



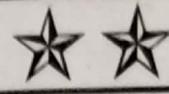
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Practical No.1

. Aim : Implement linear search to find an item in the list.

. Theory :

Linear Search

Linear search is one of the simplest algorithm in which targeted item is sequentially matched with each item in the list.

It is worst searching algorithm with worst case time complexity. It is a naive approach. On the other hand in case of an ordered list, instead of searching the list in sequence, a binary search is used which will start by examining the middle term.

✓ Linear search is a technique to compare each and every element with the key element to be found, if both of them matches, the algorithm returns that element found and its position. It is also found.

1) unsorted
Algorithm :

- Step 1 : Create an empty list and assign it to a variable

- Step 2 : Accept the total no. of elements to be ingored into the list from the user . say 'n'

- Step 3 : Use for loop for adding the elements onto the list

Step 4 : Print the list

Steps : Accept an element from the user that is to be searched into the list.

Steps : Use for loop in a range from '0' to the total no. of elements to search the elements from the list.

Step 7 Use for loop that the elements in the list is equal to the accepted from user

Step 8 : If the element is found then print the statement that the element is found along with the elements present in.

Step 9 : Use the condition at the loop to print the statement that the element is not found if the element which is accepted from user is not there in the list

```
X = [4, 3, 5, 8, 9, 2]
a = int(input("Enter a number : "))
for i in range(len(x)):
    if a == x[i]:
        print("Element found at : ", i)
        break
    else:
        print("Element not found")
```

Sorted Program

```
s = list(input("Enter the numbers : "))  
s.sort()  
print(s)  
a = int(input("Enter the Number to be searched : "))  
for i in range(len(s)):  
    if (a == s[i]):  
        print("Number found in position", i)  
        break  
    if (a != s[i]):  
        print("Number not found")
```

Output

```
Enter the Numbers : 2, 5, 8, 9, 7  
[1, 2, 5, 8, 9]  
Enter the Numbers to be searched : 5  
Number found in position 2
```

```
Enter the Numbers : 2, 3, 9, 5, 1  
[1, 2, 3, 5, 9]
```

```
Enter the Numbers to be searched : 7  
Number not found.
```

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Step 1 : Draw the output of given algorithm.

Step 2 : Sorted Arrays Search :

Sorting means to arrange the elements in increasing or decreasing order.

Algorithm :

Step 1 : Create Empty list and assign it to variable.

Step 2 : Accept total no. of elements to be inserted into the list from user, say 'n'

Step 3 : Use for loop for using append() method to add element in the list.

Step 4 : Use sort() method to sort the accepted element and assign in increasing order the list from point the list.

Step 5 : Use If statement to give the range in which element is not found in given range then display "Element not found".

Step 6 : To use else statement if element is not found in range then satisfy the given condition.

Step 7 : Use for loop for range from 0 to the top no. of elements to be searched before doing this accept or search no. from user using Input Statement

Step 8 : Use for loop that the elements in the list equal to the element accepted from user

Step 9 : If the element is found then print the statement that the element is found along with the element position.

Step 10 : Use another if loop to print that the element is not found if the element which is accepted from user is not there in the list.

Step 11 : ~~Print the input and output of above algorithm.~~

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29/11/19

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Source code

```

a = 108 + C.append("Element " + str(i) + ": ")
a = sorted(a)
C = len(a)
S = input("Enter search no.: ")
if C > 0:
    if S < a[0] or S > a[-1]:
        print("Not found")
    else:
        first, last = 0, C - 1
        for i in range(C):
            mid = (first + last) // 2
            print(f"found at {mid}")
            if S == a[mid]:
                print(f"Number found")
                break
        else:
            if S < a[mid]:
                last = mid - 1
            else:
                first = mid + 1

```

✓
mr

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Practical - 2 NO. 2

Aim : Implement Binary Search to find an searched no. in the list

Theory :

Binary Search

Binary Search is also known as half-interval search, logarithmic search or binary chop is a search algorithm, search or binary chop is a search algorithm that finds a position of a target value within a sorted array. If you are looking for the number which is at the end of the list then you need to search entire list in linear search. Which is time consuming this can be avoided by using Binary Search.

