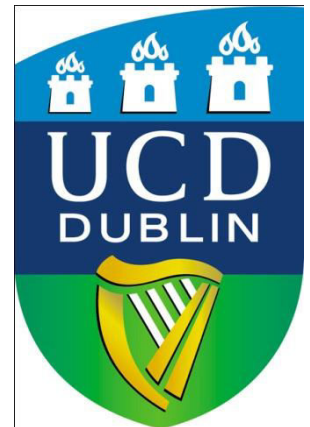


# COM307000 - Software

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# Why Software?

- ❑ Why is software as important to security as crypto, access control, protocols?
- ❑ Virtually all information security features are implemented in software
- ❑ If your software is subject to attack, your security can be broken
  - Regardless of strength of crypto, access control, or protocols
- ❑ Software is a poor *foundation* for security

# **Software Flaws and Malware**

# Bad Software is Ubiquitous

- ❑ NASA Mars Lander (cost \$165 million)
  - Crashed into Mars due to...
  - ...error in converting English and metric units of measure
  - Believe it or not
- ❑ Denver airport
  - Baggage handling system — very buggy software
  - Delayed airport opening by 11 months
  - Cost of delay exceeded \$1 million/day
  - What happened to person responsible for this fiasco?
- ❑ MV-22 Osprey
  - Advanced military aircraft
  - Faulty software can be fatal

# Software Issues

## Alice and Bob

- ❑ Find bugs and flaws by accident
- ❑ Hate bad software...
- ❑ ...but they learn to live with it
- ❑ Must make bad software work

## Trudy

- ❑ Actively looks for bugs and flaws
- ❑ Likes bad software...
- ❑ ...and tries to make it misbehave
- ❑ Attacks systems via bad software

# Complexity

- "Complexity is the enemy of security", Paul Kocher, Cryptography Research, Inc.

System	Lines of Code (LOC)
Netscape	17 million
Space Shuttle	10 million
Linux kernel 2.6.0	5 million
Windows XP	40 million
Mac OS X 10.4	86 million
Boeing 777	7 million

- A new car contains more LOC than was required to land the Apollo astronauts on the moon

# Lines of Code and Bugs

- ❑ Conservative estimate: 5 bugs/10,000 LOC
- ❑ **Do the math**
  - Typical computer: 3k exe's of 100k LOC each
  - Conservative estimate: 50 bugs/exe
  - Implies about 150k bugs per computer
  - So, 30,000-node network has 4.5 billion bugs
  - Maybe only 10% of bugs security-critical and only 10% of those remotely exploitable
  - Then "only" 45 million critical security flaws!

# Software Security Topics

- ❑ Program flaws (unintentional)
  - Buffer overflow
  - Incomplete mediation
  - Race conditions
- ❑ Malicious software (intentional)
  - Viruses
  - Worms
  - Other breeds of malware



# Program Flaws

- ❑ An **error** is a programming mistake
  - To err is human
- ❑ An error may lead to incorrect state: **fault**
  - A fault is internal to the program
- ❑ A fault may lead to a **failure**, where a system departs from its expected behavior
  - A failure is externally observable

**error** → **fault** → **failure**

# Example

```
char array[10];  
for(i = 0; i < 10; ++i)  
    array[i] = `A`;   
array[10] = `B`;
```

- ❑ This program has an **error**
- ❑ This error might cause a **fault**
  - Incorrect internal state
- ❑ If a fault occurs, it might lead to a **failure**
  - Program behaves incorrectly (external)
- ❑ We use the term **flaw** for all of the above

# Secure Software

- ❑ In software engineering, try to ensure that a program does what is intended
- ❑ *Secure* software engineering requires that software **does what is intended...**
- ❑ **...and nothing more**
- ❑ Absolutely secure software? Dream on...
  - But, absolute security *anywhere* is impossible
- ❑ How can we manage software risks?

