Bellman-Ford Algorithm

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History

- The algorithm was first proposed by Alfonso Shimbel in 1955
- Named after Richard Bellman and Lester Ford Jr who published it in 1956 and 1958
- Edward F. Moore also published the same algorithm in 1957, and for this reason it is also sometimes called the Bellman–Ford–Moore algorithm

Bellman-Ford

- > The Bellman-Ford algorithm operates on an input graph, G, with | V | vertices and | E | edges
- ➤ Bellman-Ford algorithm is a single-source shortest path algorithm
- > Bellman-Ford can also detect negative edge cycles which is a useful feature.
- Returns a Boolean value indicating whether or not there is a negative-weight cycle that is reachable from the source.
- If there is such cycle the algorithm indicates nor solution exists, If there is no such cycle, the algorithm produces the shortest paths and their weights.
- > The main contribution of this algorithm is in its ordering of relaxations.

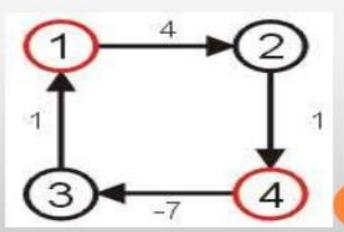
Bellman-Ford

Input: Graph and a source vertex src

Output: Shortest distance to all vertices from src. If there is a negative weight cycle, then shortest distances are not calculated, negative weight cycle is reported.

NEGATIVE CYCLES:

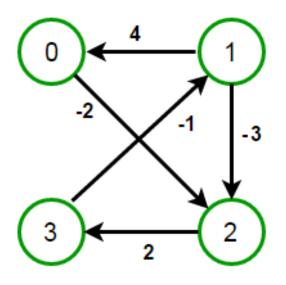
- Clearly, if we have negative vertices, it may be possible to end up in a cycle whereby each pass through the cycle decreases the total *length*
- Thus, a shortest length would be undefined for such a graph
- Consider the shortest path from vertex 1 to 4...
- We will only consider nonnegative weights.

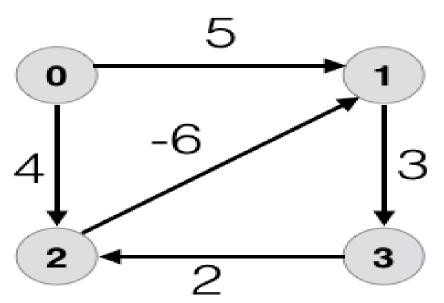


Negative Weight Cycle Example

Negative weight cycle: A **negative-weight cycle** is a **cycle** in a graph whose edges sum to a **negative** value.

For example, consider the following graph – It has one **negative-weight cycle**, 1-2-3-1 with sum -2 .

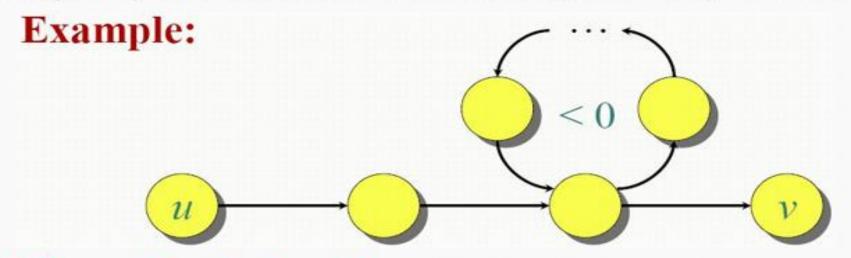




Soution = $-\infty$

Negative-weight cycles

Recall: If a graph G = (V, E) contains a negative-weight cycle, then some shortest paths may not exist.



Bellman-Ford algorithm: Finds all shortest-path lengths from a **source** $s \in V$ to all $v \in V$ or determines that a negative-weight cycle exists.

