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Characterization and complexity of
Thin Strip Graphs

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“You may also include one or more general quotes related to your topic.”

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“Another quote.”

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Acknowledgment

I want to thank ...

Contents

1	Background	2
1.1	Graphs and intersections	2
1.1.1	Graphs	2
1.1.2	Intersection graphs	4
1.2	Complexity	6
1.2.1	P vs NP	7
1.2.2	$\exists\mathbb{R}$ complexity class	8
1.3	Geometry	11
1.3.1	Stabbing	11
1.3.2	Coin graphs	12
2	Interval graphs	13
2.1	Interval graphs	13
2.1.1	Unit interval graphs	14
2.2	Mixed unit interval graphs	14
2.2.1	Characterization	16
2.3	Unfettered unit interval graphs	18
2.3.1	Recognition	19
3	Thin strip graphs	21
3.0.1	Interval graphs	21
3.1	Characterization of thin strip graphs	22
3.2	Recognition	23

4	Thin two-level graphs	24
4.1	Characterization of TTL	24
4.1.1	Relation with MUIG	25
4.2	Induced forbidden subgraphs	26
5	Complexity	27
5.1	Recognizing Thin Strip Graphs	27
	Appendices	29
A	Graph classes hierarchy	30
B	Problems in inclusion	32
C	Problems in forbidden induced subgraph characterization	33
D	Problems in complexity	34

Introduction

Talk about what this is etc...

Chapter 1

Background

1.1 Graphs and intersections

1.1.1 Graphs

A graph G is defined as $G = (V, E)$, where V is the set of vertices and E the set of edges, where $E \subseteq \binom{V}{2}$. The vertices $v, w \in V$ such that $e = vw \in E$ links are called the *endpoints* of e .

Definition 1.1.1. An embedding of a graph G into a surface Σ is a mapping of G in Σ where the vertices correspond to distinct points and the edges correspond to simple arcs connecting the images of their endpoints. [GF17].

A graph G is planar if there is an embedding of this graph that does not have any crossing between the edges.

Definition 1.1.2. Let $G = (V, E)$ and $S \subset V$, an induced subgraph is a graph H of G whose vertex set is S and its edge set $F = \{vw : v, w \in S, vw \in E\}$.

Definition 1.1.3. Let $G = (V, E)$ its complement graph \overline{G} is the graph such that its edge set is defined as: $\{vw : v, w \in V, vw \notin E\}$.

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Definition 1.1.4. H is called a *minor* of G if H can be constructed by deleting edges and vertices, or contracting edges.

Theorem 1.1.5 (Kuratowski). *A graph G is planar if and only if it does not contain K_5 or $K_{3,3}$ as a minor or a induced subgraph.*

Definition 1.1.6. A path P_n in a graph G is a sequence of vertices $v_1 v_2 v_3 \dots v_n$ such that $v_i v_{i+1} \in E$.

Definition 1.1.7. A cycle C_n in a graph $G = (V, E)$ is a path $v_1 \dots v_n$ such that $v_1 = v_n$.

Definition 1.1.8. A chord of a cycle C_n with $n \geq 4$ is an edge that connects two non consecutive vertices of C_n .

Definition 1.1.9. A triangular chord of a cycle is a chord that will create a new triangle (C_3).

Definition 1.1.10. A graph $G = (V, E)$ is complete if every pair of distinct $v_1, v_2 \in V$ are adjacent. This is denoted K_n with n the size of the graph. If G is an induced graph of H then G is a clique of H .

Definition 1.1.11. A graph G is bipartite if there exist two disjoint subsets $A, B \subset V$ such that $A \cup B = V$ and each edge $e \in E$ has an endpoint on A and the another on B .

Definition 1.1.12. A bipartite graph G with bipartitions A and B is complete bipartite if every pair of vertices $v \in A, w \in B$ are adjacent. It is denoted as $K_{n,m}$, being n and m the size of each bipartition.

Some graphs can be characterized with properties. A property of a graph is a property that is preserved under all its isomorphisms. These properties are called *hereditary* if they are also preserved under all its induced subgraphs; they are called *minor-hereditary* if they are also preserved under its minors (e.g. Kuratowski's planar graph characterization [1.1.5]).

Definition 1.1.13. An forbidden induced subgraph (minor) of a graph class X is a graph such that if it is the induced subgraph (minor) of a graph G , we know that $G \notin X$.

The coloration of a graph is a color assignment to each vertex such that the color of the two endpoints of every edge of the graph is different.

Definition 1.1.14. The chromatic number of a graph $\chi(G)$ is the smallest number of colors needed to have an acceptable coloration of G .

Definition 1.1.15. The clique number of a graph $\omega(G)$ is the size of the biggest clique of G . We can observe that for every graph: $\chi(G) \geq \omega(G)$.

Definition 1.1.16. A perfect graph is a graph that respects this condition for every induced subgraph:

$$\omega(G) = \chi(G)$$

Theorem 1.1.17 (Lovasz). *G is perfect if and only if \overline{G} is perfect.*

1.1.2 Intersection graphs

Definition 1.1.18. The *intersection graph* of a collection ζ of objects is the graph (ζ, E) such that $c_1 c_2 \in E \Leftrightarrow c_1 \cap c_2 \neq \emptyset$.

An intersection can be seen as a relationship between two objects. In this thesis, it will be important to define these relations more formally to characterize intersection graphs.

Definition 1.1.19. A partial order is a binary relation \leq over a set A satisfying these axioms:

- if $a \leq b$ and $b \leq a$ then $a = b$ (antisymmetry).
- if $a \leq b$ and $b \leq c$ then $a \leq c$ (transitivity).
- $a \leq a$ (reflexivity).

Definition 1.1.20. A total order is a partial order where the reflexivity order is replaced by the connex property:

$$a \leq b \text{ or } b \leq a$$

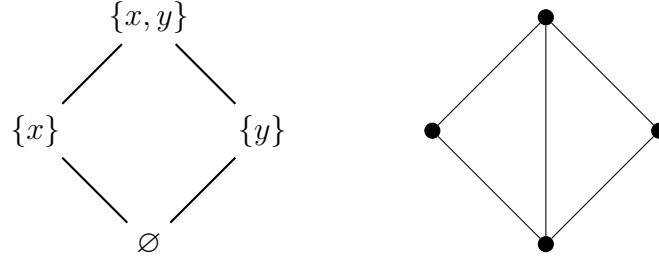


Figure 1.1: On the left, Hasse diagram of a poset of the power set of 2 elements ordered by inclusion. On the right, the comparability graph of this poset.

