

Data Structure & Algorithms

Sunbeam Infotech

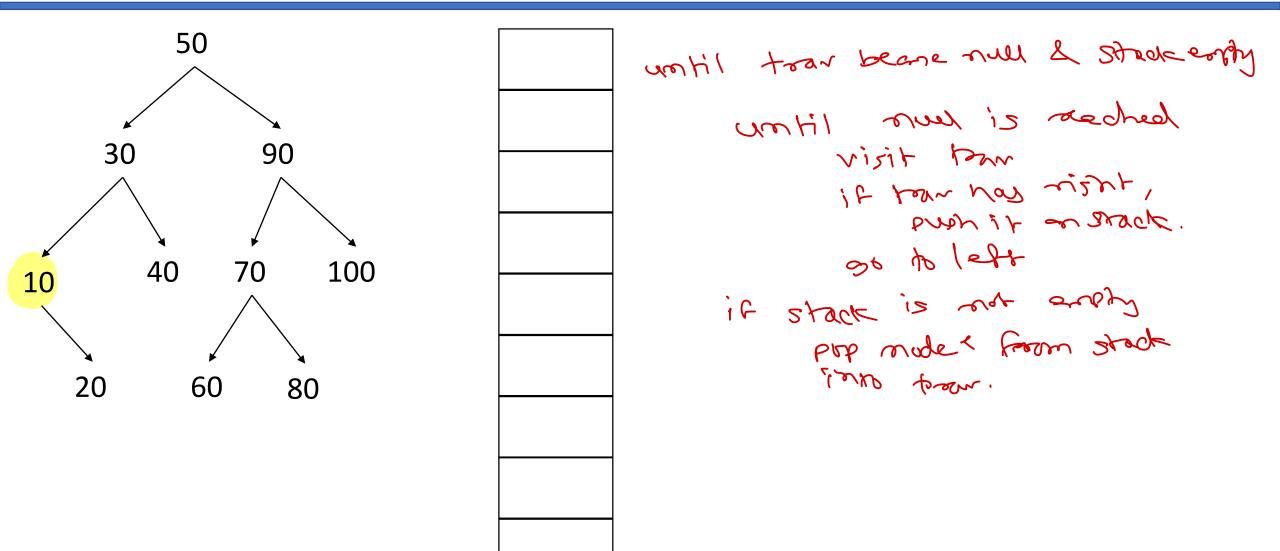


Agenda

- BST In-order traversal (non-recursive)
- BST Post-order traversal (non-recursive) /
- BST BFS & DFS (non-recursive) /
- BST Delete node /
- BST Balance tree
- BST Types
 - Skewed Binary Tree
 - AVL Tree
 - R & B Tree /

