# ICPC Team Reference Material

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# 1 Setup

# 1.1 Vimrc

```
1 let mapleader = "\"
2 syntax on
3 filetype plugin on
4 set nocompatible
5 set autoread
6 set foldmethod=marker
7 set autoindent
8 set clipboard+=unnamedplus
9 set number relativenumber
10 set shiftwidth=2 softtabstop=2 expandtab
11 map <leader>c :w! && !compile %:p:r<CR>
12 vmap < <gv</pre>
13 vmap > >gv
```

# 1.2 Capslock as Escape

```
1 setxkbmap -layout us
2 xmodmap -e 'clear Lock'
3 xmodmap -e 'keycode 66 = Escape'
```

# 1.3 Compilation

# 2 Graph algorithms

#### 2.1 Adjacency list representation

```
1 class Graph {
 2 public:
      vector <int> _head, _next, _to, _cost;
      int edge_number;
     Graph() = default;
     Graph (int V, int E, bool isDirec) {
       isDirected = isDirec;
       _head.assign(V + 9, 0);
10
       _next.assign(isDirected ? E + 9 : E * 2 + 9, 0);
11
       _to.assign(isDirected ? E + 9 : E * 2 + 9, 0);
12
        // _cost.assign(isDirected ? E + 9 : E * 2 + 9, 0);
13
       edge_number = 0;
14
15
16
17
      void addEdge(int u, int v, int w = 0) {
18
        _next[++edge_number] = _head[u];
        _to[edge_number] = v;
20
           _cost[edge_number] = w;
^{21}
        _head[u] = edge_number;
24
      void addBiEdge(int u, int v, int w = 0) {
       addEdge(u, v, w);
26
        addEdge(v, u, w);
27
      void dfs(int node) {
       vis[node] = true;
        for (int i = _head[node]; i; i = _next[i]) if(!vis[_to[i]]) {
      dfs(_to[i]);
35
```

# 2.2 Depth first search (DFS)

```
1 void DFS(int node)
2 {
3    vis[node] = true;
4    for(int i = Head[node]; i; i = Next[i])
5     if(!vis[To[i]);
6     DFS(To[i]);
7 }
```

# 2.3 Breadth first search (BFS)

```
1 void BFS(int src)
2
      memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
      memset (Par, -1, sizeof (Par[0]) * (n + 2));
      queue <int> q;
      q.push(src);
      dis[src] = 0;
10
      while (q.size())
1.1
12
          u = q.front(); q.pop();
13
         for(int i = Head[u]; i; i = Next[i]) if(dis[To[i]] == 00) {
14
       dis[To[i]] = dis[u] + 1;
Par[To[i]] = u;
15
16
       q.push(To[i]);
```

#### 2.4 0-1 BFS

```
1 const int N = 1e5 + 9, M = 2e5 + 9, oo = 0x3f3f3f3f3f;
 2 11 INF = 0x3f3f3f3f3f3f3f3f3f;
   int Head[N], Par[N], Next[M], To[M], Cost[M], ne, n, m, u, v, st, tr, tax;
   void addEdge(int from, int to, int cost) {
     Next[++ne] = Head[from];
     Head[from] = ne;
     Cost[ne] = cost;
1.1
     To[ne] = to;
12
13
14 void clear() {
     memset(Head, 0, sizeof(Head[0]) * (n + 2));
1.5
17 }
19 void BFS(int src, int trg)
20
     memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
     memset(Par, -1, sizeof(Par[0]) * (n + 2));
     degue <int> 0:
24
     O.push front (src);
25
     dis[src] = 0;
26
     int node:
27
28
     while (O.size())
29
30
          node = Q.front(); Q.pop_front();
31
         if (node == trg) return;
          for(int i = Head[node]; i; i = Next[i])
     if(dis[node] + Cost[i] < dis[To[i]])</pre>
35
36
          dis[To[i]] = dis[node] + Cost[i];
37
         if (Cost[i])
38
           Q.push_back(To[i]);
39
40
           Q.push_front(To[i]);
41
42
43
```

# 2.5 0-1 BFS (grid)

```
const int N = 1e3 + 9, M = 2e5 + 9, oo = 0x3f3f3f3f3f;
   int dis[N][N], n, m, si, sj, ti, tj;
   char grid[N][N];
10 bool valid(int r, int c) {
    return r >= 1 && r <= n && c >= 1 && c <= m;
11
12
13
   int ZBFS(int sr, int sc, int tr, int tc)
     memset(dis, 0x3f, sizeof (dis)); // memset(dis, 0x3f, n * m) we don't do that here
     deque <pair <int, int> > Q;
     dis[sr][sc] = 0;
19
     Q.push_front({sr, sc});
20
21
     int r, c, nr, nc, ncost;
22
     while (O.size())
23
24
         tie(r, c) = Q.front(); Q.pop_front();
         if (r == tr && c == tc) return dis[r][c];
         for(int i = 0; i < 8; ++i)</pre>
29
       nr = r + dr[i];
30
       nc = c + dc[i];
31
32
       if(!valid(nr, nc)) continue;
33
       ncost = (i != grid[r][c]);
34
35
       if(dis[r][c] + ncost < dis[nr][nc])
36
37
           dis[nr][nc] = dis[r][c] + ncost;
38
39
           if (ncost)
40
       Q.push_back({nr, nc});
```

# 2.6 Articulation points and bridges

```
rootChildren, dfs_timer, bridgeInx;
 2 bool Art[N];
 3
   vector < pair <int, int> > bridges(M);
   void _clear() {
                   0, sizeof(Head[0]) * (n + 2));
     memset (Head,
     memset(dfs_num, 0, sizeof(dfs_num[0]) * (n + 2));
     memset(Par, -1, sizeof(Par[0]) * (n + 2));
memset(Art, 0, sizeof(Art[0]) * (n + 2));
     ne = dfs_timer = bridgeInx = 0;
10
11
12
13
   void Tarian(int node) {
14
     dfs_num[node] = dfs_low[node] = ++dfs_timer;
     for(int i = Head[node]; i; i = Next[i]) {
15
       if (dfs_num[To[i]] == 0) {
16
         if(node == root) ++rootChildren;
17
         Par[To[i]] = node;
         Tarjan(To[i]);
20
         dfs_low[node] = Min(dfs_low[node], dfs_low[To[i]]);
21
22
         if(dfs_low[To[i]] >= dfs_num[node])
23
     Art[node] = true;
24
25
         if(dfs_low[To[i]] > dfs_num[node])
26
     bridges[bridgeInx++] = make_pair(node, To[i]);
27
       else if(To[i] != Par[node])
28
         dfs_low[node] = Min(dfs_low[node], dfs_num[To[i]]);
29
30
31 }
   int main() {
     for(int i = 1; i <= n; ++i)</pre>
35
       if(dfs_num[i] == 0) {
36
37
         rootChildren = 0;
38
         Tarjan(i);
39
         Art[root] = (rootChildren > 1);
40
41
42
     cout << "Art Points :\n";
43
     for(int i = 1; i <= n; ++i)</pre>
44
       if (Art[i])
         cout << i << " ";
      cout << "\nBridges :\n";
48
      for(int i = 0; i < bridgeInx; ++i)</pre>
       cout << bridges[i].first << " - " << bridges[i].second << endl;</pre>
50
```

# 2.7 Bi-connected components

```
1 int Head[N], Next[M], To[M], Par[N], dfs_num[N], dfs_low[N], ne, n, m, u, v, root, rootChildren,
         dfs_timer, Stack[N], top, ID;
 2 bool Art[N];
3 vector < vector <int> > BiCCs(N), BiCCIDs(N);
   void addEdge(int from, int to) {
     Next[++ne] = Head[from];
      Head[from] = ne;
     To[ne] = to;
9
10
11 void clear() {
12
    memset (Head,
                         0, sizeof(Head[0])
                       0, sizeof(dfs_num[0])
-1, sizeof(Par[0])
     memset (dfs_num,
                                                 * (n + 2));
                                                  * (n + 2));
     memset (Par.
15
     memset (Art,
                        0, sizeof(Art[0])
                                                  * (n + 2));
     ne = dfs timer = top = ID = 0;
16
     BiCCs = BiCCIDs = vector < vector <int> > (N);
```

```
20
   void Tarjan (int node)
^{21}
      dfs_num[node] = dfs_low[node] = ++dfs_timer;
      Stack[top++] = node;
^{24}
25
      for(int i = Head[node]; i; i = Next[i]) {
        if(dfs_num[To[i]] == 0) {
26
27
          if(node == root) ++rootChildren;
          Par[To[i]] = node;
28
29
          Tarjan(To[i]);
30
          dfs_low[node] = Min(dfs_low[node], dfs_low[To[i]]);
31
          if(dfs_low[To[i]] >= dfs_num[node])
32
33
34
        Art[node] = true;
35
36
        for (int x = -1; x ^ To[i];)
37
38
            x = Stack[--top];
39
            BiCCIDs[x].emplace_back(ID);
40
            BiCCs[ID].emplace_back(x);
41
42
        BiCCIDs[node].emplace_back(ID);
43
        BiCCs[ID].emplace_back(node);
44
45
        else if(To[i] != Par[node])
46
47
          dfs_low[node] = Min(dfs_low[node], dfs_num[To[i]]);
48
49
50
51
52
53
      for(int i = 1; i <= n; ++i)</pre>
54
        if(dfs_num[i] == 0) {
55
          root = i;
          rootChildren = 0;
56
57
          Tarjan(i);
          Art[root] = (rootChildren > 1);
58
59
60
     for(int i = 1; i <= ID; ++i) {
  cout << "Component : " << i << " contains : ";</pre>
61
62
        for(int j = 0; j < (int)BiCCs[i].size(); ++j)</pre>
63
64
          cout << BiCCs[i][j] << " \n"[j == BiCCs[i].size() - 1];</pre>
65
```

# 2.8 Bipartite graph

```
1 bool checkBiPartite(int node, int par = 0) {
     if(vis[nodel)
        return color[par] != color[node];
     color[node] = color[par] ^ 1;
      vis[node] = true;
     bool ok = true;
      for(int i = Head[node]; i; i = Next[i])
       if(To[i] != par)
10
         ok &= checkBiPartite(To[i], node);
1.1
     return ok;
12 3
13
14 int main() {
1.5
     bool isBiPartite = true:
     for(int i = 1; i <= n; ++i)
16
17
       if(!vis[i])
18
         isBiPartite &= checkBiPartite(i);
     cout << (isBiPartite ? "YES" : "NO") << endl;</pre>
19
20
```

# 2.9 Bellman ford

```
return true;
       return false;
10
    bool Bellman_Ford(int src)
13
14
       memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
15
       memset(Par, -1, sizeof(Par[0]) * (n + 2));
16
17
18
       bool newRelaxation = true;
19
20
       for(int i = 2; i <= n && newRelaxation; ++i) {</pre>
         newRelaxation = false;
for(int i = 1; i <= n; ++i)</pre>
21
       for(int j = Head[i]; j; j = Next[j])
if(dis[i] < INF && dis[i] + Cost[j] < dis[To[j]]) {</pre>
         dis[To[j]] = dis[i] + Cost[j];
26
27
         newRelaxation = true;
28
29
30
       return hasNC();
31 3
```

# 2.10 Connected components

```
1  void DFS(int node) {
2   visited[node] = true;
3   for(int e = Head[node]; e; e = Next[e])
4   if(!visited[To[e]])
5   DFS(To[e]);
6  }
7  
8  int main() {
9   for(int node = 1; node <= n; ++node)
10   if(!visited[node])
11   ++CCS, DFS(node);
12   cout << CCS << endl;
13 }</pre>
```

# 2.11 Cycle detection

```
void DFS(int node, int parent = -1)
      if (hasCycle |= visited[node])
       return:
     visited[node] = true;
      for(int i = Head[node]; i; i = Next[i])
       if(To[i] != parent)
         DFS(To[i], node);
10
11
12 int main() {
13
     for(int i = 1; i <= n; ++i)
       if(!visited[i])
15
         DFS(i);
      cout << (hasCycle ? "YES" : "NO") << endl;</pre>
16
17
```