Algorithm :

Step 1 : Create empty list and assign it to a variable

Step 2 : Using input method accept the range of given list

Step 3 : use for loop, add elements in list, using append() method,

Step 4 : Use `sort()` to sort the accepted elements and assign it in increasing order of 1st + 2nd + ... + last after sorting.

Step 5 : Use if loop to give the range in which element is found in given range then display message "Element + NOT found".

Step 6 : Then use the else statement, if statement is found in range then satisfy the below condition

Step 7 : Accept an argument & key of the element that element has to be sorted.

Step 8 : Initialize first to 0 and last to last element of list & array is starting from 0 & it is initialized 1 less than the serial count.

Step 9 : use for loop & assign to the given loop.

Step 10 : If statement in 1st + and still the element to be searched is not found then find the middle element (cm)

Step 11 : Else if the from to be searched is less than middle term then Initialize (last - cm) = mid cm - 1

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else

Initialize first (el) = mid (cm) - 1

Step 12 : Repeat till you found the element & print the input
2. Output of above algorithm.

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```
# Bubble sort
a = list(input("Enter Numbers : "))
for p in range(len(a)-1):
    for c in range(len(a)-1-p):
        if a[c] > a[c+1]:
            t = a[c]
            a[c] = a[c+1]
            a[c+1] = t
print(a)
```

Output :

Enter Numbers : 2, 4, 8, 1, 9, 5, 6
[1, 2, 4, 5, 6, 8, 9]

Enter Numbers : 0.43, 0.31, 0.78, 0.19
[0.19, 0.32, 0.43, 0.78]

Enter Numbers : 232, 543, 123, 901
[123, 232, 543, 901]

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Practical No. 3

Bubble Sort

- Aim : Implementation of bubble sort program on given list
- Theory : Bubble sort is based on the idea of repeatedly comparing pairs of adjacent elements and then swapping their position if they are in wrong order. This is the simplest form of sorting available. In this, we sort the given elements in ascending and descending order by comparing two adjacent elements at a time.

Algorithm

Step 1 : Bubble sort starts by comparing the first two elements of an array and swapping it if necessary.

Step 2 : If we want to sort the elements of array in ascending order then first element is greater than second then if we need to swap elements.

Step 3 : If the element is smaller than second then we do not swap the elements.

Step 4 : Again second and third elements are compared and swapped if it is necessary and this process goes on till last and second last element is compared and swapped.

Rein + C" $\xrightarrow{\text{Oxid.}}$ $\text{GO}_2 + \text{H}_2\text{O}$

help Ca

