

# Real Time Operating systems (RTOS) concepts

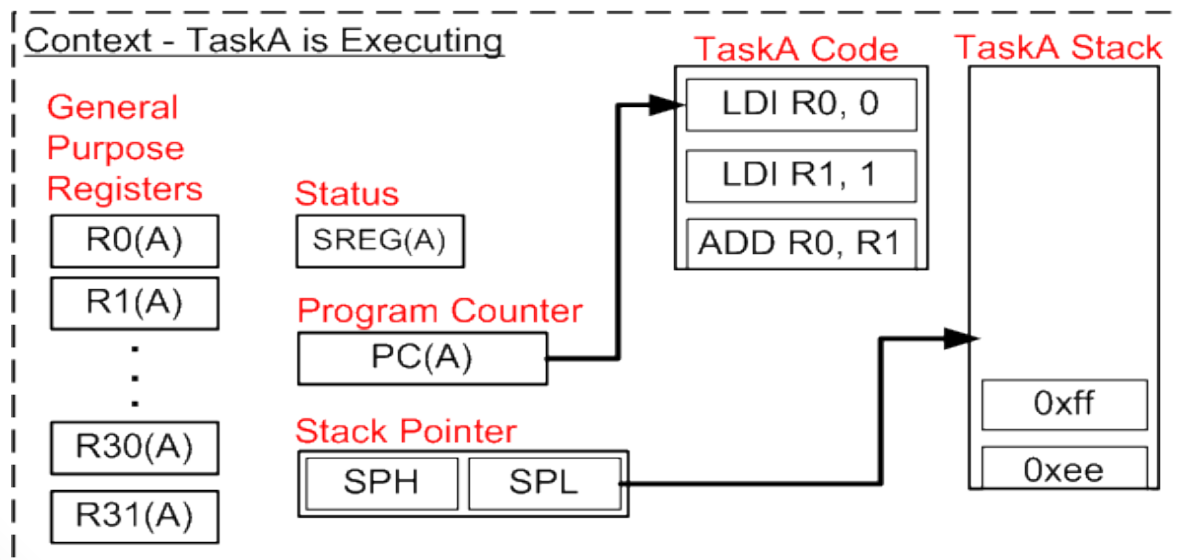
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# Content:

- Task Context.
- Context Switching between tasks.
- Shared Data Problem.
- *Non Reentrant Function.*
- Reentrant Function.
- Gray area of reentrancy.
- How to protect Shared Data?
- References and Read more

# Task Context

- Every task has it's own context (it's own Data).
- Every Task Created has it's own data structure called Task Control Block (TCB).
- Task saves it's data like: tasks status, ID , priority, stack pointer, Pointer to function( task itself) ....., in it's **TCB**.
- Task context also saved in it's own stack and CPU registers.

