# National University of Computer & Emerging Sciences CS 3001 - COMPUTER NETWORKS

Lecture 9
Chapter 3

19th September, 2023

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Office Hours: 02:30 pm till 06:00 pm (Every Tuesday & Thursday)

# Chapter 3 Transport Layer

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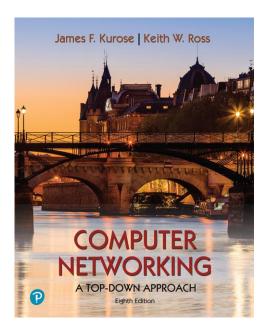
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# Computer Networking: A Top-Down Approach

8<sup>th</sup> edition Jim Kurose, Keith Ross Pearson, 2020

#### Transport layer: overview

#### Our goal:

- understand principles behind transport layer services:
  - multiplexing, demultiplexing
  - reliable data transfer
  - flow control
  - congestion control

- learn about Internet transport layer protocols:
  - UDP: connectionless transport
  - TCP: connection-oriented reliable transport
  - TCP congestion control

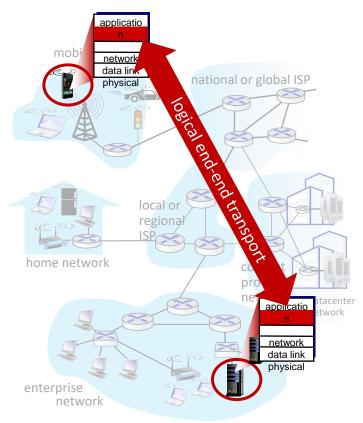
#### Transport layer: roadmap

- Transport-layer services
- Multiplexing and demultiplexing
- Connectionless transport: UDP
- Principles of reliable data transfer
- Connection-oriented transport: TCP
- Principles of congestion control
- TCP congestion control
- Evolution of transport-layer functionality

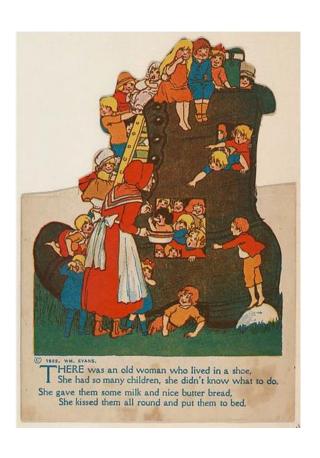


#### Transport services and protocols

- provide logical communication between application processes running on different hosts
- transport protocols actions in end systems:
  - sender: breaks application messages into segments, passes to network layer
  - receiver: reassembles segments into messages, passes to application layer
- two transport protocols available to Internet applications
  - TCP, UDP



#### Transport vs. network layer services and protocols



#### household analogy:

- 12 kids in Ann's house sending letters to 12 kids in Bill's house:
- hosts = houses
- processes = kids
- app messages = letters in envelopes

#### Transport vs. network layer services and protocols

- transport layer: communication between processes
  - relies on, enhances, network layer services
- network layer: communication between hosts

#### household analogy:

- 12 kids in Ann's house sending letters to 12 kids in Bill's house:
- hosts = houses
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- app messages = letters in envelopes

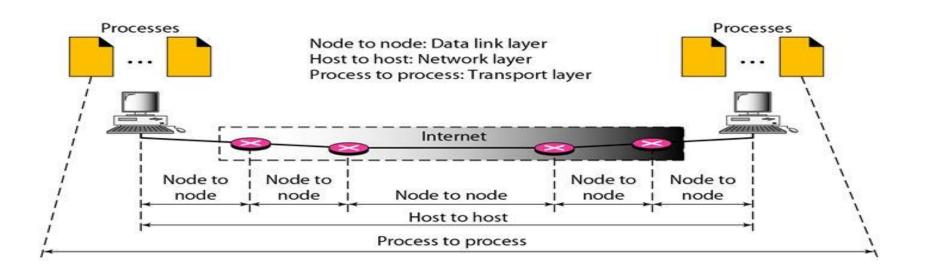
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#### Transport Layer vs Network Layer & Data Link Layer

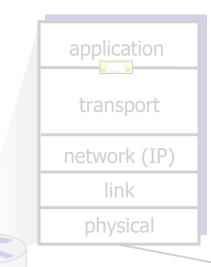
- The Data Link Layer is responsible for delivery of frames between two neighboring nodes over a link. This can be called node-to-node delivery.
- The network layer is responsible for delivery of datagrams between two hosts. This can be called host-to-host delivery.

