Faculty of Computer & Information Sciences

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Research Topic (2) DIET and MIXING Model

Introduction

What's OR (Operation Research) ? (Ref. 1)

Operation Research is a systematic strategy for problem-solving and decision-making that is valuable in the management of organizations. In operations research, problems are separated into fundamental parts and then solved in defined steps by mathematical analysis. What's more, It's an interdisciplinary part of Applied Mathematics and formal Sciences that utilizes methods such as mathematical modeling, Statistics and Algorithm to show up optimal or near optimal solutions to complex problems. It's typically concerned with optimizing the maximum (benefit, profit, assembly line performance, crop yield, production, band width, etc.) or minimum (loss, cost, risk, etc.) of some objective function operations research helps management achieve its goals using strategies and methods.

What're the phases of Operation Research ? (Ref. 2)

- 1. Defining the problem and collecting data
- 2. Creating a mathematical model
- 3. Selecting appropriate data input
- 4. Checking the model and its solutions
- 5. Preparing to the execute the model (Model Validity)
- 6. Final Execution or Implementation

1- Defining the problem:

Identify the principle that mentioned before; Decision variables, Objectives, Constraints.

2- <u>Creating a mathematical model :</u>

Translate the problem definition into mathematical relationships. If the resulting model fits one of the standard mathematical models, such as LP. We can usually reach a solution by using available algorithm.

3- Selecting appropriate data input

No model will work appropriately if data input isn't proper. The purpose behind this step is to have adequate contribution to work and test the model.

4- Checking the model and its solutions

It is far the simplest of all OR phases. We should find a solution. If the model doesn't behave properly, then updating and modification is considered at this stage is a must.

5- Preparing to the execute the model (Model Validity)

checks whether or not the proposed model does what it purport to do. A model is supposed to be valid if it can give a solid and reliable prediction of the system's performance. A model must be appropriate for a longer time and can be updated every now and then taking into consideration the past, present and future aspects of the problem.

6- Final Execution or Implementation

The translation of the results into understandable operating instructions to be issued to the people who will administer the recommended system.

Formulate the Mathematical Linear Programming for the problem?

[1] Decision Variables:

Let x1,x2,x3 and x4 denote the number of units of type 1,2,3 and 4 respectively.

X1 the number of units of food of type **1**

X2 the number of units of food of type 2

X3 the number of units of food of type 3

X4 the number of units of food of type 4

[2] The Objective:

is to minimize the cost i.e., Minimize Z=Rs (45x1+40x2+85x3+65x4)

[3] Constraints:

Constraints are on the fulfillment of the daily requirements of the various constituents i.e.,

proteins	$3x1 + 4x2 + 8x3 + 6x4 \ge 800$
fats	$2x1 + 2x2 + 7x3 + 5x4 \ge 200$
carbohydrates	$6x1 + 4x2 + 7x3 + 4x4 \ge 700$

where $x1, x2, x3, x4 \ge 0$

What's the meaning of basic feasible solution of a general Linear Programming Problem ?

In linear programming theory, a basic feasible solution (BFS) is a solution with an insignificant arrangement of non-zero variables. Geometrically, each BFS corresponds to a corner of the polyhedron of feasible and attainable solutions. If there exists an optimal solution, then there exists an optimal BFS. Hence, If you want to find an optimal solution, it is sufficient to consider the BFS-s. This fact is utilized by the simplex algorithm, which basically goes from some BFS to another until an optimal and ideal one is found.

What's the solution to the problem using Regular Simplex method?

Modeling: Z=45 x1+40 x2+85 x3+65 x4. Subject to $3x1+4x2+8x3+6x4 \ge 800$, for proteins $2x1+2x2+7x3+5x4 \ge 200$, for fats $6x1+4x2+7x3+4x4 \ge 700$, for carbohydrates where $x1,x2,x3,x4 \ge 0$ hidden. Add 3 surplus and 3 artificial.

- 1. As the constraint -1 is of type '≥' we should subtract surplus variable S1 and add artificial variable A1.
- 2. As the constraint -2 is of type '≥' we should subtract surplus variable S2 and add artificial variable A2.
- **3.** As the constraint -3 is of type '≥' we should subtract surplus variable S3 and add artificial variable A3.

