



# COMBINATIONAL LOGIC

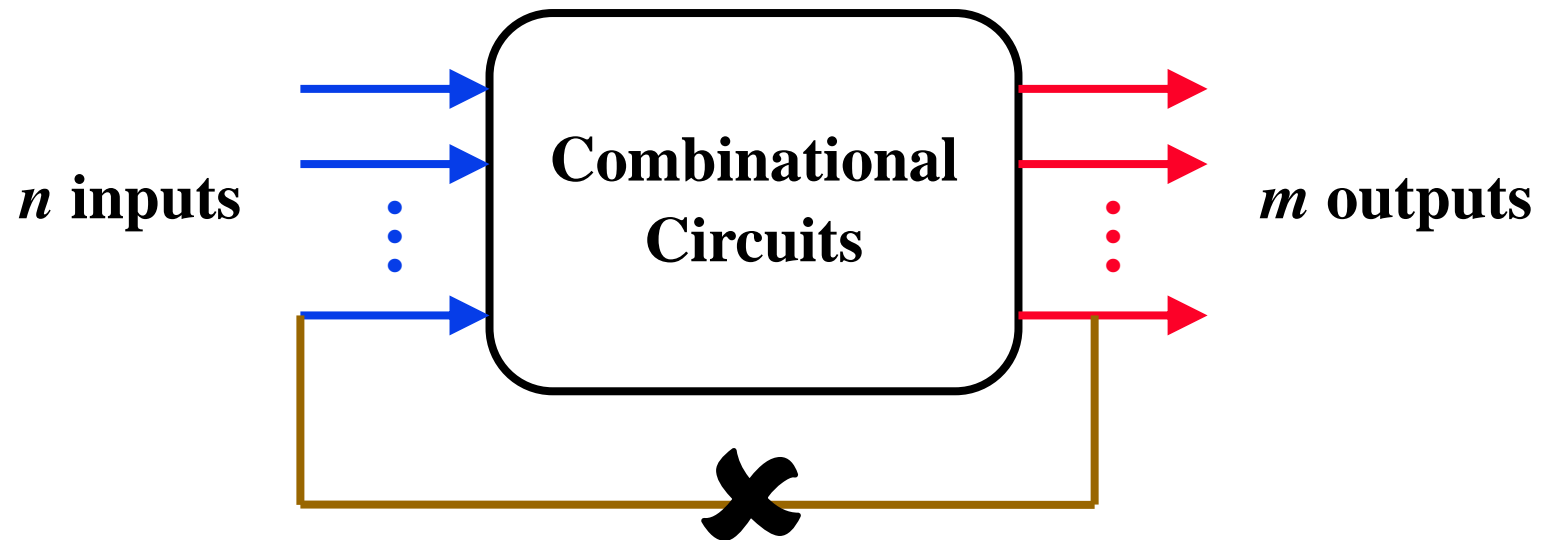
## DIGITAL LOGIC DESIGN

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# Combinational Circuits

★ Output is function of input only

i.e. no feedback

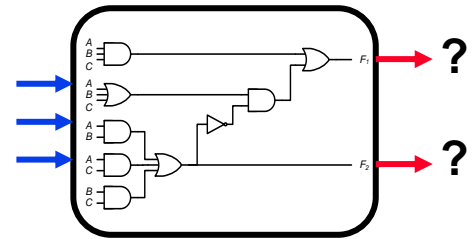


When **input** changes, **output** may change (after a delay)

# Combinational Circuits

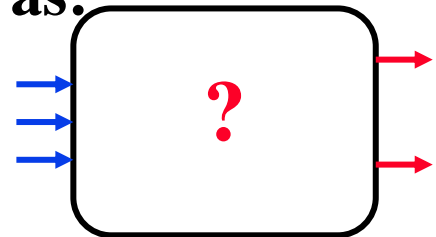
## ★ Analysis

- Given a circuit, find out its *function*
- Function may be expressed as:
  - ◆ Boolean function/equation
  - ◆ Truth table