# Program Flaws

- ❑ Program flaws are **unintentional**
  - But can still create security risks
- ❑ We'll consider 3 types of flaws
  - Buffer overflow (smashing the stack)
  - Incomplete mediation
  - Race conditions
- ❑ These are the most common flaws

# Buffer Overflow



# Attack Scenario

- ❑ Users enter data into a Web form
- ❑ Web form is sent to server
- ❑ Server writes data to array called buffer, without checking length of input data
- ❑ Data “overflows” buffer
  - Such overflow might enable an attack
  - If so, attack could be carried out by anyone with Internet access

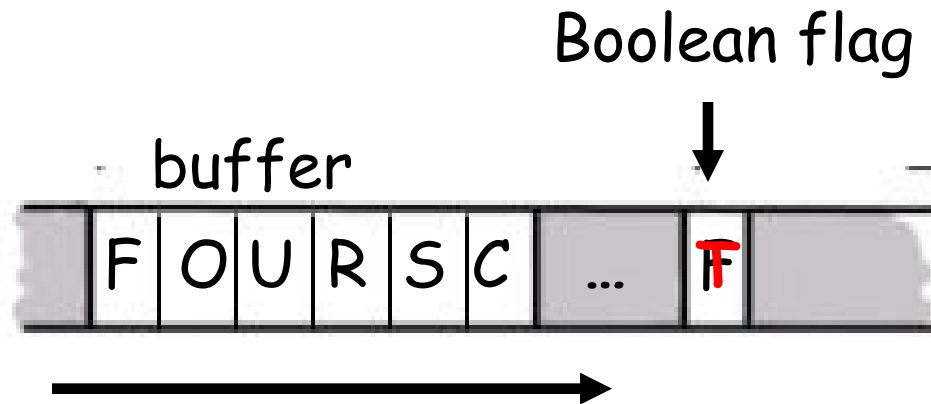
# Buffer Overflow

```
int main() {  
    int buffer[10];  
    buffer[20] = 37; }
```

- ❑ **Q:** What happens when code is executed?
- ❑ **A:** Depending on what resides in memory at location “buffer[20]”
  - Might overwrite **user** data or code
  - Might overwrite **system** data or code
  - Or program could work just fine

# Simple Buffer Overflow

- ❑ Consider boolean flag for authentication
- ❑ Buffer overflow could overwrite flag allowing anyone to authenticate