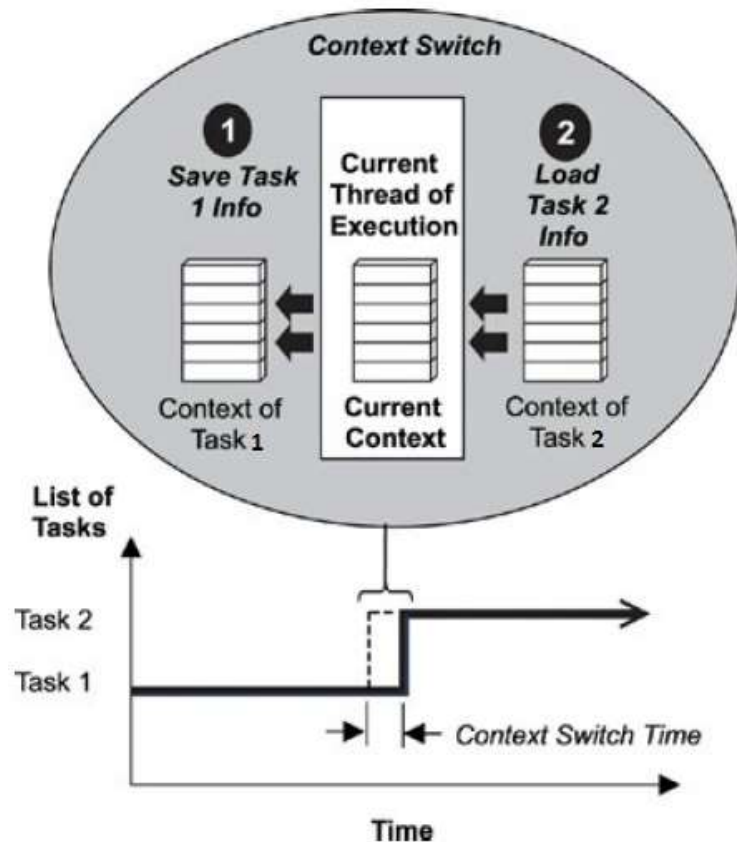


# Context Switching between tasks

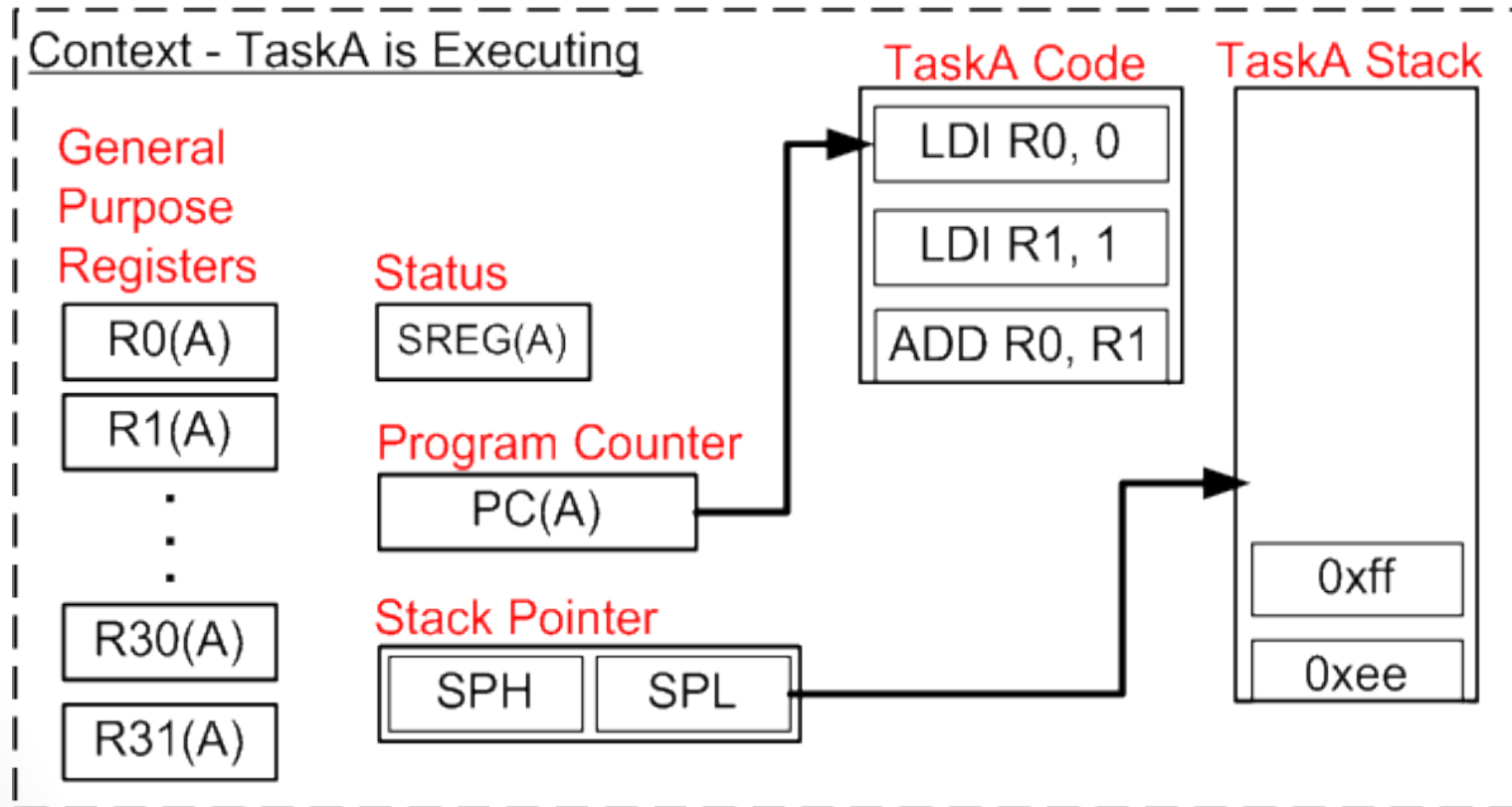
- **Context switching** : Is how to switch the processor between the context of one task to the another, so the system must:
  - Save the state of the old process,
  - Then load the saved state for the new process,
  - The new process continues from where it left off just before the context switch.
- When the task is not running, its context is frozen within the TCB, to be restored the next time the task runs.
- **The Dispatcher** :
  - Is the part of the scheduler that performs context switching.

