

# Todays Content

1. Recursion Intro
2. Steps to Recursive codes

## Use of Recursion:

Sorting

Trus : BT & BST

Backtracking

Dynamic Programming

Graphs

## Rules for Function:

1. Every time a function call is made, it is stored in top of stack
2. When function returns or it's complete execution it will exit stack
3. Even if function return type is void, we can write return;

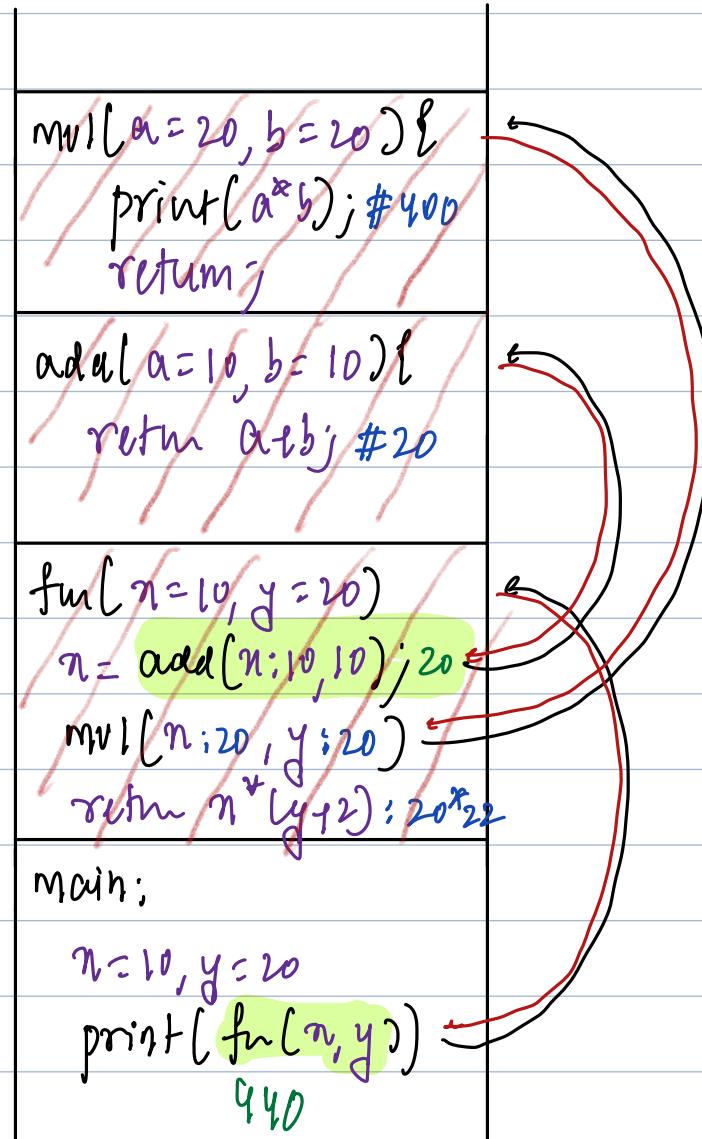
## Function Calls

```
void mul(int a, int b) {
    print(a*b);
    return;
}
```

```
int add(int a, int b) {
    return a+b;
}
```

```
int fun(int x, int y) {
    x = add(x, 10);
    mul(x, y);
    return x * (y + 2);
}
```

```
main() {
    int x = 10, y = 20;
    print(fun(x, y));
}
```



Recursion: Function calling itself

Solving a Problem using SubProblem

Same Problem smaller Input

$$\text{sum}(5) = \underline{1+2+3+4+5}$$

$$\text{sum}(5) = \text{sum}(4) + 5 \quad \# \text{ Sub Problem} = \text{sum}(4)$$

$$\text{sum}(N) = \underline{1+2+\dots+N-1+N}$$

$$\text{sum}(N) = \text{sum}(N-1) + N \quad \# \text{ Sub Problem} = \text{sum}(N-1)$$

Steps to write Recursive Code:

Assumption: Decide what your function does  $\# \{\text{input, does, return}\}$

Main Logic: Solve problem using subproblems  $\#$  Recursive step

Bare Condition: Input for which recursion needs to stop

Ass: Given  $N$ , calculate & return sum of first  $N$  natural numbers

```
int sum(int N){  
    if(N==1){ return 1; }  
    return sum(N-1) + N  
}
```

```
main() {  
    print(sum(4));  
}
```

```
int sum(N=4) {  
    if (N==1) { return 1; }  
    return sum(N-1) + N  
}
```

```
int sum(N=3) {  
    if (N==1) { return 1; }  
    return sum(N-1) + N  
}
```

```
int sum(N=2) {  
    if (N==1) { return 1; }  
    return sum(N-1) + N  
}
```

```
int sum(N=1) {  
    if (N==1) { return 1; }  
    return sum(N-1) + N  
}
```

sum(N=4):

\*if N>=1: return  
return sum(N-1) + N

sum(N=3):

\*if N>=1: return  
return sum(N-1) + N

sum(N=2):

\*if N>=1: return  
return sum(N-1) + N

sum(N=1):

\*if N>=1: return  
return sum(N-1) + N

main()

print(sum(4)): 10

## Importance of Base Conditions:

```
int sum(int n){  
    return sum(n-1) + n;  
}
```

If no base conditions code goes  $\infty$  loop = TLE : Time Limit Exceeded.

