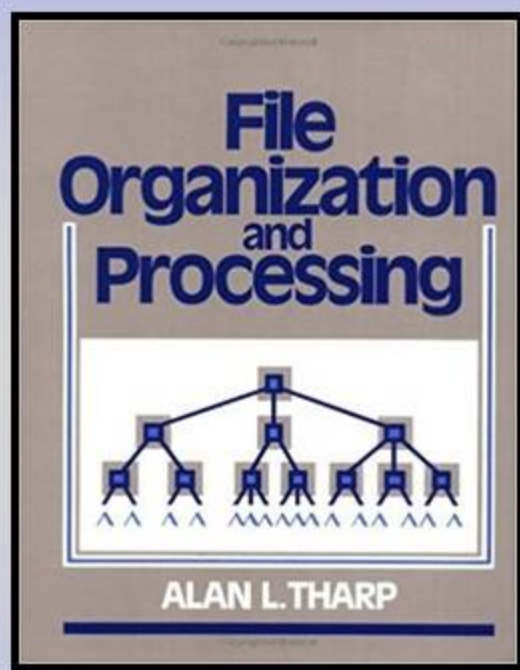


File Organization & Processing

CS2202

Instructor:

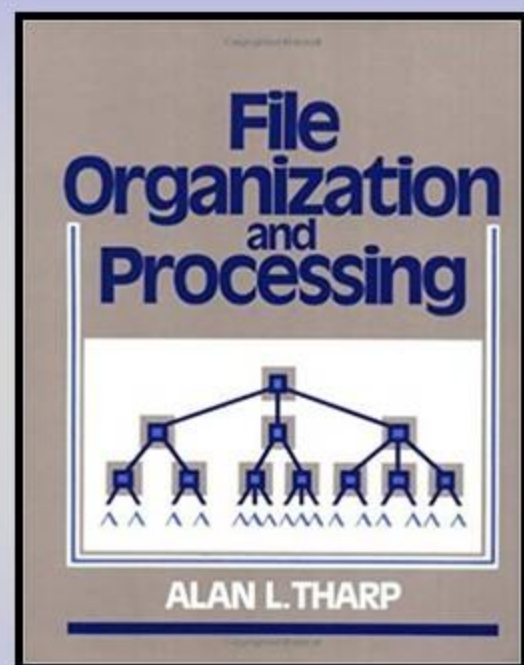
Dr. Mohamed Hassan





Chapter 2:

- Secondary Storage
-



Secondary Storage Devices

➤ **Two** major types of secondary storage devices:

1. **Direct Access Storage Devices (DASDs)**

- **Magnetic Discs**

 - Hard disks (high capacity, low cost, fast)

 - Floppy disks (low capacity, lower cost, slow)

- **Optical Discs**

 - CD-ROM = (Compact disc, read-only memory)

2. **Serial Devices**

- **Magnetic tapes** (very fast sequential access)

Storage and Files

- **Storage has major implications for DBMS design!**
 - **READ**: transfer data from disk to main memory (RAM).
 - **WRITE**: transfer data from RAM to disk.
 - Both operations are **high-cost** operations, relative to in-memory operations, so DB must be planned carefully!



Storage and Files

- **Why Not Store Everything in Main Memory?**
 - **Costs** too much: Cost of RAM about 100 times the cost of the same amount of disk space, so relatively small size.
 - Main memory is **volatile**.
 - **Typical storage hierarchy:**
 - **Main memory** (RAM) (**primary** storage) **for currently used data.**
 - **Disk** for **the main database** (**secondary** storage).
 - **Tapes** for **archiving older versions** of the data (**tertiary** storage).



Storage Hierarchy

- **Primary storage** : random access memory (RAM)
 - typical **capacity** a number of **GB**
 - **cost** per **MB \$2-3.00**
 - typical **access time** **5ns to 60ns**
- **Secondary storage**: magnetic disk/ optical devices/ tape systems
 - typical **capacity** a number of **100GB** for **fixed media**; ∞ for removable
 - **cost** per **MB \$0.01** for **fixed media**, more for **removable**
 - typical **access time** 8ms to **12ms** for fixed media, larger for removable



Units of Measurement

Spatial units:

- o **byte**: 8 bits
- o **kilobyte (KB)**: 1024 or 2^{10} bytes
- o **megabyte (MB)**: 1024 kilobytes or 2^{20} bytes
- o **gigabyte (GB)**: 1024 megabytes or 2^{30} bytes

Time units:

- o **nanosecond (ns)** one- billionth (10^{-9}) of a second
- o **microsecond (μ s)** one- millionth (10^{-6}) of a second
- o **millisecond (ms)** one- thousandth (10^{-3}) of a second

Primary versus Secondary Storage

- Primary storage costs **several hundred times** as much per unit as secondary storage,
- but has access times that are **250,000 to 1,000,000 times** faster than secondary storage.