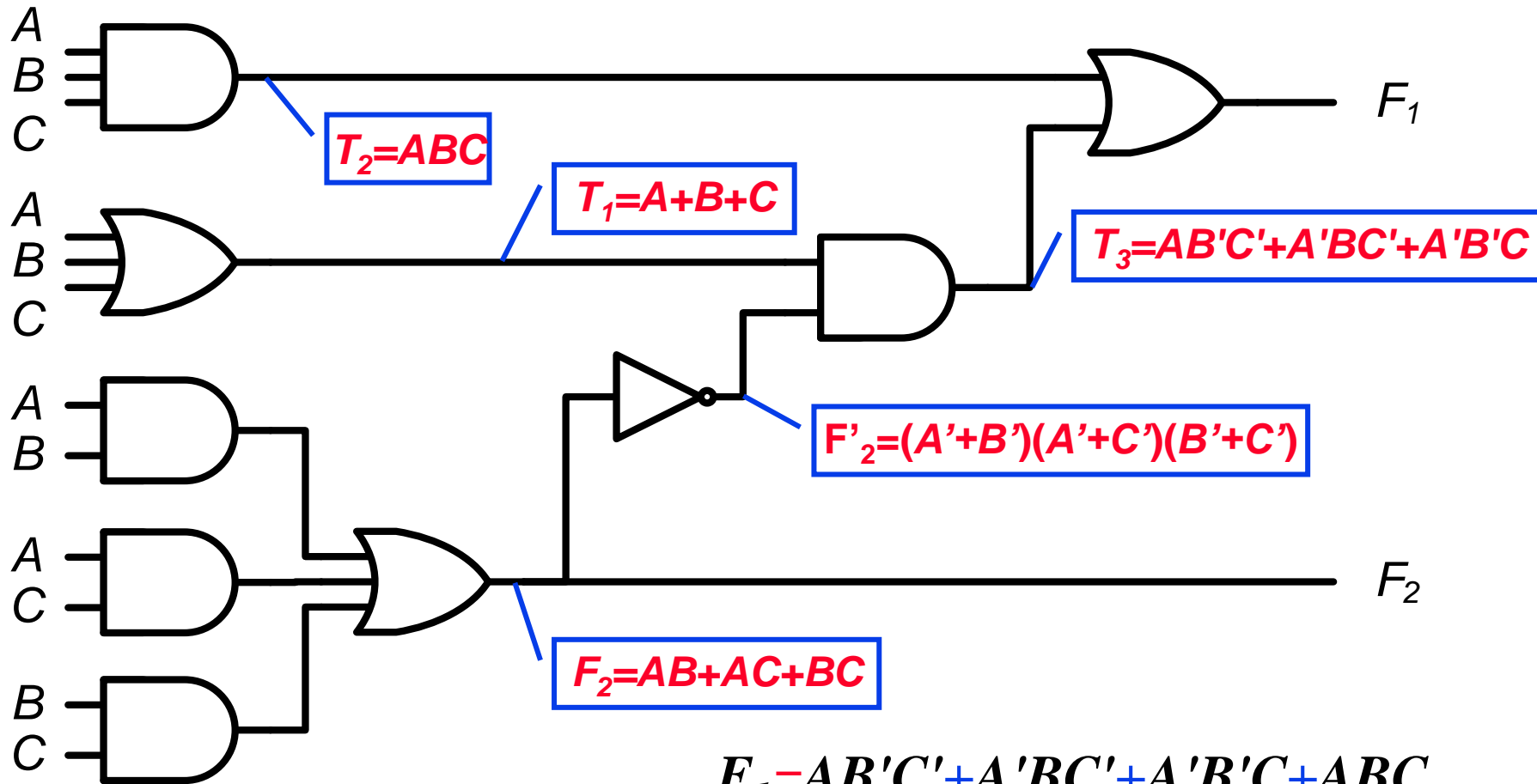
## ★ Design

- Given a desired function/specification, determine its *circuit*
- Function/specification may be expressed as:
  - ◆ Boolean function
  - ◆ Truth table
  - ◆ Statement



# Analysis Procedure

## ★ Boolean Expression Approach



$$F_1 = AB'C' + A'BC' + A'B'C + ABC$$

$$F_2 = AB + AC + BC$$

# Boolean Expression Approach

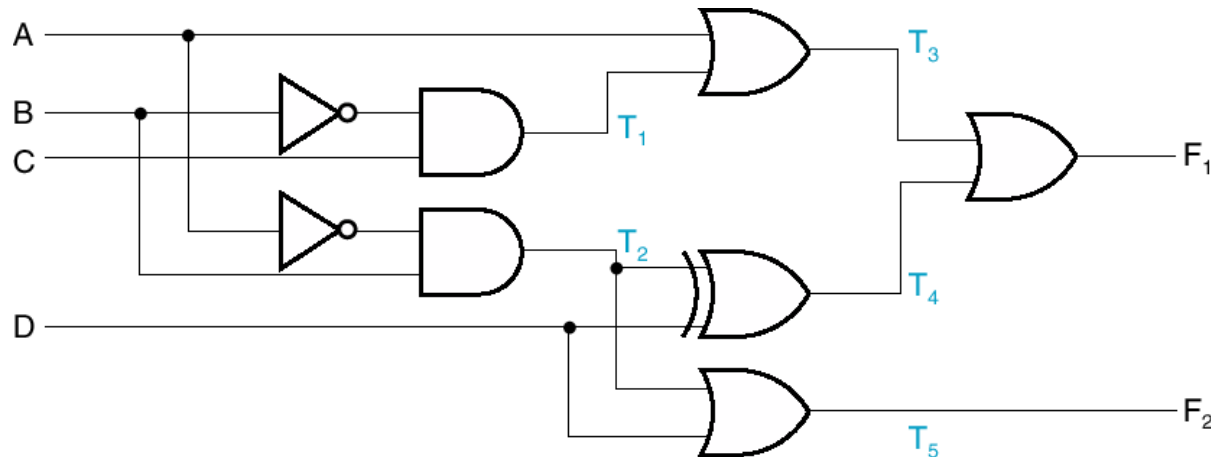
Label gate outputs of input variables

Determine Boolean functions or values

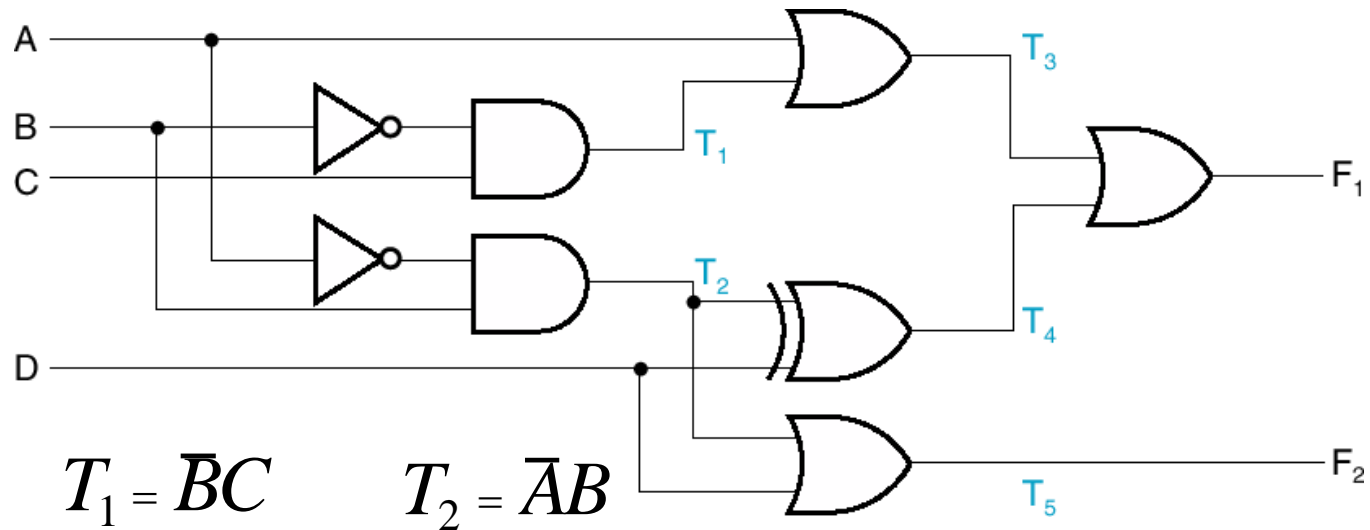
Label outputs of gates fed by previously labeled gates

