

LPP

I Canonical form : $\rightarrow \text{Max } Z = \sum_{i=1}^n c_i x_i$ Subject to $\sum_{j=1}^m a_{ij} x_j \leq b_i \quad i=1, 2, 3, \dots, n$
 $x_{ij} \geq 0 \quad j=1, 2, \dots, m$

Characteristics are as follows \rightarrow

(a) If Objective function is of minimization type
 $\sum a_i x_i \geq b_i$

(b)

maximization

$\sum a_i x_i \leq b_i$

(c) If the constraint are of equality type then write in form of inequality

$a_1 x_1 + a_2 x_2 = b$, then $a_1 x_1 + a_2 x_2 \leq b$, & $a_1 x_1 + a_2 x_2 \geq b$.

(d) If any variable is unrestricted eg x_i , then write $x_i - x_i' = x_i$, where $x_i, x_i' \geq 0$

II Standard form : \rightarrow In standard form introduce slack variables & express objective function (Z) as well as all constraints in form of equalities

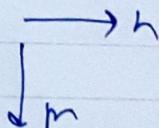
$$Z = c_1 x_1 + c_2 x_2 + \dots + c_n x_n + 0 s_1 + 0 s_2 + \dots + 0 s_m$$

$$\text{Subject to } a_{11} x_1 + a_{12} x_2 + \dots + a_{1n} x_n + s_1 + 0 s_2 + 0 s_3 + \dots + 0 s_m = b_1$$

$$\text{i.e. } \sum_{j=1}^n a_{ij} x_j + s_i = b_i$$

All the conditions are expressed in equations using slack variables

Right hand side of all constraint must be non-negative
~~ie~~ \geq must be of maximizer
 $x_1, x_2, \dots, x_n, s_1, \dots, s_m \geq 0$
 else multiply by -1



① $Z = -3x_1 + 2x_2 - x_3$

 $x_1 - 3x_2 + 2x_3 \geq -6$
 $3x_1 + 4x_3 \leq 3$
 $-3x_1 + 5x_2 \leq 4$
 $x_1, x_2 \geq 0, x_3 \text{ unrestricted}$

Since x_3 is unrestricted
 $x_3 = x_3' - x_3'' \quad x_3', x_3'' > 0$

Maximize $Z' = 3x_1 - 2x_2 + x_3 + 0s_1 + 0s_2 + 0s_3$
 subject to

$$-x_1 + 3x_2 - 2x_3 + s_1 = -6 \quad + 0s_2 + 0s_3$$

$$3x_1 + 4x_2 + 4x_3 + s_2 + s_3 = 3$$

$$-3x_1 + 5x_2 + 0x_3 + 0s_1 + 0s_2 + s_3 = 4$$

$$x_i, s_i \geq 0$$

→ Definitions →

- i) set of values which satisfy constraint is solution
- ii) solution of LPP which satisfy non-negativity condition is feasible solution
- iii) feasible solution which optimizes feasible objective function is called optimal feasible solution.
- iv) Basic feasible solution is solution obtained by putting any n variables to zero & finding values of remaining m variables then variables which are not equal to zero are called basic solution variables & other n zero valued are non basic variables.

No of Basic Solution is $\binom{n+m}{m}$

v) Basic feasible solution → satisfy non negative condition & basic solution (BFS)

Two types of BFS

i) Non degenerate BFS → if in basic feasible solution, all m values of basic solution variable are +ve, non zero then it is non degenerate basic feasible solution

ii) Degenerate BFS → if one or more m basic variables are 0 then degenerate BFS

① Which of the following are degenerate, basic, & solution?

$$\text{Maximize } Z = x_1 + 3x_2 + 3x_3$$

optimal

$$x_1 + 2x_2 + 3x_3 \leq 4$$

$$2x_1 + 3x_2 + 5x_3 = 7$$

→ No of Variables 3, No of constraint 2

∴ No of basic solution can be obtained by 3 - 2 = 1 variable to 0

No of ^① Basic solutio NonBasic ^② equals to 0 ^③ Basic variables ^④ Equation ^⑤ Values

