

# 1334. Find the City With the Smallest Number of Neighbors at a Threshold Distance

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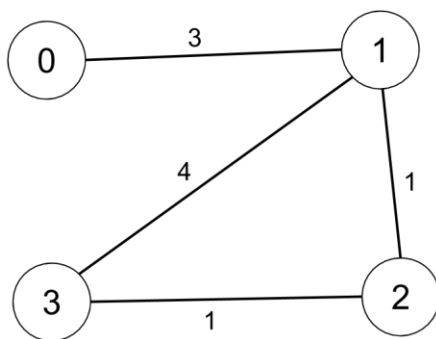
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## Problem Statement:

There are  $n$  cities numbered from 0 to  $n-1$ . Given the array `edges` where `edges[i] = [fromi, toi, weighti]` represents a bidirectional and weighted edge between cities `fromi` and `toi`, and given the integer `distanceThreshold`. Return the city with the smallest number of cities that are reachable through some path and whose distance is at most `distanceThreshold`. If there are multiple such cities, return the city with the greatest number. Notice that the distance of a path connecting cities  $i$  and  $j$  is equal to the sum of the edges' weights along that path.

## Examples

Example 1:



Input:  $n = 4$ , `edges = [[0,1,3],[1,2,1],[1,3,4],[2,3,1]]`,

`distanceThreshold = 4`

Output: 3

Explanation: The figure above describes the graph. The neighboring cities at a distanceThreshold = 4 for each city are:

City 0 -> [City 1, City 2]

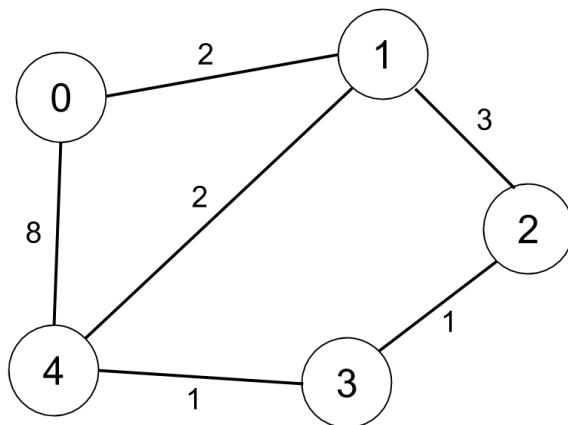
City 1 -> [City 0, City 2, City 3]

City 2 -> [City 0, City 1, City 3]

City 3 -> [City 1, City 2]

Cities 0 and 3 have 2 neighboring cities at a distanceThreshold = 4, but we have to return city 3 since it has the greatest number.

Example 2:



Input: n = 5,

edges = [[0,1,2],[0,4,8],[1,2,3],[1,4,2],[2,3,1],[3,4,1]],

distanceThreshold = 2

Output: 0

Explanation: The figure above describes the graph. The neighboring cities at a distanceThreshold = 2 for each city are:

City 0 -> [City 1]

City 1 -> [City 0, City 4]

City 2 -> [City 3, City 4]

City 3 -> [City 2, City 4]

City 4 -> [City 1, City 2, City 3]

The city 0 has 1 neighboring city at a distanceThreshold = 2.

## Approach for Dijkstra's algorithm:

To solve the problem of finding the city with the smallest number of reachable cities within a given distance threshold using Dijkstra's algorithm, follow these steps:

### Step-by-Step Approach

#### 1) Graph Representation:

- a) Use an adjacency list to store the graph. Each city will have a list of pairs representing the connected cities and the distances to them.

#### 2) Dijkstra's Algorithm Implementation:

- a) Initialize a distance array to store the shortest distance from the starting city to each other city. Set all distances to infinity (INT\_MAX), except the starting city which is set to 0.
- b) Use a priority queue (min-heap) to process the cities based on the shortest known distance.
- c) For the current city, update the distances to its neighboring cities if a shorter path is found.

#### 3) Iterate and Count Reachable Cities:

- a) For each city, run Dijkstra's algorithm to get the shortest distances to all other cities.
- b) Count how many cities have a distance less than or equal to the distance threshold.

#### 4) Track Minimum Reachable Cities:

- a) Keep track of the minimum number of reachable cities and the corresponding city index. If there is a tie, choose the city with the larger index.

