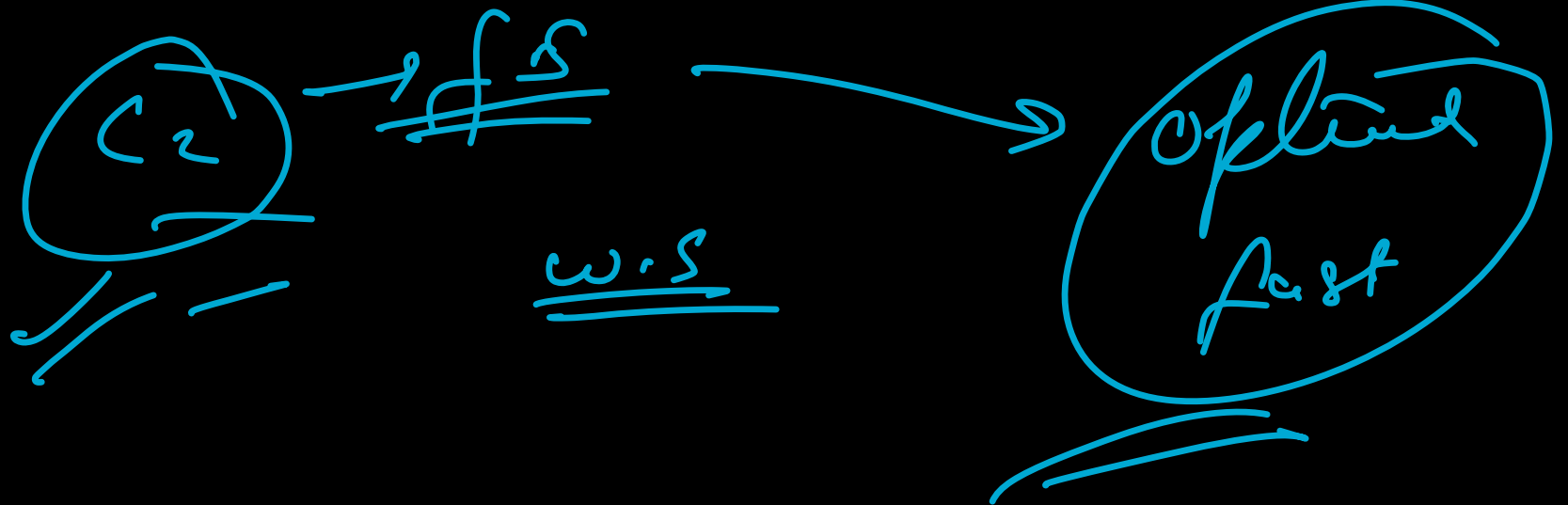
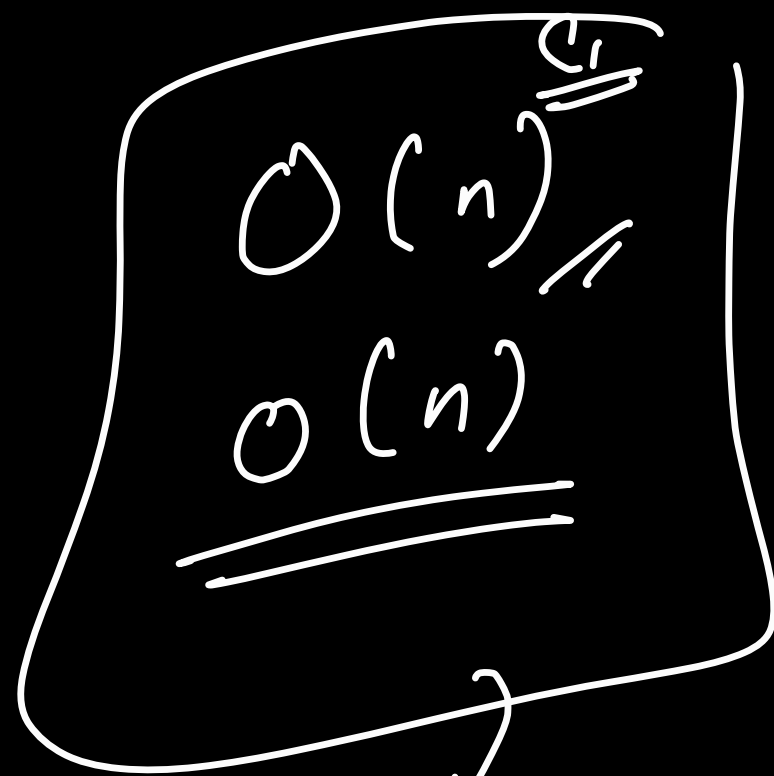
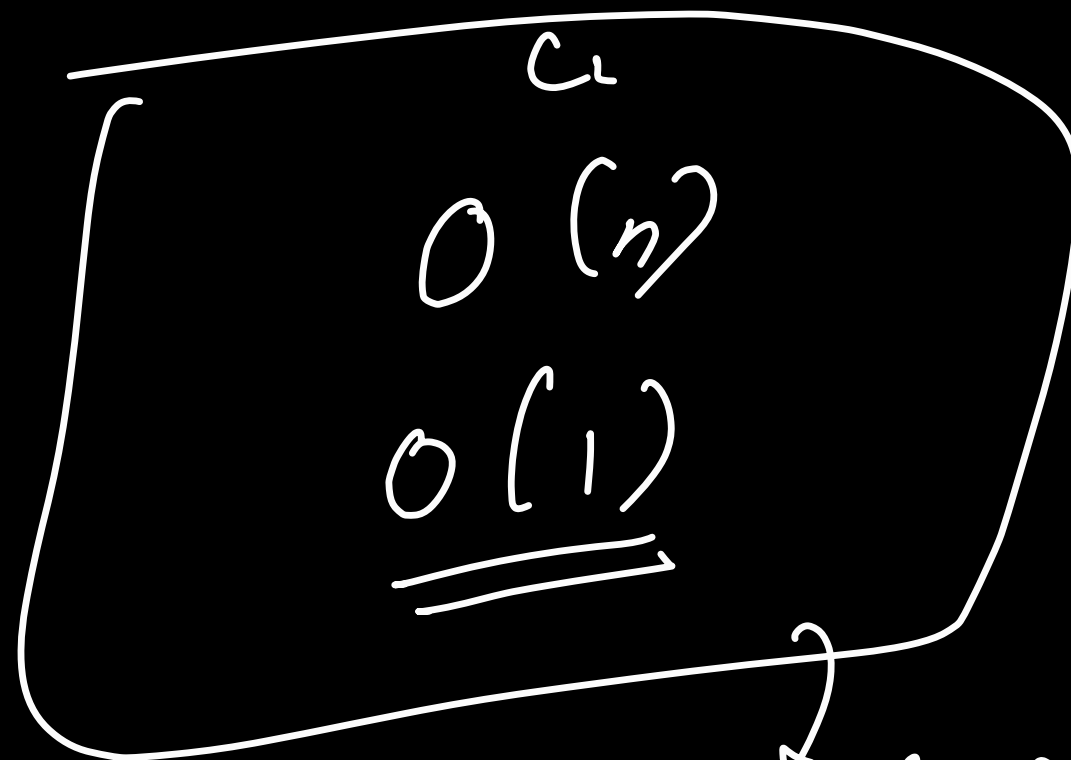


