

Today's Content

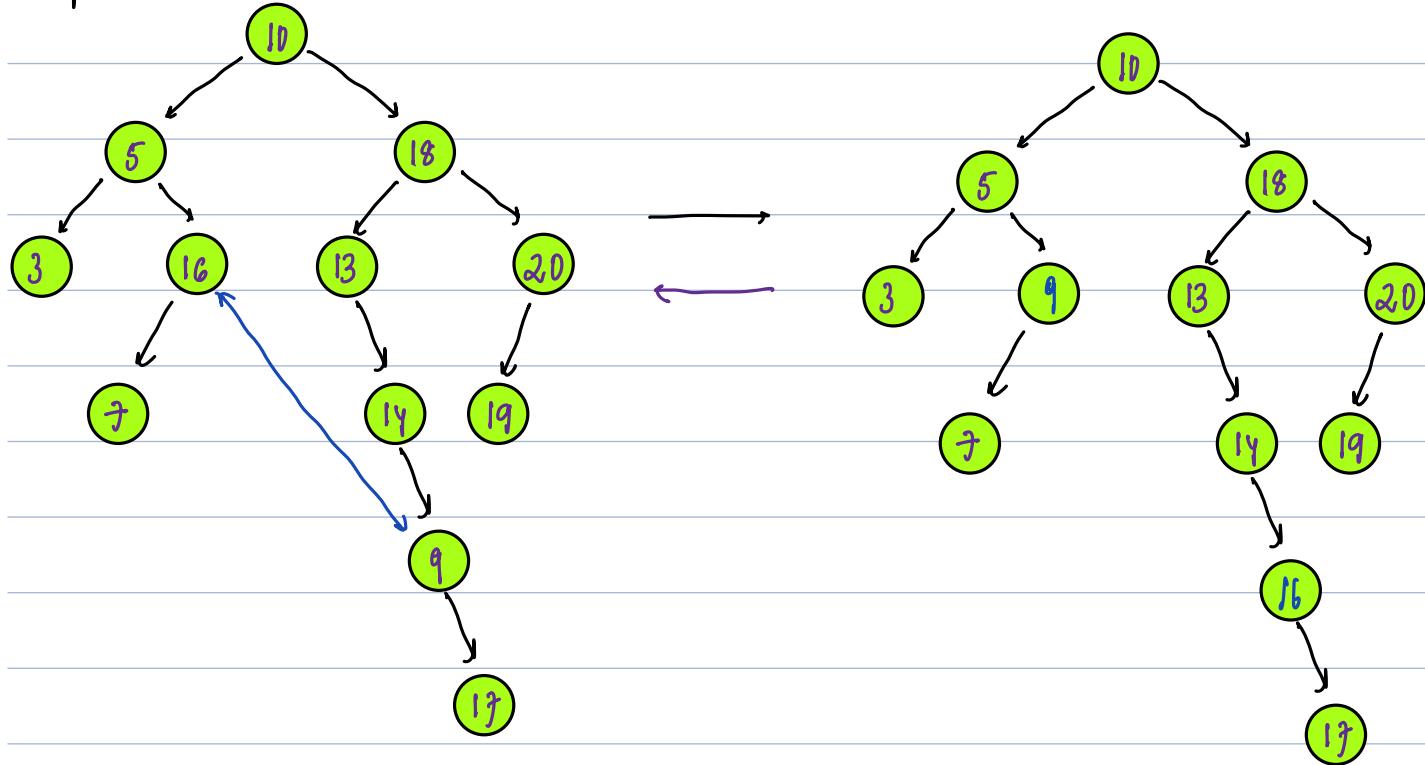
1. ReOrder BST

2. Insert BST in given BT

Recover BST:

Given a BT, which is formed by swapping 2 distinct nodes data in BST
recover original BST, by swapping 2 nodes data back.