Memory Hierarchy

- At the **primary storage level**, the memory hierarchy includes, at **the most expensive** end' **cache memory**, which is a **static RAM** (Random Access Memory).
- The **next level of primary storage** is **DRAM (Dynamic RAM)**, The advantage of DRAM is its **low cost, lower speed compared with static RAM**.
- Programs normally reside and execute in DRAM.



Memory Hierarchy-flash memory

Flash memory,

- Since 1988 it has become common, particularly because it is **nonvolatile**, using **EEPROM** (**E**lectrically **E**rasable **P**rogrammable **R**ead-**O**nly **M**emory) technology.
- Its life is **10,000-1,000,000** times erase...
- **Read/write** is **fast**, but **erase** is **slow**...
- Therefore special arrangements are made for the file system, regarding file delete or update.
- **Capacities** up to **128 GB** has been realized to date.

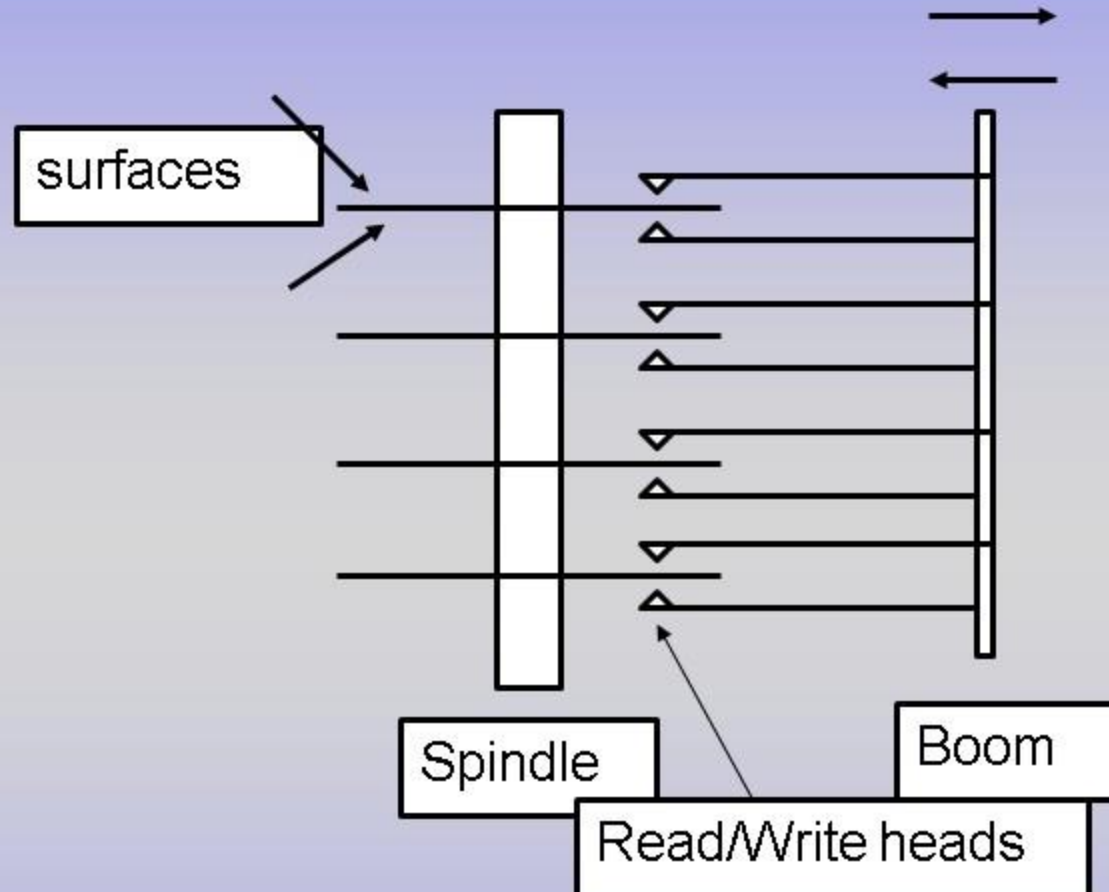


Magnetic Disks

- Bits of data (0's and 1's) are stored on **circular magnetic platters** called **disks**.
- A disk rotates rapidly (& **never stops**).
- A **disk head** reads and writes bits of data as they pass under the head.
- Often, **several platters** are organized into a **disk pack (or disk drive)**.