## Solution Code for Dijkstra's algorithm:

```
#include <iostream>
#include <vector>
#include <queue>
#include <limits.h>
#include <algorithm>

using namespace std;

typedef pair<int, int> pii;

class Solution {
public:
    vector<vector<pii>> graph;

    vector<int> dijkstra(int n, int start, int distanceThreshold) {
        vector<int> distances(n, INT_MAX);
        distances[start] = 0;
        priority_queue<pii, vector<pii>, greater<pii>> pq;
        pq.push({0, start});

        while (!pq.empty()) {
            int u = pq.top().second;
            int dist = pq.top().first;
            pq.pop();

            if (dist > distances[u]) continue;

            for (const auto& edge : graph[u]) {
                int v = edge.first;
                int weight = edge.second;

                if (distances[u] + weight < distances[v]) {
                    distances[v] = distances[u] + weight;
                    pq.push({distances[v], v});
                }
            }
        }

        return distances;
    }

    int findTheCity(int n, vector<vector<int>>& edges, int distanceThreshold) {
        graph = vector<vector<pii>>(n);

        for (const auto& edge : edges) {
```

```

        int u = edge[0];
        int v = edge[1];
        int weight = edge[2];
        graph[u].emplace_back(v, weight);
        graph[v].emplace_back(u, weight);
    }

    int minReachableCount = INT_MAX;
    int resultCity = -1;

    for (int i = 0; i < n; ++i) {
        vector<int> distances = dijkstra(n, i, distanceThreshold);
        int reachableCount = 0;

        for (int j = 0; j < n; ++j) {
            if (i != j && distances[j] <= distanceThreshold) {
                reachableCount++;
            }
        }

        if (reachableCount < minReachableCount || (reachableCount ==
minReachableCount && i > resultCity)) {
            minReachableCount = reachableCount;
            resultCity = i;
        }
    }

    return resultCity;
}
};

```

## Approach for Bellman-Ford algorithm:

To solve the problem of finding the city with the smallest number of reachable cities within a given distance threshold using the Bellman-Ford algorithm, follow these steps:

### Step-by-Step Approach

#### 1) Graph Representation:

- a) Use an edge list to store the graph. Each edge will be represented as a triplet (u, v, w) where u and v are the cities and w is the weight (distance).

## 2) Bellman-Ford Algorithm Implementation:

- a) Initialize a distance array to store the shortest distance from the starting city to each other city. Set all distances to infinity (INT\_MAX), except the starting city, which is set to 0.
- b) Relax all edges  $|V| - 1$  times (where  $|V|$  is the number of vertices).
- c) Optionally, a final pass through all edges can detect negative weight cycles, but this is unnecessary for this problem since distances are non-negative.

## 3) Iterate and Count Reachable Cities:

- a) For each city, run the Bellman-Ford algorithm to get the shortest distances to all other cities.
- b) Count how many cities have a distance less than or equal to the distance threshold.

## 4) Track Minimum Reachable Cities:

- a) Keep track of the minimum number of reachable cities and the corresponding city index. If there is a tie, choose the city with the larger index.

## Solution Code for Bellman-Ford's algorithm:

```
#include <vector>
#include <limits.h>
using namespace std;

class Solution {
public:
    vector<int> bellmanFord(int n, int start, const vector<vector<int>>& edges) {
        vector<int> dist(n, INT_MAX);
        dist[start] = 0;

        for (int i = 0; i < n - 1; ++i) {
            for (const auto& edge : edges) {
                int u = edge[0];
                int v = edge[1];
                int weight = edge[2];

                if (dist[u] != INT_MAX && dist[u] + weight < dist[v]) {
                    dist[v] = dist[u] + weight;
                }
            }
        }

        return dist;
    }
};
```

```

    }

    int findTheCity(int n, vector<vector<int>>& edges, int distanceThreshold) {
        int minReachableCount = INT_MAX;
        int resultCity = -1;

        for (int i = 0; i < n; ++i) {
            vector<int> distances = bellmanFord(n, i, edges);
            int reachableCount = 0;

            for (int j = 0; j < n; ++j) {
                if (i != j && distances[j] <= distanceThreshold) {
                    reachableCount++;
                }
            }

            if (reachableCount < minReachableCount || (reachableCount ==
minReachableCount && i > resultCity)) {
                minReachableCount = reachableCount;
                resultCity = i;
            }
        }

        return resultCity;
    }
};

```