## 11.01.2025

# **Problem Name: word subsets**

#### **Problem Statement:**

```
You are given two string arrays <a href="words1">words1</a> and <a href="words2">words2</a>.

A string <a href="b">b</a> is a subset of string <a href="a if every letter in b">a</a> occurs in <a href="a including multiplicity">a including multiplicity</a>.

A string <a href="a from words1">a from words1</a> is <a href="words1">universal</a> if for every string <a href="b in words2">b</a> in words2</a>, <a href="b in words2">b</a> is a subset of <a href="a a.">a</a>.
```

Return an array of all the universal strings in words1.

#### **Example 1:**

```
Input: words1 = ["amazon","apple","facebook","google","leetcode"], words2 =
["e","o"]
Output: ["facebook","google","leetcode"]
```

#### Example 2:

```
Input: words1 = ["amazon", "apple", "facebook", "google", "leetcode"], words2 =
["lc", "eo"]
Output: ["leetcode"]
```

#### Example 3:

```
Input: words1 = ["acaac","cccbb","aacbb","caacc","bcbbb"], words2 =
["c","cc","b"]
Output: ["cccbb"]
```

## **Approach: Using Frequency Count (Optimized)**

#### Key Idea:

- To determine if a word is **universal**, it must satisfy the condition of being a **superstring** for all strings in words2.
- · Convert the problem into frequency counting.

### Steps:

- 1. Calculate the Maximum Requirement for words2:
  - a. Compute the **maximum frequency** of each character across all words in words2.
- 2. Check each word in words1:
  - a. Count the frequency of characters for each word in words1.
  - b. Compare with the **maximum required frequency** from words2.

#### C++ Code:

```
class Solution {
public:
    // Helper function to check if word1 can cover the requirements
    bool isSubSet(vector<int>& charRequired, vector<int>& temp) {
       for (int i = 0; i < 26; i++) {
            if (temp[i] < charRequired[i])</pre>
               return false;
        7
       return true;
    }
    vector<string> wordSubsets(vector<string>& words1, vector<string>& words2) {
        vector<string> res;
       vector<int> charRequired(26, 0); // Store the maximum frequency needed
for each character
        // Step 1: Calculate the maximum frequency requirement across words2
        for (string& word : words2) {
            vector<int> temp(26, 0);
            for (char ch : word) {
                temp[ch - 'a']++;
                charRequired[ch - 'a'] = max(charRequired[ch - 'a'], temp[ch -
'a']);
           3
        7
        // Step 2: Check each word in words1 against the required frequency
        for (string& word : words1) {
            vector<int> temp(26, 0);
            for (char ch : word) {
               temp[ch - 'a']++;
            3
            if (isSubSet(charRequired, temp)) {
               res.push_back(word);
       return res;
   3
3;
```

### **Explanation:**

1. Input:

```
a. words1 = ["amazon", "apple", "facebook", "google", "leetcode"]
b. words2 = ["e", "o"]
```

2. **Step 1:** Calculate the maximum frequency requirement from words2.

```
a. e -> 1, o -> 1 \rightarrow charRequired = [0,0,0,...1,0,...1,...0]
```

3. Step 2: Check each word in words1.

```
a. "facebook" has both "e" and "o" → valid.
```

b. "amazon" → doesn't satisfy "o" requirement.

### **Complexity Analysis:**

• Time Complexity:

```
a. 0(n * 1 + m * k)
b. n = number of words in words1.
c. 1 = average length of words in words1.
d. m = number of words in words2.
e. k = average length of words in words2.
```

Space Complexity:

a. 0(26) for storing the frequency count (constant space).

# Edge Cases Handled:

- If words1 or words2 is empty.
- If all words in words1 are subsets of words2.

#### **Test Cases:**

```
Input: words1 = ["a","b","c"], words2 = ["a"]
Output: ["a"]

Input: words1 = ["abc","def","ghi"], words2 = ["z"]
Output: []
```

### **Summary:**

· Use a frequency count approach for efficiency.

- Count the maximum required frequency for each character in words2.
- Validate each word in words1 using the **subset** condition.

Would you like me to provide more test cases or further optimize this? 😊