

# Sorting Algorithms I



Mark Matthews PhD

# Summary

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- Why sorting is important
- Sorting in java
- Selection Sort
- Types of sorting algorithms
- Insertion Sort (maybe next time)



# Why is sorting important?

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- Sorted data powers almost everything
- More efficient to work with a sorted dataset
- More human-friendly also
- Enables much more efficient search algorithms (which we'll see soon!)

# Why is sorting important?

---

Find '1'.

Life without sorting

[38, 2, 8, 48, 27, 43, 23, 3, 17, 33, 1]

Worse case scenario we have to look through the whole list

Time complexity:  $O(N)$  or linear



# What do we mean by sorting?

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Organising a set of similar items in a collection by some property

Items need to be homogenous  
e.g., of the same type

Can be in increasing or decreasing order  
e.g., 1,2,3,4 or 4,3,2,1

Sort by shared property  
e.g., lexicographical (alphabetical),  
numerical order, date etc.



## Sorting problem

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Ex. Student records in a university.

item	→	Chen	3	A	991-878-4944	308 Blair
Rohde		2	A	232-343-5555	343 Forbes	
Gazsi		4	B	766-093-9873	101 Brown	
Furia		1	A	766-093-9873	101 Brown	
Kanaga		3	B	898-122-9643	22 Brown	
Andrews		3	A	664-480-0023	097 Little	
Battle	→	4	C	874-088-1212	121 Whitman	

Sort. Rearrange array of  $N$  items into ascending order.

Andrews	3	A	664-480-0023	097 Little
Battle	4	C	874-088-1212	121 Whitman
Chen	3	A	991-878-4944	308 Blair
Furia	1	A	766-093-9873	101 Brown
Gazsi	4	B	766-093-9873	101 Brown
Kanaga	3	B	898-122-9643	22 Brown
Rohde	2	A	232-343-5555	343 Forbes

# Sorting applications

KAYAK  
Part of Booking.com

Flights Hotels Cars Holidays More

Return 1 Adult Economy 0 bags

Dublin (DUB) → Honolulu (HNL)

Fri 14/2 Sun 16/2

**Cheapest** €1425 • 23h 24m

**Best** ⓘ

**Quickest** €6391 • 20h 15m

Other sort

Track prices OFF

7 of 527 flights

**Fare assistant** ⓘ

- Cabin bag 0 +
- Checked bag 0 +
- Payment method Visa Debit

**Stops**

- Direct
- 1 stop
- 2+ stops €1425

**Times**

Take-off Landing

Take-off from DUB Fri 05:00 – 21:00

Take-off from HNL Sun 07:00 – 23:59

**Airlines**

Select all | Clear all

- Aer Lingus €1340
- American Airlines €1340
- British Airways €1302

**Fly to Honolulu with Aer Lingus**

Save time on your trip by pre-clearing U.S. Immigration in Dublin Airport. Book now at aerlingus.com

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**View Deal**

**Rating: 7** 0 1 2 3 4 5 6 7 8 9 10

**€1425**  
American Airlines  
Main Cabin

**View Deal**

**Rating: 2** 0 1 2 3 4 5 6 7 8 9 10

**€2045**  
United Airlines  
Premium Economy

**View Deal**

**Rating: 2** 0 1 2 3 4 5 6 7 8 9 10

**€2133**  
Delta

**View Deal**

**Rating: 1** 0 1 2 3 4 5 6 7 8 9 10

**€2149**  
KLM

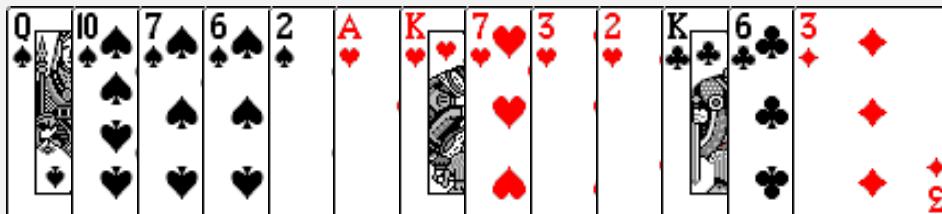
**View Deal**

# Sorting applications

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Postal packages



playing cards

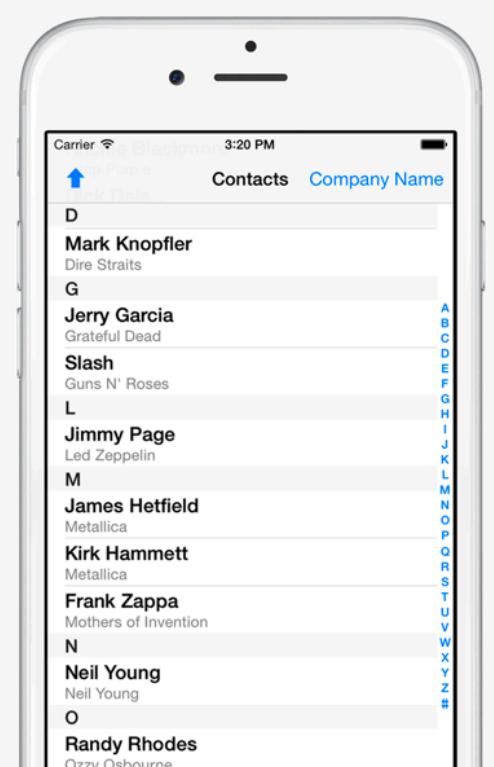


Databases, large datasets



Hogwarts houses

Contacts



# Sorting in Java



## Sample sort client 1

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Goal. Sort **any** type of data.

Ex 1. Sort random real numbers in ascending order.

seems artificial (stay tuned for an application)

```
public class Experiment
{
    public static void main(String[] args)
    {
        int N = Integer.parseInt(args[0]);
        Double[] a = new Double[N];
        for (int i = 0; i < N; i++)
            a[i] = StdRandom.uniform();
        Insertion.sort(a);
        for (int i = 0; i < N; i++)
            StdOut.println(a[i]);
    }
}
```

```
% java Experiment 10
0.08614716385210452
0.09054270895414829
0.10708746304898642
0.21166190071646818
0.363292849257276
0.460954145685913
0.5340026311350087
0.7216129793703496
0.9003500354411443
0.9293994908845686
```

## Sample sort client 2

---

Goal. Sort **any** type of data.

Ex 2. Sort strings in alphabetical order.

```
public class StringSorter
{
    public static void main(String[] args)
    {
        String[] a = StdIn.readAllStrings();
        Insertion.sort(a);
        for (int i = 0; i < a.length; i++)
            StdOut.println(a[i]);
    }
}
```

% more words3.txt  
bed bug dad yet zoo ... all bad yes

% java StringSorter < words3.txt  
all bad bed bug dad ... yes yet zoo  
[suppressing newlines]

## Sample sort client 3

---

Goal. Sort **any** type of data.

Ex 3. Sort the files in a given directory by filename.

```
import java.io.File;

public class FileSorter
{
    public static void main(String[] args)
    {
        File directory = new File(args[0]);
        File[] files = directory.listFiles();
        Insertion.sort(files);

        for (int i = 0; i < files.length; i++)
            StdOut.println(files[i].getName());
    }
}
```

```
% java FileSorter .
Insertion.class
Insertion.java
InsertionX.class
InsertionX.java
Selection.class
Selection.java
Shell.class
Shell.java
ShellX.class
ShellX.java
```

# Sorting in Java

---

Goal. Sort **any** type of data (for which sorting is well defined).

Q. How can `sort()` know how to compare data of type `Double`, `String`, and `java.io.File` without any information about the type of an item's key?

Callback = reference to executable code.

- Client passes array of objects to `sort()` function.
- The `sort()` function calls object's `compareTo()` method as needed.

Implementing callbacks.

- Java: interfaces.
- C: function pointers.
- C++: class-type functors.
- C#: delegates.
- Python, Perl, ML, Javascript: first-class functions.



# Callbacks: roadmap

## client

```
public class StringSorter
{
    public static void main(String[] args)
    {
        String[] a = StdIn.readAllStrings();
        Insertion.sort(a);
        for (int i = 0; i < a.length; i++)
            StdOut.println(a[i]);
    }
}
```

## Comparable interface (built in to Java)

```
public interface Comparable<Item>
{
    public int compareTo(Item that);
}
```

key point: no dependence  
on String data type

## data-type implementation

```
public class String
implements Comparable<String>
{
    ...
    public int compareTo(String b)
    {
        ...
        return -1;
        ...
        return +1;
        ...
        return 0;
    }
}
```

## sort implementation

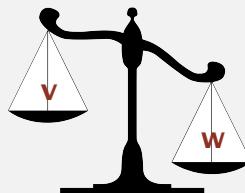
```
public static void sort(Comparable[] a)
{
    int N = a.length;
    for (int i = 0; i < N; i++)
        for (int j = i; j > 0; j--)
            if (a[j].compareTo(a[j-1]) < 0)
                exch(a, j, j-1);
            else break;
}
```

## Comparable API

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Implement `compareTo()` so that `v.compareTo(w)`

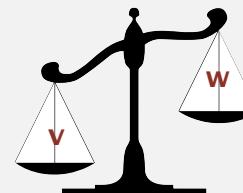
- Defines an order (total order)
- Returns a negative integer, zero, or positive integer if v is less than, equal to, or greater than w, respectively.
- Throws an exception if incompatible types (or either is null).



less than (return -1)



equal to (return 0)



greater than (return +1)

Built-in comparable types. Integer, Double, String, Date, File, ...

