

Today's Content

a. Sum() / Size() / Max() / Height()

b. Balanced Tree()

c. Diameter Tree()

Implement below functions using recursion

```
int size(Node *root)
```

```
int sum(Node *root)
```

```
int max(Node *root)
```

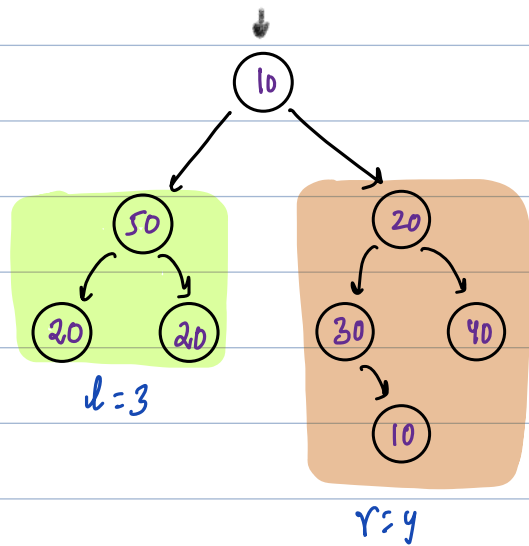
```
int height(Node *root)
```

Implement below functions

Ass: Given root of BT: calculate & return size of BT

TODO: do dry run

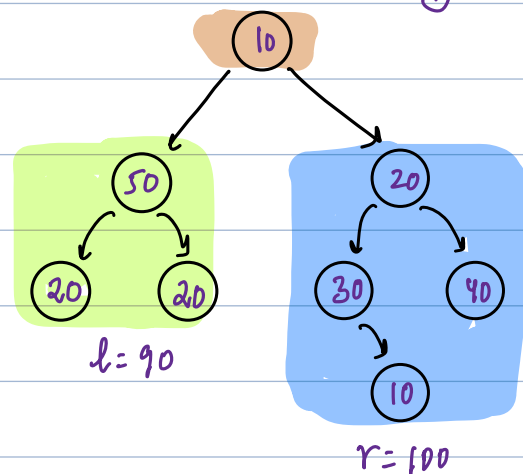
```
int size(Node root){  
    if(!root == nullptr) {  
        return 0;  
    }  
    int l = size(root->left);  
    int r = size(root->right);  
    return l+r+1;  
}
```



Ass: Given root of BT: calculate & return sum of BT

TODO: do dry run

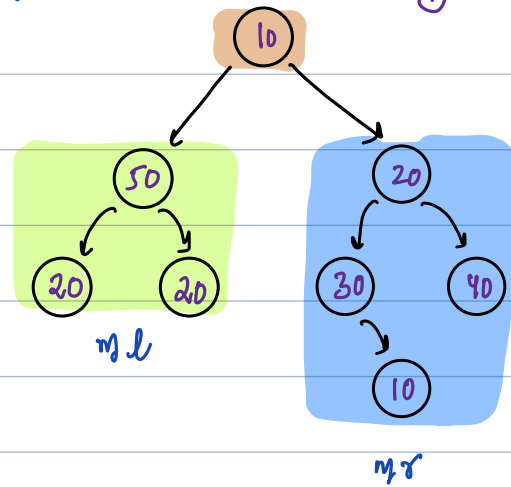
```
int sum(Node root){  
    if(!root == nullptr) {  
        return 0;  
    }  
    int l = sum(root->left);  
    int r = sum(root->right);  
    return l+r+root->data;  
}
```



Ass: Given root of BT: calculate & return min of BT

TODO: do dry run

```
int min(Node root) {
    if (root == null) {
        return INT_MIN;
    }
    int ml = min(root.left);
    int mr = min(root.right);
    return min(ml, mr, root->data);
}
```



