

Today's Content

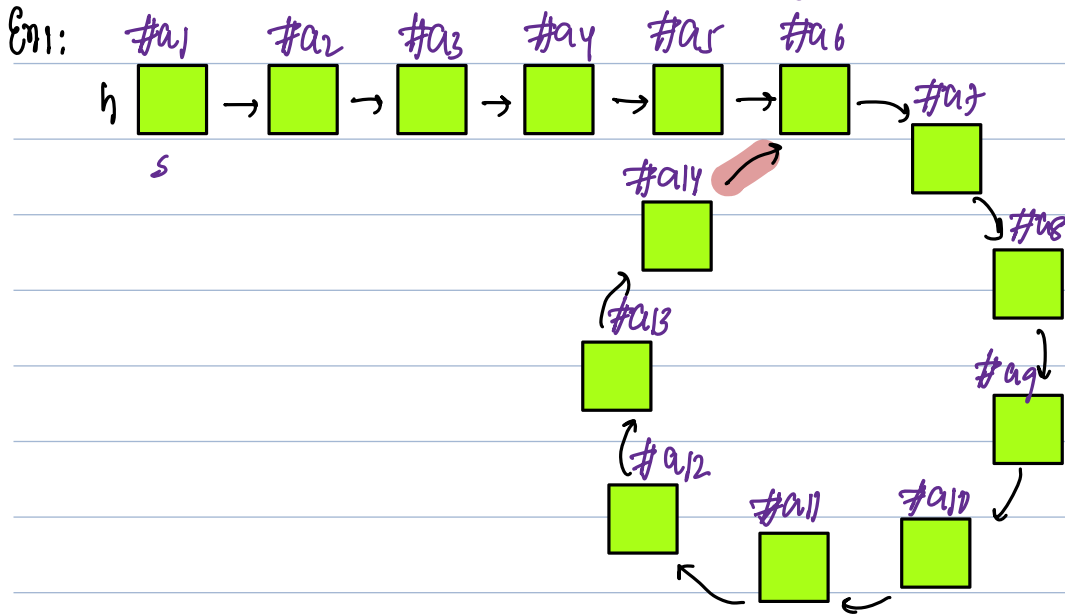
1. Cycle Detection
2. If cycle exists remove cycle
3. Intersection of 2 Linked Lists
4. Proof for cycle detection

Q1 Given a head node of Linked List, check for cycle detection?

#Note1: When last node points to any prev node, that's when cycle formed

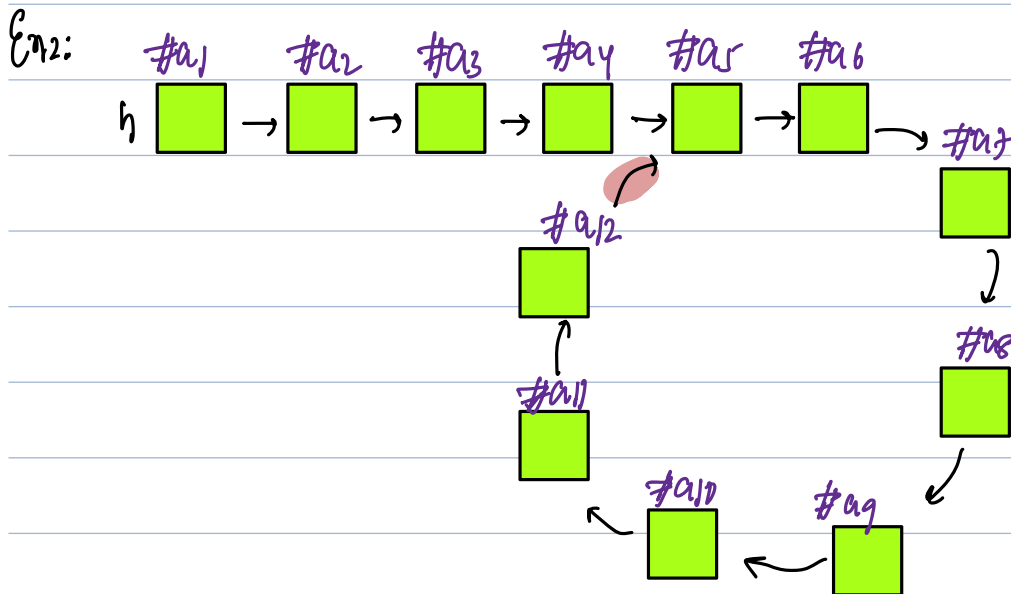
#Note2: if no cycle return null.