# 2.12 Directed cyclic graph into acyclic

```
13
     HeadDAG[from] = neDAG;
14
     CostDAG[ne] = cost;
15
     ToDAG[neDAG] = to;
     ++out[from];
17
19
   void _clear() {
     20
21
22
     memset(compID, 0, sizeof(compID[0]) * (n + 2));
23
     memset(compSize, 0, sizeof(compSize[0]) * (n + 2));
24
     memset(HeadDAG, 0, sizeof(HeadDAG[0]) * (n + 2));
memset(out, 0, sizeof(out[0]) * (n + 2));
25
     ne = dfs_timer = top = neDAG = ID = 0;
26
27
29
   void Tarjan (int node)
30
      dfs_num[node] = dfs_low[node] = ++dfs_timer;
32
      in_stack[Stack[top++] = node] = true;
33
34
     for(int i = Head[node]; i; i = Next[i]) {
35
       if(dfs_num[To[i]] == 0)
36
         Tarian(To[i]):
37
       if(in_stack[To[i]])
38
         dfs_low[node] = Min(dfs_low[node], dfs_low[To[i]]);
39
40
41
     if(dfs_num[node] == dfs_low[node]) {
42
43
44
       for(int cur = -1; cur ^ node;) {
45
         in_stack[cur = Stack[--top]] = false;
46
         compID[cur] = ID;
47
         ++compSize[ID];
48
49
50 }
51
52 void Tarjan() {
     for(int i = 1; i <= n; ++i)
53
54
       if(dfs_num[i] == 0)
55
         Tarjan(i);
56 }
57
58
   void DFS(int node)
59
      dfs_num[node] = 1;
     for(int i = Head[node]; i; i = Next[i]) {
62
        if(compID[node] != compID[To[i]])
63
     addEdgeDAG(compID[node], compID[To[i]]);
64
         if(dfs_num[To[i]] == 0)
     DFS(To[i]);
65
66
67
68
69
   void construct dag() {
     memset(dfs_num, 0, sizeof(dfs_num[0]) * (n + 2));
70
     for(int i = 1; i <= n; ++i)
       if(dfs_num[i] == 0)
74
75
```

# 2.13 Dijkstra (dense graph)

```
1 /** Diikstra on dense graphs
       complexity : O(n^2 + m)
   int Head[N], Par[N], Next[M], To[M], Cost[M], ne, n, m, u, v, st, tr, tax;
   ll dis[N];
   void Dijkstra(int src, int V)
      memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
     memset(Par, -1, sizeof(Par[0]) * (n + 2));
      vector <bool> mark(V + 1, false);
13
14
      dis[src] = 0;
1.5
      for(int i = 1; i <= V; ++i) {</pre>
16
        int node = 0;
        for (int j = 1; j <= V; ++j)
17
         if(!mark[j] && dis[j] < dis[node])</pre>
18
19
     node = j;
20
        if(dis[node] == INF) break;
```

# 2.14 Dijkstra (grid)

```
const int dr[] = { 1, -1, 0, 0, 1, 1, -1, -1 };
const int dc[] = { 0, 0, 1, -1, 1, -1, 1, -1 };
const char dir[] = {'D', 'U', 'R', 'L'};
    const int N = 1e3 + 9, M = 2e5 + 9, oo = 0x3f3f3f3f3f;
    int grid[N][N], dis[N][N], n, m;
    bool valid(int r, int c) {
     return r >= 1 && r <= n && c >= 1 && c <= m;
11
13
    void Dijkstra(int sr, int sc)
14
15
      memset(dis, 0x3f, sizeof(dis)); // memset(dis, 0x3f, n * m) we don't do that here
16
      priority_queue <tuple <int, int, int> > Q;
17
18
      dis[sr][sc] = grid[sr][sc];
19
      Q.push({-grid[sr][sc], sr, sc});
20
21
      int cost, r, c, nr, nc;
       while (O.size())
           tie(cost, r, c) = Q.top(); Q.pop();
25
          if((-cost) > dis[r][c]) continue; // lazy deletion
26
27
           for (int i = 0; i < 4; ++i)
28
        nr = r + dr[i];
29
30
        nc = c + dc[i];
31
        if(!valid(nr, nc)) continue;
32
33
        if(dis[r][c] + grid[nr][nc] < dis[nr][nc])
34
35
36
             dis[nr][nc] = dis[r][c] + grid[nr][nc];
37
             Q.push({-dis[nr][nc], nr, nc});
41 }
```

# 2.15 Dijkstra (negative weighted graph)

```
1 int Head[N], Par[N], Next[M], To[M], Cost[M], ne, n, m, u, v, st, tr, tax;
    void Dijkstra(int src)
      memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
     memset(Par, -1, sizeof(Par[0]) * (n + 2));
      priority_queue <pair <11, int> > Q;
      dis[src] = 0;
10
     Q.push({-dis[src], src});
      int node;
      while(Q.size()) {
        tie(cost, node) = Q.top(); Q.pop();
17
        if((-cost) > dis[node]) continue;
18
19
        for(int i = Head[node]; i; i = Next[i])
20
         if(dis[node] + Cost[i] < dis[To[i]])</pre>
21
        dis[To[i]] = dis[node] + Cost[i];
       Q.push({-dis[To[i]], To[i]});
        Par[To[i]] = node;
```

#### 2.16 Dijkstra (sparse graph)

```
/** Dijkstra on sparse graphs
       - complexity : O(n + m)logn -> O(nlogn + m)
       - Single Source Single Destination Shortest Path Problem
       - Positive Weight Edges only
       Subpaths of shortest paths from u to v are shortest paths!
   int Head[N], Par[N], Next[M], To[M], Cost[M], ne, n, m, u, v, st, tr, tax;
 8
   ll dis[N]:
10 void Dijkstra(int src. int trg)
11 {
     memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
     memset(Par, -1, sizeof(Par[0]) * (n + 2));
     priority_queue <pair <11, int> > Q;
16
      dis[src] = 0;
17
     Q.push({-dis[src], src});
18
19
     int node:
20
21
      while(Q.size()) {
22
       tie(cost, node) = Q.top(); Q.pop();
23
24
       if((-cost) > dis[node]) continue; // lazy deletion
25
       if (node == trg) return;
                                         // cheapest cost in case of positive weight edges
26
27
       for(int i = Head[node]; i; i = Next[i])
         if(dis[node] + Cost[i] < dis[To[i]])</pre>
29
30
        dis[To[i]] = dis[node] + Cost[i];
31
       Q.push({-dis[To[i]], To[i]});
32
       Par[To[i]] = node;
33
34
35
```

# 2.17 Edge classification

```
1 int Head[N], Next[M], To[M], Par[N], in_time[N], ne, n, m, u, v, dfs_timer;
 2 char dfs_num[N];
    void edgeClassification(int node)
      dfs num[node] = EXPLORED:
      in_time[node] = ++dfs_timer;
      for(int i = Head[node]; i; i = Next[i])
10
          if (dfs num[To[i]] == UNVISITED)
11
12
13
        cout << "Tree Edge : " << node << " -> " << To[i] << endl;</pre>
        Par[To[i]] = node;
15
        edgeClassification(To[i]);
16
17
18
          else if (dfs_num[To[i]] == VISITED)
19
        /** Cross Edges only occur in directed graph */
20
        if(in_time[To[i]] < in_time[node])
  cout << "Cross Edge : " << node << " -> " << To[i] << endl;</pre>
21
22
23
24
          cout << "Forward Edge : " << node << " -> " << To[i] << endl;
25
26
          else if (dfs_num[To[i]] == EXPLORED)
27
29
          cout << "Bi-Directional Edge : " << node << " -> " << To[i] << endl;</pre>
30
31
          cout << "Backward Edge : " << node << " -> " << To[i] << " (Cycle)" << endl;
32
33
34
35
     dfs_num[node] = VISITED;
36
37
38 int main() {
     for(int i = 1; i <= n; ++i) if(!dfs num[i])</pre>
39
            edgeClassification(i);
41
```

#### 2.18 Eulerian tour tree

```
1 \quad \textbf{int} \ \ \text{Head[N], To[M], Next[M], Cost[M], ne, n, m, u, v, w, Last[N], First[N], euler\_tour[1 + N << 1];}
2 ll Height[1 + N << 1];
3 int euler_timer;
5 void clear() {
 6
     memset (Head,
                          0, sizeof(Head[0])
                                                     * (n + 2));
                                                     * (n + 2)):
      memset (Last.
                          0, sizeof(Last[0])
      memset (First.
                          0, sizeof(First[0])
                                                     * (n + 2));
     ne = euler_timer = 0;
10
11
12
       euler\_tour[1 .. n * 2 - 1] = which records the sequence of visited nodes
13
       Height[1..n * 2 - 1] = which records the depth of each visited node
                           = records the index of the first occurrence of node i in euler_tour
17
                           = records the index of the last occurrence of node i in euler_tour
18
19
   void EulerianTour(int node, ll depth = 0)
20
21
      euler tour[++euler_timer] = node;
23
      Height[euler_timer] = depth;
      First[node] = euler_timer;
      for(int i = Head[node]; i; i = Next[i])
       if(First[To[i]] == 0)
      EulerianTour(To[i], depth + Cost[i]);
31
      euler_tour[++euler_timer] = node;
      Height[euler_timer] = depth;
33
34
35
      Last[node] = euler_timer;
36
37
38
   void show() {
      for(int i = 1; i < (n << 1); ++i) cout << euler tour[i] << " ";cout << endl;</pre>
      for(int i = 1; i < (n << 1); ++i) cout << Height[i] << " "; cout << endl; for(int i = 1; i <= n; ++i) cout << First[i] << " "; cout << endl;
      for(int i = 1; i <= n; ++i)</pre>
                                          cout << Last[i] << " ";
46
      EulerianTour(1);
48
      show();
49
```

#### 2.19 Floodfill

```
1 /** check if there is a path from (0, 0) to (n - 1, m - 1) using '.' only **/
 3 int dr[4] = \{1, -1, 0, 0\};
   int dc[4] = \{0, 0, 1, -1\};
   char grid[N][M];
 6 int n, m;
   bool valid(int r, int c) {
     return r >= 0 && r < n && c >= 0 && c < m && grid[r][c] == '.';
10 }
11
12 bool isDis(int r, int c) {
13
     return r == n - 1 && c == m - 1;
14
   bool FloodFill(int r, int c) {
     if(!valid(r, c)) return false;
     if(isDis(r, c)) return true;
19
20
      grid[r][c] = '#';
21
      for(int i = 0; i < 4; ++i)
       if(FloodFill(r + dr[i], c + dc[i])) return true;
23
     return false:
25
26
27
     cout << (FloodFill(0, 0) ? "YES" : "NO") << endl;</pre>
```

#### 2.20 Floyd warshall (all-pairs shortest path)

```
1 /** -The graph has a 'negative cycle' if at the end of the algorithm,
     the distance from a vertex v to itself is negative.
      - before k-th phase the value of d[i][j] is equal to the length of
     the shortest path from vertex i to the vertex j,
     if this path is allowed to enter only the vertex with numbers smaller than k
     (the beginning and end of the path are not restricted by this property).
10 int Par[N][N], n, m, u, v, tax;
11 ll adj[N][N], dis[N][N];
12
13
   vector <int> restorePath(int st, int tr)
      vector <int> path;
      if(dis[st][tr] == INF) return path;
      for(int i = tr; st ^ i; i = Par[st][i])
19
       path.push_back(i);
20
21
     path.push_back(st);
22
      reverse(path.begin(), path.end());
23
     return path:
24
25
26
    void Flovd Warshall()
27
      for(int i = 1; i <= n; ++i)</pre>
       for (int j = 1; j <= n; ++j)
         Par[i][j] = i;
31
       for(int i = 1; i <= n; ++i)
for(int j = 1; j <= n; ++j)</pre>
33
34
35
     if(dis[i][k] + dis[k][j] < dis[i][j])</pre>
36
         dis[i][j] = dis[i][k] + dis[k][j];
Par[i][j] = Par[k][j];
37
38
39
40 }
```