# Context Switching between tasks Cont.

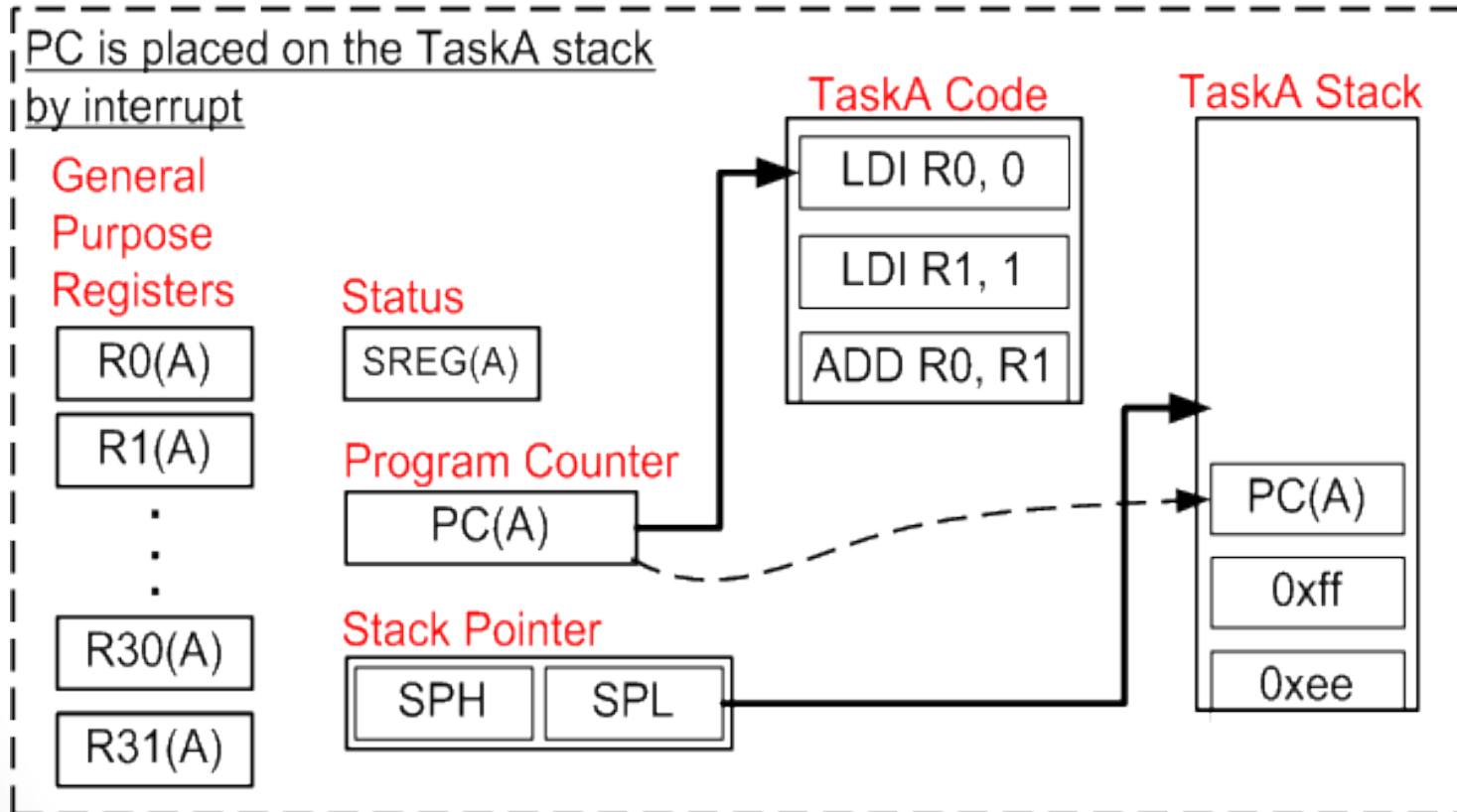
- **Context Switch Time:** Is the time it takes for the scheduler to switch from one task to another.
- frequent context switching makes a performance overhead.



