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SER-501-A2

Q1) a) max subarray by Brute force

```
def find_maximum_subarray_brute(A):
  Return a tuple (i,j) where A[i:j] is the maximum subarray.
  time complexity = O(n^2)
  max = 0
  start = 0
  end = 0
  if (len(A) == 0):
    return None
  # loop running till the length of the array
  for i in range(0, len(A)):
    sum = 0
    j = i
    sum = sum + A[i]
    # loop for taking in account the size of array we are considering
    for j in range(j, len(A)):
      sum = sum + A[j]
      if sum > max:
         max = sum
         start = i
         end = i
  tup = (start, end)
  return tup
```

Q1) b) maximum crossing subarray

```
def find_maximum_crossing_subarray(A, low, mid, high):
    """
    Find the maximum subarray that crosses mid
    Return a tuple ((i, j), sum) where sum is the maximum subarray of A[i:j].
    """
    max_left = 0
    max_right = 0
    left_sum = 0
    sum = 0
    i = mid
```

```
# taking sum of left side
  while i \ge low:
    sum = sum + A[i]
    if (sum > left_sum):
       left sum = sum
       max_left = i
    i = i - 1
  j = mid + 1
  right sum = 0
  sum = 0
  # taking sum of right side
  while (i \le high):
    sum = sum + A[i]
    if (sum > right sum):
       right\_sum = sum
       max_right = i
    j = j + 1
  # returning sum of both sides
  tup = (max_left, max_right)
  tup1 = (tup, left_sum + right_sum)
  return tup1
# The recursive method to solve max subarray problem
def find_maximum_subarray_recursive_helper(A, low=0, high=-1):
  Return a tuple ((i, j), sum) where sum is the maximum subarray of A[i:j].
  # base case
  if (low == high):
    tup = (low, high)
    return (tup, A[low])
  # finding midpoint of the array
  mid = ((low + high) // 2)
  # calling function for left array
  ((left_low, left_high),
  left_sum) = find_maximum_subarray_recursive_helper(A, low, mid)
  # calling function for right array
  ((right_low, right_high),
  right_sum) = find_maximum_subarray_recursive_helper(A, mid + 1, high)
  # calling function for crossing array
  ((cross_low, cross_high),
  cross_sum) = find_maximum_crossing_subarray(A, low, mid, high)
  print(cross_sum)
  # if the sum of left array is greater
  if (left_sum >= right_sum) and (left_sum >= cross_sum):
     tup_left = (left_low, left_high)
```

```
tup_left1 = (tup_left, left_sum)
    return (tup_left1)
  # if the sum of right array is greater
  elif (right_sum >= left_sum) and (right_sum >= cross_sum):
    tup_right = (right_low, right_high)
    tup_right1 = (tup_right, right_sum)
    return (tup_right1)
  # if the sum of crossing array is greater
    tup_cross = (cross_low, cross_high)
    tup_cross1 = (tup_cross, cross_sum)
    return (tup_cross1)
# The recursive method to solve max subarray problem
def find_maximum_subarray_recursive(A):
  Return a tuple (i,j) where A[i:j] is the maximum subarray.
  if (len(A) == 0):
    return None
  # if the array is not of length of power2 we apply padding to left
  A = PadLeft(A)
  return find maximum subarray recursive helper(A, 0, len(A) - 1)[0]
def PadLeft(A):
  # checking the closest power of 2 to the len(A)
  nextPower = NextPowerOfTwo(len(A))
  deficit = int(math.pow(2, nextPower) - len(A))
  # concatenate 0 with the original array
  for x in range(0, deficit):
    A = np.concatenate((A, [0]))
  return A
def NextPowerOfTwo(number):
  # Returns next power of two following 'number'
  return math.ceil(math.log(number, 2))
Q1) c) max subarray by iterative method
def find_maximum_subarray_iterative(A):
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```
Return a tuple (i,j) where A[i:j] is the maximum subarray.
  arr sum = 0
  arr_val = 0
  left = -1
  right = -1
  left1 = 0
  if (len(A) == 0):
    return None
  # loop runing till the length of A
  for i in range(0, (len(A))):
     # keeping sum after every index
     arr\_sum = arr\_sum + A[i]
     # comparing two sums
     if (arr_sum > arr_val):
       arr_val = arr_sum
       left = left1
       right = i
     # if arr_sum get less then 0 we reset arr_sum to zero
    if arr_sum < 0:
       arr\_sum = 0
       left1 = i + 1
  tup = (left, right)
  return tup
Q2) a) matrix multiplication
def square_matrix_multiply(A, B):
  Return the product AB of matrix multiplication.
  A = asarray(A)
  B = asarray(B)
  assert A.shape == B.shape
  assert A.shape == A.T.shape
  n = len(A)
  # intializing an n*n matrix c
  C = [[0 \text{ for } i \text{ in } range(n)] \text{ for } j \text{ in } range(n)]
  # assigning values to c
  for i in range(n):
     for j in range(n):
       for k in range(n):
          C[i][j] = C[i][j] + (A[i][k] * B[k][j])
  C = np.array(C)
  return C
```