# 2.21 Minimum spanning tree (Kruskal)

```
1 int n, m, u, v, w;
 2 vector < tuple <int, int, int> > edges;
   UnionFind uf:
   pair < ll, vector < pair <int, int> > > Kruskal()
     sort(edges.begin(), edges.end());
     vector < pair <int, int> > mstEdges;
     int from, to, cost;
     ll minWieght = 0;
      for(tuple <int, int, int> edge : edges)
14
15
          tie(cost, from, to) = edge;
16
         if(uf.union_set(from, to))
17
18
       minWieght += cost;
19
       mstEdges.push_back(make_pair(from, to));
20
21
22
     if(mstEdges.size() == n - 1)
23
       return make_pair(minWieght, mstEdges);
26
      return make_pair(-1, vector < pair <int, int> > ());
27
```

# 2.22 Kth ancestor and lowest common ancestor (binary lifting)

```
1 int Head[N], To[M], Next[M], Par[N], up[N][LOG + 1], Log[N], Level[N], ne, n, u, v, q;
3
     memset(Head, 0, sizeof(Head[0]) * (n + 2));
      memset(Par, 0, sizeof(Par[0]) * (n + 2));
memset(Level, 0, sizeof(Level[0]) * (n + 2));
10 int lastBit(int a) {
11
     return (a & -a);
12
13
14
   void logCalc()
15
      for(int i = 2; i < N; ++i)</pre>
        Log[i] = Log[i >> 1] + 1;
19
20
21
    void DFS(int node, int depth = 0)
22
23
      Level[node] = depth:
24
      up[node][0] = Par[node]; // Par[root] = root
25
      for(int i = 1; i <= LOG; ++i) {</pre>
26
       up[node][i] = up[up[node][i - 1]][i - 1];
27
28
29
30
      for(int i = Head[node]; i; i = Next[i]) if(To[i] != Par[node]) {
31
          Par[To[i]] = node;
          DFS(To[i], depth + 1);
33
34
36
   int KthAncestor(int u, int k)
37
      if(k > Level[u]) return -1;
39
      for(int i = lastBit(k); k; k -= lastBit(k), i = lastBit(k))
40
41
       u = up[u][Log[i]];
42
43
      return u:
44
45
46
    int LCA(int u, int v)
47
      if(Level[u] < Level[v]) swap(u, v);</pre>
49
      int k = Level[u] - Level[v];
50
51
      u = KthAncestor(u, k);
52
      if(u == v) return u;
53
54
      for(int i = LOG; i >= 0; --i)
55
        if(up[u][i] ^ up[v][i])
56
57
     u = up[u][i];
     v = up[v][i];
      return up[u][0];
63
65
      logCalc();
      for(int i = 1; i <= n; ++i) if(Par[i] == 0) {</pre>
67
         Par[i] = i;
68
69
          DFS(i);
       }
70
71
      cin >> q;
72
      while (q--)
73
74
75
76
          cout << LCA(u, v) << endl;</pre>
77
78
```

# 2.23 Lowest common ancestor (euler tour)

```
1 int Head[N], To[M], Next[M], Cost[M], ne, n, m, u, v, w, q;
2 int Last[N], First[N], euler_tour[N << 1], Height[N << 1], euler_timer;
3
4 void EulerianTour(int node, int depth = 0)
5 {
6 euler_tour[++euler_timer] = node;</pre>
```

```
Height[euler_timer] = depth;
     First[node] = euler_timer;
      for(int i = Head[node]; i; i = Next[i])
       if(First[To[i]] == 0)
11
^{12}
13
      EulerianTour(To[i], depth + Cost[i]);
14
15
      euler_tour[++euler_timer] = node;
16
     Height[euler_timer] = depth;
17
18
19
     Last[node] = euler_timer;
20 }
21
   int main()
25
      SparseTable <int> st(Height + 1, Height + euler_timer + 1, [&] (int a, int b) { return a <= b; });
27
28
      while (q--)
29
30
         cin >> 1 >> r:
31
         int left = Last[1];
32
33
          int right = Last[r]:
         if(left > right) swap(left, right);
34
35
36
          cout << euler_tour[ st.query(left, right) ] << endl;</pre>
37
38 }
```

#### 2.24 Minimum vertex cover

```
1 bool DFS(int node, int par = -1) {
2
3 bool black = false;
4 for(int e = Head[node]; e; e = Next[e])
5 if(To[e] != par)
6 black |= DFS(To[e], node);
7
8 MVC += black;
9 return !black;
10 }
```

# 2.25 Shortest path faster algorithml (SPFA)

```
1 /** Shortest Path Faster Algorithm :
       - This algorithm runs in O(kE) where k is a number depending on the graph.
        - The maximum k can be V (which is the same as the time complexity of Bellman Fords).
        - However in practice SPFA (which uses a queue) is as fast as Dijkstras (which uses a priority
             queue).
        - SPFA can deal with negative weight edge. If the graph has no negative cycle, SPFA runs well on
        - If the graph has negative cycle(s), SPFA can also detect it as there must be some vertex (those
              on the negative cycle)
       that enters the queue for over V dash 1 times.
10 int Head[N], Par[N], Next[M], To[M], Cost[M], Cnt[N], ne, n, m, u, v, st, tax;
11 ll dis[N];
12 bool Inq[N];
13
14 void set() {
     memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
     memset(Par, -1, sizeof(Par[0]) * (n + 2));
     memset (Cnt,
                   0, sizeof(Cnt[0]) * (n + 2));
     memset (Inq,
                   0, sizeof(Inq[0]) * (n + 2));
19
20
21
   bool SPFA(int src)
22
23
     _set();
24
25
     deque <int> 0:
26
     Q.push_front (src);
27
28
     dis[src] = 0;
     Cnt[src] = 1;
29
     Inq[src] = 1;
```

```
int node;
      while(Q.size()) {
        node = Q.front(); Q.pop_front(); Inq[node] = 0;
        for(int i = Head[node]; i; i = Next[i])
          if(dis[node] + Cost[i] < dis[To[i]]) {
      dis[To[i]] = dis[node] + Cost[i];
Par[To[i]] = node;
40
41
      if(!Inq[To[i]])
42
43
          if(++Cnt[To[i]] == n)
44
            return true; // graph has a negative weight cycle
45
          if(Q.size() && dis[To[i]] > dis[Q.front()])
46
            Q.push_back(To[i]);
49
            Q.push_front(To[i]);
51
          Inq[To[i]] = true;
53
55
      return false:
56
```

# 2.26 Single source shortest path

```
1 int Head[N], Par[N], Next[M], To[M], ne, n, m, u, v, st, tr;
 2 ll dis[N];
   void BFS(int src)
      memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
     memset(Par, -1, sizeof(Par[0]) * (n + 2));
     queue <int> Q;
10
     O.push (src):
11
     dis[src] = 0;
12
13
      while(Q.size()) {
       node = Q.front(); Q.pop();
      for(int i = Head[node]; i; i = Next[i])
         if (dis[To[i]] == INF) {
     dis[To[i]] = dis[node] + 1;
     Q.push(To[i]);
21
22
23
```

# 2.27 Single source shortest path (grid)

```
1 const int dr []
                        = {-1, 0, 1, 0};
 2 const int de []
   const int dc [] = {0, 1, 0, -1};
const char dir [] = {'U', 'R', 'D', 'L'};
   map <char, int> inv = { {'U', 0}, {'R', 1}, {'D', 2}, {'L', 3}};
   int dis[N][N], n, m;
   char Par[N][N];
   bool valid(int r, int c) {
     return r >= 1 && r <= n && c >= 1 && c <= m && dis[r][c] == oo;
10
11
13
    void BFS (int sr, int sc)
      memset(dis, 0x3f, sizeof(dis));
      memset(Par, -1, sizeof(Par));
18
      queue < pair <int, int> > Q;
      dis[sr][sc] = 0;
20
      Q.push({sr, sc});
21
      while (Q.size())
       tie(r, c) = Q.front(); Q.pop();
        for(int i = 0; i < 4; ++i) {
         nr = r + dr[i];
          nc = c + dc[i];
```

# 2.28 Tarjan (strongly connected components)

```
1 int Head[N], To[M], Next[M], Cost[M];
 2 int dfs_num[N], dfs_low[N];
 3 int Stack[N], compID[N], compSize[N];
 4 int ne, n, m, u, v, w;
 5 int dfs_timer, top, ID;
 6 bool in_stack[N];
    void _clear() {
      memset(Head, 0, sizeof(Head[0]) * (n + 2));
memset(dfs_num, 0, sizeof(dfs_num[0]) * (n + 2));
10
1.1
      memset(compID, 0, sizeof(compID[0]) * (n + 2));
     memset(compSize, 0, sizeof(compSize[0]) * (n + 2));
ne = dfs_timer = top = ID = 0;
12
13
14 }
15
    void Tarjan (int node)
16
17
      dfs_num[node] = dfs_low[node] = ++dfs_timer;
      in_stack[Stack[top++] = node] = true;
20
21
      for(int i = Head[node]; i; i = Next[i]) {
^{22}
        if(dfs_num[To[i]] == 0)
23
          Tarjan(To[i]);
24
25
        if (in_stack[To[i]])
          dfs_low[node] = Min(dfs_low[node], dfs_low[To[i]]);
26
27
28
      if(dfs_num[node] == dfs_low[node]) {
29
         for(int cur = -1; cur ^ node;) {
32
          in_stack[cur = Stack[--top]] = false;
33
          compID[cur] = ID;
34
          ++compSize[ID];
35
36
37 }
38
39 void Tarjan() {
     for (int i = 1; i <= n; ++i) if (dfs_num[i] == 0)
40
41
            Tarjan(i);
42
```

# 2.29 Tree diameter

# 2.30 Tree diameter (weighted tree)

```
1 void DFS(int node, long long cost, int par = -1) {
2    if(cost > diameter) diameter = cost, at = node;
3    for (int e = Head[node]; e; e = Next[e])
4    if(To[e] != par)
5    DFS(To[e], cost + Cost[e], node);
6 }
7
```

```
8 int main() {
9    DFS(1, 011);
10    from = at, diameter = 0;
11    DFS(from, 011);
12    to = at;
13    cout << diameter << endl;
14 }</pre>
```

# 2.31 Topological sort (kahns algorithm)

```
1 vector <int> kahn(int n)
 2
      vector <int> ready, ret;
      for(int i = 1; i <= n; ++i)</pre>
       if(!in[i])
         ready.push_back(i);
10
      while(!ready.empty())
11
         node = ready.back(); ready.pop_back();
12
         ret.push_back(node);
14
          for(int i = Head[node]; i; i = Next[i])
      if(--in[To[i]] == 0)
       ready.push_back(To[i]);
18
     return ret:
20
21
22 int main() {
      vector <int> v = kahn(n);
      if((int)v.size() == n)
       for(int i : v)
         cout << i << ' ';
        cout << "not a DAG!" << endl;</pre>
29
```

#### 2.32 Problems

#### 2.32.1 Restoring the path

```
1 const int dr []
                        = \{-1, 0, 1, 0\};
2 const int dc [] = {0, 1, 0, -1};
3 const char dir [] = {'U', 'R', 'D', 'L'};
 4 map <char, int> inv = { {'U', 0}, {'R', 1}, {'D', 2}, {'L', 3}};
 6 /** Implicit Graphs
        - in BFS, Dijkstra or Bellman-Ford function write -> Par[nr][nc] = dir[i ^ 2]
        - char Par[N][N] initialize with -1
        - si start i
        - sj strat j
       - fi target :
14
       - char dir and its map inv
16
17
18 string restorePath(int si, int sj, int fi, int fj) {
      string s:
      if (Par[ei][ej] == -1) return s;
20
      for(char i = Par[fi][fj]; (si ^ fi) || (sj ^ fj); i = Par[fi][fj]) {
       s += dir[inv[i] ^ 2];
27
      reverse(s.begin(), s.end());
30
      return s:
31
    /** Explicit Graphs (BFS, Dijkstra or Bellman-Ford)
33
34
35
        - int Par[N] initialize with -1
        - 11 dis[N] initialize with 0x3f
        - 11 INF = 0 \times 3f3f3f3f3f3f3f3f3f
```

```
38
   vector <int> restorePath(int dest) {
     vector <int> path;
     if(dis[dest] == INF) return path;
44
     for(int i = dest; ~i; i = Par[i])
       path.push_back(i);
45
46
47
     reverse(path.begin(), path.end());
48
     return path;
49
50
51
   /** in case of Floyd-Warshall:
52
       - 11 dis[N][N] initialize with 0x3f
       - 11 INF = 0 \times 3f3f3f3f3f3f3f3f3f3f
55
       - int Par[N][N] initialize with
       - in Floyd-Warshall function write -> Par[i][j] = Par[k][j];
57
58
59 vector <int> restorePath(int st, int tr) {
60
     vector <int> path;
61
     if(dis[st][tr] == INF) return path;
     for(int i = tr; st ^ i; i = Par[st][i])
63
       path.push_back(i);
64
65
66
     path.push back(st);
67
     reverse(path.begin(), path.end());
     return path:
```

#### 2.32.2 Tree distances

```
1 int Head[N], Next[M], To[M], Par[N], ne, n, m, u, v, diameter, At, From;
   int E[N << 1], H[N << 1], F[N], L[N], timer, SP[N << 1][LOG + 1], Log[N << 1];
 5
   void clear() {
     memset(Head, 0, sizeof(Head[0]) * (n + 2));
     memset(Par, -1, sizeof(Par[0]) * (n + 2));
     timer = 0;
10
12 void EulerTour(int node, int depth = 0, int par = -1) {
14
     H[timer] = depth;
15
     F[node] = timer;
16
     for(int i = Head[node]; i; i = Next[i])
17
18
       if(To[i] != par) {
19
         EulerTour(To[i], depth + 1, node);
          E[++timer] = node;
20
21
         H[timer] = depth;
22
23
     L[node] = timer;
    void dfs(int node, int depth = 0, int par = -1) {
      if (depth > diameter) diameter = depth, At = node;
29
     for(int i = Head[node]; i; i = Next[i])
if(To[i] != par)
30
         dfs(To[i], depth + 1, node);
31
32
33
34
   void bulid()
35
37
     dfs(1); From = At; diameter = 0; dfs(From);
39
40
      for(int i = 2; i <= (n << 1); ++i)
41
       Log[i] = Log[i >> 1] + 1;
42
      for(int i = 1; i < (n << 1); ++i)</pre>
44
       SP[i][0] = i;
45
46
     int MaxLog = Log[(n << 1)];</pre>
     for(int j = 1, k, h; j <= MaxLog; ++j) {
  k = (1 << j);</pre>
47
48
        h = (k >> 1);
49
        for (int i = 1; i + k - 1 < (n << 1); ++i)
50
51
     const int & x = SP[i][j - 1];
```

```
const int & y = SP[i + h][j - 1];
54
55
      SP[i][j] = H[x] <= H[y] ? x : y;
57
58
59
60
    int query(int 1, int r)
61
      int d = r - 1 + 1:
      int lg = Log[d];
int k = (1 << lg);</pre>
63
64
65
      const int & x = SP[1][1g];
const int & y = SP[1 + d - k][1g];
66
67
      return (H[x] <= H[y] ? x : y);
70
72
    int LCA(int u, int v) {
      return query (u, v);
74
75
76
   int distance(int u, int v) {
      int 1 = F[u];
      int r = F[v]:
78
79
      if(1 > r) swap(1, r);
80
81
      int ix = LCA(1, r);
      return (H[1] + H[r] - H[ix] - H[ix]);
83
85
86
87
      for(int i = 1; i <= n; ++i)</pre>
89
        cout << max(distance(i, At), distance(i, From)) << " \n"[i == n];</pre>
90
```

#### 2.32.3 Subtree sizes

```
1 int Head[N], Next[M], To[M], Par[N], sbtree_size[N], ne, n, m, u, v, w;
2
3 void dfs(int node, int par = -1) {
4     sbtree_size[node] = 1;
5     for(int i = Head[node]; i; i = Next[i])
6     if(To[i] != par) {
7         dfs(To[i], node);
8         sbtree_size[node] += sbtree_size[To[i]];
9     }
10 }
11
2 int main()
13 {
14     dfs(1);
15     for(int i = 1; i <= n; ++i)
16     cout << sbtree_size[i] - 1 << " \n"[i == n];
17 }</pre>
```

#### 2.32.4 Shortest cycle

```
1 /** for each node run BFS and minmize the cycle length
2
   int BFS (int src)
      memset(dis, 0x3f, sizeof(dis[0]) * (n + 2));
      memset(Par, -1, sizeof(Par[0]) * (n + 2));
10
      Q.push(src);
11
      dis[src] = 0;
12
13
      int node, ret = oo;
14
      while(Q.size())
15
16
          node = Q.front(); Q.pop();
          for(int i = Head[node]; i; i = Next[i])
17
18
19
        if (dis[To[i]] != oo) {
         if(Par[node] != To[i]) {
   if(dis[node] + 1 + dis[To[i]] < ret)</pre>
20
^{21}
        ret = dis[node] + 1 + dis[To[i]];
```

```
23
^{24}
          continue;
25
26
27
        dis[To[i]] = dis[node] + 1;
28
        Par[To[i]] = node;
29
        Q.push(To[i]);
30
31
32
     return ret;
33
```

#### 2.32.5 Path sum for each node

```
1 int Head[N], Next[M], To[M], ne, u, v, n, m, subtree_size[N], level[N];
2 11 dis[N];
   void dfs(int node, int par = -1) {
      subtree_size[node] = 1;
      for(int i = Head[node]; i; i = Next[i]) if(To[i] != par) {
          level[To[i]] = level[node] + 1;
          dfs(To[i], node);
          subtree_size[node] += subtree_size[To[i]];
10
11 }
12
13 void reRoot(int node, ll pd, int par = -1) {
14
    dis[node] = pd;
for(int i = Head[node]; i, i = Next[i]) if(To[i] != par) {
          reRoot(To[i], pd - subtree_size[To[i]] + (n - subtree_size[To[i]]), node);
16
17
18 }
20
   void get_dis()
21
23
      11 pd = 0;
24
      for(int i = 1; i <= n; ++i)
25
       pd += level[i];
26
     reRoot(1, pd);
27
     for(int i = 1; i <= n; ++i)
  cout << dis[i] << " \n"[i == n];</pre>
28
29
30 }
```

# 3 Data structures

# 3.1 Segment tree

```
template <class T, class F = function <T(const T &, const T &)> >
   class SegmentTree
      vector <T> _A;
      vector <T> ST;
      vector <T> LT;
     int _N;
10 public :
     template <class iter>
11
     SegmentTree(iter _begin, iter _end, const F _func = [](T a, T b) {return a <= b ? a : b;}) : func(
12
            func)
13
14
       _N = distance(_begin, _end);
        _N = (1 << (int)ceil(log2(_N)));
15
        _A.assign(_N + 1, 0);
        ST.assign(_N << 1, 0);
19
        LT.assign(_N << 1, 0);
20
21
          _typeof(_begin) i = _begin;
22
        for(int j = 1; i != _end; ++i, ++j)
   _A[j] = *i;
23
24
25
        build(1, 1, _N);
26
27
      void build(int p, int 1, int r)
```

```
30
         if(1 == r) {
31
           ST[p] = A[1];
           return;
 34
         int mid = (1 + r) >> 1;
36
         build(p + p, l, mid);
37
         build(p + p + 1, mid + 1, r);
38
 39
         const T & x = ST[p + p];
const T & y = ST[p + p + 1];
 40
 41
 42
         ST[p] = func(x, y);
 43
 44
       void update_range(int ul, int ur, int delta) {
 47
         update_range(ul, ur, delta, 1, 1, _N);
 48
 49
 50
       T query(int ql, int qr) {
51
         return query(ql, qr, 1, 1, _N);
 52
53
       void update_point(int inx, int delta)
 54
 55
 56
         inx += N - 1:
         ST[inx] = delta;
 57
 58
         while (inx > 1) {
           const T & x = ST[inx + inx];
           const T & y = ST[inx + inx + 1];
 64
 65
           ST[inx] = func(x, y);
66
67
68
69 private :
       void update_range(int ul, int ur, int delta, int p, int l, int r)
70
 72
         if(r < ul || ur < 1)
 73
           return:
 74
 75
         if(ul <= 1 && r <= ur) {
          ST[p] += delta;
           LT[p] += delta;
           return;
 79
 80
81
         propagate(p);
         int mid = (1 + r) >> 1;
 83
 84
         update_range(ul, ur, delta, p + p, l, mid);
update_range(ul, ur, delta, p + p + 1, mid + 1, r);
 85
 86
 87
         const T & x = ST[p + p];
         const T & y = ST[p + p + 1];
 90
 91
         ST[p] = func(x, y);
92
 93
94
       T query(int ql, int qr, int p, int l, int r)
95
96
         if(r < ql || qr < 1)
97
           return INT MAX;
98
 99
         if (gl <= 1 && r <= gr)
100
          return ST[p]:
101
102
         propagate(p);
103
         int mid = (1 + r) >> 1;
         const T & x = query(ql, qr, p + p, l, mid);
107
         const T & y = query(q1, qr, p + p + 1, mid + 1, r);
108
109
         return func(x, y);
110
111
       void propagate(int p) {
112
         if(LT[p]) {
113
                         += LT[p];
114
           ST[p + p]
115
           ST[p + p + 1] += LT[p];
           LT[p + p] += LT[p];
LT[p + p + 1] += LT[p];
           LT[p] = 0;
120
```

# 3.2 Merge sort tree

```
1 class SegmentTree
 2
     vector <vector <int> > sTree;
      vector <int> localArr;
     int NP2, oo = 0x3f3f3f3f3f;
     template <class T>
     SegmentTree(T begin, T end)
10
11
        int n = _end - _begin;
12
        while (NP2 < n) NP2 <<= 1;
13
        sTree.assign(NP2 << 1, vector <int> ());
16
        localArr.assign(NP2 + 1, 0);
17
18
        __typeof(_begin) i = _begin;
for(int j = 1; i != _end; i++, ++j)
19
20
         localArr[j] = *i;
21
       build(1, 1, NP2);
22
23
24
      void build(int p, int 1, int r)
25
26
27
28
        sTree[p].push_back(localArr[l]);
29
         return;
30
31
32
        build(left(p), l, mid(l, r));
33
       build(right(p), mid(l, r) + 1, r);
34
35
36
37
38
     int query(int ql, int qr, int k) {
      return query(q1, qr, k, 1, 1, NP2);
39
40
41
42 private :
     int query(int ql, int qr, int k, int p, int l, int r)
        if(isOutside(ql, qr, l, r))
46
47
48
        if(isInside(ql, qr, l, r)) {
         return sTree[p].end() - upper_bound(sTree[p].begin(), sTree[p].end(), k);
49
50
51
        return query(ql, qr, k, left(p), l, mid(l, r)) +
52
         query(q1, qr, k, right(p), mid(1, r) + 1, r);
53
54
55
56
      void merge(int p)
57
        vector <int> & L = sTree[left(p)];
59
        vector <int> & R = sTree[right(p)];
60
61
        int l_size = L.size();
62
        int r_size = R.size();
63
        int p_size = l_size + r_size;
        I. push back (INT MAX) :
65
        R.push_back(INT_MAX);
66
67
68
        sTree[p].resize(p_size);
69
        for(int k = 0, i = 0, j = 0; k < p_size; ++k)</pre>
70
        if(L[i] <= R[j])
     sTree[p][k] = L[i], i += (L[i] != INT_MAX);
      sTree[p][k] = R[j], j += (R[j] != INT_MAX);
75
76
        L.pop_back();
77
       R.pop_back();
78
79
80
     inline bool isInside(int ql, int qr, int sl, int sr) {
81
       return (gl <= sl && sr <= gr);
82
83
     inline bool isOutside(int ql, int qr, int sl, int sr) {
       return (sr < ql || qr < sl);
```

```
87
88    inline int mid (int 1, int r) {
89        return ((1 + r) >> 1);
90    }
91
92    inline int left(int p) {
93        return (p << 1);
94    }
95
    inline int right(int p) {
96        inline int right(int p) {
97        return ((p << 1) | 1);
98    }
99    };</pre>
```

#### 3.3 Merge sort

```
1 ll inversions;
    template <class T>
    void merge(T localArr [], int l, int mid, int r)
      int 1 size = mid - 1 + 1;
     int r size = r - mid:
      T L[1_size + 1];
10
     T R[r_size + 1];
      for(int i = 0; i < l_size; ++i) L[i] = localArr[i + 1];</pre>
      for(int i = 0; i < r_size; ++i) R[i] = localArr[i + mid + 1];</pre>
14
15
16
      if(sizeof(T) == 4) Mx = INT_MAX;
17
      else Mx = LONG_MAX;
18
19
      L[l_size] = R[r_size] = Mx;
20
      for(int k = 1, i = 0, j = 0; k <= r; ++k)
if(L[i] <= R[j])</pre>
21
23
          localArr[k] = L[i], i += (L[i] != Mx);
         localArr[k] = R[j], j += (R[j] != Mx), inversions += l_size - i;
29
    void merge_sort(T localArr [], int l, int r)
30
31
32
33
          int mid = (1 + r) >> 1;
34
          merge_sort(localArr, 1,
35
          merge_sort(localArr, mid + 1, r);
36
         merge(localArr,

    mid, r);

37
    template <class T>
    void merge_sort(T _begin, T _end)
42
      const int sz = _end - _begin;
44
     __typeof(*_begin) localArray[sz];
45
46
        _typeof(_begin) k = _begin;
47
      for(int i = 0; k != _end; ++i, ++k)
48
       localArray[i] = *k;
49
50
     merge_sort(localArray, 0, sz - 1);
51
52
      for(int i = 0; k != _end; ++i, ++k)
53
        *k = localArray[i];
```

# 3.4 Sparse table

```
1 template <class T, class F = function <T(const T&, const T&)>>
2 class SparseTable
3 {
4   int _N;
5   int _Log;
6   vector <T> _A;
7   vector < vector <T> > ST;
8   vector <int> Log;
```

```
F func;
10
     SparseTable() = default;
15
      SparseTable(iter _begin, iter _end, const F _func = less <T> ()) : func(_func)
16
17
        _N = distance(_begin, _end);
18
19
        Log.assign(N + 1, 0);
20
        for (int i = 2; i <= _N; ++i)</pre>
21
          Log[i] = Log[i >> 1] + 1;
22
        LOG = Log[ N1:
23
25
         _A.assign(_N + 1, 0);
26
        ST.assign(_N + 1, vector <T> (_LOG + 1, 0));
27
28
          _typeof(_begin) i = _begin;
29
        for(int j = 1; i != _end; ++i, ++j)
30
          _{A[j]} = *i;
31
32
        build();
33
34
      void build()
35
36
37
        for (int i = 1; i <= _N; ++i)</pre>
          ST[i][0] = i;
38
39
40
        for(int j = 1, k, d; j <= _LOG; ++j)  // the two nested loops below have overall time complexity</pre>
41
42
      k = (1 << j);
43
     d = (k >> 1);
44
45
      for(int i = 1; i + k - 1 <= _N; ++i)</pre>
46
          T const & x = ST[i][j - 1]; // starting subarray at index = i with length = 2^{j-1} T const & y = ST[i + d][j - 1]; // starting subarray at index = i + d with length = 2^{j-1}
47
48
49
50
          ST[i][j] = func(A[x], A[y]) ? x : y;
51
52
53
54
55
      T query(int 1, int r) // this query is O(1)
56
57
58
        T const & x = ST[1][Log[d]];
59
        T const & y = ST[1 + d - (1 << Log[d])][Log[d]];
60
61
        return func(_A[x], _A[y]) ? x : y;
62
63 };
```

# 3.5 Union find disjoint sets

```
Maintain a set of elements partitioned into non-overlapping subsets. Each
      partition is assigned a unique representative known as the parent, or root. The
       following implements two well-known optimizations known as union-by-size and
      path compression. This version is simplified to only work on integer elements.
        find\_set(u) returns the unique representative of the partition containing u.
       same\_set(u, v) returns whether elements u and v belong to the same partition.
       union set(u, v) replaces the partitions containing u and v with a single new
      partition consisting of the union of elements in the original partitions.
10
11
13
       O(a(n)) per call to find_set(), same_set(), and union_set(), where n is the
       number of elements, and a(n) is the extremely slow growing inverse of the Ackermann function
       (effectively a very small constant for all practical values of n).
16
17
18
       O(n) for storage of the disjoint set forest elements.
19
       O(1) auxiliary for all operations.
20
21
22 class UnionFind
23
24
     vector <int> par:
      vector <int> siz;
25
     int num sets;
26
     size t sz;
```

```
UnionFind() : par(1, -1), siz(1, 1), num_sets(0), sz(0) {}
      UnionFind(int n): par(n + 1, -1), siz(n + 1, 1), num_sets(n), sz(n) {}
33
34
35
        assert(u <= sz);
36
37
38
        for(leader = u; ~par[leader]; leader = par[leader]);
39
40
        for(int next = par[u]; u != leader; next = par[next]) {
41
         par[u] = leader;
42
         u = next;
43
        return leader;
46
47
      bool same_set(int u, int v) {
48
        return find_set(u) == find_set(v);
49
50
51
      bool union_set(int u, int v) {
52
       if(same_set(u, v)) return false;
53
54
        int x = find set(u);
55
       int y = find_set(v);
56
57
        if(siz[x] < siz[y]) swap(x, y);</pre>
59
        par[v] = x:
       siz[x] += siz[y];
63
        return true;
64
65
66
      int number_of_sets() {
67
       return num_sets;
68
69
70
      int size of set(int u) {
71
       return siz[find_set(u)];
72
73
      size_t size() {
75
       return sz;
      void clear() {
79
       par.clear();
80
        siz.clear();
       sz = num_sets = 0;
81
82
83
84
      void assign(size t n) {
85
       par.assign(n + 1, -1);
       siz.assign(n + 1, 1);
87
       sz = num_sets = n;
89
      map < int, vector <int> > groups(int st) {
91
       map < int, vector <int> > ret;
93
        for(size_t i = st; i < sz + st; ++i)</pre>
         ret[find_set(i)].push_back(i);
95
96
        return ret;
97
98 };
```

#### 3.6 Bubble sort

```
1  /**
2  Bubble sort consists of n rounds. On each round, the algorithm iterates
3  through the elements of the array. Whenever two consecutive elements are found
4  that are not in correct order, the algorithm swaps them. The algorithm can be
5  implemented as follows:
6  **/
7  template <class T>
9  void bubble_sort(T_begin, T_end, int round) {
10  const int sz = _end - _begin;
11  int localArray[sz];
12
13  _typeof(_begin) k = _begin;
14  for(int j = 0; k != _end; ++k, ++j)
```

```
15
       localArray[j] = *k;
16
17
     round = min(round, sz);
     for(int i = 0; i < round; ++i) /* n rounds -> n_th element **/
       for(int j = 0; j < sz - 1; ++j) if(localArray[j] > localArray[j + 1])
20
                 swap(localArray[j], localArray[j + 1]);
^{21}
22
     for(int j = 0; k != _end; ++k, ++j)
23
24
       *k = localArray[j];
25
26
27
     After the first round of the algorithm, the largest element will be in the correct
28
      position, and in general, after k rounds, the k largest elements will be in the
      correct positions.
```

# 4 Mathematics

# 4.1 Pisano periodic sequence

```
template <class T>
    using matrix = vector < vector <T> >;
    template <class T> string to_string(T x) {
      int sn = 1;
     if (x < 0) sn = -1, x *= sn;
      string s = "";
       s = "0123456789"[x % 10] + s, x /= 10;
10
      } while(x);
      return (sn == -1 ? "-" : "") + s;
1.1
12 3
13
14 auto str_to_int(string x) {
15    uil28 ret = (x[0] == '-' ? 0 : x[0] = '0');
16    for(int i = 1; i < (int)x.size(); ++i) ret = ret * 10 + (x[i] = '0');</pre>
     return (x[0] == '-' ? -1 * (i128)ret : ret);
17
18 }
20 istream & operator >> (istream & in, i128 & i) noexcept {
      in >> s;
23
24
      return in;
25
26
27 ostream & operator << (ostream & os, const i128 i) noexcept {
28
     os << to string(i);
29
     return os;
30 }
31
32 void Fast() {
     cin.sync_with_stdio(0);
      cin.tie(0);
      cout.tie(0);
36
37
38 ll n;
39 vector <int> primes;
40 matrix <11> fibMatrix = \{\{1, 1\},
41
           {1.0}
42 };
43
44 i128 gcd(i128 a, i128 b) {
45
     while (a && b)
        a > b ? a %= b : b %= a;
     return a + b;
50 i128 lcm(i128 a, i128 b) {
     return a / gcd(a, b) * b;
52
54 vector < array <11, 2> > factorize(11 x) {
      vector < array <11, 2> > ret;
56
      for(int i = 0; 111 * primes[i] * primes[i] <= x; ++i) {</pre>
57
        if(x % primes[i]) continue;
58
59
        int cnt = 0;
        while (x % primes[i] == 0) {
60
61
          cnt++;
          x /= primes[i];
```

```
ret.push_back({primes[i], cnt});
 65
67
       if(x > 1) ret.push_back({x, 1});
69
 70
71
    matrix <11> MatMul (matrix <11> A, matrix <11> B, 11 mod) {
      int ra = A.size(), cb = B[0].size(), ca = A[0].size();
73
       matrix <i128> C(ra, vector <i128> (cb));
74
75
       for(int i = 0; i < ra; ++i)</pre>
         for (int j = 0; j < cb; ++j) {
    C[i][j] = 0;
    for(int k = 0; k < ca; ++k)
76
 77
78
       C[i][j] = (C[i][j] + (i128)A[i][k] * B[k][j]);
 82
       matrix <11> ret(ra, vector <11> (cb));
 83
       for(int i = 0; i < ra; ++i)</pre>
         for (int j = 0; j < cb; ++j)
  ret[i][j] = C[i][j] % mod;</pre>
 84
 85
 86
 87
       return ret:
 88
 89
90 matrix <11> MatPow(matrix <11> A, 11 p, 11 mod) {
91
       int r = A.size(), c = A[0].size();
       assert (r == c && p);
       matrix <11> result = A;
 95
97
         if(p & 111) result = MatMul(result, A, mod);
98
         A = MatMul(A, A, mod);
aa
         p >>= 111;
100
101
       return result;
102
103
104 i128 ModExp(i128 a, ll p) {
       i128 result = 1;
105
106
       while (p)
        if (p & 111) result = result * a;
107
108
109
        p >>= 111;
111
       return result;
112
113
114 ll nthFib(ll n, ll mod) {
       return MatPow(fibMatrix, n, mod)[0][1];
115
116
117
118 bool is_period(ll n, ll mod) {
       return nthFib(n, mod) == 0 && nthFib(n + 1, mod) == 1;
119
120
121
122 ll solver(ll x, ll mod) {
       vector < array <11, 2> > factors = factorize(x);
       for(int i = 0; i < (int) factors.size(); ++i) {</pre>
125
         while(x % factors[i][0] == 0 && is_period(x / factors[i][0], mod))
126
          x /= factors[i][0];
127
128
       return x;
129
130
131 ll pisano_prime(ll val) {
       if(val == 2) return 3;
132
       if (val == 5) return 20;
133
       if (val % 10 == 1 || val % 10 == 9)
134
         return solver(val - 1, val);
135
136
137
       return solver(2 * (val + 1), val);
139
    const int N = 1e7 + 9;
140
141 bitset <N> isPrime;
143 void Precomputation_Sieve() {
144
       isPrime.set();
       int _sqrt = sqrtl(N);
145
146
       for(int i = 5; i <= _sqrt; i += 6) {</pre>
147
        if(isPrime[i]) for (int j = i * i; j < N; j += i + i) isPrime.reset(j);</pre>
148
149
         if(isPrime[i]) for (int j = i * i; j < N; j += i + i) isPrime.reset(j);</pre>
152
153
```

```
155 vector <int> Primes(int n) {
      vector <int> _Primes;
      if(n >= 2) _Primes.push_back(2);
159
      if(n >= 3) _Primes.push_back(3);
161
      for (int i = 5; i <= n; i += 6) {
162
       if(isPrime[i]) _Primes.push_back(i);
163
164
        if(isPrime[i]) _Primes.push_back(i);
165
        i -= 2;
166
167
      return _Primes;
168
169
    void initialize()
      primes = Primes(N);
174
175
176 void Solve() {
177
      initialize():
178
179
      cin >> n:
      vector < array <11, 2> > factors = factorize(n);
180
181
182
183
      for (int i = 0; i < (int) factors.size(); ++i) {</pre>
        ans = lcm(ans, (i128)pisano_prime(factors[i][0]) * ModExp(factors[i][0], factors[i][1] - 1));
184
185
      cout << ans << endl;</pre>
187
```

