

Q1)

$$\begin{aligned}
\text{a) } T(n) &= 5T(n/3) + n \log n, \quad n=3^k \\
&= 5T(3^{k/3}) + 3^k \log 3^k \\
T(n/3) &= 5(5T(n/9) + n/3 \log n/3) + n \log n \\
T(n/9) &= 5^2 T(n/3^2) + 5n/3 \log n/3 + n \log n \\
&= 5^3 T(n/3^3) + 5^3 n/3^3 \log n/3^3 + 5n/3 \log n/3 + n \log n \\
&\cdot \\
&\cdot \\
&= 5^k T(n/3^k) + \sum_{i=0}^{k-1} (5/3)^k n \log (n/3^k), \quad (k = \log_3 n) \\
&= 5^{(\log_3 n)} T(1) + \sum_{i=0}^{\log_3 n} (5/3)^i n \log (n/3^i) \\
&= 5^{(\log_3 n / \log_3 5)} + \sum_{i=1}^k (5/3)^i n \log (n/3^i) \\
&= n^{(\log_3 5)} + \sum_{i=1}^k (5/3)^i n \log (n/3^i) \\
&\leq n^{(\log_3 5)} + \sum_{i=0}^k (5/3)^i n \log n \\
&\leq n^{(\log_3 5)} + (\log_3 n) n^{(\log_3 5)} n (\log n)^2 \\
&= O(n^{(\log_3 5)})
\end{aligned}$$

$$T(n) = T(n-1) + n^2, \quad T(n-1) = T(n-2) + (n-1)^2 + n^2$$

$$T(n-2) = T(n-3) + (n-2)^2 + (n-1)^2 + n^2$$

We observe that:

$$T(n) = T(n-k) + (n-k+1)^2 + (n-k+2)^2 + (n-k+3)^2 + \dots + (n-1)^2 + n^2$$

Now assuming $k = n-1$, we have:

$$T(n) = T(1) + 1^2 + 2^2 + \dots + (n-1)^2 + n^2$$

$$T(n) = 1 + 1^2 + 2^2 + \dots + n^2 \quad (\text{using } T(1) = 1)$$

$$T(n) = 1 + (n(n+1)(2n+1)) / 6 \quad (\text{sum of squares})$$

$$O(n) = n^3$$

b) [44, 937, 13, 69, 37, 80, 472, 49, 300, 183]

Merge sort

And refers to the fact that the elements were divided in the previous step.

[44, 937, 13, 69, 37, 80, 472, 49, 300, 183] Divide(#1)

[44, 937, 13, 69, 37] and [80, 472, 49, 300, 183] D(#2)

[44, 937, 13] and [69, 37] || [80, 472, 49] and [300, 183] D(#3)

[44, 937] and [13] || [69] and [37] || [80, 472] and [49] || [300] and [183] D(#4)

[44] and [937] || [13] || [69] || [37] || [80] and [472] || [49] || [300] || [183] D(#5)

[44, 937]||[13] || [37,69] || [80, 472] || [49] || [183, 300] Compare and Merge (#1)
 [13, 44, 937] || [37,69] || [49, 80, 472] || [183, 300] C&M(#2)
 [13,37, 44, 69, 937] || [49, 80, 83, 300, 472] C&M(#3)
 [13, 37, 44, 49, 69, 80, 83, 300, 472, 937] C&M(#4)

Insertion sort

(||) divides the sorted from unsorted, ({}) the current element.

All shifts were performed to the left and were shown as 1 only

[44||, {937}, 13, 69,37, 80, 472, 49, 300, 183] do nothing
 [44, 937||, {13}, 69,37, 80, 472, 49, 300, 183] copy 13 and shift 2 steps
 [13, 44, 937||, {69},37, 80, 472, 49, 300, 183] copy 69 and shift one step
 [13, 44, 69, 937||, {37}, 80, 472, 49, 300, 183] copy 37 and shift 3 steps
 [13, 37, 44, 69, 937||, {80}, 472, 49, 300, 183] copy 80 and shift 1 steps
 [13, 37, 44, 69, 80, 937||, {472}, 49, 300, 183] copy 472 and shift 1 steps
 [13, 37, 44, 69, 80, 472, 937||, {49}, 300, 183] copy 49 and shift 5 steps
 [13, 37, 44, 49,69, 80, 472, 937||, {300}, 183] copy 300 and shift 2 steps
 [13, 37, 44, 49,69, 80, 300, 472, 937||, {183}] copy 183 and shift 4 steps
 [13, 37, 44, 49,69, 80, 183, 300, 472, 937,||]

c)

Worst case happens when the pivot is less the all the other elements resulting in recursive calls on $n-1$, where n is the number of elements in the previous step.