⑥

$Z \geq 0$

⑦

Degenerate

⑧

Value of Z

⑨

To optimal

⑩

⑪

⑫

⑬

⑭

⑮

⑯

⑰

⑱

⑲

⑳

5
4
3
2
1

$$\begin{array}{l} \text{Row 1: } x_1 + 2x_2 + 4x_3 = 4 \\ \text{Row 2: } 2x_1 + 3x_2 + 3x_3 = 7 \\ \text{Row 3: } 2x_1 + 3x_2 + 5x_3 = 7 \end{array}$$

$$\begin{array}{l} \text{Row 1: } x_1 + 2x_2 + 4x_3 = 4 \\ \text{Row 2: } x_2 + x_3 = 1 \\ \text{Row 3: } 2x_2 + 5x_3 = 0 \end{array}$$

$$\begin{array}{l} \text{Row 1: } x_1 + 2x_2 + 4x_3 = 4 \\ \text{Row 2: } x_2 + x_3 = 1 \\ \text{Row 3: } 2x_2 + 5x_3 = 0 \end{array}$$

Optimal solution
Basic

$$x_3 = 0$$

$$x_1 = 0$$

$$x_2 = 0$$

$$x_1 = 0$$

Optimal solution
Basic

Simplex Method

Step 1. Write in standard form

Step 2. Write initial basic solution

Step 3. Make initial simplex table

Step 4. Are all entries in \geq row non negative?

i) If Yes then current solution is optimal basic feasible solution stop here if 2 has more than 5 zeros, else continue

ii) If No select most negative or smallest entry in \geq row. The corresponding column is key column or pivot column & the corresponding variable enters in the basis.

Step 5. Obtain replacement ratio by dividing solution column by key column

Step 6. Are all ratios infinite / And/or negative?

i) If Yes CPP has unbounded solution. Stop Here

ii) If No select minimum finite non negative ratio

Corresponding row is called key row & corresponding variable leaves the basis. In case of tie to select any one.

If any ratio is zero then consider that ratio which has denominator positive in pivot column

Mark

Step 7. Make the key element as intersection of key row & key column

Step 8. Make New simplex method by elementary row transformations

i) Make key element 1 by dividing key row by key element

ii) Make all other elements of key column zero.

Step 9. Go to step 4 & continue



① Maximize $Z = 3x_1 + 2x_2$ subject to $3x_1 + 2x_2 \leq 18$
 $0 \leq x_1 \leq 4$
 $0 \leq x_2 \leq 6$

$\rightarrow Z - 3x_1 - 2x_2 + 0s_1 + 0s_2 + 0s_3 = 0$
 S.T.

$$\begin{aligned} 3x_1 + 2x_2 + s_1 &= 18 \\ &\quad + 0s_2 + 0s_3 \\ x_1 + 0s_1 + s_2 + 0s_3 &= 4 \\ x_2 + 0s_1 + 0s_2 + s_3 &= 6 \end{aligned}$$

$$x_1, x_2 \geq 0, s_1, s_2, s_3 \geq 0$$

Initial basic feasible solution let $x_1 = 0, x_2 = 0$

$$s_1 = 18, s_2 = 4, s_3 = 6$$

$$Z = 0$$

Iteration	Basic Var	Coeff					RHS	Ratio
		x_1	x_2	s_1	s_2	s_3		
0	Z	-3	-2	0	0	0	0	$0/-3 = 0$
	s_1	3	2	1	0	0	18	$18/3 = 6$
	s_2	1	0	0	1	0	4	$4/1 = 4 \rightarrow$
	s_3	0	1	0	0	1	6	$6/0 = \infty$

All entries Not Non Negative. Choose Most small entry
 $\therefore x_2$ is entering basic, column is key column
 \therefore write ratio