#### 4.2 Euler totient function

```
1 int lp[N], Primes[664580], pnx; /** size of Primes = n / (ln(n) - 1.08) */
 3
   void linear sieve (int n) {
     for (int i = 2; i <= n; ++i) {
       if (lp[i] == 0) {
         lp[i] = Primes[pnx++] = i;
       for (int j = 0, comp; j < pnx && Primes[j] <= lp[i] && (comp = i * Primes[j]) <= n; ++j) {
         lp[comp] = Primes[j];
10
11
12 }
13
14 ll Phi(ll a) { // for Queries
1.5
     ll ret = a, p;
     for (int i = 0; i < pnx && (p = Primes[i], true); ++i) {</pre>
16
17
      if (p * p > a) break;
       if (a % p) continue;
18
       ret -= ret / p;
       while (a % p == 0) a /= p;
     if (a > 1) ret -= ret / a;
23
   return ret;
24
```

#### 4.3 Extended wheel factorization

```
1 /*
    Constraints:
 3
      1 <= n <= 1e7
     2 <= a <= 1e{14}
      linear sieve takes O(n)
      Factorization takes O(n / (ln(n) - 1.08))
10
     Space Complexity:
1.1
     O(MaxN + n / (ln(n) - 1.08)
12 */
13
14 int lp[N];
15 int Primes[664580], pnx; /** size of Primes = n / (ln(n) - 1.08) */
16
17 void linear sieve(int n) {
     for (int i = 2; i <= n; ++i) {
  if (lp[i] == 0) {</pre>
18
```

```
20
         lp[i] = Primes[pnx++] = i;
^{21}
        for (int j = 0, comp; j < pnx && Primes[j] <= lp[i] && (comp = i * Primes[j]) <= n; ++j) {
         lp[comp] = Primes[j];
26
28
   vector<pair<ll, int>> Factorization(ll a) {
      vector<pair<11, int> > ret;
30
31
      for (int i = 0, cnt; i < pnx && (p = Primes[i], true) && p * p <= a; ++i) {</pre>
32
       if (a % p) continue;
33
        cnt = 0;
        while (a % p == 0) a /= p, ++cnt;
34
       ret.emplace_back(p, cnt);
37
      if (a > 1) ret.emplace_back(a, 1);
39
```

#### 4.4 Least prime factorization

```
1 /*
     Constraints:
     1 <= n <= 1e7
     Time Complexity:
      linear_sieve takes O(n)
     Factorization takes O(log(n))
10
    O(MaxN + n / (ln(n) - 1.08)
11 */
12
13 int lp[N];
14 int Primes[664580], pnx; /** size of Primes = n / (ln(n) - 1.08) */
16
   void linear sieve(int n) {
     for (int i = 2; i <= n; ++i) {
17
18
        if (lp[i] == 0) {
19
         lp[i] = Primes[pnx++] = i;
        for (int j = 0, comp; j < pnx && Primes[j] <= lp[i] && (comp = i * Primes[j]) <= n; ++j) {
          lp[comp] = Primes[j];
^{24}
^{25}
26
27
   vector<pair<int, int>> Factorization(int n) {
      vector<pair<int, int>> ret;
      while (n > 1) {
       int p = leastPrime[n], cnt = 0;
while (n % p == 0) n /= p, ++cnt;
30
31
        ret.emplace_back(p, cnt);
      return ret;
35
```

# 4.5 Mobius function

```
1 /**
        Constraints:
        1 <= x <= 1e7
        2 <= n <= 10^{14}
        Time Complexity:
        linear_sieve takes O(x)
        mobius takes O(n / (ln(n) - 1.08))
10
11
        O(MaxN + n / (ln(n) - 1.08))
12 */
13
14 int lp[N], Primes[664580], pnx; /** size of Primes = n / (ln(n) - 1.08) */
15
   void linear_sieve(int x) {
  for (int i = 2; i <= x; ++i) {</pre>
17
        if (lp[i] == 0) {
18
          lp[i] = Primes[pnx++] = i;
19
20
```

```
for (int j = 0, comp; j < pnx && Primes[j] <= lp[i] && (comp = i * Primes[j]) <= x; ++j) {
23
         lp[comp] = Primes[j];
25
26
27
28 int mobius(ll n) {
29
30
31
     for (int i = 0; i < pnx && (p = Primes[i], pp = p * p, true); ++i) {
32
      if (pp > n) break;
33
       if (n % p) continue;
34
       if (n % pp == 0) return 0;
35
       n /= p;
36
       mob = -mob:
     if (n > 1) mob = -mob;
39
     return mob;
40 }
```

#### 4.6 Phi factorial

```
2
       Constraints:
 3
       1 <= x <= 1e7
       2 <= n <= 1e7
       Time Complexity:
       linear sieve takes O(x)
       phi_factorial takes O(n)
10
       Space Complexity:
11
       O(MaxN + n / (ln(n) - 1.08))
12 */
13
14 int lp[N], Primes[664580], pnx; /** number of primes = n / (ln(n) - 1.08) **/
   void linear_sieve(int x) {
17
     for (int i = 2; i \le x; ++i) {
18
       if (lp[i] == 0) {
19
         lp[i] = Primes[pnx++] = i;
20
21
       for (int j = 0, comp; j < pnx && Primes[j] <= lp[i] && (comp = i * Primes[j]) <= x; ++j) {
22
23
         lp[comp] = Primes[j];
24
25
26
27
28 ll phi_factorial(int n) {
    for (int i = 2; i <= n; ++i) {
31
       ret = ret * (lp[i] == i ? i - 1 : i);
32
33
     return ret;
34 1
```

#### 4.7 Linear sieve

```
1 <= n <= 1e7
       Time Complexity:
       linear_sieve takes O(n)
       Space Complexity:
       O(MaxN + n / (ln(n) - 1.08))
10 **/
12 int lp[N];
13 int Primes[664580], pnx; /** size of Primes = n / (ln(n) - 1.08) */
15 void linear_sieve(int n) {
16
     for (int i = 2; i <= n; ++i) {
17
       if (lp[i] == 0) {
18
         lp[i] = Primes[pnx++] = i;
19
       for (int j = 0, comp; j < pnx && Primes[j] <= lp[i] && (comp = i * Primes[j]) <= n; ++j) {
20
         lp[comp] = Primes[j];
21
22
23
```

# 4.8 Segmented sieve

```
2 int Primes[664580], pnx; /** size of Primes = n / (ln(n) - 1.08) */
 3 bool isPrime[N];
    void linear sieve(int n) {
 6
      for (int i = 2; i <= n; ++i) {
  if (lp[i] == 0) {</pre>
           lp[i] = Primes[pnx++] = i;
10
         for (int j = 0, comp; j < pnx && Primes[j] <= lp[i] && (comp = i * Primes[j]) <= n; ++j) {
           lp[comp] = Primes[j];
11
^{12}
13
14
16
    vector<ll> segmented_sieve(ll 1, ll r) {
17
      1 += 1 == 1;
18
      int limit = r - 1 + 1;
19
      vector<ll> ret:
       memset(isPrime, true, sizeof(isPrime));
20
21
22
       for (int i = 0; i < pnx && (p = Primes[i], true); ++i) {
  for (ll j = max(p * p, (l + p - l) / p * p); j <= r; j += p)
    isPrime[j - l] = false;</pre>
23
24
       for (int i = 0; i < limit; ++i)</pre>
29
30
           ret.emplace_back(i + 1);
31
32
```

#### 4.9 Miller Rabin test

```
1 ll ModExp(ll base, ll e, ll mod)
2 {
      base %= mod;
      for(result = 1; e; e >>= 111)
      result = ((i128) result * base) % mod;
         base = ((i128)base * base) % mod;
10
11
     return result;
12
13
14 bool CheckComposite(ll n, ll p, ll d, int r)
15
     11 a = ModExp(p, d, n);
if(a == 1 || a == n - 1)
16
17
18
        return false;
      for(int i = 1; i < r; ++i)</pre>
^{21}
          a = ((i128)a * a) % n;
          if(a == n - 1)
23
      return false;
      return true:
26
   bool Miller(ll n)
      if(n < 2) return false;</pre>
      for(r = 0, d = n - 1; (d & 111) == 0; d >>= 111, ++r);
34
35
      for(int p: {2, 3, 7, 11, 13, 17, 19, 23, 29, 31, 37})
36
          if(n == p)
37
38
      return true:
39
         if(CheckComposite(n, p, d, r))
40
      return false;
41
      return true:
43
   int main()
```

```
17   cout << (Miller(n) ? "Yes, it is Prime" : "No, it is not a prime") << endl;
18 }
```

# 4.10 Stable marriage problem

```
1 const int N = 1e3+ 9, M = 1e3 + 9, oo = 0x3f3f3f3f3f;
 2 queue <int> 0;
    \textbf{int} \ \text{husband[N], wife[N], Next[N], order[N][N], pref[N][N], n, v;}
    void clear() {
     memset(wife, 0, sizeof(wife[0]) * (n + 2));
      memset(husband, 0, sizeof(husband[0]) * (n + 2));
     memset(Next, 0, sizeof(Next[0]) * (n + 2));
    void engage(int man, int woman) {
     int exWife = wife[man];
14
     wife[man] = woman;
15
     husband[woman] = man;
     if (exWife)
17
       Q.push(exWife);
18
19
20
    void Solve()
      _clear();
      for(int i = 1; i <= n; ++i)</pre>
27
        for(int j = 1; j <= n; ++j)</pre>
28
         cin >> pref[i][j];
29
30
     for(int i = 1; i <= n; ++i)</pre>
31
       for(int j = 1; j <= n; ++j) {
32
         cin >> v:
33
         order[i][v] = j;
34
35
36
     for(int i = 1; i <= n; ++i)</pre>
37
       Q.push(i);
39
40
      while (Q.size())
41
42
          woman = Q.front(); Q.pop();
43
         man = pref[woman][++Next[woman]];
44
45
         if(!wife[man] || order[man][woman] < order[man][wife[man]])</pre>
46
     engage(man, woman);
47
          else
48
     Q.push (woman);
49
      for(int i = 1; i <= n; ++i)</pre>
        cout << husband[i] << endl;</pre>
53
54
55
56
57
      3 2 1
59
      2 1 3
60
61
62
63
      2 3 1
64
65
66
```

# 4.11 Euler totient sieve

```
Space Complexity:
       MaxN
10
    int EulerPhi[N];
13
14
    void Phi_sieve(int n) {
15
       for (int i = 1; i <= n; ++i) {
16
          EulerPhi[i] = i;
17
18
       for (int i = 2; i <= n; ++i) {
   if (EulerPhi[i] == i)</pre>
19
       for (int j = i; j <= n; j += i) {
EulerPhi[j] -= EulerPhi[j] / i;</pre>
20
21
```