After adding surplus and artificial variables

Min
$$Z = 45x1 + 40x2 + 85x3 + 65x4 + 051 + 052 + 053 + MA1 + MA2 + MA3$$

subject to

$$3x1+4x2+8x3+6x4-S1$$
 + A1 = 800
 $2x1+2x2+7x3+5x4$ - S2 + A2 = 200
 $6x1+4x2+7x3+4x4$ - S3 + A3=700

and $x1,x2,x3,x4,S1,S2,S3,A1,A2,A3 \ge 0$

Table 1													
Iteration-1 n-1		Cj	45	40	85	65	0	0	0	М	М	М	
В	Св	Хв	<i>x</i> 1	<i>x</i> 2	<i>x</i> 3	<i>x</i> 4	S1	S2	S3	A1	A2	A3	Min Ratio XBx3
<i>A</i> 1	M	800	3	4	8	6	-1	0	0	1	0	0	8008=100
A2	M	200	2	2	(7)	5	0	-1	0	0	1	0	2007=28.5714→
A3	M	700	6	4	7	4	0	0	-1	0	0	1	7007=100
Z=1700M		$\mathbf{Z}j$	11 <i>M</i>	10M	22M	15 <i>M</i>	-M	-M	-M	M	M	M	
		Z j-C j	11 <i>M</i> -45	10 <i>M</i> -40	22 <i>M</i> -85↑	15M-65	-M	-M	-M	0	0	0	

Enter =x3, Drop =A2, Key Element =7.

Positive maximum Z j-C j is 22M-85 and its column index is 3. So, the entering variable is x3.

Minimum ratio is 28.5714 and its row index is 2. So, the leaving basis variable is A2.

 \therefore The pivot element is 7.

$$R2(new)=R2(old)/7$$

$$R1(new)=R1(old) - 8R2(new)$$

$$R3(new)=R3(old) - 7R2(new)$$

Table 2													
Iteration-2		C_j	45	40	8 5	65	0	0	0	М	M	M	
В	C B	Хв	<i>x</i> 1	<i>x</i> 2	<i>x</i> 3	<i>x</i> 4	<i>S</i> 1	S2	<i>S</i> 3	A 1	A2	A 3	Min Ratio <i>XBx</i> 1
<i>A</i> 1	М	571. 428 6	0.7143	1.7143	0	0.2857	- 1	1.1429	0	1	-1.1429	0	571.42860 .7143= 800
<i>x</i> 3	8 5	28.5 714	(0.2857)	0.2857	1	0.7143	0	-0.1429	0	0	0.1429	0	28.57140. 2857=100 →
A3	M	500	4	2	0	-1	0	1	1	0	-1	1	5004=125
Z=1071.428 6M+2428.57 14		Z_j	4.7143 M+24.2 857	3.7143 M+24.2 857	8 5	0.7143 <i>M</i> +60.714 3	- M	2.1429 <i>M</i> - 12.1429	- M	M	- 2.1429 <i>M</i> +12.142 9	M	
		Z j- C j	4.7143 <i>M</i> - 20.7143	3.7143 <i>M</i> - 15.7143	0	- 0.7143 <i>M</i> -4.2857	- M	2.1429 <i>M</i> - 12.1429	- M	0	3.1429 <i>M</i> +12.142 9	0	

Enter = x1, Drop =x3, Key Element =0.2857.

Positive maximum Z j-C j is 4.7143M-20.7143 and its column index is 1. So, the entering variable is x.

Minimum ratio is 100 and its row index is 2. So, the leaving basis variable is x3.

 \therefore The pivot element is 0.2857.

$$R2(\text{new})=R2(\text{old})\div0.2857$$

$$R1(\text{new}) = R1(\text{old}) - 0.7143R2(\text{new})$$

$$R3(\text{new})=R3(\text{old}) - 4R2(\text{new})$$

Table 3													
Iteration-3		Cj	4 5	40	85	65	0	0	0	M	М	M	
В	C B	ХB	<i>x</i> 1	<i>x</i> 2	<i>x</i> 3	<i>x</i> 4	<i>S</i> 1	S2	<i>S</i> 3	A 1	A2	A 3	Min Ratio XBS2
<i>A</i> 1	M	500	0	1	-2.5	-1.5	-1	1.5	0	1	-1.5	0	5001.5=333
<i>x</i> 1	4 5	100	1	1	3.5	2.5	0	-0.5	0	0	0.5	0	
A3	M	100	0	-2	-14	-11	0	(3)	-1	0	-3	1	1003=33.33 33→
Z=600M+ 4500		Z_j	4 5	- M+ 45	- 16.5 <i>M</i> +1 57.5	12.5 <i>M</i> +1 12.5	- М	4.5 <i>M</i> - 22.5	- М	M	- 4.5 <i>M</i> +2 2.5	M	
		Z j - C j	0	- M+5	16.5 <i>M</i> +7 2.5	12.5 <i>M</i> +4 7.5	- М	4.5 <i>M</i> - 22.5↑	- М	0	5.5 <i>M</i> +2 2.5	0	

Enter = S2, Drop =A3, Key Element =3.

