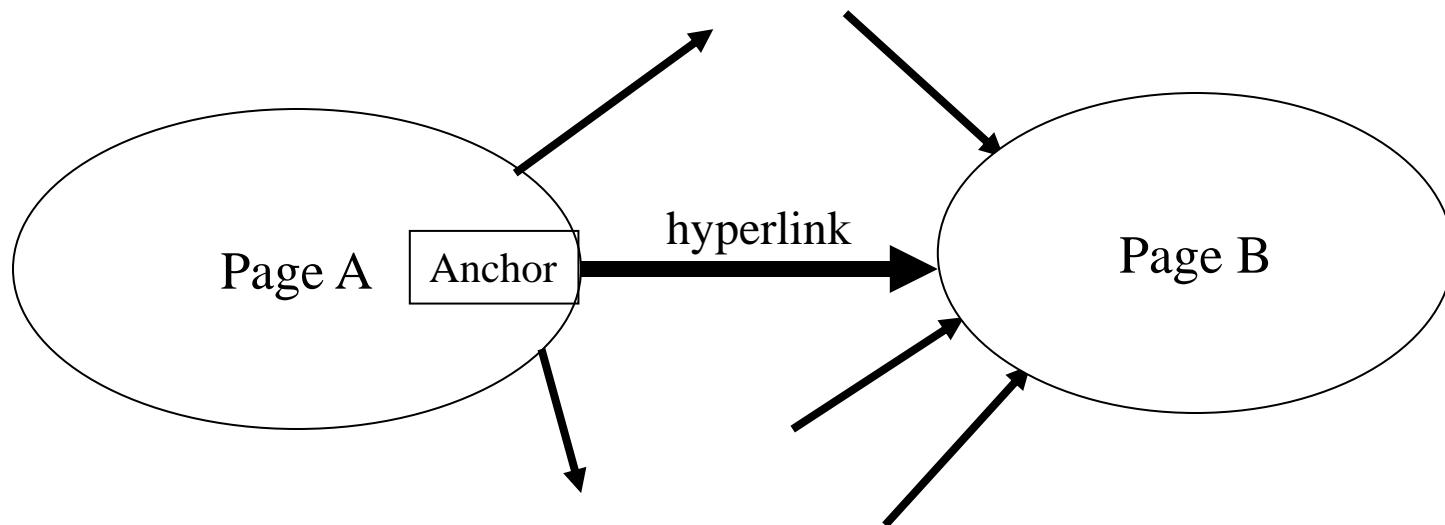


Introduction to
Information Retrieval

Link analysis

The Web as a Directed Graph

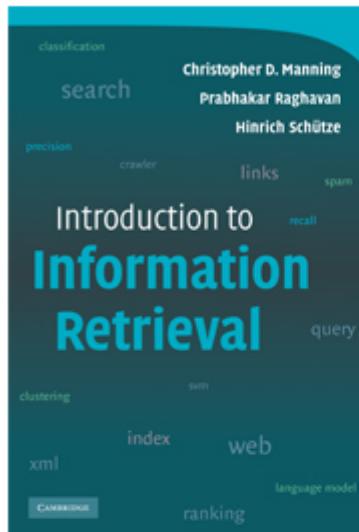


Our study of link analysis builds on two intuitions:

1. The anchor text pointing to page B is a good description of page B.
2. The hyperlink from A to B represents an endorsement of page B, by the creator of page A

Assumption 1: reputed sites

Introduction to Information Retrieval



This is the companion website for the following book.

Christopher D. Manning, Prabhakar Raghavan and Hinrich Schütze, *Introduction to Information Retrieval*

You can order this book at CUP, at your local bookstore or on the internet. The best search

The book aims to provide a modern approach to information retrieval from a computer science perspective. It is available at Cambridge University and at the University of Stuttgart.

We'd be pleased to get feedback about how this book works out as a textbook, what is missing, etc. Please send comments to: informationretrieval (at) yahoogroups (dot) com

Assumption 2: annotation of target

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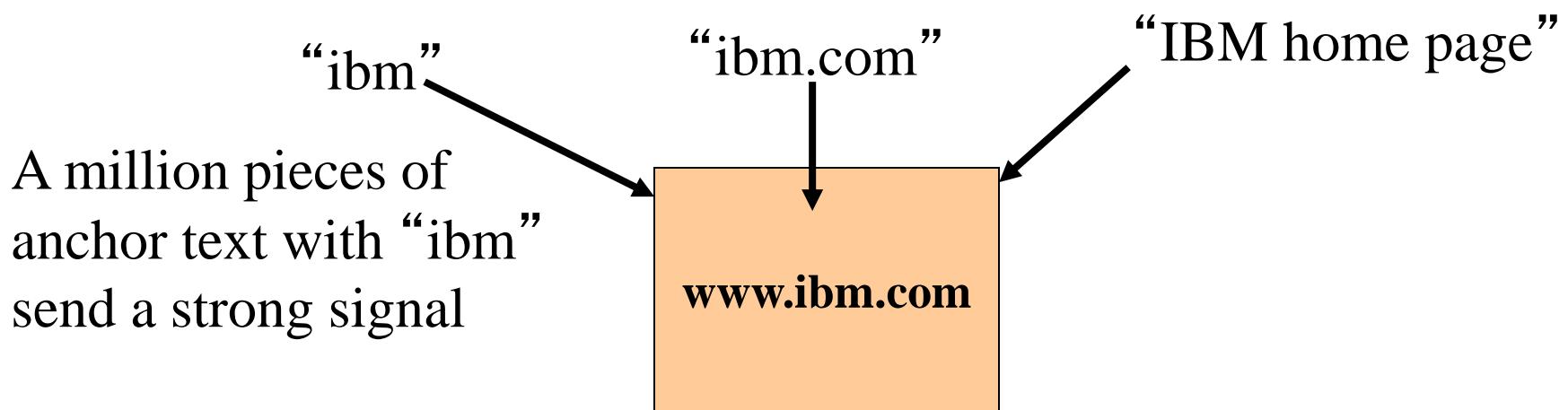
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Anchor Text

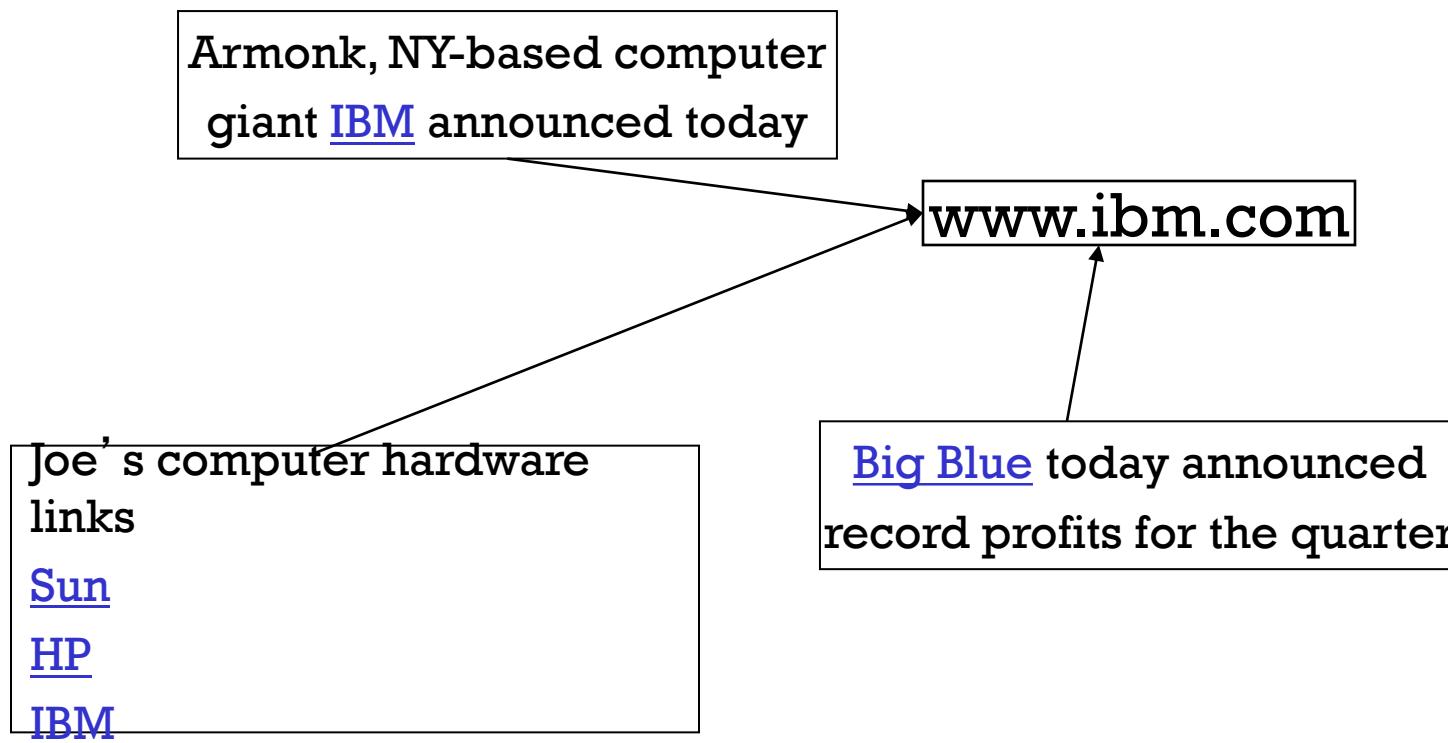
WWW Worm - McBryan [Mcbr94]

- For ***ibm*** how to distinguish between:
 - IBM's home page (mostly graphical)
 - IBM's copyright page (high term freq. for 'ibm')
 - Rival's spam page (arbitrarily high term freq.)



Indexing anchor text

- When indexing a document D , include (with some weight) anchor text from links pointing to D .



Indexing anchor text

- Can sometimes have unexpected effects, e.g., spam, **miserable failure**
- Can score anchor text with weight depending on the authority of the anchor page's website
 - E.g., if we were to assume that content from cnn.com or yahoo.com is authoritative, then trust (more) the anchor text from them
 - Increase the weight of off-site anchors (non-nepotistic scoring)

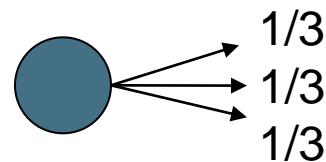
Link analysis: Pagerank

PageRank

- Every node in the web graph is assigned a numerical score between 0 and 1, known as its *PageRank*.
- The PageRank of a node will depend on the link structure of the web graph.
- Given a query, a web search engine computes a composite score for each web page that combines hundreds of features such as cosine similarity and term proximity , together with the PageRank score.

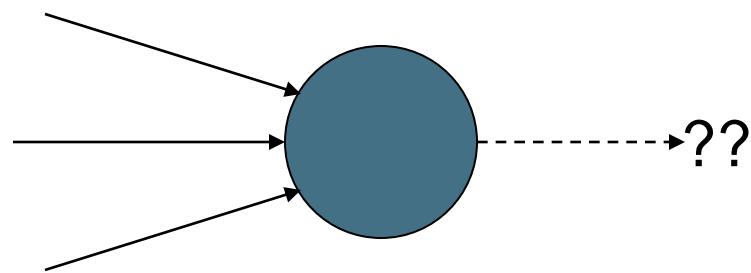
Pagerank scoring – Random walk

- Imagine a user doing a random walk on web pages:
 - Start at a random page
 - At each step, go out of the current page along one of the links on that page, equiprobably
- “In the long run” each page has a long-term visit rate
 - use this as the page’s score.



