

# COMP20230: Data Structures & Algorithms

## Lecture 17: Graphs (2)

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## Today

- Graph Traversal: DFS and BFS Complexity
- Weighted Graphs
- Shortest Path – Dijkstra's Algorithm

## Take home message

Graphs can be weighted along edges to better represent and model some scenarios.

Dijkstra's algorithm is a shortest path algorithm.

# Recap: Graphs

A collection of vertices (nodes), joined together by edges.  
No definite beginning or end.

## Representations of graphs

Two common methods: adjacency lists and adjacency matrices

## Searching

Many algorithms begin by searching (e.g. DFS, BFS) their input graph to obtain this structural information.

Worst case is you need to traverse and explore **every** vertex and **every** edge.

Time Complexity for  $|V|$  vertices and  $|E|$  edges

$$\mathcal{O}(|V| + |E|)$$

where, depending on sparsity,  $\mathcal{O}(|E|)$  is between  $\mathcal{O}(1)$  and  $\mathcal{O}(|V|^2)$

Breadth-first search is **complete**, but depth-first search is not, i.e. the input to breadth-first search is assumed to be a finite graph, represented explicitly (e.g. by an adjacency list).

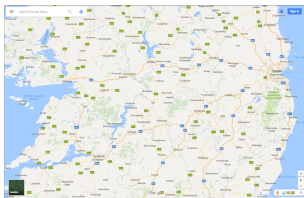
# Weighted Graphs

Sometimes graphs have a weights on edges.

## Examples

A road network with the distance between two cities

A computer network – hops between routers for a streaming video packet with time between nodes (e.g. think ping and traceroute).



# Routing

The screenshot displays the Google Maps routing interface. On the left, a sidebar shows the start and end points: Mountshannon, Co. Clare and Arklow, Co. Wicklow. Below this, there are three route options listed:

- via M7**: 2 h 56 min, 273 km. Fastest route, the usual traffic. ▲ This route has tolls. [DETAILS](#)
- via M7 and R747**: 3 h, 216 km
- via M6**: 3 h 1 min, 270 km

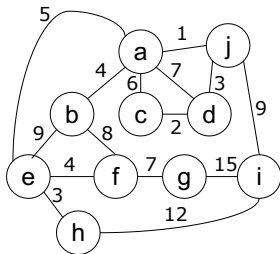
The main map area shows the route highlighted in blue. The route starts at Mountshannon, goes south to Limerick, then east through Ennis, Roscrea, and Thurles, passing through the M7 and M6 motorways, before reaching Arklow. A 'Sign in' button is visible in the top right corner of the map area. The bottom of the map shows the Google logo, map data copyright (©2017 Google), and a scale bar (20 km).

**GREAT BRITAIN AND IRELAND with 2 JERSEY COASTS**

Latitude and Longitude coordinates are provided for numerous locations across the British Isles. The table lists locations such as London, Edinburgh, Glasgow, and many others, along with their respective coordinates. The map includes a compass rose and a scale bar.

Place	Latitude	Longitude
London	51° 30' N	0° 07' W
Edinburgh	55° 55' N	3° 07' W
Glasgow	55° 55' N	4° 15' W
Belfast	54° 38' N	5° 50' W
Cardiff	53° 20' N	3° 08' W
Manchester	53° 20' N	2° 15' W
Birmingham	52° 28' N	1° 50' W
Liverpool	53° 20' N	3° 00' W
Sheffield	53° 28' N	1° 08' W
Nottingham	52° 55' N	1° 05' W
Leeds	53° 45' N	1° 05' W
Bristol	51° 30' N	2° 35' W
Exeter	50° 45' N	3° 35' W
Plymouth	50° 45' N	4° 15' W
Southampton	50° 45' N	1° 00' W
Portsmouth	50° 45' N	1° 05' W
Woolwich	51° 28' N	0° 05' W
Canterbury	51° 10' N	0° 55' E
London	51° 30' N	0° 07' W
Edinburgh	55° 55' N	3° 07' W
Glasgow	55° 55' N	4° 15' W
Belfast	54° 38' N	5° 50' W
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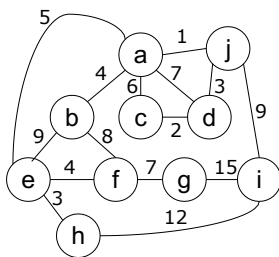
# Weighted Graph Adjacency List



a	b(4), c(6), d(7), e(5), j(1)
b	a(4), e(9), f(8)
c	a(6), d(2)
d	a(7), c(2), j(3)
e	a(5), b(9), f(4), h(3)
f	b(8), e(4), g(7)
g	f(7), i(15)
h	e(3), i(12)
i	g(15), h(12), j(9)
j	a(1), d(3), i(9)