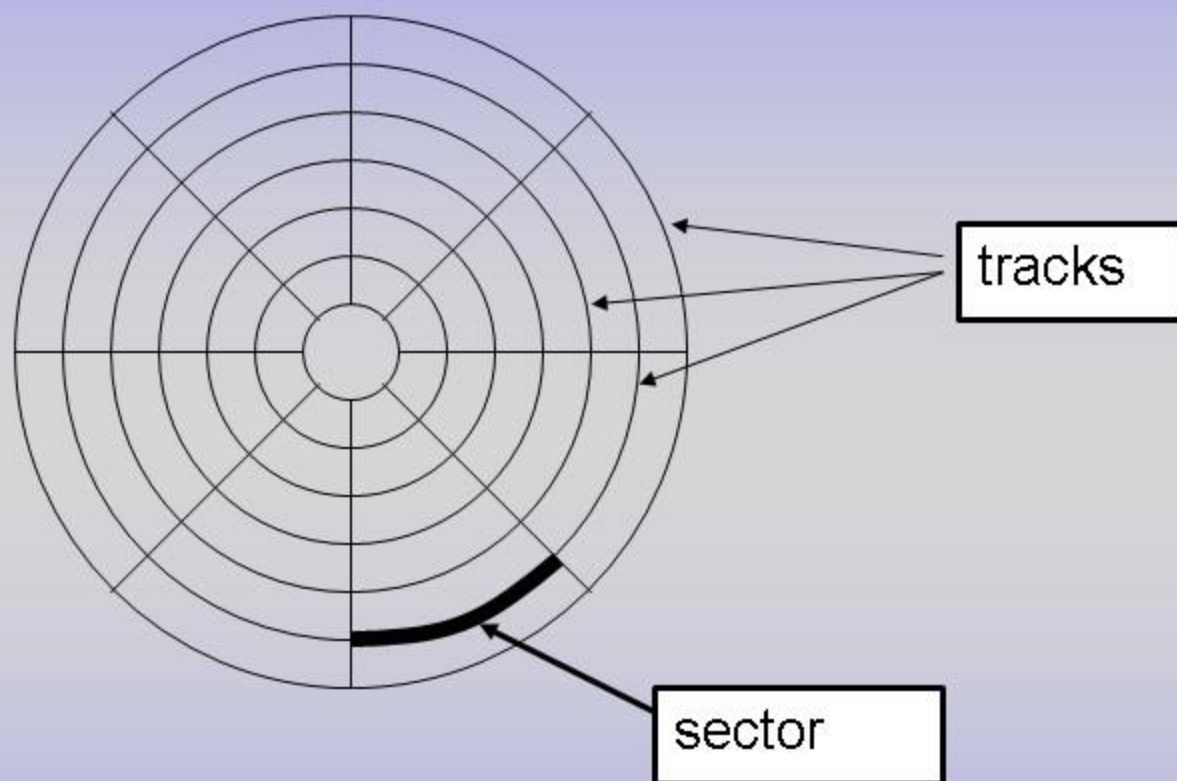


A Disk Drive



Disk drive with 4 platters and 8 surfaces and 8 RW heads

Looking at a surface



Surface of disk showing tracks and sectors

Organization of Disks

- Disk contains concentric **tracks**.
- Tracks are divided into **sectors**
- A **sector** is the smallest addressable unit in a disk.



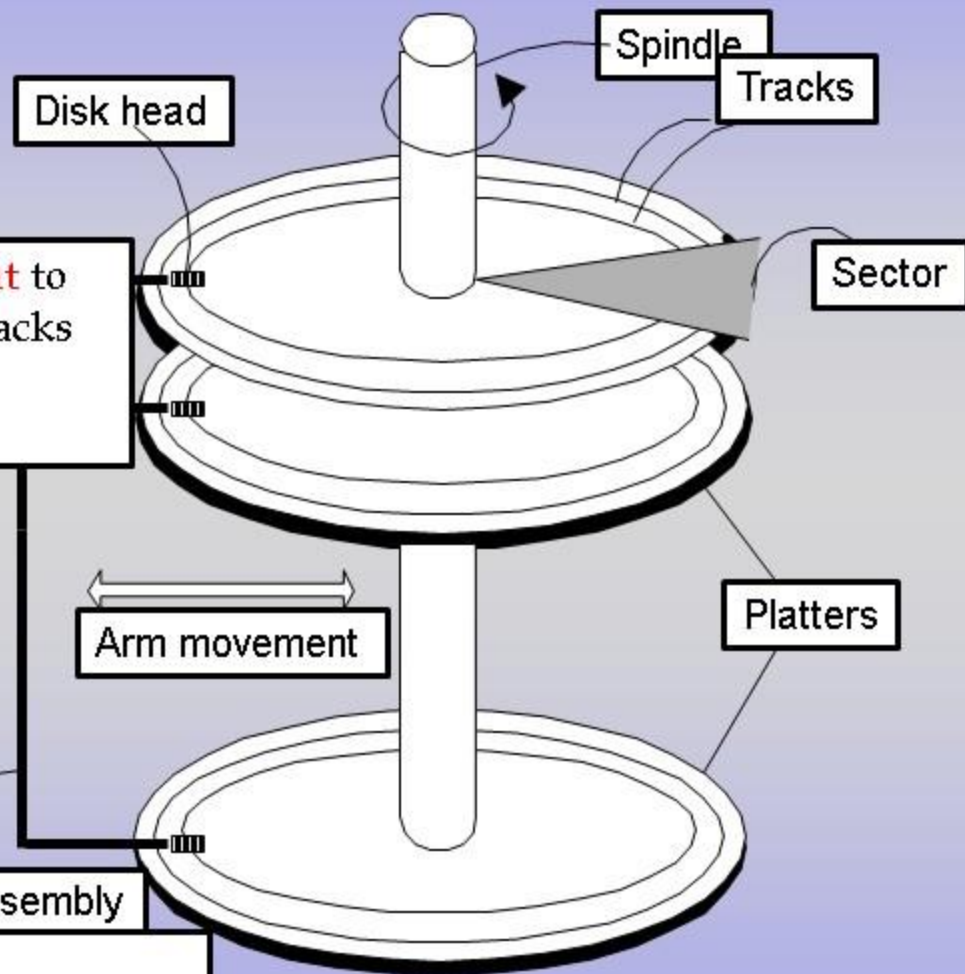
Components of a Disk

❖ The platters spin (say, 90rps).

