

Today's Content

1. Count of subsets with sum = k
2. Calculate & return count of Quadruplets with sum = k

Given $arr[N]$ & k calculate &

Check if there exists a subset with $sum = k$

Constraints

$$1 \leq N \leq 40$$

$$\# 2^{10} = 1024 \approx 10^3$$

$$1 \leq arr[i] \leq 10^7$$

$$\# (2^{10})^4 = (10^3)^4 = 10^{12}$$

$$1 \leq k \leq 10^9$$

0 1 2 3 4 5 6

Ex: $arr[] = \{3, 7, 2, 9, 2, 5, 7\}$ $k = 10$

Idea 1: Generate all subsets:

For every subset, iterate & calculate $sum = k$

$$T.C: O(2^N * N) = O(N * 2^N) \quad S.C: O(1)$$

$$\# N \leq 40, N = 40 \quad 2^{40} * 40 \gg 10^8 \text{ TLE.}$$

Idea 2:

$arr[] = \{3, 7, 4, 9, 2, 5\}$ $k = 18$ 2^6 Subsets

$a[] = \{3, 7, 4\} : 2^3$

$\{ \}$

$\{3\}$

$\{7\}$

$\{4\}$

$\{3, 7\}$

$\{7, 4\}$

$\{3, 4\}$

$\{3, 7, 4\}$

$b[] = \{9, 2, 5\}$

$\{ \}$

$\{9\}$

$\{2\}$

$\{5\}$

$\{9, 2\}$

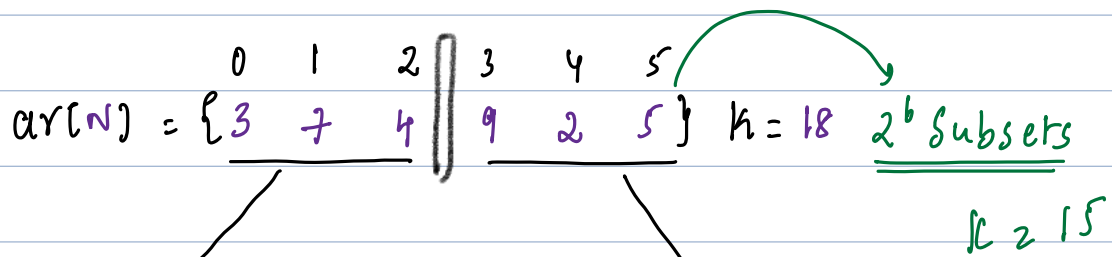
$\{2, 5\}$

$\{9, 5\}$

$\{9, 2, 5\}$

#obs1: If we combine 1 subset in $a[]$ with 1 subset in $b[]$
 can get subset in $c[]$.

#obs2: Can get all subsets in $c[]$
 If we combine 1 subset in $a[]$ with 1 subset in $b[]$



$$a[N/2] = \{ 3 \ 7 \ 4 \} : 2^3$$

$$b[N/2] = \{ 9 \ 2 \ 5 \}$$

$$2^{N/2}$$

$$n$$

{ }	:	0
{ 3 }	:	3
{ 7 }	:	7
{ 4 }	:	4
{ 3 7 }	:	10
{ 7 4 }	:	11
{ 3 4 }	:	7
{ 3 7 4 }	:	14

$$sum1[] = [0 \ 3 \ 4 \ 7 \ 7 \dots]$$

$$sum2[] = [0 \ 2 \ 5 \ 7 \ 9 \dots]$$

0 :	{ }
9 :	{ 9 }
2 :	{ 2 }
5 :	{ 5 }
11 :	{ 9 2 }
7 :	{ 2 5 }
17 :	{ 9 5 }
16 :	{ 9 2 5 }

$2^{N/2}$
 y

#Q: Check if there exists 2 elements x, y such that $x + y = k$

x is subset sum of $a[]$ & y is subset sum of $b[]$

Idea: Generate all pairs of x & y .

TC: $O(2^N)$

Idea2: Store all subset sums of $b[]$ in array $sums[]$

Sort $sums[]$

Take each subset sum[] of $a[]$ as n .

For every n , $y = k - n$, search in sorted array $sums[]$

TC: $O(2^{N/2} + N/2 * 2^{N/2} + 2^{N/2} * N/2)$ SC: $O(2^{N/2})$

TC: $O(N * 2^{N/2})$ SC: $O(2^{N/2})$

Idea3: Store all subset sums of $a[]$ in $sum_a[]$ & sort it

Store all subset sums of $b[]$ in $sum_b[]$ & sort it

Given 2 sorted arrays, check if there exists a pair with $sum = k$.

→ 2 pointers for pair sum.

TC: $O(N/2 * 2^{N/2} + N/2 * 2^{N/2} + 2 * 2^{N/2})$

TC: $O(N * 2^{N/2})$ SC: $O(2 * 2^{N/2}) = O(2^{N/2})$

Idea4: Store all subset sums of $b[]$ in hashset hs .

Take each subset sum[] of $a[]$ as n .

For every n , $y = k - n$, search in hashset hs .

TC: $O(2^{N/2} * 1 + 2^{N/2} * 1)$ SC: $O(2^{N/2})$

TC: $O(2^{N/2})$ SC: $O(2^{N/2})$

Note: Talk about when to apply meet in the middle

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Given $A[]$, $B[]$, $C[]$, $D[]$ no. of quadruplets (i, j, k, l) such that
 $A[i] + B[j] + C[k] + D[l] = S$
 $0 \leq i, j, k, l < N$

 $k=15$

	0	1	2	3	4	5
$A[] =$	{3	4	2	7	3	9}

$B[] =$	{6	3	4	2	6	4}
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$C[] =$	{2	7	6	4	3	5}
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$D[] =$	{6	4	3	8	9	4}
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$AB \quad \#N^2$

$CD \quad \#N^2$

 $\#N^2$