

# CS341

# Artificial Intelligence

## Lecture 5

---

**DR. HEBA MOHSEN**

# Heuristic Search

---

- On the average they improve the quality of the paths that are explored.
- Using Heuristics, we can hope to get good ( though possibly non-optimal ) solutions
- There are good general purpose heuristics that are useful in a wide variety of problem domains.
- Special purpose heuristics exploit domain specific knowledge
- Heuristic search uses **Heuristic Function**: This is a function that maps from problem state descriptions to measures of desirability, usually represented as numbers.

# Example (1): 8-puzzle

- $f_1(T)$  = the number correctly placed tiles on the board:

$$f_1 \left( \begin{array}{|c|c|c|} \hline 1 & 3 & 2 \\ \hline 8 & & 4 \\ \hline 5 & 6 & 7 \\ \hline \end{array} \right) = 4$$

- $f_2(T)$  = number of incorrectly placed tiles on board:

- gives (rough!) estimate of how far we are from goal

$$f_2 \left( \begin{array}{|c|c|c|} \hline 1 & 3 & 2 \\ \hline 8 & & 4 \\ \hline 5 & 6 & 7 \\ \hline \end{array} \right) = 4$$

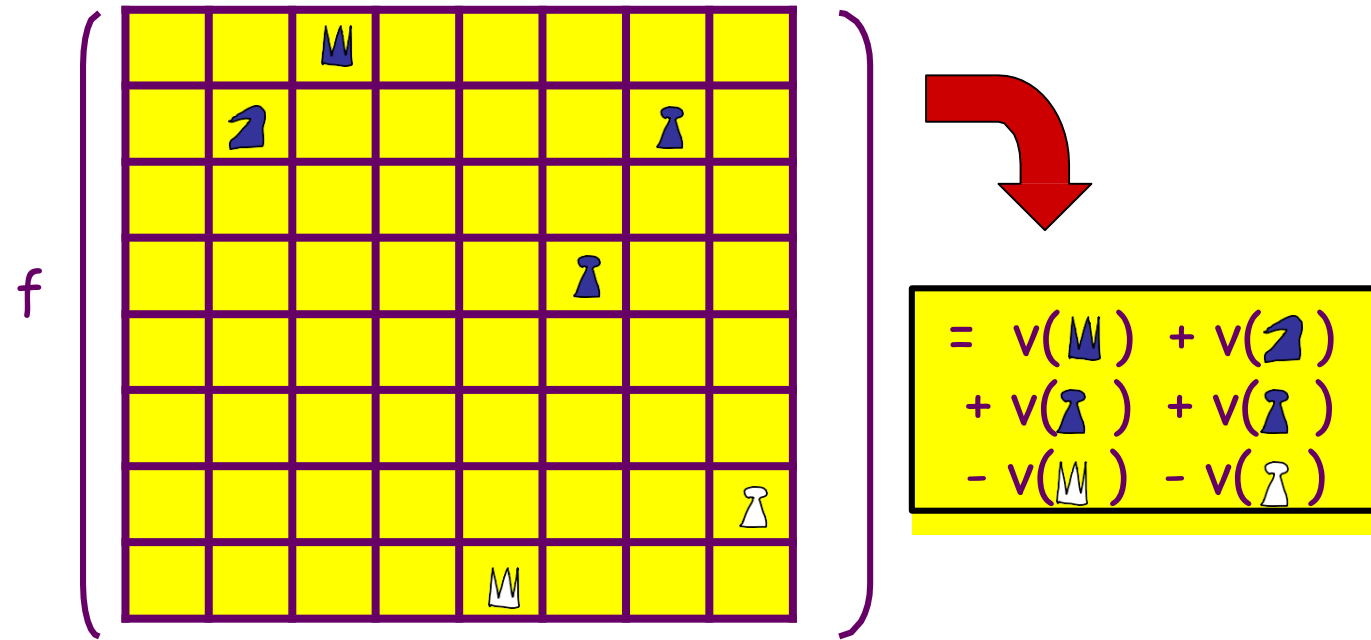
Most often, 'distance to goal' heuristics are more useful !

- **$f_3(T)$  = the sum of ( the horizontal + vertical distance that each tile is away from its final destination):**
  - gives a better estimate of distance from the goal node

$$f_3 \left( \begin{array}{|c|c|c|} \hline 1 & 3 & 2 \\ \hline 8 & & 4 \\ \hline 5 & 6 & 7 \\ \hline \end{array} \right) = 1 + 1 + 2 + 2 = 6$$

# Examples (2): Chess:

- $F(T) = (\text{Value count of black pieces}) - (\text{Value count of white pieces})$



# Best-first search

---

Idea: use an **evaluation function**  $f(n)$  for each node

- $f(n)$  provides an estimate for the total cost.
- Expand the node  $n$  with smallest  $f(n)$ .

Implementation:

Order the nodes in fringe increasing order of cost.

Special cases:

- greedy best-first search
- $A^*$  search

# Example

We start from source "S" and search for goal "I" using given costs and Best First search.

pq (priority queue) initially contains S We remove s from and process unvisited neighbors of S to pq. pq now contains {A, C, B} (C is put before B because C has lesser cost)

We remove A from pq and process unvisited neighbors of A to pq. pq now contains {C, B, E, D}

We remove C from pq and process unvisited neighbors of C to pq. pq now contains {B, H, E, D}

We remove B from pq and process unvisited neighbors of B to pq. pq now contains {H, E, D, F, G}

We remove H from pq. Since our goal "I" is a neighbor of H, we return.