We know the $T(0) = T(1) = 0$

$T(n) = n + T(n-1)$ Solving the relation:

$$T(n) = n + T(n-1)$$

$$T(n-1) = n + (n-1) + T(n-2)$$

$$T(n-2) = n + (n-1) + (n-2) + T(n-3)$$

$$T(n) = n + (n-1) + \dots + 3 + 2 + 1 \text{ (partial sum of the series is } n(n+1))$$

$$T(n) = n^2 + n$$

$$O(n) = n^2$$

Q2)

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File	Edit	View	Search	Terminal	Help				
1	17	43	57	58	92	93	99	100	
1	17	43	57	58	92	93	99	100	
1	17	43	57	58	92	93	99	100	

Part a - Time analysis of Quick Sort

Array Size	Time Elapsed (ms)	compCount	moveCount
2000	1.344	13927	47141
4000	3.668	25893	88355
6000	5.19	41685	141111
8000	9.544	52453	178551
10000	15.58	76845	257143
12000	22.44	96399	321057
14000	32.564	117679	390561
16000	41.888	131531	437265
18000	52.902	144345	481127
20000	70.4	187006	614290

Part b - Time analysis of Insertion Sort

Array Size	Time Elapsed (ms)	compCount	moveCount
2000	9.582	1015628	1019628
4000	62.876	3976005	3984005
6000	232.278	8961802	8973802
8000	522.104	16107587	16123587
10000	1029.45	24962621	24982621
12000	1751.63	35821519	35845519
14000	2828.13	48810292	48838292
16000	4146.18	63852925	63884925
18000	5860.64	81365184	81401184
20000	8121.6	100295387	100335387

Part c - Time analysis of Hybrid Sort

Array Size	Time Elapsed (ms)	compCount	moveCount
2000	0.532	14686	44986
4000	2.136	27402	84582
6000	5.016	44077	135145
8000	9.248	55601	171189
10000	14.58	80897	247607
12000	21.792	101295	310037
14000	30.184	123167	376995
16000	40.336	137885	421957
18000	52.398	151348	464666
20000	67.06	194829	596107

Q3)