#Note3: If cycle is there, remove cycle {last node \rightarrow null ptr}
Return start / first node of cycle.



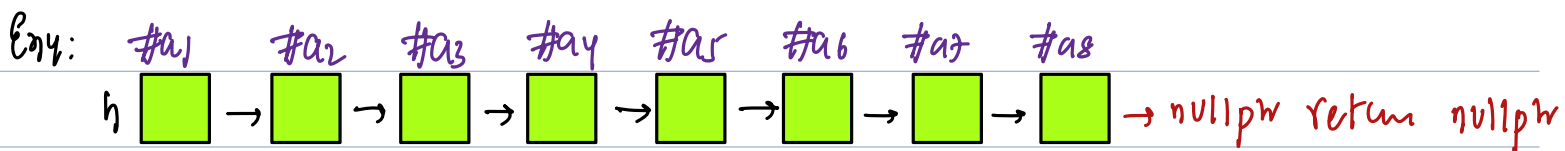
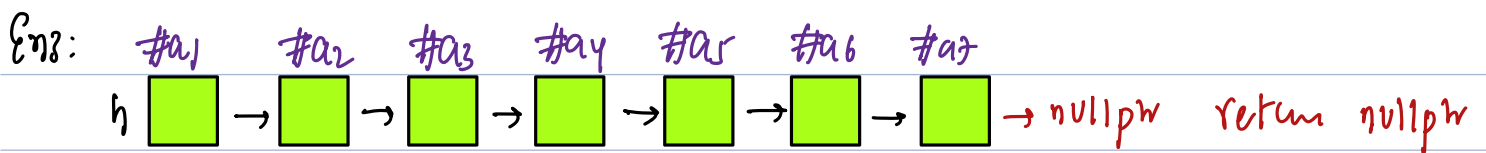
Cycle exist

1. break cycle
2. return a6



Cycle exist

1. break cycle
2. return a5

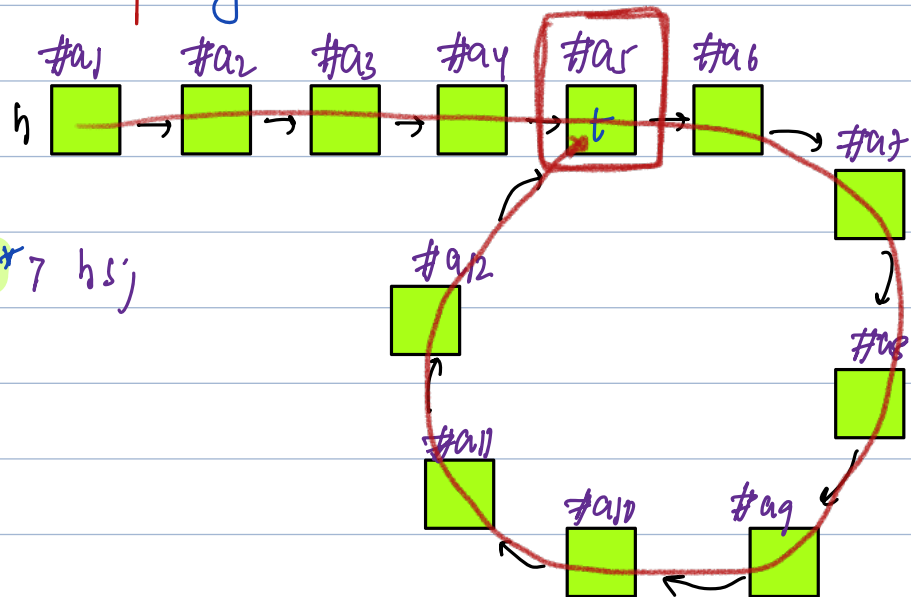


Idea 1: Create a hashset $\langle \text{Node}^* \rangle$ TC: $O(N)$ SC: $O(N)$ \rightarrow hashset

