

Lecture 50

(Linked List in C++ (Part 1))

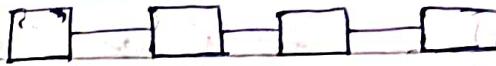
→ What is a linked list? → collection of Nodes [Info | link]
 linear data structure used to store a list of values

Array



a single memory
Block with partitions

linked list



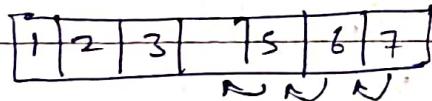
memory blocks linked
to each other

→ Challenges of array

→ static size

→ Contiguous memory
allocation

→ Inserting and Deleting is
costly ($O(n)$)



100 bytes

20 bytes

80 bytes

But No ~~element~~ arry
can be inserted

→ Advantages of a linked list over an array

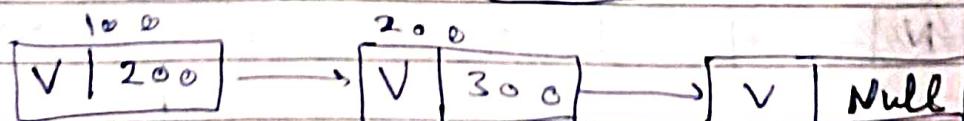
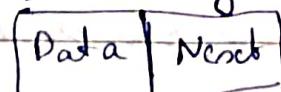
→ Dynamic Size

→ Non contiguous memory allocation

→ Insertion and Deletion is not expensive.

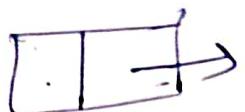
→ List node

blocks of memory → node



↑
Head

↑ Tail



→ Types of linked lists

→ Singly linked list

every node points to its successor node



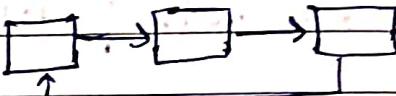
→ Doubly linked list

every node is connected to its previous & next node.



→ Circular linked list

the last node point to head node.



★ Implementation of a listnode in a singly linked list

For C

```
struct node {
    int value;
    struct node* next;
};
```

```
class Node { public:
    int value;
    Node* next;
```

```
Node (int data) :
```

```
    value = data;
```

```
    next = null;
```

```
int main () {
```

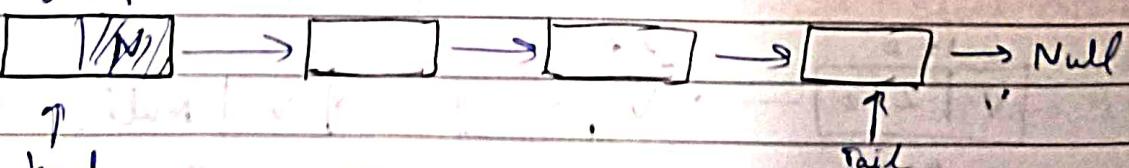
```
    Node* n = new Node(1)
```

```
    cout << n->value << " " << n->next << endl;
```

```
    return 0;
```

→ Traversal in a Singly linked list

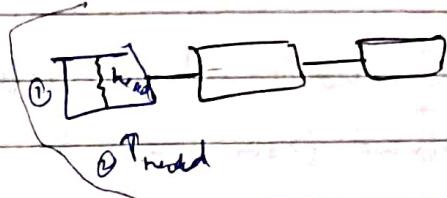
Create a
new pointer
temp to traverse
the linked
list



```
temp = temp->next
```

⇒ Insertion at k^{th} position in a Singly linked list.

• Add a node at the start



// Continued to previous code

(1) `Void InsertAtHead (Node* &head, int val) {`

`Node* new-node = new Node(val);`

`new-node->next = head;`

`head = new-node;`

`Void display (Node* head)`

`Node* temp = head;`

`while (temp != Null) {`

`cout << temp->val << " -> ";`

`temp = temp->next;`

`} cout << " Null " << endl;`

`}`

`int main () {`

`Node* head = Null;`

`InsertAtHead (head, 2)`

`display (head)`

`InsertAtHead (head, 1)`

`display (head)`

`return 0;`

Output

`2 -> Null`

`1 -> 2 -> Null`

>Add node to end

(2) `Void insertAtTail (Node*&head, int Val) {`

`Node* new-node = new Node (Val);`

`Node* temp = head;`

`while (temp->next != Null) {`

`temp = temp->next;`

`}`

`temp->next = new-node;`

`int main () {`

`insertAtTail (head, 3);`

`display (head);`

`return 0;`

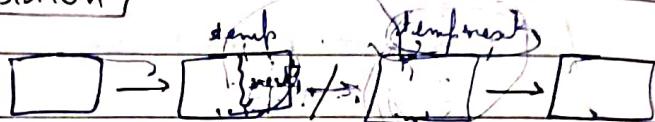
`1 -> 2 -> 3 -> Null`

(A) Time Complexity:- $O(n)$

$O(1)$ {when tail is given}

Add a node at an arbitrary position

// Continued from previous



③ void InsertAtPosition (Node* &head, int val, int pos) {

if (pos == 0) {

 InsertAtHead (head, value)

 return;

}

 Node* new-node = new Node (val);

~~temp~~

 Node* temp = head;

 int current-position = 0;

 while (current-position != pos-1) {

 temp = temp->next;

 current-position++;

} // temp is pointing to node at pos-1

// new-node->next = temp->next;

 temp->next = new-node;

}

int main () {

 InsertAtPosition

 (head, 4, 1);

 display (head)

(B) Time complexity $\Rightarrow O(pos)$

$O(n)$ (Worst case)

→ Updation at k^{th} position in a Singly linked list

(continued)

Void UpdateAtPos (Node* &head, int k, int val) {

Node* temp = head;

int currpos = 0;

while (currpos != k) {

temp → temp → next;

(currpos++)

// temp is pointing to node at k

temp → value = value;

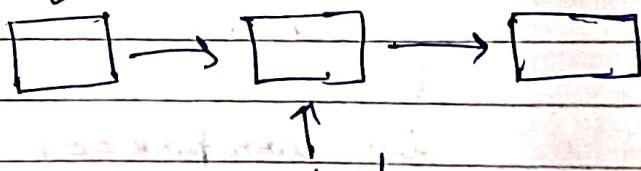
int main () {

UpdateAtPosition (head, 2, 5)

display (head)

→ Deletion at k^{th} Position in a Singly linked List.

