

Todays Content

1. Cycle Detection
2. If you Ennts remove you
3. Intersection of 2 Linked List
4. Proof for cycle detection

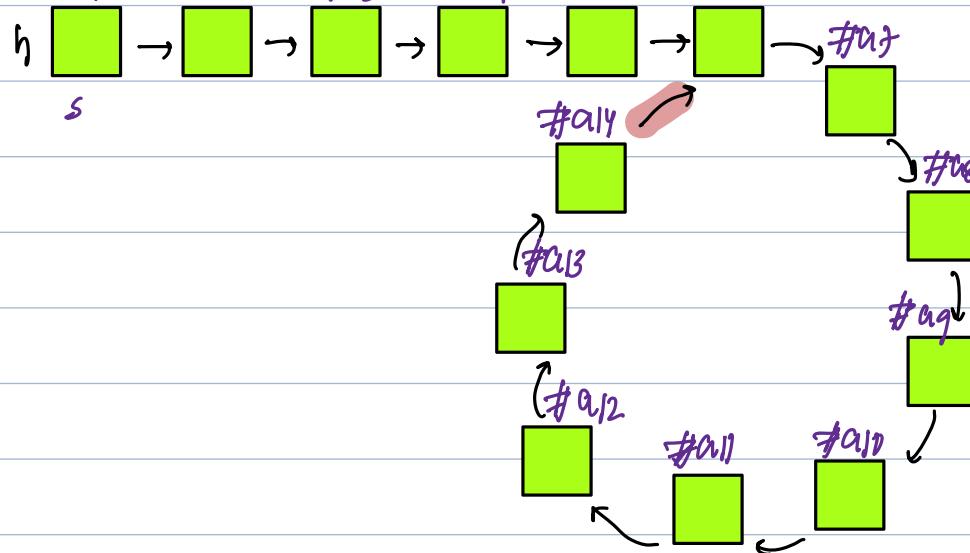
Q1 Given a head node of Linked List, check for cycle detection?

#Note1: When last node points to any prev node, that's when cycle formed

#Note2: If no cycle return null.

#Note3: If cycle is there, remove cycle {lastnode → nullptr}
 Return start/first node if cycle

