

## Practical 1:

Aim: Linear Search

### Q1] Unsorted Array:

Algorithm:-

Step 1: Define a function with 2 parameters. Use for cond<sup>n</sup> statement with range i.e. length of array to find index.

Step 2: Now use if conditional statement to check whether the given no. by user is equal to the elements in the arrays.

Step 3: If the condition in step 2 is satisfied then return index of given array. If the condition satisfies then get out of loop.

Step 4: Now initialize a variable to enter elements in the arrays from user. Now use split() to split values.

Step 5:- Now initialise a variable as empty array.

*theory:-*  
Linear search is one of the simplest searching algorithm in which each item is sequentially matched with each item in the list. 36

It is worst searching algorithm with worst case time complexity. It is a four approach on the other hand instead of the ordered list instead of searching list in sequence. A binary search is used which will start by examining the middle term. Linear search is a technique to compare each and every element with the key element to be found if both of them matches then its position is found.

```

# sorted array
# code:
def linear(arr, n):
    for i in range(len(arr)):
        if arr[i] == n:
            return i
inp = input("Enter elements in array:").split()
array = []
for ind in inp:
    array.append(int(ind))
print("Elements in array are:", array)
array.sort()
n1 = int(input("Enter element to be searched:"))
n2 = linear(array, n1)
if n2 == n1:
    print("Element found at position", n2)
else:
    print("Element not found")

```

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Step 6: Now use for conditional statement to append the element given as input by user in the empty array.

Step 7: Now again initialise another variable to ask user to find the element in the array.

Step 8: Again initialise a variable to call the defined function.

Step 9: Use if conditional statement to check if variable in step 8 matches with the element you want to find the position. If condition satisfied print element not found.

### Q3] Unsorted Array:

Algorithm:-

Step 1: Define a function with 2 parameters  
use for conditional statements  
with range i.e. length of array  
to find index.

Step 2: Now use if conditional statement  
to check if whether the given  
statement is equal to the element  
in arrays.

Step 3: If the condn in step 2 is satisfied  
then the index no of given  
array. If the condn does not satisfy  
then get out of the loop.

Step 4: Now initialise a variable to  
enter elements in the arrays  
from user. Now use split()  
to split the values.

Unsorted array

(odd)

linear (arr, n):

for i in range (len(arr)):

if arr[i] == n:

return i

inp = input ("Enter elements in array: ") .split()

array = []

for ind in inp:

array.append (int(ind))

print ("Elements in array are: ", array)

x1 = int (input ("Enter element to find: "))

x2 = linear (array, n)

if x2 == x1:  
print ("Element found at location: ", x2)

else:  
print ("Element not found: ")

Output for sorted array :-

>>> Element in array : 1 2 3 5

>>> Elements in array are : [1, 2, 3, 5]

>>> Enter element to search : 2

Element found at position : 1

>>> Element in array are : [1, 2, 3, 4, 5]

>>> Element to search : 7

Element not found.

Output for unsorted array:

>>> Elements in array are : 3, 2, 4, 5, 1

>>> Elements in array are : [3, 2, 4, 5, 1]

>>> Element to be searched : 4

found at location 2.

>>> Element in array : 2 4 5 3 1

>>> Elements in arrays are [2, 4, 5, 3, 1]

>>> Element to be searched : 6

Not found.

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Step:- Now initialise a variable as array  
i.e. empty

Step:- Now use for conditional loop  
statement to append the element as  
input by user in empty array

Step:- Now again initialise another  
variable to as k for user to  
find element in the array

Step:- Again initialise a variable to  
call the variable function.

Step:- To check if user using if  
condn statement to see if variable  
in step 8 matches the elements  
in the list. If not satisfied  
print not found.

## Practical 2:

Aim: Binary Search

Algorithm:

Step 1:- Define a function with 2 parameters  
Now initialise variable with 0 in  
while loop to find the mid value

Step 2:- Use if condn statement to determine  
at which position the mid  
value should point

Step 3:- If the cond doesn't satisfies  
return -1

Step 4:- Now initialise a variable to  
enter elements in array

Step 5:- Use for condition statement  
to append the element in  
empty array.