```
def helpCalcs + stimsets
```

help(Calib +>SPRINT10108 +)

23 Sept Caligariopsis +, larva +, -

`pivot = 0 | first [first]`

$$L = \rho_0 r s + \gamma$$

$$g = \log E$$

while not done:

while $i < p$ and $\text{ans} + [4]^{2-p} \neq 0$ do $i = i + 1$

$$1 + \pi = \pi$$

where $\alpha_{pq} + [r] = p^{\circ} \wedge q$ and $r = 1$.

$$T - \lambda = \delta$$

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卷之三

$$t = \alpha^{1g} + \zeta^1$$

$$t' = \alpha^{10} S + \sum F_i P_i S + \sum Q_i P_i S + \sum R_i P_i S + \sum T_i P_i S = \alpha^{10} S + \sum S^2 T_i$$

$\alpha = [r \text{ (temp)} + C \text{ (envelope)}] \text{ range} / 100 \text{ first } - 11$

$\text{change}(\omega, x) = \text{entropy}(x)$

Ones + \Rightarrow inner element + 0])
n = \Rightarrow append (b) + 0])
quick call(s +)
Done

Doktrin SOR

Q4: Implement quick sort to sort the given list.

Theory : the ocrk sort is a recursive algorithm based on the divide and conquer the technique

Step 1: Quick sort first selects a value which is called pivot value. First element serve as pivot value since we know that largest element ends up as last in that list.

Step 2 : The partition process will happen next . It will find one split point and at the same time move other elements to the appropriate side of the list .
either less than or greater than pivot value .

Step 3. Partitioning begins by locating two position markers. Let's call them leftmargin & rightmark at the beginning of the list and a remaining item in the list. The goal of partition process is to move items that are on comparing side with respect to pivot value. Also converging on the split point

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Step 5: We begin partitioning leftmost unit. A value that is greater than the pivot value document rightmost until we find value less than the pivot value. At this point we have 4000 items that are not yet partitioned. Current split point.

Step 5: In the pivot value which becomes less than current we update position of rightmark to now split. New split point is now in place.

Step 6: The pivot value can be exchanged with the current to split point. Pivot value (is now in place)

Step 7: In addition, all the items to left of split position less than pivot will the items to the right of split. If pivot is greater than pivot, the can now be divided at point & hence sort can be invoked recursively on 2 halves.

Step 8: The quicksort function invokes a recursive function,

QuickSortRecap

Step 9: quicksort pop begins with some base case as follows

Step 10: If length of list is less than 0, Quicksort already sorted

Step 11: If it is greater, then it can be partitioned and split

Step 12: The partition function implements the process described till

Step 13: Display and stuck the resulting and output of above algorithm.

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Step 7: Once element satisfies the condition
that is lesser than stack's greater than
than it will be pushed into the stack.

Step 8: After pushing elements as lesser than the element
it will check and initialize the value.

Step 9: Compared with to check the elements but pop method
used to delete the element from the stack.

Step 10: If the pop method value is less than, then re-
move the stack is empty or else delete the
smallest item stack at topmost position.

Step 11: This condition checks whether the
bottom of elements are zero while the
second case another top is assigned any
value is acknowledged any value then we can
say the stack is empty.

Step 12: Assign the elements values in push method
to add and then the given value is
poped off.

Step 13: Finally the final output of above
algo is given.

Output:
push(10)
push(20)

>>> S.push (20)
>>> S.pop
 $\Sigma 20, 0, 0, 0, 0$

>>> S.pop C
Data = 20

>>> S.pop
 $\Sigma 0, 0, 0, 0, 0$
>>> S.push (10)
>>> S.push (20)
>>> S.push (30)
>>> S.push (40)
>>> S.push (50)
>>> S.pop

S.pop()

$\Sigma 10, 20, 30, 40, 50$

Mr
24/01/2020

47

Program

```

class Queue:
    global r
    global f
    global a
    def __init__(self):
        self.r = f = 0
        self.a = [0, 0, 0, 0]
    def enqueue(self, value):
        self.r += 1
        if self.r == len(self.a):
            print("Queue is full")
        else:
            self.a[self.r] = value
            self.r += 1
    def dequeue(self):
        if self.r - f == len(self.a):
            print("Queue is empty")
        else:
            value = self.a[f]
            self.a[f] = None
            print("Queue Element Deleted", value)
            self.r -= 1
b = Queue()
b.enqueue(1)
b.enqueue(2)
b.enqueue(3)
b.dequeue()
b.dequeue()
b.dequeue()
b.dequeue()

```

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PRACTICAL NO. 5

• Aim: Implementing a queue using python list

• theory: Queue is a linear data structure which has 2 reference front and rear.

Implementing a queue using python list is the simplest as the python list provides inbuilt functions to perform the specified operations of the queue. It is based on the principle that a new element is inserted after rear and element of queue is deleted which is at front. In simple terms, a queue can be described as a data structure based on first in first out FIFO principle.

- queue(): Creates a new empty queue
- enqueue(): Insert an element at the rear of the queue and similar to that of insertion of linked using tail.
- dequeue(): Returns the element which was at the front. The front is moved to the successive element. A dequeue operation cannot remove element if the queue is empty.

Output:

