

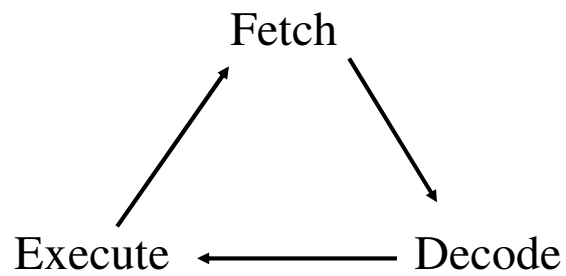
ECE 344

Microprocessors

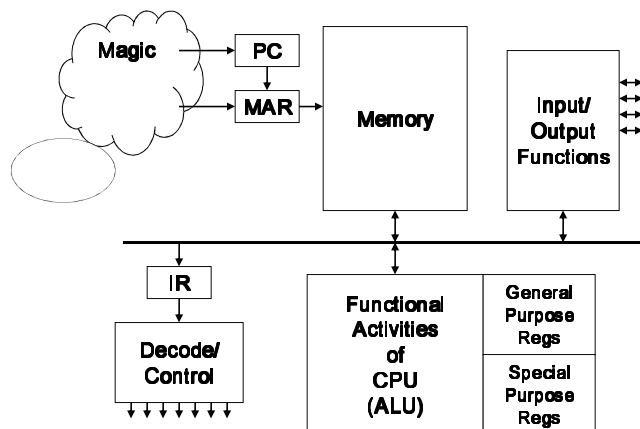
How to Do Work in Processor

- Instruction set defines what's possible
- Work accomplished by repeating instructions on data until desired result has been obtained

Basic Premise of All Computers



Basic Computer Organization



Computer Architecture

- Work is done in computers by moving information from one place to another, sometimes with a transformation
 - Register to register
 - Register to memory
 - Memory to register
 - Memory to memory

Computer Architecture

- The instruction set identifies the work (movement, transformation) that can be done in a particular architecture
- Any work to be done in a machine must be accomplished by a series of individual instructions
- Complexity of instruction set determines how much work can be done by an instruction

Complex Instruction Set Computer : CISC

- Instructions complex to allow for extended work
- Memory to memory activities allowed
- Instructions may take multiple words to store
- Complexity may limit pipelining
- Hooks to help with data structures

Reduced Instruction Set Computer : RISC

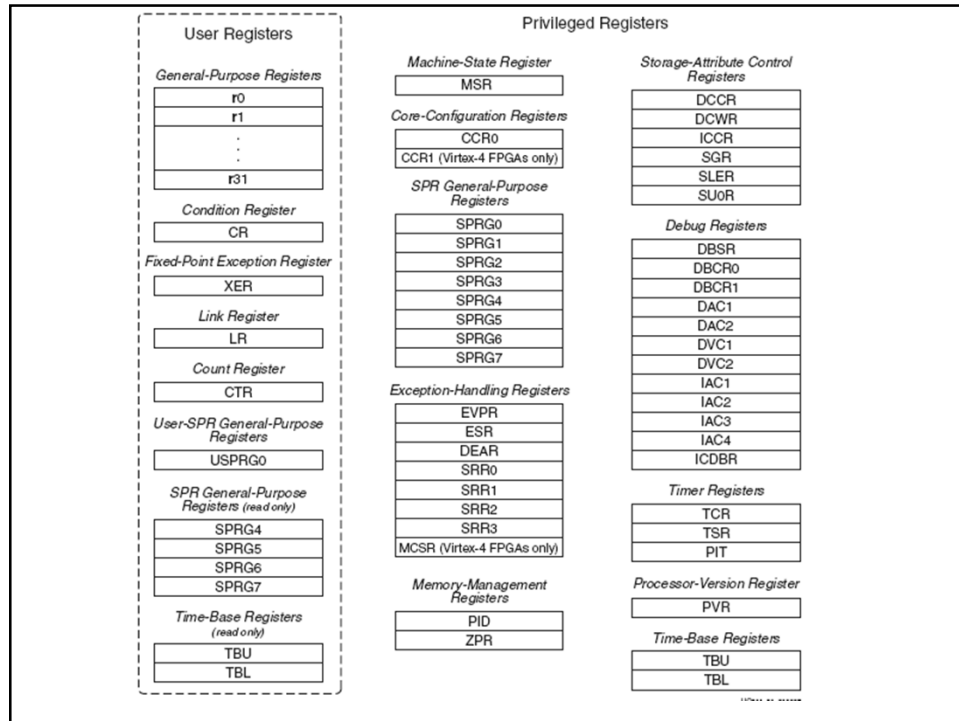
- Minimal number of instructions and addressing modes
- Fixed instruction format
- Hardwired instruction decoding
- Single cycle execution
- Load/store organization (work in registers)
- Use of register windows

Common Registers in Processors

- Program counter (PC) – where is program executing
- Memory Address Register (MAR) – location of interest in memory
- Accumulator (ACC) – for single address machine
- General purpose registers (Rn) – for data and addresses

Common Registers in Processors

- Address registers (An) – point to location in memory
- Instruction Register (IR) – holds instruction to be executed
- Status Register (SR) – holds information about status of system
- System Control Registers – hold information about overall system operation



Processor Function in Register Form

Example: add r10, r11, r12

fetch: $MAR \leftarrow PC$
 $IR \leftarrow M[MAR]$
 $PC \leftarrow PC + ??$

- decode -

execute: $r10 \leftarrow r11 + r12$

Instruction Types

- Work: add, subtract, mult, div, AND, OR, EXOR
- Movement: reg to reg, reg to mem, mem to reg
- Program control: jump, jump to subroutine, return from subroutine, conditional jump
- System control: set control bits, modify control registers, interrupt system

System Memory

- Used to hold Instructions, Data, OS info
- Organized by byte (8 bits), halfword (16 bits), or word (32 bits)
- Instructions can move bytes, halfwords, or words to or from memory
- Registers identify location in memory (either PC or Rn)

Memory Organization - Bytes

	00	01	02	03	04	05	06	07
0000	00	01	02	03	04	05	06	07
0008	08	09	0A	0B	0C	0D	0E	0F
0010	10	11	12	13	14	15	16	17
0018	18	19	1A	1B	1C	1D	1E	1F

