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; Mechanically Proving Approximation Bounds on Lattice-Linear Parallel
; Algorithms
;
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; 1. Introduction
;
; Lattice-Linear Predicates (LLPs) describe a class of predicates
; defined over lattices of vectors for which the problem of finding a
; minimal satisfying vector has a simple parallel/distributed algorithm.
; LLPs have been shown to be effective at parallelizing a variety of
; well-known optimization problems improving upon existing solutions. Garg's
; algorithm for detecting LLPs runs in time proportional to the height of
; the lattice, offering a large cut-down on the solution space. For
; lattices with very large height, however, it is simply infeasible to
; parallelize any solution. Many such problems are NP-hard, and thus if  $P \neq$ 
; NP, then there does not exist any algorithm for them in NC.
;
;
; To combat the hardness of such problems, a common remedy is to
; resort to approximations. While LLPs describe optimization problems,
; it is hard to encode approximation algorithms as LLPs because of
; their optimality. In this project, we discover a method for
; exploiting the lattice via abstraction to implicitly achieve
; approximations with LLPs. Furthermore, we represent this method in ACL2 to
; mechanically construct robust proofs of correctness.
;
; 2. Background
;
; A lattice is a partially-ordered set where each pair of elements has a
; well-defined least upper bound (join) and greatest lower bound (meet).
; The lattice is said to be distributive if join distributes over meet
; (equivalent to meet distributes over join).
;
; A Galois connection between two lattices A and B is a pair of functions
;  $\alpha : A \rightarrow B$  and  $\gamma : B \rightarrow A$  such that  $\gamma(\alpha(a)) \geq a$  and
;  $\alpha(\gamma(b)) \leq b$ . A Galois connection is called a Galois insertion
; (GI) if  $\alpha(\gamma(b)) = b$ . That is, B is exactly representable in A. We
; will assume only GIs throughout this project since it is well-known that a
; Galois connection may be reduced to a GI. It is also known that  $\alpha$ 
; uniquely determines  $\gamma$ , and vice versa. Finally, it is known that for
; order-preserving  $f : A \rightarrow A$   $(\alpha \circ f \circ \gamma) \leq g$  pointwise for all  $g : B \rightarrow B$ 
; such that  $\alpha \circ f \leq g \circ \alpha$ . These properties allow us to
; simplify the project, and we do not verify them in ACL2.
;
; d is a quasi-metric on A if for all  $x, y, z$  in A,  $d(x, y) = 0$  iff  $x = y$ ,
;  $d(x, y) > 0$  iff  $x > y$ , and  $d(x, y) + d(y, z) \geq d(x, z)$ . For the purposes of
; this problem, a specific measure and metric is important, but that is
; unimportant for this project.
;

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; Let L be a distributive lattice of n -dimensional vectors and B be a
; predicate over L . Then, we define $\text{forbidden}(G, i, B) := \text{forall } H \text{ in } L \text{ such}$
; that $G \leq H, G[i] = H[i] \implies \text{not } B(H)$. This states that component i is
; forbidden in vector G in L if, no matter how G is progressed, fixing i
; keeps the predicate false. We then say that B is lattice-linear if for any
; G that makes B false, G has a forbidden component. For the purposes of
; minimality, we also need to know how much to progress a forbidden
; component so as to make it no longer forbidden while not jumping over the
; optimal solution. Let $\text{delta}(G, i, B)$ be the greatest value such that for all
; H in L such that $H \geq G, H[i] < \text{delta}(G, i, B) \implies \text{not } B(H)$ whenever
; $\text{forbidden}(G, i, B)$. That is, it is the least value we need to bring
; component i up to in order to make it no longer forbidden. Advancing a
; vector G with a forbidden component i is done by setting the i th component
; to $\text{delta}(G, i, B)$.

; Distributivity in the lattice is only necessary for certain performance
; improvements, so for the purposes of this project we do not worry
; ourselves with it. We later note the potential for future work with
; distributive lattices. For the purposes of this paper, and as is typically
; and classically presented in LLPs, we will consider vectors ordered
; component-wise throughout this project and report.