All are Not Unchanged \therefore key entry smallest is 1 (Not 0)
 s_3 leaves key element is intersection of key row & key column is 1

1	Z	0	-2	0	3	1	0	12
	s_3	0	2	-1	-3	0	6	$6/2 = 3$
	x_1	1	0	0	1	0	4	$4/1 = 4$
	s_3	0	1	0	0	1	6	$6/1 = 6$

key element already! Make col 0 row 0
 $R_1 \rightarrow 3R_2 + R_1$
 $R_2 \rightarrow R_2 - 3R_1$
 $R_3 \rightarrow R_3 - R_2$

Repeat procedure. x_2 enters S, Leaves.
 key element 2 ∵ Need to make col one.

	$R_2 \rightarrow R_2/2$	$R_1 \rightarrow R_1 + 2R_2$	$R_3 \rightarrow R_3 - R_2$
Z	0	0	1
x_2	0	1	$\frac{1}{2}$
x_4	1	0	0
s_3	0	0	- $\frac{1}{2}$
			3
			4
			3

All entries Non-Negative hence achieved optimal solution
 basic feasible

$$x_1 = 4$$

$$x_2 = 3$$

$$x_4 = 0$$

$$s_2 = 0$$

$$s_3 = 3$$

Q Why is s_1 extra?
 as s_2 is non variable

Entry for non basic variable s_1 is 0. Hence LPP has alternate solution.

The alternate optimum solution is obtained by taking s_2 as entering variable
 ∵ s_2 enters

	0	0	1	0	0	18	Ratio
x_2	0	1	$\frac{1}{2}$	$-\frac{3}{2}$	0	3	$\frac{3}{\frac{1}{2}} = 6$
x_1	1	0	0	0	0	4	$\frac{4}{1} = 4$
s_3	0	0	$-\frac{1}{2}$	$+\frac{7}{2}$	1	3	$\frac{3}{-\frac{1}{2}} = 27$

All entries not non-negative Minimum finite is +2 ∵ s_2 leaves,

$$R_3 \rightarrow \frac{2}{3}R_3 \text{ to make key element 1}$$

$$R_1 \rightarrow \frac{2}{3}R_1 + R_2 \quad R_2 \rightarrow R_2 - R_3 \text{ to make other 0}$$

x_1	x_2	s_1	s_2	s_3	RHS	rate
7	0	0	1	0	0	10
x_2	0	1	0	0	1	6
x_1	1	0	$\frac{1}{3}$	0	$-\frac{2}{3}$	2
s_2	0	0	$-\frac{1}{3}$	1	$\frac{2}{3}$	2

All entries is ≥ 0 & No Negative
None Values

$$\begin{aligned} x_1 &= 2 \\ x_2 &= 6 \\ s_1 &= 0 \\ s_2 &= 2 \end{aligned}$$

$\therefore \infty$ No of nonbasic optimal feasible solution are $X = \lambda x_1 + (1-\lambda)x_2$ where $\lambda \in [0, 1]$

$$X = \begin{bmatrix} x_1 - 4 \\ x_2 - 3 \\ x_3 - 6 \\ s_1 = 0 \\ s_2 = 0 \\ s_3 = 3 \end{bmatrix} \quad X = \begin{bmatrix} 2 \\ 6 \\ 0 \\ 0 \\ 3 \end{bmatrix}$$

$$X = \begin{bmatrix} 2+2\lambda \\ 6-3\lambda \\ 0 \\ 2-2\lambda \\ 3\lambda \end{bmatrix}$$

Since X has only 1 zero \therefore Non basic solution
As it does not have 2 zeros $\Rightarrow 2$ zeros

