

## Todays Content

1. Sorting Intro
2. Bubble Sort
3. Selection Sort
4. Insertion Sort

Sorting: Arranging data in inc/dec order based on parameter

0 1 2 3 4 5

Ex1:  $A[] = \{2 3 9 12 17 19\}$  : Inc based on value

0 1 2 3 4 5

Ex2:  $A[] = \{19 6 5 2 -1 -19\}$  : Dec based on value

0 1 2 3 4 5

Ex3:  $A[] = \{2 7 4 9 6 10\}$  : Inc based on factors

#factors 2 2 3 3 4 4

Sorting library?  $T.C: O(N \log N)$  To sort  $N$  elements.

1. To Sort int arr[]:

sort(arr, arr+n); # It will sort arr[] from arr[0] to arr[N-1] In Inc

sort(arr, arr+n, greater<int>()); # It will sort arr[] In Dec

#include <functional> #for greater

2. To sort vector<int> v;

sort(v.begin(), v.end()); #Sort in increasing Order

sort(v.begin(), v.end(), greater<int>()); #Sort in Decreasing Order

#include <functional> #for greater

Stable Sorting: When 2 data points have same value, their relative order should be same, before & after sorting, that type of sorting algo is called **stable sorting Algo.**

Ex:  $arr[5] = \{1 5 2 6 2\}$  #Sort inc based on value;

#Sort1 =  $\{1 2 2 5 6\}$  #Sort1 is stable

#Sort2 =  $\{1 2 2 5 6\}$  #Sort2 is unstable

Inplace Sorting:

Sorting Algo with  $SC: O(1)$  It is considered as Inplace Sorting.

# All Sorting Algorithms

#Basic: SelectionSort  $\leftarrow$  BubbleSort  $\leftarrow$  InsertionSort

TC :  $O(N^2)$   $O(N^2)$   $O(N^2)$

StableSort: \*

✓

✓

UnstableSort: ✓ ✓ ✓

#Optimized: MergeSort QuickSort HeapSort

TC :

StableSort:

UnstableSort:

#Linear: CountSort RadixSort

TC :

StableSort:

UnstableSort:

## Selectin Sort:

At every step:

Iterate in arr[] select min q bring to its correct position.

0 1 2 3 4 5

arr[] = { 9 8 4 10 6 2 }      mini: index with min value

0 1 2 3 4 5      MinInd      Swap(arr[0] arr[5]);

i=0 j:[0 5] { 9 8 4 10 6 2 }  
mini=0 1 2 2 2 5

swap(arr[0] arr[5])

0 1 2 3 4 5  
i=1 j:[1 5] { 2 8 4 10 6 9 }  
mini=1 2 2 2 2 2

swap(arr[1] arr[2])

0 1 2 3 4 5  
i=2 j:[2 5] { 2 9 8 10 6 9 }  
mini=2 2 2 4 4 4

swap(arr[2] arr[4])

0 1 2 3 4 5  
i=3 j:[3 5] { 2 4 6 10 8 9 }  
mini=3 4 4 4 4

swap(arr[3] arr[4])

0 1 2 3 4 5  
i=4 j:[4 5] { 2 4 6 8 10 9 }  
mini=4 5 5 5

swap(arr[4] arr[5])

0 1 2 3 4 5  
i=5 j:[5 5] { 2 4 6 8 9 10 }  
mini=5 5 5

swap(arr[5] arr[5])

#Pseudo Code:

i: 0 to N-1:

Iterate j: i to N-1 q calculate mini;

swap arr[mini] q arr[i];

void SelectionSort(int arr[], int N) { TC: O(N^2) SC: O(1)}

```
for(int i=0; i<N; i++) {  
    # i: j: [i..N-1] & calculate minj;  
    int mini = i;  
    for(int j=i; j<N; j++) {  
        # arr[mini] && arr[j]  
        if(arr[j] < arr[mini]) {  
            mini = j;  
        }  
    }  
    swap(arr[i] & arr[mini]);  
}
```

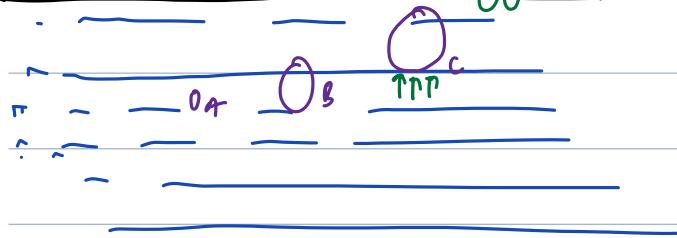
#Selection Sort:

	0	1	2	3	4	5
i=0	j:[0 5]	{ 6 → 5 → 6 → 3 → 8 → 7 }	MinInd	Swap(arr[0] arr[3]);		
	mini =	0 1 1 3 3 3	3	swap(arr[0] arr[3])		

	0	1	2	3	4	5
i=1	j:[1 5]	{ 3 → 5 → 6 → 6 → 8 → 7 }				
	mini =	1 1 1 1 1	1	swap(arr[1] arr[1])		

	0	1	2	3	4	5
i=	j:[ ]	{ 3 5 6 6 8 7 }	# Issue: Relation order between mini =	6 6 is changed		

Idea Bubble Sort: Bigger bubble will come out on top.



