# ICPC Team Reference Material

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# 1 Setup

## 1.1 Vimrc

```
1 let mapleader = "\"
2 syntax on
3 filetype plugin on
4 set nocompatible
5 set autoread
6 set foldmethod=marker
7 set autoindent
8 set clipboard+=unnamedplus
9 set number relativenumber
10 set shiftwidth=2 softtabstop=2 expandtab
11 map <leader>c :w! && !compile %:p:r<CR>
12 vmap < <gv</pre>
13 vmap > >gv
```

# 1.2 Capslock as Escape

```
1 setxkbmap -layout us
2 xmodmap -e 'clear Lock'
3 xmodmap -e 'keycode 66 = Escape'
```

# 1.3 Compilation

# 2 Graph algorithms

#### 2.1 Adjacency list representation

```
1 class Graph {
 2 public:
      vector <int> _head, _next, _to, _cost;
      int edge_number;
      Graph() = default;
      Graph (int V, int E, bool isDirec) {
       isDirected = isDirec;
       _head.assign(V + 9, 0);
10
11
       _next.assign(isDirected ? E + 9 : E \star 2 + 9, 0);
       _to.assign(isDirected ? E + 9 : E * 2 + 9, 0);
12
        // _cost.assign(isDirected ? E + 9 : E * 2 + 9, 0);
13
       edge_number = 0;
14
15
16
17
      void addEdge(int u, int v, int w = 0) {
18
        _next[++edge_number] = _head[u];
        _to[edge_number] = v;
20
           _cost[edge_number] = w;
^{21}
        _head[u] = edge_number;
24
      void addBiEdge(int u, int v, int w = 0) {
       addEdge(u, v, w);
26
        addEdge(v, u, w);
27
      void dfs(int node) {
       vis[node] = true;
        for (int i = _head[node]; i; i = _next[i]) if(!vis[_to[i]]) {
      dfs(_to[i]);
35
```

## 2.2 Depth first search (DFS)

```
1  void DFS(int node)
2  {
3     vis[node] = true;
4     for(int i = Head[node]; i; i = Next[i])
5     if(!vis[To[i]);
6     DFS(To[i]);
7  }
```

## 2.3 Breadth first search (BFS)

```
1 void BFS(int src)
2
      memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
      memset (Par, -1, sizeof (Par[0]) * (n + 2));
      queue <int> q;
      q.push(src);
      dis[src] = 0;
10
1.1
      while (q.size())
12
          u = q.front(); q.pop();
13
         for(int i = Head[u]; i; i = Next[i]) if(dis[To[i]] == 00) {
14
       dis[To[i]] = dis[u] + 1;
Par[To[i]] = u;
15
16
       q.push(To[i]);
```

#### 2.4 0-1 BFS

```
1 #include <bits/stdc++.h>
 2 using namespace std;
   typedef int64_t 11;
 4 int main() {}
 6 const int N = 1e5 + 9, M = 2e5 + 9, oo = 0x3f3f3f3f3f;
   11 INF = 0x3f3f3f3f3f3f3f3f3f;
 9 int Head[N], Par[N], Next[M], To[M], Cost[M], ne, n, m, u, v, st, tr, tax;
10 11 dis[N];
11
12 void addEdge(int from, int to, int cost) {
        Next[++ne] = Head[from];
13
        Head[from] = ne;
14
        Cost[ne] = cost;
To[ne] = to;
1.5
17 }
19 void _clear() {
20
        memset(Head, 0, sizeof(Head[0]) * (n + 2));
21
22 }
23
24 void BFS(int src, int trg)
25 (
        memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
26
27
        memset(Par, -1, sizeof(Par[0]) * (n + 2));
28
29
        degue <int> 0:
        Q.push_front(src);
30
31
        dis[src] = 0;
34
        while(Q.size())
35
36
            node = Q.front(); Q.pop_front();
37
            if(node == trg) return;
38
39
            for(int i = Head[node]; i; i = Next[i])
40
                if(dis[node] + Cost[i] < dis[To[i]])</pre>
41
                    dis[To[i]] = dis[node] + Cost[i];
42
43
                    if (Cost [il)
                        Q.push_back(To[i]);
44
45
46
                        Q.push_front(To[i]);
47
49
```

## 2.5 0-1 BFS (grid)

```
1 #include <bits/stdc++.h>
 2 using namespace std;
   typedef int64_t 11;
    int main() {}
 6 const int dr[] = { -1, -1, 0, 1, 1, 1, 0, -1 };
7 const int dc[] = { 0, 1, 1, 1, 0, -1, -1, -1 };
8 const char dir[] = {'D', 'U', 'R', 'L'};
10 const int N = 1e3 + 9, M = 2e5 + 9, oo = 0x3f3f3f3f3f;
12 int dis[N][N], n, m, si, sj, ti, tj;
13 char grid[N][N];
14
15 bool valid(int r, int c) {
        return r >= 1 && r <= n && c >= 1 && c <= m;
16
17 }
18
19 /**
20 7 0 1
21 \//
22 6-*-2
24 5 4 3
25 **/
26
27 int ZBFS(int sr, int sc, int tr, int tc)
28
29
        memset(dis, 0x3f, sizeof (dis)); // memset(dis, 0x3f, n * m) we don't do that here
30
        deque <pair <int, int> > Q;
31
32
        dis[sr][sc] = 0;
33
34
        Q.push_front({sr, sc});
```

```
36
        int r, c, nr, nc, ncost;
37
        while(Q.size())
38
39
            tie(r, c) = Q.front(); Q.pop_front();
40
            if(r == tr && c == tc) return dis[r][c];
41
42
            for(int i = 0; i < 8; ++i)
43
44
                nr = r + dr[i];
45
                nc = c + dc[i];
46
47
                if(!valid(nr, nc)) continue;
48
                ncost = (i != grid[r][c]);
49
50
                if (dis[r][c] + ncost < dis[nr][nc])
                    dis[nr][nc] = dis[r][c] + ncost;
53
54
55
                        Q.push_back({nr, nc});
                    else
57
                        Q.push_front({nr, nc});
59
60
61
        return oo:
62
```

## 2.6 Articulation points and bridges

```
1 int Head[N], Next[M], To[M], Cost[M], Par[N], dfs_num[N], dfs_low[N], ne, n, m, u, v, w, root,
          rootChildren, dfs_timer, bridgeInx;
   bool Art[N];
 3
   vector < pair <int, int> > bridges(M);
 5
   void _clear() {
      memset(Head, 0, sizeof(Head[0]) * (n + 2));
 6
      memset(dfs_num, 0, sizeof(dfs_num[0]) * (n + 2));
     memset(Par, -1, sizeof(Par[0])
memset(Art, 0, sizeof(Art[0])
                                           * (n + 2));
* (n + 2));
     ne = dfs_timer = bridgeInx = 0;
10
11 }
12
13
   void Tarjan (int node) {
      dfs_num[node] = dfs_low[node] = ++dfs_timer;
      for(int i = Head[node]; i; i = Next[i]) {
16
        if(dfs_num[To[i]] == 0) {
17
          if(node == root) ++rootChildren;
          Par[To[i]] = node;
18
19
          Tarjan(To[i]);
20
          dfs_low[node] = Min(dfs_low[node], dfs_low[To[i]]);
21
          if(dfs_low[To[i]] >= dfs_num[node])
22
23
     Art[node] = true;
24
25
          if(dfs_low[To[i]] > dfs_num[node])
26
      bridges[bridgeInx++] = make_pair(node, To[i]);
27
        else if(To[i] != Par[node])
29
          dfs_low[node] = Min(dfs_low[node], dfs_num[To[i]]);
31
32
33
   int main() {
      for(int i = 1; i <= n; ++i)</pre>
35
        if(dfs_num[i] == 0) {
36
          root = i;
37
          rootChildren = 0;
38
          Tarjan(i);
          Art[root] = (rootChildren > 1);
39
40
41
      cout << "Art Points :\n";</pre>
      for(int i = 1; i <= n; ++i)
        if (Art[i])
         cout << i << " ";
46
48
      for(int i = 0; i < bridgeInx; ++i)
  cout << bridges[i].first << " - " << bridges[i].second << endl;</pre>
49
50
```

### 2.7 Bi-connected components

```
1 int Head[N], Next[M], To[M], Par[N], dfs_num[N], dfs_low[N], ne, n, m, u, v, root, rootChildren,
          dfs_timer, Stack[N], top, ID;
   bool Art[N];
   vector < vector <int> > BiCCs(N), BiCCIDs(N);
   void addEdge(int from, int to) {
     Next[++ne] = Head[from];
     Head[from] = ne;
     To[ne] = to;
9 }
10
11 void clear() {
                        0, sizeof(Head[0])
12
     memset (Head.
                                                 * (n + 2));
     memset (dfs_num,
13
                        0, sizeof(dfs_num[0]) * (n + 2));
     memset (Par,
                       -1, sizeof(Par[0])
                                                * (n + 2)):
14
                        0, sizeof(Art[0])
     memset (Art.
                                                 * (n + 2)):
1.5
     ne = dfs_timer = top = ID = 0;
16
     BiCCs = BiCCIDs = vector < vector <int> > (N);
17
18
19
20
    void Tarjan(int node)
21
22
      dfs_num[node] = dfs_low[node] = ++dfs_timer;
23
     Stack[top++] = node;
24
      for(int i = Head[node]; i; i = Next[i]) {
25
26
       if(dfs_num[To[i]] == 0) {
         if(node == root) ++rootChildren;
27
         Par[To[i]] = node;
28
29
         Tarian(To[i]);
30
31
          dfs_low[node] = Min(dfs_low[node], dfs_low[To[i]]);
         if (dfs_low[To[i]] >= dfs_num[node])
        Art[node] = true;
35
        for (int x = -1; x ^ To[i];)
36
37
38
            x = Stack[--top];
           BiCCIDs[x].emplace_back(ID);
39
40
           BiCCs[ID].emplace_back(x);
41
42
        BiCCIDs[node].emplace back(ID);
43
        BiCCs[ID].emplace_back(node);
44
45
46
        else if(To[i] != Par[node])
47
         dfs_low[node] = Min(dfs_low[node], dfs_num[To[i]]);
48
49
50
   int main()
52
53
     for(int i = 1; i <= n; ++i)
       if(dfs_num[i] == 0) {
54
55
         root = i;
56
          rootChildren = 0:
57
          Tarjan(i):
58
         Art[root] = (rootChildren > 1);
59
     for(int i = 1; i <= ID; ++i) {
       cout << "Component : " << i << " contains : ";</pre>
63
        for(int j = 0; j < (int)BiCCs[i].size(); ++j)</pre>
          cout << BiCCs[i][j] << " \n"[j == BiCCs[i].size() - 1];</pre>
65
66
```

# 2.8 Bipartite graph

```
1 bool checkBiPartite(int node, int par = 0) {
2    if(vis[node])
3    return color[par] != color[node];
4
5    color[node] = color[par] ^ 1;
6    vis[node] = true;
7    bool ok = true;
8    for(int i = Head[node]; i; i = Next[i])
9    if(To[i] != par)
10    ok &= checkBiPartite(To[i], node);
11    return ok;
12    }
13
```

```
14 int main() {
15    bool isBiPartite = true;
16    for(int i = 1; i <= n; ++i)
17     if(!vis[i])
18     isBiPartite &= checkBiPartite(i);
19    cout << (isBiPartite ? "YES" : "NO") << endl;
20 }</pre>
```

#### 2.9 Bellman ford

```
1 // Bellman Ford Algorithm : In programming contests, the slowness of Bellman Fords and its negative
cycle detection feature causes it to be used only to solve the SSSP problem on small graph
             which is not guaranteed to be free from negative weight cycle.
    bool hasNC() {
        for(int i = 1; i <= n; ++i)</pre>
          for(int j = Head[i]; j; j = Next[j])
  if(dis[i] < INF && dis[i] + Cost[j] < dis[To[j]])</pre>
        return true:
        return false:
10
11
12
    bool Bellman Ford(int src)
13
       memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
14
        memset (Par, -1, sizeof (Par[0]) * (n + 2));
15
        bool newRelaxation = true;
20
        for(int i = 2; i <= n && newRelaxation; ++i) {
21
          newRelaxation = false;
       for(int i = 1; i <= n; ++i)
  for(int j = Head[i]; j; j = Next[j])
if(dis[i] < INF && dis[i] + Cost[j] < dis[To[j]]) {</pre>
22
24
          dis[To[j]] = dis[i] + Cost[j];
Par[To[j]] = i;
25
26
          newRelaxation = true;
       return hasNC();
```

