

# Graphs

## Graph:

- ✚ **Graph** is a mathematical representation of a network and it describes the relationship between lines and points.
- ✚ Graph is a way of expressing information in pictorial form.
- ✚ A graph is a collection of nodes and edges. A graph is also called a network.
- ✚ A node is whatever you are interested in: person, city, team, project, computer, etc.
- ✚ An edge represents a relationship between nodes.

## Daily life examples of Graphs

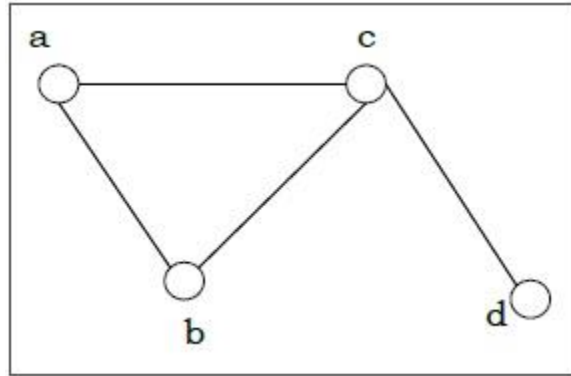
- ✚ Map of the world, or world
  - Pakistan
  - Provinces
  - Islamabad
  - Sector H-9
  - NUML
  - Ghazali Block
- ✚ Internet
- ✚ WAN, LAN
- ✚ Our Knowledge
- ✚ Our friendship circle
- ✚ Human body

And many more.....

## What is a Graph?

**Definition** – A graph (denoted as  $G=(V,E)$   $G=(V,E)$ ) consists of a non-empty set of vertices or nodes  $V$  and a set of edges  $E$ .

**Example** – Let us consider, a Graph  
is  $G=(V,E)$   $G=(V,E)$  where  $V=\{a,b,c,d\}$   $V=\{a,b,c,d\}$  and  $E=\{\{a,b\},\{a,c\},\{b,c\},\{c,d\}\}$



**Degree of a Vertex** – The degree of a vertex  $V$  of a graph  $G$  (denoted by  $\deg(V)$ ) is the number of edges incident with the vertex  $V$ .

Vertex	Degree	Even / Odd
a	2	even
b	2	even
c	3	odd
d	1	odd

**Even and Odd Vertex** – If the degree of a vertex is even, the vertex is called an even vertex and if the degree of a vertex is odd, the vertex is called an odd vertex.

**Degree of a Graph** – The degree of a graph is the largest vertex degree of that graph. For the above graph the degree of the graph is 3.

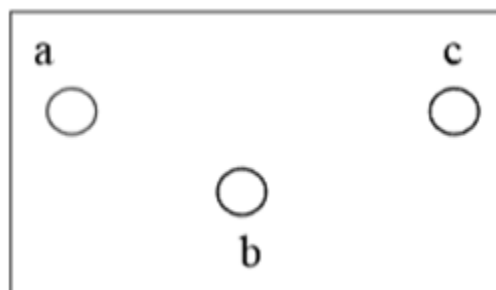
**The Handshaking Lemma** – In a graph, the sum of all the degrees of all the vertices is equal to twice the number of edges.

## Types of Graphs

There are different types of graphs, which we will learn in the following section.

