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RECAP

Transport Layer

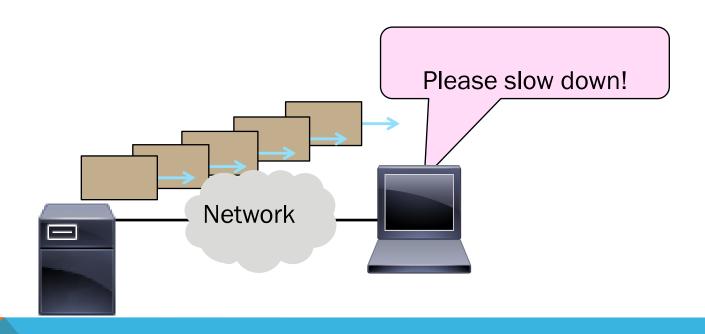
- End-to-End Communication
- Ports
- Protocols
 - UDP
 - TCP
 - Reliabilty
- Improving Performance
 - Sliding Window

TODAY'S PLAN

- Improving Performance
 - Flow Control
- Application Layer
 - DNS

Adding flow control to the sliding window algorithm

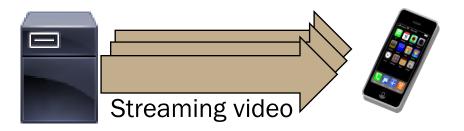
To slow the over-enthusiastic sender



PROBLEM

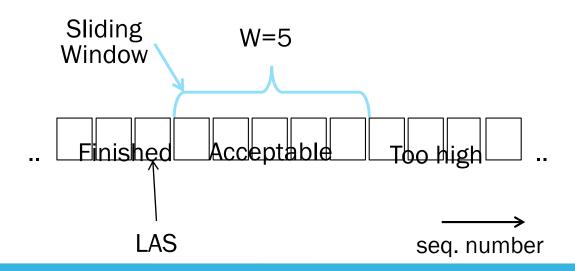
Sliding window uses pipelining to keep the network busy

What if the receiver is overloaded?

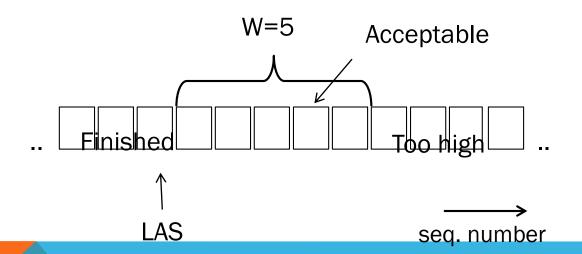


Consider receiver with W buffers

 LAS=LAST ACK SENT, app pulls in-order data from buffer with recv() call

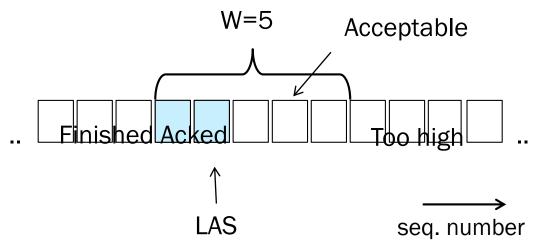


Suppose the next two segments arrive but app does not call recv()



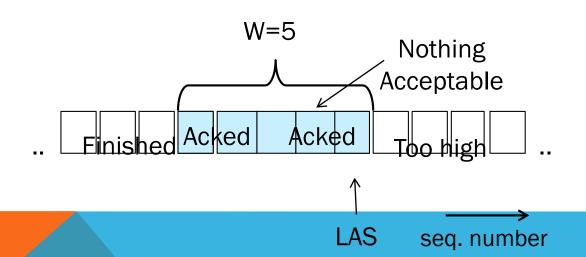
Suppose the next two segments arrive but app does not call recv()

LAS rises, but we can't slide window!



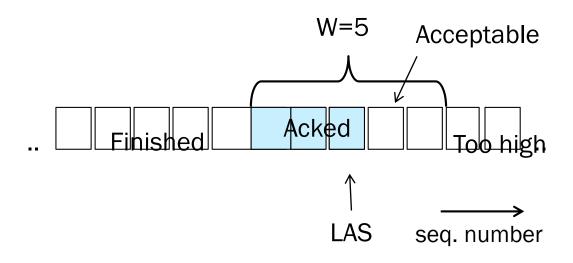
If further segments arrive (even in order) we can fill the buffer

• Must drop segments until app recvs!



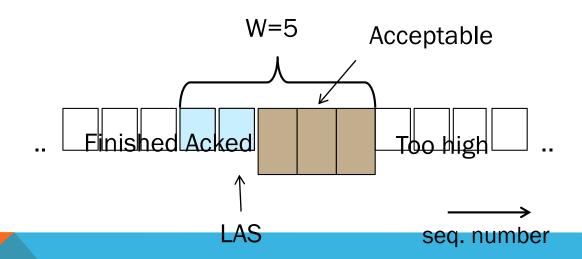
App recv() takes two segments

Window slides

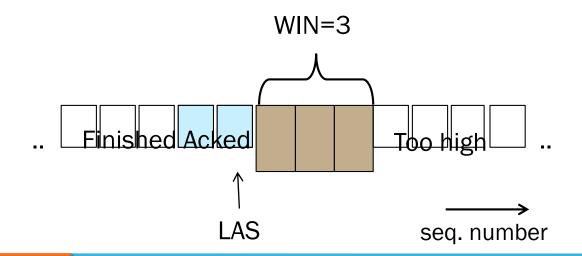


Avoid loss at receiver by telling sender the available buffer space

• WIN=#Acceptable, not W (from LAS)

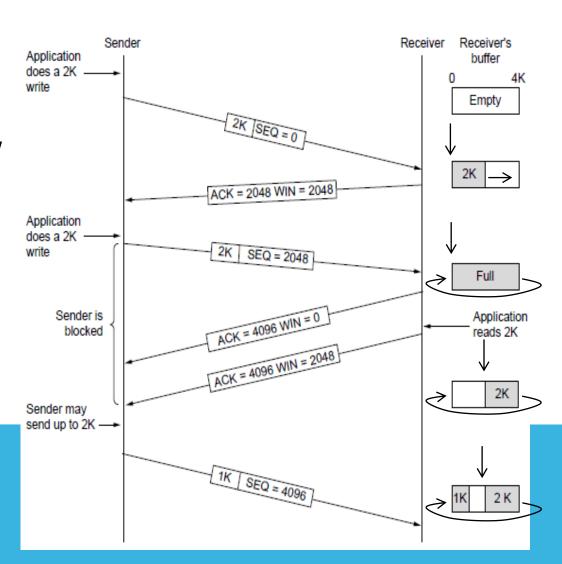


Sender uses the lower of the sliding window and <u>flow</u> <u>control window</u> (WIN) as the effective window size



TCP-style example

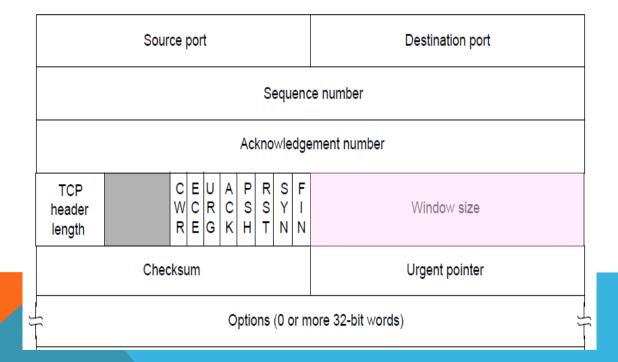
- SEQ/ACK sliding window
- Flow control with win
- SEQ + length < ACK+WIN</p>
- 4KB buffer at receiver
- Circular buffer of bytes