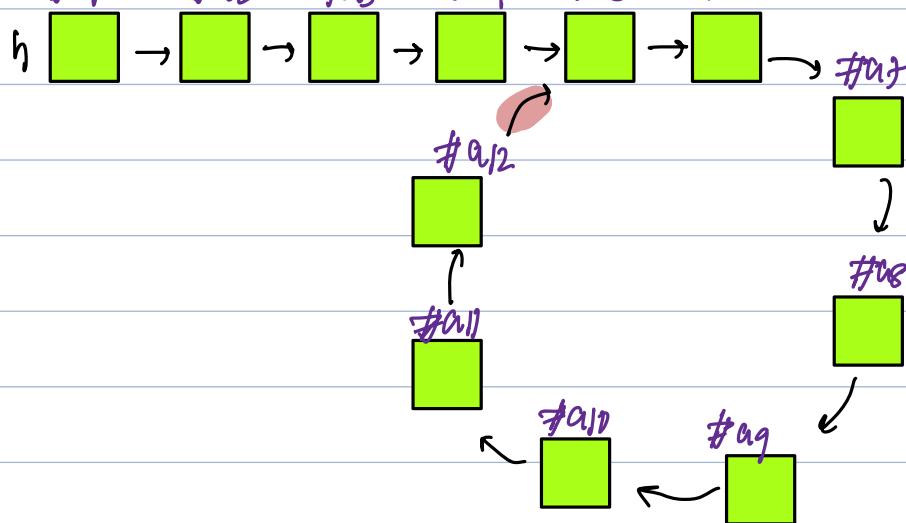
Ex1: #a1 #a2 #a3 #a4 #a5 #a6



Cycle exit

1. break cycle
2. return a6

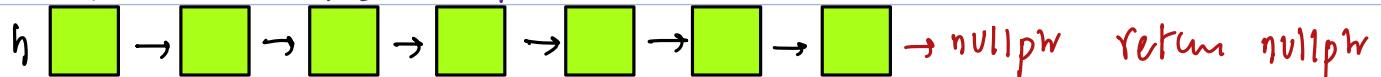
Ex2: #a1 #a2 #a3 #a4 #a5 #a6



Cycle exit

1. break cycle
2. return a6

Ex3: #a1 #a2 #a3 #a4 #a5 #a6 #a7



Ex4: #a1 #a2 #a3 #a4 #a5 #a6 #a7 #a8

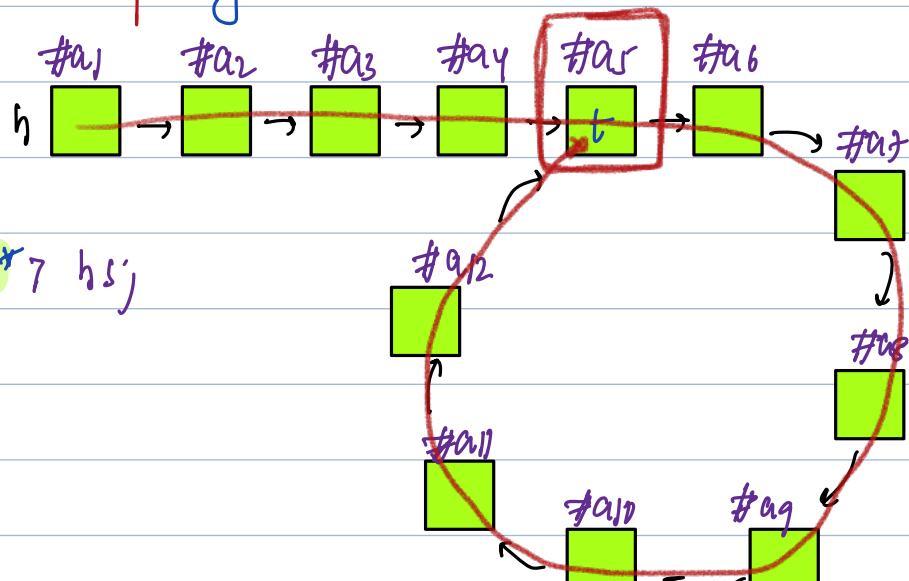


Ideas: Create a HashSet Node^* $T.C: O(N)$ $S.C: O(N) \rightarrow$ HashSet

Iterate on Linked List & Insert each node address.

Obs1 # if a node address already exists: cycle exists

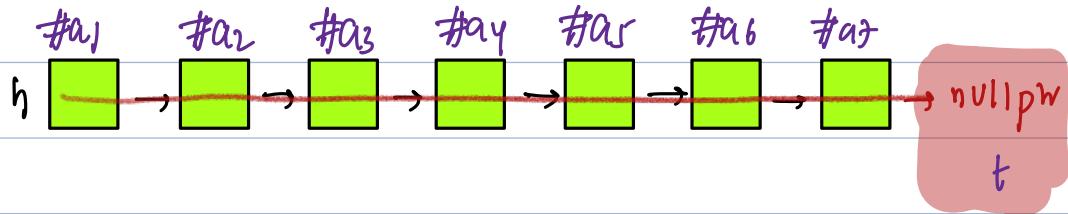
Obs2 # if $t = \text{nullptr}$ you don't exist



unordered_set<Node*> hs;

```
#a1 #a2 #a3  
#a4 #a5 #a6  
#a7 #a8 #a9  
#a10 #a11 #a12
```

Ex2:

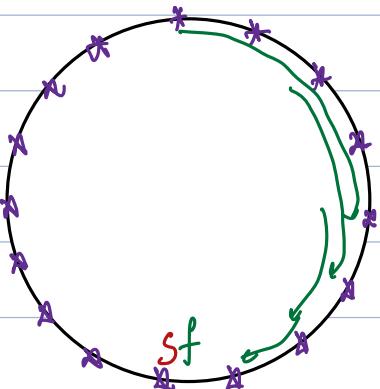


```
#a1 #a2 #a3  
#a4 #a5 #a6  
#a7
```