#### 4.12 Mobius sieve

```
1 /*
      Constraints:
      1 <= n <= 1e7
      Time Complexity:
      mu_sieve takes O(n)
      Space Complexity:
10
11
12 int mu[N], lp[N], Primes[78522], pnx;
13
14
    void mu_sieve(int n) {
15
      mu[1] = 1;
      fill(mu, mu + N, 1);

for (int i = 2; i <= n; ++i) {

   if (lp[i] == 0) {

      lp[i] = Primes[pnx++] = i;
17
18
19
20
           mu[i] = -1;
^{21}
         for (int j = 0, nxt; j < pnx && Primes[j] <= lp[i] && (nxt = i * Primes[j]) <= n; ++j) {
           lp[nxt] = Primes[j];
           mu[nxt] = (lp[i] == Primes[j] ? 0 : -mu[i]);
26
```

#### 4.13 Wheel sieve

```
1 /**
       Constraints:
       1 <= n <= 1e9
       2 <= x <= 9700000
       Time Complexity:
       wheel_sieve takes O(n / ln(ln(n)))
       coPrimes takes O(x * ln(ln(x)))
10
       Space Complexity:
11
       O(MaxN / 32 + n / (1n(n) - 1.08) + x)
12 **/
13
14 bitset<N> isPrime;
15
16 int Primes[50908031], pnx; /** size of Primes = n / (ln(n) - 1.08) */
19
     int basis[5] = {3, 5, 7, 11, 13};
20
21
     bitset<30100> isCoprime;
23
      isCoprime.set();
24
25
      for (int b : basis)
       for (int d = b * b; d <= x; d += b << 1)
26
27
         isCoprime.reset(d);
28
      for (int i = 17; i <= x; i += 2)
       if (isCoprime[i]) ret.push_back(i);
```

```
ret.push_back(x + 1);
     ret.push_back(x + 17);
     return ret;
35
37 void wheel_sieve(int n) {
     int basis[6] = {2, 3, 5, 7, 11, 13};
39
     vector<int> wheel = coPrimes(2 * 3 * 5 * 7 * 11 * 13);
40
     int sz = wheel.size();
41
42
     for (int k = 0; k < sz; ++k)
43
       inx[wheel[k]] = k;
44
45
     isPrime.set():
     inx[1] = sz - 2
46
     int inc[sz - 1];
49
     for (int i = 1; i < sz; ++i)</pre>
50
       inc[i - 1] = wheel[i] - wheel[i - 1];
51
52
      for (int p : basis) {
53
       if (n >= p)
         Primes[pnx++] = p;
54
55
56
57
     int c = 0;
     for (11 i = 17; i <= n; i += inc[c++]) {
58
59
      if (isPrime[i]) {
60
        Primes[pnx++] = i;
        int d = inx[i % 30030];
61
         for (ll j = i * i; j <= n; j += i * inc[d++]) {
     isPrime.reset(j);
     if (d == sz - 1) d = 0;
65
66
67
       if (c == sz - 1) c = 0;
68
69
```

# 5 String Processing

#### **5.1** Trie

```
Trie* children[26]; // Pointer = 8 Byte; 8*26 = 208 Byte
      int prefixs, words; // 8 Byte
      bool iseow; // 1 Byte
      char cur_letter; // 1 Byte
      vector <string> lex;
      priority_queue <pair <int, string>, vector <pair <int, string>>, greater <pair <int, string>>> occurrence; // small at top
10 public:
      Trie(char lett = '\0') {
11
12
        memset (children, 0, sizeof (children));
13
        prefixs = words = 0;
        iseow = false;
        cur_letter = lett;
17
18
      void insert(string &str) { // O(1)
19
        Trie* cur = this;
20
        int inx, strsz = str.size();
for(int i = 0; i < strsz; ++i) {</pre>
21
          inx = str[i] - 'a';
23
          if(cur->children[inx] == nullptr)
      cur->children[inx] = new Trie(str[i]);
           cur = cur->children[inx];
27
          cur->prefixs++;
28
29
        cur->iseow = true;
30
        cur->words++;
31
32
33
      int search_word(string &str) { // O(1)
34
        Trie* cur = this;
        int inx, strsz = str.size();
for(int i = 0; i < strsz; ++i) {</pre>
35
36
          inx = str[i] - 'a';
37
           if(cur->children[inx] == nullptr) {
38
      return 0;
```

```
41
           cur = cur->children[inx];
 42
         return cur->words;
       int search_prefix(string &str) { // O(1)
         Trie* cur = this;
         int inx = 0, strsz = str.size();
         for(int i = 0; i < strsz; ++i) {
  inx = str[i] - 'a';</pre>
 49
 51
           if(cur->children[inx] == nullptr) {
       return 0;
54
           cur = cur->children[inx];
55
         return cur->prefixs;
58
 59
       bool erase(string &str) {
60
         if (!search_word(str))
 61
           return false:
62
         Trie* cur = this:
64
         int inx, strsz = str.size();
 65
         for(int i = 0; i < strsz; ++i) {</pre>
          inx = str[i] - 'a';
 66
           if(--cur->children[inx]->prefixs == 0) {
 67
       cur->children[inx] = nullptr;
 69
       return true:
 70
           cur = cur->children[inx];
         if (--cur->words == 0) {
           cur->iseow = false;
 75
 76
         return true;
 77
        \textbf{void} \ \text{dfs} (\texttt{Trie*} \ \text{node, string s}) \ \ \{ \ // \ \text{lex order dfs -> traverse all the strings starting from root } \\ 
 81
         if (node->iseow) {
           lex.emplace_back(s);
 82
 83
 84
         for (int j = 0; j < 26; ++j)
           if(node->children[j] != nullptr) {
       dfs(node->children[j], s + string(1, node->children[j]->cur_letter));
 89
90
91
       void dfs2(Trie* node, string s) { // autocomplete dfs -> traverse all the strings starting from the
             end of the given prefix
92
         if (node->iseow)
93
          if(occurrence.size() < 10) {
       occurrence.push(make_pair(node->words, s));
 95
           else (
 96
       if(node->words > occurrence.top().first) {
         occurrence.pop():
         occurrence.push (make_pair(node->words, s));
100
102
103
         for(int i = 0; i < 26; ++i) if(node->children[i] != nullptr) {
104
       dfs2(node->children[i], s + string(1, node->children[i]->cur_letter));
105
106
107
108 public:
       vector <string> lex_order() { // all strings in lexicographical order
109
110
         lex.clear():
         Trie* cur = this;
111
         for(int i = 0; i < 26; ++i) if(cur->children[i] != nullptr) {
112
       dfs(cur->children[i], string(1, cur->children[i]->cur_letter));
113
114
         return lex;
117
118
       void autocomplete(string &pref) { // suggest top ten words with max frequency
119
         if(!search_prefix(pref))
120
121
         Trie* cur = this:
122
         int inx, presz = pref.size();
123
         for(int i = 0; i < presz; ++i) {
  inx = pref[i] - 'a';</pre>
           cur = cur->children[inx];
         for(int i = 0; i < 26; ++i) if(cur->children[i] != nullptr) {
       dfs2(cur->children[i], string(1, cur->children[i]->cur_letter));
130
```

```
132
133
         vector <string> st;
134
         while (!occurrence.empty()) {
          st.emplace_back(pref + occurrence.top().second);
136
           occurrence.pop();
138
         if(cur->iseow) {
139
           st.emplace_back(pref);
140
141
         while(!st.empty()) {
142
           cout << st.back() << endl;</pre>
143
           st.pop_back();
144
145
146 }:
```

# 5.2 Knuth Morris Pratt (KMP)

```
2 * KMP(Knuth-Morris-Pratt) Algorithm
   ** Longest Prefix
   *** proper prefix = all prefixes except the whole string
   *** propre suffix = all suffixes except the whole string
   ** Prefix Function = Failure Function
   *** Given String P of len m, Find F[m];
    *** let t = P[0...i]
    *** f[i] = length of the longest proper prefix of t that is suffix of t
   *** calculating i different ways
   *** match the pattern against itself
    *** O(m) for failure function
   *** O(n) for KMP
14
    **/
16 vector <int> LongestPrefix(string &p) {
      int psz = p.size();
      vector <int> longest_prefix(psz, 0);
19
      for(int i = 1, k = 0; i < psz; ++i) {</pre>
20
        while(k && p[k] != p[i]) k = longest_prefix[k - 1];
longest_prefix[i] = (p[k] == p[i] ? ++k : k);
21
22
23
^{24}
      return longest_prefix;
    vector <int> KMP(string &s, string &p) {
      int ssz = s.size(), psz = p.size();
      vector <int> longest_prefix = LongestPrefix(p), matches;
      for(int i = 0, k = 0; i < ssz; ++i) {
   while(k && p[k] != s[i]) k = longest_prefix[k - 1]; // Fail go back</pre>
32
33
34
        k += (p[k] == s[i]);
35
36
        if(k == psz) {
37
         matches.emplace_back(i - psz + 1);
           k = longest_prefix[k - 1]; // faill safe and find another pattern
39
40
41
42
```

# 6 Geometry

#### 6.1 Point

```
15
      bool operator == (point other) const {
16
        return ((fabs(x - other.x) < EPS) && (fabs(y - other.y) < EPS)); // " < EPS " equal to " == zero "
17
19
      bool operator > (point other) const {
20
        if(fabs(x - other.x) > EPS)
^{21}
          return x > other.x;
        return y > other.y;
23
24
      ld dist(point other) { // Euclidean distance
25
       ld dx = this->x - other.x;
ld dy = this->y - other.y;
27
        return sqrtl(dx * dx + dy * dy);
28
29
31
      ld DEG_to_RAD(ld theta) {
32
        return theta * PI / 180.0;
33
34
35
      ld RAD_to_DEG(ld theta) {
36
        return theta * 180.0 / PI;
37
38
      point rotate(ld theta) {
39
40
        ld rad = DEG_to_RAD(theta);
        return point (cos (theta) * x - sin (theta) * y,
41
         sin(theta) * x + cos(theta) * v);
43
44 };
```

# 7 More advanced topics

# 7.1 A\*-Algorithm

```
1 const int dr []
                        = \{-1, 0, 1, 0\};
   const int dr [] = {-1, 0, 1, 0,
const int dc [] = {0, 1, 0, -1};
const char dir [] = {'U', 'R', 'D', 'L'};
map <char, int> inv = { {'U', 0}, {'R', 1}, {'D', 2}, {'L', 3}};
2 const int dc []
   const int N = 1e3 + 9, M = 2e5 + 9, oo = 0x3f3f3f3f3f;
   const 11 INF = 0x3f3f3f3f3f3f3f3f3f;
 9 char grid[N][N];
10 int dis[N][N], n, m, si, sj, ti, tj;
11 char Par[N][N];
    vector < pair <int, int> > restorePath(int sr, int sc, int tr, int tc)
13
14
15
      vector < pair <int, int> > ret;
16
      if(dis[tr][tc] == oo) return ret;
17
18
      for(char i = Par[tr][tc]; (sr ^ tr) || (sc ^ tc); i = Par[tr][tc])
19
20
          ret.push_back({tr, tc});
^{21}
          tr += dr[inv[i]];
          tc += dc[inv[i]];
      ret.push_back({sr, sc});
      reverse(ret.begin(), ret.end());
28
   bool valid(int r, int c) {
30
      return r >= 0 && r < n && c >= 0 && c < m && grid[r][c] != '%';
31
    /** admissible heuristic **/
    int manhattanDistance(int x1, int y1, int x2, int y2) {
     return (abs(x1 - x2) + abs(y1 - y2));
37
39 int Astar(int sr, int sc, int tr, int tc)
40
41
      memset(dis, 0x3f, sizeof (dis));
42
      memset(Par, -1, sizeof (Par));
43
      priority_queue <tuple <int, int, int> > Q;
      dis[sr][sc] = 0;
46
47
      O.push({-manhattanDistance(sr, sc, tr, tc), sr, sc});
      int hoost, r, c, nr, nc;
```

```
while(Q.size())
 51
            tie(hcost, r, c) = Q.top(); Q.pop();
            if (r == tr && c == tc) return dis[r][c];
 55
            for (int i = 0; i < 4; ++i)
 56
          nr = r + dr[i];
 57
 58
          nc = c + dc[i];
 59
 60
          if(!valid(nr, nc)) continue;
 61
 62
          if(dis[r][c] + 1 < dis[nr][nc])
 63
              dis[nr][nc] = dis[r][c] + 1;
Par[nr][nc] = dir[i ^ 2];
 64
              Q.push({-dis[nr][nc] -manhattanDistance(nr, nc, tr, tc), nr, nc});
 67
 68
 69
 70
       return -1;
 71
 72
 73
     void Solve()
 74
       cin >> si >> sj >> ti >> tj >> n >> m;
 75
 76
       for(int i = 0; i < n; ++i)</pre>
 77
         for(int j = 0; j < m; ++j)
  cin >> grid[i][j];
 78
 79
       cout << Astar(si, sj, ti, tj) << endl;
vector < pair <int, int> > path = restorePath(si, sj, ti, tj);
 83
 84
          cout << point.first << " " << point.second << endl;</pre>
 85
 86
 87
 88 /**
 89
       P -> strat
        . -> target
 90
 91
 92
        input:
0 2 2 3 5 5
 93
        88P8-
        -8---
        88888
 99
100
        output:
101
        0.2
102
103
        1 2
104
105
        2 3
106
107 **/
```

