Syllabus of Operating System

1) Basic Introduction → types, process Diagram
   System call

2) ** Process Scheduling → FIFO
   → SJF
   → Pr
   → Round Robin

3) Process Synchronization → Semaphore

4) Deadlock & threads → Banker

5) ** Memory Management → Virtual Mem.
   → Paging
   → Segmentation
   → Page Replacement algo
   → fragmentation

6) ** Disk Scheduling
   → SCAN, CSCAN, FCFS

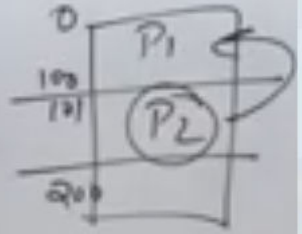
7) UNIX Commands → ls, mkdir, CD, Chmod, Open, Seq, Random, Linked

8) File Mgmt and Secur

Operating System and its Functions

→ Windows (82%)

→ Primary goal



→ Convenience

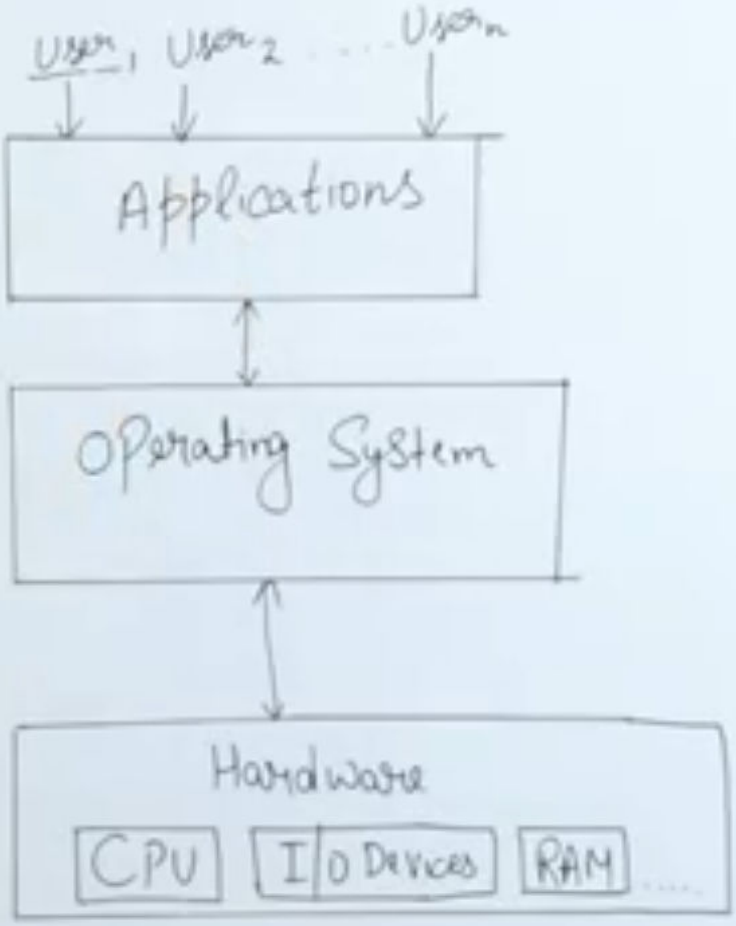
(95%)

→ throughput (Linux)

5) Security

U → OS → H

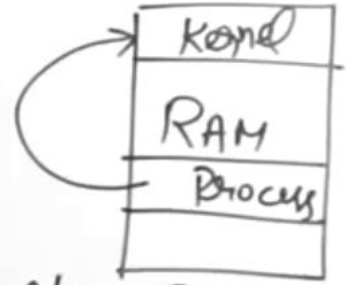
1) Resource Mgmt.
2) Process Mgmt. (CPU Scheduling)
3) Storage mgmt (HD) → file System
4) Memory Mgmt (RAM)



"System Call"

- File Related ⇒ Open(), Read(), Write(), Close(), Create file etc.
- Device Related ⇒ Read, Write, Reposition, ioctl, fcntl
- Information ⇒ get Pid, attributes, get System time and data
- Process Control ⇒ Load, Execute, abort, Fork, Wait, Signal, Allocate etc.
- Communication ⇒ Pipe(), Create/delete Connections, Shmget()

Printf.



