

Stacks -2

Lecture-46

Raghav Garg

COLLEGE
WALLAH

Ques: Balanced Brackets

'(', ')'

'Easy'

→ string s = "()()()"; true

string s = "(()())"; true

string s = "()(())"; false

string s = ")()("; false

string s = "((()()))"; true

pop()
push()
top()

COLLEGE
WALLAH

Ques: Balanced Brackets

String s = "()_i()()";



st

Steps

- 1) If you see opening bracket
→ push
- 2) If you see a closing bkt,
st → top ko dekho , agar → '('
st.pop()

→ If st.size() == 0 , true

Ques: Balanced Brackets

string s = "(())()"; true
i



Steps

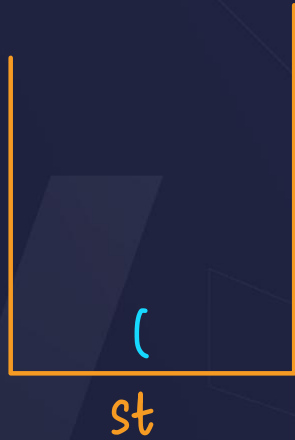
1) If you see opening bracket
→ push

2) If you see a closing bkt,
st → top ko dekho , agar → '('
st.pop()

→ If st.size() == 0, true

Ques: Balanced Brackets

string s = "() (() () ";



Steps

1) If you see opening bracket
→ push

2) If you see a closing bkt,
st → top ko dekho , agar → '('
st.pop()

→ If st.size() == 0 , true
else false

Ques: Balanced Brackets

string s = "())()(";
i



st

Steps

1) If you see opening bracket
 → push

2) If you see a closing bkt,
 st → top ko dekho , agar → '('

st.pop() , agar st empty → false

→ If st.size() == 0 , true

else false

Ques: Remove Consecutive Duplicates in a string

string s = "aaabbcddaa bffg";

str = "gfbadcb a"

rev ↪ "a b c d a b f g"



st
char

COLLEGE
WALLAH

Ques: Next greater element

Problem Statement :

arr	3	1	2	5	4	6	2	3
ans	5	2	5	6	6	-1	3	-1

Brute Force Solution : T.C. S.C.
 $O(n^2)$ $O(1)$

$O(n)$

COLLEGE
WALLAH

Method-2 'Using stack'

	<i>i</i>							
arr	3	1	2	5	4	6	2	3
ans	5	2	5	6	6	-1	3	-1

```
ans[n-1] = -1;
st.push(arr[n-1])
```



```
while (st.top() <= arr[i]) st.pop(); POP
```

```
ans[i] = st.top;    ANS
```

```
st.push(arr[i]);    PUSH
```

pop, ans, push

Example

arr = 4 1 2 5 4 3 4 8 2 7
ans = 5 2 5 8 8 4 8 -1 7 -1



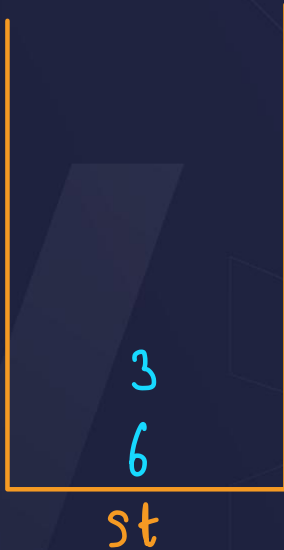
pop

ans

push

Ques: Previous greater element

arr	3	1	2	5	4	6	2	3
pge	-1	3	3	-1	5	-1	6	6



pop
ans
push

COLLEGE
WALLAH

Ques: Stock Span problem

Input - {100,80,60,70,60,75,85}

Output - {1 , 1 , 1, 2 , 1, 4 , 6}

span , pehle se lekar
continuously

	0	1	2	3	4	5	6
arr	100	80	60	70	60	75	85
ans	1	1	1	2	1	4	6

	0	1	2	3	4	5	6	7
arr	100	80	60	81	70	60	75	85
pge	-1	100	80	100	81	70	81	100
ans	1	1	1	3	1	1	3	7



Previous Greater
Element

	0	1	2	3	4	5	6	7	i
arr	100	80	60	81	70	60	75	85	
pgi	-1	0	1	0	3	4	3	0	
span	1	1	1	3	1	1	3	7	



pop \rightarrow while (arr[st.top()] \leq arr[i])