❖ The arm assembly is moved **in or out** to position a head on a desired **track**. Tracks under heads make a **cylinder** (imaginary!).

❖ Only one head reads/writes at any one time.

❖ **Block size** is a multiple of **sector size** (which is often fixed).



Disk Controller

- **Disk controllers**: typically embedded in the disk drive, which **acts as an interface between the CPU and the disk hardware**.
- The controller has **an internal cache** (typically a number of MBs) that it uses to **buffer data for read/write requests**.



Accessing Data

- When a program reads a byte from the disk, the operating system locates the surface, track and sector containing that byte, and reads the entire sector into a special area in main memory called **buffer**.
- The **bottleneck** of a disk access is moving the read/write arm.
- So it makes sense to store a file in tracks that are below/above each other on different surfaces, rather than in several tracks on the same surface.

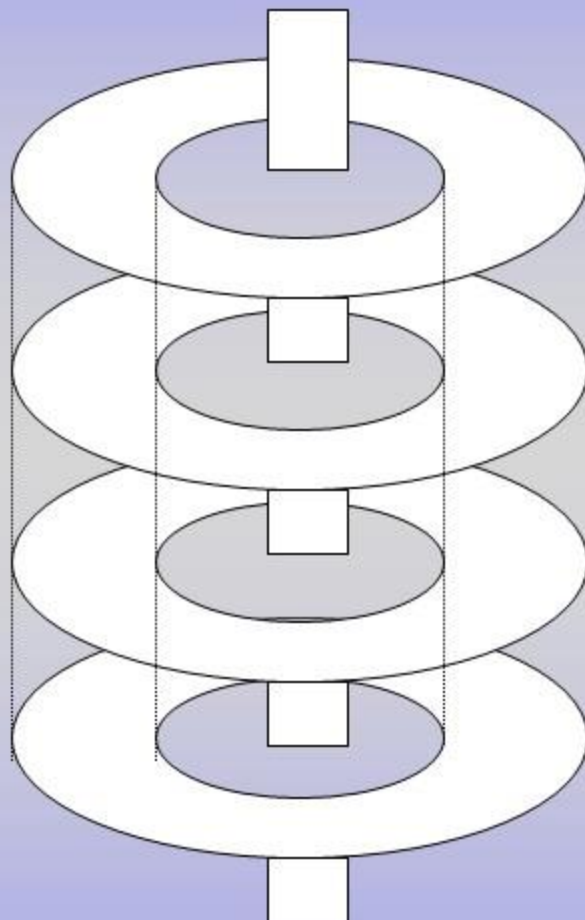


Cylinders

- A **cylinder** is the set of tracks at a given radius of a disk pack.
 - i.e. a cylinder is the set of tracks that can be accessed without moving the disk arm.
- All the information on a cylinder can be accessed without moving the read/write arm.



Cylinders



Estimating Capacities

- **Track capacity** = # of sectors/track * bytes/sector
- **Cylinder capacity** = # of tracks/cylinder * track capacity
- **Drive capacity** = # of cylinders * cylinder capacity
- **Number of cylinders** = # of tracks in a surface

