

Week 7

Chapter 6: Conditional Processing

What's Next

- Boolean and Comparison Instructions
- Conditional Jumps
- **Conditional Loop Instructions**
- Conditional Structures
- Application: Finite-State Machines
- Conditional Control Flow Directives

Conditional Loop Instructions

- LOOPZ and LOOPE
- LOOPNZ and LOOPNE

LOOPZ and LOOPE

- Syntax:
 LOOPE *destination*
 LOOPZ *destination*
- Logic:
 - $ECX \leftarrow ECX - 1$
 - if $ECX > 0$ and $ZF=1$, jump to *destination*
- Useful when scanning an array for the first element that does **not** match a given value.

In 32-bit mode, ECX is the loop counter register. In 16-bit real-address mode, CX is the counter, and in 64-bit mode, RCX is the counter.

LOOPNZ and LOOPNE

- LOOPNZ (LOOPNE) is a conditional loop instruction
- Syntax:
 - LOOPNZ *destination*
 - LOOPNE *destination*
- Logic:
 - $ECX \leftarrow ECX - 1$;
 - if $ECX > 0$ and $ZF=0$, jump to *destination*
- Useful when scanning an array for the first element that matches a given value.

LOOPNZ Example

The following code finds the first positive value in an array:

```
.data
array SWORD -3,-6,-1,-10,10,30,40,4
sentinel SWORD 0
.code
    mov esi,OFFSET array
    mov ecx,LENGTHOF array
next:
    test WORD PTR [esi],8000h ; test sign bit
    pushfd ; push flags on stack
    add esi,TYPE array
    popfd ; pop flags from stack
    loopnz next ; continue loop
    jnz quit ; none found
    sub esi,TYPE array ; ESI points to value
quit:
```

Locate the first nonzero value in the array. If none is found, let ESI point to the sentinel value:

```
.data
array  SWORD 50 DUP(?)
sentinel SWORD 0FFFFh
.code
    mov esi,OFFSET array
    mov ecx,LENGTHOF array
L1: cmp WORD PTR [esi],0           ; check for zero
```

(fill in your code here)

```
quit:
```

... (solution)

```
.data
array SWORD 50 DUP(?)
sentinel SWORD 0FFFFh
.code
    mov esi,OFFSET array
    mov ecx,LENGTHOF array
L1: cmp WORD PTR [esi],0           ; check for zero
    pushfd                        ; push flags on stack
    add esi,TYPE array
    popfd                         ; pop flags from stack
    loope L1                      ; continue loop
    jz quit                       ; none found
    sub esi,TYPE array            ; ESI points to value
quit:
```


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Conditional Structures

- Block-Structured IF Statements
- Compound Expressions with AND
- Compound Expressions with OR
- WHILE Loops
- Table-Driven Selection

Block-Structured IF Statements

Assembly language programmers can easily translate logical statements written in C++/Java into assembly language. For example:

```
if( op1 == op2 )  
    X = 1;  
else  
    X = 2;
```

```
mov  eax,op1  
cmp  eax,op2  
jne  L1  
mov  X,1  
jmp  L2  
L1:  mov  X,2  
L2:
```

Implement the following pseudocode in assembly language. All values are unsigned:

```
if( ebx <= ecx )  
{  
    eax = 5;  
    edx = 6;  
}
```

```
cmp ebx,ecx  
ja  next  
mov eax,5  
mov edx,6  
next:
```

(There are multiple correct solutions to this problem.)

Implement the following pseudocode in assembly language. All values are 32-bit signed integers:

```
if( var1 <= var2 )
    var3 = 10;
else
{
    var3 = 6;
    var4 = 7;
}
```

```
mov eax,var1
cmp eax,var2
jle L1
mov var3,6
mov var4,7
jmp L2
L1: mov var3,10
L2:
```

(There are multiple correct solutions to this problem.)

