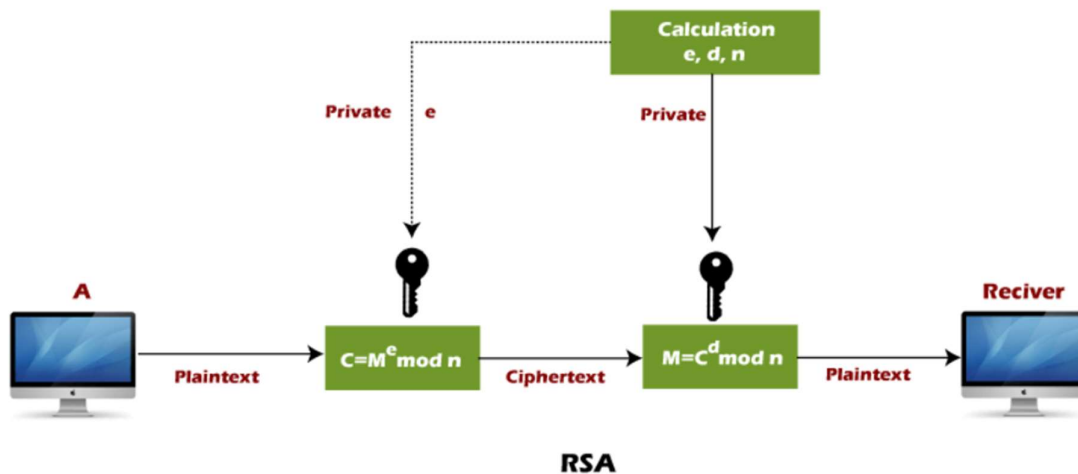


**Final Year B. Tech, Sem VII 2022-23**  
**PRN – 2020BTECS00211**  
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**Batch: B4**  
**Practical No – 11**

**Title:** Implementation of RSA Algorithm.

**Theory:**

RSA algorithm is a public key encryption technique and is considered as the most secure way of encryption. It was invented by Rivest, Shamir and Adleman in year 1978 and hence name RSA algorithm.



**RSA algorithm uses the following procedure to generate public and private keys:**

- Select two large prime numbers,  $p$  and  $q$ .
- Multiply these numbers to find  $n = p \times q$ , where  $n$  is called the modulus for encryption and decryption.
- Choose a number  $e$  less than  $n$ , such that  $n$  is relatively prime to  $(p - 1) \times (q - 1)$ . It means that  $e$  and  $(p - 1) \times (q - 1)$  have no common factor except 1. Choose "e" such that  $1 < e < \phi(n)$ ,  $e$  is prime to  $\phi(n)$ ,  
 **$\gcd(e, \phi(n)) = 1$**
- If  $n = p \times q$ , then the public key is  $\langle e, n \rangle$ . A plaintext message  $m$  is encrypted using public key  $\langle e, n \rangle$ . To find ciphertext from the plain text following formula is used to get ciphertext  $C$ .  
 **$C = m^e \bmod n$**   
Here,  $m$  must be less than  $n$ . A larger message ( $> n$ ) is treated as a concatenation of messages, each of which is encrypted separately.
- To determine the private key, we use the following formula to calculate the  $d$  such that:  
 **$D_e \bmod \{(p - 1) \times (q - 1)\} = 1$**   
**Or**  
 **$D_e \bmod \phi(n) = 1$**
- The private key is  $\langle d, n \rangle$ . A ciphertext message  $c$  is decrypted using private key  $\langle d, n \rangle$ . To calculate plain text  $m$  from the ciphertext  $c$  following formula is used to get plain text  $m$ .  
 **$m = c^d \bmod n$**

**Example:**

## RSA Example

1. Select primes:  $p=17$  &  $q=11$
2. Compute  $n = pq = 17 \times 11 = 187$
3. Compute  $\phi(n) = (p-1)(q-1) = 16 \times 10 = 160$
4. Select  $e$ :  $\gcd(e, 160) = 1$ ; choose  $e=7$
5. Determine  $d$ :  $de=1 \bmod 160$  and  $d < 160$   
Value is  $d=23$  since  $23 \times 7 = 161 = 10 \times 160 + 1$
6. Publish public key  $KU = \{7, 187\}$
7. Keep secret private key  $KR = \{23, 17, 11\}$

