

**Final Year B. Tech, Sem VII 2022-23**  
**PRN – 2020BTECS00211**  
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**Batch: B4**  
**Practical No – 11**

**Title:** Implementation of Diffie-Hellman Key Exchange.

**Theory:**

The Diffie-Hellman key exchange was one of the most important developments in public-key cryptography and it is still frequently implemented in a range of today's different security protocols.

It allows two parties who have not previously met to securely establish a key which they can use to secure their communications. In this article, we'll explain what it's used for, how it works on a step-by-step basis, its different variations, as well as the security considerations that need to be noted in order to implement it safely.

The Diffie-Hellman algorithm is being used to establish a shared secret that can be used for secret communications while exchanging data over a public network using the elliptic curve to generate points and get the secret key using the parameters.

- For the sake of simplicity and practical implementation of the algorithm, we will consider only 4 variables, one prime  $P$  and  $G$  (a primitive root of  $P$ ) and two private values  $a$  and  $b$ .
- $P$  and  $G$  are both publicly available numbers. Users (say Alice and Bob) pick private values  $a$  and  $b$  and they generate a key and exchange it publicly. The opposite person receives the key and that generates a secret key, after which they have the same secret key to encrypt.

## Step by Step Explanation

Alice

Bob

Public Keys available = P, G    Public Keys available = P, G

Private Key Selected = a

Private Key Selected = b

Key generated =

$$x = G^a \bmod P$$

Key generated =

$$y = G^b \bmod P$$

Exchange of generated keys takes place

Key received = y

key received = x

Generated Secret Key =

$$k_a = y^a \bmod P$$

Generated Secret Key =

$$k_b = x^b \bmod P$$

Algebraically, it can be shown that

$$k_a = k_b$$

Users now have a symmetric secret key to encrypt

### Example:

Step 1: Alice and Bob get public numbers  $P = 23$ ,  $G = 9$

Step 2: Alice selected a private key  $a = 4$  and  
Bob selected a private key  $b = 3$

Step 3: Alice and Bob compute public values

Alice:  $x = (9^4 \bmod 23) = (6561 \bmod 23) = 6$

Bob:  $y = (9^3 \bmod 23) = (729 \bmod 23) = 16$

Step 4: Alice and Bob exchange public numbers

Step 5: Alice receives public key  $y = 16$  and  
Bob receives public key  $x = 6$

Step 6: Alice and Bob compute symmetric keys

Alice:  $k_a = y^a \bmod p = 65536 \bmod 23 = 9$

Bob:  $k_b = x^b \bmod p = 216 \bmod 23 = 9$

Step 7: 9 is the shared secret.

### Code Snapshots:

```
#include <bits/stdc++.h>
#define ll long long
#define ull unsigned long long
#define pb emplace_back
#define po pop_back
#define vi vector<ll>
#define vii vector<vector<ll>>
using namespace std;

vector<int> primeNums;
vector<bool> prime(10000001,1);

void SieveOfEratosthenes(int n){
    for(int p=2; p*p<=n; p++){
        if(prime[p] == true){
            for (int i = p * p; i <= n; i += p)
                prime[i] = false;
        }
    }
}
```

```

        for(int i=3;i<n;i+=2){
            if(prime[i]) primeNums.push_back(i);
        }
    }

ll power(ll a, ll b, ll p){
    if (b == 1)
        return a;
    else
        return (((long long int)pow(a, b)) % p);
}

void findPrimefactors(unordered_set<int> &s, int n){
    while (n%2 == 0){
        s.insert(2);
        n = n/2;
    }
    for (int i = 3; i <= sqrt(n); i = i+2){
        while (n%i == 0){
            s.insert(i);
            n = n/i;
        }
    }
    if (n > 2)
        s.insert(n);
}

int primitiveRoot(int n){
    unordered_set<int> s;
    int phi = n-1;
    findPrimefactors(s, phi);
    for (int r=2; r<=phi; r++){
        bool flag = false;
        for (auto it = s.begin(); it != s.end(); it++){
            if (power((ll)r, (ll)phi/(*it),(ll)n) == 1)
            {
                flag = true;
                break;
            }
        }
        if (flag == false)
            return r;
    }
    return -1;
}

```

```

int main(){
    // prime number till 100000000
    SeiveOfEratosthenes(100000000);
    int privateNumberA, privateNumberB;
    cout<<"Enter the privateNumber of A and B respectively : ";
    cin>>privateNumberA>>privateNumberB;
    cout<<"\nFinding prime Number and a primitive root ...\n";

    srand(time(0));
    int p = primeNums[rand() % primeNums.size()];
    int g = primitiveRoot(p);

    cout<<"\tPrime Number : "<<p<<"\n";
    cout<<"\tPrimitive Root : "<<g<<"\n";

    // calculating the private key for a
    ll x = power(g,privateNumberA,p);
    if(x<0) x = p + x;
    cout<<"\nThe private key a for A is : "<<x<<"\n";

    // calculate private key for b
    ll y = power(g,privateNumberB,p);
    if(y<0) y = p + y;
    cout<<"The private key b for B is : "<<y<<"\n";

    ll ka = power(y, privateNumberA, p); // Secret key for A
    if(ka<0) ka = p + ka;
    ll kb = power(x, privateNumberB, p); // Secret key for B
    if(kb<0) kb = p + kb;

    cout<<"\n\nSecret key for the A is : "<<ka;
    cout<<"\nSecret key for the B is : "<<kb<<endl;
    return 0;
}

```

### Output Snapshots:

```
PROBLEMS  OUTPUT  TERMINAL  GITLENS  DEBUG CONSOLE

PS C:\Users\Ashitra\OneDrive\Desktop\7th sem\Practicals\CNS\Programs> cd "c:\Users\Ashitra\OneDrive\Desktop\7th sem\Practicals\CNS\Programs\" ; if ($?) { g++ DiffieHelman.cpp -o DiffieHelman } ; if ($?) { .\DiffieHelman }
    Prime Number : 12589
    Primitive Root :2

The private key a for A is : 4334
The private key b for B is : 512

Secret key for the A is :10937
Secret key for the B is : 10937
PS C:\Users\Ashitra\OneDrive\Desktop\7th sem\Practicals\CNS\Programs> 
```

### Conclusion:

1. The Diffie Hellman key Exchange has proved to be a useful key exchange system due to its advantages.
2. While it is really tough for someone snooping the network to decrypt the data and get the keys, it is still possible if the numbers generated are not entirely random.