

Today's Content

- a. Intro to Trees
- b. Terminology
- c. Tree Traversals

Recursive Code

1. Assumption = Decide what your function does
2. Main Logic = Solving Problems using Subproblems

Note: When we call subproblem
It works according to assumption.

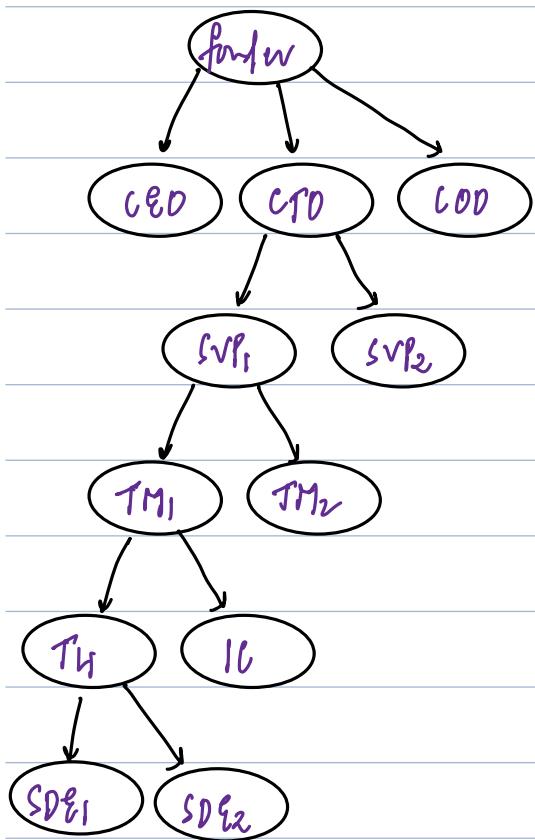
3. Base Condition = Input for which code stops.

Linear Data Structure:

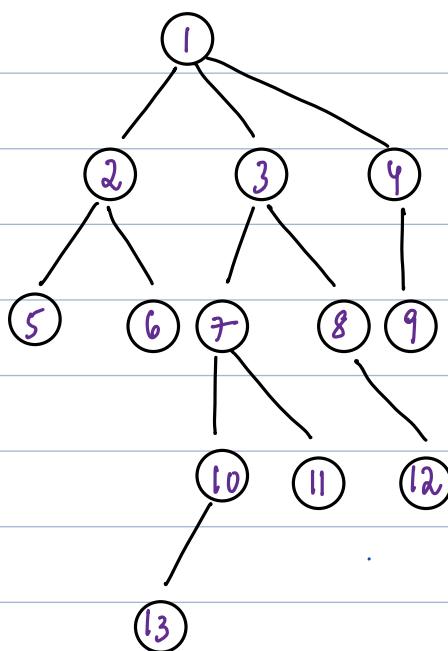
- a. Arrays
- b. Linked List
- c. Stacks / Queue

Hierarchical Data Structure: Tree

- a. Company position



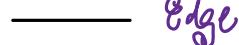
level:



Naming Convention:



Node



Edge

Parent: If node has child, it's a parent node

Ex: 1 2 3 4 ..

Child: If node has parent it's a child node

Ex: 2 3 4 ..

Root: Node without any parent

Ex: 1

Leaf: Node without any child

Ex: 5 6 9

height(Node): length of longest path from node to one of its leaf node.