Iterate on LinkedList & Insert each node address.

obs 1 # of a node address already exists: Cycle exists

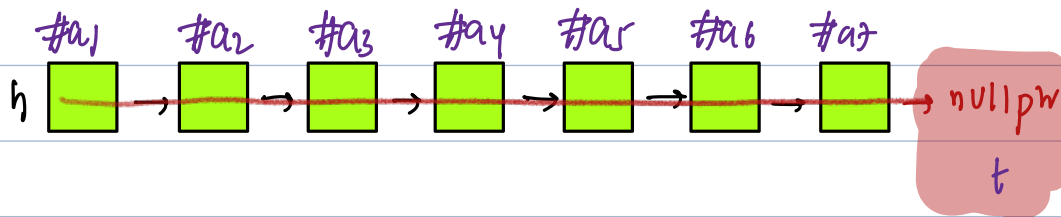
obs 2 # of $t = \text{null}$ cycle doesn't exist



unordered_set $\langle \text{Node}^* \rangle$ & hsj

#a1	#a2	#a3
#a4	#a5	#a6
#a7	#a8	#a9
#a10	#a11	#a12

Ex 2:

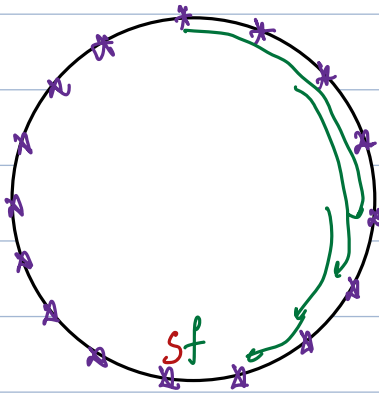


#a1	#a2	#a3
#a4	#a5	#a6
#a7		

Node* yu(Node* h) {

}