Not quite enough

- The web is full of dead-ends.
 - Random walk can get stuck in dead-ends.
 - Makes no sense to talk about long-term visit rates.



Teleporting

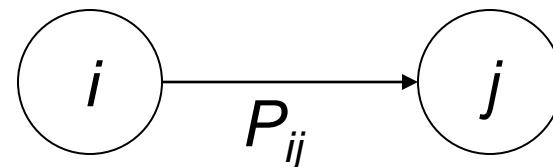
- What if the current location of the surfer, the node A, has no out-links?
- To address this we introduce an additional operation for our random surfer: the *teleport* operation.
- In the teleport operation the surfer jumps from a node to any other node in the web graph.
- If N is the total number of nodes in the web graph¹, the teleport operation takes the surfer to each node with probability $1/N$.
- The surfer would also teleport to his present position with probability $1/N$.

-
- In assigning a PageRank score to each node of the web graph, we use the teleport operation in two ways:
 - (1) When at a node with no out-links, the surfer invokes the teleport operation.
 - (2) At any node that has outgoing links, the surfer invokes the teleport operation with probability $0 < \alpha < 1$ and the standard random walk with probability $1 - \alpha$, where α is a fixed parameter chosen in advance.
- Typically, α might be 0.1.

Markov chains

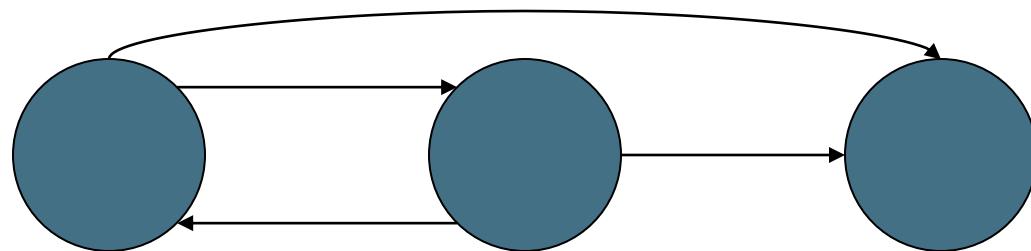
- A Markov chain consists of n states, plus an $n \times n$ transition probability matrix \mathbf{P} .
- At each step, we are in one of the states.
- For $1 \leq i, j \leq n$, the matrix entry P_{ij} tells us the probability of j being the next state, given we are currently in state i .

$P_{ii} > 0$
is OK.



Markov chains

- Clearly, for all i , $\sum_{j=1}^n P_{ij} = 1$.
- **Markov chains are abstractions of random walks.**
- *Exercise:* represent the teleporting random walk from 3 slides ago as a Markov chain, for this case:



Ergodic Markov chains

- For any *ergodic* Markov chain, there is a unique long-term visit rate for each state.
 - *Steady-state probability distribution.*
- Over a long time-period, we visit each state in proportion to this rate.
- It doesn't matter where we start.

-
- We can readily derive the transition probability matrix P for our Markov chain from the $N \times N$ matrix A :
 1. If a row of A has no 1's, then replace each element by $1/N$. For all other rows proceed as follows.
 2. Divide each 1 in A by the number of 1's in its row. Thus, if there is a row with three 1's, then each of them is replaced by $1/3$.
 3. Multiply the resulting matrix by $1 - \alpha$.
 4. Add α/N to every entry of the resulting matrix, to obtain P .

Probability vectors

- A probability (row) vector $\mathbf{x} = (x_1, \dots x_n)$ tells us where the walk is at any point.
- E.g., (000...1...000) means we're in state i .