Determine Boolean function or values

Repeat 2 until done



# Boolean Expression Approach



$$T_3 = A + T_1 = A + \bar{B}C$$

# Boolean Expression Approach :Cont

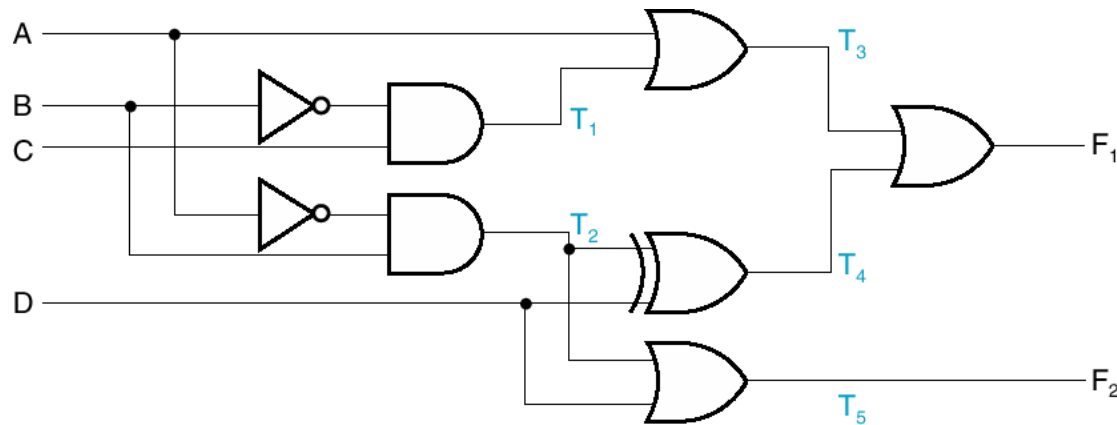
## .. Analysis Example

$$T_4 = T_2 \oplus D = (A\bar{B}) \oplus D = A\bar{B}\bar{D} + AD + B\bar{D}$$

$$T_5 = T_2 + D = A\bar{B} + D$$

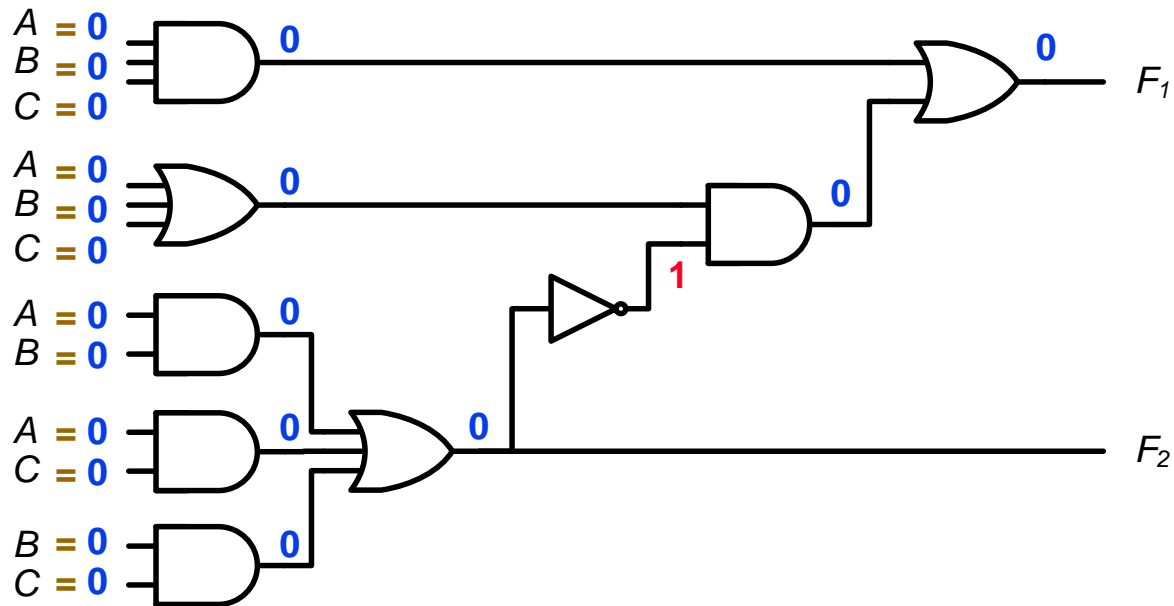
$$F_1 = T_3 + T_4 = A + BC + ABD + AD + B\bar{D}$$