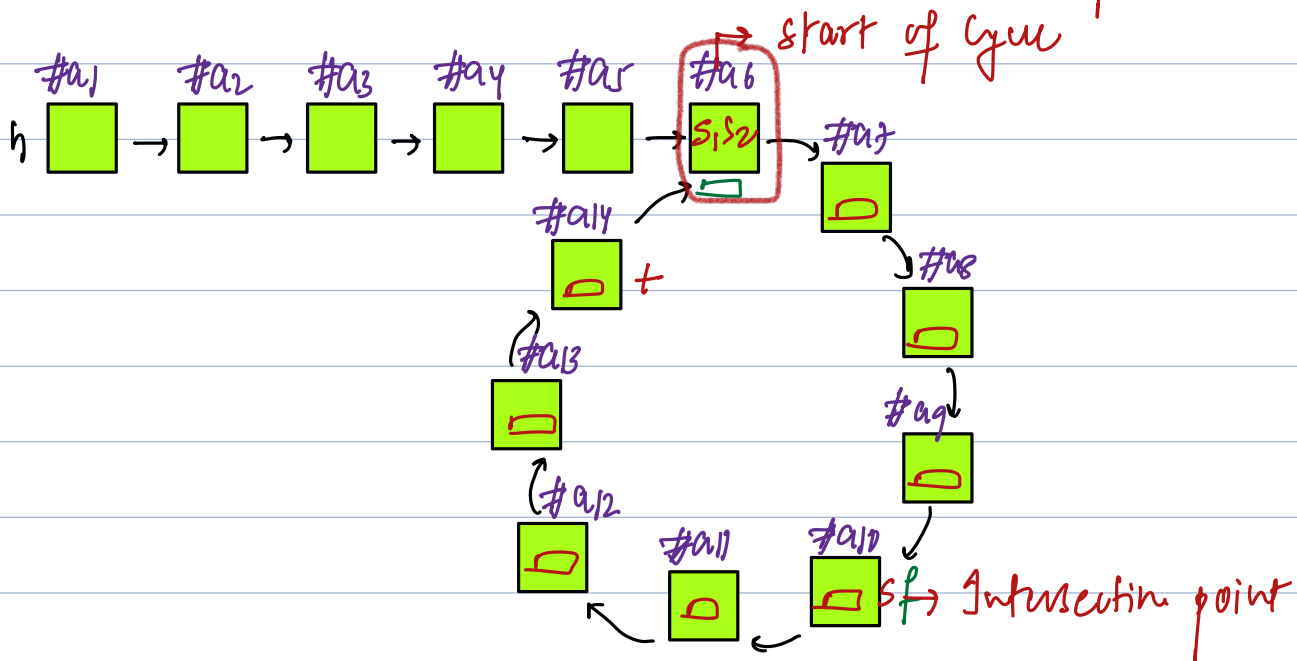
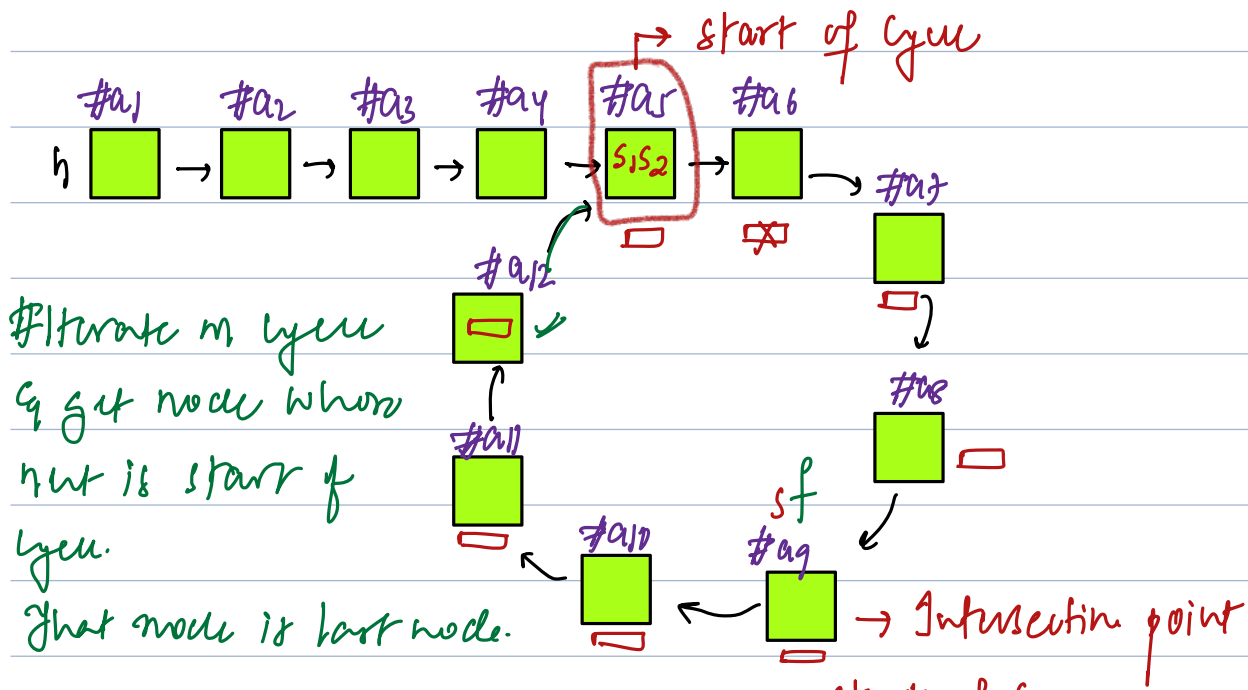
Idea 2:



Distance between them

Initial: $4 \rightarrow 3 \rightarrow 2 \rightarrow 1 \rightarrow 0$

In a circle, if $\text{slow} : 1 \text{ step}$ $\text{fast} : 2 \text{ steps}$, irrespective of initial distance, they will meet, because distance between will decrease by 1 step after each iteration in circle {d < 0, meeting each other}