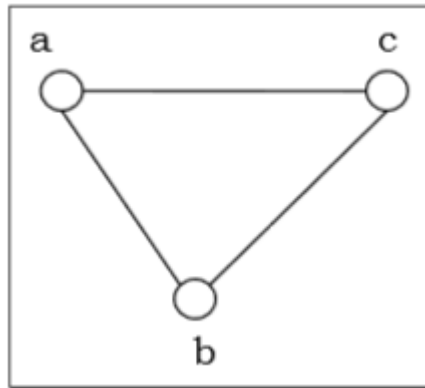
### Null Graph

A null graph has no edges. The null graph of  $n$  vertices is denoted by  $N_n$



### Simple Graph

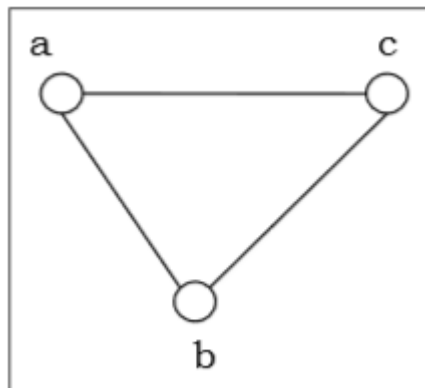
A graph is called simple graph/strict graph if the graph is undirected and does not contain any loops or multiple edges.



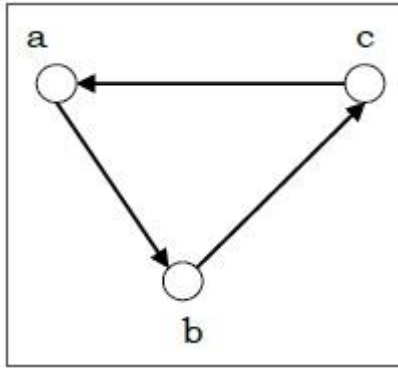
### Directed and Undirected Graph

A graph  $G=(V,E)$  is called a directed graph if the edge set is made of ordered vertex pair (a), and a graph is called undirected if the edge set is made of unordered vertex pair (b), and if directed edges called directed graph (b)

a)

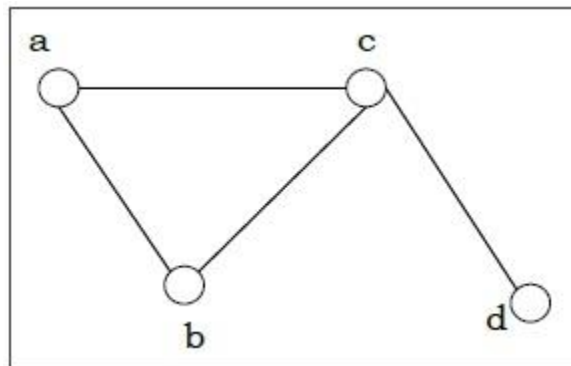


b)

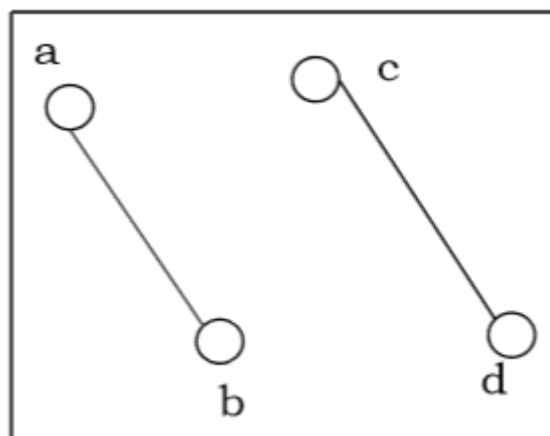


### Connected and Disconnected Graph

A graph is connected if any two vertices of the graph are connected by a path; while a graph is disconnected if at least two vertices of the graph are not connected by a path. If a graph  $G$  is disconnected, then every maximal connected sub graph of  $G$  is called a connected component of the graph  $G$ .



