

# Todays Content

1. Peak Element

2. Find Unique element

Given a unsorted arr[] with all distinct elements return any one local maxima

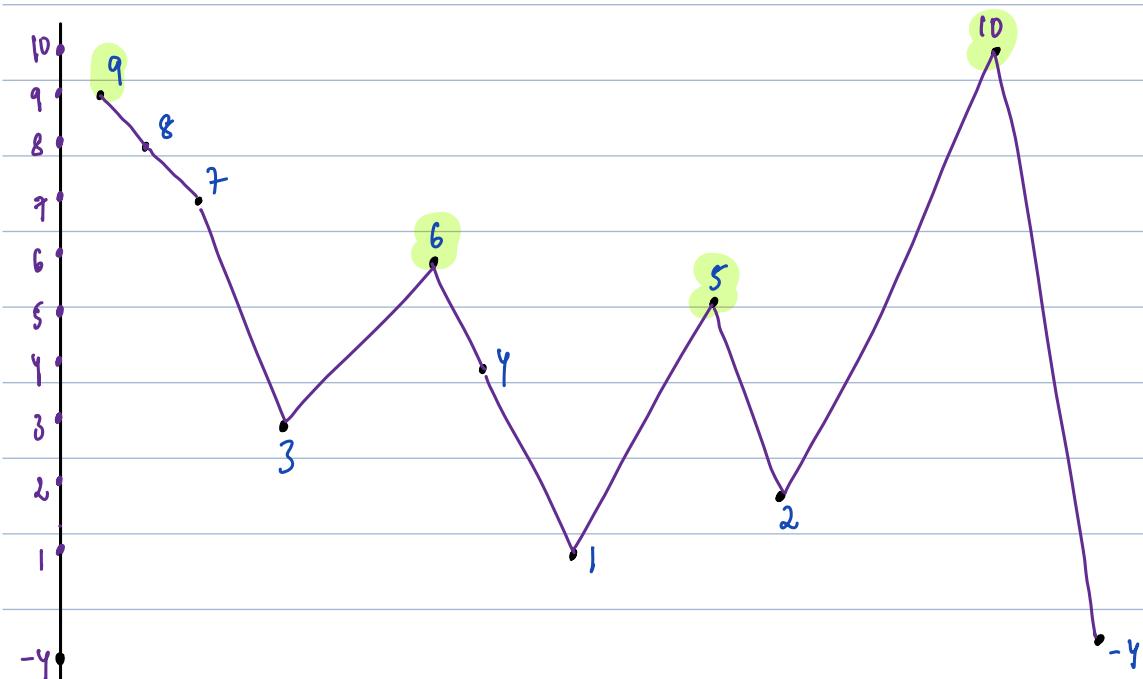
local maxima: An element is said to be local maxima, if  $>$  than  
its adjacent elements {immediate left & right}

#  $\text{arr}[i]$  is local maxima if =  $\text{arr}[i-1] \leftarrow \text{arr}[i] \rightarrow \text{arr}[i+1]$

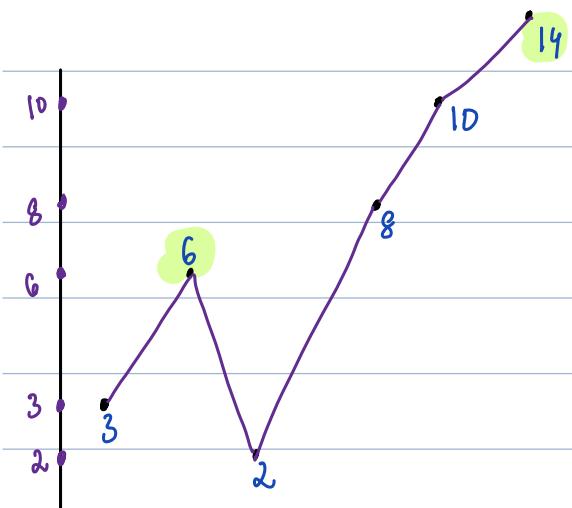
#  $\text{arr}[0]$  is local maxima if =  $\text{arr}[0] > \text{arr}[1]$

#  $\text{arr}[n-1]$  is local maxima if =  $\text{arr}[n-2] < \text{arr}[n-1]$

Ex:  $\text{arr}[11] = \{9, 8, 7, 3, 6, 4, 1, 5, 2, 10, -4\}$  Note: Return any one.



Ex:  $\text{arr}[6] = \{3, 6, 2, 8, 10, 14\}$  local maxima = 6 14



Note: Return any of them.

Obs: At least 1 local maxima exists,  
because max will be  $>$  its  
adjacent elements.