$$F_2 = T_5 = A\bar{B} + D$$



# Analysis Procedure

## ★ Truth Table Approach

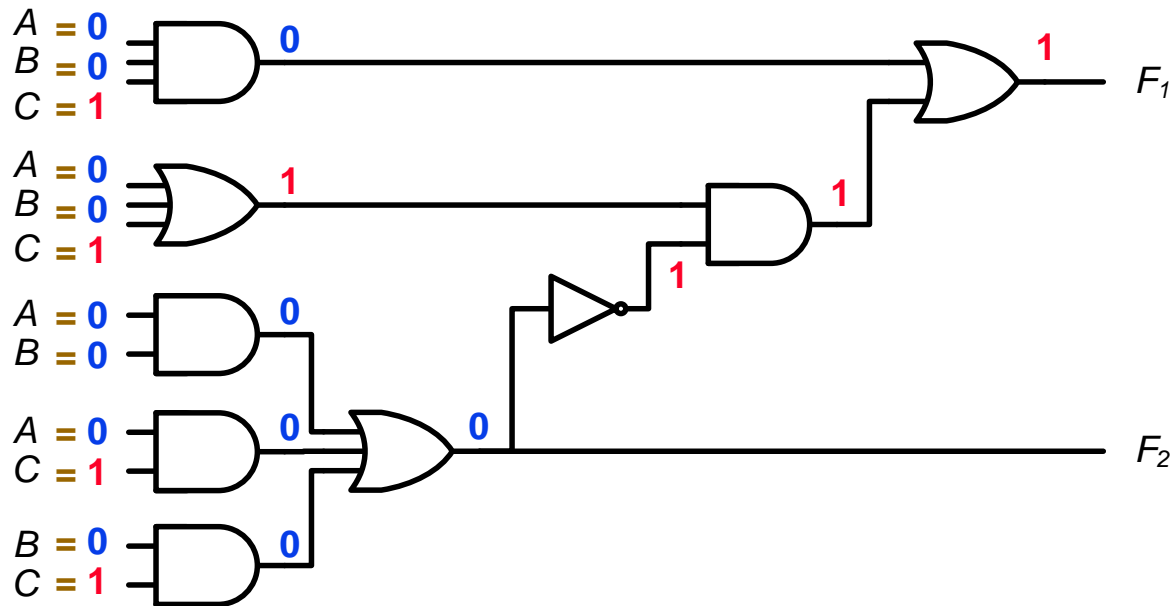


A	B	C	F <sub>1</sub>	F <sub>2</sub>
0	0	0	0	0



# Analysis Procedure

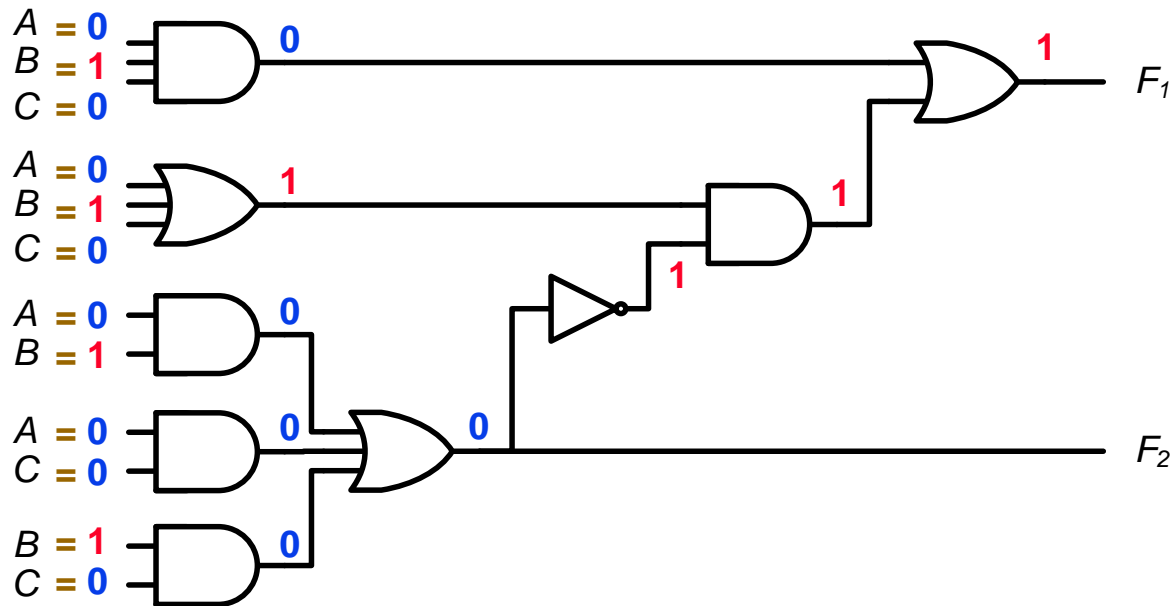
## ★ Truth Table Approach



A	B	C	F <sub>1</sub>	F <sub>2</sub>
0	0	0	0	0
0	0	1	1	0

## Analysis Procedure

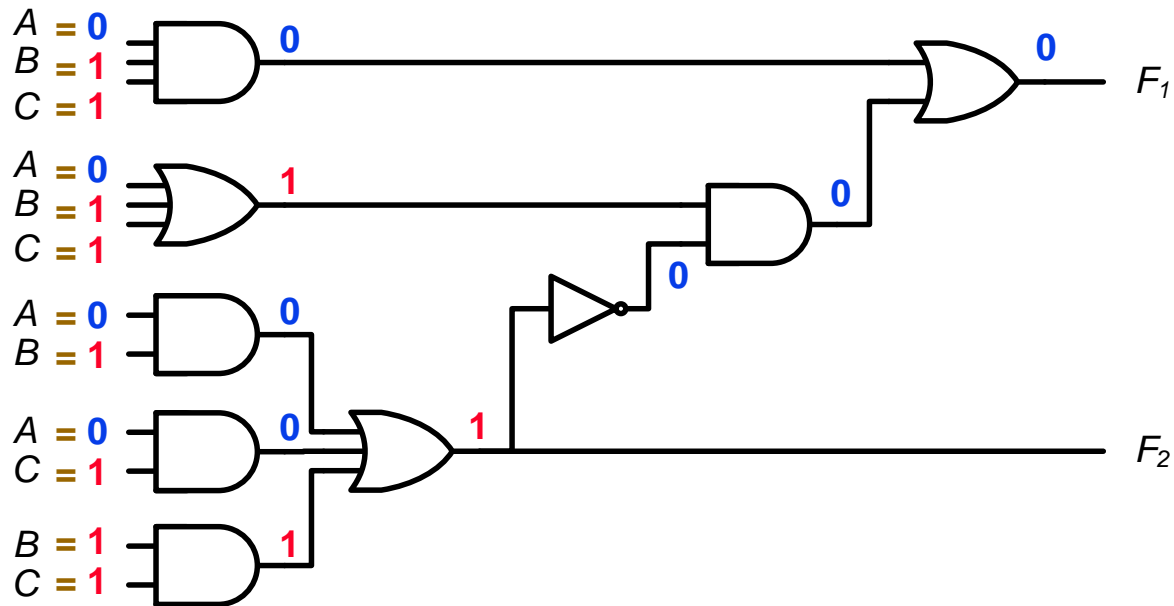
## ★ Truth Table Approach



$A$	$B$	$C$	$F_1$	$F_2$
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0

# Analysis Procedure

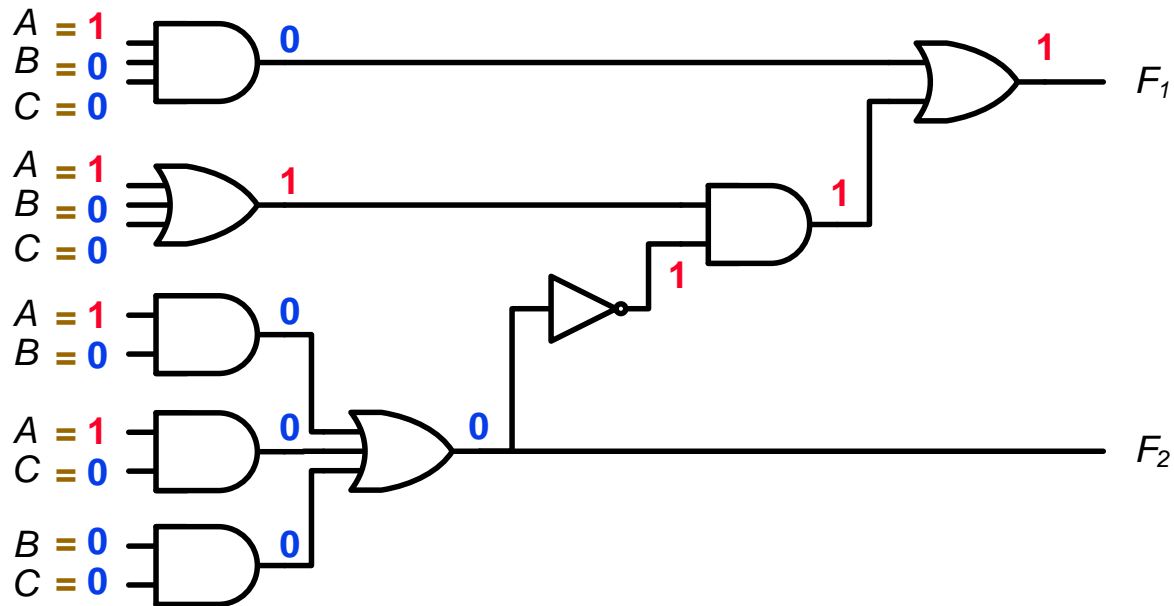
## ★ Truth Table Approach



A	B	C	$F_1$	$F_2$
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1

# Analysis Procedure

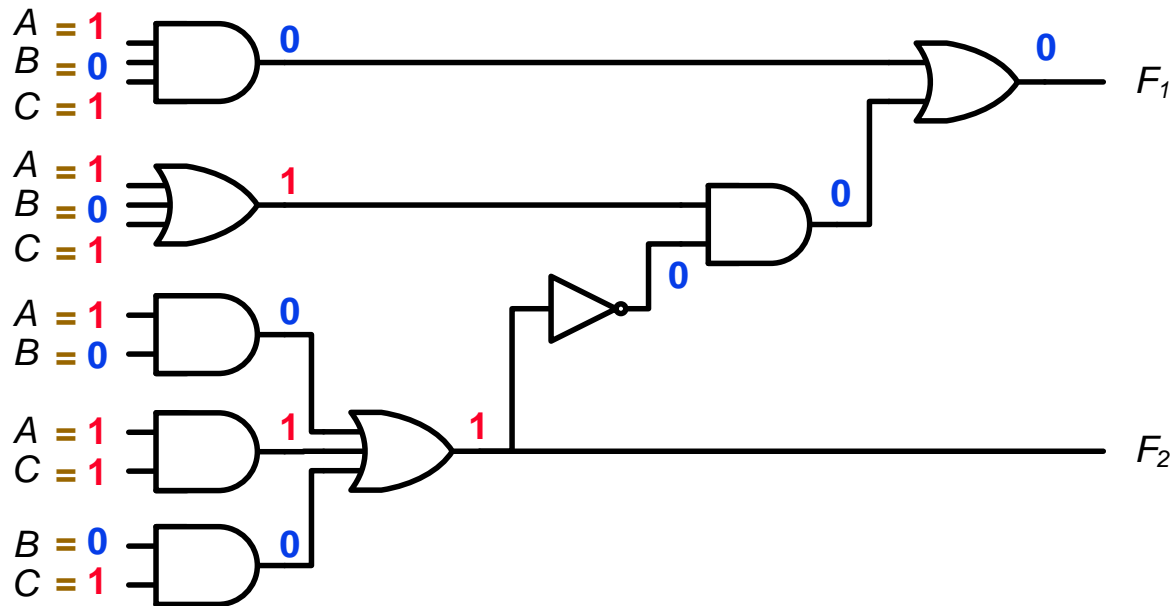
## ★ Truth Table Approach



$A$	$B$	$C$	$F_1$	$F_2$
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0

