

## Today's Content

1. Recursion Intro
2. Steps to Recursive codes

### Use of Recursion:

1. Sorting
2. Trees : BT/BST
3. BackTracking
4. Dynamic Programming
5. Graphs

## Rules for function:

1. Every time a function call is made, it is stored on top of stack
2. When function returns or its complete execution it will exist stack
3. Even if function return type is void, we can write return;

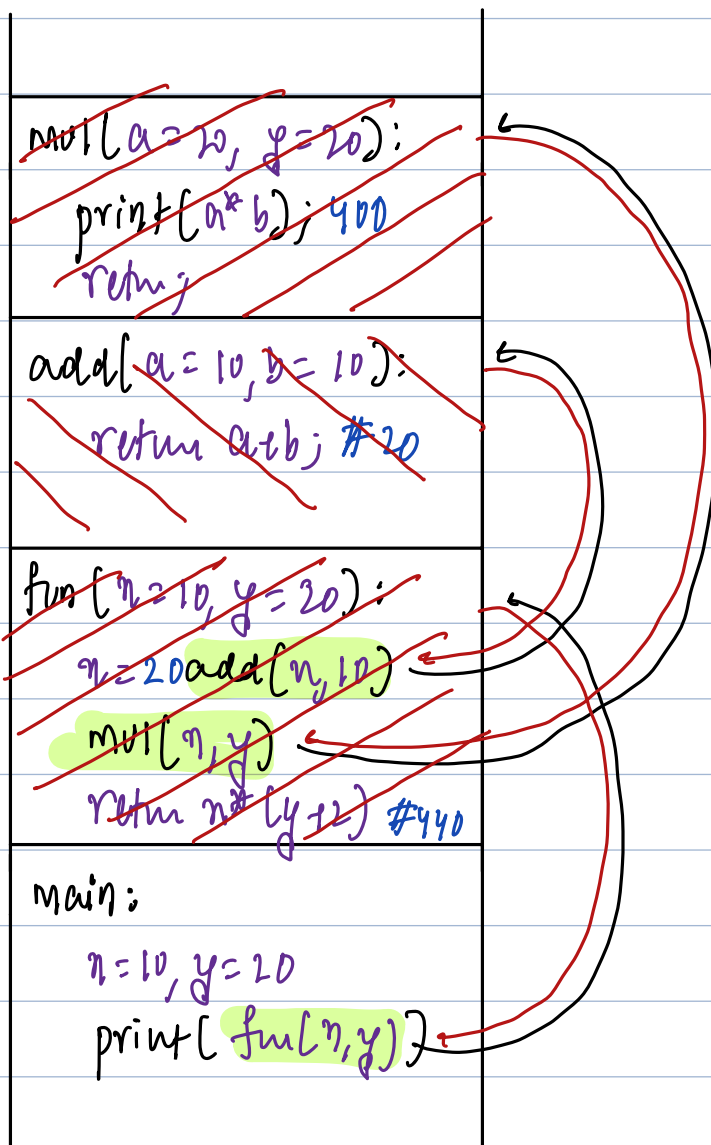
## function calls

```
void mul(int a, int b) {  
    print(a*b);  
}
```

```
int add(int a, int b) {  
    return a+b;  
}
```

```
int fun(int x, int y) {  
    x = add(x, 10);  
    mul(x, y);  
    return x*(y+2);  
}
```

```
main() {  
    int x=10, y=20;  
    print(fun(x, y))  
}
```



Recursion: Function calling itself

Solving Problem with Subproblem

smaller instance of same problem

$$\text{sum}(5) = 1+2+3+4+5$$

$$\text{sum}(5) = \text{sum}(4) + 5 \quad \# \text{ Sub Problem} = \text{sum}(4)$$

$$\text{sum}(N) = 1+2+3+\dots N-1+N$$

$$\text{sum}(N) = \text{sum}(N-1) + N \quad \# \text{ Sub Problem} = \text{sum}(N-1)$$

Steps to write Recursive Code:

Assumption: Decide what your function does # {Input, Does, Return}

Main logic: Solve problem with subproblems. # {Recursive step}

Base Condition: Input for which recursion stops # {stop}

Ass: Calculate & return sum of 1<sup>st</sup> N natural numbers

```
int sum(int N){  
    if(N==1) { return 1; }  
    return sum(N-1) + N  
    # 1+2+3+... N-1 + N  
}
```

```
main() {  
    print(sum(4));  
}
```

```
int sum(N=4) { : a  
    if(N==1) { return 1; }  
    return sum(N-1) + N  
}
```

```
int sum(N=3) { : b  
    if(N==1) { return 1; }  
    return sum(N-1) + N  
}
```

```
int sum(N=2) { : c  
    if(N==1) { return 1; }  
    return sum(N-1) + N  
}
```

```
int sum(N=1) { : d  
    if(N==1) { return 1; }  
    return sum(N-1) + N  
}
```

```
sum(N=1):  
    if(N==1) { return 1; }  
    return sum(N-1) + N
```

```
sum(N=2):  
    if(N==1) { return 1; }  
    return sum(N-1) + 1 + N
```

```
sum(N=3):  
    if(N==1) { return 1; }  
    return sum(N-1) + 3 + N
```

```
sum(N=4):  
    if(N==1) { return 1; }  
    return sum(N-1) + 6 + N
```

```
main()  
    print(sum(4));
```

