

#Linear Search Unsorted

SOURCE CODE:

```
print("abhay singh")
print("Roll No.1750")
found=False
a=[21,31,41,51,61,71,81,91]
print(a)
search=int(input("Enter the number to be searched: "))
for i in range(len(a)):
    if(search==a[i]):
        print("The number is found at" ,i+1)
        found=True
        break;
if(found==False):
    print("The number is not found")
```

OUTPUT:

```
[Python 3.4.3 Shell]
File Edit Shell Debug Options Window Help
Python 3.4.3 (v3.4.3:9b73f1c3e601, Feb 24 2015, 22:43:06
) ] on win32
Type "copyright", "credits" or "license()" for more information
```

Poochad 1

Ans:- Search a number from the list using linear unsorted

Theory: The process of identifying or finding a particular record is called searching.

There are two types of search.

i) Linear search

ii) Binary search

The linear search is further classified as UNSORTED. Here we'll look at the UNSORTED Linear search, also known as Sequential search, also as a process that checks every element in the list sequentially until the desired element is found.

When the elements we searched are not already arranged in ascending or descending order. They are arranged in random that is what it calls Unsorted Linear search.

UNSORTED Linear search.

- The data is entered in random manner.
- User needs specified the element to be searched in a entered list.
- Check the condition that whether the entered number matches or not matches then displaying the location plus increment 1 as data is stored from location zero.

- If all element are checked one by one and not found then prompt message number not found

#linear search sorted

```
found=False

a=[12,45,66,86,98]

print("Abhay singh")

print("Roll No. 1750")

search=int(input("Enter a Number to be searched :"))

if(search<a[0] or search>a[len(a)-1]):

    print("Number does not exists")

else:

    for x in range(len(a)):

        if(search==a[x]):

            print("The number is found at",x,"index number")

            found=True

            break

    if(found==False):

        print("Number not in the list")


```

(c) DataCamp
In [1]: from array import array
array([12, 45, 66, 86, 98])
Python 3.4.3 (v3.4.3:9b7etting, Feb 24 2015, 22:43:06) [MSC v.1600 32 bit (I)
Type "copyright", "credits" or "license()" for more information.

Practical 2.

Aim: To search a number from the list using linear sorted method.

Theory :- Searching and sorting are different modes are type of data structure.

SORTING - To lexicically sort the inputed data in ascending or descending manner.

SEARCHING - To search element and to display the same.

In searching that too in Linear sorted seq.
The data is arranged in ascending and descending
that is all what it meant by searching through 'seq'
that is well arranged data.

SORTED linear search.

The user is supposed to enter data in sorted manner.
User has to give an element for searching through sorted
if element is found display with an updation as val
is stored from iteration '0'.
If data or element not found print the same.

1.E

Practical 3

Aim: To search a number from the given sorted list using binary search.

Theory: A linear search also known as a half interval search, is an algorithm used in computer science to locate the search to the binary - the array must be sorted in either ascending or descending order.

At each step of the algorithm a comparison mode and the procedure branches into one of two directions.

Typically the key value is compared to the middle element of the array. If the key value is less than or greater than this middle element, the algorithm know which half of the array to continues searching as the array is sorted.

This process is represented on progressively smaller segments of the array until the value is located.

Because each step in the algorithm divides array size is half a binary search will complete successfully in logarithmic time.

#BINARY SEARCH:

```
print("Abhay singh")
print("Roll No.1750")
a=[12,25,89,57,99]
print(a)

search=int(input("Enter a number to be searched :"))

l=0

r=len(a)-1

while(True):
    m=(l+r)//2

    if(l>r):
        print("Number is not found")
        break

    if(search==a[m]):
        print("Number found at:",m+1)
        break

    elif(search<a[m]):
        r=m-1

    else:
        l=m+1
```

```
[Python 3.4.3 Shell]
File Edit Shell Debug Options Window Help
Python 3.4.3 (v3.4.3:9b73f1c3e601, Feb 24 2015, 22:43:06) [MSC v.1600 32
Type "copyright", "credits" or "license()" for more information.
>>> ===== RESTART =====
>>>
Abhay singh
Roll No.1750
[12, 25, 89, 57, 99]
Enter a number to be searched : 99
Number found at: 5
>>> ===== RESTART =====
>>>
Abhay singh
Roll No.1750
[12, 25, 89, 57, 99]
Enter a number to be searched : 66
Number is not found
>>>
```

Q.