Node* detectCycle(Node* h) { Tc: $O(N)$ Sc: $O(1)$

Node* s = h, *f = h;

bool cycleExist = false;

while(f != nullptr && f->next != nullptr) {

s = s->next;

f = f->next->next;

if(s == f) {

cycleExist = true;

break;

}

if(cycleExist == false) {

return nullptr;

#cycle detection

Node* s1 = h, *s2 = s;

while(s1 != s2) {

s1 = s1->next;

s2 = s2->next;

#start of cycle

Node* start = s1;

Node* temp = start;

while(temp->next != start) {

temp = temp->next;

temp->next = nullptr;

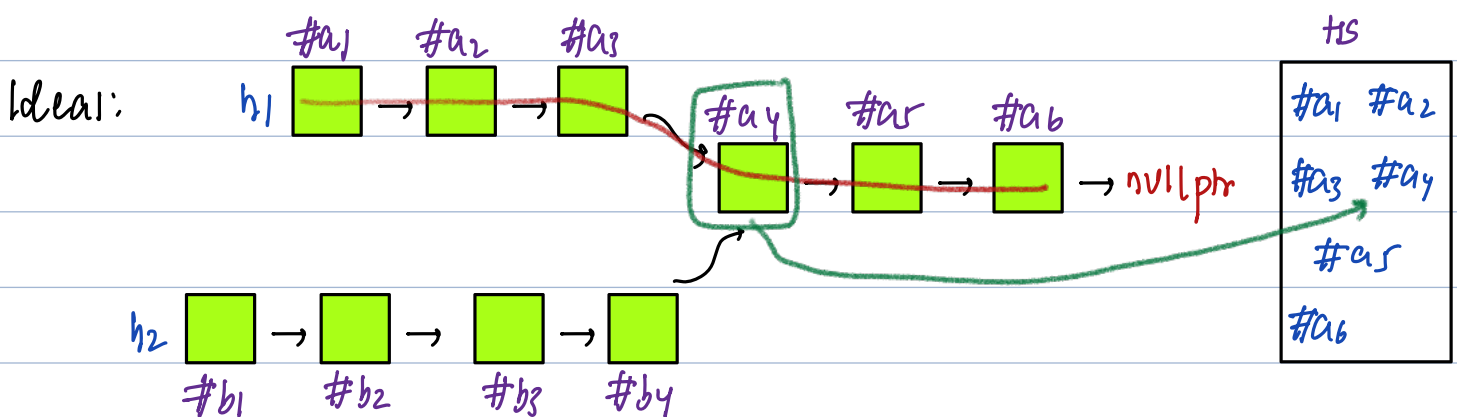
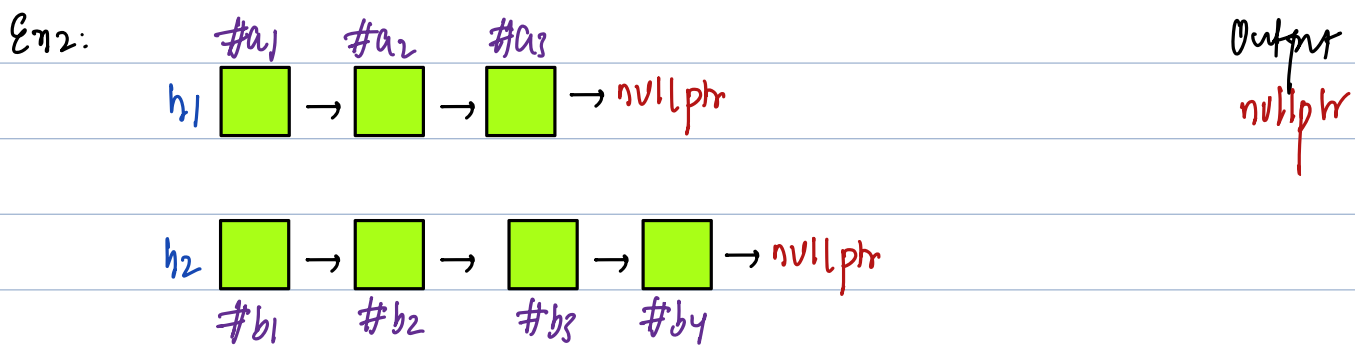
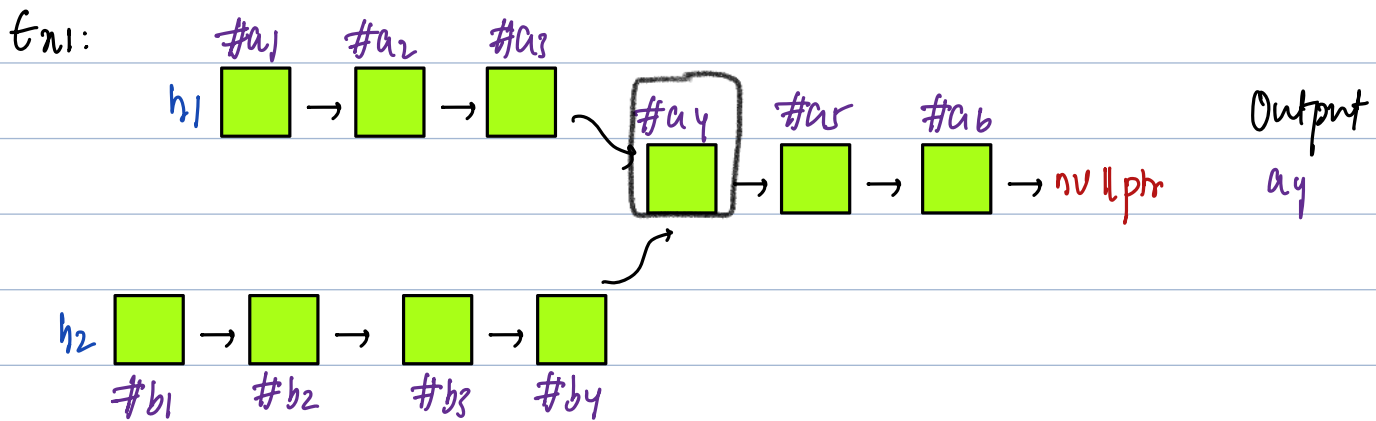
#breaking cycle