User-defined comparable types. Implement the Comparable interface.

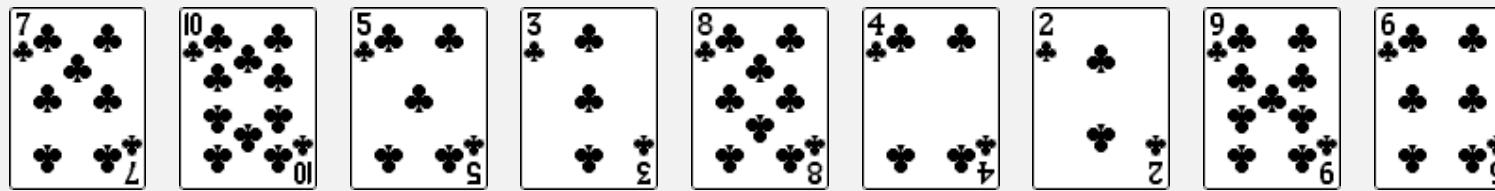
# Selection Sort



## Selection sort demo

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- In iteration  $i$ , find index  $\min$  of smallest remaining entry.
- Swap  $a[i]$  and  $a[\min]$ .

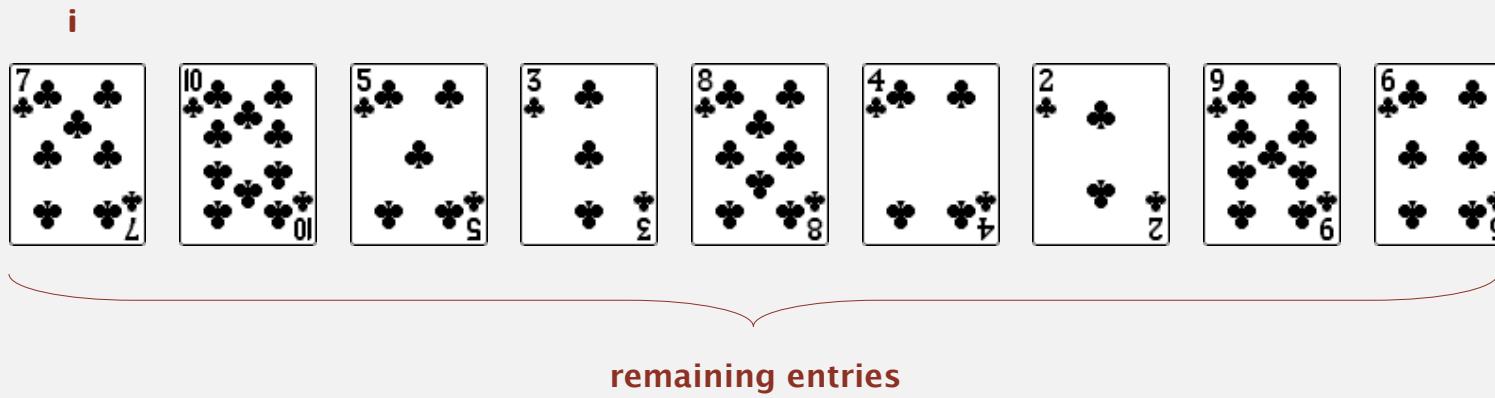


initial

## Selection sort demo

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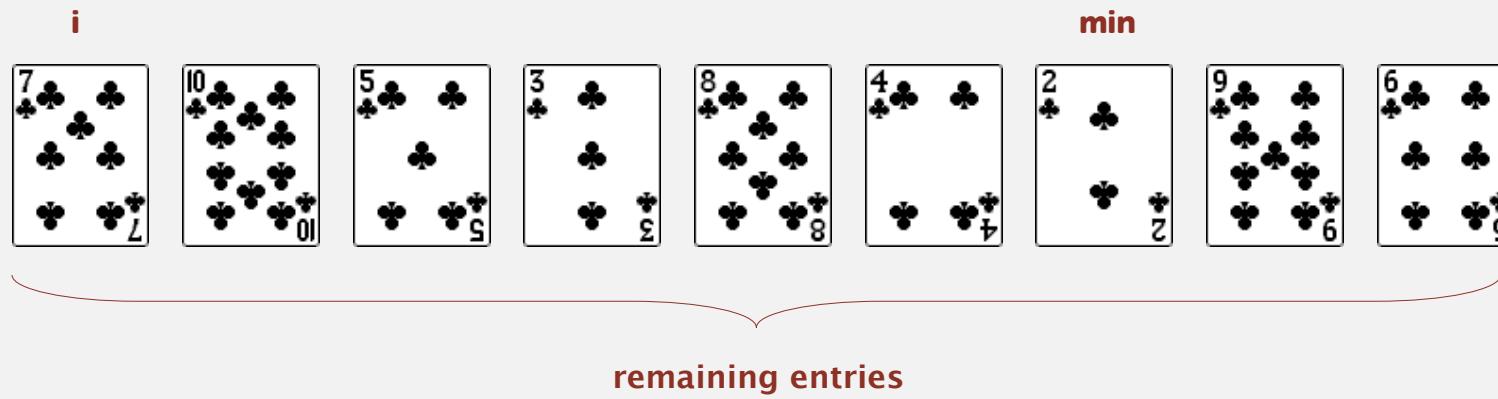
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## Selection sort demo

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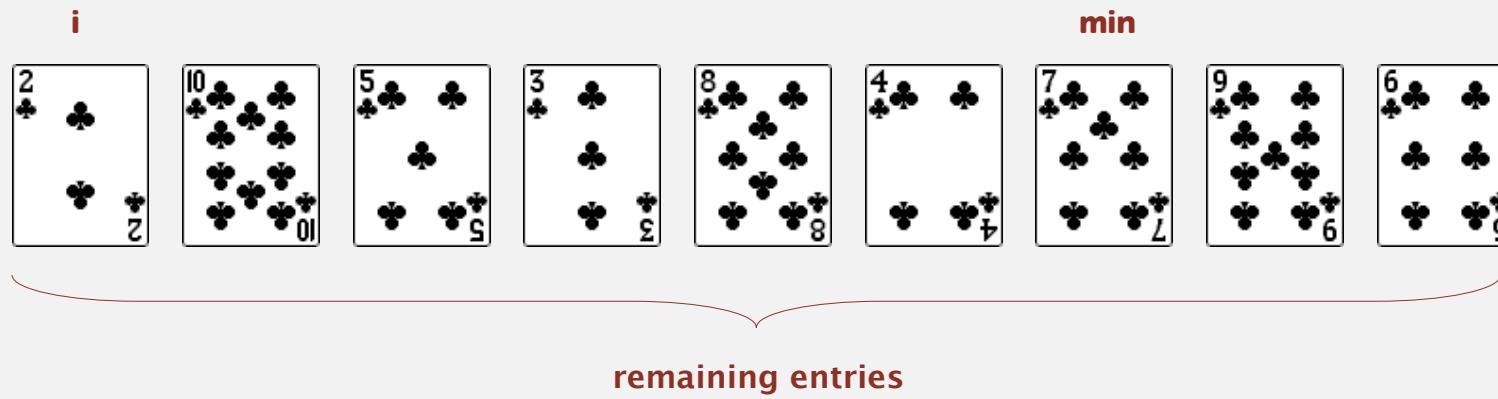
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## Selection sort demo