Definition 1.1.21. A partially ordered set or poset (S, \leq) where S a set and \leq a partial order on S .

Definition 1.1.22. A spanning order $(V, <)$ of a graph $G = (V, E)$ is a total order on V such that for any three vertices $u < v < w$:

$$uw \in E \rightarrow uv \in E \text{ or } vw \in E$$

Definition 1.1.23. A graph $G = (V, E)$ is a comparability graph if there exists a partial order \leq such that $vw \in E \Leftrightarrow v \leq w$ or $w \leq v$. Equivalently, G is a comparability graph if it is the comparability graph of a poset. For example, the Hasse diagram (figure 1.1) is a comparability graph where the relation is inclusion.

Definition 1.1.24. A graph $G = (V, E)$ is a co-comparability graph if its complement is a comparability graph.

There are multiple characterizations for the co-comparability graph class; we will see one that uses a poset to characterize it:

Theorem 1.1.25 (Damaschke [Dam92]). *A graph G is a co-comparability graph if and only if it has a spanning order.*

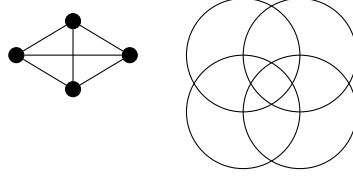


Figure 1.2: Realization of a UDG (Unit Disk Graph).

Disks

A disk graph G is a graph that is an intersection graph of disks on the plane, when the size of the disk is unitary, we talk about unit disk graphs. This class of graphs is important for this thesis, as thin strip graphs are a sub-class of unit disk graphs.

We will refer to the unit disk graph class as UDG and an example of a realization can be found in the figure 1.2.

Induced forbidden subgraphs The characterization of this class with respect to its induced forbidden subgraphs has been studied [AZ16].

Theorem 1.1.26 (Atminas-Zamaraev). *For every integer $k > 1$, $\overline{K_2 + C_{2k+1}}$ is a minimal induced subgraph of UDG.*

Theorem 1.1.27 (Atminas-Zamaraev). *For every integer $k > 4$, $\overline{C_{2k}}$ is a minimal induced subgraph of UDG.*

1.2 Complexity

Complexity theory has the objective to establish lower bounds on how efficient an algorithm can be for a given problem [Sip06]. This approach let us have a reference point to establish the difficulty of a problem.

Definition 1.2.1. Let Σ be a finite alphabet, Σ^* every word derived from Σ , $L \subseteq \Sigma^*$ is a decision problem.

Definition 1.2.2. A decider for a decision problem A is an deterministic algorithm V where

$$A = \{w | V \text{ accepts } w\}$$

A is polynomially decidable if it has a polynomial time decider [Sip06].

Definition 1.2.3. A verifier for a decision problem A is an deterministic algorithm V where

$$A = \{w | V \text{ accepts } \langle w, c \rangle \text{ for some string } c\}$$

A is polynomially verifiable if it has a polynomial time verifier [Sip06].

1.2.1 P vs NP

Definition 1.2.4. A problem $L \in \mathcal{P}$ if L is polynomially decidable.

Definition 1.2.5. A problem $L \in \mathcal{NP}$ if L is polynomially verifiable. Thus, $\mathcal{P} \subseteq \mathcal{NP}$.

To prove a bound of complexity on an unknown problem L we have to find another problem with already known complexity and find equivalences between those two. This can be achieved through *reductions*.

Definition 1.2.6. A reduction of a problem L to a problem M is a mapping of an instance of L (I_L) to an instance of M (I_M) such that I_L is true for the problem L if and only if I_M is true for the problem M . This is noted $L \leq M$ and $L \leq_P M$ if the reduction is done in polynomial time.

With this concept we can define new complexity classes. \mathcal{NP} -hard is the set of problems so that we can reduce every \mathcal{NP} problem to. The set of problems that are both \mathcal{NP} -hard and \mathcal{NP} are called \mathcal{NP} -complete. This is generalized to every complexity class (\mathcal{P} , $\exists\mathbb{R}$, RP , etc...)

Satisfiability problem The satisfiability problem (SAT) is to decide the satisfiability of a CNF formula ϕ . A CNF formula is a boolean formula that is a conjunction of multiple clauses c_k . A clause is a disjunction of multiple literals. A literal may be a variable or a negation of a variable.

Theorem 1.2.7 (Cook-Levin). *SAT is \mathcal{NP} -complete.*

Clique problem The clique problem is to find a maximum clique of a graph G .

Theorem 1.2.8. *CLIQUE is \mathcal{NP} -complete. [Kar72]*

Theorem 1.2.9. *CLIQUE is QPTAS when applied to disk graphs. [BGK⁺17]*

Theorem 1.2.10 (Clark-Colbourn). *CLIQUE is \mathcal{P} when applied to unit disk graphs. [CCJ90]*

1.2.2 $\exists\mathbb{R}$ complexity class

$\exists\mathbb{R}$ is the class that describes the problems that can be reduced to *the existential theory of the reals*[Exi06a]. The existential theory of the reals is the problem of deciding if a sentence of this form is true:

$$(\exists X_1 \dots \exists X_n) : F(\exists X_1, \dots, \exists X_n)$$

where F is a quantifier-free formula in the reals. In other words, it is a conjunction of clauses where each clause is a real polynomial inequality where each variable X_k is a real number. We can see that ETR is NP-hard because SAT can be reduced to it.

Proof. Let's take an instance of SAT ϕ_{SAT} with clauses c_k and variables x_k , we can construct an instance of ETR ϕ_{ETR} where we can construct variables in the domain $\{0, 1\}$ with this equality, so for each variable X_k :

$$X_k - X_k^2 = 0$$

Each literal of each clause will be positive or negative depending if the literal is cancelled in ϕ_{SAT} :

$$x_k \rightarrow l = X_k$$

$$\neg x_k \rightarrow l = (1 - X_k)$$

Then for each clause we can have a polynomial for which the sum of the values of every literal in the clause must be greater than one, so that at least one literal is true:

$$\sum_{l \in c_k} l \geq 1$$

With this proof, it is easy to see that ϕ_{ETR} is valid if and only if ϕ_{SAT} is also valid. □

□

This result can show us that $P \subseteq NP \subseteq \exists\mathbb{R}$.

