CS112: Data Structures

Lecture 02

linked lists

CS112: Data Structures

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Class Web Page

- **http://sakai.rutgers.edu/**

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  - **Policies**
  - **Syllabus**
  - **Assignments**
  - **Lecture notes**
  - **etc....**

- **You are assumed to know anything posted.**

Review

What is a Data Structure

- **A way to store multiple pieces of data**
- **Stores some relationship among the pieces**

What to know about a DS

- **What operations can we do?**
- **What do they cost?**
  - **Time**
  - **Memory space**

Review
Asymptotic Costs

- **Problem: actual cost depends on many details**
- **We want a measure of cost that does not depend on these**

Review: Solutions

- **Count operations, not time**
- **Op count = f(input size)**
- **Among inputs of the same size, use worst or average op count**
- **Abstract away details of f: O(f)**
  - **If O(f) > O(g), if n gets big enough f(n) will be larger than k * g(n)**

Example

Arrays a1, a2 in increasing order, length=n

Do they have any common element?

```
int i1 = 0;  int i2 = 0;
while (i1 < n && i2 < n && a1[i1] != a2[i2]){
   if (a1[i1]<a2[i2]){i1++;}
      else {i2++;}}
if (i1 = = n || i2 = = n){
   System.out.println("no");} else {
   System.out.println("yes ");
```

Which Operations to Count?

Count should model time of algorithm

- **Most frequent / inner loop**
- **Most time consuming**
- **Inherent in algorithm, not language**

- **Count a1[i1] != a2[i2]**

Size of input

- **Number of ops = f(input size)**
- **How do we define size of input?**
  - **For this example: n**

Worst / Average Case

- **Worst case: 2n-2 = O(n)**
- **Average case**
  - **Assume will find a match**
  - **Assume each sum (i1 + i2) equally likely to be location of first match**
  - **Sum over all cases prob(case)*cost(case)**

$$\sum_{s=0}^{2n-2} 1/(2n-2) * s$$

$$= (1/(2n-2)) * (2n-2)(2n-1)/2 = O(n)$$

Review: Rules for Big O

- **k is O(1)**

  341 is O(1)
- **f+g = max (O(f), O(g))**

  n + 1 is max(O(n), O(1)) = O(n)
- **k * f = O(f)**

  $O(4*n^4) = O(n^4)$
- **$O(n^A) < O(n^B)$ if A < B**

  $O(n^3) < O(n^4)$
- **O(polynomial) is O(highest exponent term)**

  $5\ n^4 + 44\ n^2 + 55\ n + 12$ is $O(n^4)$

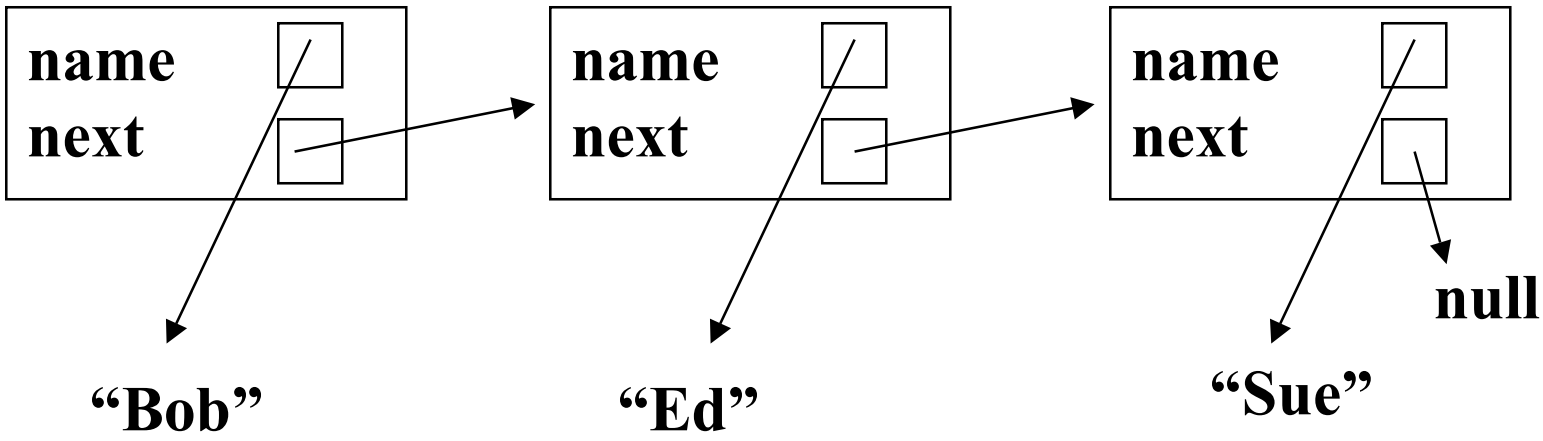
Review: Names for Big O

- **O(1) is constant**
- **O(n) is linear**
- **O(n^2) is quadratic**
- **O(k^n) is exponential**