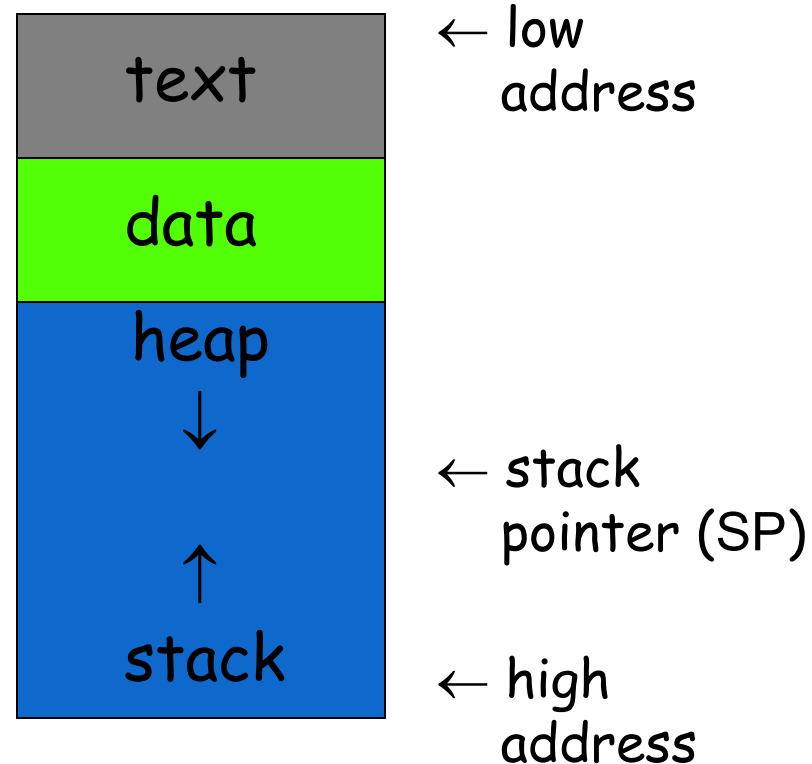


- ❑ In some cases, Trudy need not be so lucky as in this example



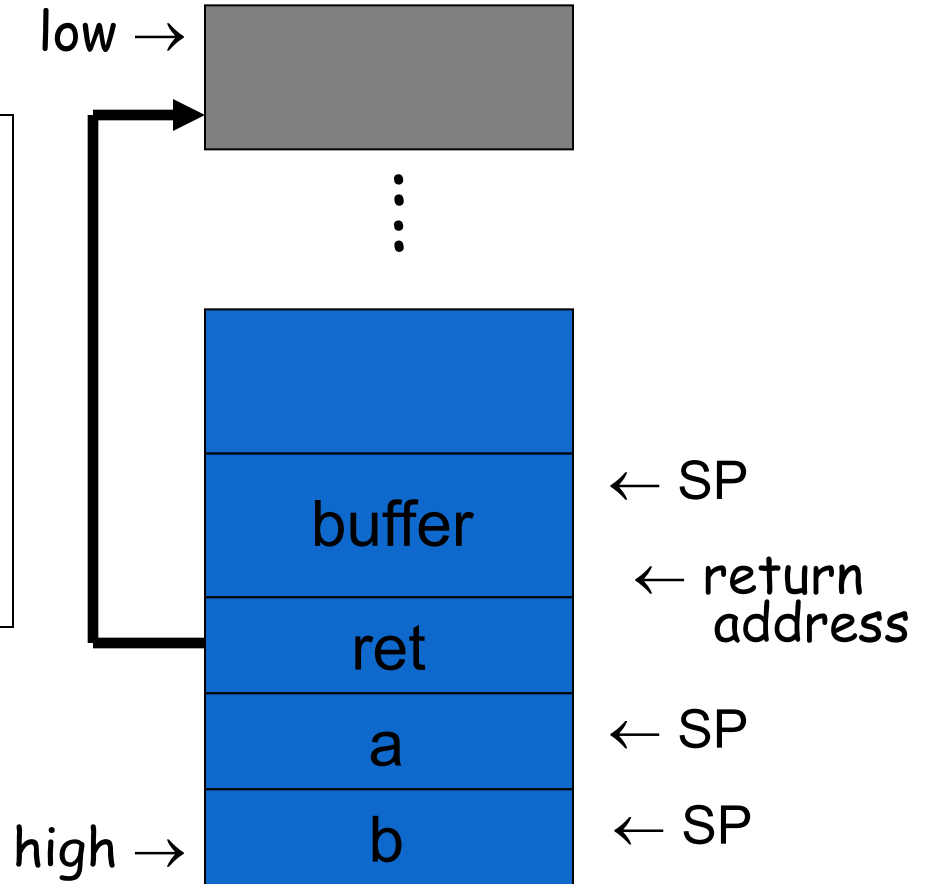
# Memory Organization

- ❑ **Text** — code
- ❑ **Data** — static variables
- ❑ **Heap** — dynamic data
- ❑ **Stack** — “scratch paper”
  - Dynamic local variables
  - Parameters to functions
  - Return address



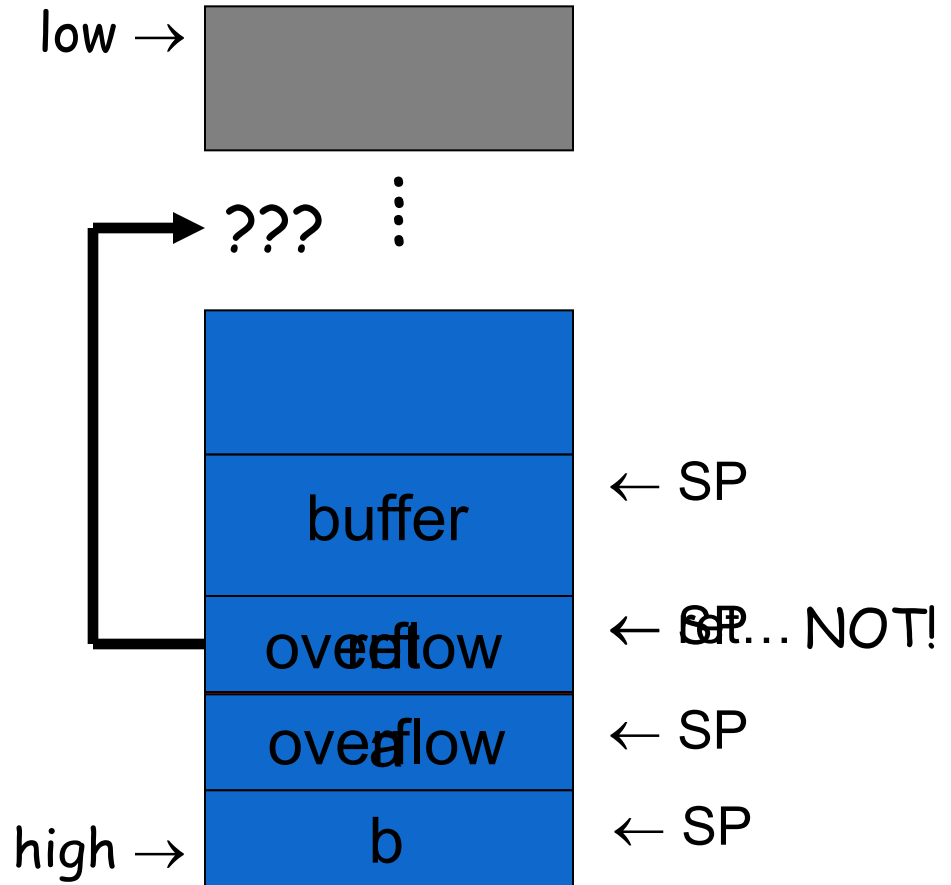
# Simplified Stack Example

```
void func(int a, int b){  
    char buffer[10];  
}  
void main(){  
    func(1,2);  
}
```



