

## Today's Content

1. Rotate array by  $k$  times
2. Vector pass by value & pass by reference.
3. Count no: of distinct elements

Q: Given an arr[], l & r reverse entire array from l..r

Ex: arr[12]  
 l=2 r=9

0	1	2	3	4	5	6	7	8	9	10	11
10	20	30	40	50	60	70	80	90	100	110	120
		100	90	80	70	60	50	40	30		

$p_1 \rightarrow p_1 \rightarrow p_1 \rightarrow p_1$   $p_2 \leftarrow p_2 \leftarrow p_2 \leftarrow p_2$

Idea:

	0	1	2	3	4	5	6	7	8	9	10	11
$p_1 < p_2$	10	20	100	90	80	70	60	50	40	30	110	120
2 < 9	swap(arr[p <sub>1</sub> ], arr[p <sub>2</sub> ]); p <sub>1</sub> ++; p <sub>2</sub> --;											
3 < 8	swap(arr[p <sub>1</sub> ], arr[p <sub>2</sub> ]); p <sub>1</sub> ++; p <sub>2</sub> --;											
4 < 7	swap(arr[p <sub>1</sub> ], arr[p <sub>2</sub> ]); p <sub>1</sub> ++; p <sub>2</sub> --;											
5 < 6	swap(arr[p <sub>1</sub> ], arr[p <sub>2</sub> ]); p <sub>1</sub> ++; p <sub>2</sub> --;											
6 < 5	$p_1 > p_2$ : stop arr[p <sub>1</sub> ..p <sub>2</sub> ] is reversed.											

void reverseRange(int arr[], int l, int r){

int p<sub>1</sub>=l, p<sub>2</sub>=r;

while(p<sub>1</sub> < p<sub>2</sub>){

//swap arr[p<sub>1</sub>] & arr[p<sub>2</sub>]

int tmp = arr[p<sub>1</sub>];

arr[p<sub>1</sub>] = arr[p<sub>2</sub>];

arr[p<sub>2</sub>] = tmp;

p<sub>1</sub>++; p<sub>2</sub>--;

}

}

Iteration: Worst Case = Entire Array = N/2

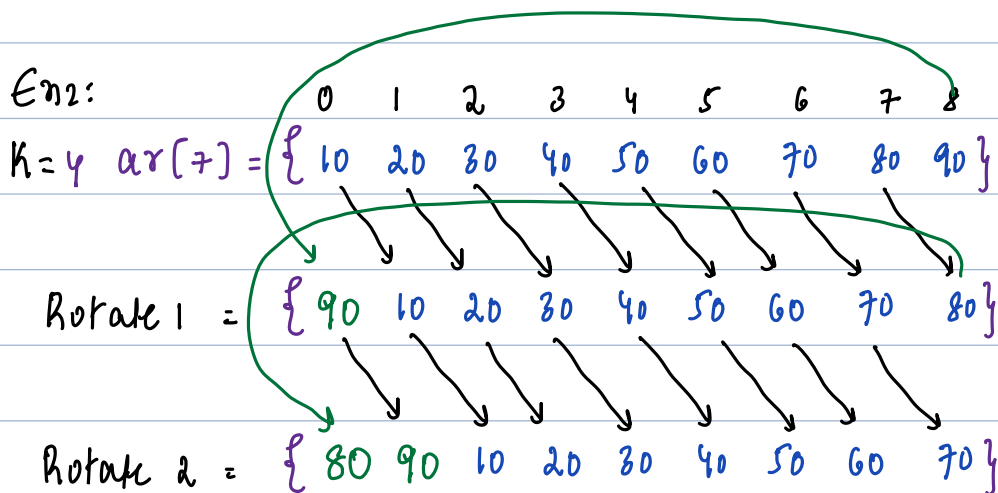
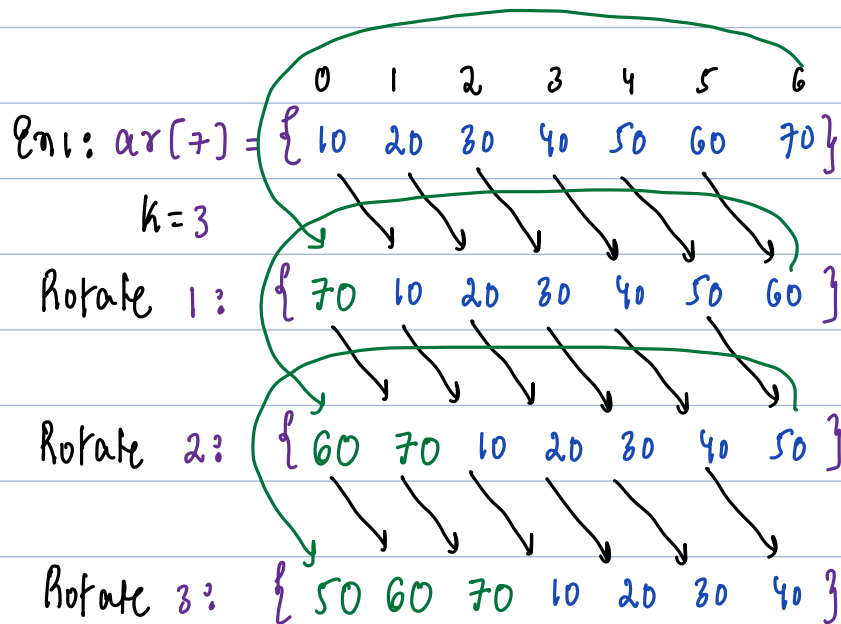
TC: O(N) SC: O(1)