Problems in $\exists\mathbb{R}$

In this section we will describe some problems that are $\exists\mathbb{R}$ -complete and will give an overview of the proof.

The art gallery problem Given a simple polygon P (without crossings between every side), we introduce *guards*. A guard g is a point such that every point of the polygon is watched by a guard. A point p is watched by a point q if the segment pq is contained in P . The subset G , being G the set of guards and $G \subseteq P$, is optimum if it has the minimal cardinality covering the whole polygon.

The art gallery problem is to decide, given a polygon P and a number of guards k , whether there exists a configuration of k guards in G guarding the whole polygon. The art gallery problem is $\exists\mathbb{R}$ -complete [AAM17].

Proof idea First of all, we can see that the art gallery problem is in $\exists\mathbb{R}$ if we reduce this problem to ETR. If we have an instance (P, k) of the art gallery problem we can have a formula [EH06] like this:

$$\phi = \{\exists x_1 y_1, \dots, x_k y_k \forall p_x p_y : \text{INSIDE-POLYGON}(p_x, p_y) \rightarrow \bigvee_{1 \leq i \leq k} \text{SEES}(x_i, y_i, p_x, p_y)\}$$

Where INSIDE-POLYGON returns 1 if $(p_x, p_y) \in P$ and SEES returns 1 if the segment $(x, y)(p_x, p_y) \in P$. ϕ is not a ETR formula, so we would like to construct a quantifier-free formula with the idea of ϕ . To achieve this, the main idea is to have a small set of points $Q \subseteq P$ such that if these points are watched, the whole polygon is watched. This subset Q is called the *witness set*. The only thing is now to create a polynomial for each point that ensures that the point is watched by a guard.

To finish the proof we have to prove that the art gallery problem is $\exists\mathbb{R}$ -hard. For this part an $\exists\mathbb{R}$ -complete problem has been deducted from ETR. For the problem ETR-INV we have a set of variables $\{x_1, \dots, x_n\}$ and a set of equations of this form:

$$x = 1, \quad x + y = z, \quad x \cdot y = 1$$

and the problem decides if it exists a solution to this set of equations such that the value of each variable is real in $[\frac{1}{2}, 2]$.

A reduction of ETR-INV is found to the art gallery problem by constructing a polygon P and finding a number g for that polygon such that the instance of ETR-INT is true if and only if P is covered by at most g guards.

Stretchability A pseudoline is a simple closed curve in the plane. The stretchability problem is to decide if given a pseudoline arrangement, it is equivalent to an arrangement of straight lines.

Proof idea ETR can be reduced to STRETCHABILITY due to Mnev's universality theorem. [Sch10]

Unit disk graph recognition The unit disk graph recognition is the problem that decides if a graph G is a unit disk graph. Unit disk graph recognition is $\exists\mathbb{R}$ -complete. [Sch13a]

Proof idea UDG recognition is a corollary of deciding whether a graph with a given length is realizable. This problem is $\exists\mathbb{R}$ -complete.

The reduction is done from STRETCHABILITY [Sch13a]. The reduction is done by adding a vertex to V for each pseudoline intersection. For each three consecutive points u_1, u_2, u_3 along a pseudoline a widget will be added that will be only realizable if and only if the pseudoline can be stretched with the same arrangement.

1.3 Geometry

Definition 1.3.1. $\text{dist}(a, b)$ denotes the distance between the points a and b and is calculated with:

$$\text{dist}(a, b) = \sqrt{(a_x - b_x)^2 + (a_y - b_y)^2}$$

The intersection of convex objects is a matter well studied for multiple subjects. In our case, it is interesting to know some properties about the intersection of disks, those ones being convex objects.

A set S is convex if:

$$\forall p, q \in S \ \forall \lambda \in [0, 1] : (1 - \lambda)p + \lambda q \in S$$

1.3.1 Stabbing

A *stabbing* is a point that traverses a set of intersecting objects. A lot of research has been done [Sch13b] on the minimal amount of stabbings to cover every object in a set. Stabbings can also be done with more complex structures than points, in that case we are talking about *coverings*.

Theorem 1.3.2 (Helly). *Given a set Q of objects in \mathbb{R}^d , if for each subset of Q of size $d + 1$ their intersection is non empty, then $\bigcap_{q \in Q} q \neq \emptyset$. [Hel23]*

Theorem 1.3.3. *The problem that for a set of n disks whether there exists a regular n -gon whose vertices stab every disk of the set can be decided in $O(n^{10.5}/\sqrt{\log(n)})$ [Sch13b]*

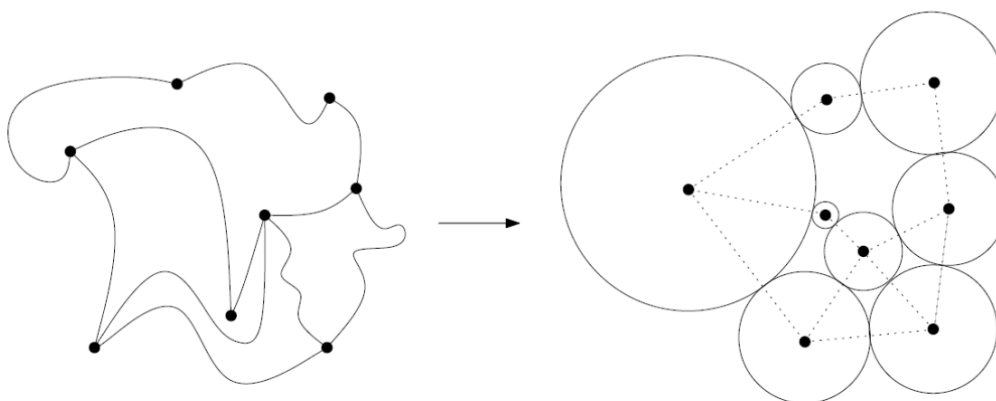


Figure 1.3: Circle packing of a planar graph. [Nac16]

1.3.2 Coin graphs

Penny graphs can be defined as disk graphs where the disks can just touch each other without overlapping. A famous theorem is derived from this class of graphs: the circle packing theorem.

