

# **Wireless Sensor Networks**

# Wireless sensor network

- ❖ A wireless sensor network (WSN) is a wireless network consisting of **spatially distributed autonomous devices** using sensors to **cooperatively monitor physical or environmental conditions**, such as temperature, sound, vibration, pressure, motion or pollutants, at different locations
- ❖ Wireless sensors are limited in memory, computation power, bandwidth, and energy. Due to their small physical size, they can be embedded in the physical environment
- ❖ Low cost & energy implies low power CPU, radio with minimum bandwidth and range
- ❖ Ad-hoc deployment implies no maintenance or battery replacement

# Sensor network

## ❖ Sensor

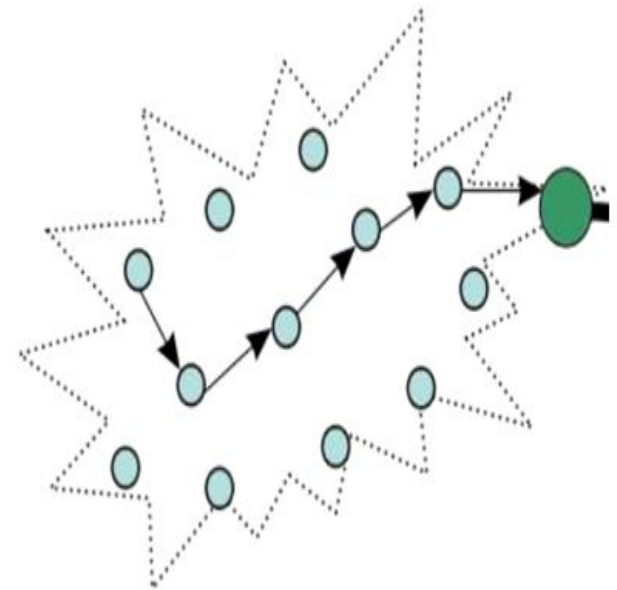
- A transducer
- converts physical phenomenon e.g. heat, light, motion, vibration, and sound into electrical signals

## ❖ Sensor node

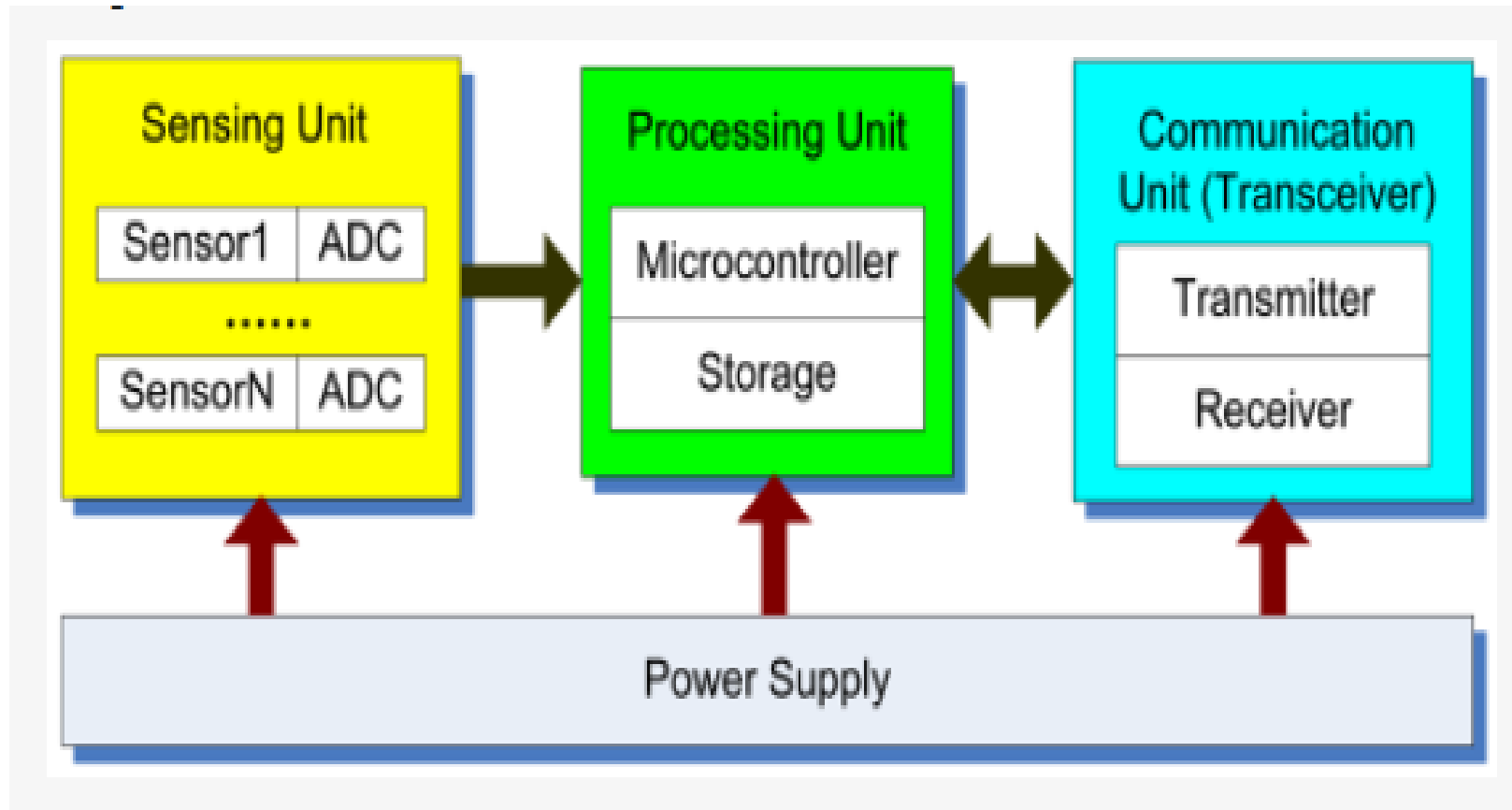
- basic unit in sensor network
- contains on-board sensors, processor, memory, transceiver, and power supply

## ❖ Sensor network

- consists of a large number of sensor nodes
- nodes deployed either inside or very close to the sensed phenomenon



# Hardware Architecture of sensor node



# Characteristics of WSN

- ❖ Sensing and data processing are essential
- ❖ WSNs have many more nodes and are more densely deployed
- ❖ Hardware must be cheap; nodes are more prone to failures
- ❖ WSNs operate under very strict energy constraints
  - Sleep as much as possible.
  - Acquire data only if indispensable.
  - Use data fusion and compression.
  - Transmit and receive only if necessary.  
Receiving is just as costly as sending.
- ❖ WSN nodes are typically static
- ❖ The communication scheme is many-to-one (data collected at a base station) rather than peer-to-peer

# Data Collection

- ❖ Centralized data collection puts extra burden on nodes close to the base station. Clever routing can alleviate that problem
- ❖ **Clustering:** data from groups of nodes are fused before being transmitted, so that fewer transmissions are needed
- ❖ Often getting measurements from a particular area is more important than getting data from each node
- ❖ Security and authenticity should be guaranteed. However, the CPUs on the sensing nodes cannot handle fancy encryption schemes.

# Infrastructure vs. Ad-Hoc Networks

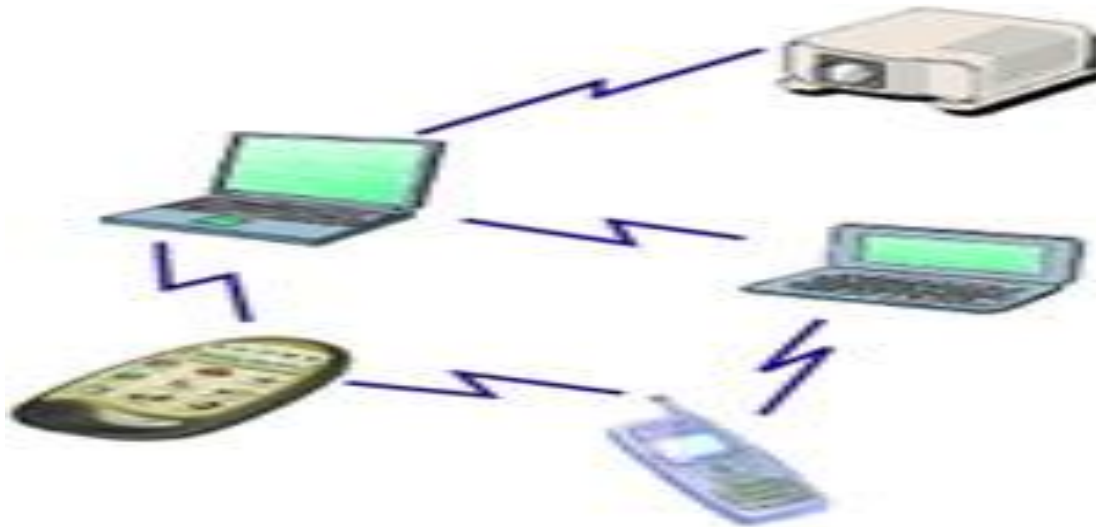
- ❖ Infrastructure network consists of **wired network or fixed connectivity**

**Ad-hoc means *'for this purpose'***

- ❖ No need for infrastructure (like routers, cell towers, etc.)
- ❖ It can accommodate new devices at any time.
- ❖ MANET: Mobile Ad-Hoc Network
- ❖ WSN: Wireless sensor network

# MANET

- ❖ MANET – (Mobile Ad-Hoc NETwork) a system of mobile nodes (laptops, sensors, etc.) interfacing without the assistance of centralized infrastructure (access points, bridges, etc.)





# Ad Hoc Wireless Networks

- ❖ Large number of self-organizing static or mobile nodes that are possibly randomly deployed
- ❖ Near(est)-neighbor communication
- ❖ Wireless connections
  - Links are fragile, possibly asymmetric
  - Connectivity depends on power levels and fading
  - Interference is high for omnidirectional antennas
- ❖ WSNs are ad hoc networks (wireless nodes that self-organize into an infrastructure-less network).