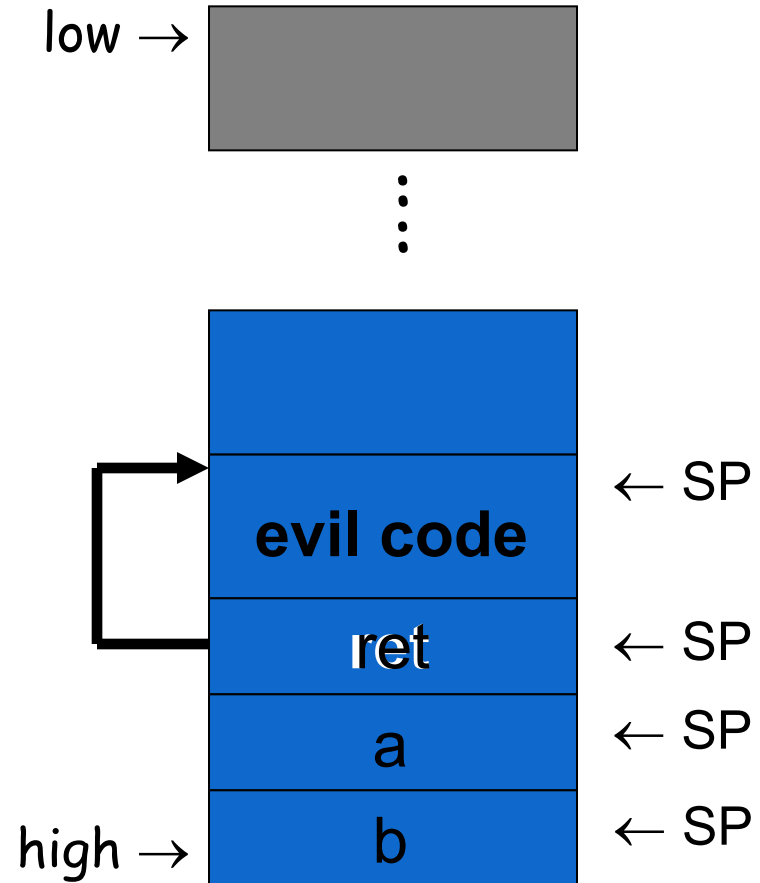
# Smashing the Stack

- ❑ What happens if buffer overflows?
- ❑ Program “returns” to wrong location
- ❑ A crash is likely



# Smashing the Stack

- ❑ Trudy has a better idea...
- ❑ **Code injection**
- ❑ Trudy can run code of her choosing...
  - ...on your machine



