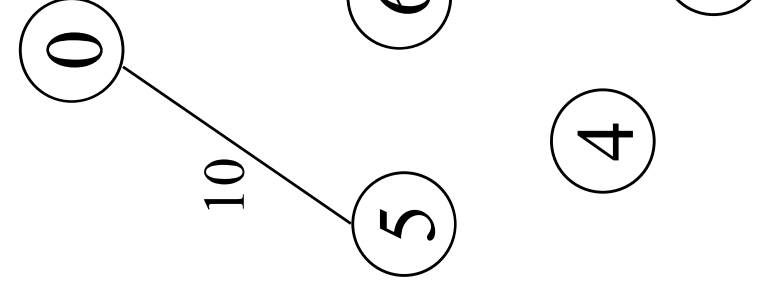
3-~~18~~-6

3-~~22~~-4

4-~~24~~-6

4-~~25~~-5

0-~~28~~-1



↑  
+ 3-~~5~~-6  
cycle



0-~~10~~-5

2-~~12~~-3

1-~~14~~-6

1-~~16~~-2

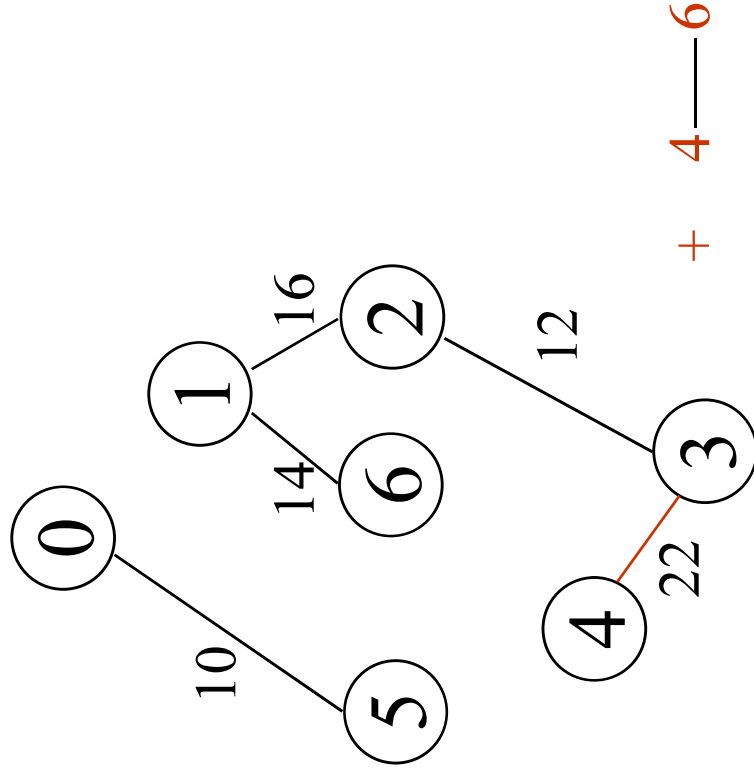
3-~~18~~-6

3-~~22~~-4

4-~~24~~-6

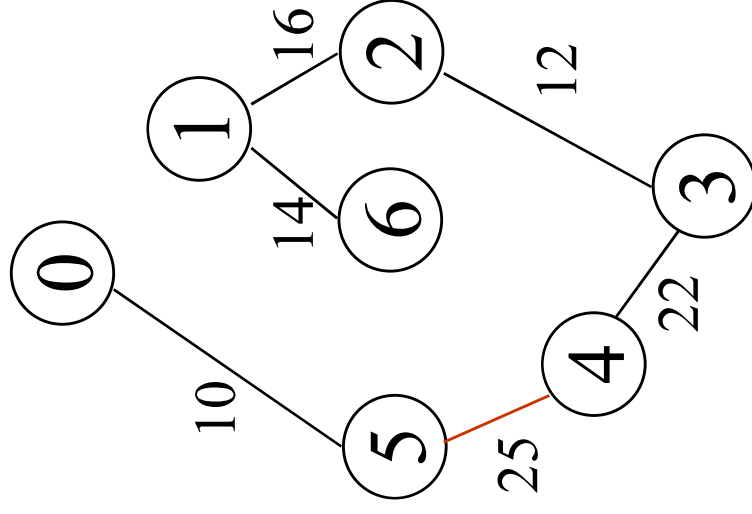
4-~~25~~-5

0-~~28~~-1



+ 4-6

cycle



cost = 10 + 25 + 22 + 12 + 16 + 14

# Kruskal Algorithm for MST

```
MST-KRUSKAL( $G, w$ )
1   $A \leftarrow \emptyset$ 
2  for each vertex  $v \in V[G]$ 
3      do MAKE-SET( $v$ )
4  sort the edges of  $E$  into nondecreasing order by weight  $w$ 
5  for each edge  $(u, v) \in E$ , taken in nondecreasing order by weight
6      do if FIND-SET( $u$ )  $\neq$  FIND-SET( $v$ )
7          then  $A \leftarrow A \cup \{(u, v)\}$ 
8              UNION( $u, v$ )
9  return  $A$ 
```

# Prim's Algorithm

```
MST-PRIM( $G, w, r$ )
1  for each  $u \in V[G]$ 
2      do  $\text{key}[u] \leftarrow \infty$ 
3       $\pi[u] \leftarrow \text{NIL}$ 
4   $\text{key}[r] \leftarrow 0$ 
5   $Q \leftarrow V[G]$ 
6  while  $Q \neq \emptyset$ 
7      do  $u \leftarrow \text{EXTRACT-MIN}(Q)$ 
8      for each  $v \in \text{Adj}[u]$ 
9          do if  $v \in Q$  and  $w(u, v) < \text{key}[v]$ 
10             then  $\pi[v] \leftarrow u$ 
11                  $\text{key}[v] \leftarrow w(u, v)$ 
```

# Examples for Prim's Algorithm

