

Name :- Aashi Ashokbhai Goyani

ID :- 1002205247

HANDS ON :- 12

## Ques 2 Defining the dynamic table's behaviour

When inserting an element into the table, if the table is full, the table doubles in size.

Doubling the size of the table takes  $O(n)$  time, where  $n$  is the current size of the table.

After resizing the new element is inserted into the table.

### i) Aggregate Method:

In this method, we calculate the total time for  $n$  insertions and divide it by  $n$  to find average time per insertion.

Let  $n_i$  be the size of the table when inserting  $i$ th element.

Initially  $n_1 = 1$

when resizing  $n_i = 2 \times n_{i-1}$

The time taken to insert an element into a table of size  $n$  is proportional to  $n$ .

$T_a \rightarrow$  amortized time per insertion.

$$T_a = \frac{T}{n} \rightarrow \text{total time for } n \text{ insertion.}$$

The first insertion takes  $O(1)$  time (without resizing).

The second insertion takes  $O(1)$  time (resizing needed  
doubling the size)

The third insertion takes  $O(2)$  time (resizing,  
doubling the size)

The  $k^{th}$  insertion takes  $O(2^{k-1})$  time (resizing,  
doubling size  $k-1$  times)

The  $n^{th}$  insertion takes  $O(2^{k-1})$  times

$$T = O(1) + O(1) + O(2) + O(4) + \dots + O(2^{k-1})$$

$$T = O(2^0) + O(2^1) + O(2^2) + O(2^3) + \dots + O(2^{k-1})$$

This is geometric series with common ratio  $= 2$

Sum of geometric series :-

$$T = O(2^{k-1}) = O(2^{\log_2 n} - 1) = O(n-1) = O(n)$$

Thus,  $T_a = O(1)$

## ii) Accounting method.

In this method, we assign each insertion a "charge" that covers the cost of insertion and resizing.

$O(1)$  charge  $\rightarrow$  for each insertion operation

$O(n)$  charge  $\rightarrow$  for each resizing operation.

$C_i$  be the total cost of inserting the  $i^{\text{th}}$  element.

$$A_i = C_i + C_{i-1}$$

Where  $A_i$  is the amortized cost for insertion.

$$C_i = O(1) \quad \text{for } i^{\text{th}} \text{ insertion}$$

$$C_{i-1} = O(n) \quad \text{for } i^{\text{th}} \text{ resizing operation}$$

The total charge  $T$  for  $n$  insertion can be calculated as,

$$T = n \times O(1) + \underbrace{n \times O(n)}_2$$

$\underbrace{\qquad\qquad\qquad}_{\substack{\text{Cost of} \\ n \text{ insertions}}}$        $\underbrace{\qquad\qquad\qquad}_{\substack{\text{Cost of resizing} \\ \text{operations.}}}$

$$T = O(n) + O\left(\frac{n^2}{2}\right) = O(n^2)$$

The amortized cost per operation is then;

$$\boxed{T_a = \frac{T}{n} = O(n)}$$