

Today's Activity

Interesting TC based questions

Match following complexities

1. Linear	A: $N^{k+h} \rightarrow O(N^{k+h})$	(1)
2. Logarithmic	B: $5^{N*2} \rightarrow (5^2)^N = 25^N$	(3)
3. Exponential	C: $N/4 \log_2^N \rightarrow O(N \log_2^N)$	(5)
4. Polynomial	D: $2^N N + 10^5 \rightarrow O(N)$	(2)
5. Linear Logarithmic	E: $10N + 9N/100 + 340N^2 \rightarrow O(N^2)$	(6)
6. Quadratic	F: $10^3 \log_2^{4N} \xrightarrow{\hspace{2cm}} O(\log_2^N)$	(4)

$$\left[\log_2^a + \log_2^b = \log_2^{ab} \right]$$

$$\log_2^{4N} = \left\{ \log_2^4 + \log_2^N \right\} = 2 + \log_2^N$$

Note:

Exponential: Number raised to power N.

Tricky Questions.

```
void fun(int n){
```

```
    int c=0;
```

```
    for(int i=N/2; i<=N; i++)
```

```
        int j=N;
```

```
        while(j>1)
```

```
            j=j/2;
```

```
    }
```

```
}
```

i	j	
$N/2$	$j: [N, N/2, N/4, \dots]$	$\log_2^N + 1$
$N/2+1$	$j: [N, N/2, N/4, \dots]$	$\log_2^N + 1$
$N/2+2$	$j: [N, N/2, N/4, \dots]$	$\log_2^N + 1$
\vdots		\vdots
N	$j: [N, N/2, N/4, \dots]$	$\log_2^N + 1$
$N+1$		

Outer loop = $[N/2 \dots N] = N - N/2 + 1 \Rightarrow N/2 + 1$

Inner loop = $[N/2 + 1] \log_2^N$

Total iterations:

$$= N/2 + 1 + [N/2 + 1] \log_2^N$$

$$= N/2 + 1 + \frac{N}{2} \log_2^N + \log_2^N$$

$$= O(N \log_2^N)$$

```
void fun(int N){
```

```
    int c=0;
```

```
    for(int i=N; i>0; i/=2){
```

```
        for(int j=0; j<i; j++){
```

```
            c++;
```

```
        }
```

```
    }
```

Table

i	j = [0..i-1]
N	j : [0..N-1] = N
N/2	j : [0..N/2-1] = N/2
N/4	j : [0..N/4-1] = N/4
N/8	j : [0..N/8-1] = N/8
⋮	⋮
1	j : [0..0] = 1

Outer loop = $\log_2 N$

Inner loop = $N + N/2 + N/4 + N/8 + \dots - 1 = 2N - 1$

$$S = N + N/2 + N/2^2 + N/2^3 + N/2^4 + \dots$$

$$\left\{ \begin{array}{l} 2S = 2N + N + N/2 + N/2^2 + N/2^3 + N/2^4 + \dots - 2 \\ S = N + N/2 + N/2^2 + N/2^3 + N/2^4 + \dots - 2 + 1 \end{array} \right.$$

$$S = 2N - 1$$

$$\text{Total} = \log_2 N + 2N - 1 \Rightarrow O(N)$$

```
void fun(int N){
```

```
  for(int i=0; i<N; i++){
```

```
    for(int j=N; j>i; j--){
```

\Rightarrow

$j = i+1 > i:$

$j = i > i: \text{stop}$

i	j: [N... i+1]
0	j: [N... 1] = N
1	j: [N... 2] = N-1
2	j: [N... 3] = N-2
...	...
N-1	j: [N... N] = 1
N: stop	

$a \quad b \quad b-a+1$

Outer loop: $i = [0, N-1] = N-1-0+1 = N$

Inner loop: $j = \frac{(N)(N+1)}{2}$

Total iterations = $N + \frac{N^2 + N}{2} = \frac{N^2}{2} + \frac{3N}{2} = O(N^2)$

Q3

```
int fun(int N, int k) {
```

```
    int ans = 0;
```

```
    for(int i = 1; i <= N; i++) {
```

```
        int p = pow(i, k);
```

```
        for(int j = 1; j <= p * i^k; j++) {
```

```
            ans = ans + p;
```

```
        }
```

```
    }
```

```
}
```

Table

i	j: [1, i ^k]
1	[1, 1 ^k] = 1 ^k
2	[1, 2 ^k] = 2 ^k
3	[1, 3 ^k] = 3 ^k
⋮	
N	[1, N ^k] = N ^k