② Minimize $Z = x_1 - 3x_2 + 3x_3$ S.T. $3x_1 - x_2 + 2x_3 \leq 7$
 $2x_1 + 4x_2 \geq -12$
 $-4x_1 + 3x_2 + 8x_3 \leq 10$

standard form $3x_1 - x_2 + 2x_3 + s_1 + 0s_2 + 0s_3 = 7$

$$-2x_1 + 4x_2 - 0x_3 + 0s_1 + s_2 + 0s_3 = -12$$

$$-4x_1 + 3x_2 + 8x_3 + 0s_1 + 0s_2 + s_3 = 10$$

(Mat Z^T) $Z^T + x_1 + -3x_2 + 3x_3 + 0s_1 + 0s_2 + 0s_3 = 0$
 $x_i, s_i \geq 0$

Initial basic feasible solution $x_1 = 0$ $x_2 = 0$ $x_3 = 0$

$$\therefore S_1 = 7, S_2 = 12, S_3 = 10$$

① Basis		x_1	x_2	x_3	S_1	S_2	S_3	PMS	Ratio
Z	Z	1	-3	+3	0	0	0	0	-0
S_1		3	-1	2	1	0	0	7	-7
S_2		-2	-4	0	0	1	0	12	-3
S_3		-4	3	8	0	0	1	10	$10/3$

x_2 enter	S_3 leave	Z	x_2	S_1	S_2	S_3	PMS	Z_{max}
		-3	0	11	0	0	1	$3/5$
		$5/3$	0	$14/3$	0	0	$1/3$	$34/3$
		$-22/3$	0	$32/3$	0	1	$4/3$	$76/3$
		$4/3$	1	$8/3$	0	0	$1/3$	$10/3$

x_1 enter	S_3 leave	Z	x_1	x_2	x_3	S_1	S_2	S_3	Z_{max}
		0	0	$97/5$	$9/5$	0	$8/5$	$147/5$	
		1	0	$14/5$	$3/5$	0	$7/5$	$37/5$	
		0	0	$468/15$	$66/15$	$140/15$	$160/15$	$106/15$	
		0	1	$96/15$	$12/15$	0	$9/15$	$174/15$	

Thus we have optimal feasible solution as all rows non negative

$$x_1 = 3/5 \quad x_3 = 0 \quad S_2 = 1062/15 \quad Z_{max} = 147/5$$

$$x_2 = 174/15 \quad S_1 = 0 \quad S_3 = 0$$

$$Z_{min} = -Z_{max}$$

$$\therefore Z_{min} = -147/5$$

S.F.

③ $Z = 2x_1 - 2x_2 + 4x_3 - 5x_4 + 0/s_1 + 10s_2$
 $x_1 + 2x_2 + 3x_3 + 0/s_1 + 1/s_2 - 7$
 $3x_1 + 4x_2 + 6x_3 + 0/s_1 + 4/s_2 = 15$

No of variables = 3 No of constraint = 2

∴ No of basic solution can be obtained by $3-2-1$ Variable to 0

$$x_1 = 0, \quad \begin{cases} 2x_2 + 3x_3 = 7 \\ 4x_2 + 6x_3 = 15 \end{cases} \quad \begin{cases} x_2 = ? \\ x_3 = ? \end{cases} \text{ Math error ie. unbounded solution}$$

$$x_2 = 0, \quad \begin{cases} x_1 + 3x_3 = 7 \\ 3x_1 + 6x_3 = 15 \end{cases} \quad \begin{cases} x_1 = ? \\ x_3 = ? \end{cases}$$

$$x_3 = 0, \quad \begin{cases} x_1 + 2x_2 = 7 \\ 3x_1 + 4x_2 = 15 \end{cases} \quad \begin{cases} x_1 = ? \\ x_2 = ? \end{cases}$$