1 i n

More generally, the vector $\mathbf{x} = (x_1, \dots x_n)$ means the walk is in state i with probability x_i .

$$\sum_{i=1}^n x_i = 1.$$

Change in probability vector

- If the probability vector is $\mathbf{x} = (x_1, \dots x_n)$ at this step, what is it at the next step?
- Recall that row i of the transition prob. Matrix \mathbf{P} tells us where we go next from state i .
- So from \mathbf{x} , our next state is distributed as $\mathbf{x}\mathbf{P}$
 - The one after that is $\mathbf{x}\mathbf{P}^2$, then $\mathbf{x}\mathbf{P}^3$, etc.
 - (Where) Does this converge?

How do we compute this vector?

- Let $\mathbf{a} = (a_1, \dots, a_n)$ denote the row vector of steady-state probabilities.
- If our current position is described by \mathbf{a} , then the next step is distributed as $\mathbf{a}\mathbf{P}$.
- But \mathbf{a} is the steady state, so $\mathbf{a}=\mathbf{a}\mathbf{P}$.
- Solving this matrix equation gives us \mathbf{a} .
 - So \mathbf{a} is the (left) eigenvector for \mathbf{P} .
 - (Corresponds to the “principal” eigenvector of \mathbf{P} with the largest eigenvalue.)
 - Transition probability matrices always have largest eigenvalue 1.

Link analysis: HITS

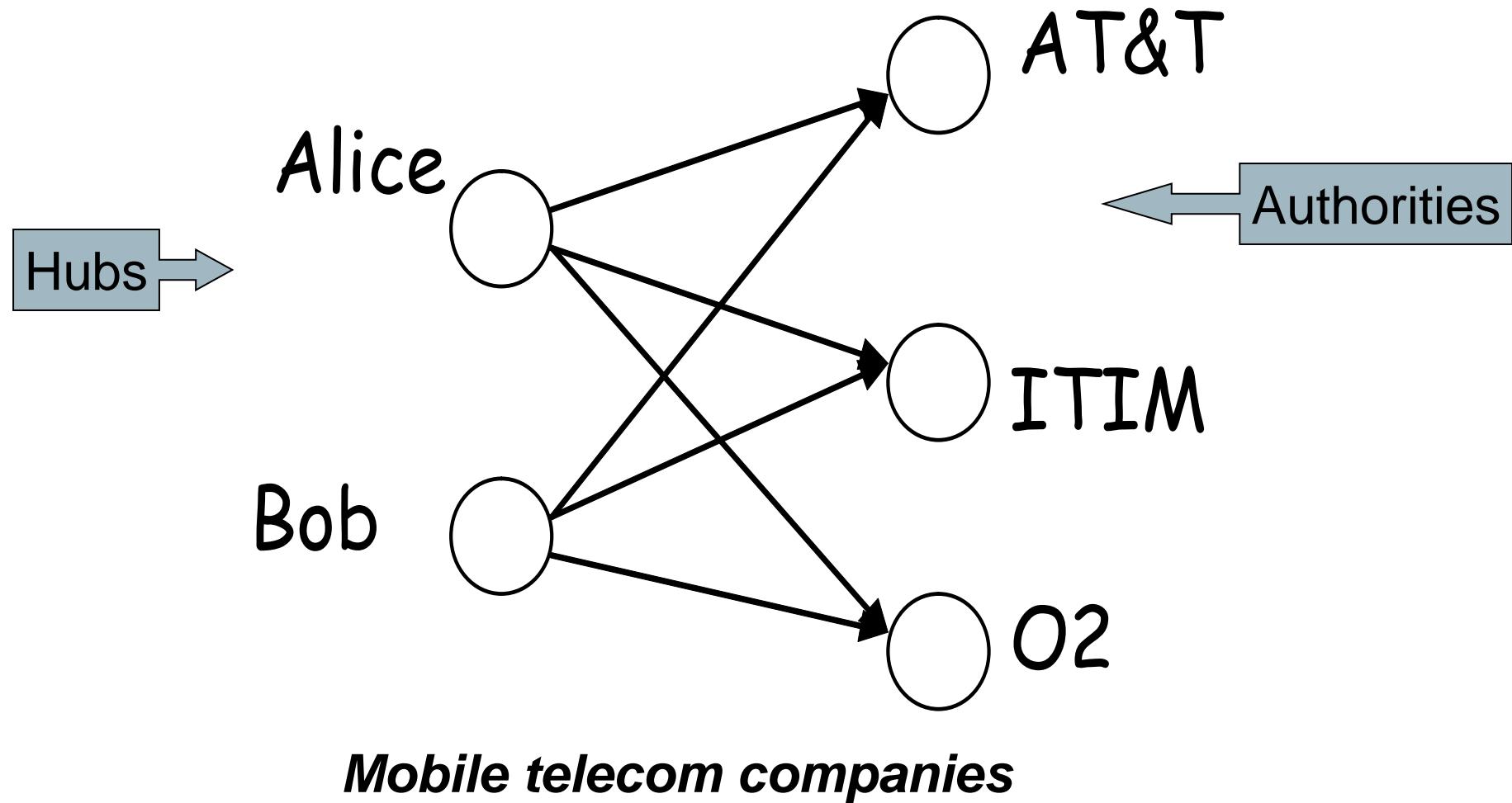
Hyperlink-Induced Topic Search (HITS)