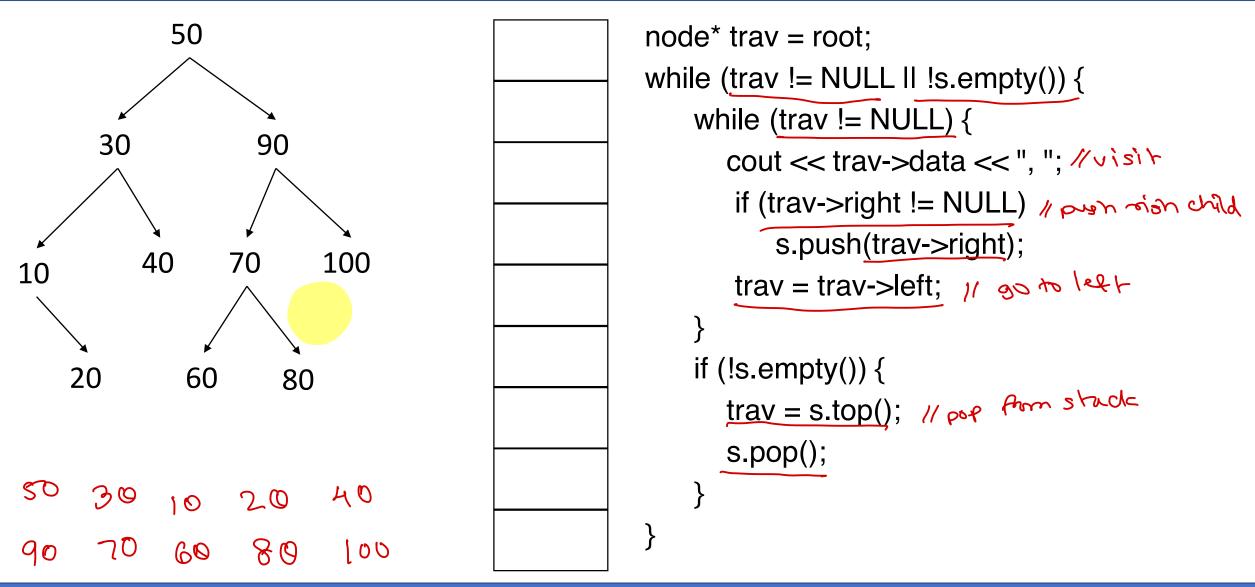


BST – Non-Recursive Algorithm – PreOrder





BST – Non-Recursive Algorithm – PreOrder

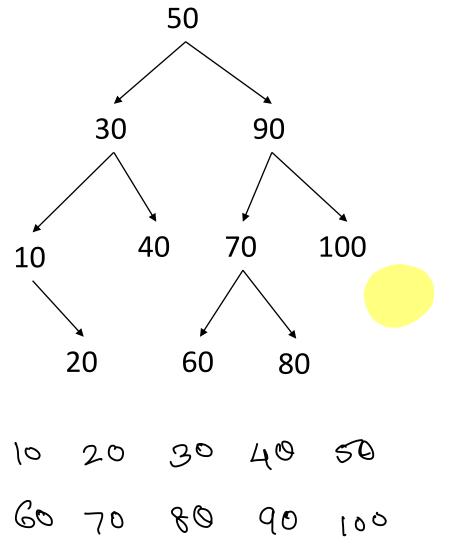


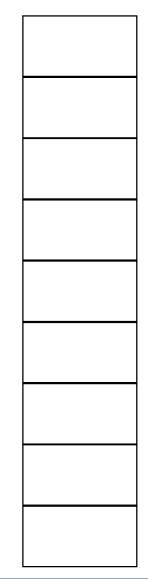


BST – Non-Recursive Algorithm – InOrder





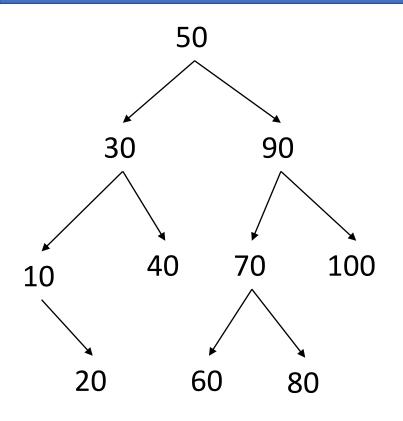




trav = robor; until toar is null as stage is end; until nul is reached: push trav on stack go to the left if stack is not empty: pop town four stade visit toav go to the risar



BST – Non-Recursive Algorithm – InOrder



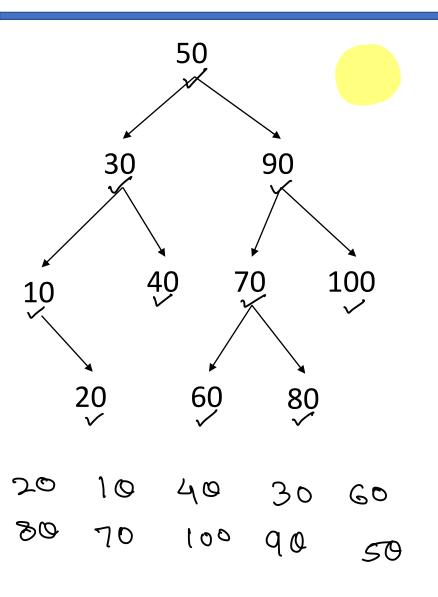
```
node* trav = root;
while (trav != NULL II !s.empty()) {
   while (trav != NULL) {
       s.push(trav);
       trav = trav->left;
    if (!s.empty()) {
      trav = s.top();
      s.pop();
      cout << trav->data << ", ";
       trav = trav->right;
```