Q: binary search  
code

```
# def binary (arr, key):  
    start = 0  
    end = len(arr)  
    while start < end:  
        mid = (start + end) // 2  
        if arr[mid] > key:  
            end = mid  
        elif arr[mid] < key:  
            start = mid + 1  
        else:  
            return mid
```

return -1

arr = input('Enter elements in sorted:').split()

arr = []

for ind in array:

arr.append(int(ind))

key = int(input('Enter element to search:'))

index = binary (arr, key)

if index < 0:

print ('Element not found:')

else:

print ('Element not found: ', index)

>>> Enter the elements in array:

3 5 10 12 15 20

>>> Element to be searched : 12  
12 element was found at index 3

>>> Enter the elements in array  
-3, 0, 1, 5, 6, 7, 8

>>> Element to search : 2  
Element was not found

/ ✓

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Step:- Now initialise available to find element in array

Step:- Now calls the function

Step:-

Theory:- Binary search is also known as half interval search algorithm that finds the position of the target. If you are looking for no. which is at end of list then you need to search at the entire list.

## Practical 3:

Aim: Implementation of Bubble sort program on given list.

Theory:- Bubble sort is based on the idea of repeatedly comparing pairs of adjacent elements and then swapping their position. If the sort is simplest form of sorting. In this we sort the given array in ascending order or descending order by comparing adjacent elements at a time.

### Algorithms:-

Step 1:- Bubble sort algorithm starts by comparing 1<sup>st</sup> & 2<sup>nd</sup> elements of an array and swapping them.

Step 2:- If we want to sort the elements of array in ascending order then 1<sup>st</sup> element is greater than second then we need to swap the element.

Step 3:- If the 1<sup>st</sup> element is smaller than second element we do not swap the element.

#BubbleSort  
#ode  
inp = input('Enter Elements').split()  
arr = []  
for ind in inp:  
 arr.append(int(ind))  
print('Element of array before sorting:', arr)  
n = len(arr)  
  
for i in range(0, n):  
 for j in range(n-1):  
 if arr[i] > arr[j]:  
 temp = arr[i]  
 arr[i] = arr[j]  
 arr[j] = temp  
print('Element of array after bubblesort:', arr)

7) Enter element : 2 3 6 1 8 5

7) Enter Elements of arrays before sorting:[2, 3, 6, 1, 8, 5]

7) Sorted list[1, 2, 3, 5, 6]

~~no swap~~

Step 4: Again second and  $\underline{3^{\text{rd}}}$  third element are compared and swapped if it is necessary and this process go on until last but second second last element is compared and swapped.

Step 5: If there are  $n$  elements to be sorted then the process mentioned above should be mentioned to get required result.

Step 6: Stick the output of input of above algorithm of bubble sort stepwise.

Output after each step

After first iteration = 5 4 3 2  
After second iteration = 4 3 2 1  
After third iteration = 3 2 1 4  
After fourth iteration = 2 1 3 4

After fifth iteration = 1 2 3 4

After sixth iteration = 1 2 3 4

## Practical 4: Stack

Aim:- Implementation of Stack using Python list

Theory:- A stack is linear data structure that can be represented in a word form by physical stack or pile. The elements in the form of stack at the top most. Therefore it works in LIFO (Last in first out). It has 3 basic operations namely: push, pop, peek.

Algorithm:-

Step1:- Create a class stack with instance variable items

Step2:- Define the init method with sl argument and initialize the initial value and then initialize the initial value and then initialize the empty value.

Step3:- Define methods push & pop under class stack

Step4:- Use if statement to give the condition if length of given list is greater than range of list then print stack is full

# code  
class

Stack:

def \_\_init\_\_(self):

self.l = [0, 0, 0, 0]

self.tos = -1

def push(self, data):

n = len(self.l)

if self.tos == n - 1:

print("The stack is full")

else:

self.tos += 1

self.l[self.tos] = data

def pop(self):

if self.tos < 0:

print("The stack is empty")

self.l[self.tos] = 0

self.tos -= 1

n = stack()

output  
=>n.push(10)

>>>n.push(9)

>>>n.push(8)

>>>n.push(7)

>>>n.push(6)

>>>n.push(5)

>>>n.push(4)

>>>n.push(3)

>>>n.push(2)

>>>n.push(1)

>>>n.push(0)

>>>n.pop()

$\Rightarrow n \cdot \text{pop}$

$\Rightarrow n \cdot 2 + 2 \cdot 0^2 - 1 \cdot 0^2$

$[10, 9, 8, 6]$

~~first step~~

~~first pop = 10~~

~~first = 9, 8, 6~~

~~first = 8, 6~~

~~first = 6~~

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Step 5:- Else print Enter element.

