

Todays Content

1. Peak Element
2. Find Unique Element

Google:

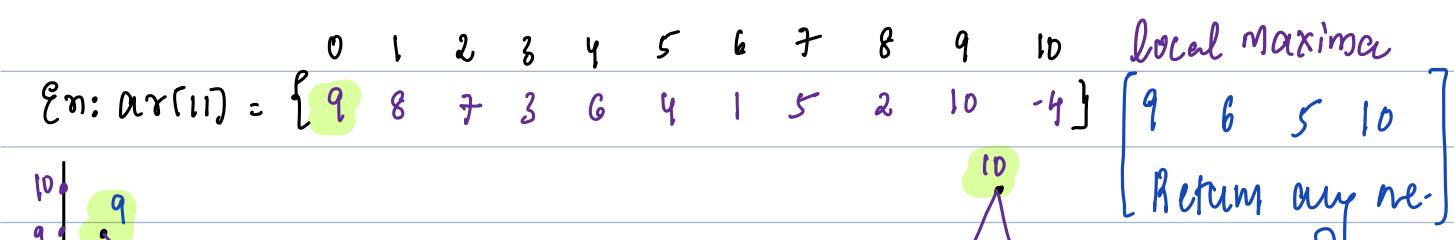
Given a unsorted arr[] with all distinct elements return anyone local maxima

local maxima: An element is said to be local maxima, if > than its adjacent elements [immediate left & right]

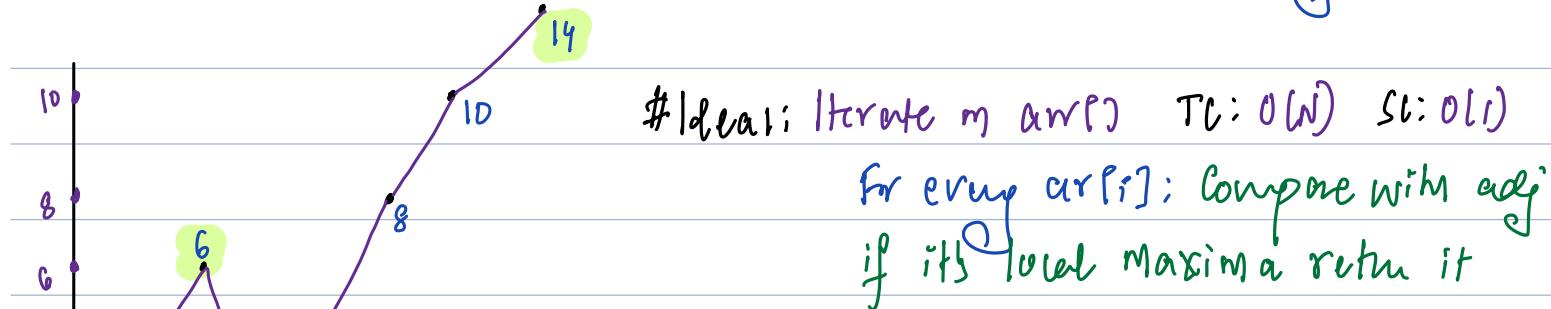
$arr[i]$ local maxima: $arr[i] > arr[i-1] \text{ & } arr[i] > arr[i+1]$

$arr[0]$ local maxima: $arr[0] > arr[1]$

$arr[N-1]$ local maxima: $arr[N-1] > arr[N-2]$



Ex: $arr[6] = \{3, 6, 2, 8, 10, 14\}$ Local Maxima = 6, 14 Return anyone



Idea: Iterate in arr[]: TC: O(N) SC: O(1)
Calculate & return max of arr[]

Ideas:

Target = Any One Local Maxima Search Space = In arr[]

Case-1

$\text{ar}[m]$ if [$\text{ar}[m-1] < \text{ar}[m]$ & $\text{ar}[m] > \text{ar}[m+1]$] {
 return $\text{ar}[m]$;
 } else
 $\text{ar}[m-1]$ $\text{ar}[m+1]$

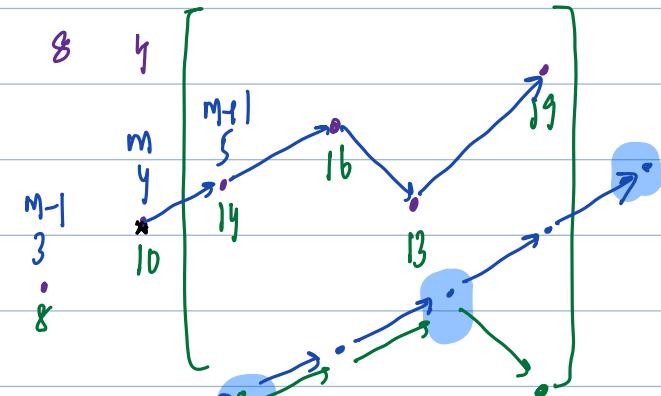
Case-2 if [$\text{ar}[m] < \text{ar}[m+1]$] { # goto right, it's a guarantee, local maxima }

$\text{ar}[m+1]$ 0 1 2 3 4 5 6 7 8 counts

6 4 12 8 10 14 16 13 19

$l \quad h \quad m$

0 8 4
3 10
8



Case-3 if [$\text{ar}[m-1] > \text{ar}[m]$] {

goto left

$\text{ar}[m-1]$

$\text{ar}[m]$

$\text{ar}[m+1]$

Case-IV: if Both sides are inc, go anywhere to get local maxima

$\text{ar}[m-1]$

$\text{ar}[m+1]$

$\text{ar}[m]$

obs: Goto side where data increases.

