

Today's content

1. Peak Element

2. Find Unique element

Given a unsorted arr[] with all distinct elements return any one local maxima

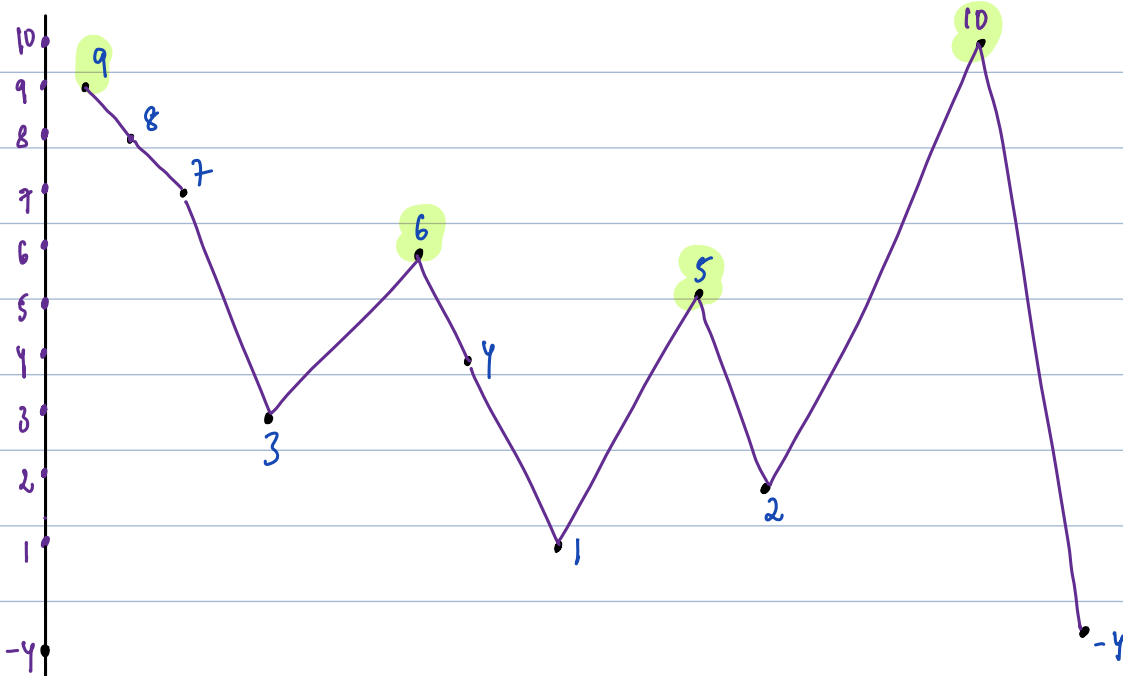
local maxima: An element is said to be local maxima, if $>$ than its adjacent elements [immediate left & right]

$arr[i]$ is local maxima if $= arr[i-1] < arr[i] > arr[i+1]$

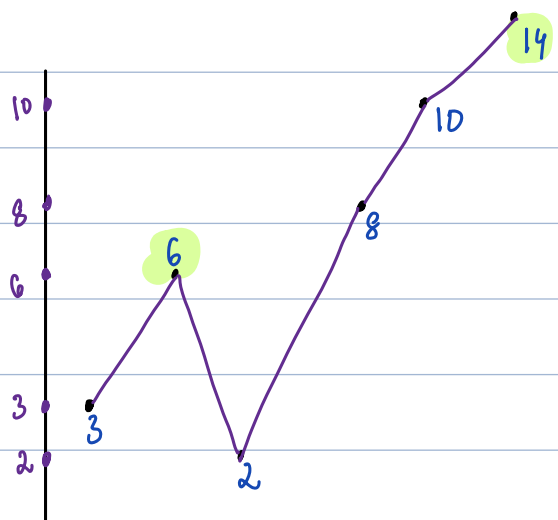
$arr[0]$ is local maxima if $= arr[0] > arr[1]$

$arr[n-1]$ is local maxima if $= arr[n-2] < arr[n-1]$

Ex: $arr[] = \{9, 8, 7, 3, 6, 4, 1, 5, 2, 10, -4\}$ local maxima: 9 6 5 10
Note: Return any one.



Ex: $arr[] = \{3, 6, 2, 8, 10, 14\}$ local maxima = 6 14



Note: Return any of them.

obs 1: Atleast 1 local maxima exists, because max ele will be $>$ its adjacent elements.

Idea: Iterate & return max of $arr[]$
Tc: $O(N)$ Sc: $O(1)$

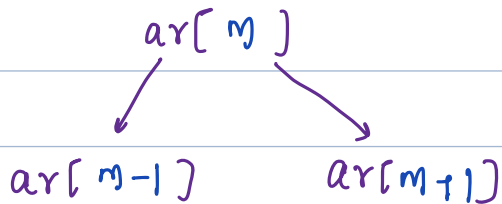
Idea: Iterate on $arr[]$; Tc: $O(N)$ Sc: $O(1)$

for every $arr[i]$, check if it's local maxima by comparing with adjacent elements, if it is return $arr[i]$;

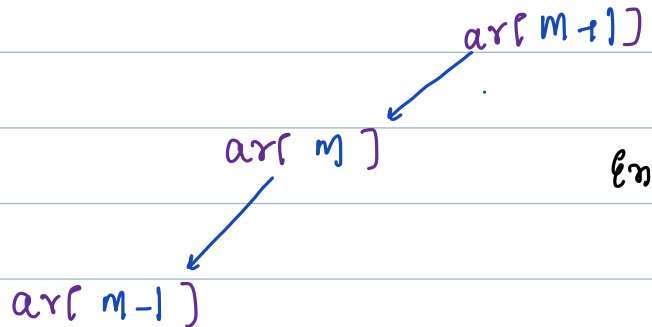
Idea:

Target = Any no local maxima Search Space = $\ln \text{arr}[]$

Case-1 if ($\text{arr}[m] > \text{arr}[m-1]$ && $\text{arr}[m] > \text{arr}[m+1]$) { return $\text{arr}[m]$ }

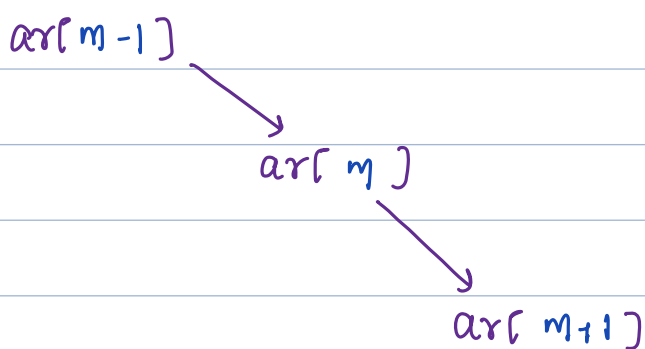


Case-2 if ($\text{arr}[m] < \text{arr}[m+1]$) { #go to right, because it's a guarantee we can find at least 1 local maxima on right }

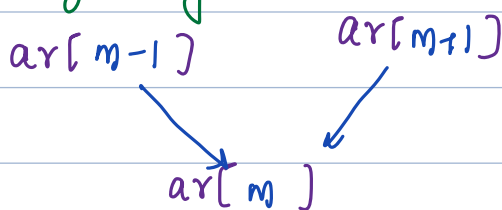


	$m-1$ m $m+1$		
l	0	1	2
h	6	4	12
m	8	10	14

Case-3 #go to left:
guarantee local maxima exists

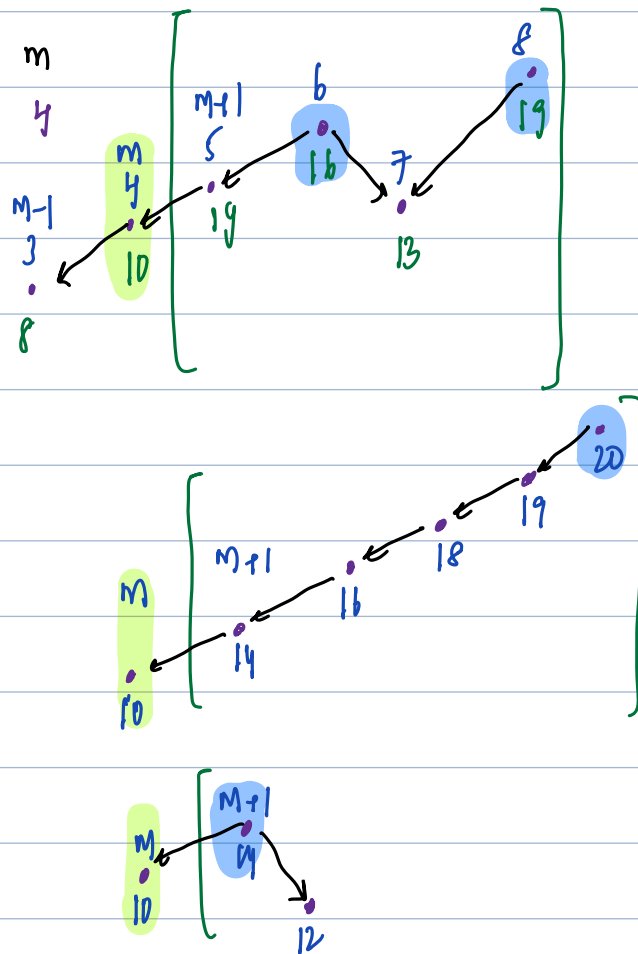


Case-IV: #go to any side.