```

>>> b.enqueue(4)
queue element inserted 4
>>> b.enqueue(5)
queue element inserted 5
>>> b.enqueue(6)
queue element inserted 6
>>> b.enqueue(7)
queue element inserted 7
>>> b.enqueue(8)
queue element inserted 8
>>> b.enqueue(9)
queue is full
>>> b.a
[4,5,6,7,8]
>>> b.dequeue()
queue element deleted 4
>>> b.dequeue()
queue element deleted 5
101110
>>> b.dequeue()
queue element deleted 6
>>> b.dequeue()
queue element deleted 7
>>> b.dequeue()
queue element deleted 8
>>> b.dequeue()
queue is empty
>>> b.a
[0,0,0,0,0,0]

```

- Algorithm:
- Step 1: Define a class queue and assign global variables then define `push()` method with self argument in `init()` - assign or initialize the initial value with the help of self argument.
- Step 2: Define an empty list and define an `enqueue()` method with 2 arguments, assign the length of empty list.
- Step 3: Use if statement that length is equal to less than queue is full or else insert the element in empty list or display that queue element added successfully and increment by 1.

- Step 4: Define `dequeue()` with self argument. If this is true, use if statement that front is equal to length of list then display queue is empty or else give that front is at zero and using while delete the present from result and increment front by 1.

- Step 5: Now call the `queue()` function and give the elements that has to be added in the empty list using `enqueue()` and print the list after adding queue is empty and same for deleting and displaying the list after deleting the element from the list.

source code:

```
def evaluate(s):
    k = s.split(' ')
    n = len(k)
    stack = []
    for i in range(n):
        if k[i].is digit():
            stack.append(int(k[i]))
        elif k[i] == '+':
            a = stack.pop()
            b = stack.pop()
            stack.append(a + b)
        elif k[i] == '-':
            a = stack.pop()
            b = stack.pop()
            stack.append(b - a)
        elif k[i] == '*':
            a = stack.pop()
            b = stack.pop()
            stack.append(b * a)
    return stack.pop()

s = "8 6 9 * +"
n = evaluate(s)
print("the value is:", n)
```

Practical No. 4

49

Aim: Program of Evaluation of given string by using stack in python environment i.e postfix.

Theory:

The postfix expression is free of any parenthesis. Further we took care of the priorities of the operation in the program. A given postfix expression can easily be evaluated using stack. Reading the expression is always from left-right in postfix.

Algorithm:

Step 1: Define evaluate as function, then create an empty stack in Python.

Step 2: Convert the string to a list by using the string method 'split'

Step 3: Calculate the length of string & print it

Step 4: Use for loop to assign the range 0 to string then give condition using if statement

Step 5: Scan the taken list from left to right. If taken in as operand convert it from a string to an integer & push the value onto the 'P'.

Ex

Step 6 : If the token is an operator & it will need two operands, pop the 1st token, the first pop is second operand, and the second pop is the right operand.

Step 7 : Perform the arithmetic operation. Push the result back on the stack.

Step 8 : When the input expression has been completed processed, the result is on the stack. Pop the 'P' & return the value.

Step 9 : Print the result or string after the evaluation of postfix.

Step 10 : Attach output exp of above algorithm.

Output:

>> evaluated value is : 462

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Source code

Printed 02/05/

```

class node:
    global data
    global next
    def __init__(self, item):
        self.data = item
        self.next = None