Q2) b) matrix multiplication by Strassen method

```
def square_matrix_multiply_strassens(A, B):
  Return the product AB of matrix multiplication.
  Assume len(A) is a power of 2
  A = asarray(A)
  B = asarray(B)
  assert A.shape == B.shape
  assert A.shape == A.T.shape
  assert (len(A) & (len(A) - 1)) == 0, " A is not a power of 2 "
  # base case
  if len(A) == 1:
    base = [[0]]
    base[0][0] = A[0][0] * B[0][0]
    return base
  else:
    # spliting matrices into 4
    A11, A12, A21, A22 = split(A)
    B11, B12, B21, B22 = split(B)
    # Finding the values of M to be applied
    # M1=(A11+A22)*(B11+B22)
    M1 = square_matrix_multiply_strassens(np.add(A11, A22),
                           np.add(B11, B22))
    \# M2 = (A21 + A22)*(B11)
    M2 = square_matrix_multiply_strassens(np.add(A21,
                               A22), B11)
    # M3=(A11)*(B12-B22)
    M3 = square matrix multiply strassens(A11,
                           np.subtract(B12, B22))
    # M4=(A22)*(B21-B11)
    M4 = square_matrix_multiply_strassens(A22, np.subtract(B21, B11))
    \# M5 = (A11 + A12)*(B22)
    M5 = square_matrix_multiply_strassens(np.add(A11, A12), B22)
    # M6=(A21-A11)*(B11+B12)
    M6 = square\_matrix\_multiply\_strassens(np.subtract)
                           (A21, A11), np.add(B11, B12))
    # M7=(A12-A22)*(B21+B22)
    M7 = square matrix multiply strassens(np.subtract
                           (A12, A22), np.add(B21, B22))
    # assingning values of the above to the final matrix
    X11 = np.add(np.subtract(np.add(M1, M4), M5), M7)
    X12 = np.add(M3, M5)
    X21 = np.add(M2, M4)
    X22 = \text{np.add(np.subtract(np.add(M1, M3), M2), M6)}
    c = [[0 \text{ for row in range}(len(X11) * 2)] \text{ for col in range}(len(X11) * 2)]
```

```
for i in range(len(X11)):
         for j in range(len(X11)):
           c[i][j] = X11[i][j]
           c[i][j + len(X11)] = X12[i][j]
           c[i + len(X11)][j] = X21[i][j]
           c[i + len(X11)][j + len(X11)] = X22[i][j]
      c = np.array(c)
      return c
   pass
def split(matrix):
   # split matrix into quarters
   m = len(matrix) // 2
   a = [[0 \text{ for } i \text{ in } range(m)] \text{ for } j \text{ in } range(m)]
   b = [[0 \text{ for i in range}(m)] \text{ for j in range}(m)]
   c = [[0 \text{ for i in range}(m)] \text{ for j in range}(m)]
   d = [[0 \text{ for i in range}(m)] \text{ for j in range}(m)]
   for i in range(0, m):
      for j in range(0, m):
         a[i][j] = matrix[i][j]
         b[i][j] = matrix[i][m + j]
        c[i][j] = matrix[i + m][j]
        d[i][j] = matrix[i + m][m + j]
        j = j + 1
     j = j + 1
   return a, b, c, d
```

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<u>Output</u>

```
In [42]: runfile('I:/MS/SER501/Assignment2/assignment_2.py', wdir='I:/MS/SER501/Assignment2')
Q1a) maximum subarray by brute force
(7, 10)
Q1b) maximum subarray by recursive method
(7, 10)
Q1c) maximum subarray by iterative method
(7, 10)
Q2a) matrix multiplication
[[ 96 68 69 69]
[ 24 56 18 52]
[ 58 95 71 92]
[ 90 107 81 142]]
Q2b) matrix multiplication by strassen method
[[ 96 68 69 69]
[ 24 56 18 52]
[ 58 95 71 92]
[ 90 107 81 142]]
In [43]:
```

Results after running flake8 and McCabe complexity command

```
C:\WINDOWS\system32\cmd.exe

(I:\python) C:\Users\Abdul Samad Khan\Documents>flake8 I:\MS\SER501\Assignment2

(I:\python) C:\Users\Abdul Samad Khan\Documents>flake8 --max-complexity 10 I:\MS\SER501\Assignment2

(I:\python) C:\Users\Abdul Samad Khan\Documents>
```