

Understand Security Capabilities of Information Systems

The security capabilities of information systems include memory protection, virtualization, Trusted Platform Module (TPM), interfaces, and fault tolerance. It is important to carefully assess each aspect of the infrastructure to ensure that it sufficiently supports security. Without an understanding of the security capabilities of information systems, it is impossible to evaluate them, nor is it possible to implement them properly.

Memory Protection

Memory protection is a core security component that must be designed and implemented into an operating system. It must be enforced regardless of the programs executing in the system. Otherwise instability, violation of integrity, denial of service, and disclosure are likely results. Memory protection is used to prevent an active process from interacting with an area of memory that was not specifically assigned or allocated to it.

Memory protection is discussed throughout Chapter 9 in relation to the topics of isolation, virtual memory, segmentation, memory management, and protection rings.

Meltdown and Spectre

In late 2017, two significant memory errors were discovered. These issues were given the names Meltdown and Spectre. These problems arise from the methods used by modern CPUs to predict future instructions to optimize performance. This can enable a processor to seemingly make reliable predictions about what code to retrieve or process even before requested. However, when the speculative execution is wrong, the procedure is not completely reversed (i.e., not every incorrect predicted step is undone). This can result in some data remnants being left behind in memory in

an unprotected state.

Meltdown is an exploitation that can allow for the reading of private kernel memory contents by a nonprivileged process. Spectre can enable the wholesale theft of memory contents from other running applications. An astoundingly wide range of processors are vulnerable to one or both of these exploits. While two different issues, they were discovered nearly concurrently and made public at the same time. By the time of the publication of this book, patches are likely to be available to address these issues in existing hardware, and future processors should have native mechanisms to prevent such exploitations.

For a thorough discussion of these concerns, please listen to the Security Now podcast or read the show notes of episodes #645, “The Speculation Meltdown”; #646, “InSpectre”; and #648, “Post Spectre?” at <https://www.grc.com/securitynow.htm>.

Virtualization

Virtualization technology is used to host one or more operating systems within the memory of a single host computer. This mechanism allows virtually any OS to operate on any hardware. It also allows multiple OSs to work simultaneously on the same hardware. Common examples include VMware Workstation Pro, VMware vSphere and vSphere Hypervisor, VMware Fusion for Mac, Microsoft Hyper-V, Oracle VirtualBox, XenServer, and Parallels Desktop for Mac.

Virtualization has several benefits, such as being able to launch individual instances of servers or services as needed, real-time scalability, and being able to run the exact OS version needed for a specific application. Virtualized servers and services are indistinguishable from traditional servers and services from a user’s perspective. Additionally, recovery from damaged, crashed, or corrupted virtual systems is often quick, simply consisting of replacing the virtual system’s main hard drive file with a clean backup version and then relaunching it. (Additional coverage of virtualization and

some of its associated risks are covered in Chapter 9 along with cloud computing.)

Trusted Platform Module

The *Trusted Platform Module (TPM)* is both a specification for a cryptoprocessor chip on a mainboard and the general name for implementation of the specification. A TPM chip is used to store and process cryptographic keys for the purposes of a hardware supported/implemented hard drive encryption system. Generally, a hardware implementation, rather than a software-only implementation of hard drive encryption, is considered to be more secure.

When TPM-based whole-disk encryption is in use, the user/operator must supply a password or physical Universal Serial Bus (USB) token device to the computer to authenticate and allow the TPM chip to release the hard drive encryption keys into memory. While this seems similar to a software implementation, the key difference is that if the hard drive is removed from its original system, it cannot be decrypted. Only with the original TPM chip can an encryption be decrypted and accessed. With software-only hard drive encryption, the hard drive can be moved to a different computer without any access or use limitations.

A *hardware security module (HSM)* is a cryptoprocessor used to manage/store digital encryption keys, accelerate crypto operations, support faster digital signatures, and improve authentication. An HSM is often an add-on adapter or peripheral or can be a Transmission Control Protocol/Internet Protocol (TCP/IP) network device. HSMs include tamper protection to prevent their misuse even if physical access is gained by an attacker. A TPM is just one example of an HSM.

HSMs provide an accelerated solution for large (2,048+ bit) asymmetric encryption calculations and a secure vault for key storage. Many certificate authority systems use HSMs to store certificates; ATM and POS bank terminals often employ proprietary HSMs; hardware SSL accelerators can include HSM support; and Domain Name System Security Extensions (DNSSEC)—compliant Domain

Name System (DNS) servers use HSM for key and zone file storage.

Interfaces

A *constrained or restricted interface* is implemented within an application to restrict what users can do or see based on their privileges. Users with full privileges have access to all the capabilities of the application. Users with restricted privileges have limited access.

Applications constrain the interface using different methods. A common method is to hide the capability if the user doesn't have permissions to use it. Commands might be available to administrators via a menu or by right-clicking an item, but if a regular user doesn't have permissions, the command does not appear. Other times, the command is shown but is dimmed or disabled. The regular user can see it but will not be able to use it.

The purpose of a constrained interface is to limit or restrict the actions of both authorized and unauthorized users. The use of such an interface is a practical implementation of the Clark-Wilson model of security.

Fault Tolerance

Fault tolerance is the ability of a system to suffer a fault but continue to operate. Fault tolerance is achieved by adding redundant components such as additional disks within a redundant array of inexpensive disks (RAID) array, or additional servers within a failover clustered configuration. Fault tolerance is an essential element of security design. It is also considered part of avoiding single points of failure and the implementation of redundancy. For more details on fault tolerance, redundant servers, RAID, and failover solutions, see Chapter 18, "Disaster Recovery Planning."

Summary

Secure systems are not just assembled; they are designed to support security. Systems that must be secure are judged for their ability to support and enforce the security policy. This process of evaluating the effectiveness of a computer system is certification. The certification process is the technical evaluation of a system's ability to meet its design goals. Once a system has satisfactorily passed the technical evaluation, the management of an organization begins the formal acceptance of the system. The formal acceptance process is accreditation.

The entire certification and accreditation process depends on standard evaluation criteria. Several criteria exist for evaluating computer security systems. The earliest, TCSEC, was developed by the U.S. Department of Defense. TCSEC, also called the Orange Book, provides criteria to evaluate the functionality and assurance of a system's security components. ITSEC is an alternative to the TCSEC guidelines and is used more often in European countries. In 2005, TCSEC was replaced by the Common Criteria. Regardless of which criteria you use, the evaluation process includes reviewing each security control for compliance with the security policy. The better a system enforces the good behavior of subjects' access to objects, the higher the security rating.

When security systems are designed, it is often helpful to create a security model to represent the methods the system will use to implement the security policy. We discussed several security models in this chapter. The Bell-LaPadula model supports data confidentiality only. It was designed for the military and satisfies military concerns. The Biba model and the Clark-Wilson model address the integrity of data and do so in different ways. These models are often used as part of the foundation when designing security infrastructure for commercial applications.

All of this understanding must culminate into an effective system security implementation in terms of preventive, detective, and corrective controls. That's why you must also know the access control

models and their functions. This includes the state machine model, Bell-LaPadula, Biba, Clark-Wilson, the information flow model, the noninterference model, the Take-Grant model, the access control matrix model, and the Brewer and Nash model.