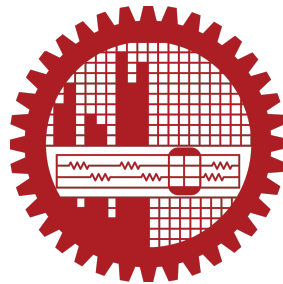


CSE 306_Assignment1

4-Bit Arithmetic and Logic Unit

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1 Design of Arithmetic Unit

Truth Table : Arithmetic Operations

CS_2	CS_1	$CS_0(C_{in})$	X_i	Y_i	F	Arithmetic Operation
0	0	0	A_i	$\overline{B_i}$	$A - B - 1$	Subtract with borrow
0	0	1	A_i	$\overline{B_i}$	$A - B$	Subtract
0	1	0	A_i	1	$A - 1$	Decrement A
0	1	1	A_i	1	A	Transfer A

Subtract with borrow :

$$\begin{aligned}
 &= A - B - 1 \\
 &= (A + \overline{B} + 1) - 1 \\
 &= A + \overline{B} \\
 &= A + (2^n - 1 - B) \\
 &= 2^n + (A - B - 1)
 \end{aligned}$$

Subtract :

$$\begin{aligned}
 &= A - B \\
 &= A + \overline{B} + 1 \\
 &= A + (2^n - 1 - B) + 1 \\
 &= 2^n + (A - B)
 \end{aligned}$$

Decrement A :

$$\begin{aligned}
 &= A - 1 \\
 &= A + 2^n - 1 \\
 &= 2^n + (A - 1)
 \end{aligned}$$

Transfer A :

$$\begin{aligned}
 &= A \\
 &= A + 2^n - 1 + 1 \\
 &= 2^n + A
 \end{aligned}$$

So we have,

$$X_i = A_i$$

$$Y_i = \overline{CS_1} \cdot \overline{B_i} + CS_1 = \overline{(CS_1 + B_i)} + CS_1$$

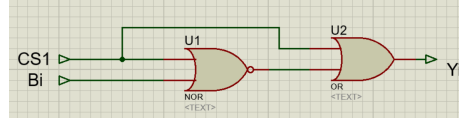


Figure 1: Design of Y_i

2 Design of Logic Unit

Truth Table : Logical Operations

CS_2	CS_1	$CS_0(c_{in})$	X_i	Y_i	$F_i = X_i \oplus Y_i$	Operation
1	0	0	$A_i + \overline{B_i}$	B_i	$A_i.B_i$	AND
1	0	1	$A_i + \overline{B_i}$	B_i	$A_i.B_i$	AND
1	1	0	$\overline{A_i}$	1	$\overline{A_i}$	Complement A
1	1	1	$\overline{A_i}$	1	$\overline{A_i}$	Complement A

Explanation: We can't modify Y_i because that would change the arithmetic operations and neither can omit A_i in any input, So we change X_i .

Let,

$$X_i = A_i + K_i$$

$$F_i = X_i \oplus Y_i$$

$$F_i = (A_i \oplus K_i) \oplus \overline{B_i}$$

$$F_i = (A_i \oplus K_i)B_i + \overline{(A_i \oplus K_i)}.\overline{B_i}$$

$$F_i = A_iB_i + K_iB_i + \overline{A_i}.\overline{K_i}.\overline{B_i}$$

Here the desired operation is A_iB_i . So putting $K_i = \overline{B_i}$

$$F_i = A_iB_i$$

$$F_i = X_i \oplus Y_i$$

$$F_i = X_i \oplus 1$$

$$F_i = \overline{X_i}$$

$$F_i = \overline{A_i + K_i}$$

Here the desired output is $\overline{A_i}$. So putting $K_i = 0$

$$F_i = \overline{A_i}$$

So we need $K_i = \overline{B_i}$ when we will do AND operation and $K_i = 0$ for NOT operation.

We get,

$$X_i = A_i + CS_2 \cdot \overline{CS_1} \cdot \overline{B_i} = A_i + CS_2 \overline{(CS_1 + B_i)}$$

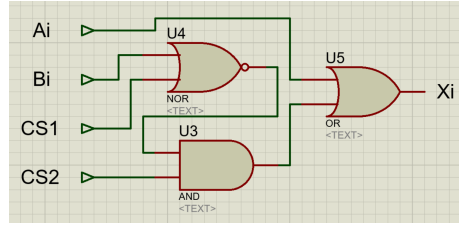


Figure 2: Design of X_i

So finally we have,

$$X_i = A_i + CS_2 \overline{(CS_1 + B_i)}$$

$$Y_i = \overline{(CS_1 + B_i)} + CS_1$$