

Today's Content

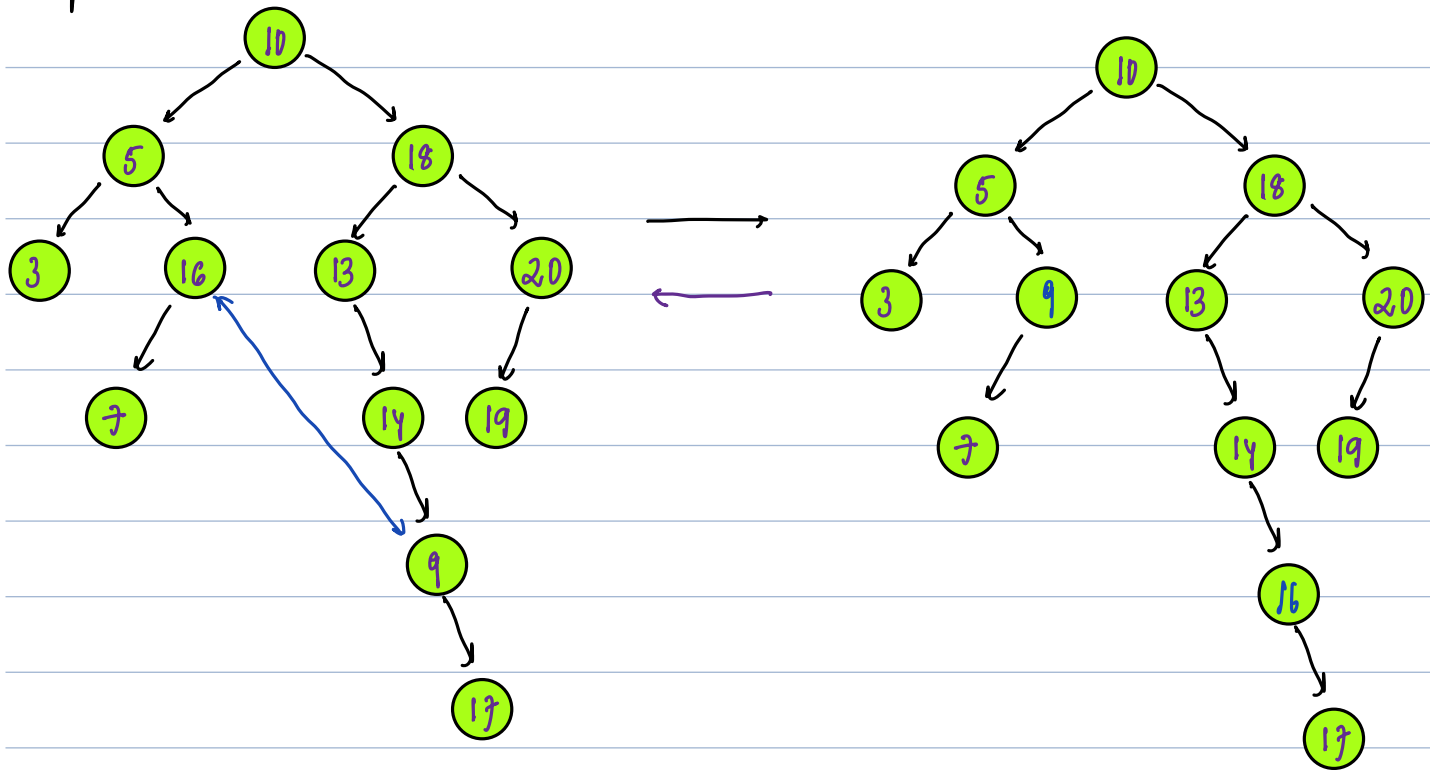
1. ReOrder BST

2. Longest BST in given BT

Recover BST:

Given a BT, which is formed by swapping 2 distinct nodes data in BST
recover original BST, by swapping 2 nodes data back.