25

Aim : To demonstrate the use of stack.

Theory: In computer science, a stack is an abstract data structure that serves as a collection of elements with two principal operations push, which adds an element to the collection and pop, which removes the most recently added element that was not yet removed. The order may be LIFO (Last in First Out) or FILO (First In Last Out).

Three basic operations are performed in the stack:

- Push: Adds an item in the stack if the stack full then it is said to be overflow condition.
- Pop: Removes an item from the stack in the order in which they are pushed. If the stack is empty, then it is said to be a underflow condition.
- Peek or top: Returns top element of stack.
- Is empty: Returns true if stack is empty else false

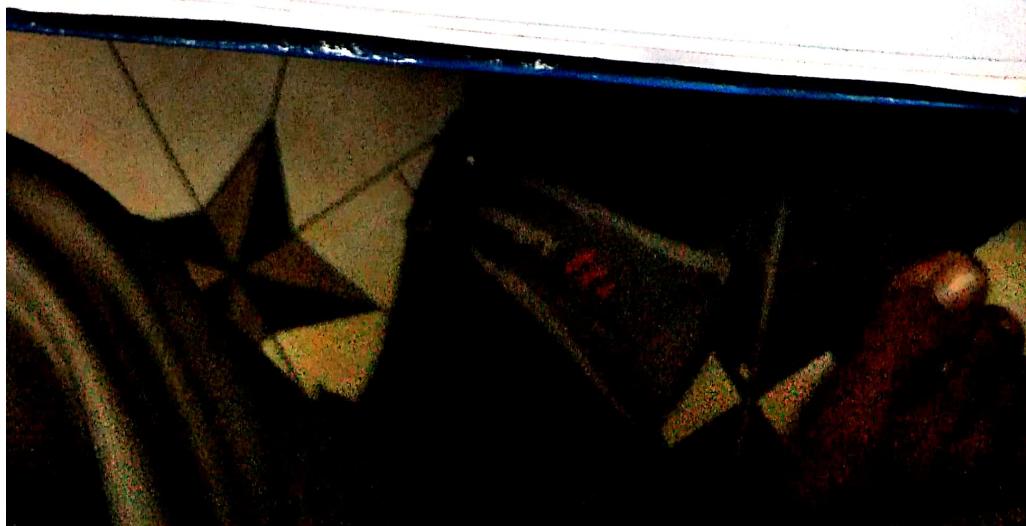
Push

data
data
data
stack

Pop

data
" " " "