int localmaxima(vector<int> &arr) { Tc: $\Theta(\log_2^n)$ Sc: $\Theta(1)$

int N = arr.size();

if(N == 1) { return arr[0]; }

if(arr[0] > arr[1]) { return arr[0]; }

if(arr[N-1] > arr[N-2]) { return arr[N-1]; }

int l = 1, h = N-2;

while(l <= h) {

int m = (l+h)/2;

if(arr[m-1] < arr[m] && arr[m] > arr[m+1]) { return arr[m]; }

else if(arr[m] < arr[m+1]) { # goto right

 l = m+1;

else { # arr[m-1] > arr[m] goto left

 h = m-1;

}

}

Note: Comparing only elements,
take care of corner elements

if $m == 0$: arr[-1] < arr[0]: Err.

if $m == N-1$: arr[N-1] > arr[N]: Err.

Q: Google

Given a unsorted arr[] return any one local maxima

local maxima: An element is said to be local maxima, if \geq than its adjacent elements [immediate left & right]

$$\underline{arr[i] \geq arr[i-1]} \text{ & } \underline{arr[i] \geq arr[i+1]}$$

If $i=0$: $arr[0]$ is local maxima $\Rightarrow arr[0] \geq arr[1]$

If $i=N-1$: $arr[N-1]$ is local maxima $\Rightarrow arr[N-1] \geq arr[N-2]$

int localmaximaRepeating (vector<int> &arr) {

Q. Every element occurs twice, except for 1, find a unique element

Note: Duplicates are adjacent to each other

Ex1: 0 1 2 3 4 5 6 7 8 9 10

$\text{arr}[] = \{ 6, 6, 2, 2, 7, 9, 9, 4, 4, 10, 10 \}$ ans = 7

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14

$\text{arr}[] = \{ 3, 3, 1, 1, 8, 8, 10, 10, 19, 6, 6, 2, 2, 4, 4 \}$ ans = 19

Idea1: 1. Calculate XOR of all elements

2. Iterate in arr[]: Compare adj elements & return unique el

TC: $O(N)$ SC: $O(1)$

Idea2:

Target: Unique element SearchSpace: In arr[]

Discard?

$\text{arr}[] = \{ 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14 \}$

if m is here:

goto right

if 1st occurrence is Even:

On left

goto right

if m is here

goto left

if 1st occurrence is Odd:

On right

goto left

Note: It is possible that m can find m 2nd occurrence

Con: if m is m 2nd occ: Bring to 1^m

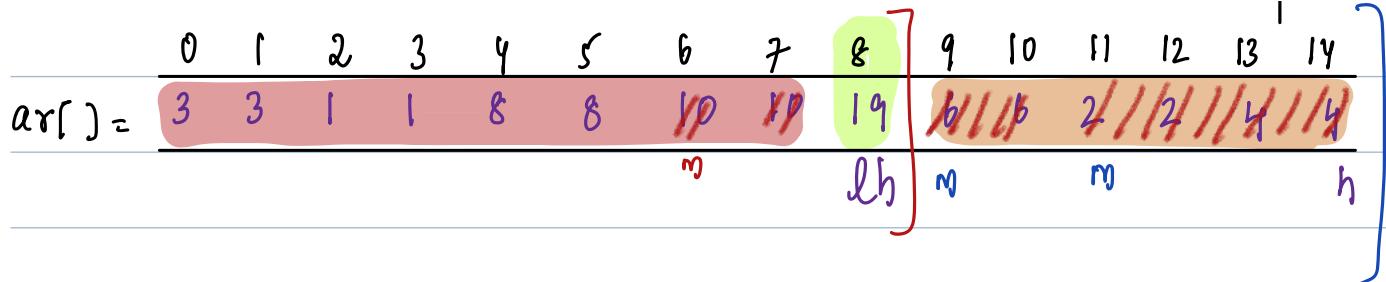
if ($\text{arr}[m-1] == \text{arr}[m]$) {

$m = m - 1$

$\text{arr}[m-1] \leftarrow \text{arr}[m]$

1^m 2nd

3



$l \quad h \quad m \quad l^m \text{ or } m^h \text{ occurrence}$

Even or Odd

- | | | | | |
|---|----|----|--|--|
| 0 | 14 | 7 | $ar[m] == ar[m-1]$; $m = m-1$, $m = 6$ | M: Even : goto right, $l = m+2$; |
| 8 | 14 | 11 | 1 st occurrence | $m = 11$ M: Odd : goto left, $h = m-1$; |
| 8 | 10 | 9 | 1 st occurrence | $m = 9$ M: Odd : goto left, $h = m-1$; |
| 8 | 8 | 8 | Unique ele; return $ar[m]$; | |

int unique(vector<int> &ar){ TC: $O(\log_2 N)$ SC: $O(1)$

int N = ar.size();

if(N == 1) { return ar[0]; }

if(ar[0] != ar[1]) { return ar[0]; }

if(ar[N-1] != ar[N-2]) { return ar[N-1]; }

int l = 1, h = N-2

while(l <= h)

int m = (l+h)/2;

if(ar[m-1] != ar[m] && ar[m] != ar[m+1]) { return ar[m]; }

if $m == 0$: $ar[-1] \neq ar[0]$

if $m == N-1$: $ar[N-1] \neq ar[N]$

if($ar[m-1] == ar[m]$): # m m 2nd occ: Bring to 1st occ

$m = m-1;$

if($m \% 2 == 0$): # On left

$l = m+2;$ # goto right

else if($m \% 2 == 1$): m right

$h = m-1;$