Input:



#Ideal: Apply InOrder in Pre

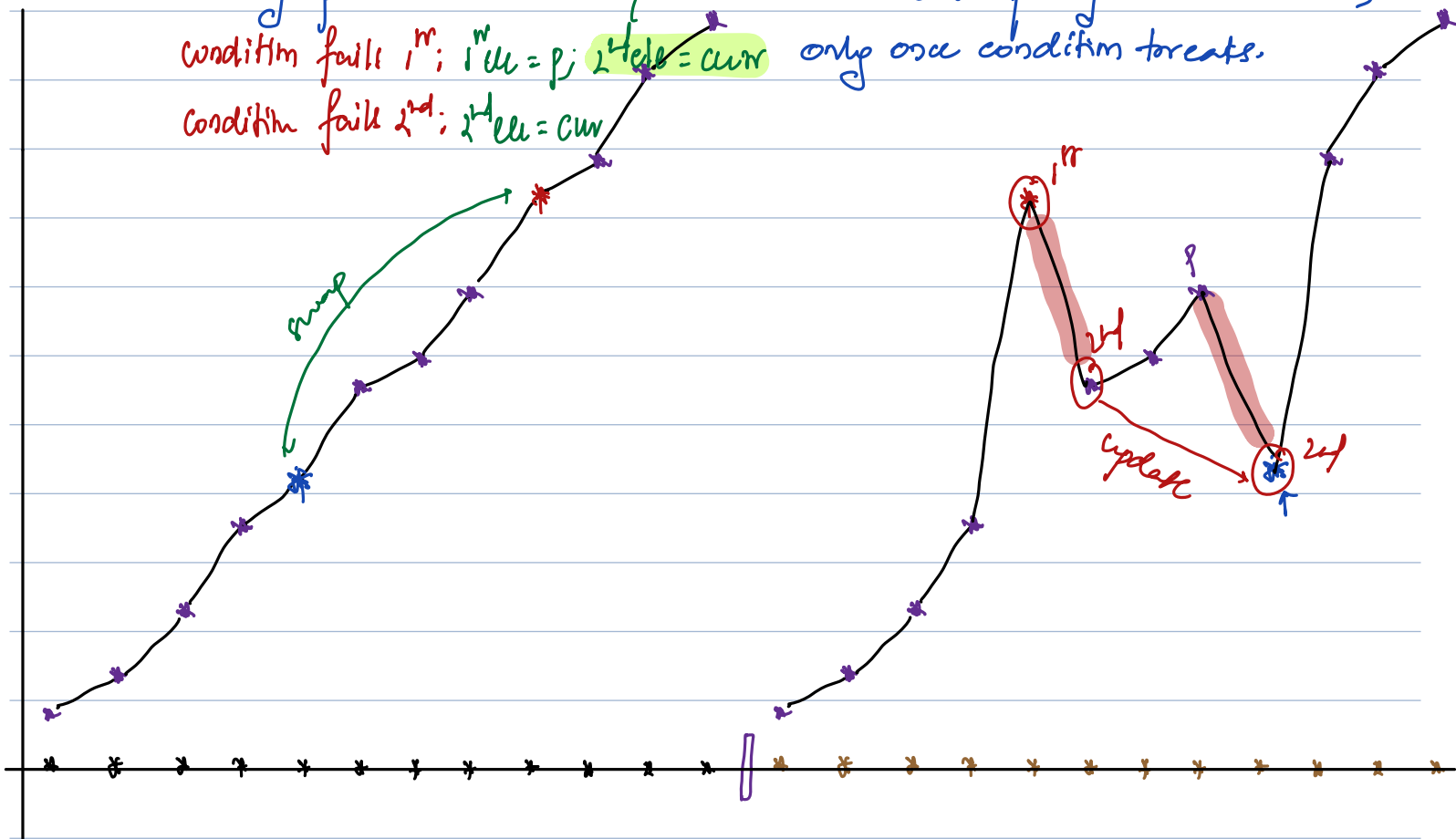
Ideally: prev & cur;