## 2.10 Connected components

```
1  void DFS(int node) {
2   visited[node] = true;
3   for(int e = Head[node]; e; e = Next[e])
4   if(!visited[To[e]])
5   DFS(To[e]);
6  }
7  
8  int main() {
9   for(int node = 1; node <= n; ++node)
10   if(!visited[node])
11   ++CCs, DFS(node);
12   cout << CCs << endl;
13 }</pre>
```

## 2.11 Cycle detection

```
1  void DFS(int node, int parent = -1)
2  {
3    if(hasCycle |= visited[node])
4     return;
5    visited[node] = true;
6
7    for(int i = Head[node]; i; i = Next[i])
8     if(To[i] != parent)
9     DFS(To[i], node);
10  }
11
12  int main() {
13    for(int i = 1; i <= n; ++i)
14    if(!visited[i])
15    DFS(i);</pre>
```

```
16   cout << (hasCycle ? "YES" : "NO") << endl;
17 }
```

## 2.12 Directed cyclic graph into acyclic

```
1 int HeadDAG[N], ToDAG[M], NextDAG[M], CostDAG[M], neDAG, Head[N], To[M], Next[M], Cost[M], dfs_num[N],
           dfs_low[N], out[N], Stack[N], compID[N], compSize[N], ne, n, m, u, v, w. dfs_timer, top, ID;
   void addEdge(int from, int to, int cost = 0) {
     Next[++ne] = Head[from];
Head[from] = ne;
     Cost[ne] = cost:
     To[ne] = to;
   void addEdgeDAG(int from, int to, int cost = 0) {
     NextDAG[++neDAG] = HeadDAG[from];
     HeadDAG[from] = neDAG;
     CostDAG[ne] = cost;
15
     ToDAG[neDAG] = to;
16
    ++out[from];
17 }
18
19 void clear() {
                       0, sizeof(Head[0])
                                             * (n + 2));
     memset (Head.
     memset(dfs_num, 0, sizeof(dfs_num[0]) * (n + 2));
21
      memset(compID, 0, sizeof(compID[0]) * (n + 2));
     memset(compSize, 0, sizeof(compSize[0]) * (n + 2));
      memset(HeadDAG, 0, sizeof(HeadDAG[0]) * (n + 2));
      memset (out,
                      0, sizeof(out[0])
26
     ne = dfs_timer = top = neDAG = ID = 0;
27
28
29
   void Tarjan (int node)
30
31
      dfs_num[node] = dfs_low[node] = ++dfs_timer;
32
      in_stack[Stack[top++] = node] = true;
33
34
      for(int i = Head[node]; i; i = Next[i]) {
        if(dfs_num[To[i]] == 0)
35
36
          Tarjan(To[i]);
37
38
        if (in_stack[To[i]])
39
          dfs_low[node] = Min(dfs_low[node], dfs_low[To[i]]);
40
41
42
     if(dfs_num[node] == dfs_low[node]) {
43
44
        for(int cur = -1; cur ^ node;) {
45
         in_stack[cur = Stack[--top]] = false;
46
          compID[cur] = ID;
47
          ++compSize[ID];
48
49
50 }
51
   void Tarjan() {
     for(int i = 1; i <= n; ++i)
       if(dfs_num[i] == 0)
55
          Tarjan(i);
56
57
58
   void DFS (int node)
59
60
      dfs_num[node] = 1;
61
      for(int i = Head[node]; i; i = Next[i]) {
     if(compID[node] != compID[To[i]])
addEdgeDAG(compID[node], compID[To[i]]);
62
63
         if(dfs_num[To[i]] == 0)
64
     DFS(To[i]);
65
66
67
   void construct_dag() {
     memset(dfs_num, 0, sizeof(dfs_num[0]) * (n + 2));
      for(int i = 1; i <= n; ++i)</pre>
73
       if(dfs_num[i] == 0)
74
         DFS(i);
75
```

### 2.13 Dijkstra (dense graph)

```
1 /** Dijkstra on dense graphs
       complexity : O(n^2 + m)
   int Head[N], Par[N], Next[M], To[M], Cost[M], ne, n, m, u, v, st, tr, tax;
   void Dijkstra(int src. int V)
      memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
10
     memset(Par, -1, sizeof(Par[0]) * (n + 2));
11
12
      vector <bool> mark(V + 1, false);
13
14
      dis[src] = 0;
      for(int i = 1; i <= V; ++i) {</pre>
        int node = 0;
        for (int j = 1; j <= V; ++j)
         if(!mark[j] && dis[j] < dis[node])
19
20
21
        if(dis[node] == INF) break;
        mark[node] = true;
for(int i = Head[node]; i; i = Next[i])
24
          if(dis[node] + Cost[i] < dis[To[i]])</pre>
25
26
        dis[To[i]] = dis[node] + Cost[i];
Par[To[i]] = node;
27
29
```

## 2.14 Dijkstra (grid)

```
const int dr[] = { 1, -1, 0, 0, 1, 1, -1, -1 };
const int dc[] = { 0, 0, 1, -1, 1, -1, 1, -1 };
const char dir[] = {'D', 'U', 'R', 'L'};
    const int N = 1e3 + 9, M = 2e5 + 9, oo = 0x3f3f3f3f3f;
    int grid[N][N], dis[N][N], n, m;
 9
    bool valid(int r, int c) {
10
      return r >= 1 && r <= n && c >= 1 && c <= m;
11
12
13
    void Dijkstra(int sr, int sc)
14
       memset(dis, 0x3f, sizeof (dis)); // memset(dis, 0x3f, n * m) we don't do that here
17
       priority_queue <tuple <int, int, int> > Q;
       dis[sr][sc] = grid[sr][sc];
19
       Q.push({-grid[sr][sc], sr, sc});
20
21
       int cost, r, c, nr, nc;
22
       while (Q.size())
23
24
           tie(cost, r, c) = 0.top(); 0.pop();
25
          if((-cost) > dis[r][c]) continue; // lazy deletion
26
27
          for(int i = 0; i < 4; ++i)
        nr = r + dr[i];
30
        nc = c + dc[i];
32
         if(!valid(nr, nc)) continue;
33
34
         if (dis[r][c] + grid[nr][nc] < dis[nr][nc])
35
36
             dis[nr][nc] = dis[r][c] + grid[nr][nc];
37
             Q.push({-dis[nr][nc], nr, nc});
38
39
40
```

# 2.15 Dijkstra (negative weighted graph)

```
1 int Head[N], Par[N], Next[M], To[M], Cost[M], ne, n, m, u, v, st, tr, tax;
```

```
2 ll dis[N];
   void Dijkstra(int src)
     memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
     memset(Par, -1, sizeof(Par[0]) * (n + 2));
     priority_queue <pair <11, int> > Q;
10
     Q.push({-dis[src], src});
11
12
13
     int node;
14
     11 cost:
1.5
     while(Q.size()) {
       tie(cost, node) = 0.top(); 0.pop();
16
       if((-cost) > dis[node]) continue;
19
       for(int i = Head[node]; i; i = Next[i])
20
         if(dis[node] + Cost[i] < dis[To[i]])</pre>
21
22
       dis[To[i]] = dis[node] + Cost[i];
23
       Q.push({-dis[To[i]], To[i]});
       Par[To[i]] = node;
25
26
27 1
```

## 2.16 Dijkstra (sparse graph)

```
/** Dijkstra on sparse graphs
      - complexity : O(n + m)logn -> O(nlogn + m)
        - Single Source Single Destination Shortest Path Problem
        - Positive Weight Edges only
        Subpaths of shortest paths from u to v are shortest paths!
 7 int Head[N], Par[N], Next[M], To[M], Cost[M], ne, n, m, u, v, st, tr, tax;
 8 11 dis[N];
10 void Dijkstra(int src, int trg)
11 {
     memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
13
     memset(Par, -1, sizeof(Par[0]) * (n + 2));
     priority_queue <pair <11, int> > Q;
     Q.push({-dis[src], src});
17
18
19
     int node;
20
21
      while(Q.size()) {
22
        tie(cost, node) = Q.top(); Q.pop();
23
        if((-cost) > dis[node]) continue; // lazy deletion
if(node == trq) return; // cheapest cost in case of positive weight edges
24
25
        for(int i = Head[node]; i; i = Next[i])
         if (dis[node] + Cost[i] < dis[To[i]])
29
30
        dis[To[i]] = dis[node] + Cost[i];
31
        Q.push({-dis[To[i]], To[i]});
32
        Par[To[i]] = node;
33
34
35
```

## 2.17 Edge classification

```
edgeClassification(To[i]);
16
17
18
          else if(dfs_num[To[i]] == VISITED)
         /** Cross Edges only occur in directed graph */
        if(in_time[To[i]] < in_time[node])</pre>
          cout << "Cross Edge : " << node << " -> " << To[i] << endl;</pre>
^{24}
          cout << "Forward Edge : " << node << " -> " << To[i] << endl;
25
26
          else if(dfs_num[To[i]] == EXPLORED)
27
28
        if(Par[node] == To[i])
          cout << "Bi-Directional Edge : " << node << " -> " << To[i] << endl;</pre>
29
30
          cout << "Backward Edge : " << node << " -> " << To[i] << " (Cycle)" << endl;
33
35
      dfs_num[node] = VISITED;
36
37
38
    int main() {
39
      for(int i = 1; i <= n; ++i) if(!dfs_num[i])</pre>
40
            edgeClassification(i);
41 }
```

#### 2.18 Eulerian tour tree

```
1 \quad \textbf{int} \; \texttt{Head[N], To[M], Next[M], Cost[M], ne, n, m, u, v, w, Last[N], First[N], euler\_tour[1 + N << 1];}
 2 ll Height[1 + N << 1];
 3 int euler_timer;
 5
    void _clear() {
      memset (Head.
                          0. sizeof (Head[0])
                                                     * (n + 2));
      memset (Last,
                                                    * (n + 2));
                          0, sizeof(Last[0])
      memset (First.
                          0, sizeof(First[0])
                                                    *(n + 2));
      ne = euler_timer = 0;
10
11
12
       euler\_tour[1 .. n * 2 - 1] = which records the sequence of visited nodes
13
14
       Height[1 ... n * 2 - 1] = which records the depth of each visited node
       First[1 .. n]
                         = records the index of the first occurrence of node i in euler_tour
                           = records the index of the last occurrence of node i in euler_tour
18
19
20
    void EulerianTour(int node, 11 depth = 0)
21
      euler_tour[++euler_timer] = node;
      Height[euler_timer] = depth;
24
      First[node] = euler timer;
      for(int i = Head[node]; i; i = Next[i])
        if(First[To[i]] == 0)
29
      EulerianTour(To[i], depth + Cost[i]);
30
31
       euler_tour[++euler_timer] = node;
       Height[euler_timer] = depth;
33
34
35
      Last[node] = euler_timer;
36
37
38
    void show() {
39
      for(int i = 1; i < (n << 1); ++i) cout << euler_tour[i] << " ";cout << endl;</pre>
      for(int i = 1; i < (n << 1); ++i) cout << Height[i] << " "; cout << endl; for(int i = 1; i <= n; ++i) cout << First[i] << " "; cout << endl;
      for(int i = 1; i <= n; ++i)
                                       cout << Last[i] << " ";
                                                                         cout << endl;
44
46
      EulerianTour(1);
48
      show();
49
```

## 2.19 Floodfill

```
1 /** check if there is a path from (0, 0) to (n - 1, m - 1) using '.' only **/
```

```
3 int dr[4] = \{1, -1, 0, 0\};
 4 int dc[4] = \{0, 0, 1, -1\};
 5 char grid[N][M];
 6 int n, m;
 8 bool valid(int r, int c) {
     return r >= 0 && r < n && c >= 0 && c < m && grid[r][c] == '.';
10 }
11
12 bool isDis(int r, int c) {
13
     return r == n - 1 && c == m - 1;
14
1.5
16 bool FloodFill(int r, int c) {
      if(!valid(r, c)) return false;
     if(isDis(r, c)) return true;
20
     grid[r][c] = '#';
21
     for(int i = 0; i < 4; ++i)
       if(FloodFill(r + dr[i], c + dc[i])) return true;
23
24
     return false:
25
26
27 int main() {
     cout << (FloodFill(0, 0) ? "YES" : "NO") << endl;</pre>
29
```