④ No of variables =

④ Max $Z = 10x_1 + x_2 + 2x_3$

$$14x_1 + x_2 - 6x_3 + 3x_4 = 7$$

$$16x_1 + 1/2x_2 - 6x_3 \leq 5$$

$$3x_1 + x_2 - x_3 \leq 0$$

$$x_1, x_2, x_3, x_4 \geq 0$$

S.F. $\rightarrow Z - 10x_1 - x_2 - 2x_3 = 0$

$$14/3x_1 + 1/3x_2 - 2x_3 + 1x_4 = 7/3 \quad \text{Here } x_4 \text{ is slack variable}$$

$$16x_1 + 1/2x_2 - 6x_3 + s_1 = 5$$

$$3x_1 - x_2 - x_3 + s_2 = 0$$

$$x_1, s_1 > 0$$

Initial basic feasible solution $x_1, x_2, x_3 = 0$

$$x_4 = 7/3$$

$$s_1 = 5$$

$$s_2 = 0$$

Iteration	Basis	x_1	x_2	x_3	x_4	s_1	s_2	s_3	RHS Ratio
0	Z	-107	-1	-2	0	0	0	0	-
x_1 enters	x_4	$\frac{14}{3}$	$\frac{1}{3}$	-2	1	0	0	$\frac{4}{3}$	$\frac{1}{2}$
s_1 leaves	s_1	16	$\frac{1}{2}$	-6	0	1	0	5	$\frac{5}{16}$
	(3)		-1	-1	0	0	1	0	0
x_3 enters	Z	0	$-\frac{10}{3}$	$-\frac{17}{3}$	0	0	$\frac{10}{3}$	0	-
x_1 leaves	x_4	0	$\frac{1}{3}$	$-\frac{4}{3}$	1	0	$-\frac{14}{3}$	$\frac{7}{3}$	$-\frac{7}{16}$
s_1	s_1	0	$\frac{35}{6}$	$-\frac{7}{3}$	0	1	$-\frac{16}{3}$	5	$\frac{10}{3}$
	x_1	1	$-\frac{1}{3}$	$-\frac{1}{3}$	0	0	$\frac{1}{3}$	0	$\frac{1}{3}$

All ratios $\rightarrow \infty \therefore$ Unbounded solution
 No variable will leave basis, x_3 cannot enter hence unbounded solution

Big M (Penalty Method)

$$\begin{aligned} \textcircled{1} \quad \text{Maximize } Z &= 6x_1 + 4x_2 \\ 2x_1 + 3x_2 &\leq 30 \\ 3x_1 + 2x_2 &\leq 24 \\ x_1 + x_2 &\geq 3 \\ x_1, x_2 &\geq 0 \end{aligned}$$

Is the solution unique? if not find another solution

$$\begin{aligned} \rightarrow \text{Max } Z &= 6x_1 + 4x_2 - MA_3 \\ 2x_1 + 3x_2 + s_1 &= 30 \\ 3x_1 + 2x_2 + s_2 &= 24 \\ x_1 + x_2 - s_3 + A_3 &= 3 \\ x_1, x_2, x_3, s_1, s_2, s_3, A_3 &\geq 0 \end{aligned}$$

$$\begin{aligned} A_3 \rightarrow \text{Penalty} \\ \text{Remove Penalty by } M \text{ (3)} + \textcircled{1} \\ \therefore Z = (6+M)x_1 + (4+M)x_2 + s_3 \\ -MS_3 - 3M \\ \therefore Z = (6+M)x_1 + (4+M)x_2 + MS_3 - 3M \end{aligned}$$

$$\begin{aligned} \text{Initial QFS Part } 6-3=3 \text{ variable } 0 \\ x_1 = 0 \quad \therefore s_1 = 30 \\ x_2 = 0 \quad \therefore s_2 = 24 \\ x_3 = 0 \quad \therefore A_3 = 3 \end{aligned}$$

Mis van big NO

Basis	x_1	x_2	s_1	s_2	s_3	A_3	RHS	rate
Z	-6-M	-4-M	0	0	M	0	-3M	
S_1	2	3	1	0	0	0	30	15
S_2	3	2	0	1	0	0	24	8
$\leftarrow A_3$	$\textcircled{3} \textcircled{1}^*$	1	0	0	-1	1	3	3

x_1 enters	Z	0	2	0	0	-6	-	18	12
A_3 leaves	S_1	0	1	1	0	2	-	24	12
\therefore leaves	S_2	0	-1	0	1	$\textcircled{3}$ *	-	15	+5
	x_1	1	1	0	0	-1	-	3	-3