# Weighted Graph Adjacency Matrix



	a	b	c	d	e	f	g	h	i	j
a		4	6	7	5					1
b	4				9	8				
c	6			2						
d	7		2							3
e	5	9				4		3		
f		8			4		7			
g						7			15	
h					3				12	
i							15	12		9
j	1			3					9	

## Shortest Path – Simple Case

How would we calculate the shortest path if there were no weights?

# Unweighted Graph Shortest Path

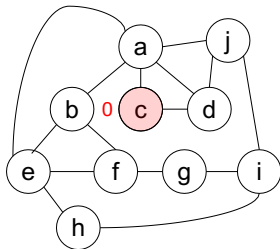
## Shortest Path – Simple Case

How would we calculate the shortest path if there were no weights?

Ignore weights on edges

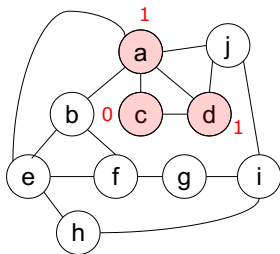
Breadth First Search will give us shortest path

# Breadth First Search – Unweighted Shortest Path



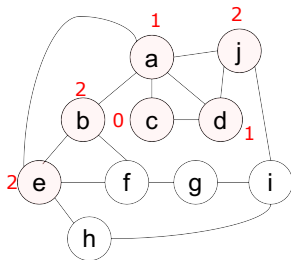
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# Breadth First Search – Unweighted Shortest Path



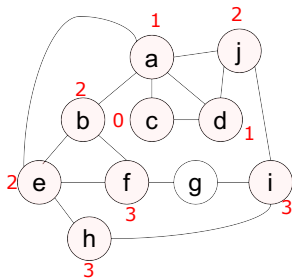
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# Breadth First Search – Unweighted Shortest Path



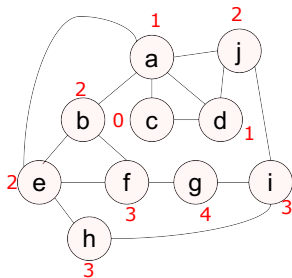
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# Breadth First Search – Unweighted Shortest Path



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# Breadth First Search – Unweighted Shortest Path



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# Searching a weighted graph

## Dijkstra's algorithm

Find the shortest paths between nodes in a graph.

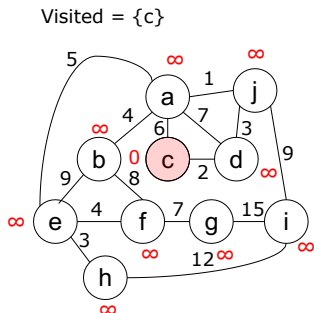
## Dijkstra Variations

Original is to find the shortest path between two given nodes  
Variant fixes a vertex as the origin and finds shortest paths to all other nodes in the graph (a shortest-path tree)

# Dijkstra's Algorithm

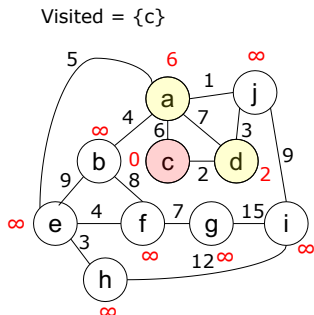
- Pick start and end vertices (from c to g)
- Mark start vertex as 0 distance and all other vertices at  $\infty$
- Repeat until we reach end vertex:
  - ① Calculates total distance to each unvisited neighbour for the current vertex
  - ② Update neighbour's distance if smaller
  - ③ Pick the vertex with the lowest marked distance and set as current
  - ④ Mark current as visited

# Dijkstra's Algorithm



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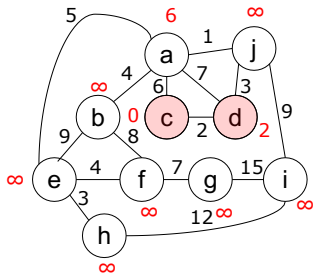
# Dijkstra's Algorithm