Bubble Sort: At each iteration

Compare adj elements, if not in correct order swap

Note: 0 1 2 3 4 5 Increasing

$arr[] = \{ 9, 8, 4, 10, 6, 2 \}$  #small option

$i=0: \{ \overset{8}{\uparrow}, 4, 6, 2, 10, 2 \}$   $i=0, j \leftarrow 5; j = 5 \text{ stop}$

$i=1: \{ 4, \overset{8}{\uparrow}, 6, 2, 10, 2 \}$   $i=1, j \leftarrow 4, j \leftarrow 5-1; j=4$

$i=2: \{ 4, 6, \overset{2}{\uparrow}, 8, 9, 10 \}$   $i=2, j \leftarrow 3; j \leftarrow 5-2; j=3$

$i=3: \{ 4, 2, 6, \overset{8}{\uparrow}, 9, 10 \}$   $i=3, j \leftarrow 2; j \leftarrow 5-3; j=2$

$i=4: \{ 2, 4, 6, 8, \overset{9}{\uparrow}, 10 \}$   $i \rightarrow j \leftarrow N-i-1; j=1$

$i=5: \{ 2, 4, 6, 8, 9, \overset{10}{\uparrow} \}$

void bubbleSort(int[] arr, int n) { TC:  $O(N^2)$  SC:  $O(1)$

for (int i=0; i < n; i++) {

    for (int j=0; j < N-i-1; j++) {

        if (arr[j] > arr[j+1]) { # j = N-1: arr[N-1] > arr[N]: Err.

            swap arr[j] & arr[j+1];

}

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# Stability: # Relative order maintained in small elements

$$ar[] = \{ 6 \ 3 \ 5 \ 6 \ 4 \}$$

$$i=0 \quad \{ 3 \ 5 \ 6 \ 4 \ 6 \}$$

$$i=1 \quad \{ 3 \ 5 \ 4 \ 6 \ 6 \}$$

$$i=2 \quad \{ 3 \ 4 \ 5 \ 6 \ 6 \}$$

3

TODD: Check in place or not.

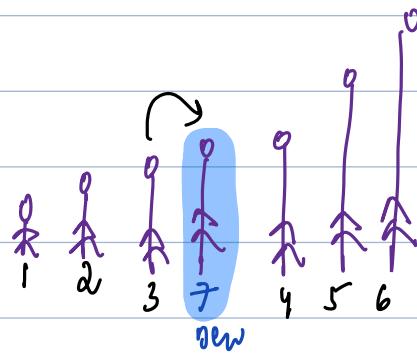
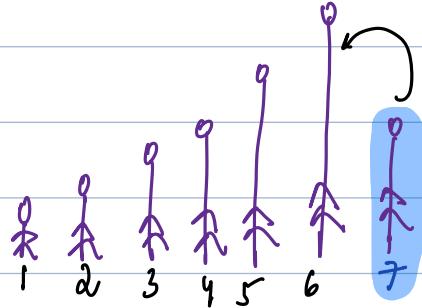
## Insertion Sort:

We insert 1 element in existing sorted data to make entire data sorted.

# while (prev person > new person) {  
    if no prev we stop }  
    We swap



if new person at start position,  
no prev person so we can stop



Insertion Sort:

1st Run:    Sort    Insert     $i = 0$      $\{ 10 \} \quad \{ 9, 4, 8, 6, 2 \}$

$i = 1$      $[ ] \quad \{ 9 \}$

$i = 2$      $[ ] \quad \{ 9, 4 \}$

$i = 3$      $[ ] \quad \{ 9, 4, 8 \}$

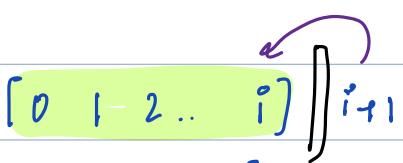
$i = 4$      $[ ] \quad \{ 9, 4, 8, 6 \}$

$i = 5$      $\underline{[ ]}$

Void Insertion(int ar[], int N) { TC: O(N) SC: O(1) }

for (int i=0; i < N-1; i++) {

# i = [0..i] is sorted



int j = i+1; // Insert in sorted data [0...i];

while (j > 0 & ar[j-1] > ar[j]) { # j=0 {ar[-1] > ar[0]}

swap ar[j-1] & ar[j]

j--;

}

}

}