# GRAPH SEARCH ALGORITHMS

>DFS(Depth First Search)

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#### Depth First Search

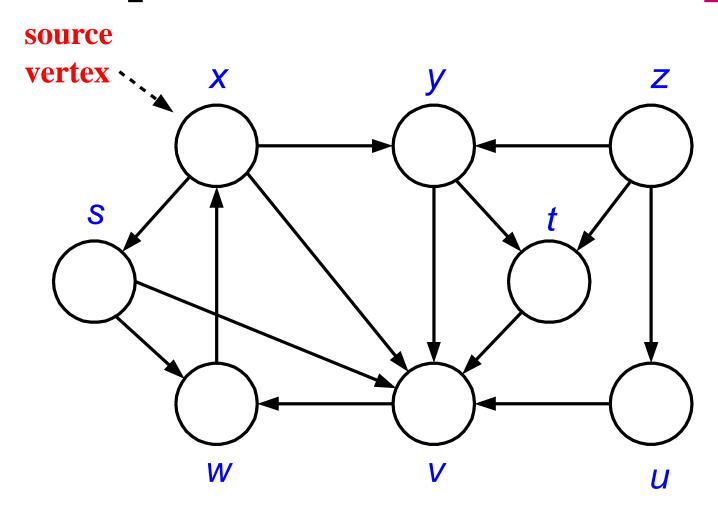
- □ DFS, go as far as possible along a single path until it reaches a dead end (that is a vertex with no edge out or no neighbor unexplored) then backtrack
- As the name implies the DFS search deeper in the graph whenever possible. DFS explores edges out of the most recently discovered vertex v that still has unexplored edges leaving it. Once all of v's edges have been explored, the search backtracks to explore edges leaving the vertex from which v was discovered.

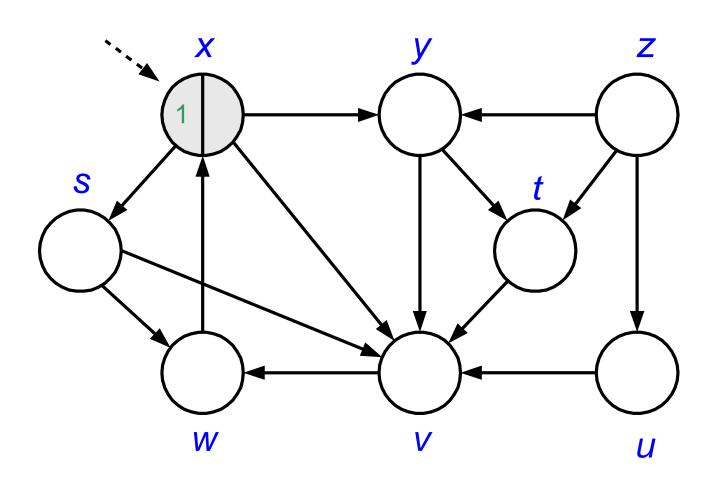
#### Depth-First Search

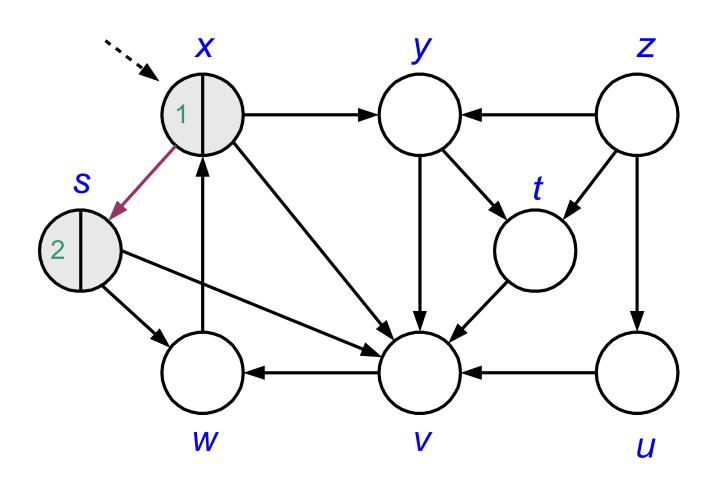
- Graph G = (V, E) directed or undirected
- Adjacency list representation
- Goal: Systematically explore every vertex and every edge
- Idea: search deeper whenever possible
  - Using a Stack(LIFO)

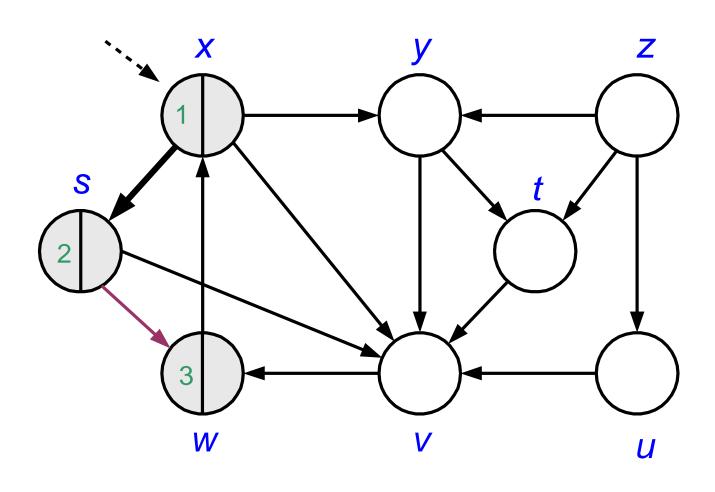
#### Information Needed to Maintain in DFS Algorithm

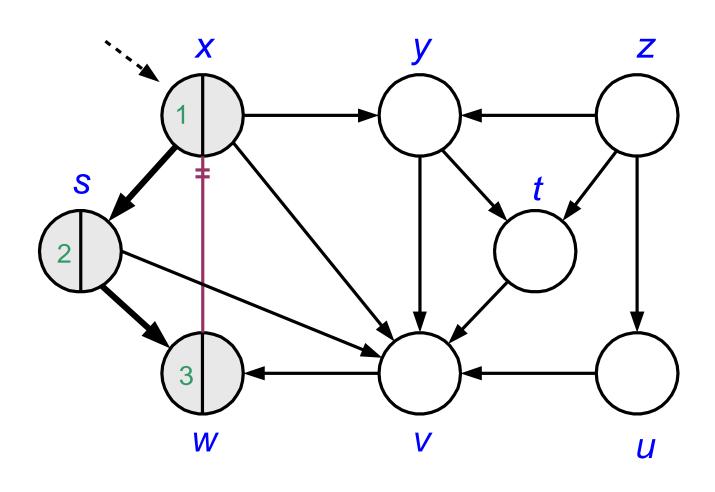
- Input graph G = (V,E) is represented using adjacency list.
- Each vertex is colored **white**, **gray or black**.
  - (1) **Vertex with White Color**: undiscovered vertex.
  - (2) **Vertex with Gray Color**: Consider to be discovered.
  - (3) **Vertex with Black Color:** Adjacency-list of Vertex has been discovered.
- Additional Data Structures
  - (1) color[u]: Store the color of each vertex  $u \in V$
  - (2)  $\Pi[u]$  : Store the predecessor of u.
- Besides creating depth first forest DFS also timestamps each vertex. Each vertex goes through two timestamps:
  - (3) d[u] : Store when u is first discovered and grayed (discovery time).
  - (4) f[u] : Store when the search finishes examining u's adjacency-list and blackens u (finishing time).
    - $\Box$  f[u] > d[u]