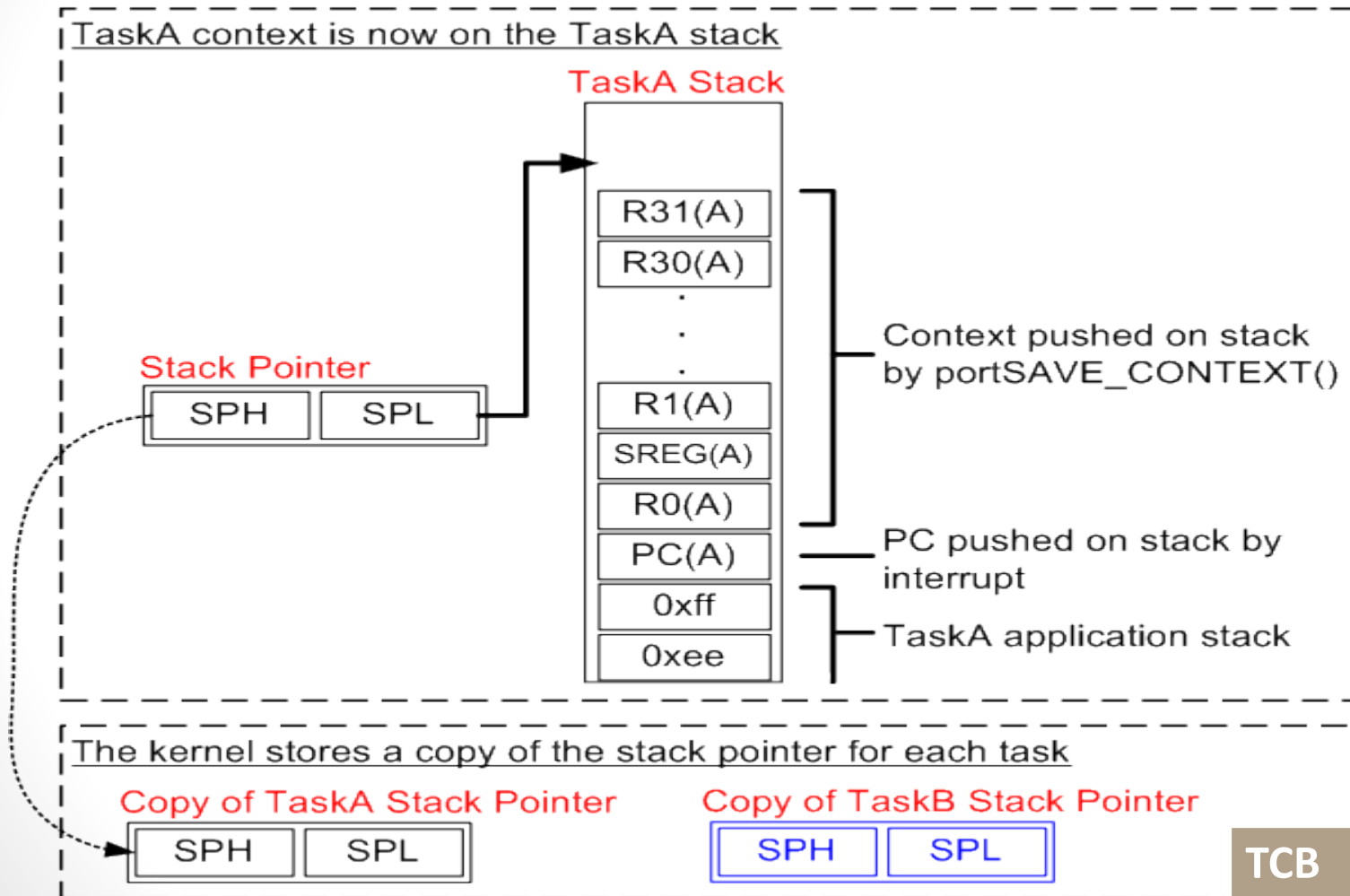
# Context Switching between tasks Cont.



