

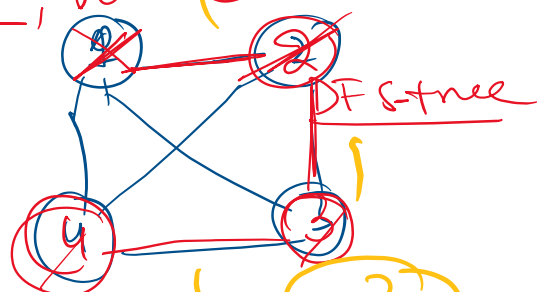
Lecture 25

Minimal Cost Spanning Tree

Spanning Tree: $G=(V, E)$, $T \subseteq G$
undirected/unweighted graph.

Sub graph of G that will contain exactly $|V| - 1$ edges and all V vertices of graph.

1 \rightarrow DFS
2 \rightarrow BFS



Cost, Sum weights of all edges

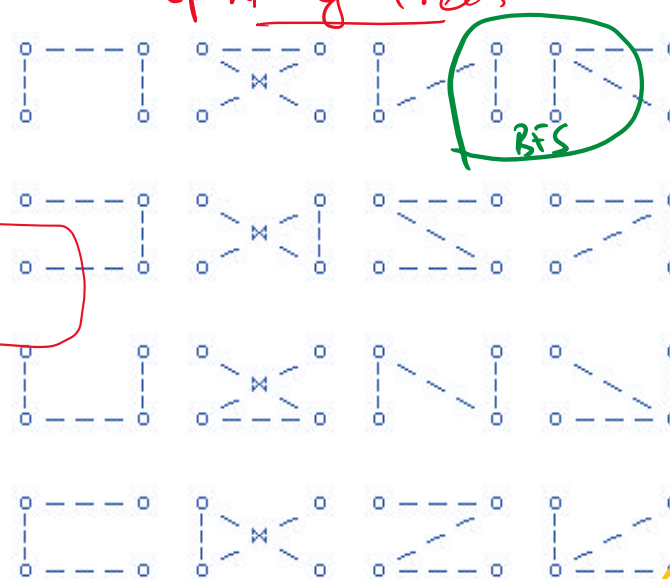
$$\text{Tree Cost} = \sum_{i=1}^{|E|} w(u,v)$$

Cost = 3

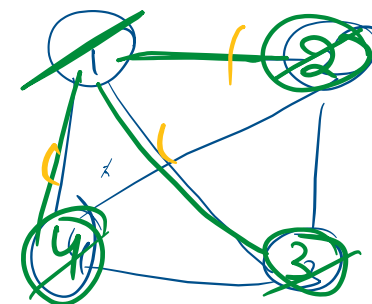
$|V|$ (subset V) $E(V) - 1$ edges

Vertices

Spanning trees



BFS



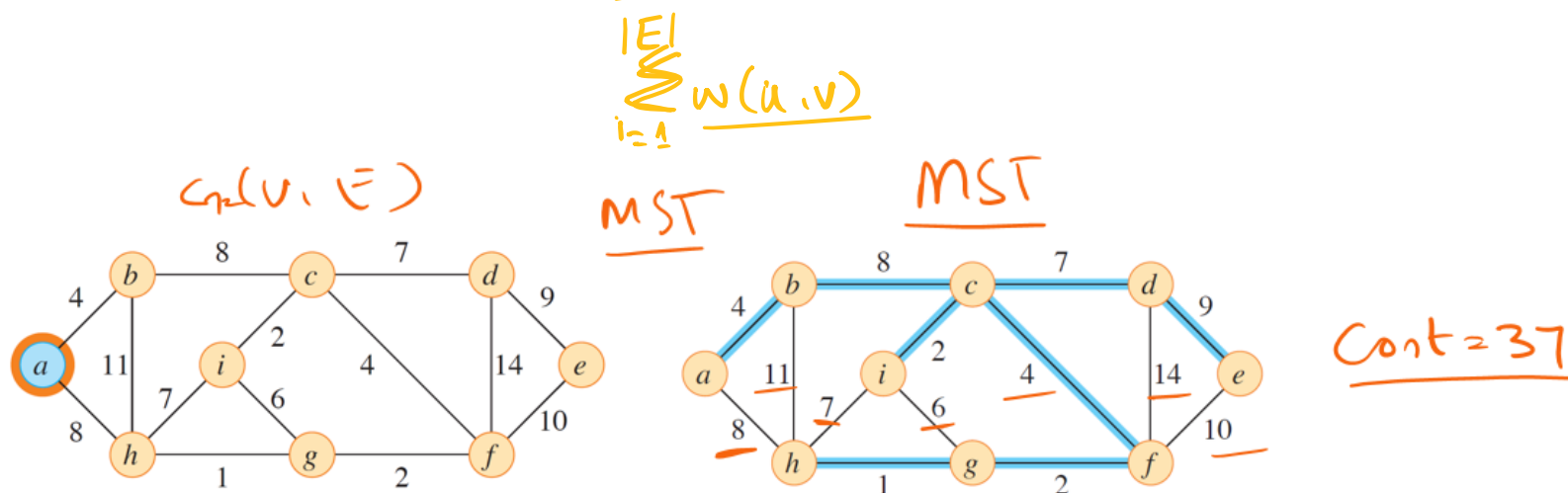
Cost = 3

\rightarrow Number of edges

16

Minimal Cost Spanning Tree (MST) :