Last-in First-Out



```
print("Abhaysingh 1750")

class stack:

    globaltos

    def __init__(self):
        self.l=[0,0,0,0,0,0]
        self.tos=-1

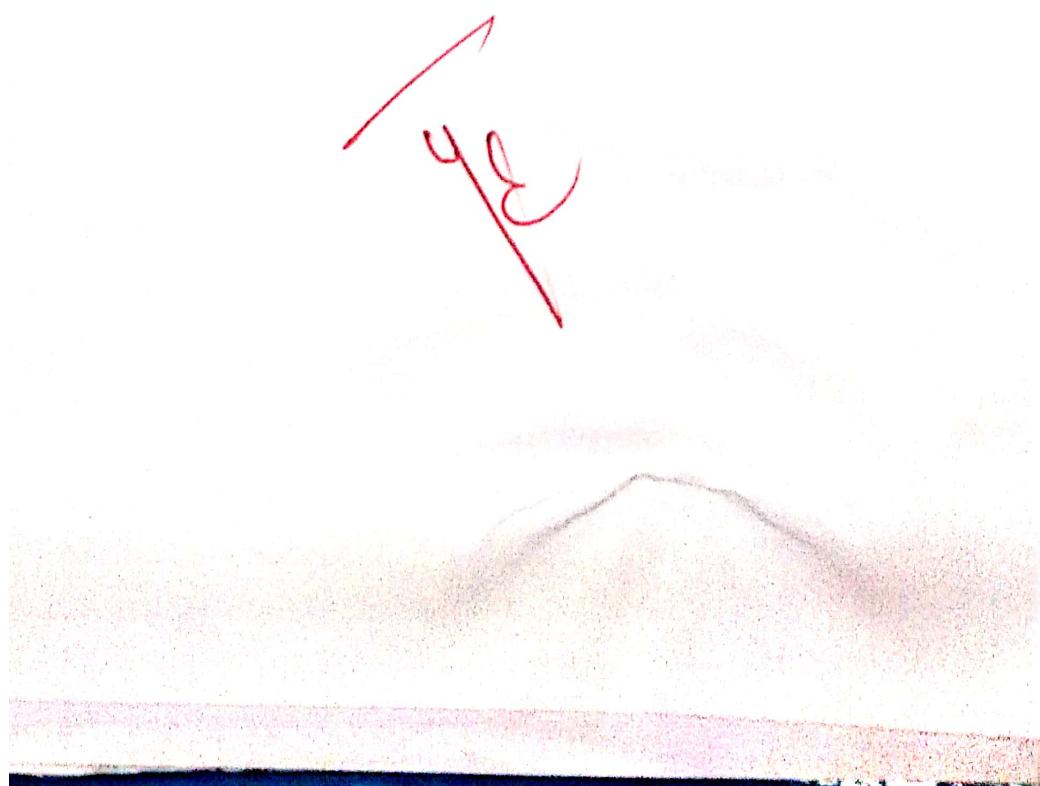
    def push(self,data):
        n=len(self.l)
        if self.tos==n-1:
            print("stack is full")
        else:
            self.tos=self.tos+1
            self.l[self.tos]=data

    def pop(self):
        if self.tos<0:
            print("stack is empty")
        else:
            k=self.l[self.tos]
            self.tos=self.tos-1

s=stack()
s.push(10)
s.push(20)
s.push(30)
s.push(40)
```

```
q.add(50)
q.add(60)
q.add(70)
q.add(80)
q.remove()
q.remove()
q.remove()
q.remove()
```

```
>>>
=====
Abhay singh 1750
queue is full
3
40
50
60
70
queue is full
>>> |
```



```
print("Abhaysingh 1750")

class queue:
    global r
    global f
    def __init__(self):
        self.r=0
        self.f=0
        self.l=[0,0,0,0,0]
    def add(self,data):
        n=len(self.l)
        if self.r<n-1:
            self.l[self.r]=data
            self.r=self.r+1
        else:
            print("queue is full")
    def remove(self):
        n=len(self.l)
        if self.f<n-1:
            print(self.l[self.f])
            self.f=self.f+1
        else:
            print("queue is full")
q=queue()
q.add(3)
q.add(40)
```

```
s.pop()
s.pop()
s.pop()
s.pop()
s.pop()
s.pop()
s.pop()
s.pop()
s.pop()

stack = []
stack.append(1780)
stack.append(18700)
stack.append(10)
stack.append(60)
stack.append(30)
stack.append(40)
stack.append(30)
```

In a queue, one end is always closed
(deque) because queue is open on both of its ends.
Enqueue c) can be termed as add to queue i.e adding

a element in queue.

Dequeue c) can be termed as delete or remove i.e deleting or removing of element

Front is use to get the front data item from a queue
Rear is used to get the last item from a queue

✓ WSC

Aim: To demonstrate the use of circular queue in data structure.