# Environment monitoring

*Zebranet*: a WSN to study the behavior of zebras



- Special GPS-equipped collars were attached to zebras
- Data exchanged with peer-to-peer info swaps
- Coming across a few zebras gives access to the data

# Medical application



- Vital sign monitoring
- Accident recognition
- Monitoring the elderly

- ❖ Intel deployed a 130-node network to monitor the activity of residents in an elder care facility.
- ❖ Patient data is acquired with wearable sensing nodes (the “watch”)

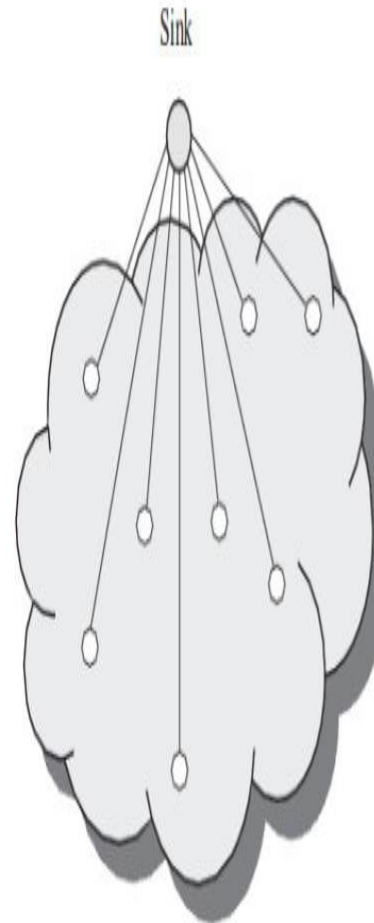
# Data-dissemination Schemes

## Direct communication with the base station

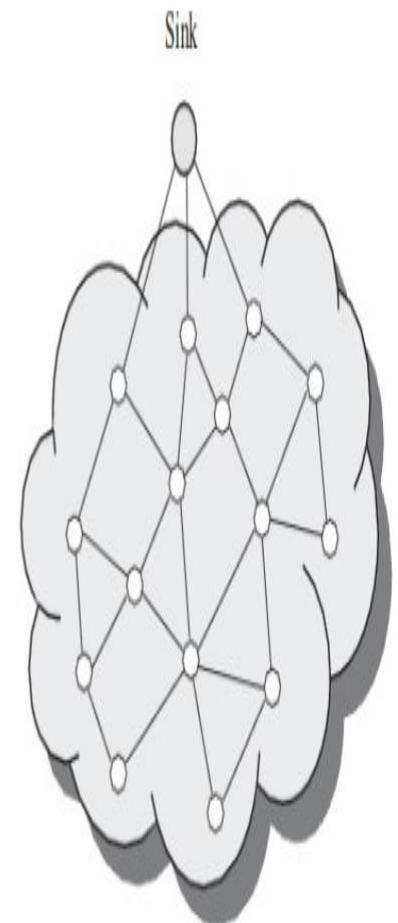
- ❖ Sensor nodes communicate with the base station directly.

## Multi-hop Scheme

- ❖ Transmit through some other intermediate nodes.



Single-hop network



Multi-hop network

# WSN network Architecture

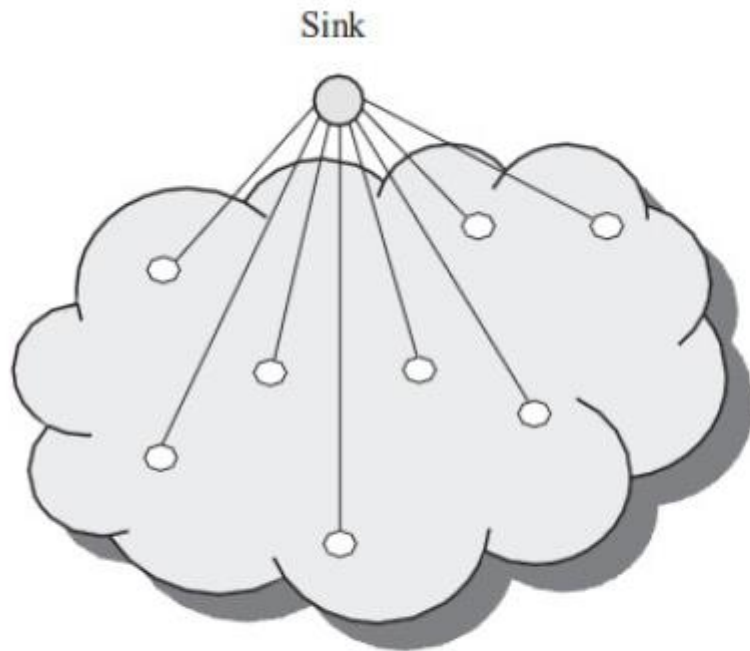
- ❖ Flat Architecture:

Each node plays the same role in performing sensing task and all sensor nodes are peers

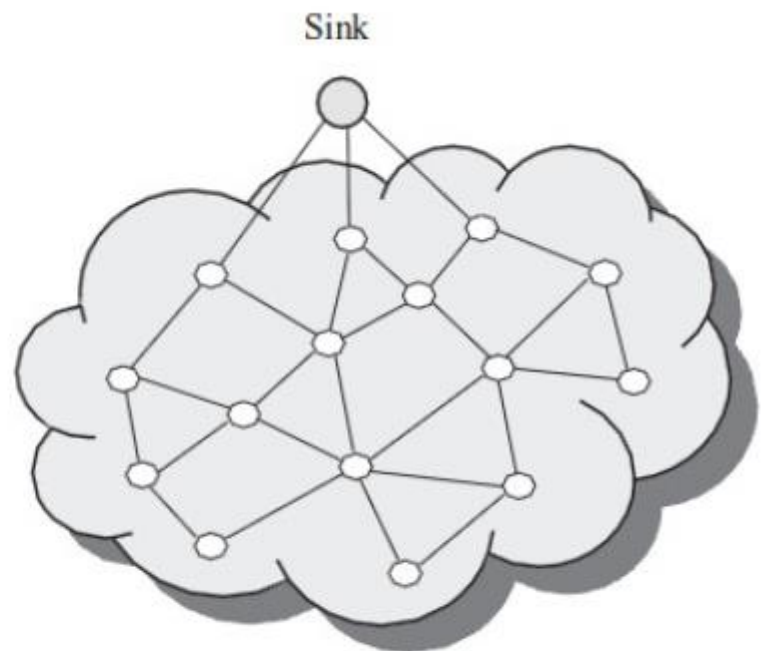
- ❖ Hierarchical Architecture:

Sensor nodes are organized clusters, where the cluster members send their data to the sink

# Flat Architecture

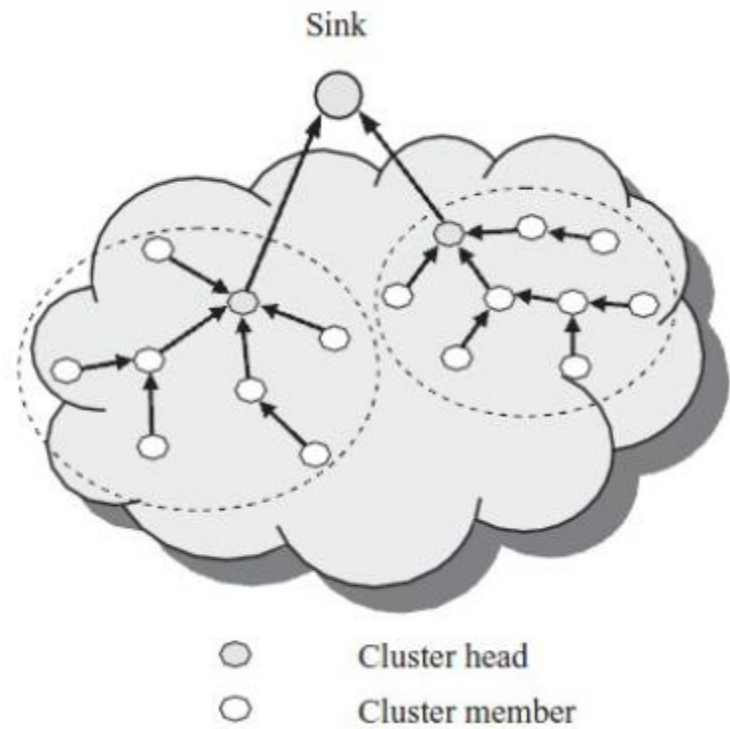
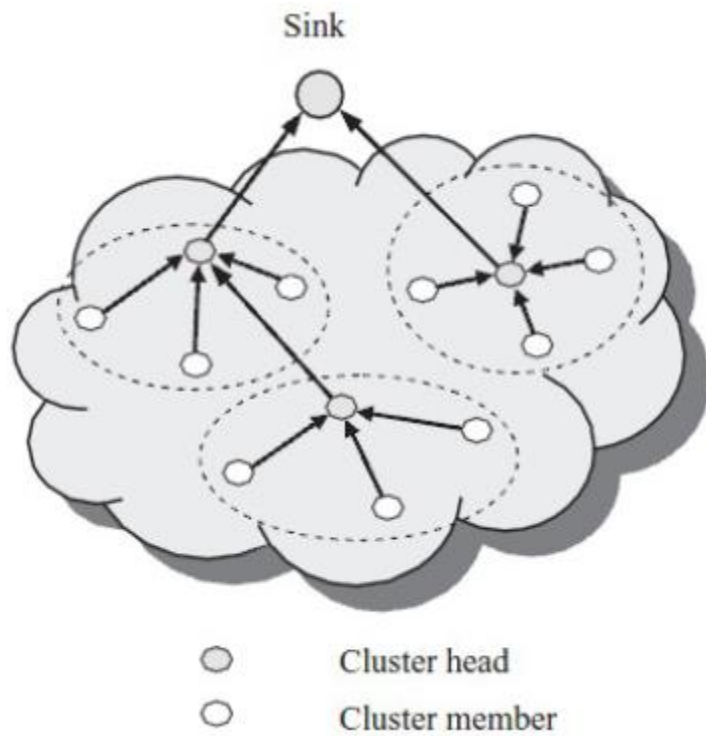