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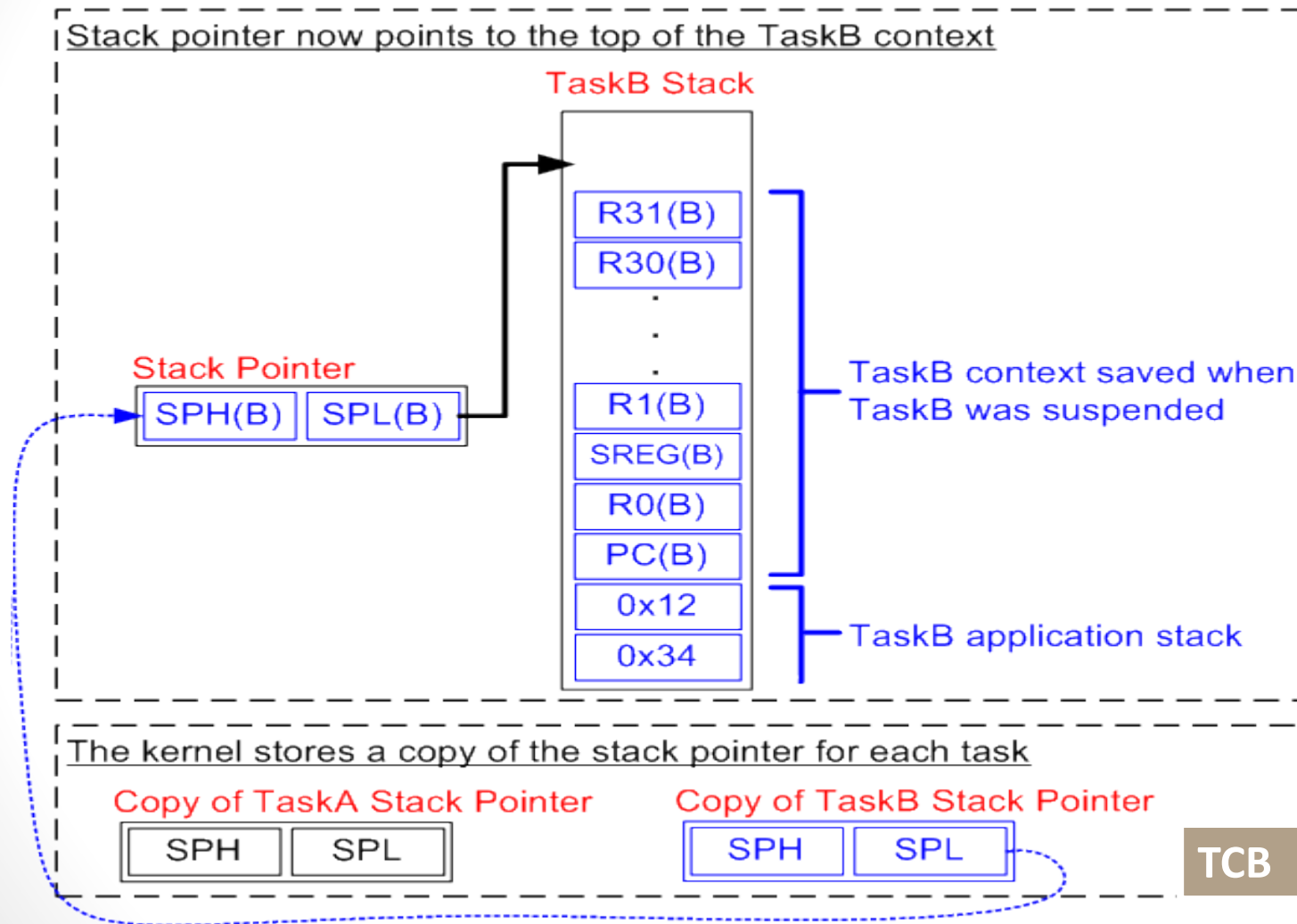