  O(k^n) is bigger than any polynomial

Review: Linked Lists

- **Class Node: instance variables for**
  - **A name**
  - **The next node in order**

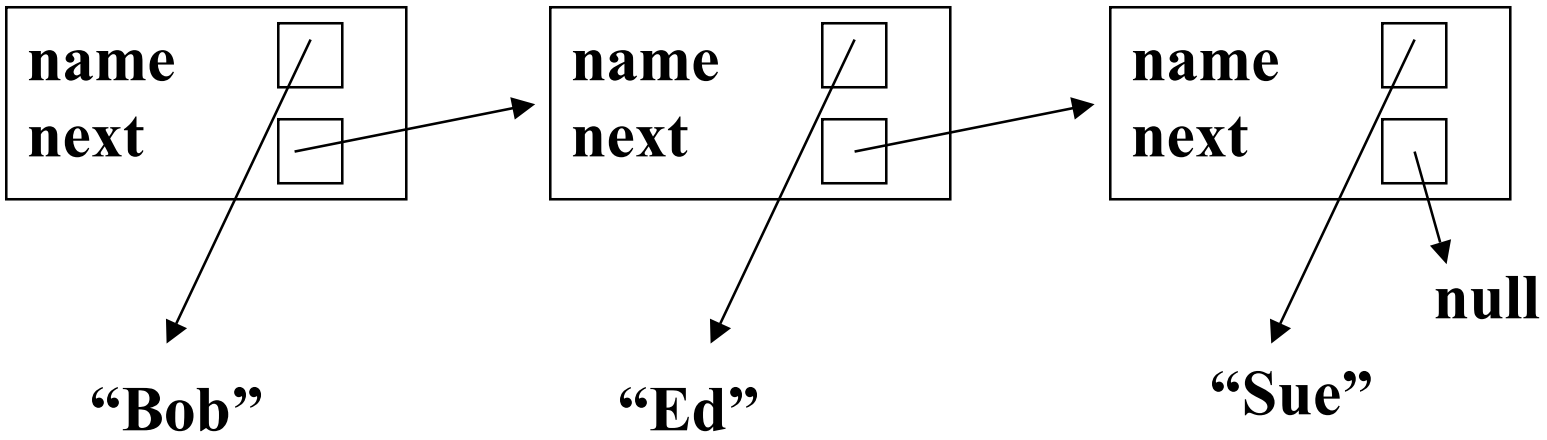


The Node Class

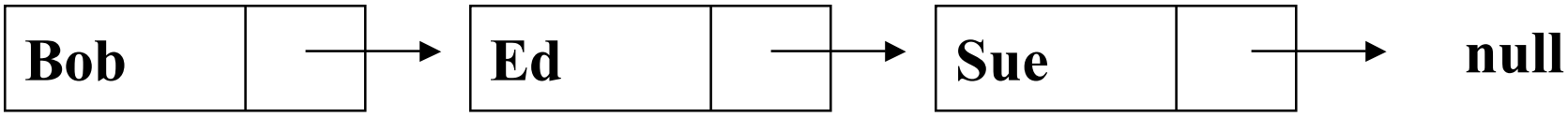
```
public class Node{
    private String name;
    private Node next;

    public Node(String nm, Node nxt){
     name = nm;
     next = nxt;
    }
 …
}
```

Storing "who is next"



- **Can also draw this way**



Operations on linked lists

- **Insert at head**
- **Remove at head**
- **Insert after given node**
- **Remove after given node**
- **Find last**
- **Insert at end**
- **Remove last**
- **Find element i**
- **Find by data**

Generic Lists

- **Problem: suppose you want to have a list of Strings and a list of ints and …**
- **Class declarations and methods are almost identical**
- **Solution in older java: list of Object**
  - **But give up ability for compiler to check**
- **Solution in java 1.5: "generics"**
  - **Class & method definitions parameterized by type**

Generic List

```
public class Node<E> {
    private E data;
    private Node<E> next;

    public Node<E>(E dat, Node<E> nxt){
     data = dat;
     next = nxt;
    }
   public E getHead(Node<E> head){
       E headData = head.data;
     return headData;} …
```

