

## Lecture - 49

### (Binary Search Algo)

#### ★ Binary Search Code:-

```

int Binary Search (vector<int> &input, int target) {
    // Define Search Space
    int lo = 0;           // Start of search space
    int hi = input.size() - 1; // End of search space

    while (lo <= hi) {
        // Calc the midpoint for the search space
        int mid = (lo + hi) / 2;

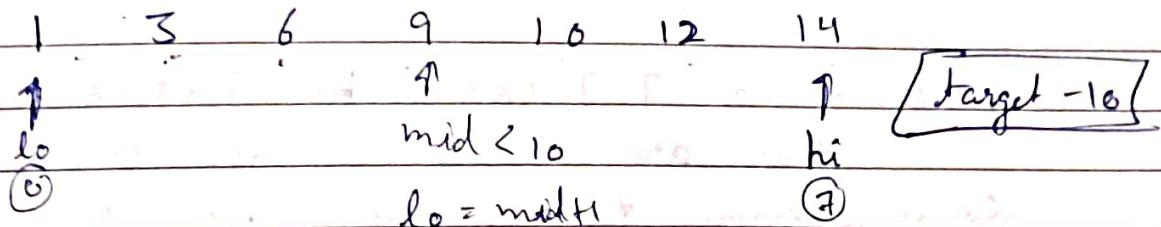
        if (target == input[mid]) {
            return mid;
        } else if (target < input[mid]) {
            hi = mid - 1; // Discard the right of mid.
        } else { // target > input[mid]
            lo = mid + 1; // Discard the left of mid.
        }
    }

    return -1;
}

int main() {
    vector<int> V{1, 3, 6, 9, 10, 12, 14};
    int n = size_of(V); // size of input
    int target = 10;
    cout << binarySearch(V, target) << "\n";
    return 0;
}

```

## ④ Time complexity Analysis



$$N + \frac{N}{2} + \frac{N}{4} + \frac{N}{8} + \dots$$

$$\frac{N}{2^0} + \frac{N}{2^1} + \frac{N}{2^2} + \frac{N}{2^3} + \dots = \frac{N}{2^K}$$

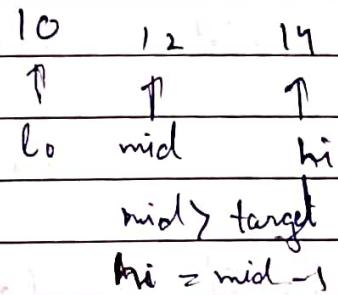
$$\frac{N}{2^K} = 1$$

$$N = 2^K$$

$$\log_2 N = \log_2 2^K$$

$$\log_2 N = K \log_2 2$$

$$K = \log_2 N$$



$$10 12$$

$$\uparrow$$

$$lo$$

$$mid = 10$$

$$mid = lo$$

target found

$O(\log N)$

## ⑤ Space complexity 2 - $O(1)$

### ⑥ Recursive Binary Search Implementation

$f(\text{input}, \text{target}, \text{lo}, \text{hi}) =$

$[\text{lo}, \text{hi}]$  denotes search space

finds the target inside  
 $\text{input} [\text{lo}, \text{hi}]$

```
int mid = (lo + hi) / 2;
if (lo > hi) return -1;
if (input[mid] == target) return mid;
if (input[mid] < target) {
    return f(input, target, mid + 1, hi);
} else {
    return f(input, target, lo, mid - 1)
}
```

→ Time complexity  $\rightarrow O(\log n)$   
 Space complexity  $\rightarrow O(1)$

→ When ;  $l_0 = \text{INT\_MAX}$ ,  $h_0 = \text{INT\_MAX}$  then  
 there are chances of overflow  
 So, to overcome this problem we use

$$\text{Use } \text{mid} = \frac{l_0 + h_0 - l_0}{2}$$

$$= \frac{2l_0 + (h_0 - l_0)}{2}$$

$$\boxed{\text{mid} = \frac{l_0 + (h_0 - l_0)}{2}} \quad // \text{Modified mid}$$

// to tackle overflow

Q-1) Find the first occurrence of a given element  $x$ , given that the given array is sorted. If no occurrence of  $x$  is found then return  $-1$ .

input :- arr = [2, 5, 5, 5, 6, 6, 8, 9, 9, 9]

$x = 5$

$\hookrightarrow m \rightarrow \text{current} = \text{target}$

output :- 1

↓  
 index  $m$  might be a possible ans.

this  $m$  might be our ans. or we can find something better than this.