5/1

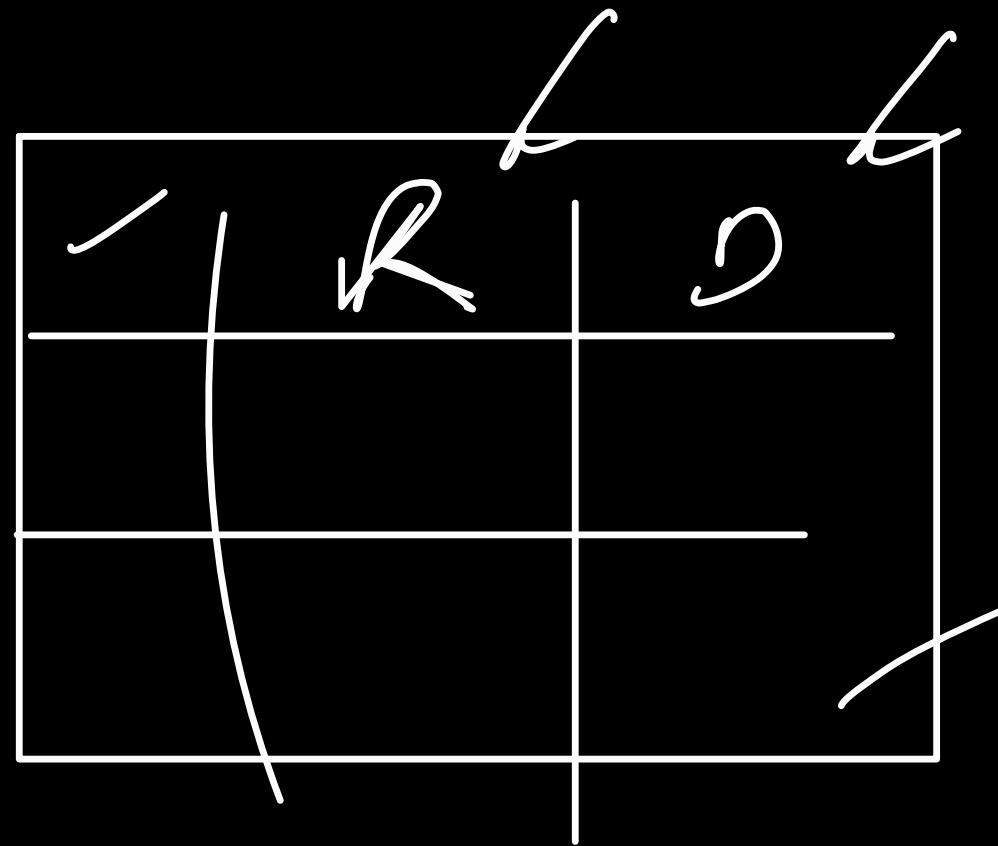




$$\underline{\underline{O(n^2)}}$$



$$\underline{\underline{O(n)}}$$



JOIN THE DARKSIDE

Prim's Algorithm → MST

↳ greedy

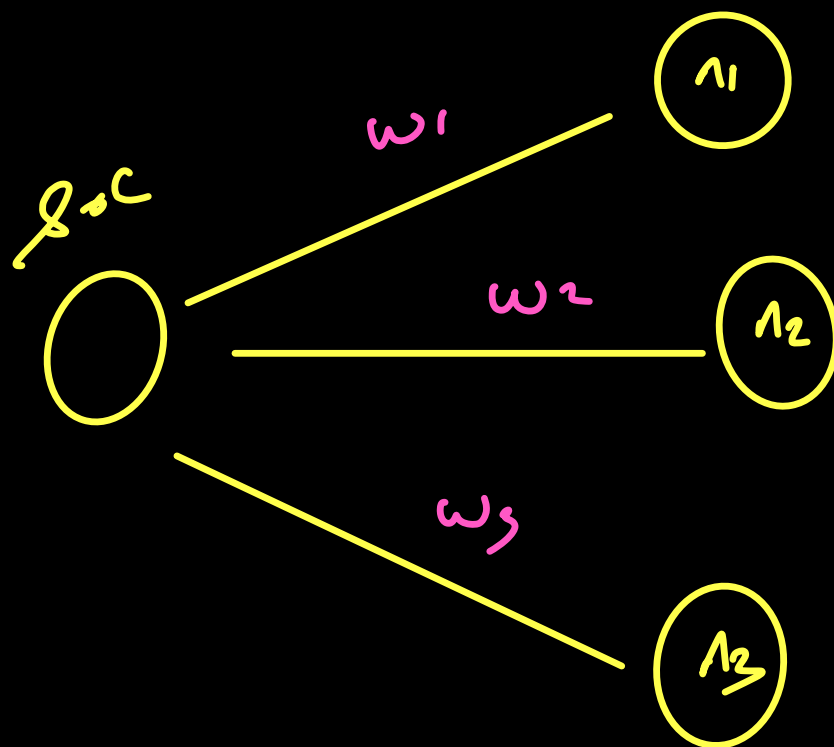
↳ Priority Queue

↳ bfs + greedy

↳ source node

sum + smallest
wt edge

src

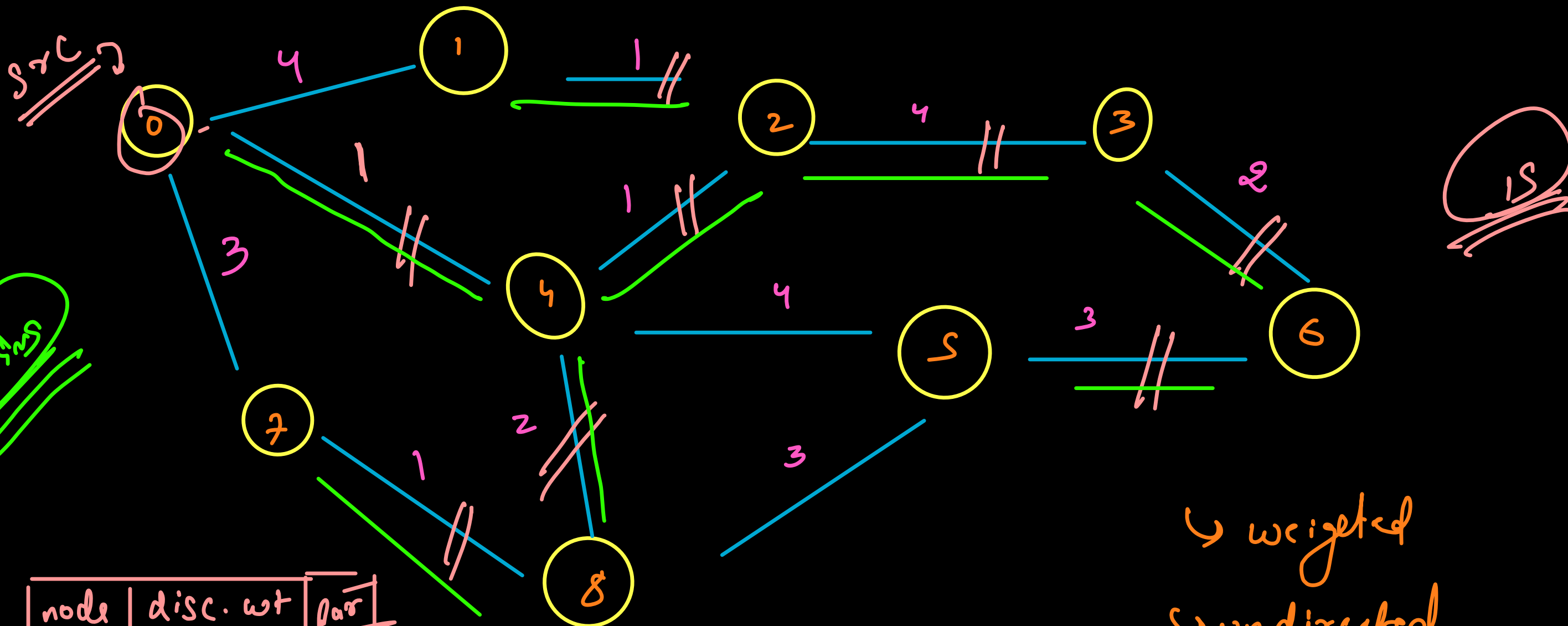


n1	n2	n3
----	----	----

vis → n1, n2, n3

greedy

priority queue → min heap



weighted
undirected

node	disc. wt	par
0	-1	-
1	1	0
2	1	1
3	4	2
4	1	0
5	3	4
6	2	3
7	1	0
8	2	4

$(0,0)$ $(1,4)$ $(1,2)$ $(2,8)$
 $(4,1)$ $(1,9)$ $(4,3)$ $(2,6)$
 $(3,5)$

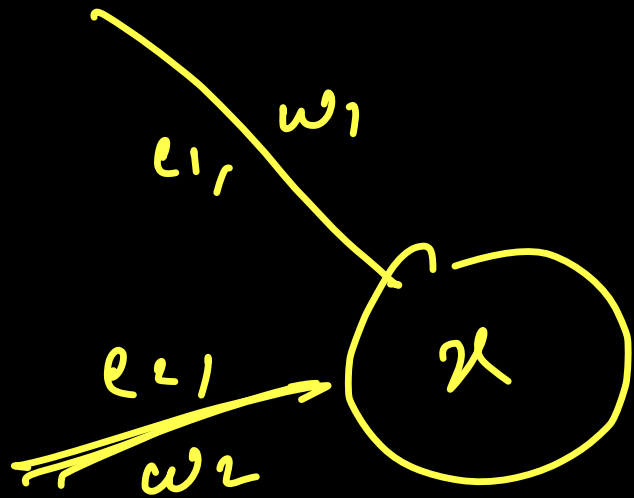
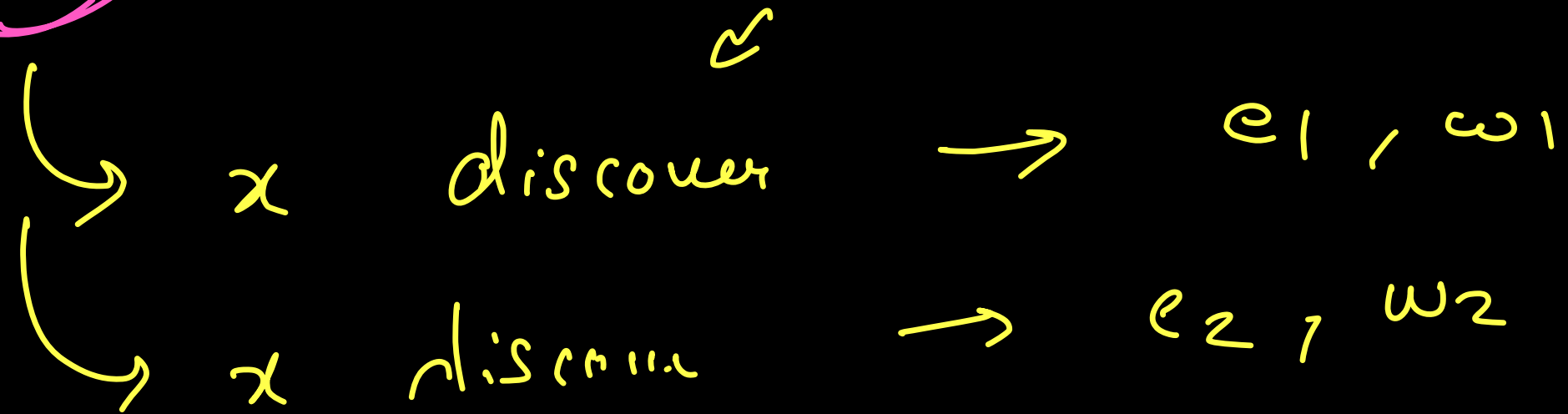
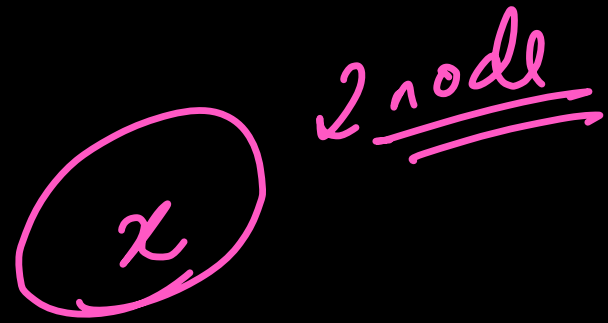
pair < wt, vertex >

