

Q: Show the CYK Algorithm with the following CNF Grammar?

$S \rightarrow AB \mid BC$

$A \rightarrow BA \mid a$

$B \rightarrow CC \mid b$

$C \rightarrow AB \mid a$

Is "bbaaa" belongs to the language of Grammar or not using CYK Algorithm. Yes!

Solution: This is because start state is available in top cell.

Substring Length ↑	$j=5$	5-1 <u>bbaaa</u> (S, A, C)				
	$j=4$	4-1 <u>bbaa</u> ϕ	4-2 <u>baaa</u> S, A, C			
	$j=3$	3-1 <u>bba</u> A	3-2 <u>baa</u> ϕ	3-3 <u>aaa</u> S, A, C		
	$j=2$	2-1 <u>bb</u> ϕ	2-2 <u>ba</u> S, A	2-3 <u>aa</u> B	2-4 <u>aa</u> B	
	$j=1$	1-1 <u>b</u> B	1-2 <u>b</u> B	1-3 <u>a</u> A, C	1-4 <u>a</u> A, C	1-5 <u>a</u> A, C
		b	b	a	a	a

$j=2$ ba BA, BC aa AA, AC, CA, CC
 $\downarrow \quad \downarrow \quad \downarrow \quad \downarrow$
 $A \quad U \quad S \quad \phi \quad \phi \quad \phi \quad B$

$j=3$ bba
 $\bullet \begin{matrix} b & ba \\ \downarrow & \downarrow \\ B & \times S, A \end{matrix} = \begin{matrix} \phi & A \\ \downarrow & \downarrow \\ BS, & BA \end{matrix}$
 $\bullet \begin{matrix} bb & a \\ \downarrow & \downarrow \\ \phi & \times \end{matrix} = \phi$
baa
 $\bullet \begin{matrix} b & aa \\ \downarrow & \downarrow \\ B & \times B \end{matrix} = \begin{matrix} \phi \\ \downarrow \\ BB \end{matrix}$
 $\bullet \begin{matrix} ba & a \\ \downarrow & \downarrow \\ S, A & \times A, C \end{matrix} = \begin{matrix} \phi & \phi & \phi & \phi \\ \downarrow & \downarrow & \downarrow & \downarrow \\ SA, & SC, & AA, & AC \end{matrix}$
aaa
 $\bullet \begin{matrix} a & aa \\ \downarrow & \downarrow \\ A, C & \times B \end{matrix} = \begin{matrix} S, C & \phi \\ \downarrow & \downarrow \\ AB, & CB \end{matrix}$
 $\bullet \begin{matrix} aa & a \\ \downarrow & \downarrow \\ B & \times A, C \end{matrix} = \begin{matrix} A & S \\ \downarrow & \downarrow \\ BA, & BC \end{matrix}$

$j=4$ bbaa
 $\bullet \begin{matrix} b & baa \\ \downarrow & \downarrow \\ B & \times \phi \end{matrix} = \phi$
 $\bullet \begin{matrix} bb & aa \\ \downarrow & \downarrow \\ \phi & \times \end{matrix} = \phi$
 $\bullet \begin{matrix} bba & a \\ \downarrow & \downarrow \\ A & \times A, C \end{matrix} = \begin{matrix} \phi & \phi \\ \downarrow & \downarrow \\ AA, & AC \end{matrix}$
baaa
 $\bullet \begin{matrix} b & aaa \\ \downarrow & \downarrow \\ B & \times S, A, C \end{matrix} = \begin{matrix} \phi & A & S \\ \downarrow & \downarrow & \downarrow \\ BS, & BA, & BC \end{matrix}$
 $\bullet \begin{matrix} ba & aa \\ \downarrow & \downarrow \\ S, A & \times B \end{matrix} = \begin{matrix} \phi & S, C \\ \downarrow & \downarrow \\ SB, & AB \end{matrix}$
 $\bullet \begin{matrix} baa & a \\ \downarrow & \downarrow \\ \phi & \times \end{matrix} = \phi$

$j=5$
 $\bullet \begin{matrix} b & baaa \\ \downarrow & \downarrow \\ B & \times S, A, C \end{matrix} = \begin{matrix} \phi & A & S \\ \downarrow & \downarrow & \downarrow \\ BS, & BA, & BC \end{matrix}$
 $\bullet \begin{matrix} bb & aaa \\ \downarrow & \downarrow \\ \phi & \times \end{matrix} = \phi$
 $\bullet \begin{matrix} bba & aa \\ \downarrow & \downarrow \\ A & \times B \end{matrix} = \begin{matrix} S, C \\ \downarrow \\ AB \end{matrix}$
 $\bullet \begin{matrix} bbaa & a \\ \downarrow & \downarrow \\ \phi & \times \end{matrix} = \phi$

Q: Show the CYK Algorithm with the following CNF Grammar?

$S \rightarrow AB \mid BC$
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 $C \rightarrow AB \mid a$

Is "bbbaa" belongs to the language of Grammar or not using CYK Algorithm.

Solution:

No, it doesn't belongs to this grammar.

5	ϕ	—	—	—	—
4	A	ϕ	—	—	—
3	ϕ	A	ϕ	—	—
2	ϕ	ϕ	A, S	B	—
1	B	B	B	A, C	A, C
	b	b	b	a	a

b b b a
 $b \text{ bba} = B A = A$
 $bb \text{ ba} = \phi$
 $bbb \text{ a} = \phi$

b b a a
 $b \text{ baa} = B \phi$
 $bb \text{ aa} = \phi$
 $bba \text{ a} = A, (AC)$
 $AA = \phi$
 $AC = \phi$

b b a
 $b \text{ ba} = B A, S$
 $bb \text{ a} = \phi$
 BA, BS

b b b
 $B \phi$
 b b
 AA
 AC
 CA
 CC

b b b a a
 $b \text{ bbaa} = B \phi$
 $bb \text{ baa} = \phi$
 $bbb \text{ aa} = \phi$
 $bbba \text{ a} = A, AC$

AA
 AC

b a a
 $ba \text{ a} = A, S$
 $b \text{ aa} = B B$
 $AA \phi$
 $AC \phi$
 $SA \phi$
 $SC \phi$
 $BB \phi$

Q: Show the CYK Algorithm with the following

$S \rightarrow AB \mid BC$
 $A \rightarrow BA \mid a$
 $B \rightarrow CC \mid b$
 $C \rightarrow AB \mid a$

Is "ababa" belongs to the language of Grammar or not using CYK Algorithm.

Solution:

\rightarrow As S is here so ababa belongs to language

j=5	S, C, A				
j=4	B	B			
j=3	B	S, C	B		
j=2	S, C	A, S	S, C	A, S	
j=1	A, C	B	A, C	B	A, C
	a	b	a	b	a

j=2

$ab \rightarrow A, C \times B$
 $ab \rightarrow AB, ACB$

$ba \rightarrow B \times A, C$
 $\rightarrow BA, BC$

$ab \rightarrow A, C \times B$
 AB, CB

$ba \rightarrow B \times A, C$
 $\rightarrow BA, BC$

V=3

$aba \rightarrow \begin{cases} a \rightarrow A, C \\ ba \rightarrow A, C \end{cases} \rightarrow \begin{cases} AA, AS \\ CA, CS \end{cases}$
 $\rightarrow ab \rightarrow S, C$
 $a \rightarrow A, C$
 $\rightarrow \begin{cases} SA, SC \\ CA, CC \end{cases}$

$S \rightarrow AB \mid BC$
 $A \rightarrow BA \mid a$
 $B \rightarrow CC \mid b$
 $C \rightarrow AB \mid a$

Is "aaaa" belongs to the language of Grammar or not using CYK Algorithm.

Solution:

j=5	S, C, A				
j=4	B	B			
j=3	S, C, A	S, C, A	S, C, A		
j=2	B	B	B	B	
j=1	A, C	A, C	A, C	A, C	A, C

$a, aaaa \rightarrow (A, C) \times B$
 $aa, aaa \rightarrow B \times (S, C, A)$
 $aaa, aa \rightarrow (S, C, A) \times B$
 $aaaa, a \rightarrow B \times (A, C)$
 AB, CB
 (S, C)
 BC, BC, BA
 $\times (S) (A)$
 SB, CB, AB
 $\times \times (S, C)$
 BA, BC
 $(A) (S)$

aaaa

Yes it belongs to the grammar as top most box has starting state in it

$a, aa \rightarrow (A, C) \times B \rightarrow AB, CB$
 $aa, a \rightarrow B \times (A, C) = (BA), (BC)$

aaa
 $a, aa \rightarrow (A, C) \times B \rightarrow AB, CB$
 $aa, a \rightarrow B \times (A, C) \rightarrow BA, BC$

aaaa

$a, aaa \rightarrow (A, C) \times (S, C, A) \rightarrow A^x S^x, AC^x, AA^x, CB^x, CC^x, CA^x$
 $aa, aa \rightarrow B \times B \rightarrow BB^x$
 $aaa, a \rightarrow (S, C, A) \times (A, C) \rightarrow SA^x, SC^x, CA^x, CC^x, AA^x, AC^x$

$aaaa$
 $a, aaa \rightarrow (A, C), (S, C, A)$
 $aa, aa \rightarrow B, B$
 $aaa, a \rightarrow (S, C, A), (A, C)$

aa
 $(A, C) \times (A, C)$
 AA, AC, CA, CC
 $\times \times \times B$

aa
 $(A, C) \times (A, C)$
 AA, AC, CA, CC
 $\times \times \times B$

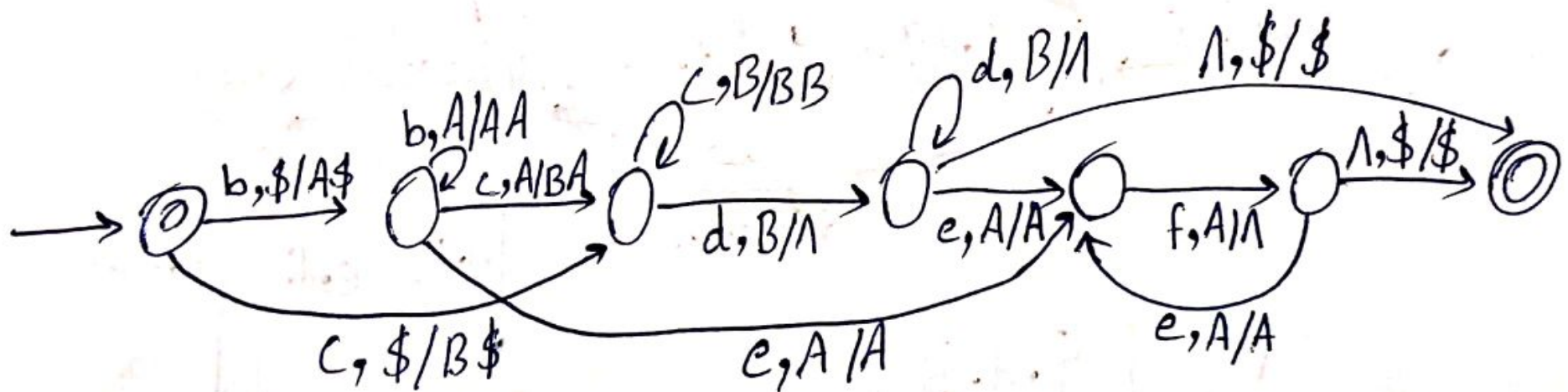
$$b^m c^n d^n (ef)^m$$

Λ

$$m \geq 1 \quad n = 0$$

$$m = 0 \quad n \geq 1$$

$$m \geq 1 \quad n \geq 1$$



Q1: Design a PDA for the following Language.