Idea: Iterate & return max of arr[]  
Tc:  $O(n)$  Sc:  $O(1)$

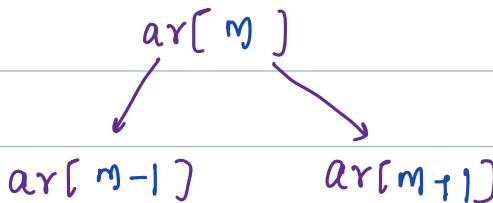
Idea: Iterate on arr[]; Tc:  $O(n)$  Sc:  $O(1)$

For every  $\text{arr}[i]$ , check if it's local maxima by comparing  
with adjacent elements, if it is return  $\text{arr}[i]$ ;

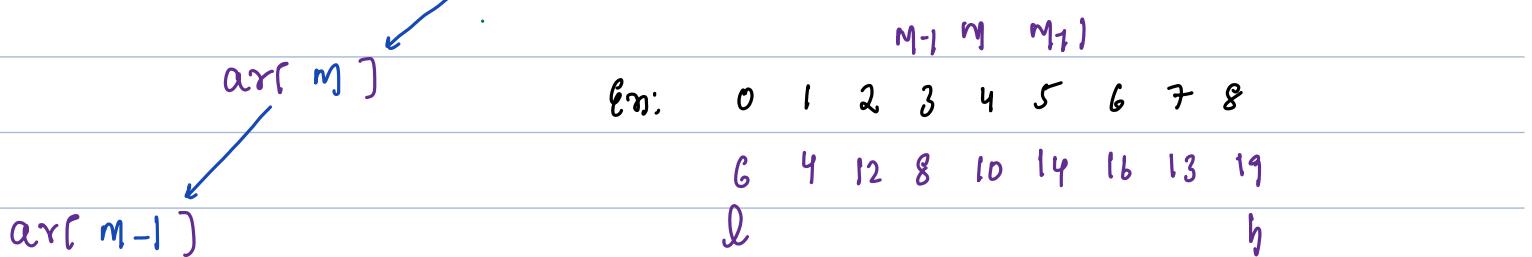
Ideas:

Target = Any no local maxima Search Space =  $\ln \text{arrP}$

Care-1 if ( $\text{ar}[m] > \text{ar}[m+1]$  &  $\text{ar}[m] > \text{ar}[m+2]$ ) { return  $\text{ar}[m]$  }

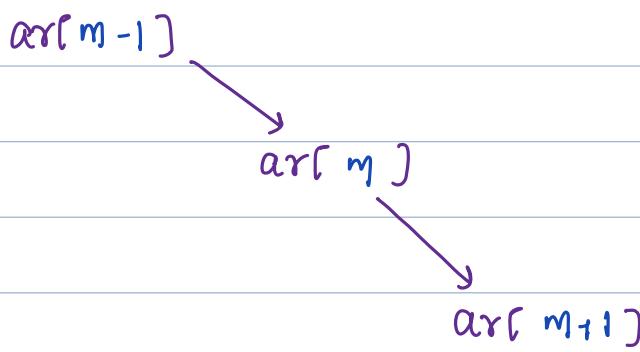


Case-2 if ( $\text{arr}[m] < \text{arr}[m+1]$ ) { # go to right, because it's a guarantee we  
                   $\text{arr}[m+1]$  can find at least 1 local maxima on right }

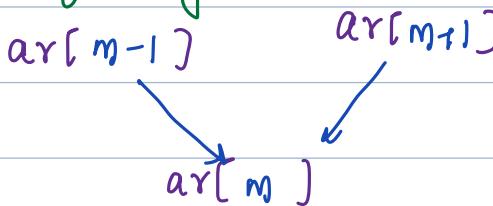


Case-3 # go to left;

guarantees local maxima exists



## Case-IV: #got my side.



#Con: After landing at mid:

On the side where data increases, it's a guarantee local maximum exists in that side.

int localmaxima( vector<int> &arr ) { TC:  $O(\log \frac{n}{2})$  SC:  $O(1)$

int n = arr.size();

if (n == 1) { return arr[0]; }

if (arr[0] > arr[1]) { return arr[0]; }

if (arr[n - 1] > arr[n - 2]) { return arr[n - 1]; }

int l = 0, h = n - 1;

Issue:

M = 0:  $arr[0] > arr[-1]$

M = N - 1:  $arr[N - 1] > arr[N - 2]$  &  $arr[N - 1] > arr[N]$

while (l <= h) {

int m = (l + h) / 2;

if (arr[m] > arr[m - 1] && arr[m] > arr[m + 1]) { return arr[m]; }

else if (arr[m] < arr[m + 1]) {

l = m + 1; # goto right;

arr[m+1]

arr[m]

else {

h = m - 1; # goto left;

arr[m-1]

arr[m]

}

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# Note: When comparing adjacent elements, be careful about corner elements