N_{eff}^* $\propto \ln(N_{\text{eff}} + b)$

3

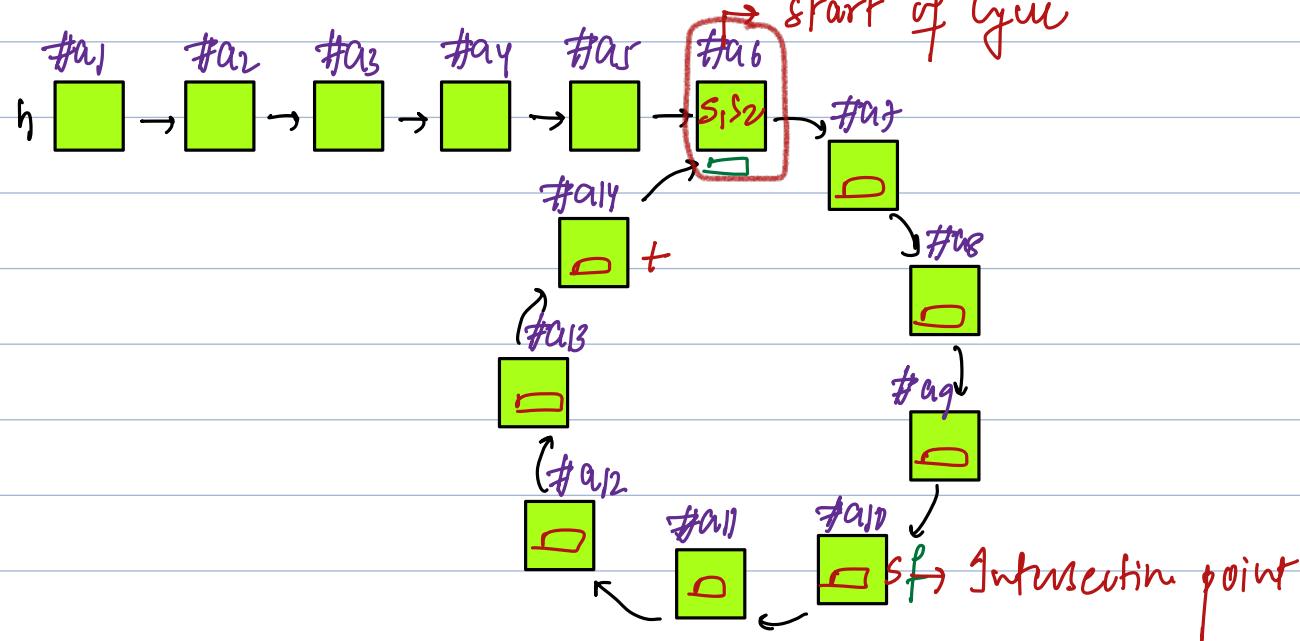
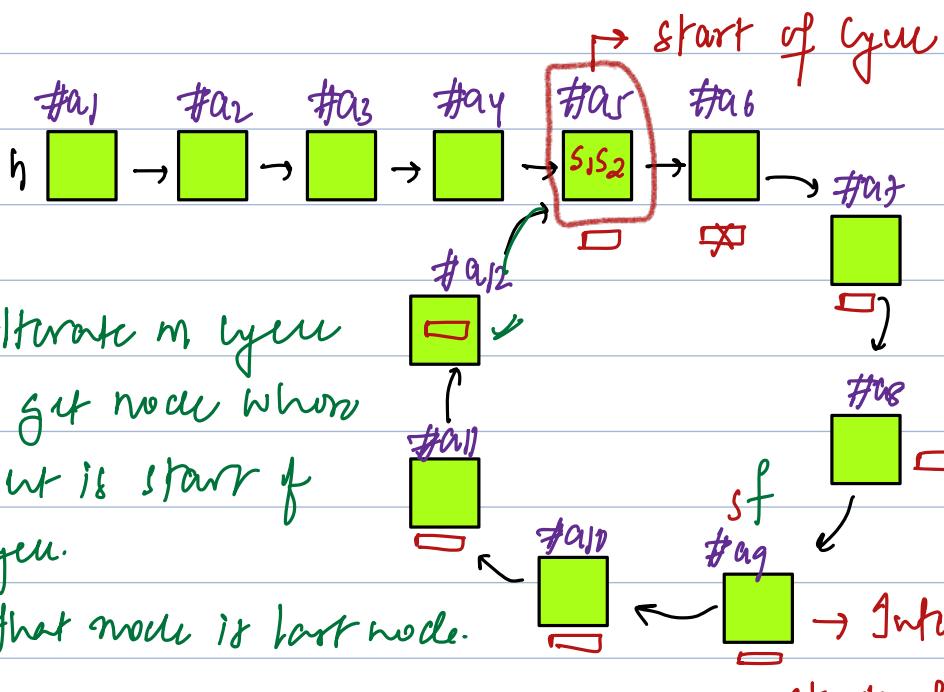
Ideas:

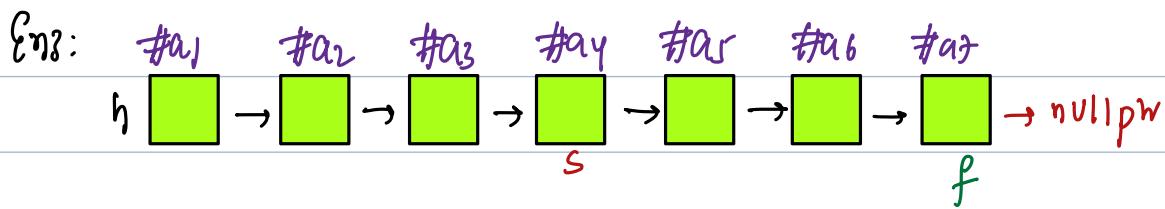


Distance between them

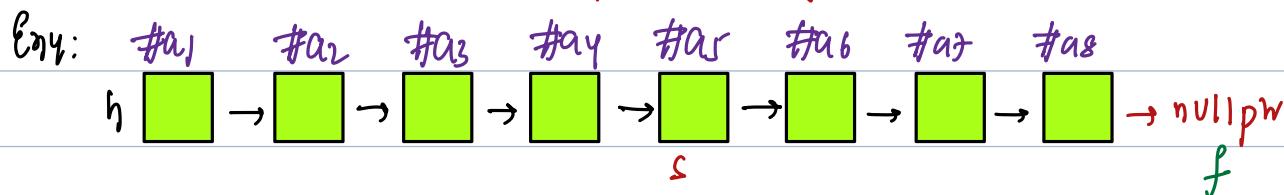
Initial: $4 \rightarrow 3 \rightarrow 2 \rightarrow 1 \rightarrow 0$

In a circle, if 1:step : 2:steps, irrespective of initial distance, they will meet, because distance between will decrease by 1 step after each iteration in circle (d=20, meeting each other)





#Cont1: if (f->next == nullptr) { No loop }



#Cont2: if (f == nullptr) { no loop }

Include both Cont1 & Cont2, because we are not sure of even or odd length

Steps:

1.

2.

3.

Q1 Given a head node of Linked List, check for cycle detection?

#Note1: When last node points to any prev node, that's when cycle formed

#Note2: If no cycle return null.

#Note3: If cycle is there, remove cycle {last node \rightarrow nullptr}

Return start / first node of cycle.

Node* detectCycle(Node *h) { TC: O(N) SC: O(1)}

Node *s = h, *f = h;

bool cycleExist = false;

while(f != nullptr && f->next != nullptr){

s = s->next

f = f->next->next;

if(s == f){

cycleExist = true;

break;

}

if(cycleExist == false){

} return nullptr;

cycle detection

Node *s₁ = h, *s₂ = s₁;

while(s₁ != s₂) {

s₁ = s₁->next;

s₂ = s₂->next;

start of cycle

Node *start = s₁;

