

Today's Content:

Attendance = 22

1. Nearest smaller element in left
2. Nearest smaller element in left

nearest element on left side:

Given $ar[N]$, for every $ar[i]$: find nearest smaller element on left side
Nearest smaller $ar[i]$, distance between indices

Ex1 0 1 2 3 4 5

$ar[6] = 4 5 2 10 3 12$

$ans[6] = -1 4 -1 2 2 3$

Ex2 0 1 2 3 4 5 6 7

$ar[8] = 4 6 10 11 7 8 3 5$

$ans[8] = -1 4 6 10 6 7 -1 3$

Ideal: for every $ar[i]$:

Iterate on left: find i^{th} element $< ar[i]$.
0 .. $i-1$ i

`vector<int> nearestElement(int ar[], int n) { TC: O(N^2) SC: O(1)`

`vector<int> ans(N, -1);`

`for (int i=0; i < N; i++) {`

for every $ar[i]$ calculate i^{th} no $< ar[i]$ on left

`for (int j=i-1; j >= 0; j--) {`

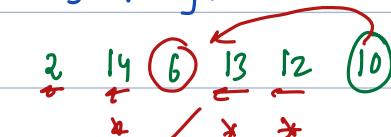
`if (ar[j] < ar[i]) {`

`ans[i] = ar[j];`

`break;`

`}`

`return ans;`



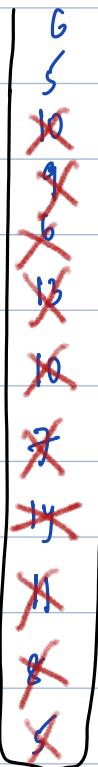
Ideas

hint: { 8 14 11 10 5 11 }

With presence of 5, 8 cannot
be ans for future elements.

obs:

	0	1	2	3	4	5	6	7	8	9	10	11
ar[12] =	5	8	11	14	7	10	13	6	9	10	5	6
ans[12] =	-1	5	8	11	5	7	10	5	6	9	-1	5



Containers:

1. Any element which can be ans, store it
2. If we are 100% sure, it cannot be ans, delete it.

Operations:

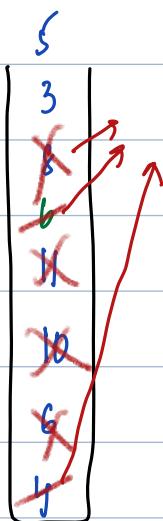
1. Insert top()
2. delete top()
3. Ans top()
4. size()

Ans: Stack data structure

By Runa: 0 1 2 3 4 5 6 7

$ar[8] = 4 6 10 11 6 8 3 5$

$ans[] = -1 4 6 10 4 6 -1 3$



stack<int> st;

element $ar[i]$ comes:

while($st.size() > 0$ $\text{and } st.top() > ar[i]$) {

} $st.pop()$;

if($st.size() > 0$) {

} $ans[i] = st.top()$;

else {

} $ans[i] = -1$

$st.push(ar[i])$

$\text{vector<int> nearestElement(vector<int> arr)}$ & $TC: O(N)$ $SC: O(N)$

stack<int> st;

int N = arr.size();

vector<int> ans(N, -1);

for(int i=0; i < N; i++) {

while($st.size() > 0$ $\text{and } st.top() > ar[i]$) {

} $st.pop()$;

if($st.size() > 0$) {

} $ans[i] = st.top()$;

else {

} $ans[i] = -1$

$st.push(ar[i])$

return ans;

Total Iterations: $N + N = 2N$

Total Outerloop + Total Innerloop.

(N)

1 pop = 1 inner iteration.

Total push = N

Total pop = N

Total Innerloop = N

Nearest smaller Index in left:

	0	1	2	3	4	5	6	7
arr[8]	4	6	10	11	7	8	3	5
ans[8]	-1	0	1	2	1			

4:7
3:11
2:10
1:6
0:4

Containers:

1. tiny element which can be ans, store its index.
2. if we are 100% sure, it cannot be ans, delete it.

```
vector<int> nearestElementIndex(vector<int> arr){
```

```
stack<int> st;
```

```
int N = arr.size();
```

```
vector<int> ans(N, -1);
```

```
for(int i=0; i<N; i++) {
```

```
    while(st.size() > 0 && arr[st.top()] >= arr[i]) {
```

```
        st.pop();
```

```
    if(st.size() > 0) {
```

```
        ans[i] = st.top();
```

```
    else {
```

```
        ans[i] = -1;
```

```
        st.push(i);
```

```
    }
```

```
}
```

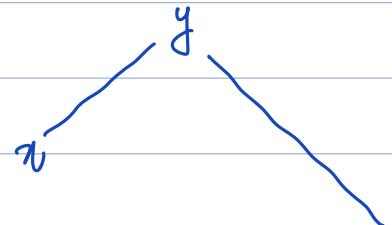
Nearest Greater on Left Side:

Ex:

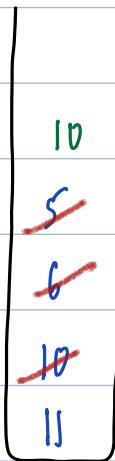
arr = 0 1 2 3 4 5 6 7

arr[8] = 11 10 6 5 10 7 4 8

ans[] = -1 11 10 6



#obs: if y is present \Rightarrow
 y cannot be ans.



Idea:

while st.top() \leq arr[i] {
 st.pop()
}

vector<int> nearestElement(vector<int> arr) {

stack<int> st;

int N = arr.size();

vector<int> ans[N, -1];

for (int i = 0; i < N; i++) {

 while (st.size() > 0 && st.top() \leq arr[i]) {

 st.pop();

 if (st.size() > 0) {

 ans[i] = st.top();

 else {

 ans[i] = -1;

 st.push(arr[i]);

return ans;

}

Next greater element in left

Ex:

0 1 2 3 4 5 6 7
arr[8] = 11 10 6 5 10 7 4 8

ans[] =



vector<int> nearestElement(vector<int> arr) {

stack<int> st;

int N = arr.size();

vector<int> ans(N, -1);

for (int i = 0; i < N; i++) {

while (st.size() > 0 && arr[st.top()] <= arr[i]) {

st.pop();

if (st.size() > 0) {

ans[i] = st.top();

else {

ans[i] = -1;

st.push(i);

return ans;

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All Versions: