

Today's Content

Strength: 22

1. Man histogram area
2. Man rectangular area

Histogram Area:

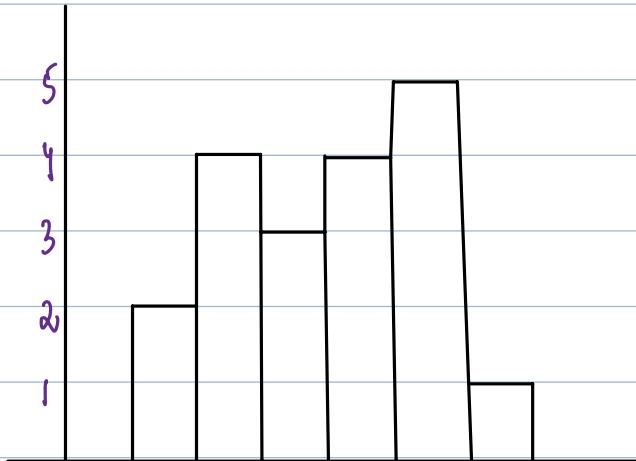
Given continuous blocks of histogram find max Rectangle Area.

Note1: Rectangle area, should be within histogram

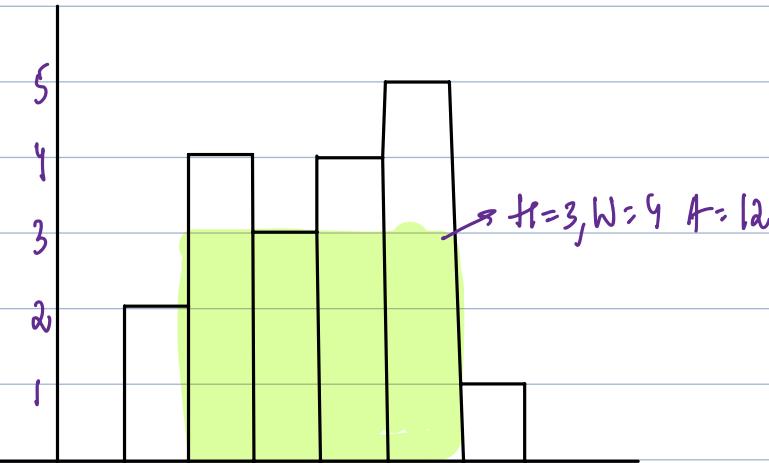
Note2: Width of each histogram is 1.

Ex1: 0 1 2 3 4 5

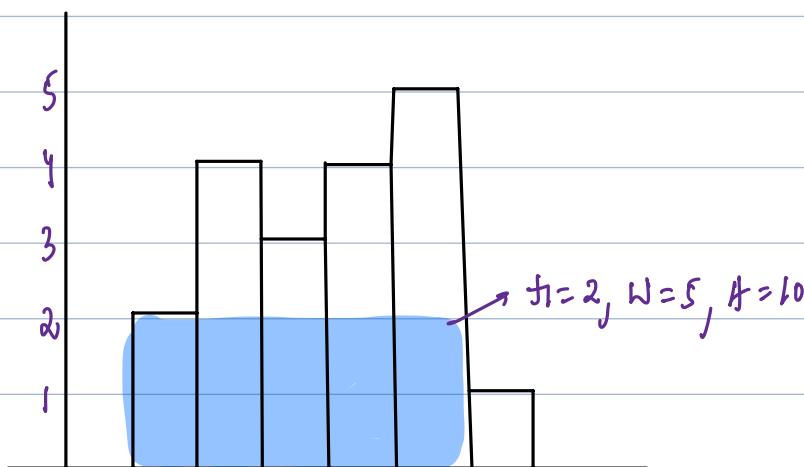
$$arr[6] = \{ 2, 4, 3, 4, 5, 1 \}$$



Rect1:

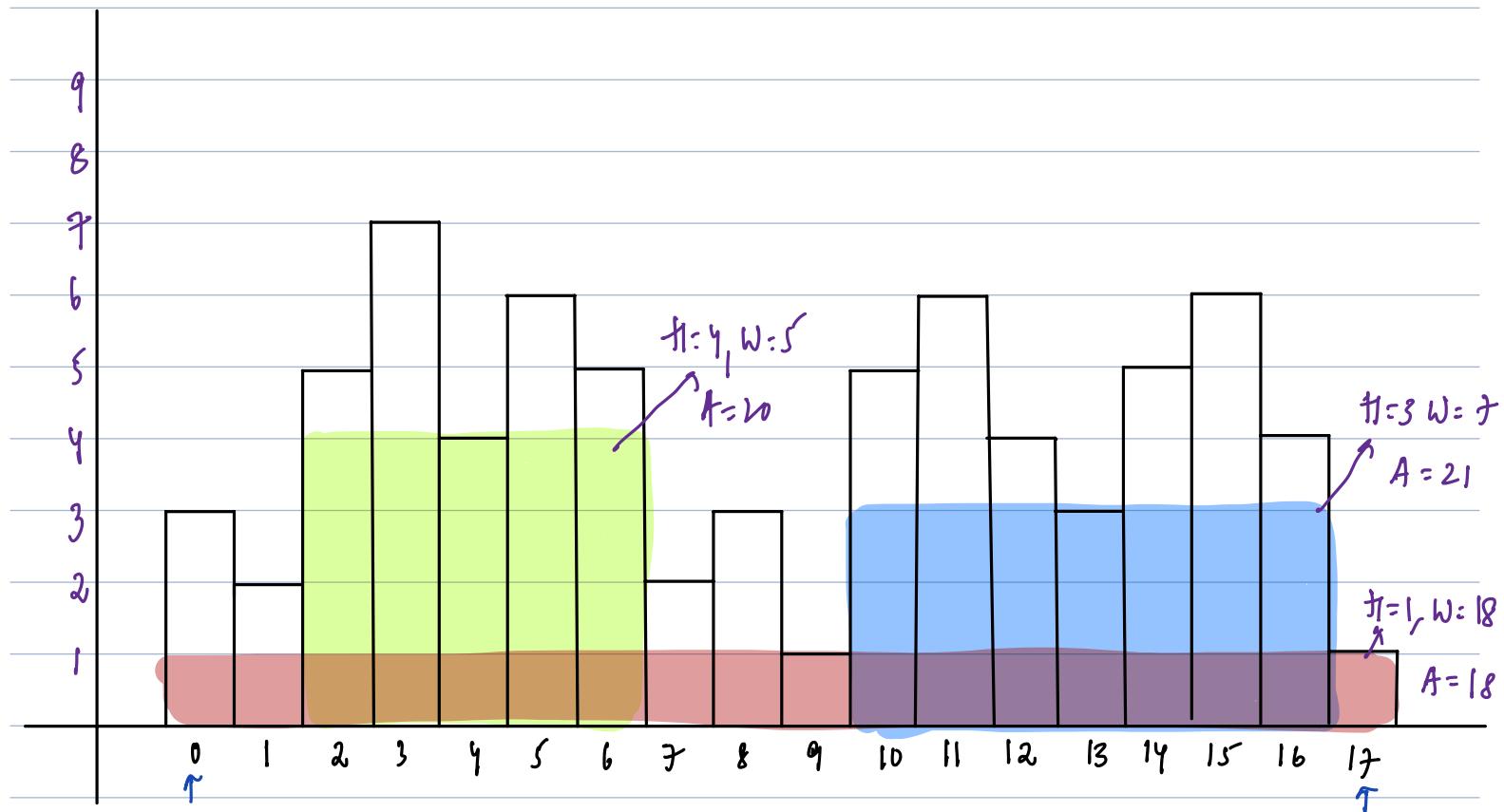


Rect2:



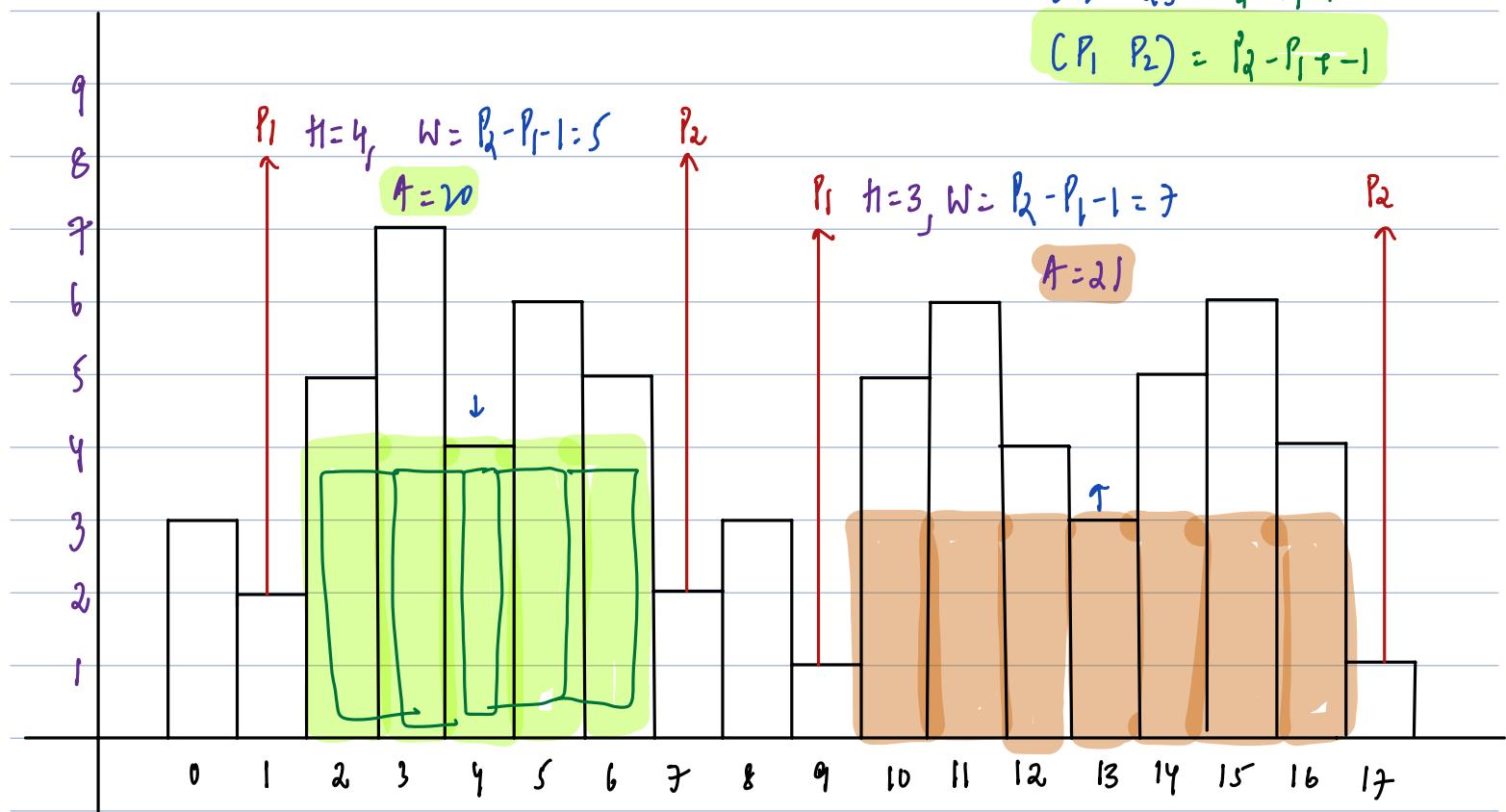
Ex2

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17
3	2	5	7	4	6	5	2	3	1	5	6	4	3	5	6	4	1



$$[P_1 \ P_2] = P_2 - P_1 + 1$$

$$[P_1 \ P_2] = P_2 - P_1 - 1$$



#Rectangle:

#Hint1: Rectangle height should match histogram height
Which?