return start;

28 Given the heads of two singly linked-lists

Return nodes at which the two lists intersect.

If the two linked list have no intersection at all, return null.



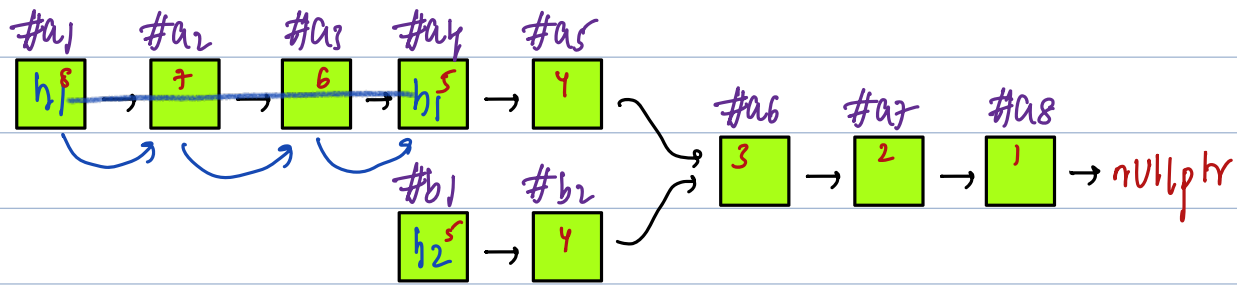
1. Iterate on h1 & Insert all addresses in hashset & store b1

2. Iterate on h2:

if node exists in bs: Return node

if not, goto next node.