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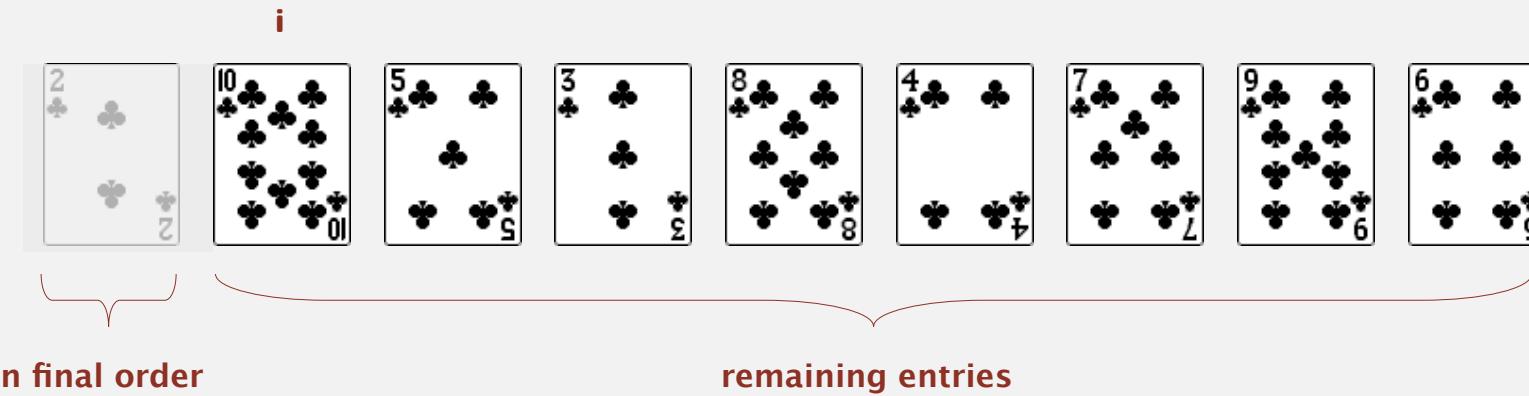
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## Selection sort demo

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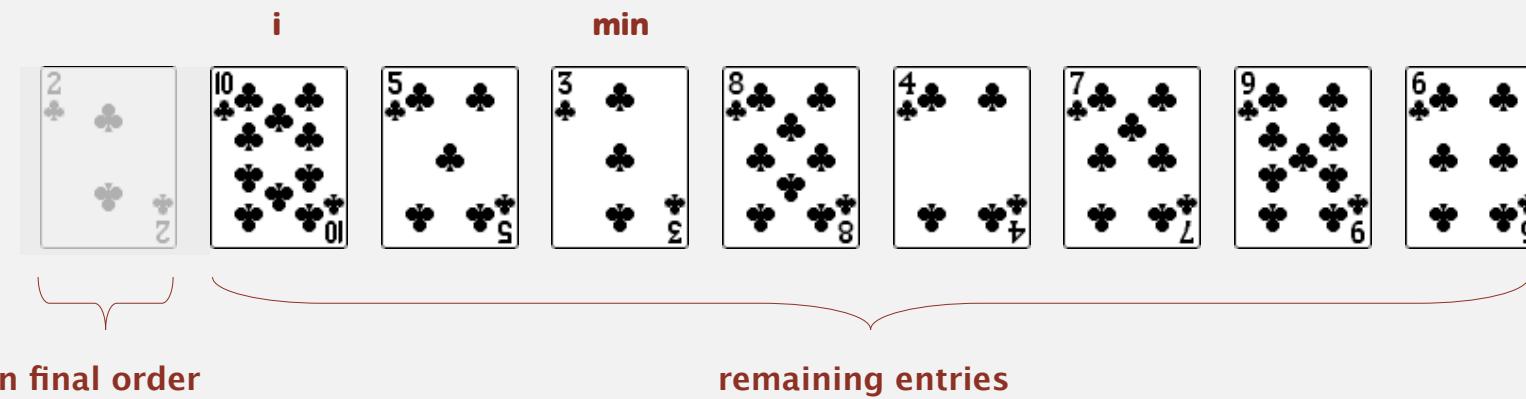
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## Selection sort demo

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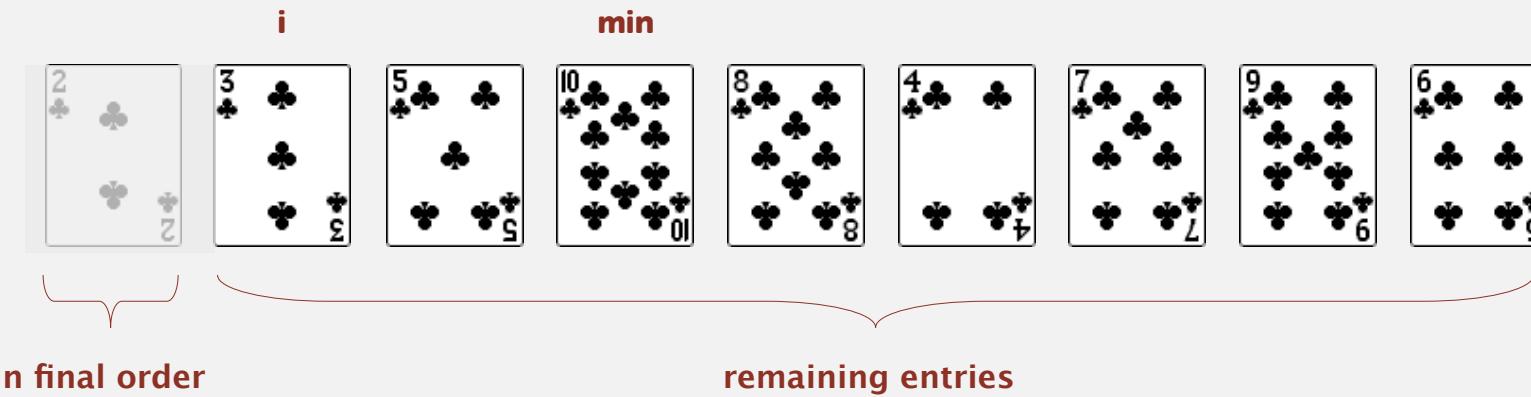
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## Selection sort demo

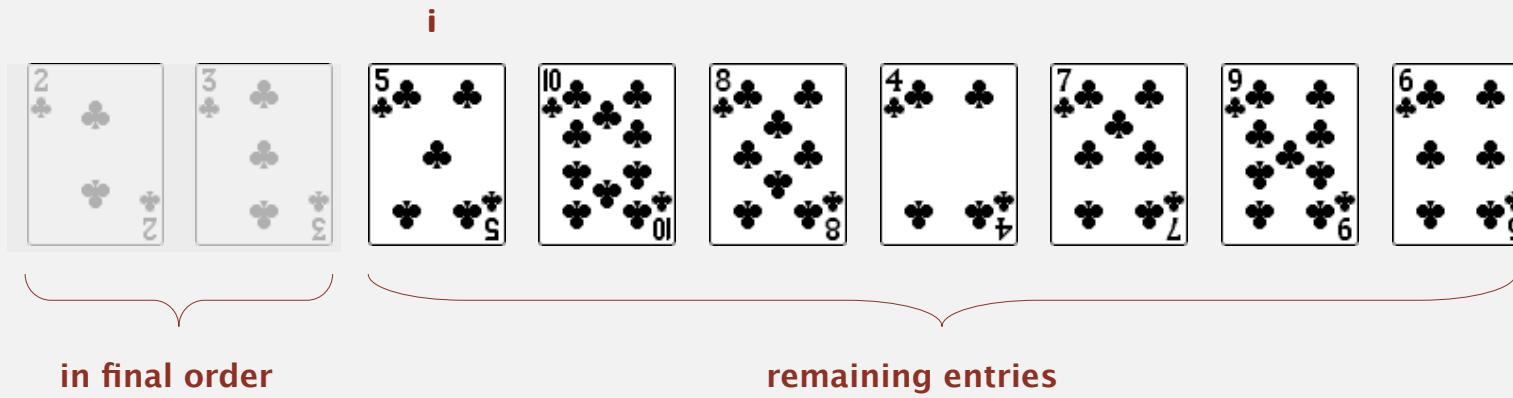
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## Selection sort demo

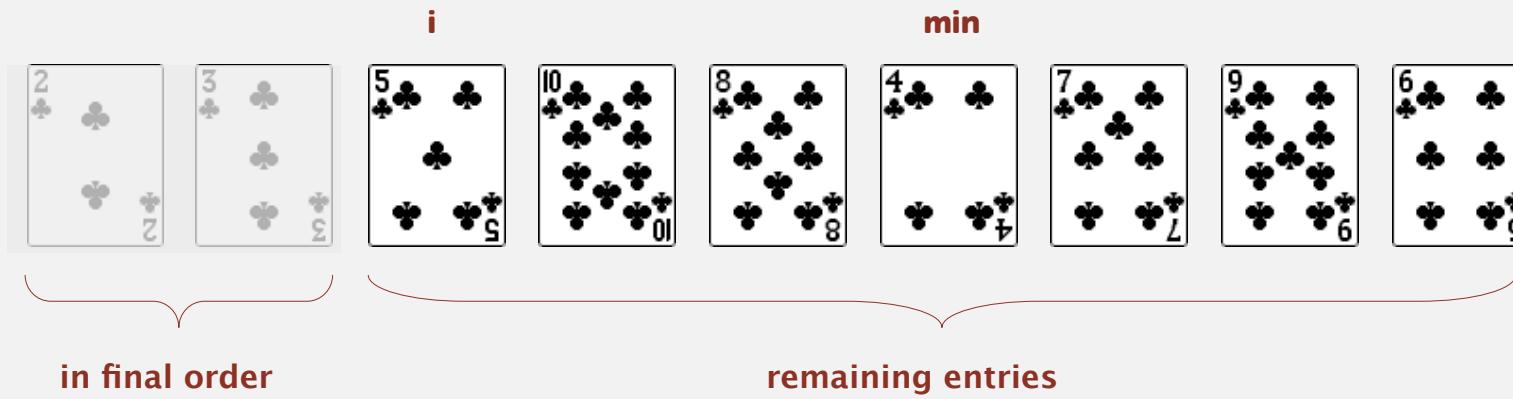
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## Selection sort demo