Single-hop network

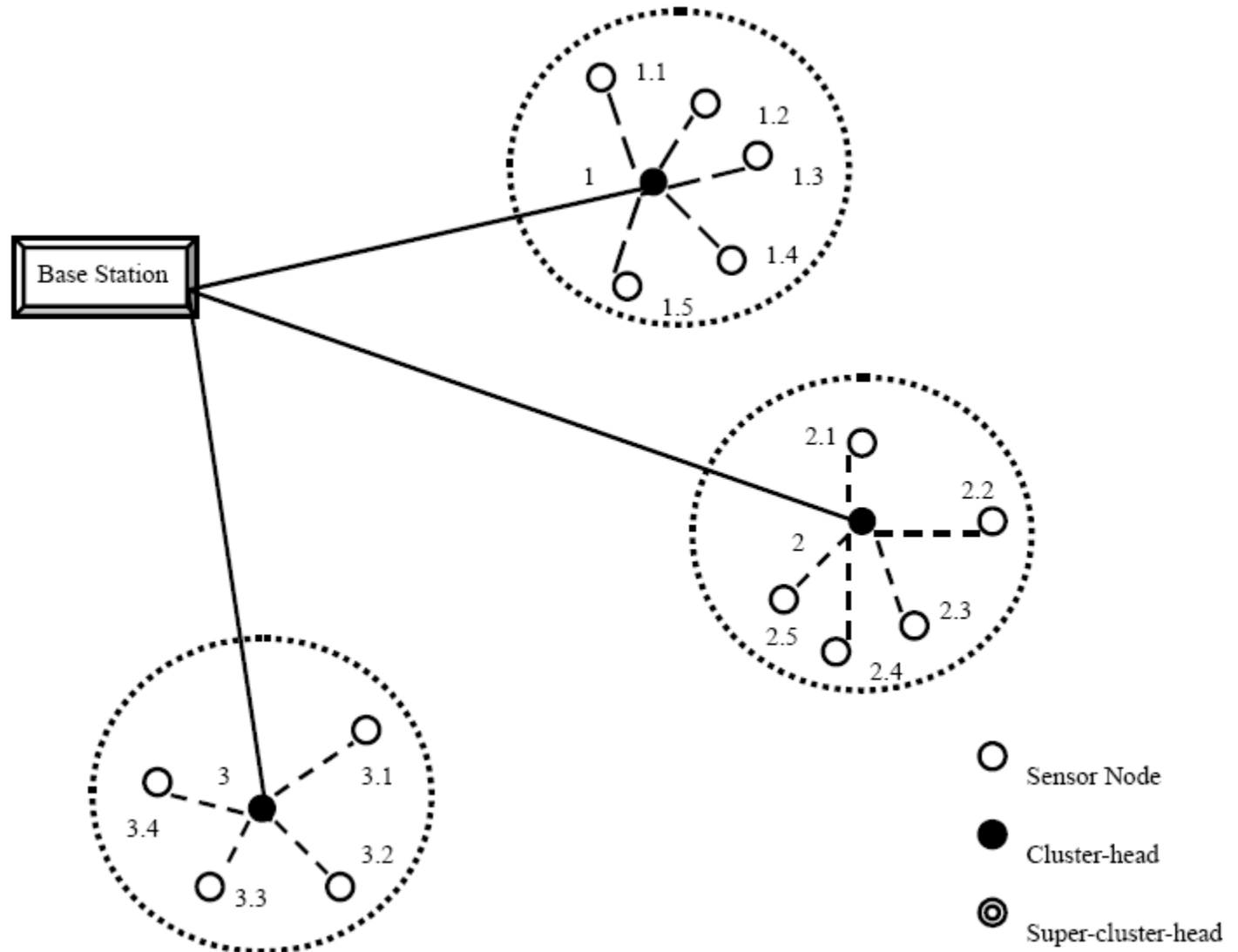


Multi-hop network

# Hierarchical Architecture



# Hierarchical Architecture



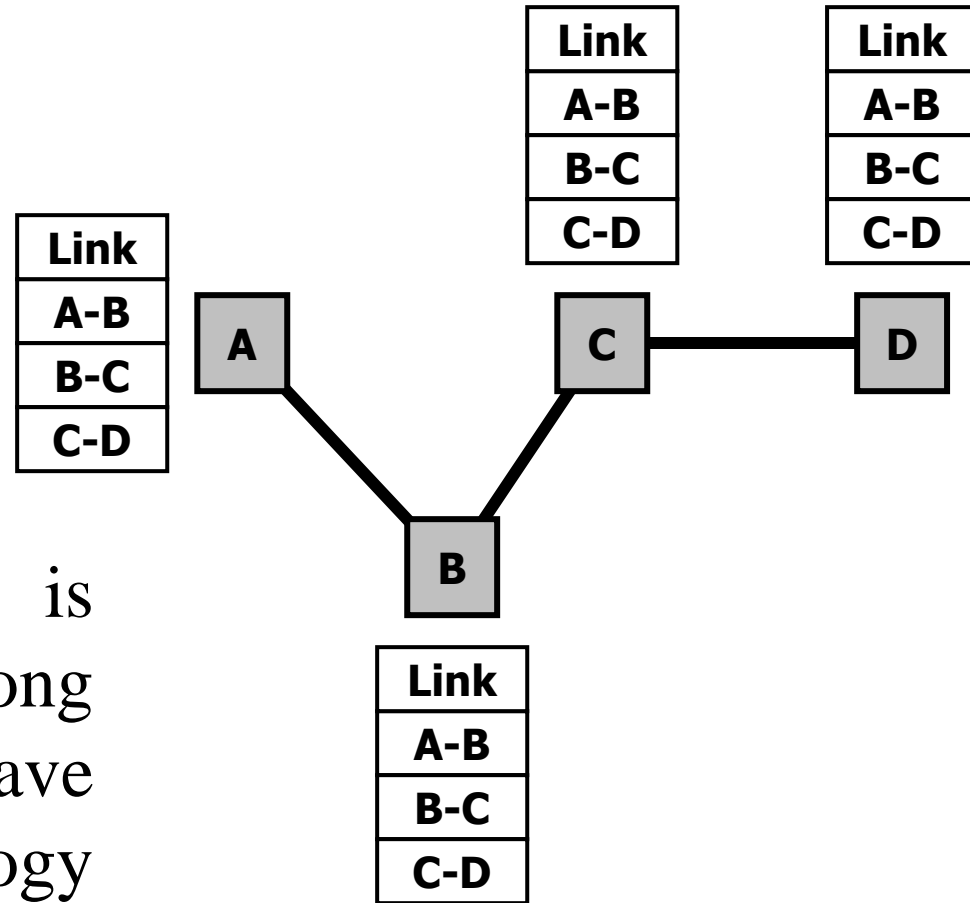


# Ad-Hoc Routing Protocol

- ❖ An ad-hoc routing protocol is a convention that controls how nodes decide which way to route packets between computing devices in a mobile ad-hoc network
- ❖ Foundation in most protocols: neighbor discovery
  - Nodes send periodic announcements as broadcast packets (beacon messages, alive messages, ...)
  - Can embed “neighbor table” into such messages; allows nodes to learn “2-hop neighborhood”
- ❖ Popular types of routing protocols:
  - Proactive
  - Reactive
  - Geographic

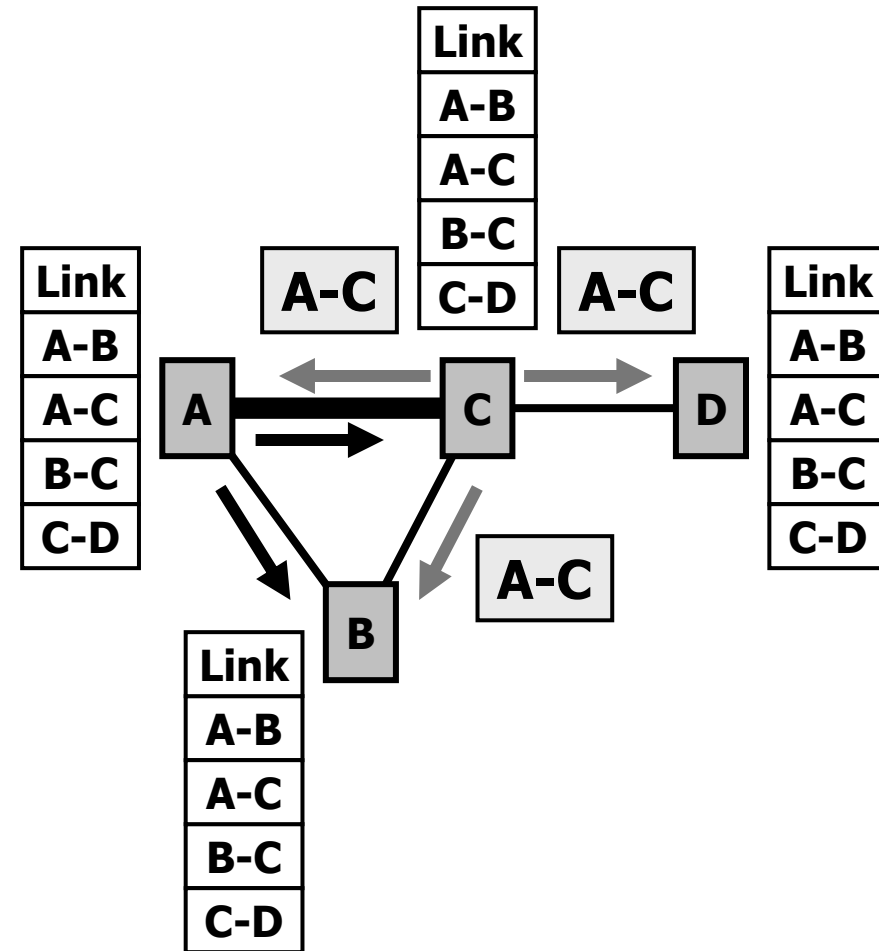
# Proactive: “Link-State” Algorithms

- ❖ Each node shares its link information so that all nodes can build a map of the full network topology
- Assuming the topology is stable for a sufficiently long period, all nodes will have the same topology information



# Proactive: “Link-State” Algorithms

- ❖ Link information is updated when a link changes state (goes up or down)
  - by sending small “hello” packets to neighbors
- ❖ Nodes A and C propagate the existence of link A-C to their neighbors and, eventually, to the entire network

