

ECE 340: PROBABILISTIC METHODS IN ENGINEERING

SOLUTIONS TO HOMEWORK #5

3.10.

Solution

- a) Note that the size of the alphabet of each bit in the password is 2. Then there is a total 2^m number of possible passwords. The sample space S can be expressed as $S = \{B, FB, FFB, FFFB, \dots\}$, where we denote an 'F' as fail and an 'B' as 'BINGO'. Note that the samples space contains a total 2^m number of outcomes.
- b) The mapping from S to S_X is simply:

S	B	FB	FFB	FFFB	...	FFF...FB
	↓	↓	↓	↓	↓	↓
S_X	1	2	3	4	...	2^m

c) $P(\{X=1\}) = \frac{1}{2^m}$

$$P(\{X=2\}) = P(FB) = P(B \text{ at } 2^{\text{nd}} \text{ time} | 1^{\text{st}} \text{ time fail}) \times P(1^{\text{st}} \text{ time fail}) = \frac{2^m-1}{2^m} \frac{1}{2^{m-1}} = \frac{1}{2^m}$$

$$P(\{X=3\}) = P(FFB)$$

$$= P(B \text{ at } 3^{\text{rd}} \text{ time} | 1^{\text{st}} \text{ time fail and } 2^{\text{nd}} \text{ time fail}) \times P(1^{\text{st}} \text{ time fail and } 2^{\text{nd}} \text{ time fail})$$

$$= P(B \text{ at } 3^{\text{rd}} \text{ time} | 1^{\text{st}} \text{ time fail and } 2^{\text{nd}} \text{ time fail}) \times P(2^{\text{nd}} \text{ time fail} | 1^{\text{st}} \text{ time fail}) \times P(1^{\text{st}} \text{ time fail}) = \frac{1}{2^{m-2}} \frac{2^{m-2}-1}{2^{m-1}} \frac{1}{2^m} = \frac{1}{2^m}$$

So, for $1 \leq n \leq 2^m$

$$P(\{X=n\}) = \frac{1}{2^m}$$

3.13. Solution

- a) By the properties of pmfs, we know that,

$$\sum_{k \in S_X} p_K = 1$$

So,

$$1 = \sum_{k \in S_X} \frac{c}{k^2}$$
$$1 = c \sum_{k \in S_X} \frac{1}{k^2}$$

Recalling the following result from the sum of infinite series: $\sum_{k=1}^{\infty} \frac{1}{k^2} = \frac{\pi^2}{6}$, we have:

$$1 = \frac{c\pi^2}{6}$$

So,

$$c = \frac{6}{\pi^2}$$

b) $P(\{X > 4\}) = \sum_{k=5}^{\infty} \frac{6}{\pi^2 k^2}$

or,

$$P(\{X > 4\}) = 1 - P(\{X \leq 4\}) = 1 - \sum_{k=1}^4 \frac{6}{\pi^2 k^2} = 1 - 0.8655 = 0.1345$$

c) $P(\{6 \leq X \leq 8\}) = P[X \leq 8] - P[X \leq 5] = \sum_{k=1}^8 \frac{6}{\pi^2 k^2} - \sum_{k=1}^5 \frac{6}{\pi^2 k^2}$

$$P(\{6 \leq X \leq 8\}) = 0.9286 - 0.8898$$

$$P(\{6 \leq X \leq 8\}) = 0.0388$$

3.16a i) We want to calculate $P(\{X > 2\})$. However, because $\{X=2\} \cup \{X=51\} = \Omega$, we know that the event of $\{X > 2\}$ is the same as the event of $\{X=51\}$. Thus, $P(\{X > 2\}) = 0.2$. (see the solution to problem 3.7).

The event $\{X > 50\}$ is also the same as $\{X=51\}$, so $P(\{X > 50\}) = 0.2$.

