



**IK2204**

# **Layer 2 Issues**

Bridges, Hubs, Switches, VLANs,  
Spanning Tree...

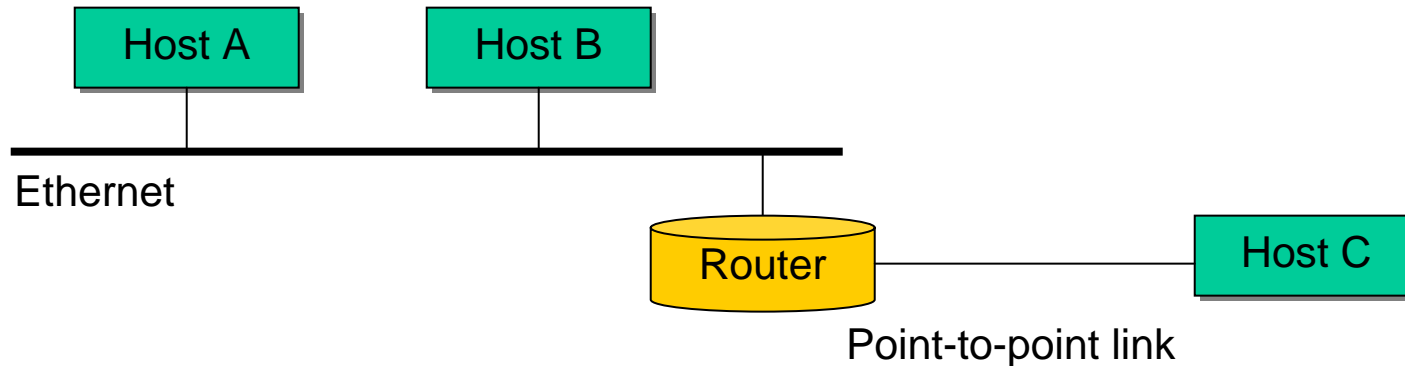
# Contents

- Perlman, chapter 3 and 5
- Video lectures no. 2-11
- And then some.....

# Different Link Types

- Point-to-point
  - A link that connects exactly two interfaces.
  - Dial-up, V.35, etc.
- Multicast (broadcast)
  - a link that supports a native mechanism at the link layer for sending packets to all (i.e., broadcast) or a subset of all neighbors.
  - Ethernet, FDDI, Wireless, etc.
- Non-Broadcast Multiple Access
  - More than two interfaces can attach, but no native form of multicast or broadcast.
  - ATM, X.25, Frame Relay, etc.

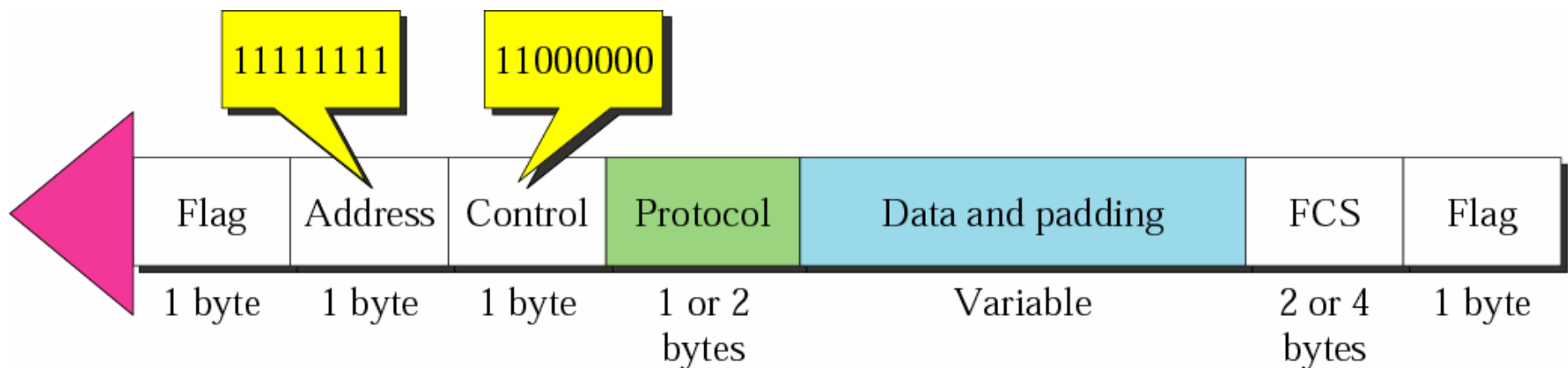
# Point-to-Point Links



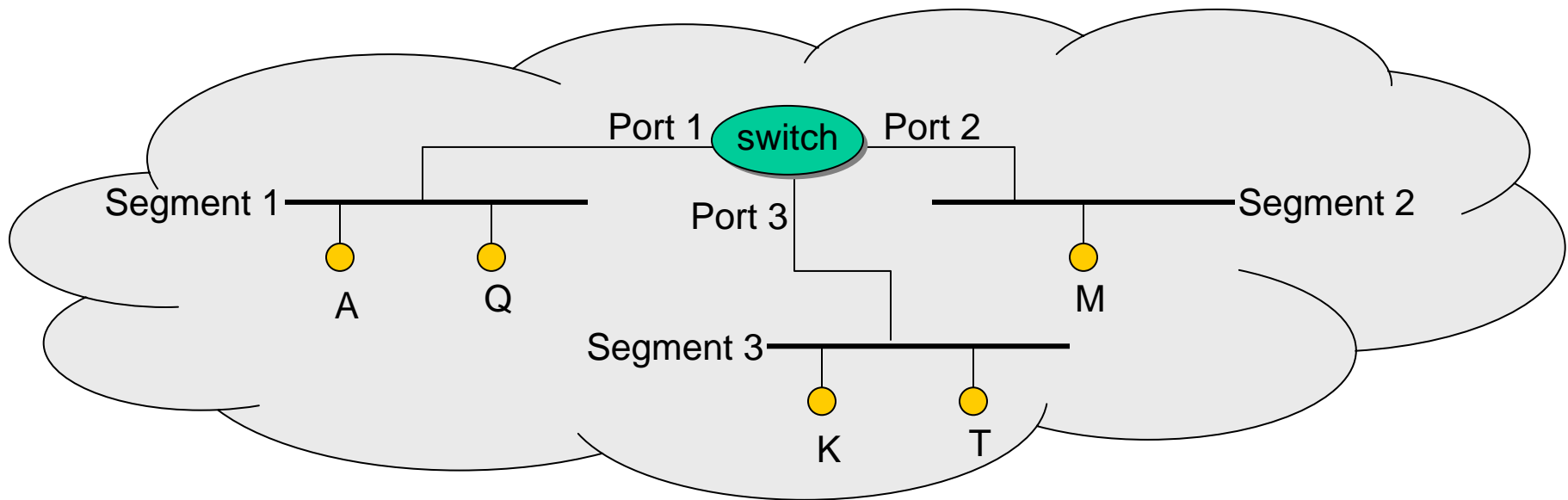
- Point-to-point link connects exactly 2 interfaces
- Widespread point-to-point protocols:
  - SLIP (CSLIP) for transmission of IP datagrams over serial lines
  - PPP (Point-to-Point Protocol)

# PPP (RFC 1331, 1332)

- Link Control Protocol (LCP) to establish, configure, and test data-links (includes option negotiation)
- Network Control Protocol (NCP) to carry data, e.g., IP (RFC 1332)
- CRC check on every frame
- Dynamic negotiation of IP address of each end



# Ethernet LAN



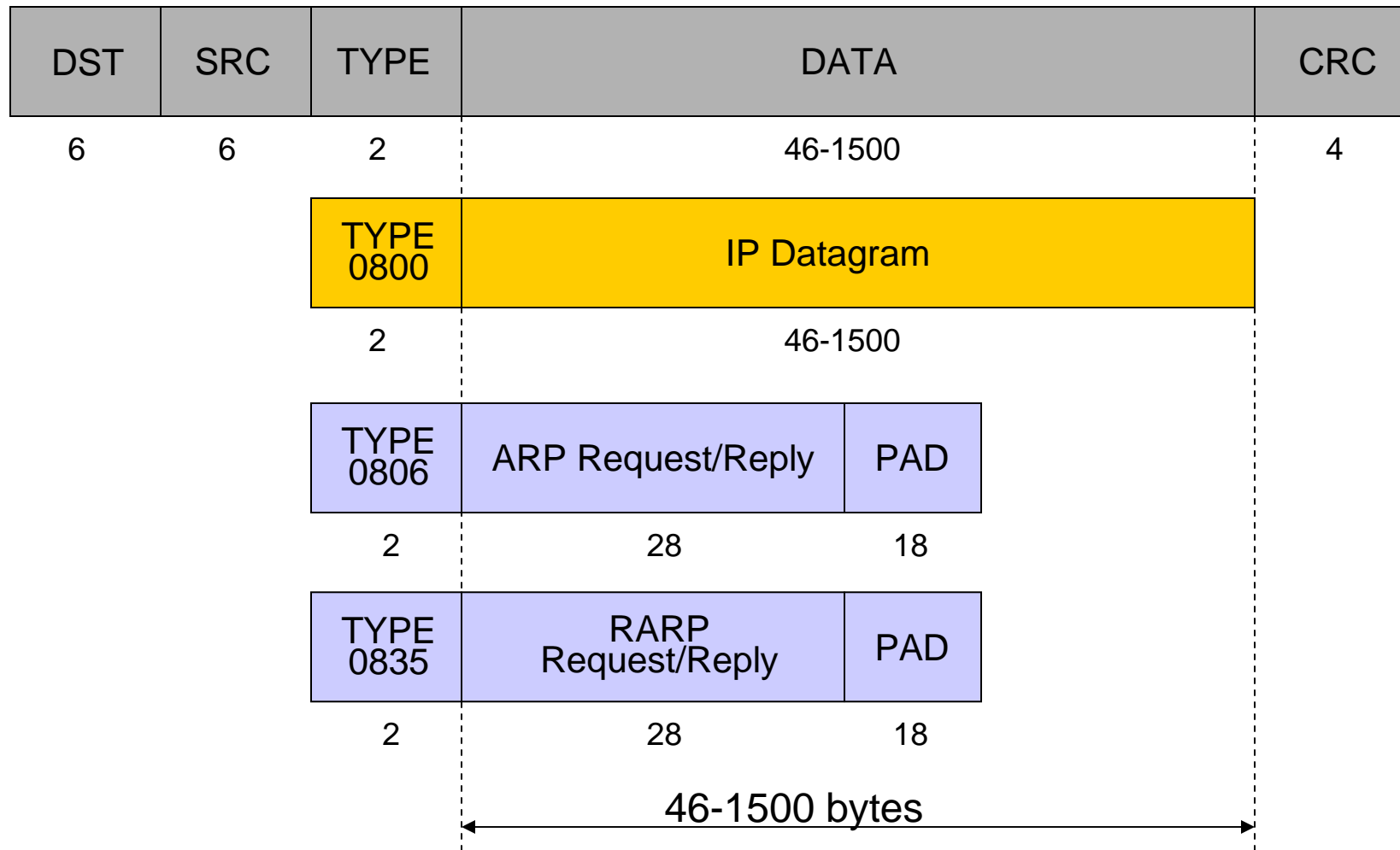
- Ethernet LAN normally built with *switches* and *hubs*
- Nodes attached to a hub are on the same *segment* (L1)
  - No hubs shown in figure
- Ethernet is a broadcast link technology
- A LAN is a *broadcast domain*