## Q: Google

Given a unsorted arr[] return any one local maxima

local maxima: An element is said to be local maxima, if  $\geq$  than its adjacent elements [immediate left & right]

$$\underline{arr[i] \geq arr[i-1]} \text{ & } \underline{arr[i] \geq arr[i+1]}$$

If  $i=0$ :  $arr[0]$  is local maxima  $\Rightarrow arr[0] \geq arr[1]$

If  $i=N-1$ :  $arr[N-1]$  is local maxima  $\Rightarrow arr[N-1] \geq arr[N-2]$

int localmaximaRepeating (vector<int> &arr) {

Q. Every element occurs twice, except for 1, find a unique element

Note: Duplicates are adjacent to each other

Ex1: 0 1 2 3 4 5 6 7 8 9 10

$\text{arr}[] = \{ 6, 6, 2, 2, 7, 9, 9, 4, 4, 10, 10 \}$  ans = 7

$\text{arr}[] = \{ 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14 \}$

$\text{arr}[] = \{ 3, 3, 1, 1, 8, 8, 10, 10, 19, 6, 6, 2, 2, 4, 4 \}$  ans = 19

# Idea1: Calculate sum of all arr[] elements & return unique element

TC:  $O(N)$  SC:  $O(1)$

# Idea2:

Target: Unique element SearchSpace: In arr[]

$\text{arr}[] = \{ 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12 \}$

if m is here: go right

On left:

1<sup>st</sup> occur of m: even

left of unique go right

if m is here: go left

On right

1<sup>st</sup> occur of m: odd

right of unique: go left

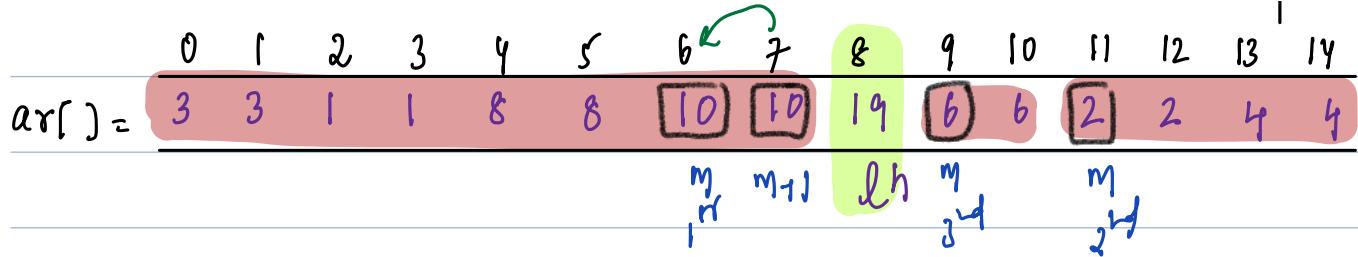
# Note: If a possible that mid can have m 2<sup>nd</sup> occurrence

if mid and m 2<sup>nd</sup> occurrence bring to 1<sup>st</sup> occurrence

if  $[\text{arr}[m]] = [\text{arr}[m-1]]$  {

$m = m - 1$ ; # bring to 1<sup>st</sup> occurrence

}



$l \ h \ m \ ar[m-1] = ar[m]$  left or right

0	14	7	2 <sup>0</sup> occ	$m = m - 1; j$	$m \& .2 = 0$ : left: goto right, $l = m + 2;$
8	14	11	1 <sup>1</sup> occ	$j$	$m \& .2 = 0$ : right: goto left, $h = m - 1;$
8	10	9	1 <sup>2</sup> occ	$j$	$m \& .2 = 0$ : right; goto left, $h = m - 1;$
8	8	8	$ar[8] != ar[7]$	$\& ar[8] != ar[9]$	return $ar[8]$

int unique(vector<int> &arr){ TC:  $O(\log \frac{n}{2})$  SC:  $O(1)$  }

int n=arr.size();

int l=0, h=n-1;  
while(l <= h){  
 int m=(l+h)/2;

Note: if m came to start or last index, in that case, that itself is ans, in this case

if( $m == 0 \&& m == n-1$ ) { return arr[m]; } problem

if( $ar[m] == ar[m-1] \&& ar[m] == ar[m+1]$ ) { return arr[m]; }

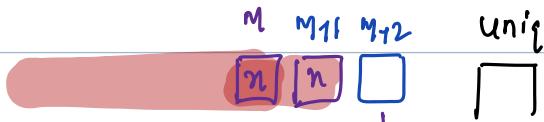
if( $ar[m-1] == ar[m]$ ) { #update m to 1<sup>st</sup> occurrence.

}  $m = m - 1;$

if( $m \& .2 == 0$ ) { #1<sup>st</sup> occurs even: left go right }

$l = m + 2$

3



else { #1<sup>st</sup> occurs odd right: go left }

$h = m - 1;$

3

