

Class Quene & Node * head; Node * tail; Quent & in style; public! Que () & this's head = NULL; this -> tail = NULL; a this > sige = 0; 3 Void enquere (in data) Void enque (int data) {
Nool * new Mode = new Nool * (data); if (this > head = = N ULL) {

this > head = this > fail = new Nod; 3 else q this->tail > next = new Node; this stail = new Moole; Void dequene () { if (Thead == NULL) { ____ sut wn; 8 elle & Node * oldhead = this -> head; Node * newhead = this > head - next: this - sheard = necestrend; of (this-shead == NULL) this bail == NULL) oldhorad - next = NULL delete oldhead; I this sing -

get int getsize()? return this > size; bool isemfly () { return this > head == NWLL; in mainty & was all well int front () & front () &

if (this > head == NULL) notwom -1;

neturn this > head > data; int main () § Brene qui g. n. engres (10). qu. cnquene(20); //10 20 qu- dogdue 1); //2030 gy. enquer (40); // 20 3040 while (not quis Empty ()) { gu. degnemb;

ren & had steat

hast

Con my

int get front () {

if (this > front = -1) return -1;

return this > V [this > front]; bool is 6 mpty 1) 50 return this front ==-1; int main () § while (not que emplys) {

conderque get front () < <" ";

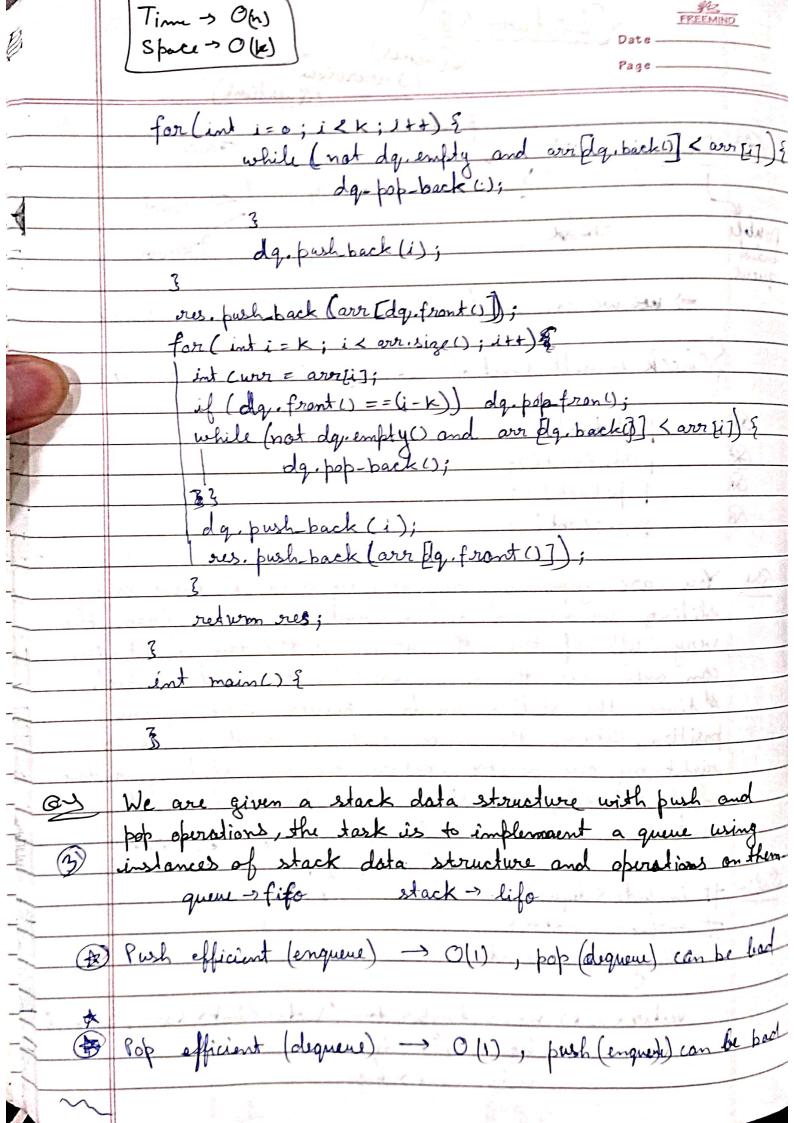
que dequeue ();