Node *temp = start;

while(temp->next != start){

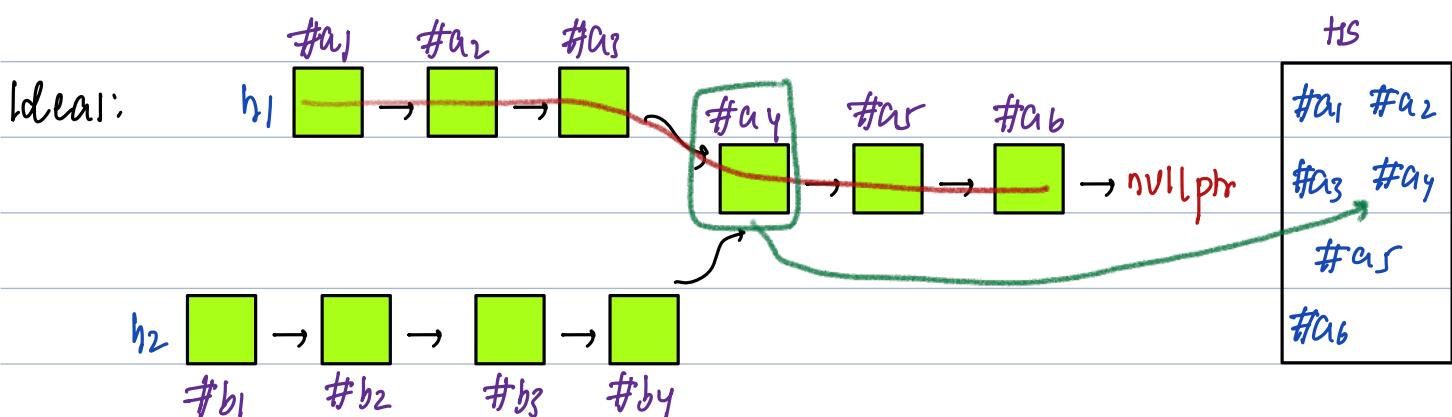
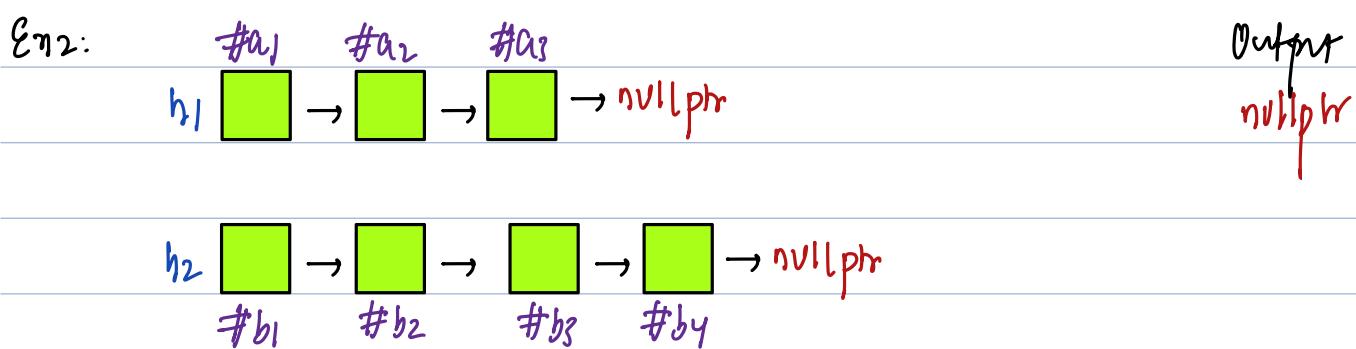
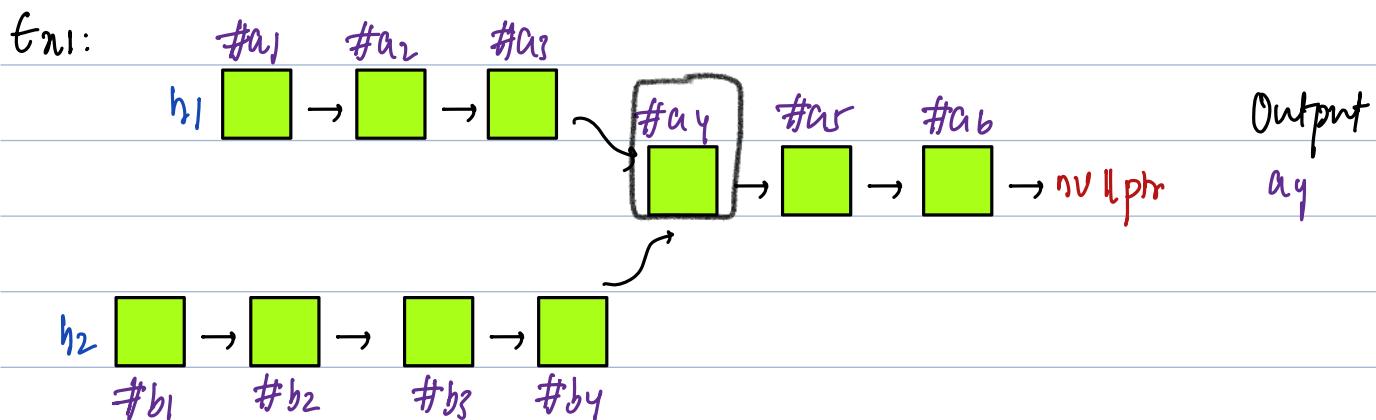
} temp = temp->next;