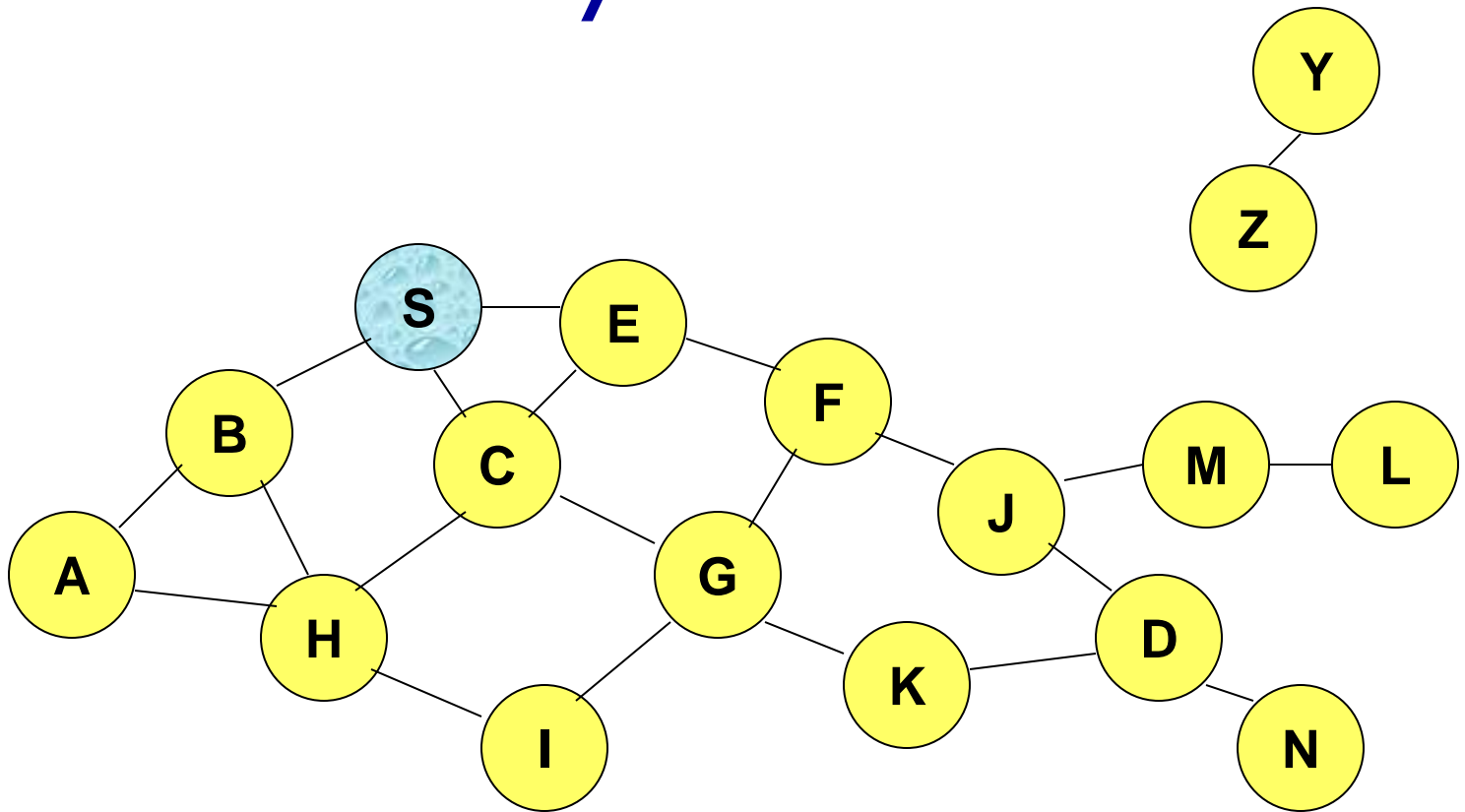


# Reactive: DSR

## Dynamic Source Routing

- ❖ Search for route when needed only
  - Search using **Route Request (RREQ)** broadcasts
  - Response using **Route Reply (RREP)** message
- ❖ Every message along route contains entire path to help intermediate nodes to decide what to do with message