→ Delete a node at the start



Void DeleteAtHead (Node* &head) {

Node* temp = head;

head = head → next;

free (temp);

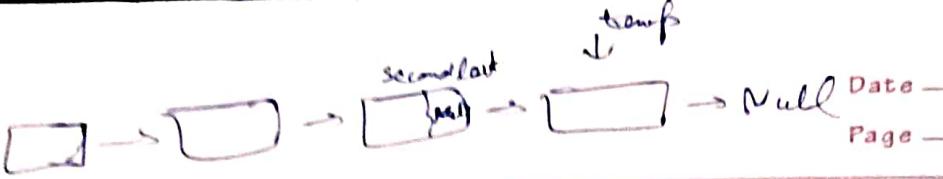
Time complexity $\rightarrow O(1)$

3

int main () {

DeleteAtHead (head)

Display (head)



Void DeleteAtTail (Node* &head) {

Node* secondLast = head;

while (secondLast->Next->Next != Null) {

SecondLast = secondLast->Next;

}

// Now SecondLast points to second last node

Node* temp = secondLast->next; // Node to be deleted

secondLast->next = Null;

free (temp);

}

int main () {

DeleteAtTail (head)

Display (head)

}

// Deletion at Arbitrary Position.

Void DeleteAtPosition (Node* &head, int pos) {

if (pos == 0) {

DeleteAtHead (head);

return;

}

int currPos = 0;

Node* prev = head;

while (currPos != pos - 1)

prev = prev->next;

currPos++

// prev is pointing to node at pos-1

Node* temp = prev->next;

prev->next = prev->next->next;

free (temp);

O(1)

Complexity

Time

Time complexity

O(n)

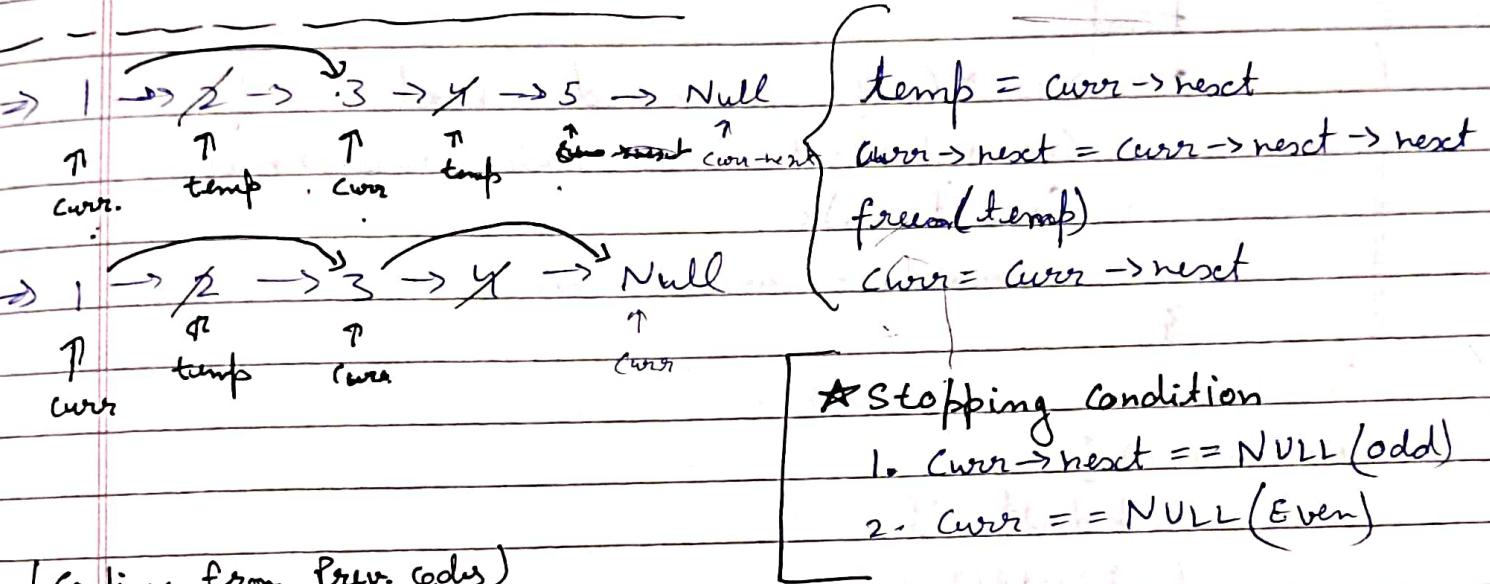
last node

Deleted Pos

(head, 1)

Display (head)

Q3 Given the head of a linked list, delete every alternate element from the list starting from the second element.



(Continue from Prev. code)

Soln: class Node { }

class linkedlist { }

public:

Node* head;

linkedlist () {

head = NULL

Void insertAtTail (int Val.)

Node* new-node = new Node (Val);

if (head == NULL) { //linked list is empty

head = new-node;

return;

}

Node* temp = head

while (temp->next != NULL) {

temp = temp->next;

}

temp->next = new-node;

}

Void display { Some func }

```

Void DeleteAlternateNodes(Node* & head) {
    Node* currNode = head;
    while (currNode != NULL && (currNode->next) != NULL) {
        Node* temp = currNode->next;
        currNode->next = currNode->next->next;
        free(temp);
        currNode = currNode->next;
    }
}

```

{

int main() {

linkedlist ll;

ll.InsertAtTail(1);

ll.InsertAtTail(2);

ll.InsertAtTail(3);

ll.InsertAtTail(4);

ll.InsertAtTail(5);

ll.display();

deleteAlternateNodes(ll.head);

ll.display();

return 0;

{

Q-21 Given the head of a sorted linked list, delete all duplicates such that each element appears only once. Return the linked list sorted as well.

while (curr != NULL) {

while (curr->next != NULL && curr->value == curr->next->value) {

 // remove curr->next Node as code of Delete
 // Alternate nodes

{ curr = curr->next }

{

$$\begin{matrix} 1 & \rightarrow & 2 & \rightarrow & 2 & \rightarrow & 3 & \rightarrow & 3 & \rightarrow & 3 \\ & & & & & & & & & & \end{matrix}$$

$$1 \rightarrow 2 \rightarrow 3$$

Continued from Prev. Ques.

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Void DeleteDuplicateNodes (Node* & head)

Node* curr-node = head;

while (curr-node) {

 while ((curr-node->next) && (curr-node->value == curr-node->next->value))

 Node* temp = curr-node->next;

 curr-node->next = curr-node->next->next;

 free (temp);

// this loop ends when current node and
// next node values are different or
// linked list ends

 curr-node = (curr-node->next);

int main () {

 ll. - -

 ll. - -

 ll. DeleteDuplicateNodes (ll. head);

 ll. display();

 return 0;

Time complexity

$O(n)$

Q3) Given the head of a singly linked list, and print the reversed list.