#Hint2: Take every histogram height $h[i]$ as rectangle height:

Iterate q calculate P_2 # P_2 is smaller index on right for $h[i]$

Iterate q calculate P_1 # P_1 is smaller index on left for $h[i]$

$$\text{Area} = h[i] * (P_2 - P_1 - 1)$$

$$\text{ans} = \max(\text{ans}, \text{Area})$$

$$\text{TC: } O(N+N) = O(N^2) \text{ Sc: } O(1)$$

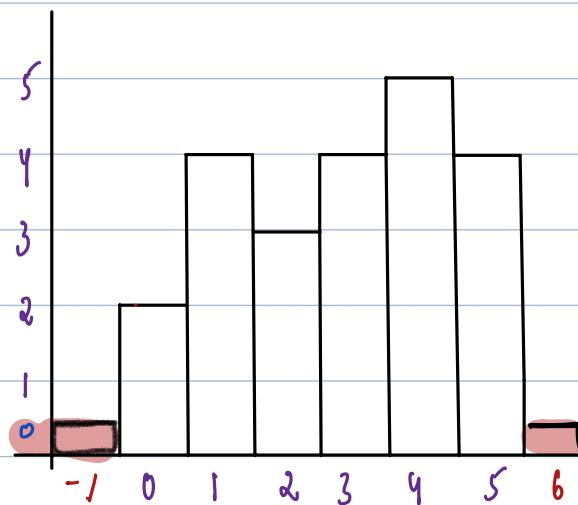
#Id 02: Optimization

In above for every $h[i]$ we need i^{th} smaller index on left & right.

Precompute i^{th} smaller index on left & right for all elements at once

Dy Run: 0 1 2 3 4 5

$ar[] = 2 4 3 4 5 1$



$sl[] P_1 : -1 0 0 2 3 2$ # Note: Default value for $P_1 = -1$

$sr[] P_2 : -1 2 6 6 5 6$ # Note: Default value for $P_2 = N$

Width $P_2 - P_1 - 1 : 6 1 5 3 1 3$

Area $h * W : 12 4 15 12 5 12$

int RectangularArea(vector<int> &arr) { TC: O(N+N+N) = O(N) SC: O(N+N) = O(N)

int N = arr.size();

10:27pm

for smaller index in left

stack<int> st1;

vector<int> p1[N, -1];

for (int i = 0; i < N; i++) {

while (st1.size() > 0 && st1.top() >= arr[i]) {

{ st1.pop(); }

if (st1.size() > 0) {

{ p1[i] = st1.top(); }

st1.push(arr[i]); }

}

for smaller index in right;

stack<int> st2;

vector<int> p2[N, N];

for (int i = N-1; i >= 0; i--) {

while (st2.size() > 0 && st2.top() >= arr[i]) {

{ st2.pop(); }

if (st2.size() > 0) {

{ p2[i] = st2.top(); }

st2.push(arr[i]); }

}

For each histogram calculate rectangular area

int ans = 0;

for (int i = 0; i < N; i++) {

{ ans = max(ans, arr[i] * (p2[i] - p1[i] - 1)); }

$\frac{\text{height}}{\text{width}}$

return ans;

3

8

Given a binary matrix $[N][M]$

Return max rectangular area containing only 1's.