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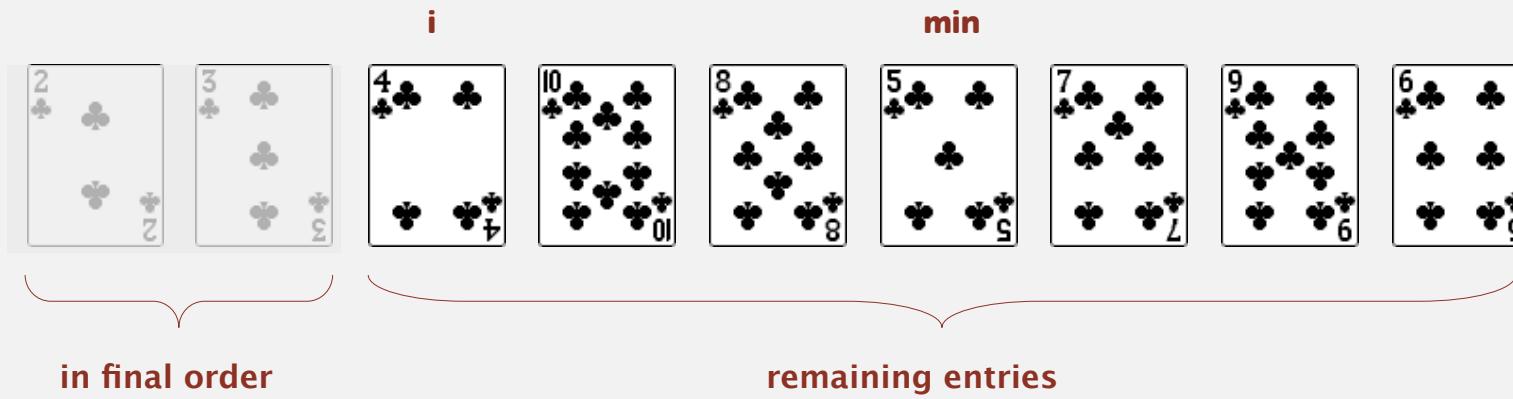
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## Selection sort demo

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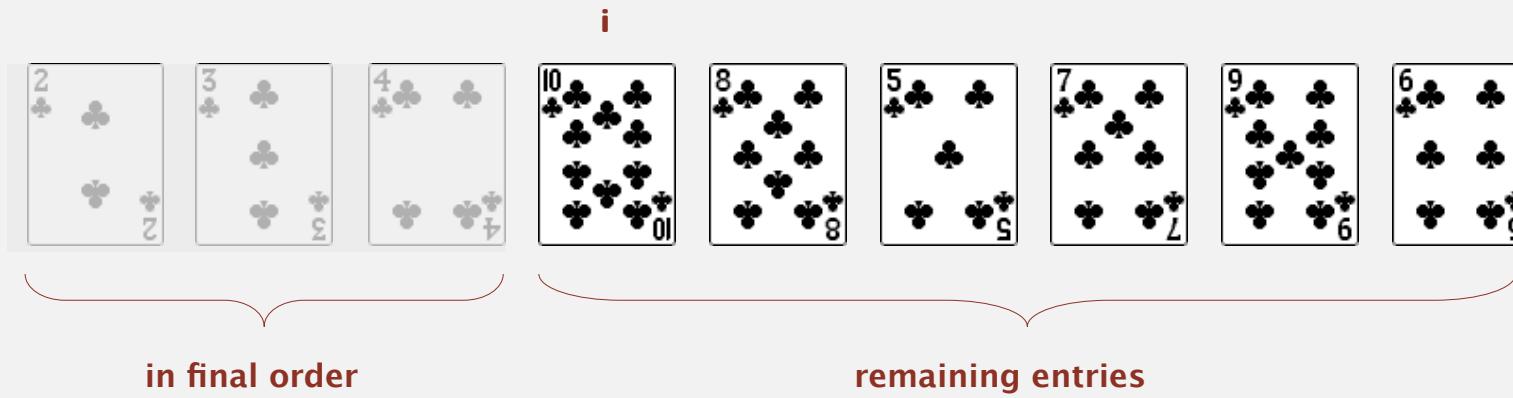
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## Selection sort demo

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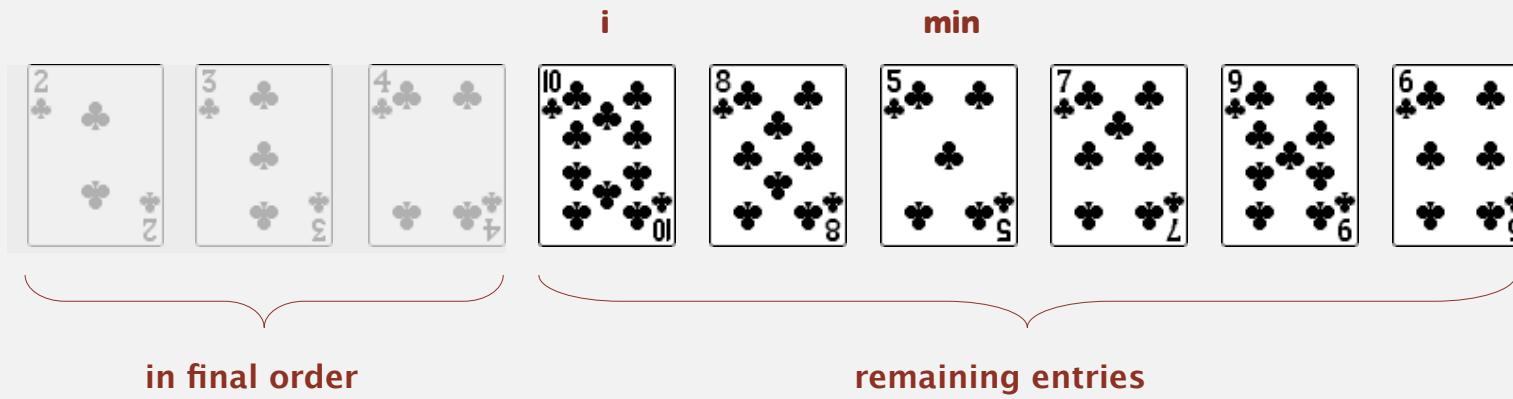
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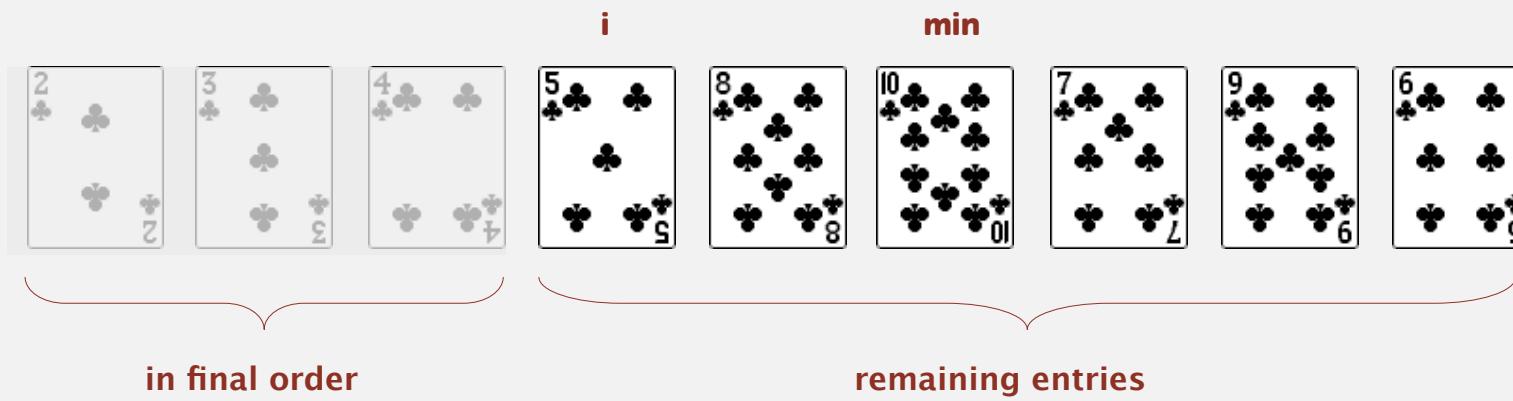
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## Selection sort demo