Memory Organization - Halfwords

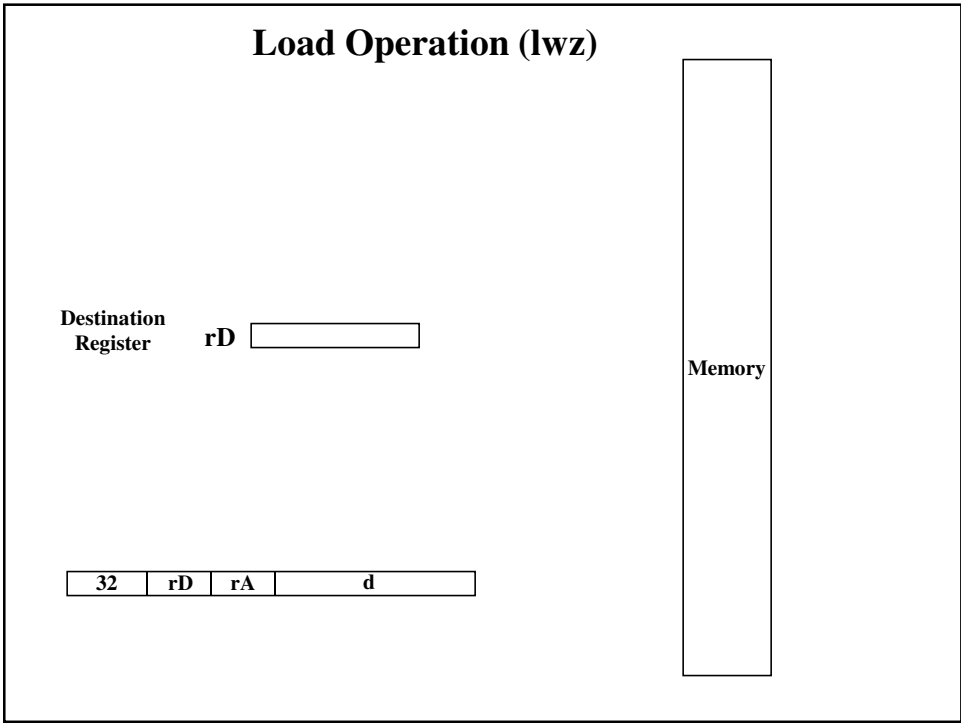
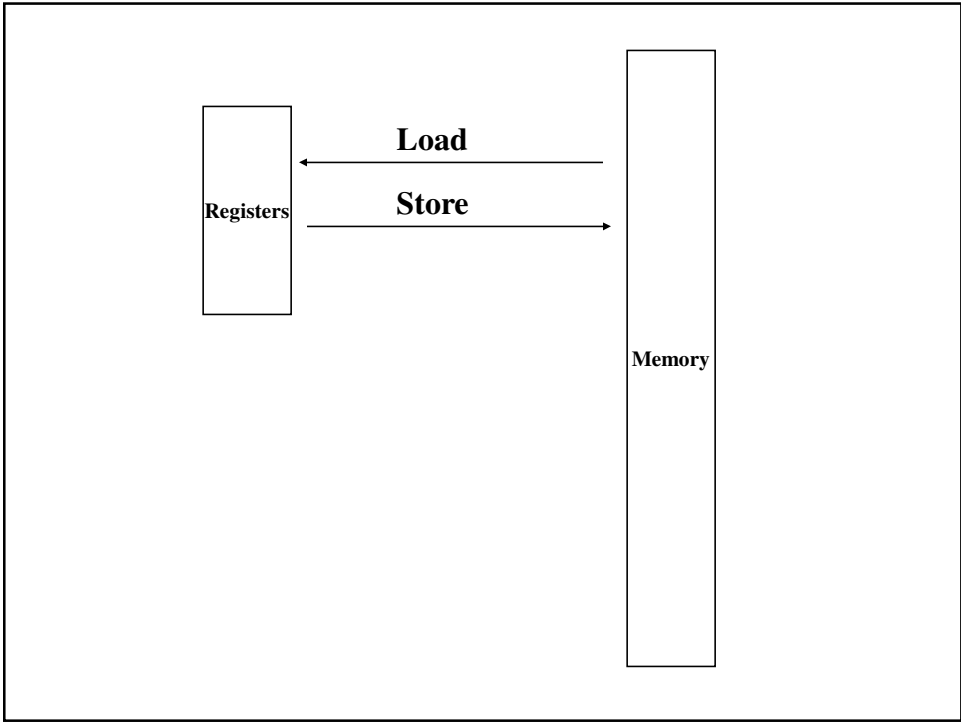
	00	02	04	06
0000	0001	0203	0405	0607
0008	0809	0A0B	0C0D	0E0F
0010	1011	1213	1415	1617
0018	1819	1A1B	1C1D	1E1F