## 2.20 Floyd warshall (all-pairs shortest path)

```
/** -The graph has a 'negative cycle' if at the end of the algorithm,
    the distance from a vertex v to itself is negative.
      - before k-th phase the value of d[i][j] is equal to the length of
     the shortest path from vertex i to the vertex j,
     if this path is allowed to enter only the vertex with numbers smaller than k
     (the beginning and end of the path are not restricted by this property).
10 int Par[N][N], n, m, u, v, tax;
13 vector <int> restorePath(int st, int tr)
      vector <int> path;
     if(dis[st][tr] == INF) return path;
17
18
      for(int i = tr; st ^ i; i = Par[st][i])
19
       path.push_back(i);
20
21
     path.push back(st);
22
     reverse(path.begin(), path.end());
23
      return path;
   void Floyd_Warshall()
27
     for(int i = 1; i <= n; ++i)</pre>
29
        for(int j = 1; j <= n; ++j)</pre>
30
         Par[i][j] = i;
31
     for(int k = 1; k <= n; ++k)</pre>
33
       for(int i = 1; i <= n; ++i)
     for(int j = 1; j <= n; ++j)
if(dis[i][k] + dis[k][j] < dis[i][j])</pre>
34
35
36
         dis[i][j] = dis[i][k] + dis[k][j];
37
         Par[i][j] = Par[k][j];
38
39
40 }
```

# 2.21 Minimum spanning tree (Kruskal)

```
1 int n, m, u, v, w;
2 vector < tuple <int, int, int> > edges;
3 UnionFind uf;
4
5 pair < ll, vector < pair <int, int> > > Kruskal()
6 {
7 sort(edges.begin(), edges.end());
```

```
vector < pair <int, int> > mstEdges;
10
      int from, to, cost;
11
      ll minWieght = 0;
      for(tuple <int, int, int> edge : edges)
15
          tie(cost, from, to) = edge;
16
          if(uf.union_set(from, to))
17
18
        minWieght += cost;
19
        mstEdges.push_back(make_pair(from, to));
20
21
22
      if(mstEdges.size() == n - 1)
23
        return make_pair(minWieght, mstEdges);
26
      return make_pair(-1, vector < pair <int, int> > ());
27
```

# 2.22 Kth ancestor and lowest common ancestor (binary lifting)

```
1 int Head[N], To[M], Next[M], Par[N], up[N][LOG + 1], Log[N], Level[N], ne, n, u, v, q;
     memset(Head, 0, sizeof(Head[0]) * (n + 2));
memset(Par, 0, sizeof(Par[0]) * (n + 2));
     memset(Level, 0, sizeof(Level[0]) * (n + 2));
10 int lastBit(int a) {
11
     return (a & -a);
12
13
   void logCalc()
14
15
      Log[1] = 0;
16
      for(int i = 2; i < N; ++i)</pre>
       Log[i] = Log[i >> 1] + 1;
18
19
    void DFS(int node, int depth = 0)
      Level[node] = depth;
      up[node][0] = Par[node]; // Par[root] = root
25
26
      for(int i = 1; i <= LOG; ++i) {
       up[node][i] = up[up[node][i - 1]][i - 1];
27
28
29
30
      for(int i = Head[node]; i; i = Next[i]) if(To[i] != Par[node]) {
31
          Par[To[ill = node:
32
          DFS(To[i], depth + 1);
33
34
    int KthAncestor(int u, int k)
38
      if(k > Level[u]) return -1;
40
      for(int i = lastBit(k); k; k -= lastBit(k), i = lastBit(k))
41
       u = up[u][Log[i]];
42
43
      return u:
44
45
   int LCA(int u, int v)
46
47
      if(Level[u] < Level[v]) swap(u, v);</pre>
      int k = Level[u] - Level[v];
51
      if (u == v) return u:
53
      for(int i = LOG; i >= 0; --i)
55
       if(up[u][i] ^ up[v][i])
57
      v = up[v][i];
59
60
61
      return up[u][0];
62
```

```
64 int main()
65
      logCalc();
     for(int i = 1; i <= n; ++i) if(Par[i] == 0) {
          Par[i] = i;
69
70
71
     cin >> q;
73
     while (q--)
74
75
          cin >> u >> v;
76
          cout << LCA(u, v) << endl;</pre>
77
78
```

## 2.23 Lowest common ancestor (euler tour)

```
int Head[N], To[M], Next[M], Cost[M], ne, n, m, u, v, w, q;
   int Last[N], First[N], euler_tour[N << 1], Height[N << 1], euler_timer;</pre>
    void EulerianTour(int node, int depth = 0)
      euler_tour[++euler_timer] = node;
      Height[euler_timer] = depth;
     First[node] = euler timer;
      for(int i = Head[node]; i; i = Next[i])
11
      if(First[To[i]] == 0)
      EulerianTour(To[i], depth + Cost[i]);
15
      euler_tour[++euler_timer] = node;
16
      Height[euler_timer] = depth;
17
18
19
     Last[node] = euler timer;
20
21
22
   int main()
      SparseTable <int> st(Height + 1, Height + euler_timer + 1, [&] (int a, int b) { return a <= b; });
27
28
      while (q--)
29
30
          cin >> 1 >> r:
31
32
         int left = Last[1];
33
         int right = Last[r];
         if(left > right) swap(left, right);
34
35
36
         cout << euler_tour[ st.query(left, right) ] << endl;</pre>
37
38 }
```

#### 2.24 Minimum vertex cover

```
1 bool DFS(int node, int par = -1) {
2
3 bool black = false;
4 for(int e = Head[node]; e; e = Next[e])
5 if(To[e] != par)
6 black |= DFS(To[e], node);
7
8 MVC += black;
9 return !black;
10 }
```

## 2.25 Shortest path faster algorithml (SPFA)

```
1 /** Shortest Path Faster Algorithm:
2 - This algorithm runs in O(kE) where k is a number depending on the graph.
3 - The maximum k can be V (which is the same as the time complexity of Bellman Fords).
4 - However in practice SPFA (which uses a queue) is as fast as Dijkstras (which uses a priority queue).
5 - SPFA can deal with negative weight edge. If the graph has no negative cycle, SPFA runs well on
```

```
- If the graph has negative cycle(s), SPFA can also detect it as there must be some vertex (those
              on the negative cycle)
        that enters the queue for over V dash 1 times.
10 int Head[N], Par[N], Next[M], To[M], Cost[M], Cnt[N], ne, n, m, u, v, st, tax;
12 bool Inq[N];
13
14 void _set() {
15
     memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
      memset(Par, -1, sizeof(Par[0]) * (n + 2));
17
      memset (Cnt,
                   0, sizeof(Cnt[0]) * (n + 2));
18
      memset(Inq, 0, sizeof(Inq[0]) * (n + 2));
19
    bool SPFA(int src)
24
      deque <int> Q;
26
      Q.push_front (src);
27
28
      dis[src] = 0;
29
      Cnt[src] = 1;
      Inq[src] = 1;
30
31
32
33
      while(Q.size()) {
34
        node = Q.front(); Q.pop_front(); Inq[node] = 0;
        for(int i = Head[node]; i; i = Next[i])
  if(dis[node] + Cost[i] < dis[To[i]]) {</pre>
      dis[To[i]] = dis[node] + Cost[i];
39
      Par[To[i]] = node;
40
41
      if(!Inq[To[i]])
42
43
          if(++Cnt[To[i]] == n)
44
            return true; // graph has a negative weight cycle
45
46
          if(Q.size() && dis[To[i]] > dis[Q.front()])
47
            Q.push_back(To[i]);
48
            Q.push_front(To[i]);
49
50
51
          Inq[To[i]] = true;
53
54
55
      return false:
56
```

## 2.26 Single source shortest path

```
1 int Head[N], Par[N], Next[M], To[M], ne, n, m, u, v, st, tr;
2 11 dis[N];
    void BFS(int src)
      memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
      memset(Par, -1, sizeof(Par[0]) * (n + 2));
      queue <int> Q;
10
     Q.push(src);
11
     dis[src] = 0;
12
13
      while(Q.size()) {
14
       node = Q.front(); Q.pop();
15
       for(int i = Head[node]; i; i = Next[i])
         if (dis[To[i]] == INF) {
     dis[To[i]] = dis[node] + 1;
Par[To[i]] = node;
     Q.push(To[i]);
^{21}
22
```

# 2.27 Single source shortest path (grid)

```
1 const int dr [] = {-1, 0, 1, 0};
```

```
2 const int dc [] = {0, 1, 0, -1};
3 const char dir [] = {'U', 'R', 'D', 'L'};
 4 map <char, int> inv = { {'U', 0}, {'R', 1}, {'D', 2}, {'L', 3}};
 6 int dis[N][N], n, m;
9 bool valid(int r, int c) {
     return r >= 1 && r <= n && c >= 1 && c <= m && dis[r][c] == oo;
11
12
13 void BFS(int sr, int sc)
14
      memset(dis, 0x3f, sizeof(dis)):
1.5
16
     memset (Par, -1, sizeof (Par));
      queue < pair <int, int> > Q;
      dis[sr][sc] = 0;
20
     Q.push({sr, sc});
21
     int r, c, nr, nc;
23
      while(Q.size()) {
        tie(r, c) = Q.front(); Q.pop();
25
        for(int i = 0; i < 4; ++i) {</pre>
26
27
         nr = r + dr[i]:
28
         nc = c + dc[i];
29
30
          if(!valid(nr, nc)) continue;
31
          dis[nr][nc] = dis[r][c] + 1;
          Par[nr][nc] = dir[i ^ 2];
          Q.push({nr, nc});
35
36 }
```

## 2.28 Tarjan (strongly connected components)

```
1 int Head[N], To[M], Next[M], Cost[M];
 2 int dfs_num[N], dfs_low[N];
 3 int Stack[N], compID[N], compSize[N];
 4 int ne, n, m, u, v, w;
 5 int dfs_timer, top, ID;
 6 bool in_stack[N];
   void clear() {
     memset (Head,
                      0, sizeof(Head[0]) * (n + 2));
     memset(dfs_num, 0, sizeof(dfs_num[0]) * (n + 2));
memset(compID, 0, sizeof(compID[0]) * (n + 2));
1.1
      memset(compSize, 0, sizeof(compSize[0]) * (n + 2));
12
13
     ne = dfs_timer = top = ID = 0;
14
15
16
    void Tarjan(int node)
17
      dfs_num[node] = dfs_low[node] = ++dfs_timer;
19
      in_stack[Stack[top++] = node] = true;
20
21
      for(int i = Head[node]; i; i = Next[i]) {
22
       if(dfs_num[To[i]] == 0)
23
          Tarjan(To[i]);
24
25
        if (in stack[To[i]])
26
          dfs_low[node] = Min(dfs_low[node], dfs_low[To[i]]);
27
28
     if(dfs_num[node] == dfs_low[node]) {
29
30
31
        for(int cur = -1; cur ^ node;) {
32
         in_stack[cur = Stack[--top]] = false;
33
          compID[cur] = ID;
34
          ++compSize[ID];
35
36
37 }
38
39 void Tarjan() {
     for(int i = 1; i <= n; ++i) if(dfs_num[i] == 0)
40
41
            Tarian(i):
42
```

#### 2.29 Tree diameter

## 2.30 Tree diameter (weighted tree)

```
1 void DFS(int node, long long cost, int par = -1) {
2    if(cost > diameter) diameter = cost, at = node;
3    for (int e = Head[node]; e; e = Next[e])
4    if(To[e] != par)
5    DFS(To[e], cost + Cost[e], node);
6  }
7
8  int main() {
9    DFS(1, 011);
10    from = at, diameter = 0;
11    DFS(from, 011);
12    to = at;
13    cout << diameter << endl;
14 }</pre>
```

## 2.31 Topological sort (kahns algorithm)

```
vector <int> kahn(int n)
 2
      vector <int> ready, ret;
      for(int i = 1; i <= n; ++i)</pre>
       if(!in[i])
         ready.push_back(i);
10
      while (!ready.empty())
11
          node = ready.back(); ready.pop_back();
         ret.push_back(node);
14
          for(int i = Head[node]; i; i = Next[i])
16
      if(--in[To[i]] == 0)
       ready.push_back(To[i]);
18
19
      return ret:
20
21
22 int main() {
      vector <int> v = kahn(n);
      if((int)v.size() == n)
       for(int i : v)
         cout << i << ' ';
        cout << "not a DAG!" << endl;
29
```