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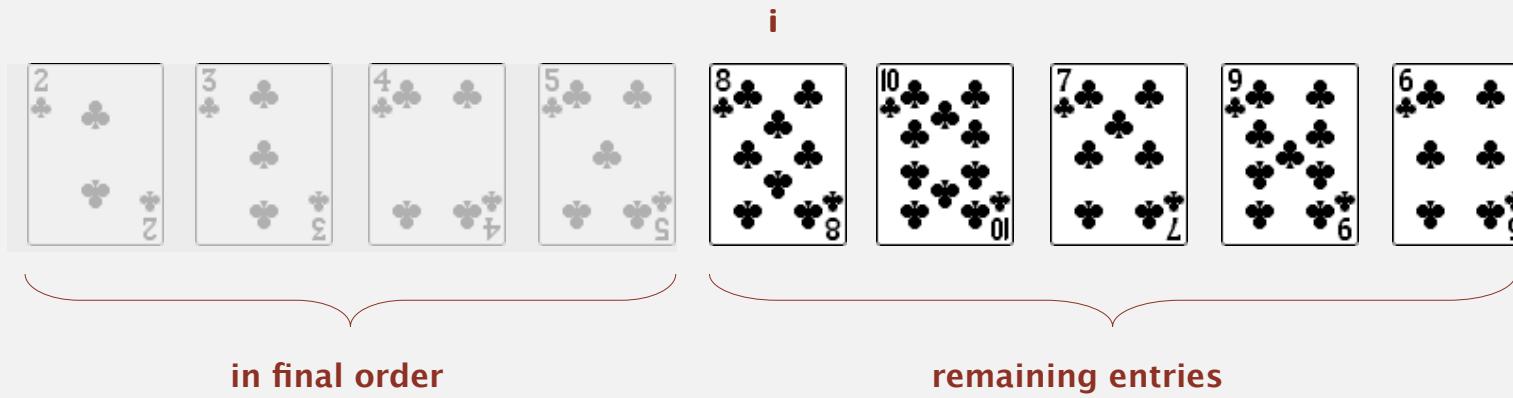
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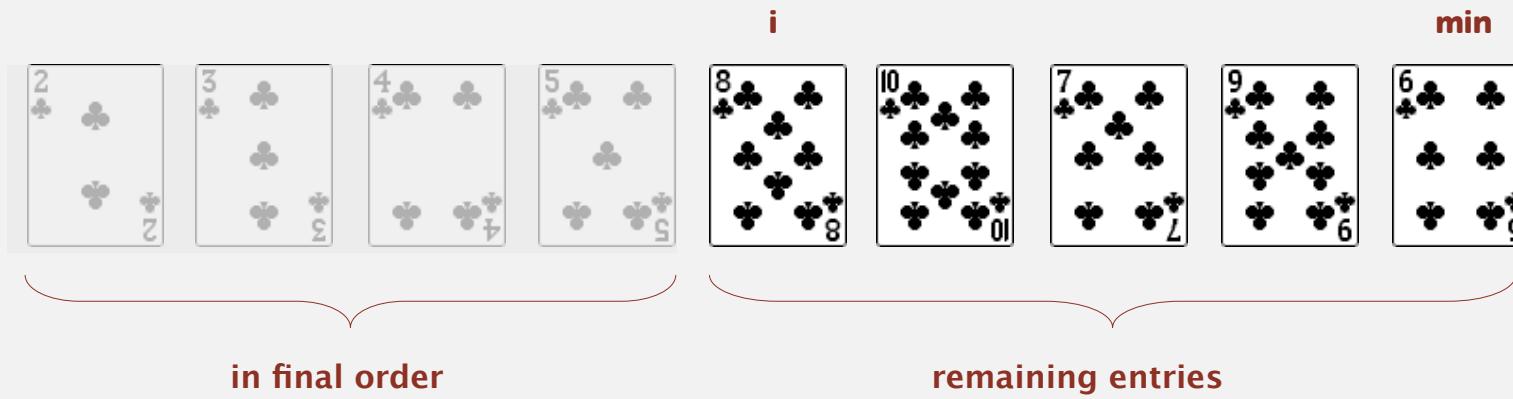
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## Selection sort demo

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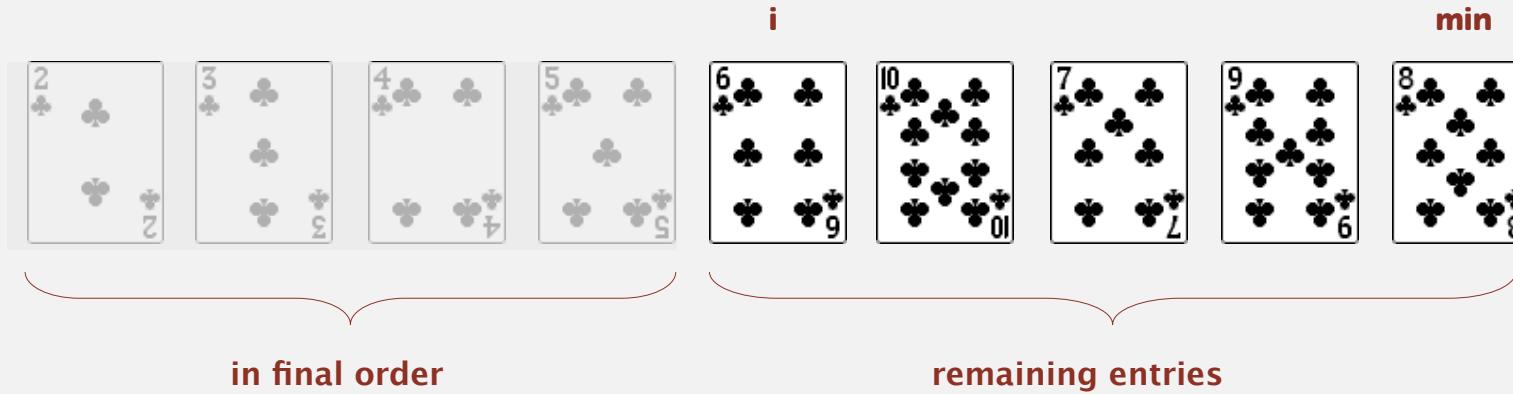
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## Selection sort demo