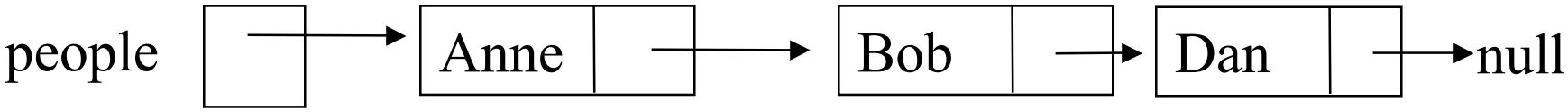
Generic List

In some method:

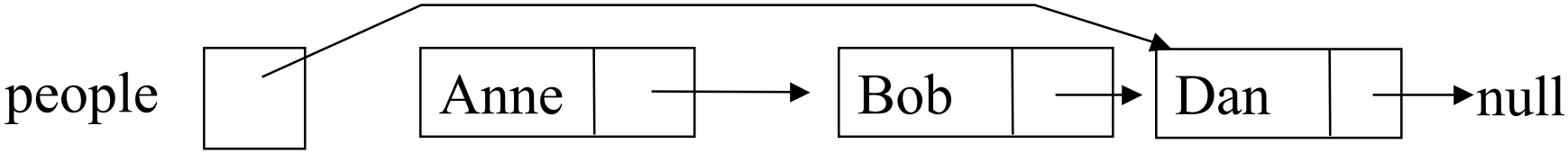
```
Node<String> n1 = new Node<String>;
  …
  String name = n1.getHead( );
```