Communication on the Internet is not defined as the exchange of data between two nodes or between two hosts. Real communication takes place between two processes.

The Transport Layer is responsible for the process-to-process delivery.

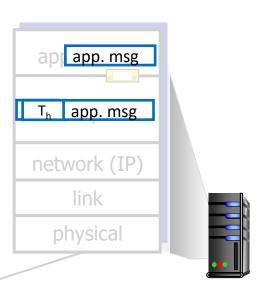


#### **Transport Layer Actions**

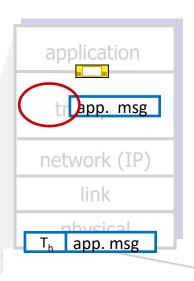


#### Sender:

- is passed an applicationlayer message
- determines segment header fields values
- creates segment
- passes segment to IP

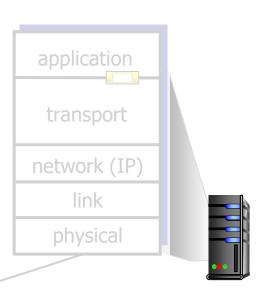


#### **Transport Layer Actions**



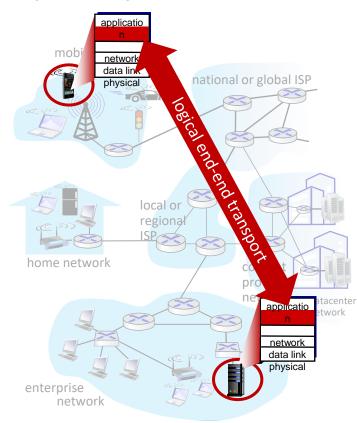
#### Receiver:

- receives segment from IP
- checks header values
- extracts application-layer message
- demultiplexes message up to application via socket



# Two principal Internet transport protocols

- TCP: Transmission Control Protocol
  - reliable, in-order delivery
  - congestion control
  - flow control
  - connection setup
- UDP: User Datagram Protocol
  - unreliable, unordered delivery
  - no-frills extension of "best-effort" IP
- services not available:
  - delay guarantees
  - bandwidth guarantees