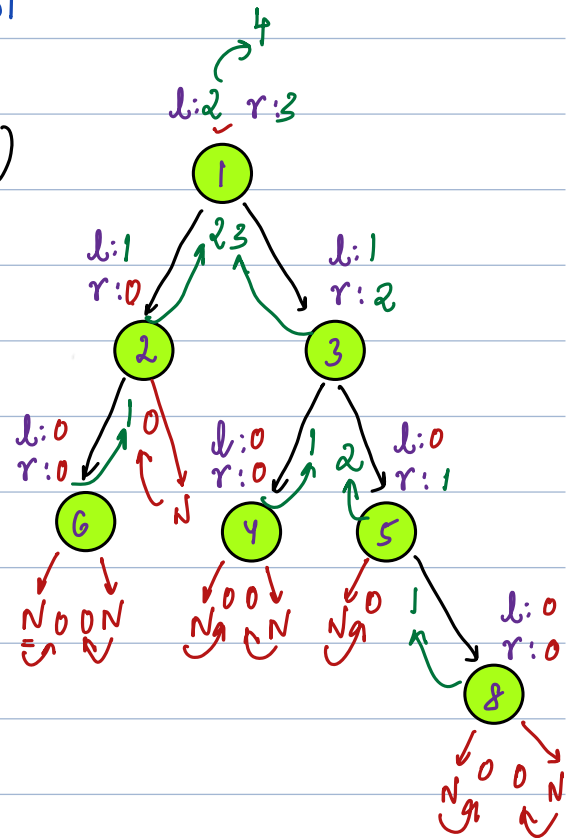
Note: Height based on nodes

Ass: Given root of BT: Calculate & return height of BT

#obs1: height of Tree = height(root)

#obs2: height(node) = 1 + max(height of child nodes)

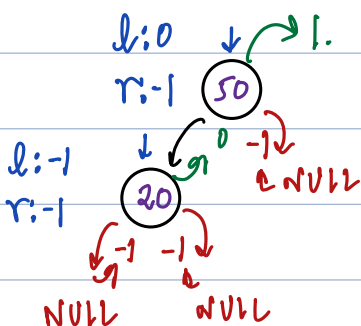
```
int height(Node root) {
    if (root == null) {
        return 0;
    }
    int lh = height(root.left);
    int rh = height(root.right);
    return max(lh, rh) + 1; // height of root node
}
```



#Note: If we want height based on edges:

```
if (root == null) { return -1; }
```

Dry Run:

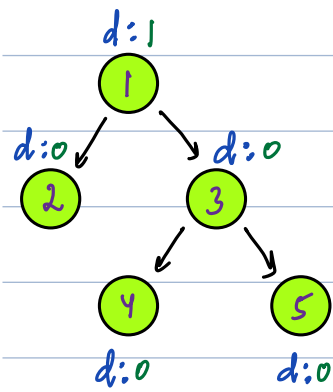


Q8 Given a BT, check if it's Balanced or Not.

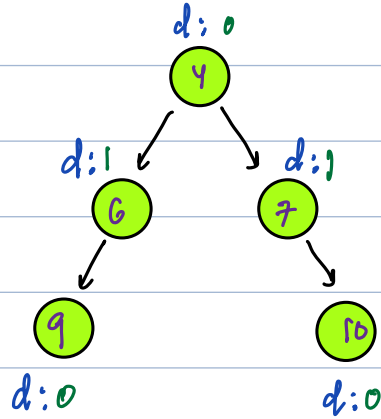
Note: A Binary Tree is said to be balanced if

for every node: $abs(height(LST) - height(RST)) \leq 1$.