$h(3)$: 3 Edges, 4 Nodes

$h(4)$: 1 Edge, 2 Nodes

$h(1)$: 0 Edges, 5 Nodes

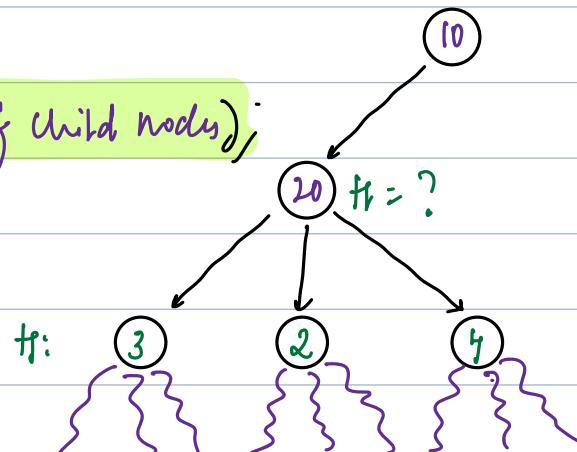
$h(12)$: 0 Edges, 1 Node

#obs

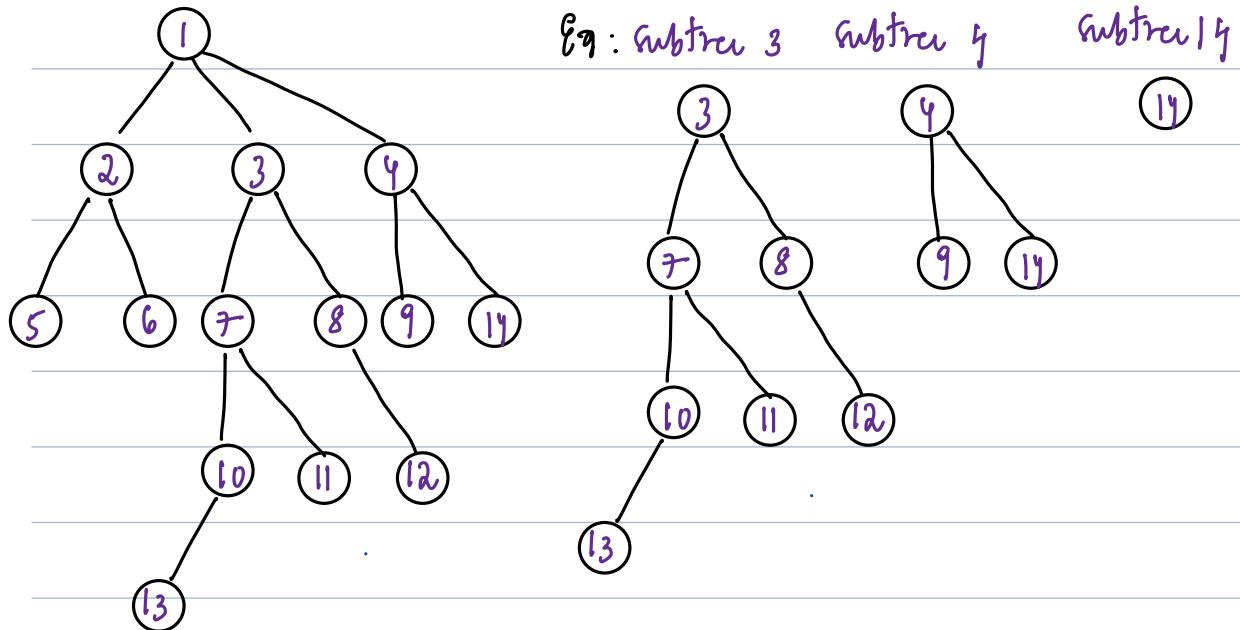
1. $h(\text{leaf})$: 0 Edge, 1 Node

2. $h(\text{Node})$: $\max(h(\text{height of child nodes}))$

3. $h(\text{Tree})$: $h(\text{root})$



#Subtree: A subtree is a subset of a tree structure that starts at a node q and contains all its descendants.



#Obs: No. of subtree is \sum # No. of nodes

Siblings: Nodes with same parents

Ancestor: All nodes along path from node to root.