#En2:



#Idea: TC: $O(N+M)$ SC: $O(1)$

Step1:

$l_1 = \text{size}(h_1), l_2 = \text{size}(h_2);$

if $(l_1 > l_2)$ { # Extra nodes $l_1 - l_2$
update h_1 $(l_1 - l_2)$ times
}

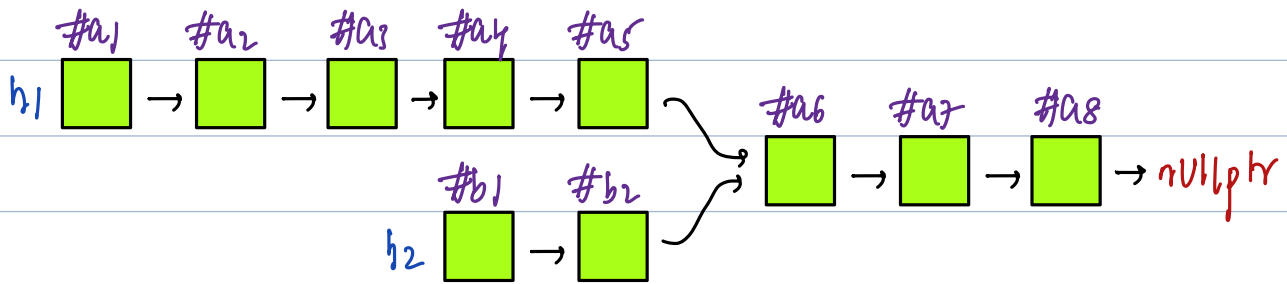
else if $(l_2 > l_1)$ { # Extra nodes $l_2 - l_1$
update h_2 $(l_2 - l_1)$ times
}

Step2:

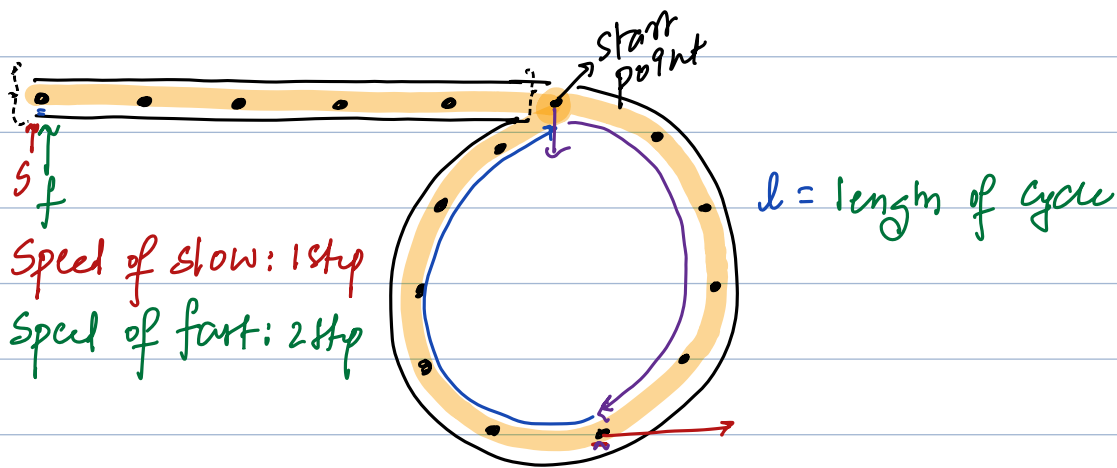
while $(h_1 \neq \text{nullptr})$ {
if $(h_1 == h_2)$ { return h_1 }
 $h_1 = h_1 \rightarrow \text{next}$
 $h_2 = h_2 \rightarrow \text{next}$
}

return nullptr

#Ideas:



Cycle Detection: Proof.



Note: Btwn s & f move same amount of time.

$ds: a + kl + d$ // k : No. of times s pointer iterating in cycle

$df: a + ml + d$ // M : No. of times f pointer iterating in cycle

$$df = 2 * ds;$$

$$a + ml + d = 2[a + kl + d]$$

$$\cancel{a} + ml + \cancel{d} = \cancel{2a} + 2kl + \cancel{2d}$$

$$ml = a + 2kl + d$$

$$ml - 2kl - d = a$$

$$a = l[M - 2k] - d;$$

$$a = ly - d;$$