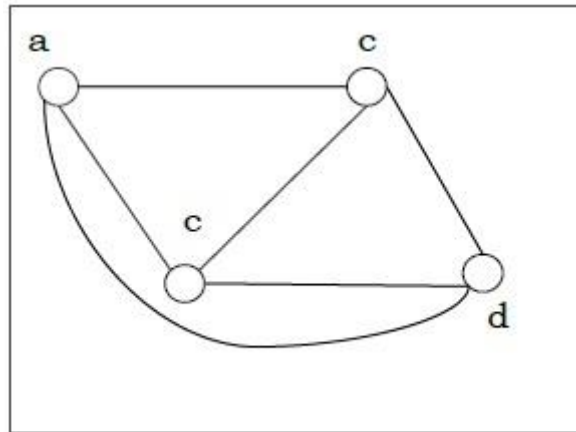
**Connected Graph**



**Disconnected Graph**

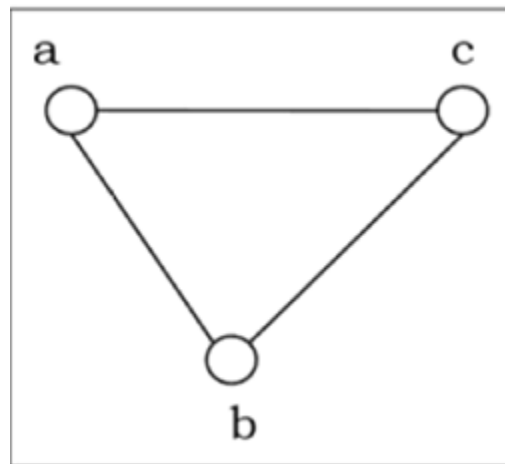
### Regular Graph

A graph is regular if all the vertices of the graph have the same degree. In a regular graph  $G$  of degree  $r$ , the degree of each vertex of  $G$  is  $r$ .



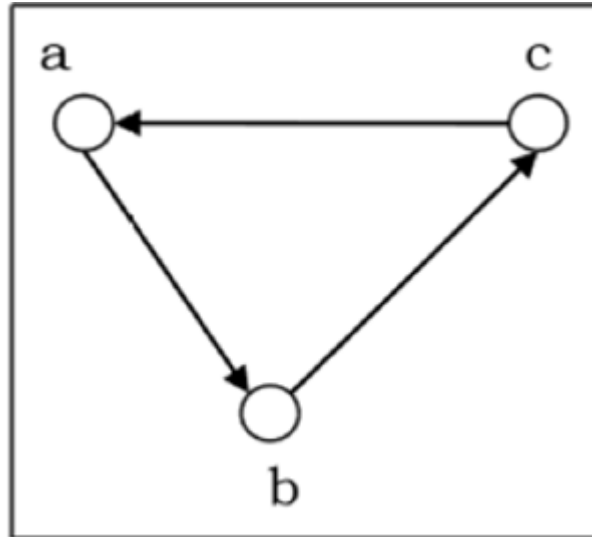
### Complete Graph

A graph is called complete graph if every two vertices pair are joined by exactly one edge. The complete graph with  $n$  vertices is denoted by  $K_n$ .



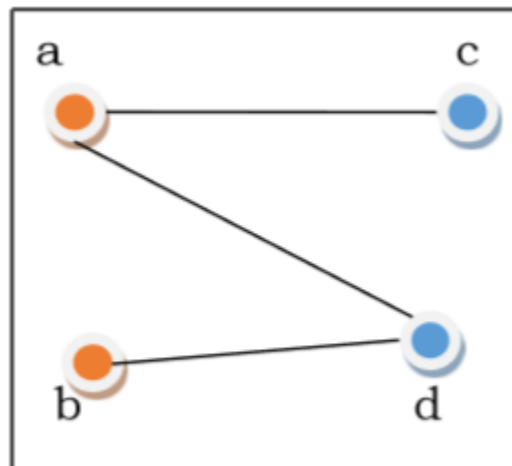
### Cycle Graph

If a graph consists of a single cycle, it is called cycle graph. The cycle graph with  $n$  vertices is denoted by  $C_n$ .



