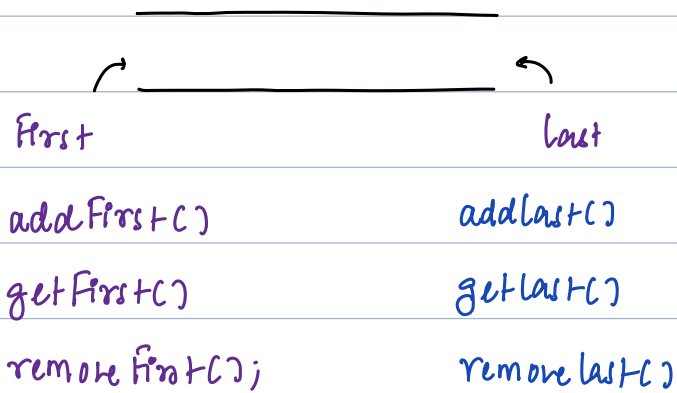


Today's Content

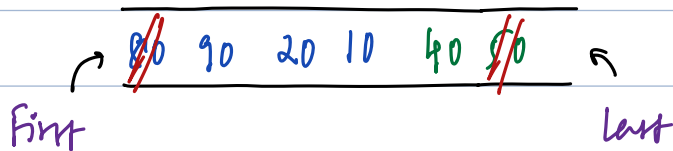
1. Deque Intro
2. Print Max of all subarrays of $len = k$
3. More questions.

Deque :



Dry Run:

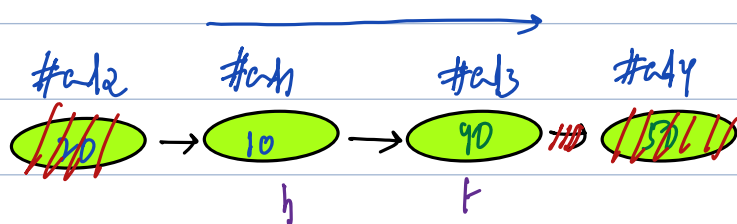
Ex: $af(10)$ ✓ $af(20)$ ✓ $al(40)$ ✓ $al(50)$ ✓ $af(80)$ ✓ $rl()$ 50 $rf()$ 80 $af(90)$ ✓



Implementation:

a. Using single linked list

Ex: $af(10)$ ✓ $af(20)$ ✓ $al(40)$ ✓ $al(50)$ ✓ $gf()$ 20 $gl()$ 50 $rf()$ 20 $rl()$ 50



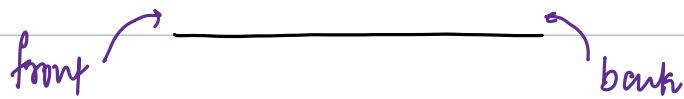
Iterate till last before node of delete of delete last node.

b. Using double linked list

In above $dlr()$ takes $O(N)$ but in dll , $dlr()$ can be done in $O(1)$ because he can traverse back in $dllr()$.

Code: TODO

Inbuilt Syntax:



C++

```
deque<type> de;
```

```
de.push-front(x);
```

```
de.front();
```

```
de.pop-front();
```

```
de.push-back(x);
```

```
de.back();
```

```
de.pop-back();
```

Java : o.c.i

```
Deque<Data-type> name = new ArrayDeque<>();
```

```
name.addFirst(x)
```

```
name.getFirst()
```

```
name.removeFirst();
```

```
name.addLast(x)
```

```
name.getLast()
```

```
name.removeLast();
```

Not sure one choice?

python

```
de = deque();
```

```
de.appendleft(x);
```

```
de.popleft();
```

```
de[0]
```

```
de.append(x); #add last
```

```
de.pop(); #delete last
```

```
de[-1]; #Access last element
```

Given $arr[N]$ & k

Print max element in every subarray of size $= k$

Ex:

	0	1	2	3	4	5	6	7	8	
$arr[] = \{$	10	1	9	3	7	6	5	11	8	$\}$

$k = 4$

Output = 10 9 9 7 11 11

Idea1: Generate all subarrays of len $= k$

Iterate on subarray & calculate max

T.C: $O(N-k+1) * O(k) = O(N^2)$ S.C: $O(1)$

Idea2: Sliding window? max

	0	1	2	3	4	5	6	7	8	
$arr[] = \{$	10	1	9	3	7	6	5	11	8	$\}$

$k = 4$

Subarray max

$\{0..3\}$ $max = 10$

$\{1..4\}$ In max : delete $arr[0]$ & add $arr[4]$.

$max += arr[4]; max -= arr[0]$ For add, inverse is sub

$max *= arr[4]; max /= arr[0]$ For mul inverse is /

$max(arr[4], max)$ For max inverse doesn't exist

Note: We can add an element in max, but we cannot remove it.

Idea2: Sliding window + TreeMap.

arr[] = { 0 1 2 3 4 5 6 7 8 }
 { 10 1 9 3 7 6 5 11 8 }

k=4

~~10~~ 1 9 3
 +

Operations	HashMap()	TreeMap
getMax()	$O(N)$	$O(\log N)$
delete(n)	$O(1)$	$O(\log N)$
insert(n)	$O(1)$	$O(\log N)$

Tc: $O((N-k+1) * \log N) \approx O(N \log N)$ Sc: $O(N)$

Idea3: Sliding window + Deque

Ex: 0 1 2 3 4 5 6 7 8 9 10 11
 arr[] = { ~~15~~ ~~12~~ ~~8~~ ~~4~~ ~~10~~ ~~9~~ ~~7~~ 12 10 7 14 3 }

k=5

Already deleted

Idea2: # All possible answers

Output

15

12

10

12

12

12

14

14

~~15~~ ~~12~~ ~~8~~ ~~4~~ ~~10~~ ~~9~~ ~~7~~ 12 10 7 14 3

Operation: deque()

a. addLast()

b. getLast()

c. popLast()

d. getFirst()

e. popFirst()

void maxSubarray (vector<int> &arr) { Tc: $O(N)$ Sc: $O(N)$

int N = arr.size();

deque<int> dq;

for (int i = 0; i < k; i++) { # 1st subarray {0..k-1}

Insert arr[i];

while (dq.size() > 0 && dq.back() < arr[i]) {

 dq.pop-back();

 dq.push-back(arr[i]);

}

print(dq.front());

int s = 1, e = k; # 2nd subarray {1..k}

while (e < N) {

Remove arr[s-1] add arr[e]

s-1 s e-1 e

if (arr[s-1] == dq.front()) {

 dq.pop-front();

while (dq.size() > 0 && dq.back() < arr[e]) {

 dq.pop-back();

 dq.push-back(arr[e]);

print(dq.front());

s++;

e++;

}

}