In recursion, function calls are stored in top of stack, before  
we go to  $\infty$  loop, stack will reach its limit: stack overflow

TLE: Memory Limit Exceeded

Online Servers:

Time limit = 1sec =  $10^8$  iterations

Space limit : TULLY

Stack limit :

Heap limit :

Array size limit in global vs local :

Issue in below code: Never reach base condition, always keep it

```
int sum(int n){  
    return sum(n-1) + n;  
}  
if (n==1){return 1;}
```

at start

$$\text{fact}(5) = 1 * 2 * 3 * 4 * 5$$

$$\text{fact}(5) = \text{fact}(4) * 5 \quad \# \text{ Sub Problem} = \text{fact}(4)$$

$$\text{fact}(n) = 1 * 2 * \dots * n - 1 * n$$

$$\text{fact}(n) = \text{fact}(n-1) * n \quad \# \text{ Sub Problem} = \text{fact}(n-1)$$

Assumption: Decide what your function does  $\# \{\text{input}, \text{docs, return}\}$

Main logic: Solve problem using subproblems  $\#$  Recursive step

Have a belief that subproblem will work as per assumption.

Bare Condition: Input for which recursion needs to stop

$$0! = 1.$$

Ass: Given  $N$ , calculate & return  $n!$

```
int fact(int n){  
    if(N == 0){ return 1; }  
    return  $\frac{\text{fact}(n-1) * n}{(n-1)! * n}$   
}
```

```
main() {
```

```
    int a = fact(3)
```

3

```
int fact(N=3) {
```

```
* if(N==0) { return 1 }
```

```
    return fact(N-1) * N
```

3

```
int fact(N=2) {
```

```
* if(N==0) { return 1 }
```

```
    return fact(N-1) * N
```

3 1

```
int fact(N=1) {
```

```
* if(N==0) { return 1 }
```

```
    return fact(N-1) * N
```

3 1

```
int fact(N=0) {
```

```
    if(N==0) { return 1 }
```

```
    return fact(N-1) * N
```

3

38 Given  $N$ , print all numbers from 1 to  $N$  in Order

$\text{Inc}(5) = 1 \ 2 \ 3 \ 4 \ 5$

$\text{Inc}(5) = \text{Inc}(4) \ # \ 1 \ 2 \ 3 \ 4 \ 5$   
 $\quad \quad \quad \text{print}[5]$

$\text{Inc}(N) = \underline{1 \ 2 \ 3 \ 4 \dots N-1 \ N}$

$\text{Inc}(N) = \text{Inc}(N-1)$   
 $\quad \quad \quad \text{print}[N]$

Ass: Given  $N$ , print numbers from 1.. $N$  & return nothing

void  $\text{Inc}(\text{int } N)$ {

    if( $N == 0$ ){  $\text{return};$ }

1 2 3 4 .. N-1 N

$\text{Inc}(N-1)$

$\text{print}[N]$

}

main() { Output: 1 2 3

    └ gnc(3)

void gnc( n=3 ) {

\* if[ n=20 ] { return; }

✓ gnc( n-1 )  
✓ print( n )

void gnc( n=2 ) {

\* if[ n=20 ] { return; }

✓ gnc( n-1 )  
✓ print( n )

void gnc( n=1 ) {

\* if[ n=20 ] { return; }

✓ gnc( n-1 )  
✓ print( n )

void gnc( n=0 ) {

if[ n=20 ] { return; }

gnc( n-1 )  
print( n )

48 Given N, print all numbers from N to 1 in the Order

Dec(5) = 5 4 3 2 1

Dec(5) = print(5)  
Dec(4)

Dec(N) = N N-1 N-2.. 1

Dec(N) = print(N)  
Dec(N-1)

Ass: Given N print N-1 in desc order & return nothing

```
void Dec(int N){  
    if(N==0){return;}  
    print(N);  
    Dec(N-1);  
}
```

58 Given  $N$ , print all numbers from  $N \ N-1 \dots 1 \ 1 \ 2 \dots N$  Dec Inc Order

$\text{DecInc}(5) = 5 \ 4 \ 3 \ 2 \ 1 \ 1 \ 2 \ 3 \ 4 \ 5$

$\text{DecInc}(5) = \text{print}(5)$

$\text{DecInc}(4) = 4 \ 3 \ 2 \ 1 \ 1 \ 2 \ 3 \ 4$

$\text{print}(5)$

$\text{DecInc}(N) = N \ \underline{N-1 \dots 1} \ \underline{1 \dots N-1} \ N$

$\text{DecInc}(N) = \text{Print}(N)$

$\text{DecInc}(N-1)$

$\text{Print}(N)$

Ass: Given  $N$ : Print  $N \ N-1 \dots 1 \ 1 \dots N-1 \ N$  & return nothing

void DecInc(int N){

if( $N == 0$ ) { return; }

$\text{Print}(N);$

$\text{DecInc}(N-1);$

$\text{Print}(N);$