Condition fail 1st; 1st ele = p; 2nd ele = cur

Condition fail 2nd; 2nd ele = cur

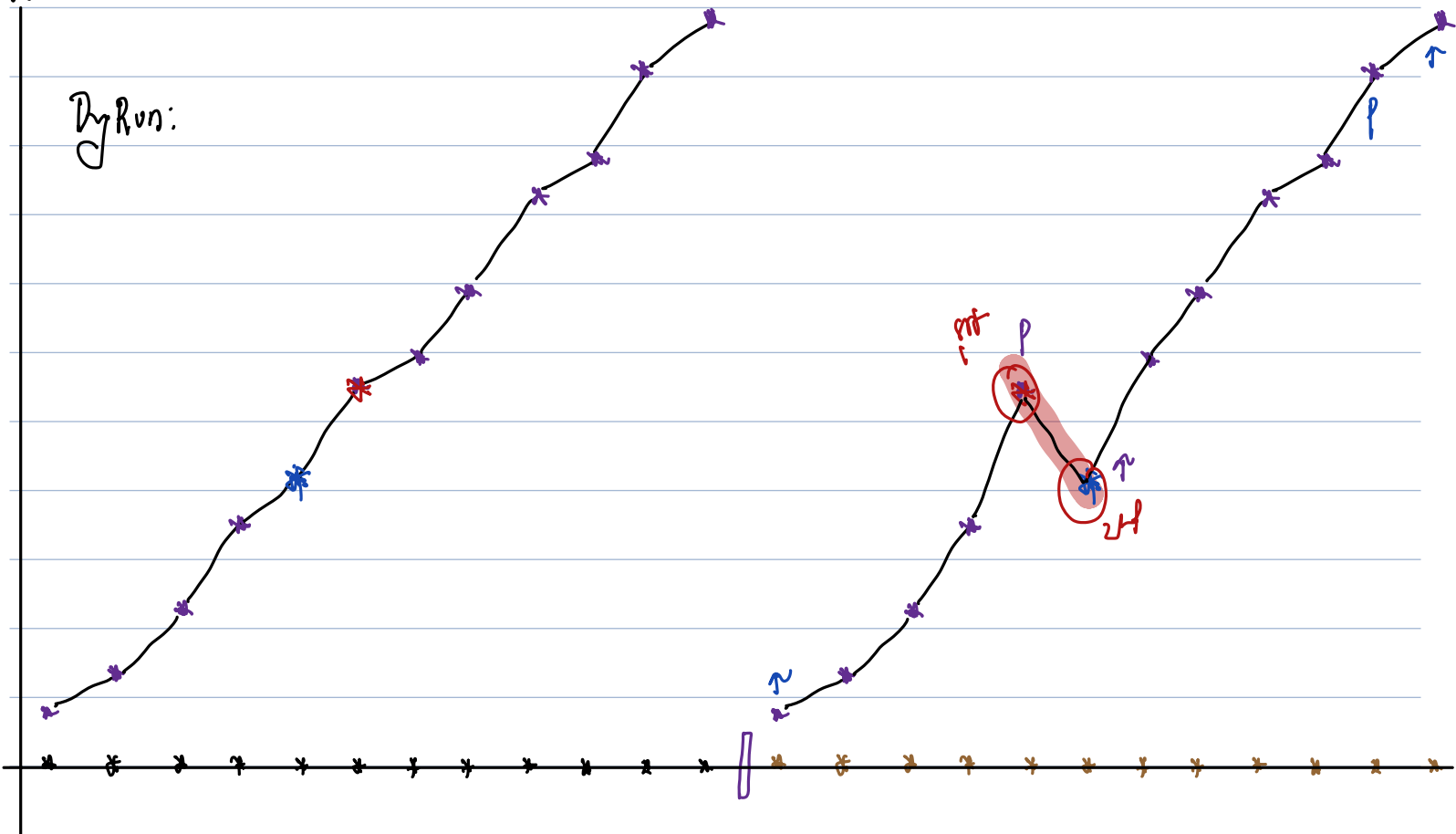
To handle edge case:

When we swap adjacent elements, only once condition breaks.



#slow:

By Run:



#Idea Apply InOrder In Tree } # Discard in checking, if inorder is sorted.
 Ideally: prev & curr;
 Condition fail 1st: $curr = prev$; 2nd: $curr = prev$
 Condition fail 2nd: $curr = prev$

Node *p = nullptr, *f = nullptr, *s = nullptr;

void InOrder(Node *root) { TC: $O(N)$

if (root == nullptr) { return; }

InOrder(root->left);

if (p != nullptr && p->data > root->data) { # p < curr->data && curr = root.

if (f == nullptr) { # fail 1st time

{ f = p; s = root;

else {

s = root;

}

p = root; # updating prev

InOrder(root->right);

Node* restoreBST(Node *root) {

*p = nullptr, *f = nullptr, *s = nullptr;