#include <iostream>

#include <vector>

Using namespace std;

→ Time complexity  $\rightarrow O(\log n)$   
→ Space "  $\rightarrow O(1)$

Void firstOccurance (vector<int> &input, int target) {

int lo = 0;  
int hi = input.size() - 1;  
int ans = -1;

while (lo <= hi) {

mid = lo + (hi - lo) / 2

if (a[mid] >= target) if (input[mid] == target) {

ans = mid;

hi = mid - 1

} else if (input[mid] > target) {

hi = mid - 1;

else {

lo = mid + 1;

}

}

return ans;

int main() {

n → enter

Vector → enter

target → enter

cout << firstOccurance (vector, target) <"\n";  
return 0;

}

Q) Find the square root of the given non-negative value  $x$ . Round it off to the nearest floor integer value.

input  $x = 4 \quad | \quad 11$

Output  $\Rightarrow 2 \quad | \quad 3$

$$\sqrt{x} \rightarrow [1, n]$$

within this range find the biggest value whose square is less than or equal to  $n$ .

Method 1

ans = -1

Linear Search

```
for (i=1; i <= n; i++) {
    if (i * i < n) {
        ans = i;
    }
    else {
        break;
    }
}
```

Time complexity

$O(\sqrt{n})$

Method 2 Binary Search      int ans = -1;      while ( $lo \leq hi$ ) {

int mid =  $lo + (hi - lo)/2$ ;      sqrt (int n)

[i, n]      if (mid \* mid < target) {

$lo = 1, hi = n$

ans = mid;

~~log~~  $lo = mid + 1$ ;

} else {

$hi = mid - 1$ ;

int main () {

cout << sqrt(24) << "\n";

## Lecture - 45

### (Problem Solving on Binary Search - 1)

Q1 Given an array of integers 'a' that is sorted in non-decreasing order. Find the first and the last position of the given 'target' element in the sorted array. Followed zero-based indexing.  
if target is not found, return  $[1, -1]$ .  
input :- arr[ ] = [1, 2, 3, 3, 3, 5, 11], target = 3

output  $\approx [2, 4]$

Sol 1 Same as Q1 For 1st occurrence  $\rightarrow$  Same as Q1 previous lecture

Q2 For last occurrence  $\rightarrow$  1st occurrence val fun. Ko hi use karke use ek bade element ki 1st position  $\xrightarrow{(x)}$  nikala.  
and  $(n-1)$  will be the last occurrence of target element.

Approach 2

$\Rightarrow$  if ( $\text{input}[\text{mid}] > \text{target}$ ) {

$\text{ans} = \text{mid};$

$\text{hi} = \text{mid} - 1;$

else  $\text{lo} = \text{mid} + 1;$

if ( $\text{a}[\text{mid}] \leq \text{target}$ ) {

$\text{lo} = \text{mid} + 1;$

if ( $\text{a}[\text{mid}] == \text{target}$ ) {

$\text{ans} = \text{mid};$

3, 3

3 else {

$\text{hi} = \text{mid} - 1;$

3

Time complexity  
 $O(\log n)$

Space complexity  
 $O(1)$

Void lower\_bound (vector<int> &input, int target) {

Same as prev. Ques.

{

Void Upperbound (vector<int> &

Same {but targets will be changed?  
        } acc. to approach ①

{

int main () {

n → enter

vector → enter

target → enter

vector<int> result;

int lb = lower\_bound (vector, target);

if (input[lb] == target) {

    result.push\_back (-1);

    result.push\_back (-1);

} else {

    result.push\_back (lb);

    result.push\_back (lb - 1);

}

cout << result[0] << " " << result[1] << "\n";

return 0;

{

Q3 A rotated sorted array is a sorted array on which rotation operation has been performed some no. of times. Given a rotated sorted array, find the index of the minimum element in the array. Follow 0-based indexing.

It is guaranteed that all the elements in the array are unique.