class linkedlist:
    global s
    def __init__(self):
        self.s = None
    def add2(self, item):
        newnode = node(item)
        if self.s == None:
            self.s = newnode
        else:
            head = self.s
            while head.next != None:
                head = head.next
            head.next = newnode
    def add3(self, item):
        newnode = node(item)
        if self.s == None:
            self.s = newnode
        else:
            newnode.next = self.s
            self.s = newnode
    def display(self):
        head = self.s
        while head.next != None:
            print(head.data)
            head = head.next
        print(head.data)

```

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Practical No. 2

- Purpose: Implementation of single linked list by adding the nodes from last position.
- Theory: A linked list is a linear data structure which stores the elements in a node in a linear fashion but no necessarily contiguous. The individual element of the linked list called a Node. Node comprises of 2 parts: (1) Data (2) Next: Data stores all the information like + the element for example full number name, address etc whereas next refers to the next node. In case of large list if we add/remove any element in the list all the elements of list has to adjust itself every time we add, it's very tedious task so linked list is used to solving this type of problem.

- Algorithm:
- Step 1: Transversing of a linked list means visiting all the nodes of the linked list in order to perform some operation on them.
- Step 2: The entire list can be accessed using the first node of linked list, the first node of the linked list in turn is referred by the head pointer of linked list.

Step 3: Thus the entire linked list can be transversed using the node which is referred by the head pointer of the linked list.

Step 4: Now that we know that we can traverse the entire linked list using the head pointer we should only use it to refer the first Node as & only.

Step 5: We should not use the head pointer to transverse the original linked list because the head pointer is our only reference to the first node. If we modify the reference of the linked list so in order to changes we cannot have the head pointer as a valid reference.

Step 6: - We may lose the reference to the first node in the linked list and hence not be able to modify the linked list. So in order to avoid making some unwanted changes to the list node we will use a temporary node to transverse the entire linked list.

Step 7: - We will use this temporary node as a copy of node we are currently transversing. Since we are making temporary node copy of current node the datatype of temporary node should also be node.

Output:
 >>> q.addB(10)
 >>> q.addB(20)
 >>> q.addB(30)
 >>> q.addB(40)
 >>> q.addL(50)
 >>> q.addL(60)
 >>> q.addL(70)
 >>> q.addL(80)
 >>> q.print()

40
30
20
10
50
60
70
80

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24/01/2022

Step 8: Now that current p is pointing to the first node
if we want to access 2nd node or 1st we can
refer p as the next node of 1st node.

Step 9: But the 3rd node is referred by current.
So we can traverse to 2nd nodes as $\text{h}=\text{h}.next$

Step 10: Similarly we can traverse rest nodes in the
linked list using same method by while loop.

Step 10.1: Our concern now is to find terminating
condition for the while loop.

Step 12: The last node in the linked list is referred
by the tail of linked list. Since the last node of linked
list does not have any next node, the value p
in the next field of last node is None.

Step 13: So we can refer the last node of
linked list as $\text{self.s}=\text{None}$.

Step 14: We have to now see to start traversing
the linked list & how to identify
whether we have reached the last node
of linked list or not.

Step 15: Mention the coding of input and output
of above algorithm.

(0.0) 0.0
(0.1) 0.0
(0.2) 0.0
(0.3) 0.0
(0.4) 0.0
(0.5) 0.0
(0.6) 0.0
(0.7) 0.0
(0.8) 0.0
(0.9) 0.0
Capable

Binary Tree

10

Class node :

def __init__(self, value):

self.left = None

self.val = value

self.right = None

Class BST:

def __init__(self):

self.root = None

def add(self, value):

p = node(value)

if self.root == None:

self.root = p

print("Root is added successfully", p.val)

while True:

if h.left == None:

h.left = p

print("Node is added successfully", p.val)

else:

h = h.left

if h.right == None:

h.right = p

print("Node is added successfully", p.val)

else:

h = h.right

if h == None:

h.right = p

print("Node is added successfully", p.val)

else:

h = h.right

h = h.right

Theory: Binary tree is a tree which supports maximum of 2 children for any node within the tree. Thus any particular node can have 0, 1 or 2 children. There is another property of binary tree that is ordered such that one child is identified as left child and another as right child.

Traverses: (1) Transverse the left subtree. Then right subtree. (2) Visit root node.

(3) Transverse the right subtree and repeat it.

Postorder: 1) Visit the root node.

2) Transverse the left subtree. The left subtree in turn might have left and right subtrees.

3) Transverse the right subtree and repeat it.

Postorder: 1) Transverse the left subtree. The left subtree in turn might have left and right subtrees.

2) Transverse the right subtree.

3) Visit the root node.

break

successively, "val")

Algorithm

Step 1: Define class Node and define ~~insertion~~ insert() with 2 arguments. Initialize the value in this method.

Step 2: Again define another BST that is binary search tree with ~~inser~~ insert() with set argument and ~~add~~ add() root is None.

Step 3: Define add() method for adding the node. Defining variable p that p=nodeC value).

Step 4: Use if statement for checking the condition that if p is None then use else statement. For if node is less than the main node then put it in left otherwise right.

Step 5: Use while loop for checking the node p is less than or greater than the main node and breaking loop if p is not satisfying.

Step 6: Use if statement within if else statement for checking that node is greater than main root then for find min right side

Step 7: After this left subtree and right subtree repeat this method & arrange the node according to binary search tree.