Step 6:- Push method to insert element &  
Pop is used to delete the element from  
stack.

Step 7:- If pop value method value is less  
than 1 then return stack is  
empty or else delete the element  
from stack at top most position

Step 8:- Assign the element values in push  
method & print given

Step 9:- Attach the imp & output of above  
algorithm

Step 10:- First condition checks whether  
the no of elements are zero while  
second condition checks whether top is  
assigned any value. If top is not  
assigned any value then  
print stack is empty.

## Practical No. 5

Q5) Aim - Implement Quick sort to sort the given

Theory:- The quick sort is a recursive algorithm based on 'divide & conquer' technique.

ALGORITHM :-

Step 1:- Quick sort first selects a value which is called as pivot value. It moves as first value element. It moves as first pivot value. Since we know first will eventually end up as last in that list.

Step 2:- The partition process will happen next. It will find the split point and at same time move other items to appropriate side of the list either less than or greater than pivot value.

Step 3:- Partition begins by locating 2 markers lets call them left mark and right mark at the beginning of the list and of remaining items in the list. The goal of partition process is to move items that are on the wrong side with respect to their values which also

# code 46  
 print ("Quick Sort")  
 def partition (arr, low, high):  
 i = low - 1  
 pivot = arr [high]  
 for j in range (low, high):  
 if arr [j] <= pivot:  
 i = i + 1  
 arr [i], arr [j] = arr [j], arr [i]  
 arr [i+1], arr [high] = arr [high], arr [i+1]  
 return i + 1  
 def quicksort (arr, low, high):  
 if low < high:  
 pi = partition (arr, low, high)  
 quicksort (arr, low, pi - 1)  
 quicksort (arr, pi + 1, high)  
 n1 = input ("Enter element in the list").split()  
 alist1 = [ ]  
 for b in n1:  
 alist1.append (int(b))  
 print ("Elements in list are:", alist1)  
 n = len (alist1)  
 quicksort (alist1, 0, n-1)  
 print ("Elements after quick sort are:", alist1)

Output:-

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Quick sort

Enter elements in the list 21, 22, 20, 30, 24, 30, 34

(i) elements in list are: [21, 22, 20, 30, 24, 30, 34]  
 elements after quicksort: [20, 21, 22, 24, 30, 34]

converge on the split point.

Step 1:- we begin by increasing the leftmark until it locates a value greater than P.V. Then decrement rightmark until we find value i.e. less than P.V. At this point we have found 2 items that are not in place.

Step 2:- At the point the right right mark becomes less than leftmark respectively. The position of rightmark is now the split point.

Step 3:- The pivot point can be split with content of split point in P.V. is now in place.

Step 4:- In addition, all the elements to the left of p.v. is less than the pivot mark. All items to the right is greater than p.v.

Step 5:- The quick () invokes a recursive ()

Step 6:- quick sort helper, begins with same base case merge sort.

Q

Step 10:- If length of list is zero or equal one it is sorted.

Step 11:- If it is greater than one it is partitioned in recursive manner.

Step 12:- The partition() process implements the described earlier.

Step 13:- Display & stick the coding of above algorithm.

Q

Aim:- Implementation of Stack using

Theory:- A stack is a linear data structure that can be represented in the world in the form of a physical stack or a pile. The elements in stack are added or removed only from one position i.e. the top position. Thus stack work LIFO principle as the element that was inserted last will be removed first. A stack can be implemented using Array as well as linked stack has 3 basic operations.