# 7.2 Mo's algorithm

```
const int N = 3e4 + 9, M = 2e5 + 9, oo = 0x3f3f3f3f, Mod = 1e9 + 7;
 2 const 11 INF = 0x3f3f3f3f3f3f3f3f3f;
   const int BLK = 256;
 5
   struct query
 6
     int 1, r, id, blk;
     query() = default;
     query(int _l, int _r, int _id) {
11
       1 = _1;
       r = _r;
        id = _id;
14
       blk = 1 / BLK;
15
16
17
     bool operator < (const query other) const {</pre>
18
       if (blk ^ other.blk)
         return blk < other.blk;
19
20
        return (blk & 1) ? r < other.r: r > other.r;
22
   } queries[M];
24 int res[M], freq[M << 3], cur;
26 void add(int id) {
```

```
27
     cur += (++freq[id] == 1);
28
29
   void remove(int id) {
     cur -= (--freq[id] == 0);
33
   int get_res() {
34
35
     return cur;
36
37
38
   int cur_1, cur_r, 1, r, n, q, a[N];
39
40
   void Solve()
41
      for(int i = 1; i <= n; ++i) cin >> a[i];
44
45
46
      for(int i = 1; i <= q; ++i) {
47
        cin >> 1 >> r;
48
        queries[i] = query(l, r, i);
49
50
5.1
      sort (queries + 1, queries + 1 + q);
52
      cur_l = 1, cur_r = 0; // assign to right invalid index
53
      for(int i = 1; i <= q; ++i)
54
55
56
          int ql = queries[i].1;
57
         int qr = queries[i].r;
59
60
          while(cur_r < qr) add(a[++cur_r]);</pre>
61
62
          while(cur_l > ql) add(a[--cur_l]);
63
          while(cur_r > qr) remove(a[cur_r--]);
65
          while(cur_l < ql) remove(a[cur_l++]);</pre>
66
67
68
         res[queries[i].id] = get_res();
69
70
      for(int i = 1; i <= q; ++i)</pre>
71
72
        cout << res[i] << "\n";</pre>
73
```

# 7.3 Square root decomposition

```
1 const int N = 5e5 + 9, M = 1e3 + 9, oo = 0x3f3f3f3f, Mod = 1e9 + 7;
 2 const 11 INF = 0x3f3f3f3f3f3f3f3f3f3f;
 3 const int BLK = 256;
 5 int n, q, a[N], type, x, y, z;
 6 vector <int> bs[M];
    int query(int 1, int r, int val)
10
      int cur_1 = 1 / BLK;
      int cur_r = r / BLK;
13
14
      if(cur_1 == cur_r) {
15
        for (int i = 1; i <= r; ++i)</pre>
          ans += (a[i] >= val);
16
      else (
17
        for(int i = 1, _end = (cur_1 + 1) * BLK; i < _end; ++i)
ans += (a[i] >= val);
18
19
        for(int i = cur_l + 1; i <= cur_r - 1; ++i)
20
          ans += bs[i].end() - lower_bound(bs[i].begin(), bs[i].end(), val);
        for(int i = cur_r * BLK; i <= r; ++i)</pre>
          ans += (a[i] >= val);
26
27
28
    void build()
29
30
      for(int i = 0; i < n; ++i)</pre>
31
        bs[i / BLK].emplace_back(a[i]);
      for (int i = 0: i < M: ++i)
34
        sort(bs[i].begin(), bs[i].end());
35
37 void update(int id, int delta)
```

```
38 {
     int pos = lower_bound(bs[id / BLK].begin(), bs[id / BLK].end(), a[id]) - bs[id / BLK].begin();
     bs[id / BLK][pos] = delta;
     sort(bs[id / BLK].begin(), bs[id / BLK].end());
     a[id] = delta;
43 }
44
45 void Solve()
46
47
48
     for(int i = 1; i <= n; ++i) cin >> a[i];
49
50
51
     cin >> a:
52
      while (q--)
55
         cin >> type >> x >> y;
56
         if(type == 0)
57
58
       cin >> z:
59
       cout << query(x, y, z) << endl;</pre>
60
61
         else
62
     update(x, y);
63
64
```

# 8 Miscellaneous

# 8.1 Double comparison

```
bool approximatelyEqual(double a, double b, double epsilon)

return fabs(a - b) <= ((fabs(a) < fabs(b) ? fabs(b) : fabs(a)) * epsilon);

bool essentiallyEqual(double a, double b, double epsilon)

return fabs(a - b) <= ((fabs(a) > fabs(b) ? fabs(b) : fabs(a)) * epsilon);

bool definitelyGreaterThan(double a, double b, double epsilon)

return (a - b) > ((fabs(a) < fabs(b) ? fabs(b) : fabs(a)) * epsilon);

return (a - b) > ((fabs(a) < fabs(b) ? fabs(b) : fabs(a)) * epsilon);

bool definitelyLessThan(double a, double b, double epsilon)

return (b - a) > ((fabs(a) < fabs(b) ? fabs(b) : fabs(a)) * epsilon);

return (b - a) > ((fabs(a) < fabs(b) ? fabs(b) : fabs(a)) * epsilon);

return (b - a) > ((fabs(a) < fabs(b) ? fabs(b) : fabs(a)) * epsilon);</pre>
```

# 8.2 Fast input output

```
Fast Input/Output method for C++:
      1. cin(with sync_with_stdio(false) & cin.tie(nullptr)):
      - In = 5e61 => 420ms
      - |n| = 1e7| => 742ms
      - |n| = 5e6| => 895ms
10
      2. read (using getchar()):
      - |n| = 5e6| => 173ms
13
      - |n| = 1e7| => 172ms
14
15
      - |n| = 5e6| => 340ms
16 **/
17
18 ll readll () {
19
     bool minus = false:
     unsigned long long result = 0;
20
21
     char ch:
22
     ch = getchar();
23
     while (true) {
      if (ch == '-') break;
```

```
if (ch >= '0' && ch <= '9') break;
27
        ch = getchar();
30
      if (ch == '-') minus = true;
      else result = ch - '0';
34
        ch = getchar();
        if (ch < '0' || ch > '9') break;
        result = result \star 10 + (ch - '0');
36
37
38
      if (minus) return - (11) result:
39
40
      return result;
43
    int readi () {
      bool minus = false;
45
      unsigned int result = 0;
47
      ch = getchar();
48
49
      while (true)
        if (ch == '-') break;
if (ch >= '0' && ch <= '9') break;
50
5.1
52
        ch = getchar();
53
54
      if (ch == '-') minus = true;
      else result = ch - '0';
        ch = getchar();
        if (ch < '0' || ch > '9') break;
60
        result = result \star 10 + (ch - '0');
62
64
      if (minus) return -(int) result;
65
      return result;
66
```

# 8.3 Gcd & Lcm

```
1 ll gcd(ll a, ll b) { // binary GCD uses about 60% fewer bit operations
 2
     int shift = __builtin_ctz(a | b);
     a >>= __builtin_ctz(a);
      while (b) {
       b >>= __builtin_ctz(b);
10
       if (a > b)
11
        swap(a, b);
       b -= a;
     return a << shift:
16
17 ll lcm(ll a, ll b) {
18
     return a / gcd(a, b) * b;
19
```

#### 8.4 Modular calculations

```
16 }
17
   // 2. Modular Multiplication
20 ll binMul(ll a, ll b, ll p) {
23
     while (b) {
24
      if (b & 111)
25
        res = (res + a) % p;
26
       a = (a + a) % p;
27
      b >>= 1;
28
29
     return res:
30 }
   // 3. Modular Multiplicative Inverse
34 ll modInv(ll b, ll p) {
35
     return binExp(b, p - 2, p); // Guaranteed that p is a Prime Number
36
```

#### 8.5 Overloaded operators to accept 128 Bit integer

```
1 typedef __uint128_t
    typedef __int128
    template <class T> string to_string(T x)
      int sn = 1; if (x < 0) sn = -1, x \neq sn; string s = "";
      do { s = "0123456789"[x % 10] + s, x /= 10; } while(x);
return (sn == -1 ? "-" : "") + s;
 9
10
11 auto str_to_int(string x)
12
      ui128 ret = (x[0] == '-' ? 0 : x[0] - '0');
for(int i = 1; i < x.size(); ++i) ret = ret * 10 + (x[i] - '0');
return (x[0] == '-' ? -1 * (i128) ret : ret);
13
14
1.5
16
    istream & operator >> (istream & in, i128 & i) noexcept { string s; in >> s; i = str_to_int(s); return
19 ostream & operator << (ostream & os, const i128 i) noexcept { os << to_string(i); return os; }
20 istream & operator >> (istream & in, uil28 & i) noexcept { string s; in >> s; i = str_to_int(s);
21 ostream & operator << (ostream & os, const ui128 i) noexcept { os << to_string(i); return os; }
```

#### 8.6 Policy based data structures

#### 8.7 stress test

```
12
       ofstream cout("input.in");
 13
       i32 t = 2;
 14
       cout << t << endl;</pre>
 16
         i32 n = int64(1, 100), m = int64(1, 100);

cout << n << " " << m << endl;
 17
 18
 19
20
21
22
23
24
25 }
26
27 i3:
           i32 u = int64(1, n), v = int64(1, n), c = int64(1, 4);

cout << u << " " << v << " " << c << endl;
    i32 main (i32 arg, char* args[]) {
       fast();
29
      i32 tc = 0;
i32 limit = 100;
30
31
       if(arg != 3) return 0;
32
33
      34
35
36
      bf = flags + args[1] + " " + args[1] + ex;
oz = flags + args[2] + " " + args[2] + ex;
char bff[bf.size() + 1];
37
38
39
       char ozz[oz.size() + 1];
       strcpy(bff, bf.c_str());
       strcpy(ozz, oz.c_str());
44
       // compile command
45
       system(bff);
```

```
46
      system(ozz);
47
      ex = ".out";
pr = "./";
49
      bf = pr + args[1] + " < input.in > " + args[1] + ex;
oz = pr + args[2] + " < input.in > " + args[2] + ex;
50
51
52
       strcpy(bff, bf.c_str());
53
       strcpy(ozz, oz.c_str());
54
55
56
       while (++tc <= limit) {</pre>
         gen();
cerr << tc << endl;
57
58
         // run command
system(bff);
59
60
          system(ozz);
62
          ifstream brute_forces("brute_force.out");
63
          ifstream optimizes("optimized.out");
64
          string brute_force, optimized;
65
          getline(brute_forces, brute_force, (char)EOF);
66
67
          getline(optimizes, optimized, (char)EOF);
          if(brute_force != optimized) {
  cerr << "Wrong Answer" << endl;</pre>
69
70
            break;
71
         } else if (tc == limit) {
  cout << "Accepted" << endl;</pre>
72
73
74
75
```