3.17.

a) First, note that $S_Y = \{-1, 0, 1, 2\}$. Given the probabilities, we obtain the pmf:

$$p_{-1} = P(\{Y = -1\}) = \frac{1}{10}$$

$$p_0 = P(\{Y = 0\}) = \frac{2}{10}$$

$$p_1 = P(\{Y = 1\}) = \frac{3}{10}$$

$$p_2 = P(\{Y = 2\}) = \frac{4}{10}$$

b) The output is equal to the input if and only if the noise is zero. Hence, the probability that the output is equal to the input is : $p_2 = P\{(Y=2)\} = \frac{4}{10}$

c) The probability that the output of the channel is positive is:

$$P\{(Y > 0)\} = P\{(Y=1)\} + P\{(Y=2)\} = \frac{3}{10} + \frac{4}{10} = \frac{7}{10}$$

3.21.

a) Compare $E[Y]$ to $E[X]$.

First, remember that X is the maximum of the number of heads that Michael gets *and* the number of heads that Carlos gets in pair of flips. The pmf for X is given below:

$$P\{(X=0)\} = P\{((T,T,T,T))\} = 1/16$$

$$P\{(X=1)\} = P\{((T,H,T,H), (H,T,H,T), (T,H,H,T), (H,T,T,H), (T,H,T,T), (H,T,T,T), (T,T,T,H), (T,T,H,T))\} = 8/16$$

$$P\{(X=2)\} = P\{((H,H,T,T), (H,H,T,H), (H,H,H,T), (H,H,H,H), (T,T,H,H), (T,H,H,H), (H,T,H,H))\} = 7/16.$$

So,

$$E\{X\} = \sum_{k=0}^2 k * p_X(k) = 0 * \frac{1}{16} + 1 * \frac{8}{16} + 2 * \frac{7}{16} = \frac{22}{16} = \frac{11}{8}$$

Y is the number of heads in two tosses of a fair coin. Its pmf is given by:

$$P\{(Y=0)\} = P\{((T,T))\} = 1/4$$

$$P\{(Y=1)\} = P\{((H,T), (T,H))\} = 2/4 = 1/2$$

$$P\{(Y=2)\} = P\{((H,H))\} = 1/4$$

$$E\{Y\} = \sum_{k=0}^2 k * p_Y(k) = 0 * \frac{1}{4} + 1 * \frac{1}{2} + 2 * \frac{1}{4} = 1$$

3.28.

$$E\{X\} = \sum_{k=-1}^2 k * p_X(k) = -1 * \frac{1}{10} + 0 * \frac{2}{10} + 1 * \frac{3}{10} + 2 * \frac{4}{10} = 1$$

$$VAR\{X\} = E\{X^2\} - E\{X\}^2 = \left((-1)^2 * \frac{1}{10} + 0^2 * \frac{2}{10} + 1^2 * \frac{3}{10} + 2^2 * \frac{4}{10} \right) - 1^2 = 1$$