Compound Expression with AND (1 of 3)

- When implementing the logical AND operator, consider that HLLs use short-circuit evaluation
- In the following example, if the first expression is false, the second expression is skipped:

```
if (a1 > b1) AND (b1 > c1)  
    x = 1;
```

Compound Expression with AND (2 of 3)

```
if (a1 > b1) AND (b1 > c1)
    X = 1;
```

This is one possible implementation . . .

```
        cmp  a1,b1                ; first expression...
        ja   L1
        jmp  next
L1:      cmp  b1,c1                ; second expression...
        ja   L2
        jmp  next
L2:      ; both are true
        mov  X,1                  ; set X to 1
next:
```

Compound Expression with AND (3 of 3)

```
if (a1 > b1) AND (b1 > c1)
    X = 1;
```

But the following implementation uses 29% less code by reversing the first relational operator. We allow the program to "fall through" to the second expression:

```
    cmp al,b1                ; first expression...
    jbe next                ; quit if false
    cmp bl,cl                ; second expression...
    jbe next                ; quit if false
    mov X,1                  ; both are true
next:
```


Implement the following pseudocode in assembly language. All values are unsigned:

```
if( ebx <= ecx
    && ecx > edx )
{
    eax = 5;
    edx = 6;
}
```

```
cmp ebx,ecx
ja  next
cmp ecx,edx
jbe next
mov eax,5
mov edx,6
next:
```

(There are multiple correct solutions to this problem.)

Compound Expression with OR (1 of 2)

- When implementing the logical OR operator, consider that HLLs use short-circuit evaluation
- In the following example, if the first expression is true, the second expression is skipped:

```
if (a1 > b1) OR (b1 > c1)  
    x = 1;
```

Compound Expression with OR (2 of 2)

```
if (a1 > b1) OR (b1 > c1)
    X = 1;
```

We can use "fall-through" logic to keep the code as short as possible:

```
    cmp al,b1                ; is AL > BL?
    ja  L1                  ; yes
    cmp bl,cl                ; no: is BL > CL?
    jbe next                ; no: skip next statement
L1: mov X,1                  ; set X to 1
next:
```

WHILE Loops

A WHILE loop is really an IF statement followed by the body of the loop, followed by an unconditional jump to the top of the loop. Consider the following example:

```
while( eax < ebx)
    eax = eax + 1;
```

This is a possible implementation:

```
top: cmp  eax, ebx          ; check loop condition
     jae  next             ; false? exit loop
     inc  eax              ; body of loop
     jmp  top              ; repeat the loop
next:
```

Implement the following loop, using unsigned 32-bit integers:

```
while( ebx <= val1)
{
    ebx = ebx + 5;
    val1 = val1 - 1
}
```

```
top: cmp ebx, val1          ; check loop condition
     ja  next              ; false? exit loop
     add ebx, 5             ; body of loop
     dec val1
     jmp top                ; repeat the loop
next:
```

Table-Driven Selection (1 of 4)

- Table-driven selection uses a table lookup to replace a multiway selection structure
- Create a table containing lookup values and the offsets of labels or procedures
- Use a loop to search the table
- Suited to a large number of comparisons

Table-Driven Selection (2 of 4)

Step 1: create a table containing lookup values and procedure offsets:

```
.data
CaseTable BYTE 'A'           ; lookup value
          DWORD Process_A     ; address of procedure
          EntrySize = ($ - CaseTable)
          BYTE 'B'
          DWORD Process_B
          BYTE 'C'
          DWORD Process_C
          BYTE 'D'
          DWORD Process_D
```

```
NumberOfEntries = ($ - CaseTable) / EntrySize
```

Table-Driven Selection (3 of 4)

Table of Procedure Offsets:

'A'	00000120	'B'	00000130	'C'	00000140	'D'	00000150
-----	----------	-----	----------	-----	----------	-----	----------

The diagram illustrates a table-driven selection process. A horizontal table contains four entries: 'A' with offset 00000120, 'B' with offset 00000130, 'C' with offset 00000140, and 'D' with offset 00000150. Below the table, a line labeled 'lookup value' points to the 'B' entry. Another line labeled 'address of Process_B' points to the offset 00000130 for entry 'B'.