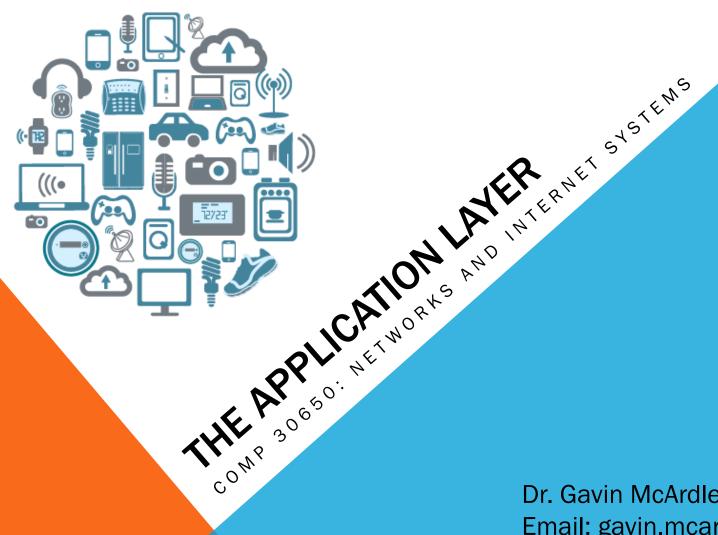


TCP HEADER

Window size for flow control

Relative to ACK, and in bytes





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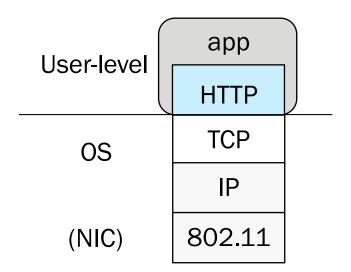
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RECALL

Application layer protocols are often part of an "app"

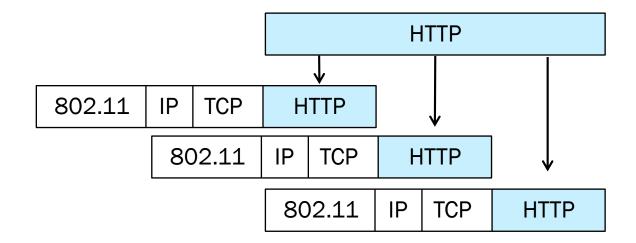
- But don't need a GUI
 - DNS
- Some have GUIs
 - Web browser



RECALL

Application layer messages are often split over multiple packets

Or may be aggregated in a packet ...



APPLICATION COMMUNICATION NEEDS Vary widely with app; must build on Transport services

Web

Series of variable length, reliable request/reply exchanges

TCP

DNS

Short, reliable request/reply exchanges

UDP

Skype

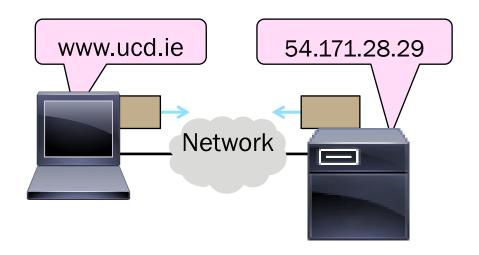
Real-time (unreliable) delivery

UDP

FILLING IN THE GAPS OF OUR BIG PICTURE

The DNS (Domain Name System)

Human-readable host names



NAMES AND ADDRESSES

Names are higher-level <u>identifiers</u> for resources <u>Addresses</u> are lower-level <u>locators</u> for resources

Multiple levels, e.g. full name → email → IP address →
 Ethernet address

Resolution is mapping a name to an address

Lookup in an address book or directory or some sort of Table!