Positive maximum Z j-C j is 4.5M-22.5 and its column index is 6. So, the entering variable is S2.

Minimum ratio is 33.3333 and its row index is 3. So, the leaving basis variable is A3.

∴The pivot element is 3.

$$R3(\text{new})=R3(\text{old})/3$$

$$R1(\text{new}) = R1(\text{old}) - 1.5R3(\text{new})$$

$$R2(\text{new}) = R2(\text{old}) + 0.5R3(\text{new})$$

Table 4													_
Iteration-4		C_j	4 5	40	85	65	0	0	0	М	М	М	
В	C B	ХВ	<i>x</i> 1	<i>x</i> 2	<i>x</i> 3	<i>x</i> 4	<i>S</i>	S 2	S3	A 1	A 2	A3	Min Ratio XBx3
<i>A</i> 1	M	450	0	2	(4.5)	4	-1	0	0.5	1	0	-0.5	4504.5=100→
<i>x</i> 1	4 5	116.66 67	1	0.666 7	1.1667	0.666 7	0	0	- 0.1667	0	0	0.1667	116.66671.166 7=100
S2	0	33.333	0	- 0.666 7	-4.6667	- 3.666 7	0	1	0.3333	0	-1	0.3333	
Z=450M+ 5250		Zj	4 5	2 <i>M</i> + 30	4.5 <i>M</i> +5 2.5	4 <i>M</i> + 30	- М	0	0.5 <i>M</i> -7.5	M	0	- 0.5 <i>M</i> + 7.5	
		<i>Z j</i> - <i>C j</i>	0	2 <i>M</i> - 10	4.5 <i>M</i> - 32.5↑	4 <i>M</i> -35	- М	0	0.5 <i>M</i> - 7.5	0	- <i>M</i>	- 1.5 <i>M</i> +7 .5	

Enter = X3, Drop = A1, Key Element = 4.5.

Positive maximum Z j-C j is 4.5M-32.5 and its column index is 3. So, the entering variable is x3.

Minimum ratio is 100 and its row index is 1. So, the leaving basis variable is A1.

 \therefore The pivot element is 4.5.

$$R1(\text{new})=R1(\text{old})\div 4.5$$

$$R2(\text{new}) = R2(\text{old}) - 1.1667R1(\text{new})$$

$$R3(\text{new}) = R3(\text{old}) + 4.6667R1(\text{new})$$

Table 5													
Iteratio n-5		Cj	4 5	40	8 5	65	0	0	0	М	М	М	
В	C B	ХB	<i>x</i> 1	<i>x</i> 2	<i>x</i> 3	<i>x</i> 4	S1	S 2	S3	<i>A</i> 1	A 2	A3	Min Ratio XBx2
<i>x</i> 3	8 5	100	0	0.444 4	1	0.888	0.222 2	0	0.111	0.2222	0	-0.1111	1000.4444=225
<i>x</i> 1	4 5	0	1	(0.148 1)	0	0.370 4	0.259	0	0.296 3	-0.2593	0	0.2963	00.1481=0→
S2	0	500	0	1.407 4	0	0.481	1.037	1	0.185	1.037	-1	-0.1852	5001.4074=355. 2632
Z=8500		Z_j	4 5	44.44 44	8 5	58.88 89	7.222 2	0	3.888 9	7.2222	0	3.8889	
		Z j- C j	0	4.444 4↑	0	- 6.111 1	7.222 2	0	3.888 9	- M+7.22 22	- М	- M+3.88 89	

Enter = X2, Drop = X1, Key Element = 4.5.

Positive maximum Z j-C j is 4.4444 and its column index is 2. So, the entering variable is x2.

Minimum ratio is 0 and its row index is 2. So, the leaving basis variable is X1.