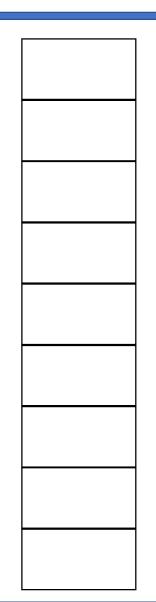


BST – Non-Recursive Algorithm – PostOrder

LRV



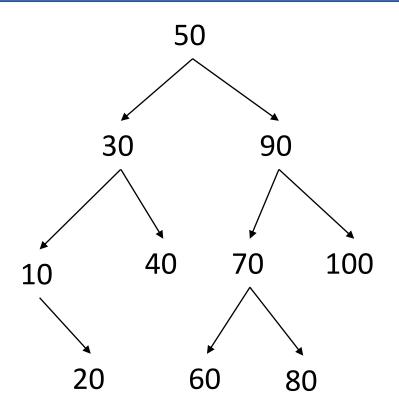




toda = 200 pt) copile four not sull as stack not subjy: until null is reached. push tour on stack goto left of tour if stade is not empty: bob moge juto toan if right node is present a visited: point node mark mode as visited on ofce toar onl. er &: push node on stack 50 to the right



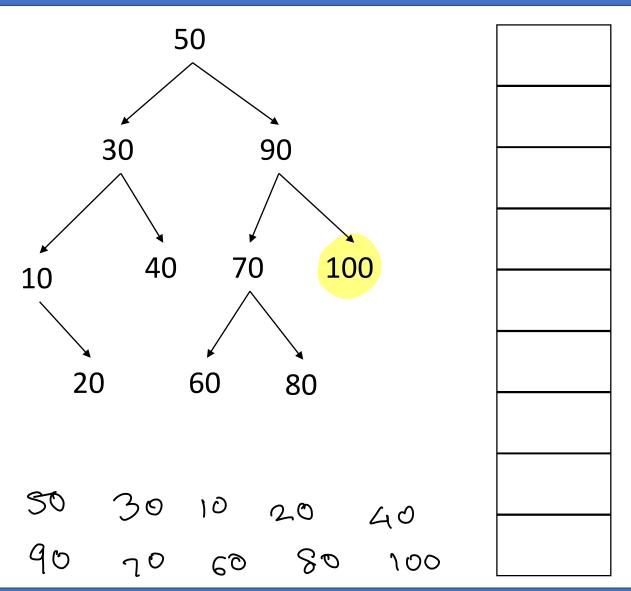
BST – Non-Recursive Algorithm – PostOrder



```
node* trav = root;
while (trav != NULL | !s.empty()) {
   while (trav != NULL) {
       s.push(trav);
       trav = trav->left;
   if (!s.empty()) {
      trav = s.top(); s.pop();
      if(trav->right±= NULL & trav->right->visited) {
          cout << trav->data << ", ";
          trav->visited = true;
          trav = NULL;
      } else {
          s.push(trav);
          trav = trav->right;
```



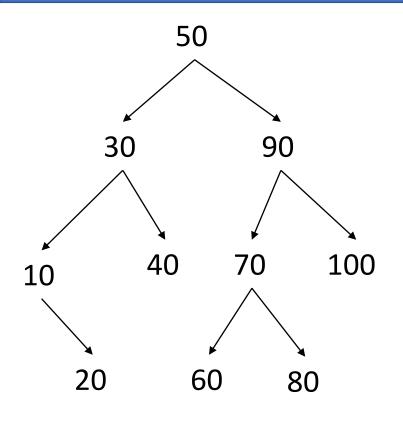
B\$T - Non-Recursive Algorithm - DFS - any broany tree (non-served)



push root an stack while stade is not empty: bob worde jugo par risit it if tour has right, it four has left, push it on stack.