# Smashing the Stack

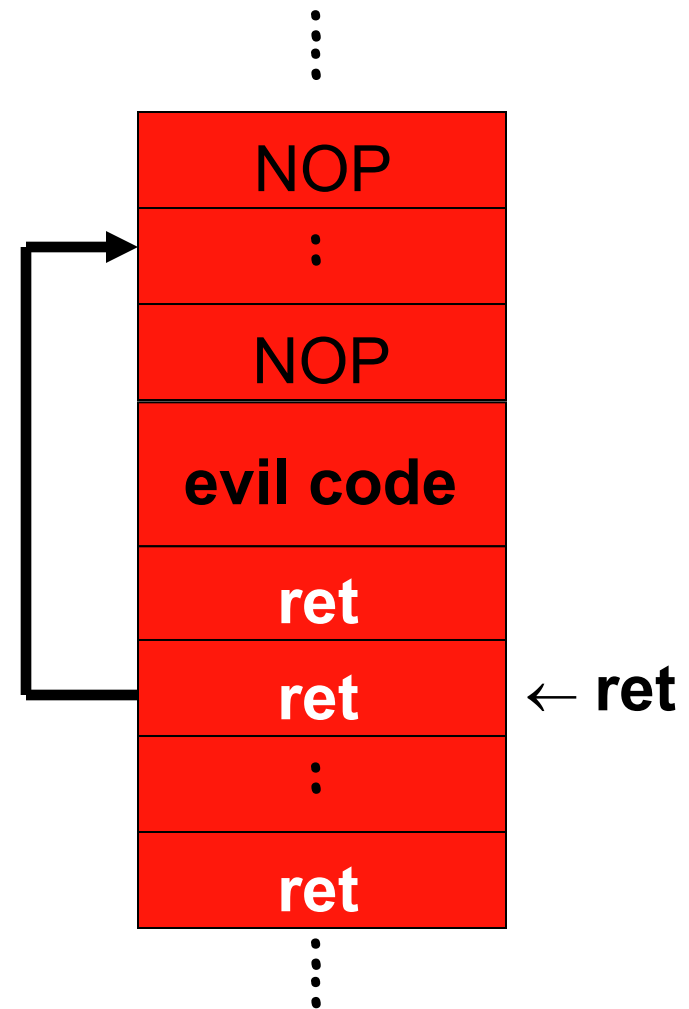
❑ Trudy may not know...

- 1) Address of evil code
- 2) Location of **ret** on stack

❑ Solutions

- 1) Precede evil code with NOP "landing pad"\*
- 2) Insert **ret** many times

\* NOP aka "no-op slide". A landing pad is just a long sequence of nop instructions, so that no matter where eip lands in that string, it will progress to the first "real" instruction.



# Stack Smashing Summary

- ❑ A buffer overflow must exist in the code
- ❑ Not all buffer overflows are exploitable
  - Things must align properly
- ❑ If exploitable, attacker can **inject code**
- ❑ Trial and error is likely required
  - Fear not, lots of help is available online
  - [Smashing the Stack for Fun and Profit](#), Aleph One
- ❑ Stack smashing is “attack of the decade” ...
  - ...for many recent decades
  - Also heap, integer overflows, format strings, etc.