Input:



#1deal: Apply InOrder In Pre

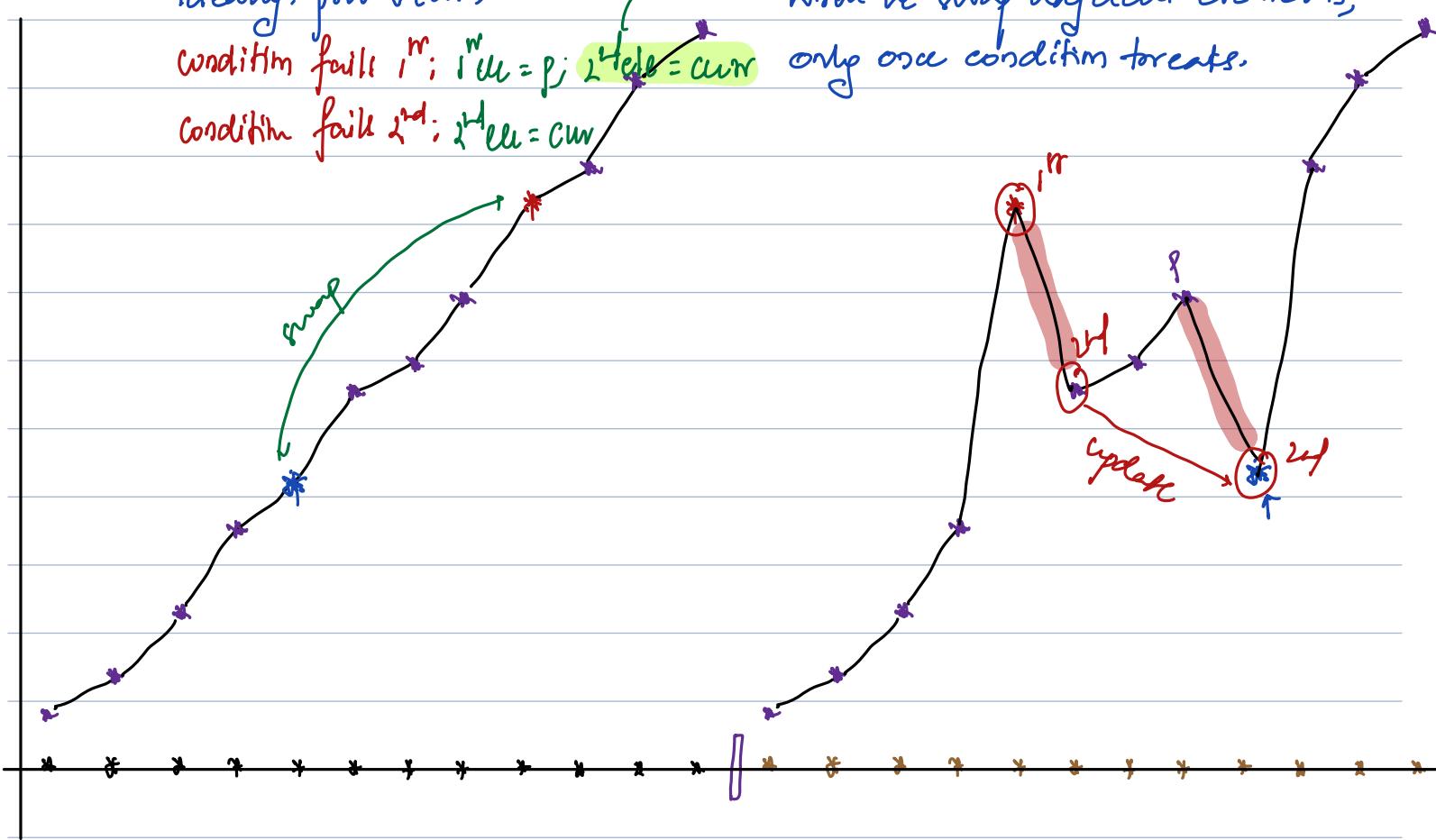
Ideally: $p_{\text{sat}} < \text{LW}$

Condition fails in \mathbb{R}^n ; $\mathbb{R}^n = p$; $\mathbb{R}^n = \text{curv}$