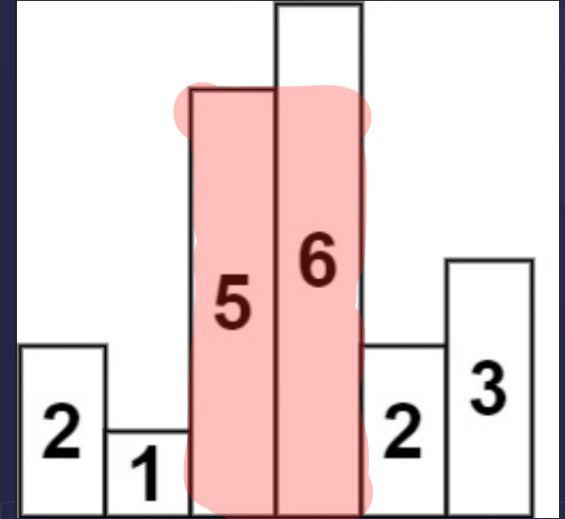
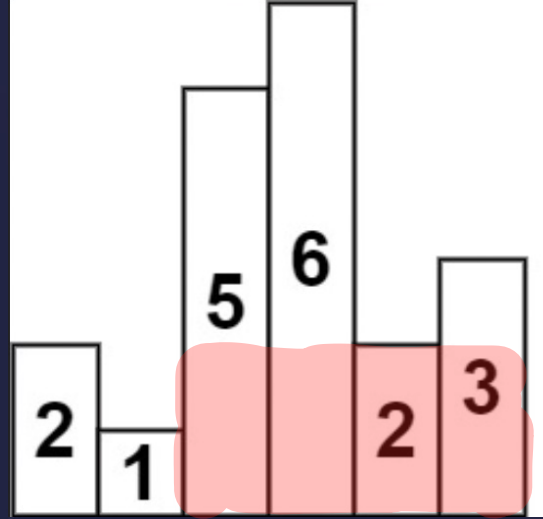
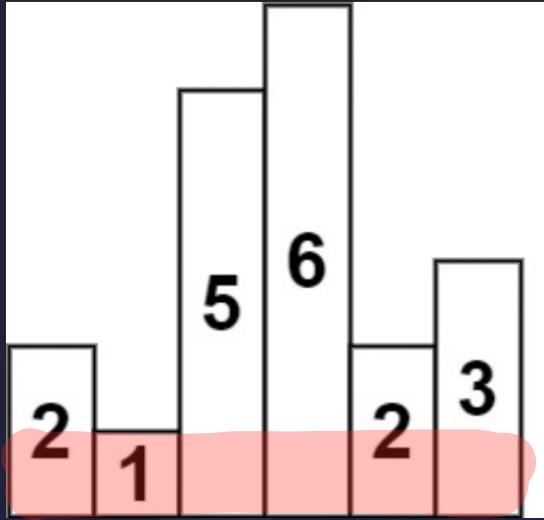
ans

push

Ques: Largest Rectangle in Histogram

[Leetcode - ~~84~~]

84

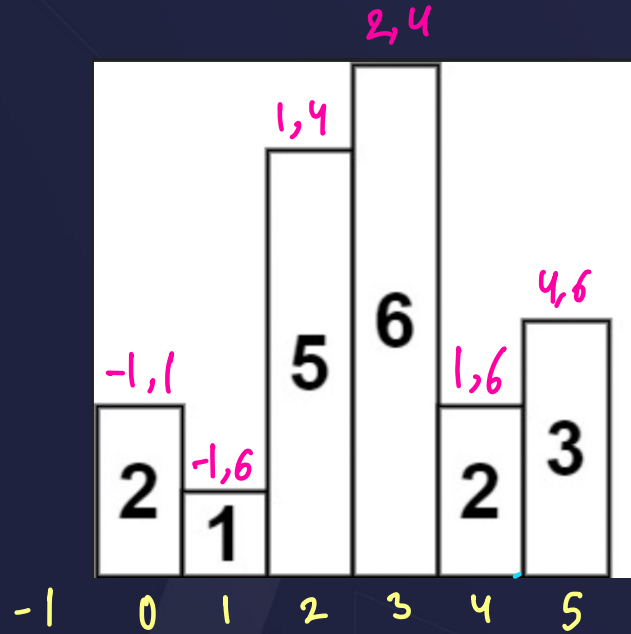


{ 2 , 1 , 5 , 6 , 2 , 3 }

COLLEGE
WALLAH

Ques: Largest Rectangle in Histogram

[Leetcode - ~~84~~]
84



2	1	5	6	2	3
2	6	10	6	8	3

Hint : Nse , Pse

Area = height * breadth

\downarrow
 $nsi[i] - psi[i] - 1$

Next
Smaller
Element

	-1	0	1	2	3	4	5	6
		2	1	5	6	2	3	
nse	1	-1	2	2	-1	-1		
nsi	1	6	4	4	6	6		

pop → after se bad element

ans

push

2
1

Previous Smaller Element

	-1	0	1	2	3	4	5	6
arr		2	1	5	6	2	3	
pse	-1	-1	1	5	1	2		
psi	-1	-1	1	2	1	4		



pop

ans

push

COLLEGE
WALLAH

```
int n = arr.size(); 2
int nsi[n];
stack<int> st;
nsi[n-1] = n;
st.push(n-1);
for(int i=n-2;i>=0;i--){
    while(st.size()>0 && arr[st.top()]>=arr[i]) st.pop();
    if(st.size()==0) nsi[i] = n;
    else nsi[i] = st.top();
    st.push(i);
}
```

Right ✓

	-1	0	1	2
arr		2	4	
nsi		2	2	
psi		-1	0	

```
int psi[n];
stack<int> gt;
nsi[0] = -1;
st.push(0);
for(int i=1;i<n;i++){
    while(gt.size()>0 && arr[gt.top()]>=arr[i]) gt.pop();
    if(gt.size()==0) psi[i] = -1;
    else psi[i] = gt.top();
    gt.push(i);
}
```

psi Right

```
int maxArea = 0; 4
for(int i=0;i<n;i++){
    int height = arr[i]; 4
    int breadth = nsi[i] - psi[i] - 1;
    int area = height * breadth; 4
    maxArea = max(maxArea,area);
}
return maxArea; 4
```

THANK YOU!

COLLEGE
WALLAH