## Assignment 2 Functions - SOLUTIONS Due Date: Wednesday, January 19

January, 2022

## 1 Preamble

This assignment covers everything taught so far. The solutions that you hand in should be **neat** and **legible**, this is an assignment, not a quiz, so I expect you to take your time and present thorough and detailed solutions.

## 2 Name and Date:

Print your name and todays date below;		
Name	Date	

Question 1. There exists a function that allows us to determine the length of a binary string. We call this function len. Here a few examples to understand how it works,

- If S = 1001, then len(S) = 4.
- If  $\mathbf{R} = 110001$ , then  $len(\mathbf{R}) = 6$ .
- If  $\mathbf{T} = \epsilon$ , then len( $\mathbf{T}$ ) = 0.

We can also define the operation of multiplication by a scalar for binary strings. So suppose  $n \in \mathbb{N}$  and **S** is some binary string, then,

$$n \cdot \mathbf{S} = \underbrace{\mathbf{S} + \dots + \mathbf{S}}_{\text{n times}}$$

Again we resort to a few examples to demonstrate how multiplication by a scalar works,

- If S = 1001, then  $2 \cdot S = S + S = 10011001$ .
- If R = 0, then  $4 \cdot R = R + R + R + R = 0000$ .
- If T = 01, then  $3 \cdot T = T + T + T = 010101$ .

Let S = 001 and T = 11, answer the following,

- (a) Let S represent the set of all binary strings. Define the length function using mapping notation.
- Solution. len:  $\mathbb{S} \to \mathbb{N}$ .

**Solution.** len(S) = len(001) = 3.

(c) Compute len(T).

(b) Compute len(S).

Solution.  $len(\mathbf{T}) = len(11) = 2$ .

(d) Compute len(S + T).

**Solution.** len(S + T) = len(001 + 11) = len(00111) = 5.

(e) Compute  $len(3 \cdot \mathbf{S})$ .

**Solution.**  $len(3 \cdot S) = len(3 \cdot 001) = len(001001001) = 9.$ 

(f) Compute  $3 \cdot \text{len}(\mathbf{S})$ .

Solution.  $3 \cdot \operatorname{len}(\mathbf{S}) = 3 \cdot 3 = 9$ .

(g) Compute  $len(4 \cdot T)$ .

**Solution.**  $len(4 \cdot T) = len(4 \cdot 11) = len(111111111) = 8.$ 

(h) Compute  $4 \cdot \text{len}(\mathbf{T})$ .

Solution.  $4 \cdot \operatorname{len}(\mathbf{T}) = 4 \cdot 2 = 8$ .

Question 2. Let F be a function. We call F linear if both of the following conditions are satisfied,

1. For all inputs x and y,

$$F(x+y) = F(x) + F(y).$$

2. For all  $c \in \mathbb{F}$ , and all inputs x,

$$F(c \cdot x) = c \cdot F(x).$$

If  $\mathbb{F} = \mathbb{N}$ , then based on your results from Question 1, do you think that the length function, len, is linear? Explain your answer.

**Solution.** From our results in parts (b),(c) we see that  $len(\mathbf{S}) + len(\mathbf{T}) = 3 + 2 = 5$ , and from part (d) we see that  $len(\mathbf{S} + \mathbf{T}) = 5$ , hence it looks like the first condition of linearity is satisfied so far.

From our result in part (e), we see that  $len(3 \cdot \mathbf{S}) = 9$  and from part (f) we see that  $3 \cdot len(\mathbf{S}) = 9$ , hence it looks like the second condition of linearity is satisfied as well.

Combining our two hypothesis, we conclude that it appears like the length function is indeed linear.

**Question 3.** Sometimes in math we would like a function that simply gets rid of trailing decimals and returns a whole number, aka an integer. This function is known as the floor function. We define it with mapping notation as floor:  $\mathbb{R} \to \mathbb{Z}$ , and it works as follows, if  $x \in \mathbb{R}$ , then floor(x) is the smallest integer that is less than or equal to x. Lets see how it works in the following examples,

- If x = 4.2, then floor(x) = floor(4.2) = 4.
- If x = -7.4, then floor(x) = floor(-7.4) = -8.
- If x = 5, then floor(x) = floor(5) = 5.
- If x = 0.4, then floor(x) = floor(0.4) = 0.
- (a) Compute floor (2.5).