Condition fails 2nd: 2nd ell = cur

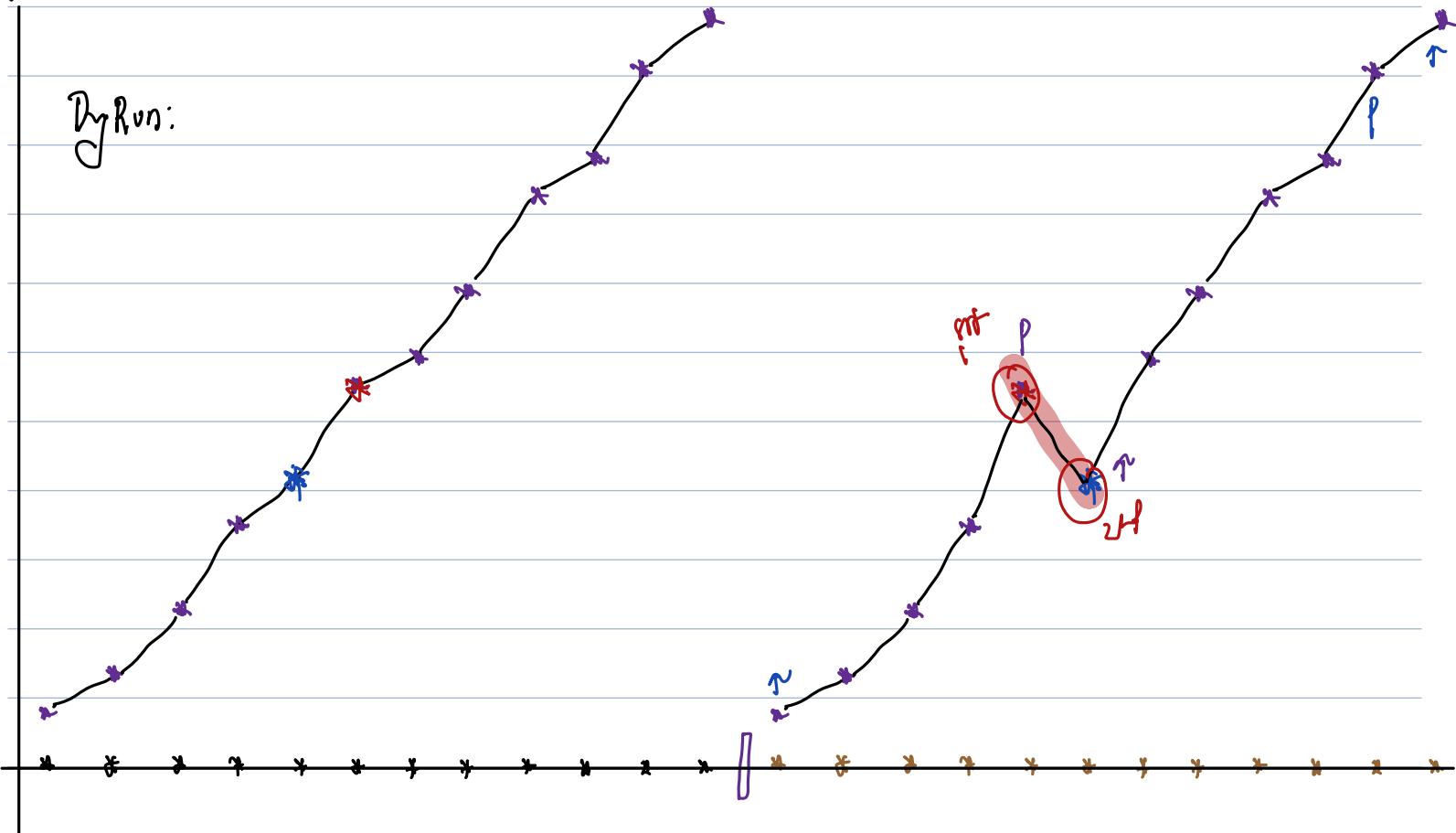
To handle edge case:-

When we swap adjacent elements, only one condition fails.



打蛇：

By Run:



#Idea Apply InOrder in Pre } # Discussed in checking, if inorder is sorted.
 Ideally: $p_{pre} \leq curr$:
 Condition fails 1st: $curr = p_j$; 2nd: $curr = curr$
 Condition fails 2nd: $curr = curr$

Node * $p = \text{nullptr}$, * $f = \text{nullptr}$, * $s = \text{nullptr}$;

void InOrder(Node *root) { TC: OCN)

if (root == nullptr) { return; }

InOrder(root->left);

if ($p \neq \text{nullptr}$ && $p \rightarrow \text{data} > root \rightarrow \text{data}$) { # $p < curr \rightarrow \text{data}$ && $curr = root$.

if ($f == \text{nullptr}$) { # failing 1st time

$f = p_j$ $s = root$;

else {

$s = root$;

$p = root$; # updating p_{pre}

3 InOrder(root->right);

Node * restoreBST(Node *root) {

* $p = \text{nullptr}$, * $f = \text{nullptr}$, * $s = \text{nullptr}$;