0, 4, 2, 8, 1, 9 | 6, 3, 5

visited

In prime we start with a src & try to discover other

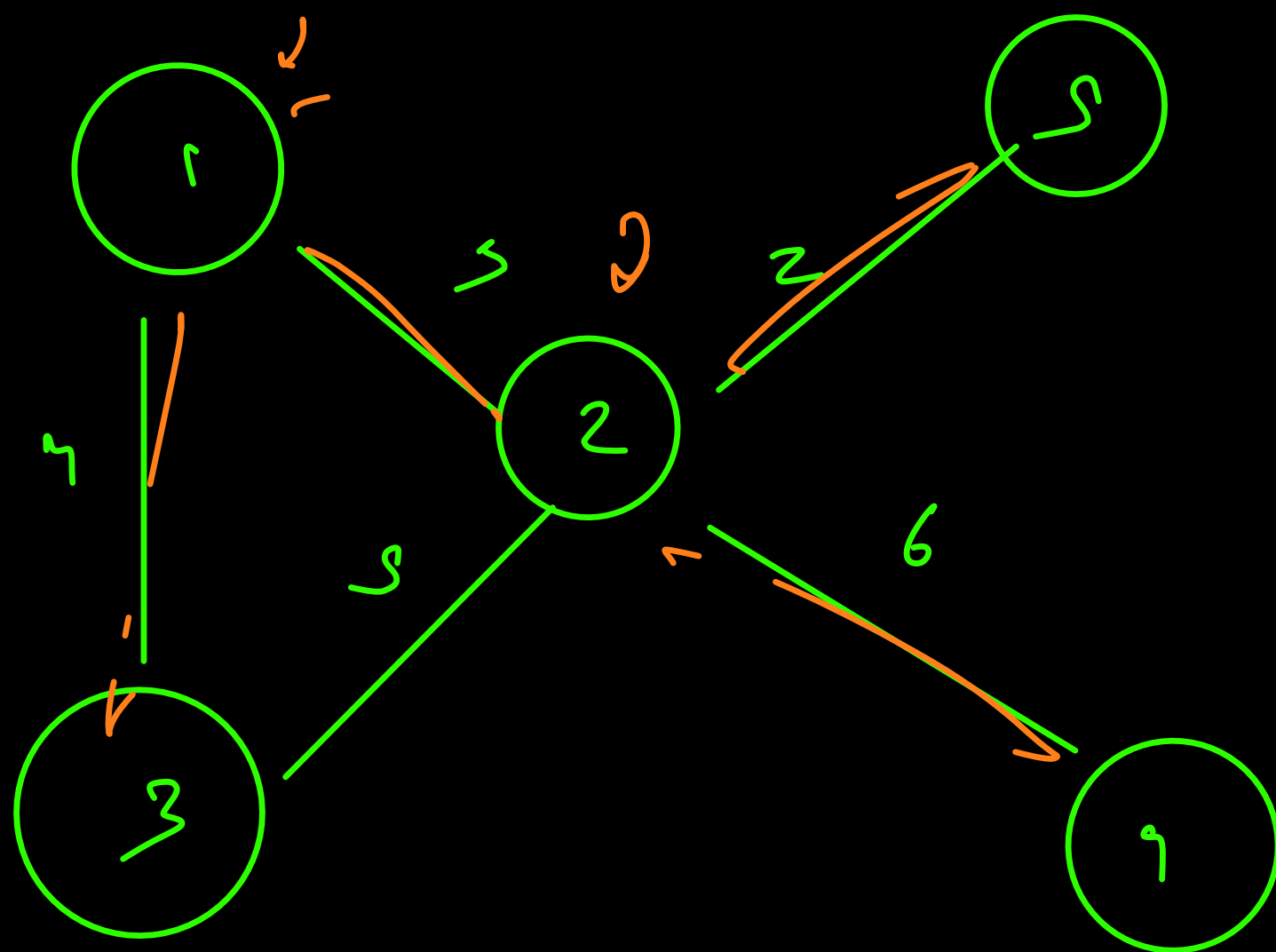
nodes.



$w_2 < w_1$

we should update the discover
wt by w_2 .

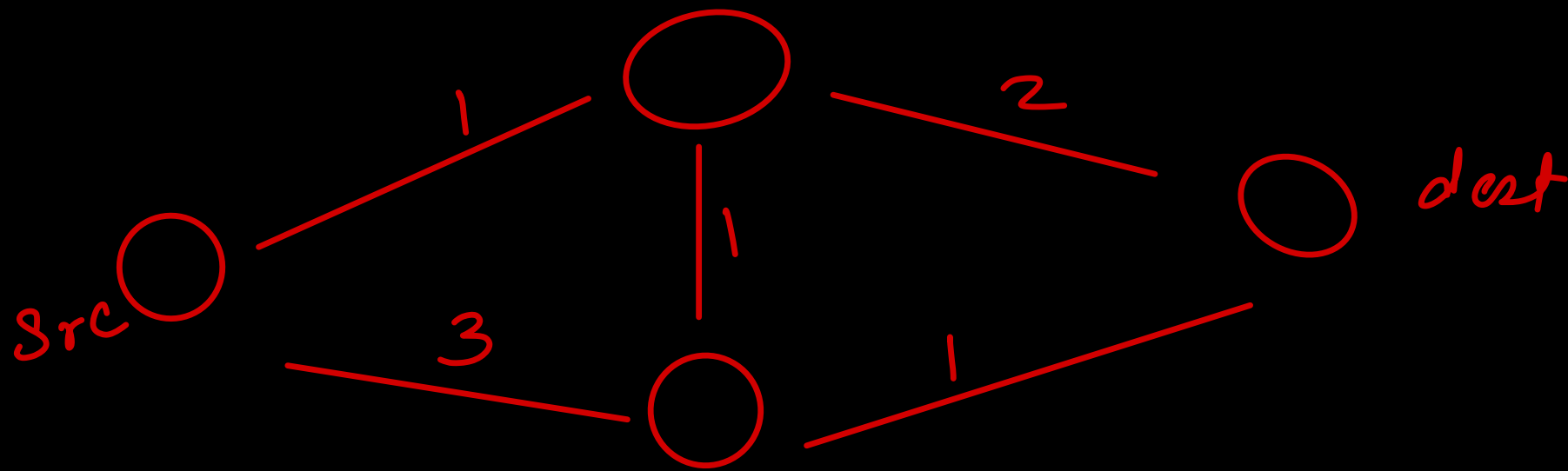
at the start of the algo, disc. wt of every node
except src is ∞ , because we not discovered
any step



$(3, 2)$
 $(2, 5)$
 $(4, 3)$

$(6, 4)$

Qn Given an undirected weighted graph, along with a src and dest. find the shortest path b/w the src & dest in terms of sum of edge costs.

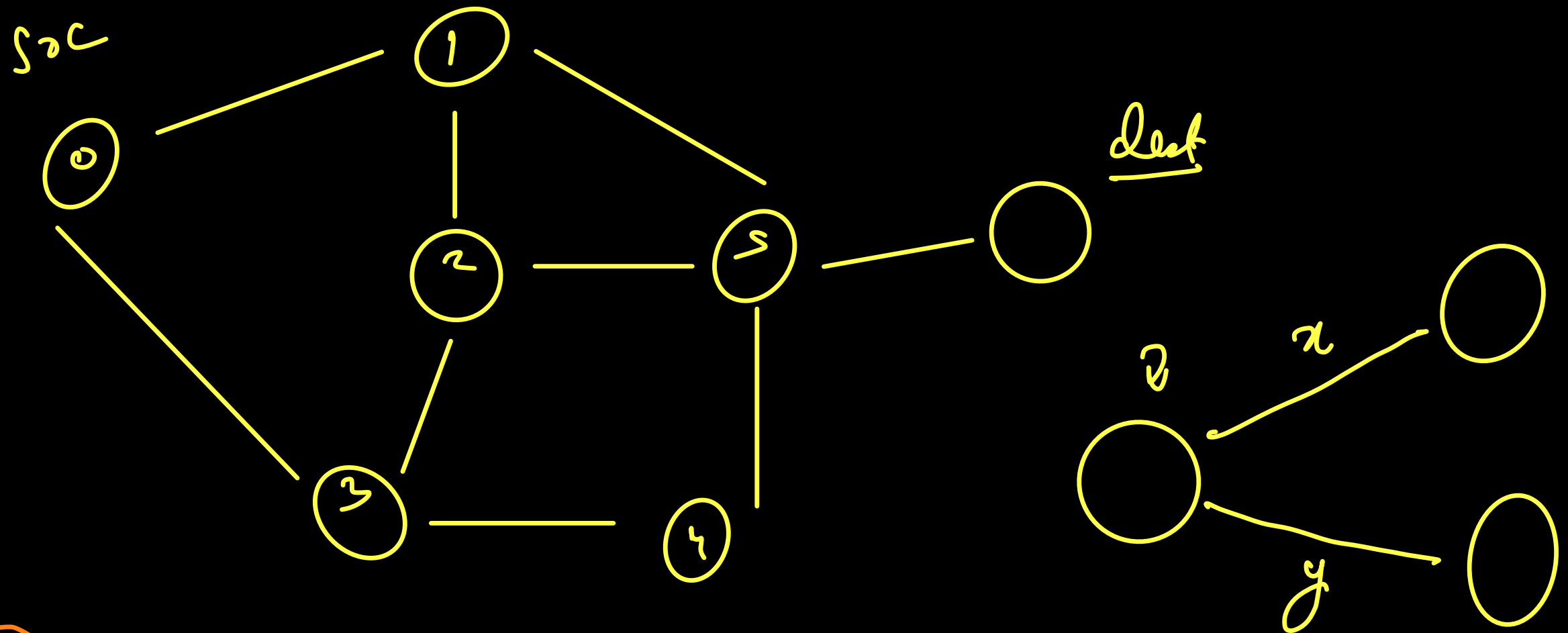


→ 3

we already know how to solve shortest path problem for unweighted graph.