Theory: The queues that use implement using an array suffer from one limitation. In that implementation there is a possibility that the queue is reported as full, even though there might be empty slots at the beginning of the queue. To overcome this limitation we can implement queue as circular queue. In Circular queue we go on adding the elements to the queue and spend the end of the array. The next element is stored in the first slot of the array.

```
print("Abhaysingh 1750")

class queue:

    global r

    global f

    def __init__(self):

        self.r=0

        self.f=0

        self.l=[0,0,0,0,0]

    def add(self,data):

        n=len(self.l)

        if self.r<n-1:

            self.l[self.r]=data

            print("data added",data)

            self.r=self.r+1

        else:

            s=self.r

            self.r=0

            if self.r<self.f:

                self.l[self.r]=data

                self.r=self.r+1

            else:

                print("queue is full")

    def remove(self):

        n=len(self.l)

        if self.f<n-1:
```

```
s=self.f  
  
self.f=0  
  
if self.f<self.r:  
    print(self.l[self.f])  
  
    self.f=self.f+1  
  
else:  
  
    print("queue is empty")  
  
self.f=s  
  
q=queue()  
  
q.add(33)  
  
q.add(44)  
  
q.add(55)  
  
q.add(66)  
  
q.add(77)  
  
q.add(88)  
  
q.remove()  
  
q.add(66)
```

```
>>>  
=====  
Abhay singh 1750  
data added 33  
data added 44  
data added 55  
data added 66  
data added 77  
queue is full  
data removed: 33  
data added 66  
>>>
```

```
### After before linkedlist(simple)###

class node:
    global data
    global next

def __init__(self,item):
    self.data=item
    self.next=None

class linkedlist:
    global s

def __init__(self):
    self.s=None

def addL(self,item):
    newnode=node(item)

    if self.s==None:
        self.s=newnode

    else:
        head=self.s

        while head.next!=None:
            head=head.next

        head.next=newnode

def addB(self,item):
    newnode=node(item)

    if self.s==None:
        self.s=newnode

    else:
        newnode.next=self.s
        self.s=newnode

def display(self):
    head=self.s

    while head.next!=None:
        print(head.data)
```

2000-0000
The 2000-0000 long-term survey was designed
specifically to evaluate the effects of the 1999-2000
flood "event" on the vegetation of the area. The
area

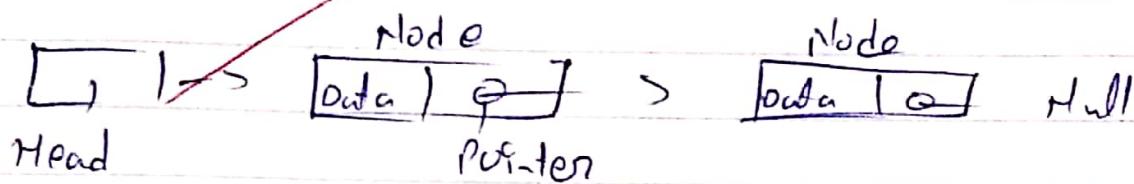
Poachiral 7

Aim: To demonstrate the use of Linkedlist in data structure.

Theory: A Linked List is a sequence of data structure
 Linked list is a sequence of links which contains items
 Each link contains a connection to another link.

- Link - Each link of a linked list can store a data called an element.
- NEXT - Each link of a linked list contain a link to the next link called next.
- Linked - A linked list contains the connection link to the first link called first.

Linked List Representation :-



Types of Linked List :

- Simple
- Doubly
- Circular

900

900

900

900

```
def evaluate(s):  
    k=s.split()  
    n=len(k)  
    stack=[]  
  
    for i in range(n):  
  
        if k[i].isdigit():  
  
            stack.append(int(k[i]))  
  
        elif k[i]=='+':  
  
            a=stack.pop()  
            b=stack.pop()  
  
            stack.append(int(b)+int(a))  
  
        elif k[i]=='-':  
  
            a=stack.pop()  
            b=stack.pop()  
  
            stack.append(int(b)-int(a))  
  
        elif k[i]=='*':  
  
            a=stack.pop()  
            b=stack.pop()  
  
            stack.append(int(b)*int(a))  
  
        else:  
  
            a=stack.pop()  
            b=stack.pop()  
  
            stack.append(int(b)/int(a))  
  
    return stack.pop()  
  
s="8 6 9 * +"
```

Python 3.6.2 Shell
File Edit Shell Debug Options Window Help
Python 3.6.2 (v3.6.2:5tc0d3b5, Jul 4 2017, 04:14:54) [MSC v.1900 32 bit (Intel)] on win32
Type "copyright", "credits" or "license()" for more information.
>>> RESTART: D:\kabhay singh 1750\evaluate.py
The Evaluate value is: 0
REHAY SINON
>>>

Practical 8

Aim: To evaluate Postfix expression using stack.

