

Programming Paradigms CSI2120 – Winter 2018

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Functional Programming in Lisp

- **Language designed by John McCarthy between 1956 - 1959 at MIT for applications related to artificial intelligence**
 - one of the oldest languages still in use
- **LISP = LISt Processor**
- **Derived from λ -calculus.**
 - λ -calculus allows functions to be the values of an expression
- **Many dialects**
 - Lisp 1.5 (1960), Scheme (1975), Common Lisp (1985)...
- **Rich language: functional, symbolic.**
- **Syntax and semantics are simple and uniform**

Creation of Lisp

- **1960: McCarthy published his paper on Lisp**
- **Lisp/Scheme has a few simple operators and a rich notation for functions.**
 - But this is combined with simple data structures
- **As a result we have a full and expressive programming language**

Nine Key Concepts

1. **Conditions (if-then-else)**
2. **Functions as data types**
3. **Recursions**
4. **Variables as pointers**
5. **Automatic garbage collection**
6. **A program is an expression (not a series of statements)**
7. **Symbols or atoms**
8. **Lists and trees**
9. **Complete language available at all times (read-eval-print)**

Pure Functional Programming

- **A program corresponds to a function call**
- **A function is a composition of functions**
- **Functions are non-causal**
 - Depend only on the parameters passed
- **No variables, no assignments**
- **No loops, no control statement**
 - Beyond the if-then-else function

Functional Programming in Practise

- **Some additions to pure functional programming**
 - Local definition of certain values
 - Assignments (lexically scoped variables)
 - Use an execution sequence (in order to break up the program).

Functional Programming in Scheme

- **Scheme is LISP dialect designed at MIT in 1975, mainly for education**
- **Initially small**
 - But it is now a complete language.
- **Standardized by ANSI / IEEE**
 - Language continues to evolve
- **Commonly used as interpreted language**
 - But may also be compiled to be executed efficiently.

Application Area of Functional Programming

- **Applications of symbolic computation, i.e., non-numerical applications**
 - Artificial intelligence (expert systems, natural language interfaces, ...)
 - Automated reasoning (theorem proving, proofs of programs ...)
 - Symbolic Computation
 - Games

Basic Concepts

- **The list is the fundamental data structure**
- **Atom: a number, a string or a symbol.**
 - All data types are equal
- **An expression is either an atom or a list**
- **A List is a series of expressions in parenthesis**
 - Including the empty list ()
- **A function is a first class object (first-class data) that can be created, assigned to variables, passed as a parameter or returned as a value.**

Evaluations of Expressions

- **Constants are evaluated to what they are.**
 - numbers
 - > 2 => 2
 - strings
 - > "Hello" => "Hello"
- **Identifiers evaluate to the value that is attributed to them.**
 - > * => #<procedure:*>
- **Lists are evaluated by**
 - first evaluating the head, i.e., the first expression; the value of this expression must be a function
 - The arguments of this function are the values obtained by evaluating the expressions contained in the rest of the list

A First Scheme Session

- In its simplest form, a Scheme interpreter session uses the interactive READ-EVAL-PRINT programming model (REPL)
- Example:

```
> (+ 3 4)
```

```
7
```

- In the example, we have a list.
 - The first entry is the function +
 - The rest of the list are constant expressions 3 and 4.
- The list is read as `(+ 3 4)`, evaluated and then the result 7 printed.

Scheme Interpreter

- **Classical MIT Scheme interpreter**
 - <http://www.gnu.org/software/mit-scheme/>
 - While available, not well supported under windows
- **Racket**
 - Another LISP dialect
- **DrRacket**
 - <http://racket-lang.org/>
 - Convenient and full-fledged programming environment to run LISP dialects

