

Todays Activity

Interesting TC based questions

Match following Complexities

1. Linear

$$A: N^{k+b} \rightarrow O(N^{k+b})$$

(1)

2. Logarithmic

$$B: S^{N*2} \rightarrow (S^2)^N = 2S^N$$

(3)

3. Exponential

$$C: N/4 \log_2^N \rightarrow O(N \log_2^N)$$

(5)

4. Polynomial

$$D: 3^0 N + 10^5 \rightarrow O(N)$$

(2)

5. Linear Logarithmic

$$E: 10N + 9N/100 + 340N^2 \rightarrow O(N^2)$$

(6)

6. Quadratic

$$F: 10^3 \log_{\frac{1}{2}}^{4N} \rightarrow O(\log_2^N)$$

(2)

$$\left[\log_2^a + \log_2^b = \log_2^{a+b} \right]$$

$$\log_2^{4N} = \left\{ \log_2^4 + \log_2^N \right\} = 4 + \log_2^N$$

Note:

Exponential: Number raised to power N.

Tricky Questions.

void fun(int N){

int c=0;

for(int i=N/2; j <= N; i++) {

 int j=N;

 while(j > 1) {

 j=j/2;

 }

}



i	j	
$N/2$	$j: [N, N/2, N/4, \dots]$	$\log_2^N + 1$
$N/2+1$	$j: [N, N/2, N/4, \dots]$	$\log_2^N + 1$
$N/2+2$	$j: [N, N/2, N/4, \dots]$	$\log_2^N + 1$
\vdots		\vdots
N	$j: [N, N/2, N/4, \dots]$	$\log_2^N + 1$
$N+1$		

$$\text{Outer loop} = [N/2 .. N] = N - N/2 + 1 \Rightarrow N/2 + 1$$

$$\text{Inner loop} = [N/2+1] \log_2^N$$

Total iterations:

$$= N/2 + 1 + [N/2+1] \log_2^N$$

$$= N/2 + 1 + \frac{N}{2} \log_2^N + \log_2^N$$

$$= O(N \log_2^N)$$

void fun(int N){

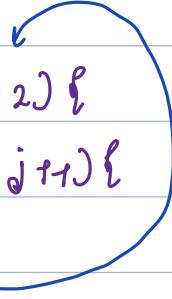
int c=0;

for (int i=N; i>0; i/=2) {

 for (int j=0; j<i; j++) {

 c+=1;

 3



Table

i	j = [0..i-1]
N	j : [0..N-1] = N
N/2	j : [0..N/2-1] = N/2
N/4	j : [0..N/4-1] = N/4
N/8	j : [0..N/8-1] = N/8
:	
i	j : [0..0] = 1

$$\text{Outer loop} = \log_2^N$$

$$\text{Inner loop} = N + N/2 + N/4 + N/8 + \dots - 1 = 2^N - 1$$

$$S = N + N/2 + N/2^2 + N/2^3 + N/2^4 + \dots$$

$$\left\{ \begin{array}{l} 2S = 2N + N + N/2 + N/2^2 + N/2^3 + N/2^4 + \dots - 2 \\ S = N + N/2 + N/2^2 + N/2^3 + N/2^4 + \dots - 2 + 1 \end{array} \right.$$

$$S = 2N - 1$$

$$\text{Total} = \log_2^N + 2N - 1 \Rightarrow O(N)$$

void fun(int N){

 for(int i=0; i < N; i++) {

 for(int j=N; j > i; j--) {

 j = i+1 > i:

 j = i > i : stop

 3

 3

i

j: [N... i+1]

0

j: [N... 1] = N

1

j: [N... 2] = N-1

2

j: [N... 3] = N-2

⋮

N-1

j: [N... N] = 1

N: stop

a b b-a+1

Outer loop: i = [0..N-1] = N-1-0+1 = N

Inner loop: j = $\frac{(N)(N+1)}{2}$

Total iterations = $N + \frac{N^2+N}{2} = \frac{N^2}{2} + \frac{3N}{2} = O(N^2)$

83

```
int fun(int N, int k) {
```

```
    int ans = 0;
```

```
    for (int i = 1; i <= N; i++) {
```

```
        int p = pow(i, k); //  $i^k$ 
```

```
        for (int j = 1; j <= p; j++) {
```

```
            ans = ans + p;
```

3

3

Tabelle

i	j: [1, i^k]
1	[1, 1^k] = 1
2	[1, 2^k] = 2
3	[1, 3^k] = 3
⋮	⋮
N	[1, N^k] = N

Outer loop = N

Inner loop = $1^k + 2^k + 3^k + \dots + N^k$

Higher Order Term

$$\text{if } k=1 : S = 1 + 2 + 3 + \dots + N = \frac{(N)(N+1)}{2} = \left[\frac{N^2}{2} \right]$$

$$\text{if } k=2 : S = 1^2 + 2^2 + 3^2 + \dots + N^2 = \frac{(N)(N+1)(2N+1)}{6} = \frac{2N^3}{6} = \left[\frac{N^3}{3} \right]$$

$$\text{if } k=3 : S = 1^3 + 2^3 + 3^3 + \dots + N^3 = \left[\frac{(N)(N+1)}{2} \right]^2 = \left[\frac{N^2(N+1)^2}{4} \right] = \left[\frac{N^6}{4} \right]$$

$$\text{if ignore } k : S = 1^k + 2^k + 3^k + \dots + N^k \rightarrow \frac{N^{k+1}}{k+1}$$