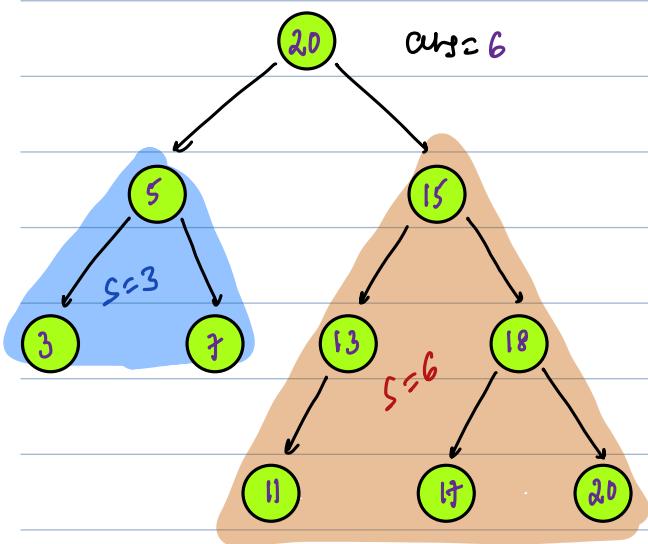
InOrder(root);

Swap $f \rightarrow \text{data}$ && $s \rightarrow \text{data}$;

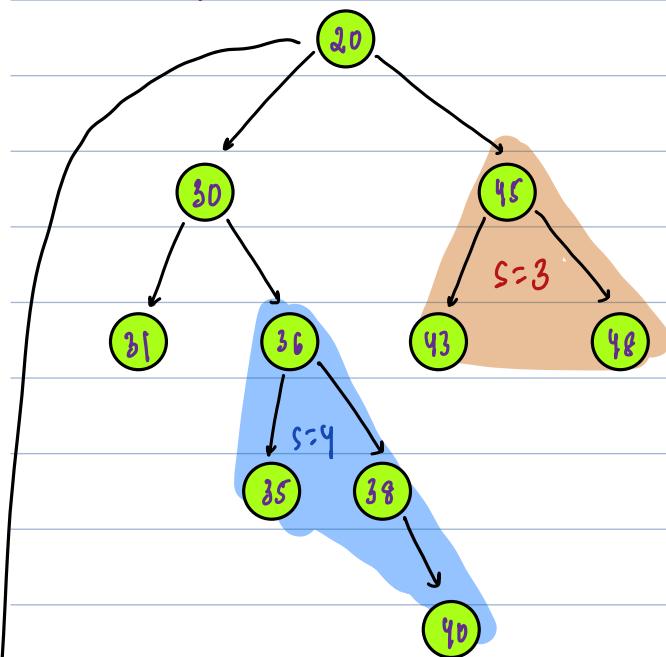
return root;

20. Return size of largest subtree in BT which is BST

fn1:



fn2: ans = 4



Ideal; *Wrong approach.

1. Apply InOrder on BT & store in arr[].

2. Iterate in arr[]:

Calculate longest part of arr[], which is increasing.

In: 31 30 35 36 38 40 20 42 43 45 48

It need not be a subtree.

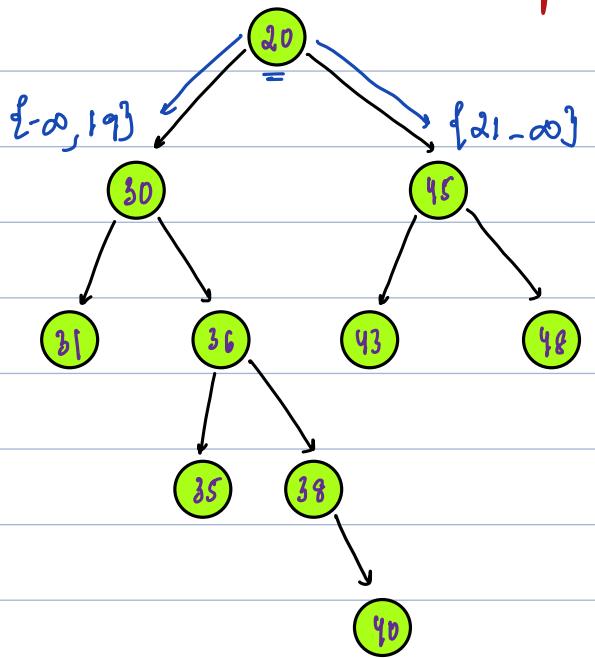
Idea2: for every node;

if C is BST(node) { # Subtree is BST
 ans = max(ans, size(node));