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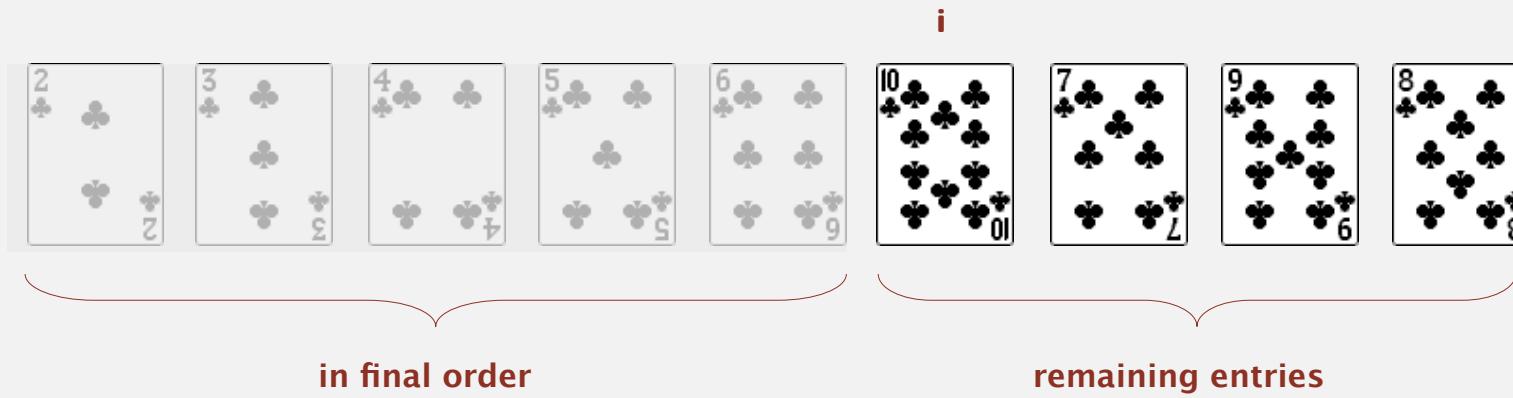
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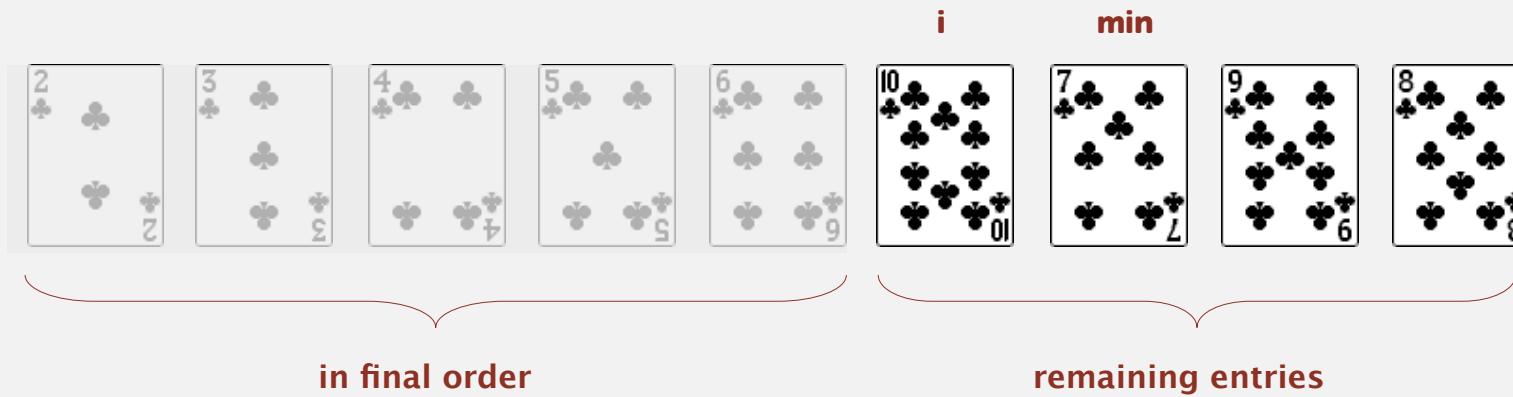
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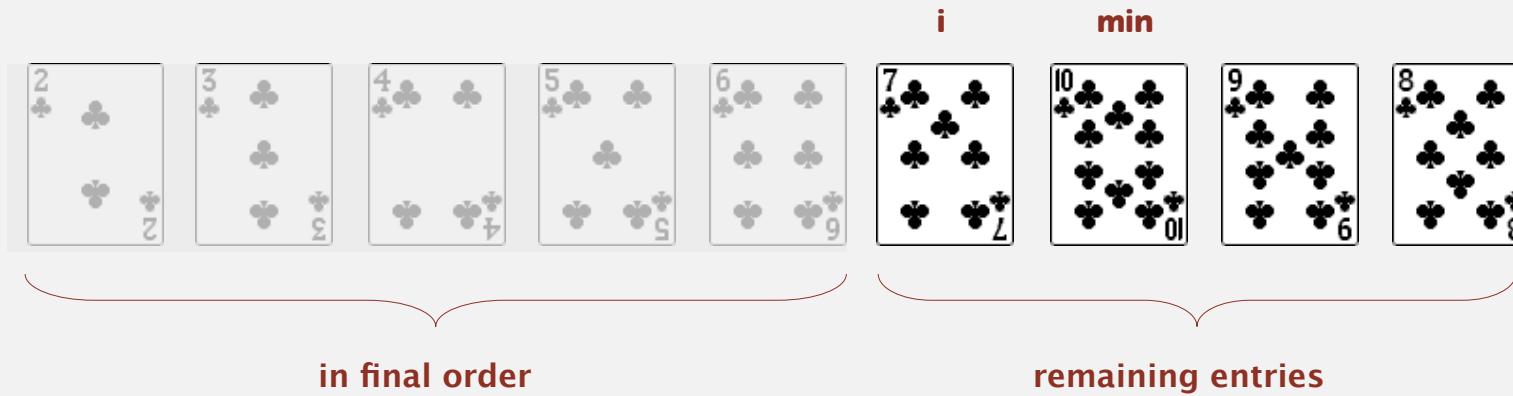
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## Selection sort demo

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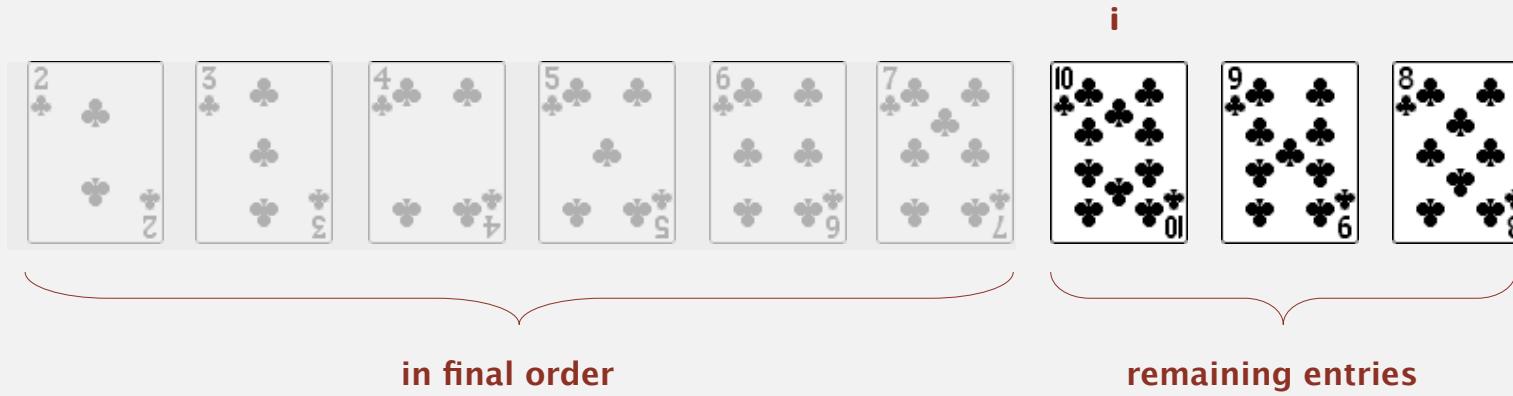
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## Selection sort demo

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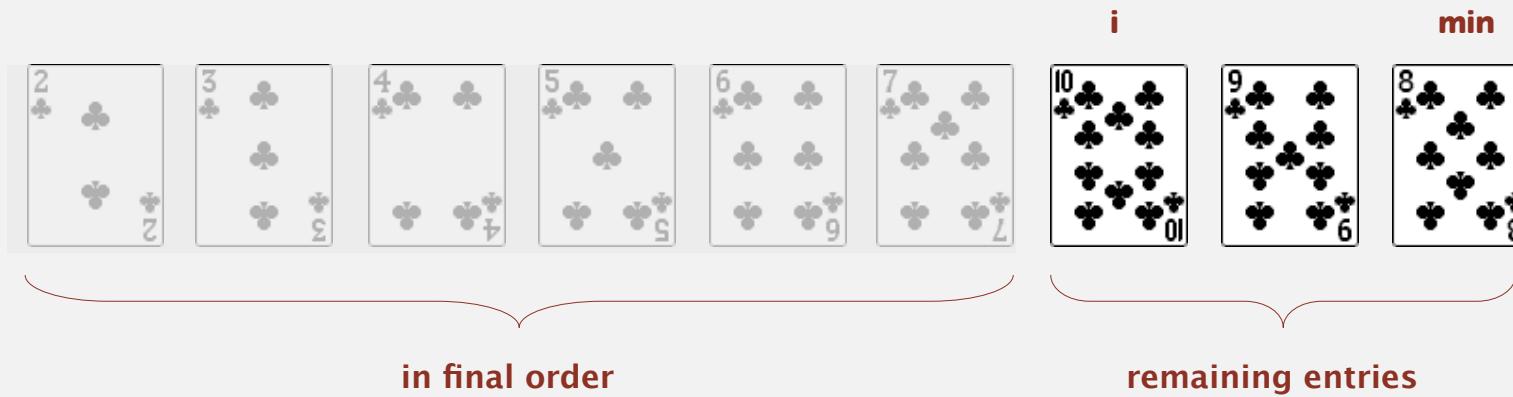
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## Selection sort demo