#### 2.32 Problems

#### 2.32.1 Restoring the path

```
10
       - si start i
11
       - sj strat j
       - fi target :
       - fj target j
       - char dir and its map inv
14
16 **/
17
18 string restorePath(int si, int sj, int fi, int fj) {
19
     if(Par[ei][ej] == -1) return s;
20
21
22
     for(char i = Par[fi][fj]; (si ^ fi) || (sj ^ fj); i = Par[fi][fj]) {
    s += dir[inv[i] ^ 2];
23
24
        fi += dr[inv[i]];
       fj += dc[inv[i]];
27
28
29
     reverse(s.begin(), s.end());
30
31
32
33
   /** Explicit Graphs (BFS, Dijkstra or Bellman-Ford)
34
        - int Par[N] initialize with -1
35
36
        - 11 dis[N] initialize with 0x3f
37
       -11 INF = 0x3f3f3f3f3f3f3f3f3f
38
39
40 vector <int> restorePath(int dest) {
     vector <int> path;
     if (dis[dest] == INF) return path;
44
     for(int i = dest; ~i; i = Par[i])
45
       path.push_back(i);
46
47
     reverse(path.begin(), path.end());
48
     return path;
49
50
51
   /** in case of Floyd-Warshall:
52
53
        - 11 dis[N][N] initialize with 0x3f
54
       -11 INF = 0x3f3f3f3f3f3f3f3f3f
55
        - int Par[N][N] initialize with
                                             Par[i][j] = i;
56
        - in Floyd-Warshall function write -> Par[i][j] = Par[k][j];
57
58
59 vector <int> restorePath(int st, int tr) {
60
     vector <int> path;
     if(dis[st][tr] == INF) return path;
61
     for(int i = tr; st ^ i; i = Par[st][i])
63
64
       path.push_back(i);
65
66
     path.push back(st);
67
     reverse(path.begin(), path.end());
68
     return path:
```

#### 2.32.2 Tree distances

```
1 int Head[N], Next[M], To[M], Par[N], ne, n, m, u, v, diameter, At, From;
   int E[N << 1], H[N << 1], F[N], L[N], timer, SP[N << 1][LOG + 1], Log[N << 1];</pre>
     memset(Head, 0, sizeof(Head[0]) * (n + 2));
     memset(Par, -1, sizeof(Par[0]) * (n + 2));
     ne = 0;
     timer = 0;
   void EulerTour(int node, int depth = 0, int par = -1) {
     E[++timer] = node;
     H[timer] = depth;
     F[node] = timer;
16
17
      for(int i = Head[node]; i; i = Next[i])
18
       if(To[i] != par) {
19
          EulerTour(To[i], depth + 1, node);
20
         E[++timer] = node;
         H[timer] = depth;
21
22
     L[node] = timer;
```

```
25 }
26
    void dfs(int node, int depth = 0, int par = -1) {
      if (depth > diameter) diameter = depth, At = node;
      for(int i = Head[node]; i; i = Next[i])
       if(To[i] != par)
31
          dfs(To[i], depth + 1, node);
32
33
34
   void bulid()
35
36
      EulerTour(1);
      dfs(1); From = At; diameter = 0; dfs(From);
37
38
39
      Log[1] = 0;
for(int i = 2; i <= (n << 1); ++i)
41
        Log[i] = Log[i >> 1] + 1;
42
43
      for(int i = 1; i < (n << 1); ++i)</pre>
44
        SP[i][0] = i;
45
46
      int MaxLog = Log[(n << 1)];</pre>
47
      for(int j = 1, k, h; j <= MaxLog; ++j) {</pre>
48
        k = (1 << j);
49
        h = (k >> 1);
        for (int i = 1; i + k - 1 < (n << 1); ++i)
50
51
      const int & x = SP[i][j - 1];
52
53
      const int & y = SP[i + h][j - 1];
54
      SP[i][j] = H[x] \le H[y] ? x : y;
57
58
59
60
   int query(int 1, int r)
61
      int d = r - 1 + 1;
63
     int lg = Log[d];
int k = (1 << lg);</pre>
64
65
      const int & x = SP[1][1g];
const int & y = SP[1 + d - k][1g];
66
67
68
69
      return (H[x] <= H[v] ? x : v);
70
   int LCA(int u, int v) {
     return query (u, v);
74
76 int distance(int u, int v) {
77
      int 1 = F[u];
      int r = F[v]:
78
79
      if(1 > r) swap(1, r);
80
81
      int ix = LCA(1, r);
      return (H[1] + H[r] - H[ix] - H[ix]);
83
85
    int main()
87
      for(int i = 1; i <= n; ++i)</pre>
89
        cout << max(distance(i, At), distance(i, From)) << " \n"[i == n];</pre>
90
```

#### 2.32.3 Subtree sizes

```
1 int Head[N], Next[M], To[M], Par[N], sbtree_size[N], ne, n, m, u, v, w;
   void dfs(int node, int par = -1) {
      sbtree_size[node] = 1;
      for(int i = Head[node]; i; i = Next[i])
       if(To[i] != par) {
         dfs(To[i], node);
          sbtree_size[node] += sbtree_size[To[i]];
 9
10 }
1.1
12 int main()
13 {
      dfs(1);
14
      for(int i = 1; i <= n; ++i)
15
        cout << sbtree_size[i] - 1 << " \n"[i == n];</pre>
16
17 }
```

#### 2.32.4 Shortest cycle

```
1 \ /** for each node run BFS and minmize the cycle length
 2
    int BFS (int src)
 5
      memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
 6
      memset(Par, -1, sizeof(Par[0]) * (n + 2));
      queue <int> 0:
10
      Q.push(src);
11
      dis[src] = 0;
15
16
           node = Q.front(); Q.pop();
17
           for(int i = Head[node]; i; i = Next[i])
18
        if(dis[To[i]] != oo) {
   if(Par[node] != To[i]) {
     if(dis[node] + 1 + dis[To[i]] < ret)
ret = dis[node] + 1 + dis[To[i]];</pre>
19
20
21
22
23
24
           continue;
25
26
27
         dis[To[i]] = dis[node] + 1;
         Par[To[i]] = node;
29
         Q.push(To[i]);
30
31
32
      return ret;
33
```

#### 2.32.5 Path sum for each node

```
1 int Head[N], Next[M], To[M], ne, u, v, n, m, subtree_size[N], level[N];
 2 11 dis[N];
   void dfs(int node, int par = -1) {
     subtree_size[node] = 1;
      for(int i = Head[node]; i; i = Next[i]) if(To[i] != par) {
          level[To[i]] = level[node] + 1;
          dfs(To[i], node);
          subtree_size[node] += subtree_size[To[i]];
10
11 }
12
13 void reRoot(int node, ll pd, int par = -1) {
     dis[node] = pd;
for(int i = Head[node]; i; i = Next[i]) if(To[i] != par) {
14
1.5
16
         reRoot(To[i], pd - subtree_size[To[i]] + (n - subtree_size[To[i]]), node);
17
18 }
19
20
   void get_dis()
^{21}
^{24}
      for(int i = 1; i <= n; ++i)</pre>
25
       pd += level[i];
26
27
      reRoot(1, pd);
28
      for(int i = 1; i <= n; ++i)</pre>
        cout << dis[i] << " \n"[i == n];</pre>
29
30
```

## 3 Data structures

## 3.1 Segment tree

```
1 template <class T, class F = function <T(const T &, const T &)>> 2 class SegmentTree 3 { vector <T> _A;
```

```
vector <T> ST;
      vector <T> LT;
      F func;
      int _N;
10 public :
      template <class iter>
      SegmentTree(iter _begin, iter _end, const F _func = [](T a, T b) {return a <= b ? a : b;}) : func(
13
14
        _N = distance(_begin, _end);
15
        N = (1 << (int)ceil(log2(N)));
16
17
        _A.assign(_N + 1, 0);
       18
21
          _typeof(_begin) i = _begin;
22
        for(int j = 1; i != _end; ++i, ++j)
23
24
25
       build(1, 1, _N);
26
27
28
      void build(int p, int l, int r)
29
30
31
         ST[p] = A[1];
32
         return:
33
34
       int mid = (1 + r) >> 1;
36
37
       build(p + p, l, mid);
38
       build(p + p + 1, mid + 1, r);
39
40
        const T & x = ST[p + p];
41
       const T & y = ST[p + p + 1];
42
43
       ST[p] = func(x, y);
44
45
46
      void update_range(int ul, int ur, int delta) {
47
        update_range(ul, ur, delta, 1, 1, _N);
48
49
50
      T query(int ql, int qr) {
51
        return query(ql, qr, 1, 1, _N);
52
53
54
      void update_point(int inx, int delta)
55
56
        inx += N - 1:
       ST[inx] = delta:
57
58
59
        while(inx > 1) {
60
         inx >>= 1:
61
         const T & x = ST[inx + inx];
         const T & y = ST[inx + inx + 1];
64
         ST[inx] = func(x, y);
66
67
68
69 private :
70
      void update_range(int ul, int ur, int delta, int p, int l, int r)
71
72
73
74
       if(r < ul || ur < 1)
75
       if(ul <= 1 && r <= ur) {
76
         ST[p] += delta;
77
          LT[p] += delta;
78
         return;
81
       propagate(p);
83
        int mid = (1 + r) >> 1;
84
85
86
        update_range(ul, ur, delta, p + p, l, mid);
        update_range(ul, ur, delta, p + p + 1, mid + 1, r);
87
88
       const T & x = ST[p + p];
89
        const T & y = ST[p + p + 1];
        ST[p] = func(x, y);
94
      T query(int ql, int qr, int p, int l, int r)
```

```
96
        if(r < ql || qr < 1)
97
          return INT_MAX;
98
99
        if(q1 <= 1 && r <= qr)
100
          return ST[p];
102
103
104
        int mid = (1 + r) >> 1;
106
        const T & x = query(q1, qr, p + p, 1, mid);
107
        const T & y = query(q1, qr, p + p + 1, mid + 1, r);
108
109
        return func(x, y);
110
       void propagate(int p) {
113
        if(LT[p]) {
          ST[p + p]
115
           ST[p + p + 1] += LT[p];
          LT[p + p] += LT[p];
         LT[p + p + 1] += LT[p];
LT[p] = 0;
117
118
119
120
121 };
```

#### 3.2 Merge sort tree

```
class SegmentTree
     vector <vector <int> > sTree;
     vector <int> localArr;
     int NP2, oo = 0x3f3f3f3f;
     template <class T>
     SegmentTree (T _begin, T _end)
10
11
       int n = _end - _begin;
12
       while (NP2 < n) NP2 <<= 1;
13
14
15
       sTree.assign(NP2 << 1, vector <int> ());
       localArr.assign(NP2 + 1, 0);
         _typeof(_begin) i = _begin;
19
       for(int j = 1; i != _end; i++, ++j)
20
        localArr[j] = *i;
21
22
       build(1, 1, NP2);
23
24
25
     void build(int p, int l, int r)
26
27
       if(1 == r) {
        sTree[p].push_back(localArr[1]);
28
29
         return;
30
       build(left(p), l, mid(l, r));
33
       build(right(p), mid(l, r) + 1, r);
34
35
36
37
38
     int query(int ql, int qr, int k) {
39
       return query(ql, qr, k, 1, 1, NP2);
40
41
    int query(int ql, int qr, int k, int p, int l, int r)
45
       if(isOutside(ql, qr, l, r))
46
47
48
       if(isInside(ql, qr, l, r)) {
49
        return sTree[p].end() - upper_bound(sTree[p].begin(), sTree[p].end(), k);
50
51
52
       53
54
55
56
     void merge(int p)
57
       vector <int> & L = sTree[left(p)];
```

```
59
        vector <int> & R = sTree[right(p)];
61
        int l_size = L.size();
        int r_size = R.size();
63
        int p_size = l_size + r_size;
65
        L.push_back(INT_MAX);
        R.push_back(INT_MAX);
67
68
       sTree[p].resize(p_size);
69
70
       for(int k = 0, i = 0, j = 0; k < p_size; ++k)</pre>
         if(L[i] <= R[j])
71
      sTree[p][k] = L[i], i += (L[i] != INT_MAX);
72
73
      sTree[p][k] = R[j], j += (R[j] != INT_MAX);
        R.pop_back();
78
79
80
      inline bool isInside(int ql, int qr, int sl, int sr) {
81
       return (q1 <= s1 && sr <= qr);
82
83
      inline bool isOutside(int ql, int qr, int sl, int sr) {
84
85
       return (sr < ql || qr < sl);</pre>
86
87
      inline int mid (int 1, int r) {
       return ((1 + r) >> 1);
91
      inline int left(int p) {
93
       return (p << 1);
94
95
     inline int right (int p) {
97
       return ((p << 1) | 1);
98
99 };
```