; 3. A Method for Approximating LLP Problems

; Suppose we have an LLP B over a lattice of n -dimensional vectors L with a
; quasi-metric d defined over it. Let L' be a lattice of m -dimensional
; vectors with (α, γ) a GI between L and L' . We say that $f : L \rightarrow L,$
; $f' : L' \rightarrow L'$ are epsilon-complete (e.c.) w.r.t the GI if $d(f(G), (\gamma$
; $\circ f' \circ \alpha)(G)) \leq \epsilon$ for all G in L .

; The basic idea is as follows: with L' being a lattice with smaller height
; than L , progress through the smaller lattice, checking whether we have
; stepped over a feasible solution in L (by simply checking $B(f(G))$ where G
; is the last position we checked). If $f(G)$ is feasible, then it must be
; optimal (proven in ACL2, described later), so $d(f(G), (\gamma \circ f' \circ$
; $\alpha)(G)) = d(\text{OPT}, (\gamma \circ f' \circ \alpha)(G)) \leq \epsilon$ where OPT is the
; optimal solution in L . We can abuse the GI property to make this more
; efficient (also proven in ACL2) since repeated application of $(\gamma \circ f'$
; $\circ \alpha)$ gives $(\gamma \circ f' \circ \alpha \circ \gamma \circ f' \circ \alpha) = (\gamma \circ f' \circ$
; $\alpha)$. Many applications of this will cancel out many instances of $(\alpha$
; $\circ \gamma)$ by the GI property, thereby saving many unnecessary computations.

; Counter-intuitively, the epsilon distance does not accumulate because we
; "fix" a comparison point in L at each step. This is the crux and novelty
; of the project, proven in ACL2.

; 4. ACL2 Representation and Results

; In this section, we will use "components" and "states" interchangeably.

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; We represent LLPs abstractly so as to apply these proofs to any concrete
; or realized instance by utilizing the "encapsulate" primitive from ACL2.
; This approach provides the definitional properties of LLPs without
; exposing the definition of the abstract functions and predicates. The
; encapsulate includes properties for the necessary theory / realizable
; functions and constants.
;
; In LLPs, lattices are typically over vectors of reals ordered
; component-wise, however for the sake of simplicity in this project, we
; consider vectors of naturals ordered similarly. This comes in the form of
; the definitions vecp and lattice-leq, both of which have a small number of
; associated theorems.
;
; From these definitions and the basic properties provided by encapsulate,
; we prove a number of important lemmas. The most important theorems proven
; are:
;   1. bounded-by-opt:
;       This proves that advancing a vector below OPT keeps it
;       below OPT.
;   2. gi-makes-stepping-efficient:
;       This proves that GI's enable an optimization on
;       the approximation algorithm.
;   3. constant-approx-dist:
;       This proves that the algorithm gives a constant approximation
;       factor (namely, epsilon) GIVEN that it only takes a finite
;       number of steps to reach OPT in the standard LLP algorithm.
;   4. lattice-leq-component & advance-is-higher-lattice:
;       These prove that advancing a vector with the aforementioned
;       advance rule moves it strictly up the lattice. This directly
;       implies (with bounded-by-opt) that it only takes a finite
;       number of steps to reach OPT in the standard LLP algorithm
;       (thus supporting constant-approx-dist).
;
; 5. Conclusion
;
; In this project, presented a novel technique for representing
; approximation algorithms with LLPs, modeled this technique in ACL2, and
; utilized the ACL2 theorem prover to prove an approximation bound on this
; technique. Using an automated theorem prover like ACL2 not only gave strong
; assurance of the technique's validity, but also offered valuable insights
; into automated deduction and the technique's underlying recursive
; structure. This project can be further extended to verify very important
; LLP optimizations utilizing lattice distributivity and Birkhoff's theorem.
; It would also be a good step to apply these ACL2 theorems to verify an
; application of this technique.

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(defun indexp (x y) (and (natp x) (natp y) (>= x 1) (<= x y)))

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(defun vecp (G n)

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(if (zp n)
    (null G)
    (and (consp G)
          (natp (car G))
          (<= 0 (car G))
          (vecp (cdr G) (- n 1)))))

(defthm vecp-length-strict
  (implies (and (natp n) (vecp G n))
            (equal (len G) n)))

(defthm vecp-is-true-listp
  (implies (and (natp n) (vecp x n))
            (true-listp x)))

(defun ith (G i)
  (if (or (atom G) (<= i 1))
      (car G)
      (ith (cdr G) (- i 1))))