performance of sorting algorithms

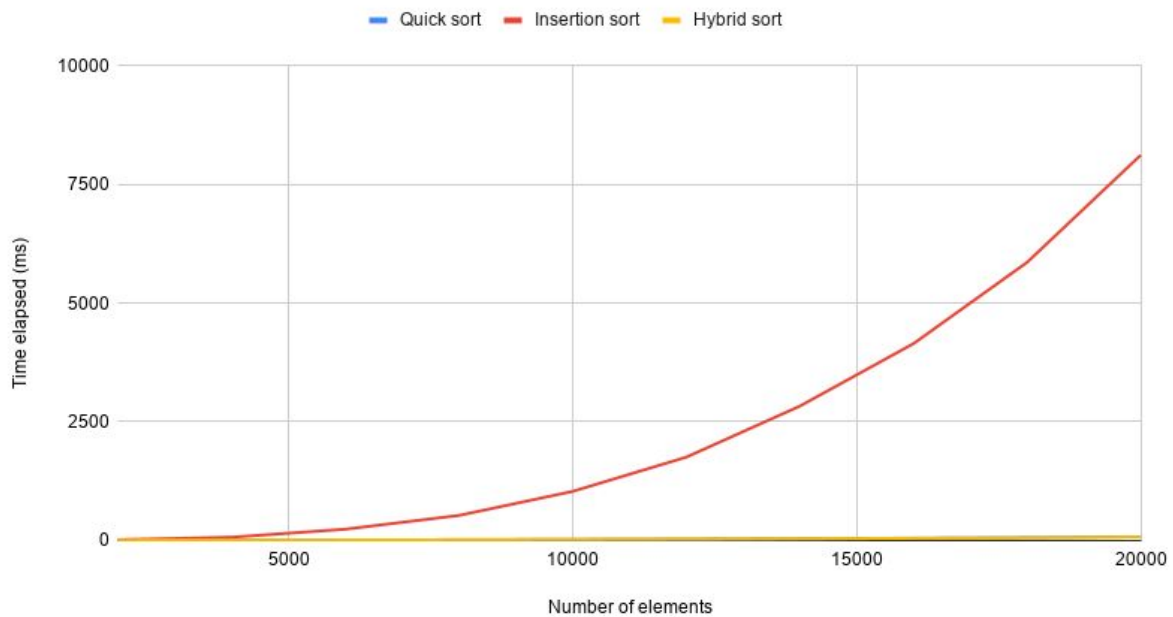


Figure 1: Sorting algorithms performance.

performance of sorting algorithms

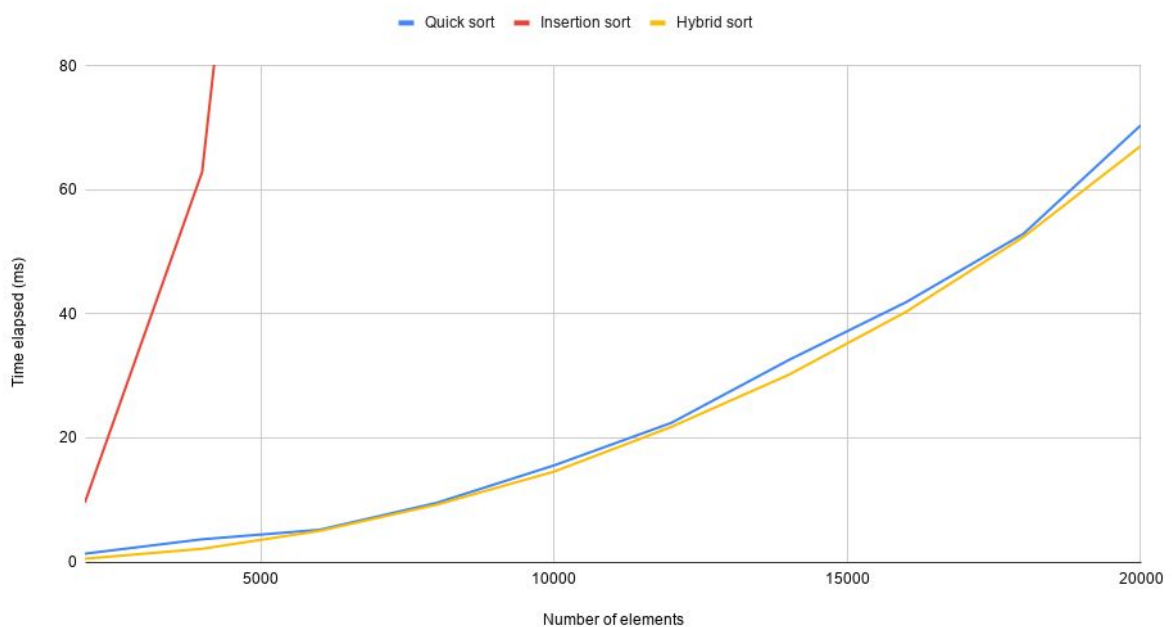


Figure 2: Close-up of the Sorting algorithms performance graph.

As we see in the graphs above Insertion Sort has significantly higher time compared to Quick sort and Hybrid Sort (figure 1). This matches the theory since worst-case for Insertion sort is n^2 and so is its average case. While for Quick Sort it is $O(n \log n)$ for average and n^2 for worst case scenario. We can

realize that the worst-case does not always happen and that's why we had a huge difference between the two algorithms (difference between 20000^2 and $20000 * \log 20000$ is enormous), Since Hybrid is Quick sort for most of the process it is expected that the time is very close to Quick sort (figure 2). Hybrid sort, has the good of both, since the best case scenario for Insertion is $O(n)$ and $O(n \log n)$ for Quick sort, using Insertion for very small arrays is more efficient in terms of time, since the number of comparisons is small, resulting in less time in Insertion compared to Quick sort where we need to partition first then recursively call on itself.