## CODE

```

class Stack:
    global tos
    def __init__(self):
        self.l = [0, 0, 0, 0, 0]
        self.tos = -1
    def push(self, data):
        n = len(self.l)
        if self.tos == n - 1:
            print("Stack is full")
        else:
            self.tos = self.tos + 1
            self.l[self.tos] = data
    def pop(self):
        if self.tos < 0:
            print("Stack is empty")
        else:
            k = self.l[self.tos]
            print("data =", k)
            self.l[self.tos] = 0
            self.tos = self.tos - 1

```

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push, pop, peek. The operation of adding  
or removing the elements is known as  
B Push & Pop.

## ALGORITHM

Step 1: Create a class with instance variable items.

Step 2: Define the init() with self argument to initialize the initial value list. Then initialise to an empty vector.

Step 3: Define methods push & pop under class stack.

Step 4: Use if condition to give the condition that if length of given list is greater than range of list then print: stack is full

Step 5: Else print statement as insert the element into the stack & initialise the value.

Step 6: Push() used to insert the element but pop() use to delete the element from stack

Q5. Step 7: If in POP() value is less than  
print The stack is empty or  
delete the element for Stack  
topmost position.

Output:-

```

s = stack()
s.push(10)
s.push(20)
s.push(30)
s.push(40)
s.push(50)
s.push(60)
s.push(70)
s.push(80)
s.pop()
s.pop()
s.pop()
s.pop()
s.pop()
s.pop()
s.pop()
s.pop()

```

Step 8:- First condition checks whether  
no of elements is zero  
second can whether tos  
is any  $\geq$  value. If tos  
is any value. If tos  
is any value. If tos  
is not any value. So we  
are sure that stack is empty.

Step 9:- Assign the elements values  
push() to add & print the  
is popped.

Step 10:- Attach the input & output  
above Algo rithm.

Output:-

```

stack is full
data = 70
data = 60
data = 50
data = 40
data = 30
data = 20
data = 10
stack is empty

```

#Code:

Class Queue:

```
Q = Queue()
global r
global f
def __init__(self):
    self.r = 0
    self.f = 0
    self.l = [0, 0, 0, 0, 0, 0]
def add(self, data):
    n = len(self.l)
    if self.r < n - 1:
        self.l[self.r] = data
        self.r += 1
    else:
        print("Queue is full")
def remove(self):
    n = len(self.l)
    if self.f < n - 1:
        print(self.l[self.f])
        self.f += 1
    else:
        print("Queue is empty")
```

Q = Queue()

Practical 6:-

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Title: Implementing a Queue using python list

Theory: Queue is a linear data structure which has 2 references front & rear. Implementing a queue using python lists is the simplest as the list as python list provides inbuilt functions to perform the specified operations of the queue. It is based on principle that a new element is inserted after rear end element of the queue is deleted which is at front. In simple term, a queue can be described as a data structure based on first in first out.

Queue(): creates a new empty queue.

Enqueue(): Insert an element at the top rear of the queue and similar to that of insertion of linked list by tail.

Dequeue(): Returns the element which was at front the front is moved to successive element & dequeue operation cannot remove element if queue is empty.

## Algorithm

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Step 1:- Define a class Queue Class with global variables then define with self argument in define assign or, initialize the initial with help of self argument.

Step 2:- Define a empty list and enqueue() method with 2 arguments assign the length of empty list

Step 3:- Define Queue() with self argument if statement that front is equal to length of list then display Queue is empty or else give that front is at 0 & else delete the element from front side & increment it by 1.

Step 4:- Use if statement that length is equal to rear then queue is full or else insert the element in empty list or display that Queue element added successfully & increase by it

Step 5:- Now all the Queue() function be given the element that has to be added in the empty list by enqueue() prior list also.

## Output :-

Q.add(30)  
Q.add(40)  
Q.add(50)  
Q.add(60)  
Q.add(70)  
Q.add(80)  
Q.add(90)

Queue is full

Q.remove()  
30  
Q.remove()  
40  
Q.remove()  
50  
Q.remove()  
60  
Q.remove()  
70  
Q.remove()  
80  
Q.remove()