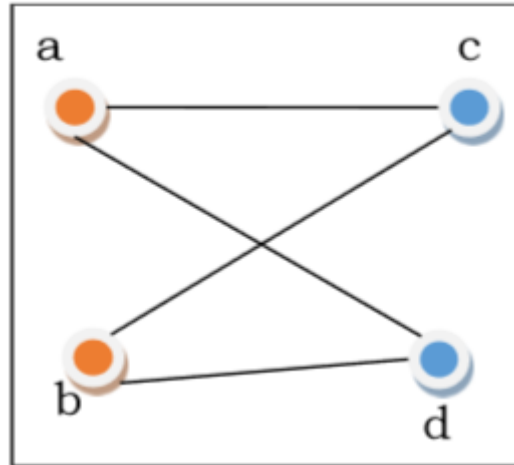
### Bipartite Graph

If the vertex-set of a graph  $G$  can be split into two disjoint sets,  $V_1$  and  $V_2$ , in such a way that each edge in the graph joins a vertex in  $V_1$  to a vertex in  $V_2$ , and there are no edges in  $G$  that connect two vertices in  $V_1$  or two vertices in  $V_2$ , then the graph  $G$  is called a bipartite graph.



### Complete Bipartite Graph

A complete bipartite graph is a bipartite graph in which each vertex in the first set is joined to every single vertex in the second set. The complete bipartite graph is denoted by  $K_{x,y}$  where the graph  $G$  contains  $x$  vertices in the first set and  $y$  vertices in the second set.



## Representation of Graphs

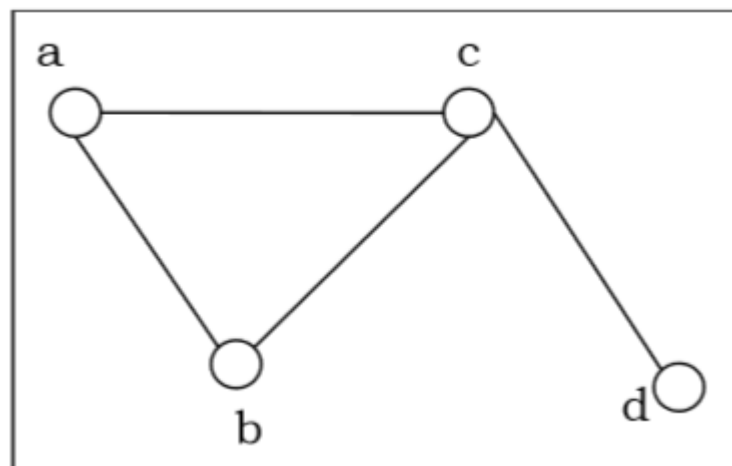
There are mainly two ways to represent a graph –

- Adjacency Matrix
- Adjacency List

### Adjacency Matrix

#### Adjacency Matrix of an Undirected Graph

Let us consider the following undirected graph and construct the adjacency matrix –

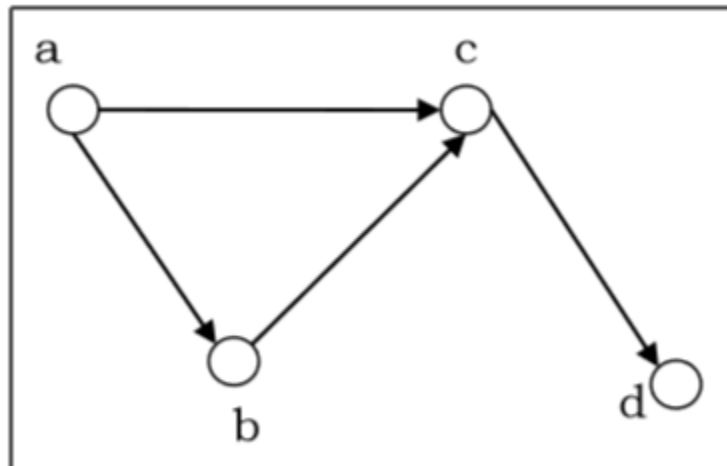


Adjacency matrix of the above undirected graph will be –

	<b>a</b>	<b>b</b>	<b>c</b>	<b>d</b>
<b>a</b>	0	1	1	0
<b>b</b>	1	0	1	0
<b>c</b>	1	1	0	1
<b>d</b>	0	0	1	0