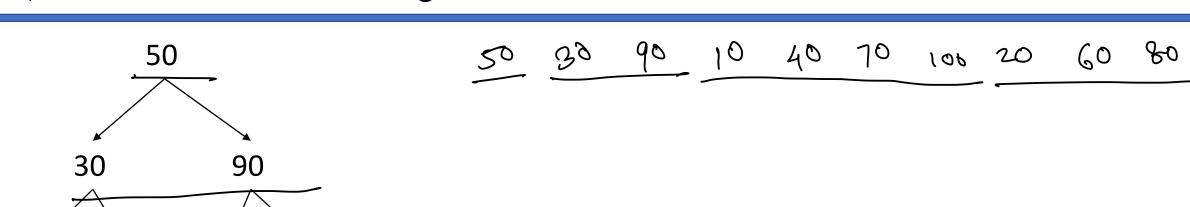
BST – Non-Recursive Algorithm – DFS

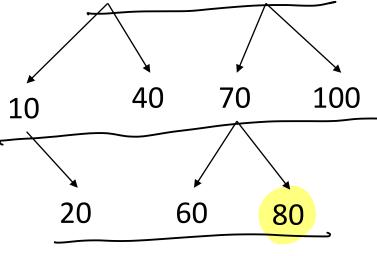


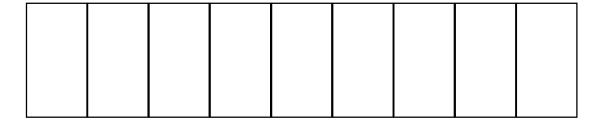
```
s.push(root);
while (!s.empty()) {
   trav = s.top();
   s.pop();
   cout << trav->data << ", ";
   if (trav->right != NULL)
       s.push(trav->right);
   if (trav->left != NULL)
       s.push(trav->left);
```





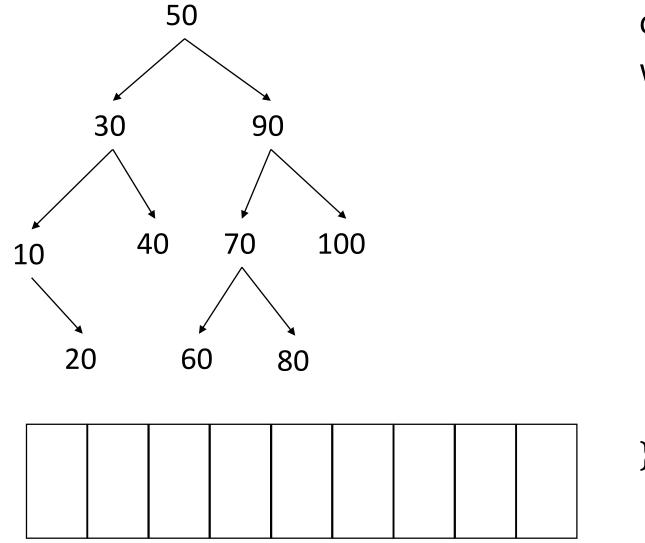








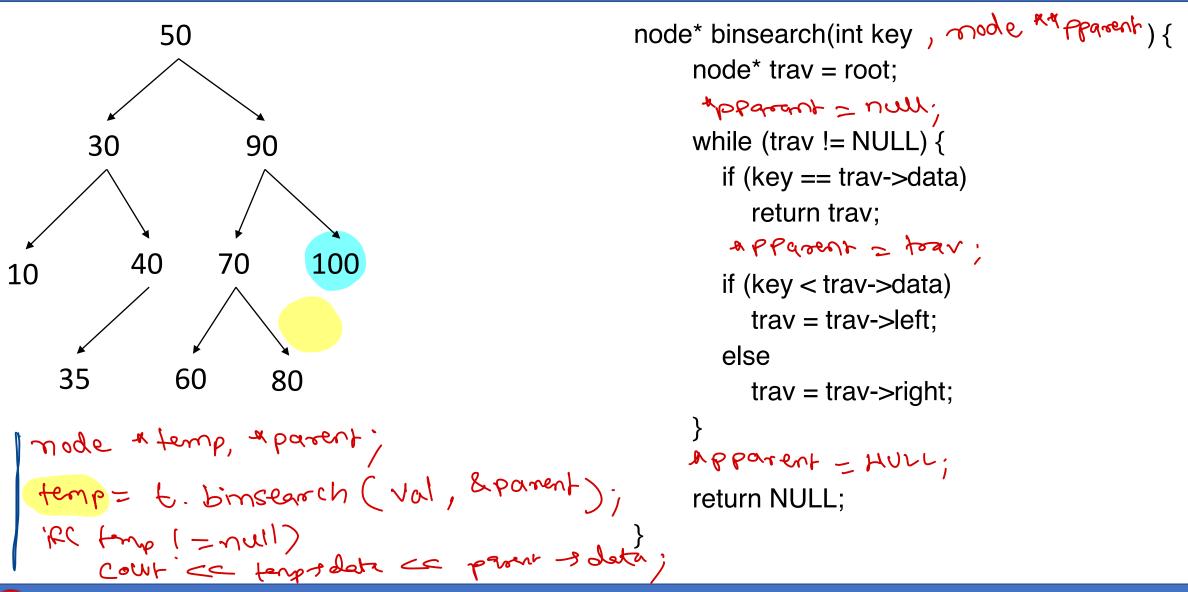
BST – Non-Recursive Algorithm – BFS



```
q.push(root);
while (!s.empty()) {
   trav = q.front();
   q.pop();
    cout << trav->data << ", ";
    if (trav->left != NULL)
       q.push(trav->left);
    if (trav->right != NULL)
       q.push(trav->right);
```

