



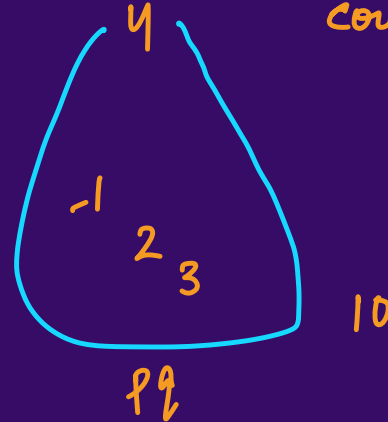
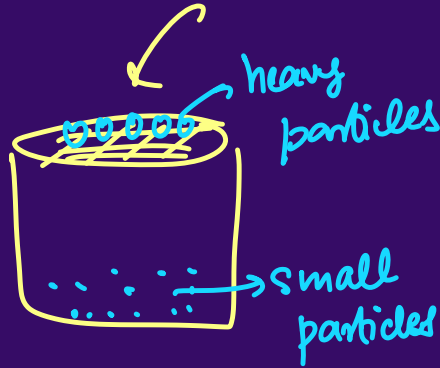
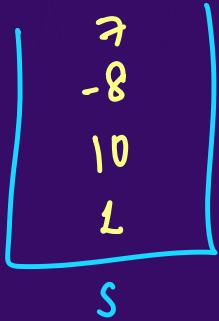
Priority Queues

What and Why?

BST \rightarrow searching $\rightarrow O(\log n)$
insertion \rightarrow

T.C

push, pop, top



pq.pop();
cout << pq.top();

What and Why?

T.C.

$\text{top}() \rightarrow O(1)$

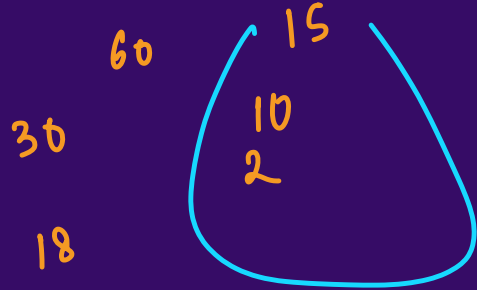
$\text{push}(x) \rightarrow O(\log n)$

$\text{pop}() \rightarrow O(\log n)$

current
size

If we want a DS in which we can always get the max^m or min^m element at any pt. of time then we use pqueues (heaps)

$n = \text{pq.size}();$



$\text{push}(10)$
 $\text{push}(2)$
 $\text{push}(18)$
 $\text{pop}()$
 $\text{push}(30)$
 $\text{push}(15)$

$\text{pop}()$
 $\text{top}() \rightarrow 15$
 $\text{push}(60)$
 pop

2 Types of heaps

max Heap (default)

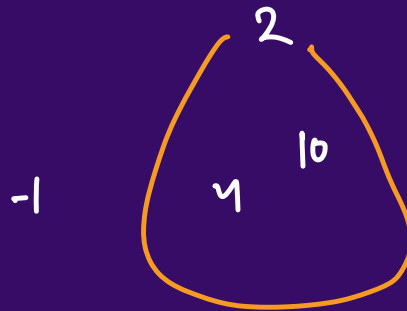


max ele is
on the top

min Heap



min element
is on top



```
push(10)
push(2)
push(4)
push(-1)
pop()
```

Priority Queue STL

include <queue>

maxHeap → priority_queue<int> pq

minHeap → priority_queue<int, vector<int>, greater<int>> pq

Problem Identification

- 1) K^{th} smallest, largest . Top k frequent elements, Closest k
- 2) At any pt. of time , minimum / max elements are required
- 3) Sorting (sometimes)

Ques:

Q1 : Find the kth smallest element in a given array.

arr = { 10, 20, -4, 6, 18, 24, 105, 118 } $k=3$

Method-1 : BIsort , arr[k-1] T.C. = $O(n \log n)$ S.C. = $O(\log n)$

Method-2 : Selection Sort \rightarrow T.C. = $O(k^*n)$ S.C. = $O(1)$

Method-3 : Quick Select \rightarrow T.C. = $O(n)$ [Not in the worst case]