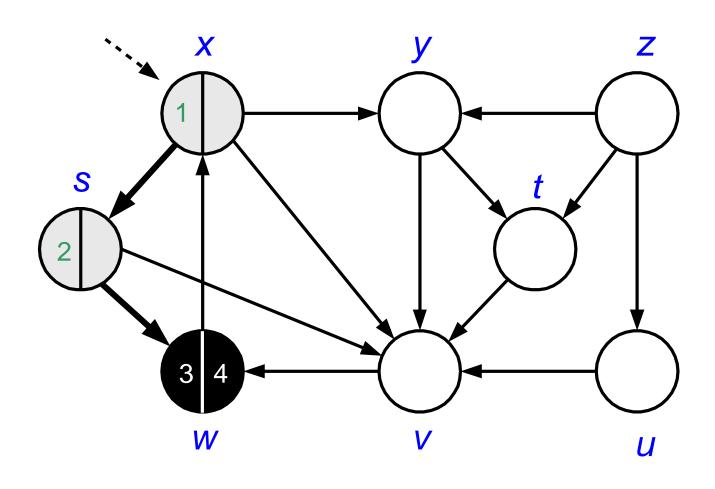


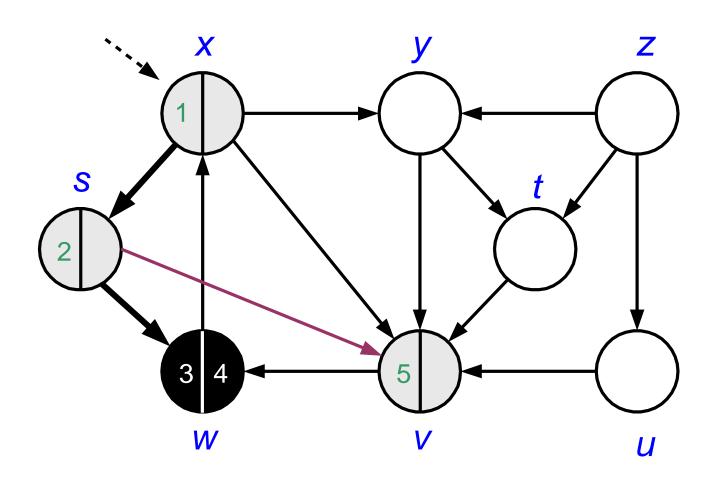


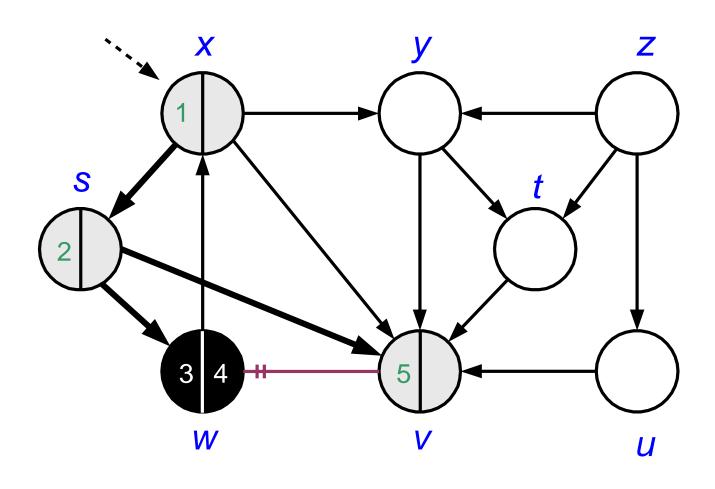


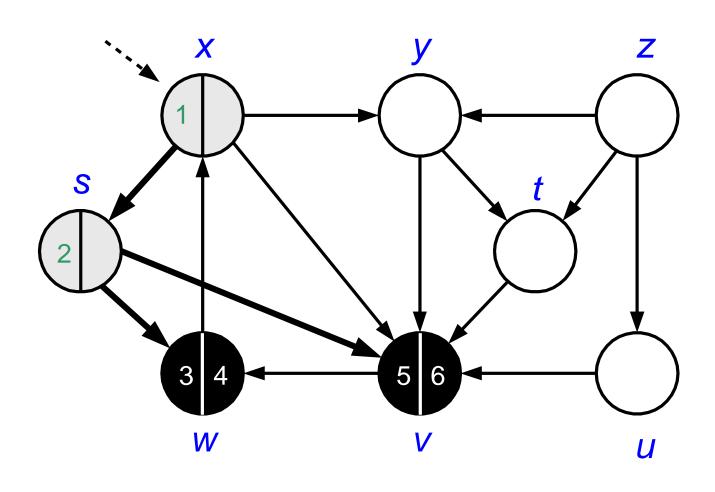


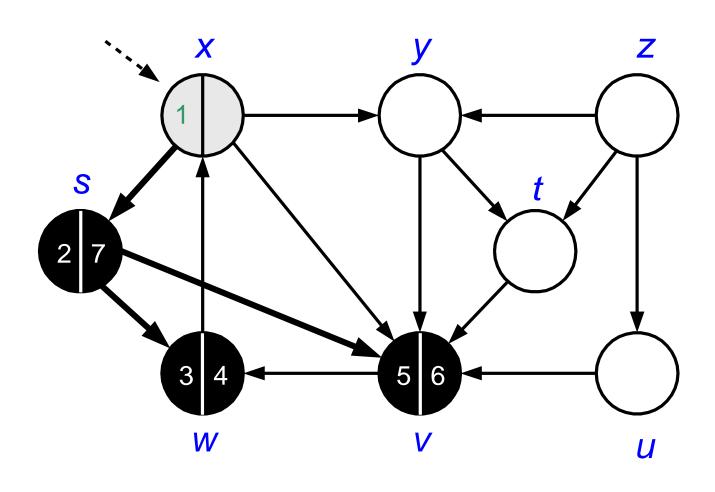


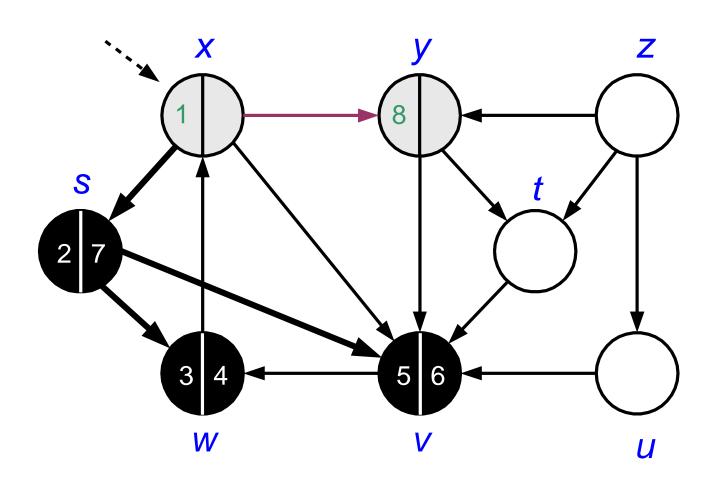


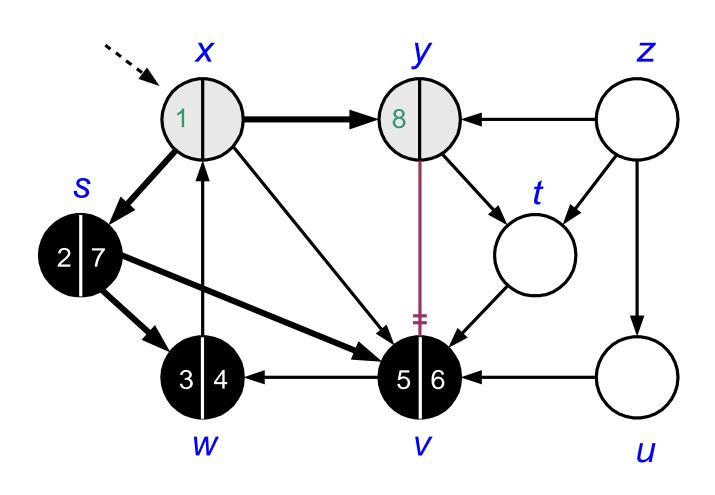


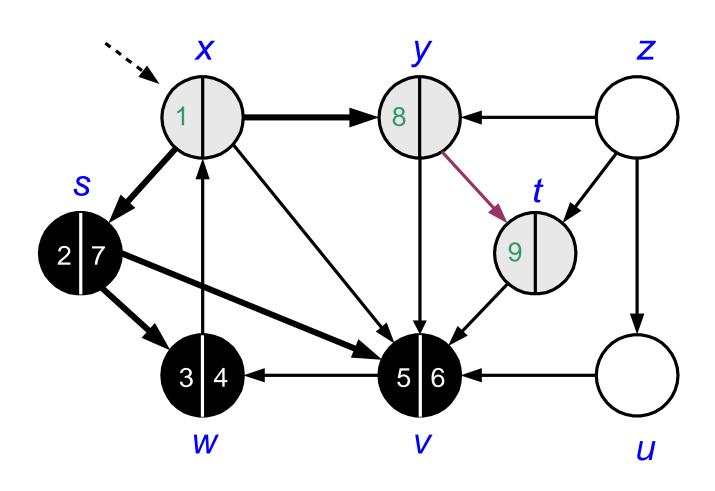


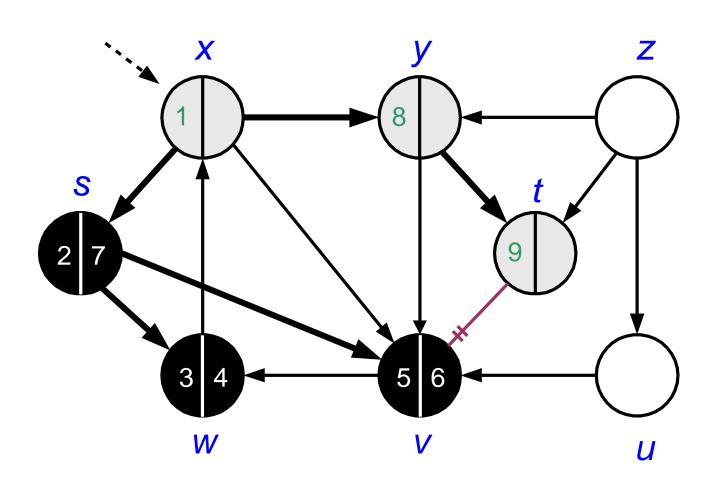


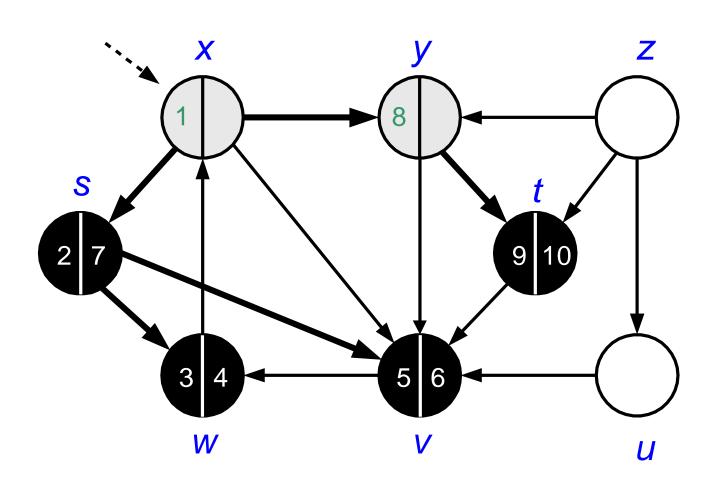


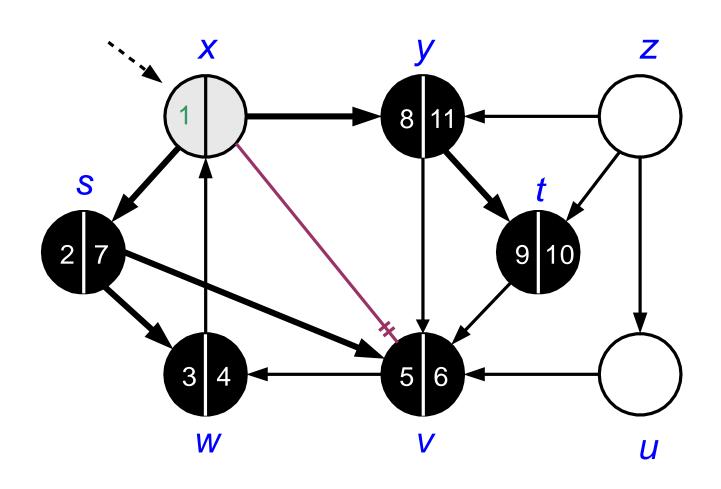


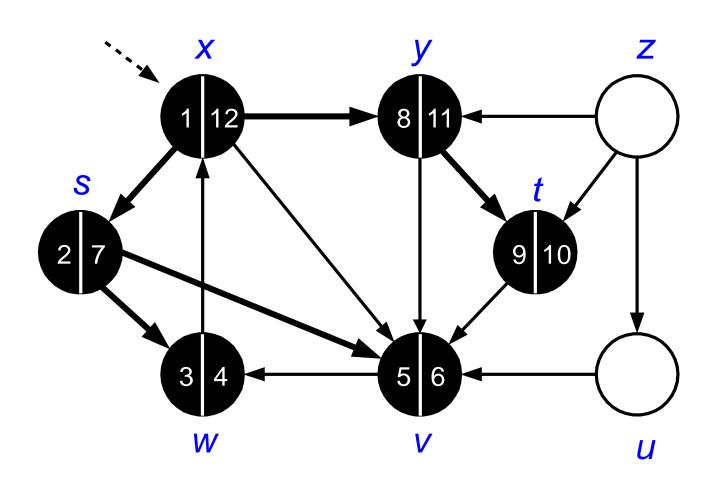


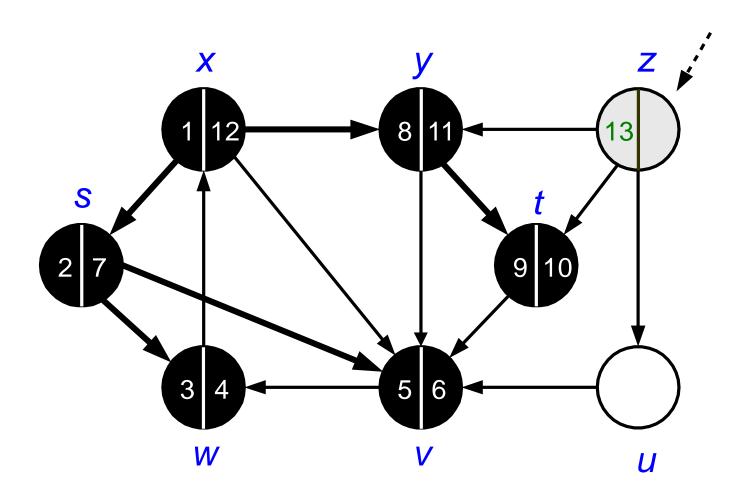


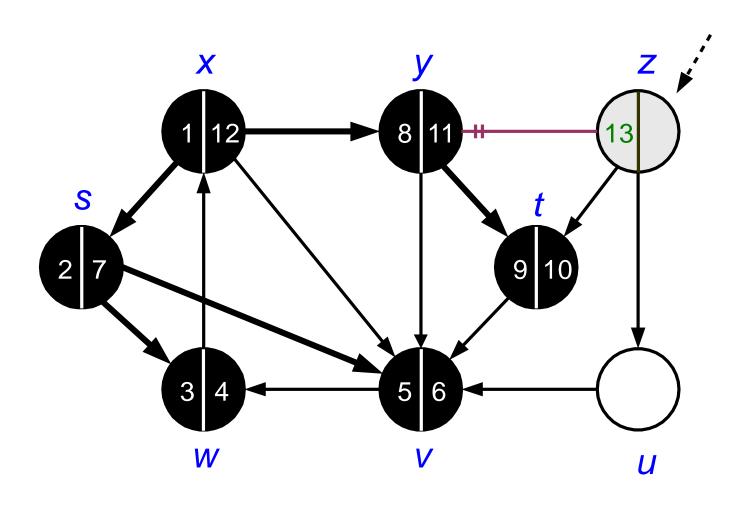