# 3.3 Merge sort

```
1 ll inversions;
    void merge(T localArr [], int l, int mid, int r)
      int l_size = mid - 1 + 1;
     int r size = r - mid;
      T L[1 size + 1];
10
     T R[r_size + 1];
11
12
      for(int i = 0; i < l_size; ++i) L[i] = localArr[i + 1];
for(int i = 0; i < r_size; ++i) R[i] = localArr[i + mid + 1];</pre>
13
14
15
      if(sizeof(T) == 4) Mx = INT_MAX;
      else Mx = LONG_MAX;
      L[l_size] = R[r_size] = Mx;
20
^{21}
      for (int k = 1, i = 0, j = 0; k \le r; ++k)
        if(L[i] <= R[j])
22
23
          localArr[k] = L[i], i += (L[i] != Mx);
24
25
          localArr[k] = R[j], j += (R[j] != Mx), inversions += l_size - i;
26
   template <class T>
    void merge_sort(T localArr [], int l, int r)
31
32
33
          int mid = (1 + r) >> 1;
          merge_sort(localArr, 1,
34
35
          merge_sort(localArr, mid + 1, r);
          merge(localArr,
                              l, mid, r);
37
38 }
40 template <class T>
41 void merge_sort(T _begin, T _end)
      const int sz = _end - _begin;
```

```
44 __typeof(*_begin) localArray[sz];
45 __typeof(_begin) k = _begin;
47    for(int i = 0; k != _end; ++i, ++k)
48    localArray[i] = *k;
49    merge_sort(localArray, 0, sz - 1);
51    for(int i = 0; k != _end; ++i, ++k)
53    for(int i = 0; k != _end; ++i, ++k)
54    *k = localArray[i];
```

## 3.4 Sparse table

```
template <class T, class F = function <T(const T&, const T&)> >
 2 class SparseTable
3
     int N;
     int _LOG;
     vector <T> A;
     vector < vector <T> > ST;
     vector <int> Log;
11 public :
12
     SparseTable() = default:
13
14
     template (class iter)
     SparseTable(iter _begin, iter _end, const F _func = less <T> ()) : func(_func)
15
16
17
       _N = distance(_begin, _end);
18
19
       Log.assign(_N + 1, 0);
for(int i = 2; i <= _N; ++i)
20
        Log[i] = Log[i >> \overline{1}] + 1;
24
25
        _A.assign(_N + 1, 0);
26
       ST.assign(N + 1, vector < T > (LOG + 1, 0));
27
28
         _typeof(_begin) i = _begin;
       for (int j = 1; i != _end; ++i, ++j)
_A[j] = *i;
29
30
31
32
       build();
33
34
35
     void build()
36
37
       for (int i = 1; i <= _N; ++i)</pre>
38
        ST[i][0] = i;
39
40
       41
42
     k = (1 << j);
     d = (k >> 1);
43
44
     for(int i = 1; i + k - 1 <= _N; ++i)</pre>
45
46
47
         T const & x = ST[i][j - 1];
                                     // starting subarray at index = i with length = 2^{j-1}
         T const & y = ST[i + d][j - 1]; // starting subarray at index = i + d with length = 2^{i}[j - 1]
49
50
         ST[i][j] = func(A[x], A[y]) ? x : y;
51
52
53
54
55
     T query(int 1, int r) // this query is O(1)
56
57
       T const & x = ST[1][Log[d]];
58
       T const & y = ST[1 + d - (1 << Log[d])][Log[d]];
59
60
61
       return func(_A[x], _A[y]) ? x : y;
63 };
```

## 3.5 Union find disjoint sets

```
Maintain a set of elements partitioned into non-overlapping subsets. Each
       partition is assigned a unique representative known as the parent, or root. The
       following implements two well-known optimizations known as union-by-size and
       path compression. This version is simplified to only work on integer elements.
        find_set(u) returns the unique representative of the partition containing u.
        same\_set(u, v) returns whether elements u and v belong to the same partition.
        union_set(u, v) replaces the partitions containing u and v with a single new
10
       partition consisting of the union of elements in the original partitions.
11
12
       Time Complexity:
        O(a(n)) per call to find_set(), same_set(), and union_set(), where n is the
13
14
       number of elements, and a(n) is the extremely slow growing inverse of the Ackermann function
       (effectively a very small constant for all practical values of n).
15
       O(n) for storage of the disjoint set forest elements.
        O(1) auxiliary for all operations.
20
21
22 class UnionFind
23
24
      vector <int> par:
25
      vector <int> siz:
26
      int num sets:
27
      size t sz:
29
   public:
      UnionFind(): par(1, -1), siz(1, 1), num_sets(0), sz(0) {}
      UnionFind(int n): par(n + 1, -1), siz(n + 1, 1), num_sets(n), sz(n) \{\}
33
      int find_set (int u)
34
35
        assert(u <= sz);
36
37
        int leader;
38
        for(leader = u; ~par[leader]; leader = par[leader]);
39
40
        for(int next = par[u]; u != leader; next = par[next]) {
         par[u] = leader;
41
42
          11 = next:
43
44
        return leader;
45
46
47
      bool same_set(int u, int v) {
        return find_set(u) == find_set(v);
49
50
51
      bool union_set(int u, int v) {
52
        if(same_set(u, v)) return false;
53
54
        int x = find set(n):
55
        int y = find_set(v);
56
57
        if(siz[x] < siz[y]) swap(x, y);</pre>
59
        par[v] = x:
        siz[x] += siz[y];
63
64
65
66
      int number_of_sets() {
67
        return num_sets;
68
69
      int size of set (int u) {
70
71
        return siz[find_set(u)];
72
73
74
      size_t size() {
75
       return sz;
76
79
       par.clear();
80
        siz.clear();
81
        sz = num\_sets = 0;
82
83
84
      void assign(size_t n) {
        par.assign(n + 1, -1);
        siz.assign(n + 1, 1);
        sz = num_sets = n;
89
      map < int, vector <int> > groups(int st) {
91
        map < int, vector <int> > ret;
```

#### 3.6 Bubble sort

```
Bubble sort consists of n rounds. On each round, the algorithm iterates
      through the elements of the array. Whenever two consecutive elements are found
      that are not in correct order, the algorithm swaps them. The algorithm can be
      implemented as follows:
   void bubble_sort(T _begin, T _end, int round) {
10
     const int sz = _end - _begin;
1.1
     int localArray[sz];
12
     __typeof(_begin) k = _begin;
for(int j = 0; k != _end; ++k, ++j)
13
14
15
       localArray[j] = *k;
16
      round = min(round, sz);
17
      for(int i = 0; i < round; ++i) /* n rounds -> n_th element **/
18
       for(int j = 0; j < sz - 1; ++j) if(localArray[j] > localArray[j + 1])
20
                 swap(localArray[j], localArray[j + 1]);
21
      for(int j = 0; k != _end; ++k, ++j)
23
24
       *k = localArray[j];
25
26
27
      After the first round of the algorithm, the largest element will be in the correct
28
      position, and in general, after k rounds, the k largest elements will be in the
29
30
      correct positions.
```

## 4 Mathematics

# 4.1 Pisano periodic sequence

```
template <class T>
    using matrix = vector < vector <T> >;
    template <class T> string to_string(T x) {
       int sn = 1;
      if (x < 0) sn = -1, x *= sn;
       string s = "";
          s = "0123456789"[x % 10] + s, x /= 10;
10
11
       return (sn == -1 ? "-" : "") + s;
12 }
13
14 auto str_to_int(string x) {
15    uil28 ret = (x[0] == '-' ? 0 : x[0] = '0');
16    for(int i = 1; i < (int)x.size(); ++i) ret = ret * 10 + (x[i] = '0');
17    return (x[0] == '-' ? -1 * (il28)ret : ret);</pre>
18
20 istream & operator >> (istream & in, i128 & i) noexcept {
     string s;
23
24
25
27 ostream & operator << (ostream & os, const i128 i) noexcept {
28
      os << to string(i);
29
      return os;
30 }
31
32 void Fast() {
      cin.sync_with_stdio(0);
      cin.tie(0);
```

```
35
       cout.tie(0);
 36
 37
    11 n;
 39
    vector <int> primes;
    matrix <11> fibMatrix = {{1, 1},
41
42 };
 43
 44
    i128 gcd(i128 a, i128 b) {
 45
       while (a && b)
 46
        a > b ? a %= b : b %= a;
 47
       return a + b;
48
 49
    i128 lcm(i128 a, i128 b) {
      return a / gcd(a, b) * b;
 52
 53
 54
    vector < array <11, 2> > factorize(11 x) {
 55
       vector < array <11, 2> > ret;
 56
       for(int i = 0; 111 * primes[i] * primes[i] <= x; ++i) {
 57
        if(x % primes[i]) continue;
 58
 59
         int cnt = 0:
        while (x % primes[i] == 0) {
 60
 61
          cnt++;
62
          x /= primes[i];
 63
 64
        ret.push_back({primes[i], cnt});
 65
 67
       if(x > 1) ret.push_back({x, 1});
69
 70
71
    matrix <11> MatMul (matrix <11> A, matrix <11> B, 11 mod) {
 72
       int ra = A.size(), cb = B[0].size(), ca = A[0].size();
73
       matrix < i128 > C(ra, vector < i128 > (cb));
74
 75
       for(int i = 0; i < ra; ++i)</pre>
        for (int j = 0; j < cb; ++j) {
  C[i][j] = 0;</pre>
 76
 77
 78
           for (int k = 0; k < ca; ++k)
 79
       C[i][j] = (C[i][j] + (i128)A[i][k] * B[k][j]);
 80
       matrix <11> ret(ra, vector <11> (cb));
       for(int i = 0; i < ra; ++i)</pre>
        for (int j = 0; j < cb; ++j)
 84
 85
         ret[i][j] = C[i][j] % mod;
 86
 87
       return ret:
 88
 89
90 matrix <11> MatPow(matrix <11> A, 11 p, 11 mod) {
91
      int r = A.size(), c = A[0].size();
       assert (r == c && p);
      matrix <11> result = A;
97
        if(p & 111) result = MatMul(result, A, mod);
        A = MatMul(A, A, mod);
98
        p >>= 111;
99
100
101
       return result;
102
103
104 i128 ModExp(i128 a, ll p) {
105
       i128 \text{ result} = 1;
106
       while (p)
        if (p & 111) result = result * a;
107
108
        a *= a;
109
        p >>= 111;
110
113
114 ll nthFib(ll n, ll mod) {
115
       return MatPow(fibMatrix, n, mod)[0][1];
116
117
118 bool is_period(ll n, ll mod) {
      return nthFib(n, mod) == 0 && nthFib(n + 1, mod) == 1;
119
120
122 ll solver(ll x, ll mod) {
       vector < array <11, 2> > factors = factorize(x);
       for(int i = 0; i < (int) factors.size(); ++i) {</pre>
125
         while(x % factors[i][0] == 0 && is_period(x / factors[i][0], mod))
          x /= factors[i][0];
```

```
127
128
      return x;
129
130
    ll pisano_prime(ll val) {
131
      if(val == 2) return 3;
133
      if (val == 5) return 20;
134
      if(val % 10 == 1 || val % 10 == 9)
135
        return solver (val - 1, val);
136
137
      return solver(2 * (val + 1), val);
138
139
140 const int N = 1e7 + 9:
141 bitset <N> isPrime:
143 void Precomputation_Sieve() {
      isPrime.set();
      int _sqrt = sqrtl(N);
146
      for(int i = 5; i <= _sqrt; i += 6) {</pre>
147
148
       if(isPrime[i]) for (int j = i * i; j < N; j += i + i) isPrime.reset(j);
149
        i += 2;
150
        if(isPrime[i]) for (int j = i * i; j < N; j += i + i) isPrime.reset(j);
151
        i -= 2;
152
153 }
154
155 vector <int> Primes(int n) {
      vector <int> Primes:
156
157
      if(n >= 2) _Primes.push_back(2);
159
      if(n >= 3) _Primes.push_back(3);
160
161
      for (int i = 5; i <= n; i += 6) {
162
       if(isPrime[i]) _Primes.push_back(i);
163
        if(isPrime[i]) _Primes.push_back(i);
164
165
        i -= 2;
166
167
      return Primes;
168
169
170
    void initialize()
171
      Precomputation_Sieve();
172
173
      primes = Primes(N);
174
175
176 void Solve() {
177
      cin >> n;
178
      vector < array <11, 2> > factors = factorize(n);
179
180
181
      for (int i = 0; i < (int) factors.size(); ++i) {</pre>
182
       ans = lcm(ans, (i128)pisano_prime(factors[i][0]) * ModExp(factors[i][0], factors[i][1] - 1));
183
184
      cout << ans << endl:
185
186
    void MultiTest(bool Tests)
189
      int tc = 1; (Tests) && (cin >> tc);
      for(int i = 1; i <= tc; ++i)</pre>
191
        Solve();
192
193
194 int main()
195 {
196
      Fast(); initialize(); MultiTest(1);
197
```