Theorem 1.3.4 (Circle packing theorem). *The circle packing theorem states that every simple connected planar graph G is a penny graph. [BS93]*

Corollary 1.3.5. *Planar graphs \subseteq disk graphs [Spi12].*

Chapter 2

Interval graphs

In this chapter we are going to introduce the class of interval graphs, which is one of the most used classes of intersection graphs at all. There are multiple types of interval graphs and the most relevant for the thesis are going to be defined below.

The mixed unit interval graphs have been introduced by Dourado et al. [DLP⁺12a] and its characterization is given by Joos [Joo13]. The proof of the characterization will not be given because of the length of it, but each family of forbidden subgraphs will be presented.

Last, we present unfettered unit interval graphs, which have been defined by Hayashi et al. [HKO⁺17] while defining thin strip graphs in their paper.

2.1 Interval graphs

An interval graph is a graph G that is the intersection graph of a collection of closed intervals in \mathbb{R} .

First we present the main characterizations of interval graphs. In the next sections we present some other subclasses of interval graphs that will help us characterize the thin strip graphs on chapter 3. There are multiple characterizations of interval graphs that are equivalent, in this thesis we are going to present only one, which is the most relevant for our research:

Theorem 2.1.1. *G is an interval graph if and only if G does not contain C_4 as an induced subgraph and \overline{G} has a transitive orientation (it is a co-comparability graph). (Gilmore and Hoffman [GH64])*

Proof. We want to prove that G is an interval graph if and only if it does not contain C_4 and it is

2.1.1 Unit interval graphs

When every interval has the same length (or *unitary*), the intersection graph of this interval set is referred to as a unit interval graph (or UIG). Roberts [Rob68] shows in his paper about indifference relations a characterization of UIG and an interesting equivalence with another class:

Theorem 2.1.2 (Roberts [Rob68]). *A graph is a unit interval graph if and only if it is a proper interval graph (an interval graph where no interval is a strict subset of another).*

Theorem 2.1.3 (Roberts [Rob68]). *An interval graph is a unit interval graph if and only if it has no induced subgraph $K_{1,3}$ (or claw).*

2.2 Mixed unit interval graphs

The next class of interval graphs that we present are mixed unit interval graphs, where each interval is unitary and can be closed, open, open-closed or closed-open.

In this paper we define four classes of unitary interval graphs:

$$\mathcal{U}^{++} = \{[x, y] : x, y \in \mathbb{R}, x \leq y\}$$

$$\mathcal{U}^{--} = \{(x, y) : x, y \in \mathbb{R}, x \leq y\}$$

$$\mathcal{U}^{+-} = \{[x, y) : x, y \in \mathbb{R}, x \leq y\}$$

$$\mathcal{U}^{-+} = \{(x, y] : x, y \in \mathbb{R}, x \leq y\}$$

where $\mathcal{U}^{++} = \text{UIG}$.

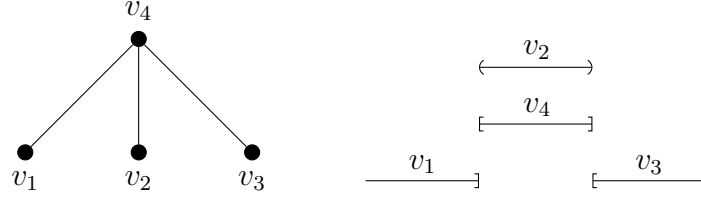


Figure 2.1: Representation of $K_{1,3}$ as a MUIG.

Theorem 2.2.1 (Dourado et al. [DLP⁺12b]). *The classes of the graphs \mathcal{U}^{--} , \mathcal{U}^{++} , \mathcal{U}^{-+} , \mathcal{U}^{+-} , and $\mathcal{U}^{-+} \cup \mathcal{U}^{+-}$ are the same.*

We can already show an equivalence between UIG and MUIG:

Theorem 2.2.2 (Dourado et al. [DLP⁺12a]). *$\text{UIG} \subsetneq \text{MUIG}$.*

Proof. The strict inclusion is straightforward: we know that $\text{UIG} = \mathcal{U}^{++} \subset \text{MUIG}$ by definition.

And we only have to find a forbidden subgraph in UIG that is a MUIG. By theorem 2.1.3 we have $K_{1,3}$ forbidden in UIG, but it is in fact in MUIG if we take an interval set a, b, c, d such that:

- $a = [x, x + 1]$
- $b = (x, x + 1)$
- $c = [x + 1, x + 2]$ (or $[x + 1, x + 2)$)
- $d = [x - 1, x]$ (or $[x - 1, x)$)

for all $x \in \mathbb{R}$ (see Figure 2.1). □

A complete characterization by induced forbidden subgraphs have been found independently by A. Schuchat et al. [SSTW14a] and F. Joos [Joo13]. However, the Schuchat paper gives a polynomial algorithm to recognize MUIGs:

Theorem 2.2.3 (Schuchat et al. [SSTW14b]). *The MUIG recognition problem is in \mathcal{P} . Moreover, there is an algorithm that solves it in $O(|V|^2)$ for V the vertex set of a graph.*

2.2.1 Characterization

In this section we will go over the characterization of MUIG given by Joos with forbidden subgraphs. We will also review each one of these forbidden subgraphs and discuss them:

Theorem 2.2.4 (Joos [Joo13]). *G is a MUIG if and only if it is a $\{F\} \cup \mathcal{R} \cup \mathcal{S} \cup \mathcal{S}'' \cup \mathcal{T}$ -free interval graph.*