Theory: Stack is an (ADT) and works on LIFO (Last-in-first-out) i.e. Push & Pop operation.

A Postfix expression is a collection of operators and operands in which the operator is placed after the operands.

Steps to be followed:

- (1) Read all the symbols one by one from left to right in the given Postfix expression.
- (2) If the reading symbol is operand then push it on the stack.
- (3) If the reading symbol is operator (+, -, *, /, etc). Then perform Two Pop operation and store the two popped operand in two different variable (operand 1 & operand 2) perform reading symbol operator using operand 1 & operand 2. Push result back onto the stack.
- (4) Finally! Perform a Pop operation and display the P value as final result.

value of postfix expression:
 $S = 12 \ 3 \ 6 \ 4 \ - + *$

stack:

4	a
6	b
13	
12	

$$b - a = 6 - 4 = 2 \text{ Store again in stack}$$

2	a
3	b
12	

$$b + a = 3 + 2 = 5 \text{ Store result in stack}$$

5	a
12	b

$$b * a = 12 * 5 = 60$$

(rt)

>>>
ABHAY SINGH 1750
[6, 5, 3, 2, 1, 4]
[1, 2, 3, 4, 5, 6]
>>> |

Practical 9

Aim: To sort given random data by using bubble sort.

Theory: SORTING is type in which any random data is sorted i.e arranged in ascending or descending order.

BUBBLE Sort sometimes referred as sorting sort.

Is a simple sorting algorithm that repeatedly steps through the list, compares adjacent elements and swaps them if they are in wrong order. The pass through the list is repeated until the list is sorted. The algorithm which a comparison sort is named for the way smaller or larger elements "bubble" to the top of the list.

Although the algorithm is simple, it is too slow as compare two element checks if condition fail that only swap otherwise goes on eg:

First Pass

$(5\ 1\ 4\ 2\ 8) \rightarrow (1\ 5\ 4\ 2\ 8)$ Here algorithm compare the first two elements and swaps $5 > 1$

$(1\ 5\ 4\ 2\ 8) \rightarrow (1\ 4\ 5\ 2\ 8)$ swap since $5 > 4$

$(1\ 4\ 5\ 2\ 8) \rightarrow (1\ 4\ 2\ 5\ 8)$ swap since $5 > 2$

$(1\ 4\ 2\ 5\ 8) \rightarrow (1\ 4\ 2\ 5\ 8)$ Now since these elements are already sorted $(8 > 5)$ algorithm does not swap them.

Q. 1

Second Pass:

$$(14258) \rightarrow (14258)$$

$$(14258) \rightarrow (12458) \quad \text{Drop } 8 \text{ as } 4 > 2$$

$$(12458) \rightarrow (12458)$$

Third Pass

(12458) Jt clock and give the data in sorted order

```
print("ABHAY SINGH 1750")
def quickSort(alist):
    quickSortHelper(alist,0,len(alist)-1)
def quickSortHelper(alist,first,last):
    if first<last:
        splitpoint=partition(alist,first,last)
        quickSortHelper(alist,first,splitpoint-1)
        quickSortHelper(alist,splitpoint+1,last)
def partition(alist,first,last):
    pivotvalue=alist[first]
    leftmark=first+1
    rightmark=last
    done=False
    while not done:
        while leftmark<=rightmark and alist[leftmark]<=pivotvalue:
            leftmark=leftmark+1
        while alist[rightmark]>=pivotvalue and rightmark>=leftmark:
            rightmark=rightmark-1
        if rightmark<leftmark:
            done=True
        else:
            temp=alist[leftmark]
            alist[leftmark]=alist[rightmark]
            alist[rightmark]=temp
    temp=alist[first]
```