Table-Driven Selection (4 of 4)

Step 2: Use a loop to search the table. When a match is found, call the procedure offset stored in the current table entry:

```
        mov ebx,OFFSET CaseTable      ; point EBX to the table
        mov ecx,NumberOfEntries      ; loop counter

L1:  cmp al,[ebx]                    ; match found?
     jne L2                          ; no: continue
     call NEAR PTR [ebx + 1]         ; yes: call the procedure
     call WriteString                ; display message
     call Crlf
     jmp L3                          ; and exit the loop
L2:  add ebx,EntrySize               ; point to next entry
     loop L1                         ; repeat until ECX = 0

L3:                                     required for
                                     procedure pointers
```

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Application: Finite-State Machines

- A finite-state machine (FSM) is a graph structure that changes state based on some input. Also called a **state-transition diagram**.
- We use a graph to represent an FSM, with squares or circles called **nodes**, and lines with arrows between the circles called **edges**.

Application: Finite-State Machines

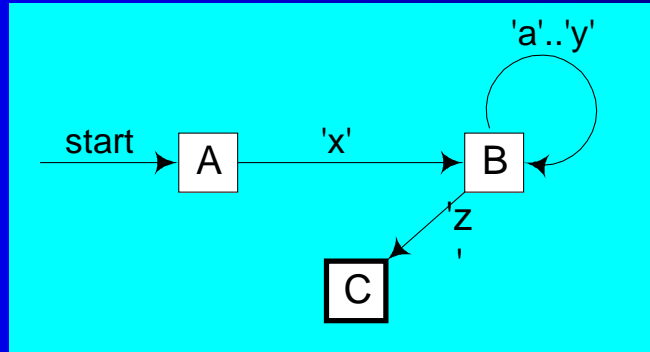
- A FSM is a specific instance of a more general structure called a **directed graph**.
- Three basic states, represented by nodes:
 - Start state
 - Terminal state(s)
 - Nonterminal state(s)

Finite-State Machine

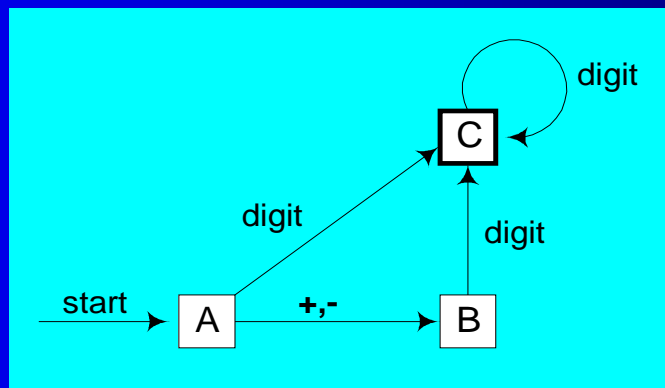
- Accepts any sequence of symbols that puts it into an accepting (final) state
- Can be used to recognize, or validate a sequence of characters that is governed by language rules (called a regular expression)
- Advantages:
 - Provides visual tracking of program's flow of control
 - Easy to modify
 - Easily implemented in assembly language

Finite-State Machine Examples

- FSM that recognizes strings beginning with 'x', followed by letters 'a'..'y', ending with 'z':

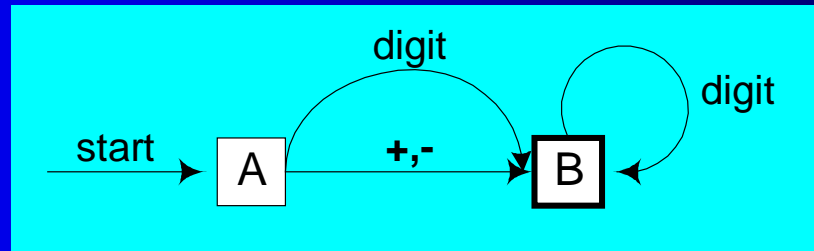


- FSM that recognizes signed integers:



Your Turn . . .