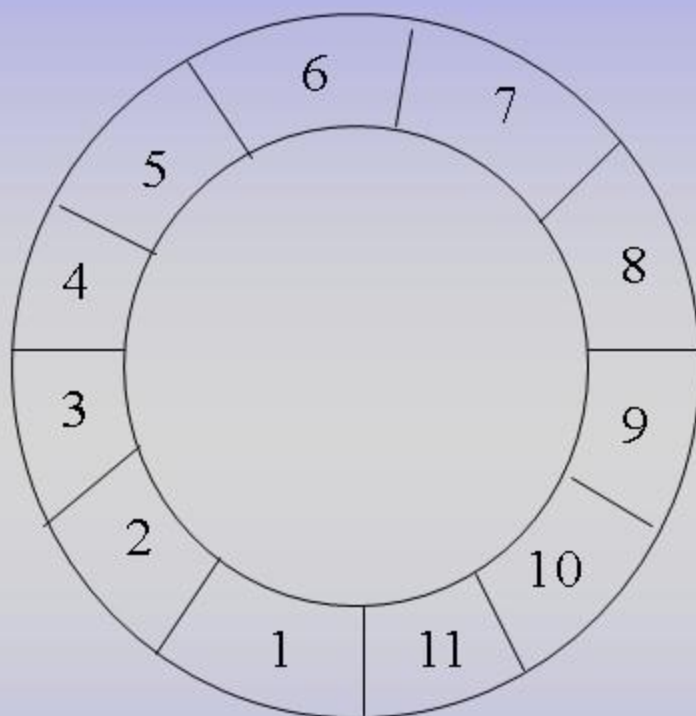


Exercise

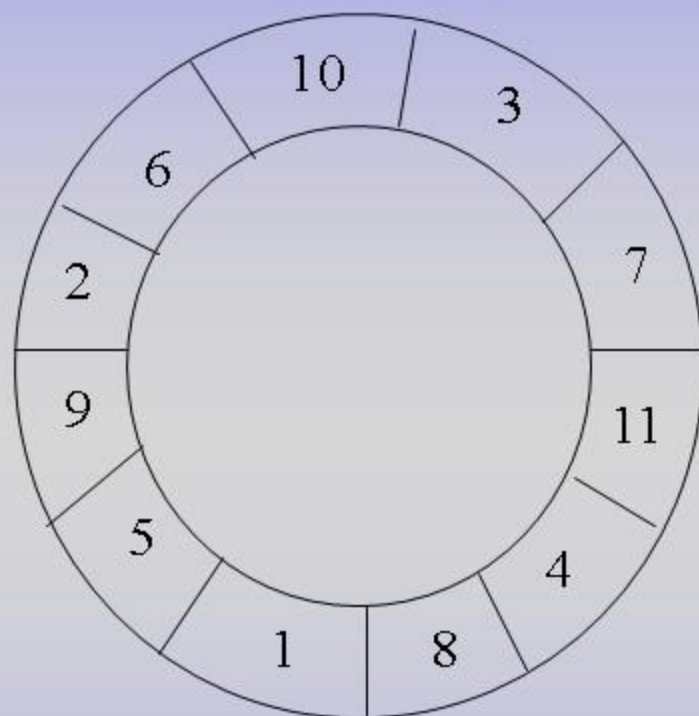
- Store a file of 20000 records on a disk with the following characteristics:
 - # of bytes per sector = 512
 - # of sectors per track = 40
 - # of tracks per cylinder = 11
 - # of cylinders = 1331
- Q1.** How many cylinders does the file require if each data record requires 256 bytes?
- Q2.** What is the total capacity of the disk?



Organizing Tracks by sector



**Physically adjacent
sectors**



**Sectors with 3:1
interleaving**

Clusters

- Usually **File manager**, under the operating system, maintains the logical view of a file.
- **File manager** views the file as a series of **clusters**, each of a number of sectors. The clusters are ordered by their logical order.
- **Files** can be seen in the form of logical sectors or blocks, which **needs to be mapped** to physical clusters.
- **File manager** uses a **file allocation table (FAT)** to **map logical sectors of the file to the physical clusters**.



Extents

- If there is **a lot of room on a disk**, it may be possible to make **a file consist entirely of contiguous clusters**. Then we say that the file is one **extent**. (**very good for sequential processing**)
- If there **isn't enough contiguous space available** to contain an entire file, the file is **divided into two or more noncontiguous** parts. **Each part is a separate extent**.



Internal Fragmentation

- **Internal fragmentation**: loss of space within a sector or a cluster.
- 1) **Due to records not fitting exactly in a sector**:
e.g. Sector size is 512 and record size is 300 bytes.
Either
 - store one record per sector, or
 - allow records *span* sectors...
 - 2) **Due to the use of clusters**: If the file size is not a multiple of the cluster size, then the last cluster will be partially used.



Organizing Tracks by Block

- Disk tracks may be divided into user-defined blocks rather than into sectors.
- Blocks can be **fixed** or **variable** length.
- A block is **usually organized to hold an integral number of logical records**.
- **Blocking Factor** = number of records stored in a block.
- No internal fragmentation, no record spanning over two blocks.
- In block-addressing scheme each block of data may be accompanied by one or more **subblocks containing extra information about the block**: record count, last record key on the block...



Non-data Overhead

- **Both blocks and sectors require non-data overhead** (written during formatting)
- **On sector addressable disks**, this information involves **sector address, track address, and condition (usable/defective)**. Also pre-formatting involves placing gaps and synchronization marks between the sectors.
- **On block-organized disk**, where a block may be of **any size**, more information is needed and the programmer should be aware of some of this information to utilize it for better efficiency...



The Cost of a Disk Access

- The time to access a sector in a track on a surface is divided into 3 components:

Time Component	Action
Seek Time	Time to move the read/write arm to the correct cylinder
Rotational delay (or latency)	Time it takes for the disk to rotate so that the desired sector is under the read/write head
Transfer time	Once the read/write head is positioned over the data, this is the time it takes for transferring data



Seek time

- **Seek time** is the time required to move the arm to the correct cylinder.
- Largest in cost.