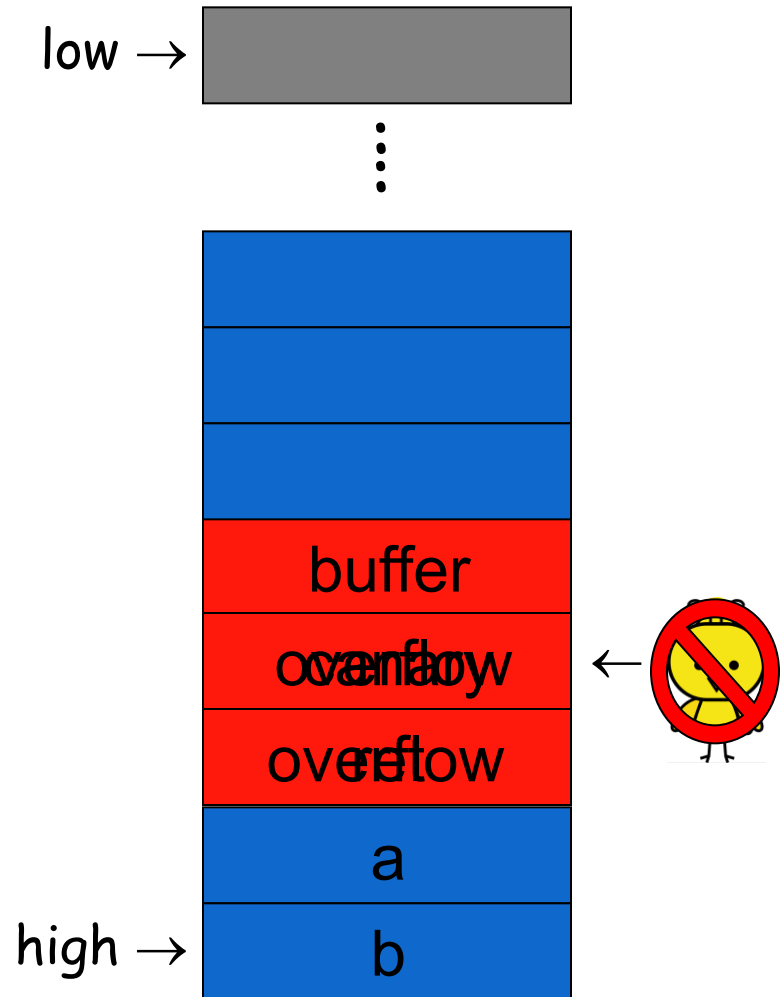
# Stack Smashing Defenses

- ❑ Employ **non-executable stack**
  - “No execute” **NX bit** (if available)
  - Seems like the logical thing to do, but some real code executes on the stack (Java, for example)
- ❑ Use a **canary**
- ❑ Address space layout randomization (**ASLR**)
- ❑ Use **safe languages** (Java, C#)
- ❑ Use **safer C functions**
  - For unsafe functions, safer versions exist
  - For example, strncpy instead of strcpy

# Stack Smashing Defenses

## □ Canary

- Run-time stack check
- Push canary onto stack
- Canary value:
  - Constant 0x000aff0d
  - Or, may depends on **ret**





# Microsoft's Canary

- ❑ Microsoft added **buffer security check** feature to C++ with /GS compiler flag
  - Based on canary (or “security cookie”)

**Q:** What to do when canary dies?

**A:** Check for user-supplied “handler”

- ❑ Handler shown to be subject to attack
  - Claimed that attacker can specify handler code
  - If so, formerly “safe” buffer overflows become exploitable when /GS is used!

# ASLR

- ❑ Address Space Layout Randomization
  - Randomize place where code loaded in memory
- ❑ Makes most buffer overflow attacks probabilistic
- ❑ Windows Vista uses 256 random layouts
  - So about 1/256 chance buffer overflow works
- ❑ Similar thing in Mac OS X and other OSs
- ❑ Attacks against Microsoft's ASLR do exist
  - Possible to “de-randomize”

# Buffer Overflow

- ❑ A major security threat yesterday, today, and tomorrow
- ❑ The good news?
  - It is possible to reduce overflow attacks (safe languages, NX bit, ASLR, education, etc.)
- ❑ The bad news?
  - Buffer overflows will exist for a long time
  - Why? Legacy code, bad development practices, clever attacks, etc.

# Incomplete Mediation



# Input Validation

- ❑ Consider: `strcpy(buffer, argv[1])`
- ❑ A buffer overflow occurs if  
 $\text{len}(\text{buffer}) < \text{len}(\text{argv}[1])$
- ❑ Software must **validate** the input by checking the length of `argv[1]`
- ❑ Failure to do so is an example of a more general problem: **incomplete mediation**

# Input Validation

- ❑ Consider web form data
- ❑ Suppose input is validated on client
- ❑ For example, the following is valid  
`http://www.things.com/orders/final&custID=112&num=55A  
&qty=20&price=10&shipping=5&total=205`
- ❑ Suppose input is not checked on server
  - Why bother since input checked on client?
  - Then attacker could send http message  
`http://www.things.com/orders/final&custID=112&num=55A  
&qty=20&price=10&shipping=5&total=25`

# Incomplete Mediation

- ❑ Linux kernel
  - Research revealed many buffer overflows
  - Lots of these due to incomplete mediation
- ❑ Linux kernel is “good” software since
  - Open-source
  - Kernel — written by coding gurus
- ❑ Tools exist to help find such problems
  - But incomplete mediation errors can be subtle
  - And tools useful for attackers too!