- input :-  $\text{Arr}[7] = \{3, 4, 5, 1, 2\}$
- output :- 3

$\rightarrow 2 \ 3 \ 4 \ 5 \ 6 \ 7$

Sols eg:-

$\begin{matrix} 7 & 2 & 3 & 4 & 5 & 6 \\ \downarrow & \leftarrow & \uparrow & & & \nwarrow \\ l_0 & & mid & & & h_i \end{matrix}$

$a[l_0] > a[h_i] \rightarrow \text{sorted Rotated}$

$\begin{matrix} 4 & 5 & 6 & 7 & 2 & 3 \\ \uparrow & \uparrow & \uparrow & \uparrow & \uparrow & \uparrow \\ l_0 & mid & & & h_i & \end{matrix}$

if  $a[mid] > a[l_0]$

$l_0 = mid + 1$ ; //min. element is at right

if  $a[mid] < a[l_0]$

$h_i = mid - 1$ ; //min. element is at left

• terminals  $\rightarrow a[mid] > a[mid + 1]$

ans = mid + 1

$a[mid] < a[mid - 1]$

ans = mid

int findMinimumInSortedRotated( vector<int>&input)

if (input.size() == 1) return input[0];

int l0 = 0, hi = input.size() - 1;

if (input[l0] < input[hi]) { //sorted array.  
return l0;

```

while (lo <= hi) {
    int mid = lo + (hi - lo) / 2
    if (input[mid] > input[mid + 1]) {
        return mid + 1;
    } else if (input[mid] < input[mid - 1]) {
        return mid;
    } else {
        if (input[mid] > input[lo]) {
            lo = mid + 1
        } else {
            hi = mid - 1
        }
    }
}
return -1;

```

int main() {
 n → enter
 vector → enter
 cout << func. x<<sup>4</sup>/n<sup>4</sup>
 return 0;
}

Q3) Given the rotated sorted array of integers, which contains distinct elements, and an integer target, return the index of target if it is in the array. Otherwise return -1.

Input:- Arr[] = {3, 4, 5, 1, 2}, target = 4

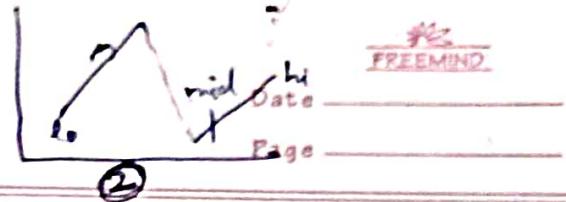
Output:- 1 ) Apply Binary search 3 times

Method 1) (1) To find min. element

(2) From [lo to min.]

(3) From [min. to hi] and get target value

method 2 :-



```
#include <iostream>
```

```
#include <vector>
```

```
using namespace std;
```

```
int binarySearchSortedRotated (vector<int>& input), int target)
```

```
int lo = 0, hi = input.size() - 1;
```

```
while (lo <= hi) {
```

```
    int mid = lo + (hi - lo) / 2;
```

```
    if (input[mid] == target) return mid;
```

```
    if (input[mid] > input[lo]) {
```

```
        if (target > input[lo] && target <= input[mid])
```

```
            if target >= lo && target <= mid
```

```
                hi = mid - 1;
```

```
        } else {
```

```
            lo = mid + 1;
```

```
}
```

```
} else {
```

```
    if (target >= input[mid] && target <= input[hi])
```

```
        lo = mid + 1;
```

```
} else {
```

```
    hi = mid - 1;
```

```
}
```

```
} else {
```

```
    return -1;
```

```
}
```

```
int main () {
```

```
n → enter; vector → enter; target → enter;
```

```
cout << funct << "\n";
```

Time complexity  $\Rightarrow O(\log n)$

Space  $\rightarrow O(1)$

H.W

FREEMIND

Date \_\_\_\_\_

Page \_\_\_\_\_

Q3 Search element in Rotated Sorted Array with duplicate elements. Return 1 if the element is found, else return -1.

input :- [0, 0, 0, 1, 1, 1, 2, 0, 0, 0], target = 2

output:- 1

Logic → Terminal elements ko remove kardeng so that lo or hi me compare kar sakin or binary search apply kar sakin (for Sorted Rotated array)

(if  $lo = hi$ )

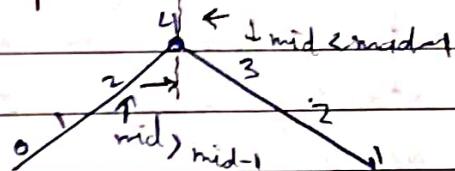
$lo++$ ,  $hi--$  until ( $lo < hi$ )

## Lecture 46

## (Binary Search Problems - 2)

Q. Given a mountain array 'arr' of length greater than 3, return the index 'i' such that  $arr[0] \geq arr[i] \geq \dots \geq arr[i-1] < arr[i] > arr[i+1] \dots > arr[arr.length - 1]$ . This index is known as the peak index.