```

def Inorder(C9004)
    if root == None:
        return
    Inorder(C9004.left)
    print(C9004.val)
    Inorder(C9004.right)

def Preorder(C9004):
    if root == None:
        return
    print(C9004.val)
    Preorder(C9004.left)
    Preorder(C9004.right)

def Postorder(C9004):
    if root == None:
        return
    Postorder(C9004.left)
    Postorder(C9004.right)
    print(C9004.val)

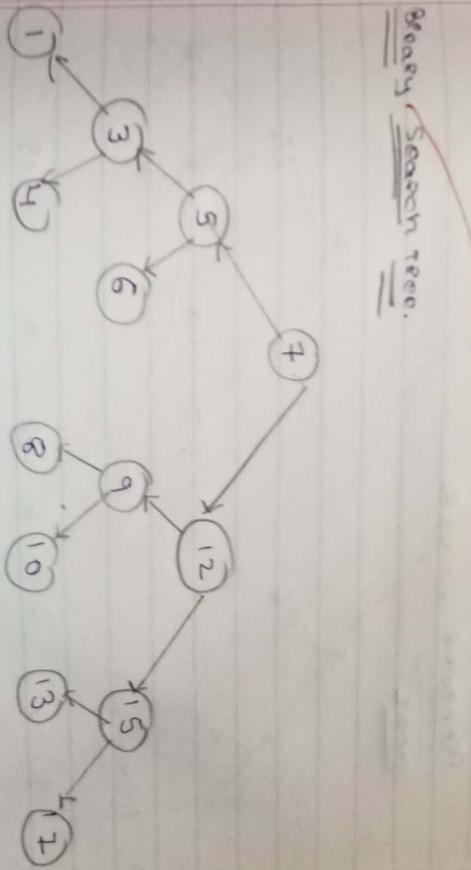
def Insertion(C9004, val):
    if C9004 == None:
        C9004 = Node(val)
    else:
        if val < C9004.data:
            C9004.left = Insertion(C9004.left, val)
        else:
            C9004.right = Insertion(C9004.right, val)
    return C9004

t = BST()
>>> t.add(12)
C9004 is added successfully: 12
>>> t.add(5)
C9004 is added successfully: 5
>>> t.add(15)
C9004 is added successfully: 15
>>> t.add(10)
C9004 is added successfully: 10
>>> t.add(20)
C9004 is added successfully: 20
>>> t.add(3)
C9004 is added successfully: 3

```

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```
>>> t.add(16)
Circular is added successfully '16'
>>> t.add(14)
Circular is added successfully '14'
>>> t.add(15)
Circular is added successfully '15'
Circular is added successfully '15')
>>> t.add(17)
Circular is added successfully '17')
>>> t.add(18)
Circular is added successfully '18')
>>> t.add(19)
Circular is added successfully '19')
Circular is added successfully '19')
>>> t.add(10)
Circular is added successfully '10')
>>> t.add(11)
Circular is added successfully '11')
Circular is added successfully '11')
>>> t.add(12)
Circular is added successfully '12')
Circular is added successfully '12')
>>> t.add(13)
Circular is added successfully '13')
Circular is added successfully '13')
>>> t.add(14)
Circular is added successfully '14')
Circular is added successfully '14')
```



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Step 8: Define Inorder(), Pre-order(), Post-order() function
Argument and use "if" statement that root must not
return true in all.

