

# Design & Analysis of Algorithms I

## Mid 1, Spring 2014

Date: 27<sup>th</sup> Feb. 2014

Time: 90 mins.

### Q1. (5+5)

Below are the pseudo codes for insertion sort and bubble sort. It is assumed that data is stored in an array  $A[1 \dots n]$ . Determine the loop invariant for the inner loops of both the sorts and prove their correctness.

Insertion Sort	Bubble Sort
<pre>for j = 2 to A.length     key = A[j]     i = j - 1     while i &gt; 0 and A[i] &gt; key         A[i+1] = A[i]         i = i - 1     A[i+1] = key</pre>	<pre>for i = 1 to A.length-1     for j = 1 to A.length - i         if ( A[j] &gt; A[j + 1] )             temp = A[j]             A[j]=A[j + 1]             A[j + 1] = temp</pre>

**Remarks: Looks good. No changes from me.**

### Modified Q2 (10)

Below is the pseudo code of count Sort. The indexes are 0-based for the C array, but 1-based for the arrays A and B. The below algorithm is stable.

If we however change the last for loop to go from 1 up to A.length, instead of A.length down to 1, it does not remain stable.

Your task is to change the Count sort code, so that with the new code, if we go from 1 to A.length in the last for loop, it still remains stable.

The modified algorithm must still be stable, and must still run in  $O(n+k)$  time.

```
COUNT-SORT (A, B , k)
//Let C[0..k] be a new array
for i = 0 to k
    C[i] = 0
for j = 1 to A.length
    C[A[j]] = C[A[j]]+1
for i = 1 to k
    C[i] = C[i] + C[i-1]
for j = A.length down to 1
    B[C[A[j]]] = A[j]
    C[A[j]] = C[A[j]] - 1
```

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### Modified Q3 (10)

Let array  $A$  be an array consisting of only zeros and ones. (0's and 1's). Suggest an algorithm to sort the records in  $O(n)$  time and  $O(1)$  additional space.