InOrder(root);

Swap f->data && s->data;

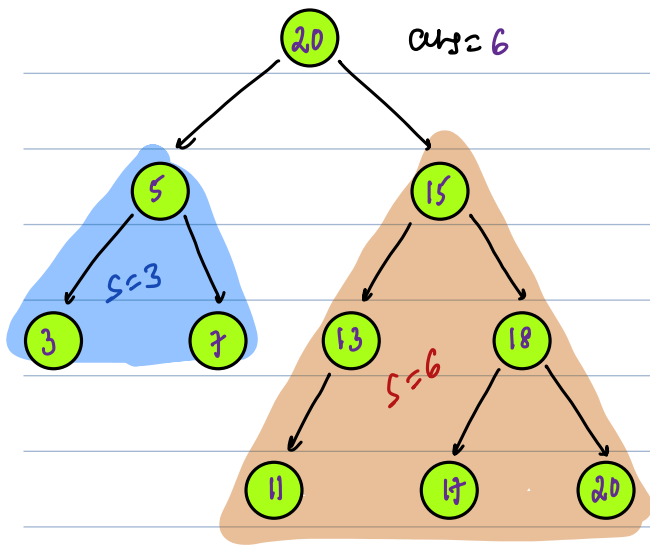
return root;

}

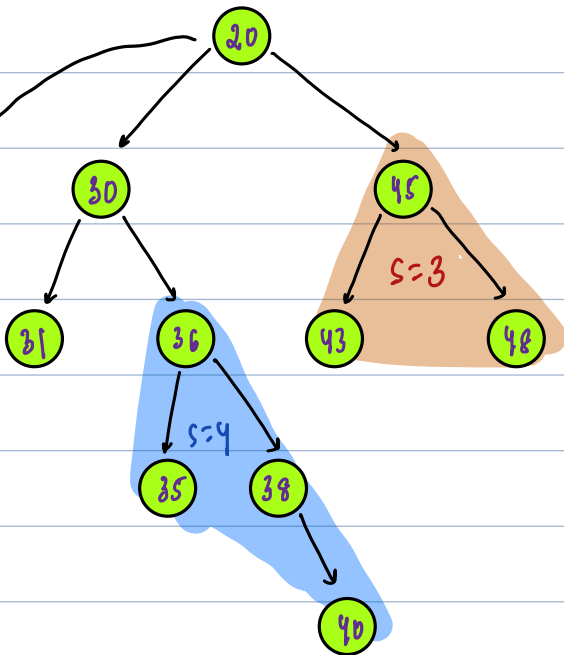


20. Return size of largest subtree in BT which is BST

fn1:



fn2: ans=4



Idea1: *Wrong approach.

1. Apply InOrder on BT & store in arr[].

2. Iterate on arr[]:

Calculate longest part of arr[], which is increasing.

In: 31 30 35 36 38 40 20 43 45 48

↓
It need not be a subtree.

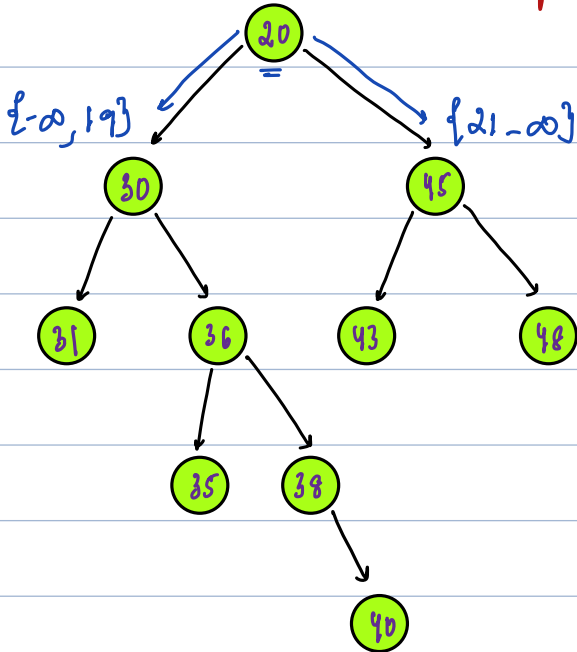
Idea2: for every node;

if [isBST(node)] { # subtree is BST
ans = max(ans, size(node));