BEFORE THE DNS - HOSTS.TXT

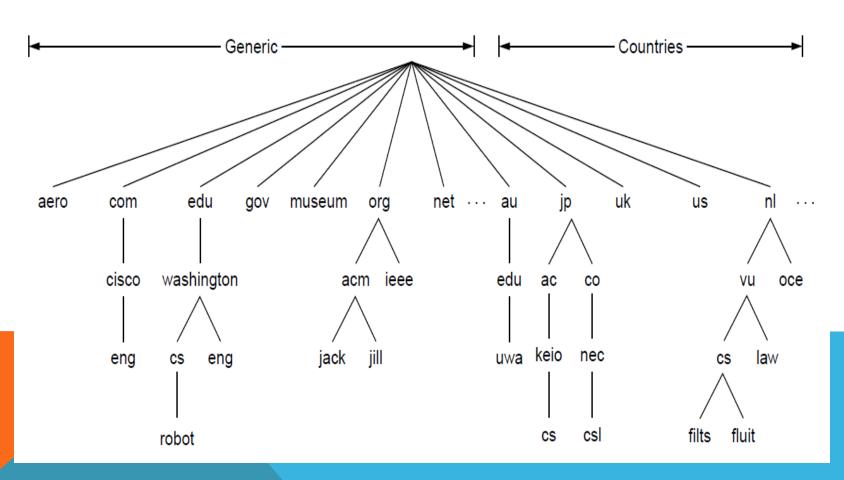
- Directory was a file HOSTS.TXT regularly retrieved for all hosts from a central machine at the NIC (Network Information Center)
- Names were initially flat, became hierarchical (e.g., lcs.mit.edu)
- Neither manageable nor efficient as the ARPANET grew ...

DNS

- A naming service to map between host names and their IP addresses
 - <u>www.ucd.ie</u> → 54.171.28.29
- Goals:
- Easy to manage (esp. with multiple parties)
- Efficient (good performance, few resources)
- Approach:
- Distributed directory based on a <u>hierarchical</u> namespace
- Automated protocol to tie pieces together

DNS NAMESPACE

Hierarchical, starting from "." (dot, typically omitted)



TLDS (TOP-LEVEL DOMAINS)

Run by ICANN (Internet Corp. for Assigned Names and Numbers)

22+ generic Top-Level Domains

Initially .com, .edu , .gov., .mil, .org, .net

~250 country code TLDs

- Two letters, e.g., ".au", plus international characters since 2010
- Widely commercialized, e.g., .tv (Tuvalu)
- Many domain hacks, e.g., instagr.am (Armenia), goo.gl (Greenland)

Other domains

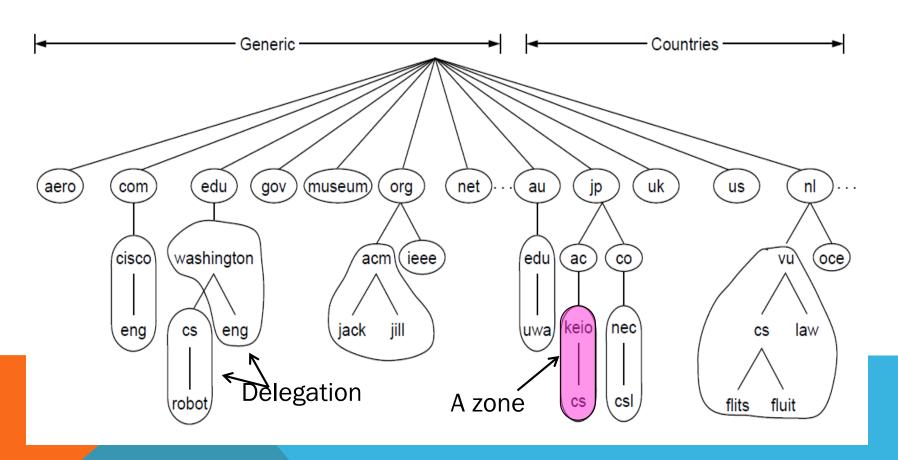
.academy to .zone and everything between.

Brand domains

google, amex, omega, etc.