Queue is empty

II Code

```

def evaluate(s):
    k = s.split()
    n = len(k)
    stack = []
    for i in range(n):
        if k[i].is digit():
            stack.append(int(k[i]))
        elif k[i] == '+':
            a = stack.pop()
            b = stack.pop()
            stack.append(int(b) + int(a))
        elif k[i] == '-':
            a = stack.pop()
            b = stack.pop()
            stack.append(int(b) - int(a))
        elif k[i] == '*':
            a = stack.pop()
            b = stack.pop()
            stack.append(int(b) * int(a))
        else:
            a = stack.pop()
            b = stack.pop()
            stack.append(int(b) / int(a))
    return stack.pop()
    r = evaluate(s)

```

### Practical No 7

Aim:- Program on Evaluation of given string by using stack in python environment ie Postfix

Theory:- The postfix expression is free of any parentheses. Further we took care of the properties of the operators in code program. A postfix expression can be easily be evaluated using stack. Reading the expression is always from left to right in postfix.

#### Algorithm:-

Step 1:- Define evaluate as a function Then Create a empty stack in python

Step 2:- Convert the string to a list by using the string method 'split'

Step 3:- Calculate the length of string in print it

Step 4:- Use for loop to assign the range of string Then give conclusion using if statement.

Q

Step 5 - Scan the token list from right. If token is an operator, convert it from a string to int & push the value onto p.

print ("The evaluated value is: ", 1) 51

Step 6 - If the token is an operator + - \* /, it will need two operands. Pop the 'P' twice. The pop is second operand & the pop is the first operand.

Step 7 - Perform the arithmetic operation. Push the result back on the 'm'.

Step 8 - When the input expression has been completely processed, the value is on the stack.

Step 9 - Print the result of string after evaluation of postfix.

Step 10 - Attach output in input of above algorithm.

```

# Code
class Node:
    global data
    global next
    def __init__(self, item):
        self.data = item
        self.next = None

class linkedlist:
    global s
    def __init__(self):
        self.s = None
    def add1(self, item):
        newnode = node(item)
        if self.s == None:
            self.s = newnode
        else:
            head = self.s
            while head.next != None:
                head = head.next
            head.next = newnode
    def addB(self, item):
        newnode = node(item)
        if self.s == None

```

### Practical 8:

Aim:- Implementation of single linked list by adding the nodes from last position

Theory:- A linked list is a linear data structure which is storing the elements in a node in a linear fashion but no necessarily continuous. The individual element of the linked list called as a Node consists of 2 parts: ① Data ② Next Data stores all the information wrt the element whereas next refers the next node.

### Algorithm:

Step1:- Traversing of a linked list means positioning all the nodes in the linked list in order to perform some operation on them.

Step2:- The entire linked list means can be accessed as the 1<sup>st</sup> node of the linked list.

Step3:- Thus the entire linked list can be traversed using the node which is referred by the Head pointer of linked list.

Step 4:- Now that we know that we can traverse the entire linked list using head pointer we should only visit/refuse the first node of list.

Step 5:- We should not use the head pointer to traverse the entire linked list because the head pointer is our only reference to 1<sup>st</sup> node.

Step 6:- We may lose the reference to in our linked list & most of our linked list. So in order to avoid making some unwanted changes to 1<sup>st</sup> node we will use temporary node to terminate the entire linked list.

Step 7:- We will use this temporary node as a copy of the node we currently have. Since we are making temporary node as copy of current node the data part of temporary node should also hold

Step 8:- Now that current is referring to the 1<sup>st</sup> node if we want to access 2<sup>nd</sup> node of list we need

```
self.s = newnode  
else:  
    newnode.next = self.s  
    self.s = newnode  
display(self):  
    head = self.s  
    while head.next != None:  
        print(head.data)  
        head = head.next  
    print(head.data)  
start = LinkedList()  
= output:
```

```
>>> start.add(80)  
>>> start.add(70)  
>>> start.add(60)  
>>> start.add(50)  
>>> start.add(40)  
>>> start.add(30)  
>>> start.add(20)
```

```
>>> start.display()  
>>> 20  
>>> 30  
>>> 40  
>>> 50  
>>> 60  
>>> 70  
>>> 80
```

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Step 5: Now that we know that we have the entire linked list using head pointer we should only visit the first node of list.

Steps: We should not use the head pointer to traverse the entire linked list because the head pointer is our only reference to 1st node.

Steps: We may lose the reference in our linked list & most of our linked list. So in order to making some unwanted changes to 1st node we will use temporary node to terminate the entire linked list.