3

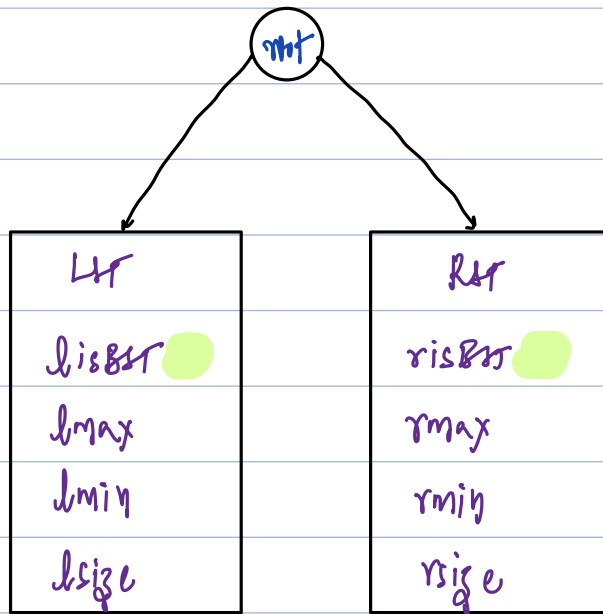
T.C: $O(N * (N + N)) = O(N^2)$

Idea3: Try Top down: Here, we cannot do top-down because we are not sure, if $[-\infty, \infty]$ root is part of BST, because if that any Inf we pass from root can result in wrong values.



Idea 4: Try Bottom up.

Case 1:



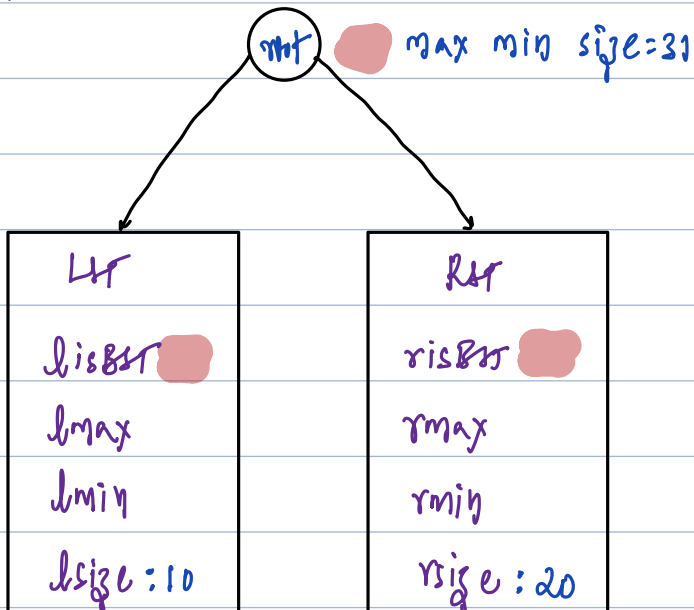
isBST: Is subtree BST or not

max: max of subtree

min: min of subtree

size: size of subtree

Case 2:



#Isbn: Finally root node will return:

isBST max min size:

↳ size of BST, it's not size of largest BST

#Coo: Take a global variable:

if a subtree is BST, compare it's size with ans.

int ans = 0;

vector<int> largestBST(Node *root) { TC: O(N)

if (root == null) {

vector<int> ans(4, 0);

ans[0] = 1; ans[1] = -∞; # Min ans[2] = ∞; # Max ans[3] = 0; # Size

} return ans;

0 1 2 3

vector<int> l = isBST(root->left); # l: isB min max size

vector<int> r = isBST(root->right); # r: isB min max size

vector<int> ans(4, 0);

if ((l[0] == 1) && (r[0] == 1) && (root->data > l[2]) && (root->data < r[1])) {

} ans[0] = 1;

ans[1] = min(l[1], r[1], root->data);

ans[2] = max(l[2], r[2], root->data);

ans[3] = l[3] + r[3] + 1;

if (ans[0] == 1) {

} ans = max(ans, ans[3]);

return ans;

}

int solve(Node *root) {

ans = 0;

largest(root);

return ans;

}

```
int solve(Node *root){
```

```
}
```

#Slight modification:

3Q Return root node of largest subtree which is BST