 \therefore The pivot element is 0.1481.

$$R2(\text{new}) = R2(\text{old}) \div 0.1481$$

$$R1(\text{new}) = R1(\text{old}) - 0.4444R2(\text{new})$$

$$R3(\text{new}) = R3(\text{old}) - 1.4074R2(\text{new})$$

Table 6													
Iteration-6		C_j	45	40	85	65	0	0	0	М	М	M	
В	Св	Хв	<i>x</i> 1	<i>x</i> 2	<i>x</i> 3	<i>x</i> 4	S1	S2	<i>S</i> 3	<i>A</i> 1	A2	A3	Min Ratio XBX4
<i>x</i> 3	85	100	-3	0	1	(2)	-1	0	1	1	0	-1	1002=50→
<i>x</i> 2	40	0	6.75	1	0	-2.5	1.75	0	-2	-1.75	0	2	
<i>S</i> 2	0	500	-9.5	0	0	4	-3.5	1	3	3.5	-1	-3	5004=125
Z=8500		$\mathbf{Z}j$	15	40	85	70	-15	0	5	15	0	-5	
		Z j - C j	-30	0	0	5↑	-15	0	5	-M+15	-M	-M-5	

Enter = X4, Drop = X3, Key Element = 2.

Positive maximum Z j-C j is 5 and its column index is 4 . So, the entering variable is X4.

Minimum ratio is 50 and its row index is 1. So, the leaving basis variable is X3.

∴The pivot element is 2.

$$R1(\text{new})=R1(\text{old})/2$$

 $R2(\text{new})=R2(\text{old}) + 2.5R1(\text{new})$
 $R3(\text{new})=R3(\text{old}) - 4R1(\text{new})$

Table 7													
Iteration-7		C_j	45	40	85	65	0	0	0	М	М	М	
В	Св	Хв	<i>x</i> 1	<i>x</i> 2	<i>x</i> 3	<i>x</i> 4	S1	S2	<i>S</i> 3	<i>A</i> 1	A2	A3	Min Ratio XBS3
<i>x</i> 4	65	50	-1.5	0	0.5	1	-0.5	0	(0.5)	0.5	0	-0.5	500.5=100→
<i>x</i> 2	40	125	3	1	1.25	0	0.5	0	-0.75	-0.5	0	0.75	
S2	0	300	-3.5	0	-2	0	-1.5	1	1	1.5	-1	-1	3001=300
Z=8250		Zj	22.5	40	82.5	65	-12.5	0	2.5	12.5	0	-2.5	
		Z_{j} - C_{j}	-22.5	0	-2.5	0	-12.5	0	2.5↑	-M+12.5	-M	-M-2.5	

Enter = S3, Drop = X4, Key Element = 0.5.

Positive maximum Z j-C j is 2.5 and its column index is 7 . So, the entering variable is S3.

Minimum ratio is 100 and its row index is 1. So, the leaving basis variable is X4.

: The pivot element is 0.5.

$$R1(\text{new}) = R1(\text{old})/0.5$$

$$R2(\text{new}) = R2(\text{old}) + 0.75R1(\text{new})$$

$$R3(\text{new})=R3(\text{old}) - R1(\text{new})$$

Table 8													
Iteration-8		Cj	45	40	85	65	0	0	0	М	М	М	
В	Св	Хв	<i>x</i> 1	<i>x</i> 2	<i>x</i> 3	<i>x</i> 4	S1	S2	<i>S</i> 3	<i>A</i> 1	A2	A 3	Min Ratio
<i>S</i> 3	0	100	-3	0	1	2	-1	0	1	1	0	-1	
<i>x</i> 2	40	200	0.75	1	2	1.5	-0.25	0	0	0.25	0	0	
S2	0	200	-0.5	0	-3	-2	-0.5	1	0	0.5	-1	0	
Z=8000		Z_j	30	40	80	60	-10	0	0	10	0	0	
		Z_{j} - C_{j}	-15	0	-5	-5	-10	0	0	-M+10	-M	-М	

Since all Zj-Cj \leq 0

Hence, optimal solution is arrived with value of variables as : x1=0, x2=200, x3=0, x4=0

then Min Z=40*200 = 8000.

Find the Dual Problem, and determine which is computationally advantageous to solve either the Primal or the Dual, then solve both problems .

1-Dual problem:

Max=
$$b^t$$
 y
 a^t y <= c y >= 0
max z= 800 y1+200 y2+700 y3.
Subject to
3 y1+2 y2+6 y3<=45,
4 y1+2 y2+4 y3<=40,
8 y1+7 y+7 y3<=85,
6 y1+5 y2+4 y3<=65.

<u>2-Determine which is computationally advantageous to solve either the Primal</u> or the Dual:

Dual problem is better than primal problem .

3- Solve the Dual problem:

	Y1	Y2	Y3	S10	S20	S30	S40	Y1	
	800	200	700					800	
S 1	0	3	2	6	1	0	0	0	45
S2	0	4	2	4	0	1	0	0	40
S 3	0	8	7	7	0	0	1	0	85
S4	0	6	5	4	0	0	0	1	68
	-800	-200	-700	0	0	0	0	-800	

Enter = y1, Drop = S2.