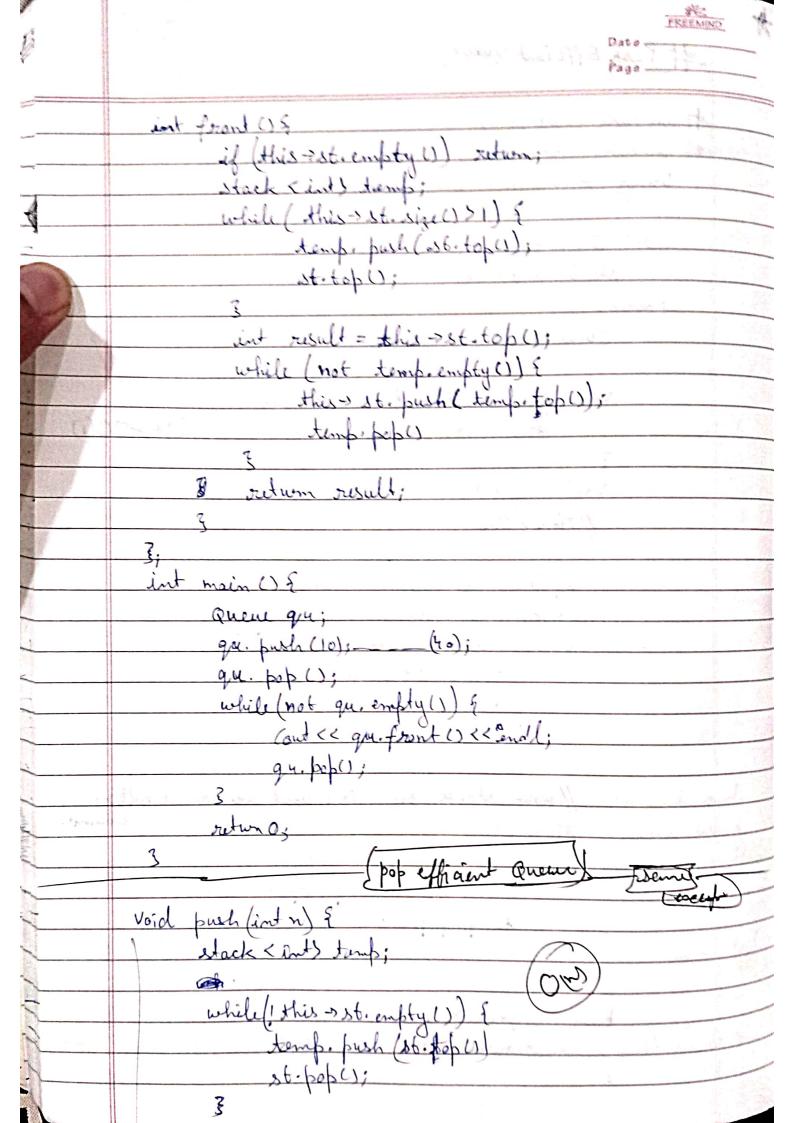
} BSimilar to stacks there is an infunit library far Queres. So, we can use it instead of making it by arrays/linked list. by arrays/linked list. # include (queue) ving namespace std; int main () { queue Lind's qu; gy. push (10); qu. push (20); gr. pop U;

while (not garenfly) & cont << que front es << "";

B	Advantage of array implementation of queues
7	ll -> space of efficiently [Good]
	ll -> (pointer) -> deletion
	June > SEL
	The state of the state of
By	Reverse the elements of a queue.
	1 /- 1 10 - 10 - 10 - 10 - 10 - 10 - 10
	1 Move elements from queue to stack 2 Move elements back from stack to green
	2 Move elements back from stack to grund
	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
Ans	# dinclude < quene>
	Using namespace eld:
	int main () {
	queue < ind > inbut;
	intput bull (10);
	input. push (20); (50); (40):
\$ 5	while (not input. empty (1) S
	st. push (input. Frant());
	input. pep();
-	3
	while (not st. empty () { input. push (st. top1);
	a input push st. tohil:
	st. pop();
`	3
₹	while (not input, empty cs) {
3	Cont ex input. front 1) see 40
N	input. pop()
	inpul. pop (); 3 returno;

