## **Insertion Sort**

Insertion Sort builds the sorted array **one element at a time**, just like sorting playing cards in your hand. You take one card (element) from the unsorted pile and **insert** it into its correct position among the already sorted ones.

Basically, we "pick" and element and assume that now its place is empty. We then check the array for a place where the element fits, and insert it there, shifting the elements, hence the name.

Algorithmic Steps (in pseudocode / algorithm format)

Algorithm: InsertionSort(A, n)

Input: Array A of n elements

Output: Sorted array A in ascending order

1. for 
$$i \leftarrow 1$$
 to n-1 do

2. 
$$key \leftarrow A[i]$$

3. 
$$j \leftarrow i - 1$$

4. while 
$$j \ge 0$$
 and  $A[j] > \text{key do}$ 

5. 
$$A[j+1] \leftarrow A[j]$$

6. 
$$j \leftarrow j - 1$$

8. 
$$A[j + 1] \leftarrow key$$

9. end for

10. return A

## Implementation in C++

```
#include <iostream>
using namespace std;
void insertionSort(int arr[], int n) {
    for (int i = 1; i < n; i++) {
        int key = arr[i];  // current element to insert
        int j = i - 1;
        // Move elements of arr[0..i-1] that are greater than key
        // to one position ahead of their current position
        while (j \ge 0 \&\& arr[j] > key) {
            arr[j + 1] = arr[j];
            j = j - 1;
        }
        arr[j + 1] = key;
    }
}
int main() {
    int arr[] = {12, 11, 13, 5, 6};
    int n = sizeof(arr) / sizeof(arr[0]);
    insertionSort(arr, n);
```

## **Time Complexity**

Case	Comparisons	Shifts	Time
Best	n-1	0	O(n)
Average	~n²/4	~n²/4	O(n²)
Worst	n²/2	n-1	O(n²)

Space complexity: O(1); in-place sorting

Stable? Ves, equal elements stay in order

Adaptive? 🗹 Yes, fewer shifts if array is nearly sorted