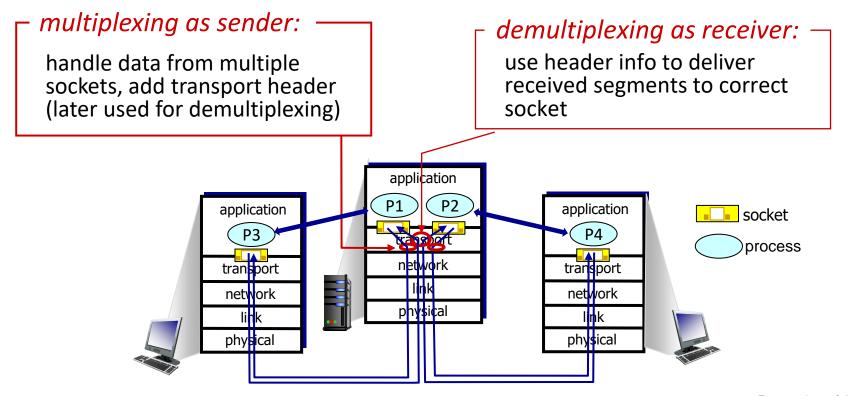


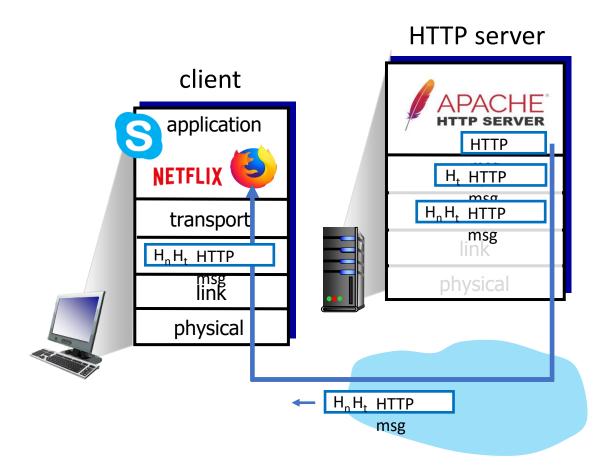
#### Chapter 3: roadmap

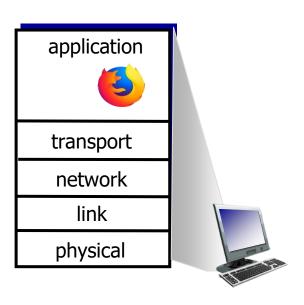
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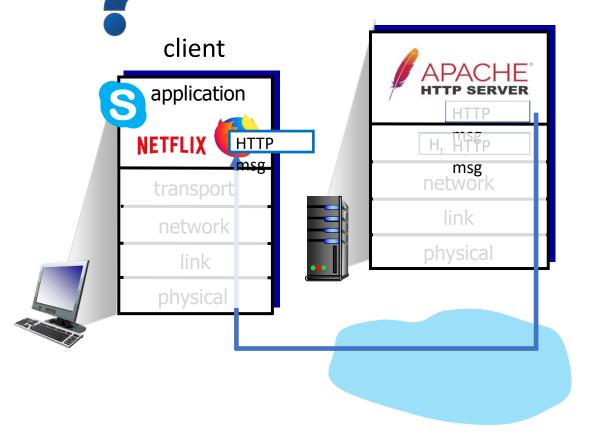
## Multiplexing/demultiplexing

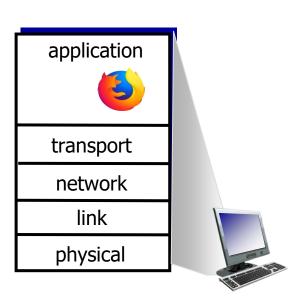


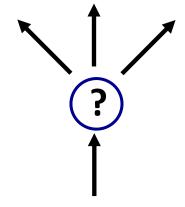




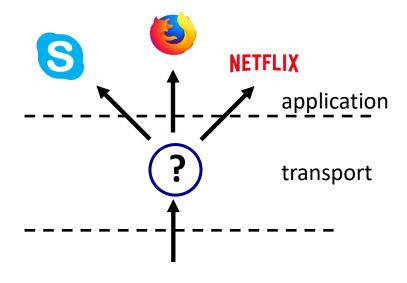
# Q: how did transport layer know to deliver message to Firefox browser process rather then Netflix process or Skype process?







de-multiplexing

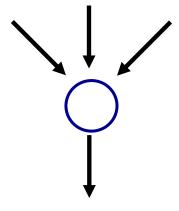


de-multiplexing

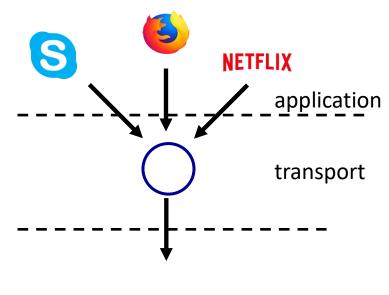








multiplexing



multiplexing



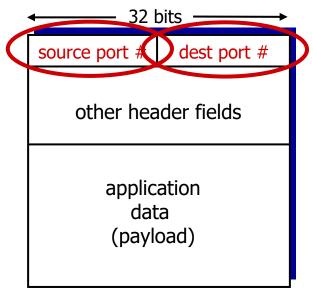
#### Port

- Port is a transport layer address / identifier to choose among the multiple processes running among the same host
- 16 bit integers ranging between 0 & 65,535

- The client side program chooses a port number randomly (ephemeral port number)
- The server side ports are not random but well known and preassigned
- IANA (Internet Assigned Number Authority) had divided port numbers into three ranges:
  - Well Known: 0 to 1023. Assigned & controlled by IANA. (Example HTTP uses Port # 80, FTP uses port # 21. List of well known ports given in RFC 1700 & updated via RFC 3232)
  - Registered: 1024 to 49,151. Not assigned / controlled by IANA, but registered with them to avoid duplication.
  - Dynamic (or Private): 49,152 to 65,535. Neither controlled, nor registered. Can be used by any process. (Ephemeral Ports.)