#### 4.2 Euler totient function

#### 4.3 Extended wheel factorization

```
1 /*
     2 <= a <= 1e{14}
      linear_sieve takes O(n)
     Factorization takes O(n / (ln(n) - 1.08))
10
     Space Complexity:
     O(MaxN + n / (ln(n) - 1.08)
11
12
13
14 int lp[N];
15 int Primes[664580], pnx; /** size of Primes = n / (ln(n) - 1.08) */
   void linear_sieve(int n) {
18
      for (int i = 2; i <= n; ++i) {
        if (lp[i] == 0) {
20
         lp[i] = Primes[pnx++] = i;
21
22
        for (int j = 0, comp; j < pnx && Primes[j] <= lp[i] && (comp = i * Primes[j]) <= n; ++j) {
         lp[comp] = Primes[j];
24
25
26
    vector<pair<11, int>> Factorization(11 a) {
      vector<pair<11, int> > ret;
31
      for (int i = 0, cnt; i < pnx && (p = Primes[i], true) && p * p <= a; ++i) {</pre>
      if (a % p) continue;
33
34
        while (a % p == 0) a /= p, ++cnt;
35
        ret.emplace_back(p, cnt);
37
      if (a > 1) ret.emplace_back(a, 1);
38
     return ret;
39
```

# 4.4 Least prime factorization

```
1 <= n <= 1e7
     linear_sieve takes O(n)
     Factorization takes O(log(n))
     Space Complexity:
10
    O(MaxN + n / (ln(n) - 1.08)
11
14
   int Primes[664580], pnx; /** size of Primes = n / (ln(n) - 1.08) */
16
   void linear_sieve(int n) {
     for (int i = 2; i <= n; ++i) {
18
       if (lp[i] == 0) {
19
         lp[i] = Primes[pnx++] = i;
20
       for (int j = 0, comp; j < pnx && Primes[j] <= lp[i] && (comp = i * Primes[j]) <= n; ++j) {
         lp[comp] = Primes[j];
```

```
27  vector<pair<int, int>> Factorization(int n) {
28    vector<pair<int, int>> ret;
29    while (n > 1) {
30        int p = leastPrime[n], cnt = 0;
31        while (n % p == 0) n /= p, ++cnt;
32        ret.emplace_back(p, cnt);
33    }
34    return ret;
35 }
```

#### 4.5 Mobius function

```
1
       Constraints:
        1 <= x <= 1e7
       2 <= n <= 10^{14}
       Time Complexity:
       linear_sieve takes O(x)
       mobius takes O(n / (ln(n) - 1.08))
10
       Space Complexity:
11
       O(MaxN + n / (ln(n) - 1.08))
12 */
13
14 int lp[N], Primes[664580], pnx; /** size of Primes = n / (ln(n) - 1.08) */
   void linear_sieve(int x) {
     for (int i = 2; i <= x; ++i) {
       if (lp[i] == 0) {
19
         lp[i] = Primes[pnx++] = i;
20
21
       for (int j = 0, comp; j < pnx && Primes[j] <= lp[i] && (comp = i * Primes[j]) <= x; ++j) {
         lp[comp] = Primes[j];
24
25
26
27
28 int mobius(ll n) {
29
30
     for (int i = 0; i < pnx && (p = Primes[i], pp = p * p, true); ++i) {</pre>
      if (pp > n) break;
       if (n % p) continue;
       if (n % pp == 0) return 0;
35
       n /= p;
36
       mob = -mob;
37
38
     if (n > 1) mob = -mob;
39
     return mob;
40 }
```

## 4.6 Phi factorial

```
1 <= x <= 1e7
        2 <= n <= 1e7
        Time Complexity:
        linear_sieve takes O(x)
        phi factorial takes O(n)
10
        Space Complexity:
11
        O(MaxN + n / (ln(n) - 1.08))
14 int lp[N], Primes[664580], pnx; /** number of primes = n / (ln(n) - 1.08) **/
16 void linear_sieve(int x) {
     for (int i = 2; i <= x; ++i) {
  if (lp[i] == 0) {</pre>
17
19
          lp[i] = Primes[pnx++] = i;
20
21
        for (int j = 0, comp; j < pnx && Primes[j] <= lp[i] && (comp = i * Primes[j]) <= x; ++j) {</pre>
         lp[comp] = Primes[j];
24
25
26 }
```

```
28  ll phi_factorial(int n) {
29     ll ret = 1;
30     for (int i = 2; i <= n; ++i) {
31         ret = ret * (lp[i] == i ? i - 1 : i);
32     }
33     return ret;
34 }</pre>
```

#### 4.7 Linear sieve

```
1 /**
 3
        1 <= n <= 1e7
        Time Complexity:
        linear_sieve takes O(n)
        Space Complexity:
       O(MaxN + n / (ln(n) - 1.08))
10 ++/
11
12 int lp[N]:
   int Primes[664580], pnx; /** size of Primes = n / (ln(n) - 1.08) */
13
14
15
   void linear sieve(int n) {
     for (int i = 2; i <= n; ++i) {
16
       if (lp[i] == 0) {
         lp[i] = Primes[pnx++] = i;
20
        for (int j = 0, comp; j < pnx && Primes[j] <= lp[i] && (comp = i * Primes[j]) <= n; ++j) {
21
22
23
24
```

## 4.8 Segmented sieve

```
1 int lp[N];
 2 int Primes[664580], pnx; /** size of Primes = n / (ln(n) - 1.08) */
    bool isPrime[N];
 5
    void linear_sieve(int n) {
 6
      for (int i = 2; i <= n; ++i) {</pre>
        if (lp[i] == 0) {
          lp[i] = Primes[pnx++] = i;
10
        for (int j = 0, comp; j < pnx && Primes[j] <= lp[i] && (comp = i * Primes[j]) <= n; ++j) {
11
          lp[comp] = Primes[j];
12
13
14
15
    vector<11> segmented_sieve(11 1, 11 r) {
16
18
      int limit = r - 1 + 1;
      vector<ll> ret;
20
      memset(isPrime, true, sizeof(isPrime));
21
22
      for (int i = 0; i < pnx && (p = Primes[i], true); ++i) {</pre>
23
24
        for (11 j = max(p * p, (1 + p - 1) / p * p); j <= r; j += p)
    isPrime[j - 1] = false;</pre>
25
26
27
      for (int i = 0; i < limit; ++i)</pre>
        if (isPrime[i])
          ret.emplace_back(i + 1);
31
```

#### 4.9 Miller Rabin test

```
using namespace std;
   typedef long long 11;
10 typedef __int128 i128;
12
13
14 void Fast() {
15
     cin.sync_with_stdio(0);
16
     cin.tie(0);cout.tie(0);
17
18
   11 ModExp(11 base, 11 e, 11 mod)
19
20
     base %= mod;
23
24
     for(result = 1; e; e >>= 111)
25
26
         if(e & 111)
27
     result = ((i128) result * base) % mod;
28
         base = ((i128)base * base) % mod;
29
30
     return result:
31 3
32
33
   bool CheckComposite(ll n, ll p, ll d, int r)
34
35
     11 a = ModExp(p, d, n);
36
     if (a == 1 || a == n - 1)
37
       return false;
     for(int i = 1; i < r; ++i)</pre>
40
41
          a = ((i128)a * a) % n;
42
         if(a == n - 1)
43
     return false;
44
45
     return true;
46
47
   bool Miller(ll n)
48
49
     if(n < 2) return false;</pre>
50
51
52
     for(r = 0, d = n - 1; (d & 111) == 0; d >>= 111, ++r);
54
55
     for(int p: {2, 3, 7, 11, 13, 17, 19, 23, 29, 31, 37})
56
57
58
     return true:
        if(CheckComposite(n, p, d, r))
59
60
     return false:
61
62
     return true:
63
64
   int main()
68
69
     11 n;
70
71
     cout << (Miller(n) ? "Yes, it is Prime" : "No, it is not a prime") << endl;</pre>
```

# 4.10 Stable marriage problem

```
18 }
19
    const int N = 1e3+ 9, M = 1e3 + 9, oo = 0x3f3f3f3f3f;
23
    int husband[N], wife[N], Next[N], order[N][N], pref[N][N], n, v;
^{24}
^{25}
                      0, sizeof(wife[0]) * (n + 2));
26
      memset (wife,
      memset(husband, 0, sizeof(husband[0]) * (n + 2));
27
      memset(Next, 0, sizeof(Next[0]) * (n + 2));
28
29
30
31 void engage (int man, int woman) {
      int exWife = wife[man];
wife[man] = woman;
      husband[woman] = man;
35
36
      if (exWife)
37
        Q.push(exWife);
38
39
40
    void Solve()
41
42
      _clear();
43
44
      cin >> n:
      for(int i = 1; i <= n; ++i)
45
        for(int j = 1; j <= n; ++j)
  cin >> pref[i][j];
46
47
      for(int i = 1; i <= n; ++i)
for(int j = 1; j <= n; ++j) {</pre>
          cin >> v;
          order[i][v] = j;
53
54
      for(int i = 1; i <= n; ++i)</pre>
56
57
        Q.push(i);
58
      int man. woman:
59
      while (Q.size())
60
61
          woman = Q.front(); Q.pop();
62
          man = pref[woman][++Next[woman]];
63
          if(!wife[man] || order[man][woman] < order[man][wife[man]])</pre>
      engage (man, woman);
66
          else
67
      Q.push (woman);
68
69
70
      for(int i = 1; i <= n; ++i)</pre>
71
        cout << husband[i] << endl;</pre>
72
73
74
    int main()
75
      Fast();
77
      int tc = 1; cin >> tc;
      for(int i = 1; i <= tc; ++i) {
80
        if(i > 1) cout << endl;</pre>
82
83
84
85
    /**
86
87
88
89
90
       2 3 1
97
```

#### 4.11 Euler totient sieve

```
1 /*
2 Constraints:
3 1 <= n <= 1e7
```

```
Time Complexity:
       Phi_sieve takes O(n * ln(ln(n)))
        Space Complexity:
10
11
12 int EulerPhi[N];
13
14 void Phi_sieve(int n) {
1.5
       for (int i = 1; i <= n; ++i) {
   EulerPhi[i] = i;</pre>
16
17
       for (int i = 2; i <= n; ++i) {
   if (EulerPhi[i] == i)</pre>
18
       for (int j = i; j <= n; j += i) {
EulerPhi[j] -= EulerPhi[j] / i;</pre>
22
23
24
```

#### 4.12 Mobius sieve

```
Constraints:
     1 <= n <= 1e7
     Time Complexity:
     mu_sieve takes O(n)
     Space Complexity:
     O (MaxN)
10 */
11
12 int mu[N], lp[N], Primes[78522], pnx;
13
14 void mu_sieve(int n) {
15
     mu[1] = 1;
     fill(mu, mu + N, 1);
16
17
      for (int i = 2; i <= n; ++i) {
18
      if (lp[i] == 0) {
         lp[i] = Primes[pnx++] = i;
mu[i] = -1;
20
21
^{22}
        for (int j = 0, nxt; j < pnx && Primes[j] <= lp[i] && (nxt = i * Primes[j]) <= n; ++j) {
23
         lp[nxt] = Primes[j];
mu[nxt] = (lp[i] == Primes[j] ? 0 : -mu[i]);
25
26
27
```