Step 9: In Inorder else statement based for giving true condition

For Preorder we have to give condition in else true if
left side and the right node

Step 10: For Preorder in else part assign left when right
and then go for root node.

Step 11: For Preorder in else part assign right when left
and then go for root node.

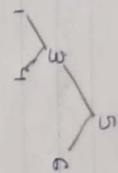
Step 12: Display the output of input of above algorithm.

Binary Search Tree

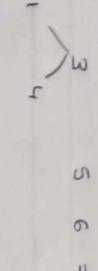
63

Inorder Clue)

Step 1:



Step 2:



Step 3 = 1 3 5 6 4 8 0 - 12 - 10 - 13 - 15 - 17

11

Preorder (Root +)

12

13

14

15

16

17

18

19

20

21

22

23

24

25

26

27

28

29

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31

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35

64

Inorder (Root +)

65

4.

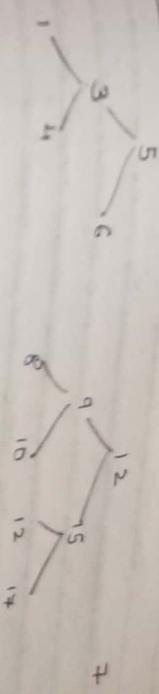
>>> Postorder (4.8002)

Step 3 : 1 0 8 - 4 6 10 2 8 10 15 13 14

6

Postorder : (204)

Step 3 :



Step 3 : 1 0 8 - 4 6 10 2 8 10 15 13 14
2 0 10 15 13 14 17 18 20 2

Step 3 :

1 0 8 - 4 6 10 2 8 10 15 13 14
2 0 10 15 13 14 17 18 20 2

Step 3 :

1 0 8 - 4 6 10 2 8 10 15 13 14
2 0 10 15 13 14 17 18 20 2

To demonstrate the use of a circular plane

168

global

$\text{SOL}_4 \cdot A = 0$
 $\text{SOL} \cdot A = 0$

tion Cadeau

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the `data` added: "data")

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— 5916 - 4) 5916 - 5916 - 5916 -

150 - 100 = 50% of water

Se 14 = 8 = 8
Defeat CII would go to

Grants-in-Aid
Research Grants
Fellowships

37

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2010-5-3

else :

s = self.tq + 5

self.tq = 0

prn((self.tq < self.tq - 2) ?

print("self.tq = " + str(self.tq - 2))

self.tq = self.tq + 1

else : print("queue is empty")

self.tq = s

t = self.tq - 1

Output

>>>q.add(100)

Data added : 100

>>>q.add(200)

Data added : 200

>>>q.add(300)

Data added : 300

>>>q.remove()

Data removed : 100

>>>q

[0, 200, 300, 0, 0]

67

computer controlled makes signal system of bus
desirable queue
CPU scheduling and memory management.

PRACTICAL NO: 11

Aim: To sort the list using merge sort

Theory: Merge sort is a divide and conquer algorithm. It divides input array in two halves and then merges the two halves. The merging operation is used for merging two halves. The merge process is key process that assumes that arr[$i..j..m..n]$ is sorted into one. The array is recursive and splits the two halves till the size becomes 1. Once the size becomes 1, the merge process comes into action and starts merging back till the complete array is merged.

Algorithm:

```
mergeSort (arr, l, r)
    if l < r
        m = l + (r - l) // 2
        mergeSort (arr, l, m)
        mergeSort (arr, m + 1, r)
        merge (arr, l, m, r)
```

Merge sort is useful for sorting linked lists and O(n log n) time complexity.

and the need of random access is low.

Inversion count problem

Used in external sorting

Merge sort is more efficient than quicksort because it only needs to be sorted can only be efficiently addressed parallelly.

Random List: [15, 89, 70, 55, 62, 99, 45, 14, 10]

Merge sorted list: [10, 14, 25, 45, 55, 62, 70, 89, 99]

Time Complexity:
 Best Case: $O(n \log n)$
 Average Case: $O(n \log n)$
 Worst Case: $O(n^2)$