

# COM3020J - Protocols

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# Protocol

- ❑ Human protocols — the rules followed in human interactions
  - Example: Asking a question in class
- ❑ Networking protocols — rules followed in networked communication systems
  - Examples: HTTP, FTP, etc.
- ❑ Security protocol — the (communication) rules followed in a security application
  - Examples: SSL, IPSec, Kerberos, etc.

# Protocols

- ❑ Protocol flaws can be very **subtle**
- ❑ Several well-known security protocols have significant flaws
  - Including WEP, GSM, and IPSec
- ❑ Implementation errors can also occur
  - Recently, IE implementation of SSL
- ❑ Not easy to get protocols right...

# Ideal Security Protocol

- ❑ Must satisfy security requirements
  - Requirements need to be precise
- ❑ Efficient
  - Minimize computational requirement
  - Minimize bandwidth usage, delays...
- ❑ Robust
  - Works when attacker tries to break it
  - Works if environment changes (slightly)
- ❑ Easy to implement, easy to use, flexible...
- ❑ Difficult to satisfy all of these!

# **Simple Security Protocols**

# Secure Entry to NSA

1. Insert badge into reader
2. Enter PIN
3. Correct PIN?
  - Yes?** Enter
  - No?** Get shot by security guard

# ATM Machine Protocol

1. Insert ATM card
2. Enter PIN
3. Correct PIN?
  - Yes?** Conduct your transaction(s)
  - No?** Machine (eventually) eats card

# **Authentication Protocols**



# Authentication

- ❑ Alice must prove her identity to Bob
  - Alice and Bob can be humans or **computers**
- ❑ May also require Bob to prove he's Bob (mutual authentication)
- ❑ Probably need to establish a **session key**
- ❑ May have other requirements, such as
  - Public keys, symmetric keys, hash functions, ...
  - Anonymity, plausible deniability, perfect forward secrecy, etc.

# Authentication

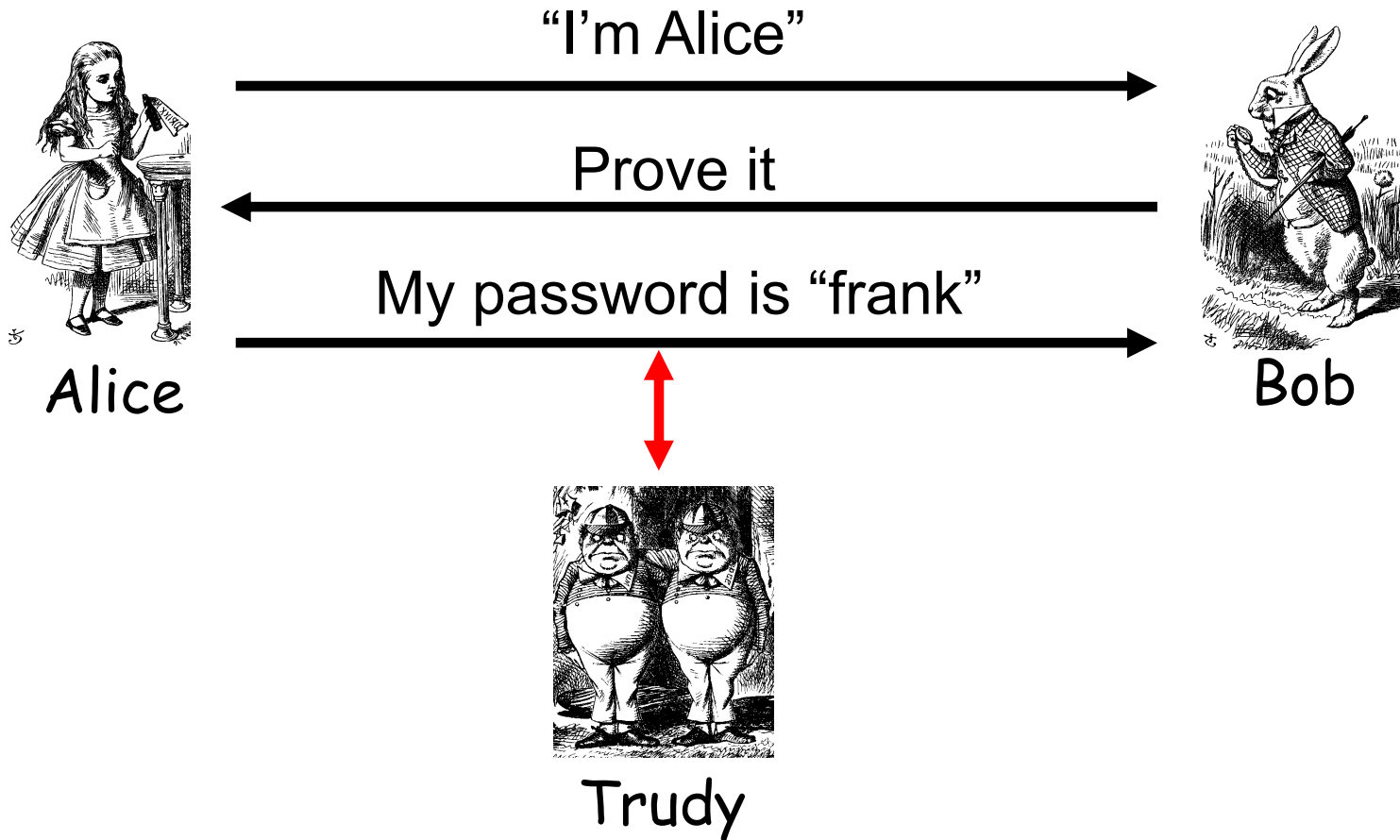
- ❑ Authentication on a stand-alone computer is relatively simple
  - Hash password with salt
  - “Secure path,” attacks on authentication software, keystroke logging, etc., can be issues
- ❑ Authentication over a network is challenging
  - Attacker can passively observe messages
  - Attacker can replay messages
  - Active attacks possible (insert, delete, change)