# Ethernet Frame (DIX)



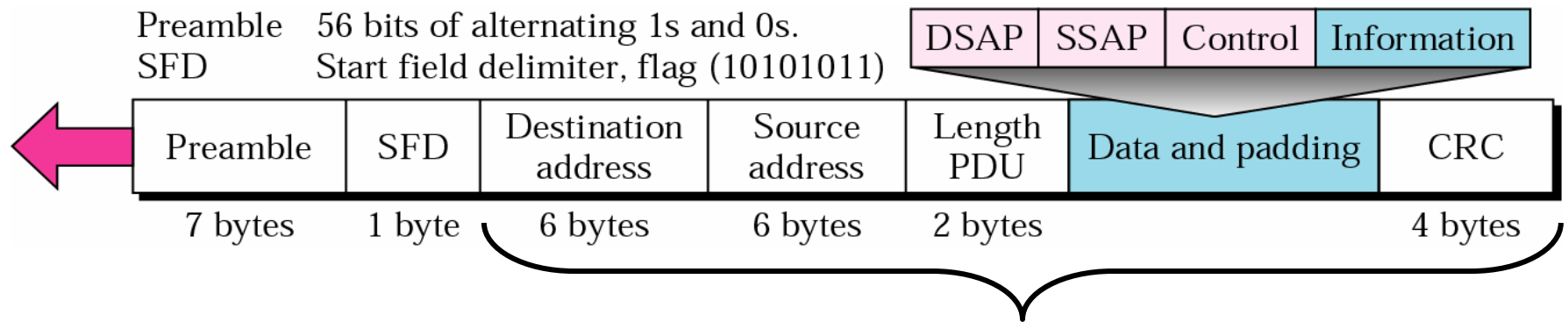
- DIX: Digital, Intel, Xerox
- DST: Destination Address, 48-bit MAC address
- SRC: Source Address, 48-bit MAC address
- Type: Type of data carried, e.g., IP packet
- Data: min size is 46 bytes, max size is 1500 bytes
- CRC: Cyclic Redundancy Check (or FCS – Frame Check Sum)
- 48-bit MAC address is normally written 12:34:56:78:9a:bc
  - Multicast (8th bit is 1), unicast, or broadcast (all bits are 1)

# Ethernet Encapsulation (RFC 894)





# IEEE 802.3 MAC Frame Format

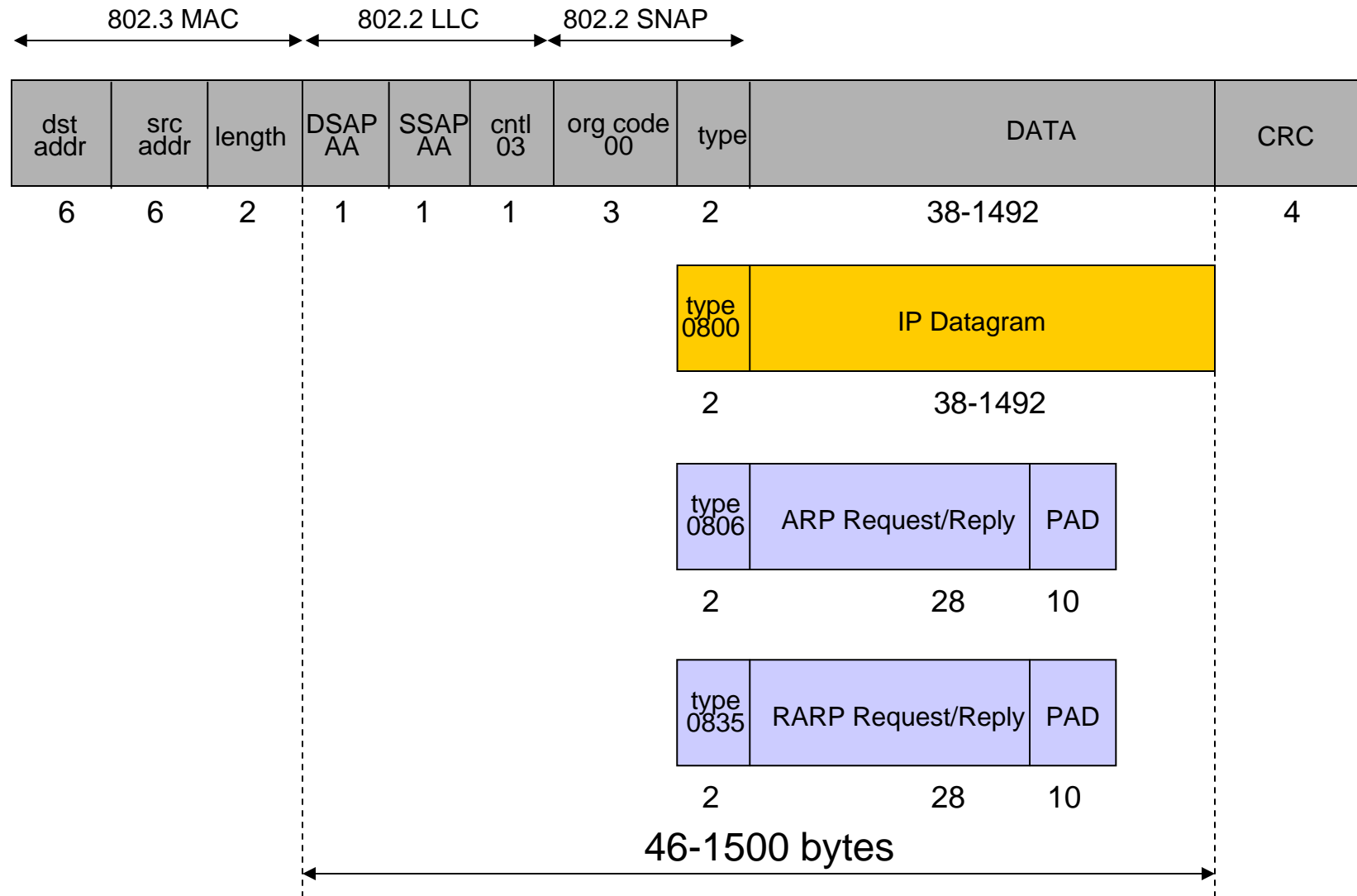


Data link layer frame

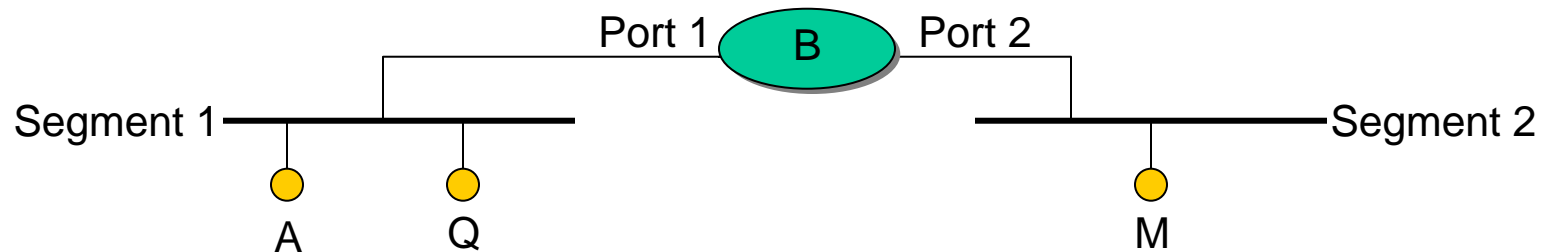
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- 18 bytes of overhead at link layer (26 bytes on the wire)
- Length/PDU
  - Length if less than 1518
    - IEEE 802.3 format
  - Otherwise PDU type
    - RFC 894

# IEEE 802.3 Encapsulation (RFC 1042)



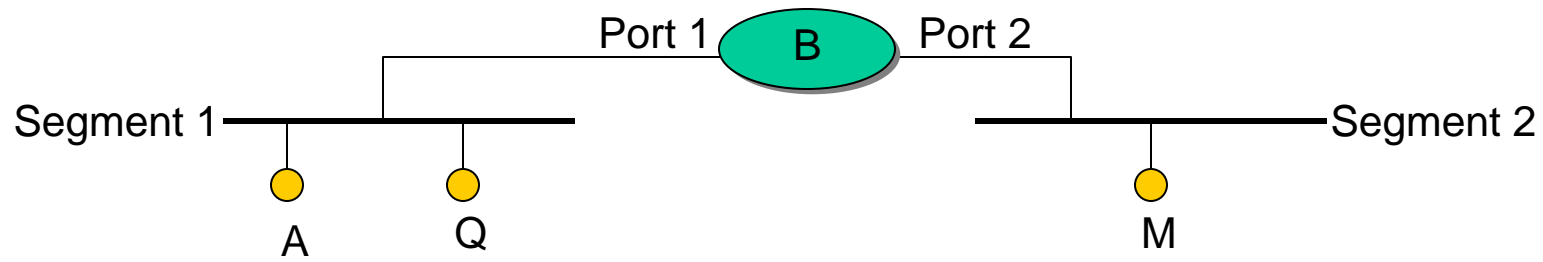
# Bridges



- A bridge operates on the link layer—buffers frames
- A bridge does not modify the structure or content of a frame
- A bridge is transparent, i.e., invisible to the clients
- A *multiport* bridge can connect many segments—a *switch*
- A bridge has a *forwarding table*—to map addresses to ports