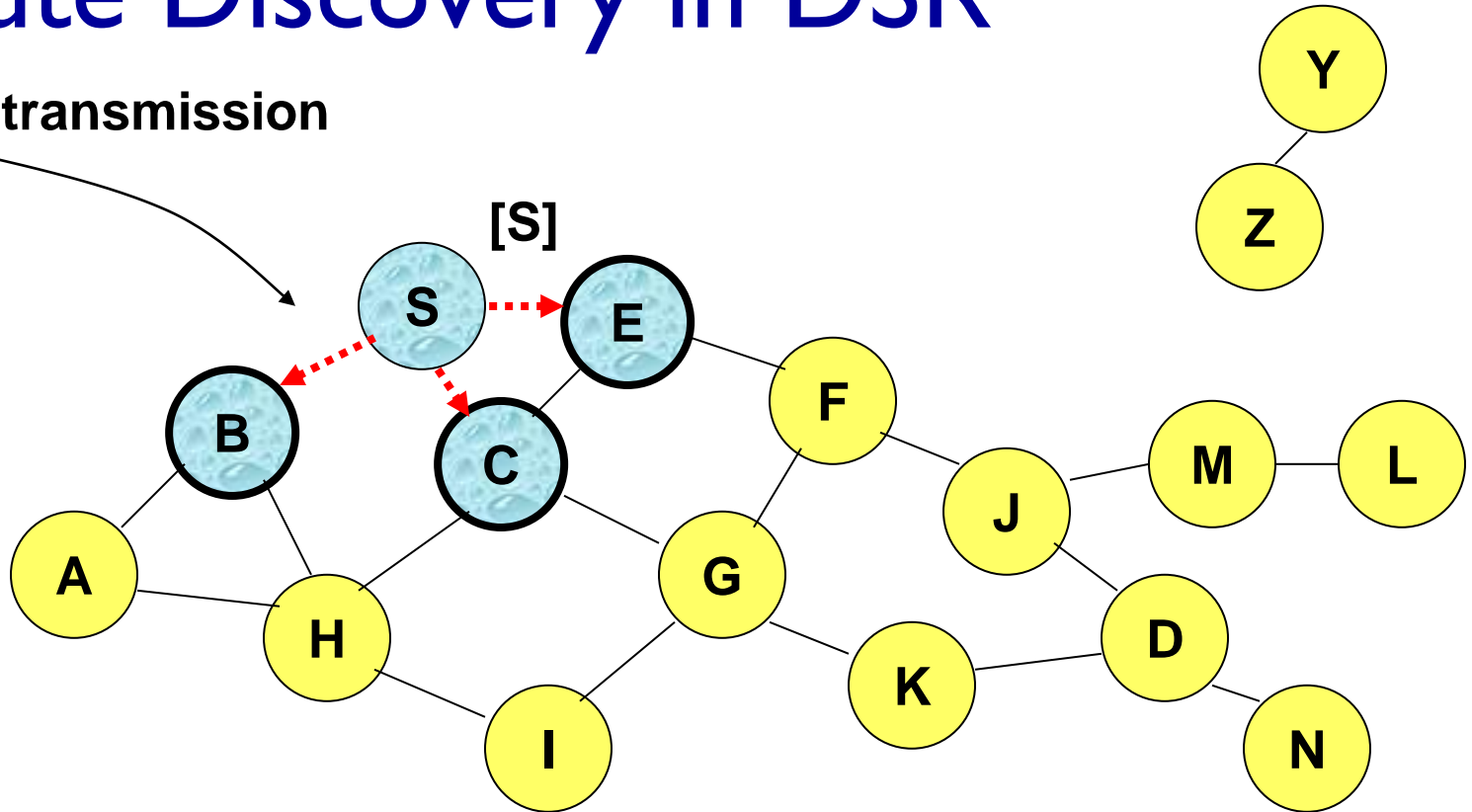
# Route Discovery in DSR



**Represents a node that has received RREQ for D from S**

# Route Discovery in DSR

Broadcast transmission

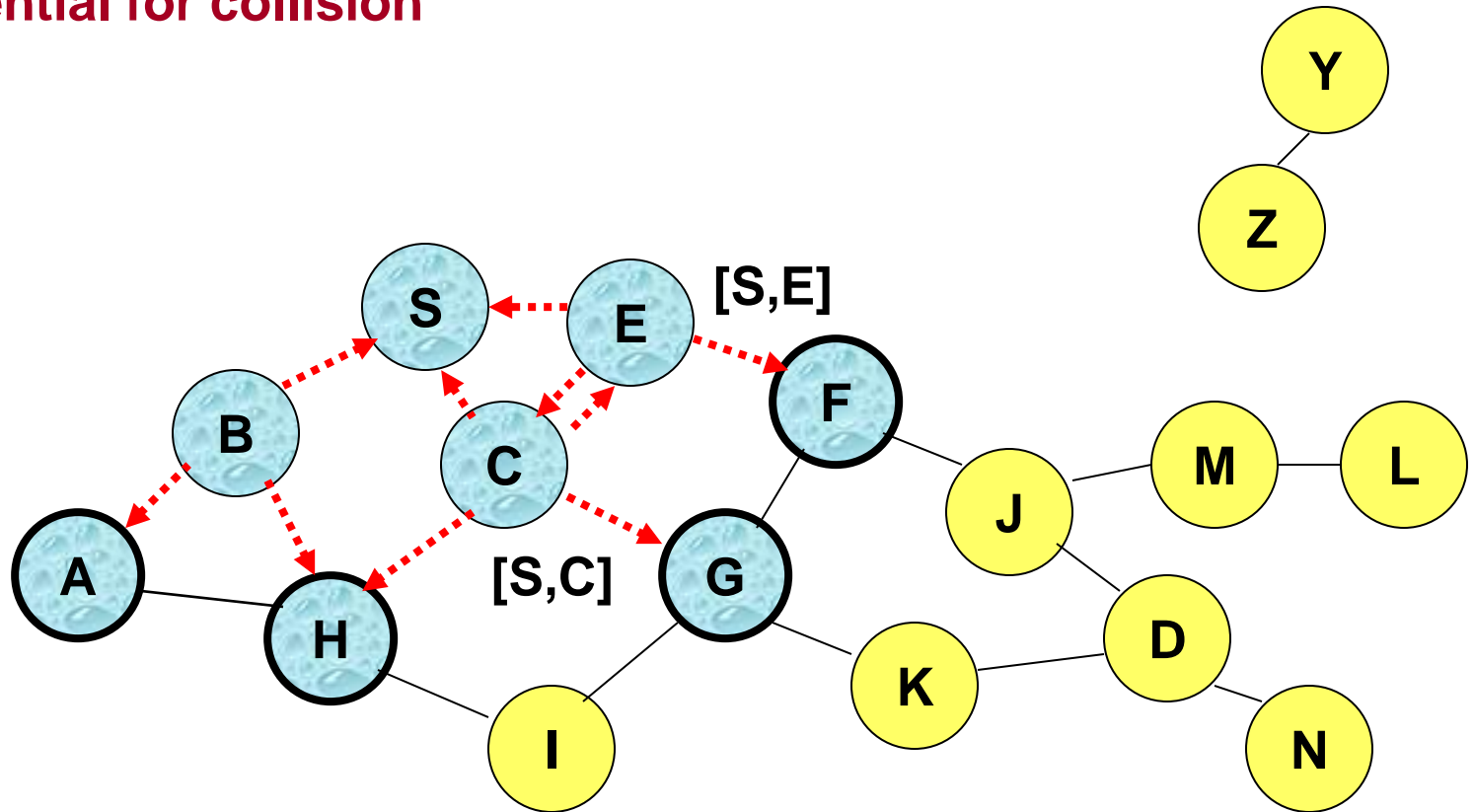


.....→ Represents transmission of RREQ

[X,Y] Represents list of identifiers appended to RREQ

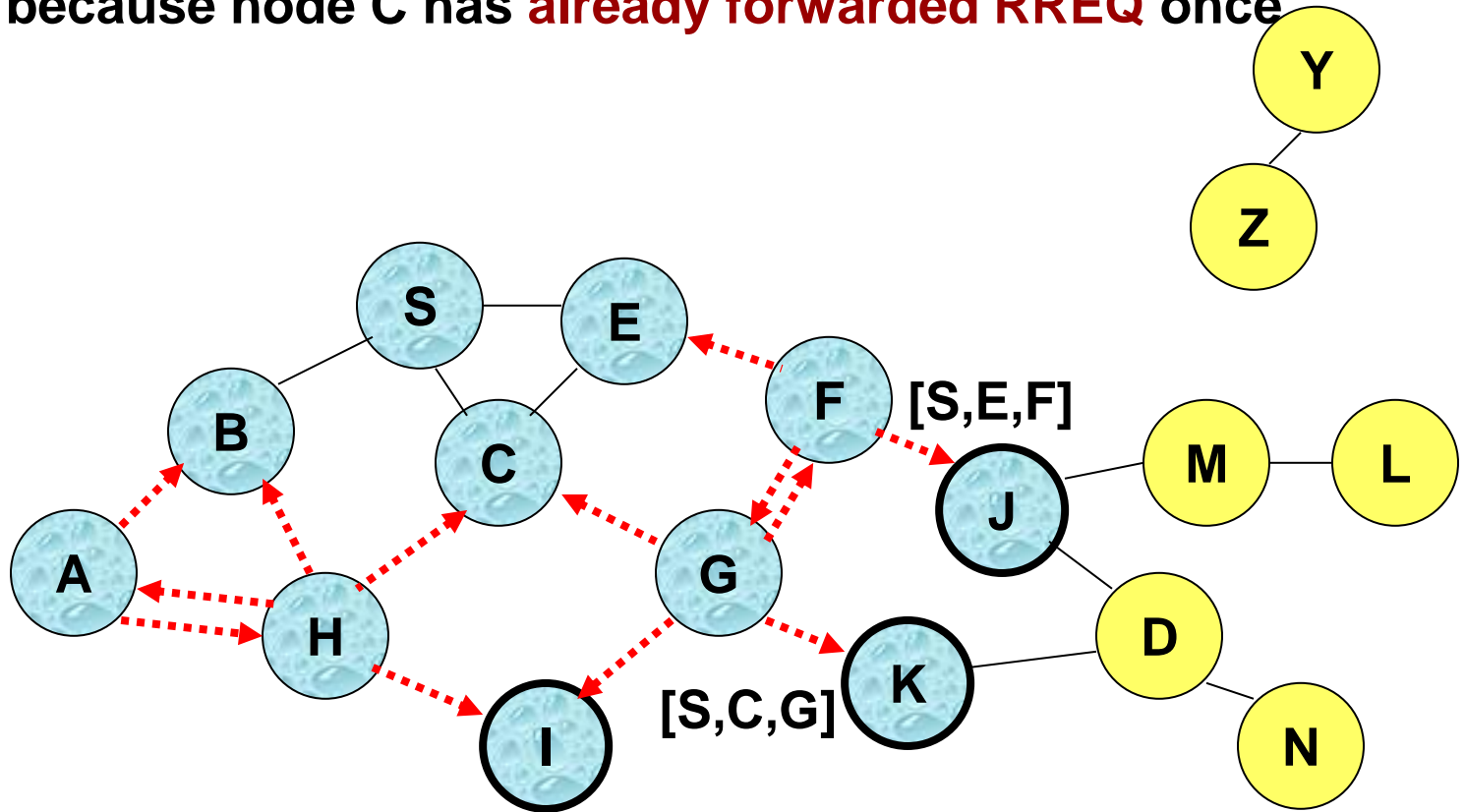
# Route Discovery in DSR

Node H receives packet RREQ from two neighbors:  
**potential for collision**



# Route Discovery in DSR

Node C receives RREQ from G and H, but does not forward it again, because node C has **already forwarded RREQ** once