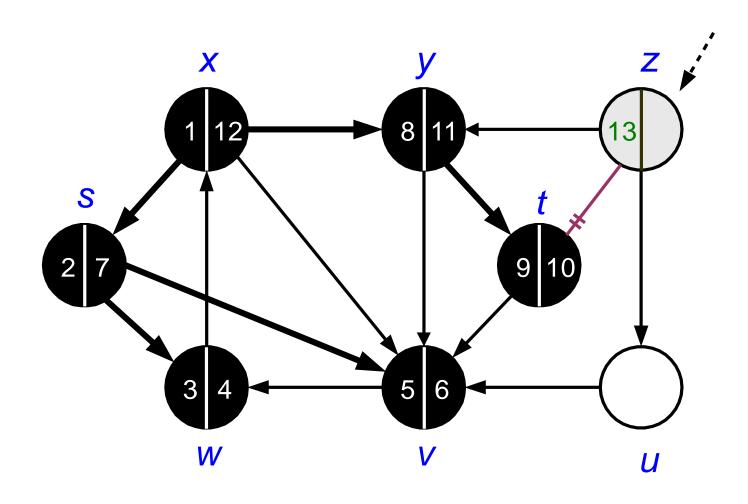


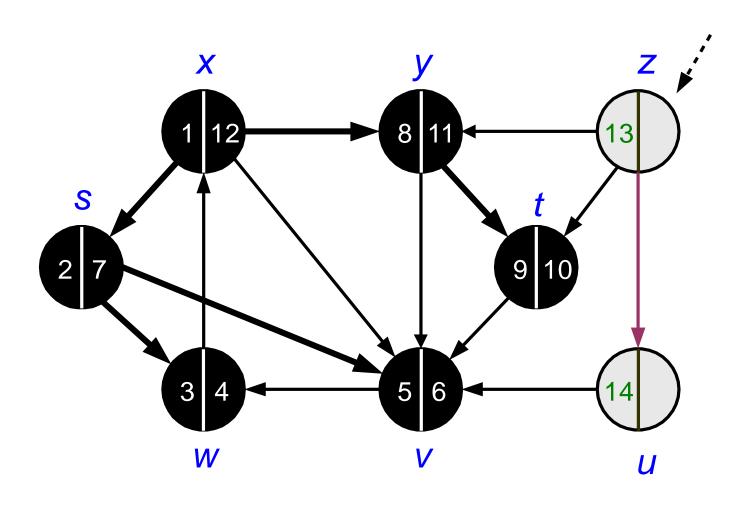


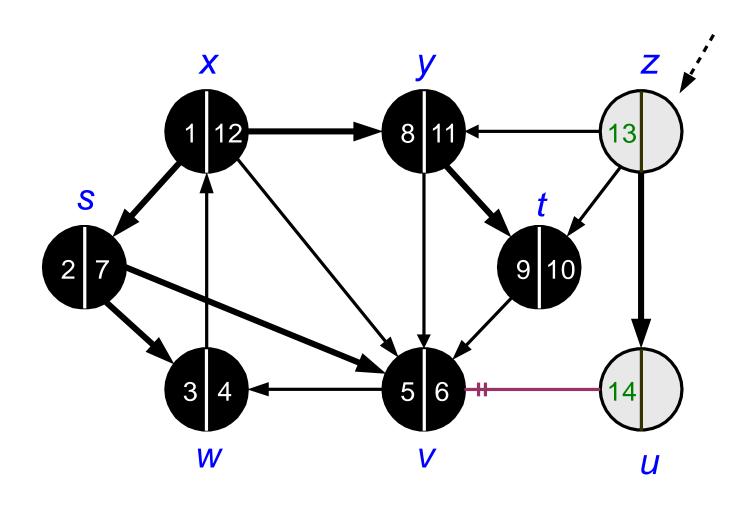


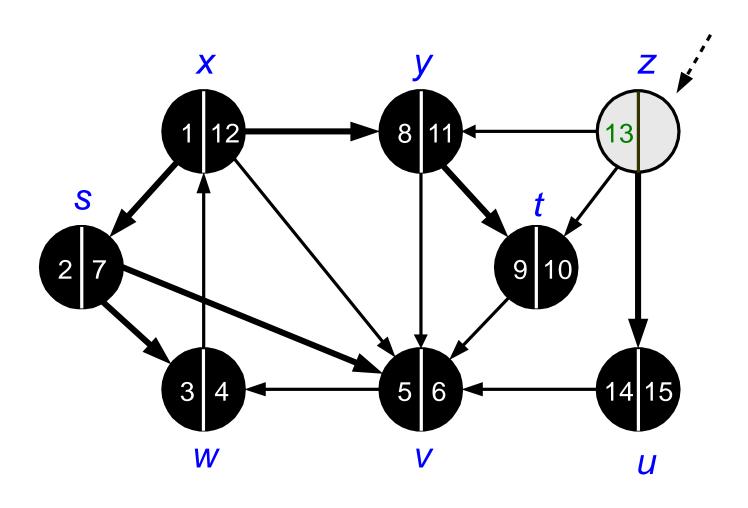


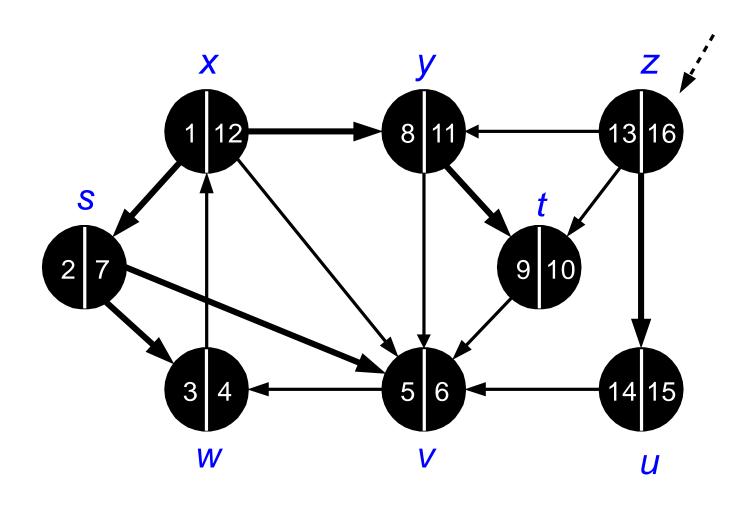


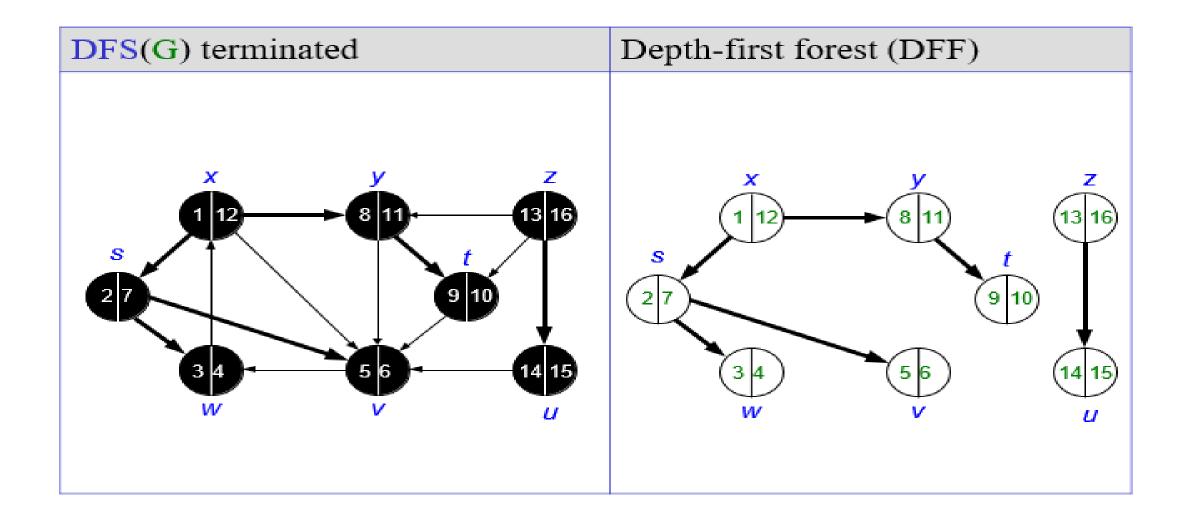












#### DFS: Algorithm

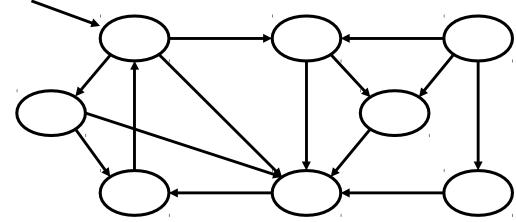
- $\Box$  DFS(G)
- 1. for each vertex u € V[G]
- 2. color[u]=white
- 3.  $\pi[u]=NIL$
- 4. time = 0
- 5. for each vertex u € V[G]
- 6. if (color[u] = = white)
- 7. DFS-VISIT(G,u)

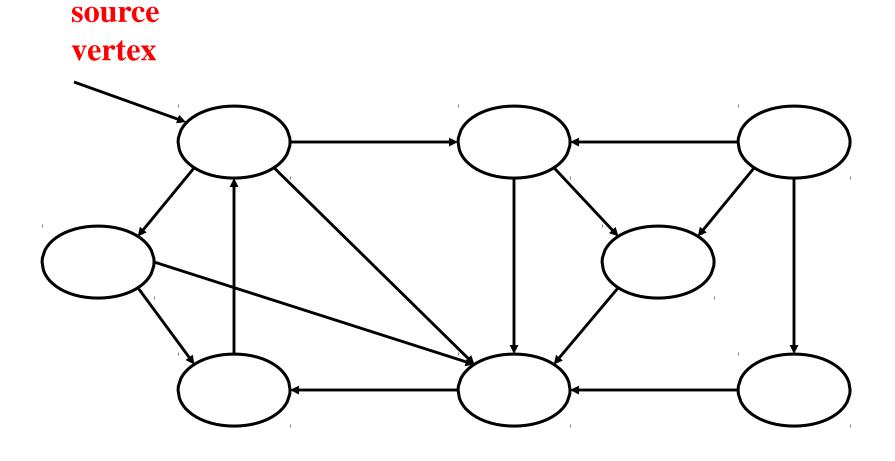
#### DFS:Algorithm (Cont.)