Bellman-Ford Steps

- 1) This step initializes distances from the source to all vertices as infinite and distance to the source itself as 0. Create an array dist[] of size |V| with all values as infinite except dist[src] where src is source vertex.
- 2) This step calculates shortest distances. Do following |V|-1 times where |V| is the number of vertices in given graph.
-a) Do following for each edge u-vIf dist[v] > dist[u] + weight of edge uv, then update dist[v]dist[v] = dist[u] + weight of edge uv
- 3) This step reports if there is a negative weight cycle in graph. Do following for each edge u-vIf dist[v] > dist[u] + weight of edge uv, then "Graph contains negative weight cycle"

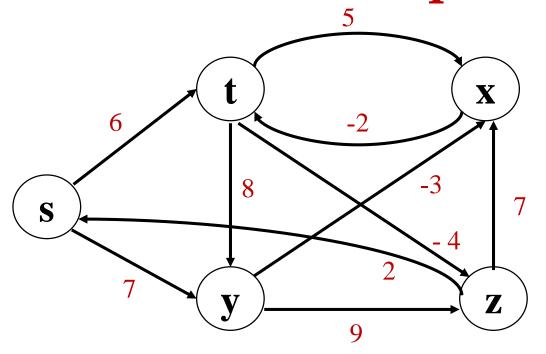
 The idea of step 3 is, step 2 guarantees the shortest distances if the graph doesn't contain a negative weight cycle. If we iterate through all edges one more time and get a shorter path for any vertex, then there is a negative weight cycle

Bellman-Ford Algorithm - Pseudocode

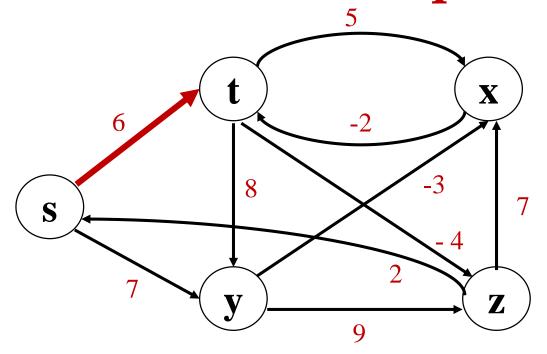
```
INITIALIZE-SINGLE-SOURCE (G, s)
BELLMAN-FORD(G, w, s)
                                                    for each vertex v \in G.V
   INITIALIZE-SINGLE-SOURCE (G, s)
                                                   v.d = \infty

\begin{array}{ccc}
3 & v.\pi \\
4 & s.d = 0
\end{array}

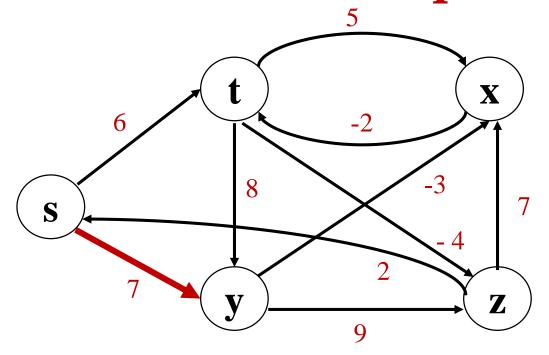
2 for i = 1 to |G.V| - 1
                                                    \nu.\pi = NIL
       for each edge (u, v) \in G.E
           RELAX(u, v, w)
                                                 Relax(u, v, w)
   for each edge (u, v) \in G.E
                                                    if v.d > u.d + w(u, v)
       if v.d > u.d + w(u, v)
                                                          v.d = u.d + w(u, v)
            return FALSE
                                                          v.\pi = u
   return TRUE
```



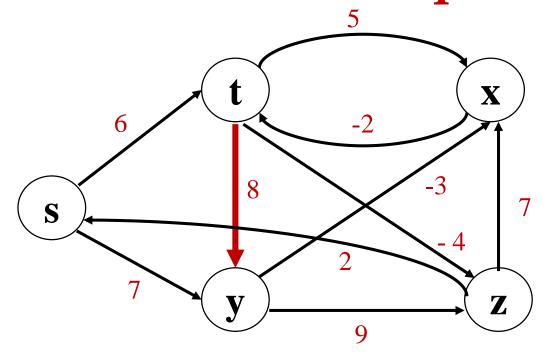
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL



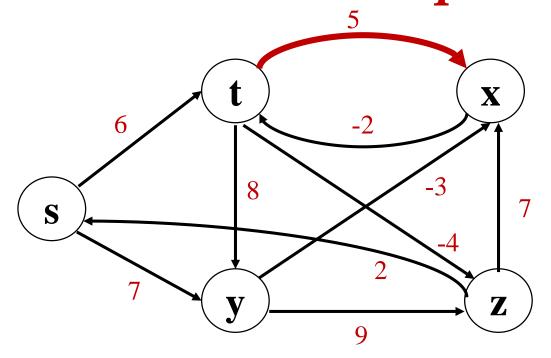
Vertex	S	t	x	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	6/s	∞/NIL	∞/NIL	∞/NIL



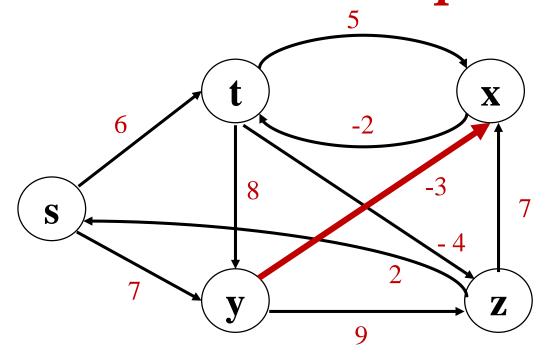
Vertex	S	t	x	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	6/s	∞/NIL	7/s	∞/NIL



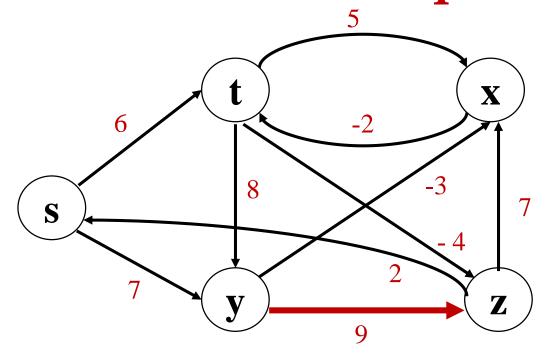
Vertex	S	t	x	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	6/s	∞/NIL	7/s	∞/NIL



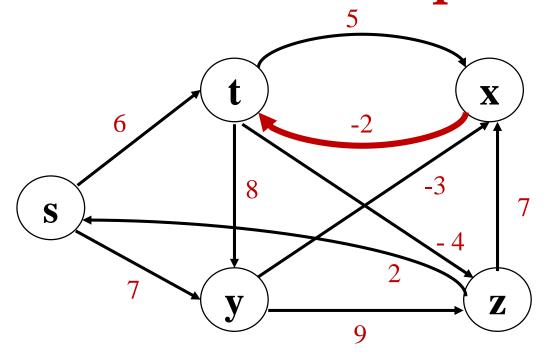
Vertex	S	t	x	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	6/s	11/t	7/s	∞/NIL



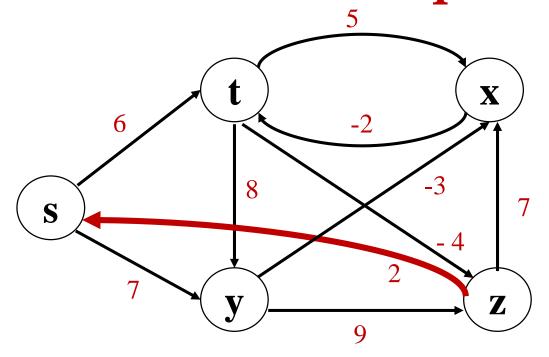
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	6/s	4/y	7/s	∞/NIL



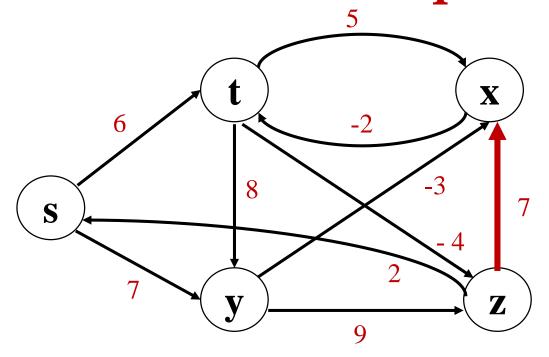
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	6/s	4/y	7/s	16/y