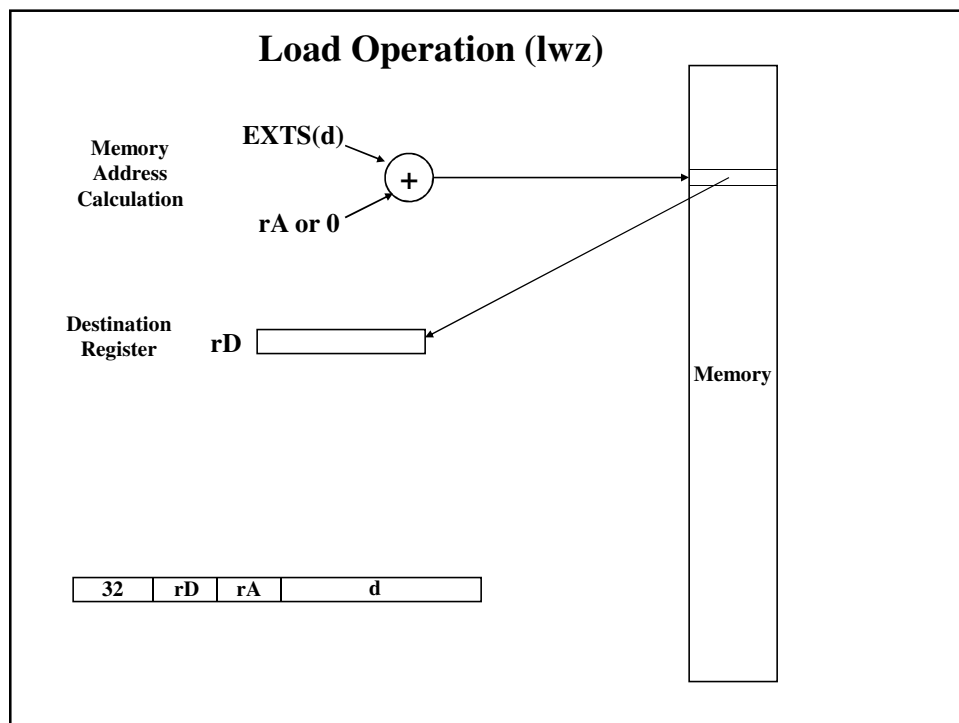
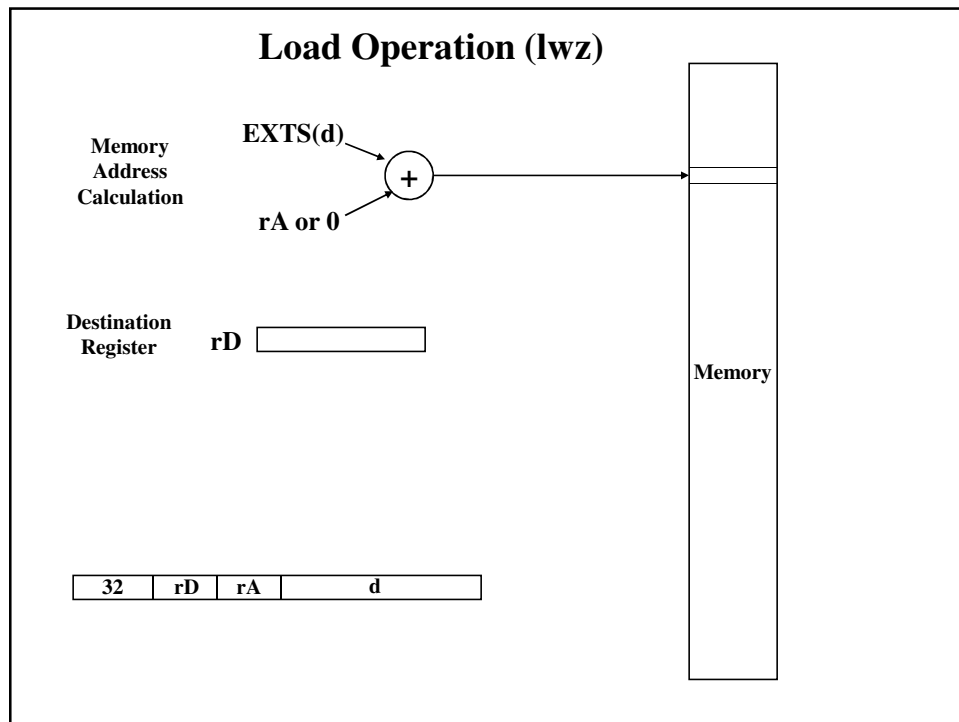
Memory Organization - Words

	00	04
0000	00010203	04050607
0008	08090A0B	0C0D0E0F
0010	10111213	14151617
0018	18191A1B	1C1D1E1F

Load & Store: Use Register as Pointer

- Goal: move data to (load) or from (store) register file
- Specify one register as pointer to memory
- Specify another register as target
- Specify offset in terms of bytes





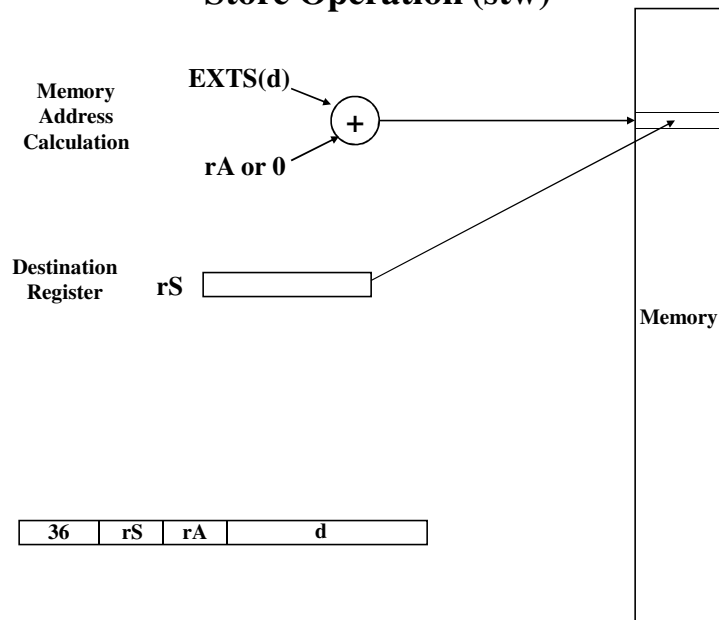
Load Operation (lwz)

Fetch: $MAR \leftarrow PC$
 $IR \leftarrow M[MAR]$
 $PC \leftarrow PC + 4$

Decode:

Execute: $EA \leftarrow (rA \text{ or zero}) + EXT(d)$
 $rD \leftarrow M[EA]$

Store Operation (stw)



Subroutine call: Program Control plus Data Movement

- Program Control: Determine target and transfer control to that location (BL, others)
- Subroutine linkage: return address → Link Register
- Parameter passing: agreement between calling and called routine
- Simple parameter passing: agree on register usage
- More complex parameter passing: stack