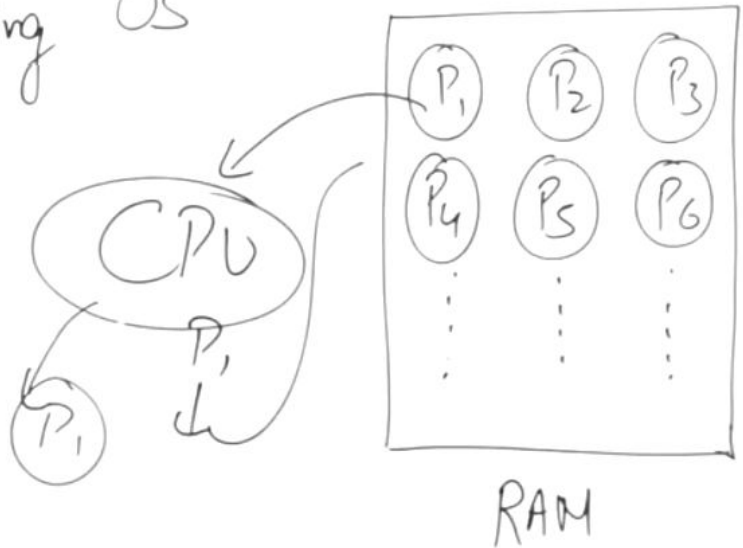
Multiprogrammed OS — Now Preemptive

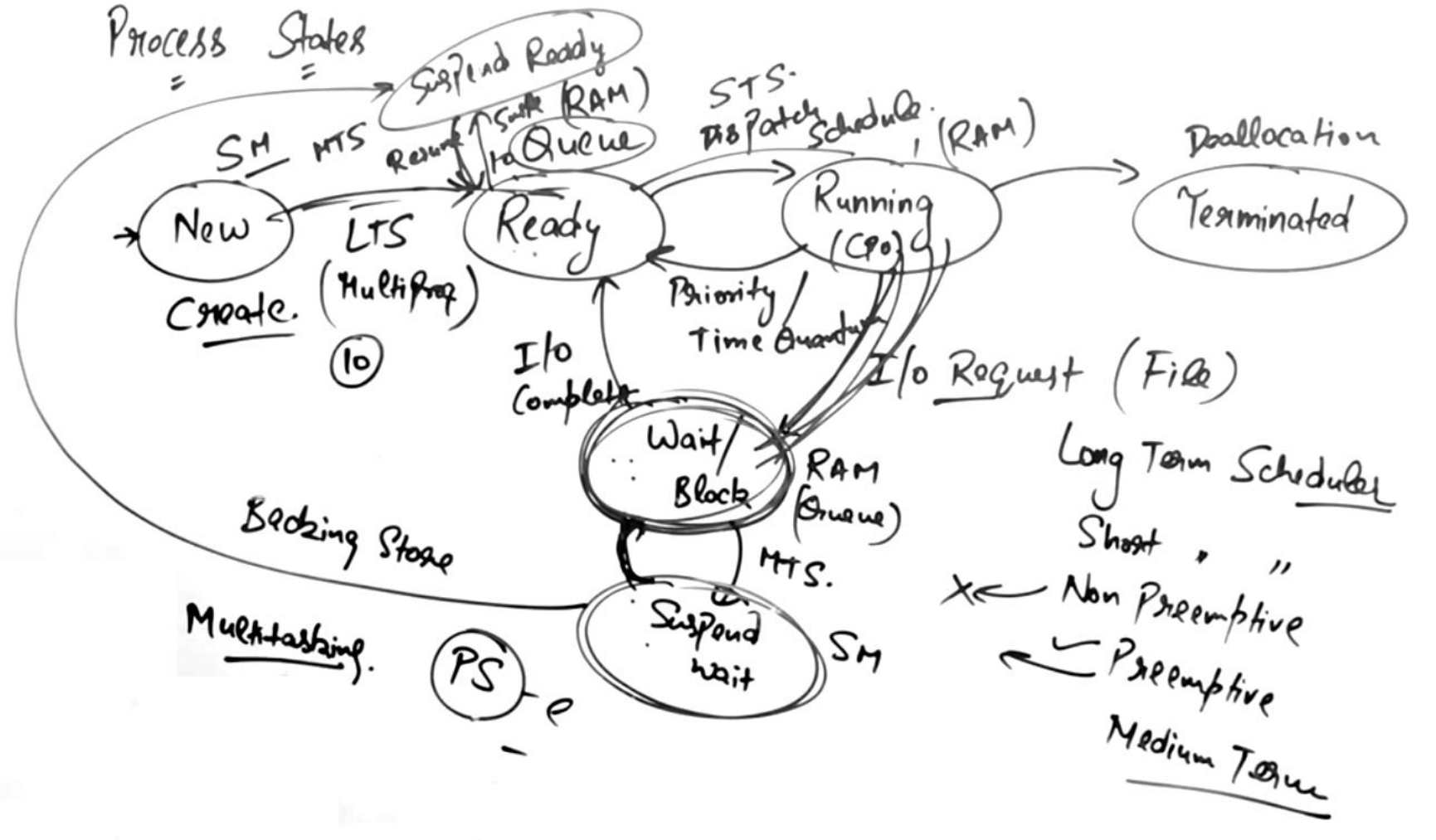
↳ IDLENESS

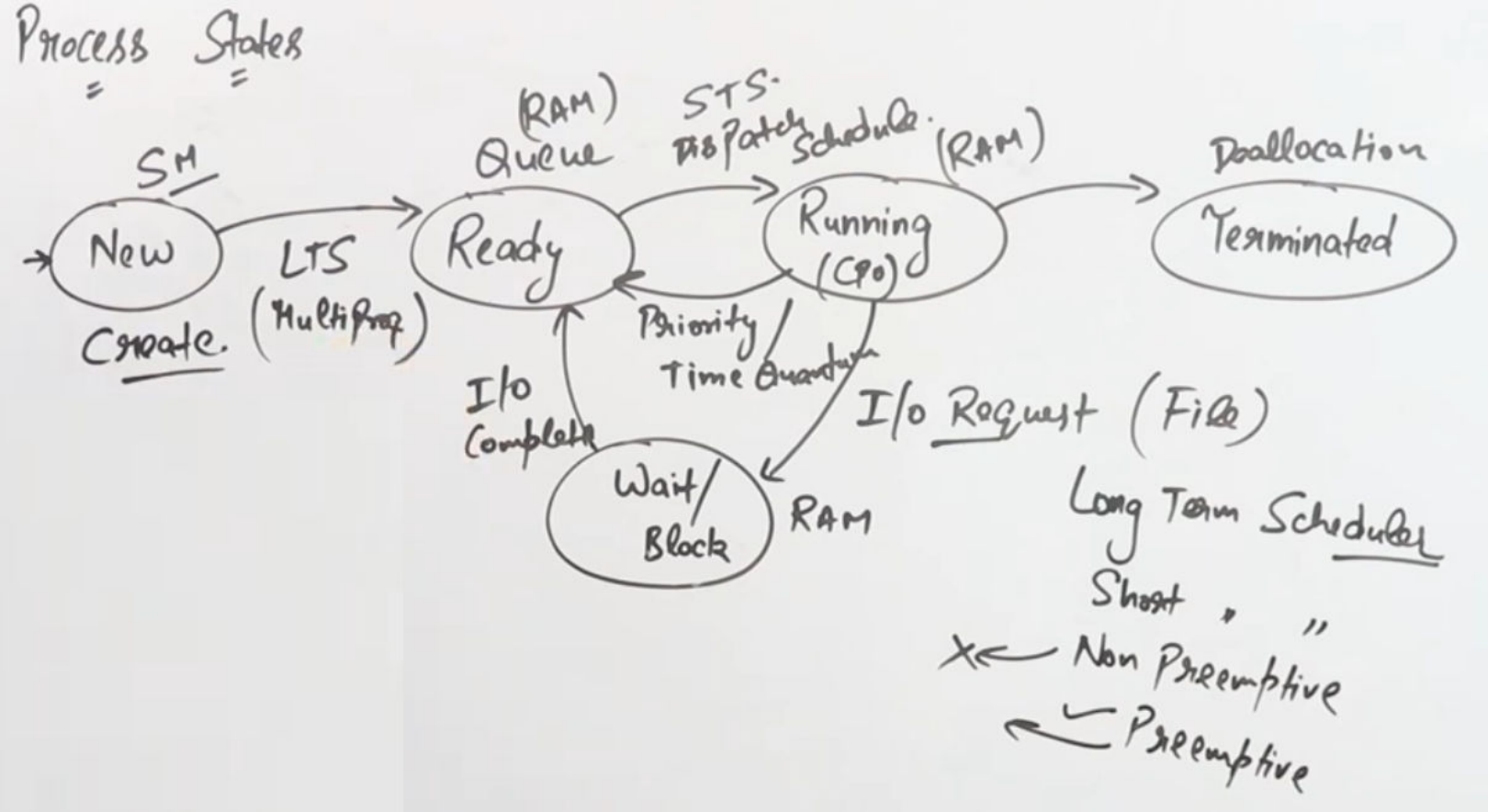
Multitasking / Time Sharing OS

↓

Responsiveness




Process States
Suspend Ready
Swap (RAM)
Queue
STS.
Dispatch
Scheduler
(RAM)
Deallocation
SM
MTS
Resume
New
LTS
(MultiProg)
Ready
Running
(CPU)
Terminated
Create.
(IO)
Priority
Time Quantum
I/O Complete
I/O Request (File)
Wait/
Block
RAM
(Queue)
Long Term Scheduler
Short " "
Non Preemptive
Preemptive
Medium Term
Backing Store
MTS.
Multitasking.
PS
Suspend
wait
SM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Short "
Non Preemptive
Preemptive

```
# include <stdio.h>
# include <unistd.h>
  int main()
  {
→   if (fork() && fork())
      fork();
    printf("Hello");
    return 0;
  }
```