Address	Port
A	1
Q	1
M	2

# Standard Bridges



Standard bridge compared to a "no-frills" bridge:

- Learn which stations that belong to each port
  - Automatically
- Avoid loops
  - Automatically through the Spanning Tree Protocol

To conform to the 802.1D standard, a bridge must implement

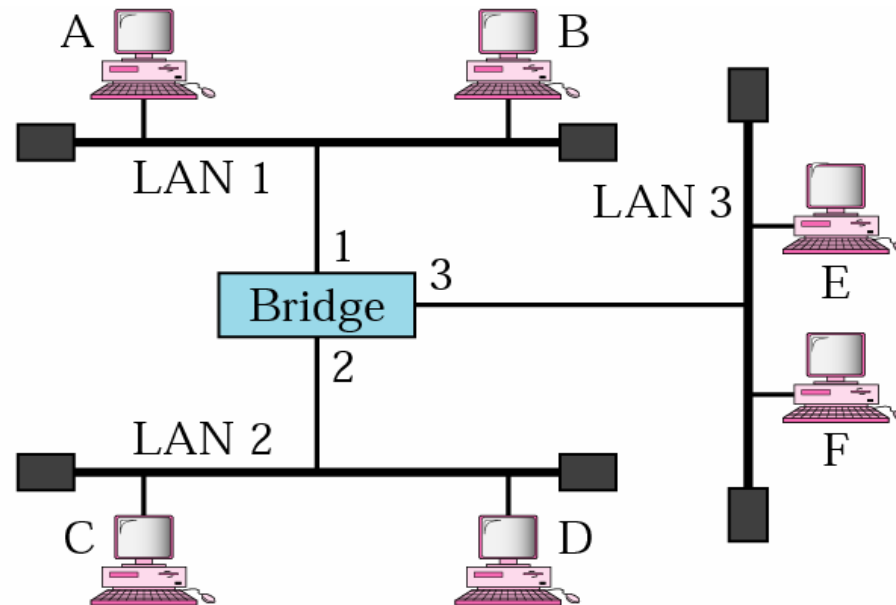
- Learning
- Spanning tree

# Learning Bridges

Strategy of a learning bridge:

1. Listen promiscuously, receiving every packet
2. Store the source address of every packet in a forwarding table (or *learning table*)
3. For each packet, check destination address with learning table
  - If not found, *flood*: forward on all (except receiving) ports
  - If addr found, forward only on port specified in learning table.  
If specified port == receiving port, drop packet
4. Bridge ages each entry in the learning table.
  - Deletes entry if no traffic is received from that source after a certain period of time

# Learning Bridges Example



Address	Port

a. Original

Address	Port
A	1

b. After A sends a frame to D

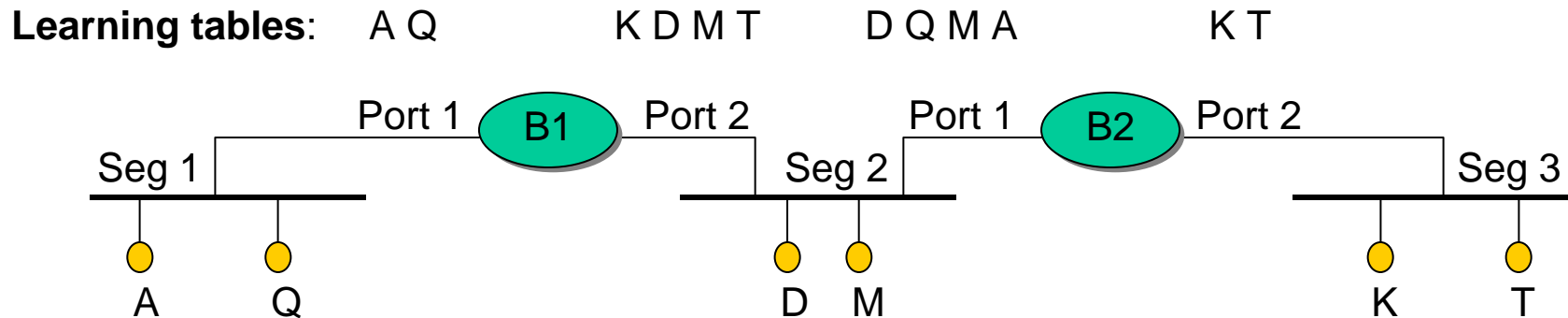
Address	Port
A	1
E	3

c. After E sends a frame to A

Address	Port
A	1
E	3
B	1

d. After B sends a frame to C

# Learning Bridges—Observations



B1 cannot distinguish between Segment 2 and Segment 3

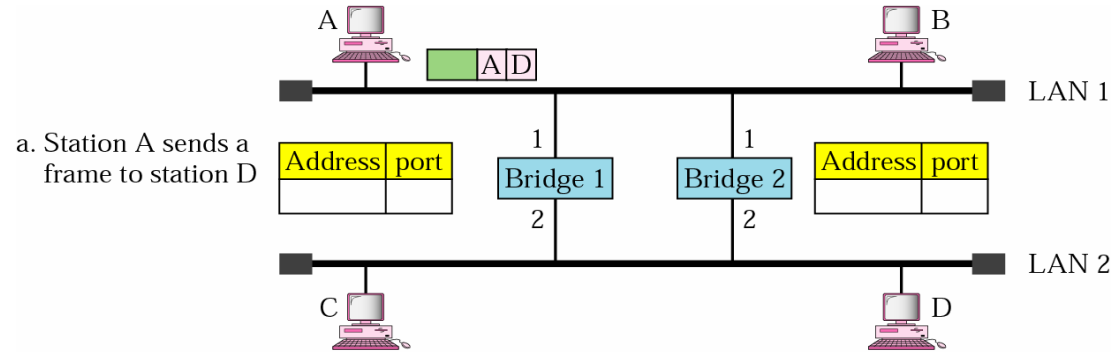
B2 cannot distinguish between Segment 1 and Segment 2

Bridges are not aware of the existence of other bridges

However,

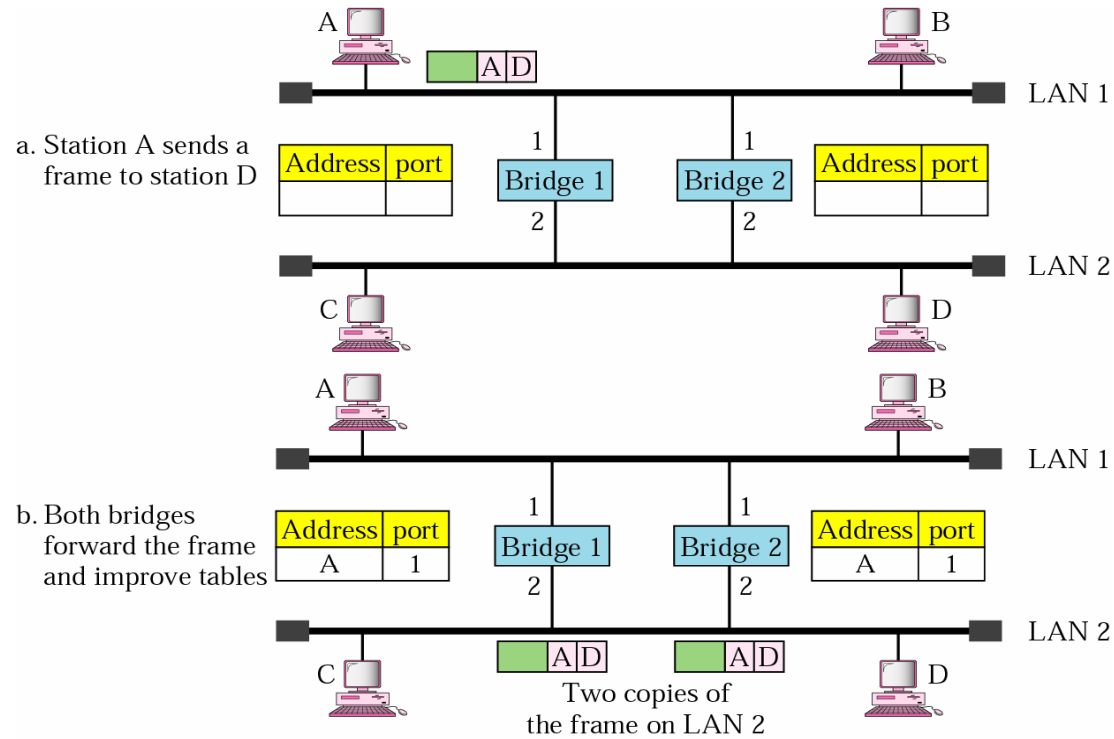
- Bridge concept works for any number of ports
- Bridge concept works for any *tree* (loop-free) topology