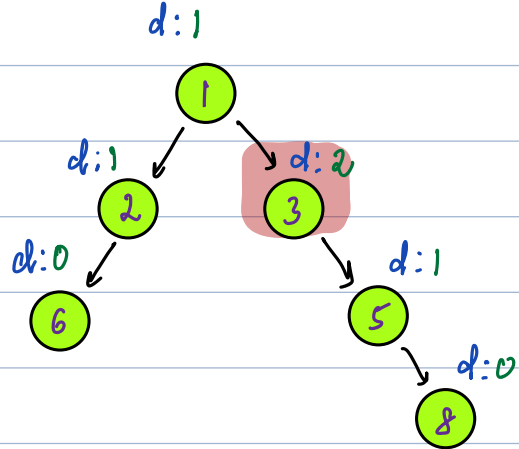
Ex1 Bal



Ex2 Bal



Ex3 Not Bal



Idea: for every node in BT:

TC: $O(N * \{N+1\}) = O(N^2)$ SC: $O(H)$

↙ stack size

$d = \text{height}(\text{node} \rightarrow \text{left});$

$r = \text{height}(\text{node} \rightarrow \text{right});$

if $(abs(d-r) > 1)$ { #Not balanced

return false;

}

return true;

Optimization: If we call $\text{height}(\text{root})$ once, it will calculate height of l & r subtree for each node?

bool isbalanced = True;

```
int height(Node root){
```

```
    if (root == nullptr) {
        return 0;
    }
```

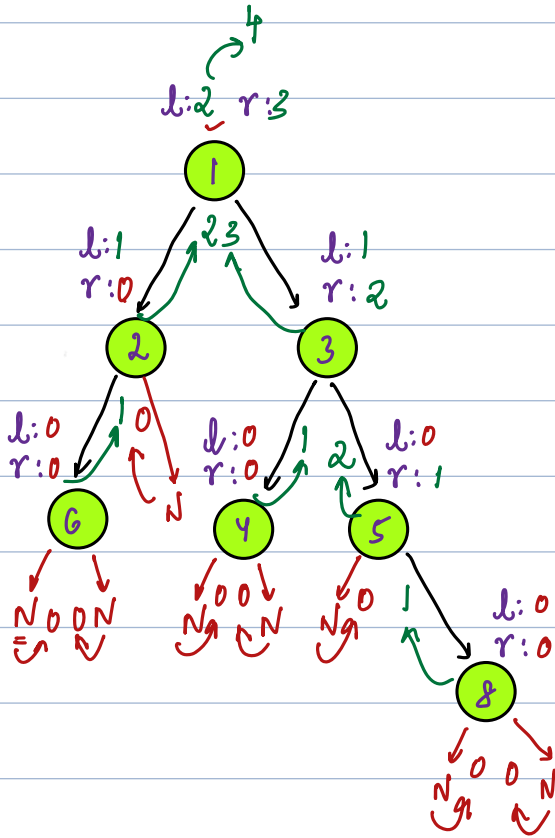
```
    int lh = height(root->left);
```

```
    int rh = height(root->right);
```

#Check if node is balanced using lh & rh

```
    if (abs(lh - rh) > 1) { isbalanced = false; }
```

```
    return max(lh, rh) + 1;
}
```



```
bool isBalancedTree(Node root){
```

```
    isbalanced = True; #reinitialize at start
```

```
    height(root);
```

```
    return isbalanced;
}
```

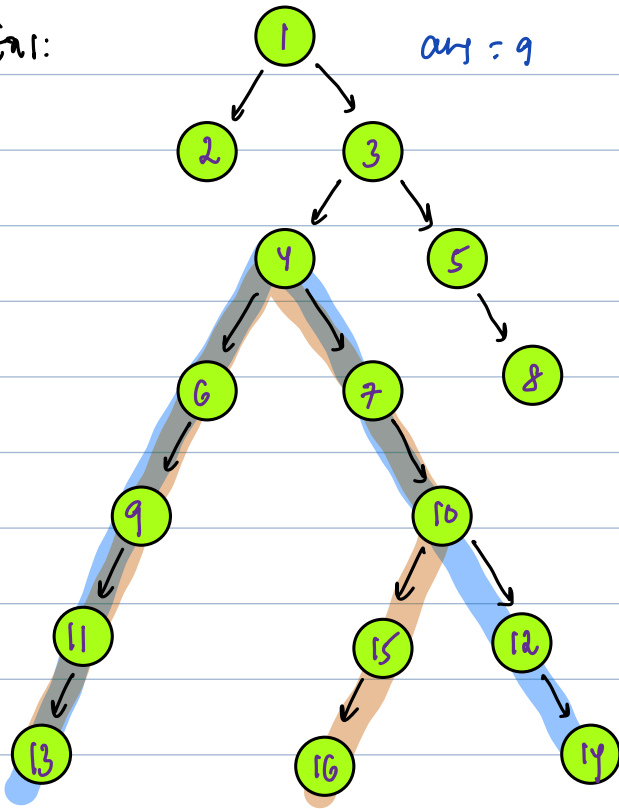
#obs:

Diameter of Binary Tree

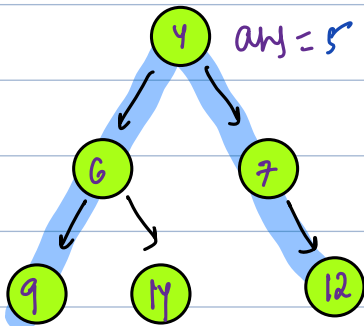
Given a binary Tree, find the length of longest path between any two nodes in the tree.

Note: Path length is calculated based on no. of nodes

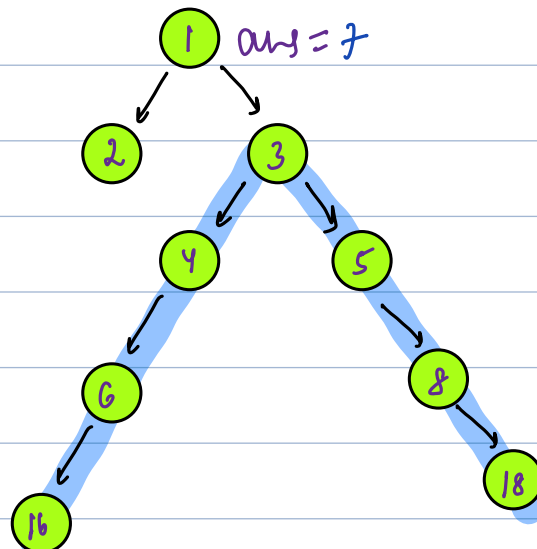
Ex1:



Ex2:

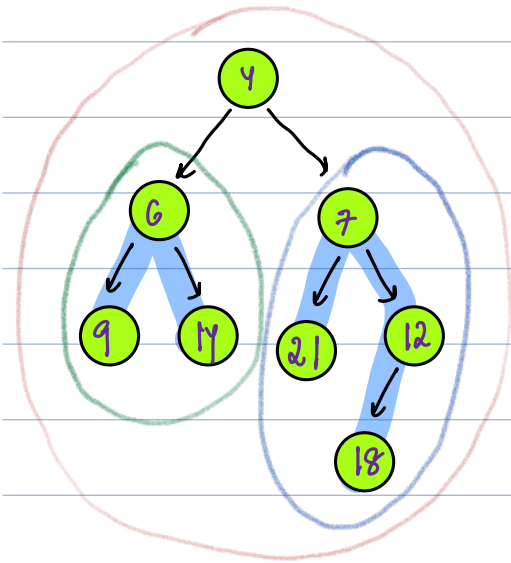


Ex3:



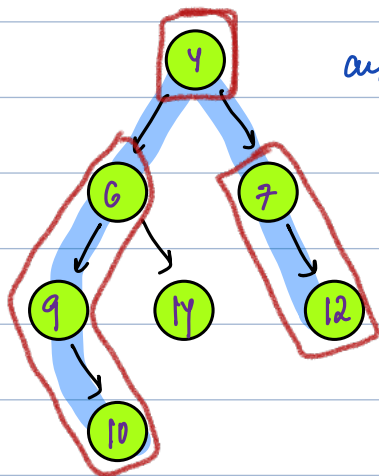
Idea: $h_{long} + \text{length of longest path in RT}$