Eq: 13: 13 10 7 3 1

↳ Node can be ancestor to itself in few problems

#Node can have only 1 parent, except root without a parent

Binary Tree: A Tree in which every node can at max have 2 children

0, 1, 2,

class Node{

int data;

Node *left, *right;

Node(int n){

data = n;

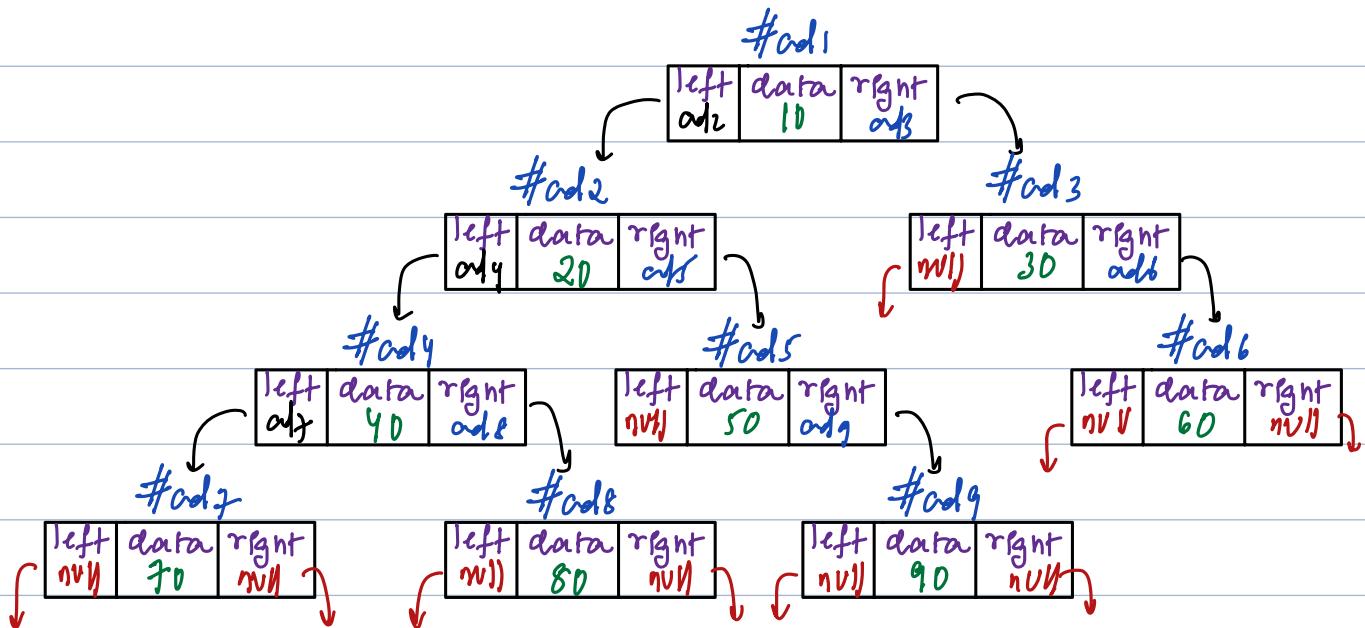
left = nullptr;

right = nullptr;