BFS

queue \rightarrow fifo

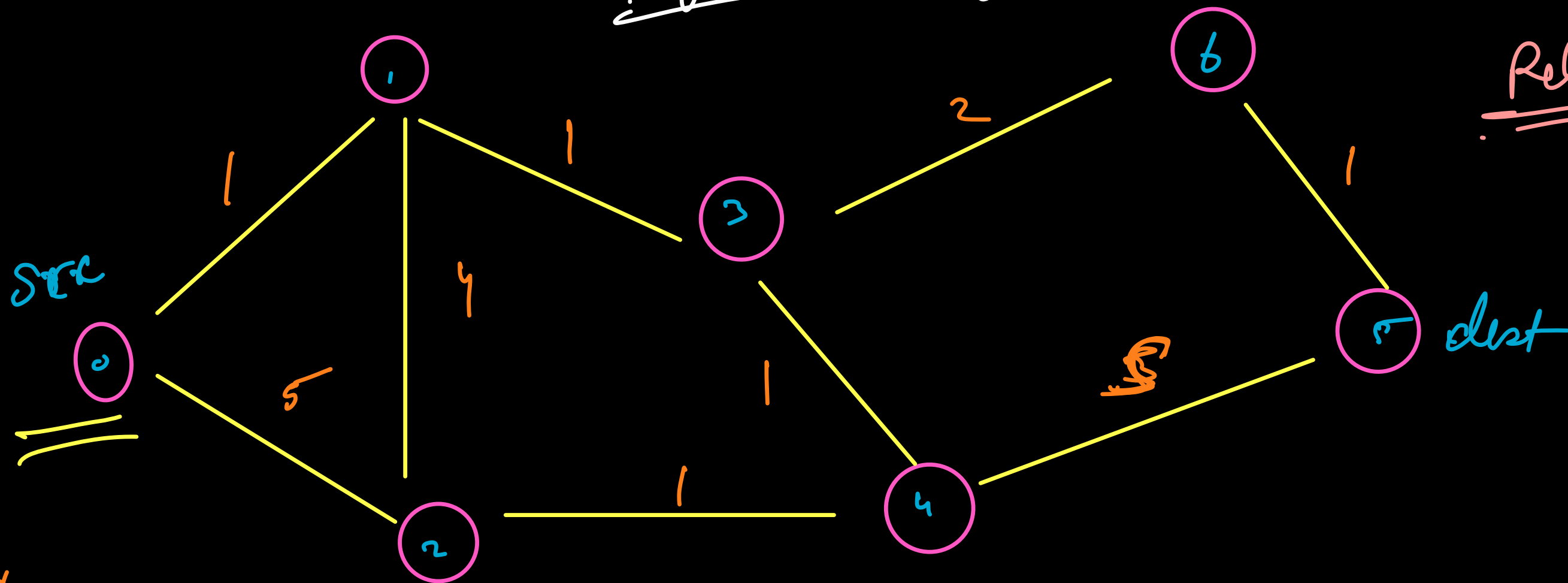


Prims

queue \rightarrow priority queue

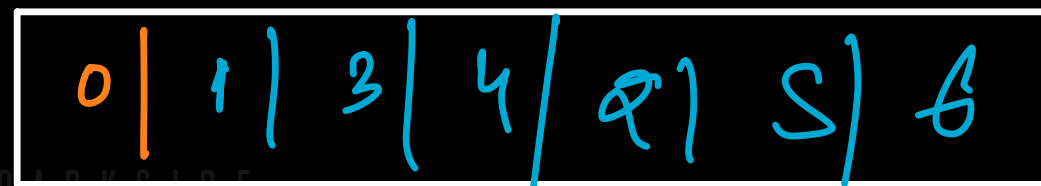
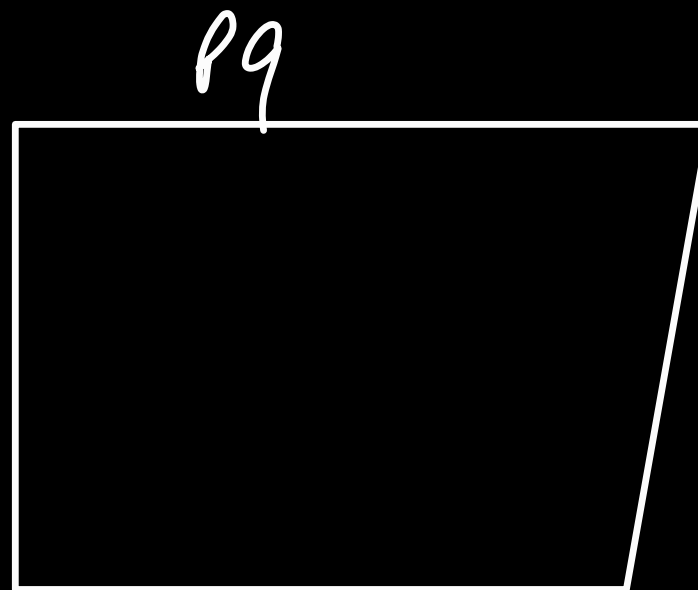
Dijkstra Algo

Relaxation



pair $\langle \text{dist}, \text{node} \rangle$

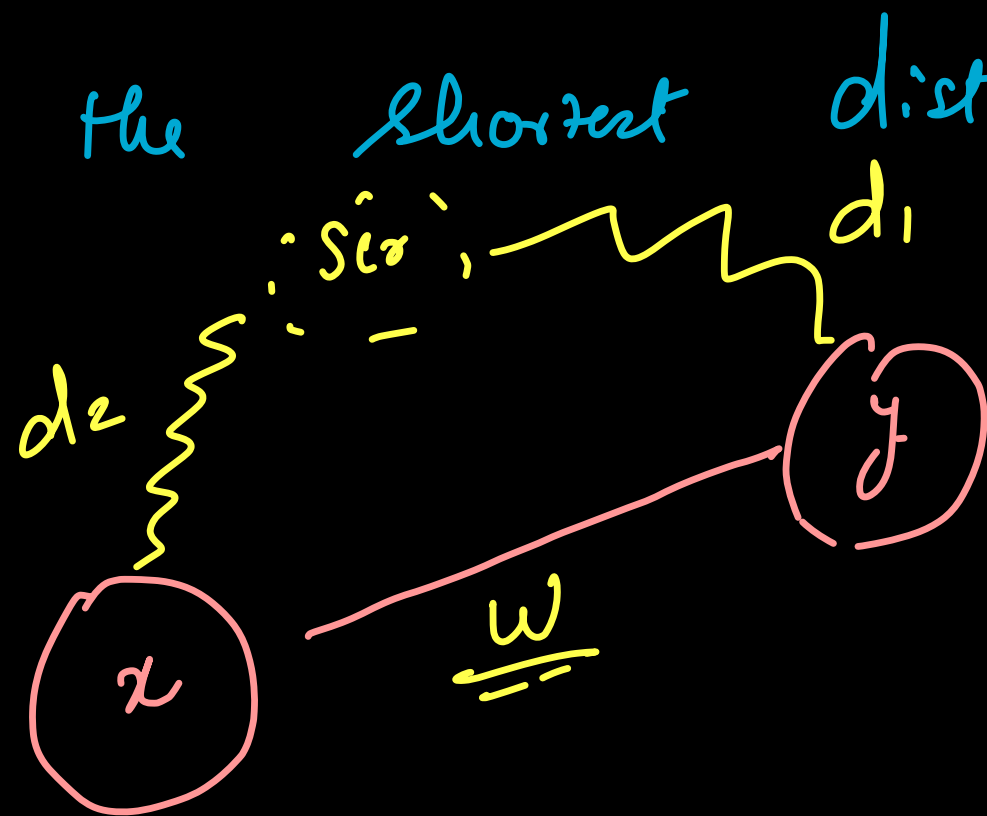
node	dist	via
0	0	xxx
1	1	0
2	5 4	0 1
3	2	1
4	3	3
5	8 5	4 6
6	4	3



vis

↓
 what is the shortest dist from src to any node.

given



node	dist
y	$d1 + d2 + w$ x
x	d2

$d1 \rightarrow$ till now what is the shortest path from
 src to y.
 $d2 \rightarrow$ till now what is the shortest path from
 src to x