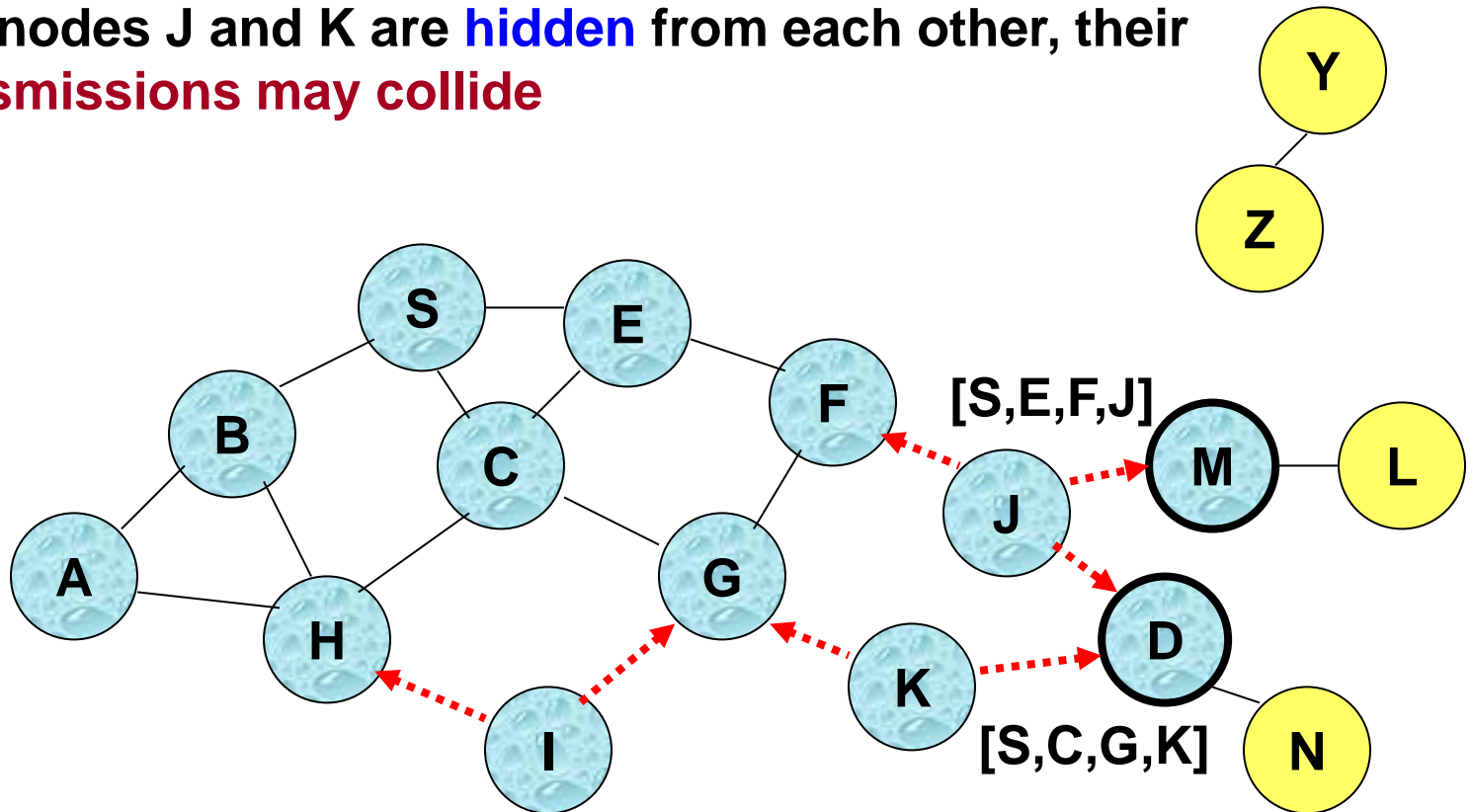




# Route Discovery in DSR

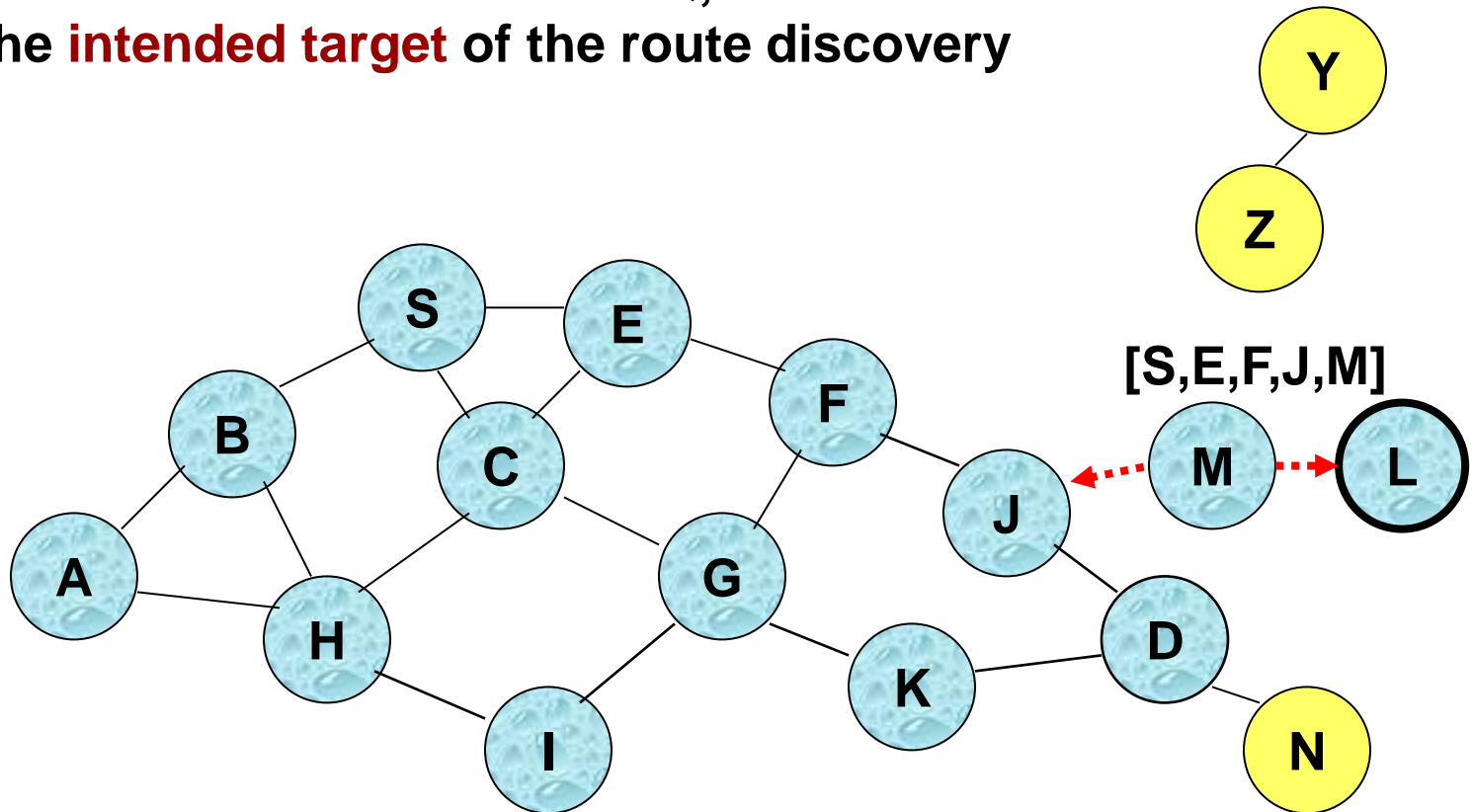
Nodes J and K both broadcast RREQ to node D

Since nodes J and K are **hidden** from each other, their transmissions may collide

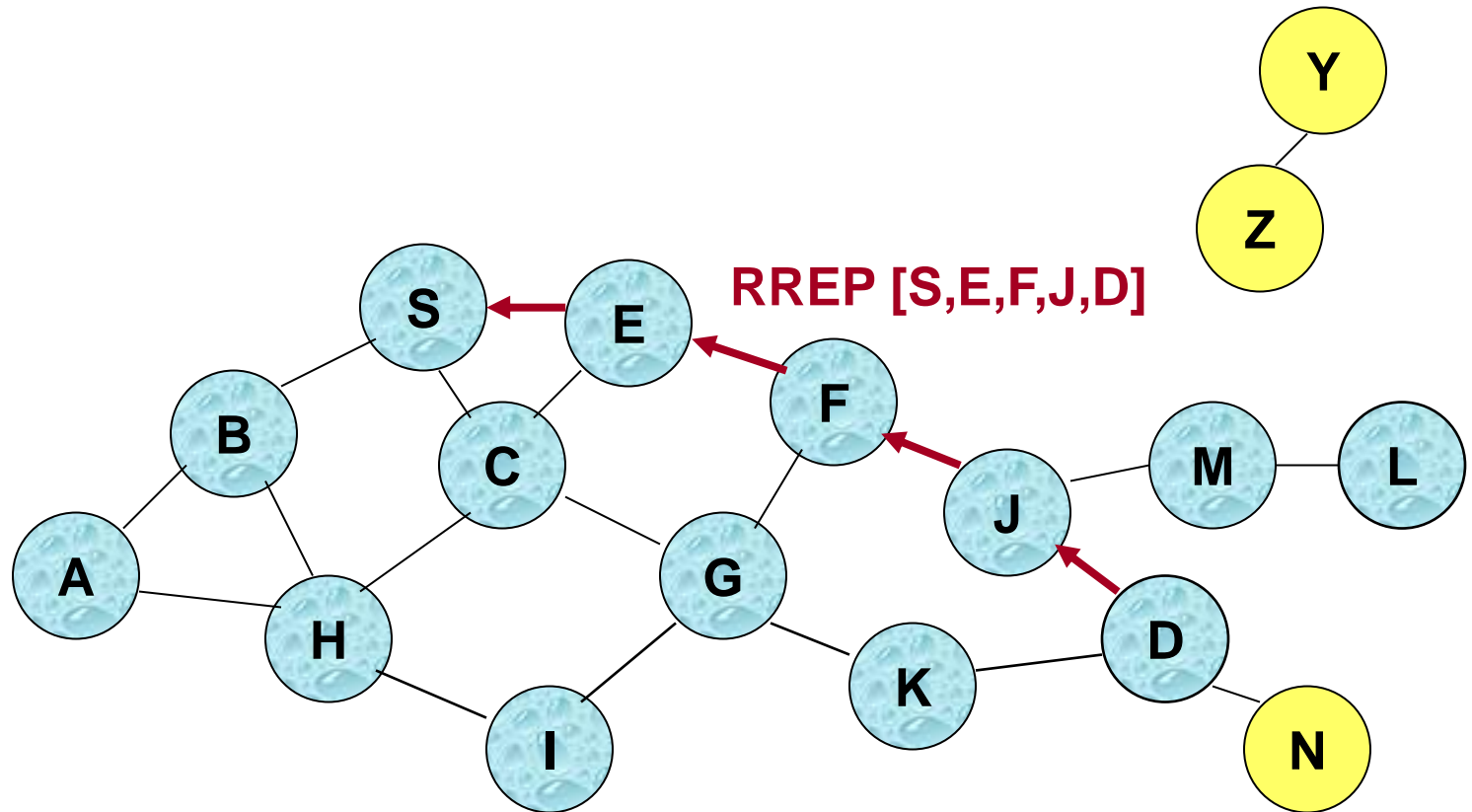


# Route Discovery in DSR

- Node D **does not forward** RREQ, because node D is the **intended target** of the route discovery



# Route Reply in DSR



← Represents RREP control message

# Route Reply in DSR

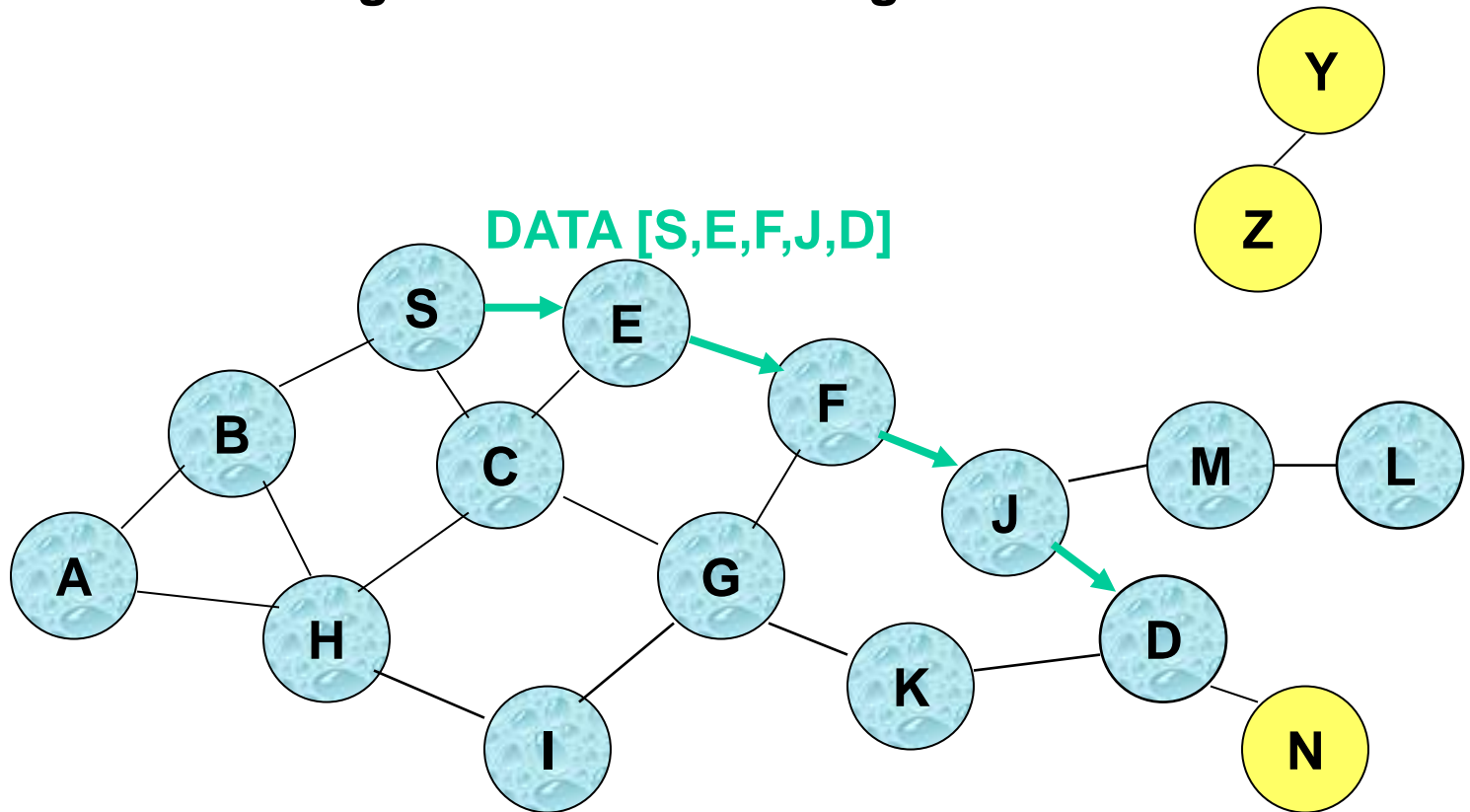
- ❖ Route Reply can be sent by reversing the route in Route Request (RREQ) only if links are guaranteed to be bidirectional
- ❖ One way to ensure this is to check, if the received RREQ was on a link that is known to be bi-directional, e.g.
  - If IEEE 802.11 MAC is used to send data, then links have to be bi-directional (since Ack is used)
- ❖ If unidirectional (asymmetric) links are allowed, then RREP may need a route discovery for S from node D
  - Route discovery not needed -> If node D already knows a route to node S
  - If a route discovery is initiated by D for a route to S, then the Route Reply is piggybacked on the Route Request from D.