Sol.) Void reverseprint (Node* head) {

 if (head == NULL) { // base case

 return;

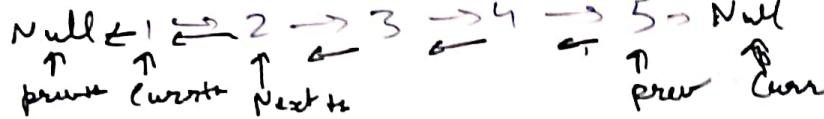
}

 reverseprint (head->next); // recursive ~~final~~ case

 cout << head->value << " "

int main () {

 Solve



(*) Reversing a linked list:-

Q3 Given the head of a singly linked list, reverse the list, and return the reversed list.

Soln Node* reverseLL(Node* &head) {

 Node* prevptr = NULL;

 Node* currptr = Head;

 while (currptr != NULL) {

 Node* Nextptr = currptr->next;

 currptr->next = prevptr;

 prevptr = currptr;

 currptr = Nextptr;

}

//~~Note~~ When this loop ends, prevptr pointing to my last node which is new head.

 Node* newhead = prevptr;

 return newhead;

}

int main() {

 ll.head = reverseLL(ll.head);

 ll.display();

 return 0;

}

Recursively

head → next → next

= head

head → next = NULL

1 → (2 → 3 → 4 → 5)

↑
head

NULL ← 1 ← [2 ← 3 ← 4 ← 5]

↑
head

↑
head → next

↑
new-head

?

Node* reverseLLRecursion (Node* &head) {

if (head == NULL || head->next == NULL) {
 return head;

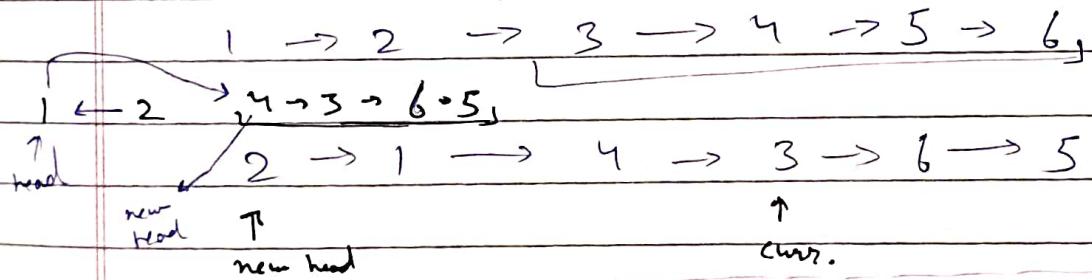
3

 Node* newHead = reverseLLRecursion (head->next);
 head->next->next = head;
 head->next = NULL;
 return newHead;

3

Q Given the head of a linked list, reverse the nodes of the list k at a time, and return the modified list.

K = 2



sol → Node* reversekLL (Node* &head, int k) {

 Node* prevptr = NULL;

 Node* currptr = head;

 int counter = 0; // Counting list k nodes

 while ((currptr != NULL && counter < k) { // reversing

 Node* nextptr = currptr->next k nodes

 currptr->next = prevptr

 prevptr = currptr;

 currptr = nextptr;

 counter++;

3

if

// currptr will give us $(k+1)^{\text{th}}$ node

if ($\text{currptr} \neq \text{NULL}$) {

 Node * newhead = reversekLL (currptr, k); // recursive call
 head \rightarrow next = newhead;

}

return currptr; // will give the new head
of connected linkedlist

int main () {

 ll.head = reversekLL (ll.head, 2);

 ll.display();

}

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8(Linked List (Part -2))

⊕ Pattern 1 > Pointers

Q-3 Given 2 linked lists, Tell if they are equal or not.
Two linked lists are equal if they have same data and arrangement of data is also the same.

Sol. bool checkEqualLinkedList(Node* head1, Node* head2)

Node* ptr1 = head1;

Node* ptr2 = head2;

while(ptr1 != NULL && ptr2 != NULL) {

if (ptr1->val != ptr2->val) {

return false;

ptr1 = ptr1->next;

ptr2 = ptr2->next;

// at this point either ptr1 is null, or

ptr2 is null, or both are null.

return (ptr1 == NULL && ptr2 == NULL);

int main() {

LinkedList ll1;

ll1.insertAtTail(1)

;

(2)

;

(3)

LinkedList ll2;

ll2.insertAtTail(1)

;

(2)

;

(3)

ll1.display(); ll2.display();

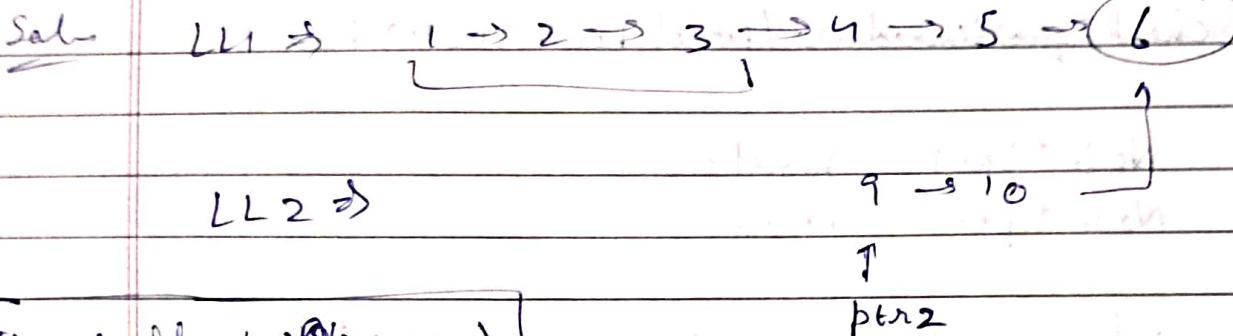
cout << checkEqualLinkedList(ll1.head(), ll2.head());

return 0;

<< endl;

Q → Given the head of two singly linked-lists head A and head B, return the node at which the two lists intersect. If the two linked lists have no intersection at all, return null.

ptr1
↓



Time Complexity = $O(n+m)$

Node * getIntersectionNode(Node * head1, Node * head2) {

Node * ptr1 = head1;
Node * ptr2 = head2;

int getLength (Node * head) {

Node * temp = head;

int length = 0;

while (temp != NULL)

length++;

temp = temp->next;

}

return length;

}

Void moveHeadByk (Node * head, int k) {

Node * ptr = head

while (k--) {

ptr = ptr->next;

return ptr;

}

~~Node^{*} void getIntersection(Node^{*} head1, ~~Node^{*}~~ head2) {~~

//Step1: calculate lengths of both linked lists.