$w \rightarrow$ edge wt of x and z

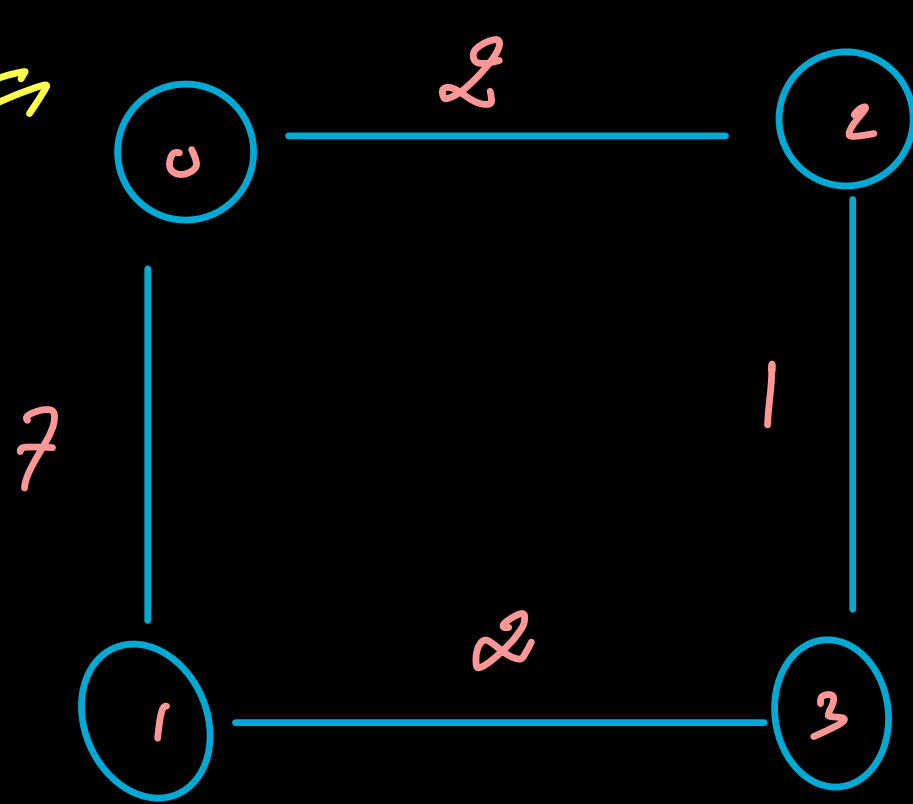
if ($d_z + w < d_1$) {

$mp[y] = d_z + w$

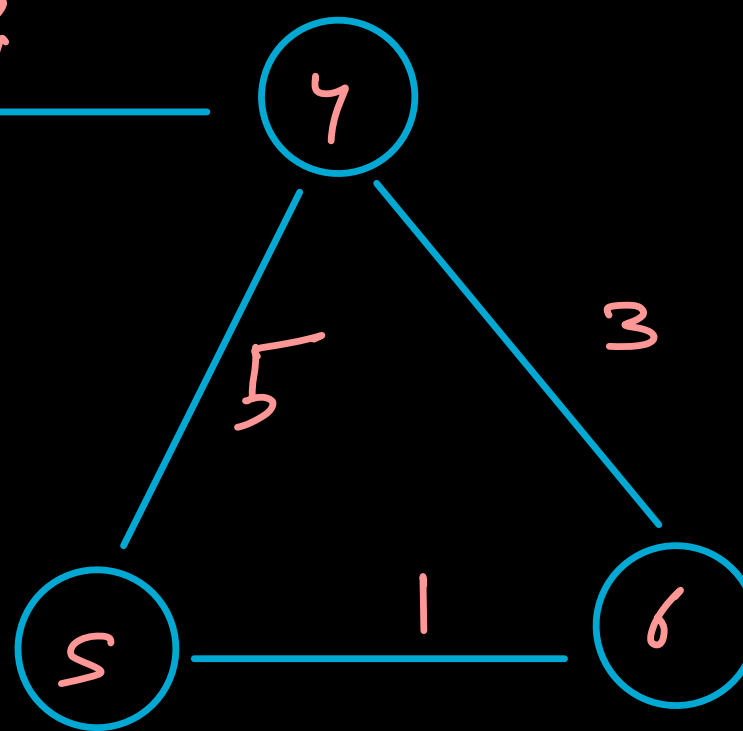
$via[y] = x$

}

~~src~~



2



dest

< dist, nodes

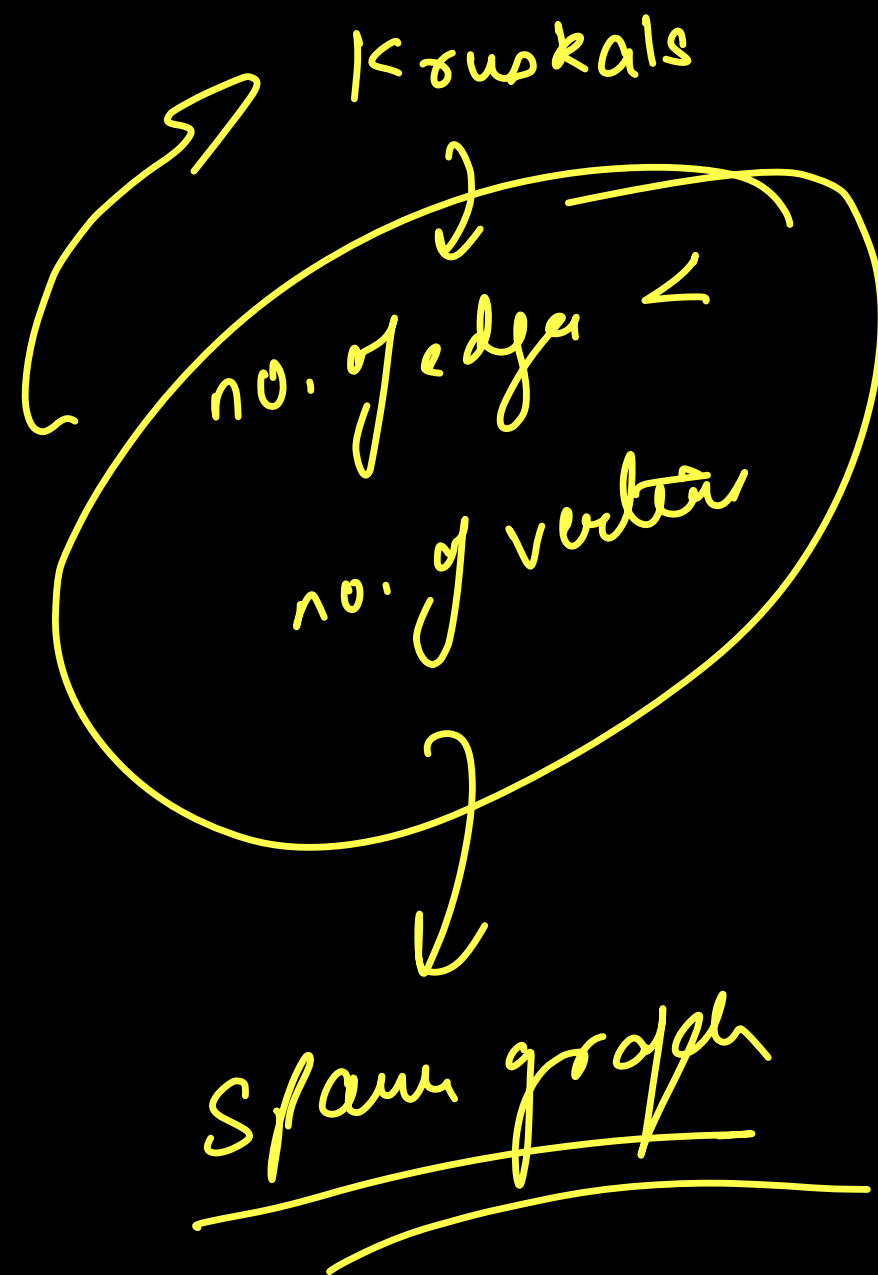
node	dist	via
0	0	xx
1	5	3
2	2	0
3	3	2
4	4	2
5	8	6
6	7	4