# Simple Authentication



- ❑ Simple and may be OK for standalone system
- ❑ But highly insecure for networked system
  - Subject to a **replay** attack (next 2 slides)
  - Also, Bob must know Alice's password

# Authentication Attack



# Authentication Attack



- ❑ This is an example of a **replay** attack
- ❑ How can we prevent a replay?

# Simple Authentication



Alice

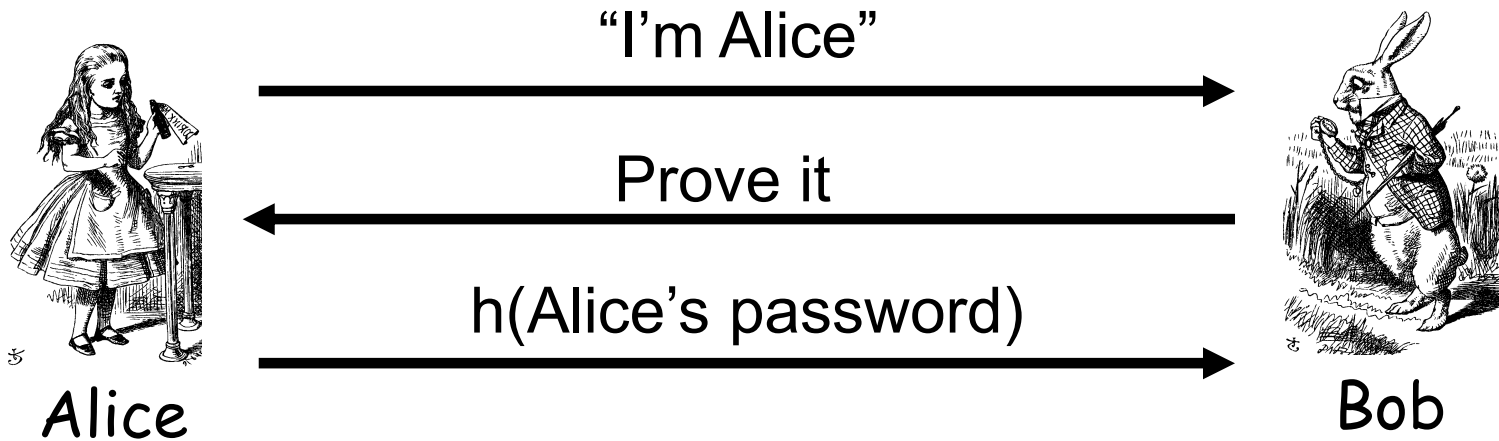
I'm Alice, my password is "frank"



Bob

- ❑ More efficient, but...
- ❑ ... same problem as previous version

# Better Authentication



- ❑ This approach hides Alice's password
  - From both Bob and Trudy
- ❑ But still subject to replay attack