DNS ZONES

A zone is a contiguous portion of the namespace



DNS ZONES

Zones are the basis for distribution

- EDU Registrar administers .edu
- UW administers washington.edu
- CS administers cs.washington.edu
 - cs.uw.edu

Each zone has a <u>nameserver</u> to contact for information about it

 Zone must include contacts for delegations, e.g., .edu knows nameserver for washington.edu

DNS RESOURCE RECORDS

A zone is comprised of DNS resource records that give information for its domain names

Туре	Meaning		
SOA	Start of authority, has key zone		
	parameters		
A	IPv4 address of a host		
AAAA	IPv6 address of a host		
("quad A")	IPVO address of a flost		
CNAME	Canonical name for an alias		
MX	Mail exchanger for the domain		
NS	Nameserver of domain or delegated		
	subdomain		

DNS RESOURCE RECORDS

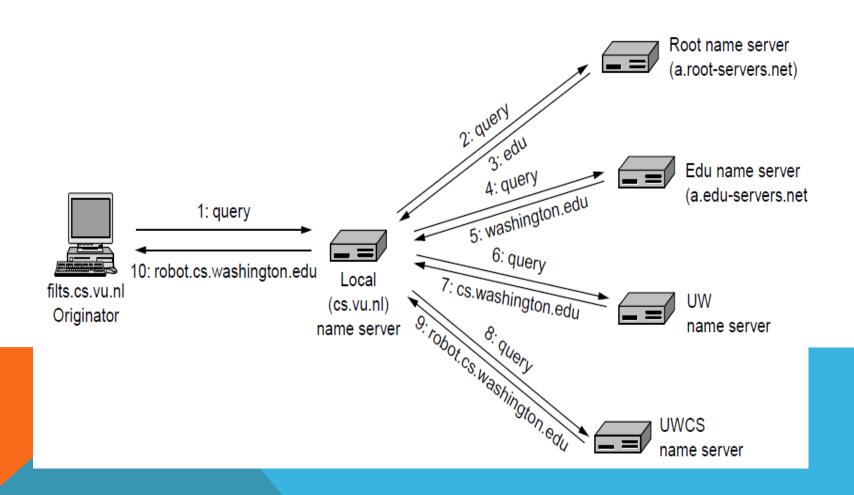
; Authoritative data cs.vu.nl. cs.vu.nl. cs.vu.nl. cs.vu.nl.	86400 86400 86400 86400 86400	u.nl IN IN IN IN	SOA MX MX NS	star boss (9527,7200,7 1 zephyr 2 top star	7200,241920,86400) Name server
star zephyr top www ftp	86400 86400 86400 86400 86400	IN IN IN IN IN	A A A CNAME CNAME	130.37.56.205 130.37.20.10 130.37.20.11 star.cs.vu.nl zephyr.cs.vu.nl	IP addresses of computers
flits flits flits flits	86400 86400 86400 86400 86400	IX I	A A MX MX MX	130.37.16.112 192.31.231.165 1 flits 2 zephyr 3 top	
rowboat		IN IN IN	A MX MX	130.37.56.201 1 rowboat 2 zephyr	Mail gateways
little-sister		IN	Α	130.37.62.23	
laserjet		IN	Α	192.31.231.216	

DNS RESOLUTION

- DNS protocol lets a host resolve any host name (domain) to IP address
- If unknown, can start with the root nameserver and work down zones

DNS RESOLUTION

flits.cs.vu.nl resolves robot.cs.washington.edu



ITERATIVE VS. RECURSIVE DNS QUERIES

Recursive query

- Nameserver completes resolution and returns the final answer
- E.g., flits → local nameserver

Iterative query

- Nameserver returns the answer or who to contact next for the answer
- E.g., local nameserver → all others

ITERATIVE VS. RECURSIVE DNS QUERIES

Recursive query

- Lets server offload client burden (simple resolver) for manageability
- Lets server cache over a pool of clients for better performance