# Learning Bridges—Loop Problem

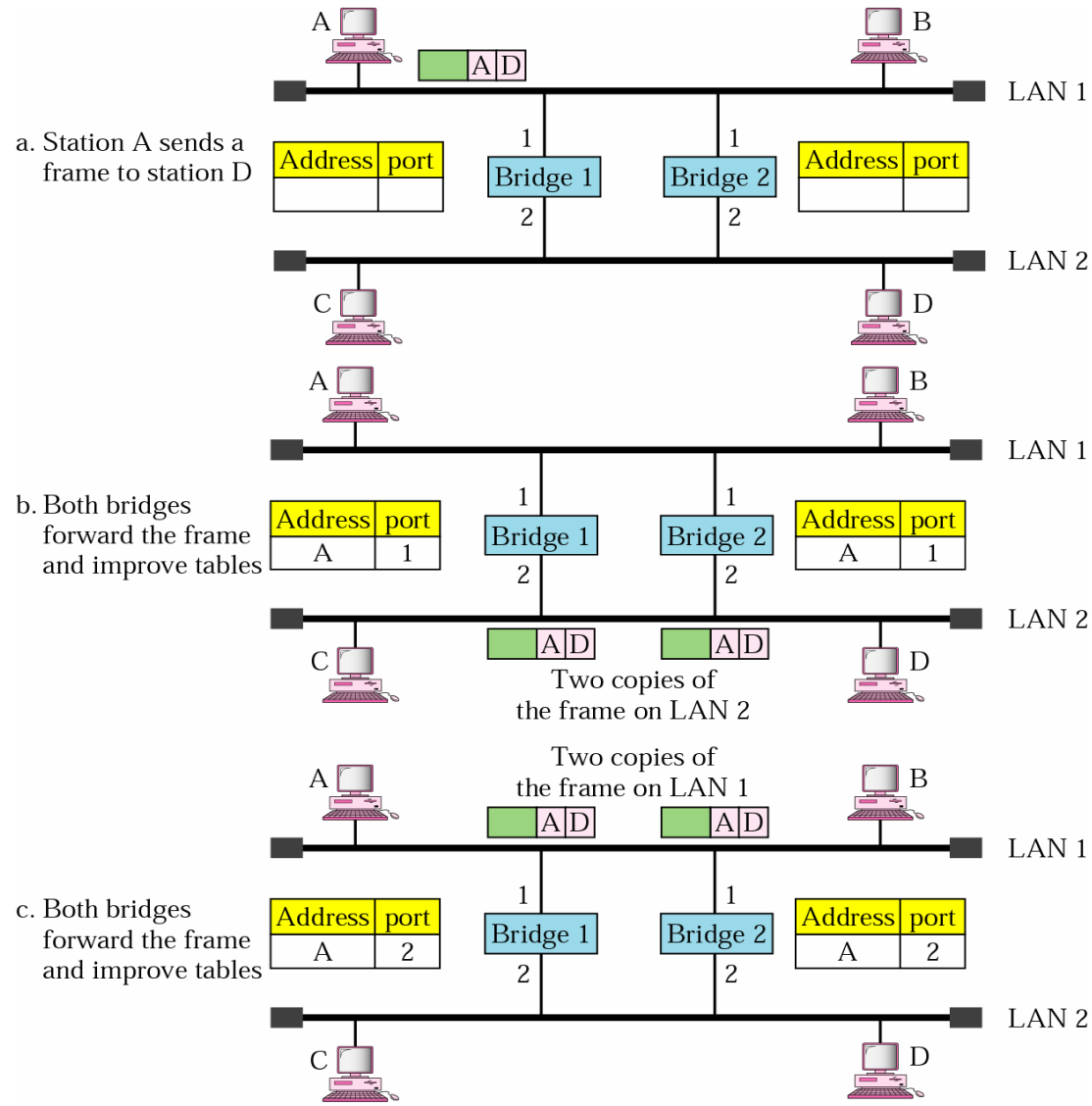




# Learning Bridges—Loop Problem

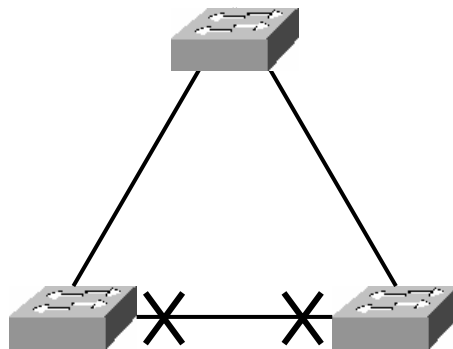


# Learning Bridges—Loop Problem



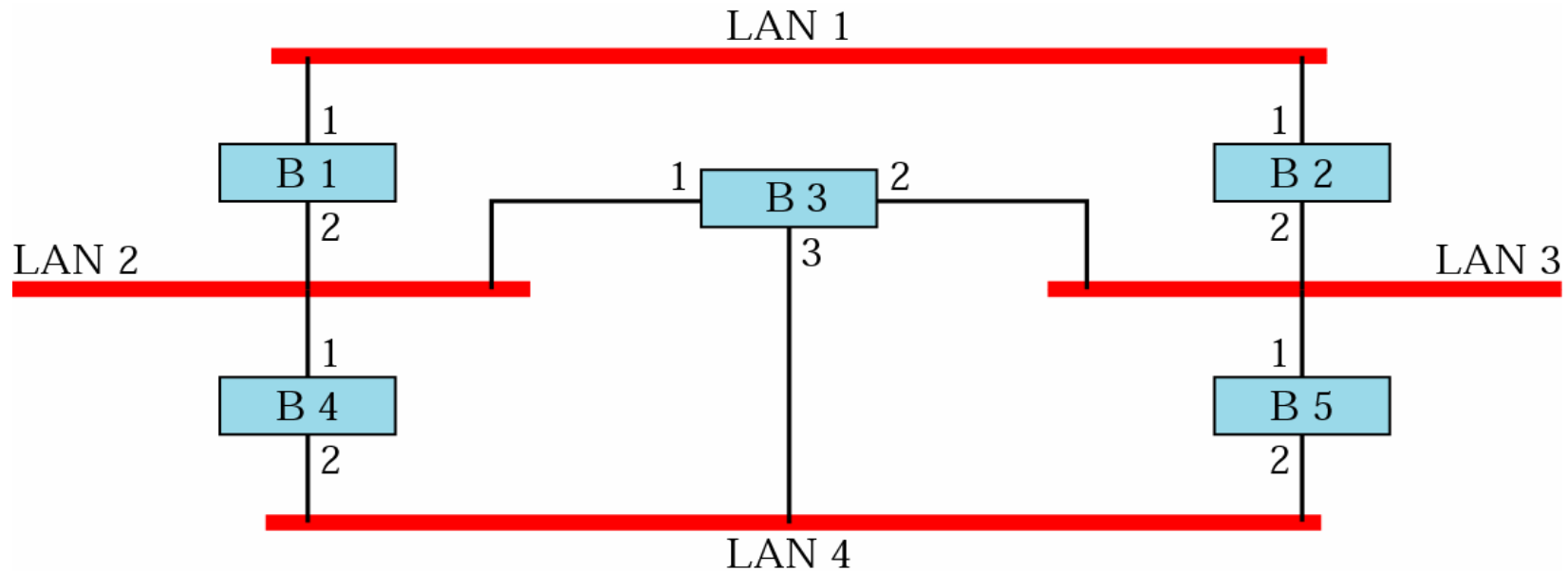
# Spanning Tree—Basic Idea

- Discover a subset of the topology that is loop-free (a tree)
- Just enough connectivity so that there is a path between every pair of segments where physically possible
  - The tree is spanning
- Disable (block) all other ports
- All bridges exchange configuration messages, called *bridge protocol data units* (BPDUs), allowing them to calculate a spanning tree



# Spanning Tree Starting Point

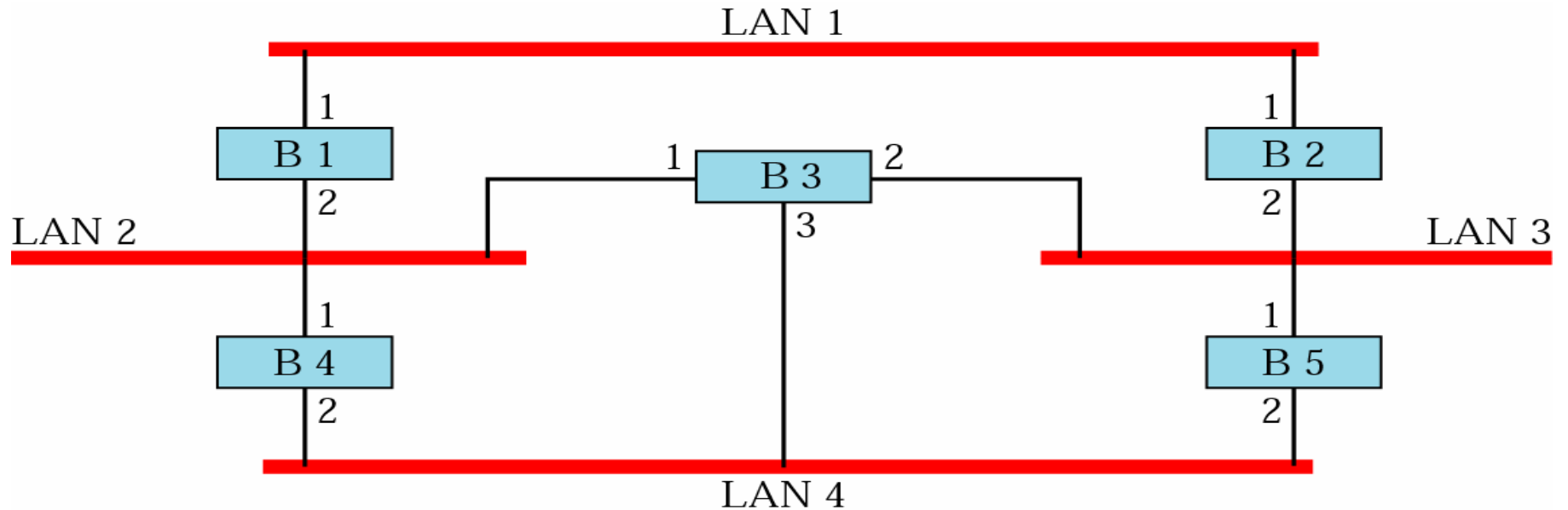
- Each bridge has a unique ID
- Each port has a unique ID within the bridge
- A *cost* can be calculated for each path between two bridges
  - A default cost value is used for each port (can be configured)



# Spanning Tree Process

1. The node with the smallest ID is selected the root bridge
2. On each bridge, select a root port
  - Port with the least cost path to the root bridge
3. On each LAN segment, select a designated bridge
  - Bridge with least cost path to root bridge
    - If two bridges have same cost, select the bridge with smallest ID
  - Mark the corresponding port as the designated port
4. Forward frames only on marked ports
  - Designated ports and root ports
  - Block on the others