(defthm vecp-nat-and-nonneg
  (implies (and (natp n) (>= n 1) (vecp G n) (indexp i n))
            (natp (ith G i)))
  :rule-classes (:type-prescription))

(defun set-ith (G i v)
  (if (atom G)
      nil
      (if (equal i 1)
          (cons v (cdr G))
          (cons (car G)
                (set-ith (cdr G) (- i 1) v)))))

(defthm set-ith-sets
  (implies (and (natp n) (indexp i n) (vecp G n))
            (equal (ith (set-ith G i v) i) v)))

(defthm set-ith-conserves
  (implies (and (natp n) (indexp i n) (indexp j n)
                (not (equal i j)) (vecp G n))
            (equal (ith (set-ith G i v) j) (ith G j))))

(defthm set-ith-type-preservation
  (implies (and (natp n) (indexp i n) (vecp G n) (natp v) (>= v 0))
            (vecp (set-ith G i v) n)))

(defun lattice-leq (x y)
  (if (consp x)
      (and (<= (car x) (car y)) (lattice-leq (cdr x) (cdr y)))
      T))

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(defthm lattice-leq-recurs-car
  (implies (and (natp n) (>= n 1) (vecp x n) (vecp y n) (lattice-leq x y))
    (<= (car x) (car y))))
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(defthm lattice-leq-recurs-cdr
  (implies (and (natp n) (>= n 1) (vecp x n) (vecp y n) (lattice-leq x y))
    (lattice-leq (cdr x) (cdr y))))
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(defthm lattice-leq-reflexive
  (implies (and (natp n) (vecp x n) (vecp y n))
    (lattice-leq x x)))
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(defthm lattice-leq-antisymmetric
  (implies (and (natp n) (>= n 1) (vecp x n) (vecp y n)
    (lattice-leq x y) (not (equal x y)))
    (not (lattice-leq y x)))
  :hints (("Goal" :use (vecp-is-true-listp)))))
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(defthm lattice-leq-transitive
  (implies (and (natp n) (vecp x n) (vecp y n)
    (lattice-leq x y) (lattice-leq y z))
    (lattice-leq x z)))
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(encapsulate
  (
    ((forbidden * *) => *)
    ((feasible *) => *)
    ((feasible-abs *) => *)
    ((alpha *) => *)
    ((gamma *) => *)
    ((advance * *) => *)
    ((advance-abs *) => *)
    ((dist * *) => *)
    ((ep) => *)
    ((delta * *) => *)
    ((get-forbidden-i * *) => *)
    ((get-forbidden *) => *)
    ((opt *) => *)
    ((stepmin) => *)
  )
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(local (defun feasible (x) (if x t t)))
(local (defun feasible-abs (x) (if x t t)))
(local (defun alpha (x) x))
(local (defun gamma (x) x))
(local (defun forbidden (G i) (if (and i G) nil nil)))
(local (defun advance (G i) (if i G G)))
(local (defun advance-abs (G) G))
(local (defun dist (x y) (if (and x y) 0 0)))
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(local (defun ep () 1))
(local (defun delta (G i) (+ 1 (ith G i))))
(local (defun get-forbidden-i (G i)
  (if (or (zp i) (atom G))
      nil
      (if (forbidden G i)
          i
          (get-forbidden-i G (- i 1))))))
(local (defun get-forbidden (G) (get-forbidden-i G (len G))))
(local (defun opt (n) (if (zp n) nil (cons 0 (opt (- n 1)))))
(local (defun stepmin () 1))

(defthm alpha-is-vecp
  (implies (and (natp n) (>= n 1) (vecp G n))
    (vecp (alpha G) n)))

(defthm gamma-is-vecp
  (implies (and (natp n) (>= n 1) (vecp G n))
    (vecp (gamma G) n)))

(defthm advance-abs-is-vecp
  (implies (and (natp n) (>= n 1) (vecp G n))
    (vecp (advance-abs G) n)))

(defthm opt-is-vecp (implies (natp n) (vecp (opt n) n)))

(defthm stepmin-is-natp (natp (stepmin)))

(defthm ep-is-natp (natp (ep)))

(defthm get-forbidden-sem-zero
  (implies (or (zp i) (atom G))
    (equal (get-forbidden G) nil)))

(defthm get-forbidden-sem-recurs
  (implies (and (natp n) (> n 0) (vecp G n) (get-forbidden G))
    (equal (get-forbidden G) (+ 1 (get-forbidden (cdr G))))))