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# Dijkstra's Algorithm

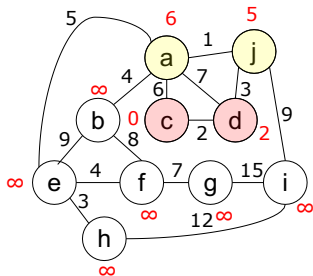
Visited = {c, d}



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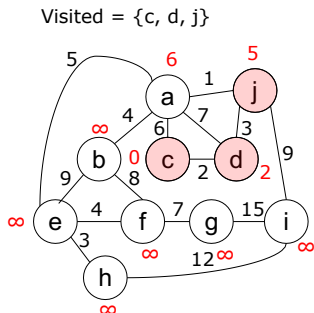
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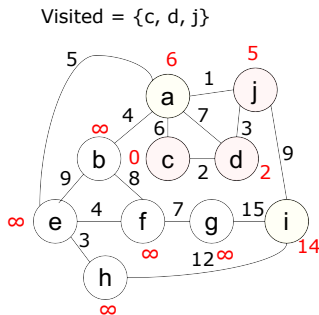
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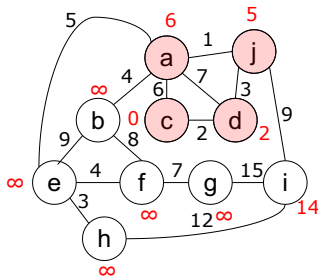


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# Dijkstra's Algorithm

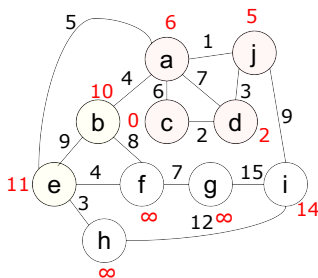
Visited = {c, d, j, a}



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# Dijkstra's Algorithm

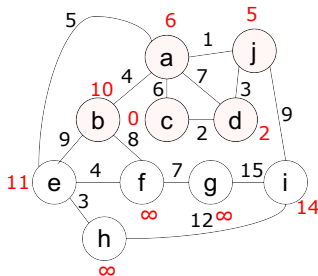
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# Dijkstra's Algorithm

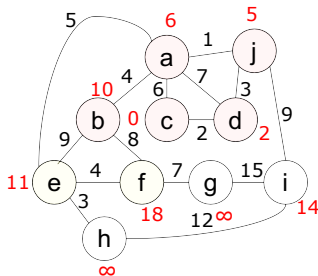
Visited = {c, d, j, a, b}



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# Dijkstra's Algorithm

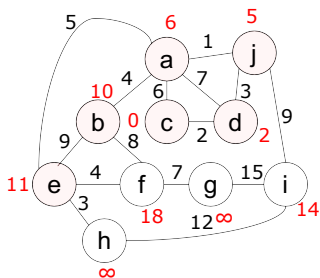
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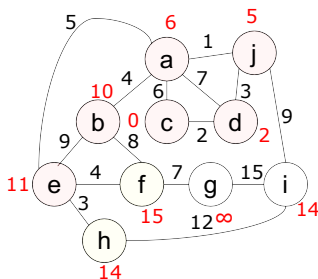
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# Dijkstra's Algorithm

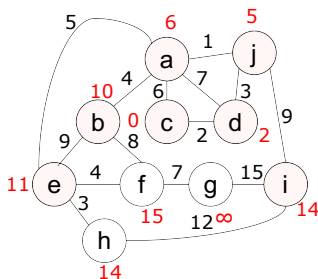
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# Dijkstra's Algorithm

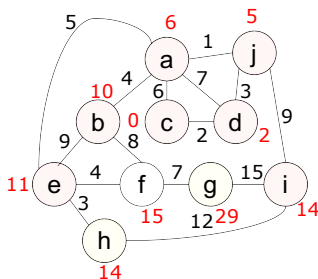
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# Dijkstra's Algorithm

Visited = {c, d, j, a, b, e, i}

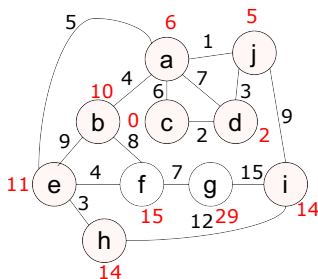


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# Dijkstra's Algorithm

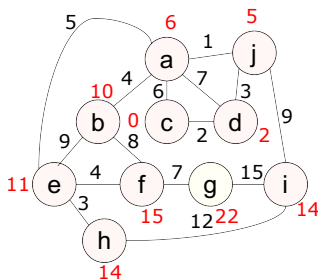
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# Dijkstra's Algorithm