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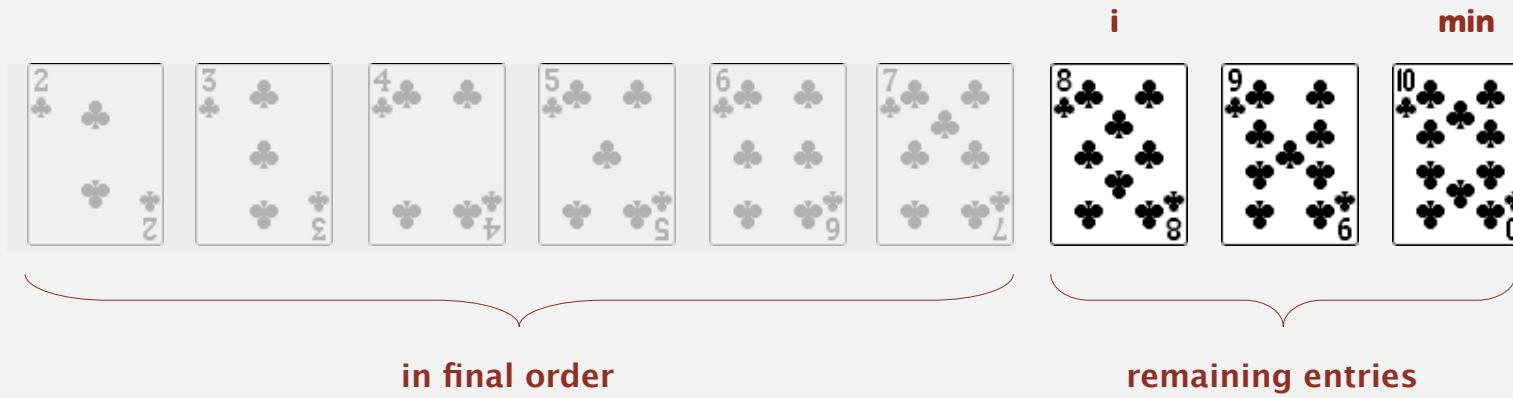
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## Selection sort demo

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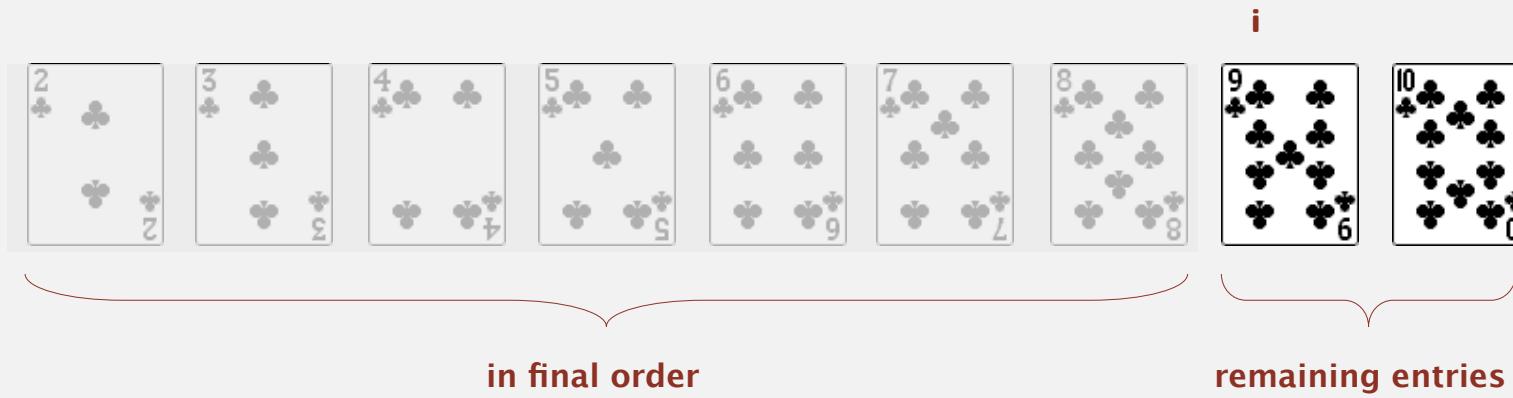
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## Selection sort demo

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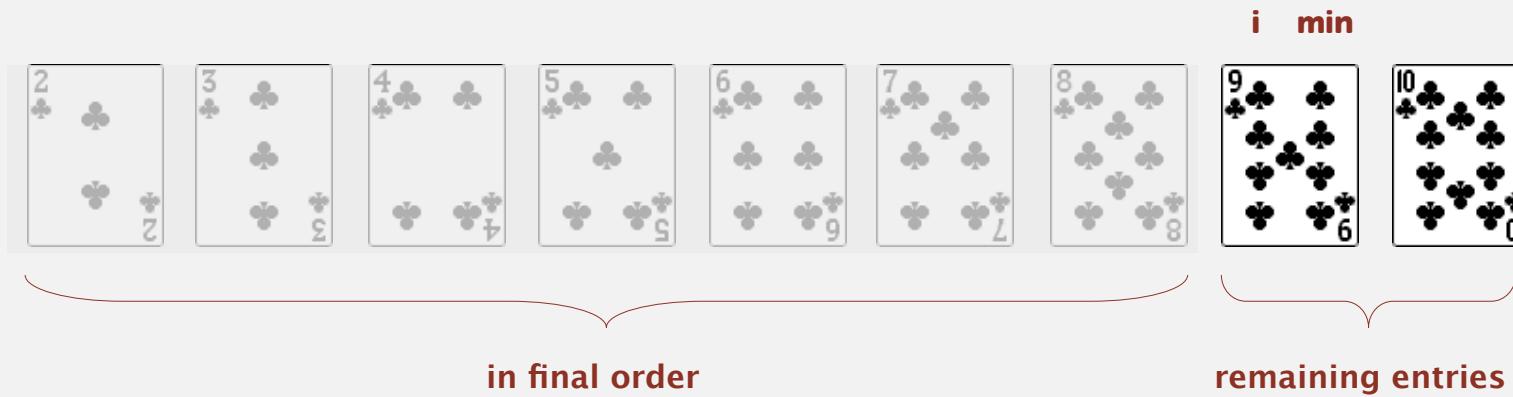
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## Selection sort demo

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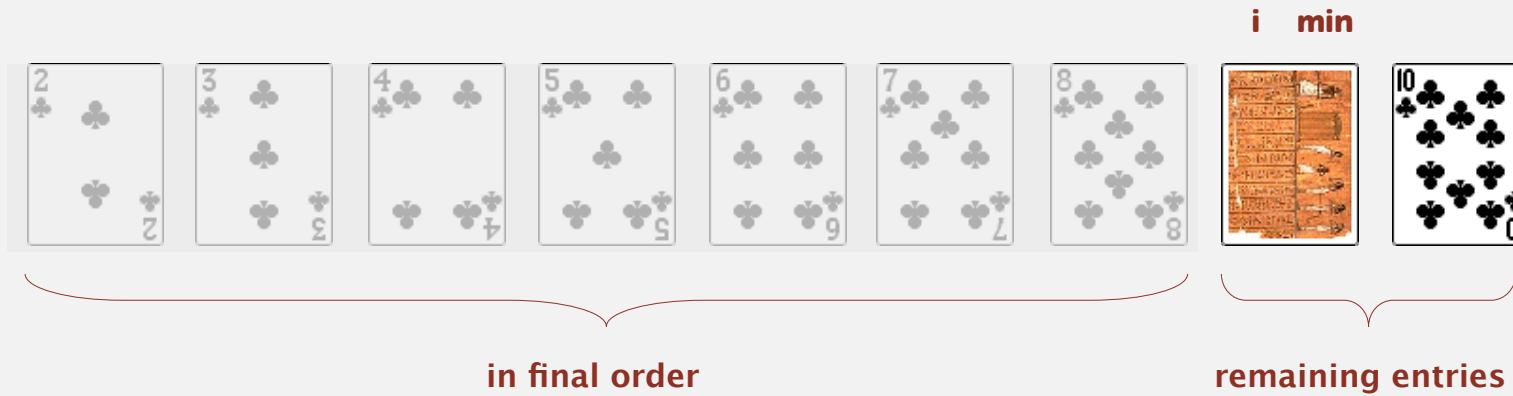
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## Selection sort demo

---

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## Selection sort demo

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## Selection sort demo

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## Selection sort demo

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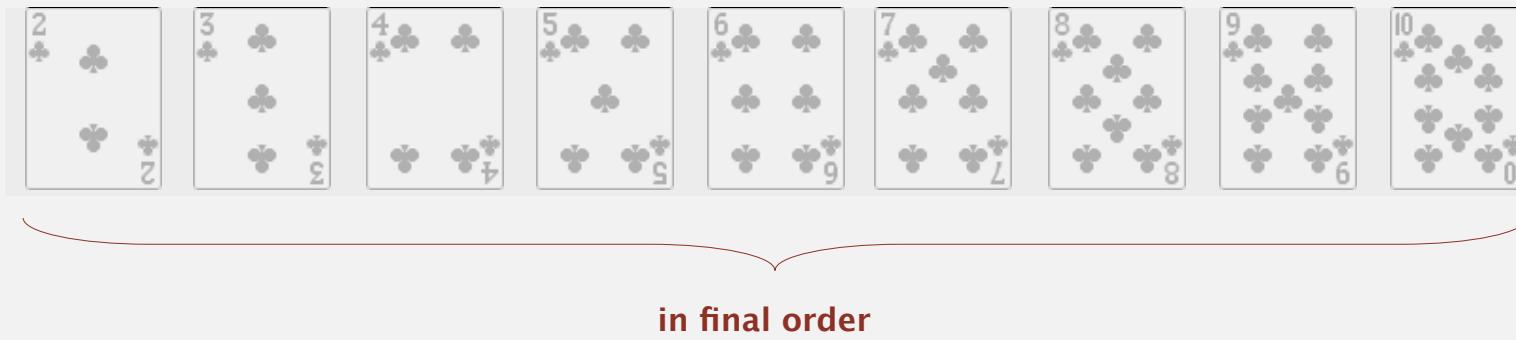
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## Selection sort demo