**Solution.** floor(2.5) = 2.

(b) Compute floor (6/3).

**Solution.** floor(6/3) = floor(2) = 2.

(c) Compute floor (19/4).

**Solution.** floor(19/4) = floor(4.75) = 4.

(d) Let f(x) = (x+1)/2 and  $g(x) = \sqrt{x-1}$ , compute floor(f(g(5))).

**Solution.** We first compute g(5),

$$q(5) = \sqrt{5-1} = \sqrt{4} = 2.$$

Next we compute f(g(5)),

$$f(g(5)) = f(2) = \frac{2+1}{2} = \frac{3}{2}.$$

Finally we compute floor(f(g(5))),

$$floor(f(g(5))) = floor(\frac{3}{2}) = floor(1.5) = 1.$$

(e) Is the floor function linear? If it is, then justify your claim. If it is not, then provide a counter example to show that it fails to be linear.

**Solution.** The floor function is **not** linear, we draw an easy counter example to show why. Notice that,

$$floor(2.5 + 1.5) = floor(4) = 4.$$

On the other hand,

$$floor(2.5) + floor(1.5) = 2 + 1 = 3.$$

Clearly floor $(2.5 + 1.5) \neq \text{floor}(2.5) + \text{floor}(1.5)$ , and hence, the floor function is not linear.

(f) Is the floor function invertible? If it is, then justify your claim. If it is not, then provide a counter example to show that it fails to be surjective or injective.

**Solution.** The floor function is **not** invertible, notice that since floor (2.5) = floor(2.4) = 2, the floor function fails to be injective. Since it fails to be injective, it fails to be invertible.

**Question 4.** Let  $S = \{1,010,00100,0001000\}$ , where each element is a binary string, and let  $\mathcal{R} = \{4,2,6,8\}$ , where each element is a natural number.

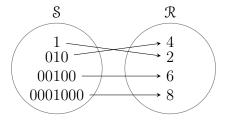
(a) Come up with an invertible function  $\Psi$  between S and  $\mathcal{R}$  and prove that your function is invertible. (**Hint:** Try using the length function)

**Solution.** I claim that,

$$\Psi \colon \mathcal{S} \to \mathcal{R}, \ \Psi(\mathbf{S}) = \operatorname{len}(\mathbf{S}) + 1.$$

is an invertible function between S and R.

*Proof:* To prove the claim from above, it is sufficient to show that the mapping diagram preserves both injectivity and surjectivity,



From the mapping diagram, we conclude that  $\Psi(\mathbf{S}) = \operatorname{len}(\mathbf{S}) + 1$  is both injective and surjective, and hence its invertible.

(b) Come up with the correct formula for the inverse function  $\Psi^{-1}$  and prove that your formula is correct using mapping tables. (**Hint:** The correct formula uses the floor function)

**Solution.** I claim that,

$$\Psi^{-1} \colon \mathcal{R} \to \mathcal{S} \,, \ \ \Psi^{-1}(r) = \mathrm{floor}\Big(\frac{r-1}{2}\Big) \cdot \mathbf{0} + \mathbf{1} + \mathrm{floor}\Big(\frac{r-1}{2}\Big) \cdot \mathbf{1}.$$

is the correct formula for the inverse function, where  $\mathbf{0}, \mathbf{1}$  are binary bits. (They are being bolded just so we don't confuse them for integers 0, 1)

*Proof:* To prove our claim, it is sufficient to show that both conditions of Definition 4.1 hold using mapping tables,

By our results from the mapping tables, we conclude that

$$\Psi^{-1}(r) = \operatorname{floor}\!\left(\frac{r-1}{2}\right) \cdot \mathbf{0} + \mathbf{1} + \operatorname{floor}\!\left(\frac{r-1}{2}\right) \cdot \mathbf{1}.$$

is indeed the inverse function for  $\Psi$ .

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