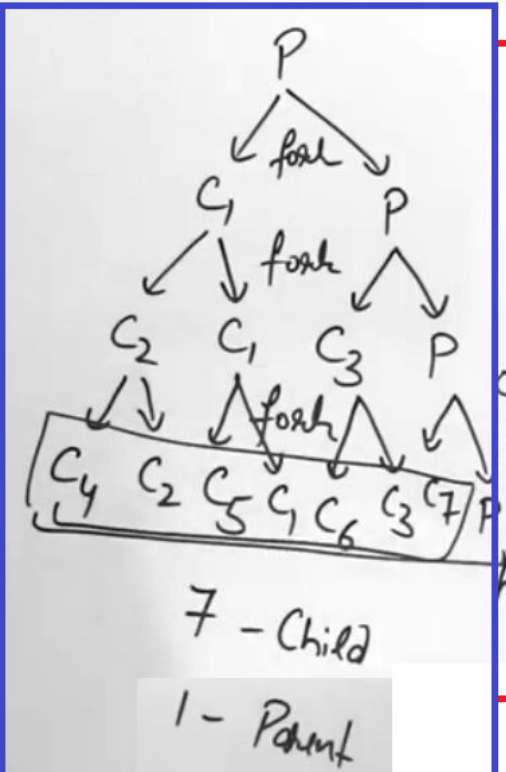
.c

running in parrellel due to fork() call - child & parent

Fork() == 0 child ✓

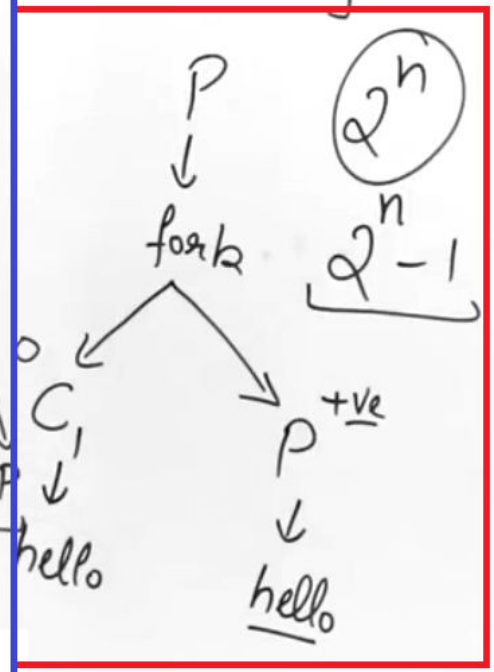
+1 +ve → Parent

-1 → Child ✗



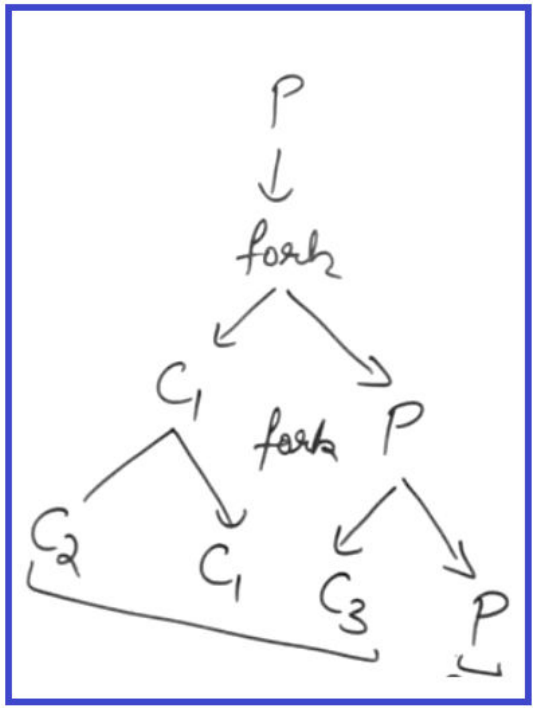
7 - Child

1 - Parent




```
main() {
fork();
fork();
Printf("hello");
}
```

fork
fork
fork
Pf("hello")



Q.
```
#include <stdio.h>
#include <unistd.h>
int main()
{
  int a;
  for (a=1; a<5; a++)
        fork();
  printf("1");
}
```

How many times it will print “1” in output?

(A) 15

(B) 16

(C) 31

(D) 32

Q.
```
#include <stdio.h>
#include <unistd.h>
int main()
{
  if (fork() && fork())
        fork();
  printf("Hello");
  return 0;
}
```

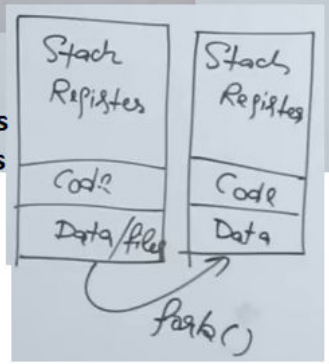
How many times it will print “Hello” in output?

(A) 2 (B) 3

(C) 4 (D) 5

Process

1) System Calls involved in Process
2) OS treats different processes differently
3) Different process have different Copies of Data, files, Code
4) Content switching is slower
5) Blocking a process will not block another
6) Independent



each process will have PID in process call. But in Threads there is 1 PID, however is this process is paused then the whole thread will be paused

content swiitching is slower as it has to save in process control block i.e. in RAM - Process

Threads (User Level)

1) There is no system Call involved
2) All User level threads treated as single task for OS
3) Threads share same copy of code and data
4) Content switching is faster
5) Blocking a thread will block entire process
6) Interdependent

blocking - in i/o request, then whole process will be blocked which includes all the Threads