743. Network Delay Time

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Problem Statement

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You are given a network of n nodes, labeled from 1 to n. You are also given times, a list of travel times as directed edges times[i] = (ui, vi, wi), where ui is the source node, vi is the target node, and wi is the time it takes for a signal to travel from source to target.

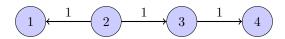
We will send a signal from a given node k. Return the minimum time it takes for all the n nodes to receive the signal. If it is impossible for all the n nodes to receive the signal, return -1.

Examples

Example 1

• Input: times = [[2,1,1],[2,3,1],[3,4,1]], n = 4, k = 2

• **Output**: 2



Iteration Table - Bellman-Ford Algorithm

Iteration	dist array
Initial	$[\infty,0,\infty,\infty,\infty]$
1	$[\infty,0,1,\infty,\infty]$
2	$[\infty,0,1,2,\infty]$
3	$[\infty, 0, 1, 2, 3]$

Iteration Table - Dijkstra's Algorithm

Iteration	dist array
Initial	$[\infty,0,\infty,\infty,\infty]$
Step 1	$[\infty,0,1,\infty,\infty]$
Step 2	$[\infty,0,1,2,\infty]$
Step 3	$[\infty, 0, 1, 2, 3]$

Example 2

• **Input**: times = [[1,2,1]], n = 2, k = 1

• Output: 1



Iteration Table - Bellman-Ford Algorithm

Iteration	dist array
Initial	$[\infty, \infty, 0]$
1	$[\infty,0,1]$

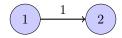
Iteration Table - Dijkstra's Algorithm

Iteration	dist array
Initial	$[\infty, \infty, 0]$
Step 1	$[\infty,0,1]$

Example 3

• Input: times = [[1,2,1]], n = 2, k = 2

• Output: -1



Iteration Table - Bellman-Ford Algorithm

Iteration	dist array
Initial	$[\infty, \infty, 0]$
1	$[\infty, \infty, 1]$

Iteration Table - Dijkstra's Algorithm

Iteration	dist array
Initial	$[\infty, \infty, 0]$
Step 1	$[\infty, \infty, 1]$

Approach for Dijkstra's Algorithm

To solve this problem, we can use Dijkstra's algorithm, which is suitable for finding the shortest paths from a single source node to all other nodes in a weighted graph.

Approach

- 1. **Graph Representation**: Represent the network using an adjacency list where each node points to a list of its neighbors and the corresponding travel times.
- 2. Dijkstra's Algorithm: Initialize a distance array 'dist' where 'dist[k] = 0' (distance to itself is zero) and 'dist[i] = infinity' for all other nodes initially. Use a priority queue (minheap) to continually extract the node with the smallest known distance from 'k' and relax (update) distances to its neighbors if a shorter path is found.

3. Result Calculation:

- After running Dijkstra's algorithm, check the 'dist' array:
 - If any node still has distance 'infinity', it means that node is unreachable from 'k', and hence it's impossible for all nodes to receive the signal. In such cases, return -1.
 - Otherwise, return the maximum value in the 'dist' array, which represents the minimum time it takes for the signal to reach all nodes.

Algorithm for Dijkstra's Algorithm

```
\overline{	ext{Algorithm}}
   Input: times: List[List[int]], n: int, k: int
   Output: Minimum time it takes for all n nodes to receive the signal, or -1 if impossible
   Build the graph from 'times' using an adjacency list;
   Initialize 'dist' array with infinity values, except 'dist[k] = 0';
   Use a priority queue (min-heap) to implement Dijkstra's algorithm starting from node k;
   while priority queue is not empty do
      Extract the node u with the smallest distance from the priority queue;
      Relax all edges from u to its neighbors in the graph;
   end
   if any node still has distance infinity then
     return-1;
   end
   return the maximum value in 'dist' array;
   end
              Algorithm 1: Network Delay Time using Dijkstra's Algorithm
```

Solution Code for Dijkstra's Algorithm

```
import heapq
   from collections import defaultdict
   def networkDelayTime(times, n, k):
      \# Step 1: Build the graph representation
      graph = defaultdict(list)
      for u, v, w in times:
         graph[u].append((v, w))
      \# Step 2: Dijkstra's algorithm to find the shortest paths from node k
      dist = \{node: float(' \setminus infty') \text{ for node in range}(1, n+1)\}
      dist[k] = 0
      min\_heap = [(0, k)] \# (distance, node)
       while min_heap:
         current\_dist, u = heapq.heappop(min\_heap)
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17
         if current\_dist > dist[u]:
             continue
19
20
          for v, weight in graph[u]:
             distance = current\_dist + weight
22
23
             if distance < dist[v]:
                 dist[v] = distance
24
                heapq.heappush(min_heap, (distance, v))
25
       # Step 3: Determine the maximum distance in dist array
27
      \max\_{dist} = \max(dist.values())
      if \max_{dist} == \frac{\text{float}(' \setminus \text{infty'})}{\text{infty'}}:
30
         return -1 # Some nodes are unreachable
         return max\_dist
```

Approach for Bellman-Ford Algorithm

To solve this problem, we can use the Bellman-Ford algorithm, which is suitable for finding shortest paths from a single source node to all other nodes in a graph with negative edge weights.

Approach

- 1. **Initialize**: Create an array 'dist' with size n+1 where $\operatorname{dist}[k]=0$ and $\operatorname{dist}[i]=\infty$ for all other nodes initially. This array will hold the minimum distance from node k to all other nodes.
- 2. **Relaxation**: Relax edges repeatedly. For each edge (u, v, w) in 'times', if dist[u] + w < dist[v], update dist[v] to dist[u] + w.
- 3. **Detect Negative Cycles**: After n-1 iterations, the 'dist' array will contain the minimum distances from k to all nodes reachable from k. If there's any further improvement in the n-th iteration, it indicates the presence of a negative cycle.
- 4. Check Unreachable Nodes: If there's no negative cycle and any node still has distance 'infinity', return -1 (some nodes are unreachable). Otherwise, return the maximum value in the 'dist' array, which represents the minimum time it takes for the signal to reach all nodes.

Algorithm for Bellman-Ford Algorithm

```
Algorithm
   Input: times: List[List[int]], n: int, k: int
   Output: Minimum time it takes for all n nodes to receive the signal, or -1 if impossible
   Initialize 'dist' array with size n+1 with all elements set to \infty, except 'dist[k] = 0';
   for i from 1 to n-1 do
       for each edge (u, v, w) in 'times' do
          if dist[u] + w < dist[v] then
           |\operatorname{dist}[v] = \operatorname{dist}[u] + w;
          end
      end
   end
   if any node still has distance \infty then
    return-1;
   end
   else
    return the maximum value in 'dist' array;
             Algorithm 2: Network Delay Time using Bellman-Ford Algorithm
```

Solution Code using Bellman-Ford Algorithm

```
import sys

def networkDelayTime(times, n, k):

# Step 1: Initialize distances

infty = sys.maxsize

dist = [\infty] * (n + 1)

dist[k] = 0

# Step 2: Relax edges repeatedly

for _ in range(n - 1):
```

```
for u, v, w in times:

if dist[u] != \infty and dist[u] + w < dist[v]:

dist[v] = dist[u] + w

# Step 3: Check for negative cycles
for u, v, w in times:

if dist[u] != \infty and dist[u] + w < dist[v]:

return -1

# Step 4: Return the maximum value in dist array
max_time = max(dist[1:])

return max_time if max_time != \infty else -1
```