- In response to a query, instead of an ordered list of pages each meeting the query, find two sets of inter-related pages:
 - *Hub pages* are good lists of links on a subject.
 - e.g., “Bob’s list of cancer-related links.”
 - *Authority pages* occur recurrently on good hubs for the subject.
- Best suited for “broad topic” queries rather than for page-finding queries.
- Gets at a broader slice of common *opinion*.

Hubs and Authorities

- Thus, a good hub page for a topic *points* to many authoritative pages for that topic.
- A good authority page for a topic is *pointed* to by many good hubs for that topic.
- Circular definition - will turn this into an iterative computation.

The hope



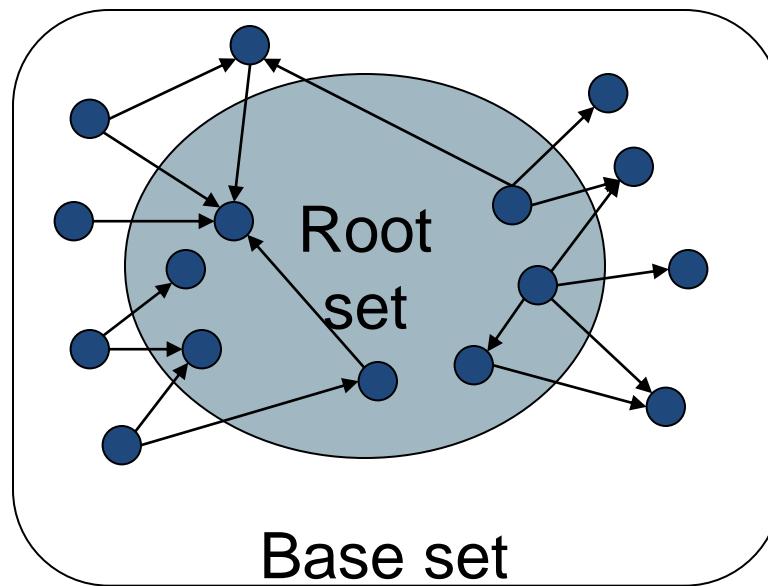
High-level scheme

- Extract from the web a base set of pages that *could* be good hubs or authorities.
- From these, identify a small set of top hub and authority pages;
→iterative algorithm.

Base set

- Given text query (say *browser*), use a text index to get all pages containing *browser*.
 - Call this the root set of pages.
- Add in any page that either
 - points to a page in the root set, or
 - is pointed to by a page in the root set.
- Call this the base set.