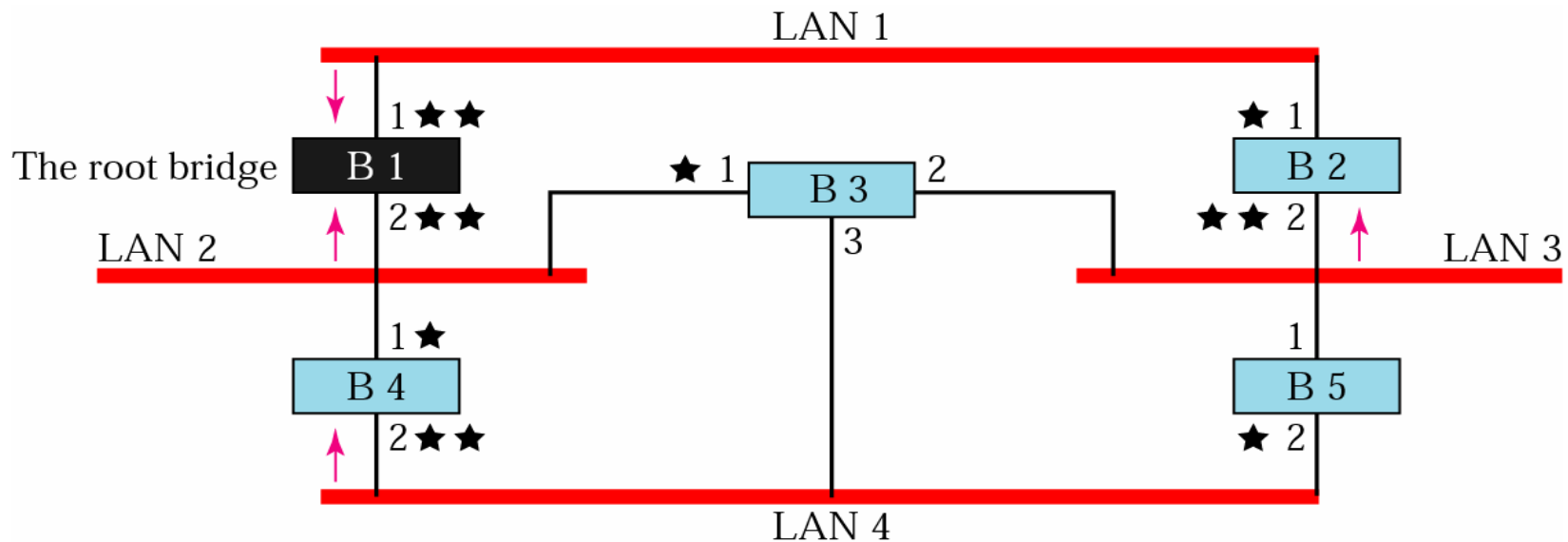
# Before Spanning Tree



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- B1 has the smallest ID
  - then assume  $B4 < B2 < B3 < B5$

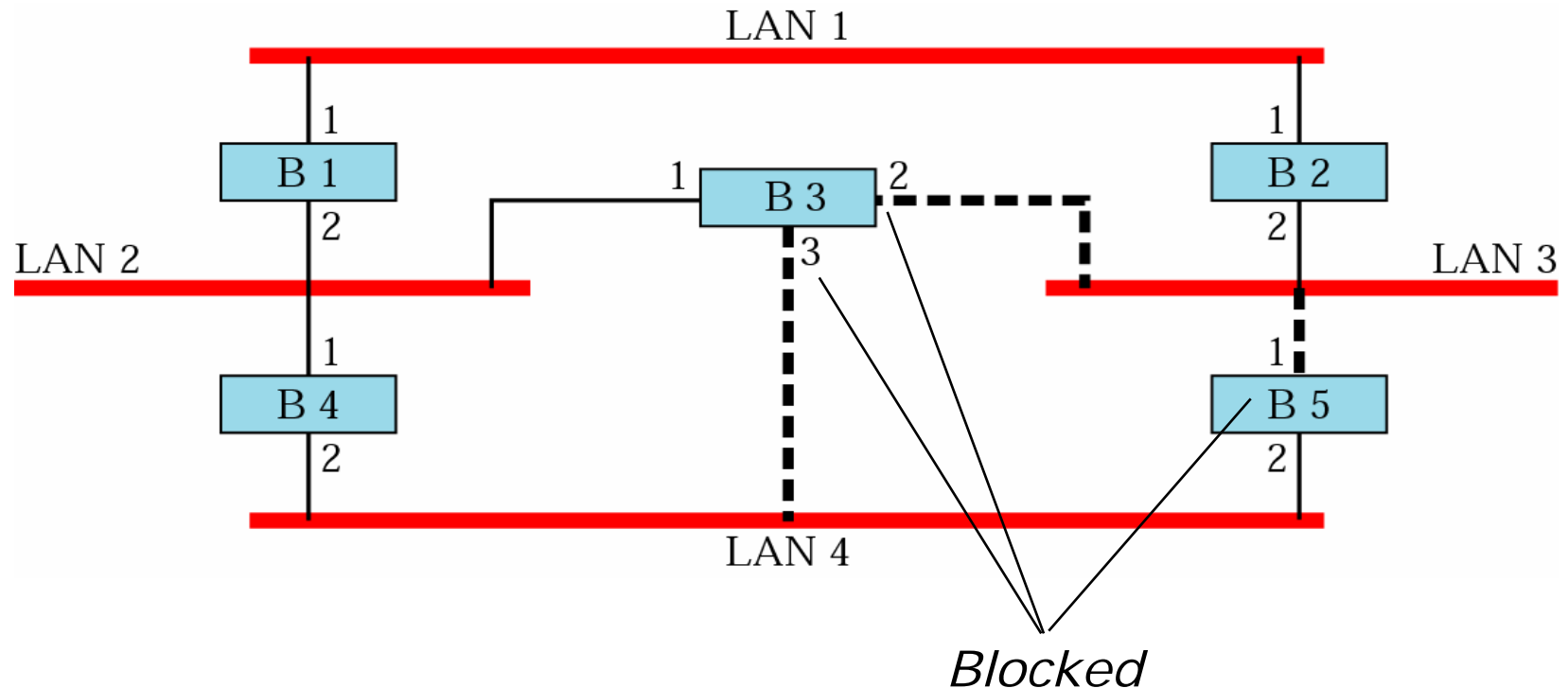
# Applying Spanning Tree



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- B1 has the smallest ID (then B4, B2, B3, B5)
- \* \* means designated bridge for a segment
- \* means root port

# Forwarding Ports and Blocking Ports



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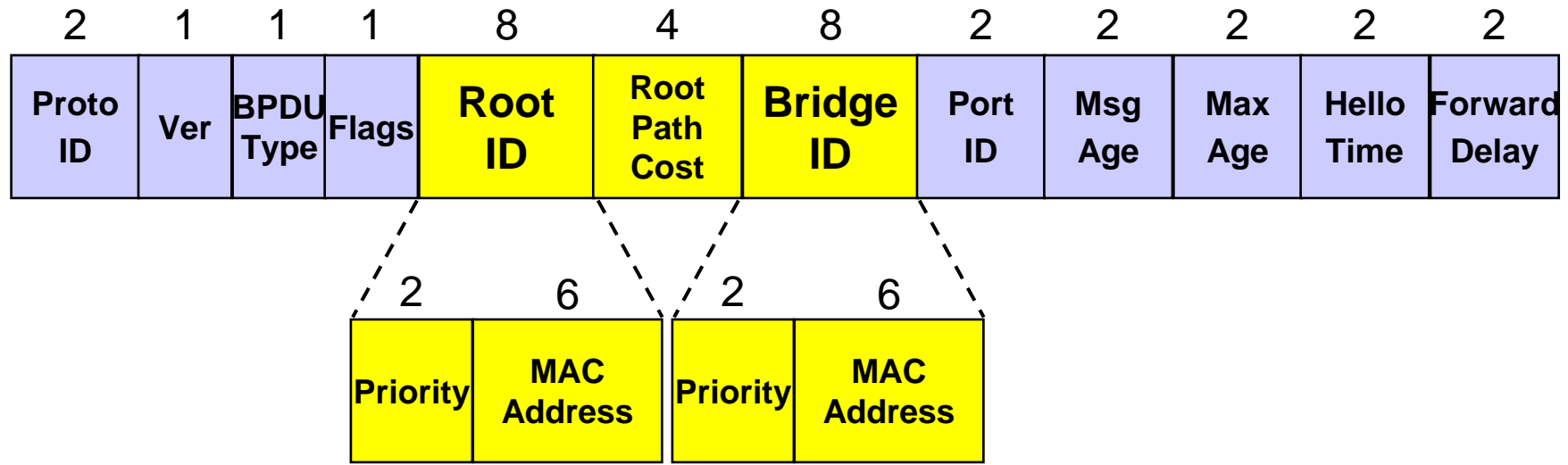
- Note that STP is not a routing protocol
  - It does not optimize routing
  - Traffic concentration towards the root



# Spanning Tree Protocol

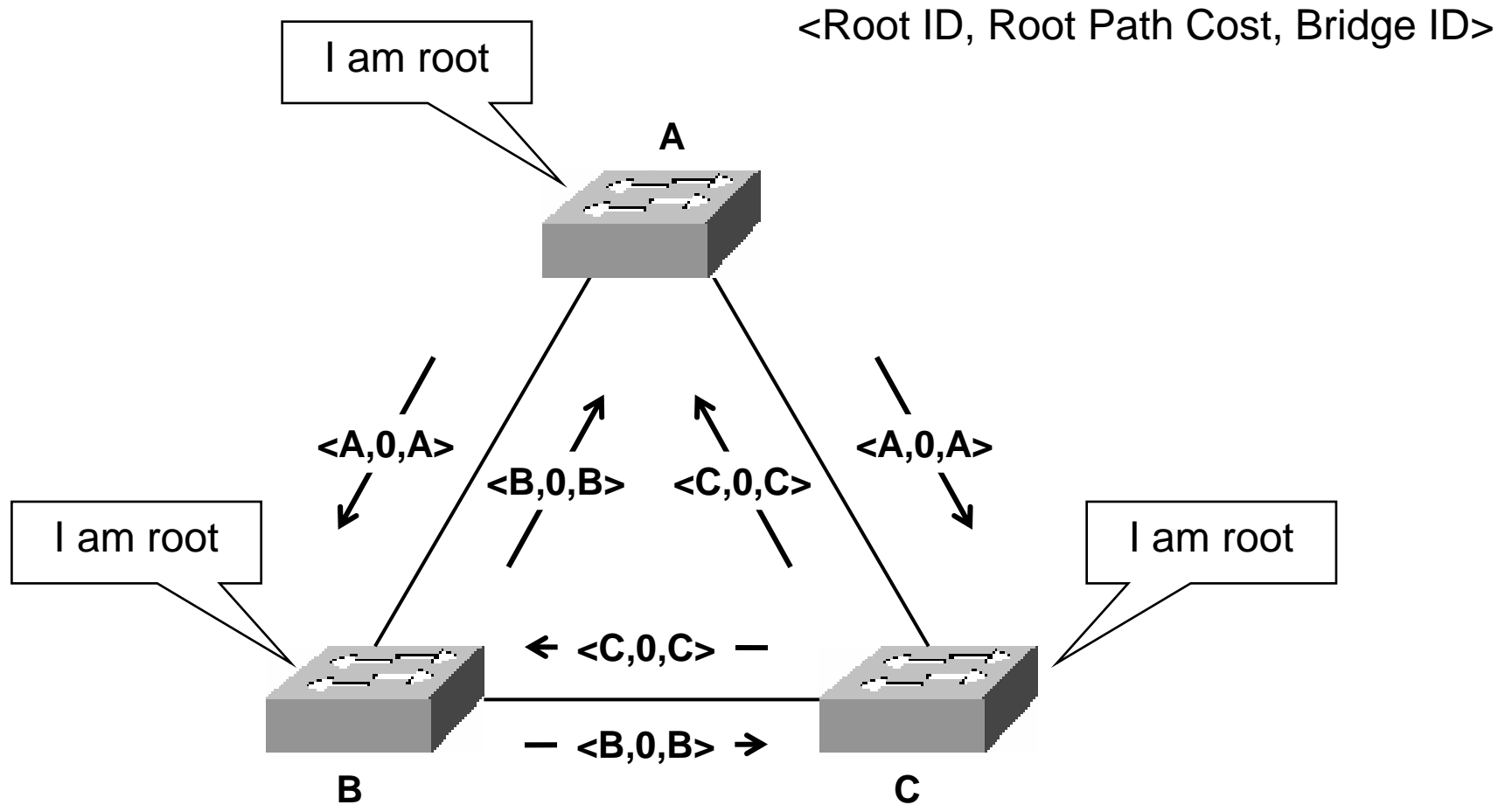
- Defined in IEEE 802.1D MAC Bridges standard
- Protocol to calculate a spanning tree
- Convergence
  - All bridges should reach a unified view of the spanning tree
- Special frames sent between neighbor switches
  - Bridge Protocol Data Units, BPDUs
  - Not forwarded!
- Based on bridge identifiers
  - Strict ordering
  - Based on MAC addresses