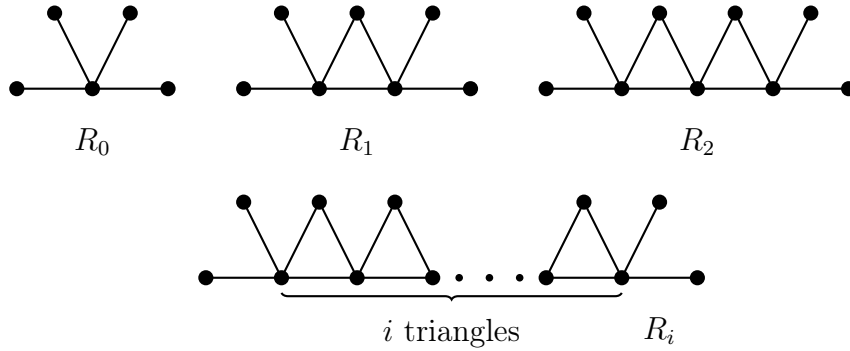


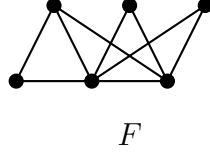
Figure 2.2: The class \mathcal{R} . [Joo13]

Except F , every family of forbidden graphs of MUIG is infinite, and is defined recursively by its precedent: then every property of these graphs has to be proved recursively. We begin first with \mathcal{R} .

Lemma 2.2.5. *\mathcal{R} is a family of co-comparability graphs.*

Proof. If we recall Theorem 1.1.25, in order to prove that \mathcal{R} is a family of co-comparability graphs we will have to find a spanning order for every R_i with $i \geq 0$. We will proceed to label our vertices with a mapping function $f : V \rightarrow \mathbb{N}$ such that $f(v) \in [1, |V|]$. This mapping will give us a spanning order by induction:

- $i = 0$: We assign the number 1 to the vertex with maximum degree v_1 . We assign then the rest of the numbers to the other vertices. We see then that $\forall u < v < w : uw \in E \rightarrow uv \in E$ because every vertex is adjacent to v_1 .

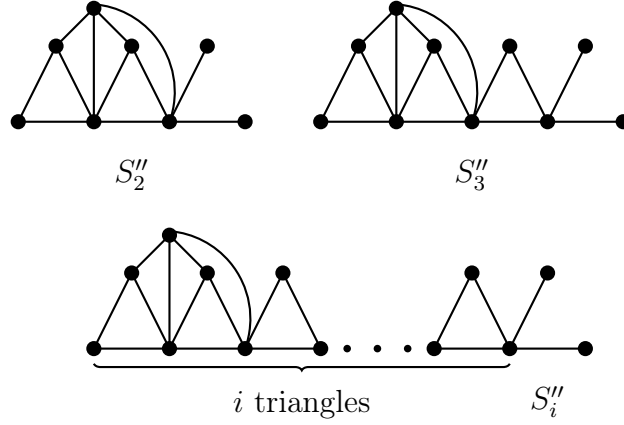
Figure 2.4: The graph F . [Joo13]

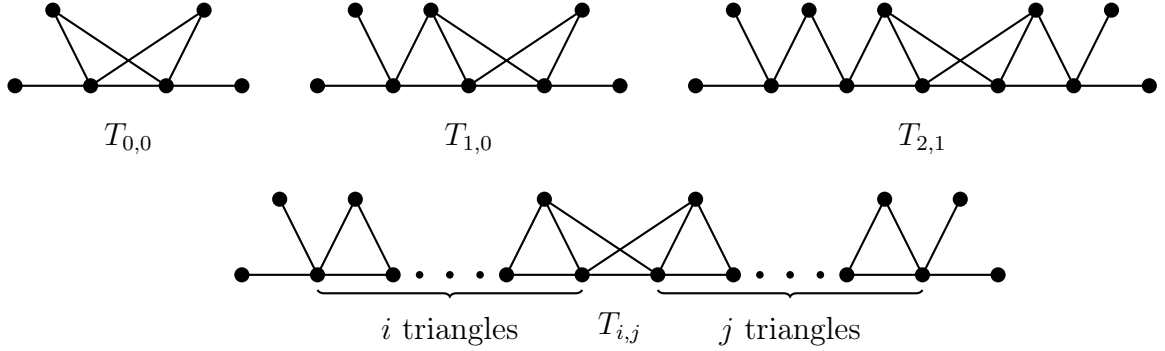
- $i = i + 1$: We define $\lambda_i = 5 + 2i$ where $\lambda_i = |V(R_i)|$. We add two vertices on each graph, where their labels are $\lambda_i + 1$ and $\lambda_i + 2$ and we also add three new edges: $v_{\lambda_i}v_{\lambda_i-1}, v_{\lambda_i}v_{\lambda_i+1}, v_{\lambda_i}v_{\lambda_i+2} \in E$.

By induction we only have to see if it holds with the new edges. We can say that it still holds with $v_{\lambda_i}v_{\lambda_i-1}$ and $v_{\lambda_i}v_{\lambda_i+1}$ because:

$$\nexists k \in \mathbb{N} : i < k < i + 1$$

Finally, we see that $v_{\lambda_i}v_{\lambda_i+2}$ is a valid edge because $v_{\lambda_i}v_{\lambda_i+1} \in E$. \square

Figure 2.3: The class \mathcal{S}'' . [Joo13]

Figure 2.5: The class \mathcal{T} . [Joo13]

2.3 Unfettered unit interval graphs

An unfettered unit interval graph can be defined as an unit interval graph such that for every touching endpoints we can chose either if they are adjacent or not.

Hayashi has characterized this class of graphs by levels. A *level structure* of a graph $G = (V, E)$ is a partition $L = \{L_i : i \in [1, t]\}$ of V such that

$$v \in L_k \rightarrow N(v) \subseteq L_{k-1} \cup L_k \cup L_{k+1}$$

where $L_0 = L_t + 1 = \emptyset$.

Theorem 2.3.1 (Hayashi et al. [HKO⁺17]). *A graph G is an unfettered unit interval graph if and only if it has a level structure where each level is a clique.*

We can clearly see that $\text{MUIG} \in \text{UUIG}$. However, we still have to see what is the location of UUIG in the higher graph classes hierarchy:

Proposition 2.3.2. *$\text{UUIG} \subset \text{co-comparability}$.*

Proof. This proposition is equivalent to say that a graph G is a UUIG if and only if it has a spanning order.