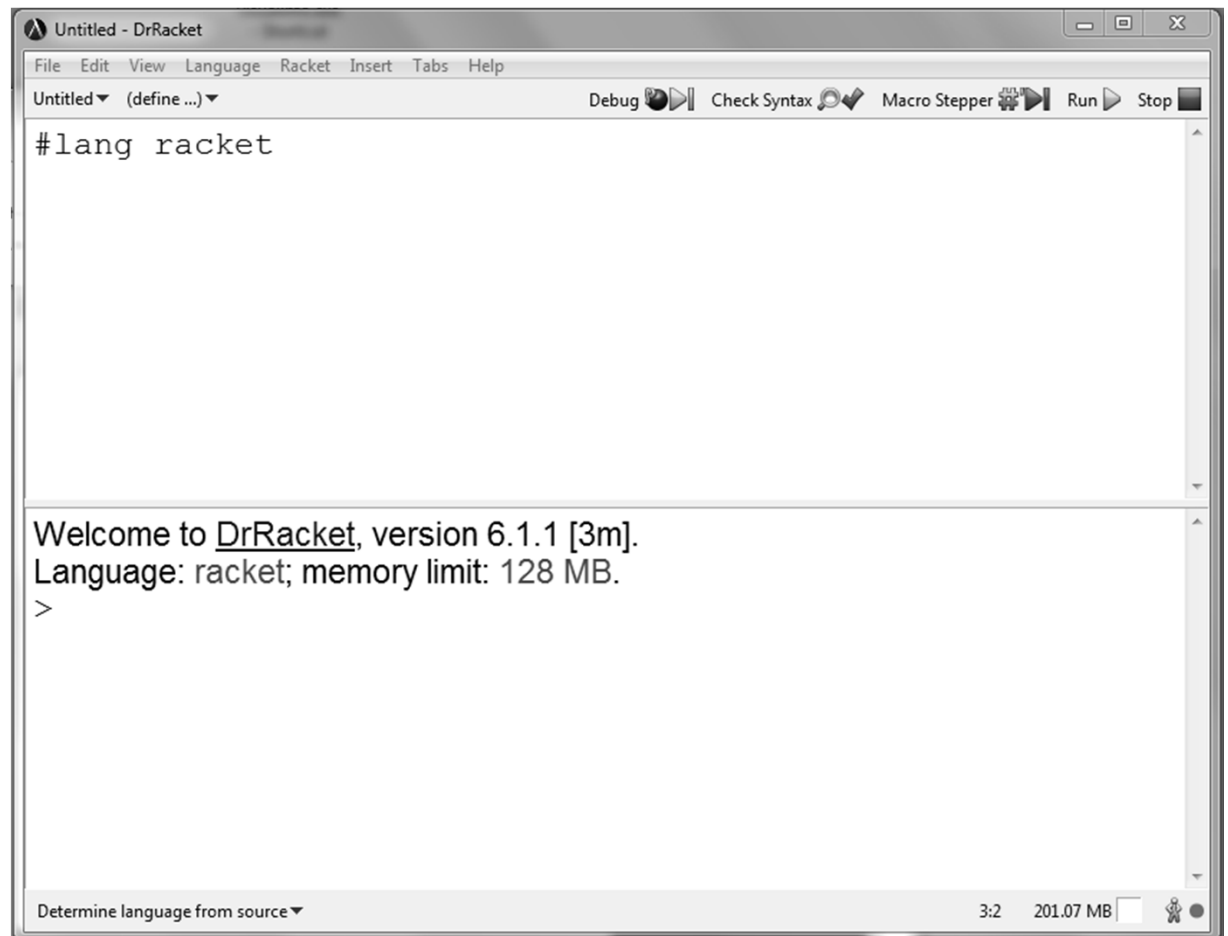
Few Remarks on the IDE

You must select a language in top of window.

Program editor

Bottom window is for running programs

This is the classic REPL Buffer



Evaluation of Expressions

- **The prefix notation is used in expressions**
 - Instead of inline operators as in $3+4*5$
 - One needs to write $(+ 3 (* 4 5))$
- **To evaluate an expression, all sub-expressions must be evaluated first.**
 - The evaluation follows normal order evaluation
 - Brackets determine the order and are not optional

$(+ 3 (* 4 5))$

$(+ 3 20)$

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Special Syntactic Forms

- Some expressions do not obey the normal order evaluation, these expressions are said to have a *special syntactic form*.
- The evaluation of their arguments is deferred until the required values are known.
- The main special syntactic forms are:
 - if statement
 - conditional branching
 - creation of local scope
 - quotation

Special Syntactic Forms: The Alternative (if statement)

```
(if (= x 0) infinity (/ 1 x))
```

- The expression following the if is evaluated first .
- If its value is true (# t) then
 - the second argument is evaluated
 - its value is returned without evaluating the third argument
- if its value is false (# f) then
 - the third argument is evaluated and returned

Special Syntactic Forms: 2.

Conditional Branching

- Conditional expressions are similar to `if` but more than two branches can be specified

```
(cond ((< x xmin) xmin) ((> x xmax) xmax) (#t x))
```

- The `cond` expression is followed by a series of lists composed of two expressions.
 - If the first of the two expressions in one of these lists evaluates to `#t` then the value of the second expression is returned
 - If the first expression evaluates to false, then evaluation proceeds to the next list.
 - If no lists evaluates to `#t` then `nil` is returned.