Software Stuff

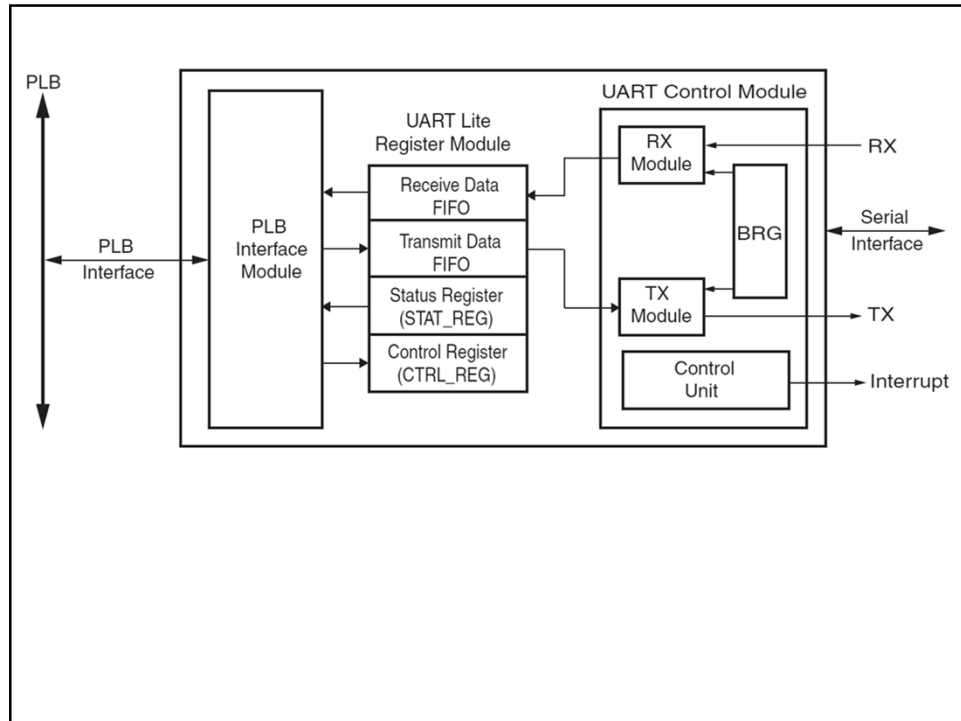
- Array implementation and utilization
 - One dimension – single index
 - Multiple dimension – calculate offset, then single index
- Other data structures
 - Stack
 - Queue
 - List
- Loops and conditional execution

Memory Mapped I/O

- I/O elements assigned addresses in memory space of processor
- I/O activity carried out by writing/reading locations in memory address space
- Work can be accomplished by bits, bytes, halfwords, words
- Understanding of I/O function imperative

UART: Serial Transmission

- Universal Serial Receiver/Transmitter (UART)
- Functionality of serial activity: specification and activity
- Interrupt activities for variety of events



XPS UART Lite Register Descriptions

Table 4 shows all the XPS UART Lite registers and their addresses.

Table 4: XPS UART Lite Registers

Base Address + Offset (hex)	Register Name	Access Type	Default Value (hex)	Description
C_BASEADDR + 0x0	Rx FIFO	Read	0x0	Receive Data FIFO
C_BASEADDR + 0x4	Tx FIFO	Write	0x0	Transmit Data FIFO
C_BASEADDR + 0x8	STAT_REG	Read	0x4	UART Lite Status Register
C_BASEADDR + 0xC	CTRL_REG	Write	0x0	UART Lite Control Register

Timer Module

- Two different modes – capture and generate
- One or two timers per module
- Three registers per counter
 - Counter itself
 - Load register
 - Control/Status register
- Interrupt capability