#### ☐ DFS-VISIT(u)

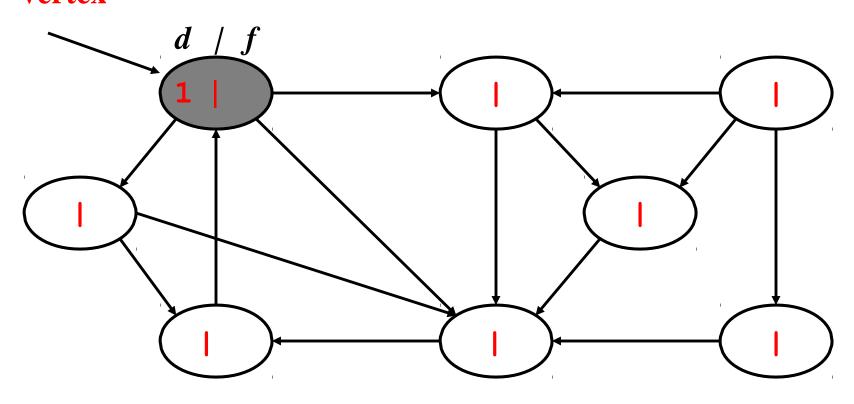
- 1. time = time + 1
- 2. d[u] = time
- 3. color[u] = gray
- 4. for each v € Adj(u) in **G** do
- 5. if (color[v] == white)
- 6.  $\pi [v] = u;$
- 7. DFS-VISIT(G,v);
- 8.  $\operatorname{color}[u] = \operatorname{black}$
- 9. time = time +1;
- 10. f[u]=time;