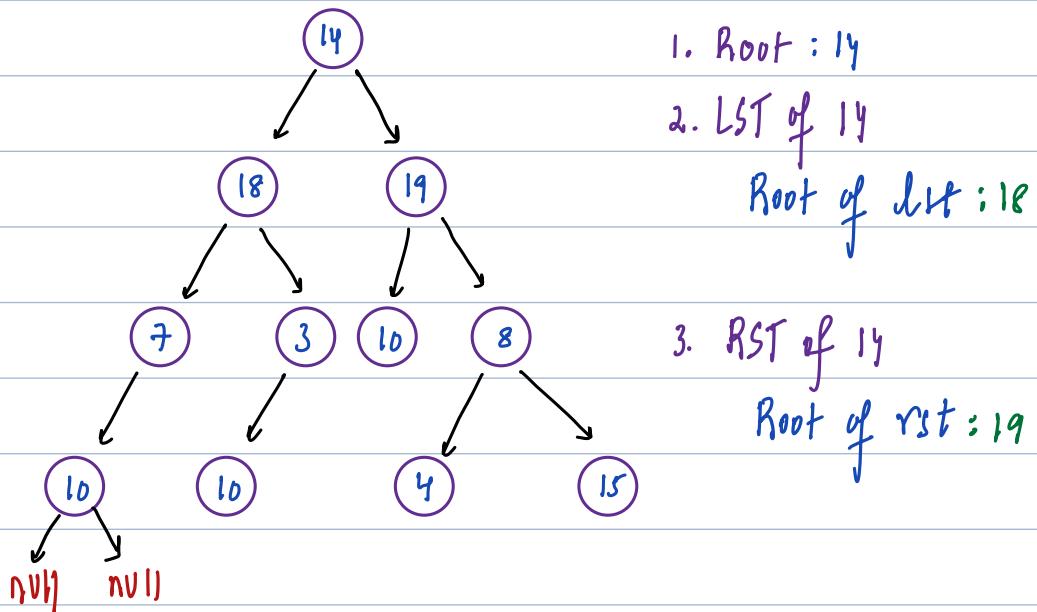
}

};

Obs: Only root node is given to access entire BT



Ex:



#Note: LST & RST concept exist for all nodes

#Note: Most of tree problems can be done with recursion.

$$\text{Ex: } \text{sum}(BT) = \text{sum}(LST) + \text{sum}(RST) + \text{root}$$

$$\text{size}(BT) = \text{size}(LST) + \text{size}(RST) + 1 \rightarrow \text{not will contribute } 1.$$

Tree Traversal

- PreOrder
- InOrder
- PostOrder

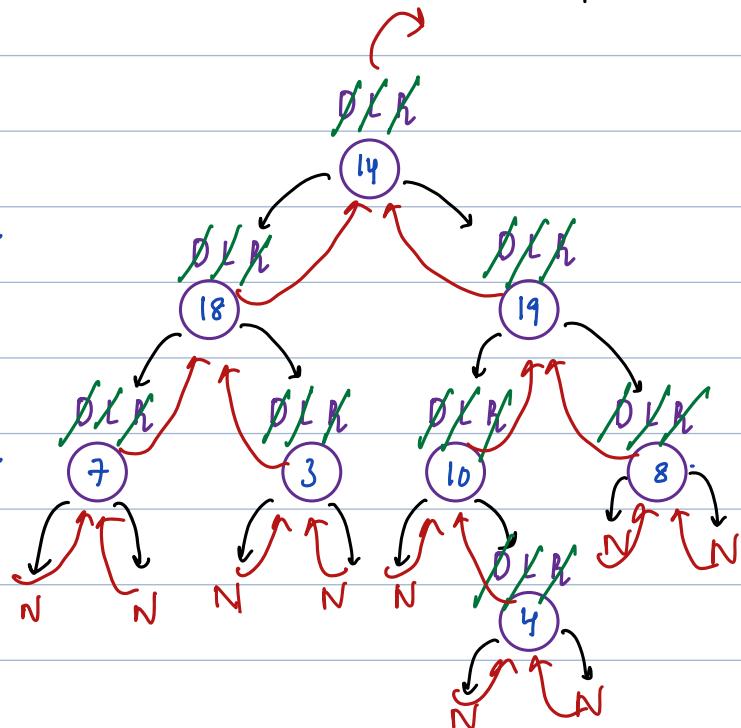
PreOrder: DLR Data Left Right

Data: Print data

left: Go to left subtree & print entire left subtree in preOrder

right: Go to right subtree & print entire right subtree in preOrder

PreTree: 19 18 7 3 19 10 4 8
root preLST preRST



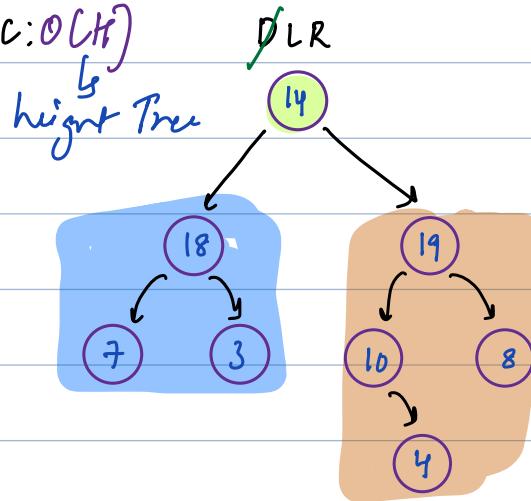
Ass: Given a root node of BT, Print entire BT in PreOrder & return nothing.