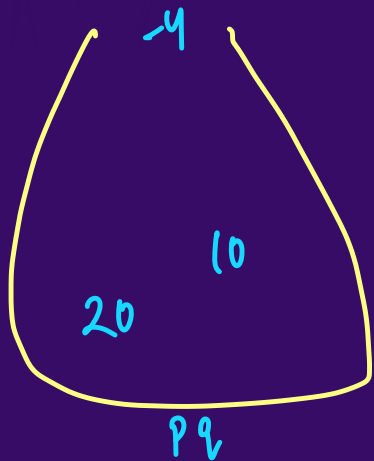
Method-4 : Using Heap \rightarrow T.C. = $O(n \cdot \log k)$

Ques:

Q1 : Find the kth smallest element in a given array.

minHeap

arr = { 10, 20, -4, 6, 18, 24, 105, 118 } $k=3$



```
for(int i=1; i<K; i++) pq.pop();  
cout << pq.top();
```

→ $K \cdot \log n$

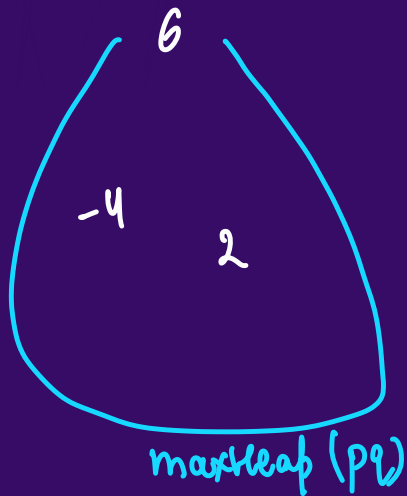
24 118 105 18
6

Ques:

maxHeap \rightarrow only k elements
not more than that
 \uparrow

Q1: Find the k th smallest element in a given array.

arr = { 10, 20, -4, 6, 18, 2, 105, 118 } $k=3$



20 18 10 105 118

Ques:

Q1 : Find the kth smallest element in a given array.

Time & Space

```
for(int i=0;i<n;i++){  
    pq.push(arr[i]); → log k  
    if(pq.size()>k) pq.pop(); → log k  
}
```

$$T.C. = O(n \log k)$$

$$S.C. = O(n) \text{ (total space)}$$

$$✓ A.S. = O(k)$$

If I insert 'n' elements in a heap,

$$\log(1) + \log(2) + \log(3) \dots \log(n)$$

$$\Rightarrow \log(n!) \xrightarrow{\sim} n \log n$$

Ques:

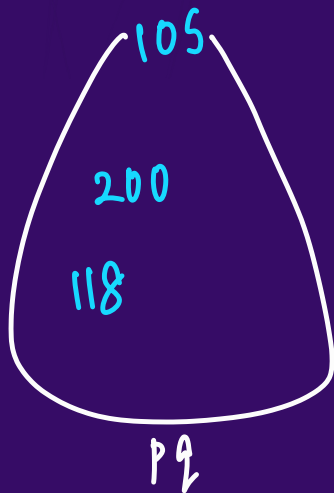
minHeap
↑

Q2: Find the kth largest element in a given array.

arr = { 10, 200, -4, 6, 18, 2, 105, 118 } k=3

$$T.C. = O(n \log k)$$

$$S.C. = O(k) \text{ (Extra)}$$



-4 6 2 10 18

Ques:

Q3 : Sort a 'k' sorted array (sort a nearly sorted array).

$\{ 4, 7, 8, 9, 10, 50, 60, 70 \}$
↑ sort

arr = $\{ 10, 9, 8, 7, 4, 70, 60, 50 \}$ $k=4$

arr = $\{ 6, 5, 3, 2, 8, 10, 9 \}$ $k=3$

↓ sort
 $\{ 2, 3, 5, 6, 8, 9, 10 \}$

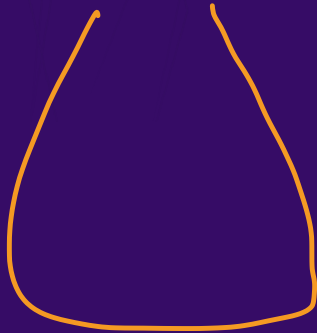
minHeap

Ques:

Q3 : Sort a 'k' sorted array (sort a nearly sorted array).