Input:- arr = [0, 4, 1, 0]  
Output:- 1



3) PeakInMountainArray (vector<int> &input) {

    lo = 0, hi = input.size() - 1;

    int ans = -1;

    while (lo <= hi) {

        int mid = lo + (hi - lo)/2;

        if (input[mid] > input[mid-1]) {

            ans = mid;

            lo = mid + 1;

        } else {

            hi = lo - mid - 1;

    }

    return ans;

}

int main () {

    vector<int> arr;

    cout << PeakInMountainArray (arr);

    return 0;

}

Q3 A peak element is an element in an array that is strictly greater than its neighbors. Given a 0-indexed integer array `nums`, find a peak element and return its index. If the array contains multiple peaks, return the index to any of the peaks. You may imagine that  $\text{nums}[-1] = \text{nums}[n] = -\infty$ , in other words an element is always considered to be strictly greater than a neighbor that is outside the array.

input:- `Arr[] = [1, 2, 1, 2, 6, 10, 3]`

output:- Either index 1 or index 5 are the correct output. At index 1, 2 is the peak element and at index 5, 10 is the peak element.

Sol. :- int findPeak (vector<int> &input) {

    int n = input.size() - 1;

    int lo = 0, hi = input.size() - 1;

    while (lo <= hi) {

        int mid = lo + (hi - lo) / 2;

        if (~~input[mid] == 0~~) {

            if (input[mid] > input[mid + 1]) {

                return 0;

            } else {

                return 1;

        }

        } else if (~~input[mid] == n - 1~~) {

            if (input[mid] > input[mid - 1]) {

                return n - 1;

            } else {

                return n - 2;

        }

```
else {
```

```
    if (input[mid] > input[mid - 1] && input[mid] > input[mid + 1])
        return mid;
```

```
    } else if (input[mid] > input[mid + 1]) {
        lo = mid + 1
```

```
    } else {
```

```
        hi = mid - 1
```

```
}
```

```
}
```

```
}
```

```
}
```

```
}
```

```
return -1;
```

```
}
```

```
int main() {
```

n → enter

vector → enter

```
cout << find_peak(vector) << "\n";
```

```
int a;
```

```
}
```

```
}
```

Q3 Search the 'target' value in a 2d integer matrix of dimensions  $n \times m$  and return 1 if found, else return 0. This matrix has the following properties:

(i) integers in each row are sorted from left to right

(ii) 1st integer of each row is greater than the last integer of previous row.

input :-

Matrix = {[1, 3, 5, 7], [10, 11, 16, 20], [23, 30, 34, 60]},

target = 3

output :- 1

1	3	5	7
10	11	16	20
23	30	34	60

Solu ① metho 1  $\rightarrow$

$i = 0 \rightarrow n$

loop  $i = 0, 1, 2, \dots, n-1$ ;  $j = 0 \rightarrow m$

linear search in 1D Array

② Method 2  $\rightarrow$

$i = 0 \rightarrow n$

loop  $i = 0, 1, 2, \dots, n-1$ ;  $j = 0 \rightarrow m$

Binary search in 1D Array

③ Method 3 copy 2D array in single 1D Array

$\{1, 3, 5, 7, 10, 11, 13, 16, 20, 23, 30, 34, 60\}$

Apply Binary Search

④ Method 4  $\rightarrow$

1	3	5	7
10	11	13	16
23	30	34	60

$3 \times 4$

$mid = 5$

$(x, y) = (1, 1) = 11$

$n = mid / m = 5 / 4 = 1$

$y = mid \% m = 5 \% 4 = 1$

#Code

Void SearchMatrix (vector<int> &<sup>a</sup>, int target)

$n = a[0].size() - 1$

$m = a[0].size() - 1$

int lo = 0, hi = ~~0~~ \* m - 1;

while (lo < hi) {

~~int~~ int mid = lo + (hi - lo) / 2

int x = mid / m

int y = mid % m

```

if (a[m][y] == target) {
    return true;
} else if (a[m][y] < target) {
    ls = mid + 1;
} else {
    hs = mid - 1;
}
return false;
}

```

```
int main() {
```

~~vector → input (A, f, &)~~  
~~target → 3~~  
~~cout << much << n;~~

```
}
```

### Lecture 47

### (Binary Search Problems - 3)

Q3 You have ' $n$ ' ( $n \leq 10^5$ ) boxes of chocolate. Each box contains a  $a[i]$  ( $a[i] \leq 10000$ ) chocolates. You need to distribute these chocolates among ' $m$ ' students such that the maximum no. of chocolates allocated to a student is minimum.

