COMP20230: Data Structures & Algorithms
Lecture 12: Hash Tables

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Outline

Last Day:

Today: Hash Tables

Hash Tables: Searchable Data Structures

Searching data structures

Example Symbol Tables

Application	Purpose of Search	Key	Value
dictionary	find word definition	word	definition
book index	find relevant pages, word occurrences	term	list of page numbers
account management	process transaction	account number	transaction details
web search	find relevant web pages	keyword	list of page titles and urls
compiler	find type and value of variable	variable name	type and value

ADT of a symbol table

For an **unordered symbol table** the ADT has the following operations:

`put(key, value)`	put key-value pair into the table
`get(key)`	value paired with key (null if key is absent)
`delete(key)`	remove key from table and value paired with key
`contains(key)`	is there a value paired with key?
`isEmpty()`	is the table empty?
`size()`	number of key-value pairs in the table
`keys()`	all the keys in the table

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Aside: Ordered Symbol Table ADT

If we want to keep our symbols ordered, we need to keep information about their rank and a number of other operations are required: `min()`, `max()`, `floor(key)`, `ceiling(key)`, `rank(key)`, `select(rank)`, `deleteMin()`, `deleteMax()`, `size(low_key,high_key)`, `keys(low_key,high_key)`

Searching data structures

Three classic data structures that can support efficient searchable symbol-table implementations:

1. Hash tables
2. Binary search trees
3. Balanced search Trees: 2–3 Trees, Red-black trees, AVL Trees

Hash figures adapted from:

Algorithms (Sedgewick & Wayne)

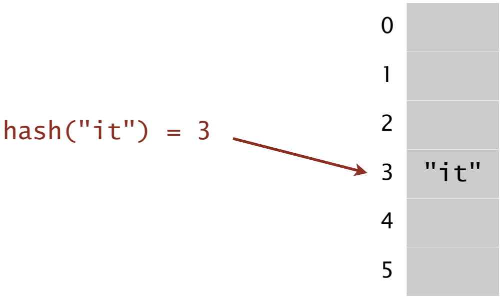
Hash Tables

Hash Tables

Save items in a key-indexed table (index is a function of the key)

Hash Function

Method for computing array index from a key.



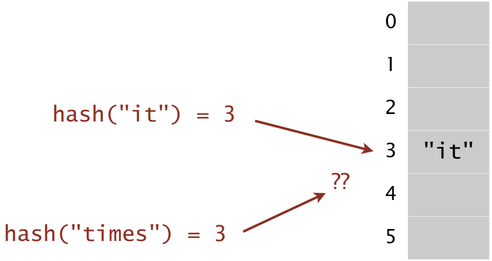
Hash Tables: Requirements and Issues

Compute the hash function

Good algorithm (i.e. fast, efficient, scalable etc.)

Collision resolution

Algorithm and data structure to handle two keys that hash to the same array index



Example: Python Dictionary

```
airports={"JFK": ("John F Kennedy Intl","United States",40.639751, -73.778925),
          "SYD": ("Sydney Intl","Australia",-33.946111,151.177222),
          "LHR": ("London Heathrow","United Kingdom",51.4775,-0.461389)}

# print a search result
print(airports["SYD"])
print("Airport Keys: ", airports.keys())

# add an airport to the dictionary
airports["AMS"]=("Schiphol","Netherlands",52.308613,4.763889)

# store the value of a search and print it
destination=airports.get("AMS")
print(destination)

# pop (search and remove) a value from dict and save it in a variable
oz_airport = airports.pop("SYD")
print("Airport Keys: ", airports.keys())

# what is the hash for key AMS?
# Does it change if I call it twice? What if I rerun the program?
print("AMS hash is:", hash("AMS"))
print("AMS hash is:", hash("AMS"))
print("DUB hash is:", hash("DUB"))
```

Output:

```
('Sydney Intl', 'Australia', -33.946111, 151.177222)
Airport Keys:  dict_keys(['JFK', 'SYD', 'LHR'])
('Schiphol', 'Netherlands', 52.308613, 4.763889)
Airport Keys:  dict_keys(['JFK', 'LHR', 'AMS'])
AMS hash is: 6708379502801481095
AMS hash is: 6708379502801481095
DUB hash is: -3052993293245237079
```

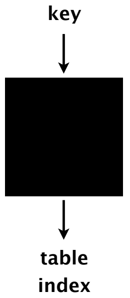
Hash Tables: Computing the Hash Function

Ideally: Scramble the keys uniformly to produce

Equally computable table index
Each table index equally likely for each key.