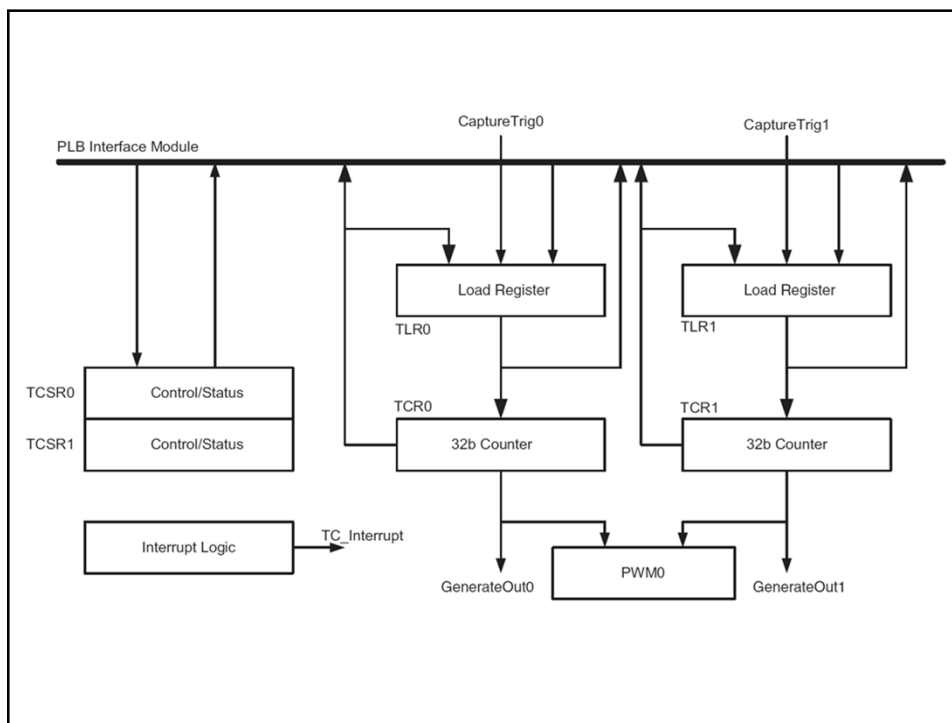


Table 4: XPS Timer/Counter Register Address Map

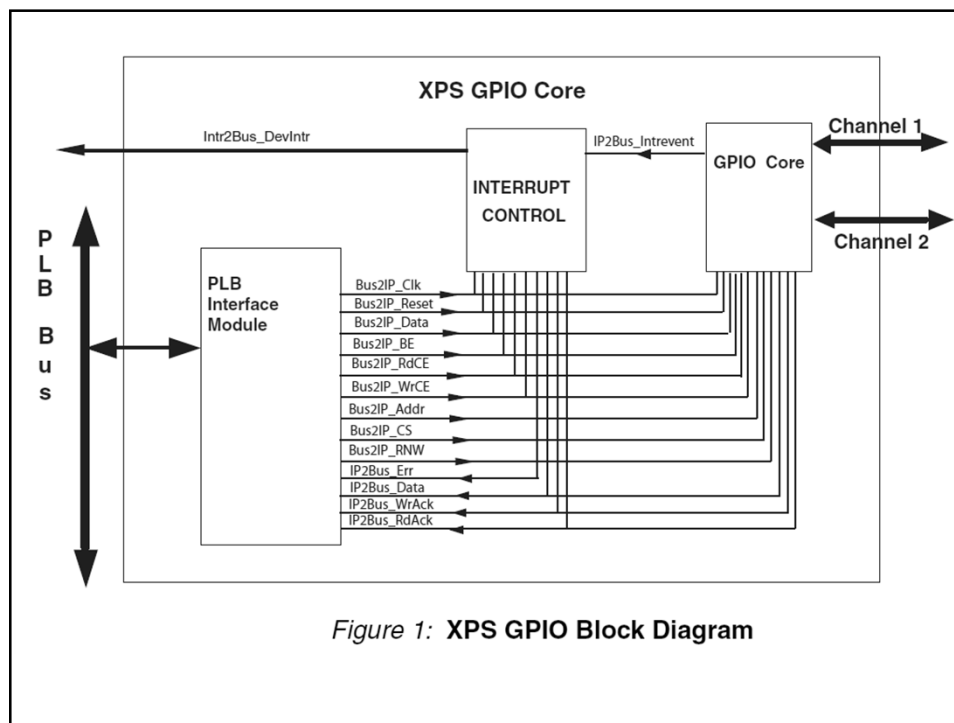
Register	Address (Hex)	Size	Type	Description
TCSR0	C_BASEADDR + 0x00	Word	Read/Write	Control/Status Register 0
TLR0	C_BASEADDR + 0x04	Word	Read/Write	Load Register 0
TCR0	C_BASEADDR + 0x08	Word	Read	Timer/Counter Register 0
TCSR1	C_BASEADDR + 0x10	Word	Read/Write	Control/Status Register 1
TLR1	C_BASEADDR + 0x14	Word	Read/Write	Load Register 1
TCR1	C_BASEADDR + 0x18	Word	Read	Timer/Counter Register 1

Pulse Width Modulation (PWM)

- Variable amount of power delivered to pin based on register settings
- Period: how often pulse happens
- Width: how much of period signal asserted
- Control: what pins used for what, when

GPIO Module

- General Purpose I/O
 - In, Out, or In & Out
- Can have one or two interfaces per module
- Can use with interrupt



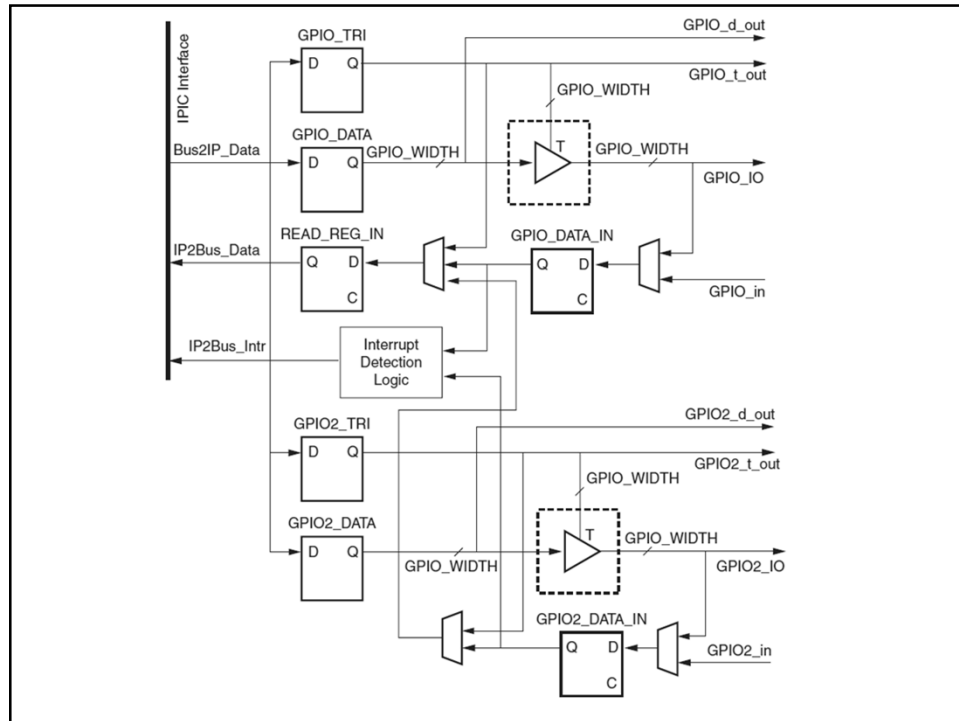


Table 4: XPS GPIO Registers

Register Name	Description	PLB Address	Access
GPIO_DATA	Channel 1 XPS GPIO Data Register	C_BASEADDR + 0x00	Read/Write
GPIO_TRI	Channel 1 XPS GPIO 3-state Register	C_BASEADDR + 0x04	Read/Write
GPIO2_DATA	Channel 2 XPS GPIO Data register	C_BASEADDR + 0x08	Read/Write
GPIO2_TRI	Channel 2 XPS GPIO 3-state Register	C_BASEADDR + 0x0C	Read/Write

Table 8: XPS GPIO Interrupt Registers

Register Name	Description	PLB Address	Access
GIE	Global Interrupt Enable Register	C_BASEADDR + 0x11C	Read/Write
IP IER	IP Interrupt Enable Register	C_BASEADDR + 0x128	Read/Write
IP ISR	IP Interrupt Status Register	C_BASEADDR + 0x120	Read/TOW ^[1]

Interrupt Controller - INTC

- Coordinates interrupt requests from modules
- Can handle up to 32 interrupt sources
- Can cascade if necessary for more than 32
- Controlled with programmed I/O techniques
- Can detect edges or levels