For each vertex of a partition L_k of UUIG (Theorem 2.3.1) we assign arbitrarily a number $i \in [\max(V(L_{k-1})) + 1, \max(V(L_{k-1})) + |V(L_k)| + 1]$;

intuitively, we assign every available number from the beginning in increasing order ($|V(L_1)|$ first numbers on the first partition and consecutively).

Because we know that each partition L_k is a clique, we can say that for each three vertices $u < v < w$, if $vw \in E \rightarrow uv \in E$ or $vw \in E$. We know this because given $u \in L_i$ and $w \in L_j$: if $uw \in E$ it means that levels L_i and L_j are adjacent, which means that $v \in L_i$ or $v \in L_j$ so v will be adjacent either to u or w . This is a spanning order. \square

If we recall the characterization of MUIG in section 2.2.1, we can see that every forbidden graph of MUIG is a UUIG (except for \mathcal{R}); which means that they are also co-comparability graphs.

2.3.1 Recognition

As we mentioned before, UUIG is a class of graphs very relevant to define TSG and that is why we have to know how this class of graphs is recognized. My first result in this thesis is to show that UUIG recognition is polynomial and can be solved with a recursive algorithm. The algorithm is based on the idea stated in the following lemma.

Lemma 2.3.3. *Let G be an UUIG with a level structure with levels L_1, \dots, L_n . $G \setminus L_i$ is a graph where each connected component is also an UUIG and the number of connected components is not bigger than two.*

Proof. By definition for a graph with a level structure, if $v \in L_i$, $N(v) = L_{i-1} \cup L_i \cup L_{i+1}$. This said, if we take off any level L_i , L_{i-1} and L_{i+1} are disconnected, but they are still connected to the other consecutive levels (L_{i-1} is connected to L_{i-2} , which is connected to L_{i-3} ... and viceversa with L_{i+1}).

And because a level is only connected to two different levels, we only have two connected components, only one if $L_i = L_1$. \square

We can then say by this lemma that we can design an algorithm that deletes a clique from the graph and then checks if the connected components of $G \setminus K_k$ are also UUIG until we arrive to the base case where $G \setminus K_k \cdots \setminus K_l$ is also

a clique (we will formalize the deleted clique by calling it H). Furthermore, that is not a necessary condition to recognize this graph. We also have to check if the vertices on the connected components ($C_1 \cup C_2 = G \setminus H$) are a clique (ie. if $v, w \in C_1$ and v and w are connected to H , then v and w also have to be adjacent).

You can read in the Algorithm 1. Our algorithm takes an arbitrary root set R as input because for each recursive call we have to begin our clique with the elements adjacent to H .

a finir
demain

Data: A graph $G = (V, E)$ and set of roots R

Result: G has a level structure where each level is a clique
not at end of this document

if *understand* **then**

 go to next section;
 current section becomes this one;

else

 go back to the beginning of current section;

end

return true;

Algorithm 1: UIG recognition

Chapter 3

Thin strip graphs

c -strip graphs are unit disk graphs such that the centers of the disks belong to $\{(x, y) : -\infty < x < \infty, 0 \leq y \leq c\}$. The class is denoted by $SG(c)$. We have $SG(0) = UIG$ and $SG(\infty) = UDG$. [HKO⁺17]

Definition 3.0.1. Thin strip graphs are defined as $TSG = \bigcap_{c>0} SG(c)$.

Remark 3.0.2. $SG(0) \neq TSG$. We can construct a $K_{1,3}$ such that we have 3 vertices with the coordinates $(1, 0)$, $(0, 0)$, $(1, 0)$ and a last one $(0, \varepsilon)$ with $\varepsilon > 0$ and arbitrarily small as seen in Figure 3.1.

Theorem 3.0.3 (Hayashi et al. [HKO⁺17]). *There is no constant t such that $SG(t) = TSG$.*

Since this class is newly defined we have to characterize it. For this purpose, some relations have been found between this class and interval graphs.

3.0.1 Interval graphs

Theorem 3.0.4 (Hayashi et al. [HKO⁺17]). $MUIG \subsetneq TSG$.

We can define a new class of graphs: unfettered unit interval graphs. These graphs are unit interval graphs where if two intersections touch, we can decide whether they intersect or not. We denote this class $UUIG$.

Introduction
of the
chapter.
Explain
Breu's
char-
acteri-
zation
of strip
graphs

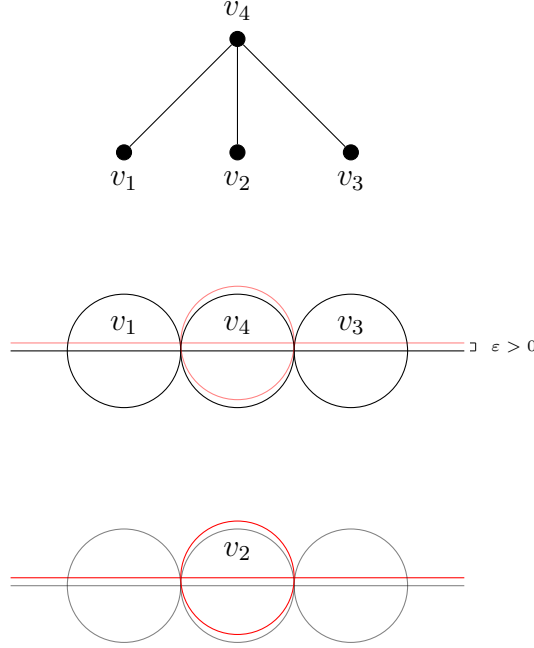


Figure 3.1: A construction of $K_{1,3}$ with a disk realization, being this graph a TSG.

Theorem 3.0.5 (Hayashi et al. [HKO⁺17]). $TSG \subsetneq UIG$.

3.1 Characterization of thin strip graphs

One of the main goals of this thesis is to characterize TSG. by forbidden induced subgraphs. To approach this, we will see how many induced forbidden subgraphs are also forbidden for TSG. We have described the families of forbidden induced subgraphs for UIG in section 2.2 and one of these families has been proven to be a forbidden induced subgraph for TSG:

Theorem 3.1.1 (Hayashi et al. [HKO⁺17]). \mathcal{R} is a forbidden induced subgraph family of TSG.