arr = { 6, 5, 3, 2, 8, 10, 9 } k=3

ans = { 2 3 5 6 8 9 10 3 }



pq (minHeap)
size $\leq k$

H.W. - leetcode 378

k^{th} smallest element in a
sorted Matrix

Ques:

pair ka usage

map

Q4 : Top K Frequent Elements

arr = { 1, 3, 2, 1, 1, 3 } K=2

ans = { 1, 3 }

sort
custom
comparator

map

1	→	3
2	→	1
3	→	2

arr = { 1, 1, 1, 3, 3, 2 }

Ques:

pair ka usage

map

Q4 : Top K Frequent Elements

K largest elements

arr = {1, 3, 2, 1, 1, 3} K=2

ans = {1, 3}

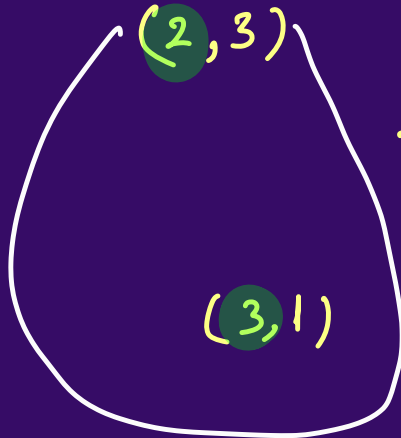
largest
[as in
with most
frequencies]

↓
minheap

map

1	→ 3
2	→ 1
3	→ 2

↓
pair < ele, freq >



→ pair < freq, ele >

(1, 2)

minheap [K] → 2

Homework:



Q : Sort Array by Increasing Frequency.

[Leetcode 1636]

Ques:

Q5: Find K Closest Elements

$$x = 3$$

$$\text{arr} = \{1, 2, 3, 4, 5, 6\}$$

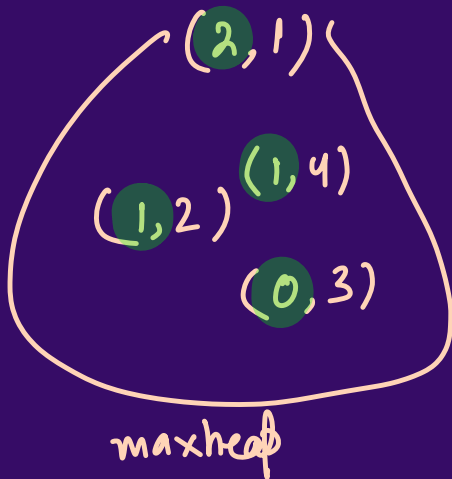
2 1 0 1 2 3

$$K = 4$$

K smallest
distances

(3, 6)

(2, 5)
VVIMP



heap \rightarrow $\langle \text{dist}, \text{ele} \rangle$

[Leetcode 658]

Ques:

Q6 : K Closest Points to Origin

arr = $\{ \underbrace{\{3, 3\}}_{\sqrt{18}}, \underbrace{\{5, -1\}}_{\sqrt{26}}, \underbrace{\{-2, 4\}}_{\sqrt{20}} \}$ $K=2$