# Analysis Procedure

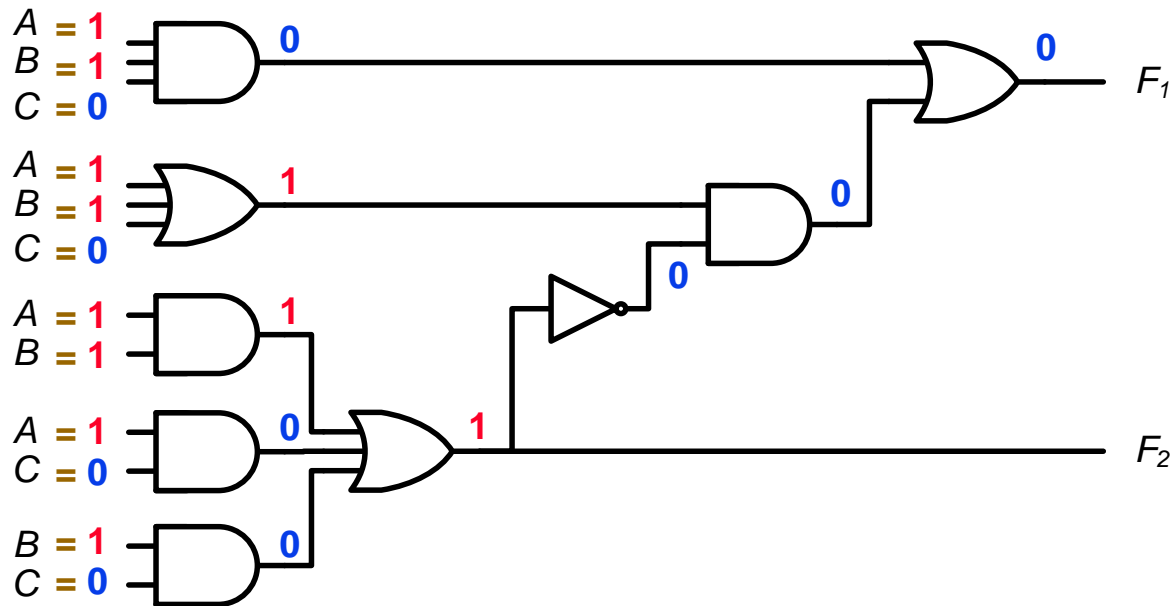
## ★ Truth Table Approach



A	B	C	$F_1$	$F_2$
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1

# Analysis Procedure

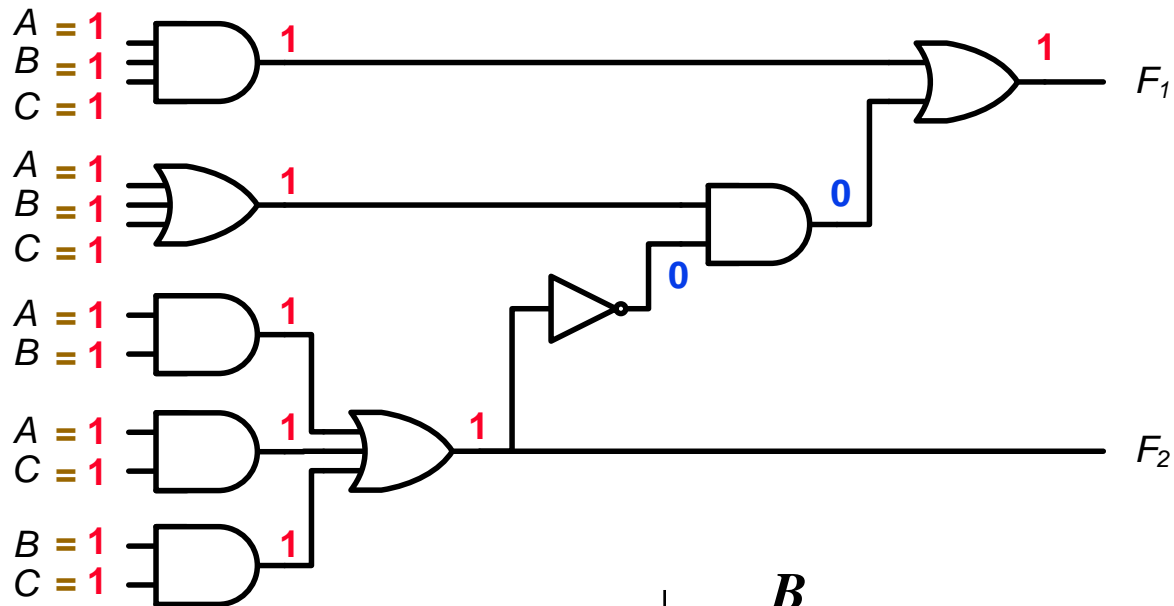
## ★ Truth Table Approach



A	B	C	$F_1$	$F_2$
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1

# Analysis Procedure

## ★ Truth Table Approach



$A$	$B$	$C$	$F_1$	$F_2$
0	0	0	0	0
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	1

	$B$			
	0	1	0	1
$A$	1	0	1	0
	$C$			

	$B$			
	0	0	1	0
$A$	0	1	1	1
	$C$			

Grouping for  $F_2$ :  
 - Blue oval around (A=0, B=1, C=0) and (A=0, B=1, C=1) → term  $AB$   
 - Green oval around (A=0, B=1, C=1) and (A=1, B=1, C=1) → term  $AC$   
 - Brown oval around (A=0, B=1, C=1) and (A=1, B=0, C=1) → term  $BC$

$$F_1 = AB'C' + A'BC' + A'B'C + ABC$$