Final = $\Theta\left(\frac{N^{k+1}}{k+1}\right)$. We cannot neglect $\frac{1}{k+1}$ because it's not a constant

Q4

```
void fun(int arr[], int N) {
```

```
    int j = 0;
```

```
    for (int i = 0; i < N; i++) {
```

```
        while ((j < N) && (arr[i] != arr[j])) {
```

```
            j++;
```

3 i=0

En: arr[N]; $\left(\begin{matrix} a_0 & a_1 & a_2 \dots a_{b_1} & a_y & a_5 \dots \dots \dots a_{N-1} \\ & & j & & \end{matrix} \right)$

i	j	iterations	
0	[0... b ₁]	$b_1 - 0 + 1 = b_1 + 1 =$	$b_1 + 1$
1	[b ₁ +1... b ₂]	$b_2 - (b_1 + 1) + 1 = b_2 - b_1 - 1 + 1 = b_2 - b_1$	b₂ - b₁
2	[b ₂ +1... b ₃]	$b_3 - (b_2 + 1) + 1 = b_3 - b_2 - 1 + 1 = b_3 - b_2$	b₃ - b₂
3	[b ₃ +1... b ₄]	$b_4 - (b_3 + 1) + 1 = b_4 - b_3 - 1 + 1 = b_4 - b_3$	b₄ - b₃
⋮			
N-1	$[b_{N-1}+1 \dots N-1]$	$= N-1 - (b_{N-1} + 1) + 1 = N-1 - b_{N-1} - 1 + 1 = N - b_{N-1} - 1$	$N - b_{N-1} - 1$

Total Outer loop = N

Total Inner loop = $N - b_{N-1}$ = N

Total Iterability = N * N = $2N = O(N^2)$

En: arr[N]; $\left(\begin{matrix} a_0 & a_1 & a_2 \dots a_{b_1} & a_y & a_5 \dots \dots \dots a_{N-1} \\ & & j=0 & & \end{matrix} \right)$

$\Rightarrow : N$

$\Rightarrow : N$

$= 2N = O(N^2)$

void fun(int n) {

```
for (int i=1; i<=2N; i++) {  
    for (int j=1; j<=i; j++) {  
        print("Hello")  
    }  
}
```

3
3

Table

i	j: [1...i]
1	j: [1..1] = 1 +
2	j: [1..2] = 2 +
3	j: [1..3] = 3 +
2 ^N	j: [1..2 ^N] = 2 ^N +

$$\text{Outer loop} = 2^N$$

$$\text{Inner loop} = \underline{1+2+3+\dots+2^N}$$

Sum of first k natural numbers = $\frac{(k)(k+1)}{2}$

$$k = 2^N \longrightarrow \frac{[2^N][2^N+1]}{2}$$

$$\frac{[2^N][2^N+1]}{2} = \frac{2^N \cdot 2^N + 2^N}{2} = \frac{2^{2N} + 2^N}{2}$$

$$\text{Total iteration} = \frac{2^{2N} + 2^N}{2} + 2^N = \frac{2^{2N} + 3 \cdot 2^N}{2} = \cancel{\frac{2^{2N}}{2}} + \frac{3 \cdot 2^N}{2}$$

$$\text{Final Big O} = O(2^{2N}) = O((2^2)^N) = O(4^N)$$

```
void fun(int N){
```

```
    for(int i=1; i<=N; i++) {
```

```
        for(int j=1; j<=2^i; j++) {
```

```
            printf("Hello");
```

```
}
```

```
3
```

i | j: [1..2ⁱ]

1 | j: [1..2¹] = 2¹

2 | j: [1..2²] = 2²

3 | j: [1..2³] = 2³

⋮

N | j: [1..2^N] = 2^N

Outer loop: N

Innner loop: 2¹ + 2² + 2³ + ... + 2^N

$$S = 2^1 + 2^2 + 2^3 + 2^4 + \dots + 2^{N-1} + 2^N$$

$$\cancel{2S} = \cancel{2^2} + \cancel{2^3} + \cancel{2^4} + \cancel{2^5} + \dots + \cancel{2^N} + 2^{N+1}$$

$$S = 2^1 + 2^2 + 2^3 + 2^4 + \dots + 2^{N-1} + 2^N$$

$$S = 2^{N+1} - 2$$

$$S = 2^1 + 2^N - 2 \Rightarrow 2 + 2^N - 2$$

$$\text{Total} = N + \boxed{2 + 2^N} - 2; = O(2^N)$$