Step 6: We will use this temporary node as a copy of the node currently being traversed. Since we are making temporary node copy of current node the data of temporary node should also be

Step 7: Now that current is referring to the 1st node if we want to access 2nd node of list we have to refer as next node of 1st node

```

self.s = newnode
newnode.next = self.s
self.s = newnode
display(self)
head = self.s
while head.next != None:
    print(head.data)
    head = head.next
print(head.data)
start = linkedlist()
= output:

```

```

>>> start.add(80)
>>> start.add(70)
>>> start.add(60)
>>> start.add(50)
>>> start.add(40)
>>> start.add(30)
>>> start.add(20)
>>> start.display()
>>> 20
>>> 30
>>> 40
>>> 50
>>> 60
>>> 70
>>> 80

```

Step 9:- But the 1<sup>st</sup> node is referred by current . so we can traverse to 2<sup>nd</sup> nodes as  $n^{\text{th}} = \text{h.next}$

Step 10:- Similarly we can traverse rest of nodes in the linked list.

Step 11:- Our concern now is to find terminating condition for the while loop.

Step 12:- The last node in the linked list is referred by tail of linked list.

Step 13:- So we can refer to lastnode of linked list.

Step 14:- We have to now see how to start traversing the linked list & how to identify whether then we have reached the last node.

Step 15:- Attach the coding or input & output of above algorithm.

### Practical 0

Exm. Implementation of merge sort by Python

Theory - Merge sort is a divide & conquer algorithm. It divides input array into 2 halves. It cuts itself to merge. The 2 sorted halves merge. A helper is used for merging.

### Algorithm:

Step 1 - The list is divided into left right in each recursive call. Adjacent elements are obtained.

Step 2 - Now begins the sorting process. This follows however in each call. The k iteration hasn't the whole lists & makes along the way.

Step 3 - If the value of  $a[i]$  is smaller than the value at  $j$ ,  $L[i]$  is assigned to the  $a[i+1]$  slot.  $i$  is incremented. If not then  $R[j]$  is chosen.

code  
sort (arr, l, m, r):  
    if n1 = m + 1 + r  
        n2 = l - m  
    L = [0] \* (n1)  
    R = [0] \* (n2)  
    for i in range (0, n1):  
        L[i] = arr [l+i]  
    for j in range (0, n2):  
        R[j] = arr [m+1+j]

    i = 0  
    j = 0  
    k = L  
    while i < n1 & j < n2:  
        if L[i] <= R[j]:  
            arr [k] = L[i]  
            i = i + 1  
        else:  
            arr [k] = R[j]  
            j = j + 1  
    k = k + 1  
    while i < n1:  
        arr [k] = L[i]  
        i = i + 1  
        k = k + 1

```

while j < n
    arr [k] = R [j]
    arr += 1
    k = k + 1
    arr (arr, i, u)

def merge sort (arr):
    if i < u:
        m = int ((l + (u - 1)) / 2)
        merge sort (arr, l, m)
        merge sort (arr, m + 1, u)
        sort (arr, l, m, u)
    print (arr)
    n = len (arr)
    mergesort (arr, 0, n - 1)
    print (arr)

```

Output:

~~[12, 23, 34, 56, 78, 45, 86, 98, 42]~~

~~(12, 23, 56, 56, 42, 45, 78, 86, 98)~~

Step 4:- This way, the values being assigned through  $arr[i+1]$  are all sorted.

Step 5:- At the end of this loop, one of the halves may not have been traversed completely leaving slots in the list.

Step 6:- Thus, the merge sort has been implemented.