$$F_2 = AB + AC + BC$$

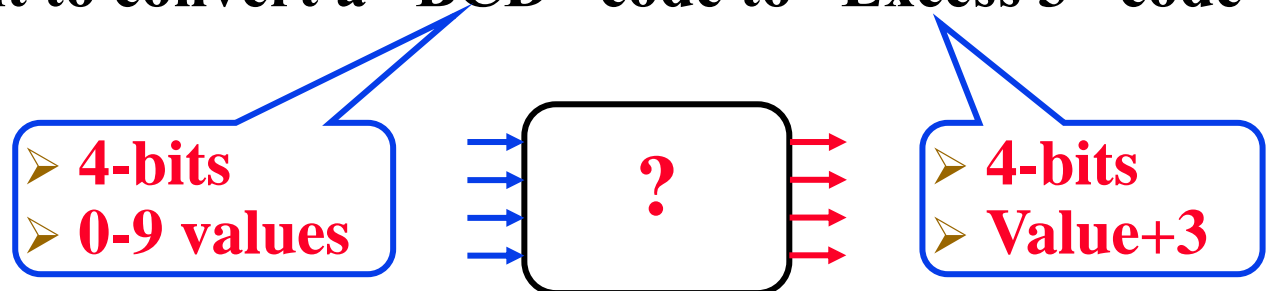
# Design Procedure

## ★ Given a problem statement:

- Determine the number of *inputs* and *outputs*
- Derive the truth table
- Simplify the Boolean expression for each output
- Produce the required circuit

## Example 1:

Design a circuit to convert a “BCD” code to “Excess 3” code

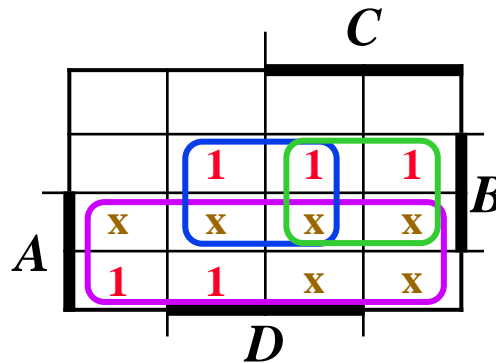




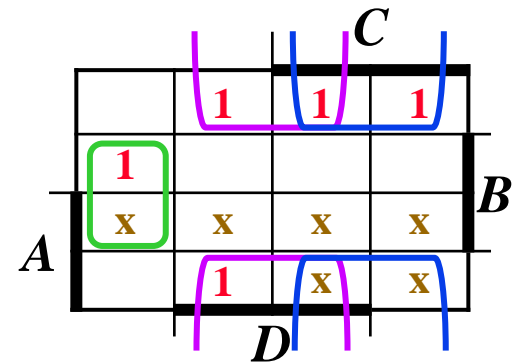
# Design Procedure

## ★ BCD-to-Excess 3 Converter

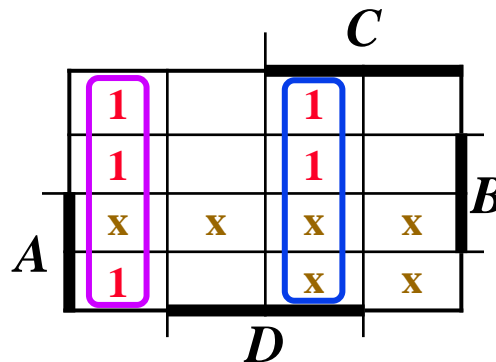
<i>A B C D</i>	<i>w x y z</i>
0 0 0 0	0 0 1 1
0 0 0 1	0 1 0 0
0 0 1 0	0 1 0 1
0 0 1 1	0 1 1 0
0 1 0 0	0 1 1 1
0 1 0 1	1 0 0 0
0 1 1 0	1 0 0 1
0 1 1 1	1 0 1 0
1 0 0 0	1 0 1 1
1 0 0 1	1 1 0 0
1 0 1 0	x x x x
1 0 1 1	x x x x
1 1 0 0	x x x x
1 1 0 1	x x x x
1 1 1 0	x x x x
1 1 1 1	x x x x