int L1 = getLength (head1)

int L2 = getLength (head2)

1) Find diff. k b/w linked-list and move longer linked list ptr by k

if ($L_1 > L_2$) {

$$\text{int } k = L_1 - L_2 ;$$

ptr = move head by k (head, k);

ptr2 = head2;

3 else 5

$$\text{int } k = L_2 - L_1;$$

ptr1 = head1

ptr2 = move head by k (head2, k);

Compare ptr1 and ptr2 nodes.

while (*ptr*) {

if ($p[tr1] == p[tr2]$) {

return p[rl];

3

ptr1 = ptr1->next;

~~ptr2 = ptr2 -> next;~~

~~return NULL;~~

7

int main() {

linkedlist ll1

III. insert All tail (1) ; (2) ; (3) ; (4) ; (5)

linker lis ll2

ll2. insertAtTail (6); (7)

ll2. head \rightarrow next \rightarrow next = ll1. head \rightarrow next \rightarrow next \rightarrow next;

```

Node* intersection = getIntersection(LL1.head, LL2.head);
if (intersection) {
    cout << intersection->val << endl;
}
else {
    cout << "-1" << endl;
}
return 0;
}

```

Q.5 Given the head of a linked list, remove the k-th node from the end of the list and return its head.

Soln void removeKthNodeFromEnd(Node* head, int k)

```

Node* ptr1 = head
Node* ptr2 = head
//move ptr2 by k steps ahead
int count = k
while (count--) {
    ptr2 = ptr2->next;
}
if (ptr2 == NULL) { // k is equal to length of linked list.
    // we have to delete head node
    Node* temp = head;
    head = head->next;
    free(temp);
    return;
}
// now ptr2 is k steps ahead of ptr1
// when ptr2 is at Null (end of list), ptr1 will be at
// node to be deleted

```

while ($\text{ptr2} \rightarrow \text{next} \neq \text{NULL}$) {

$\text{ptr1} = \text{ptr1} \rightarrow \text{next};$

$\text{ptr2} = \text{ptr2} \rightarrow \text{next};$

}

~~Node~~ // Now ptr1 is pointing to the node before k^{th} node

// Node to be deleted is $\text{ptr1} \rightarrow \text{next}$ from end

$\text{Node}^* \text{temp} = \text{ptr1} \rightarrow \text{next};$

$\text{ptr1} \rightarrow \text{next} = \text{ptr1} \rightarrow \text{next} \rightarrow \text{next};$

$\text{free}(\text{temp});$

3

int main() {

ll. -

ll. -

~~removeKthNodeFromEnd~~(ll, head, 3);

ll.display();

return 0;

}

Q3 Given 2 sorted linked lists, merge them into 1 singly linked list such that the resulting list is also sorted.

Soln → $\text{Node}^* \text{mergeTwoSortedList}(\text{Node}^* \& \text{head1}, \text{Node}^* \& \text{head2}) \{$

$\text{Node}^* \text{DummyHeadNode} = \text{new Node}(-1);$

$\text{Node}^* \text{ptr1} = \text{head1}$

$\text{Node}^* \text{ptr2} = \text{head2}$

$\text{Node}^* \text{ptr3} = \text{DummyHeadNode}$

while ($\text{ptr1} \& \& \text{ptr2}$) {

if ($\text{ptr1} \rightarrow \text{val} < \text{ptr2} \rightarrow \text{val}$) {

$\text{ptr3} \rightarrow \text{next} = \text{ptr1}$

~~else~~ $\text{ptr1} = \text{ptr1} \rightarrow \text{next}$

{ Time complexity $\rightarrow O(n+m)$

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else {

ptr3->next = ptr2

ptr2 = ptr2->next

}

3

if (ptr1 != Null) {

ptr3->next = ptr1

else {

ptr3->next = ptr2

}

return DummyHeadNode->next;

3

int main() {

ll1* ll2* ~

linkedlist ll3.

ll3.head = ~~mergeliinkedlists~~(ll1.head, ll2.head);

ll3.display();

return 0;

3;

Pattern Merging Multiple linked list

Q3 You are given an array of k linked-lists, each linked list is sorted in ascending order. Merge all the linked-lists into one sorted linked list and return it.

Expt \Rightarrow 2 Phle 2 lists ko merge kro and then uske jo list banegi usko 3rd ke saath merge karo

\Rightarrow Merge 2 lists code from prev. Question

Node* MergekLinkedLists (vector<Node*> &lists) {

if (lists.size() == 0) {

return NULL;

}

while (lists.size() > 1) {

Node* mergehead = mergeLinkedLists (lists[0], lists[1]);

lists.push_back (mergedhead);

lists.erase (lists.begin());

lists.erase (lists.begin());

}

return lists[0];

}

int main () {

LinkedList ll1;

LinkedList ll2;

LinkedList ll3;

vector<Node*> lists

lists.push_back (ll1.head);

lists.push_back (ll2.head);

lists.push_back (ll3.head);

LinkedList mergedLL;

mergedLL.head = mergekLinkedLists (lists);

mergedLL.display();

return 0;

}

Pattern: Slow Fast Pointer

Q1 Find the middle element of the given linked list.

slow \Rightarrow p₁ = p₂ \rightarrow next

fast \Rightarrow p₁ = p₂ \rightarrow next \rightarrow next

Sol:

```
Node* findMiddleElement(Node* head) {
    Node* slow = head;
    Node* fast = head;
```

```
while (fast != NULL && fast->next != NULL)
    slow = slow->next;
    fast = fast->next->next;
```

}

return slow;

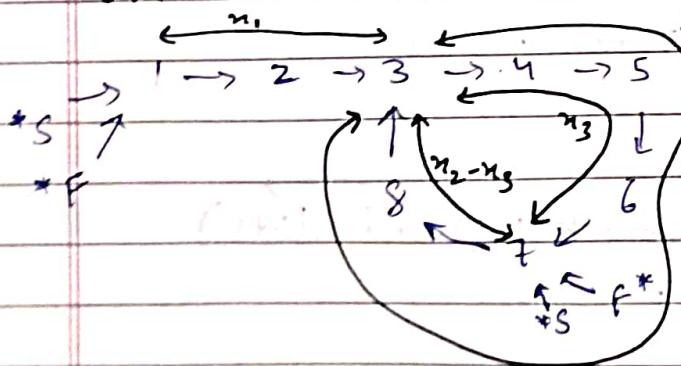
{

int

```
main() {
    Node* middleNode = findMiddleNode(ll.head);
    cout << middleNode->val << endl;
    return 0;
}
```

}

Q2 Given the head of a linked list, determine if the linked list has a cycle in it.