### Practical 10:

Time Implementation of sets using Python

#### Algorithm:

- Step 1:- Define 2 empty sets  
Now use for statements  
The range of above 2 sets  
Then print sets for addition
- Step 2:- Now add() is used  
of elements in the given for loop  
Then print sets for addition
- Step 3:- Find the union of 2 sets by using union method  
above 2 sets by using intersection method  
Print sets of union as set 3.
- Step 4:- Using if statement find out  
superset of set 3 & set 4. Display  
above set.
- Step 5:- Display No element in set 3.  
return using mathematical op.
- Step 6:- Use clear() to remove or delete  
1 print set after clearing the elements  
present in the set.
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- code
- ```
set1 = set()
set2 = set()
for i in range(0, 15):
    set1.add(i)
for i in range(1, 12):
    set2.add(i)
print("set1: ", set1)
print("set2: ", set2)
print("\n")
print("Union of set1 & set2: set3")
set3 = set1 | set2
print(set3)
print("Intersection of set1 & set2: set4")
set4 = set1 & set2
print(set4)
print("\n")
if set3 > set4:
    print("set3 is superset of set4")
else:
    print("set3 is same as set4")
if set4 < set3:
    print("set4 is subset of set3")
else:
    print("set4 is same as set3")
print("\n")
set5 = set3 - set4
print(set5)
```

```
Print("Elements in set 3 & not in  
Print("\n")  
if sets4 is disjoint(sets5):  
    Print("cause set 5 are mutually  
        exclusively")
```

Set Clear()

```
Print("After applying clear, sets is empty")  
Print("sets5 = ", sets5)
```

Output:

set1: {8, 9, 10, 11, 12, 13, 14}

set2: {1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11}

Union of set1 & set2: set3

{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14}

Intersection of set1 & set2: {8, 9, 10, 11}

set 3 is super set of set4.

Elements in set 3 & not in set4: sets5

{1, 2, 3, 4, 5, 6, 7, 12, 13, 14}

set 4 & set 5 are mutually exclusive

Set 4 & set 5 are mutually exclusive  
after applying clear set 5 is empty so

set4 = set()

### Practical!!

Aim:- Program based on binary search tree by implementing inorder, postorder traversal.

Theory:- Binary tree is a tree which supports maximum of 2 children for node in the tree. Thus any particular node have either 0 or 2 children. There is another identity of binary tree that it is ordered such that one child is identified as left child & other as right child.

- Inorder:- Traverse the left subtree in form left-right-left. i.e. left subtree inform might have left right subtrees.
- Pre-order:- Visit the root node then the left subtree & right subtree.
- Post order:- Traverse the left subtree in form right-left-right. i.e. left & right subtrees.

CODE

```
class BST:
    def __init__(self):
        self.root = None
    def add(self, value):
        p = node(value)
        if self.root == None:
            self.root = p
            print("Root is added successfully", p.val)
        else:
            n = self.root
            if p.val < n.val:
                if n.left == None:
                    n.left = p
                    print(p.val, "is added successfully")
                    break
                else:
                    n = n.left
            else:
                if n.right == None:
```

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n. My  
print(p.val, "none")  
right side successfully  
break

```

    to
    else:
        n = n.right
    def Inorder(root):
        if root == None:
            return
        else:
            Inorder(root.left)
            print(root.val)
            Inorder(root.right)
    def Preorder(root):
        if root == None:
            return
        else:
            print(root.val)
            Preorder(root.left)
            Preorder(root.right)
    def Postorder(root):
        if root == None:
            return

```

### Algorithm:

Step 1: Define class node & define init() with 2 arguments. Initialize the value in this method.

Step 2:- Again, define a class BST that is binary search tree with init() with self argument & assign the root as none

Step 3: Define add for adding the node  
Define a variable p that  
p = node(value)

Step 4:- Use if statement for checking  
the condition that root is none  
then run else statement for if node  
is less than then main node then  
put on arrange them in left side

Step 5:- Use while loop for checking the node  
is less than or greater than the  
main node & break the loop if it is  
not satisfying.

Step 6:- Use if statement within main statement for checking that node is greater than main root then put it to right side

Step 7:- After this, left side tree & subtree repeat this into arrange the node to binary search tree.

Step 8:- Define Inorder() procedure with root as argument & statement that root is none & return that's all

Step 9:- Inorder else statement used giving that condition if first root, then right node.