#### 4.13 Wheel sieve

```
1 /**
       1 <= n <= 1e9
       2 <= x <= 9700000
       Time Complexity:
       wheel sieve takes O(n / ln(ln(n)))
       coPrimes takes O(x * ln(ln(x)))
10
       Space Complexity:
       O(MaxN / 32 + n / (ln(n) - 1.08) + x)
11
12 **/
13
16 int Primes[50908031], pnx; /** size of Primes = n / (ln(n) - 1.08) */
17
18
   vector<int> coPrimes(int x) {
19
     int basis[5] = {3, 5, 7, 11, 13};
20
21
      vector<int> ret:
22
     bitset<30100> isCoprime;
23
     isCoprime.set();
24
25
     for (int b : basis)
       for (int d = b * b; d <= x; d += b << 1)
26
27
         isCoprime.reset(d);
```

```
for (int i = 17; i <= x; i += 2)</pre>
        if (isCoprime[i]) ret.push_back(i);
31
33
      ret.push_back(x + 17);
34
35
36
37
    void wheel_sieve(int n) {
      int basis[6] = {2, 3, 5, 7, 11, 13};
vector<int> wheel = coPrimes(2 * 3 * 5 * 7 * 11 * 13);
39
40
      int sz = wheel.size();
41
      for (int k = 0; k < sz; ++k)
42
        inx[wheel[k]] = k;
43
46
      inx[1] = sz - 2;
47
      int inc[sz - 1];
48
49
      for (int i = 1; i < sz; ++i)</pre>
50
        inc[i-1] = wheel[i] - wheel[i-1];
51
52
      for (int p : basis) {
53
        if (n >= p)
54
         Primes[pnx++] = p;
55
56
57
      int c = 0:
      for (11 i = 17; i <= n; i += inc[c++]) {
        if (isPrime[i]) {
         Primes[pnx++] = i;
          int d = inx[i % 30030];
          for (ll j = i * i; j <= n; j += i * inc[d++]) {
63
      isPrime.reset(j);
      if (d == sz - 1) d = 0;
65
67
        if (c == sz - 1) c = 0;
68
69
```

# 5 String Processing

#### 5.1 Trie

```
1 class Trie {
 2 private:
      Trie* children[26]; // Pointer = 8 Byte; 8*26 = 208 Byte
      int prefixs, words; // 8 Byte
      bool iseow; // 1 Byte
char cur_letter; // 1 Byte
      vector <string> lex;
      priority queue <pair <int, string>, vector <pair <int, string>>, greater <pair <int, string>>>
            occurrence; // small at top
10 public:
      Trie(char lett = '\0') {
        memset (children, 0, sizeof (children));
13
        prefixs = words = 0;
14
        iseow = false;
        cur_letter = lett;
16
17
18
      void insert(string &str) { // O(1)
19
        Trie* cur = this:
20
        int inx, strsz = str.size();
        for(int i = 0; i < strsz; ++i) {</pre>
         inx = str[i] - 'a';
          if(cur->children[inx] == nullptr)
      cur->children[inx] = new Trie(str[i]);
26
          cur = cur->children[inx];
27
          cur->prefixs++;
28
29
        cur->iseow = true;
30
        cur->words++:
31
      int search_word(string &str) { // O(1)
34
        Trie* cur = this;
35
        int inx, strsz = str.size();
        for(int i = 0; i < strsz; ++i) {
          inx = str[i] - 'a';
```

```
38
           if(cur->children[inx] == nullptr) {
 39
 40
 41
           cur = cur->children[inx];
 42
 43
         return cur->words;
 44
 45
 46
       int search_prefix(string &str) { // O(1)
 47
         Trie* cur = this;
         int inx = 0, strsz = str.size();
 48
         for(int i = 0; i < strsz; ++i) {
  inx = str[i] - 'a';</pre>
 49
 50
           if(cur->children[inx] == nullptr) {
 51
 52
       return 0:
 54
           cur = cur->children[inx];
 55
 56
         return cur->prefixs;
 57
 58
 59
      bool erase(string &str) {
         if(!search word(str))
 60
 61
           return false:
 62
         Tries cur = this:
 63
 64
         int inx. strsz = str.size();
 65
         for(int i = 0; i < strsz; ++i) {</pre>
 66
           inx = str[i] - 'a';
 67
           if (--cur->children[inx]->prefixs == 0) {
 68
       cur->children[inx] = nullptr;
       return true;
 70
 71
           cur = cur->children[inx];
 72
 73
         if(--cur->words == 0) {
 74
           cur->iseow = false;
 75
 76
         return true;
 77
 78
 79
       void dfs(Trie* node, string s) { // lex order dfs -> traverse all the strings starting from root
 80
 81
         if (node->iseow)
 82
           lex.emplace back(s);
 83
         for (int j = 0; j < 26; ++j)
 86
           if(node->children[j] != nullptr) {
 87
       dfs(node->children[j], s + string(1, node->children[j]->cur_letter));
 88
 89
 90
       {f void} dfs2(Trie* node, string s) { // autocomplete dfs -> traverse all the strings starting from the
 91
             end of the given prefix
         if(node->iseow) {
 92
 93
          if(occurrence.size() < 10) {
 94
      occurrence.push (make_pair (node->words, s));
 95
          } else {
       if (node->words > occurrence.top().first) {
         occurrence.pop();
         occurrence.push(make_pair(node->words, s));
 99
100
101
102
         for(int i = 0; i < 26; ++i) if(node->children[i] != nullptr) {
103
104
       dfs2(node->children[i], s + string(1, node->children[i]->cur_letter));
105
106
107
108
    public:
      vector <string> lex_order() { // all strings in lexicographical order
109
110
         lex.clear();
         Trie* cur = this;
         for(int i = 0; i < 26; ++i) if(cur->children[i] != nullptr) +
       dfs(cur->children[i], string(1, cur->children[i]->cur_letter));
115
116
117
118
       void autocomplete(string &pref) { // suggest top ten words with max frequency
         if(!search_prefix(pref))
119
120
           return:
121
122
         Trie* cur = this:
         int inx, presz = pref.size();
         for(int i = 0; i < presz; ++i) {
           inx = pref[i] - 'a';
           cur = cur->children[inx];
127
```

```
129
         for(int i = 0; i < 26; ++i) if(cur->children[i] != nullptr) {
130
       dfs2(cur->children[i], string(1, cur->children[i]->cur_letter));
131
132
133
         vector <string> st;
134
         while (!occurrence.empty()) {
135
           st.emplace_back(pref + occurrence.top().second);
           occurrence.pop();
136
137
138
         if(cur->iseow) {
139
           st.emplace_back(pref);
140
         while(!st.empty()) {
  cout << st.back() << endl;</pre>
141
142
143
           st.pop_back();
146
```

## 5.2 Knuth Morris Pratt (KMP)

```
* KMP (Knuth-Morris-Pratt) Algorithm
    ** Longest Prefix
   *** proper prefix = all prefixes except the whole string
    *** propre suffix = all suffixes except the whole string
   ** Prefix Function = Failure Function
    *** Given String P of len m, Find F[m];
    *** let t = P[0...i]
    *** f[i] = length of the longest proper prefix of t that is suffix of t
    *** calculating i different ways
    *** match the pattern against itself
    *** O(m) for failure function
13
    *** O(n) for KMP
14
15
16
   vector <int> LongestPrefix(string &p) {
17
     int. psz = p.size():
18
     vector <int> longest_prefix(psz, 0);
19
20
      for(int i = 1, k = 0; i < psz; ++i) {
21
        while (k &  p[k] != p[i]) k = longest_prefix[k-1];
       longest_prefix[i] = (p[k] == p[i] ? ++k : k);
^{24}
      return longest_prefix;
25
26
27
   vector <int> KMP(string &s, string &p) {
28
      int ssz = s.size(), psz = p.size();
30
      vector <int> longest_prefix = LongestPrefix(p), matches;
31
      for(int i = 0, k = 0; i < ssz; ++i) {</pre>
33
        while(k && p[k] != s[i]) k = longest_prefix[k - 1]; // Fail go back
       k += (p[k] == s[i]);
35
37
         matches.emplace_back(i - psz + 1);
38
          k = longest\_prefix[k - 1]; // faill safe and find another pattern
39
40
41
      return matches;
```

# 6 Geometry

#### 6.1 Point

```
1 class point
2 {
3  public :
4   ld x, y;
5
6   point() = default;
7   point(ld _x, ld _y) : x(_x), y(_y) {}
8
9   bool operator < (point other) const {
10   if(fabs(x - other.x) > EPS) // if(x != other.x)
11   return x < other.x;</pre>
```

```
12
        return v < other.v;
13
14
      bool operator == (point other) const {
        return ((fabs(x - other.x) < EPS) && (fabs(y - other.y) < EPS)); // " < EPS " equal to " == zero "
17
18
19
     bool operator > (point other) const {
20
       if(fabs(x - other.x) > EPS)
21
         return x > other.x;
22
        return y > other.y;
23
24
     ld dist(point other) { // Euclidean distance
25
       ld dx = this->x - other.x;
ld dy = this->y - other.y;
26
        return sqrtl(dx * dx + dy * dy);
29
30
31
      ld DEG_to_RAD(ld theta) {
32
        return theta * PI / 180.0;
33
34
35
      ld RAD to DEG(ld theta) {
        return theta * 180.0 / PI;
36
37
38
     point rotate(ld theta) {
39
40
        ld rad = DEG_to_RAD(theta);
        return point (cos (theta) * x - sin (theta) * y,
41
        sin(theta) * x + cos(theta) * y);
```

# 7 More advanced topics

# 7.1 A\*-Algorithm

```
#pragma GCC optimize("Ofast")
    #include <bits/stdc++.h>
    using namespace std;
   typedef int64_t 11;
10
11 void Fast() {
        cin.sync with stdio(0);
12
        cin.tie(0);cout.tie(0);
1.3
14
15
16 const int dr [] = {-1, 0, 1, 0};

17 const int dc [] = {0, 1, 0, -1};

18 const char dir [] = {'U', 'R', 'D', 'L'};
19 map <char, int> inv = { {'U', 0}, {'R', 1}, {'D', 2}, {'L', 3}};
^{21}
   const int N = 1e3 + 9, M = 2e5 + 9, oo = 0x3f3f3f3f3f;
22 const 11 INF = 0x3f3f3f3f3f3f3f3f3f3f;
23
\frac{24}{25}
   char grid[N][N];
   int dis[N][N], n, m, si, sj, ti, tj;
26
   char Par[N][N];
27
28
   vector < pair <int, int> > restorePath(int sr, int sc, int tr, int tc)
29
30
         vector < pair <int, int> > ret;
31
        if (dis[tr][tc] == oo) return ret;
        for(char i = Par[tr][tc]; (sr ^ tr) || (sc ^ tc); i = Par[tr][tc])
34
35
            ret.push_back({tr, tc});
36
            tr += dr[inv[i]];
37
            tc += dc[inv[i]];
38
39
40
        ret.push_back({sr, sc});
41
        reverse (ret.begin(), ret.end());
42
        return ret:
43
44
45 bool valid(int r, int c) {
        return r >= 0 && r < n && c >= 0 && c < m && grid[r][c] != '%';
```

```
47 }
 48
 49
     /** admissible heuristic **/
    int manhattanDistance(int x1, int y1, int x2, int y2) {
         return (abs(x1 - x2) + abs(y1 - y2));
 52
53
54
    int Astar(int sr, int sc, int tr, int tc)
55
56
57
58
         memset(dis, 0x3f, sizeof (dis));
         memset (Par, -1, sizeof (Par));
 59
        priority_queue <tuple <int, int, int> > Q;
 60
61
         dis[srl[scl = 0:
         Q.push({-manhattanDistance(sr, sc, tr, tc), sr, sc});
 63
64
         int hoost, r, c, nr, nc;
 65
         while(Q.size())
66
 67
             tie(hcost, r, c) = Q.top(); Q.pop();
68
             if(r == tr && c == tc) return dis[r][c];
69
70
             for(int i = 0; i < 4; ++i)
 71
 72
                 nr = r + dr[i]:
 73
                 nc = c + dc[i]:
74
 75
                 if (!valid(nr, nc)) continue;
76
 77
                 if(dis[r][c] + 1 < dis[nr][nc])
 78
 79
                     dis[nr][nc] = dis[r][c] + 1;
 80
                     Par[nr][nc] = dir[i ^ 2];
 81
                     Q.push({-dis[nr][nc] -manhattanDistance(nr, nc, tr, tc), nr, nc});
 82
83
84
85
         return -1;
86
87
 88
    void Solve()
89
90
         Fast();
91
 92
         cin >> si >> sj >> ti >> tj >> n >> m;
 93
 94
         for (int i = 0; i < n; ++i)</pre>
 95
             for(int j = 0; j < m; ++j)
96
                cin >> grid[i][j];
97
98
         cout << Astar(si, sj, ti, tj) << endl;</pre>
99
         vector < pair <int, int> > path = restorePath(si, sj, ti, tj);
100
101
         for (auto point : path)
102
            cout << point.first << " " << point.second << endl;</pre>
103 }
104
105 int main()
106
         while (t--) Solve();
108
109
110
111
112
        P -> strat
         . -> target
113
114
115
        input:
0 2 2 3 5 5
116
117
         88P8-
118
         -8---
119
        88888
120
121
123
124
^{125}
126
127
128
129
130 **/
```