# Race Conditions



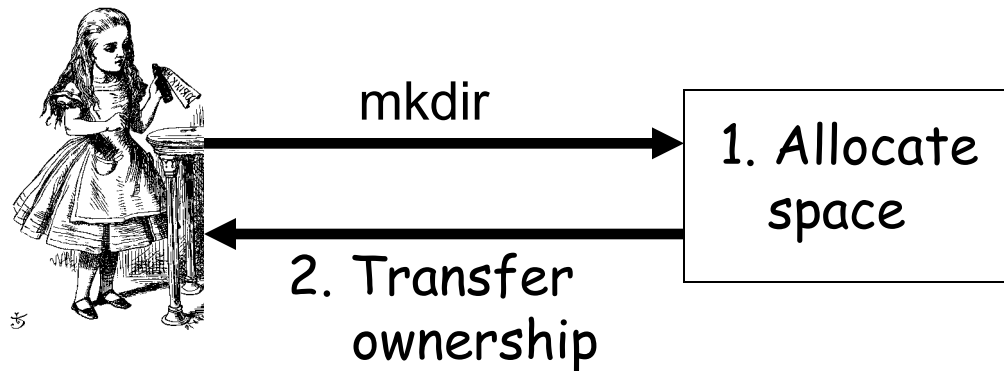


# Race Condition

- ❑ Security processes should be **atomic**
  - Occur “all at once”
- ❑ Race conditions can arise when security-critical process occurs in stages
- ❑ Attacker makes change between stages
  - Often, between stage that gives authorization, but before stage that transfers ownership
- ❑ Example: Unix mkdir

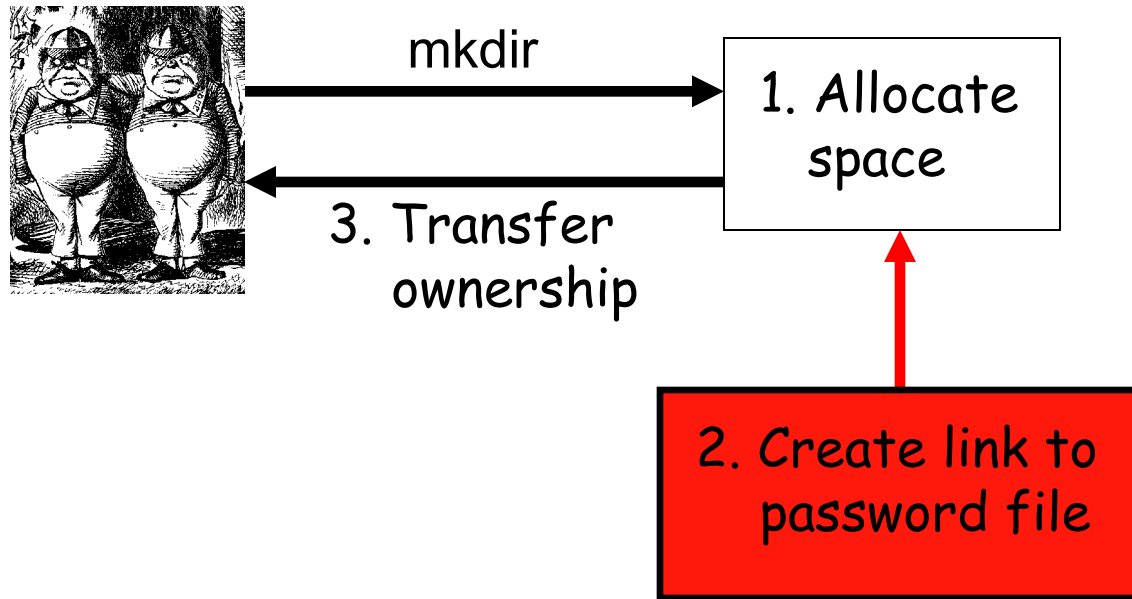
# mkdir Race Condition

- ❑ mkdir creates new directory
- ❑ How mkdir is supposed to work



# mkdir Attack

## ❑ The mkdir race condition



- ❑ Not really a “race”
  - But attacker’s timing is critical

# Race Conditions

- ❑ Race conditions are common
- ❑ Race conditions may be more prevalent than buffer overflows
- ❑ But race conditions harder to exploit
  - Buffer overflow is “low hanging fruit” today
- ❑ To prevent race conditions, make security-critical processes atomic
  - Occur all at once, not in stages
  - Not always easy to accomplish in practice

**Next...Malware**