- Explain why the following FSM does not work as well for signed integers as the one shown on the previous slide:



Implementing an FSM

The following is code from State A in the Integer FSM:

StateA:

```
    call Getnext           ; read next char into AL
    cmp al, '+'            ; leading + sign?
    je StateB             ; go to State B
    cmp al, '-'            ; leading - sign?
    je StateB             ; go to State B
    call IsDigit           ; ZF = 1 if AL = digit
    jz StateC             ; go to State C
    call DisplayErrorMsg   ; invalid input found
    jmp Quit
```

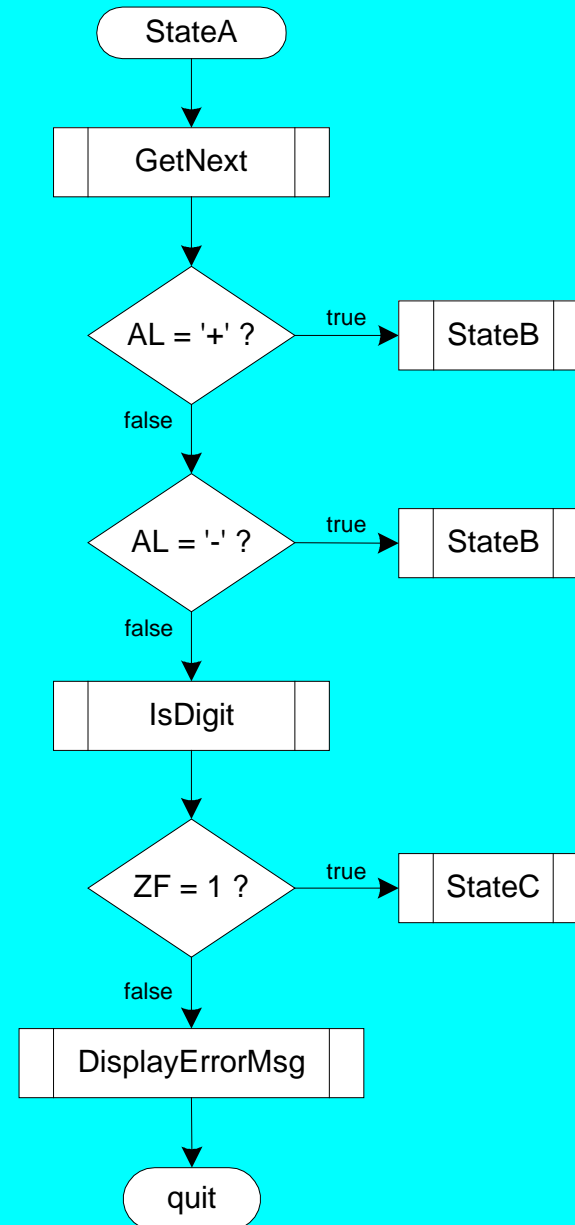

IsDigit Procedure

Receives a character in AL. Sets the Zero flag if the character is a decimal digit.

```
IsDigit PROC
    cmp     al,'0'                ; ZF = 0
    jb      ID1
    cmp     al,'9'                ; ZF = 0
    ja      ID1
    test    al,0                  ; ZF = 1
ID1: ret
IsDigit ENDP
```

Flowchart of State A

State A accepts a plus or minus sign, or a decimal digit.



- Draw a FSM diagram for hexadecimal integer constant that conforms to MASM syntax.
- Draw a flowchart for one of the states in your FSM.
- Implement your FSM in assembly language. Let the user input a hexadecimal constant from the keyboard.

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Creating IF Statements

- Runtime Expressions
- Relational and Logical Operators
- MASM-Generated Code
- .REPEAT Directive
- .WHILE Directive

Runtime Expressions

- .IF, .ELSE, .ELSEIF, and .ENDIF can be used to evaluate runtime expressions and create block-structured IF statements.
- Examples:

```
.IF eax > ebx
    mov edx,1
.ELSE
    mov edx,2
.ENDIF
```

```
.IF eax > ebx && eax > ecx
    mov edx,1
.ELSE
    mov edx,2
.ENDIF
```

- MASM generates "hidden" code for you, consisting of code labels, CMP and conditional jump instructions.


