

Todays Content

1. PSum Intro
2. Range sum queries

Thursday Class

1. Leetters
2. Count Triplets
3. O/I profision

Friday class

1. Quick C++ Transition
2. Vector in C++
3. Pass by value & Pass by reference
4. Vector & Vech1d or Array of vector
5. String in C++

Given arr[] q s₀, e₀ calculate sum of all elements from s₀ to e₀:

Ex:

arr[10] = {3 4 6 8 9 10 2 7 4 10}
s=2 e=7 : ans = 42

int sum(int arr[], int s, int e) { TC: O(N) SC: O(1)}

int ans=0;

for (int i=s; i<=e; i++) {
 ans = ans + arr[i];

return ans;

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Q1: Given $\text{ar}(n)$ elements & $\text{Qmat}[Q][2]$

In Qmat matrix, we have Q : rows & 2 : columns

Each row in Qmat represents a query.

0^{th} col in row represents : start point of query $\rightarrow s = \text{Qmat}[i][0]$

1^{st} col in row represents : end point of query $\rightarrow e = \text{Qmat}[i][1]$

for every query calculate sum of elements from index $s..e$ in $\text{ar}()$ & print

Constraints:

$$\begin{cases} 1 \leq N \leq 10^5 \\ 1 \leq \text{ar}[i] \leq 10^9 \end{cases}$$

Sum of array elements:

Min: 1

Max: $10^{14} \Rightarrow$ long variable

$$1 \leq Q \leq 10^5$$

$$\text{ar}[1] = \{1\} \quad \text{ar}[10^5] = \{10^9, 10^9, 10^9, \dots, 10^9\}$$

$$0 \leq s \leq e \leq N.$$

Ex:
 $\text{ar}[10] = \{-3, 6, 2, 4, 5, 2, 8, -9, 3, 1\}$

$\text{Qmat}[6][2]$

0:s 1:e Output

Ideas:

| | | | |
|---|---|---|----|
| 0 | 4 | 8 | 9 |
| 1 | 3 | 7 | 10 |
| 2 | 1 | 3 | 12 |
| 3 | 7 | 7 | -9 |
| 4 | 3 | 6 | 19 |
| 5 | 0 | 4 | 14 |

for every query:

Iterate from $s..e$ calculate sum & print

Expected TC: $Q \times N$

```
void RangeSum(int arr[], int N, int &mat[ ][], int s) {
```

```
    for (int i = 0; i <= s; i++) {
```

```
        int s = mat[i][0], e = mat[i][1];
```

long sum = 0; // Need long to avoid wrong answer.

```
        for (int j = s; j <= e; j++) { TC: O(N)
```

```
            sum = sum + arr[j];
```

```
        }
```

```
}
```

3

Calculated TC: $\Theta(N^2)$ SC: $\Theta(1)$

$$\left. \begin{array}{l} 1 \times N = N \\ 1 \times 1 = 1 \end{array} \right\} \Rightarrow 10^5 * 10^5 = 10^{10} \gg 10^8 \text{ TLE}$$

Optimization Idea:

Say we are given csk cricket scores for first 10 overs of batting.
After every over total score is given.

| | | | | | | | | | | |
|--------------|---|---|----|----|----|----|----|----|----|----|
| Overs: | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
| Total Score: | 2 | 8 | 14 | 29 | 31 | 44 | 65 | 79 | 88 | 97 |

$$Q1: \text{Total runs scored in } 10^{\text{th}} \text{ over} = \text{score}[10] - \text{score}[9] = 9$$

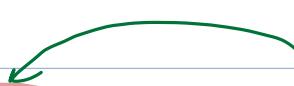
$$Q2: \text{Total runs scored in } 7^{\text{th}} \text{ over} = \text{score}[7] - \text{score}[6] = 16$$

$$Q3: \text{Total runs scored in } 6^{\text{th}} - 10^{\text{th}} \text{ over} = \text{score}[10] - \text{score}[5] = 66$$

$$Q4: \text{Total runs scored in } 3^{\text{th}} - 6^{\text{th}} \text{ over} = \text{score}[6] - \text{score}[2] = 41$$

$$Q5: \text{Total runs scored in } 4^{\text{th}} - 9^{\text{th}} \text{ over} = \text{score}[9] - \text{score}[3] = 74$$

$$\text{Con: Total runs scored in } i^{\text{th}} - j^{\text{th}} \text{ over} = \text{score}[j] - \text{score}[i-1]$$