$$m \geq 0, n \geq 0$$

$$L = \{(ab)^m c^n d^n e^m; m, n \geq 0\}$$

$$a^m b^n c^n d^n e^m = \text{least} = 1 - (m \geq 0, n \geq 0)$$

$$= abcde (m \geq 0, n \geq 0)$$

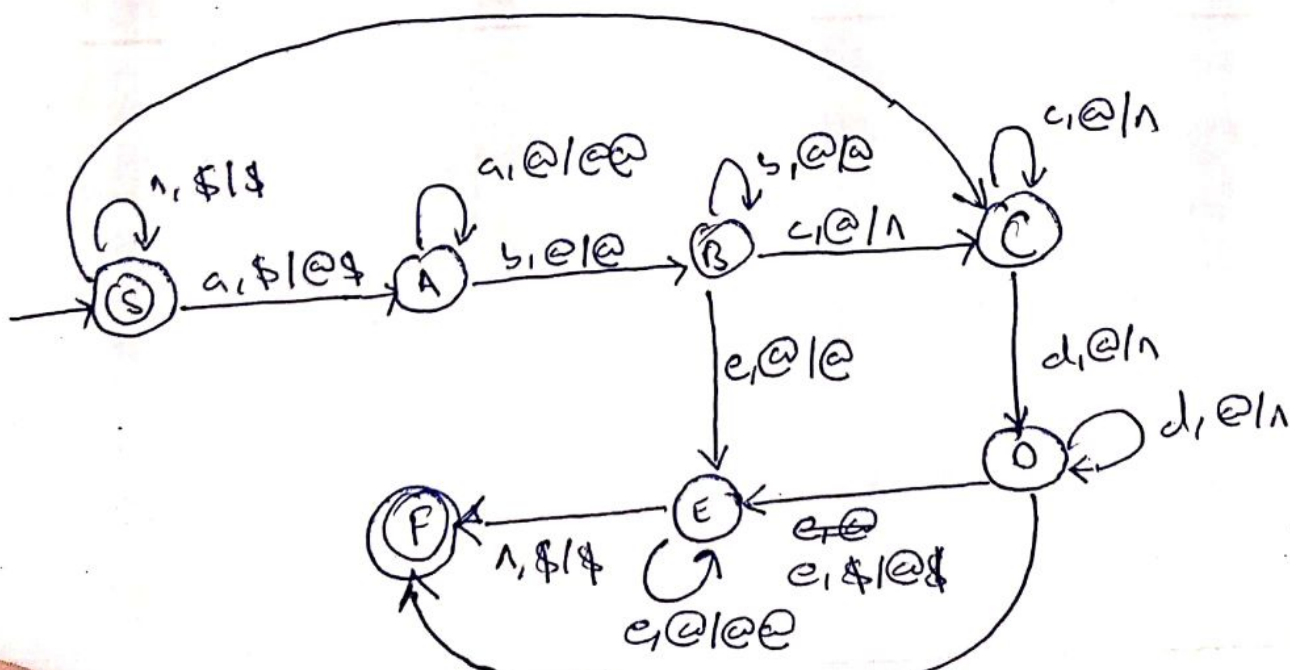
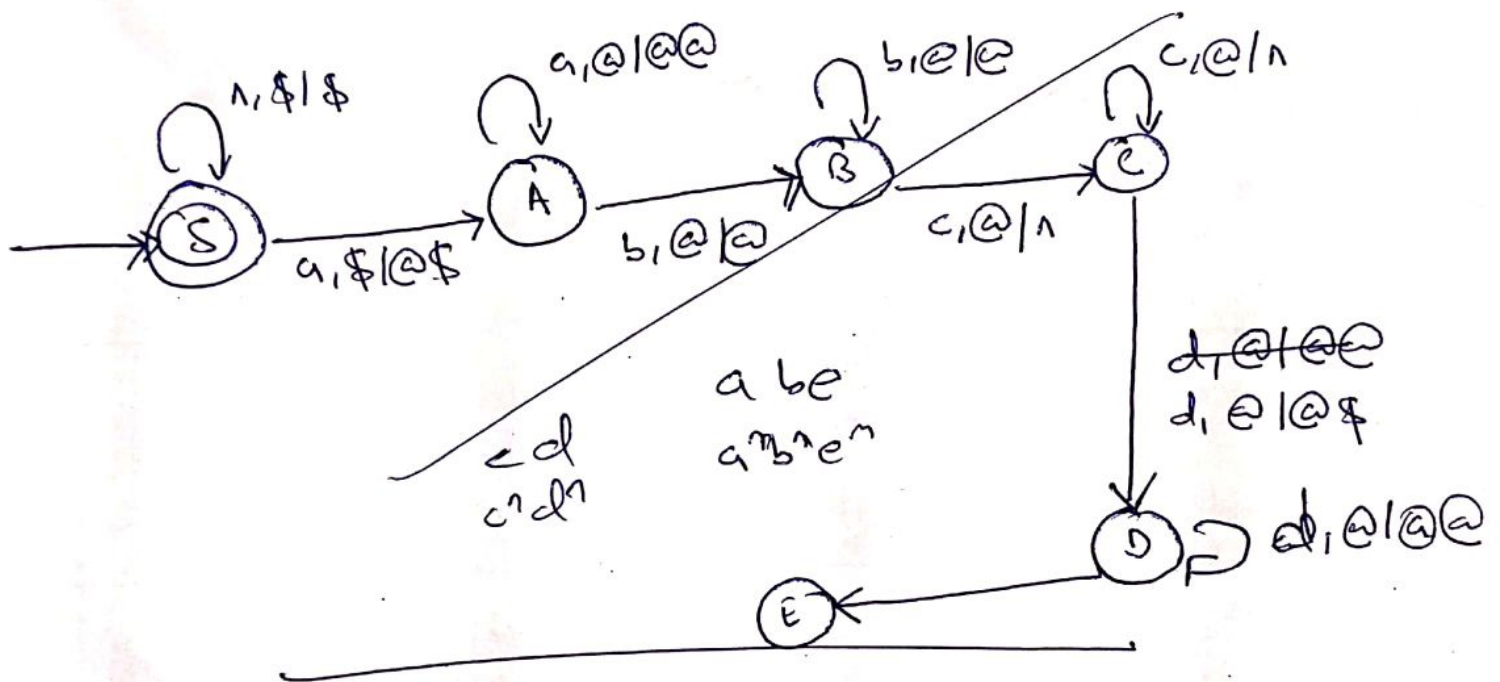
$$= cd (m \geq 0, n \geq 0)$$