Outer loop = N

Inner loop = $1^k + 2^k + 3^k + \dots + N^k$

Higher Order Term

$$\text{if } k=1: S = 1 + 2 + 3 + \dots + N = \frac{(N)(N+1)}{2} = \left[\frac{N^2}{2} \right]$$

$$\text{if } k=2: S = 1^2 + 2^2 + 3^2 + \dots + N^2 = \frac{(N)(N+1)(2N+1)}{6} = \frac{2N^3}{6} = \left[\frac{N^3}{3} \right]$$

$$\text{if } k=3: S = 1^3 + 2^3 + 3^3 + \dots + N^3 = \left[\frac{(N)(N+1)}{2} \right]^2 = \left[\frac{N^2 + N}{2} \right]^2 = \left[\frac{N^4}{4} \right]$$

$$\text{if } \text{input } k: S = 1^k + 2^k + 3^k + \dots + N^k \longrightarrow \frac{N^{k+1}}{k+1}$$

Final = $O\left(\frac{N^{k+1}}{k+1}\right)$. We cannot neglect $\frac{1}{k+1}$ because it's not a constant

Q4

```
void fun(int arr[], int N) {
```

```
    int j = 0;
```

```
    for (int i = 0; i < N; i++) {
```

```
        while ( (j < N) && (arr[i] <= arr[j]) ) {
```

```
            j++;
```

```
    }
```

Ex: arr[N]: $a_0 \ a_1 \ a_2 \dots a_{b_1} \ a_4 \ a_5 \dots a_{N-1}$

\downarrow
j

i	j	iterations
0	$[0 \dots b_1]$	$b_1 - 0 + 1 = b_1 + 1 =$ $b_1 + 1$ \uparrow
1	$[b_1 + 1 \dots b_2]$	$b_2 - (b_1 + 1) + 1 = b_2 - b_1 - 1 + 1 =$ $b_2 - b_1$ \uparrow
2	$[b_2 + 1 \dots b_3]$	$b_3 - (b_2 + 1) + 1 = b_3 - b_2 - 1 + 1 =$ $b_3 - b_2$ \uparrow
3	$[b_3 + 1 \dots b_4]$	$b_4 - (b_3 + 1) + 1 = b_4 - b_3 - 1 + 1 =$ $b_4 - b_3$ \uparrow
\vdots		
N-1	$[b_{N-1} + 1 \dots N-1]$	$= N-1 - (b_{N-1} + 1) + 1 = N-1 - b_{N-1} - 1 + 1 =$ $N - b_{N-1} - 1$

Total Outer loop = N

Total Inner loop = $N - 1 + 1 = N$

Total iterations = $N + N = 2N = O(N)$

Ex: arr[N]: $a_0 \ a_1 \ a_2 \dots a_{b_1} \ a_4 \ a_5 \dots a_{N-1}$

\downarrow
j

$\left. \begin{array}{l} i=0 \xrightarrow{\hspace{15em}} : N \\ j=0 \xrightarrow{\hspace{15em}} : N \end{array} \right\} = 2N = \underline{O(N)}$

```
void fun(int n){
```

```
    for(int i=1; i<= 2^N; i++){
        for(int j=1; j<= i; j++){
            print("hello")
        }
    }
```

Table

i	j: [1...i]
1	j: [1..1] = 1
2	j: [1..2] = 2
3	j: [1..3] = 3
⋮	
2 ^N	j: [1..2 ^N] = 2 ^N

Outer loop = 2^N

Innwr loop = $1 + 2 + 3 + \dots + 2^N$

Sum of first k natural numbers = $\frac{k(k+1)}{2}$

$k = 2^N \longrightarrow \frac{[2^N][2^N+1]}{2}$

$$\frac{[2^N][2^N+1]}{2} = \frac{2^N \times 2^N + 2^N}{2} = \frac{2^{2N} + 2^N}{2}$$

$$\text{Total iteration} = \frac{2^{2N} + 2^N + 2^N}{2} = \frac{2^{2N} + 3 \times 2^N}{2} = \frac{2^{2N}}{2} + \frac{3 \times 2^N}{2}$$

$$\text{Final BigO} = O(2^{2N}) = O((2^2)^N) = O(4^N)$$

```
void fun(int N){
```

```
    for(int i=1; i<=N; i++){
```

```
        for(int j=1; j<=2i; j++){
```

```
            print("Hello");
```

```
        }
```

```
    }
```

```
}
```

i	j: [1..2 ⁱ]
1	j: [1..2 ¹] = 2 ¹
2	j: [1..2 ²] = 2 ²
3	j: [1..2 ³] = 2 ³
⋮	
N	j: [1..2 ^N] = 2 ^N

Out loop: N

Inn loop: 2¹ + 2² + 2³ + ... 2^N

$$S = 2^1 + 2^2 + 2^3 + 2^4 + \dots + 2^{N-1} + 2^N$$

$$\begin{aligned} 2S &= 2^2 + 2^3 + 2^4 + 2^5 + \dots + 2^N + 2^{N+1} \\ S &= 2^1 + 2^2 + 2^3 + 2^4 + \dots + 2^{N-1} + 2^N \\ S &= 2^{N+1} - 2 \end{aligned}$$

$$S = 2^1 + 2^2 + 2^3 + 2^4 + \dots + 2^{N-1} + 2^N$$

$$\text{Total} = N + \boxed{2 \times 2^N - 2} = O(2^N)$$