Hash Codes

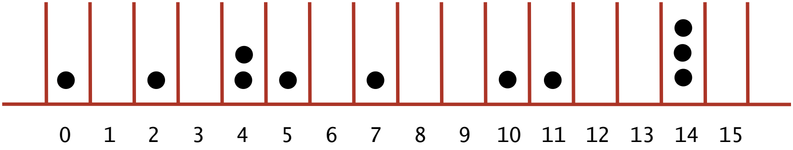
Integers, e.g.
Most significant part of a float;
Memory address of an object



Hash Tables

Uniform Hashing Assumption

Each key is equally likely to hash to an integer between 0 and $M-1$.

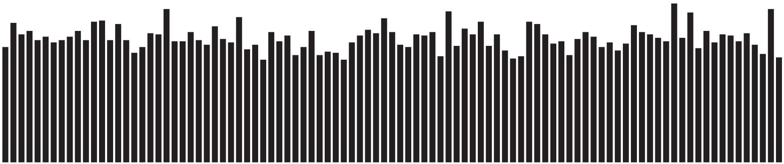


Bins and Balls

Evenly distribute balls into the slots of a hash table.
Throw balls aiming for uniform distribution at M bins.

Example Hash Table

Java hash table implementation result for distributing keys of strings (words) in Tale of Two Cities. (M=97)

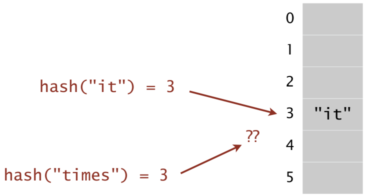


Hash value frequencies for words in Tale of Two Cities (M = 97)

Hash Tables

Collisions

Two distinct keys hashing to same index
Collisions inevitable (unless *dynamic perfect hashing* implemented – memory hungry!).



Birthday Problem

How many birthdays on the same day in a class of 70? With only 23 people, the probability that two people have same birthday is 50%

Hash Tables

Implementation

Separate Chaining Symbol Table
Linear Probing

Separate Chaining Symbol Table

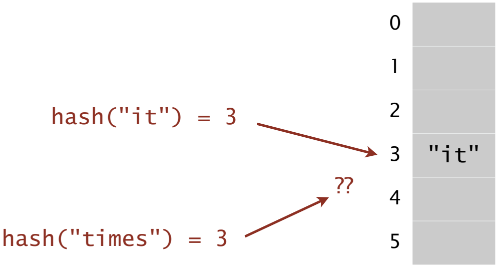
M lists and N keys.

Use an array of $M < N$ linked lists

`Hash`: Map key to integer i between 0 and $M - 1$

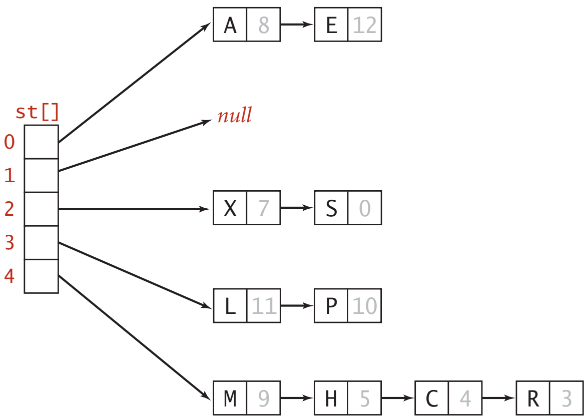
`Insert`: Put at front of ith chain (if not already there)

`Search`: Need to search only ith chain



Separate Chaining Symbol Table

key	hash	value
S	2	0
E	0	1
A	0	2
R	4	3
C	4	4
H	4	5
E	0	6
X	2	7
A	0	8
M	4	9
P	3	10
L	3	11
E	0	12



Separate Chaining Symbol Table

Getting the balance right: what size for balance between `insert` and `search`?

Analysis

Under uniform hashing assumption, prob. that the number of keys in a list is within a constant factor of N/M is extremely close to 1

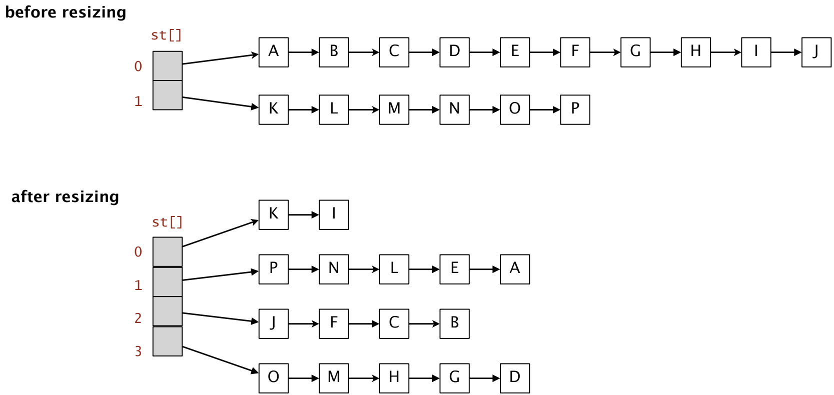
Consequences

Number of probes for `search/insert` is proportional to N/M
M too large $\Rightarrow$ too many empty chains
M too small $\Rightarrow$ chains too long
Typical choice: $M \sim N/4 \Rightarrow$ constant-time ops