# Bridge Protocol Data Unit (BPDU)

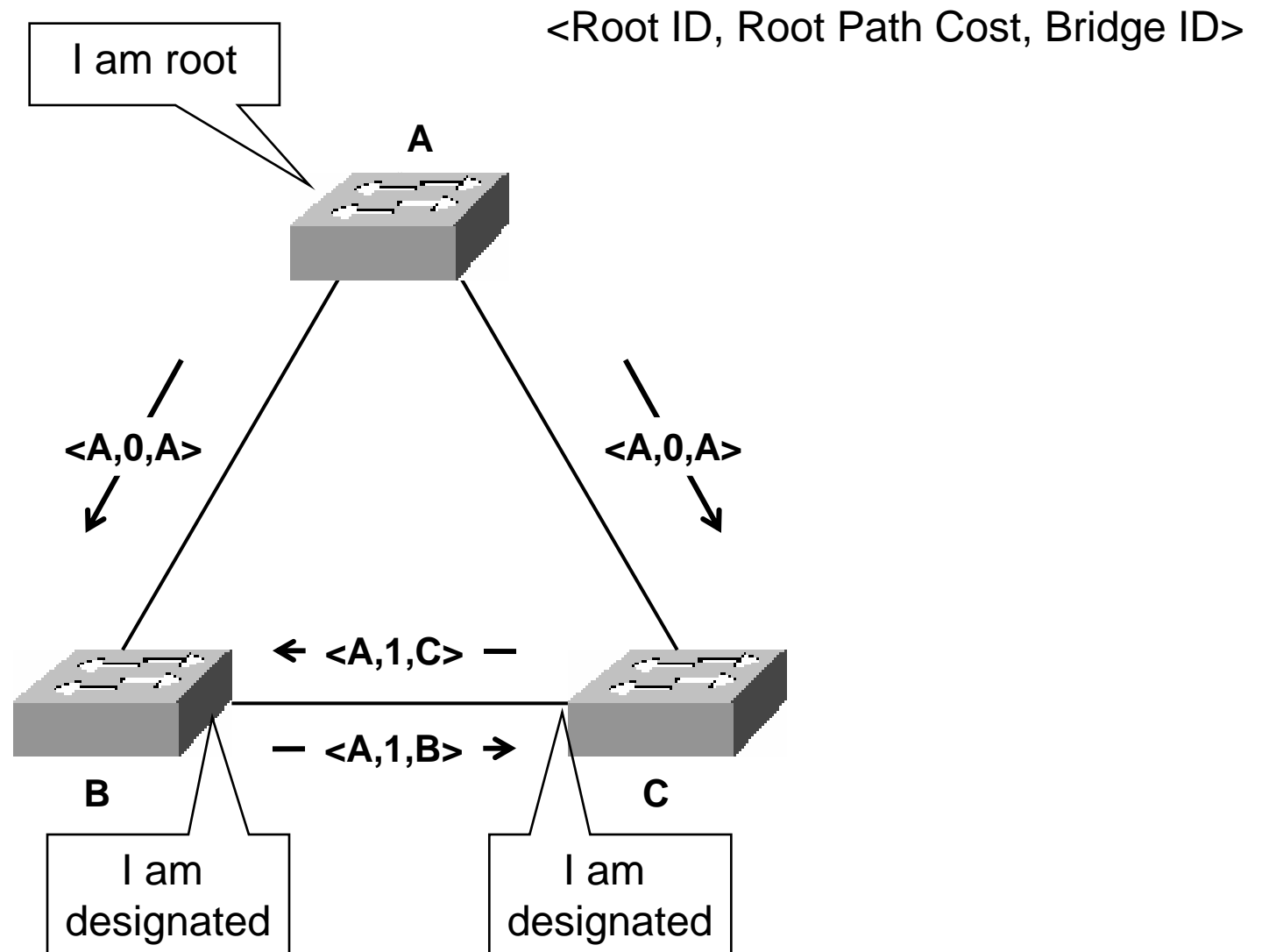


- Sent as an 802.1 frame
  - Destination MAC address 01-80-C2-00-00-00 (multicast)
  - BPDU Type 0
- Sent periodically (Hello Time) by root bridges
  - Triggers sending of BPDUs in designated bridges