# Challenge-Response

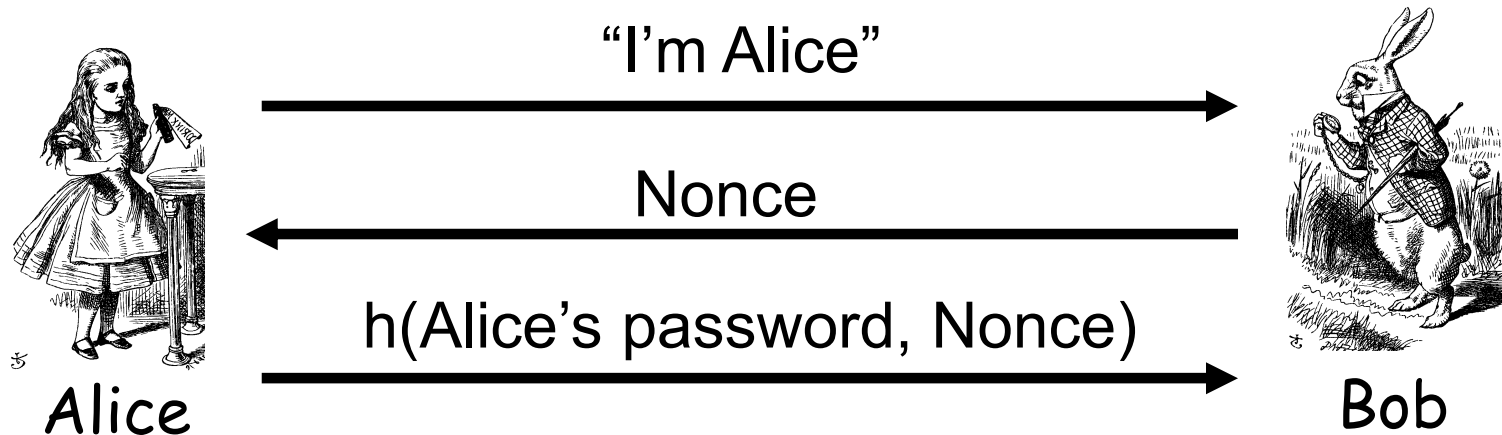
- ❑ To prevent replay, use *challenge-response*
  - Goal is to ensure “freshness”
- ❑ Suppose Bob wants to authenticate Alice
  - *Challenge* sent from Bob to Alice
- ❑ Challenge is chosen so that...
  - Replay is not possible
  - Only Alice can provide the correct *response*
  - Bob can verify the response



# Nonce

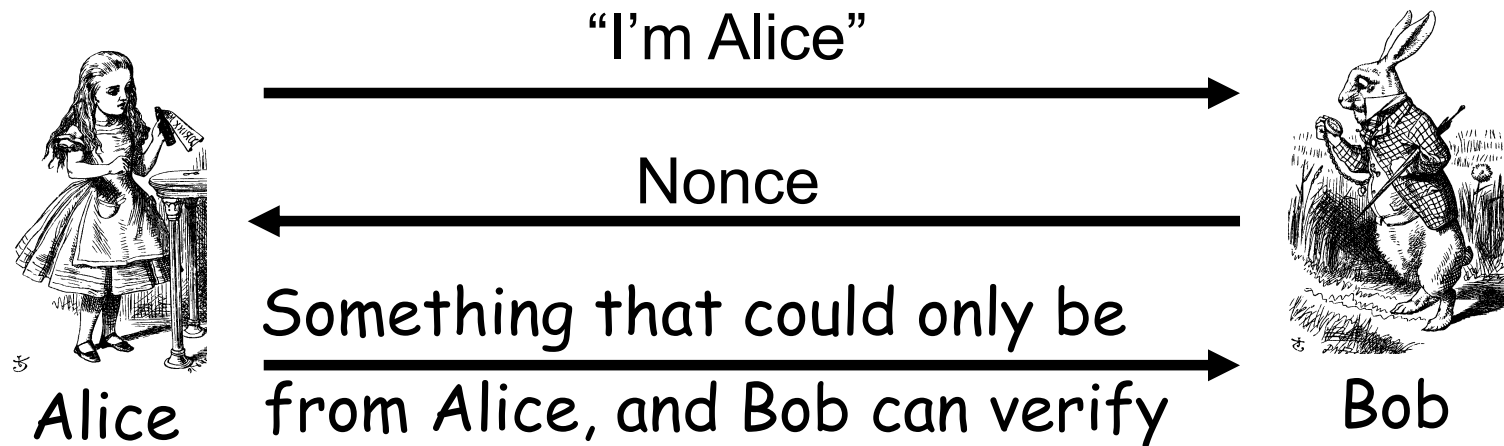
- ❑ To ensure freshness, can employ a **nonce**
  - Nonce == **n**umber used **once**
- ❑ What to use for nonces?
  - That is, what is the challenge?
- ❑ What should Alice do with the nonce?
  - That is, how to compute the response?
- ❑ How can Bob verify the response?
- ❑ Should we use passwords or keys?

