

Todays Content

1. Count of subsets with sum = k

2. Calculate & return count of Quadruplets with sum = k

Given $arr[N]$ & k calculate q

Check if there exists a subset with $\text{sum} = k$

#Constraints

$$1 \leq N \leq 40$$

$$\# 2^{10} = 1024 \approx 10^3$$

$$1 \leq arr[i] \leq 10^7$$

$$\# (2^{10})^7 = (10^3)^7 = 10^{21}$$

$$1 \leq k \leq 10^9$$

0 1 2 3 4 5 6

Ex: $arr[] = \{3, 7, 4, 9, 2, 5\}$ $k = 10$

Idea 1: Generate all subsets:

For every subset, iterate & calculate $\text{sum} = k$

$$T.C: O(2^N * N) = O(N * 2^N) \text{ SC: } O(1)$$

$$\# N = 40, N = 40 \quad 2^{40} * 40 >> 10^8 \text{ TLE.}$$

Idea 2:

$arr[] = \{3, 7, 4, 9, 2, 5\}$ $k = 18$ 2^6 Subsets

$$arr[] = \{3, 7, 4, 9, 2, 5\} : 2^6$$

$$b[] = \{9, 2, 5\}$$

$$\{\}$$

$$\{\}$$

$$\{3\}$$

$$\{9\}$$

$$\{7\}$$

$$\{2\}$$

$$\{4\}$$

$$\{5\}$$

$$\{3, 7\}$$

$$\{9, 2\}$$

$$\{7, 9\}$$

$$\{2, 5\}$$

$$\{3, 9\}$$

$$\{9, 5\}$$

$$\{3, 7, 9\}$$

$$\{9, 2, 5\}$$

#Obs1: If we combine 1 subset in $a[\ell]$ with 1 subset in $b[\ell]$
can get subset in $c[\ell]$.

#Obs2: Can get all subsets in c ,
if we combine 1 subset in $a[\ell]$ with 1 subset in $b[\ell]$

$a[\ell] = \{3, 7, 4\}$ $b[\ell] = \{9, 2, 5\}$ $k = 18$ 2^b Subsets

$$a[\ell/2] = \{3, 7, 4\} : 2^3$$

$$b[\ell/2] = \{9, 2, 5\}$$

$2^{\ell/2}$	$\{ \} : 0$	$Sum[\ell] = [0 3 4 7 \dots]$	$0 : \{ \}$	$2^{\ell/2}$
	$\{3\} : 1$	$Sum[\ell] = [0 3 4 7 \dots]$	$9 : \{9\}$	
	$\{7\} : 2$	$Sum[\ell] = [0 2 5 7 \dots]$	$2 : \{2\}$	
α	$\{4\} : 3$		$5 : \{5\}$	
	$\{3, 7\} : 10$		$11 : \{9, 2\}$	
	$\{7, 4\} : 11$		$7 : \{2, 5\}$	
	$\{3, 4\} : 7$		$14 : \{9, 5\}$	
	$\{3, 7, 4\} : 14$		$16 : \{9, 2, 5\}$	

#Q: Check if there exists 2 elements n, y such that $n+y = k$
n is subset sum of $a[\ell]$ y is subset sum of $b[\ell]$

Ideal: Generate all pairs of n & y .

$$TC: O(2^{\ell})$$

Ideas: Store all subset sums of $b[0:n]$ in array $\text{sums}[0:n]$
Sort $\text{sums}[0:n]$

Take each subset $\text{sum}[0:n]$ of $a[0:n]$ as n .

For every n , $y = k - n$, search in sorted array $\text{sums}[0:n]$

TC: $O(2^{n/2} + n/2 * 2^{n/2} + 2^{n/2} * n/2)$ SC: $O(2^{n/2})$

TC: $O(N * 2^{n/2})$ SC: $O(2^{n/2})$

Ideas: Store all subset sums of $a[0:n]$ in $\text{sums}[0:n]$ & sort it

Store all subset sums of $b[0:n]$ in $\text{sumb}[0:n]$ & sort it

Given 2 sorted arrays, check if there exists a pair with $\text{sum} = k$.

→ 2 pointer for pair sum.

TC: $O(n/2 * 2^{n/2} + n/2 * 2^{n/2} + 2 * 2^{n/2})$

TC: $O(N * 2^{n/2})$ SC: $O(2 * 2^{n/2}) = O(2^{n/2})$

Ideas: Store all subset sums of $b[0:n]$ in HashSet hs .

Take each subset $\text{sum}[0:n]$ of $a[0:n]$ as n .

For every n , $y = k - n$, search in HashSet hs .

TC: $O(2^{n/2} + 2^{n/2})$ SC: $O(2^{n/2})$

TC: $O(2^{n/2})$ SC: $O(2^{n/2})$

Note: Talk about when to apply mid in the middle

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Given $A[], B[], C[], D[]$ no. of quadruplets (i, j, k, l) such that
 $A[i] + B[j] + C[k] + D[l] = s$
0 $\leq i, j, k, l \leq N$

$k=15$

0 1 2 3 4 5

$A[] = \{3 4 2 7 3 9\}$

4B $\# N^2$

$B[] = \{6 3 4 2 6 4\}$

$\# N^2$

$C[] = \{2 7 6 4 3 5\}$

CD $\# N^2$

$D[] = \{6 4 3 8 9 4\}$