void preOrder(Node root) { TC: O(N) SC: O(H)}

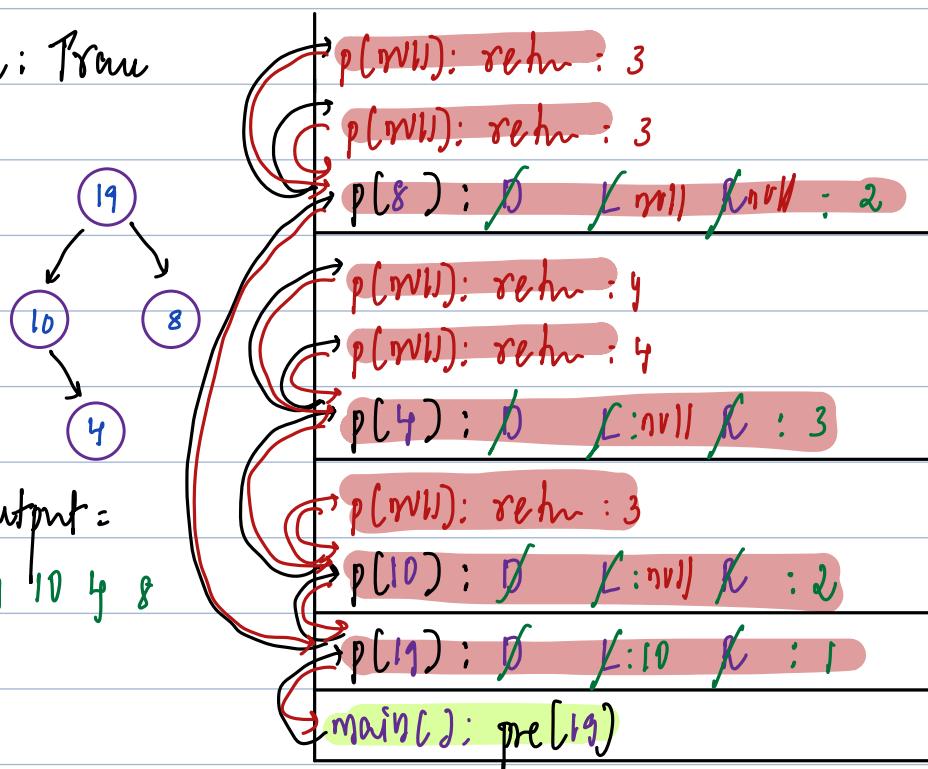
```

if[root==null] {return;}
print(root.data);
preorder(root.left); #root f LST
preorder(root.right); #root f RST
  