# Context Switching between tasks Cont.

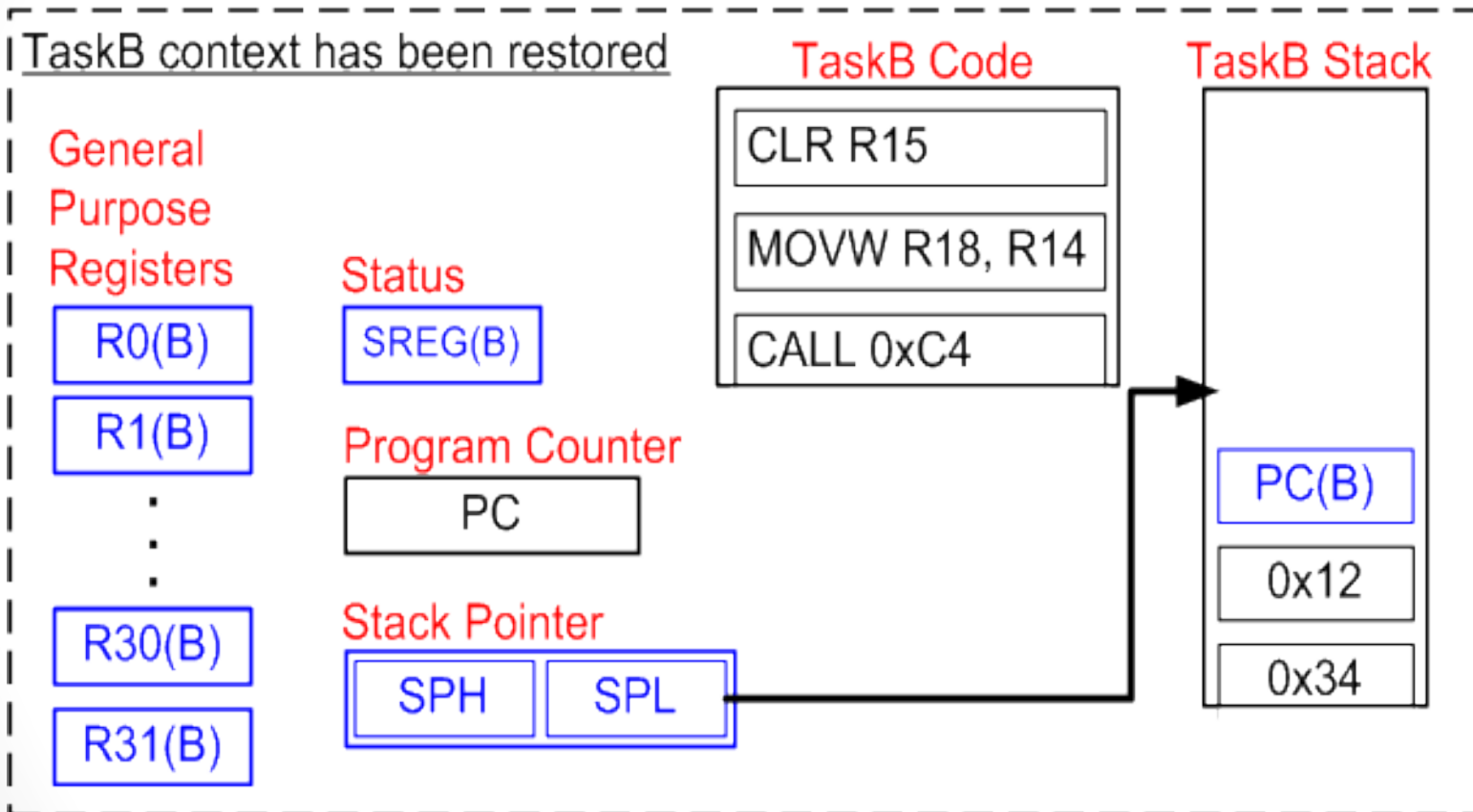