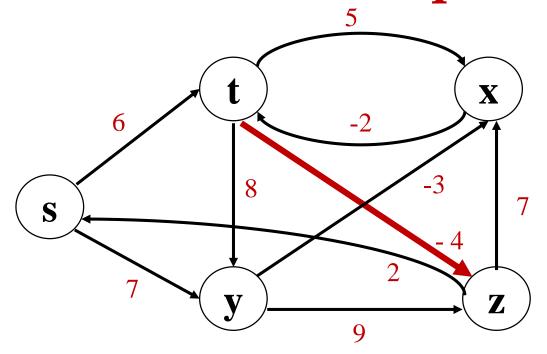
Vertex	S	t	x	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	16/y



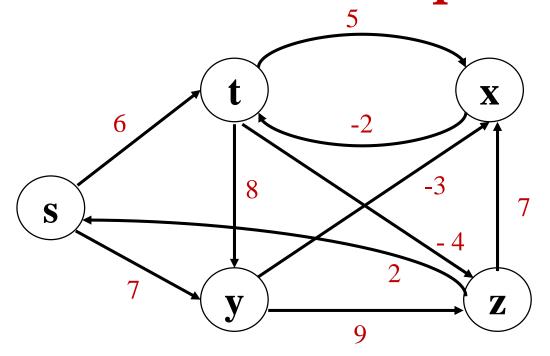
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	16/y



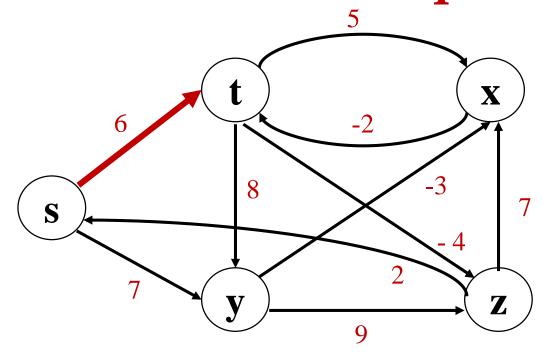
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	16/y



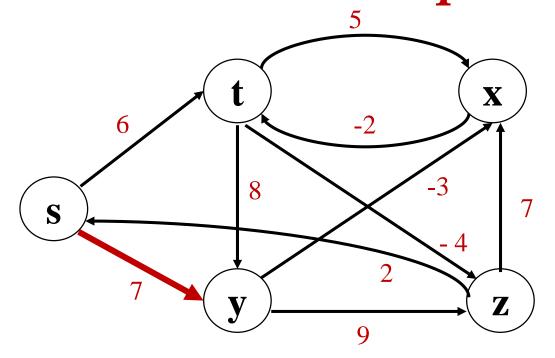
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t



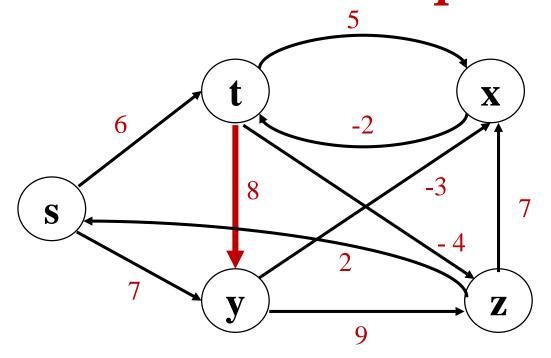
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t



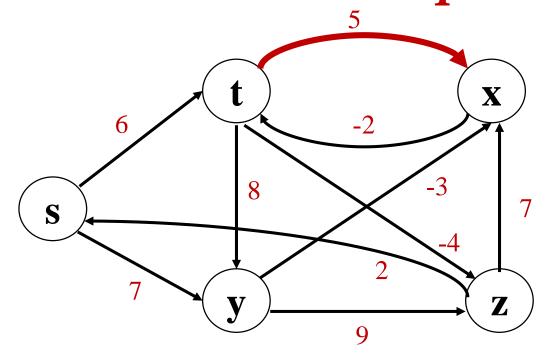
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t
2nd	0	2/x	4/y	7/s	-2/t