Separate Chaining Symbol Table

Resizing: Average length of list $N/M = constant$

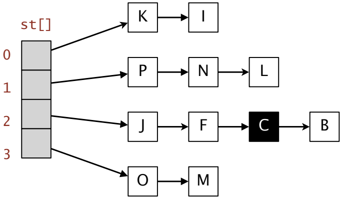
Double size of array M when $N/M \geq 8$
Halve size of array M when $N/M \leq 2$
Need to rehash all keys when resizing



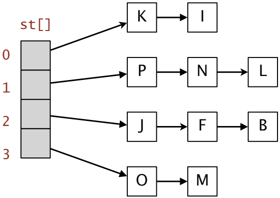
Separate Chaining Symbol Table

Deleting is straight-forward

before deleting C



after deleting C



Collision Resolution Strategy: Use Open Addressing

Open addressing

When a new key collides, find next empty slot, and put it there

st[0]	jocularly
st[1]	*null*
st[2]	listen
st[3]	suburban
⋮	*null*
st[30000]	browsing

Linear-probing Hash Table

Linear-probing

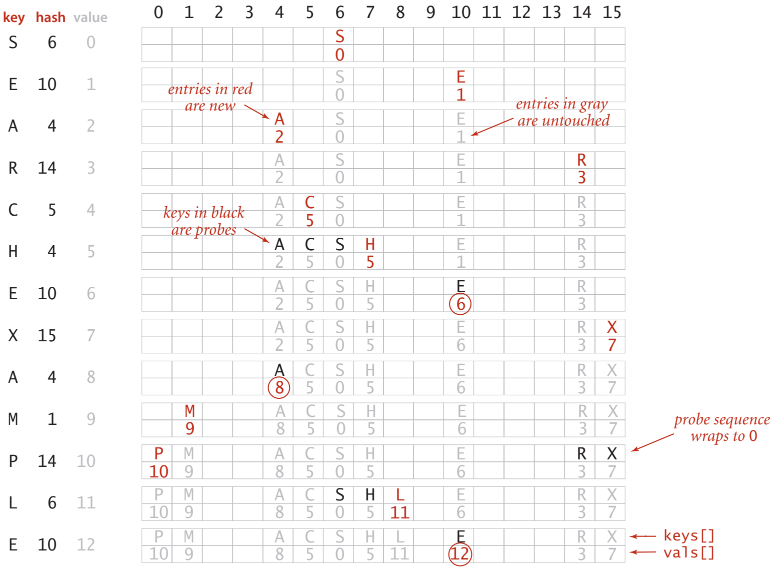
Open addressing scheme for resolving collisions in hash tables

`Hash:` Map key to integer i between 0 and $M-1$ `Insert:` Put at table index i if free; if not try $i+1$, $i+2$, etc. `Search:` Search table index i; if occupied but no match, try $i+1$, $i+2$, etc.

Note

Array size M must be greater than number of key-value pairs N

Example of Linear Probing (video on moodle)



Linear Probing Hash Table

Resizing: Average length of list $N/M \leq 1/2$

Double size of array M when $N/M \leq 1/2$
Halve size of array M when $N/M \geq 1/8$
Need to rehash all keys when resizing.

before resizing

	0	1	2	3	4	5	6	7
keys[]		E	S			R	A	
vals[]		1	0			3	2	

after resizing

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
keys[]					A		S				E				R	
vals[]					2		0				1				3	

Linear Probing Hash Table

Deletion: What happens if we delete S from hash table?

before deleting S

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
keys[]	P	M			A	C	S	H	L		E				R	X
vals[]	10	9			8	4	0	5	11		12				3	7

doesn't work, e.g., if hash(H) = 4

after deleting S ?

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
keys[]	P	M			A	C		H	L		E				R	X
vals[]	10	9			8	4		5	11		12				3	7

Linear Probing Hash Table

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vals[]	10	9			8	4		5	11		12				3	7

Cannot just leave null/None – will not find H

Need to rehash the cluster to the right of the deleted key.