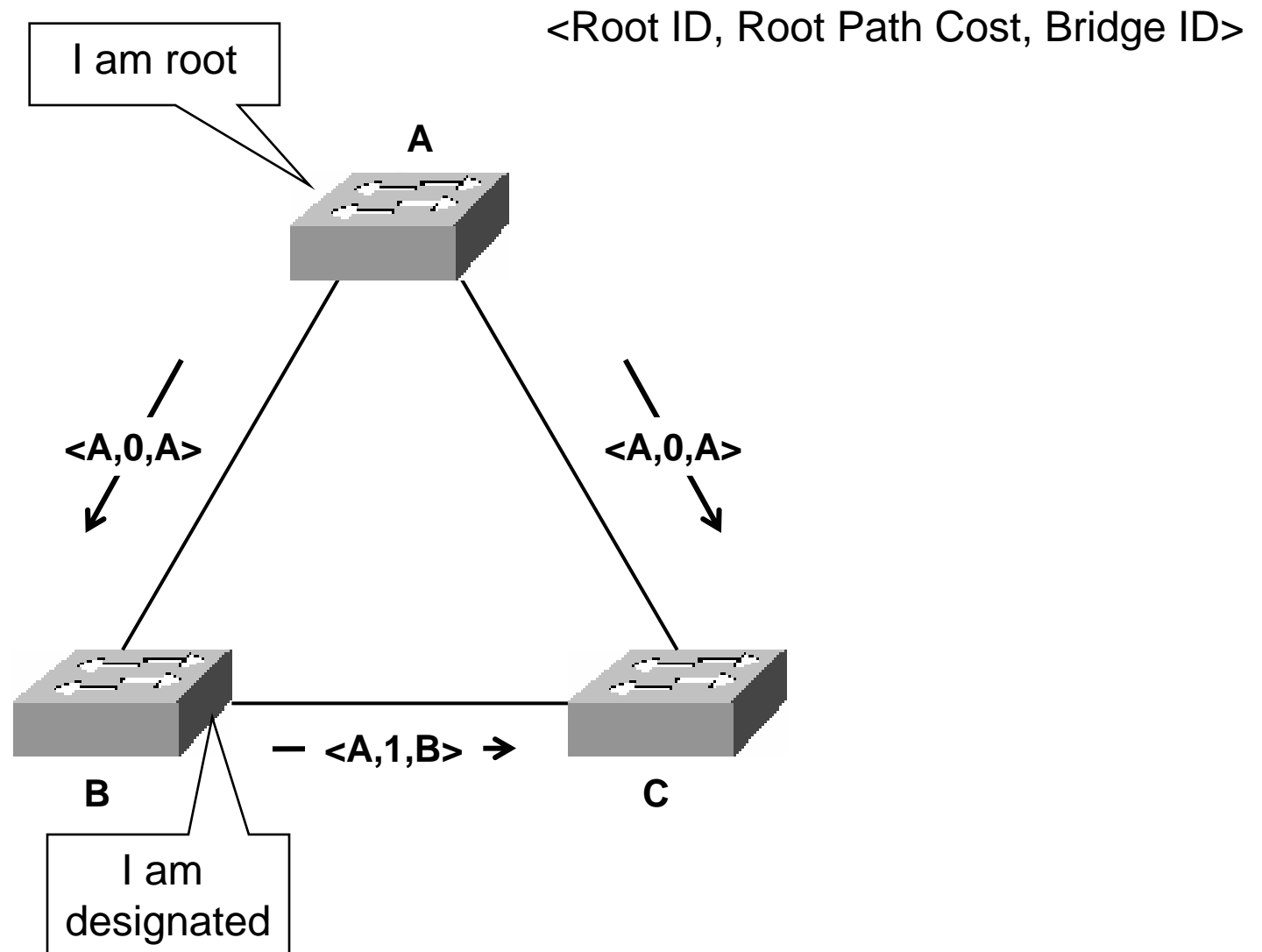
# Initial State



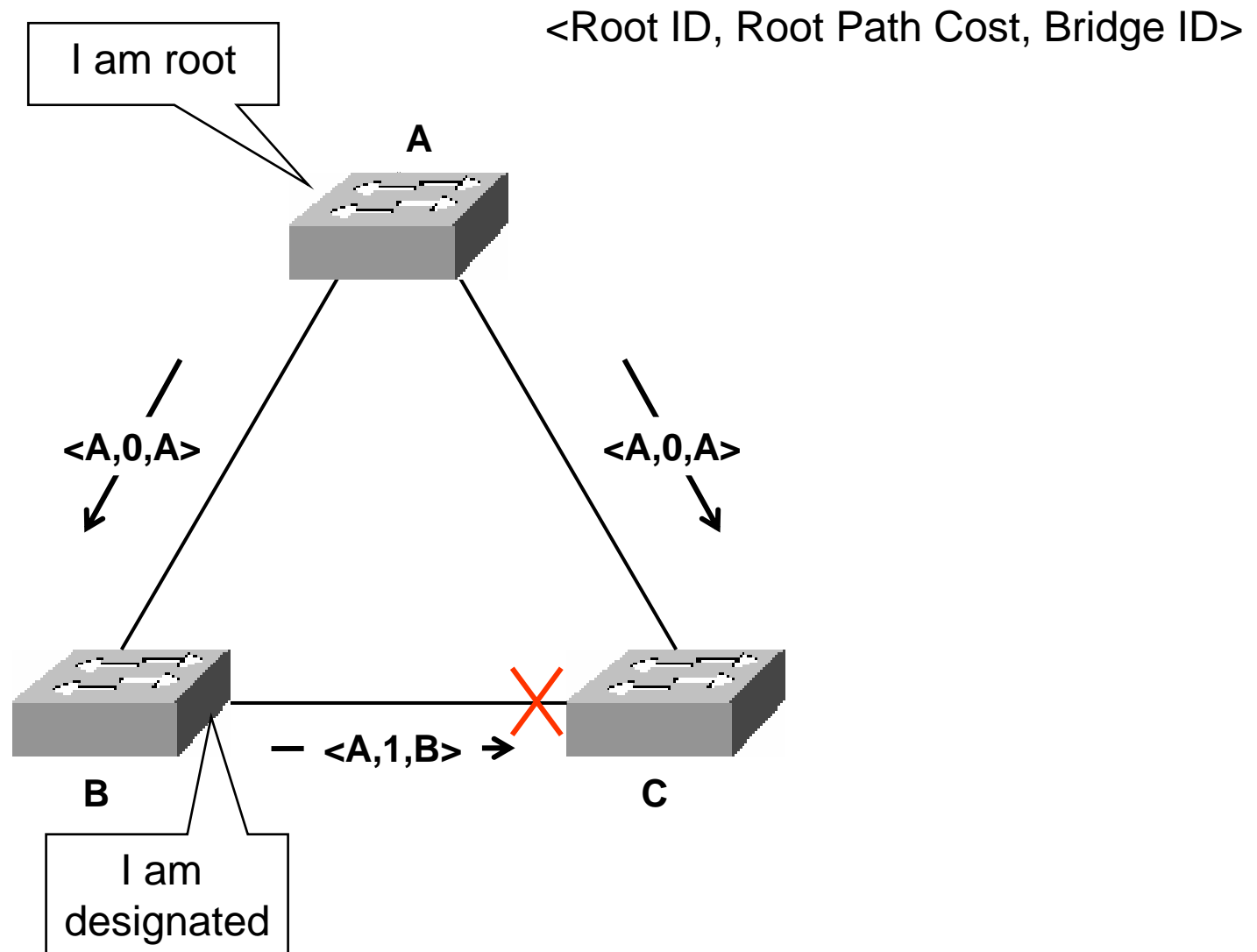
# Root Bridge Recognized



# Designated Bridge Recognized



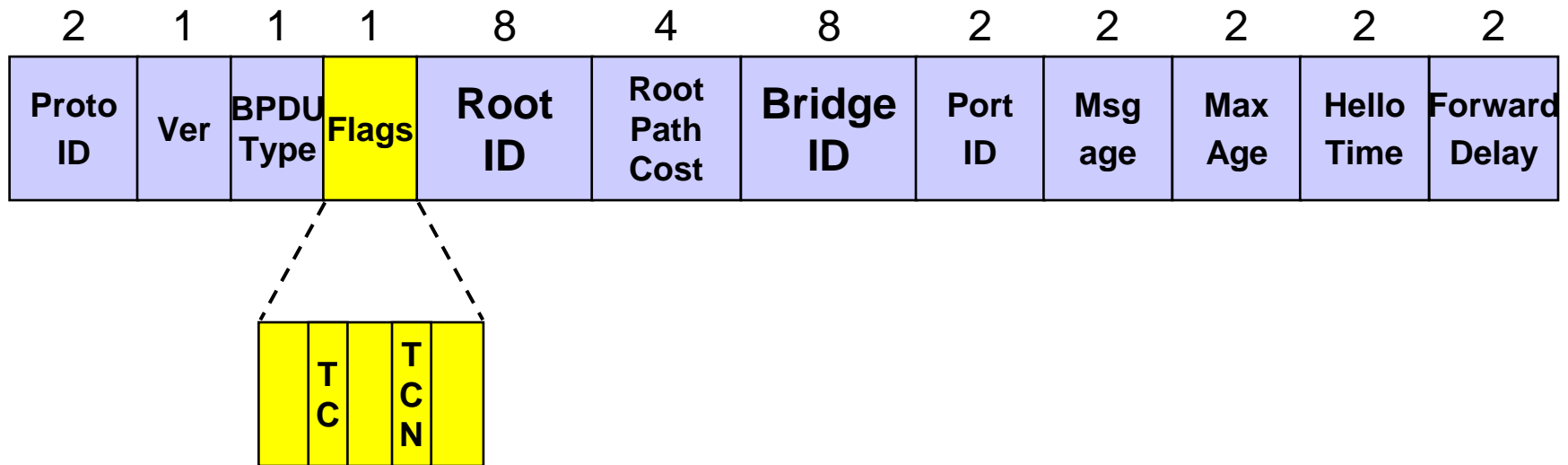
# Ports Disabled



# Topology Changes and Learning Table

- Entries in learning table expires
  - Normally after 5 minutes
  - Incorrect forwarding
    - “Black hole”
- Solution:
  - Faster expiration time when network configuration has changed
  - Topology Change Notification PDU

# Topology Change Notification



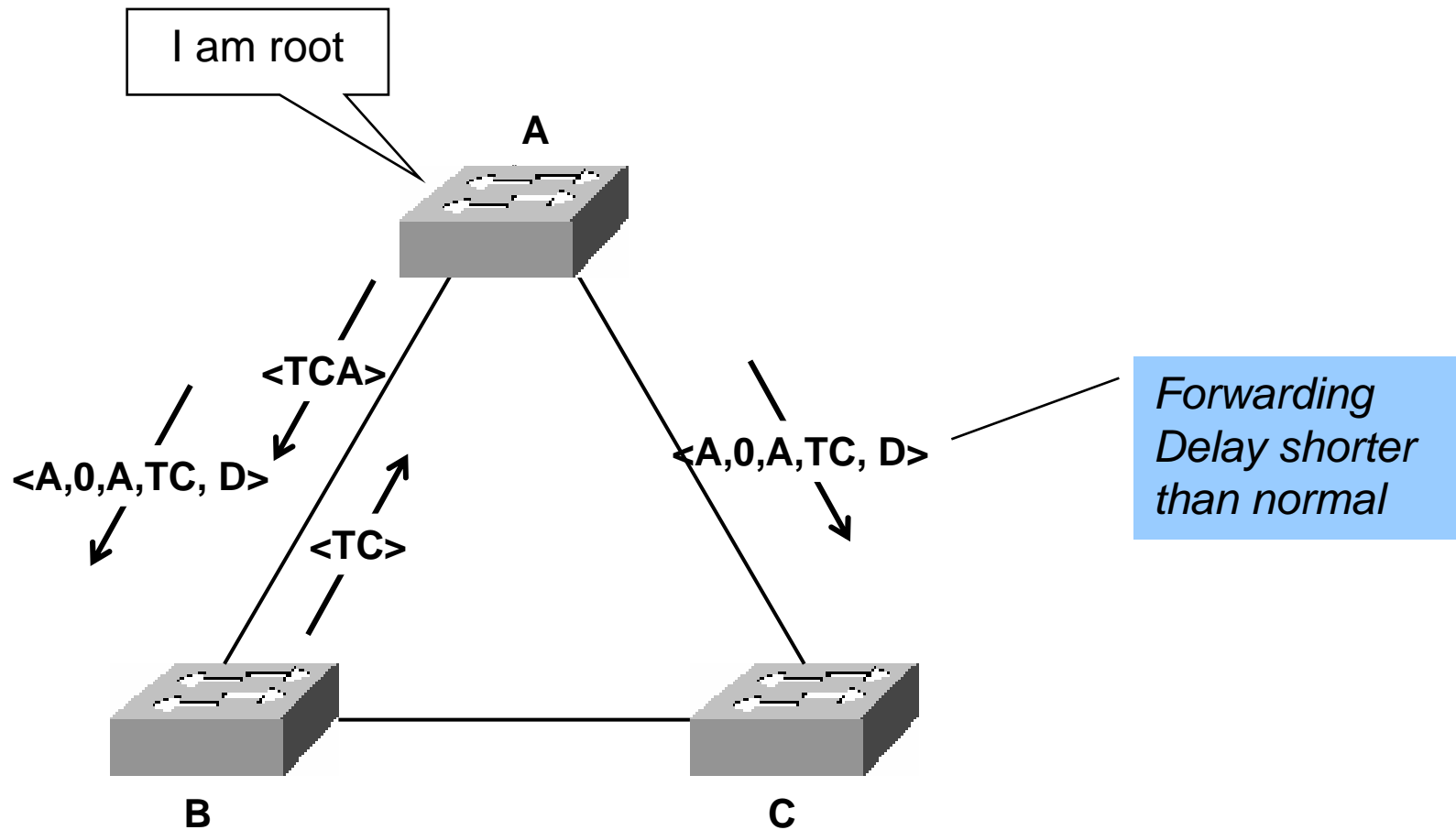
- Topology change at a bridge
  - Port failure
  - No periodic configuration BPDUs
  - Port status change
- Bridge sends spontaneous BPDUs
  - Topology Change Notification BPDUs
  - BPDUs type 0x80



# Topology Change

TCN BPDU:  
Configuration BPDU:

<Flags>  
<Root ID, Root Path Cost, Bridge ID,  
Flags, Forward Delay>



# Spanning Tree Protocol Timing

2	1	1	1	8	4	8	2	2	2	2	2
Proto ID	Ver	BPDU Type	Flags	Root ID	Root Path Cost	Bridge ID	Port ID	Msg age	Max Age	Hello Time	Forward Delay

