Todays Content

1. PfSum Intro
2. Range sum queries

Thursday Class

1. leaders
2. Count Triplets
3. 0/1 prefixSum

Friday class

1. Quick C++ Transition
2. Vector in C++
3. Pass by value & Pass by reference
4. Vector< Vector<int> > or Array of vector.
5. String in C++

Given ar[] & s_0, e_0 calculate sum of all elements from $s \to e$:

Eg:

	0	1	2	3	4	5	6	7	8	9
ar[10] = {	3	4	6	8	9	10	2	7	4	10 }

s=2 e=7 : ans = 42

```
int sum(int ar[], int s, int e){ TC: O(N) SC: O(1)
    int ans=0;
    for(int i=s; i<=e; i++){
        ans = ans + ar[i];
    }
    return ans;
}
```

Q1: Given ar[N] elements & Qmat[Q][2]

In Qmat matrix, we have Q: rows & 2: columns

Each row in Qmat represents a query.

0^{th} col in row represents : start point of query → s = Qmat[i][0]

1^{st} col in row represents : end point of query → e = Qmat[i][1]

For every query calculate sum of elements from index s..e in ar[] & print

Constraints:

$1 <= N <= 10^5$

$1 <= ar[i] <= 10^9$

$1 <= Q <= 10^5$

$0 <= s <= e < N$.

Sum = {1..10^{14}} >> int range

Min: 1 Max = $10^9 * 10^5 = 10^{14}$

ar[1] : {1} ar[10^5] : { 10^9 10^9 10^9.. 10^9 }

Ex:

	0	1	2	3	4	5	6	7	8	9
ar[10] = {	-3	6	2	4	5	2	8	-9	3	1 }

Qmat[6][2]

	0:s	1:e	Output
→ 0	4	8	9
→ 1	3	7	10
→ 2	1	3	12
→ 3	7	7	-9
→ 4	3	6	19
→ 5	0	4	14

Ideal:

For every Query:

Iterate from s..e calculate sum & print it

Expected TC: O(Q*N)

	s	e
0^{th}:	Qmat[0][0]	Qmat[0][1]
1^{st} :	Qmat[1][0]:3	Qmat[1][1]:7
2^{nd} :	Qmat[2][0]:1	Qmat[2][1]:3
i^{th} :	Qmat[i][0]:s	Qmat[i][1]:e

```
void RangeSum(int ar[], int N, int Q[][], int Q){

    for(int i=0; i<Q; i++){
        int s=Q[i][0], e=Q[i][1];
        long sum=0; // sum datatype should be long.
        for(int j=s; j<=e; j++){
            sum=sum+ar[j];
        }
        print(sum);
    }

}
```

Calculated TC: $O(Q*N)$ SC: $O(1)$

$1 \le N \le 10^5$

$1 \le Q \le 10^5$ $10^5 * 10^5 = 10^{10} >> 10^8$ TLE.

Optimization Idea:

Say we are given csk cricket scores for first 10 overs of batting. After every over total score is given.

Overs:	1	2	3	4	5	6	7	8	9	10
Total Score:	2	8	14	29	31	49	65	79	88	97

Q1: Total runs scored in 10^{th} over = score[10] - score[9] = 9

Q2: Total runs scored in 7^{th} over = score[7] - score[6] = 16

Q3: Total runs scored in 6^{th} - 10^{th} over = score[10] - score[5] = 66

Q4: Total runs scored in 3^{th} - 6^{th} over = score[6] - score[2] = 41

Q5: Total runs scored in 4^{th} - 9^{th} over = score[9] - score[3] = 74

Gen: Total runs scored in i^{th} - j^{th} over = score[j] - score[i-1]

Obs: Total sum till that point = Cummulative sum.

Cummulative sum calculate from start = prefix Sum

prefix Sum = Total sum from 0 till that index i.

String prefixSum value in array is pSum[i]

Idea: Create pf[] to optimize

1. pf[N]
2. pf[i] = Sum of all elements [0..i].

Ex:

	0	1	2	3	4	5	6	7	8	9
ar[10] = {	-3	6	2	4	5	2	8	-9	3	1 }
psum[10] :	-3	3	5	9	14	16	24	15	18	19

Output: Ans using Psum[]

	0:s	1:e		psum[e] - psum[s-1]
0	4	8	→	psum[8] - psum[3] = 18 - 9 = 9
1	3	7	→	psum[7] - psum[2] = 15 - 5 = 10
2	1	3	→	psum[3] - psum[0] = 9 - (-3) = 12
3	7	7	→	psum[7] - psum[6] = 15 - 24 = -9
4	3	6	→	psum[6] - psum[2] = 24 - 5 = 19
5	0	4	→	psum[4] - psum[-1] // Error

→ Sum[0..4] = psum[4].

Query:

```
[s..e] = if (s==0){ // [0..e]
           print( psum[e])
         else{
           print (psum[e] - psum[s-1])
```

Construct Psum[] For Given ar[N]

	0	1	2	3	4
ar[5] = {	3	-2	4	5	6 }
sum = 0 →	3	1	5	10	16
psum[5] = {	3	1	5	10	16 }

Steps: Carry forward sum L → R
1. update sum
2. storing in psum[]

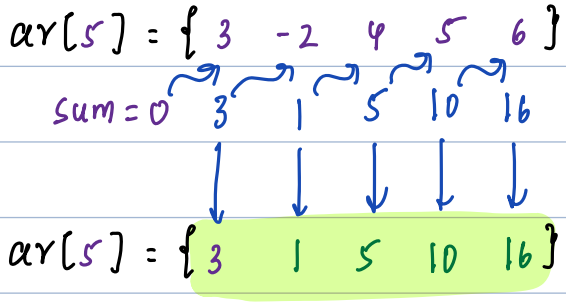
↳ pfSum array.

TC: O(N+Q) SC: O(N+5) ≈ O(N)

```
void RangeSum(int ar[], int N, int Qmat[][], int Q){
    long psum[N]; // long = Storing sum values >> int range.
    long sum=0;
    for(int i=0;i<N;i++){ ——→ TC: N
        sum=sum+ar[i]; //update
        psum[i]=sum; //Storing
    }

    for(int i=0; i<Q; i++){ →TC: Q
        int s=Qmat[i][0], e=Qmat[i][1];
        if(s==0){ //sum:[0..e]
            print(psum[e]);
        }else{
            print(psum[e]-psum[s-1]);
        }
    }
}
```

Note: Modifying same arr[] into Prefix arr[]



Issues in storing psum[] → arr[]

1. loose input arr[] information.
2. pfsum[] datatype & arr[] datatype might not match.

When can we store psum[] → arr[]

1. If arr[] is no longer needed.
2. if pfsum[] datatype & arr[] datatype is same.

Note: When multiples range queries, think in pfsum[]