Python 3.4.3 Shell
File Edit Shell Debug Options Window Help
Python 3.4.3 (v3.4.3:9b73f1c3e601, Feb 24 2015, 22:43:06)
Type "copyright", "credits" or "license()" for more information
===== RESTART =====
>>>
>>>
ABHAY SINGH 1750
[42, 45, 54, 55, 66, 67, 80, 89, 100]
>>>

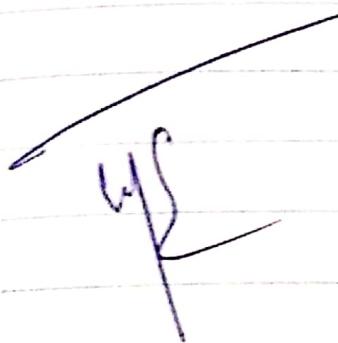
Practical 10

Ques: To evaluate i.e. sort the given data in quick sort.

Topic: Quicksort is an efficient sorting algorithm type of divide and conquer algorithm. It picks an element as pivot and partitions the given array around the picked pivot. There are many different variations of quick sort that pick pivot in different ways.

- ① Always pick first element as pivot.
- ② Always pick last element as pivot.
- ③ pick a random element as pivot.
- ④ pick median as pivot.

The key process in quicksort is Partition(). Target of partition is given an array and an element x of array as pivot, put x at correct position, put y at its current position if sorted as not all smaller element (smaller than x) before x , & put all greater (greater than x) after x . All this should be done in linear time.



1

Practical: 1 Binary Tree

Def: Binary tree is a special type of tree in which every node has either no child or a maximum of two children.

Properties:

1. Every node can have at most two children.

2. Every node can have at most one child.

3. Every node can have at most one child.

4. Every node can have at most one child.

5. Every node can have at most one child.

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94. Every node can have at most one child.

95. Every node can have at most one child.

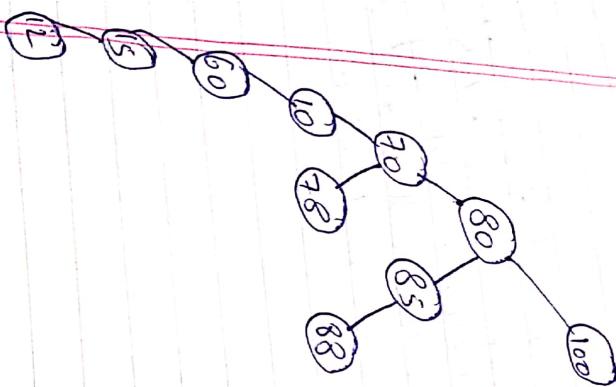
96. Every node can have at most one child.

97. Every node can have at most one child.

98. Every node can have at most one child.

99. Every node can have at most one child.

100. Every node can have at most one child.



Diagnostic representation

BINARY SEARCH TREE

```
class node:  
    global r  
    global l  
    global data  
  
    def __init__(self,l):  
        self.l=None  
        self.data=l  
        self.r=None  
  
class tree:  
    global root  
  
    def __init__(self):  
        self.root=None  
  
    def add(self,val):  
        if self.root==None:  
            self.root=node(val)  
        else:  
            newnode=node(val)  
            h=self.root  
            while True:  
                if newnode.data<h.data:  
                    if h.l==None:  
                        h=h.l  
                    else:  
                        h.l=newnode  
                        print(newnode.data," added on left of ",h.data)  
                        break  
                else:  
                    if h.r==None:  
                        h=h.r  
                    else:  
                        h.r=newnode
```

```

        print(newnode.data, " added on right of ", b.data)
        break

def preorder(self,start):
    if start!=None:
        print(start.data)
        self.preorder(start.l)
        self.preorder(start.r)

def inorder(self,start):
    if start!=None:
        self.inorder(start.l)
        print(start.data)
        self.inorder(start.r)

def postorder(self,start):
    if start!=None:
        self.postorder(start.r)
        self.postorder(start.l)
        print(start.data)

print("Name : Abhay")

t=tree()

t.add(12)

t.add(90)

t.add(13)

t.add(85)

t.add(35)

t.add(44)

t.add(71)

t.add(82)

t.add(16)

t.add(12)print("Preorder is :")

t.preorder(t.root)

print("Inorder is :")