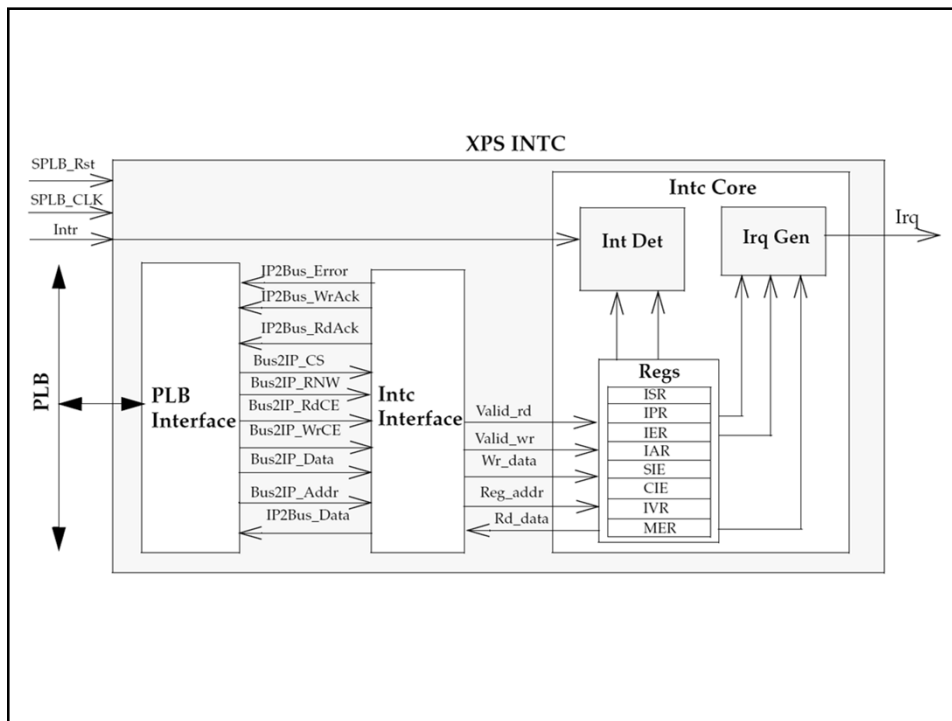


Table 4: XPS INTC Registers and Base Address Offsets

Register Name	Base Address + Offset (Hex)	Access Type	Abbreviation	Reset Value
Interrupt Status Register	C_BASEADDR + 0x0	Read / Write	ISR	All Zeros
Interrupt Pending Register	C_BASEADDR + 0x4	Read only	IPR	All Zeros
Interrupt Enable Register	C_BASEADDR + 0x8	Read / Write	IER	All Zeros
Interrupt Acknowledge Register	C_BASEADDR + 0xC	Write only	IAR	All Zeros
Set Interrupt Enable Bits	C_BASEADDR + 0x10	Write only	SIE	All Zeros
Clear Interrupt Enable Bits	C_BASEADDR + 0x14	Write only	CIE	All Zeros
Interrupt Vector Register	C_BASEADDR + 0x18	Read only	IVR	All Ones
Master Enable Register	C_BASEADDR + 0x1C	Read / Write	MER	All Zeros

PIT – Programmable Interval Timer

- User specifies period for activity by placing count in register
- Counter loaded from register – decrements to zero
- Interrupt can occur when count reaches zero
- Control of PIT: three registers (PIT, TCR, TSR)

Fixed Interval Timer - FIT

- Four periods to choose from
- Enabled/Disabled with special purpose registers

Information Representation

- Integer representation – whole numbers
 - Unsigned binary
 - 2's complement
 - Excess codes
- Floating point representation – IEEE 32 bit normalized numbers
- Other: instruction, address, parity, etc.

Unsigned Binary

$$Value = \sum_{i=0}^{n-1} b_i \times 2^i$$

$$11110000 = 240$$

$$1111000011110000 = 61,680$$

Unsigned Binary Patterns

$$0000\ 0001 = 1$$

$$0000\ 0010 = 2$$

$$0000\ 0100 = 4$$

$$0000\ 1000 = 8$$

$$0000\ 1010 = 10$$

$$0001\ 0000 = 16$$

$$0001\ 1010 = 16 + 10 = 26$$

Unsigned Binary, Fixed Point

$$Value = \sum_{i=0}^{n-1} b_i \times 2^i \times 2^{-p}$$

$$1111.0000 = 15$$

$$11110000.11110000 = 240.9375$$

$$1111.000011110000 = 15.05859375$$

$$1.111000011110000 = 1.88232421875$$

Two's Complement

$$Value = -b_{n-1} \times 2^{n-1} + \sum_{i=0}^{n-2} b_i \times 2^i$$

$$11110000 = -16$$

$$1111000011110000 = -3,856$$

Two's Complement Patterns

$$0000\ 0001 = 1$$

$$0000\ 0010 = 2$$

$$0000\ 0100 = 4$$

$$0000\ 1000 = 8$$

$$0000\ 1010 = 10$$

$$0001\ 0000 = 16$$

$$0001\ 1010 = 16 + 10 = 26$$

Two's Complement Patterns

$$1000\ 0001 = 1 + -128 = -127$$

$$1000\ 0010 = 2 + -128 = -126$$

$$1000\ 0100 = 4 + -128 = -124$$

$$1000\ 1000 = 8 + -128 = -120$$

$$1000\ 1010 = 10 + -128 = -118$$

$$1001\ 0000 = 16 + -128 = -112$$

$$1001\ 1010 = 16 + 10 + -128 = -102$$

Two's Complement, Fixed Point

$$Value = (-b_{n-1} \times 2^{n-1} + \sum_{i=0}^{n-2} b_i \times 2^i) \times 2^{-p}$$

