

# Todays Content

1. Recursion Intro
2. Steps to Recursive codes

## Use of Recursion:

1. Sorting
2. Trees : BT/BST
3. Backtracking
4. Dynamic Programming
5. Graphs

## Rules for Function:

1. Every time a function call is made, it is stored in top of stack
2. When function returns or it's complete execution it will exit stack
3. Even if function return type is void, we can write return;

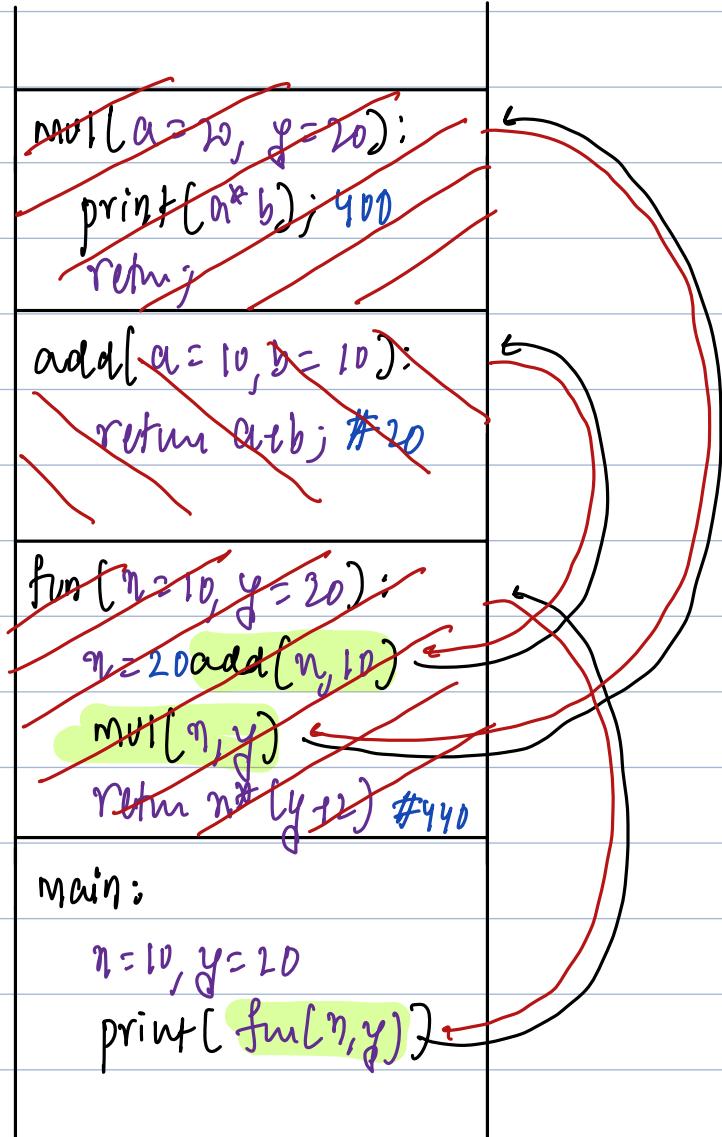
## Function Calls

```
void mul(int a, int b) {  
    print(a*b);  
    return;  
}
```

```
int add(int a, int b) {  
    return a+b;  
}
```

```
int fun(int n, int y) {  
    n = add(n, 10);  
    mul(n, y);  
    return n * (y + 2)  
}
```

```
main() {  
    int n = 10, y = 20;  
    print(fun(n, y))  
}
```



Recursion: Function calling itself

Solving Problem with Subproblem

smaller instance of same problem

$$\text{sum}(5) = \underline{1+2+3+4+5}$$

$$\text{sum}(5) = \text{sum}(4) + 5 \quad \# \text{ SubProblem} = \text{sum}(4)$$

$$\text{sum}(N) = \underline{1+2+3+\dots N-1+N}$$

$$\text{sum}(N) = \text{sum}(N-1) + N \quad \# \text{ SubProblem} = \text{sum}(N-1)$$

Steps to write Recursive Code:

Assumption: Decide what your function does  $\{\text{Input, Does, Return}\}$

Main Logic: Solve problem with subproblems.  $\{\text{Recursive Step}\}$

Bare Condition: Input for which recursion stops  $\{\text{& Stop}\}$

Ass: Calculate or return sum of  $N$  natural numbers

int sum(int N){

    if(N==1){ return 1; }

    return sum(N-1) + N

$\# 1+2+3+\dots N-1+N$

```
main() {  
    print(sum(4));
```

int sum(N=4) : a

if [N=4] { return 1; }

return sum(N-1) + N

3  
b + 4

int sum(N=3) : b

if [N=3] { return 1; }

return sum(N-1) + N

3  
3 + 3

int sum(N=2) : c

if [N=2] { return 1; }

return sum(N-1) + N

3  
1 + 2

int sum(N=1) : d

if [N=1] { return 1; }

return sum(N-1) + N

3

sum(N=1):

if N=1: return

return sum(N-1) + N

sum(N=2):

if N=2: return

return sum(N-1) + 1 + N

sum(N=3):

if N=3: return

return sum(N-1) + 3 + N

sum(N=4):

if N=4: return

return sum(N-1) + 6 + N

main()

print(sum(4)) : 10

## Importance of Base Conditions:

If base condition is not there, code will never stop : TLE Time Limit Exceeded

Before we reach TLE, will reach its memory limit, because each function call stored in stack & we will get MLE. Stack overflow.

```
int sum(int N){  
    return sum(N-1) + N  
    # 1+2+3.. N-1+N}
```