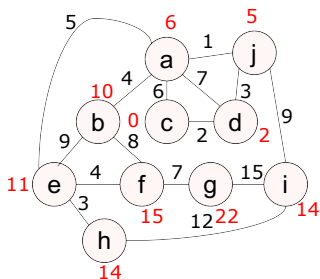
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# Dijkstra's Algorithm

Visited = {c, d, j, a, b, e, i, h, f, g}



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# Dijkstra's Algorithm

```
function dijkstra:
Input:  a graph G, a node n
Output: each node of G gets the shortest distance from n
for each node c of G not n do
    distance[c] ← infinity
endfor
distance[n] ← 0
current ← n
visited ← {n}
while visited does not contain all nodes of G do
    for each node c neighbour of current and not visited do
        # is new route to neighbour better? min(|n|+|n--c| and |c|)
        distance[c] ← min(distance[n] + weight(current, c),
distance[c])
    endfor
    # set next node to search from - lowest of all non-visited
    current ← non visited node with smallest value
    add current to visited
endwhile
```

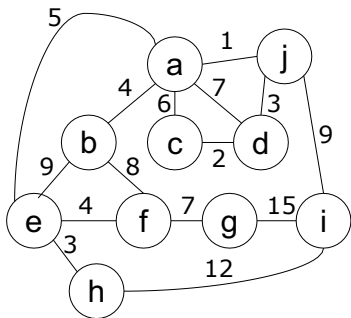
# Dijkstra Steps

## Shortest Path Tree

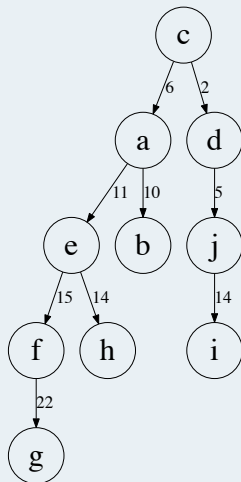
- Visited nodes are in the rows
- Distance from node c in columns for visited nodes
- Best distance from c to a column's node is highlighted in bold

	a	b	c	d	e	f	g	h	i	j
c	6(c-a)	$\infty$		2(c-d)	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	$\infty$
d	6(c-a)	$\infty$			$\infty$	$\infty$	$\infty$	$\infty$	$\infty$	5(c-d-j)
j	6(c-a)	$\infty$			$\infty$	$\infty$	$\infty$	$\infty$	14(c-d-j-i)	
a		10(c-a-b)			11(c-a-e)	$\infty$	$\infty$	$\infty$	14(c-d-j-i)	
b					11(c-a-e)	18(c-a-b-f)	$\infty$	$\infty$	14(c-d-j-i)	
e						15(c-a-e-f)	$\infty$	14(c-a-e-h)	14(c-d-j-i)	
i						15(c-a-e-f)	29(c-d-j-i)	14(c-a-e-h)		
h						15(c-a-e-f)	29(c-d-j-i)			
f							22(c-a-e-f-g)			

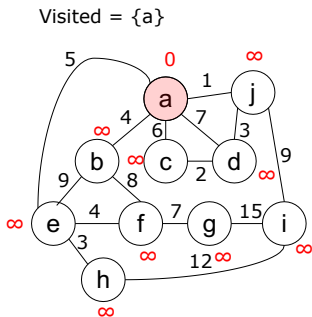
# Dijkstra: Shortest Path Tree



## Shortest Path Tree



# Exercise: From A to G?



a	b(4), c(6), d(7), e(5), j(1)
b	a(4), e(9), f(8)
c	a(6), d(2)
d	a(7), c(2), j(3)
e	a(5), b(9), f(4), h(3)
f	b(8), e(4), g(7)
g	f(7), i(15)
h	e(3), i(12)
i	g(15), h(12), j(9)
j	a(1), d(3), i(9)

## Complexity

Basic Dijkstra is  $\mathcal{O}(|V|^2)$

Can be optimised to  $\mathcal{O}(|E| + |V| \log |V|)$

Graphs can have weighted edges to better represent and model some scenarios.

Dijkstra's algorithm is a shortest path algorithm to traverse weighted edges.