$$1111.0000 = -1$$

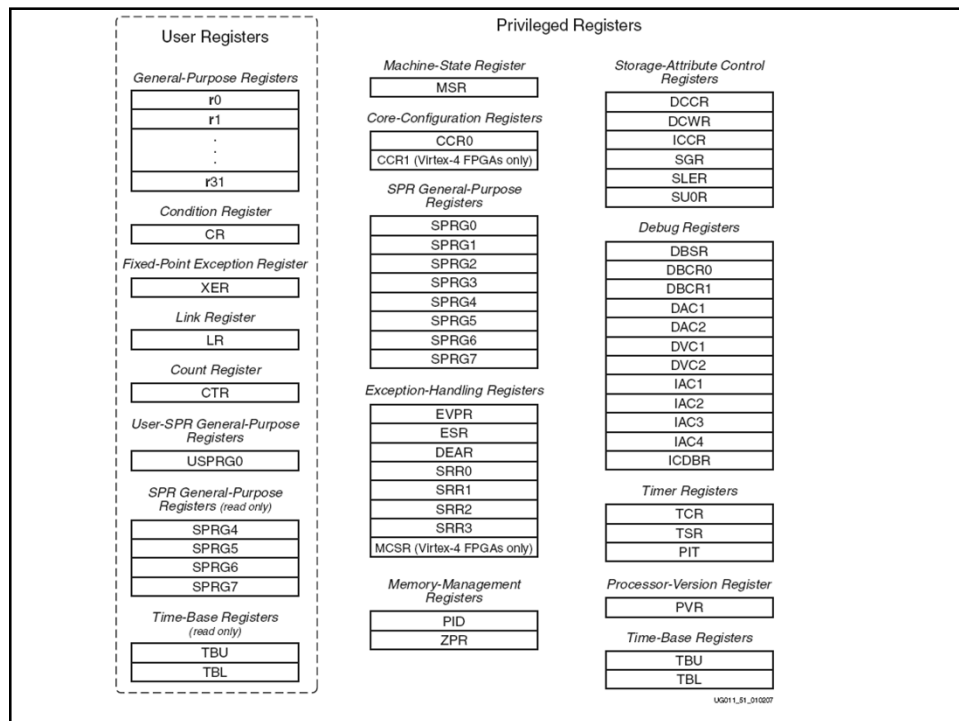
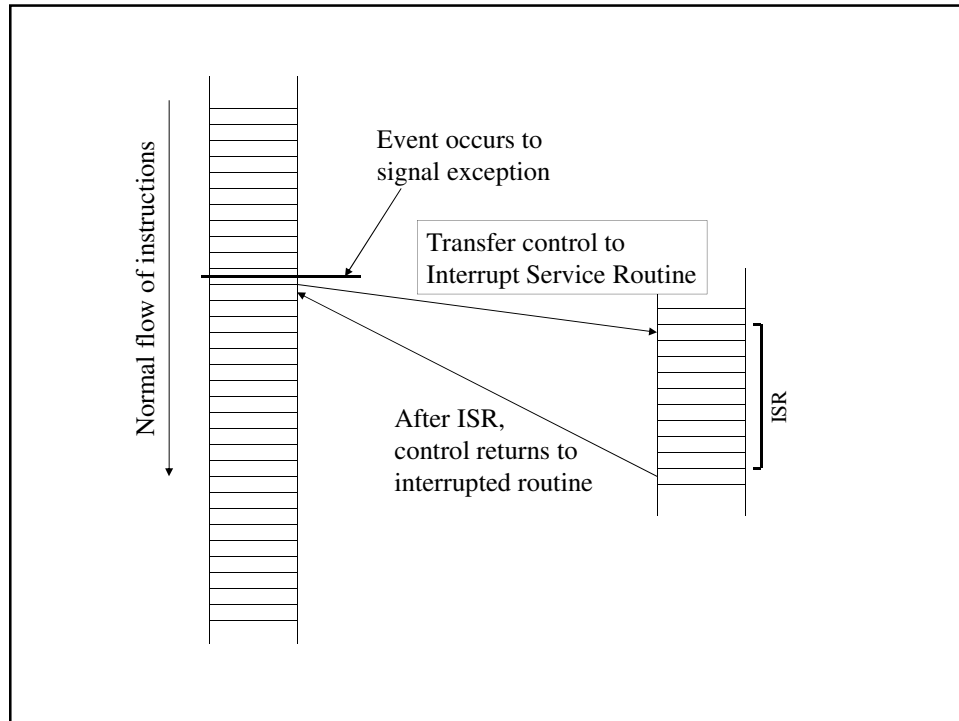
$$11110000.11110000 = -15.0625$$

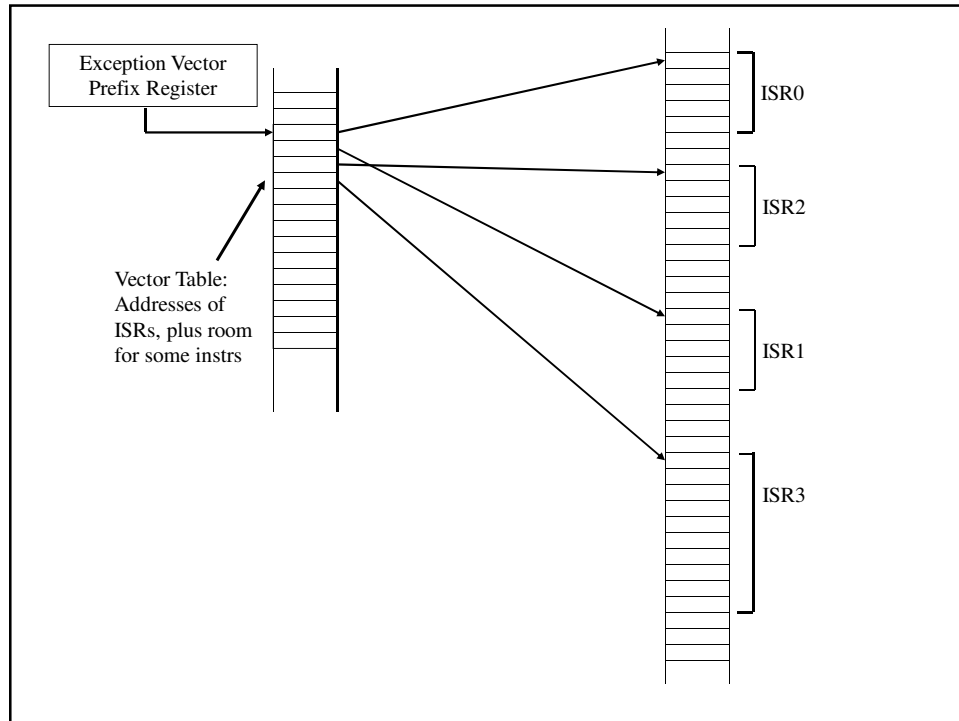
$$1111.000011110000 = -0.94140625$$

$$1.111000011110000 = -0.11767578125$$

Interrupt System (1)

- Provide timely access to processor facilities
- Mechanism for returning to program without any program change
- Caused by events needing attention
 - Errors (div by 0, illegal instr, priv violation, etc)
 - Intentional events of system (bkpt, traps, etc)
 - External events (timers, UART, GPIO, etc)
- Response activity contained in ISR





Interrupt System (2)

- Determine cause of interrupt, which ISR to invoke
 - Concept of vector for one of several ISRs
 - difference between PPC and other processors
 - Concept of polling for multiple interrupts that share a single ISR
 - Interrupt system determines which ISR to invoke; ISR looks at appropriate register for polling as needed

