

# PRINT Cipher

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**Abstract.** In this paper we prove that the One-Time-Pad has perfect security.

**Keywords:** Something · Something else

## 1 Introduction

Widely used primitives like the AES [?] do not have perfect security, and can be analysed with linear cryptanalysis [?], differential cryptanalysis [?], or differential power analysis [?]. We show that the One-Time-Pad is unconditionally secure in [Section 2](#).

## 2 Main Result

### 2.1 Sbox Analysis

The sbox for the PRINT cipher is a 3-bit to 3-bit. Since input is 3-bit so for a b-bit block, the sbox is applied  $\frac{b}{3}$  parallelly. The current state for the sbox is a  $\frac{b}{3}$  words, for each word same sbox is used and the next state is the concatenation of outputs. It is a balanced sbox and has a linear structure. The sbox is given in the following table :-

x	0	1	2	3	4	5	6	7
s[x]	0	1	3	6	7	4	5	2

#### 2.1.1 Difference Distribution Table

The sbox has a differential branch number defined as  $\min_{v, w \neq v} \{wt(v \oplus w) + wt(S(v) \oplus S(w))\}$  of 2. The difference distribution table (ddt) which is generated using Sage is as follows :-

	0	1	2	3	4	5	6	7
0	0	0	0	0	0	0	0	0
1	0	2	0	2	0	2	0	2
2	0	0	2	2	0	0	2	2
3	0	2	2	0	0	2	2	0
4	0	0	0	0	2	2	2	2
5	0	2	0	2	2	0	2	0
6	0	0	2	2	2	2	0	0
7	0	2	2	0	2	0	0	2

### 2.1.2 Linear Approximation Table

The linear branch number which is defined as  $\min_{\alpha \neq \beta, \text{LAM}(\alpha, \beta) \neq 0} \{\text{wt}(\alpha) + \text{wt}(\beta)\}$  for this sbox is **2**. The linearity of this sbox is **4**. The linear approximation table generated from Sage is as follows:-

	0	1	2	3	4	5	6	7
0	4	0	0	0	0	0	0	0
1	0	-2	0	2	0	2	0	2
2	0	0	2	2	0	0	2	-2
3	0	2	-2	0	0	2	2	0
4	0	0	0	0	2	-2	2	2
5	0	2	0	2	2	0	-2	0
6	0	0	2	-2	2	2	0	0
7	0	2	2	0	-2	0	0	2

### 2.1.3 Additional Properties of Sbox

1. The component function in 3 variables in algebraic normal form of the sbox is

$$x_0 * x_2 + x_0 + x_1 * x_2$$

2. The interpolation polynomial for the sbox is

$$(a + 1)x^6 + (a^2 + a + 1)x^5 + (a^2 + 1)x^3$$

3. The polynomials which satisfy the sbox is

- $x_0 * x_2 + x_0 + x_1 + y_1$
- $x_0 * x_1 + x_0 + x_1 + x_2 + y_2$
- $x_0 * y_1 + x_0 + x_2 + y_1 + y_2$
- $x_0 * y_2 + x_1 + y_1$

- $x_1 * x_2 + x_0 + y_0$
- $x_1 * y_0 + x_1 + x_2 + y_0 + y_2$
- $x_0 * y_0 + x_1 * y_1 + x_2 + y_2$
- $x_1 * y_2 + x_0 + x_1 + y_0$
- $x_2 * y_0 + x_1 + y_0 + y_1$
- $x_2 * y_1 + x_0 + y_0$
- $x_0 * y_0 + x_2 * y_2 + x_0 + x_1 + x_2 + y_0 + y_1$
- $y_0 * y_1 + x_2 + y_0 + y_1 + y_2$
- $y_0 * y_2 + x_1 + y_1$
- $y_1 * y_2 + x_0 + y_0 + y_1$

x - input variables y - output variables

4. Maximum degree of component function - 2
5. Minimum degree of component function - 2
6. Maximal differential probability - 0.25
7. Absolute maximal linear bias - 2
8. Relative maximal linear bias - 0.25

## References