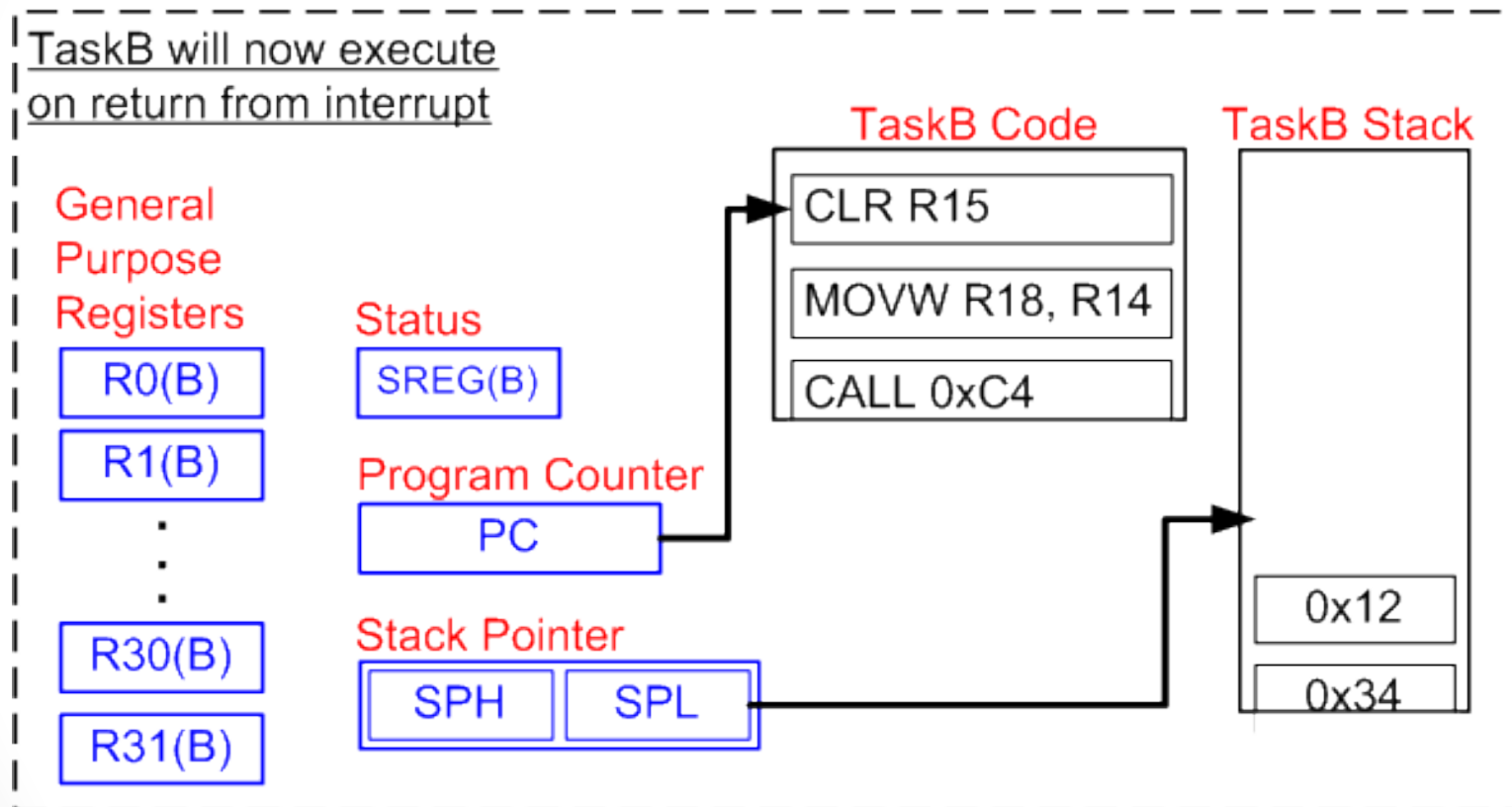
# Context Switching between tasks Cont.