#### 7.2 Mo's algorithm

```
#pragma GCC optimize ("Ofast")
    #include <bits/stdc++.h>
    using namespace std;
   typedef int64_t ll;
10 typedef __int128 i128;
11
12 void Fast() {
     cin.sync_with_stdio(0);
13
14
      cin.tie(0);cout.tie(0);
15
17 void File() {
     freopen("input.in", "r", stdin);
freopen("output.out", "w", stdout);
19
20
21
22 const int N = 3e4 + 9, M = 2e5 + 9, oo = 0x3f3f3f3f, Mod = 1e9 + 7;
23 const 11 INF = 0x3f3f3f3f3f3f3f3f3f3;
24 const int BLK = 256:
25
26 struct query
27
28
     int 1, r, id, blk;
29
      query() = default;
30
31
      query(int _l, int _r, int _id) {
        1 = _1;
        r = _r;
id = _id;
33
34
35
        blk = 1 / BLK;
36
37
38
      bool operator < (const query other) const {
39
        if(blk ^ other.blk)
40
          return blk < other.blk;</pre>
41
        return (blk & 1) ? r < other.r : r > other.r;
42
43 } queries[M];
44
45 int res[M], freq[M << 3], cur;
46
47
   void add(int id) {
     cur += (++freq[id] == 1);
49
50
51 void remove(int id) {
52
     cur -= (--freq[id] == 0);
53
54
55 int get res() {
56
     return cur:
57
59 int cur_1, cur_r, 1, r, n, q, a[N];
    void Solve()
      for(int i = 1; i <= n; ++i) cin >> a[i];
65
      cin >> q;
66
      for(int i = 1; i <= q; ++i) {
67
68
        cin >> 1 >> r;
69
        queries[i] = query(1, r, i);
70
71
72
      sort (queries + 1, queries + 1 + q);
73
      cur_1 = 1, cur_r = 0; // assign to right invalid index for(int \ i = 1; \ i \le q; \ ++i)
74
75
76
77
          int ql = queries[i].1;
78
          int qr = queries[i].r;
79
80
          // Add right
81
          while(cur_r < qr) add(a[++cur_r]);</pre>
82
83
          while(cur_l > ql) add(a[--cur_l]);
84
           // Remove right
85
          while(cur_r > qr) remove(a[cur_r--]);
86
             Remove left
          while(cur_l < ql) remove(a[cur_l++]);</pre>
87
89
          res[queries[i].id] = get_res();
90
91
      for(int i = 1; i <= q; ++i)
```

# 7.3 Square root decomposition

```
#pragma GCC optimize ("Ofast")
          #include <bits/stdc++.h>
         #define endl
         using namespace std;
         typedef int64_t l1;
 10 typedef __int128 i128;
11
12
         void Fast() {
            cin.sync with stdio(0);
              cin.tie(0);cout.tie(0);
14
15
16
         const int N = 5e5 + 9, M = 1e3 + 9, oo = 0x3f3f3f3f3f, Mod = 1e9 + 7;
         const 11 INF = 0x3f3f3f3f3f3f3f3f3f;
         const int BLK = 256;
20
21 int n, q, a[N], type, x, y, z;
22
        vector <int> bs[M];
24
         int query(int 1, int r, int val)
25
              int cur 1 = 1 / BLK;
26
               int cur r = r / BLK:
              int ans = 0;
               if(cur_l == cur_r) {
                   for (int i = 1; i <= r; ++i)
                        ans += (a[i] >= val);
33
                   for(int i = 1, _end = (cur_l + 1) * BLK; i < _end; ++i)</pre>
                   for(int i = 1, _c.ma - (cm_1 - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1, _ - 1
35
37
                   for (int i = cur_r * BLK; i <= r; ++i)</pre>
39
                      ans += (a[i] >= val);
40
41
               return ans:
42
44
         void build()
               for(int i = 0; i < n; ++i)
                   bs[i / BLK].emplace_back(a[i]);
48
 49
               for(int i = 0; i < M; ++i)</pre>
50
                   sort(bs[i].begin(), bs[i].end());
51
53
         void update(int id, int delta)
54
               int pos = lower_bound(bs[id / BLK].beqin(), bs[id / BLK].end(), a[id]) - bs[id / BLK].beqin();
              bs[id / BLK][pos] = delta;
               sort(bs[id / BLK].begin(), bs[id / BLK].end());
               a[id] = delta;
59
61
          void Solve()
63
               for(int i = 1; i <= n; ++i) cin >> a[i];
              build();
67
               cin >> q;
               while (q--)
70
71
                         cin >> type >> x >> y;
72
                        if(type == 0)
```

```
cin >> z;
75
        cout << query(x, y, z) << endl;</pre>
76
     update(x, y);
79
80 }
81
82 int main()
83
     Fast():
84
85
86
     int tc = 1;
     for(int i = 1; i <= tc; ++i)
87
88
       Solve();
```

## 8 Miscellaneous

## 8.1 Double comparison

```
bool approximatelyEqual(double a, double b, double epsilon)

return fabs(a - b) <= ((fabs(a) < fabs(b) ? fabs(b) : fabs(a)) * epsilon);

bool essentiallyEqual(double a, double b, double epsilon)

return fabs(a - b) <= ((fabs(a) > fabs(b) ? fabs(b) : fabs(a)) * epsilon);

bool definitelyGreaterThan(double a, double b, double epsilon)

return (a - b) > ((fabs(a) < fabs(b) ? fabs(b) : fabs(a)) * epsilon);

return (a - b) > ((fabs(a) < fabs(b) ? fabs(b) : fabs(a)) * epsilon);

bool definitelyLessThan(double a, double b, double epsilon)

return (b - a) > ((fabs(a) < fabs(b) ? fabs(b) : fabs(a)) * epsilon);

return (b - a) > ((fabs(a) < fabs(b) ? fabs(b) : fabs(a)) * epsilon);</pre>
```

# 8.2 Fast input output

```
Fast Input/Output method for C++:
      1. cin(with sync_with_stdio(false) & cin.tie(nullptr)):
      - |n| = 5e6| => 420ms
     - |n = 1e7| => 742ms
      - |n| = 5e6| => 895ms
      read (using getchar()):
12
      - /n = 5e6/ => 173ms
13
      - |n = 1e7| => 172ms
14
      - |n| = 5e6| => 340ms
1.5
16 **/
17
   ll readll () {
18
     bool minus = false;
     unsigned long long result = 0;
24
25
       if (ch == '-') break;
       if (ch >= '0' && ch <= '9') break;
26
       ch = getchar();
28
29
     if (ch == '-') minus = true;
30
     else result = ch - '0';
31
32
33
     while (true)
34
       ch = getchar();
       if (ch < '0' || ch > '9') break;
35
       result = result * 10 + (ch - '0');
```

```
37
      if (minus) return -(11) result;
41
43 int readi () {
     bool minus = false;
45
      unsigned int result = 0;
46
47
      ch = getchar();
48
     while (true) {
  if (ch == '-') break;
49
50
        if (ch >= '0' && ch <= '9') break;
51
        ch = getchar();
54
      if (ch == '-') minus = true;
56
      else result = ch - '0';
58
59
        ch = getchar();
        if (ch < '0' || ch > '9') break;
60
        result = result * 10 + (ch - '0');
61
62
63
64
      if (minus) return - (int) result:
65
      return result:
66
```

#### 8.3 Gcd & Lcm

```
1 ll gcd(ll a, ll b) { // binary GCD uses about 60% fewer bit operations
2
     if (!a) return b:
      int shift = __builtin_ctz(a | b);
     a >>= builtin ctz(a);
      while (b) {
       b >>= __builtin_ctz(b);
        swap(a, b);
14
      return a << shift;
15
16
17 ll lcm(ll a, ll b) {
18
     return a / gcd(a, b) * b;
19
```

#### 8.4 Modular calculations

```
- It also has important applications in many tasks unrelated to arithmetic, since it can be used
           with any operations that have the property of associativity:
 3 */
   // 1. Modular Exponentiation
   ll binExp(ll a, ll b, ll p) {
     ll res = 1;
     while (b) {
       if (b & 111)
        res = res * a % p;
       a = a * a % p;
16
18
   // 2. Modular Multiplication
20 ll binMul(ll a, ll b, ll p) {
     ll res = 0;
     a %= p:
      while (b)
      if (b & 111)
        res = (res + a) % p;
       a = (a + a) % p;
```

```
27 b >>= 1;

28 }

29 return res;

30 }

31 

32 // 3. Modular Multiplicative Inverse

33 

34 ll modInv(ll b, ll p) {

35 return binExp(b, p - 2, p); // Guaranteed that p is a Prime Number 16 

36 }
```

#### 8.5 Overloaded operators to accept 128 Bit integer

```
typedef __uint128_t
    typedef __int128
    template <class T> string to_string(T x)
      int sn = 1; if (x < 0) sn = -1, x *= sn; string s = "";
     do { s = "0123456789"[x % 10] + s, x /= 10; } while(x);
return (sn == -1 ? "-" : "") + s;
10
11 auto str_to_int(string x)
12 {
      ui128 ret = (x[0] == '-' ? 0 : x[0] - '0');
13
     for(int i = 1; i < x.size(); ++i) ret = ret * 10 + (x[i] - '0');
return (x[0] == '-' ? -1 * (i128)ret : ret);</pre>
15
16
    istream & operator >> (istream & in, i128 & i) noexcept { string s; in >> s; i = str_to_int(s); return
19 ostream & operator << (ostream & os, const i128 i) noexcept { os << to_string(i); return os; }
    istream & operator >> (istream & in, uil28 & i) noexcept { string s; in >> s; i = str_to_int(s);
21 ostream & operator << (ostream & os, const uil28 i) noexcept { os << to_string(i); return os; }
```

# 8.6 Policy based data structures

```
1 #if __cplusplus >= 20140ZL
2 #include <ext/pb_ds/assoc_container.hpp>
3 #include <ext/pb_ds/ree_policy.hpp>
4 #endif
5
6 #if __cplusplus >= 20140ZL
vsing namespace __gnu_cxx;
8 using namespace __gnu_pbds;
9 #endif
10
11 template <class T, typename Comp = less <T> >
11 template <class T, typename Comp = less <T> >
12 using indexed_set = tree <T, null_type, Comp, rb_tree_tag, tree_order_statistics_node_update>;
13
14 template <typename K, typename V, typename Comp = less <K>>
15 using indexed_map = tree <K, V, Comp, rb_tree_tag, tree_order_statistics_node_update>;
```

#### 8.7 stress test

```
1 mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().count());
    /** 64-bit signed int Generator
    i64 int64(i64 a, i64 b) {
      return uniform_int_distribution <i64> (a, b) (rng);
    /** Customize your Generator depending on the input
11 void gen () {
     ofstream cout ("input.in");
14
      cout << t << endl;
1.5
16
      while (t--) {
        i32 n = int64(1, 100), m = int64(1, 100);
cout << n << " " << m << endl;
17
18
19
20
           i32 u = int64(1, n), v = int64(1, n), c = int64(1, 4);

cout << u << " " << v << " " << c << endl;
^{21}
```

```
23
   i32 main (i32 arg, char* args[]) {
    fast();
29
     i32 tc = 0;
i32 limit = 100;
30
31
32
     if(arg != 3) return 0;
33
    34
35
36
37
     bf = flags + args[1] + " " + args[1] + ex;
oz = flags + args[2] + " " + args[2] + ex;
39
      char bff[bf.size() + 1];
40
     char ozz[oz.size() + 1];
41
     strcpy(bff, bf.c_str());
42
     strcpy(ozz, oz.c_str());
43
44
45
46
47
48
     // compile command
     system(bff);
     system(ozz);
   ex = ".out";
pr = "./";
bf = pr + args[1] + " < input.in > " + args[1] + ex;
49
```

```
oz = pr + args[2] + " < input.in > " + args[2] + ex;
51
52
      strcpy(bff, bf.c_str());
      strcpy(ozz, oz.c_str());
55
       while (++tc <= limit) {</pre>
56
57
         cerr << tc << endl;
58
         // run command
59
         system(bff);
60
         system(ozz);
61
         ifstream brute_forces("brute_force.out");
ifstream optimizes("optimized.out");
63
64
         string brute_force, optimized;
getline(brute_forces, brute_force, (char)EOF);
getline(optimizes, optimized, (char)EOF);
65
67
68
69
         if(brute_force != optimized) {
70
           cerr << "Wrong Answer" << endl;
71
           break;
72
         } else if (tc == limit) {
73
74
           cout << "Accepted" << endl;</pre>
75 }
76 }
```