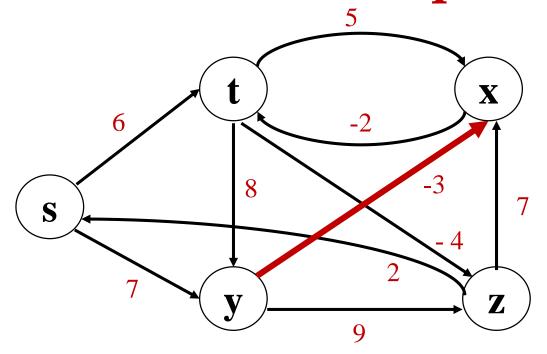
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t
2nd	0	2/x	4/y	7/s	-2/t



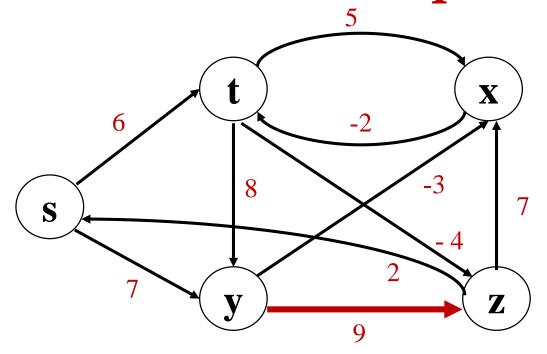
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t
2nd	0	2/x	4/y	7/s	-2/t



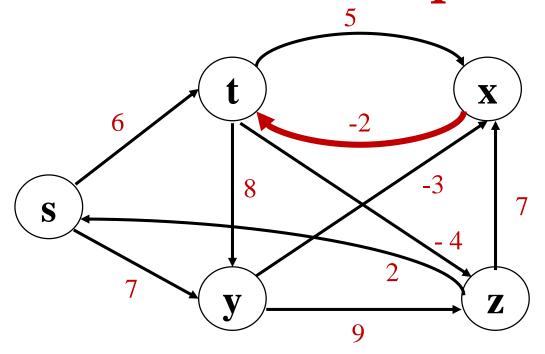
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t
2nd	0	2/x	4/y	7/s	-2/t



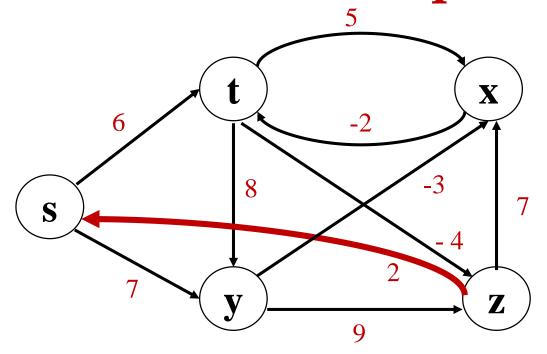
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t
2nd	0	2/x	4/y	7/s	-2/t



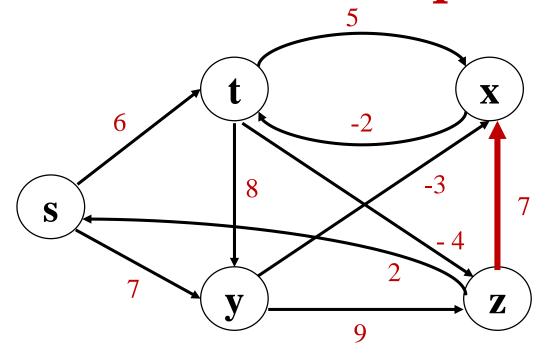
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t
2nd	0	2/x	4/y	7/s	-2/t



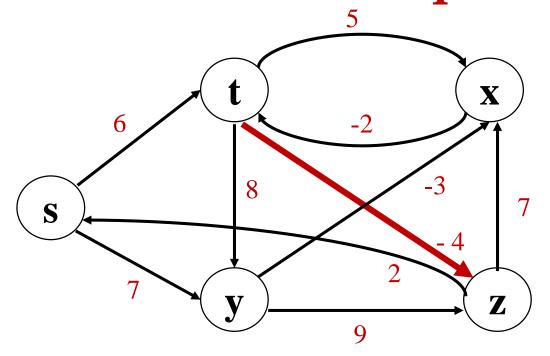
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t
2nd	0	2/x	4/y	7/s	-2/t