Example: Conditional and Top-level Define

Definition of the function `showIt` taking an argument `pts` and evaluating a conditional

```
(define (showIt pts)
  (cond ((or (<= pts 3) (>= pts 65)) 0)
        ((<= 4 pts 6) 0.5)
        ((<= 7 pts 12) 1.0)
        ((<= 13 pts 15) 1.5)
        ((<= 16 pts 18) 1.8)
        (else 2.0)))
```

=> `showIt`

```
(showIt 25)
```

=> 2.0

Special Syntactic Forms: 3. Creating Local Scope

- **Let Expressions**

```
(let ((a 3) (b 4)) (* a b))
```

=> 12

- The first argument of a let expression is a list of links created between an identifier and a value
- These links are only valid for the evaluation of the following expression(s)
 - There can be several to allow the execution of a sequence.
 - However, it evaluates to the last expression.

```
(let ((a 3) (b 4)) (* a b) (+ a b))
```

=> 7

Special Syntactic Forms: 4. Quotations

- The quote function ensures that an argument list is not evaluated.

```
(quote (1 2 3))
```

```
=> (1 2 3)
```

- But the list is rather returned as is.
- Quotation is necessary here, otherwise the first expression of a list needs to evaluate to a function.

```
`(+ 3 4) => (+ 3 4)
```

```
(+ 3 4) => 7
```

- The quote function can simply written with a ' :

```
'(1 2 3)
```

```
=> (1 2 3)
```

Quotation Example

```
(let ((a '(1 2 3)) (b '(3 4 5))) (cons a b))
```

equates to

```
(cons '(1 2 3) '(3 4 5))
```

⇒ ((1 2 3) 3 4 5)

- The function `cons` is the dot operator for lists, i.e., it puts the first expression as the head of the second list expression
- The list `(1 2 3)` becomes the first element in the combined list `((1 2 3) 3 4 5)`
- (Much) more on list processing soon.

Example: Building a List with the Function `list`

```
(list `a `b `c)
```

```
⇒ (a b c)
```

```
(list `(a b c))
```

```
⇒ ((a b c))
```

Lambdas

- **Lambda expressions are “local functions”**

The expression `(lambda (var1, var2, ...) exp1 exp2 ...)` returns a function and applies it to the variables that are the parameters to the function expressions, e.g.,

```
((lambda (x) (* x x)) 3)
```

⇒ 9

Multiple variables and multiple expression (result of last expression is the evaluation result)

```
((lambda (f x y) (f x x) (f x y) (f y y)) +  
2 3 )
```

⇒ 6

Function Definitions

A definition associates a function expression to a name:

```
(define square (lambda (x) (* x x)))
```

or equivalently:

```
(define (square x) (* x x))
```

Use of define, here procedure square:

```
(square 2)
```

$\Rightarrow 4$

Example: Factorial

Top-level Function Definition

```
(define (fact n)
  (if (> n 0)
      (* n (fact (- n 1)))
      1)
  )
```

=> fact

```
(fact 35)
```

=> 10333147966386144929666651337523200000000

Another Example Function

Conversion from degrees Fahrenheit to Celsius

```
(define (F-a-C temperature)
  (/ (- temperature 32) 1.8))
```

=> F-a-c

```
(F-a-C 95)
```

=> 35

```
(define freezing 32)
```

=> freezing

```
(F-a-C freezing)
```

=> 0

Function Definitions with Lambdas

- We have seen Lambdas can be combined with top-level defines

```
(define fct (lambda (f x) (f x x)))
```

```
(fct + 13)
```

```
=> 26
```

```
(fct * 4)
```

```
=> 16
```

- Combine with let binding: `x` is a let-bound variable in the enclosing scope

```
(let ((x `a))
```

```
  (let ((f (lambda (y) (list x y))))
```

```
    (f `b)))
```

```
=> (a b)
```

Lambda Expression and Let-Bound Variables

```
(let ((x 2) (y 3)) (+ x y))
```

is equivalent to

```
((lambda (x y) (+ x y)) 2 3)
```

In general:

```
((let (var val) ...) expr...)  $\equiv$  ((lambda (var ...)
  expr...) val...)
```

Example: Greatest Common Divisor (GCD)

- top level define for gcd function

```
(define gcd
  (lambda (a b)
    (if (= a b)
        a
        (if (> a b)
            (gcd (- a b) b)
            (gcd a (- b a))))))
```

=> gcd

```
(gcd 12 15)
```

=> 3

Summary

- **Introduction to Functional Programming**
- **Basic Scheme**
- **Special Syntactic Forms**
 - Alternative (if then else)
 - Conditional
 - Local Scope
 - Quotation
- **Let-bound variables**
- **Top-level (function) definitions**
- **Lambda expressions**