pg

0 | 2 | 3 | 4 | 1 | 6 | 5

vis



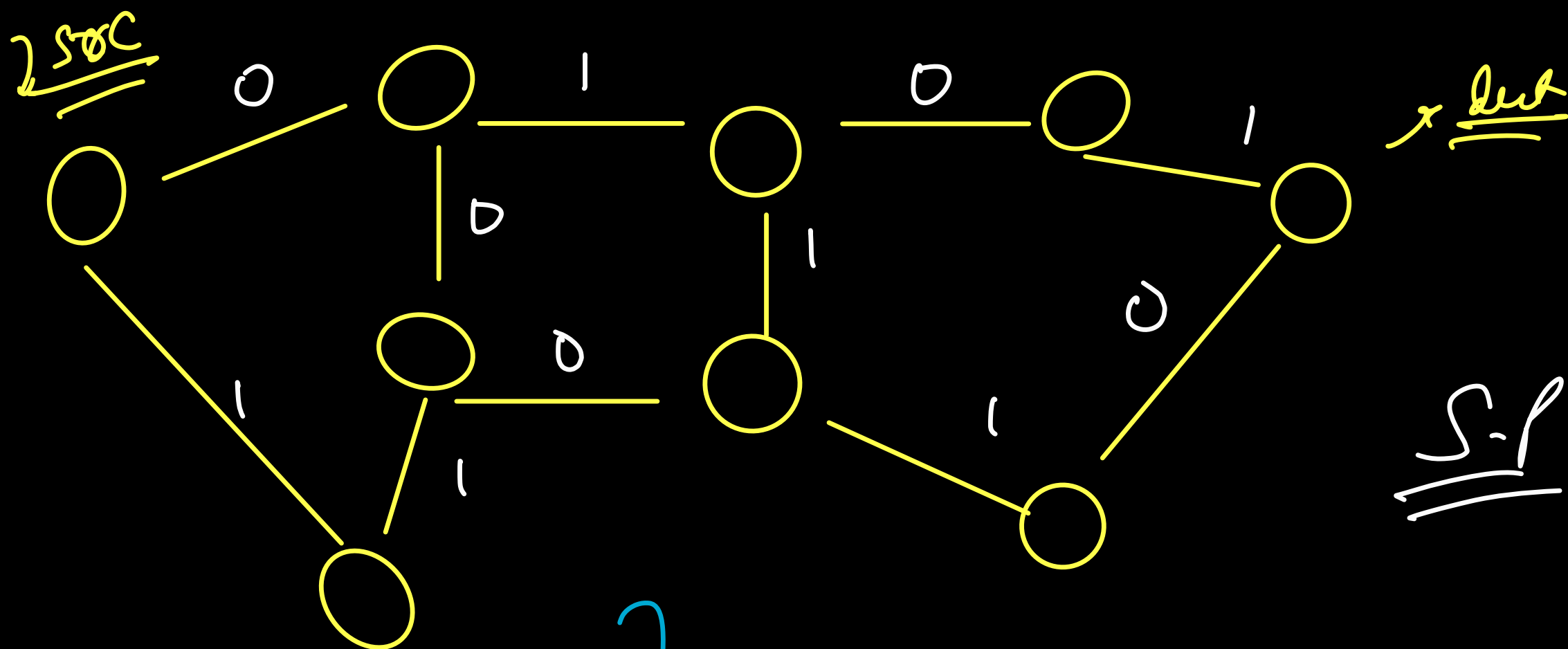
Points

no. of edges > no. of vertices

spann graph

The diagram shows a condition for a spanning graph. It starts with the word "Points" at the top. Below it, the text "no. of edges > no. of vertices" is written. An arrow points down from this text to the text "spann graph".

5

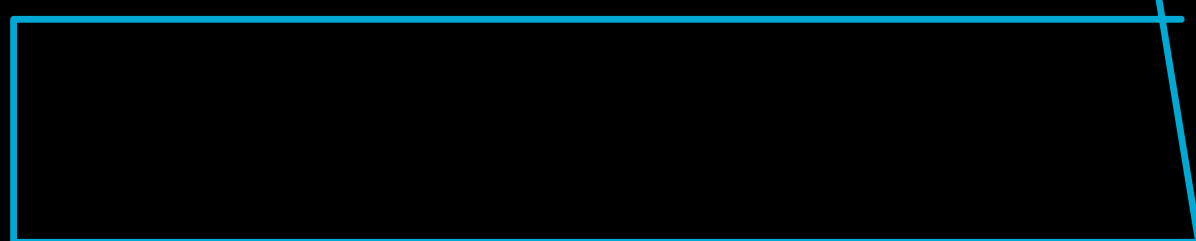


dist

S-p → Dist

0-1 BFS

google

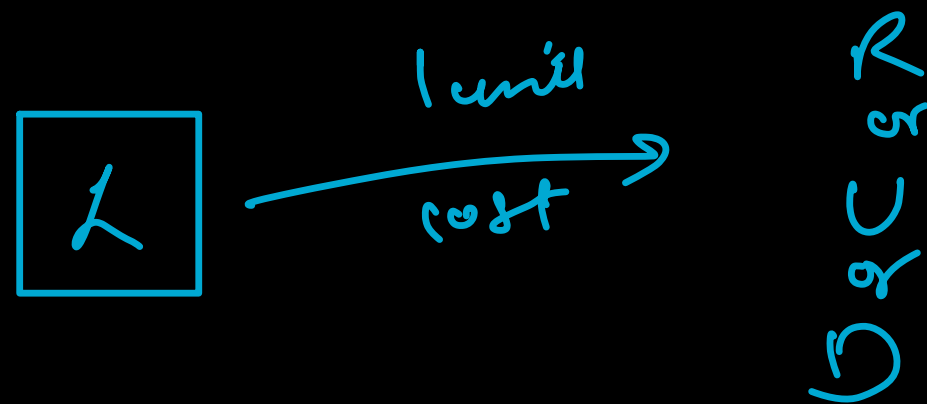


queue

fifo

deque

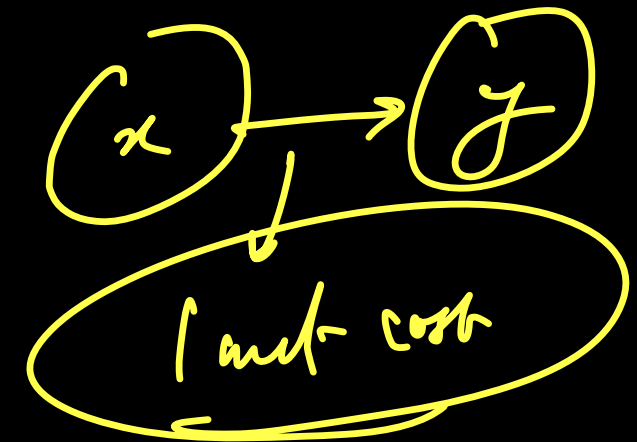
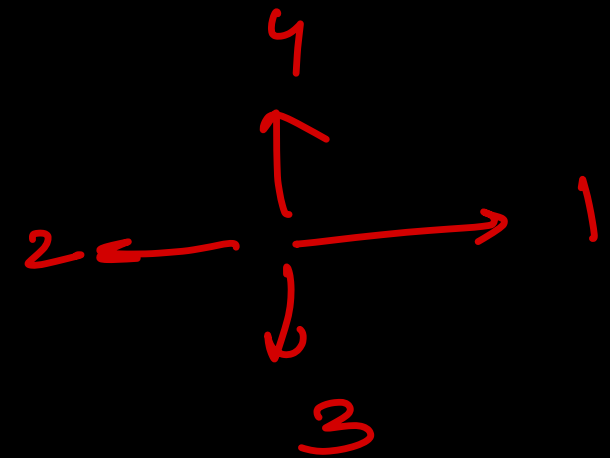
Q. Given a $m \times n$ grid, with $['L', 'R', 'U', 'D']$ chars mentioned on every cell. The char represent which cell we can move to from the cur cell. You're starting from top-left. Now by spending 1 unit cost we can replace the char of the cur cell to other chars.