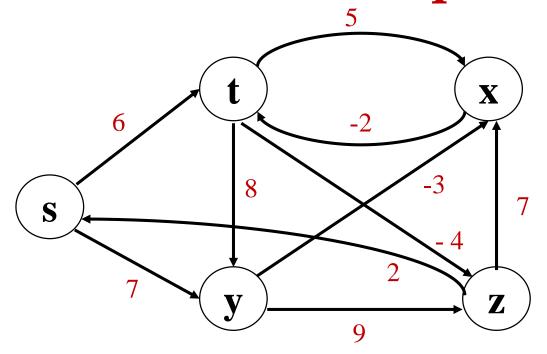
Vertex	S	t	x	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t
2nd	0	2/x	4/y	7/s	-2/t



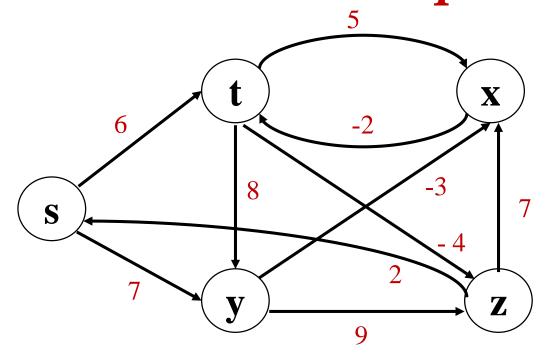
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t
2nd	0	2/x	4/y	7/s	-2/t



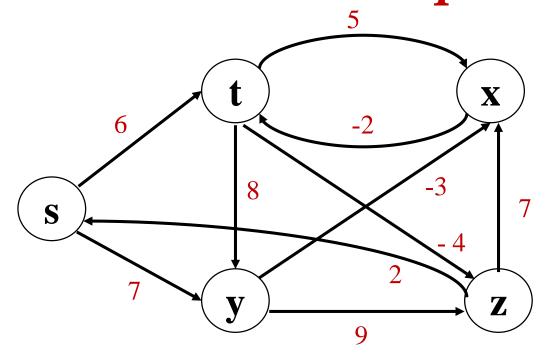
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t
2nd	0	2/x	4/y	7/s	-2/t



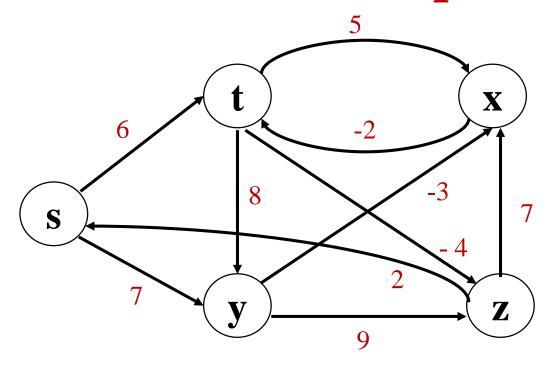
Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t
2nd	0	2/x	4/y	7/s	-2/t



Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t
2nd	0	2/x	4/y	7/s	-2/t
3rd	0	2/x	4/y	7/s	-2/t



Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t
2nd	0	2/x	4/y	7/s	-2/t
3rd	0	2/x	4/y	7/s	-2/t
4th	0	2/x	4/y	7/s	-2/t