### Code Snapshots:

```
#include <bits/stdc++.h>
using namespace std;

// void file()
// {
// #ifndef ONLINE_JUDGE
//     freopen("input.txt", "r", stdin);
//     freopen("output.txt", "w", stdout);
// #endif
// }

// Function for extended Euclidean Algorithm
int ansS, ansT;
int findGcdExtended(int r1, int r2, int s1, int s2, int t1, int t2)
{
    // Base Case
    if (r2 == 0)
    {
        ansS = s1;
        ansT = t1;
        return r1;
    }

    int q = r1 / r2;
    int r = r1 % r2;

    int s = s1 - q * s2;
    int t = t1 - q * t2;

    cout << q << " " << r1 << " " << r2 << " " << r << " " << s1 << " " << s2
    << " " << s << " " << t1 << " " << t2 << " " << t << endl;

    return findGcdExtended(r2, r, s2, s, t2, t);
}

int modInverse(int A, int M)
{
    int x, y;
    int g = findGcdExtended(A, M, 1, 0, 0, 1);
    if (g != 1) {
        cout << "Inverse doesn't exist";
        return 0;
    }
    else {
        // m is added to handle negative x
    }
}
```

```

        int res = (ansS % M + M) % M;
        cout << "inverse is" << res << endl;
        return res;
    }
}

long long powM(long long a, long long b, long long n)
{
    if (b == 1)
        return a % n;
    long long x = powM(a, b / 2, n);
    x = (x * x) % n;
    if (b % 2)
        x = (x * a) % n;
    return x;
}

int findGCD(int num1, int num2)
{
    if (num1 == 0)
        return num2;
    return findGCD(num2 % num1, num1);
}

// Code to demonstrate RSA algorithm
int main()
{
    //file();

    // Two random prime numbers
    long long p, q, e, msg;
    //17 31 7 2

    cout << "Please enter 2 prime number and e and Message to Encrypt" <<
endl;
    cin >> p >> q >> e >> msg;

    cout << "2 random prime numbers selected are " << p << " " << q << endl;

    // First part of public key:
    long long n = p * q;
    cout << "Product of two prime number n is " << n << endl;

    // Finding other part of public key.
    // e stands for encrypt

    cout << "Taken e is " << e << endl;

```

```

long long phi = (p - 1) * (q - 1);
cout << "phi is " << phi << endl;

while (e < phi) {
    // e must be co-prime to phi and
    // smaller than phi.
    if (findGCD(e, phi) == 1)
        break;
    else
        e++;
}

cout << "Final e value is " << e << endl;

// Private key (d stands for decrypt)

long long d = modInverse(e, phi);
cout << "d is " << d << endl;

cout << "\nso now our public key is " << "<" << e << "," << n << ">" <<
endl;
cout << "\nso now our private key is " << "<" << d << "," << n << ">" <<
endl << endl;

// Message to be encrypted

cout << "Message date is " << msg << endl;

// Encryption  $c = (msg^e) \% n$ 
long long c = powM(msg, e, n);
cout << "Encripted Message is " << c << endl;

// Decryption  $m = (c^d) \% n$ 
long long m = powM(c, d, n);
cout << "original Message is " << m << endl;

return 0;
}

```

## Output Snapshots:

```
PROBLEMS  OUTPUT  TERMINAL  GITLENS  DEBUG CONSOLE

PS C:\Users\Ashitra\OneDrive\Desktop\7th sem\Practicals\CNS\Programs> cd "c:\Users\Ashitra\OneDrive\Desktop\7th sem\Practicals\CNS\Programs\" ; if ($?) { g++ RSAAlgo.cpp -o RSAAlgo } ; if ($?) { .\RSAAlgo }
}
Please enter 2 prime number and e and Message to Encrypt
5 7 3 28
2 random prime numbers selected are 5 7
Product of two prime number n is 35
Taken e is 3
phi is 24
Final e value is 5
0 5 24 5 1 0 1 0 1 0
4 24 5 4 0 1 -4 1 0 1
1 5 4 1 1 -4 5 0 1 -1
4 4 1 0 -4 5 -24 1 -1 5
inverse is5
d is 5

so now our public key is <5,35>

so now our private key is <5,35>

Message date is 28
Encrypted Message is 28
original Message is 28
PS C:\Users\Ashitra\OneDrive\Desktop\7th sem\Practicals\CNS\Programs> 
```

## Conclusion:

1. It is concluded that, while establishing RSA key pairs, usage keys and general-purpose keys are integrated.
2. In usage RSA keys, two key pairs are used for encryption and signatures.