- Protocol is timer driven
- Too short timers can give loops and instabilities
- Too long timers can give long convergence times
  - Until network reaches a stable spanning tree configuration

# Rapid Spanning Tree Protocol

- Ordinary STP takes 30 – 50 seconds to converge, with default settings
- Rapid Spanning Tree Protocol (RSTP)
  - First defined in IEEE 802.1w
  - Replaces STP after 2004 in 802.1D
  - Full-duplex mode
    - No shared links
  - Backwards compatible with STP
  - Converges in around one second

# RSTP vs STP

- RSTP has two more port designations
  - Alternate Port—backup for Root Port
  - Backup port—backup for Designated Port on the segment
- In RSTP, all bridges send BPDUs automatically
  - While in STP, the root triggers BPDUs
- In RSTP, bridges act to bring the network to convergence
  - While in STP, bridges passively wait for time-outs before changing port states

# Bridging Limitations

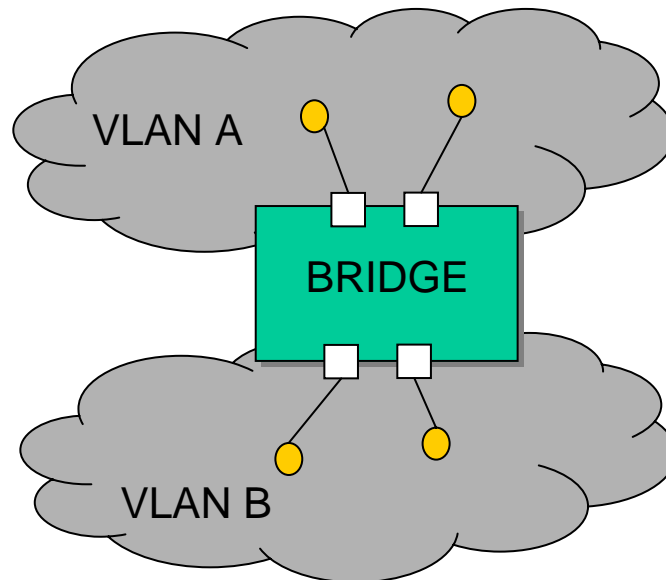
Can a bridged network be of an arbitrary size?

No, it is not scaling because of:

- Addressing
  - No structure in the MAC addresses → no aggregation can be used
- Table sizes (size of the learning table)
- Traffic congestion
  - Broadcast and multicast spans the whole broadcast domain
  - Broadcast domains must be kept to a reasonable size

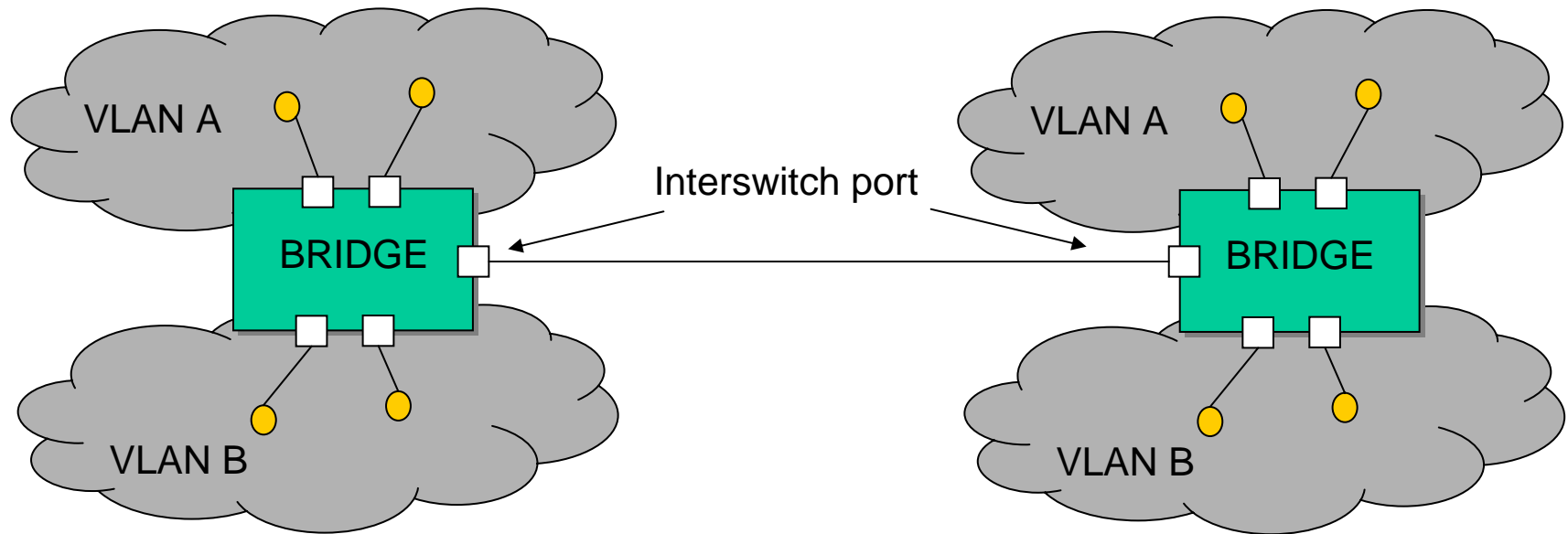
# Virtual LAN (VLAN)

- A single bridge can act as multiple bridges by grouping the ports of the bridge into separate VLANs.
- Like a LAN, each VLAN is a layer 2 realm separated from other VLANs (separate broadcast domains) by separating ports on a bridge.
- Routing between VLANs requires a router (just like LANs)

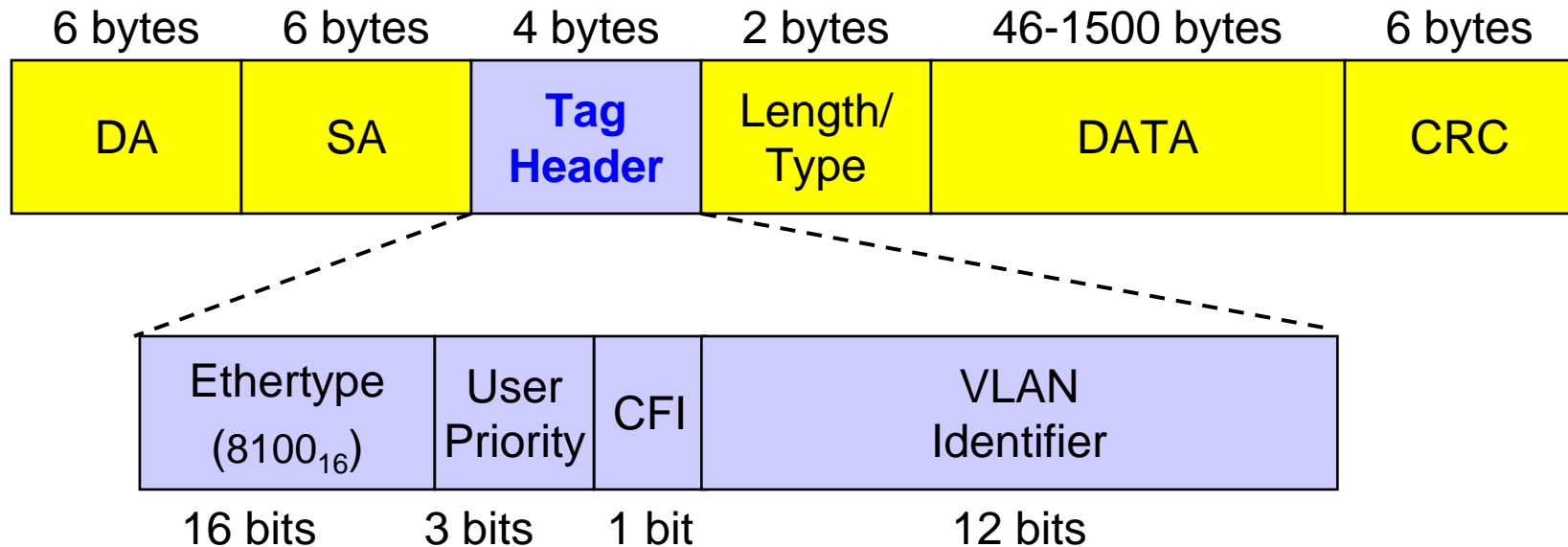


# VLAN Grouping and Frame Tagging

- VLAN membership determined through port number
  - Ports 1, 2: VLAN 1; Ports 3, 4: VLAN 2, etc
- Frames belonging to different VLANs can be sent on the same link by VLAN encapsulation using frame tagging
  - Interswitch ports belong to several VLANs



# IEEE 802.1Q Frame Tagging

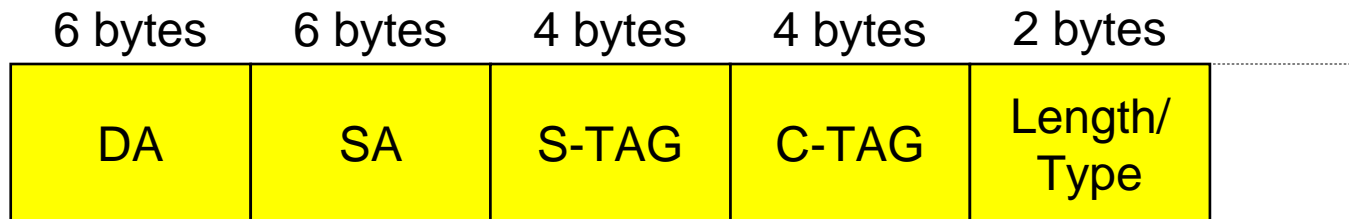


- Tag header added to Ethernet header
  - Ethertype 8100<sub>16</sub>
- 12-bit VLAN ID allows for 4096 VLANs



# VLAN Extensions

- Multiple spanning tree protocol
  - Original STP computes one spanning tree for all VLANs
  - Multiple STP computes one spanning tree per VLAN
- VLAN Stacking (Q-in-Q)
  - IEEE 802.1ad "Provider Bridge" standard
  - Extends tag space
  - Allows for providers to transport customer frames with tags
    - Outer tag: S-TAG or "provider tag" (IEEE 802.1ad)
    - Inner tag: C-TAG or "customer tag" (IEEE 802.1Q)



# Layer 2 Summary

- Point-to-point and Ethernet
- Bridging
  - Learning bridges
- Spanning Tree Protocol
- VLANs