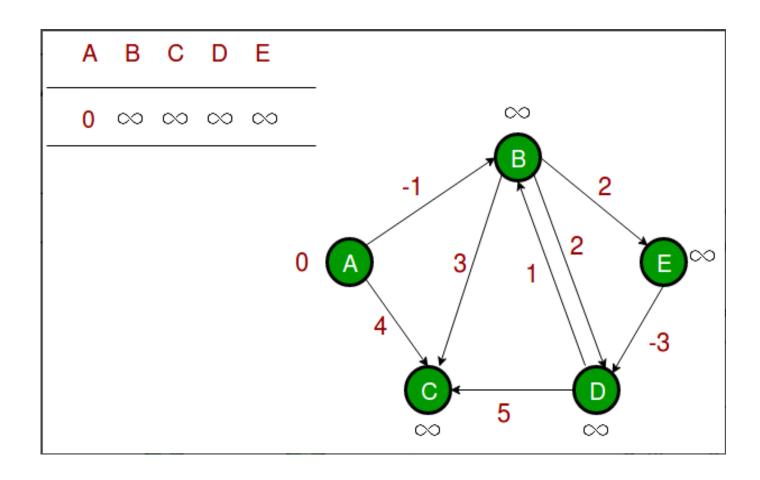


Vertex	S	t	х	У	Z
Initial	0/NIL	∞/NIL	∞/NIL	∞/NIL	∞/NIL
1st	0	2/x	4/y	7/s	-2/t
2nd	0	2/x	4/y	7/s	-2/t
3rd	0	2/x	4/y	7/s	-2/t
4th	0	2/x	4/y	7/s	-2/t
5th					

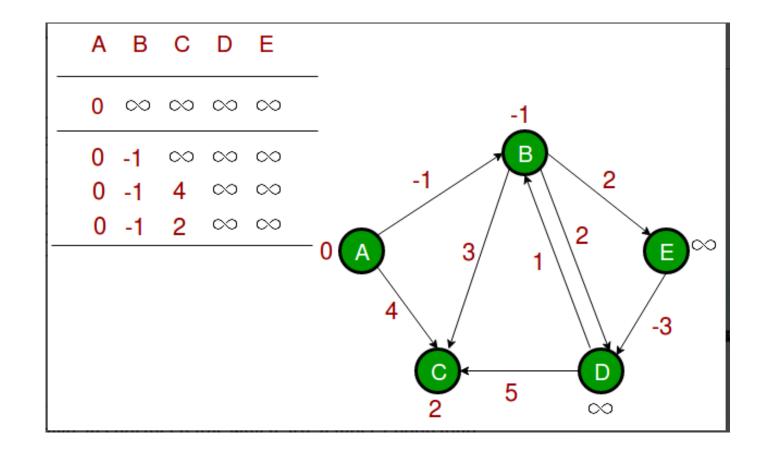
Edge List: (s,t), (s,y), (t,y), (t,x), (y,x), (y,z), (x,t), (z,s), (z,x), (t,z)

If distance is decreased(relaxed) in 5th iteration, it means there exist a negative weight cycle.

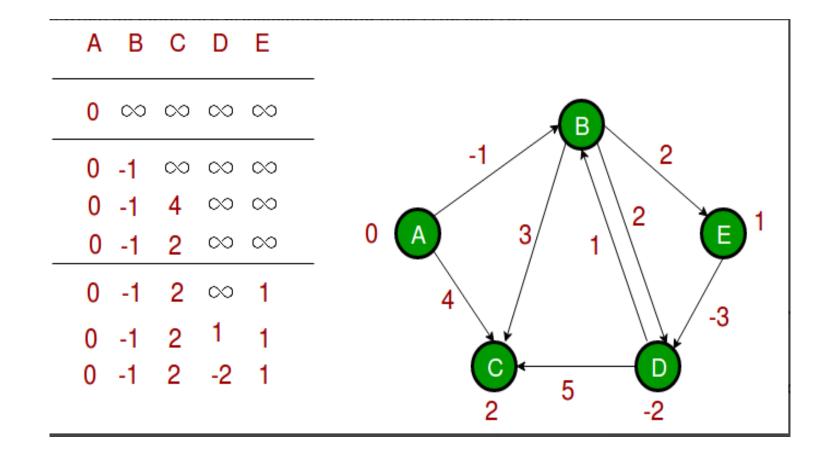
Example



Example



Example



Complexity

The Bellman-Ford algorithm runs in time O(VE), since the initialization in line 1 takes $\Theta(V)$ time, each of the |V|-1 passes over the edges in lines 2–4 takes $\Theta(E)$ time, and the for loop of lines 5–7 takes O(E) time.

Applications

- For the Internet specifically, there are many protocols that use Bellman-Ford Information protocol.
- ➤ for example the Routing Information Protocol(RIP).
- A second example is the interior gateway routing protocol to help machines exchange routing data within a system.

Thank