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Homework Assignment № 3

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Homework assignment 3

1. Problem 8.11

(a) First-fit

P1 => Mem(1) with 185 left

P2 => Mem(2) with 100 left

P3 => Mem(5) with 392 left

P4 => Mem(3) with 150 left

 $P5 \Rightarrow Mem(5)$

(b) Best-fit

P1 => Mem(6) with 10 left

P2 => Mem(2) with 100 left

P3 => Mem(5) with 392 left

P4 => Mem(4) with 0 left

 $P5 \Rightarrow Mem(5)$

(c) Worst-fit

 $P1 \Rightarrow Mem(4)$

 $P2 \Rightarrow Mem(4)$

 $P3 \Rightarrow Mem(2)$

 $P4 \Rightarrow Mem(1)$

P5 => Fragmentation as it waits for Mem(2) or Mem(4)

2. Problem 8.25

A)
$$2 * 50 \text{ ns} = 100 \text{ ns}$$

B)
$$2 + (.75 * 50) + (.25 * 100) = 64.5 \text{ ns}$$

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3. Problem 9.30

- (a) Define a page-replacement algorithm
 - i. Starting counter value: 0
 - ii. Counters are increase whenever a new page is associated with that frame
 - iii. Counters are decreased whenever one of the pages of the frame is no longer needed
 - iv. The page replacement is selected by breaking FIFO for tie breaking and finding frame with the smallest counter.
- (b) Page faults for the string. = 13

1					5															
	1	2	3	4	4	4	5	1	6	7	7	7	8	8	8	9	5	5	5	4
		1	2	3	3	3	4	5	1	6	6	6	7	7	7	8	9	9	9	5
			1	2	2	2	3	4	5	1	1	1	6	6	6	7	8	8	8	9
P	P	P	P	P			P	P	P	P			P			P	P			P

(c) Minimum possible page faults = 11

1	1	1	1	1	1	1	1	6	6	8	8	8	8	8	8	8	8	8	2
	2	2	2	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5
		3	3	3	3	3	3	3	7	7	7	7	7	7	7	7	7	4	4
			4	4	4	4	4	4	4	4	4	4	9	9	9	9	9	9	9
P	P	P	P	P				P	P	P			P					P	P

4. Problem 9.35 The 240-byte request is assigned a 256-byte segment

The 120-byte request is assigned a 128-byte segment

The 60-byte request is assigned a 64-byte segment

The 130-byte request is assigned a 256-byte segment.

After the allocation, the following segment sizes are available:

64-bytes, 256-bytes, 1K, 2K, 4K, 8K, 16K, 32K, 64K, 128K, 256K, & 512K.

After the releases of memory, the only segment in use would be a 256-byte segment containing 130 bytes of data.

The following segments will be free:

256 bytes, 512 bytes, 1K, 2K, 4K, 8K, 16K, 32K, 64K, 128K, 256K, & 512K.

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5. Problem 12.9

(a) All extens are of the same size

The amount of free space needed for a file must be known in advance, pre-allocation might be inefficient; a file the grows over time must have enough space allocated for its final size. Regardless of how often the file is used. This may invoke internal fragmentation.

The advantage is foe sequential access; the file system remembers the disk address of the block last referenced and thus can read the next block guickly.

- (b) Extents can be of any size and are allocated dynamically. Dynamic allocation involves how to satisfy a request of variable size from a list of free holes. This advantage means that there will be no internal fragmentation. But the first fit and best fit methods for selecting a free hole suffers from external fragmentation.
- (c) Extents can be of a few predetermined sizes. External fragmentation becomes a problem when the largest contiguous chunks are unable to handle a request; storage is fragmented into holes too small to handle the data, depending on the file and disk size.

6. Problem 12.11

The advantage is that while accessing a block that is stored at the middle of a file, its location can be determined by chasing the pointers stored in the FAT as opposed to accessing all of the individual blocks of the file in a sequential manner to find the pointer to the target block. Typically, most of the FAT can be cached in memory and therefore the pointers can be determined with just memory accesses instead of having to access the disk blocks

7. Problem 10.11

- (a) FCFS The schedule is 2150, 2069, 1212, 2296, 2800, 544, 1618, 356, 1523, 4965, 3681. The total distance is 13011.
- (b) SSTF The schedule is 2150, 2069, 2296, 2800, 3681, 4965, 1618, 1523, 1212, 544, 356. The total distance is 7586.
- (c) SCAN The schedule is 2150, 2069, 1618, 1523, 1212, 544, 356, 0, 2296, 2800, 3681, 4965. The total distance is 7115.
- (d) LOOK The schedule is 2150, 2069, 1618, 1523, 1212, 544, 356, 2296, 2800, 3681, 4965. The total distance is 6403.
- (e) C-SCAN The schedule is 2150, 2296, 2800, 3681, 4965, 4999, 0, 356, 544, 1212, 1523, 1618, 2069. The total distance is 9917.
- (f) C-LOOK The schedule is 2150, 2296, 2800, 3681, 4965, 356, 544, 1212, 1523, 1618, 2069. The total distance is 9137.

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8. Problem 10.17

A a) A write of one block of data requires the following: read of the parity block, read of the old data stored in the target block, computation of the new parity based on the differences between the new and old contents of the target block, and write of the parity block and the target block.

B Assume that the seven contiguous blocks begin at a four-block boundary. A write of seven contiguous blocks of data could be performed by writing the seven contiguous blocks, writing the parity block of the first four blocks, reading the eight block, computing the parity for the next set of four blocks and writing the corresponding parity block onto disk.

9. Problem 15.3

When a user creates a password, the system generates a random number (which is the salt) and appends it to the user-provided password, encrypts the resulting string and stores the encrypted result and the salt in the password file. When a password check is to be made, the password presented by the user is first concatenated with the salt and then encrypted before checking for equality with the stored password. Since the salt is different for different users, a password cracker cannot check a single candidate password, encrypt it, and check it against all of the encrypted passwords simultaneously

10. Problem 15.13

It means that the message is encrypted using the public key and then decrypted using the private key. This scheme is not sufficient to guarantee authentication since any entity can obtain the public keys and therefore could have fabricated the message. However, the only entity that can decrypt the message is the entity that owns the private key, which guarantees that the message is a secret message from the sender to the entity owning the private key; no other entity can decrypt the contents of the message.