Special Problem

The code to generate X n times and obtain an estimate of the mean of X is provided as follows:

```
clc
clear all
close all

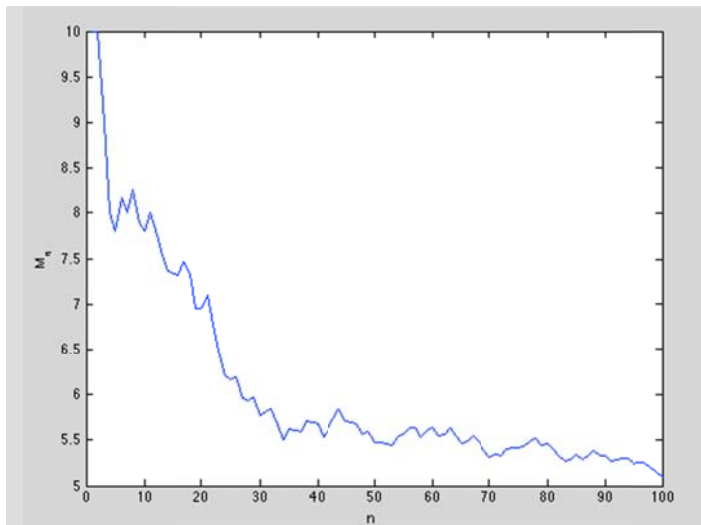
n=input('Give a value for n: ');

for i=1:n
    outcome=ceil(4.*rand(1));           %generate a random integer
                                        %between 1 and 4 with 1/4 of
                                        %probability each

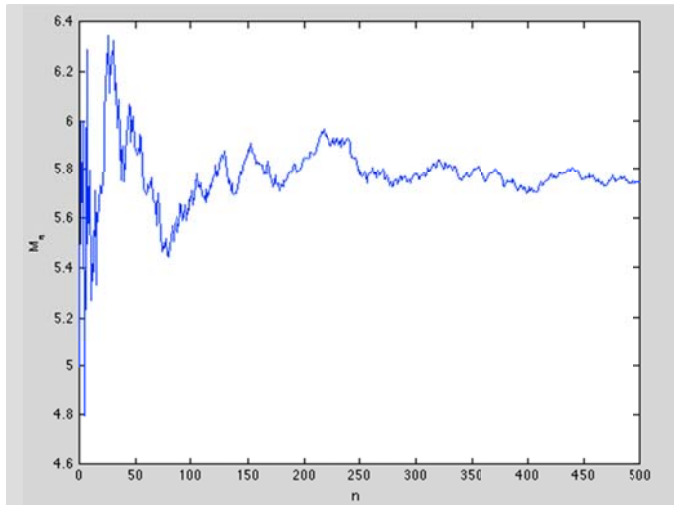
    switch(outcome)                    %assign the value to X according
    case (1)                           %to the outcome of the experiment
        X(i)=10;                      % (H,H)
    case (2)                           % (H,T)
        X(i)=5;
    case (3)                           % (T,H)
        X(i)=7;
    case (4)                           % (T,T)
        X(i)=0;
    end
    meanX(i) = sum(X)/i;
end

plot(1:n,meanX);
xlabel('n')
ylabel('M_n')
```

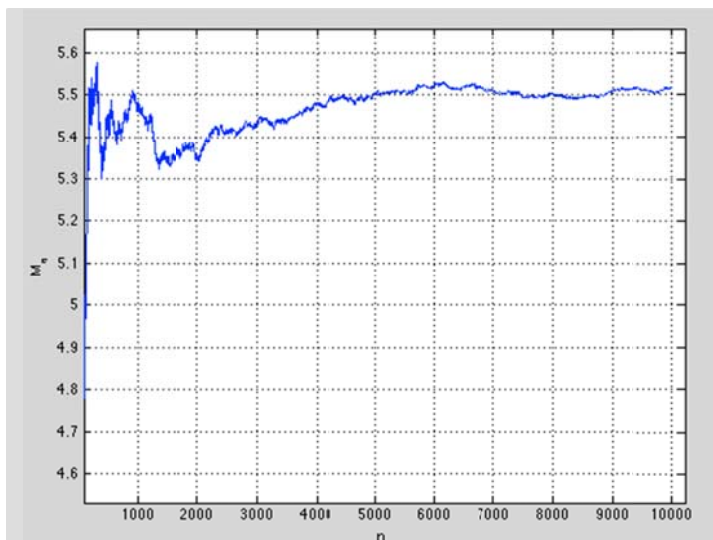
For $n = 100$,



For $n = 5000$,



For $n = 10000$,



It seems that $E\{X\}$ is converging to a value around 5.5.

Calculating the actual mean, we can see that this estimate is correct:

$$E\{X\} = 10 * \frac{1}{4} + 5 * \frac{1}{4} + 7 * \frac{1}{4} + 0 * \frac{1}{4} = 5.5$$

The first observation is that if we use small values of n , we cannot obtain a reliable estimate of the expected value of X , $E[X]$; however, as we increase the number of repetitions of the experiment, the estimate of the expected value of X obtained from the measurements (also called the *sample mean*) starts approaching the true value of $E[X]$.

In summary, if n is sufficiently large, we can use the sample mean as an estimate of $E[X]$. The larger n , the better the estimate of $E[X]$ is. Later in the course we shall see that this is a result of the so-called Law of Large Numbers.