Lecture-58 (Queur)
(3 nterview
(Questions) => we we can add/remove elements from both sides. Scheck inbuilt funt of different libraries at oppreformely bushback (): push front (); pop back (); pop-front (); By You are given an array of integers nums, there is a sliding window of size k which is moving from the very left of the the array to the very right. You can only see the k numbers in the window. Each I time the sliding window moves right by one position, Return the mase sliding window moves which parically contains the max element in each [1,2,1,-3,5,8,9,6,1] Salis Hinclude (deque) # include < vector) using namespace std; Vector < int > max Sliding Window (Vector < int > Larr, int k) & dequeux cint > dq; vector < int > res;





this set push (a);

while (not temp. empty ()) }

this -> st. push (temp. top())
temp. pop;

Void pop () {
if of this -set empty) return;

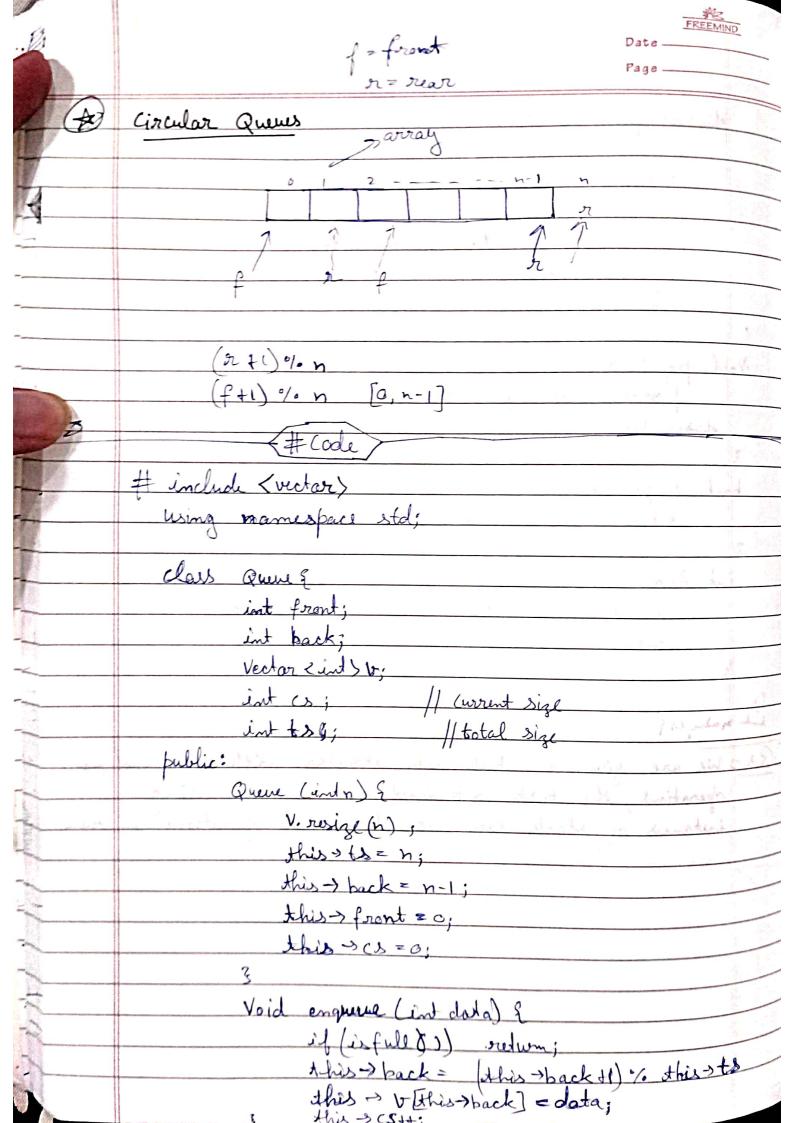
bool emply () ?

return this > st. empty ();

this > st. empty()) reducem, INT-MIN; um this > st. top();

int maky US

By s we are given a stack data structure with push and pop operations, the task is to implement a queue using instances of stack data structure and operations on them



Void dequeue () {
if (is Empty ()) return.
if (is 5 infrty ()) return; this -> front = (this > front +1) % this -> ts;
this > (s;
3
int get front () { if (this > front = = -1) return -1; return this > v [this > front];
if (this > front = = -1) return -1:
return this -> v This -> front ?:
 3
 This has a brown the
 bool isempty () {
 return this > (s == 0;
 3
bool is Full () {
 return this -> (s = 2 this -> 68;
3
And the second of the second o