} if slow/fast pointer meets at a same node then linkedlist has a cycle in it.
 $s = s \rightarrow \text{next}$
 $f = f \rightarrow \text{next} \rightarrow \text{next}$

By drawing means, cycle ka starting point head se aur pointers ke meeting point se equidistant hain.

$n_1 = n_2 = n_3$

front root node

bool detectCycle (Node* head)

```
if (!head) {
    return false;
}
```

Node* slow = head;

Node* fast = head;

```
while (fast && fast->next) {
    slow = slow->next;
    fast = fast->next->next;
```

slow = slow->next;

Fast = Fast->next->next;

```
if (slow == fast) {
    cout << slow->val << endl;
    return true;
}
```

}

return false;

}

int main () {

linkedlist ll;

ll.insertAtTail (1); -- (8);

ll.display();

ll.head->next->next->next->next->next->next->next->next
= ll.head->next->next.

cout << detectCycle (ll.head) << endl;

return 0;

}



Jab bhi hamne Question diya, cycle remove
karnے k. toh hum Slow/Fast me se ek ko head
par le jayenge and then dono ko ek ek step se
aage badhayenge. Tis point par dono k. next me
Same value hogi wha Slow ke next me null point
kar denge.

Void removecycle(Node* head) {

Node* slow = head;

Node* Fast = head;

do {

slow = slow → next;

fast = fast → next → next;

} while (slow != fast);

fast = head

while (slow → next != fast → next) {

slow = slow → next;

fast = fast → next;

}

slow → next = NULL;

}

int main() {

// Continued from prev.

removecycle(ll, head);

cout << "After removing cycle: " << endl;

cout << detectcycle(g_ll, head) << endl;

ll.display();

return 0;

}

Q5 Given the head of a linked list, determine if the ll is palindrome or not.

- steps ↗
- 1) Find Middle element → slow / fast pointer
 - 2) Break linked list into 2.
 - 3) Reverse 2nd half of ll. → (prev, curr, next)
 - 4) Compare the 2 parts of ll. (head1 { } head2)
 (start) { } (end)
 { compare }

bool isPalindrome (Node* head) {

① // Find Middle element

Node* slow = head;

Node* fast = head;

while (Fast && Fast->next) {

slow = slow->next;

fast = Fast->next->next;

} // now slow is pointing to middle element

② // break the linked list in the middle.

Node* curr = slow->next; Node* prev = slow;

slow->next = NULL;

③ // reverse the 2nd half of ll.

while (curr) {

Node* nextNode = curr->next;

curr->next = prev;

prev = curr

curr = nextNode

} // full list is reversed

④ // check if the 2 ll. are equal.

Node* head1 = head;

Node* head2 = prev;

while (head2) {

if (head1->val != head2->val) {

return false;

} // end of loop

head1 = head1->next;

head2 = head2->next;

} // end of loop

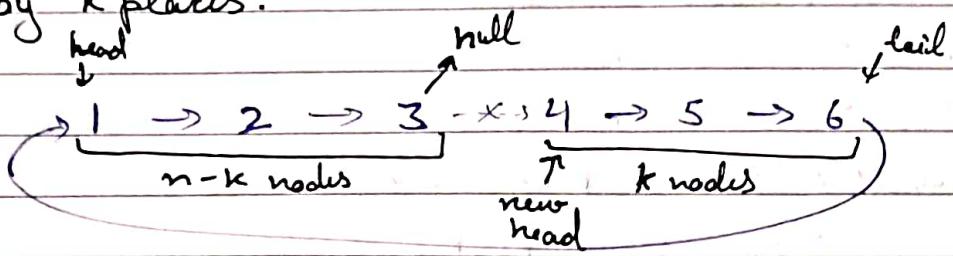
return true;

} // don't mind if list is odd

cout << isPalindrome (ll, head) << endl;

② Pattern: → Rearrangement of Nodes in a list.

Q3 Given the head of a ll., rotate the list to the right by k places.



steps:

① Find n

② Find tail node,

tail \rightarrow next = head

③ Traverse $n-k$ nodes, $(n-k)^{th}$ node \rightarrow next = NULL
 $(n-k+1)^{th}$ = new head

Node * rotateByK (Node *head, int k) {

//① Find length of ll.

int n=0;

//② Find tail node

Node * tail = head

while (tail \rightarrow next) {

n++;

tail = tail \rightarrow next;

}

n++ ; // for including last node

K = k % n;

if (k == 0) {

return head;

}

tail->next = head

Node * temp = head; // traverse $n-k$ nodes

for (i=0; i < n-k; i++) {

temp = temp \rightarrow next;

③

// temp is now pointing to $(n-k)^{th}$ node

$\text{Node}^* \text{newhead} = \text{temp} \rightarrow \text{next};$

$\text{temp} \rightarrow \text{next} = \text{NULL};$

return newhead

3

int main() {

linkedlist ll;

ll.insertAtTail(1); —————— (6),

ll.display();

ll.head = rotateByK(ll.head, 3)

ll.display();

return 0;

3

Q Given the head of a Singly linked list, group all the nodes with odd indices together followed by nodes with even indices, and return the reordered list.

Sol \hookrightarrow Node* oddEvenLinkedList (Node* head) {

if (!head) {

return head;

3

~~Node* evenhead = head->next;~~

~~Node* oddptr = head;~~

~~Node* evenptr = evenhead;~~

while (evenptr && evenptr->next) {

~~oddptr->next = oddptr->next->next;~~

~~evenptr->next = evenptr->next->next;~~

~~oddptr = oddptr->next;~~

~~evenptr = evenptr->next;~~

3

// temp is now pointing to $(n-k)^{th}$ node

$\text{Node}^* \text{newhead} = \text{temp} \rightarrow \text{next};$
 $\text{temp} \rightarrow \text{next} = \text{NULL};$
 return newhead

3

```
int main() {
    linkedlist ll;
    ll.insertATTail(1); - - - - - (6),
    ll.display();
    ll.head = rotateByK(ll.head, 3)
    ll.display();
    return 0;
}
```

Q2

Given the head of a Singly linked list, group all the nodes with odd indices together followed by nodes with even indices, and return the reordered list.

Sol: Node* oddEvenLinkedList (Node* head) {

if (!head) {
 return head;

3

~~Node* evenhead = head->next;~~
~~Node* oddptr = head;~~
~~Node* evenptr = evenhead;~~
~~while (evenptr && evenptr->next) {~~
 ~~oddptr->next = oddptr->next->next;~~
 ~~evenptr->next = evenptr->next->next;~~
 ~~oddptr = oddptr->next;~~
 ~~evenptr = evenptr->next;~~

?