Typically:

- **5 ms** (milliseconds) to move from one track to the next (track-to-track)
- **50 ms** maximum (from inside track to outside track)
- **30 ms** average (from one random track to another random track)



Latency (Rotational Latency)-1

- **Latency** is the time needed for the disk to rotate so the sector we want is under the read/write head.
- Hard disks usually **rotate** at about **5000-7000 rpm**,
 - 12-8 msec per **revolution**.
- Note:
 - **Min latency** = 0
 - **Max latency** = Time for one disk revolution
 - **Average latency (r)** = $(\text{min} + \text{max}) / 2$
= $\text{max} / 2$
= time for $\frac{1}{2}$ disk revolution
- Typically 6 – 4 ms, at average



Transfer Time-1

- **Transfer time** is the time for the read/write head to pass over a block.
- The transfer time is given by the formula:

$$\text{Transfer time} = \frac{\text{number of sectors}}{\text{track capacity in number of sectors}} \times \text{rotation time}$$

- e.g. if there are S_t sectors per track, the time to transfer one sector would be $1/S_t$ of a revolution.



Transfer Time-2

- The transfer time depends only on the speed at which the spindle rotates, and the number of sectors that must be read.

- **Sequential Reading**

- Given the following disk:
 - Avg. Seek time $s = 16$ ms
 - Avg. Rot. Latency $r = 8.3$ ms
 - Block transfer time = 8.4 ms
- a) Calculate the time to read 10 sequential blocks, on the same track.
- a) $16 + 8.3 + 8.4 * 10$



Secondary Storage Devices: Magnetic Tapes



Characteristics

- **No direct access**, but **very fast sequential access**.
- Resistant to different environmental conditions.
- Easy to **transport, store, cheaper than disk**.
- Before it was widely used to store application data; nowadays, **it's mostly used for backups or archives**.



MT Characteristics-2

- A sequence of bits are stored on magnetic tape.
- For storage, the tape is wound on a reel.
- To access the data, the tape is unwound from one reel to another.
- As the tape passes the head, bits of data are read from or written onto the tape.