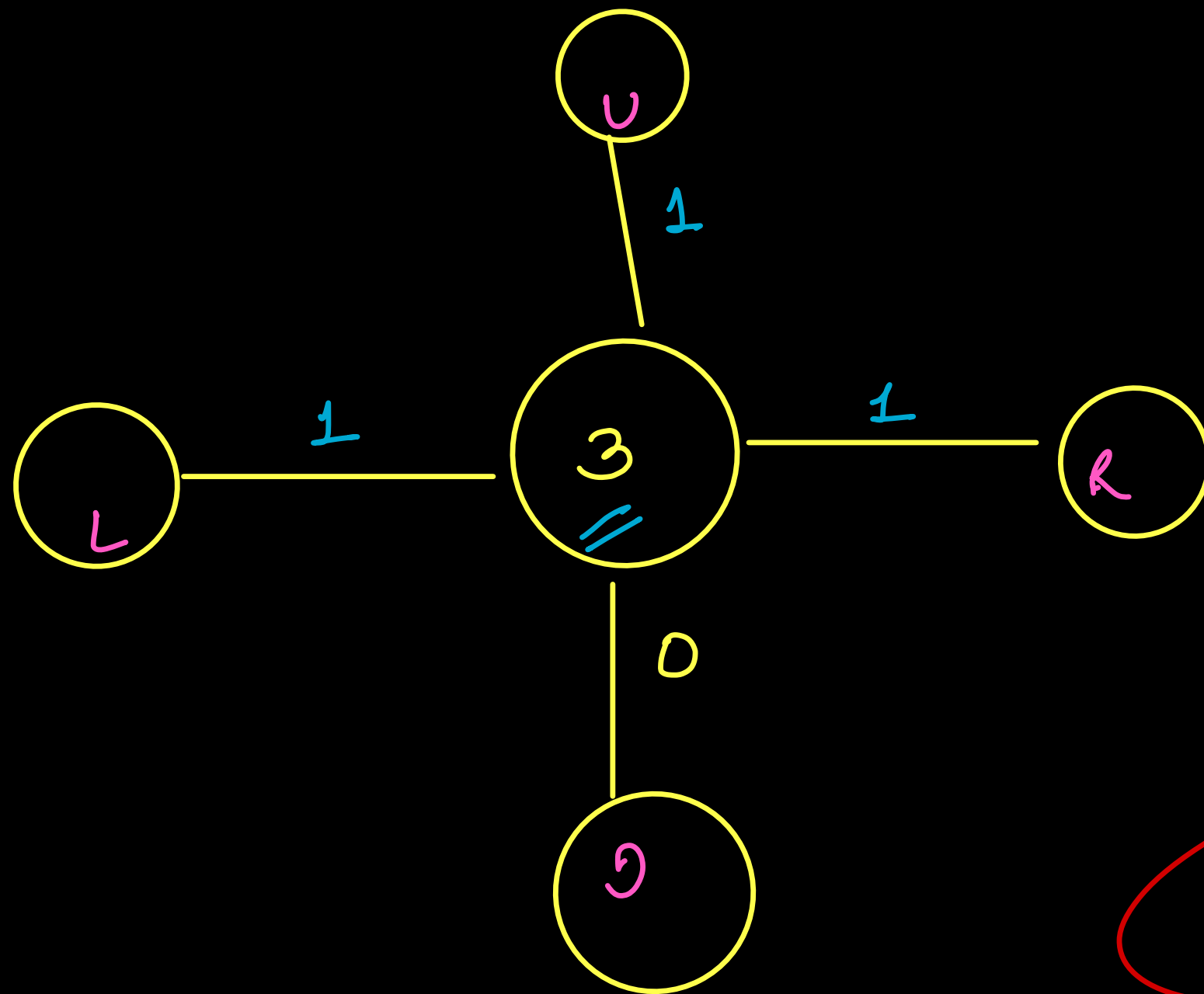


In what min cost we can reach the bottom Right

min cost

1	1	3	4	4
4	3	3	4	1
1	2	3	3	2
3	1	1	1	2
1	3	2	1	3





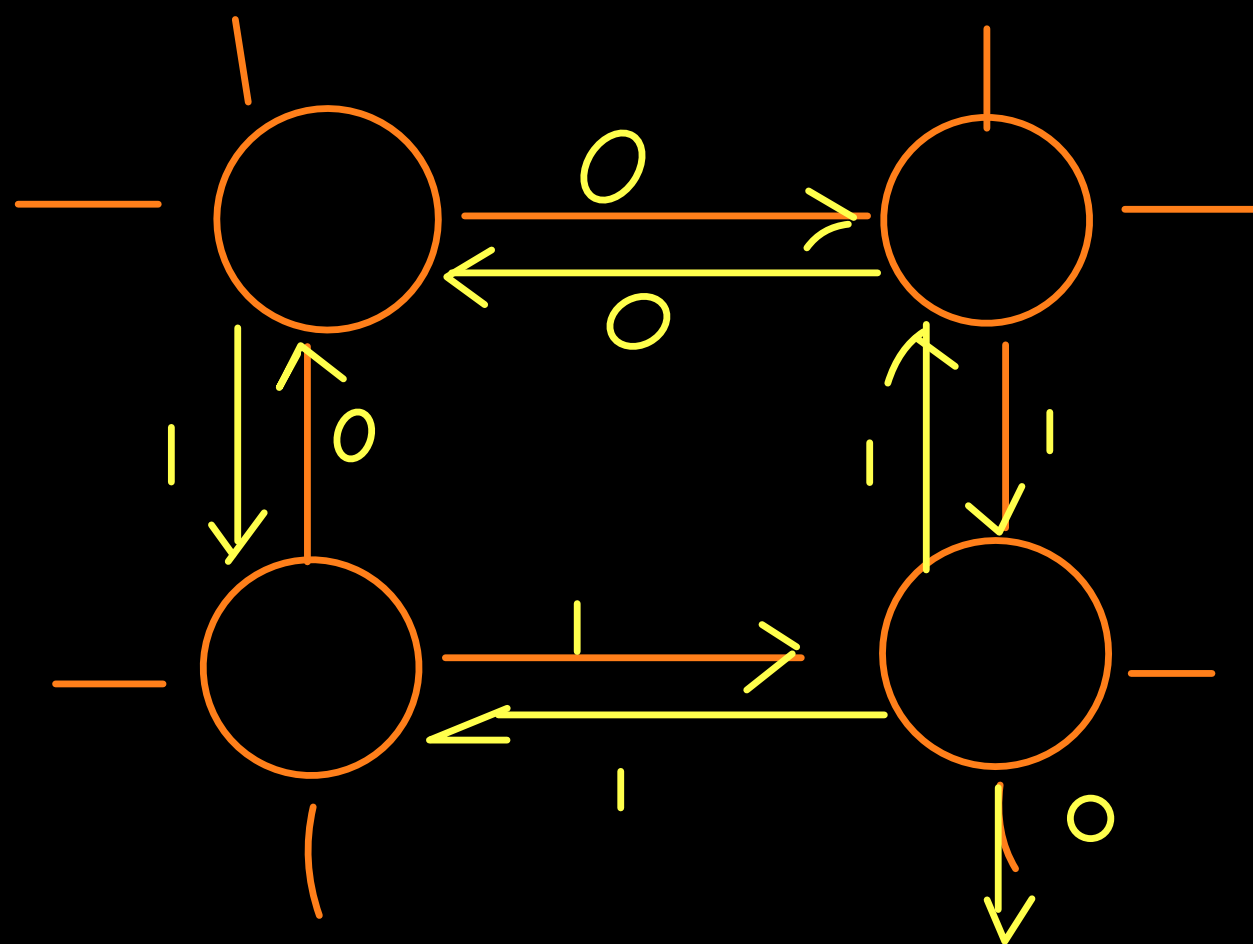
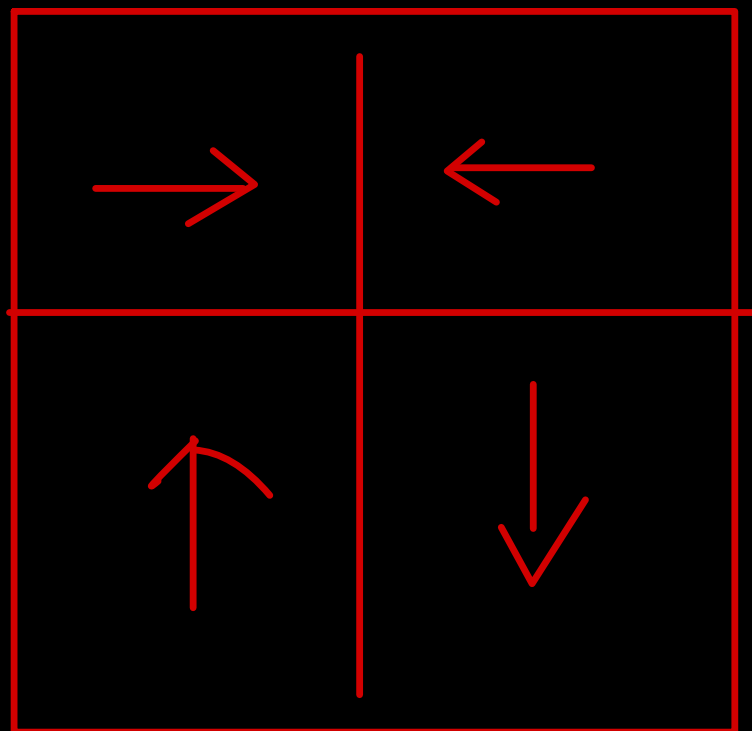
//