# Challenge-Response



- ❑ Nonce is the **challenge**
- ❑ The hash is the **response**
- ❑ Nonce prevents replay (ensures freshness)
- ❑ Password is something Alice knows
- ❑ Note: Bob must know Alice's pwd to verify

# Generic Challenge-Response



- ❑ In practice, how to achieve this?
- ❑ Hashed password works, but...
- ❑ ...encryption is much better here (why?)

# Symmetric Key Notation

- ❑ Encrypt plaintext  $P$  with key  $K$

$$C = E(P, K)$$

- ❑ Decrypt ciphertext  $C$  with key  $K$

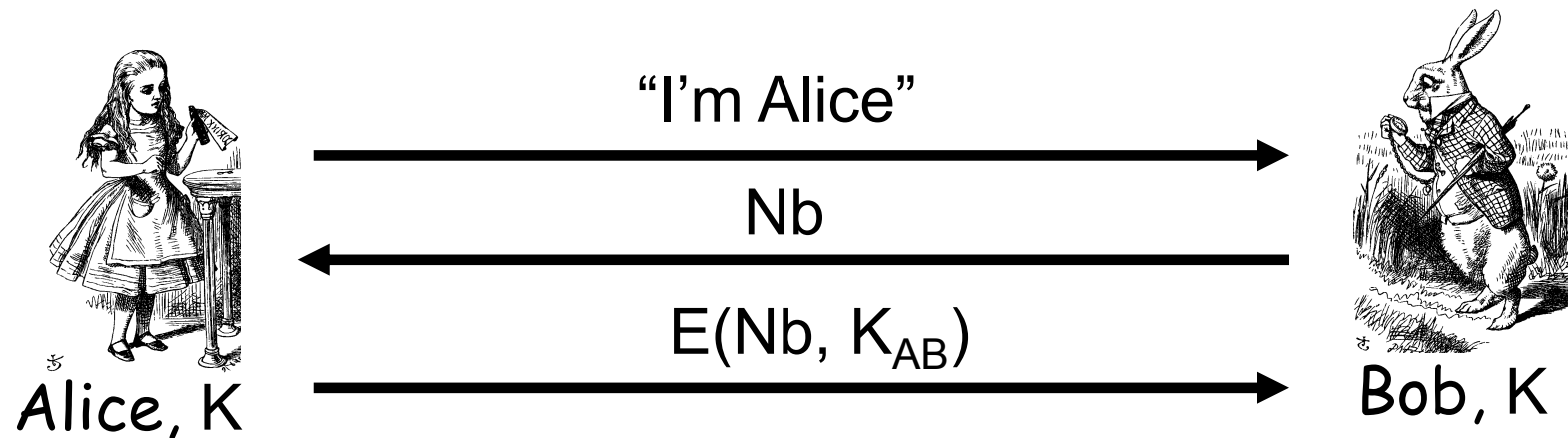
$$P = D(C, K)$$

- ❑ Here, we are concerned with attacks on protocols, **not** attacks on cryptography
  - So, we assume crypto algorithms are secure

# Authentication: Symmetric Key

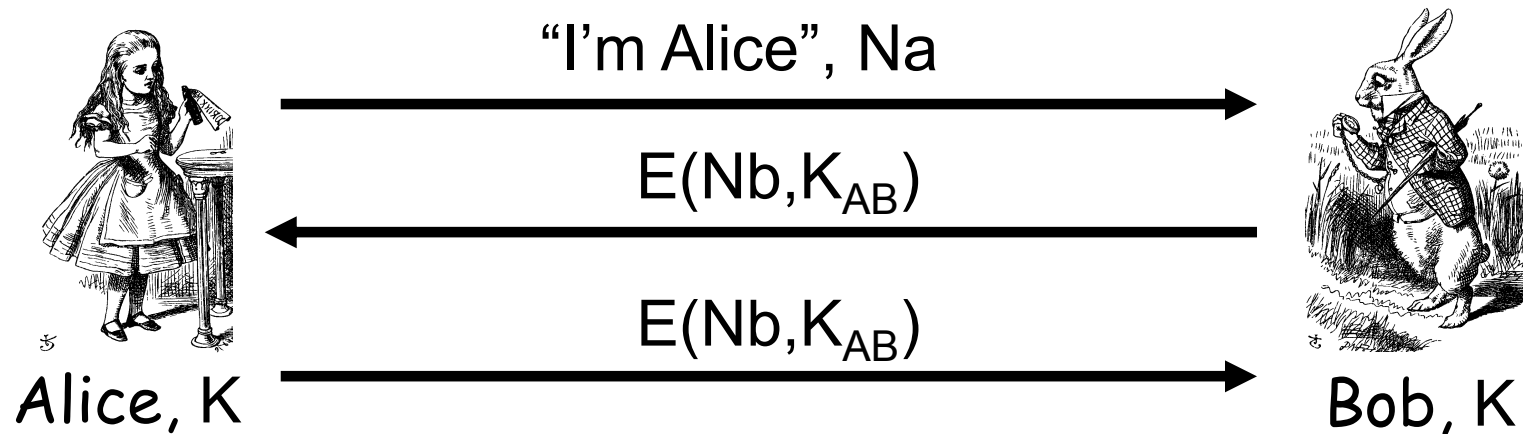
- Alice and Bob share symmetric key  $K_{AB}$
- Key  $K_{AB}$  known only to Alice and Bob
- Authenticate by proving knowledge of shared symmetric key
- How to accomplish this?
  - Cannot reveal key, must not allow replay (or other) attack, must be verifiable, ...

