

Today's Content

Sum of all subarrays man

Q Given $arr[N]$ sum of max of all subarrays

Ex:

$arr = \{ \overset{0}{2} \overset{1}{5} \overset{2}{3} \}$

#Subarrays #Max $A=25$

$\{2\}$

2

$\{2, 5\}$

5

$\{2, 5, 3\}$

5

$\{5\}$

5

$\{5, 3\}$

5

$\{3\}$

3

#Q1: launch in class

Idea 1

for every subarray:

iterate & calc max & add in ans.

$$TC: O(N^2 * N) = O(N^3) \quad SC: O(1)$$

Idea 2

0 1 2 3 4 5
 $arr[] = \{ 3, 2, 4, 5, 7, 6 \}$

$$max = 3 + 3 + 4 + 5 + 7 + 7 = ?$$

for every subarray:

calc max using carry forward & add in ans

$$TC: O(N^2 * 1) = O(N^2) \quad SC: O(1)$$

Idea 3: Contribution Technique:

Add contribution of each $arr[i]$.

Contribution of $arr[i] = arr[i] * \{ \# \text{No. of subarrays in which } arr[i] \text{ is max} \}$

0 1 2
 $arr[] = \{ 2, 5, 3 \}$ ans =

Subarrays

Max

$\{ 2 \}$

2

$$2 * 1 + 5 * 4 + 3 * 1 = 25$$

$\{ 2, 5 \}$

5

$\{ 2, 5, 3 \}$

5

$\{ 5 \}$

5

$\{ 5, 3 \}$

5

$\{ 3 \}$

3

Q Calculate in how many subarrays ele $arr[5] = 9$ is max?

DryRun1:

$arr[] = \{ 2, 3, 13, 7, 6, 9, 5, 3, 2, 11, 2 \}$

Subarray:

s e # Subarray = 12

3 {3 5} {3 6} {3 7} {3 8}

4 {4 5} {4 6} {4 7} {4 8}

5 {5 5} {5 6} {5 7} {5 8}

8

DryRun2:

$arr[] = \{ 2, 3, 13, 7, 6, 9, 5, 3, 2, 11, 2 \}$

#start $[P_1..i]$ #end $= [i..P_2]$

$s = i - P_1 = 3$ $e = P_2 - i = 4$

#Subarray = $s * e = 12$

P_1 = 1st greater index in left for $arr[i]$ # Default = -1

P_2 = 1st greater index in right for $arr[i]$ # Default = n

$[a, b] = \# b - a + 1$

$(a, b] = \# b - a$

$[a, b) = \# b - a$

Q find in how many subarrays ele $arr[2] = 5$ is max, based on above approach.

Dry Run 1:

i^* j P_2^*
 0 1 2 3 4 5 6 7 8
 $arr[] = \{ 7, 3, 5, 1, 2, 5, 4, 5, 9 \}$

#start $[P_1 \dots i]$	#end $[i \dots P_2]$	# subarrays
1	2	$\{1, 2\}$ $\{1, 3\}$ $\{1, 4\}$ $\{1, 5\}$ $\{1, 6\}$ $\{1, 7\}$
2	3	$\{2, 2\}$ $\{2, 3\}$ $\{2, 4\}$ $\{2, 5\}$ $\{2, 6\}$ $\{2, 7\}$
	4	
	5	
	6	
	7	

Q find in how many subarrays ele $arr[5] = 5$ is max, based on above approach.

Dry Run 1:

P_1^* i P_2^*
 0 1 2 3 4 5 6 7 8
 $arr[] = \{ 7, 3, 5, 1, 2, 5, 4, 5, 9 \}$

#start $[P_1 \dots i]$	#end $[i \dots P_2]$	
1	5	$\{1, 5\}$ $\{1, 6\}$ $\{1, 7\}$
2	6	$\{2, 5\}$ $\{2, 6\}$ $\{2, 7\}$
3	7	$\{3, 5\}$ $\{3, 6\}$ $\{3, 7\}$
4		$\{4, 5\}$ $\{4, 6\}$ $\{4, 7\}$
5		$\{5, 5\}$ $\{5, 6\}$ $\{5, 7\}$

#obs:

			[*] p_1			i			[*] p_2		
	0	1	2	3	4	5	6	7	8		
arr[] =	{	7	3	5	1	2	5	4	5	9	}

#obs1:

p_1 = 1st greater or equal index on left for arr[i]

p_2 = 1st greater index on right for arr[i].

#start [$p_1 \dots i$]

#end = [$i \dots p_2$]

3

5

{3 5} {3 6} {3 7}

4

6

{4 5} {4 6} {4 7}

5

7

{5 5} {5 6} {5 7}

Q find in how many subarrays ele arr[7]=5 is max, based on above approach

						p_1		i	p_2		
	0	1	2	3	4	5	6	7	8		
arr[] =	{	7	3	5	1	2	5	4	5	9	}

#start [$p_1 \dots i$]

#end = [$i \dots p_2$]

6

7

= {6 7} {7 7}

7

#Con: p_1 = 1st greater or equal index on left for arr[i]

p_2 = 1st greater index on right for arr[i].

or

p_1 = 1st greater index on left for arr[i]

p_2 = 1st greater or equal index on right for arr[i].

```
int main(vector<int> &arr) {
```

```
    int N = arr.size();
```

```
    # Nearest greater or equal elem left.
```

```
    stack<int> st;
```

```
    vector<int> p1(N, -1);
```

```
    for(int i=0; i<N; i++) {
```

```
        while(st.size() > 0 && arr[st.top()] <= arr[i])
```

```
            st.pop();
```

```
        if(st.size() > 0)
```

```
            p1[i] = st.top();
```

```
        st.push(i);
```

```
    }
```

```
    # Nearest greater elem right
```

```
    stack<int> st;
```

```
    vector<int> p2(N, N);
```

```
    for(int i=0; i<N; i++) {
```

```
        while(st.size() > 0 && arr[st.top()] <= arr[i])
```

```
            st.pop();
```

```
        if(st.size() > 0)
```

```
            p2[i] = st.top();
```

```
        st.push(i);
```

```
    }
```

```
    int A = 0;
```

```
    for(int i=0; i<N; i++) {
```

```
        A = A + arr[i] * (i - p1[i]) * (p2[i] - i)
```

```
    }
```

```
    return A;
```

```
}
```