K smallest
distances

maxheap

ans = $\{ \{3, 3\}, \{-2, 4\} \}$

heap \rightarrow

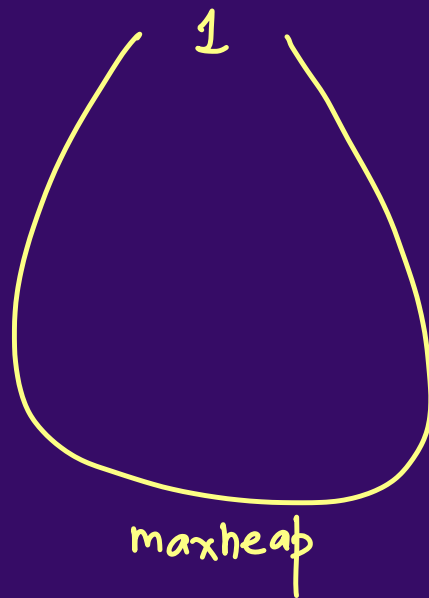
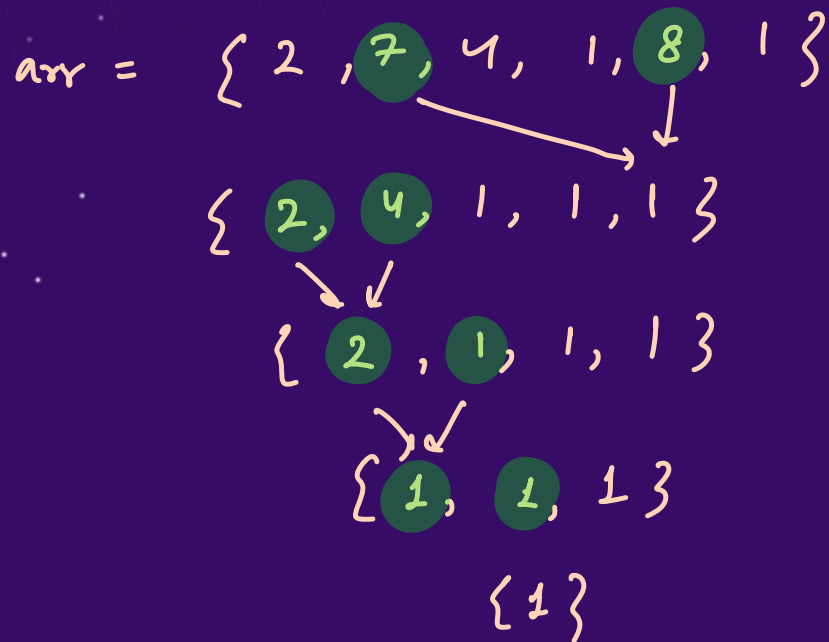
$\langle \text{dist}^2, \text{point} \rangle$

pair \langle int , $\underbrace{\text{pair} \langle \text{int}, \text{int} \rangle}_{\substack{\downarrow \\ \text{vector} \langle \text{int} \rangle}} \rangle$

$\{18, \{3, 3\}\}$

Ques:

Q7 : Last Stone Weight



$$x = 1$$

$$y = 1$$

$$x - y = 0$$

Ques:

Q7 : Last Stone Weight

arr = {2, 7, 4, 1, 8, 1}
sort

{1, 1, 2, 4, 7, 8}

{1, 1, 2, 4, 1}
sort

{1, 1, 1, 2, 4}

{1, 1, 1, 2}
sort

{1, 1, 1, 2}

{1, 1, 1} → sort

$n \log n +$

$(n-1) \log(n-1)$

$+ (n-2) \log(n-2)$

$\sum_{r=0}^{n-1} (n-r) \log(n-r)$

{1}

↑

{1, 1, 1}

[Leetcode 1046]

Ques: 'Very Interesting'

Q8: Minimum Cost to Connect all Ropes

arr = { 2, 7, 4, 1, 8 }

9, 4, 1, 8

13, 1, 8

14, 8

22

$$\text{cost} = 9 + 13 + 14 + 22 = 58$$

task is to join all the ropes with the min cost possible

You can connect only 2 ropes at a time with the cost being sum of length of those ropes

Ques:

Q8 : Minimum Cost to Connect all Ropes

$$\text{arr} = \{ 2, 7, 4, 1, 8 \}$$

$$\text{total} \rightarrow {}^nC_2 \times {}^{n-1}C_2 \times \dots {}^2C_2$$

$$6, 7, 1, 8$$

$$6, 8, 8$$

$$6, 16$$

$$22$$

$$\text{cost} = 6 + 8 + 16 + 22 = 52$$

Ques:

Q8 : Minimum Cost to Connect all Ropes

$$\text{arr} = \{2, 7, 4, 1, 8\}$$

$$3, 7, 4, 8$$

$$7, 7, 8$$

$$14, 8$$

$$\text{cost} = 3 + 7 + 14 + 22 = 46$$

Ques:

Q8 : Minimum Cost to Connect all Ropes

arr = { 6, 5, 3, 2, 8, 10, 9 }

6, 5, 5, 8, 10, 9

6, 10, 8, 10, 9

14, 10, 10, 9

14, 10, 19

24, 19

43

cost = 5 + 10 + 14 + 19 + 24 + 43

◀ **THANK YOU** ▶