# Authenticate Alice Using Symmetric Key



- ❑ Secure method for Bob to authenticate Alice
- ❑ But, Alice does not authenticate Bob
- ❑ So, can we achieve mutual authentication?

# Mutual Authentication?



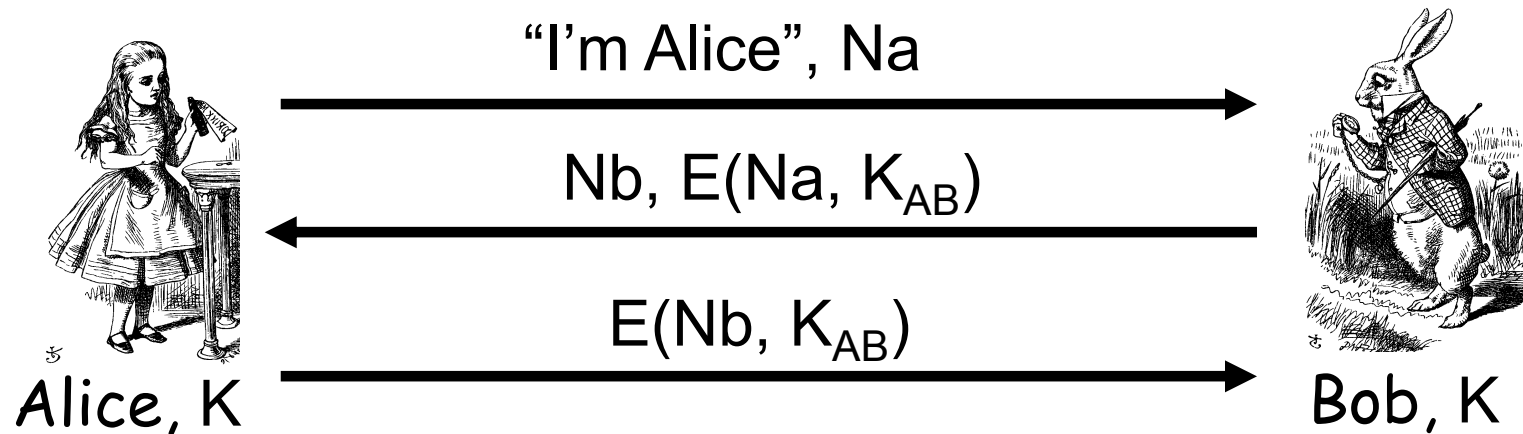
- ❑ What's wrong with this picture?
- ❑ "Alice" could be Trudy (or anybody else)!

# Mutual Authentication

- ❑ Since we have a secure one-way authentication protocol...
- ❑ The obvious thing to do is to use the protocol twice
  - Once for Bob to authenticate Alice
  - Once for Alice to authenticate Bob
- ❑ This has got to work...

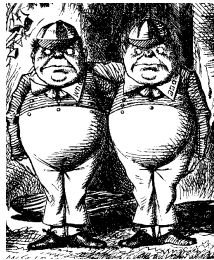


# Mutual Authentication



- ❑ This provides mutual authentication...
- ❑ ...or does it? See the next slide

# Mutual Authentication Attack



Trudy

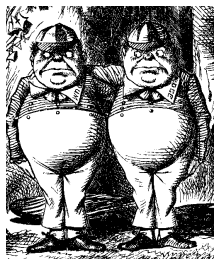
1. "I'm Alice",  $N_a$

2.  $N_b$ ,  $E(N_a, K_{AB})$

5.  $E(N_b, K_{AB})$



Bob,  $K$



Trudy

3. "I'm Alice",  $N_b$

4.  $N_c$ ,  $E(N_b, K_{AB})$



Bob,  $K$