# Route Reply in DSR

- ❖ Node S on receiving RREP, caches the route included in the RREP
- ❖ When node S sends a data packet to D, the entire route is included in the packet header
  - hence the name source routing
- ❖ Intermediate nodes use the source route included in a packet to determine to whom a packet should be forwarded

# Data Delivery in DSR

Packet header size grows with route length



# DSR Optimization: Route Caching

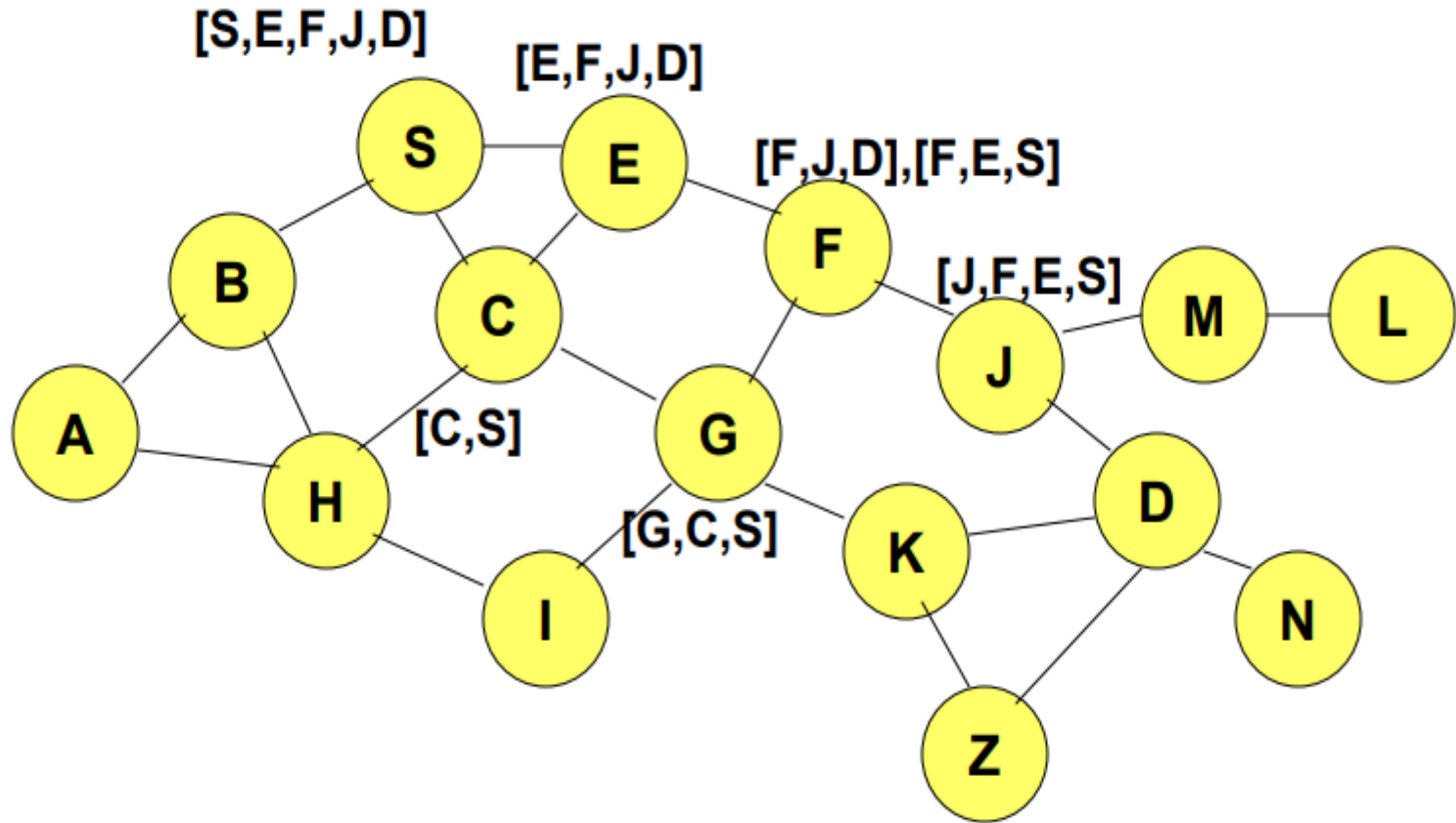
- ❖ Each node caches a new route it learns by any means
- ❖ When node S finds route [S,E,F,J,D] to node D, node S also learns route [S,E,F] to node F
- ❖ When node K receives Route Request [S,C,G] destined for node D, node K learns route [K,G,C,S] to node S
- ❖ When node F forwards Route Reply RREP [D,J,F,E,S], node F learns route [F, J, D] to node D
- ❖ When node E forwards Data [S,E,F,J,D] it learns route [E,F,J,D] to node D
- ❖ A node may also learn a route when it overhears Data packets

# DSR Optimization: Route Caching

- ❖ When node S learns that a route to node D is broken,
  - Can use another route from its local cache, if such a route to D exists in its cache.
  - Otherwise, node S initiates route discovery by sending a route request
- ❖ Node X on receiving a Route Request for some node D can send a Route Reply if node X knows a route to node D
- ❖ Use of route cache
  - can speed up route discovery
  - can reduce propagation of route requests