With the properties presented in this chapter, we can begin to state our first hypothesis:

Hypothesis 3.1.2. $F \in (UDG \cap UIG) \setminus TSG$.

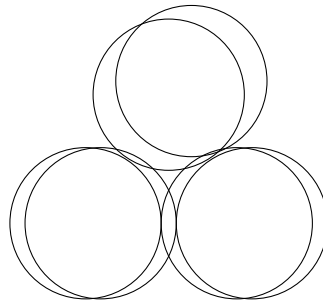


Figure 3.2: A construction of F with a unit disk realization.

Claim 3.1.3. $F \in (UDG \cap UUIG)$

Proof. We can see that $F \in UDG$ because we can have a unit disk realization (Figure 3.2) and also has a level structure $L = \{K_2, K_3, K_1\}$. \square

Theorem 3.1.4. $F \notin TSG$.

Proof. The distance between a and d has to be at most 2, because there is at least one vertex that is adjacent to both. And clearly, the other points have to be between them because they are adjacent to both, so if they were not.

this is
bullshit,
F is in
fact in
TSG...

3.2 Recognition

The recognition of this class of graphs is stated by Breu in his thesis .

explain
every-
thing
about
dan-
gerous
cycles
and com-
plement
oriented
graphs
in Breu's
paper

Chapter 4

Thin two-level graphs

Breu [Bre96] has presented in his thesis a similar class of constrained unit disk graphs where the disks are placed on k horizontal parallel lines. More formally: a disk (x, y) can be placed in $x \in (-\infty, \infty)$ and $y \in L$ with $|L| = k$.

In this chapter I define thin two-level graphs as a two-level graph where $L = \{0, \epsilon\}$ and ϵ is an arbitrarily small real number.

4.1 Characterization of TTL

Proposition 4.1.1. *An σ -SG(c) graph G (with $c < 1$) can be characterized by computing $\delta : A \times B \rightarrow E$ where $A, B \subseteq G$, A and B are UIG, and $A \cup B = \emptyset$:*

$$\delta(x, y) = \begin{cases} xy & \text{if } \text{dist}(x, y) \leq 1 \\ \emptyset, & \text{otherwise} \end{cases}$$

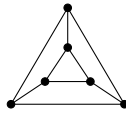


Figure 4.1: Forbidden graph in TTL

this proposition is ugly, re-do this with Breu's notations.

This class of graphs is close to our main class TSG. But we have to know at what point we can rely in this class of graphs to study TSG:

Lemma 4.1.2 (Breu [Bre96]). *Let $abcd$ be a chordless 4-cycle in a two-level graph $G = (V, E)$. Then ad and bc are level edges (they are adjacent in the same level), and the others are cross edges for every realization ϕ of G for which $\phi(a)_x < \phi(c)_x$ and $\phi(b)_x < \phi(d)_x$.*

With this preliminary result, we proceed to one of our main results:

Theorem 4.1.3. $TTL \subsetneq TSG$

Proof. By definition, we know that $TTL \subset TSG$ because the area where the disks can be placed in TTL is included in the area in TSG.

We can prove that $TTL \neq TSG$ because we can construct a graph G such that $G \in TSG$ and $G \notin TTL$. This graph D is a net* graph as described in Figure 4.1.

Part 1. D is a TSG because we can realize it as a TSG if we take as center of disks $(0, 0)$, $(0, z)$, $(0, \epsilon)$, $(1, 0)$, $(1, z)$, $(1, \epsilon)$ such that $0 < z < \epsilon$.

Part 2. Now we have to prove that D is a forbidden induced subgraph of TTL. We will try to construct it by taking a induced subgraph that is realizable: we take $D_{-1} = D - x$ with $x \in V(D)$. We notice that $V(D_{-1})$ is a chordless C_4 ($abcd$) with a vertex e adjacent to any two consecutive vertices $x, y \in V(C_4)$ creating the triangle xye .

By Lemma 4.1.2 we know that $abcd$ is a cycle if ab and cd are level edges. We can classify these vertices in two sets: $\ell(V) = a, d$ and $r(V) = b, c$ where $\forall u \in \ell(V) v \in r(V) : \phi(u)_x < \phi(v)_x$.

To realize D_{-1} we have to add a vertex e to C_4 . In our case, e cannot be added in two

Finish
proof
here

4.1.1 Relation with MUIG

A big question that was asked during this research is: What is the relationship between TTL and MUIG? We know that the net*-graph is forbidden for every two-level graph, this graph is also forbidden for MUIG because it includes

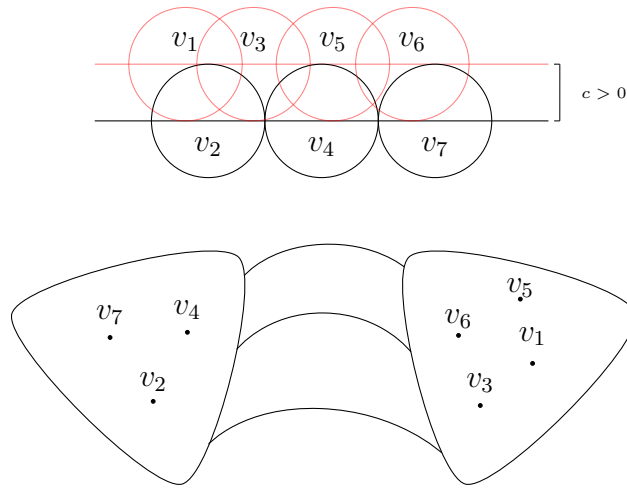


Figure 4.2: A representation of a $TL(c)$

an induced C_4 ; we can also see that every forbidden induced subgraph for MUIG also is for TTL.

However, C_4 is realizable for two-level graphs with $c = 1$ in general, which means $TTL \neq \text{interval} = \text{MUIG}$. We then know that $TTL \cup \text{MUIG} \subseteq \text{TSG}$ with $TTL \neq \text{MUIG}$.

Rewrite this better.

4.2 Induced forbidden subgraphs

Chapter 5

Complexity

5.1 Recognizing Thin Strip Graphs

Conclusions

The conclusions are to be written with care, because it will be sometimes the part that could convince a potential reader to read the whole document.