Step 10:- For pre-order, we have to condition in else that first root, left & right node

Step 11:- for post order in else part left & right & root

else:

Post order (root, left)  
Post order (root, right)  
Print (root, val)

t = BST()

- output:

t.add(1)

root is added successfully!

t.add(2)

node is added successfully to the right

t.add(4)

node is added successfully to the right

t.add(5)

node is added successfully to the right

t.add(3)

node is added to left side successfully

print ("Inorder : ", inorder(t.root))

Inorder:

$\Rightarrow \text{print}("In Preorder:", \text{Preorder}(\text{root}))$

Preorder

1  
2  
3  
4  
5

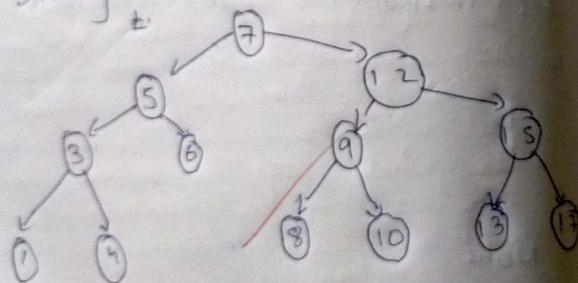
Preorder: None

$\Rightarrow \text{Print}("Post order:", \text{Postorder}(\text{t}, \text{root}))$

3  
5  
4  
2

Postorder: None

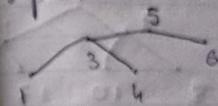
\* Binary search tree:



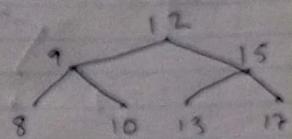
65

In Order (LVR):

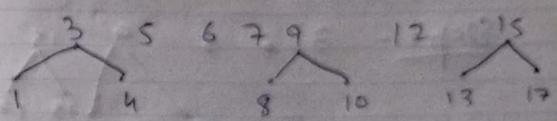
Step 1:



7



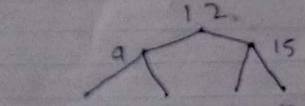
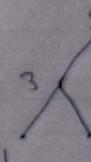
Step 2:



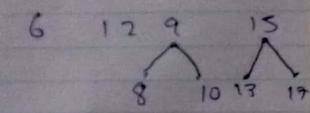
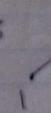
Step 3: 1, 3, 4, 5, 6, 7, 8, 9, 10, 12, 13, 15, 17.

Preorder: (VLR)

Step 1: 7



Step 2:- 7 5 3

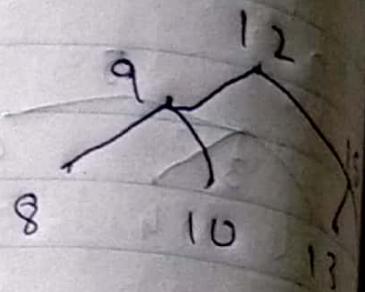
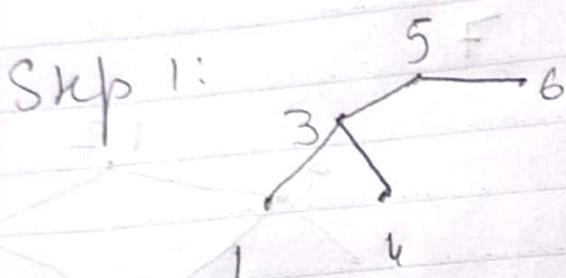


Step 3: 7, 5, 3, 1, 4, 6, 12, 9, 8, 10, 13, 15, 17,

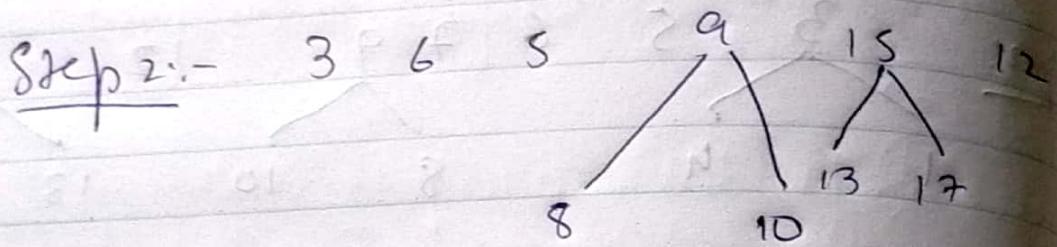
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• POST ORDER : (LRV)

Step 1:



Step 2:-



Step 3: 1, 4, 3, 6, 5, 8, 10, 9, 13, 17