(defthm get-forbidden-i-is-indexp
  (implies (and (natp n) (>= n 1) (vecp G n) (indexp i n)
    (get-forbidden-i G i))
    (indexp (get-forbidden-i G i) n)))

(defthm get-forbidden-i-is-forbidden
  (implies (and (natp n) (>= n 1) (vecp G n) (indexp i n)
    (get-forbidden-i G i))
    (forbidden G (get-forbidden-i G i))))

(defthm get-forbidden-is-indexp
  (implies (and (natp n) (>= n 1) (vecp G n) (get-forbidden G))

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      (indexp (get-forbidden G) n)))

(defthm get-forbidden-is-forbidden
  (implies (and (natp n) (>= n 1) (vecp G n) (get-forbidden G))
    (forbidden G (get-forbidden G))))

(defthm get-forbidden-gives-nil
  (implies (and (natp n) (>= n 1) (vecp G n)
    (not (get-forbidden G)) (indexp i n))
    (not (forbidden G i))))

(defthm get-forbidden-nil-is-feasible
  (implies (and (natp n) (>= n 1) (vecp G n) (not (get-forbidden G)))
    (feasible G)))

(defthm advance-goes-up
  (implies (and (natp n) (>= n 1) (indexp i n) (vecp G n))
    (lattice-leq G (advance G i))))

(defthm advance-makes-unforbidden
  (implies (and (natp n) (>= n 1) (indexp i n) (vecp G n)
    (forbidden G i))
    (not (forbidden (advance G i) i))))

(defthm advance-progresses-by-delta
  (implies (and (natp n) (>= n 1) (vecp G n) (indexp i n)
    (forbidden G i))
    (equal (advance G i) (set-ith G i (delta G i)))))

(defthm ag-galois-connection
  (implies (and (natp n) (>= n 1) (vecp G n) (vecp H n))
    (lattice-leq G (gamma (alpha G)))))

(defthm ag-galois-insertion
  (implies (and (natp n) (>= n 1) (vecp G n))
    (equal (alpha (gamma G)) G)))

(defthm llp-abs-soundness
  (implies (and (natp n) (vecp G n) (feasible-abs G))
    (feasible (gamma G))))

(defthm forbidden-def
  (implies (and (natp n) (>= n 1) (indexp i n) (vecp G n) (vecp H n)
    (forbidden G i) (lattice-leq G H)
    (equal (ith G i) (ith H i)))
    (not (feasible H))))

(defthm llp-def
  (implies (and (natp n) (>= n 1) (indexp i n) (vecp G n) (feasible G))
    (not (forbidden G i))))

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(defthm ep-completeness
  (implies (and (natp n) (>= n 1) (indexp i n) (vecp G n)
    (forbidden G i))
    (<= (dist (advance G i)
      (gamma (advance-abs (alpha G))))
      (ep))))

(defthm dist-triangle-inequality
  (implies (and (natp n) (>= n 1) (vecp G n) (vecp H n) (vecp A n))
    (<= (dist G A) (+ (dist G H) (dist H A)))))

(defthm dist-zp-is-eq
  (implies (and (natp n) (>= n 1) (vecp G n))
    (equal (dist G G) 0)))

(defthm delta-forbidden
  (implies (and (natp n) (>= n 1) (indexp i n) (vecp G n) (vecp H n)
    (forbidden G i) (lattice-leq G H)
    (< (ith H i) (delta G i)))
    (not (feasible H))))

(defthm no-forbidden-is-feasible
  (implies (and (natp n) (>= n 1) (vecp G n) (not (get-forbidden G)))
    (feasible G)))

(defthm opt-is-feasible
  (implies (and (natp n) (>= n 1))
    (feasible (opt n))))

(defthm opt-is-optimal
  (implies (and (natp n) (>= n 1) (vecp G n) (lattice-leq G (opt n))
    (not (equal G (opt n))))
    (not (feasible G))))

(defthm opt-is-optimal-contrapositive
  (implies (and (natp n) (>= n 1) (vecp G n) (feasible G))
    (lattice-leq (opt n) G)))

(defthm delta-type
  (implies (and (natp n) (>= n 1) (vecp G n) (indexp i n))
    (natp (delta G i)))
  :hints (("Goal" :use (vecp-nat-and-nonneg))))