## Importance of Base Conditions:

If base condition is not there, code will never stop : TLE Time Limit Exceeded

Before we reach TLE, will reach its memory limit, because each function call stored in stack & we will get MLE. Stack overflow.

```
int sum(int N){  
    return sum(N-1) + N  
}
```

#  $1+2+3 \dots N-1 + N$

Q: On online servers

Improve it

$$\text{Iterations} = 10^8$$

$$\text{Space limit} = 1MB = 10^6 \text{ Bytes}$$

$$\text{int} = 4B = 10^5 \text{ int}[] = 4 \times 10^5$$

$$\text{long} = 8B = 10^5 [ ] = 8 \times 10^5$$

$$\text{bool} = 1B = 10^6 \approx 1 \times 10^6$$

$$\text{Array size} = [10^5 - 10^6]$$

How many function calls can stack hold at once?

1. Each function call = approx 10B

2. We can have  $10^5$  function calls at once in stack

$$\text{fact}(5) = 5 * 4 * 3 * 2 * 1$$

$$\text{fact}(5) = 5 * \text{fact}(4) \quad \# \text{ Sub Problem} = \text{fact}(4)$$

$$\text{fact}(N) = 1 * 2 * \dots * N-1 * N$$

$$\text{fact}(N) = \text{fact}(N-1) * N \quad \# \text{ Sub Problem} = \text{fact}(N-1)$$

Assumption: Decide what your function does # {Input, Does, Return}

Main logic: Solve problem with subproblems. # {Recursive step}

Note: Subproblems will work according to assumption.

Base Condition: Input for which recursion stops # {stop}

$$0! = 1.$$

Ass: Given  $N$ , calculate & return  $N!$

```
int fact(int N){  
    if (N == 0) { return 1; }  
    return fact(N-1) * N  
}
```

main() {

int a = fact(3);

3

int fact(N=3) {

if (N == 0) { return 1; }

return fact(N-1) \* N;

3

2

3

int fact(N=2) {

if (N == 0) { return 1; }

return fact(N-1) \* N;

3

1

2

int fact(N=1) {

if (N == 0) { return 1; }

return fact(N-1) \* N;

3

1

1

int fact(N=0) {

if (N == 0) { return 1; }

return fact(N-1) \* N;

3

3Q Given  $N$ , print all numbers from 1 to  $N$  in incOrder

$\text{Inc}(5) = 1 \ 2 \ 3 \ 4 \ 5$

$\text{Inc}(5) = \text{Inc}(4)$   
 $\text{print}(5)$

$\text{Inc}(N) = 1 \ 2 \ 3 \ 4 \dots N-1 \ N$

$\text{Inc}(N) = \text{Inc}(N-1)$   
 $\text{print}(N)$

Ass: Given  $N$  print all numbers from 1..  $N$  & return nothing

```
void Inc(int N){  
    if (N == 0) { return; }  
    Inc(N-1)  
    print(N)  
}
```

}



main() {

Output: 1 2 3

inc(3)

void inc(N=3) {

\* if(N==0) { return; }

✓ inc(N-1)

✓ print(N)

void inc(N=2) {

\* if(N==0) { return; }

✓ inc(N-1)

✓ print(N)

void inc(N=1) {

\* if(N==0) { return; }

✓ inc(N-1)

✓ print(N)

void inc(N=0) {

if(N==0) { return; }

inc(N-1)

print(N)

48 Given  $N$ , print all numbers from  $N$  to 1 in decOrder

$\text{Dec}(5) = 5 \ 4 \ 3 \ 2 \ 1$

$\text{Dec}(5) = \text{Print}(5)$

$\text{Dec}(4)$

$\text{Dec}(N) = N \ N-1 \ N-2 \ \dots \ 1$

$\text{Dec}(N) = \text{Print}(N)$

$\text{Dec}(N-1)$

Ass: Given  $N$  Print  $N$  to 1 return nothing

```
void Dec(int N){
```

```
    if(N==0){return;}
```

```
    print(N)
```

```
    Dec(N-1)
```

```
}
```

58 Given  $N$ , print all numbers from  $N\ N-1\ \dots\ 1\ 1\ 2\ \dots\ N$  Dec Inc Order

DecInc(5) = 5 4 3 2 1 1 2 3 4 5

```
DecInc(5) = print(5)
            DecInc(4)
            print(5)
```

DecInc( $N$ ) =  $N\ N-1\ N-2\ \dots\ 1\ 1\ \dots\ N-2\ N-1\ N$

```
DecInc(N) = print(N)
            DecInc(N-1)
            print(N)
```

Ass: Given  $N$  print  $N\ N-1\ \dots\ 1\ 1\ \dots\ N-1\ N$  & return nothing

```
void DecInc(int N){
    if(N==0) {return;}
    print(N)
    DecInc(N-1)
    print(N)
}
```