### Adjacency Matrix of a Directed Graph

Let us consider the following directed graph and construct its adjacency matrix –

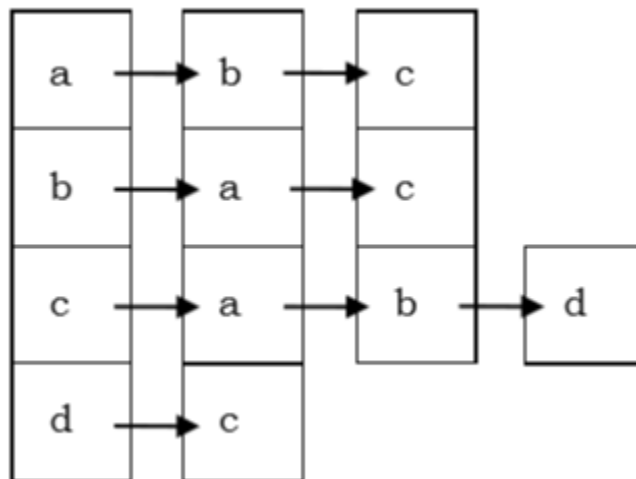




Adjacency matrix of the above directed graph will be –

	<b>a</b>	<b>b</b>	<b>c</b>	<b>d</b>
<b>a</b>	0	1	1	0
<b>b</b>	0	0	1	0
<b>c</b>	0	0	0	1
<b>d</b>	0	0	0	0

### Adjacency List



## Isomorphism

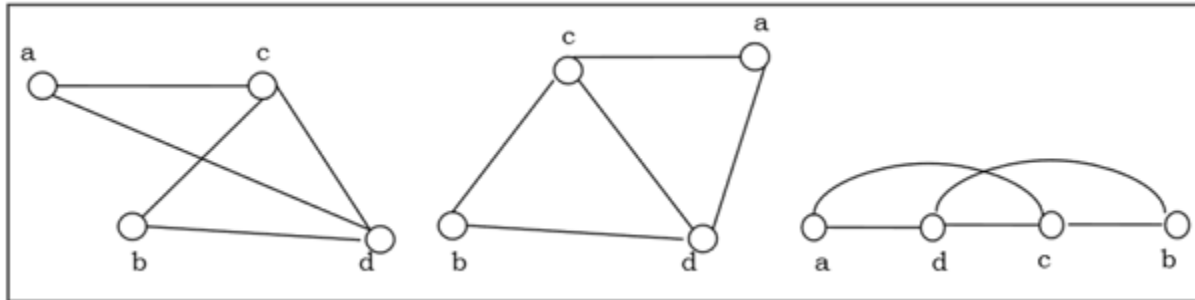
If two graphs G and H contain the same number of vertices connected in the same way, they are called isomorphic graphs (denoted by  $G \cong H$ ).

It is easier to check non-isomorphism than isomorphism. If any of these following conditions occurs, then two graphs are non-isomorphic –

- The number of connected components are different
- Vertex-set cardinalities are different
- Edge-set cardinalities are different
- Degree sequences are different

### Example

The following graphs are isomorphic –



## Euler Graphs

A connected graph  $G$  is called an Euler graph, if there is a closed trail which includes every edge of the graph  $G$ . An Euler path is a path that uses every edge of a graph exactly once. An Euler path starts and ends at different vertices.

An Euler circuit is a circuit that uses every edge of a graph exactly once. An Euler circuit always starts and ends at the same vertex. A connected graph  $G$  is an Euler graph if and only if all vertices of  $G$  are of even degree, and a connected graph  $G$  is Eulerian if and only if its edge set can be decomposed into cycles.

