

Today's Content

1. Rotate arr[] by k times
2. Vector pass by value & pass by reference.
3. Count no. of distinct elements

Q: Given an arr[], how to reverse entire array from l..r

Ex: arr[12]

$l=2 \ r=9 \ \{ 10 \ 20 \ 30 \ 40 \ 50 \ 60 \ 70 \ 80 \ 90 \ 100 \ 110 \ 120 \}$

100 90 80 70 60 50 40 30

$P_1 \rightarrow P_1 \rightarrow P_1 \rightarrow P_1 \quad P_2 \leftarrow P_2 \leftarrow P_2 \leftarrow P_2$

Idea:

$P_1 \leftarrow P_2 \ \{ 10 \ 20 \ 100 \ 90 \ 80 \ 70 \ 60 \ 50 \ 40 \ 30 \ 110 \ 120 \}$

2 \leftarrow 9 swap(arr[P₁], arr[P₂]); P_1++ ; P_2--

3 \leftarrow 8 swap(arr[P₁], arr[P₂]); P_1++ ; P_2--

4 \leftarrow 7 swap(arr[P₁], arr[P₂]); P_1++ ; P_2--

5 \leftarrow 6 swap(arr[P₁], arr[P₂]); P_1++ ; P_2--

6 \leftarrow 5 $P_1 > P_2$: Stop arr[P₁..P₂] is reversed.

void reverseRange(int arr[], int l, int r){

int P₁=l, P₂=r;

while(P₁ < P₂) {

//swap arr[P₁] & arr[P₂]

int tmp = arr[P₁];

arr[P₁] = arr[P₂];

arr[P₂] = tmp;

P_1++ ; P_2--

}

3

Iteration: Worst Case: Entire Array = $N/2$

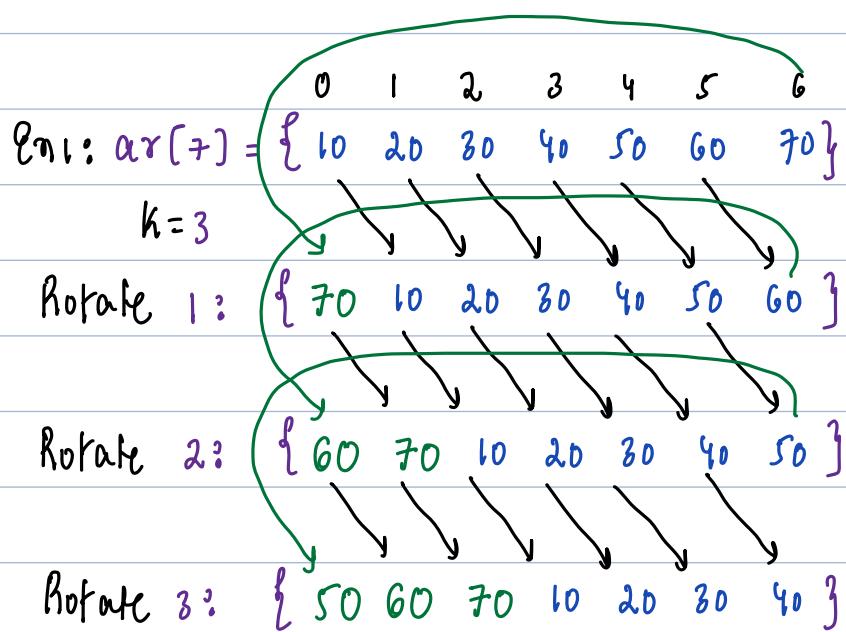
TC: $O(N)$ SC: $O(1)$

Q: Given an arr[] rotate it by k times by Right \rightarrow Left

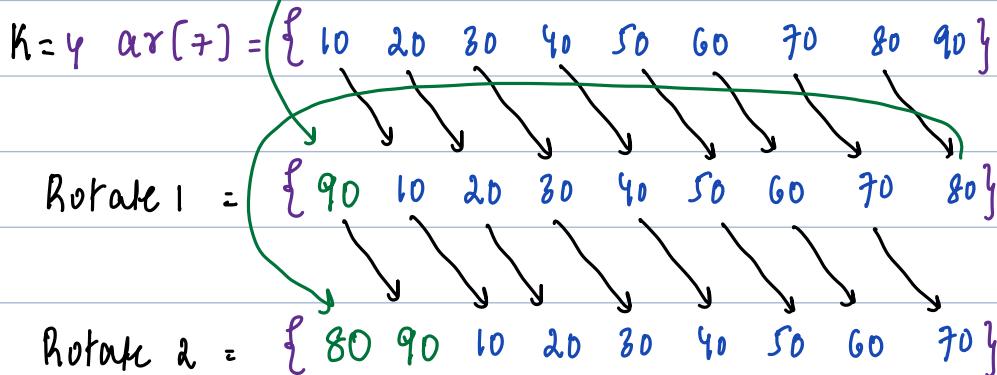
Constraints:

$$1 \leq N \leq 10^5$$

$$1 \leq K \leq 10^5$$



Ex2:



Ideas: For k times:

Estimated TC: $O(k \cdot n)$

Rotate arr[] right to left.