R2 = R2/4

R1=R1-3R2

R3=R3-8R2

R4=R4-6R2

	Y1	Y2	Y3	S10	S20	S30	S40	Y1	
	800	200	700					800	
S 1	0	0	1/2	3	1	-3/4	0	0	15
Y 1	800	1	1/2	1	0	1/4	0	0	10
S 3	0	0	3	-1	0	-2	1	0	5
S4	0	0	2	2	0	-3/2	0	1	5
	0	200	100	0	200	0	0	0	

So

S1=x1=0,s2=x2=200,s3=x3=0,s4=x4=0;

Then Z=200*40=**8000**.

4-solve primal with regular simplex methods (in number 4)

I solved this problem previously in this research , to open click **Here.**

Solve the Problem Using Dual Simplex Method, if possible.

The condition is met min+ so we can solve it but Need to adjust the constrains 1- constrains hit *-1 to <=

Then add slacks

Min Z = 45 x1+40 x2+85 x3+65 x4+0s1+0s2+0s3

-3x1-4x2-8x3-6x4+s1=-800, for proteins

-2x1-2x2-7x3-5x4-s2=-200, for fats

-6x1-4x2-7x3-4x4+s3=-700, for carbohydrates

 $x1,x2,x3,x4,s1,s2,s3 \ge 0$ hidden

		X1	X2	X3	X4	S 1	S2	S3	
		45	40	85	65	0	0	0	
S 1	0	-3	-4	-8	-6	1	0	0	-800
S2	0	-2	-2	-7	-5	0	1	0	-200
S3	0	-6	-4	-7	-4	0	0	1	-700
		45	40	85	65	0	0	0	

Enter x2,Drop s1 R1=R1*-1THEN R1=R1/4 R2=R2-2R1 R3=R3-4R1

		X1	X2	X3	X4	S 1	S2	S3	
		45	40	85	65	0	0	0	
X2	40	3/4	1	2	3/2	-1/4	0	0	200
S2	0	-1/2	0	-3	-2	-1/2	1	0	200
S 3	0	-3	0	1	2	-1	0	1	100
		15	0	5	5	10	0	0	

X2=200,x1=0,x3=0,x4=0 Then z=40*200=8000.

By using any computer software (e.g. TORA, MATLAB, ...) solve this optimization problem, or write down a program that solve the Problem using any programming language you know (program one method ONLY; Regular method, Dual Simplex method, or Dual Problem).

The Code of Solving the problem using **Regular method** with **MatLab**:

```
prob = optimproblem('ObjectiveSense', 'min');

x = optimvar('x',4,3,2,1,'LowerBound',0);

prob.Objective = 45*x(1) + 40*x(2) + 85*x(3) + 65*x(4);

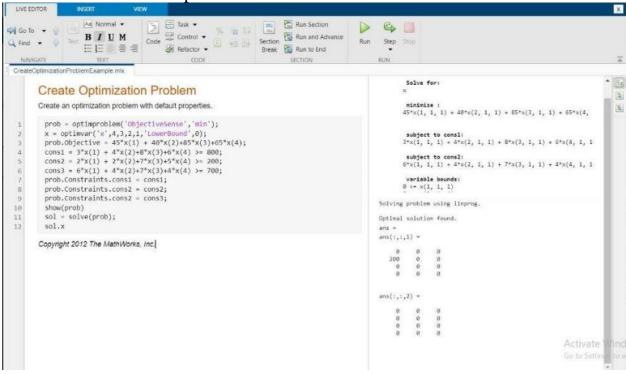
cons1 = 3*x(1) + 4*x(2) + 8*x(3) + 6*x(4) >= 800;

cons2 = 2*x(1) + 2*x(2) + 7*x(3) + 5*x(4) >= 200;

cons3 = 6*x(1) + 4*x(2) + 7*x(3) + 4*x(4) >= 700;
```

```
prob.Constraints.cons1 = cons1;
prob.Constraints.cons2 = cons2;
prob.Constraints.cons2 = cons3;
show(prob)
sol = solve(prob);
sol.x
```

A Screenshot for the output of Matlab:



Reference:

- [1] Tech Target (Website): Link
- [2] **Tutors Globe** (Website): <u>Link</u>
- [2] "Schaum's Outline of Operations Research 2nd Edition" (Book): Link
- [4] Article shared by Nikhila C (Article): Link
- [5] MatLab (Website): Link