```

3



Ex: Print



Output =

19 10 4 8

Max Stack Size \approx height of Tree

InOrder LDR Left Data Right

left: Go to left subtree & print entire left subtree in InOrder

Data: Print data

right: Go to right subtree & print entire right subtree in InOrder

InPre: 7 18 3 14 10 4 19 8

InList not InRST

Ass: Given a root node of BT, Print entire BT in Inorder

void InOrder(Node root){

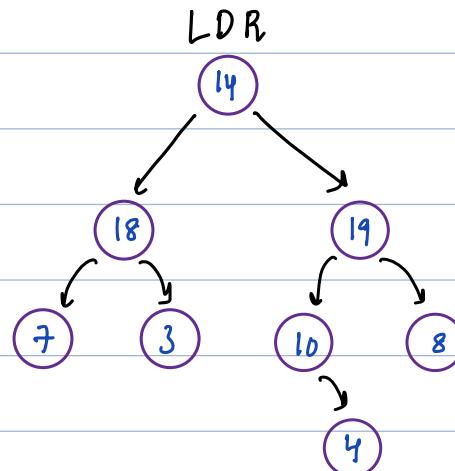
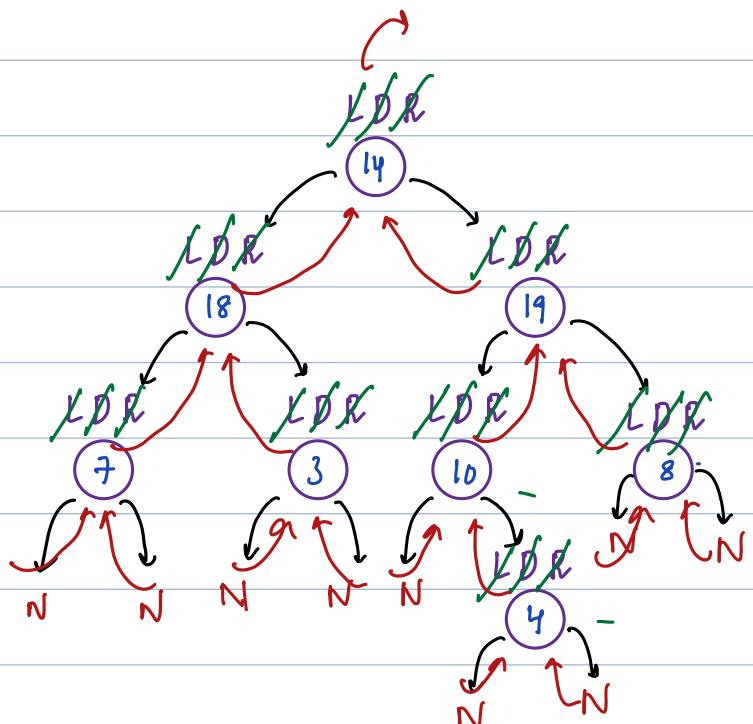
if(root==null){return;}

InOrder(root.left);

print(root.data);

InOrder(root.right);

3

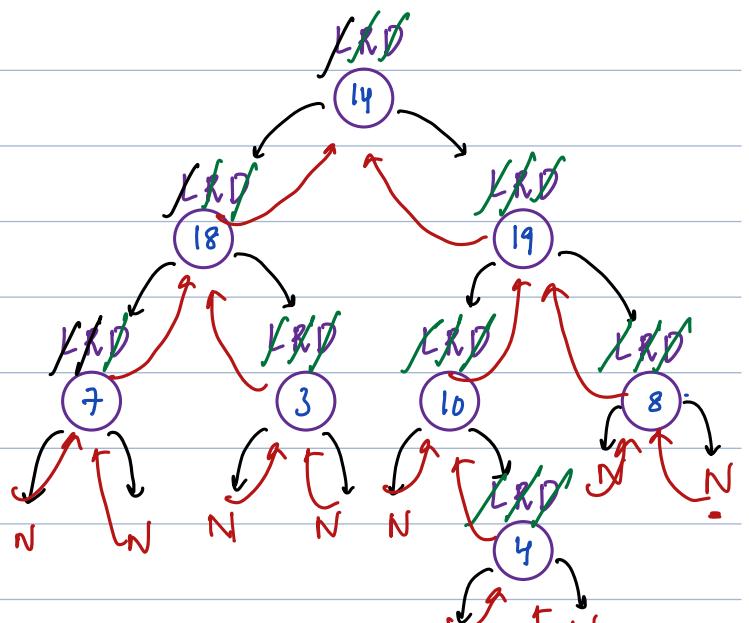


PostOrder LRD Left Right Data

Left: Go to left subtree & print entire left subtree in PostOrder

Right: Go to right subtree & print entire right subtree in PostOrder

Data: Print data



PostPre: 7 3 18 4 10 8 19 14
PostLST PostRST not

Ass: Given a root node of BT, Print entire BT in PostOrder

```
void PostOrder(Node root){  
    if(root==null){return;}  
    PostOrder(root.left);  
    PostOrder(root.right);  
    print(root.data);  
}
```