Single Rotate:

$arr[7] = \{ 10, 20, 30, 40, 50, 60, 70 \}$

$arr[7] = \{ 70, 10, 20, 30, 40, 50, 60 \}$

$int tmp = arr[6]; // 70 \leftarrow$

Shifting ↘

i i-1

$arr[6] = arr[5]$

$arr[5] = arr[4]$

$arr[4] = arr[3]$

$arr[3] = arr[2]$

$arr[2] = arr[1]$

$arr[1] = arr[0]$

$arr[0] = tmp;$

void rotateth(int arr[], int N, int k) {

for (int l=0; l < k; l++) {

int tmp = arr[N-1]; // tmp : last element

for (int i=N-1; i>=1; i--) {

arr[i] = arr[i-1];

arr[0] = tmp;

Innwlloop = NK Outterloop = K

TC: $O(N \cdot k)$ SC: $O(1)$

$$1 \times = N \times = 10^5$$

$$1 \times = k \times = 10^5$$

$$O(10^5 \times 10^5) = 10^{10} \gg 10^8 \text{ TLE}$$

Optimized Approach

0 1 2 3 4 5 6

$$\text{Ex1: } k=3 \text{ arr[7]} = \{10 20 30 40 \text{ (blue box)} 50 60 70 \text{ (green box)}\}$$

$$\text{Rotate 1} = \{70 10 20 30 40 50 60\}$$

$$\text{Rotate 2} = \{60 70 10 20 30 40 50\}$$

$$\text{Rotate 3} = \{50 60 70 10 20 30 40\}$$

0 1 2 3 4 5 6 7 8

$$k=4 \text{ arr[7]} = \{10 20 30 40 50 \text{ (blue box)} 60 70 80 90 \text{ (green box)}\}$$

$$\text{Rotate 1} = \{90 10 20 30 40 50 60 70 80\}$$

$$\text{Rotate 2} = \{80 90 10 20 30 40 50 60 70\}$$

$$\text{Rotate 3} = \{70 80 90 10 20 30 40 50 60\}$$

$$\text{Rotate 4} = \{60 70 80 90 10 20 30 40 50\}$$

Obs: Rotate arr[N] by k times

1. last k ele will shift to first

2. Remaining ele will shift to back

Ex4:

0 1 2 3 4 5 6 7 8 9 10 11 12

$$k=5 \text{ arr[13]} = \{a_0 a_1 a_2 a_3 a_4 a_5 a_6 a_7 a_8 a_9 a_{10} a_{11} a_{12}\}$$

Step1:

$$\text{Reverse entire arr[7]} \{a_{12} a_{11} a_{10} a_9 a_8 \text{ (green box)} a_7 a_6 a_5 a_4 a_3 a_2 a_1 a_0 \text{ (blue box)}\}$$

Step2:

Reverse k elements

Step3:

Reverse rem ele

0 1 2 3 4 5 6 7 8 9 10 11 12

$$\text{After Rotation} = \{a_8 a_9 a_{10} a_{11} a_{12}\} \quad | \quad \{a_0 a_1 a_2 a_3 a_4 a_5 a_6 a_7\}$$

Idea: Given arr[N] q rotate k times

1. Reverse entire arr[] : reverseRange(arr, 0, N-1) : N/2

2. Reverse first k ele : reverseRange(arr, 0, k-1) : k/2

3. Reverse remaining ele : reverseRange(arr, k, N-1) : (N-k)/2

} Iterations : N

} TC : O(N)

```

void reverseRange(int arr[], int l, int r) {
    int p1 = l, p2 = r;
    while (p1 < p2) {
        int tmp = arr[p1];
        arr[p1] = arr[p2];
        arr[p2] = tmp;
        p1++;
        p2--;
    }
}

```

Void rotatkh (int arr[], int N, int k) {

$k = k \% N$; // Edge Cases

reverseRange (arr, 0, N-1) → reverseRange (arr, 0, 3)

reverseRange (arr, 0, k-1) → reverseRange (arr, 0, 5) ? Errr. *

reverseRange (arr, k, N-1)

0 1 2 3

Ex: $N=4$ $arr[4]=\{10 20 30 40\}$ $k=6$?

$p_1=0, p_2=5$

swap $arr[0]$ w/ $arr[5]$ // RTE.

Observability

0 1 2 3

Ex: $arr[4] = 10 20 30 40$

Rufahim observing

k_i :

$0 \rightarrow 4 \rightarrow 8 \rightarrow 12 \rightarrow 16 \rightarrow 20 \dots$

$1 \rightarrow 5 \rightarrow 9 \rightarrow 13 \rightarrow 17 \rightarrow 21 \dots$

$2 \rightarrow 6 \rightarrow 10 \rightarrow 14 \rightarrow 18 \rightarrow 22 \rightarrow 26$

$3 \rightarrow 7 \rightarrow 11 \rightarrow 15 \rightarrow 19 \rightarrow 23 \rightarrow 27 \rightarrow 31 \rightarrow 35$

rotate 0 = $\underline{\underline{10 20 30 40}}$

rotate 1 = $\underline{\underline{40 10 20 30}}$

rotate 2 = $\underline{\underline{30 40 10 20}}$

rotate 3 = $\underline{\underline{20 30 40 10}}$

rotate 4 = $\underline{\underline{10 20 30 40}}$

rotate 5 = $\underline{\underline{40 10 20 30}}$

rotate 6 = $\underline{\underline{30 40 10 20}}$

rotate 7 = $\underline{\underline{20 30 40 10}}$

rotate 8 = $\underline{\underline{10 20 30 40}}$

K $N=4 \neq k \% N \Rightarrow$ In general $k \% N$

20	0
17	1
26	2
30	2
35	3

Con: Rotating k is same as Rotating $k \% N$

```
void rotate(int arr[], int N, int k){
```

0 1 2 3

$k = k \% N$; // Edge Cases

Ex: $N=4$ $arr[4] = \{10, 20, 30, 40\}$ $k=6 \Rightarrow 6 \% 4 = 2$

reverseRange(arr, 0, N-1) → reverseRange(arr, 0, 3) ✓

reverseRange(arr, 0, k-1) → reverseRange(arr, 0, 1) ✓

reverseRange(arr, k, N-1) → reverseRange(arr, 2, 3) ✓