Circular Lists

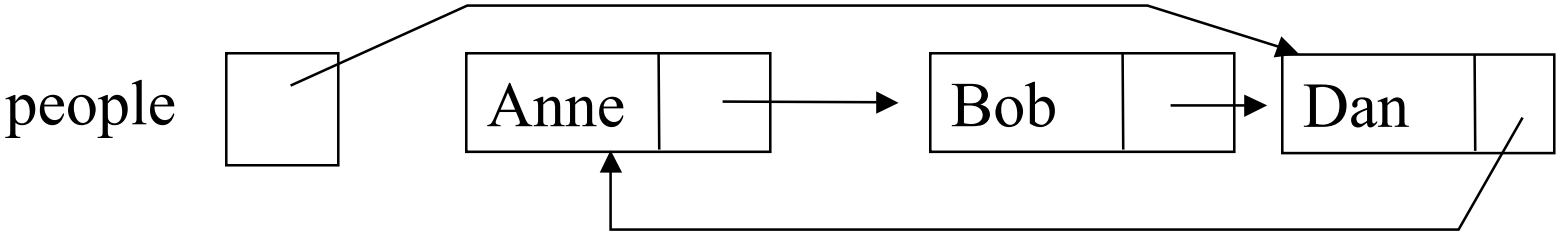
- **Problem: Cost to Access Tail**



- **Solution: point to tail, not head**

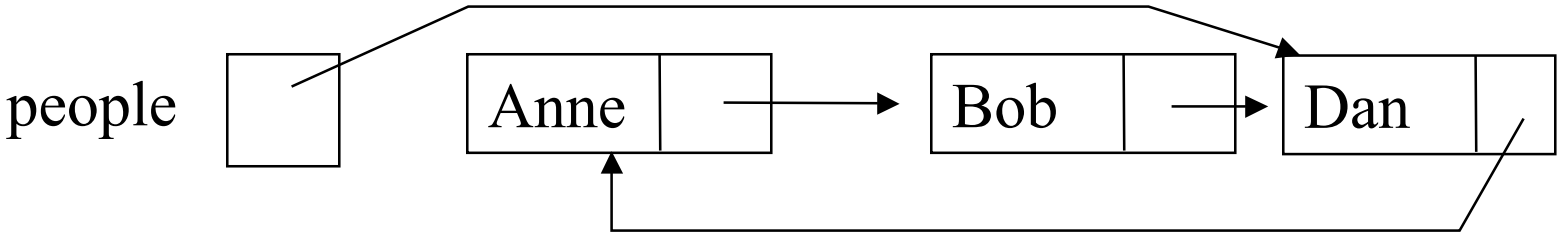


- **Problem: access to head. Solution:**



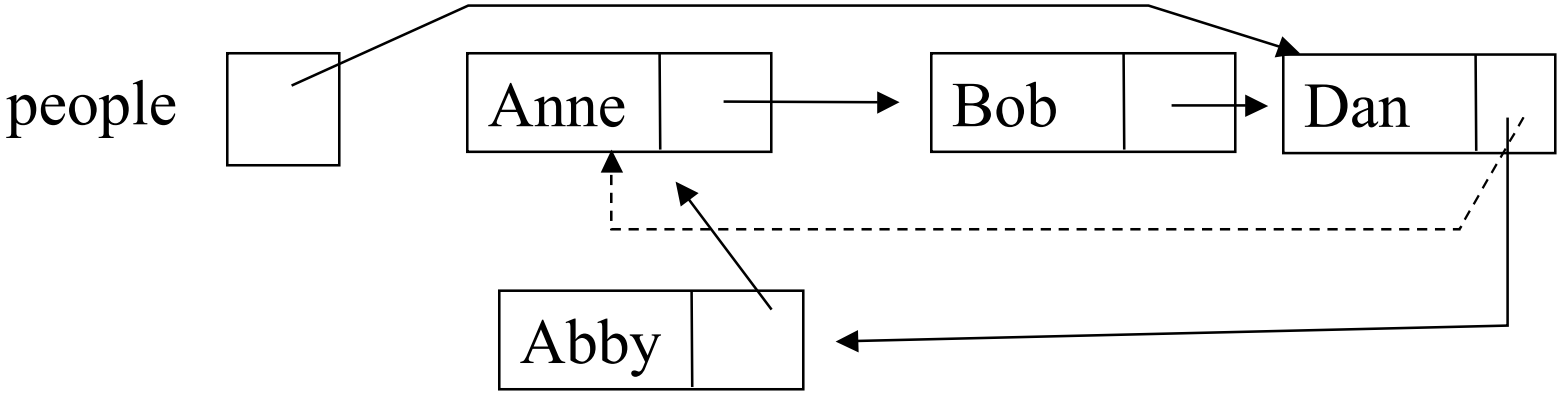
Circular List

- **First node is ??**
- **Variable place points at last node when??**

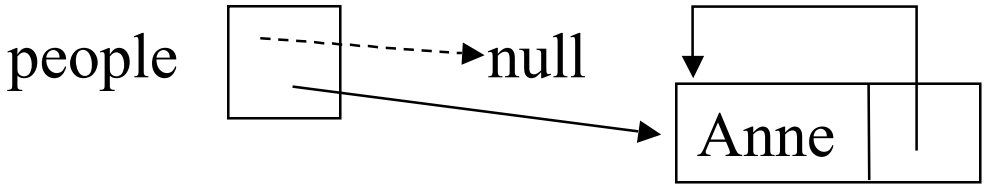


Insert at Head

- **List not empty**



- **List empty**



Insert at Head

```
if(people == null){
    people =  new Node(newName, null);
    people.next = people;
   } else {
    Node newNode= new Node(newName,
                                        people.next);
    people.next = newNode;
   }
```

Delete Head

- **You write it for circular lists**
- **Hint: 3 cases:**
  - **people empty**
  - **one node**
  - **more nodes**

Other CLL Methods

- **See resources => Java examples => fancy lists**

DLL Methods

- **See resources => Java examples => fancy lists**

Dummy Headers

- **Problem: delete head is different that delete elsewhere in list**
  - **Change pointer to list as a whole vs change the next field of some node**
- **Solution: Keep an extra "dummy" node at the head of the list**

Iterators

- **Abstract data type: a container**
  - **E.g. array or linked list**
  - **Can do mostly the same things with them, main difference is cost**
  - **Problem: one of the things I want to do is go through the data items one by one**

Processing Data Items

- **Normal way to handle same-process-different-structure problem is with a method that is defined appropriately for each class**
  - **Same name and abstract behavior**
  - **Different code**
- **But what would be the interface?**
  - **What changes from call to call? Pieces of progam** ☹

Solution

- **Instead of a method to do whole loop, have methods you can use to build the loop**
  - **hasNext**
  - **getNext**
- **State: an object**
  - **Represents a particular instance of iteration**
  - **Initialized by new**

Abstract List Traversal

- **while (list.hasNext()) {**
  print(list.getNext().data);
  }

- **list could be an Array:**
  hasNext() { return (i != list.length) }
  getNext() { i++; return list[i]; }

Abstract List Traversal

- **while (list.hasNext()) {**
  print(list.getNext().data);
  }

- **list could be a LinkedList:**
  hasNext() { return (curr != null) }
  getNext() { curr = curr.next;
  return curr; }

Iterators

- **See StringList.java and StringListIterator.java**