Graph

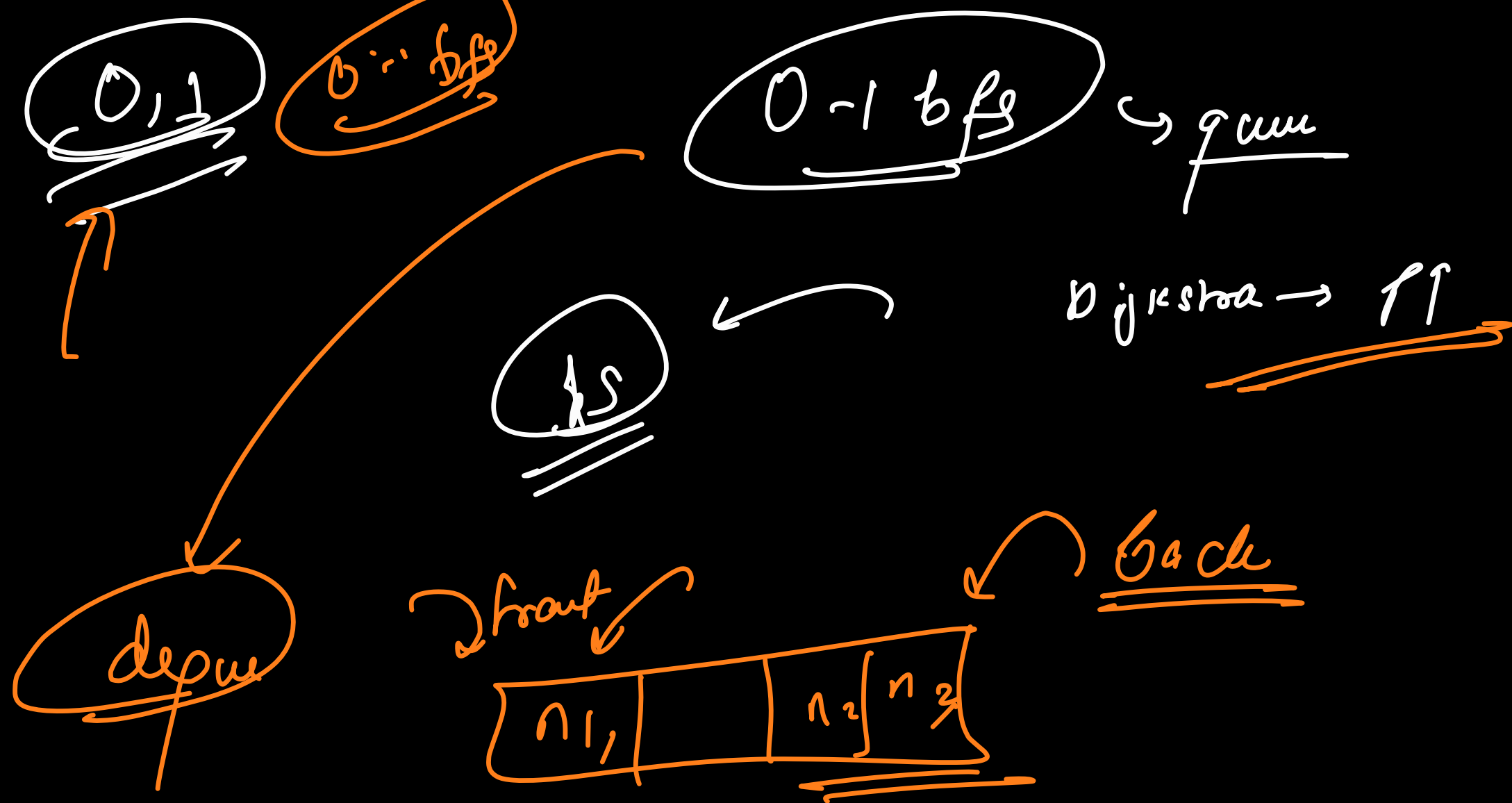
Shater job

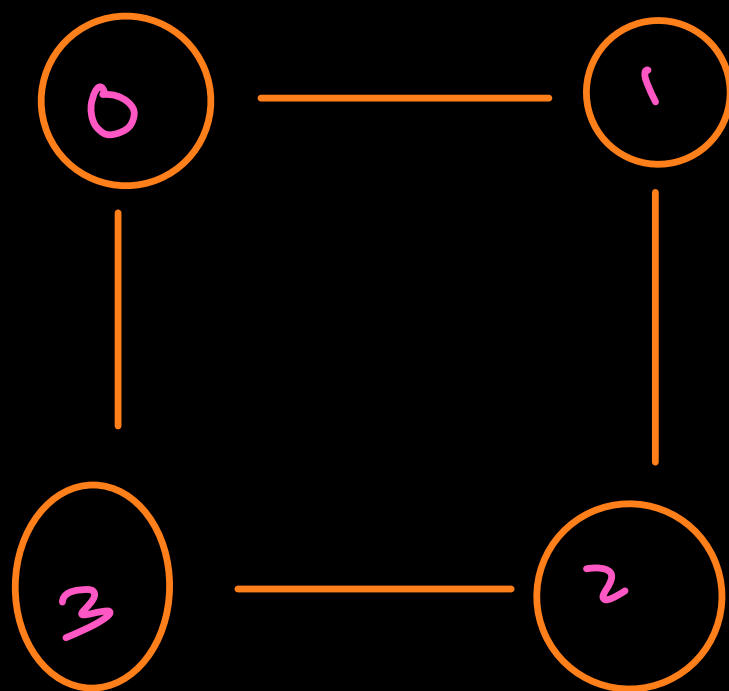
↓

Dijkstra



JOIN THE DARKSIDE





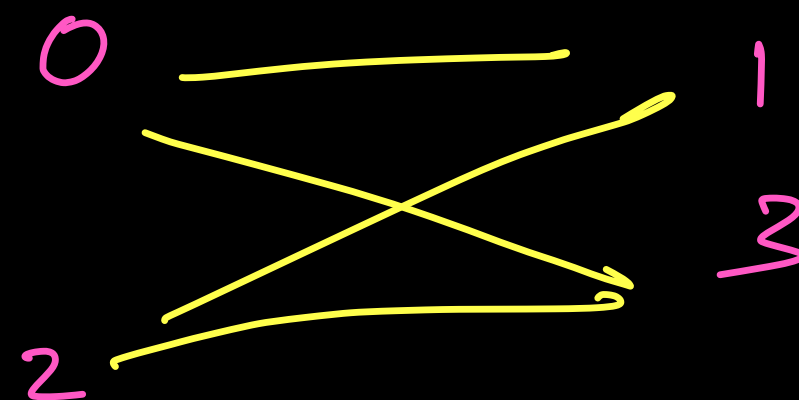
Set A

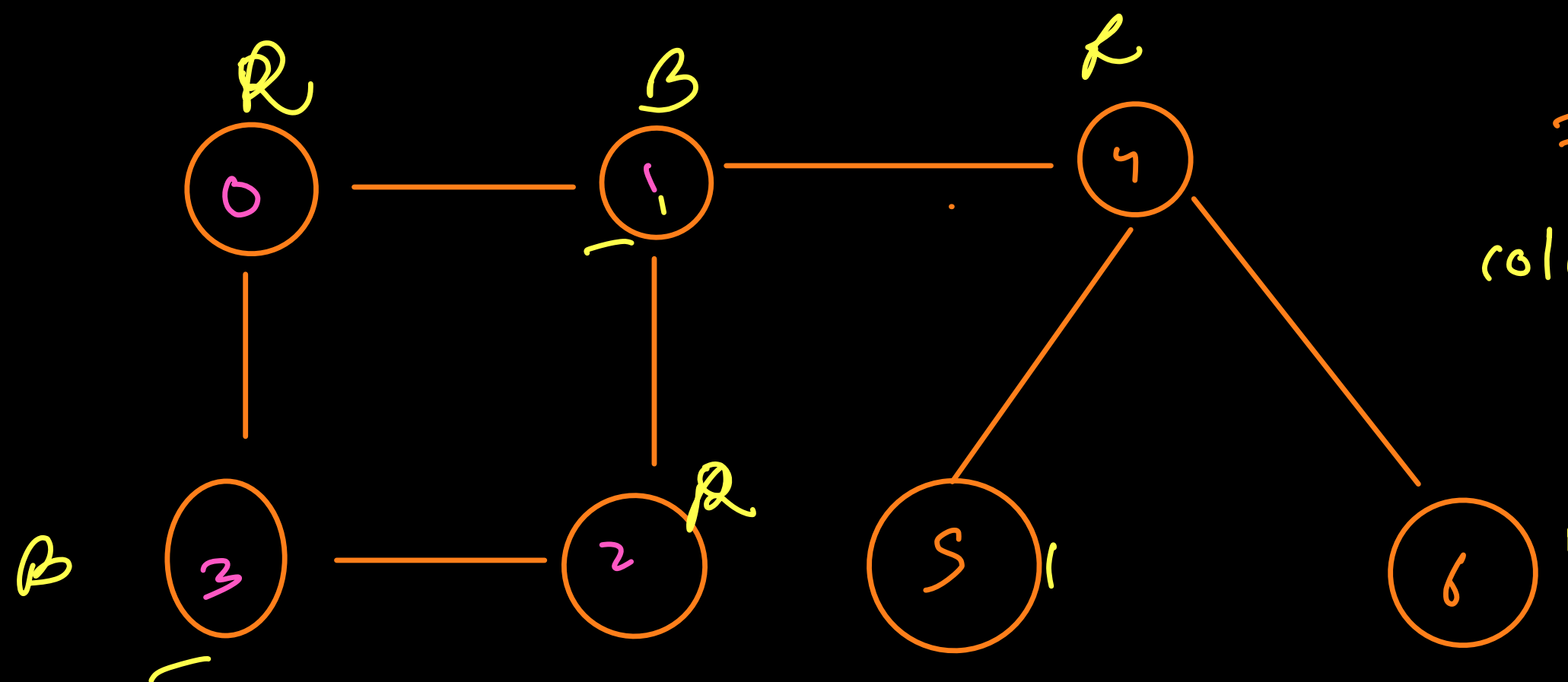
 0, 2

Set B

 1, 3

2-colors
 2 colors
 R, B





BFS
 $color[neigh] =$
 $1 \text{ if } color[con]$
 0

0 1 3 4 2 5 6

q

1	1	1	1	1	1	1
---	---	---	---	---	---	---

vis

0, 1, 3,

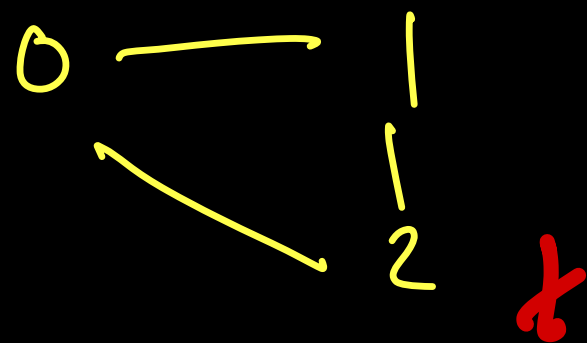
4, 2, 5, 6

0	1	2	3	4	5	6
0	1	0	1	0	1	1

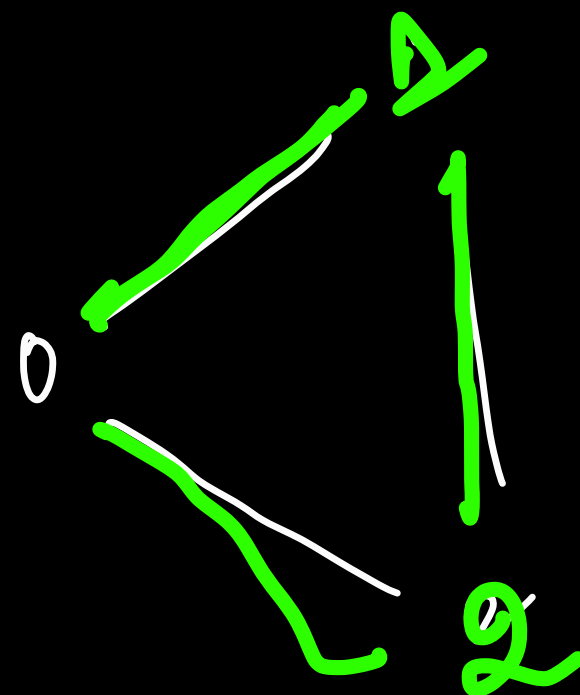
$color[i] = 1$
 Blue

$color[i] = 0$
 Green

color



1-2



func

3 things

↓

- ① —
- ② —
- ③ —