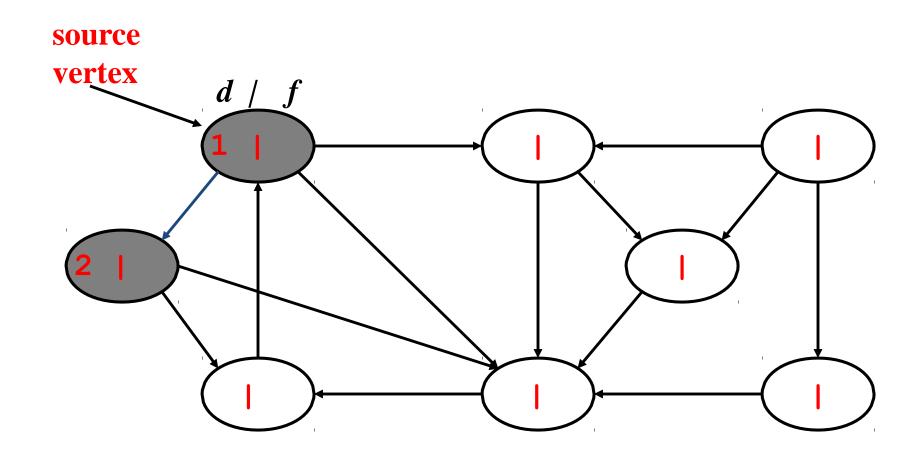


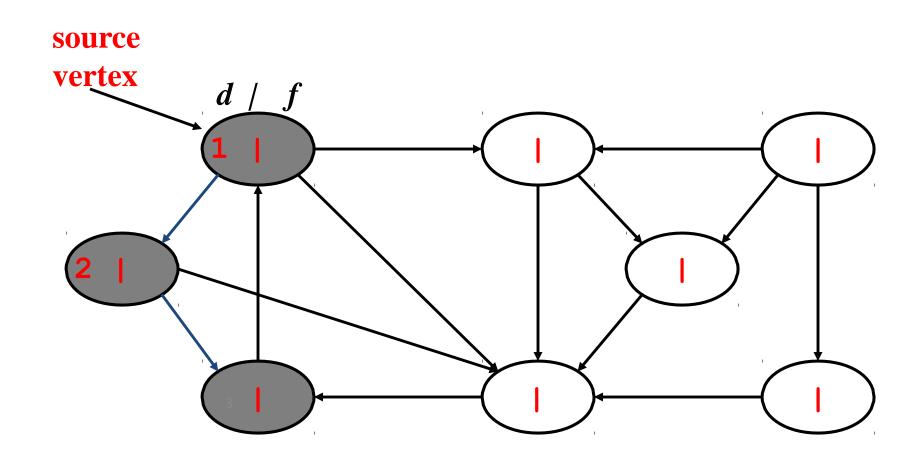


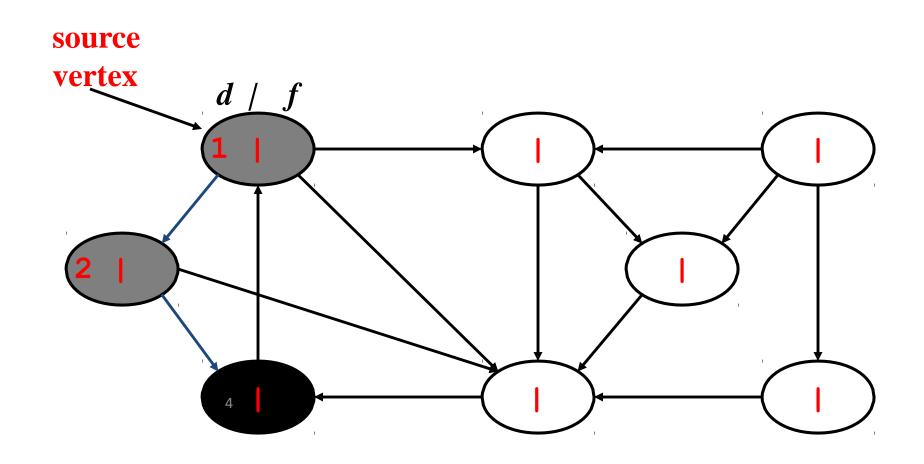


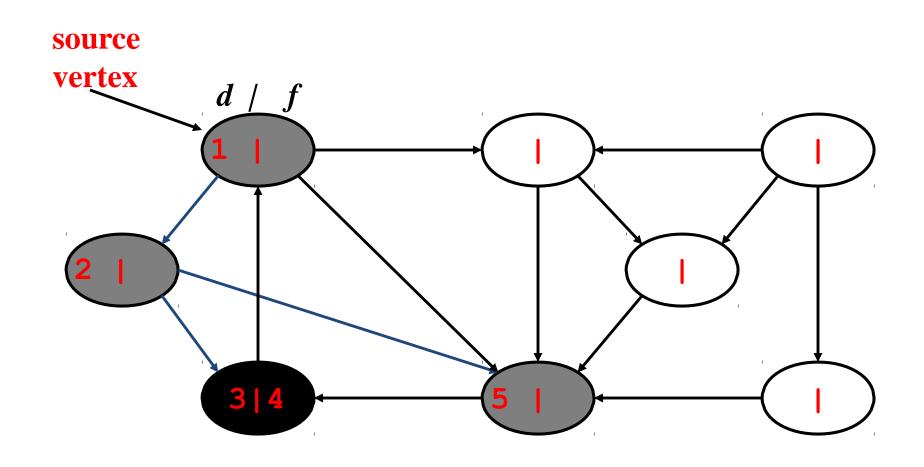
#### source vertex

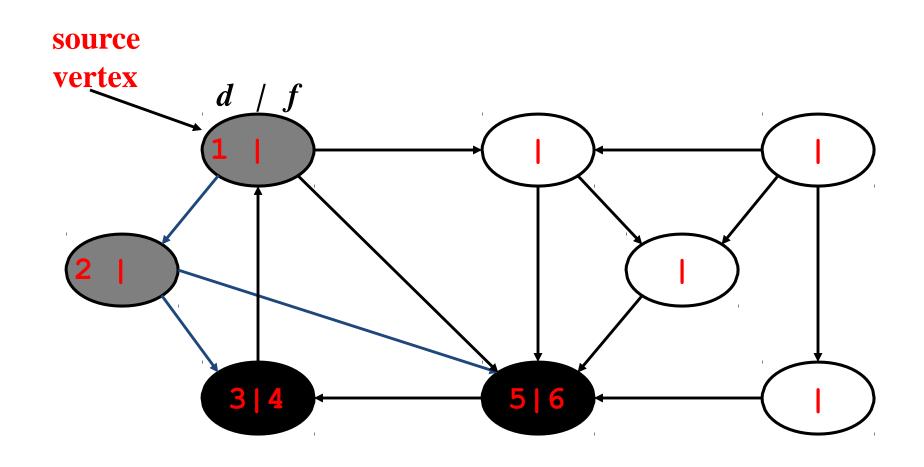


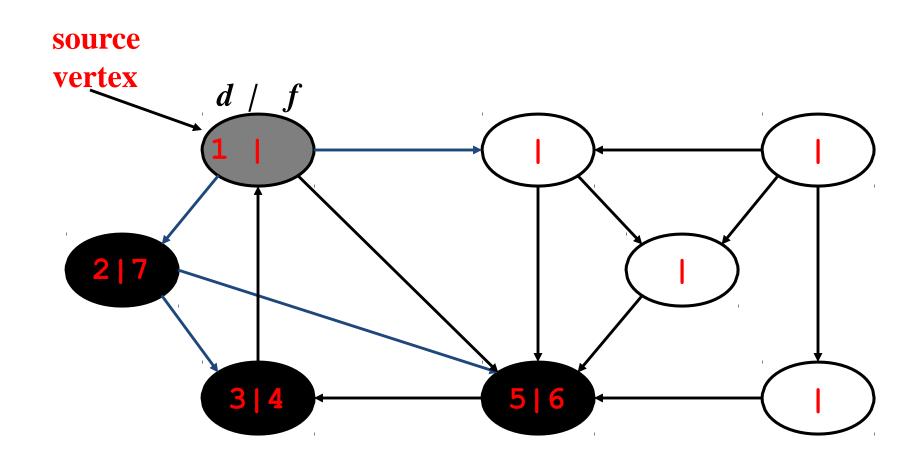


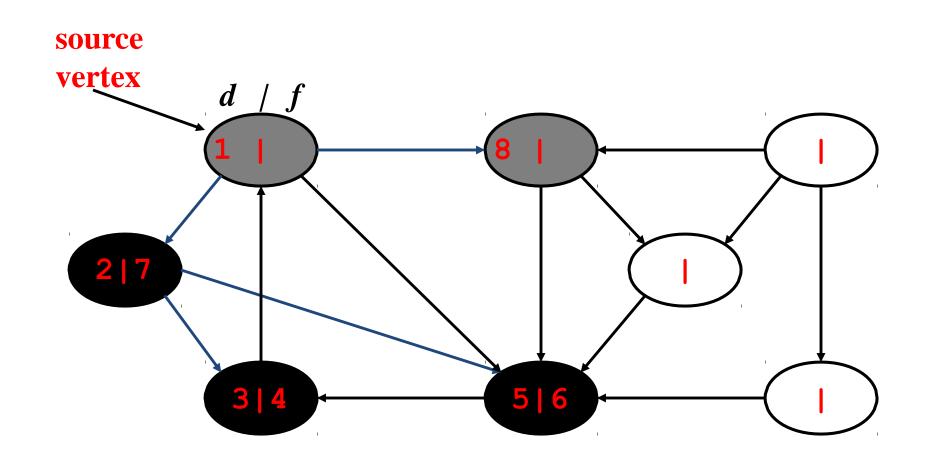


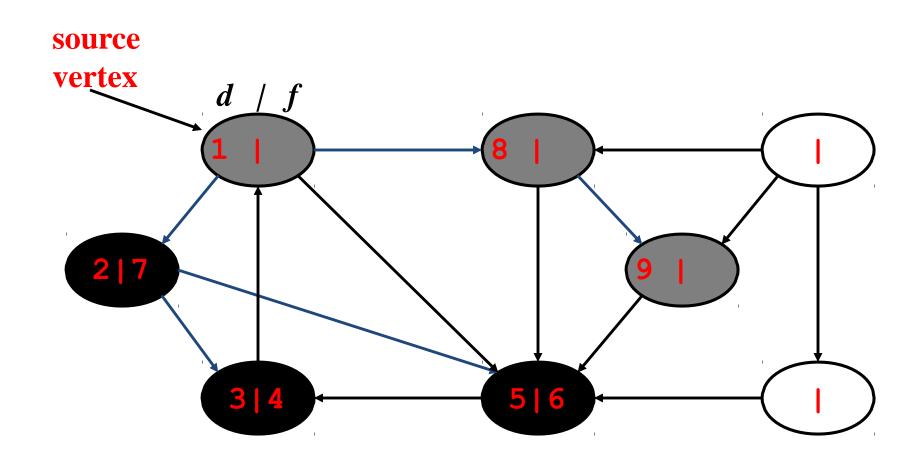


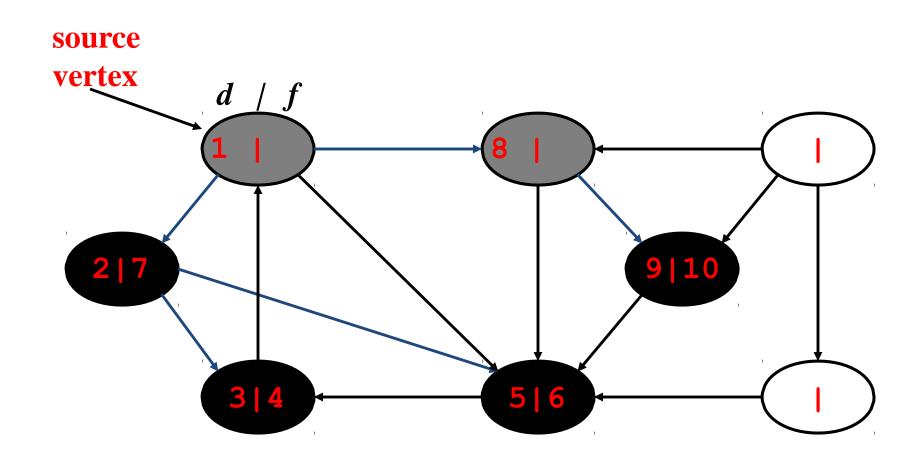


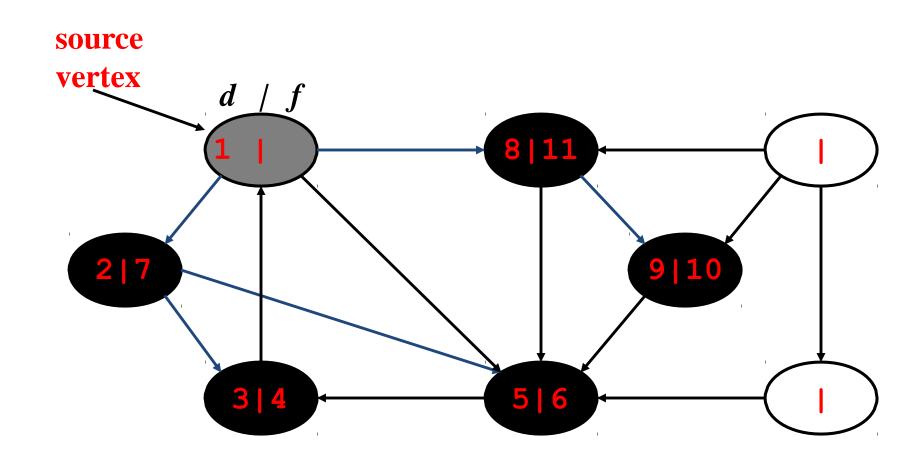


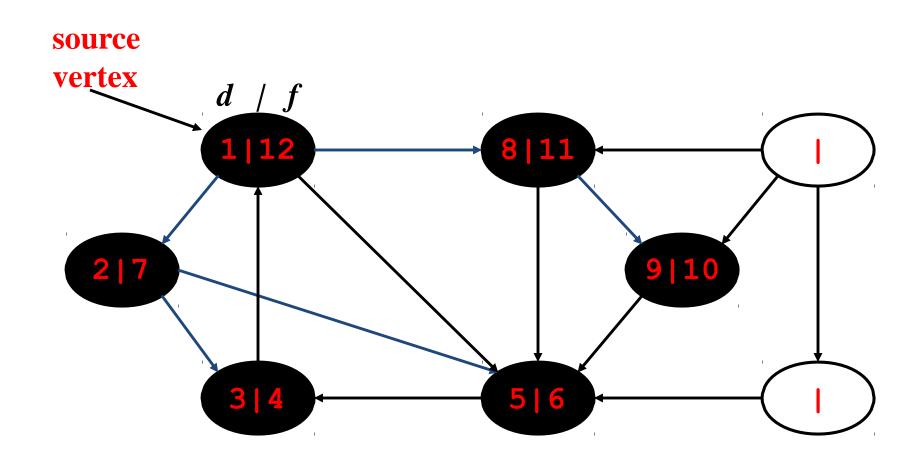


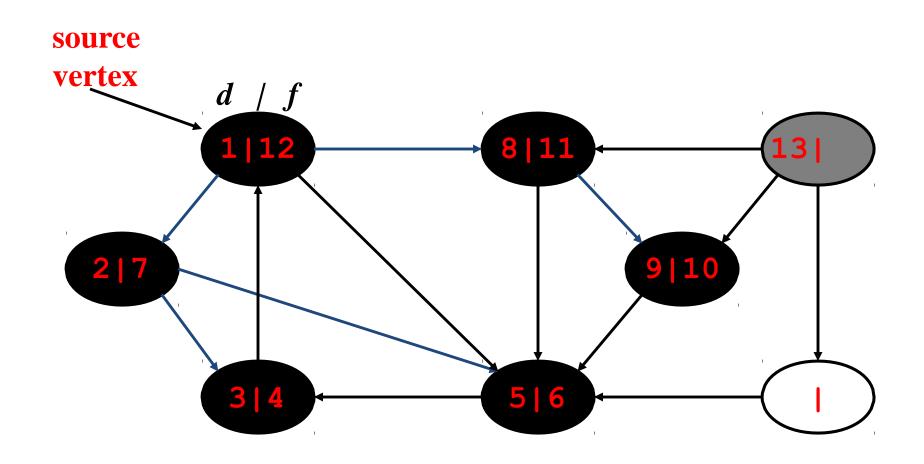


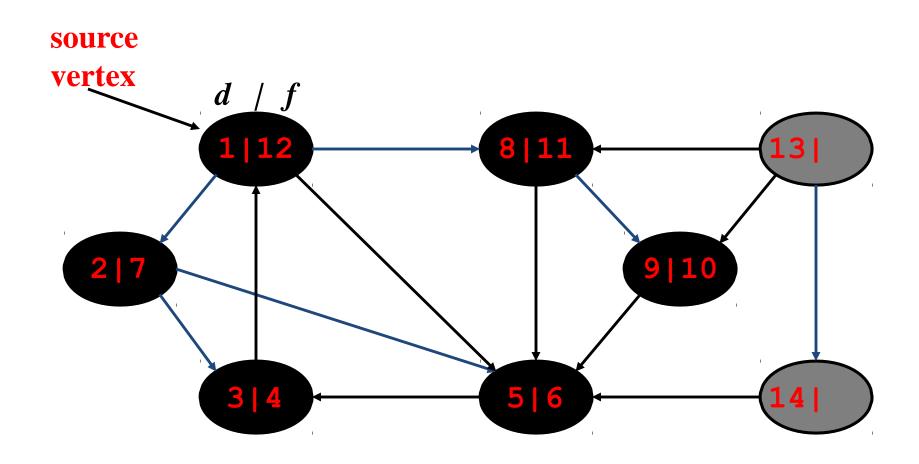


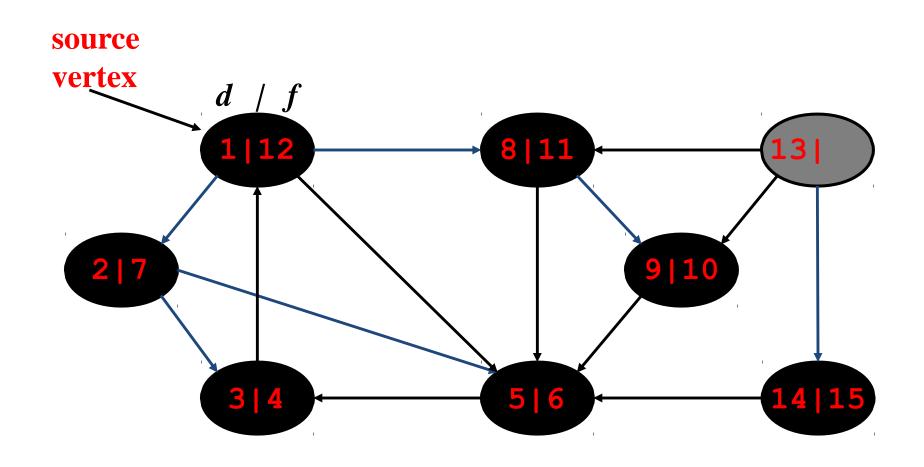


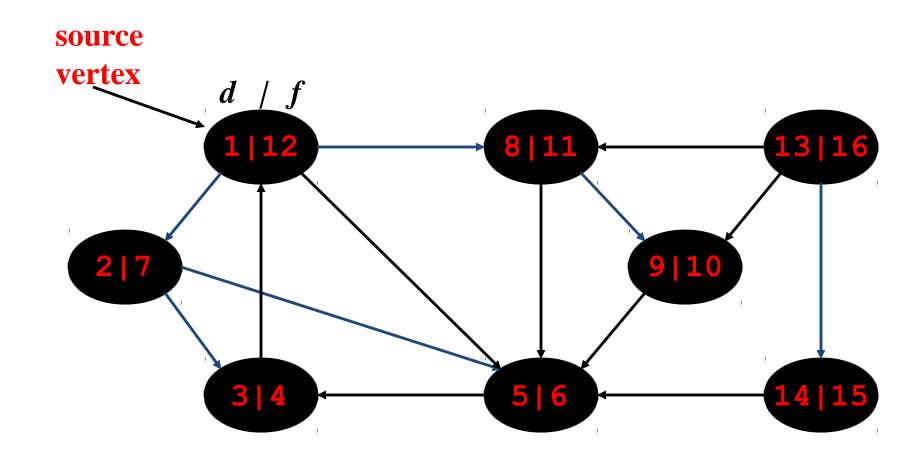




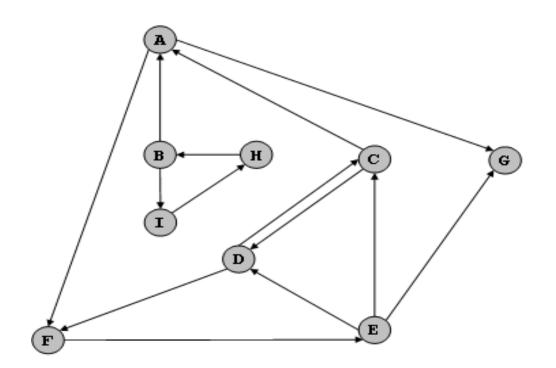








#### Directed Depth First Search



#### Adjacency Lists

A: FG

B: A H

C: AD

D: CF

E: CDG

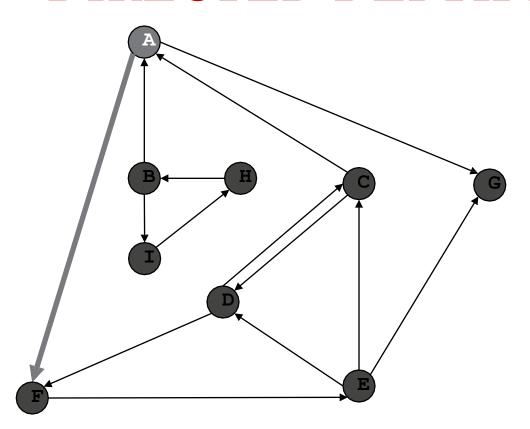
F: E

G:

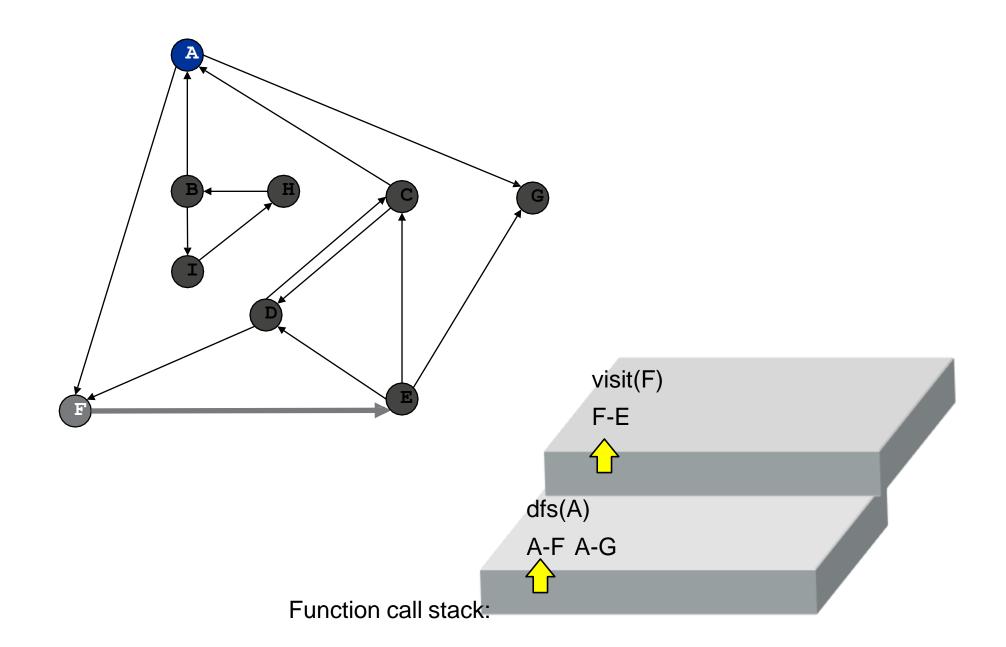
H: B

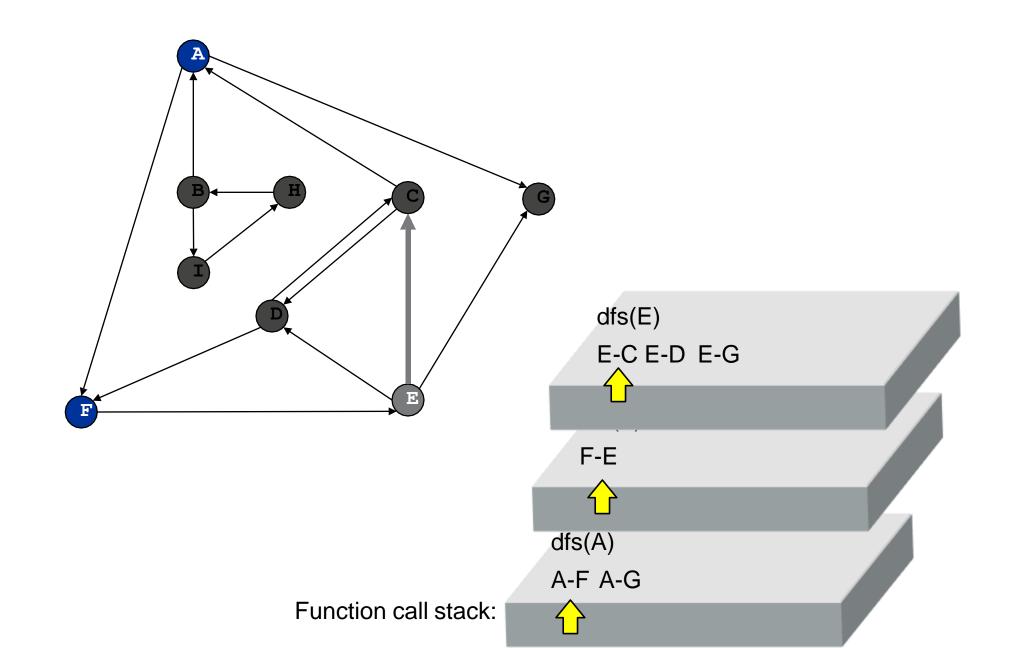
I: H

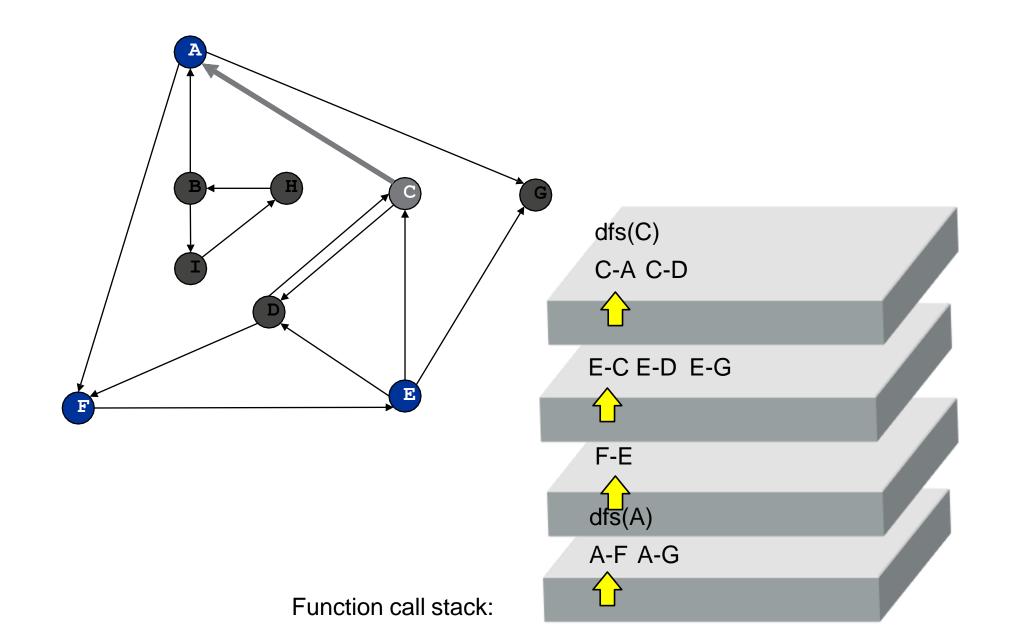
#### DIRECTED DEPTH FIRST SEARCH

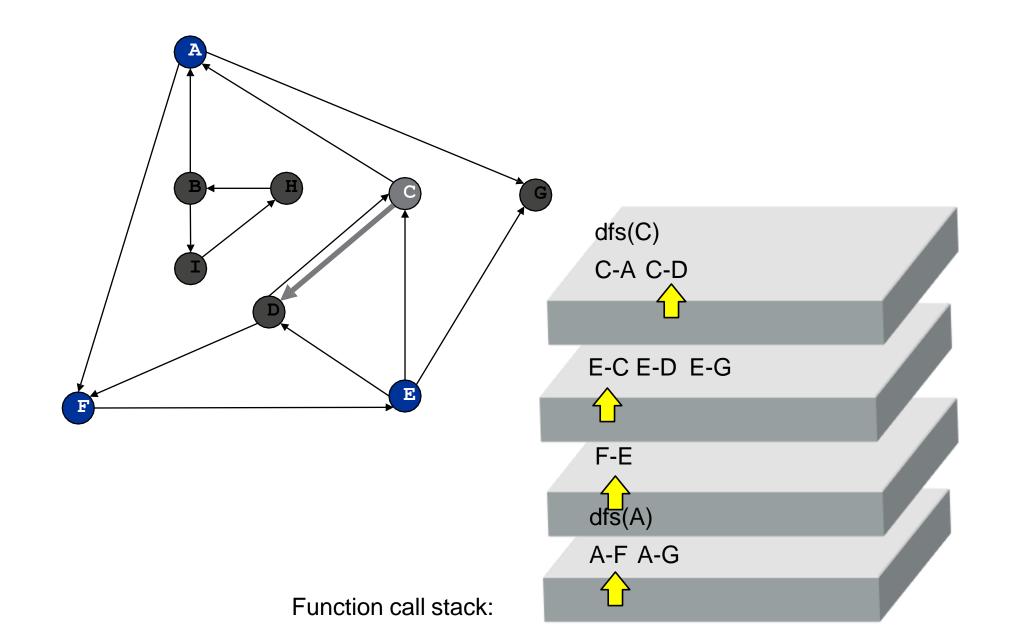


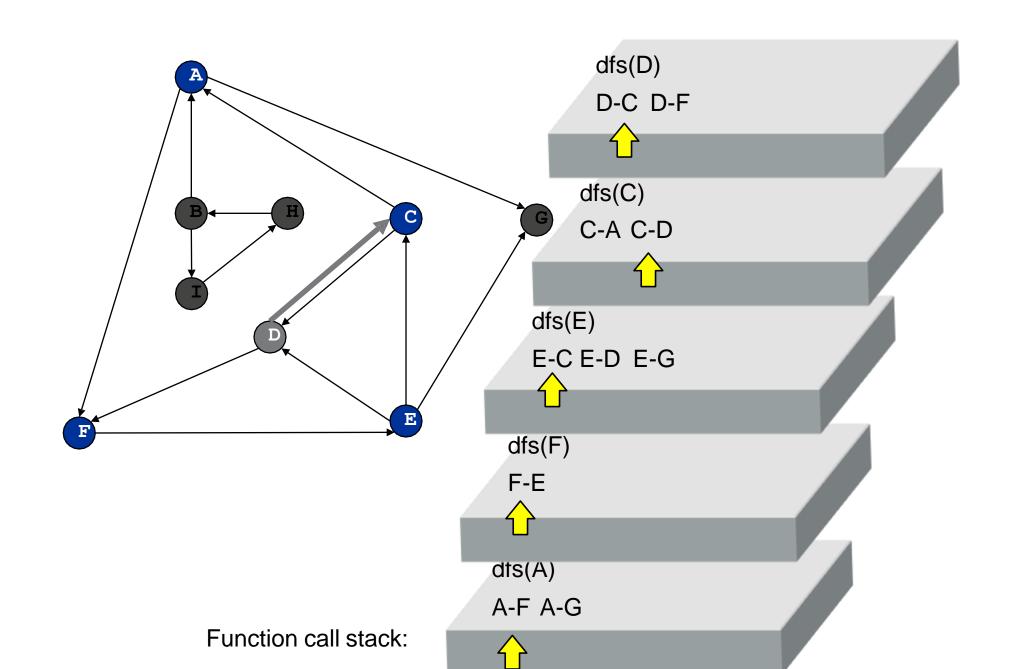
Function call stack:

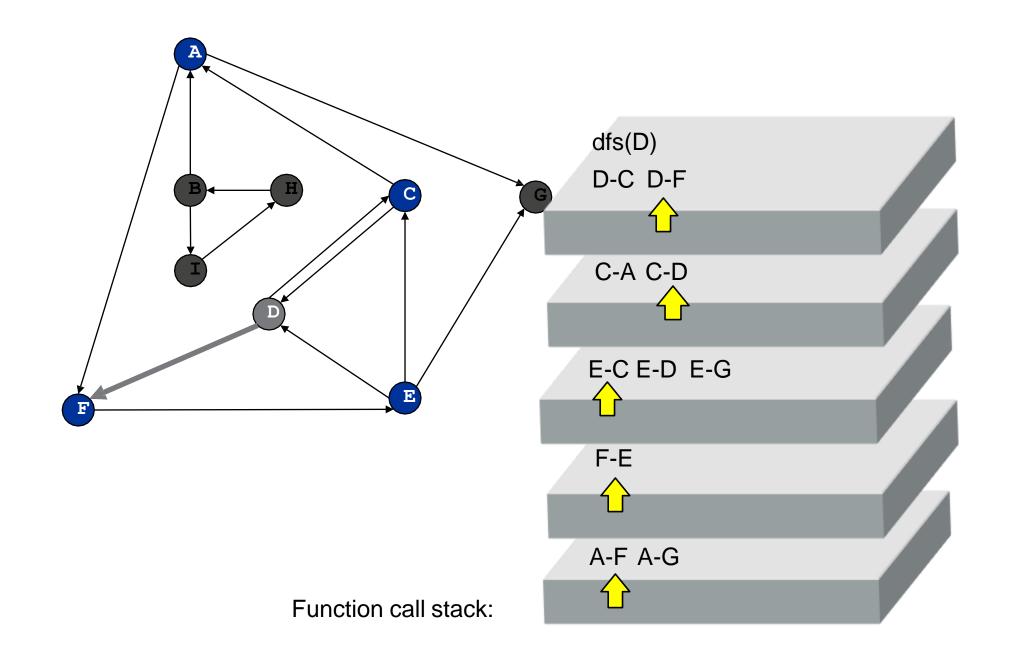


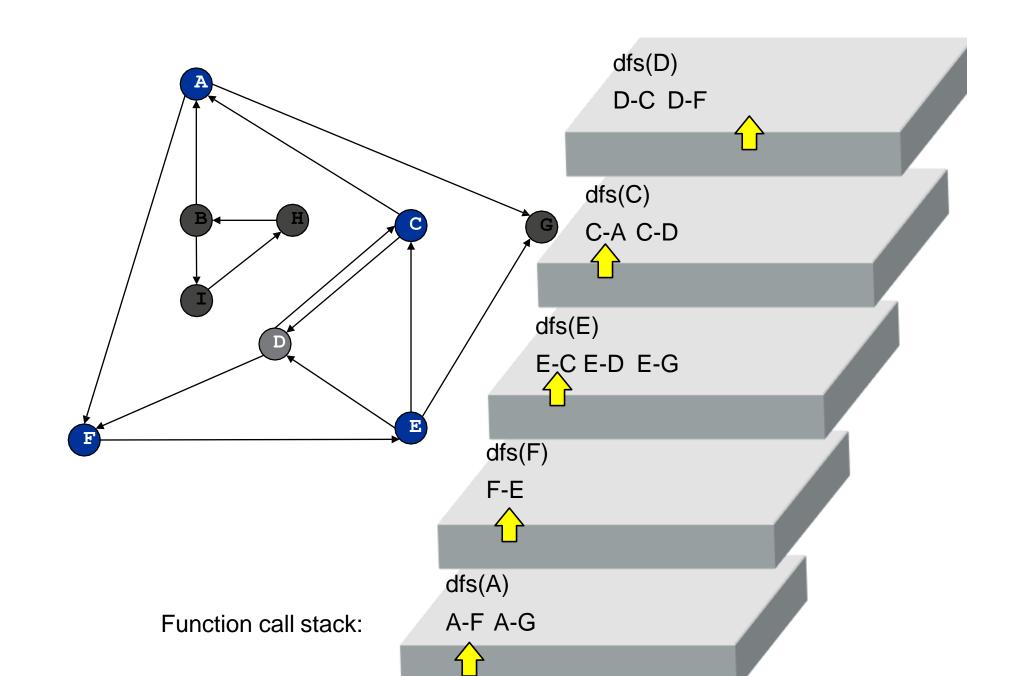


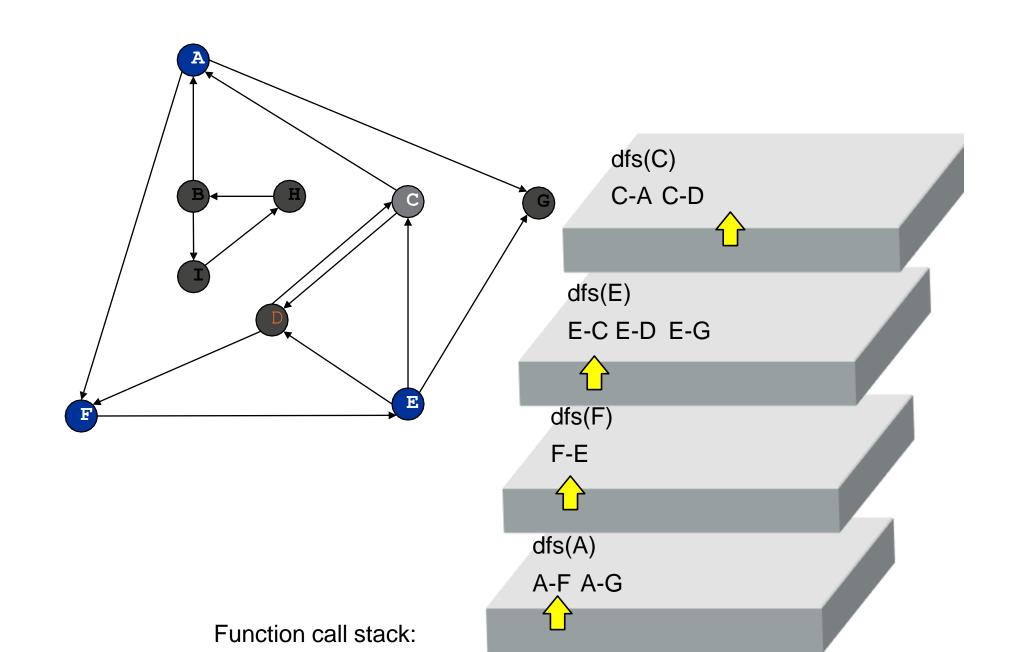


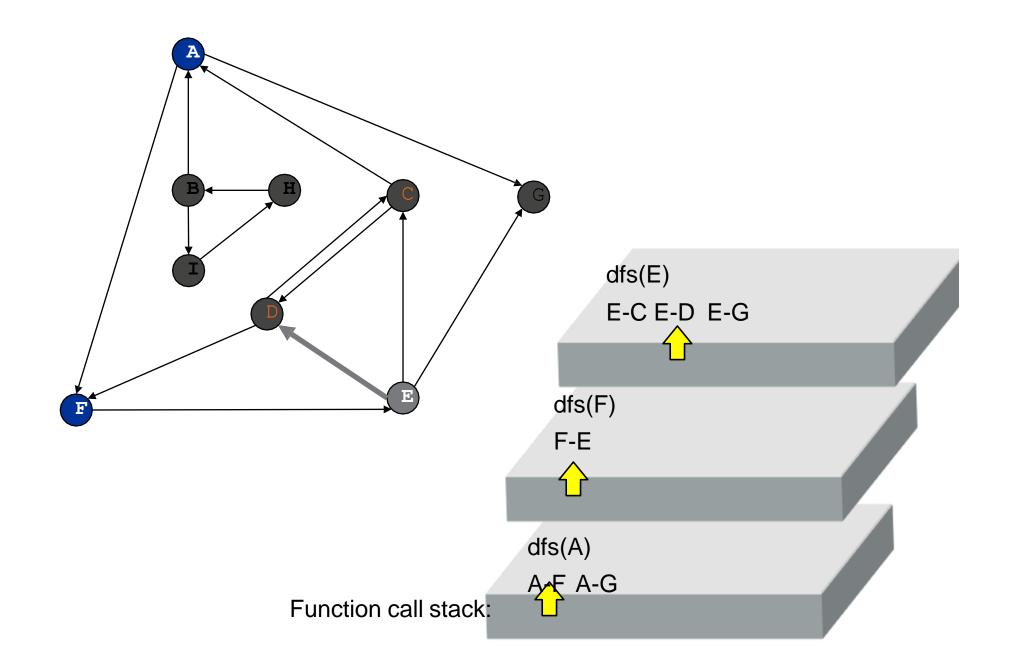


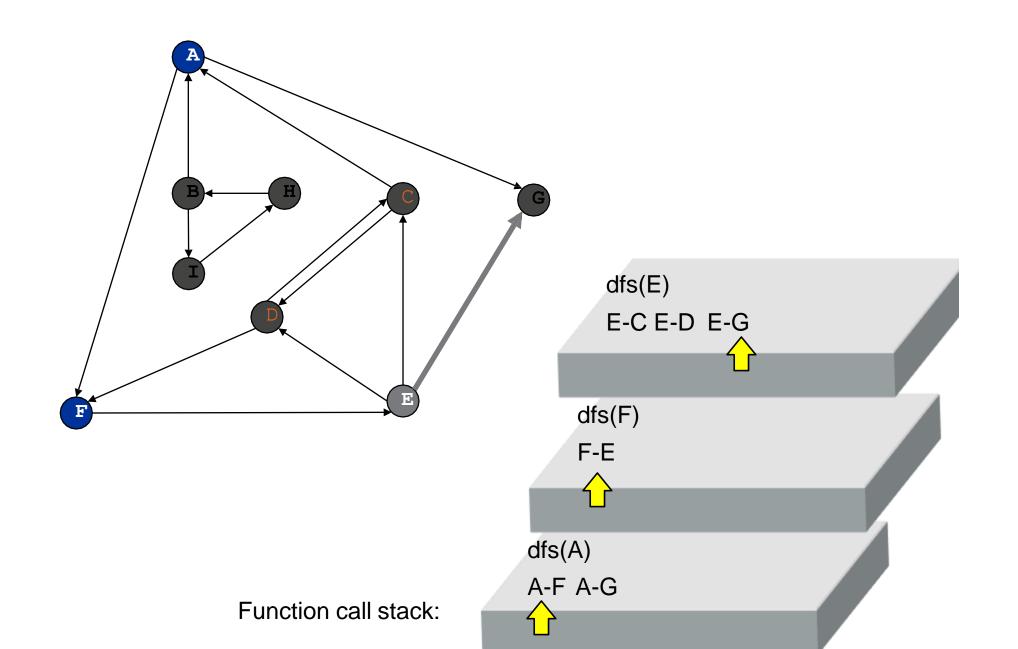


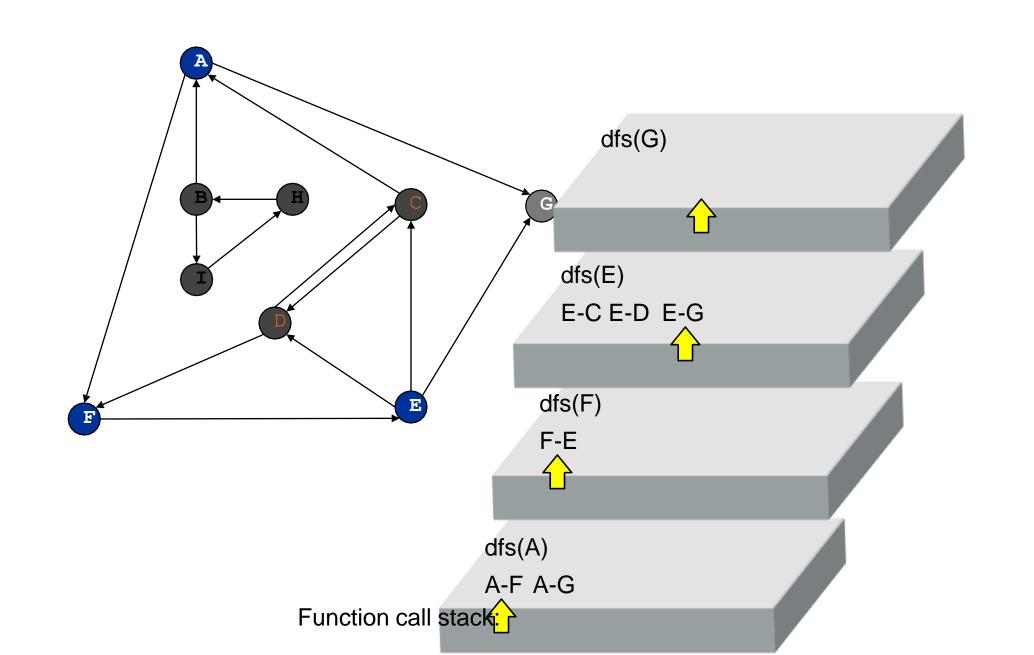


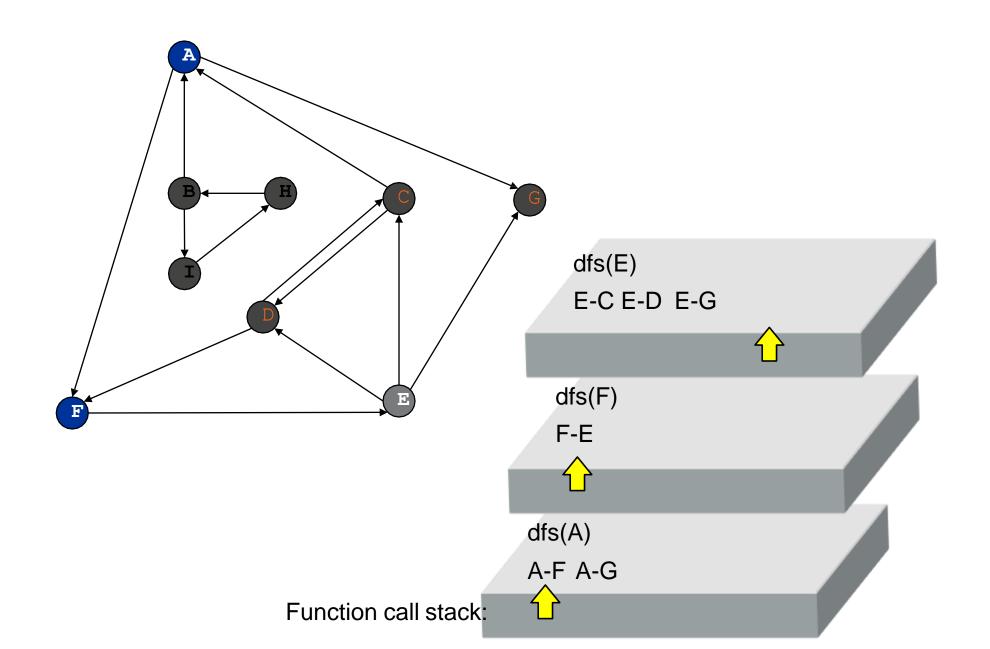


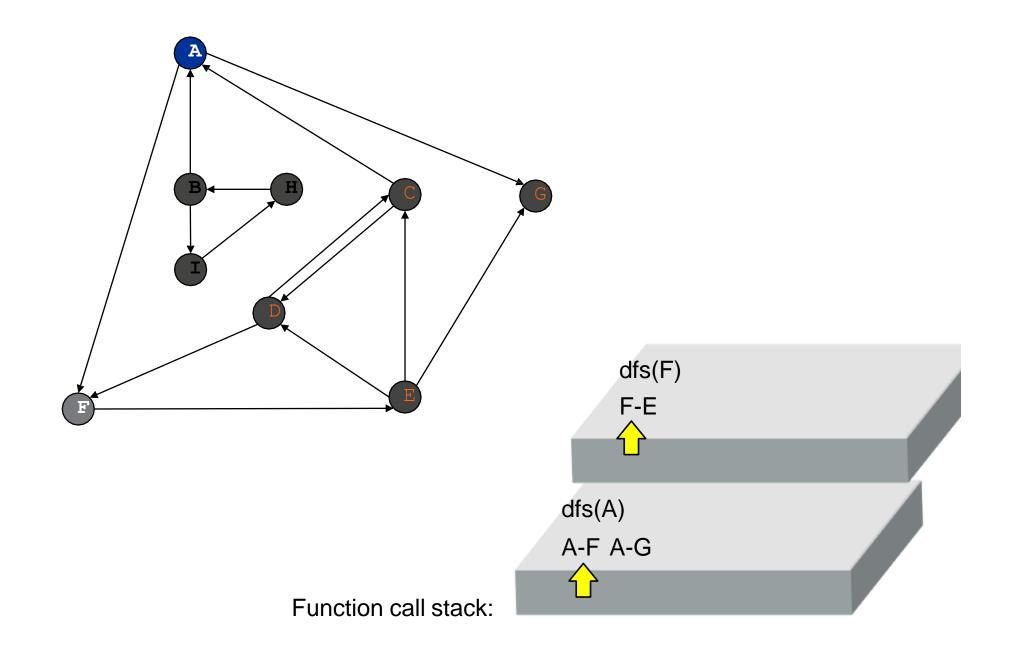


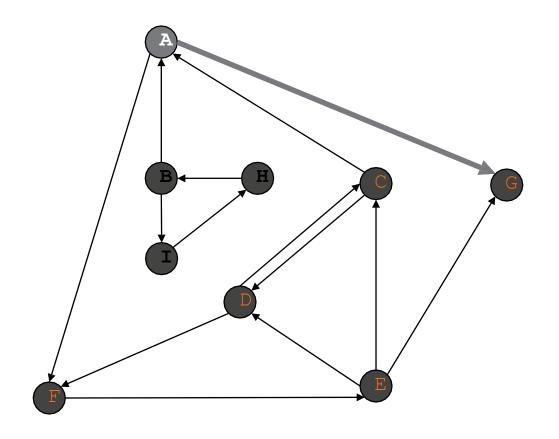






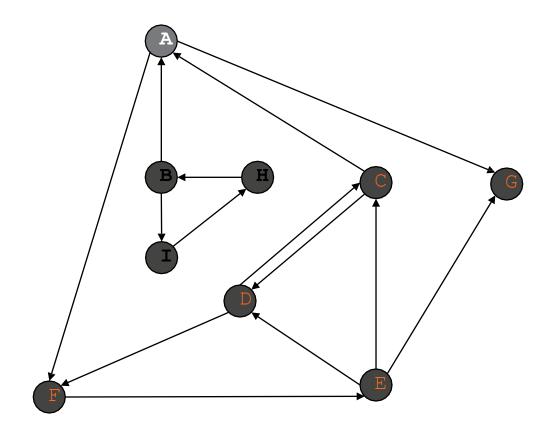






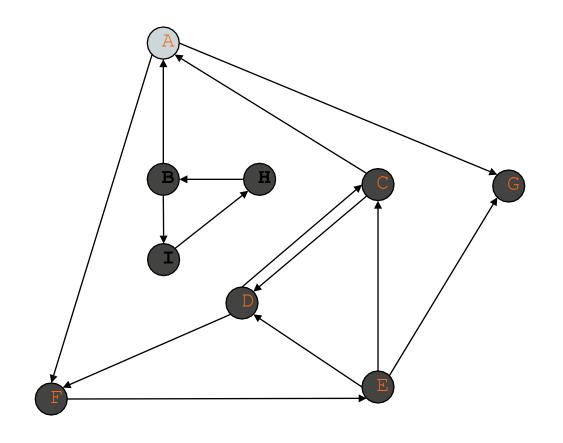
dfs(A) A-F A-G

Function call stack:

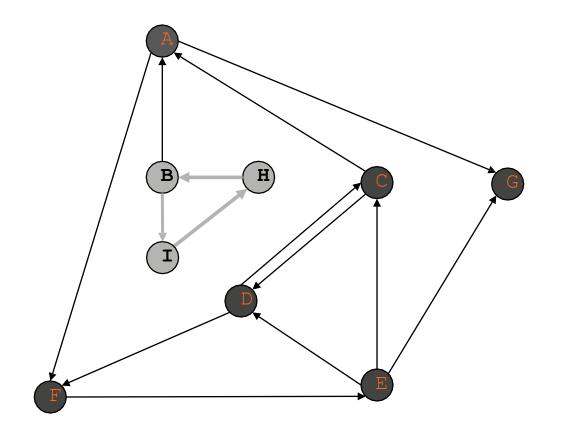


dfs(A)
A-F A-G

Function call stack:



Nodes reachable from A: A, C, D, E, F, G



*Nodes reachable from A:* A, C, D, E, F, G

#### Running Time of DFS Algorithm

Lines 1-2 and lines 4-5 of DFS take time  $\Theta(V)$ , exclusive of the time to execute the calls to DFS-VISIT.

The procedure DFS\_VISIT is called exactly once for each vertex  $v \in V$ . During an execution DFS\_VISIT(v), the loop on lines 4-7 is executed |Adj[v]| times. Since  $\Sigma |Adj[v]| = \Theta(E)$ , the total cost of executing lines 2-6 of DFS-VISIT is  $\Theta(E)$ .