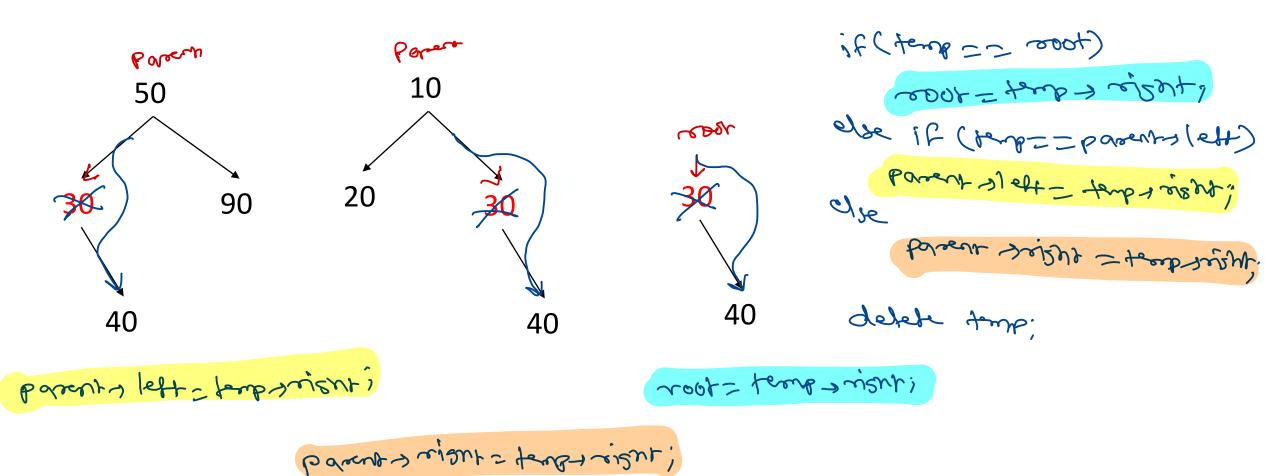


BST – Find node with its Parent

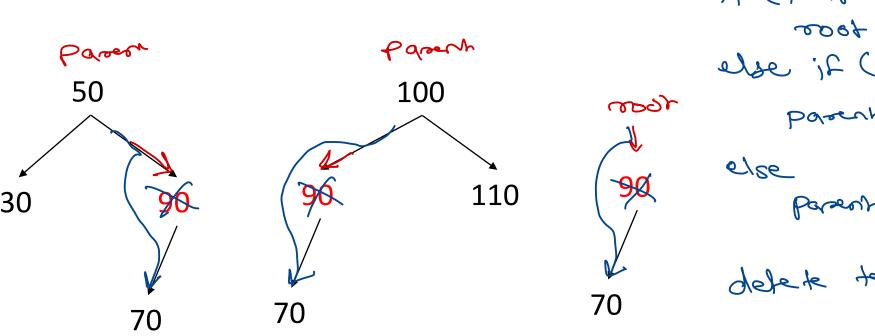


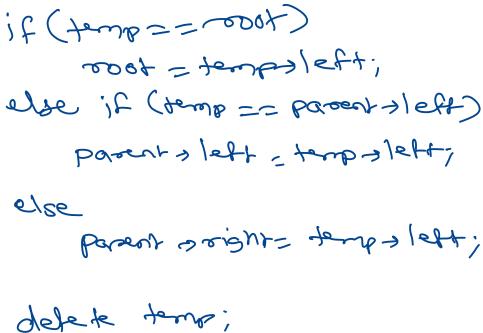


BST - Delete Node - whose left child is out. (temp)