Q: Given an arr[] rotate it by  $k$  times by Right  $\rightarrow$  Left

Constraints:

$$1 \leq N \leq 10^5$$

$$1 \leq K \leq 10^5$$

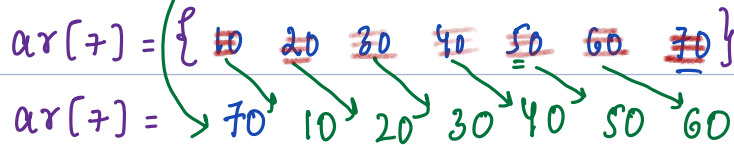


Idea: for k times:

Estimated TC:  $O(k \cdot N)$

Rotate arr[] right to left.

Single Rotate:



`int tmp = arr[6]; // 70 ✓`

Shifting ✓

`i i-1`

`arr[6] = arr[5]`

`arr[5] = arr[4]`

`arr[4] = arr[3]`

`arr[3] = arr[2]`

`arr[2] = arr[1]`

`arr[1] = arr[0]`

`arr[0] = tmp;`

`void rotateK(int arr[], int N, int k) {`

`for (int d = 0; d < k; d++) {`

`int tmp = arr[N-1]; // tmp = last element`

`for (int i = N-1; i >= 1; i--) {`

`arr[i] = arr[i-1];`

`arr[0] = tmp;`

`}`

`Inner loop =  $N \cdot k$  Outer loop =  $k$`

`TC:  $O(N \cdot k)$  SC:  $O(1)$`

`$1 \text{ s} = 10^5$`

`$1 \text{ s} = 10^5$`

`$O(10^5 \times 10^5) = 10^{10} > 10^8 \text{ TLE}$`

## Optimized Approach

Ex1:  $k=3$   $arr[] = \{10, 20, 30, 40, 50, 60, 70\}$

Rotate 1 =  $\{70, 10, 20, 30, 40, 50, 60\}$

Rotate 2 =  $\{60, 70, 10, 20, 30, 40, 50\}$

Rotate 3 =  $\{50, 60, 70, 10, 20, 30, 40\}$

$k=4$   $arr[] = \{10, 20, 30, 40, 50, 60, 70, 80, 90\}$

Rotate 1 =  $\{90, 10, 20, 30, 40, 50, 60, 70, 80\}$

Rotate 2 =  $\{80, 90, 10, 20, 30, 40, 50, 60, 70\}$

Rotate 3 =  $\{70, 80, 90, 10, 20, 30, 40, 50, 60\}$

Rotate 4 =  $\{60, 70, 80, 90, 10, 20, 30, 40, 50\}$

Obs: Rotate  $arr[N]$  by  $k$  times

1. Last  $k$  ele will shift to front

2. Remaining ele will shift to back

Ex4:  $k=5$   $arr[13] = \{a_0, a_1, a_2, a_3, a_4, a_5, a_6, a_7, a_8, a_9, a_{10}, a_{11}, a_{12}\}$

Step 1: Reverse entire  $arr[]$

0	1	2	3	4	5	6	7	8	9	10	11	12
$a_{12}$	$a_{11}$	$a_{10}$	$a_9$	$a_8$	$a_7$	$a_6$	$a_5$	$a_4$	$a_3$	$a_2$	$a_1$	$a_0$

Step 2:

Reverse  $k$  elements

Step 3:

Reverse rem ele