So the total running time of DFS is  $\Theta(V+E)$ .

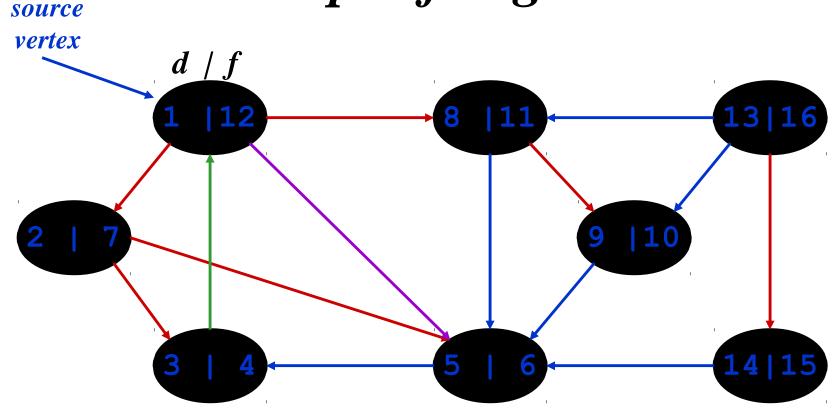
#### DFS: Application

- Topological Sort
- Strongly Connected Component

#### Classification of Edges

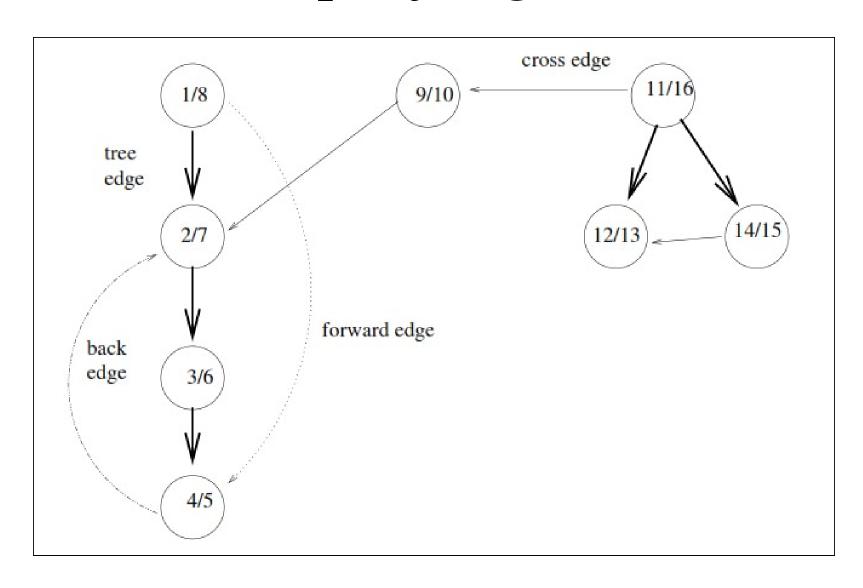
- ☐ Another interesting property of DFS is that the DFS can be used to classify the edges of the input graph G=(V,E). The DFS creates depth first forest which can have four types of edges:
- ☐ **Tree edge:** Edge (u,v) is a tree edge if **v** was first discovered by exploring edge (u,v). **White** color indicates tree edge.
- **Back edge:** Edge (u,v) is a back edge if it connects a vertex **u** to a ancestor **v** in a depth first tree. **Gray** color indicates back edge.
- □ **Forward edge:** Edge (u,v) is a forward edge if it is non-tree edge and it connects a vertex **u** to a descendant **v** in a depth first tree. **Black** color indicates forward edge.
- ☐ Cross edge: rest all other edges are called the cross edge. Black color indicates forward edge.

#### Example of Edges



Tree edges Back edges Forward edges Cross edges

# Example of Edges



Thank