Scores: 0 1 2 .. i-1  i i+1 .. j

Obs: Cumulative Sum = Total sum till that point

Cumulative Sum from 0th index = PfSum

String PfSum information in array = Pf[]

If we have Pf[] we can answer queries.

Idea: 1. Let's create pfin

2. $pfin[i] = \text{Sum of array elements from } [0..i]$

Ex:

$$\text{ar}[10] = \{ -3, 6, 2, 4, 5, 2, 8, -9, 3, 1 \}$$
$$psum[10] = \{ -3, 3, 5, 9, 14, 16, 24, 15, 18, 19 \}$$

Output: Ans using Psum()

$0:s \quad 1:e \quad psum[e] - psum[s-1]$

| | | | |
|---------------------|---|---|--|
| $i=0 \rightarrow 0$ | 4 | 8 | $psum[8] - psum[3] = 18 - 9 = 9$ |
| $i=1 \rightarrow 1$ | 3 | 7 | $psum[7] - psum[2] = 15 - 5 = 10$ |
| $i=2 \rightarrow 2$ | 1 | 3 | $psum[3] - psum[0] = 9 - (-3) = 12$ |
| $i=3 \rightarrow 3$ | 7 | 7 | $psum[7] - psum[6] = 15 - 24 = -9$ |
| 4 | 3 | 6 | $psum[6] - psum[2] = 24 - 5 = 19$ |
| 5 | 0 | 4 | $psum[4] - psum[-1] : \underline{\text{Error.}}$ |

$\rightarrow \text{Sum}[0..4] = psum[4] = 14$

Query:

```
[s..e] = if (s==0){ // [0..e]
            {
                print(psum[e])
            }
        } else {
            {
                print(psum[e] - psum[s-1])
            }
        }
```

Construct Psum[] for Given arr[]

Step1: To construct pfsum[] we carry forward sum from l→R

0 1 2 3 4

$ar[5] = \{ 3 -2 4 5 6 \}$

$sum = 0 \xrightarrow{\quad} 3 \xrightarrow{\quad} 1 \xrightarrow{\quad} 5 \xrightarrow{\quad} 10 \xrightarrow{\quad} 16$ 1. update sum
 ↓ ↓ ↓ ↓ ↓ 2. store in psum[]

$psum[5] = \{ 3 1 5 10 16 \}$

$\xrightarrow{\quad} pfsum[]$
 $TC: O(N+8) \quad SC: O(N+5) = O(N)$

void RangeSum(int ar[], int N, int &mat[][], int &s){

```
long psum[N];
long sum=0;
for (int i=0; i<N; i++) { → N
    sum = sum + ar[i]
    psum[i] = sum;
}
```

for (int i=0; i<8; i++) { // iterate in each row ⇒ query.

```
int s = mat[i][0], e = mat[i][1];
```

```
if (s==0) { // [0..e]
```

```
    print( psum[e] )
```

```
else {
```

```
    print( psum[e] - psum[s-1] )
```

```
}
```

```
}
```

```
3
```

Note: Modifying same arr[] to Prefix arr[]

$$arr[5] = \{ 3 \ -2 \ 4 \ 5 \ 6 \}$$

$$\text{sum} = 0 \quad 3 \quad 1 \quad 5 \quad 10 \quad 16$$



$$arr[5] = \{ 3 \ 1 \ 5 \ 10 \ 16 \}$$

Issues with above?

1. We will lose arr[] information
2. Datatype of arr[] & psum[] can be different

We can update arr[] with psum value

1. We are no longer using arr values
2. Datatype of arr[] & psum[] should be same.

Note: When multiple range queries given: Think of prefix sum.