Ex1:

$\text{mat}[4][5]$

	0	1	2	3	4
0	1	0	1	0	0
1	1	0	1	1	1
2	1	1	1	1	1
3	1	0	0	1	0

$\text{ans} = 6$

Ex2:

$\text{mat}[8][5]$

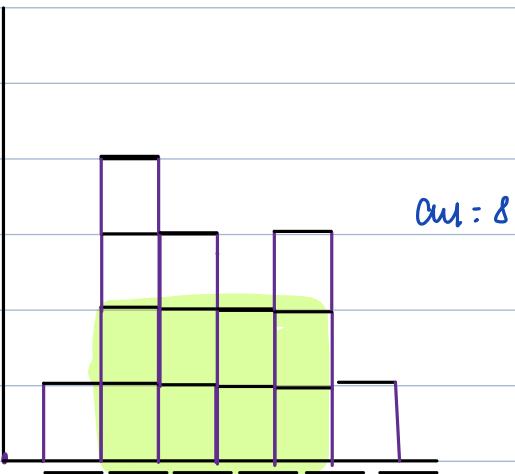
0	1	1	1	1	1	1
1	0	1	1	1	0	1
2	1	1	1	1	1	1
3	1	1	1	1	1	0
4	0	0	1	1	1	0
5	0	1	0	1	1	1
6	1	1	1	1	1	0
7	1	1	1	1	0	1

$\text{ans} = 12$

#Hint: Max rectangular area, whose base at last row.

Ex1: mat[4][6]:

	0	1	2	3	4	5
0	1	1	0	1	0	0
1	1	1	1	0	1	1
2	0	1	1	1	1	0
3	1	1	1	1	1	1



#Tracing Idea2:

Take every row base:

For each column in a row, calculate no: of consecutive 1s along Top.

For each row, apply previous rectangle area logic.

mat[8][5]

	0	1	2	3	4	5
0	1	1	1	1	1	1
1	0	1	1	1	0	1
2	1	1	1	1	1	1
3	1	1	1	1	1	0
4	0	0	1	1	1	0
5	0	1	0	1	1	1
6	1	1	1	1	1	0
7	1	1	1	1	0	1

for every Row: #N

for each col: #M

Iterate along col, to get consecutive 1s #N

Apply Above rectangular area #M

Total TC: $O(N * M * N + N * M) = O(N^2M)$

	0	1	2	3	4	5
0	0	1	0	6	9	1
1	1	2	1	7	5	0
2	3	2	8	0	1	

Rut = 8

Rut = 10

Rut = 8

Optimize Idea:

For each column: Iterate from top to down & calculate consecutive is.

mat[8][5]

	0	1	2	3	4	5
0	1↓	1↓	1	1	1	1
1	0↓	1↓	1	1	0	1
2	1↓	1↓	1	1	1	1
3	1↓	1↓	1	1	1	0
4	0↓	0↓	1	1	1	0
5	0↓	1↓	0	1	1	1
6	1↓	1↓	1	1	1	0
7	1↓	1↓	1	1	0	1

mat[8][5]

	0	1	2	3	4	5
0	1	1	1	1	1	1
1	0	2	2	2	0	2
2	1	3	3	3	1	3
3	2	4	4	4	2	0
4	0	0	5	5	3	0
5	0	1	0	6	4	1
6	1	2	1	7	5	0
7	2	3	2	8	0	1

Approach: $Tc: O(N \cdot M)$

For each col:

Iterate & store consecutive is information at each cell

For every row:

Apply rectangular area in histogram logic.

int maxRectangle(vector<vector<int>> &mat) { Tc: O(N*M) Sc: O(M) }

int N = mat.size(), M = mat[0].size();

for (int j = 0; j < M; j++) {

Iterate in j^th row & get consecutive 1's count q & store it.

int c = 0;

for (int i = 0; i < N; i++) {

mat[i]: i^th vector or i^th row

if (mat[i][j] == 1) {

 c++;

else {

 c = 0;

 mat[i][j] = c;

 3

 3

int A = 0;

for (int i = 0; i < N; i++) {

 A = max(A, RectangularArea(mat[i]));

return A;

3