Interrupt System (3): Initialization

- Activities that occur only once, to allow participation in interrupt system or identify details of interrupt process
 - Creation of ISRs to deal with interrupt action
 - Setup of registers in individual functional units
 - Building Int Table, setup of EVPR
 - Setup of MSR (EE; others as needed)

Machine State Register: Defines Processor State

- Bit 13 - WE - Wait State Enable
- Bit 14 - CE - Critical Interrupt Enable
- Bit 16 - EE - External Interrupt Enable
- Bit 17 - PR - Privilege Level
- Bit 19 - ME - Machine Check Enable
- Bit 21 - DWE - Debug Wait Enable
- Bit 22 - DE - Debug Interrupt Enable
- Bit 26 - IR - Instruction Relocate
- Bit 27 - DR - Data Relocate

Interrupt System (4): Steady State

- Activity defined by ISR (user supplied)
- In PowerPC, interrupt controller to address offset 0x500
 - Other addresses as appropriate
 - User must deal with possible events efficiently
 - User resets flags, re-establish setup condition
- Return from interrupt – return PC, MSR

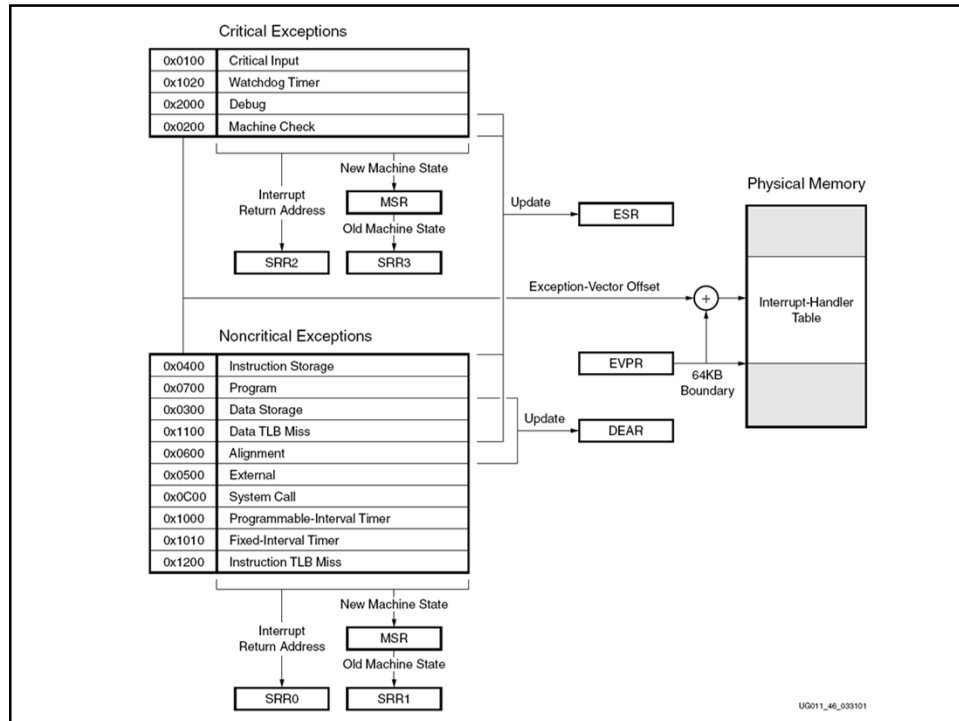
Machine Characteristic: Precise vs Imprecise

- Pipelining: overlap of operations
- Completion of tasks in order vs out of order
- In order: precise, out of order: imprecise

Another classification: Critical vs Noncritical

- Processor responds to Critical Exceptions before Noncritical Exceptions
- Critical Exceptions
 - Critical input exception (external)
 - Machine check exception
 - Watchdog timer exception
 - Debug exception
- All other exceptions (PPC405) non-critical

Exception	Vector Offset	Classification			Cause
Critical Input	0x0100	Critical	Asynchronous	Precise	External critical-interrupt signal.
Machine Check	0x0200	Critical	Asynchronous	Imprecise	External bus error.
Data Storage	0x0300	Noncritical	Synchronous	Precise	Data-access violation.
Instruction Storage	0x0400	Noncritical	Synchronous	Precise	Instruction-access violation.
External	0x0500	Noncritical	Asynchronous	Precise	External noncritical-interrupt signal.
Alignment	0x0600	Noncritical	Synchronous	Precise	Unaligned operand of <code>dcread</code> , <code>lwarx</code> , <code>stwcx</code> , <code>dcbz</code> to non-cacheable or write-through memory.
Program	0x0700	Noncritical	Synchronous	Precise	Improper or illegal instruction execution. Execution of trap instructions.
FPU Unavailable	0x0800	Noncritical	Synchronous	Precise	Attempt to execute an FPU instruction when FPU is disabled.
System Call	0x0C00	Noncritical	Synchronous	Precise	Execution of <code>sc</code> instruction.
APU Unavailable	0x0F20	Noncritical	Synchronous	Precise	Attempt to execute an APU instruction when APU is disabled.
Programmable-Interval Timer	0x1000	Noncritical	Asynchronous	Precise	Time-out on the programmable-interval timer.
Fixed-Interval Timer	0x1010	Noncritical	Asynchronous	Precise	Time-out on the fixed-interval timer.
Watchdog Timer	0x1020	Critical	Asynchronous	Precise	Time-out on the watchdog timer.
Data TLB Miss	0x1100	Noncritical	Synchronous	Precise	No data-page translation found.
Instruction TLB Miss	0x1200	Noncritical	Synchronous	Precise	No instruction-page translation found.
Debug	0x2000	Critical	Asynchronous and synchronous	Precise	Occurrence of a debug event.



Interrupt Activity

- Initialization – prepare system for action
 - Initialize individual module to be able to cause interrupt
 - Initialize Interrupt Controller to allow individual interrupts to reach processor
 - Initialize EVPR to identify ISRs
 - Initialize MSR to permit processor to respond to interrupts

Interrupt Activity

- Before processor initialization – system initialization (planning, preparing)
- Determine desired activity for each module
- Prepare ISR for each module/submodule
- Prepare Table pointing to ISRs

Interrupt Activity

- Steady State
 - Interrupt breaks fetch-decode-execute cycle
 - Vector activity selects appropriate ISR
 - ISR must:
 - Determine and carry out appropriate activity
 - Take steps to reset/clear interrupt request flag(s)
 - Correctly restore system to previous state