Visualization



Get in-links (and out-links) from a *connectivity server*

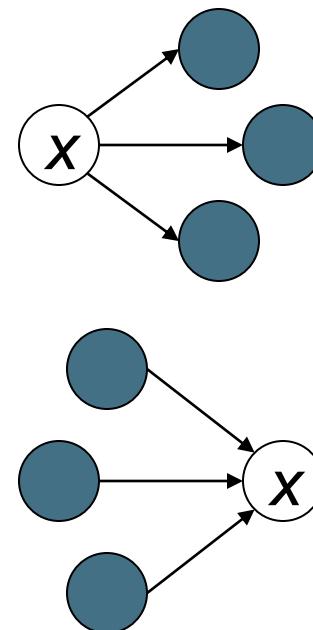
Distilling hubs and authorities

- Compute, for each page x in the base set, a hub score $h(x)$ and an authority score $a(x)$.
- Initialize: for all x , $h(x) \leftarrow 1$; $a(x) \leftarrow 1$;
- Iteratively update all $h(x)$, $a(x)$;  Key
- After iterations
 - output pages with highest $h()$ scores as top hubs
 - highest $a()$ scores as top authorities.

Iterative update

- Repeat the following updates, for all x :

$$h(x) \leftarrow \sum_{x \mapsto y} a(y)$$



$$a(x) \leftarrow \sum_{y \mapsto x} h(y)$$

Scaling

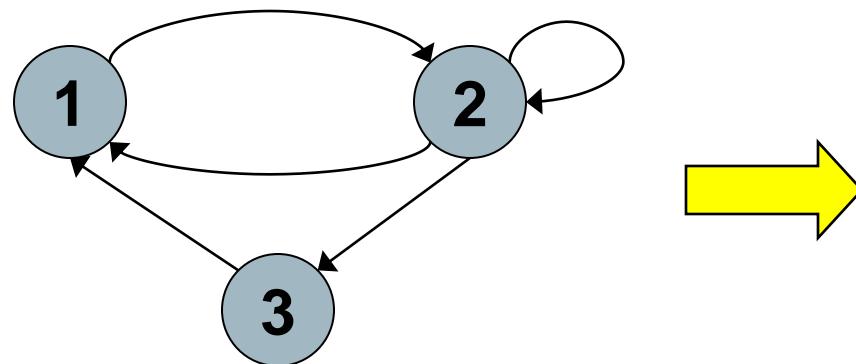
- To prevent the $h()$ and $a()$ values from getting too big, can scale down after each iteration.
- Scaling factor doesn't really matter:
 - we only care about the *relative* values of the scores.

How many iterations?

- Claim: relative values of scores will converge after a few iterations:
 - in fact, suitably scaled, $h()$ and $a()$ scores settle into a steady state!
 - proof of this comes later.
- In practice, ~5 iterations get you close to stability.

Proof of convergence

- $n \times n$ adjacency matrix A:
 - each of the n pages in the base set has a row and column in the matrix.
 - Entry $A_{ij} = 1$ if page i links to page j , else = 0.



	1	2	3
1	0	1	0
2	1	1	1
3	1	0	0

Hub/authority vectors

- View the hub scores $h()$ and the authority scores $a()$ as vectors with n components.
- Recall the iterative updates

$$h(x) \leftarrow \sum_{x \mapsto y} a(y)$$

$$a(x) \leftarrow \sum_{y \mapsto x} h(y)$$

Rewrite in matrix form

- $\mathbf{h} = \mathbf{A}\mathbf{a}$.
- $\mathbf{a} = \mathbf{A}^t\mathbf{h}$.

Recall \mathbf{A}^t is the transpose of \mathbf{A} .

Substituting, $\mathbf{h} = \mathbf{A}\mathbf{A}^t\mathbf{h}$ and $\mathbf{a} = \mathbf{A}^t\mathbf{A}\mathbf{a}$.

Thus, \mathbf{h} is an eigenvector of $\mathbf{A}\mathbf{A}^t$ and \mathbf{a} is an eigenvector of $\mathbf{A}^t\mathbf{A}$.

Further, our algorithm is a particular, known algorithm for computing eigenvectors: the *power iteration* method.

Guaranteed to converge.