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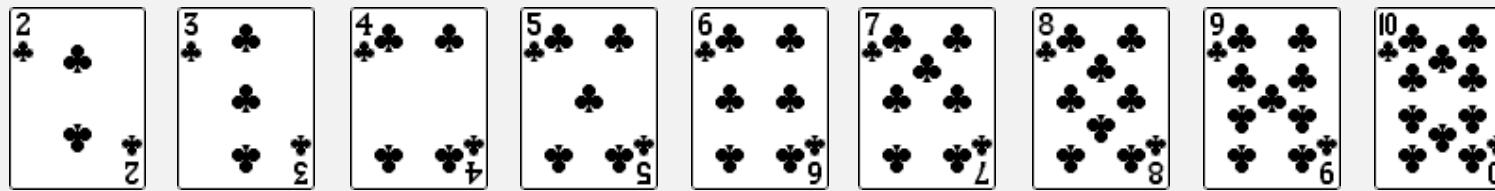
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## Selection sort demo

---

- In iteration  $i$ , find index  $\min$  of smallest remaining entry.
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sorted

## Selection sort

---

Algorithm.  $\uparrow$  scans from left to right.

Invariants.

- Entries to the left of  $\uparrow$  (including  $\uparrow$ ) fixed and in ascending order.
- No entry to right of  $\uparrow$  is smaller than any entry to the left of  $\uparrow$ .



## Human Selection Sort

---

1. Look in the array for the smallest element
2. Swap with the first element from your search
  - \*Note not always the first element in the array
3. Don't bother with the last element, logically that will take care of itself

## Selection sort: Gypsy folk dance

---



## Selection sort inner loop

---

To maintain algorithm invariants:

- Move the pointer to the right.

```
i++;
```



- Identify index of minimum entry on right.

```
int min_index = i;  
for (int j = i+1; j < arr.length; j++) {  
    if (arr[min_index] > arr[j])  
        min_index = j;  
}
```



- Exchange into position.

```
arr[i] = arr[min_index];
```



## Selection sort: Java implementation

---

```
void sort(int arr[])
{
    int temp;
    int min_index;

    // One by one move boundary of unsorted subarray
    for (int i = 0; i < arr.length-1; i++){
        min_index = i;

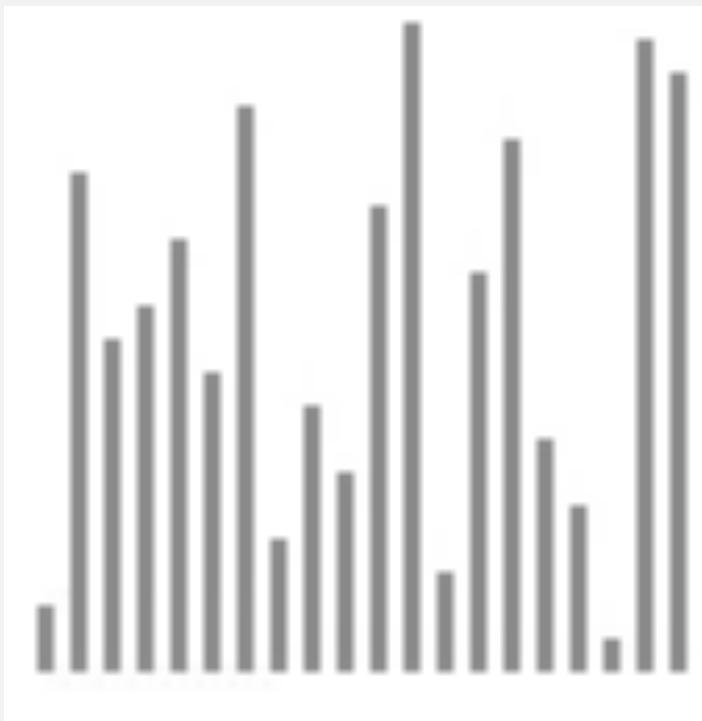
        // Find the minimum element in unsorted array
        for (int j = i+1; j < arr.length; j++) {
            if (arr[min_index] > arr[j])
                min_index = j;
        }

        // Swap the found minimum element with the first element
        temp = arr[i];
        arr[i] = arr[min_index];
        arr[min_index] = temp;
    }
}
```

## Selection sort: animations

---

20 random items

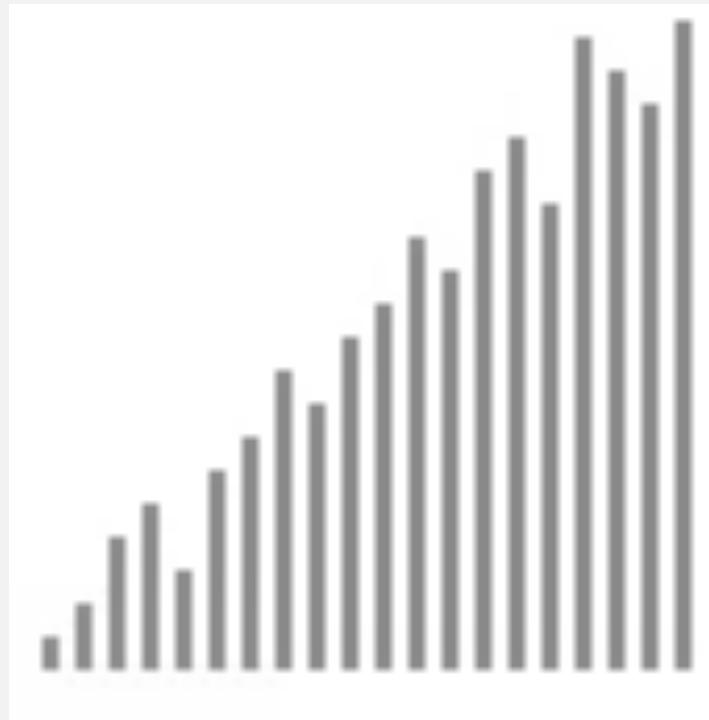


- ▲ algorithm position
- █ in final order
- ▒ not in final order

## Selection sort: animations

---

20 partially-sorted items



- ▲ algorithm position
- in final order
- ▬ not in final order

## Selection sort: analysis

**Proposition.** Selection sort uses  $(N-1) + (N-2) + \dots + 1 + 0 \sim N^2/2$  compares and  $N$  exchanges.

		a[]										
i	min	0	1	2	3	4	5	6	7	8	9	10
		S	O	R	T	E	X	A	M	P	L	E
0	6	S	O	R	T	E	X	A	M	P	L	E
1	4	A	O	R	T	E	X	S	M	P	L	E
2	10	A	E	R	T	O	X	S	M	P	L	E
3	9	A	E	E	T	O	X	S	M	P	L	R
4	7	A	E	E	L	O	X	S	M	P	T	R
5	7	A	E	E	L	M	X	S	O	P	T	R
6	8	A	E	E	L	M	O	S	X	P	T	R
7	10	A	E	E	L	M	O	P	X	S	T	R
8	8	A	E	E	L	M	O	P	R	S	T	X
9	9	A	E	E	L	M	O	P	R	S	T	X
10	10	A	E	E	L	M	O	P	R	S	T	X
		A	E	E	L	M	O	P	R	S	T	X

Trace of selection sort (array contents just after each exchange)

**Running time insensitive to input.** Quadratic time, even if input is sorted.

Data movement is minimal. Linear number of exchanges.

# Classifying Sorting Algorithms

---

## 1. Comparison vs. Non-comparison

- Comparison sorts compare two items at a time
- Non-comparison sorts simply ones that do NOT compare two items

## 2. Time complexity

- How much time it takes the algorithm based on the size of the input

## 3. Space complexity (in-place vs. out-of-place)

- How much memory / space the algorithm will need based on size of the input

## 4. Stability

- Stable sorting algorithms preserve the existing relative order of elements when comparing equal keys

## 5. Internal vs. External

- Internal sort is one where the items being sorted can be kept in main memory / RAM
- External sort is one where the items being sorted need to use external memory

## 6. Recursive vs. Non-recursive

- Recursive sorting uses a divide & conquer approach splitting up dataset into smaller inputs
- Non-recursive simply does not use recursion / does not split the dataset up



## Selection sort: summary

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<b>Comparison sort?</b>	Comparison
<b>Time Complexity</b>	$O(N^2)$
<b>Space Complexity</b>	In place
<b>Internal or External?</b>	Internal
<b>Recursive / Non-recursive?</b>	Non-recursive