Find a tree T of a given graph G that contains all the vertices of G and has the minimum total weight of the edges of G over all other such trees

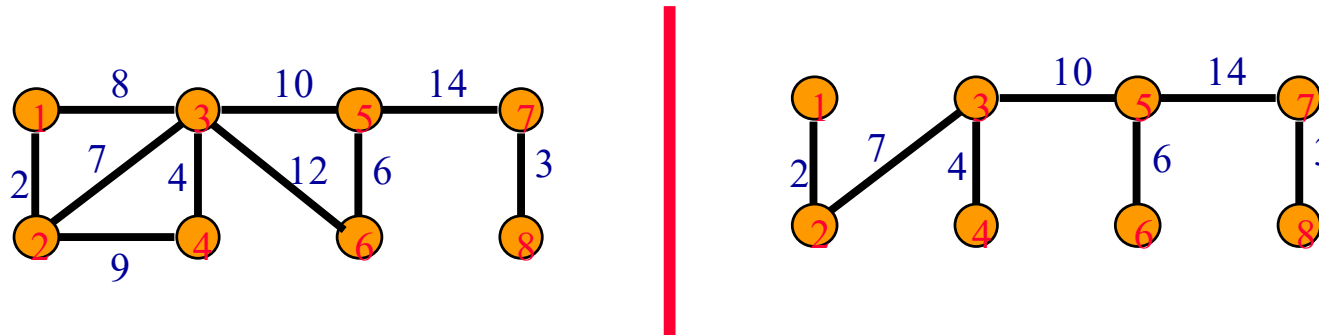
**MST Algorithms:**

1. Prim's Algorithm
2. Kruskal's Algorithm

→ Select edges with Cheapest/minimum cost
+ Not creating the cycle.