# Context Switching between tasks Cont.



# Context Switching between tasks Cont.



# Shared Data Problem.

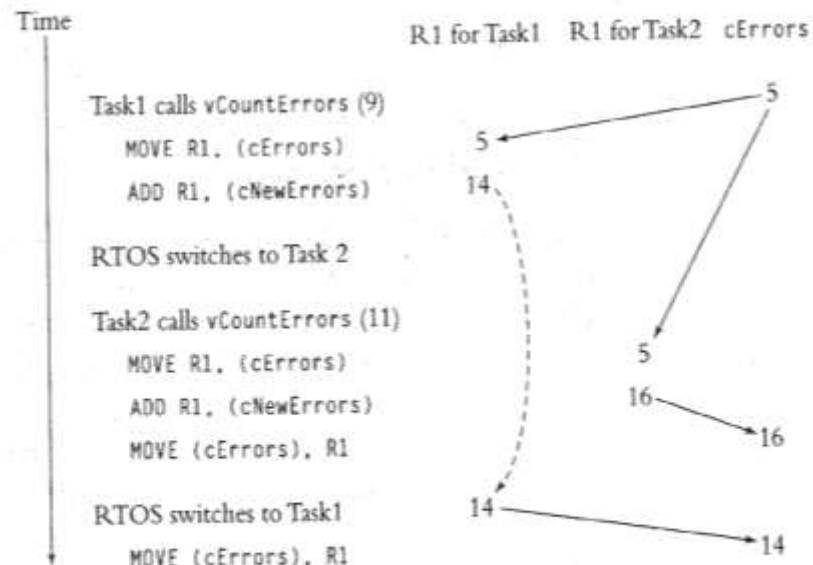
```
void Task1 (void)
{
    :
    :
    vCountErrors (9);
    :
    :
}
```

```
void Task2 (void)
{
    :
    :
    vCountErrors (11);
    :
    :
}
```

```
static int cErrors;

void vCountErrors (int cNewErrors)
{
    cErrors += cNewErrors;
}
```

```
: Assembly code for vCountErrors
:
: void vCountErrors (int cNewErrors)
:{
:   cErrors += cNewErrors;
:   MOVE R1, (cErrors)
:   ADD R1, (cNewErrors)
:   Move (cErrors), R1
:   RETURN
:}
```



# *Non Reentrant Function.*

- Is a function that can't be used between more than one task.
- Can't be interrupted or a data loss will happen.
- Example:

```
static int cErrors;  
  
void vCountErrors (int cNewErrors)  
{  
    cErrors += cNewErrors;  
}
```

# Reentrant Function.

- Is a function that can be used between more than one task with out fear of data corruption.
- Can be interrupted at any time and resumed without loss of data.
- Reentrant functions either:
  - Use local variables, Or use protected global variables.
  - Can't Call other non reentrant function.

# Gray area of reentrancy

- **Some Functions and some operation on a shared data are processor and compiler dependent.**
- **Example:**
  - **printf() Function Is reentrant or non reentrant?**
    - The answer is it's depend on the processor and on the compiler.
  - **A++**

# How to protect Shared Data?

- **Using Mutual Exclusion access.**
- **Examples on Mutual Exclusion methods are:**
  - Disable and enable interrupts,
  - Disabling Scheduling, and
  - Using Semaphores.



# References and Read more:

- **Real-Time Concepts for Embedded Systems book** by Qing Li and Carolyn.
  - <http://www.e-reading.club/book.php?book=102147>
- **An Embedded Software Primer** by David E. Simon.
  - <http://www.amazon.com/Embedded-Software-Primer-David-Simon/dp/020161569X>
- **Linux Kernel Embedded Systems Building Blocks 2e** by Jean J. Labrosse.
  - <http://www.amazon.com/Embedded-Systems-Building-Blocks-Ready/dp/0879306041>
- **FreeRTOS website.**
  - <http://www.freertos.org>