Time Complexity $O(n)$

odd \rightarrow next = evenhead;
otherwise head;

3
int main() {

 ll. head = oddEvenLinkedList(ll. head);

 ll. display();

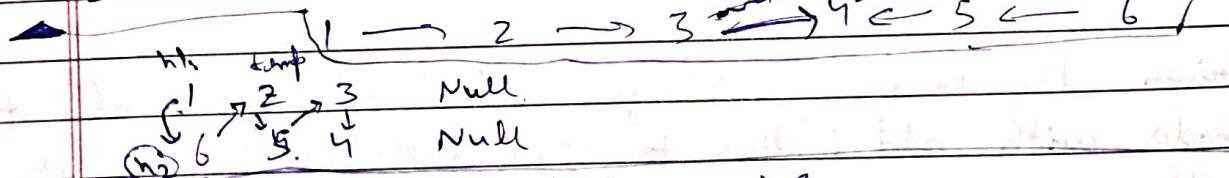
 return 0;

Q3 You are given the head of a singly linked list. The list can be represented as:

② $L_0 \rightarrow L_1 \rightarrow \dots \rightarrow L_{n-1} \rightarrow L_n$

Reorder the list to be in the following form:-

③ $L_0 \rightarrow L_n \rightarrow L_1 \rightarrow L_{n-1} \rightarrow L_2 \rightarrow L_{n-2} \rightarrow \dots$



Sol. \rightarrow Node* reorderedList (Node* &head) {

 // 1. Finding the middle element

 Node* slow = head;

 Node* fast = head;

 while (fast && fast->next) {

 slow = slow->next;

 fast = fast->next->next;

 3

 // now slow is pointing to the middle element

 // 2. separate the linked list and reverse the second half.

 Node* curr = slow->next;

 slow->next = NULL;

 Node* prev = slow

```
while (curr) {
```

```
    Node* nextptr = curr->next;
```

```
    curr->next = prev;
```

```
    prev = curr;
```

```
    curr = nextptr;
```

```
}
```

// 3. merging 2 halves of the linked list

```
.Node* ptr1 = head; // LL 1st half.
```

```
Node* ptr2 = prev; // LL 2nd half.
```

```
while (ptr1 != ptr2) {
```

```
    Node* temp = ptr1->next;
```

```
    ptr1->next = ptr2;
```

```
    ptr1 = ptr2;
```

```
    ptr2 = temp;
```

```
}
```

```
return head;
```

```
}
```

```
int main() {
```

```
    ll.head = reorderedLinkedList(ll.head)
```

```
    ll.display();
```

```
    return 0;
```

```
}
```

Q3 Given a linked list, swap every two adjacent nodes and return its head. You may not modify the values in the list nodes. Only nodes themselves may be changed.

^{head} ^{second Node}

$1 \rightarrow 2 \rightarrow [3 \rightarrow 4 \rightarrow 5 \rightarrow 6]$ Reverses

$2 \rightarrow 1 \rightarrow 4 \rightarrow 3 \rightarrow 6 \rightarrow 5$

new head

Sols

```
Node* swapAdjacent(Node* head) {
```

```
//base case
```

```
if (head == NULL || head->next == NULL) {
```

```
    return head;
```

```
}
```

```
//recursive case
```

```
Node* secondNode = head->next;
```

```
head->next = swapAdjacent(secondNode->next);
```

```
secondNode->next = head; //reversing the link  
//between 1st and 2nd Node.
```

```
return secondNode;
```

```
}
```

```
int main() {
```

```
    ll.head = swapAdjacent(ll.head);
```

```
    ll.display();
```

```
    return 0;
```

```
}
```

Lecture - 52

→ What is a Doubly linked list?



Date _____
struct node {

int info;
struct node *prev;
struct node *next;

} node;

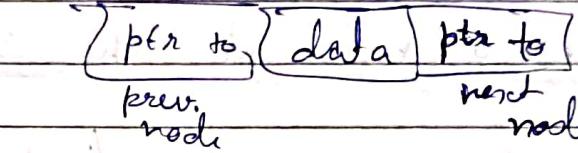
→ Advantages

→ traversal → both ways

→ insertion / deletion becomes more efficient

→ Disadvantage

→ Extra space for previous - ptr.



#include <iostream>

using namespace std;

class Node {

public:

int val;

Node* prev;

Node* next;

Node (int data) {

val = data;

prev = null;

next = null;

};

};

class doublyLinkedList {

public:

Node* head,

Node* tail;

doublyLinkedList () {

head = NULL

tail = NULL

[#Code]

int main () {

Node* new-node = new Node(3)

doublyLinkedList dll;

dll.head = new-node;

dll.tail = new-node;

cout << dll.head->val << endl;

return 0;

}

void display () {

same

};

(*) Traversal in doubly linked list

⇒ Insertion at k^{th} position in a Doubly linked list.

① Add a node at the start

Void insertAtStart (int Val)

Node* new-node = new Node (Val)

if (head == NULL)

head = new-node;

tail = new-node

return;

}

new-node → next = head

head → prev = new-node

head = new-node;

return;

}

int main

dll.insertAtStart();

;

{

 dll.display();

}

↑
head

↑
tail

↑
head

↑
tail

↑
head

↑
tail

② Add a node at end

Void insertAtEnd (int Val) {

Node* new-node = new Node (Val);

if (tail == NULL) {

head = new-node;

tail = new-node;

return;

}

tail → next = new-node;

new-node → prev = tail;

tail = new-node;

return;

}

$O(n)$ $O(k)$

⑥ Add a node at an arbitrary position

Void insertAtPosition (int Val, int k)

// assuming k is less or equal to length of dll

Node * temp = head;

int count = 1;

while (count < (k-1)) {

temp = temp -> next;

count ++;

}

int main () {

dll.insert -> pos1 (3);

dll.display ();

return 0;

}

Node * new-node = new Node (val);

new-node -> next = temp -> next;

temp -> next = new-node;

new-node -> prev = temp;

new-node -> next -> prev = new-node;

return

}

→ Deletion at kth Position in a doubly linked list.