S_2 enters	Z	0	0	0	2	0	-	48	42/5
S_1	0	$\textcircled{S_3}$	1	$-2/3$	0	-	14	$+48/5$	
S_3	0	$-1/3$	0	$1/3$	1	-	$15/3-5$	-15	
x_1	1	$2/3$	0	$1/3$	0	-	8	12	

Optimal solution as all solutions are non-negative.

$$x_1 = 8 \quad S_1 = 14 \quad Z_{\max} = 48$$

$$x_2 = 0 \quad S_2 = 0$$

$$x_3 = 5 \quad S_3 = 5$$

There are 3 basic solutions, S_1, S_3 might be 0 in Z but x_2 is also 0
Hence there is an alternative solution. x_1 enters $\therefore S_1$ leaves
 \therefore Infinitely many solutions $x_1 = \begin{bmatrix} 8 \\ 14 \\ 5 \end{bmatrix}$

Z	0	0	0	2	0	-	48	
x_2	0	1	$3/5$	$-2/5$	0	-	$42/5$	
S_3	0	0	$1/5$	$1/5$	1	-	$39/5$	
x_1	1	0	$-2/5$	$3/5$	0	-	$12/5$	

Alternate solution is $X_2 = \begin{bmatrix} 12/5 \\ 42/5 \\ 39/5 \\ 0 \\ 0 \end{bmatrix}$

$$X = \lambda X_1 + (1-\lambda)X_2 = \begin{bmatrix} 8\lambda + (1-\lambda)42/5 \\ 14\lambda + (1-\lambda)42/5 \\ 5\lambda + (1-\lambda)39/5 \\ 0 \\ 0 \end{bmatrix}$$

Non Basic solution
as $n-m-c-3 \Rightarrow$ zeros not present

① Minimize $Z = x_1 + 2x_2 + x_3$

$$x_1 + \frac{1}{2}x_2 + \frac{1}{2}x_3 \leq 1$$

$$\frac{3}{2}x_1 + 2x_2 + x_3 \geq 8$$

$$x_1, x_2, x_3 \geq 0$$

S.F. \rightarrow

$$\text{Max } Z' = -x_1 - 2x_2 - x_3 - M A_2$$

$$\text{Subt } x_1 + \frac{1}{2}x_2 + \frac{1}{2}x_3 + s_1 = 1$$

$$\frac{3}{2}x_1 + 2x_2 + x_3 + -s_2 + A_2 = 8$$

$$x_1, x_2, x_3, s_1, s_2, A_2 \geq 0$$

$$\text{Max } Z \rightarrow Z'$$

$$Z' = (-1 + \frac{3}{2}M)x_1 - (-2 + M)x_2 - (1 - M)x_3 - Ms_2 - 8M$$

Initial basic solution $s_1, s_2 = 0$ Variables to 0

$$x_1, x_2, x_3 = 0$$

$$\therefore s_1 = 1 \quad A_2 = 8$$

Basic	x_1	x_2	x_3	s_1	s_2	A_2	RHS	Ratio
Z'	$1 - \frac{3}{2}M$	$2 - 2M$	$1 - M$	0	M	0	$-8M$	
s_1	1	$\frac{1}{2}$	$\frac{1}{2}$	1	0	0	1	2
A_2	$\frac{3}{2}$	2	1	0	-1	1	8	4

x_2 enters	Z	$-2 + M$	0	$-1 + M$	$-4 + 4M$	M	0	$-4 + 4M$
s_1 leaves	x_2	2	1	1	2	0	0	2
	A_2	$-\frac{3}{2}$	0	$\frac{1}{2}$	$-\frac{1}{2}$	-1	1	4

All entries Non Negative but A_2 remains in basis with positive value

$A_2 > 0 \quad \therefore \text{Pseudo solution}$