After Rotation =  $\{a_8, a_9, a_{10}, a_{11}, a_{12}, a_0, a_1, a_2, a_3, a_4, a_5, a_6, a_7\}$

Idea: Given  $arr[N]$  rotate  $k$  times

- Reverse entire  $arr[]$ : reverseRange( $arr, 0, N-1$ ):  $N/2$
  - Reverse first  $k$  ele: reverseRange( $arr, 0, k-1$ ):  $k/2$
  - Reverse remaining ele: reverseRange( $arr, k, N-1$ ):  $(N-k)/2$
- } Iterative:  $N$   
Tc:  $O(N)$

void reverseRange (int arr[], int l, int r) {

int p1 = l, p2 = r;

while (p1 < p2) {

int tmp = arr[p1];

arr[p1] = arr[p2];

arr[p2] = tmp;

p1++; p2--;

void rotateK (int arr[], int N, int k) {

k = k % N; // Edge Cases

reverseRange (arr, 0, N-1) → reverseRange (arr, 0, 3) ✓

reverseRange (arr, 0, k-1) → reverseRange (arr, 0, 5) ? Error. ✗

reverseRange (arr, k, N-1)

p1 = 0, p2 = 5

swap arr[0] & arr[5] // RTE.

Observation

Ex: arr[4] = 

0	1	2	3
10	20	30	40

rotate 0 = 

0	1	2	3
10	20	30	40

rotate 1 = 

0	1	2	3
40	10	20	30

rotate 2 = 

0	1	2	3
30	40	10	20

rotate 3 = 

0	1	2	3
20	30	40	10

rotate 4 = 

0	1	2	3
10	20	30	40

rotate 5 = 

0	1	2	3
40	10	20	30

rotate 6 = 

0	1	2	3
30	40	10	20

rotate 7 = 

0	1	2	3
20	30	40	10

rotate 8 = 

0	1	2	3
10	20	30	40

Rotation observation

k:

0 → 4 → 8 → 12 → 16 → 20...

1 → 5 → 9 → 13 → 17 → 21...

2 → 6 → 10 → 14 → 18 → 22 → 26

3 → 7 → 11 → 15 → 19 → 23 → 27 → 31 → 35

k N=4 ≠ k/4 ⇒ In general k % N

k	N=4
20	0
17	1
26	2
30	2
35	3

Con: Rotating k is same as Rotating k % N

void rotateK(int arr[], int N, int k){

$k = k \% N$ ; // Edge Cases

Ex:  $N=4$   $arr[4] = \{10, 20, 30, 40\}$   $k=6 \Rightarrow 6 \% 4 = 2$

reverseRange(arr, 0, N-1)  $\rightarrow$  reverseRange(arr, 0, 3) ✓

reverseRange(arr, 0, k-1)  $\rightarrow$  reverseRange(arr, 0, 1) ✓

reverseRange(arr, k, N-1)  $\rightarrow$  reverseRange(arr, 2, 3) ✓