Iterative query

- Lets server "file and forget"
- Easy to build high load servers

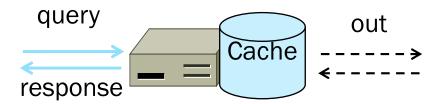
CACHING

Resolution latency should be low

Adds delay to web browsing

Cache query/responses to answer future queries immediately

- Including partial (iterative) answers
- Responses carry a TTL for caching

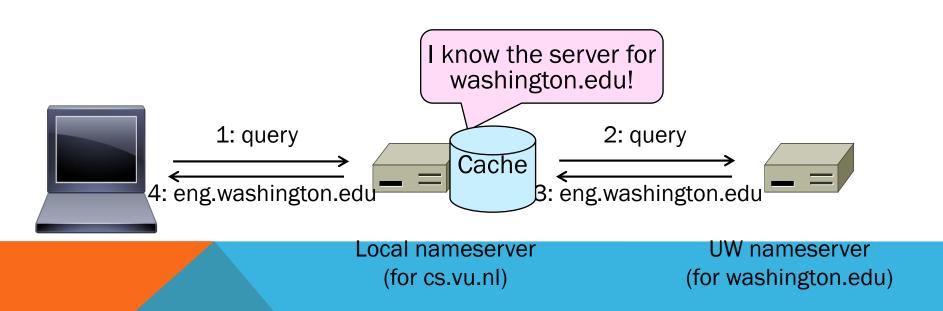


Nameserver

CACHING

flits.cs.vu.nl now resolves eng.washington.edu

And previous resolutions cut out most of the process



LOCAL NAMESERVERS

Local nameservers typically run by IT (enterprise, ISP)

- But may be your host or AP
- Or alternatives e.g., Google public DNS

Clients need to be able to contact their local nameservers

Typically configured via DHCP

ROOT NAMESERVERS

Root (dot) is served by 13 server names

- a.root-servers.net to m.root-servers.net
- All nameservers need root IP addresses
- Handled via configuration file (named.ca)

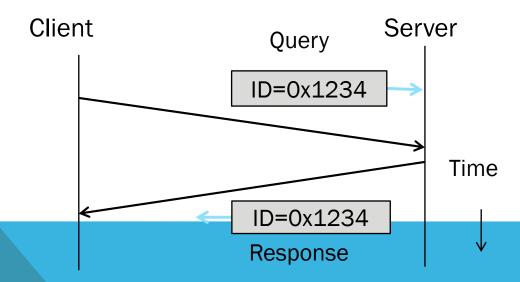
There are >250 distributed server instances

- Highly reachable, reliable service
- Most servers are reached by <u>IP anycast</u> (Multiple locations advertise same IP!)
- Routers take client to the closest one.
- Servers are IPv4 and IPv6 reachable

DNS PROTOCOL

Query and response messages

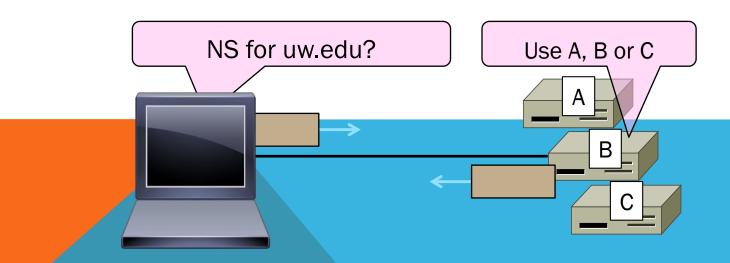
- Built on UDP messages, port 53
- ARQ for reliability; server is stateless!
- Messages linked by a 16-bit ID field



DNS PROTOCOL

Service reliability via replicas

- Run multiple nameservers for domain
- Return the list; clients use one answer
- Helps distribute load too



DNS PROTOCOL

Security is a major issue

- Compromise redirects to wrong site!
- Not part of initial protocols ...

DNSSEC (DNS Security Extensions)