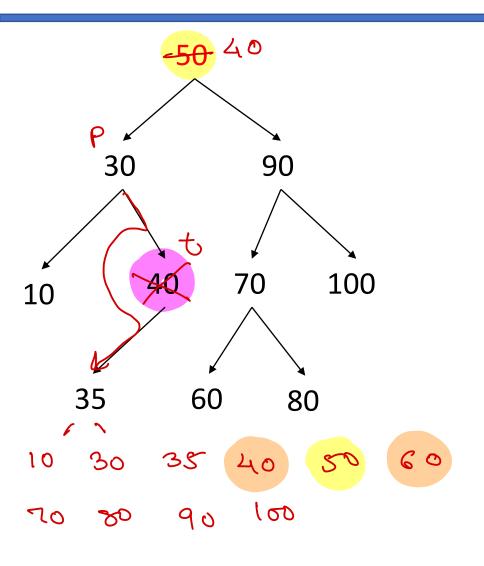






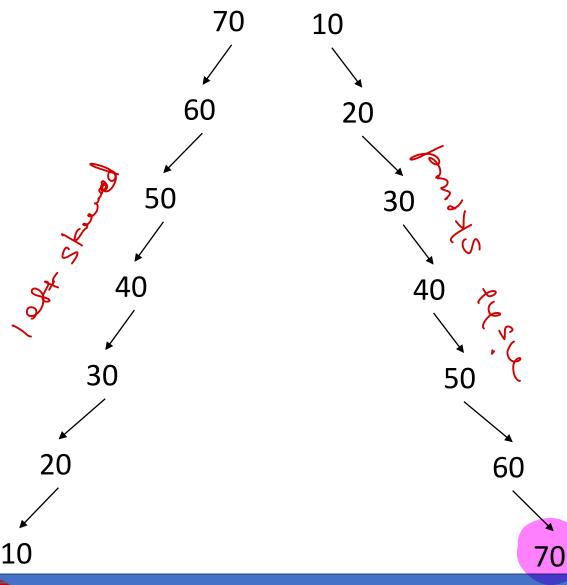






baseur = temp; proed = temp > left; while (pred = risnr! = onul) { based = begi pred= pred - right; temp > deta = poed > data. temp = poed;

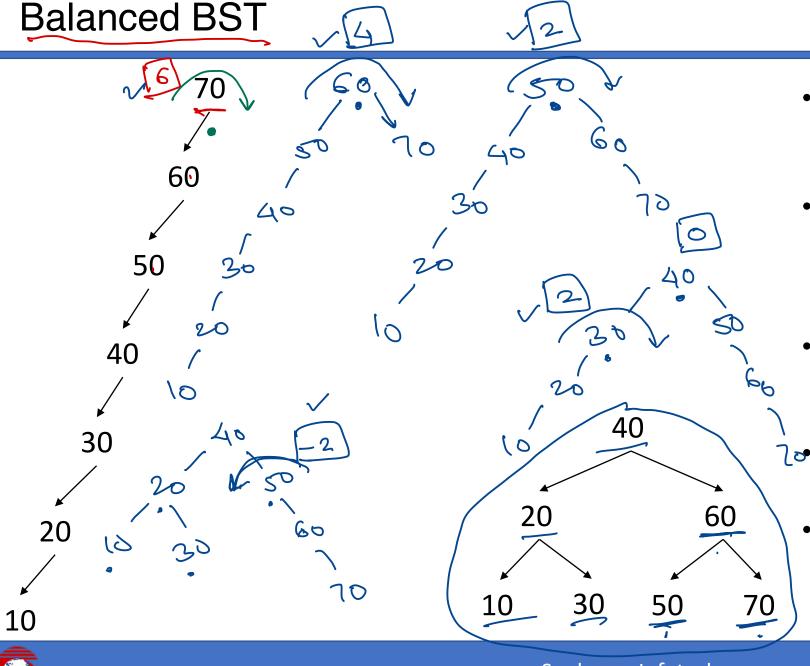
Skewed Binary Tree



- In Binary tree if only left or only right links are used, tree grows only on one side. Such tree is called as skewed binary tree.
 - Left skewed binary tree
 - Right skewed binary tree
- Time complexity of any BST is O(h).
- Such tree have maximum height i.e. same as number of elements.
- Time complexity of searching in skewed BST is O(n).

Is like linked list

heishh



- To speed up searching, height of BST should minimum as possible.
- If nodes in BST are arranged so that its height is kept as less as possible, is called as Balanced BST.
- Balance factor
 - = <u>Height of left sub tree</u> <u>Height</u> of left sub tree
- In balanced BST, BF of each node is -1, 0 or +1.
- A tree can be balanced by applying series of left or right rotations on unbalanced nodes.





Thank you!

Nilesh Ghule <nilesh@sunbeaminfo.com>