⑦ Void deleteAtStart () {

if (head == NULL) {

return;

}

Node * temp = head;

head = head -> next

if (head == NULL) {

head -> prev = NULL; }

}

free (temp);

return;

}

if (head == NULL) {

tail = NULL;

else {

head -> prev = NULL

}

① Delete At End

```

Void deleteAtEnd () {
    if (head == NULL) {
        return;
    }

    Node* temp = tail;
    tail = tail->prev;
    if (tail == NULL) {
        head = NULL;
    } else {
        tail->next = NULL;
    }

    free (temp);
    return;
}

```

② Add a node at Arbitrary Position

(C(k))

```

Void deleteAtPosition ($ int k) {
    // k <= length of ll.
    Node* temp = head;
    int count = 1;

    while (count < k) {
        temp = temp->next;
        count++;
    }
}

```

// Now temp is pointing to node at k^{th} position.

```

temp->prev->next = temp->next;
temp->next->prev = temp->prev;
free (temp);
return;
}

```

Q13 Reverse a doubly linked list. \Rightarrow

Sol: Void reverseLL (Node*&head, Node*&tail) {

```
    Node* currptr = head;
    while (currptr) {
        Node* nextptr = currptr->next;
        currptr->next = currptr->prev;
        currptr->prev = nextptr;
        currptr = nextptr;
    }
```

// Swapping head and tail pointers

Node* newhead = ~~head~~.tail;

~~tail~~ = ~~head~~.head;

head = newhead;

Q21 Find it's a palindrome or not. (DLU)

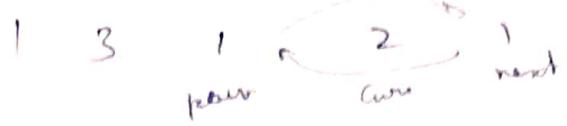
Sol: Bool isPalindrome (Node* head, Node* tail) {

```
    while (head != tail && tail != head->prev) {
        if (head->val != tail->val) {
            return false;
        }
    }
```

head = head->next;

tail = tail->prev;

return true;



Q3→ Given head (DLL), delete the nodes whose neighbours has the same value. Traverse the list from right to left.

Sol→ Void deleteSameNeighbourNode (Node^{*} &head, Node^{*} &tail) {

 Node^{*} currptr = tail → prev;

 while ((currptr) == head) {

 Node^{*} prevNode = currptr → prev;

 Node^{*} nextNode = currptr → next;

 if (prevNode == nextNode) {

 prevNode → next = nextNode;

 nextNode → prev = prevNode;

 free(currptr);

 }

 currptr = prevNode;

 }

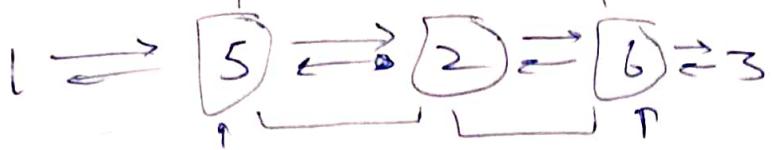
 return;

}

Q4

A critical point in a linked list is defined as either a local maxima or a local minima. Given a linked list tail, return an array of length 2 containing [min distance, max distance] where minDistance is the minimum distance b/w any two distinct critical points and maxDistance is the maximum distance b/w any two distinct critical points. If there are fewer than two critical points, return [-1, -1].

Note that a node can only be a local maxima/minima if there exists both a previous node and a next node.



Sol: bool iscriticalPoint (Node* &currNode) {

if (currNode->prev->val < currNode->val & currNode->next->val < currNode->val)
return true;

if (currNode->prev->val > currNode->val & currNode->next->val > currNode->val)
return true;

}

Vector<int> distance b/w Critical Points (Node* head, Node* tail) {

Vector<int> ans (2, INT_MAX);

int firstCP = -1, lastCP = -1;

Node* currNode = tail->prev;

if (currNode == NULL) {

ans[0] = ans[1] = -1;

return ans;

} int currPos = 0;

while (currNode->prev != NULL) {

if (iscriticalPoint (currNode)) {

if (firstCP == -1) {

firstCP = lastCP = currPos;

} else {

ans[0] = min (ans[0], currPos - firstCP);

ans[1] = currPos - firstCP;

lastCP = currPos;

}

currPos++;

currNode = currNode->prev;

}

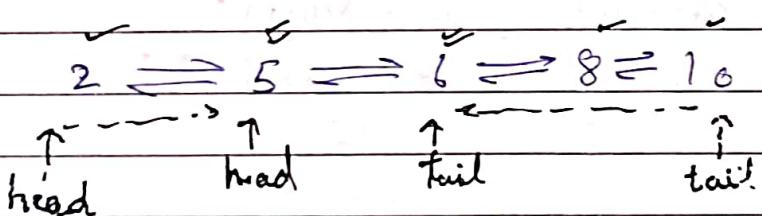
```

if (ans[0] == INT_MAX) {
    ans[0] = ans[1] = -1;
}
return ans;
}

int main() {
    vector<int> ans = distance1ToCriticalPoints(dll.head, 1);
    cout << ans[0] << " " << ans[1] << endl;
    return 0;
}

```

Q → Given the head of a DLL. The values of the linked list are sorted in non-decreasing order. Find if there exists a pair of distinct nodes such that the sum of their values is x . Return the pair in the form of a vector $[l, r]$, where l and r are the values stored in the two nodes pointed by the pointers. If there are multiple such pairs return any of them. If there is no such pair return $[-1, -1]$.

 $(x=11)$ 

Sol → Vector<int> pairSumDLL (Node* head, Node* tail, int x) {

```

vector<int> ans(2, -1);
do while (head != tail) {
    int sum = head->val + tail->val;
    if (sum == x) {
        ans[0] = head->val;
        ans[1] = tail->val;
    }
    return ans;
}

```

if ($\text{sum} > n$) { // need smaller values, i will move tail = tail \rightarrow prev; tail back.}

{ else { // need bigger values, i will move head forward. head = head \rightarrow next; }

{}

{}

return ans; }

{ }

int main() {

vector<int> ans = pairSumDLL(dll.head, dll.tail), 11);

(out << ans[0] << " " << ans[1] << endl;

return 0;

{}

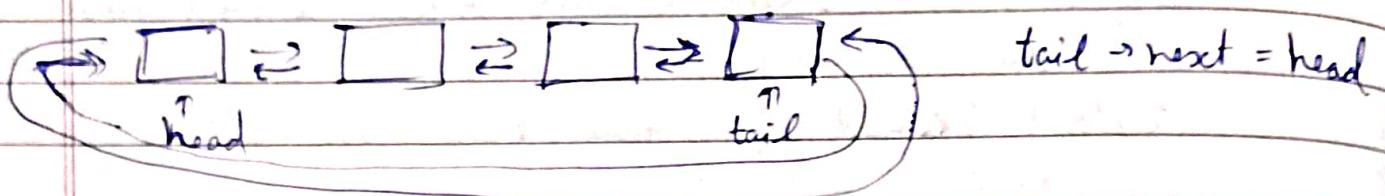
Lecture → 53

(Linkedlist (Part-4))

FREEMIND

Date _____
Page _____

Circular Linked List.