(defthm delta-bound
  (implies (and (natp n) (>= n 1) (vecp G n) (indexp i n)
    (forbidden G i))
    (>= (- (delta G i) (ith G i)) (stepmin))))

(defthm increase-stepmin

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        set-ith-type-preservation))))

(defthm advance-is-higher-lattice
  (implies (and (natp n) (>= n 1) (vecp G n) (indexp i n) (forbidden G i))
    (not (equal (advance G i) G)))
  :hints (("Goal" :use (advance-is-higher
    advance-is-vecp
    (:instance this-is-so-simple
      (H (advance G i)))))))

(defthm advance-conserves-other
  (implies (and (natp n) (>= n 1) (vecp G n)
    (indexp i n) (indexp j n)
    (forbidden G i) (not (equal i j)))
    (equal (ith (advance G i) j) (ith G j)))
  :hints (("Goal" :use (advance-progresses-by-delta set-ith-conserves))))

(defthm equal-implies-≤
  (implies (equal x y)
    (≤ x y))
  :rule-classes (:rewrite))

(defthm <-implies-≤
  (implies (< x y)
    (≤ x y))
  :rule-classes (:rewrite))

(defthm lattice-leq-component
  (implies (and (natp n) (>= n 1) (vecp G n) (indexp i n) (indexp j n)
    (forbidden G i))
    (≤ (ith G j) (ith (advance G i) j)))
  :hints (("Goal"
    :cases ((equal i j)
      (not (equal i j)))
    :use ( advance-progresses-by-delta
      set-ith-sets
      advance-conserves-other
      advance-is-higher ))))

(defthm no-forbidden-below-opt-is-opt
  (implies (and (natp n) (>= n 1) (vecp G n) (not (get-forbidden G))
    (lattice-leq G (opt n)))
    (equal (opt n) G))

(defthm advance-keeps-other-states-fixed
  (implies (and (natp n) (>= n 1) (vecp G n) (indexp j n) (indexp i n)
    (not (equal i j)) (forbidden G i))
    (equal (ith (advance G i) j) (ith G j)))
  :hints (("Goal" :use (advance-progresses-by-delta set-ith-conserves))))

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(defthm forbidden-is-infeasible
  (implies (and (natp n) (>= n 1) (vecp G n) (indexp i n) (forbidden G i))
    (not (feasible G))))

(defthm bounded-by-opt
  (implies (and (natp n) (>= n 1) (vecp G n) (indexp i n) (forbidden G i)
    (lattice-leq G (opt n)))
    (lattice-leq (advance G i) (opt n)))
  :hints (("Goal" :use (opt-is-feasible opt-is-vecp
    (:instance detect-llp-step-stays-below
      (M (opt n))))
    :in-theory (disable advance-progresses-by-delta))))

(defthm apply-gi
  (implies (and (natp n) (>= n 1) (vecp G n)
    (equal (alpha (gamma (advance-abs (alpha G))))
      (advance-abs (alpha G))))
  :hints (("Goal" :use ((:instance ag-galois-insertion
    (G (advance-abs (alpha G))))))))

(defun abstract-step (G k)
  (if (zp k)
    G
    (abstract-step (advance-abs G) (- k 1))))

(defun transit-step (G k)
  (if (zp k)
    G
    (transit-step (gamma (advance-abs (alpha G))) (- k 1))))

(defthm gi-makes-stepping-efficient
  (implies (and (natp n) (>= n 1) (natp k) (>= k 1) (vecp G n)
    (equal (gamma (abstract-step (alpha G) k)) (transit-step G k)))
  :hints (("Goal" :use ((:instance ag-galois-insertion
    (G (advance-abs (alpha G))))))))

(defthm step-to-opt
  (implies (and (natp n) (>= n 1) (vecp G n) (indexp i n) (forbidden G i)
    (lattice-leq G (opt n))
    (lattice-leq (opt n) (advance G i)))
    (equal (opt n) (advance G i)))

(defthm constant-approx-dist
  (implies (and (natp n) (>= n 1) (vecp G n) (indexp i n) (forbidden G i)
    (lattice-leq G (opt n))
    (lattice-leq (opt n) (advance G i)))
    (<= (dist (opt n) (gamma (advance-abs (alpha G)))
      (ep)))
  :hints (("Goal" :use (step-to-opt ep-completeness))))

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