- (a) One box will be allocated to exactly one student.
  - (b) All the boxes should be allocated.
  - (c) Each student has to be allocated at least one box.
  - (d) Allotment should be in contiguous order, for instance a student cannot be allocated box 1 and box 3, skipping box 2.
- calculate and return that minimum possible no.

Assume that it is always possible to distribute the chocolates

→ The first line of input will contain the <sup>value</sup> ~~line~~ of  $n$ ,

the no. of boxes.

- The second line of input will contain the n no.s denoting the no. of chocolates in each box.
- The third line will contain the m no. of students.

- input :-

4

12 34 67 90

2

Time complexity :-  $O(n \times \log(\sum a_i))$

113

- Output :-

Sol is

```
#include <iostream>
```

```
#include <vector>
```

```
Using namespace std;
```

```
bool canDistChoco(vector<int> arr, int mid, int s) {
```

```
    int n = arr.size();
```

```
    int studentRequired = 1;
```

```
    int currSum = 0;
```

```
    for (int i = 0; i < n, i++) {
```

```
        if (arr[i] > mid) {
```

```
            return false;
```

```
}
```

```
        if ((currSum + arr[i]) > mid) {
```

```
            studentRequired++;
```

```
            currSum = arr[i];
```

```
        } else {
```

```
            currSum += arr[i];
```

```
}
```

```
} return true;
```

```
}
```

```
int distChoco (Vector<int>&arr, int s) {
```

```
    int n = arr.size();
```

```
    int lo = arr[0];
```

```
    int hi = 0;
```

```
    for (i=0; i<n; i++) {
```

```
        hi += arr[i];
```

```
}
```

```
    int ans = -1;
```

```
    while (lo <= hi) { int mid = lo + (hi-lo)/2;
```

```
        if (condistChoco (arr, mid, s)) {
```

```
            ans = mid;
```

```
            hi = mid-1;
```

```
} else {
```

```
    lo = mid+1;
```

```
}
```

```
} return ans;
```

```
}
```

```
int main () {
```

```
    int n;
```

```
    Vector<int> arr;
```

```
    cin >> n >> arr;
```

```
    cout << distChoco (arr, n);
```

```
}
```

Q24 A new racing track for kids is being built in New York with 'n' starting spots. The spots are located along a straight line at positions  $x_1, x_2, \dots, x_n$  ( $x_i \in \mathbb{R}$ ). For each  $i$ ,  $n_{i+1} > n_i$ . At a given time only 'm' children are allowed to enter the race. Since the race track is for kids, they may run into each other while running. To prevent this, we want to choose the starting spots such that the minimum distance b/w any two of them is as large as possible. What

is the largest minimum distance?

→ The first line of input will contain the value of  $n$ , the no. of starting spots. The second line of input will contain the  $n$  nos. denoting the location of each spot. The third line will contain the value of  $m$ , no. of children.

• input

5

1 2 4 8 9

3

output

3

Time complexity =  $O(n \log(x_n - x_1))$

Sol → bool canPlaceStudents(vector<int>& pos, int s, int mid) {

    int studentReqd = 1;

    int lastPlaced = pos[0];

    for (i=1; i<pos.size(); i++) {

        if (pos[i] - lastPlaced >= mid) {

            studentReqd++;

            lastPlaced = pos[i];

        if (studentReqd == s) {

            return true;

}

    return false;

    int race (vector<int>& pos; int s) {

        int n = pos.size();

        int lo = 1;

        int hi = pos[n-1] - pos[0];

int ans = -1;  
while (lo ≤ hi) {

int mid = lo + (hi - lo) / 2;

if (canPlaceStudents(pas, s, mid)) {

ans = mid;

lo = mid + 1;

} else {

hi = mid - 1;

}

return ans;

}

int main() {

n → enter

vector <→ enter

s → enter      sta (No. of Students)

cout << result(pas, s);

}

(H-W) 3 Variations of these two problems

(ii) Aggressive Cows

[From Q2]

(i) Book allocation

[From Q1]