$= \text{length of longest path in LST} +$
 $\text{length of longest path in RST} + 1$



#obs: longest path will pass through rootnode of a subtree

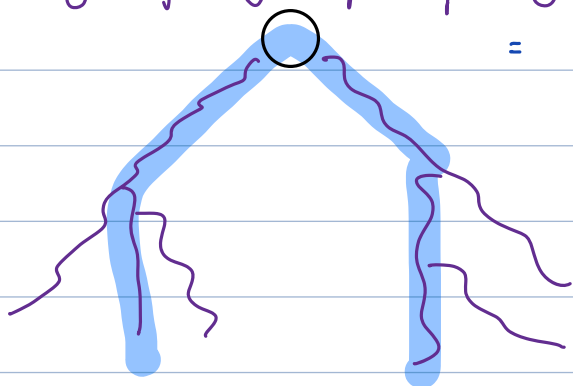
Ex: length of longest path passing through 4?



$$ans = 6 = h_L + h_R + 1 = \underline{3 + 2 + 1}$$

#con: length of longest path passing through rootnode of a subtree?

$$= \underline{\text{height(LST)} + \text{height(RST)} + 1}$$



Idea 2:

for every node:

Calculate length of longest path through node as root of subtree

$l = \text{height}(\text{root} \rightarrow \text{left})$

$r = \text{height}(\text{root} \rightarrow \text{right})$

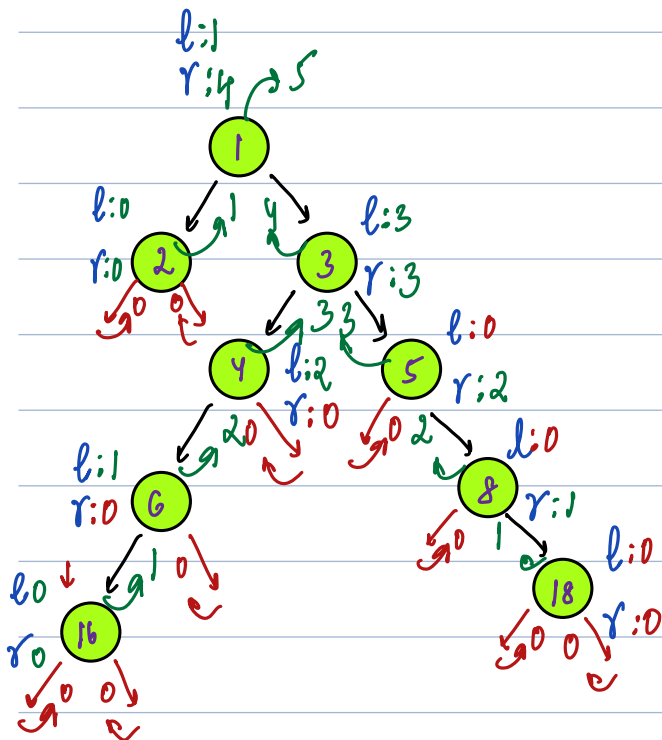
$\text{ans} = \max(\text{ans}, l+r+1);$

return ans;

T.C: $O(N*N) = O(N^2)$ S.C: $O(1)$

Ideas: If we calculate $\text{height}(\text{tree})$:

For every node we will calculate height of LST & RST.



int ans=0; T.C: $O(N)$ S.C: $O(H)$

int Height(Node root) {

if [root == NULL] { return 0; }

int l = Height(root->left);

int r = Height(root->right);

length of longest path passing through node as root of subtree l+r+1

ans = max(ans, l+r+1);

return max(l, r) + 1;

int diameter(Node root) {

ans = 0; # Reinitialize at start

Height(root);

return ans;

#obs