Relational and Logical Operators

Operator	Description
<i>expr1</i> == <i>expr2</i>	Returns true when <i>expression1</i> is equal to <i>expr2</i> .
<i>expr1</i> != <i>expr2</i>	Returns true when <i>expr1</i> is not equal to <i>expr2</i> .
<i>expr1</i> > <i>expr2</i>	Returns true when <i>expr1</i> is greater than <i>expr2</i> .
<i>expr1</i> >= <i>expr2</i>	Returns true when <i>expr1</i> is greater than or equal to <i>expr2</i> .
<i>expr1</i> < <i>expr2</i>	Returns true when <i>expr1</i> is less than <i>expr2</i> .
<i>expr1</i> <= <i>expr2</i>	Returns true when <i>expr1</i> is less than or equal to <i>expr2</i> .
! <i>expr</i>	Returns true when <i>expr</i> is false.
<i>expr1</i> && <i>expr2</i>	Performs logical AND between <i>expr1</i> and <i>expr2</i> .
<i>expr1</i> <i>expr2</i>	Performs logical OR between <i>expr1</i> and <i>expr2</i> .
<i>expr1</i> & <i>expr2</i>	Performs bitwise AND between <i>expr1</i> and <i>expr2</i> .
CARRY?	Returns true if the Carry flag is set.
OVERFLOW?	Returns true if the Overflow flag is set.
PARITY?	Returns true if the Parity flag is set.
SIGN?	Returns true if the Sign flag is set.
ZERO?	Returns true if the Zero flag is set.

Signed and Unsigned Comparisons

```
.data
val1    DWORD 5
result  DWORD ?
.code
mov eax,6
.IF eax > val1
    mov result,1
.ENDIF
```

Generated code:




```
mov eax,6
cmp eax,val1
jbe @C0001
mov result,1
@C0001:
```

MASM automatically generates an unsigned jump (JBE) because **val1** is unsigned.

Signed and Unsigned Comparisons

```
.data
val1    SDWORD 5
result  SDWORD ?
.code
mov eax,6
.IF eax > val1
    mov result,1
.ENDIF
```

Generated code:



```
mov eax,6
cmp eax,val1
jle @C0001
mov result,1
@C0001:
```

MASM automatically generates a signed jump (JLE) because **val1** is signed.

Signed and Unsigned Comparisons

```
.data
result DWORD ?
.code
mov ebx,5
mov eax,6
.IF eax > ebx
    mov result,1
.ENDIF
```

Generated code:

```
mov ebx,5
mov eax,6
cmp eax,ebx
jbe @C0001
mov result,1
@C0001:
```

MASM automatically generates an unsigned jump (JBE) when both operands are registers . . .

Signed and Unsigned Comparisons

```
.data
result SDWORD ?
.code
mov ebx,5
mov eax,6
.IF SDWORD PTR eax > ebx
    mov result,1
.ENDIF
```

Generated code:

```
mov ebx,5
mov eax,6
cmp eax,ebx
jle @C0001
mov result,1
@C0001:
```

... unless you prefix one of the register operands with the SDWORD PTR operator. Then a signed jump is generated.

.REPEAT Directive

Executes the loop body before testing the loop condition associated with the .UNTIL directive.

Example:

```
; Display integers 1 - 10:

mov eax,0
.REPEAT
    inc eax
    call WriteDec
    call Crlf
.UNTIL eax == 10
```

.WHILE Directive

Tests the loop condition before executing the loop body The .ENDW directive marks the end of the loop.

Example:

```
; Display integers 1 - 10:

mov eax,0
.WHILE eax < 10
    inc eax
    call WriteDec
    call Crlf
.ENDW
```

Summary

- Bitwise instructions (AND, OR, XOR, NOT, TEST)
 - manipulate individual bits in operands
- CMP – compares operands using implied subtraction
 - sets condition flags
- Conditional Jumps & Loops
 - equality: JE, JNE
 - flag values: JC, JZ, JNC, JP, ...
 - signed: JG, JL, JNG, ...
 - unsigned: JA, JB, JNA, ...
 - LOOPZ, LOOPNZ, LOOPE, LOOPNE
- Flowcharts – logic diagramming tool
- Finite-state machine – tracks state changes at runtime