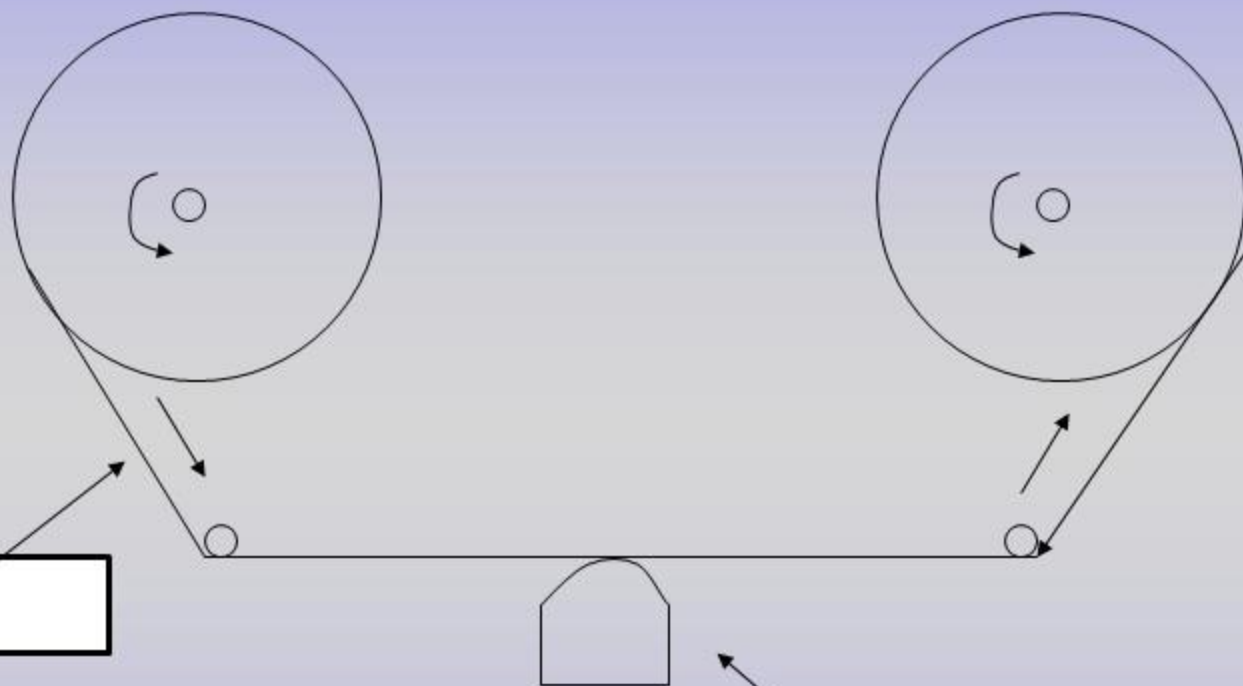


Reel 1

Reel 2

tape

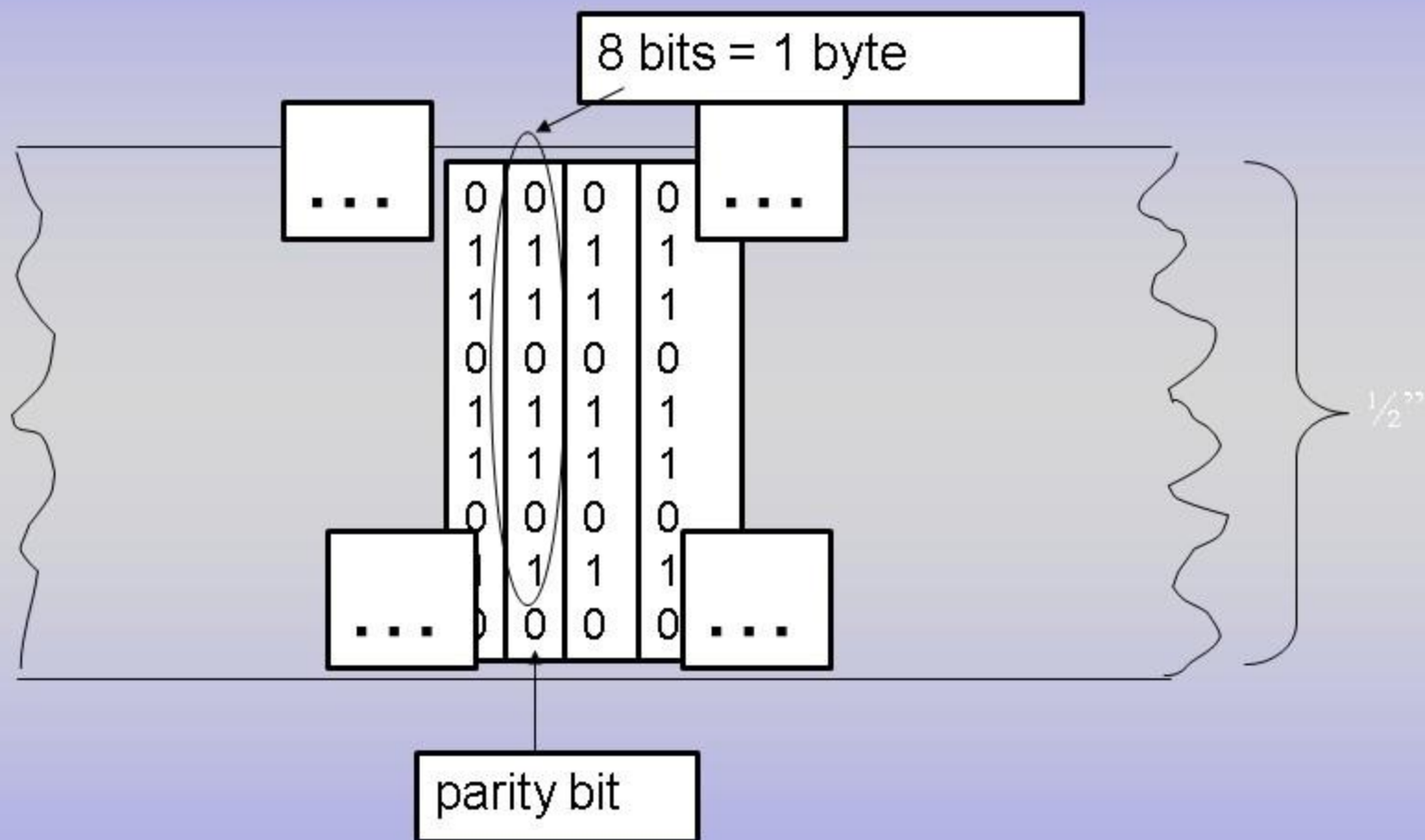
Read/write head



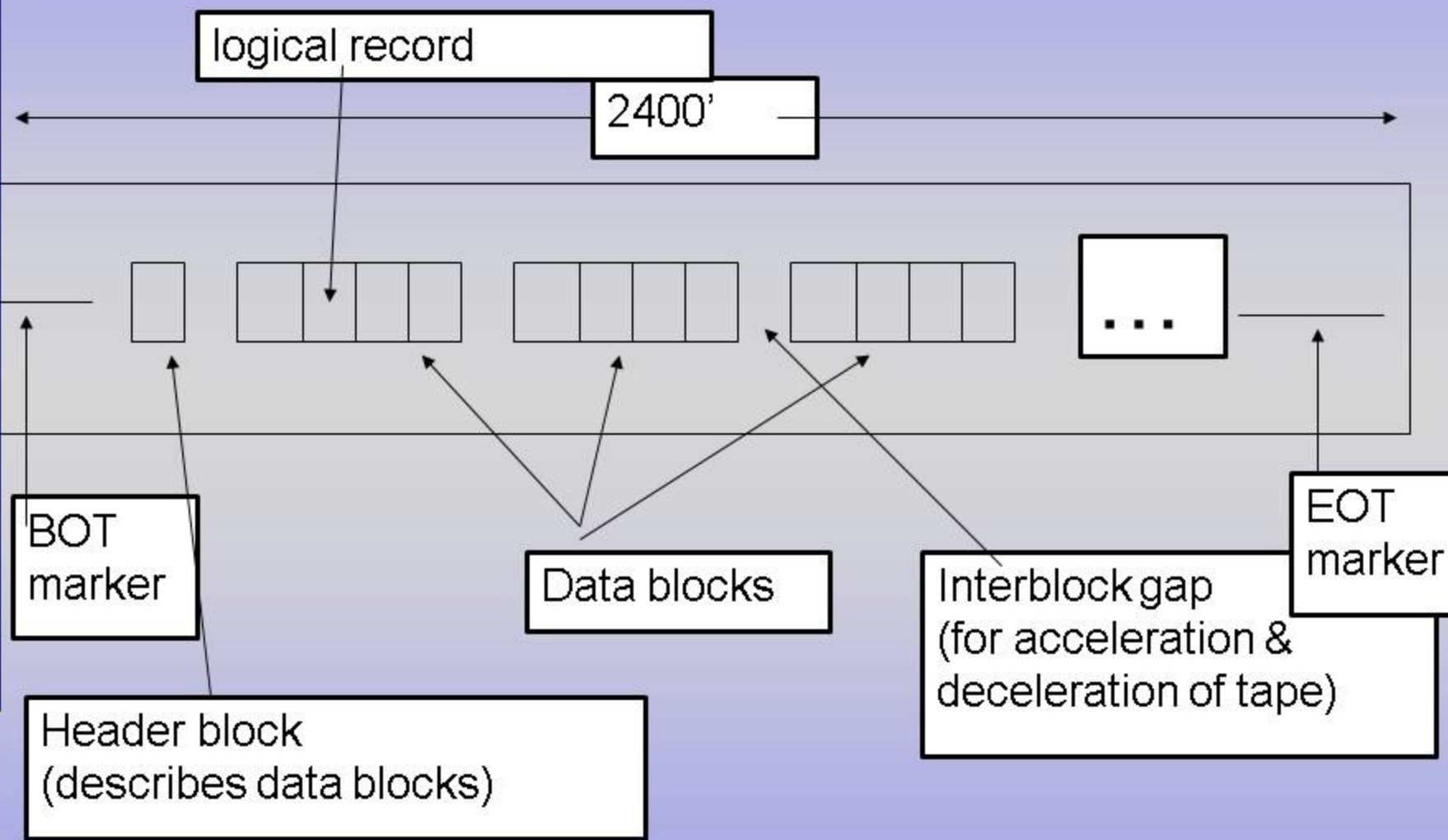
Tracks

- Typically data on tape is **stored** in **9 separate bit streams**, or **tracks**.
- Each track is a sequence of bits.
- **Recording density = # of bits per inch (bpi)**. Typically
800 or 1600 bpi.
30000 bpi on some recent devices.

MT recording in detail



Tape Organization



Secondary Storage Devices: CD-ROM



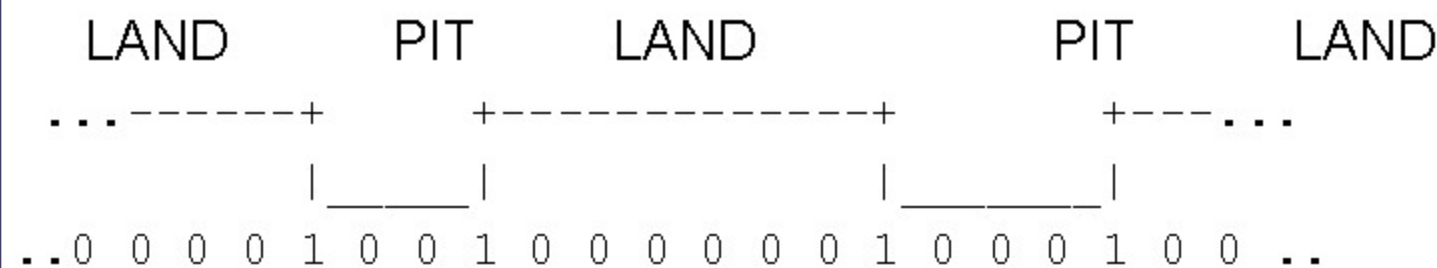
Physical Organization of CD-ROM

- **Compact Disk** – read only memory (write once), R/W is also available.
- Data is encoded and **read optically with a laser**
- Can store around +600MB data
- **Digital data** is represented as a **series of Pits and Lands:**
 - **Pit** = a little depression, forming a lower level in the track
 - **Land** = the flat part between pits, or the upper levels in the track



Organization of data

- Reading a CD is done by shining a laser at the disc and detecting changing reflections patterns.
 - **1 = change in height (land to pit or pit to land)**
 - **0 = a “fixed” amount of time between 1’s**



- Note : **we cannot have two 1's in a row!**



CD-ROM

- While the speed of CD-ROM readers is relatively higher, such as 24X(24 times CD audio speed), the speed of writing is much slower, as low as 2X.

.**The DVD** (**Digital Video Disc or Digital Versatile Disc**) technology is based on CD technology with increased storage density.

- **The DVD technology allows two-side medium**, with a storage capacity of up to 10GB.