Qn: On Online servers

Ignore it

Iterating =  $10^8$

Spec limit = 1MB =  $10^6$  bytes

Int = 4B =  $10^5$  int() =  $4 \times 10^5$

long = 8B =  $10^5$  long() =  $8 \times 10^5$

Bool = 1B =  $10^6$   $\approx 1 \times 10^6$

Array size =  $[10^5 - 10^6]$

How many function calls can stack hold at once?

1. Each function call = approx 10B

2. We can have  $10^5$  function calls at once in stack

$$\text{fact}(5) = 5 * 4 * 3 * 2 * 1$$

$$\text{fact}(5) = 5 * \text{fact}(4)$$

$$\text{fact}(n) = 1 * 2 * \dots * n - 1 * n$$

$$\text{fact}(n) = \text{fact}(n-1) * n$$

# Sub Problem = fact(n-1)

Assumption: Decide what your function does # {Input, Does, Return}

Main Logic: Solve problem with subproblems. # {Recursive Step}

Note: Subproblems will work according to assumption.

Bare Condition: Input for which recursion stops # {Stop}

$$0! = 1.$$

Ass: Given N, calculate & return  $n!$

```
int fact(int n){  
    if (n == 0) { return 1; }  
    return fact(n-1) * n  
}
```

main() {

    int a = fact(3); // 6

    int fact(N=3) { : a

        if (N == 0) { return 1; }

        return fact(N-1) \* N // 3

    }

    int fact(N=2) { : b

        if (N == 0) { return 1; }

        return fact(N-1) \* N // 2

    }

    int fact(N=1) { : c

        if (N == 0) { return 1; }

        return fact(N-1) \* N // 1

    }

    int fact(N=0) {

        if (N == 0) { return 1; }

        return fact(N-1) \* N

    }

}

38 Given N, print all numbers from 1 to N in Order

Inc( $s$ ) = 1 2 3 4 5

Inc( $s$ ) = Inc( $4$ )

    print( $s$ )

Inc( $N$ ) = 1 2 3 4 .. N-1 N

Inc( $N$ ) = Inc( $N-1$ )

    print( $N$ )

Ass: Given N print all numbers from 1.. N & return nothing

void Inc(int N){

    if [N=20] { return; }

    Inc( $N-1$ )

    print( $N$ )

}

```
main() {  
    }  
    inc(3)  
    } //
```

Output: 1 2 3

```
void func( N=3 ) {
```

\* if [ N == 0 ] { return; }

- ✓ `func(N-1)`
- ✓ `print(N)`

3

```
void  fnc( n=2 ) {
```

\* if [ N == 0 ] { return; }

$$\checkmark \text{ } \mathfrak{A}_{nC}(N-1)$$

✓ print(N)

3

```
void Inc( N=1 ) {
```

\* if [ N == 0 ] { return; }

✓  $\mathfrak{A}_{nC}(N-1)$

3

void func( n=0 ) {

if [ N==20 ] { return; }

$\mathfrak{g}_{n-1} \subset \mathfrak{g}_n$

print(N)

3

48 Given  $N$ , print all numbers from  $N$  to 1 in decOrder

Dec(5) = 5 4 3 2 1

Dec(5) = Print(5)

Dec(4)

Dec(N) = N N-1 N-2 ... 1

Dec(N) = Print(N)      ↑  
                Dec(N-1)      ↘

Ass: Given  $N$  Print  $N$  to 1 return nothing

```
void Dec(int N){  
    if(N==0){return;}  
    print(N)  
    Dec(N-1)}
```

3

58 Given  $N$ , print all numbers from  $N \ N-1 \dots 1 \ 1 \ 2 \dots N$  Dec Inc Order

$\text{DecInc}(5) = 5 \ 4 \ 3 \ 2 \ 1 \ 1 \ 2 \ 3 \ 4 \ 5$

$\text{DecInc}(5) = \text{print}(5)$

$\text{DecInc}(4)$

$\text{print}(5)$

$\text{DecInc}(N) = N \ N-1 \ N-2 \dots 1 \ 1 \dots N-2 \ N-1 \ N$

$\text{DecInc}(N) = \text{print}(N)$

$\text{DecInc}(N-1)$

$\text{print}(N)$

Ass: Given  $N$  print  $N \ N-1 \dots 1 \ 1 \dots N-1 \ N$  q return nothing

void DecInc(int N){

    if( $N==0$ ) { return; }

    print(N)

    DecInc(N-1)

    print(N)