Appendices

Appendix A

Graph classes hierarchy

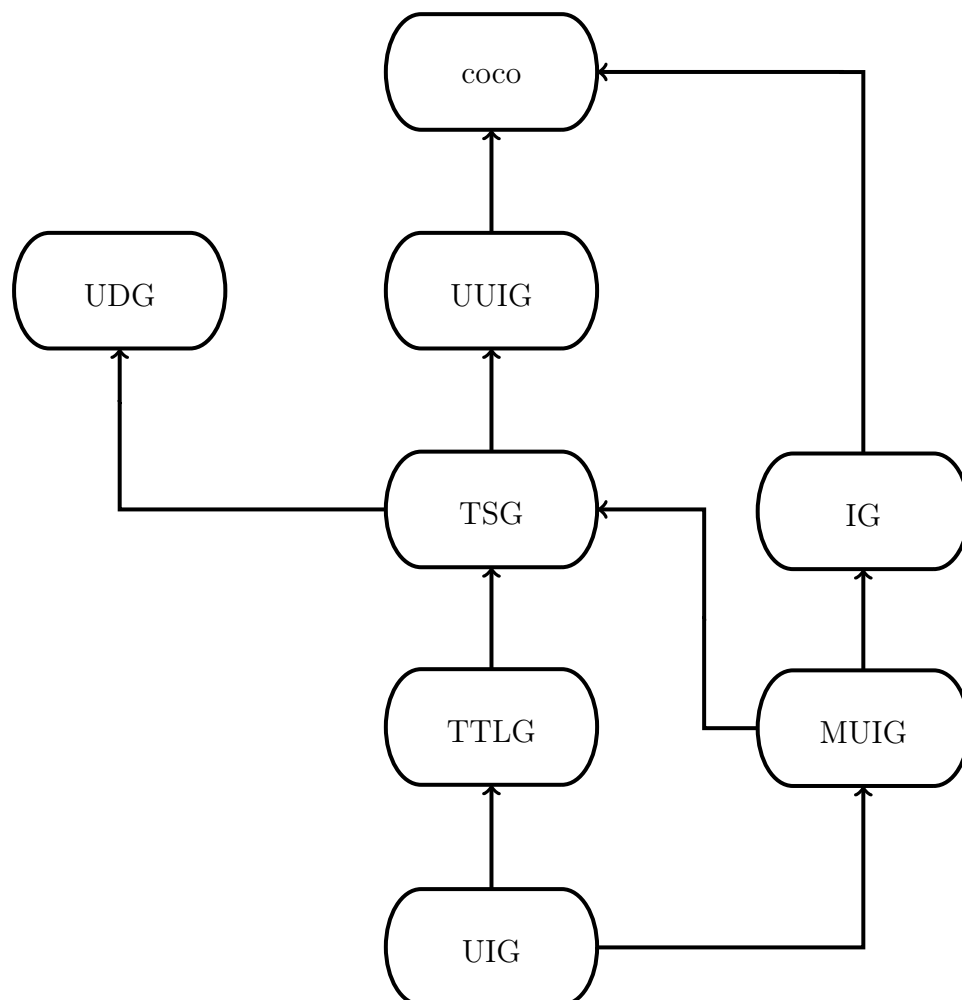


Figure A.1: A hierarchy of every relevant graph of this document. The relation $\text{class}_1 \rightarrow \text{class}_2$ means that $\text{class}_1 \subset \text{class}_2$.

Appendix B

Problems in inclusion

- **MUIG \subsetneq TSG \subsetneq UUIG** : Hayashi [HKO⁺17]
- **MUIG \neq TTLG (Open)**: To prove that MUIG \subsetneq TSG, Hayashi [HKO⁺17] could simulate MUIGs with 4 different levels. Having only two levels, I conjecture that this is not possible. However, MUIG can have C_4 , so an inclusion between these two classes is impossible (it has to be rewritten).
- **TTLG \subsetneq TSG (Open)**: This problem has been solved in my thesis by finding a forbidden graph for TTLG, theorem 4.1.3.
- **TLG \subset TSG (Open)**: This is a plausible stronger statement than the one before. However, this result could make the study of TTLG less relevant. Thus, this result would imply:

$$G \in \text{TLG}(j) \rightarrow G \in \text{SG}(k) : j, k \in \mathbb{R}$$

Appendix C

Problems in forbidden induced subgraph characterization

- **MUIG**: Joos [Joo13] gives us a complete characterization of forbidden graphs.
- **TSG (Open)**: Hayashi [HKO⁺17] says that MUIG's forbidden induced subgraphs also are in TSG. He claims that finding a graph $F \in (\text{UDG} \cap \text{UUIG}) \setminus \text{TSG}$ could be a good starting point. In my thesis I show that a forbidden induced subgraph for MUIG is in $\text{UDG} \cap \text{UUIG}$.
- **TTLG (Open)**: There are many properties about these graphs in Breu's thesis [Bre96].
- **UDG (Open)**: There is no complete characterization of UDG. Can the results of this thesis help find new ones?U

Appendix D

Problems in complexity

- **UIG/IG recognition:** Both of these problems are polynomial.
- **MUIG recognition:** Schuchat et al. give a linear algorithm ($O(|V|^2)$) to recognise MUIGs [SSTW14b].
- **UDG recognition:** $\exists\mathbb{R}$ -complete [Exi06b].
- **SG(c) recognition (Open):** Breu [Bre96] states that SG(c) recognition is polynomial if a complement edge orientation and a mapping $\phi : V \rightarrow [0, c]$ is polynomial as an input of the decision problem.
- **TSG recognition (Open):** Can we get rid of the mapping as input to recognise TSGs? In that case the problem would be at least NP.
- **UUIG recognition (Open):** Informally the recognition of this class of graphs **cannot** be polynomial because we have to find all the cliques of the graph; the CLIQUE problem is NP-complete.

Index

Care, 28

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Todo list

Text about background intro	2
a finir demain	20
Introduction of the chapter.	21
Explain Breu's characterization of strip graphs	21
this is bullshit, F is in fact in TSG...	23
explain everything about dangerous cycles and complement oriented graphs in Breu's paper	23
this proposition is ugly, re-do this with Breu's notations.	24
Finish proof here	25
Rewrite this better.	26