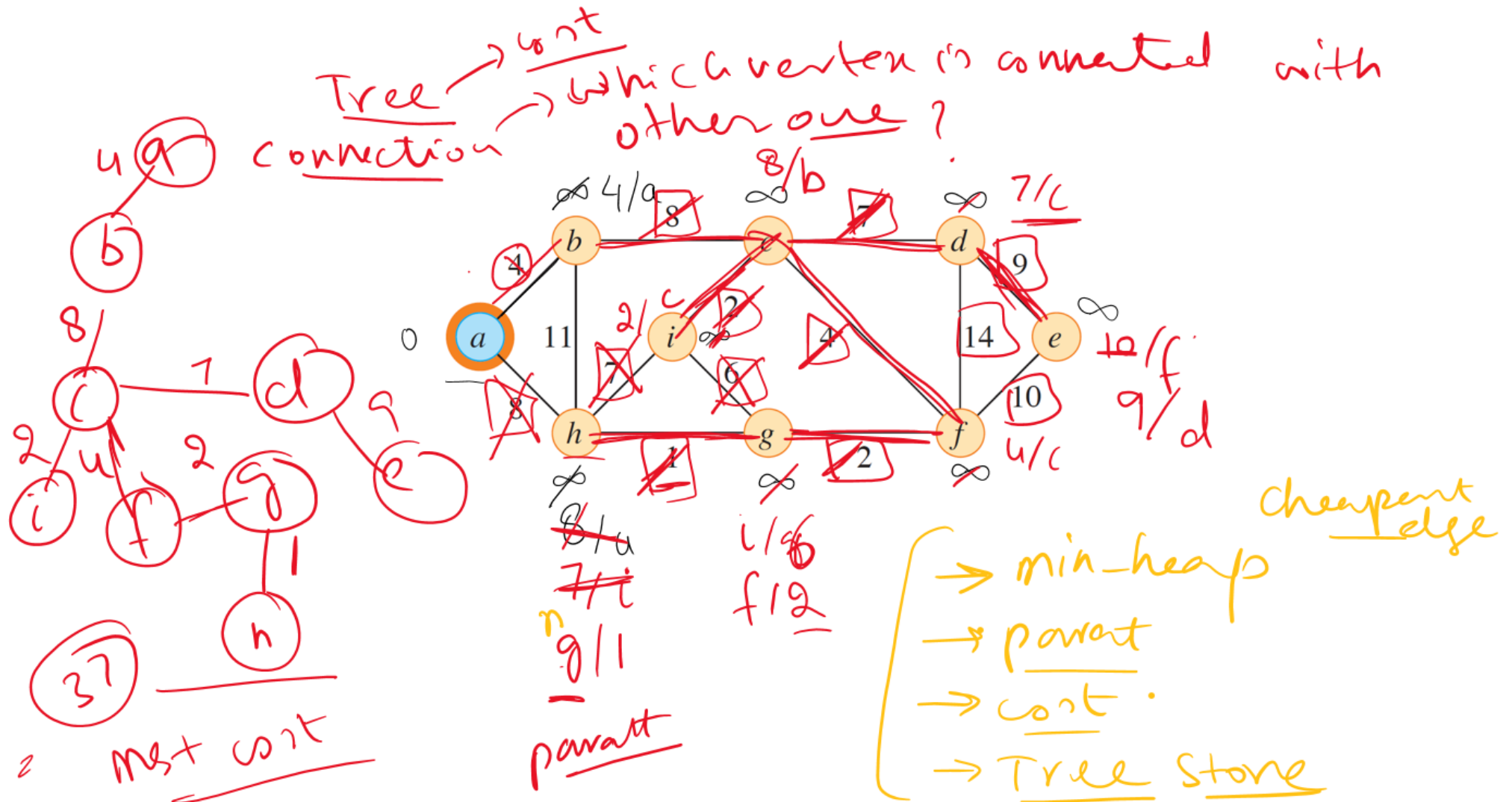
→ Greedy Approach

Prim's Method



- Start with any single vertex tree.
- Get a **2**-vertex tree by adding a cheapest edge.
- Get a **3**-vertex tree by adding a cheapest edge.
- Grow the tree one edge at a time until the tree has **n - 1** edges (and hence has all **n** vertices).

Prim's MST Algorithm



MST-PRIM(G, w, r)

```

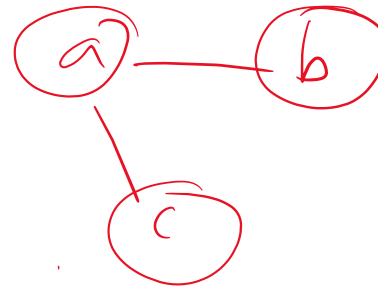
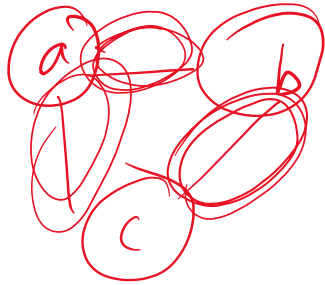
1  for each vertex  $u \in G.V$ 
2       $u.key = \infty$ 
3       $u.\pi = \text{NIL}$ 
4   $r.key = 0$ 
5   $Q = \emptyset$ 
6  for each vertex  $u \in G.V$ 
7      INSERT( $Q, u$ )
8  while  $Q \neq \emptyset$ 
9       $u = \text{EXTRACT-MIN}(Q)$ 
10     for each vertex  $v$  in  $G.Adj[u]$ 
11         if  $v \in Q$  and  $w(u, v) < v.key$ 
12              $v.\pi = u$ 
13              $v.key = w(u, v)$ 
14             DECREASE-KEY( $Q, v, w(u, v)$ )

```

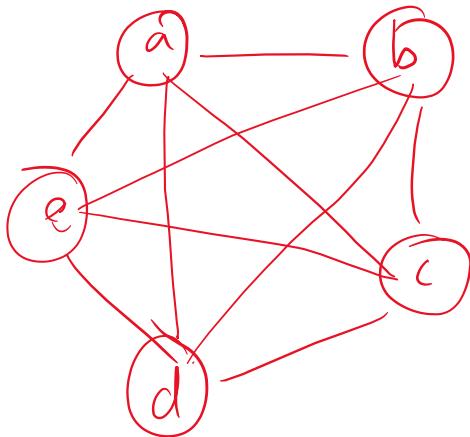
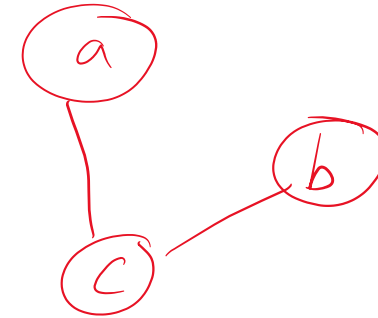
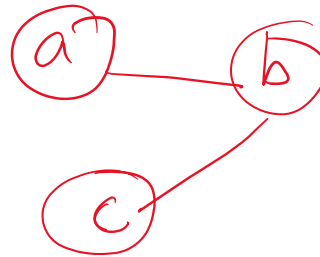
Handwritten annotations and diagrams:

- Red arrows point from $u.key = \infty$ to "cont" (content).
- Red arrows point from $u.\pi = \text{NIL}$ to "parent / connection".
- Red arrow points from $r.key = 0$ to "min heap".
- Red arrow points from the loop starting at line 6 to "insert all vertices".
- Red arrows point from $w(u, v) < v.key$ to "add u to the tree" and "update keys of u 's non-tree neighbors".

Diagram illustrating the MST-PRIM algorithm:



Scale up / 5



2⁴ / Vertices:
4x4
edges 6

