

CSE - 322

Computer Networks Sessional

NS2 Project Proposal

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Section: A1

DiffQ: Practical Differential Backlog Congestion Control for Wireless Networks

Link to the Paper

[DiffQ: Practical Differential Backlog
Congestion Control for Wireless Networks |
IEEE Conference Publication | IEEE Xplore](#)

Authors: A.Warrier; S. Janakiraman; S. Ha; I. Rhee

Published In: [IEEE INFOCOM 2009](#)

[PDF](#)

Motivations

- Tackle the flow in the middle problem
 - An initial flow gets hindered at a node due to flows generated later on
 - Ensures fairness by proper distribution of the throughput in all of the flows
 - Few flows get disproportionately large amount of bandwidth, many flows are starving
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DiffQ Design

Maintain a queue for each destination whose packet a node forwards

A node delivers a packet or inserts it in FIFO order into the queue of destination

A node also keeps track of the sizes of the destination queues of its neighbouring nodes

Queue is updated twice: if a node receives a packet or it overhears a reception

A queue differential is defined as:
 $QDi(d) = |Qi(d)| - |Qj(d)|$
 $Qi(d)$ is destination queue of d at node i

Set a priority based on this queue differential for the HOL packet of each destination queue

Sender selects the HOL packet of highest priority and transmits it

Pseudocode

Source Rate Control()

1. F = Destination of flow originating at this node
2. $qlen \leftarrow |Q_i(F)|$;
3. if $qlen > \text{QUEUE THRESH}$
 4. $rate = rate/\beta$;
5. else
 6. $rate = rate + \alpha$;

Forwarder Algorithm()

1. $\Delta \leftarrow$ Number of priority levels supported by MAC;
2. $D \leftarrow$ Maximum per-dest queue size;
3. **Flow Scheduling**
4. $F \leftarrow \text{argmax}_d QD_i(d)$;
5. $P \leftarrow$ HOL packet of $Q_i(F)$;
6. $P.\text{priority} \leftarrow \text{MAX}(QD_i(F) - D, 0)$; $P.qlen \leftarrow |Q_i(F)|$;
7. Transmit P ;

8. On receiving packet P from local application

9. Encapsulate P with DiffQ header;
10. if P is the first packet
 11. Create flow entry for P 's destination;
12. $F \leftarrow$ Destination of P ;
13. Enqueue P into $Q_i(F)$;
14. **On reception of packet P from node j**

15. $F \leftarrow$ Destination of P ;
16. if F is this node
 17. Decapsulate DiffQ Header;
 18. Send it up to the application;
19. else
 20. if No flow entry exists for F
 21. Create flow entry for F ;
 22. if node j is the routing next-hop for F
 23. $QD_i(F) \leftarrow |Q_i(F)| - |Q_j(F)|$;
 24. else
 25. Enqueue P into $Q_i(F)$;