```

```
t.postorder(t.root)

Name : Abhay
90 added on right of 12
13 added on left of 90
85 added on right of 13
35 added on left of 85
44 added on right of 35
71 added on right of 44
82 added on right of 71
16 added on left of 35
12 added on left of 13
Preorder is :
12
90
13
12
85
35
16
44
71
82
Inorder is :
12
12
13
16
35
44
71
82
85
90
Postorder is :
12
16
82
71
44
35
85
13
90
12
>>>
```

Aim: merge sort.

Theory: Merge sort is a sorting technique based on divide and conquer technique. With worst-case time complexity being $O(n \log n)$, it is one of the most popular and most-desired algorithms.

Merge sort first divides the array into two halves and then combines them in a sorted manner.

It divides an input array in two halves itself for the two halves and then merges the two divided halves. The merge() function is used for merging 2 halves. The merge(arr, l, m, r) is a recursive function that assumes that arr[l...m] & arr[m+1...r] are sorted and merges the two sorted sub-arrays into one.

```
print("ABHAY SINGH 1750")
```

```
def sort(arr,l,m,r):
```

```
    n1=m-l+1
```

```
    n2=r-m
```

```
    L=[0]*n1
```

```
    R=[0]*n2
```

```
    for i in range(0,n1):
```

```
        L[i]=arr[l+i]
```

```
    for j in range(0,n2):
```

```
        R[j]=arr[m+1+j]
```

```
    i=0
```

```
    j=0
```

```
    k=l
```

```
    while i<n1 and j<n2:
```

```
        if L[i]<=R[j]:
```

```
            arr[k]=L[i]
```

```
            i+=1
```

```
        else:
```

```
            arr[k]=R[j]
```

```
            j+=1
```

```
            k+=1
```

```
    while i<n1:
```

```
        arr[k]=L[i]
```

```
        i+=1
```

```
        k+=1
```

```
arr[k]=R[j]

j+=1

k+=1

def mergesort(arr,l,r):

    if l<r:

        m=int((l+(r-1))/2)

        mergesort(arr,l,m)

        mergesort(arr,m+1,r)

        sort(arr,l,m,r)

arr=[12,23,34,56,78,45,86,98,42]

print(arr)

n=len(arr)

mergesort(arr,0,n-1)

print(arr)
```

```
File Edit Shell Debug Options Window Help
Python 3.4.3 (v3.4.3:9b73f1c3e601, Feb 24 2015, 22:43:06) [Mi
Type "copyright", "credits" or "license()" for more information
```

```
print("ABHAY SINGH 1750")
a=[23,22,18,96,56,60]
print("Beore sorting\n",a)
for i in range(len(a)-1):
    for j in range(len(a)-1):
        if(a[j]>a[i+1]):
            t=a[j]
            a[j]=a[i+1]
            a[i+1]=t
print("After selection sort\n",a)

output:
Python 3.4.3 (v3.4.3:9b73fbc3e601, Feb 24 2015, 22:43:06) [MSC v.1600 32 bit (Intel)]
Type "copyright", "credits" or "license()" for more information.
>>> ===== RESTART =====
>>>
>>> ABHAY SINGH 1750
Beore sorting
[23, 22, 18, 96, 56, 60]
After selection sort
[60, 96, 96, 96, 96, 96]
>>> |
```

Practical 13

Aim: To evaluate > o to sort given random data by using Selection sort.

Theory: Selection Sort is a simple sorting algorithm. The sorting algorithm is an in-place comparison-based algorithm in which the list is divided into two parts - the sorted part at the left and arr. Initially the sorted part is empty and the unsorted part is the entire list. The smallest element is selected from the unsorted array & swapped with the left most element so that array becomes a part of sorted array. This process continues moving unsorted array boundary by one element to the right.

This algorithm is not suitable for large data sets average and worst case complexity are of $O(n^2)$, where n is the number of items.