### How demultiplexing works

- host receives IP datagrams
  - each datagram has source IP address, destination IP address
  - each datagram carries one transport-layer segment
  - each segment has source, destination port number
- host uses IP addresses & port numbers to direct segment to appropriate socket



TCP/UDP segment format

#### Connectionless demultiplexing

#### Recall:

when creating socket, must specify host-local port #:

```
DatagramSocket mySocketI
= new
DatagramSocket(12534);
```

- when creating datagram to send into UDP socket, must specify
  - destination IP address
  - destination port #

when receiving host receives *UDP* segment:

- checks destination port # in segment
- directs UDP segment to socket with that port #



IP/UDP datagrams with same dest. port #, but different source IP addresses and/or source port numbers will be directed to same socket at receiving host

(thus UDP socket identified by 2-tuple:

- dest IP address
- dest port number)

#### Connectionless demultiplexing: an example

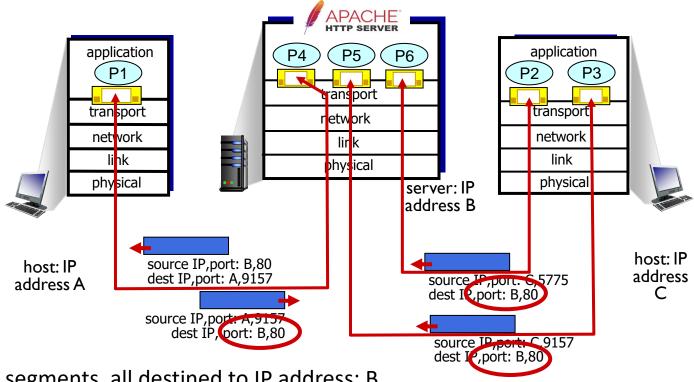
```
mySocket =
                                socket(AF INET,SOCK DGRAM)
                              mySocket.bind(myaddr,6428);
mySocket =
                                                                  mySocket =
 socket(AF INET, SOCK STREAM)
                                                                    socket(AF INET,SOCK STREAM)
mySocket.bind(myaddr, 9157);
                                                                  mySocket.bind(myaddr,5775);
                                            application
                                                                           application
              application
                 P3
                                             transport
              transport
                                                                           transport
                                              network
               network
                                                                            network
                                               link
                 lirk
                                                                              lihk
                                              physical
                                                                            physical
               physical
                              source port: 6428
                                                            source port: ?
                              dest port: 9157
                                                              dest port: ?
                                                      source port: ?
               source port: 9157
                                                      dest port: ?
                 dest port: 6428
```

### Connection-oriented demultiplexing

- TCP socket identified by 4-tuple:
  - source IP address
  - source port number
  - dest IP address
  - dest port number
- demux: receiver uses all four values (4-tuple) to direct segment to appropriate socket

- server may support many simultaneous TCP sockets:
  - each socket identified by its own 4-tuple
  - each socket associated with a different connecting client

#### Connection-oriented demultiplexing: example



Three segments, all destined to IP address: B, dest port: 80 are demultiplexed to *different* sockets

# Summary

- Multiplexing, demultiplexing: based on segment, datagram header field values
- UDP: demultiplexing using destination port number (only)
- TCP: demultiplexing using 4-tuple: source and destination IP addresses, and port numbers
- Multiplexing/demultiplexing happen at all layers

# Assignement # 2 (Chapter - 2)

- 2<sup>nd</sup> Assignment already uploaded in Google Classroom Stream Section
- Due Date: Thursday, 21<sup>st</sup> September, 2023 (During the lecture)
- Hard copy of the handwritten assignment to be submitted directly to the Instructor during the lecture.
- Please read all the instructions carefully in the uploaded Assignment document, follow & submit accordingly

# Quiz # 2 (Chapter - 2)

- Quiz # 2 for Chapter 2 to be taken in the class on Thursday, 21st September, 2023 during the lecture time
- Quiz to be take for **OWN Section** only

No Retake

Be on time