#Con: After landing at mid:

Go to side where data increases, it's a guarantee local maxima exists on that side.



```
int localmaxima(vector<int> &arr) { TC:  $O(\log N)$  SC:  $O(1)$ 
```

```
int n = arr.size();
```

```
if (n == 1) { return arr[0]; }
```

```
if (arr[0] > arr[1]) { return arr[0]; }
```

```
if (arr[n-1] > arr[n-2]) { return arr[n-1]; }
```

```
int l = 0, h = n-1;
```

```
while (l < h) {
```

```
int m = (l+h)/2;
```

```
if (arr[m] > arr[m-1] && arr[m] > arr[m+1]) { return arr[m]; }
```

```
else if (arr[m] < arr[m+1]) {
```

```
l = m+1; // goto right;
```

```
else {
```

```
h = m-1; // goto left;
```

```
}
```

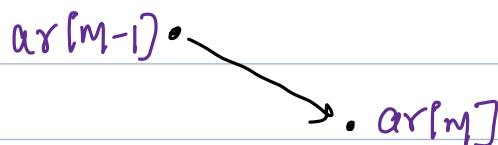
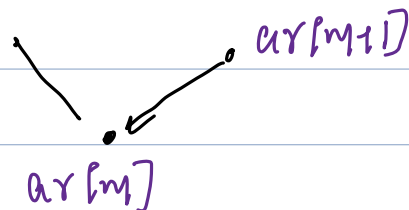
```
}
```

```
}
```

Issues:

$m=0$: $arr[0] > arr[-1]$

$m=n-1$: $arr[n-1] > arr[n-2]$ && $arr[n-1] > arr[n]$



Note: When comparing adjacent elements, be careful about corner elements

Given a unsorted arr[] return any one local maxima

local maxima: An element is said to be local maxima, if \geq than its adjacent elements [immediate left & right]

$$\underline{arr[i] \geq arr[i-1]} \text{ \& \& } \underline{arr[i] \geq arr[i+1]}$$

If $i=0$: $arr[0]$ is local maxima $\Rightarrow arr[0] \geq arr[1]$

If $i=N-1$: $arr[N-1]$ is local maxima $\Rightarrow arr[N-1] \geq arr[N-2]$

```
int localMaximaRepetition (vector<int> &arr) {
```

```
}
```

18. Every element occurs twice, except for 1, find a unique element

Note: Duplicates are adjacent to each other

Ex1: 0 1 2 3 4 5 6 7 8 9 10

arr[] = { 6 6 2 2 7 9 9 4 4 10 10 } ans = 7

arr[] = { 3 3 1 1 8 8 10 10 19 6 6 2 2 4 4 }

ans = 19

Idea1: Calculate num of all arr[] elements & return unique element

TC: $O(N)$ SC: $O(1)$

Idea2:

Target: Unique element Searchspace: In arr[]

arr[] = { 0 1 2 3 4 5 6 7 8 9 10 11 11
 { 12 12 6 6 21 21 7 9 9 4 4 10 10 }

if m is here: go right

On left:

1st occurenc of ele: even

left of unique go to right

if m is here: go left

On right

1st occurenc of ele: odd

right of unique: go to left

Note: It's possible that mid can have 2nd occurrence

if mid has 2nd occurrence bring to 1st occurrence

if (arr[mid] == arr[mid-1]) {

mid = mid - 1; # bring to 1st occurrence

}

$ar[] =$

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
3	3	1	1	8	8	10	10	19	6	6	2	2	4	4

 m $m+1$ l h m m

l h m $ar[m-1] == ar[m]$ left or Right

0 14 7 2^{nd} occ $m = m-1$; $m \% 2 == 0$: left: goto right, $l = m+2$;
 8 14 11 1^{st} occ ; $m \% 2 == 0$: right: goto left, $h = m-1$;
 8 10 9 1^{st} occ ; $m \% 2 == 0$: right: goto left, $h = m-1$;
 8 8 8 $ar[8] != ar[7]$ && $ar[8] != ar[9]$ return $ar[8]$

int unique(vector<int> &ar) { T.C: $O(\log_2 N)$ S.C: $O(1)$

int $n = ar.size()$

int $l = 0$, $h = n-1$

while ($l < h$) {

int $m = (l+h)/2$;

if ($m == 0$ || $m == n-1$) { return $ar[m]$ } problem

if ($ar[m] != ar[m-1]$ && $ar[m] != ar[m+1]$) { return $ar[m]$ }

if ($ar[m-1] == ar[m]$) { #update m to 1^{st} occurrence.

} $m = m-1$;

if ($m \% 2 == 0$) { # 1^{st} occurrence even: left go right

} $l = m+2$

else { # 1^{st} occurrence odd right: go left

} $h = m-1$

}

Note: If m came to start or last index, in that case, that itself is ans, in this ex above