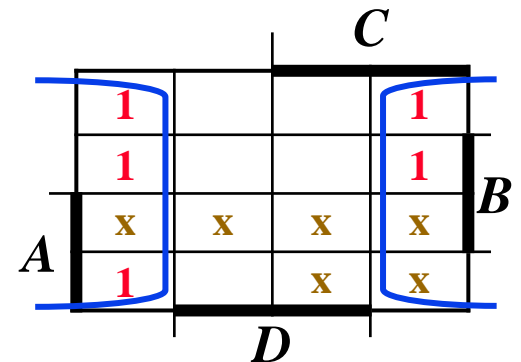
$$w = A + BC + BD$$



$$x = B'C + B'D + BC'D'$$



$$y = C'D' + CD$$

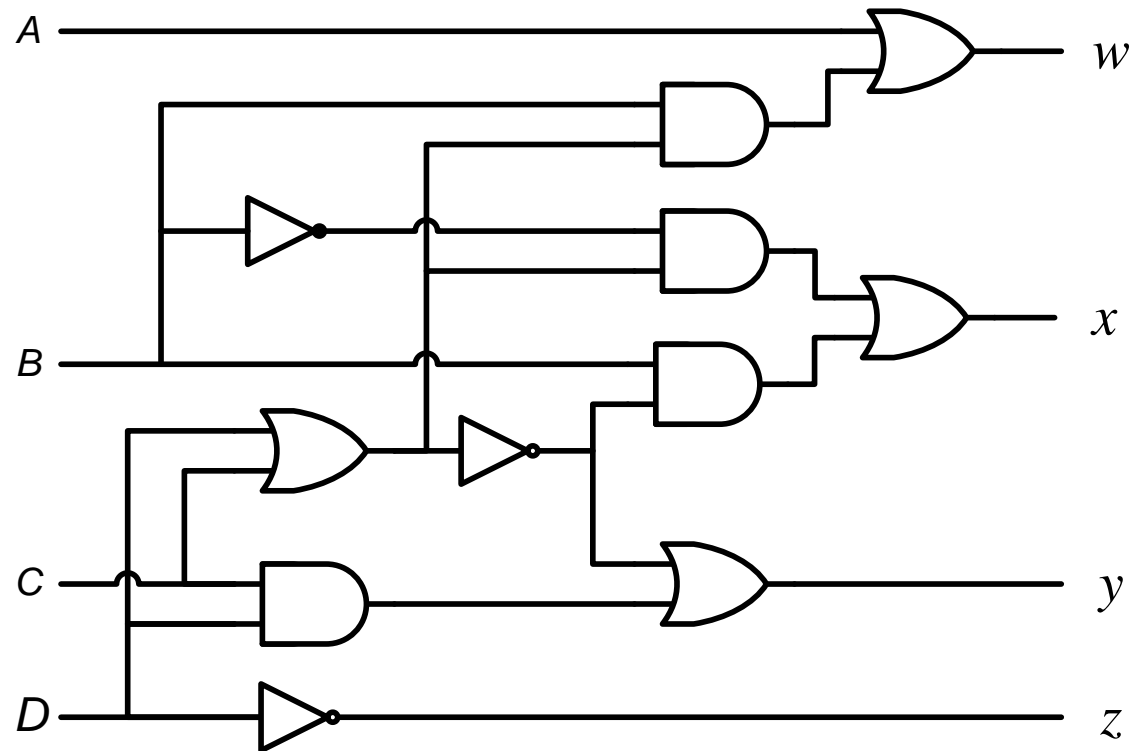


$$z = D'$$

# Design Procedure

## ★ BCD-to-Excess 3 Converter

<i>A B C D</i>	<i>w x y z</i>
0 0 0 0	0 0 1 1
0 0 0 1	0 1 0 0
0 0 1 0	0 1 0 1
0 0 1 1	0 1 1 0
0 1 0 0	0 1 1 1
0 1 0 1	1 0 0 0
0 1 1 0	1 0 0 1
0 1 1 1	1 0 1 0
1 0 0 0	1 0 1 1
1 0 0 1	1 1 0 0
1 0 1 0	x x x x
1 0 1 1	x x x x
1 1 0 0	x x x x
1 1 0 1	x x x x
1 1 1 0	x x x x
1 1 1 1	x x x x



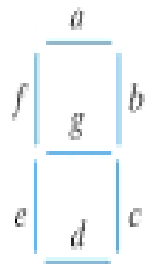
$$w = A + B(C+D)$$

$$y = (C+D)' + CD$$

$$x = B'(C+D) + B(C+D)'$$

$$z = D'$$

## Example 2: Seven-Segment Decoder



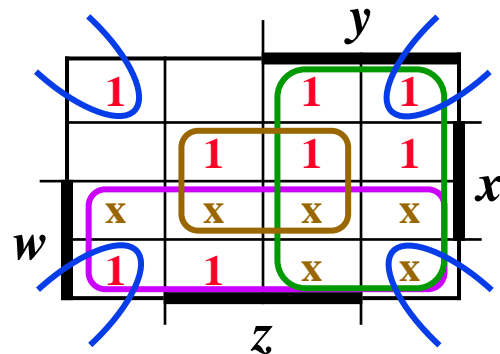
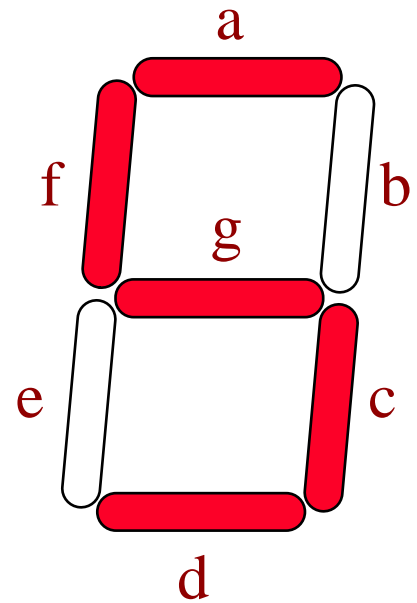
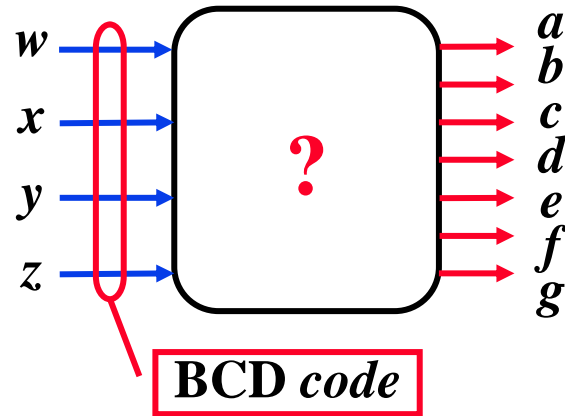
(a) Segment designation



(b) Numerical designation for display

# Example 2: Seven-Segment Decoder

<i>w x y z</i>	<i>a b c d e f g</i>
0 0 0 0	1 1 1 1 1 1 0
0 0 0 1	0 1 1 0 0 0 0
0 0 1 0	1 1 0 1 1 0 1
0 0 1 1	1 1 1 1 0 0 1
0 1 0 0	0 1 1 0 0 1 1
0 1 0 1	1 0 1 1 0 1 1
0 1 1 0	1 0 1 1 1 1 1
0 1 1 1	1 1 1 0 0 0 0
1 0 0 0	1 1 1 1 1 1 1
1 0 0 1	1 1 1 1 0 1 1
1 0 1 0	x x x x x x x
1 0 1 1	x x x x x x x
1 1 0 0	x x x x x x x
1 1 0 1	x x x x x x x
1 1 1 0	x x x x x x x
1 1 1 1	x x x x x x x



$$a = w + y + xz + x'z'$$

$$b = \dots$$

$$c = \dots$$

$$d = \dots$$

# Example 3: Implement circuit that can convert 3bit gray code into binarycode

$$B_2 = \sum m(4, 5, 6, 7)$$

$$B_1 = \sum m(2, 3, 4, 5)$$

$$B_0 = \sum m(1, 2, 4, 7)$$

		G1'G0'		G1'G0		G1G0		G1G0'	
		00		01		11		10	
G2'	0	0		1		3		2	
G2	1	1 <sub>4</sub>		1 <sub>5</sub>		1 <sub>7</sub>		1 <sub>6</sub>	

$$B_2 = G_2$$

		$G_1'G_0'$	$G_1'G_0$	$G_1G_0$	$G_1G_0'$
		00	01	11	10
$G_2'$	0	0	1	1 <sub>3</sub>	1 <sub>2</sub>
$G_2$	1	1 <sub>4</sub>	1 <sub>5</sub>	7	6

$$B_1 = G_2'G_1 + G_2G_1'$$

$$B_1 = G_2 \oplus G_1$$

		G1'G0'	G1'G0	G1G0	G1G0'
		00	01	11	10
G2'	0	0	1 <sub>1</sub>	3	1 <sub>2</sub>
G2	1	1 <sub>4</sub>	5	1 <sub>7</sub>	6

$$B_0 = G_2 \oplus G_1 \oplus G_0$$

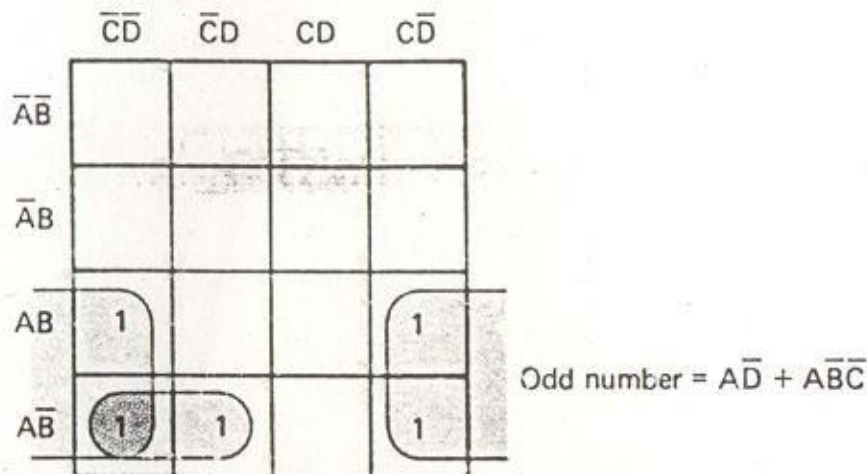
Decimal Equivalent	Gray Code			Binary Output		
	G2	G1	G0	B2	B1	B0
0	0	0	0	0	0	0
1	0	0	1	0	0	1
3	0	1	1	0	1	0
2	0	1	0	0	1	1
6	1	1	0	1	0	0
7	1	1	1	1	0	1
5	1	0	1	1	1	0
4	1	0	0	1	1	1

## Example 4: Implement a binary to gray converter

Decimal Equivalent	Binary			Gray Code		
	B2	B1	B0	G2	G1	G0
0	0	0	0	0	0	0
1	0	0	1	0	0	1
2	0	1	0	0	1	1
3	0	1	1	0	1	0
4	1	0	0	1	1	0
5	1	0	1	1	1	1
6	1	1	0	1	0	1
7	1	1	1	1	0	0

# Example 5: System Design Application

Design a circuit that can be built using an AOI and inverters that will output a HIGH (1) whenever the 4-bit hexadecimal input is an odd number from 0 to 9.



Hex Truth Table Used to Determine the Equation for Odd Numbers<sup>a</sup> from 0 to 9

D	C	B	A	DEC	
0	0	0	0	0	
0	0	0	1	1	$\leftarrow A\bar{B}\bar{C}\bar{D}$
0	0	1	0	2	
0	0	1	1	3	$\leftarrow AB\bar{C}\bar{D}$
0	1	0	0	4	
0	1	0	1	5	$\leftarrow A\bar{B}C\bar{D}$
0	1	1	0	6	
0	1	1	1	7	$\leftarrow ABC\bar{D}$
1	0	0	0	8	
1	0	0	1	9	$\leftarrow A\bar{B}\bar{C}D$

<sup>a</sup> Odd number =  $A\bar{B}\bar{C}\bar{D} + AB\bar{C}\bar{D} + A\bar{B}C\bar{D} + ABC\bar{D} + A\bar{B}\bar{C}D$ .

# Example 5: System Design Application

implementation of the "odd-number decoder" using an AOI