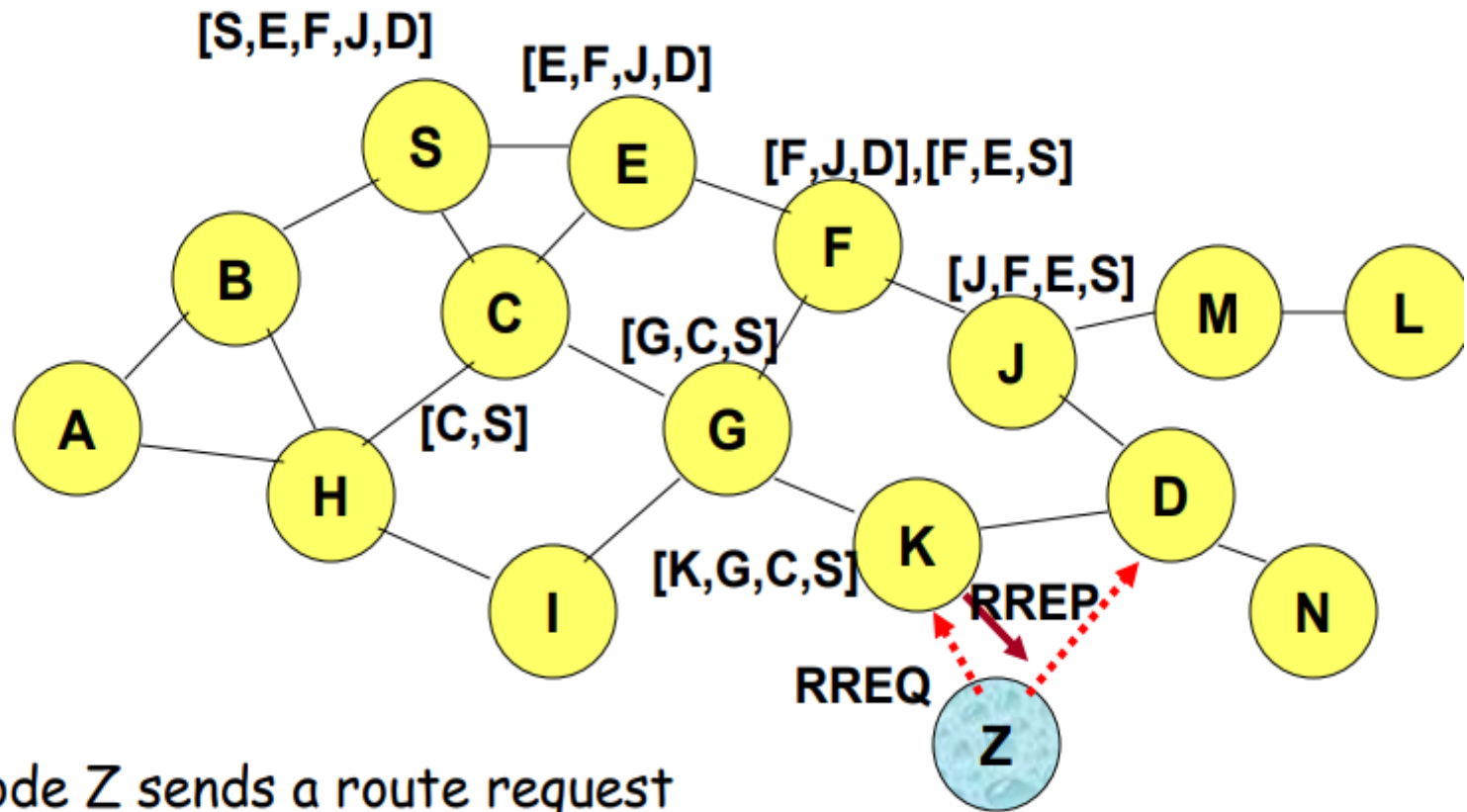


# Use of Route Caching: Example



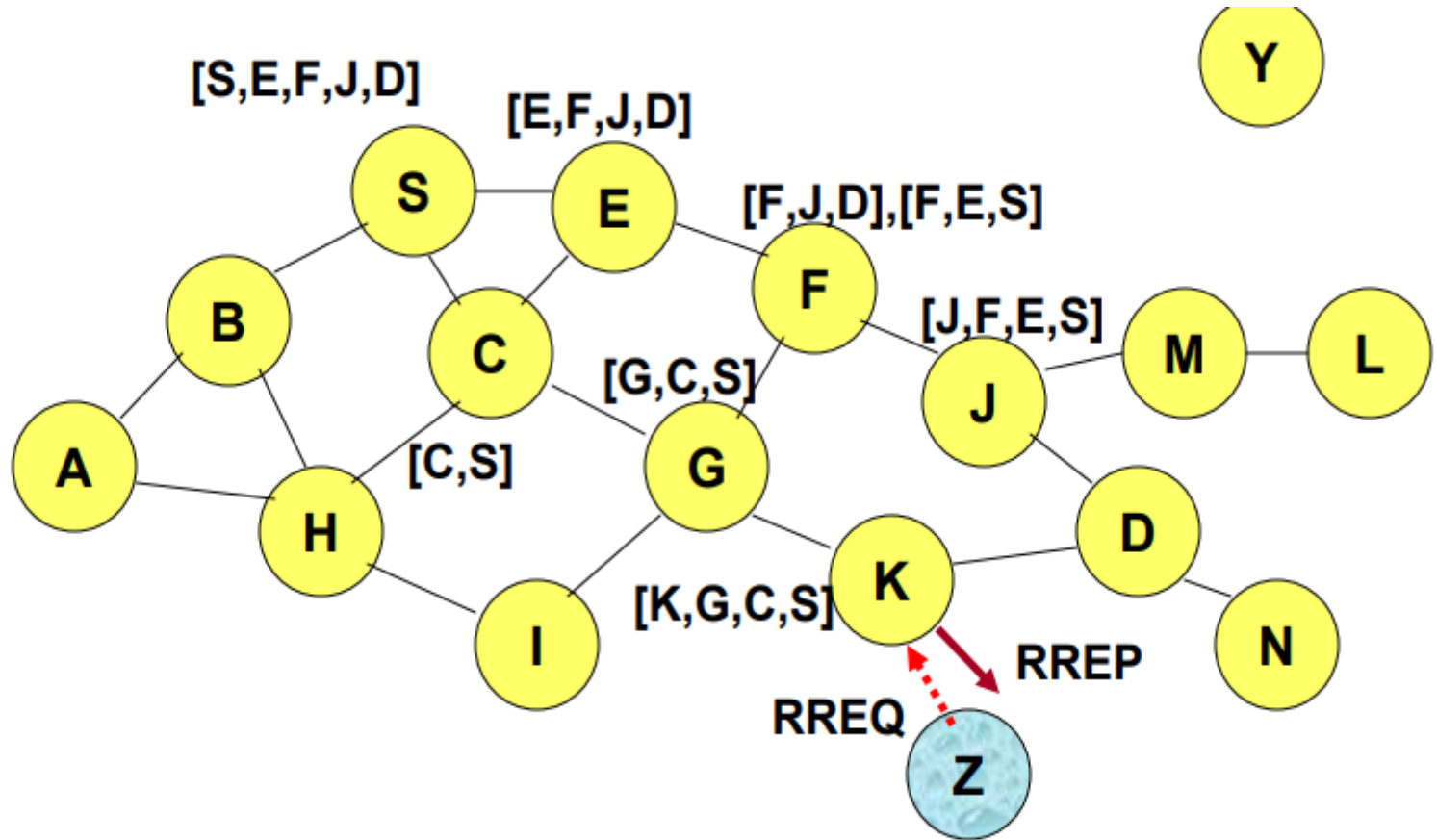
[P,Q,R] Represents cached route at a node  
(DSR maintains the cached routes in a tree format)

# Route Caching benefits: speed up of Route Discovery



When node Z sends a route request for node C, node K sends back a route reply [Z,K,G,C] to node Z using a locally cached route

# Route Caching benefits: Reduction in propagation of RREQs



- ❑ Route Reply (RREP) from node K **limits flooding** of RREQ.

# DSR : Explore the following

- ❖ Duplication of Route hops
- ❖ Route Maintenance
  - Preventing Route Reply Storms
  - Route Request hop limits
  - Packet Salvaging
  - Automatic Route Shortening
  - Increased spreading of Route Error messages

# DSR : Advantages

- ❖ Routes maintained only between nodes who need to communicate
- ❖ reduces overhead of route maintenance
- ❖ Route caching can further reduce route discovery overhead
- ❖ A single route discovery may yield many routes to the destination, due to intermediate nodes replying from local caches

# DSR : Disadvantages

- ❖ Packet header size grows with route length due to source routing
- ❖ Flood of route requests may potentially reach all nodes in the network
- ❖ Care must be taken to avoid collisions between route requests propagated by neighboring nodes
  - insertion of random delays before forwarding RREQ
- ❖ Increased contention if too many route replies come back due to nodes replying using their local cache
  - Route Reply Storm problem

# Proactive vs Reactive

## ❖ Reactive:

- Only establish/maintain routes between nodes needed them (in contrast: tables store ALL routes)
- Store entire route in each message; message size grows with route length
- Route requests cause “flooding”

## ❖ Proactive:

- Route information always available; no need to search for route (but route information can be outdated)
- Continuous exchange of route change updates

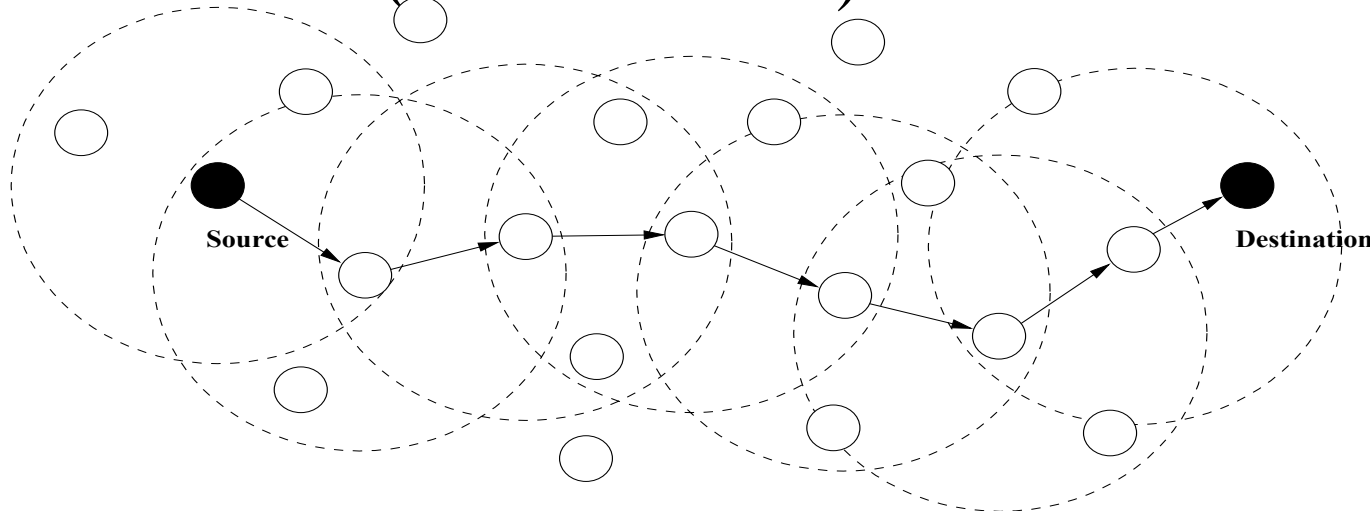
# Geographic Routing

- ❖ Nodes use location information to make routing decisions
  - sender must know the locations of itself, the destination, and its neighbors
  - location information can be queried or obtained from a **location broker**
  - location information can come from GPS (Global Positioning System) or some other form of positioning technology



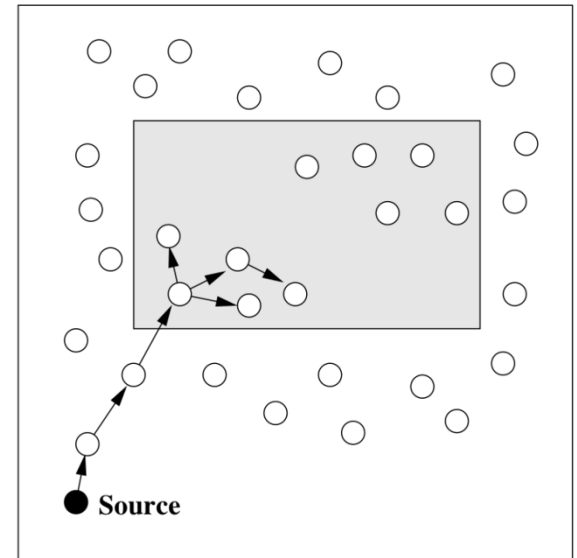
# Unicast Location-Based Routing

- ❖ One single destination
- ❖ Each forwarding node makes localized decision based on the location of the destination and the node's neighbors (**greedy forwarding**)
- ❖ **Challenge:** packet may arrive at a node without neighbors that could bring packet closer to the destination (**voids or holes**)



# Geocasting

- ❖ Packet is sent to all or some nodes within specific geographic region
- ❖ **Example:** query sent to all sensors within geographic area of interest
- ❖ **Routing challenge:**
  - propagate a packet near the target region (similar to unicast routing)
  - distribute packet within the target region (similar to flooding)



# Design Challenges

## ❖ Heterogeneity

- The devices deployed may be of various types and need to collaborate with each other.

## ❖ Distributed Processing

- The algorithms need to be centralized as the processing is carried out on different nodes.

## ❖ Low Bandwidth Communication

- The data should be transferred efficiently between sensors

# Design Challenges

## ❖ Large Scale Coordination

- The sensors need to coordinate with each other to produce required results.

## ❖ Utilization of Sensors

- The sensors should be utilized in a ways that produce the maximum performance and use less energy.

## ❖ Real Time Computation

- The computation should be done quickly as new data is always being generated.

# Applications of WSN

- ❖ Military and national security application
- ❖ Environment monitoring
- ❖ Medical application
- ❖ Home and Office Applications
- ❖ Automotive Applications