$$= abe - (m \geq 0, n \geq 0)$$

$$\{a, b, c\} = \text{push}$$

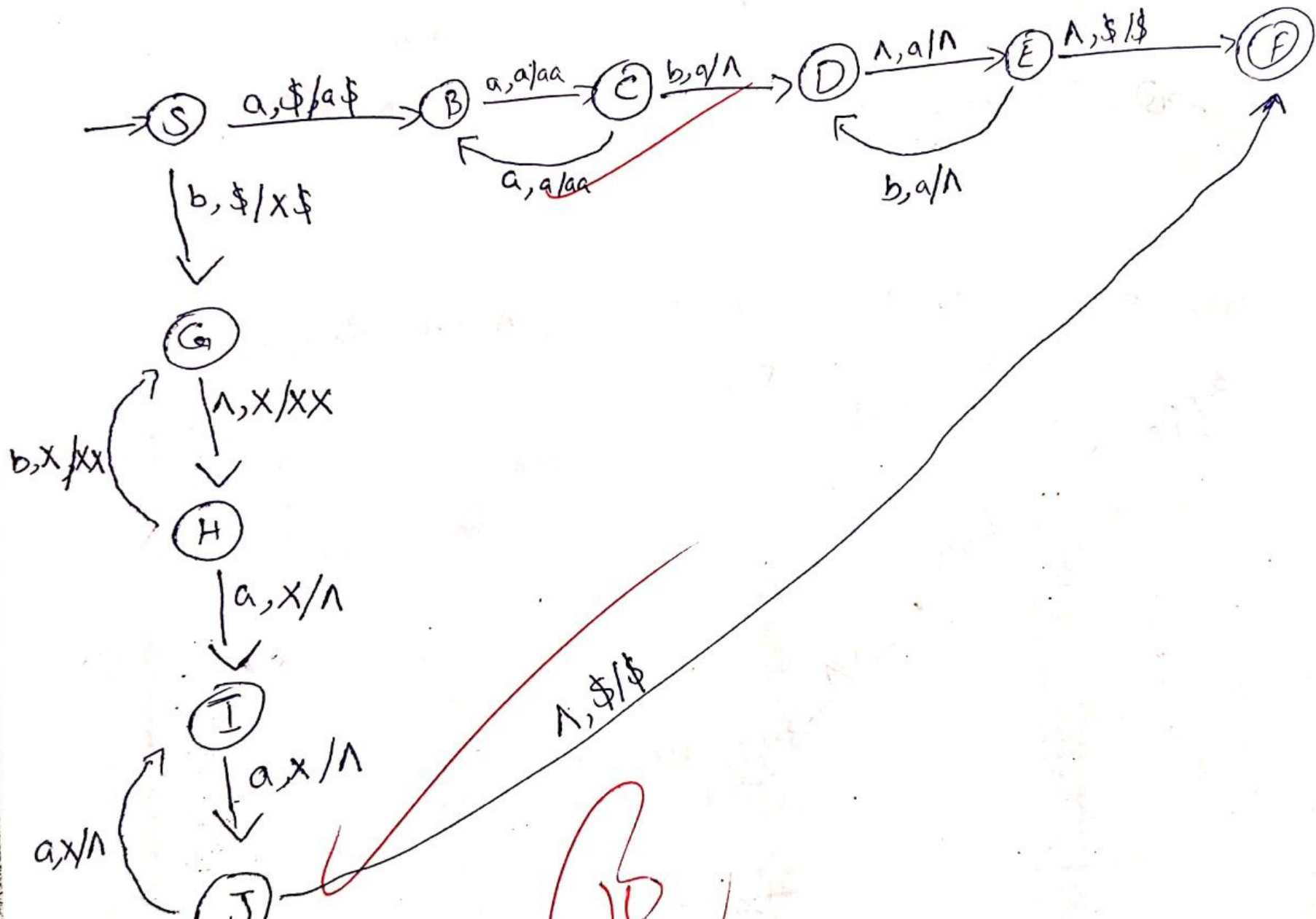
$$\{c, d\} = \text{pop}$$

$$\Gamma = \{\$, @\}$$



Q1: Design a PDA for the following Language.

L contains twice more a 's than b 's language over $\{a, b\}$



L contains 2 more a's than b's language over $\{a, b\}$

$$\Gamma = \{A, B, Z_0\}$$

Z_0 is our stack element