⇒ Advantages

- We can start traversing from any node to print all nodes until we reach visited nodes.
- When we have to traverse linked list in circular fashion multiple times.

⇒ Implementation of Node → Same

⇒ Traversal

```
temp = head;  
while (temp->next != head) {  
    temp = temp->next;  
}
```

⇒ Insertion at kth position in a circular linked list.

① Add a node at the start

Close Node {

};

class Circularlinkedlist { // Same as Preve

};

Void insertAtStart (int val) {

Node* new-node = new Node(val);

if (head == NULL) {

head = new-node;

new-node->next = head; //Circular linked list

return;

These all
functions are
passed
inside class
so, head
is not
given in function.

use do {
} while (temp != head)
at time of
display

}

`Node* tail = head;`

`while (tail->next != head) {`

`tail = tail->next;`

`}`

`// now tail is pointing to the last node`

`tail->next = new-node;`

`new-node->next = head;`

`head = new-node;`

`}`

`void printCircular() {`

`Node* temp = head;`

`for (i=0; i < 15; i++) {`

`cout << temp->val << " -> ";`

`temp = temp->next;`

`} cout << endl;`

`}`

`int main()`

`circularlinkedlist CLL;`

`CL.insertAtStart(3);`

`CL.printCircular();`

`return 0;`

`}`

① Add a node at end

Same as Prev.

`tail->next = new-node;`

`new-node->next = head;`

`// head will remain same.`

② Add Node at Arbitrary Position

Same as Singly LL.

⇒ Deletion at kth Position

① void deleteAtStart () {

 if (head == NULL) {
 return;

 }

 Node * temp = head;

 Node * tail = head;

 while (tail->next != head) {
 tail = tail->next;

 }

 head = head->next;

 tail->next = head;

 free (temp);

 }

② void DeleteAtEnd () {

 if (head == NULL) {
 return

 }

 Node * tail = head;

 while (tail->next->next != head) {

 tail = tail->next

 }

// here tail is pointing to second last

Node * temp = tail->next; // to be deleted

tail->next = head;

free (temp);

 }

③ Delete at arbitrary Position.

Some or Singly LL.

★ Template Classes :-

→ they allow us to pass data type as parameter

```
#include <iostream>
```

```
using namespace std;
```

```
template <typename T>
```

```
class Node { // Template class
```

```
public:
```

```
    T val;
```

```
    Node* next;
```

```
    Node(T data) {
```

```
        val = data;
```

```
        next = NULL;
```

```
}
```

```
};
```

```
int main() {
```

```
    Node<int>* node1 = new Node<int>(3);
```

```
    cout << node1->val << endl;
```

```
    Node<char>* node2 = new Node<char>('a');
```

```
    cout << node2->val << endl;
```

```
    return 0;
```

```
}
```

★ STL :-

(Standard Template Library)

→ set of template classes for implementing commonly used data structures & functions.

eg: vector<int>
 <char>

list

② 3 major components

set

→ containers

map

→ Standard

queue

→ Algorithms

stack

⇒ Containers: used to hold data of the same type

e.g.: vector - class that defines a dynamic array.

list → class for doubly linked list.

map → used to store unique key value pair

set ⇒ used to store unique values

⇒ Iterators: → used to iterate/traverse the container.
→ pointer like entities.

`vector<int> v = {1, 2, 3, 4, 5}`

`vector<int> :: iterator itr;`

`itr.n. begin() = 1`

⇒ Algorithms:

functions that are used to manipulate data in the containers

e.g.

`sort()` → `sort(v.begin(), v.end())`

`min()` → `min(a, b)`

`max()` → `max(a, b)`

C++ .com for
these func./template

★ What is a list?

→ template class in STL for implementing doubly linked list.

⇒ declaration of a list

`#include <list>`

`list<datatype> listname`

`list<int> rollNo;`

`list<int> list1: {1, 2, 3, 4}`

`list<int> list2 {5, 6, 7, 8}`

→ Advantage of list in C++ STL :-
 → implementation becomes easy

→ iterator functions

`list.begin()` → iterator for 1st element $\begin{bmatrix} 1, 2, 3, 4 \\ \pi \uparrow \end{bmatrix}$

`list.end()` → iterator for the position after last element. $\begin{bmatrix} 1, 2, 3, 4 \end{bmatrix}_\pi$

`list.rbegin()` → iterator for the first element in reverse iteration $\begin{bmatrix} 1, 2, 3, 4 \end{bmatrix}_{\pi \uparrow}$

`list.rend()` → iterator for the position after last element in reverse iteration $\begin{bmatrix} 1, 2, 3, 4 \end{bmatrix}_{\pi \uparrow}$

`advance(itr, n)` → advances the iter by n places.

$(itr, 1) \rightarrow \begin{bmatrix} 1, 2, 3, 4 \end{bmatrix}_{\pi \uparrow}$

#include <iostream>

#include <list>

Using raw

int main() {

list<int> L1 = {1, 2, 3, 4};

auto itr = L1.begin();

cout << *itr << endl;

// 1

auto rev_itr = L1.rbegin();

cout << *rev_itr << endl;

// 4

advance(itr, 2);

cout << *itr << endl;

// 3

// Using Range based loop.

for (auto num : L1) {

// traversal of list.

cout << num << endl;

}

// Using iterators for (auto itr = L1.begin(); itr != L1.end(); itr++) {

cout << *itr << " ";

} cout << endl;

// reverse traversal

```
for (auto itr = L1.begin(); itr != L1.end(); itr++) {
    cout << *itr << " ";
```

}

}

★ Inserting elements into a list

list.insert(itr, value) → insert value before the position of the itr.

list.insert(itr, count, value) → insert value count no. of times before itr.

list.insert(itr, str_itr, end_itr) → insert values from str_itr to end_itr before itr.

e.g. // from previous --

→ auto itr = L1.begin();

advance(itr, 2);

auto l = L1.begin();

auto r = L1.begin();

advance(r, 2);

L1.insert(itr, l, r); // [1, 2, 1, 2, 3, 4]

for (auto itr = L1.begin(); itr != L1.end(); itr++) {

cout << *itr << " ";

cout << endl;

[1, 2, 3, 4]

↑ ↑
l or itr

(l included
or not included)

★ Deleting elements from a list.

list.erase(itr) → delete the element pointed to by the itr.

list.erase(str_itr, end_itr) → delete elements from str_itr to end_itr.

→ Other member function of a list container.

push-front (Val) → insert val in the front of list

pop-front () → removes val from front of list

push-back (Val) → insert val in the back of the list.

popback () → removes val from back of list.