3

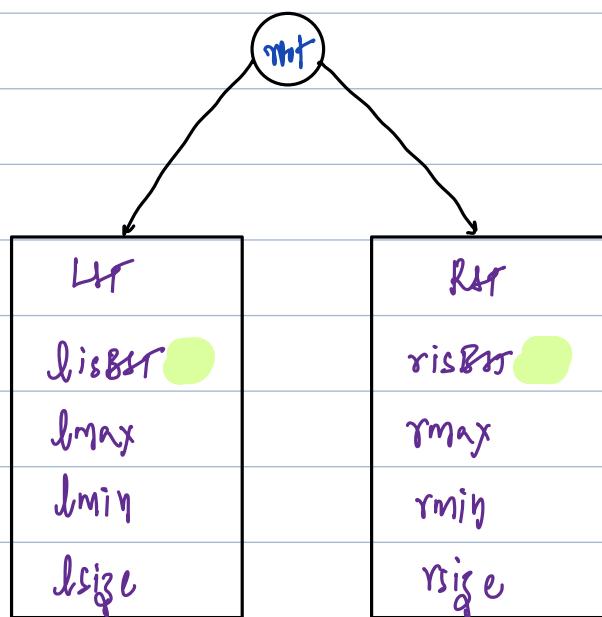
Tc: $O(N \times (N + N)) = O(N^2)$

Idea3: Try Top down: Here, we cannot do top-down because we are not sure, if
 $[-\infty, \infty]$ subtree is part of BST, because if that any inf we pass
from root can result in wrong values.



Ideas: Try bottom up.

Case1:



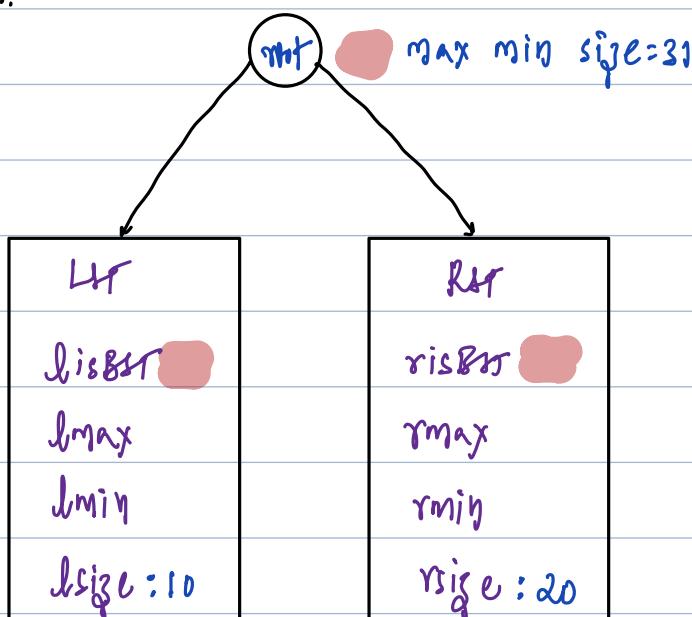
isBST: Is subtree BST or not

max: max f subtree

min: min f subtree

size: size f subtree

Case2:



if isBST: finally root node will return:

isBST max min size:

↳ size of BS, It's not size of largest BST

for loop: Take a global variable;

if a subtree is BST, compare its size with ans.

int sans = 0;

vector<int> largestBST(Node *root) { TC: O(n)}

if (root == null) {

vector<int> ans(4, 0);

ans[0] = 1; ans[1] = +∞; # Min ans[2] = -∞; # Max ans[3] = 0; # Size

} return ans;

0 1 2 3

vector<int> l = isBST(root->left); # l: isB min max size

vector<int> r = isBST(root->right); # r: isB min max size

vector<int> ans(4, 0);

if ((l[0] == 1) && (r[0] == 1) && (root->data > l[2]) && (root->data < r[2])) {

} ans[0] = 1;

ans[1] = min(l[1], r[1], root->data);

ans[2] = max(l[2], r[2], root->data);

ans[3] = l[3] + r[3] + 1;

if (ans[0] == 1) {

} sans = max(sans, ans[3]);

return ans;

}

int solve(Node *root) {

sans = 0;

largest(root);

return sans;

}

```
int solve(Node *root){
```

}

#Slight modification:

30 Return root node of largest subtree which is BST