temp->next = nullptr;

breaking cycle

} return start;

Q Given the heads of two singly linked-lists
 Return nodes at which the two lists intersect.
 If the two linked list have no intersection at all, return null.



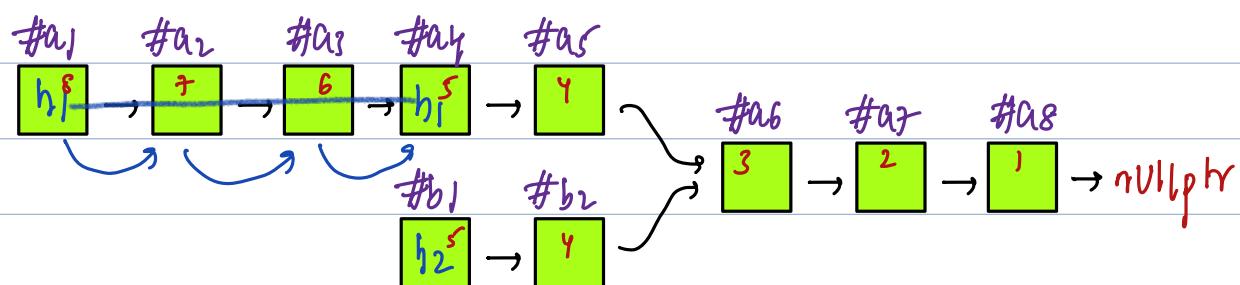
1. Iterate on h_1 & Insert all addresses in hashset $hs \rightarrow b_1$

2. Iterate on h_2 :

If node exists in hs : Return node

If not, goto next node.

#En2:



#Idea: TC: $O(N \cdot M)$ SC: $O(1)$

Step1:

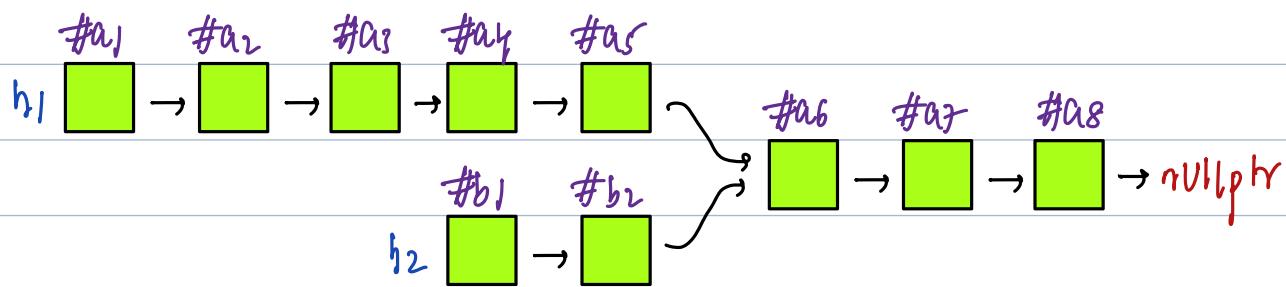
```
l1 = size(h1), l2 = size(h2);  
if [ l1 > l2 ] { #Extra nodes l1-l2  
}   #update h1 (l1-l2) times  
else if [ l2 > l1 ] { #Extra nodes l2-l1  
}   #update h2 (l2-l1) times
```

Step2:

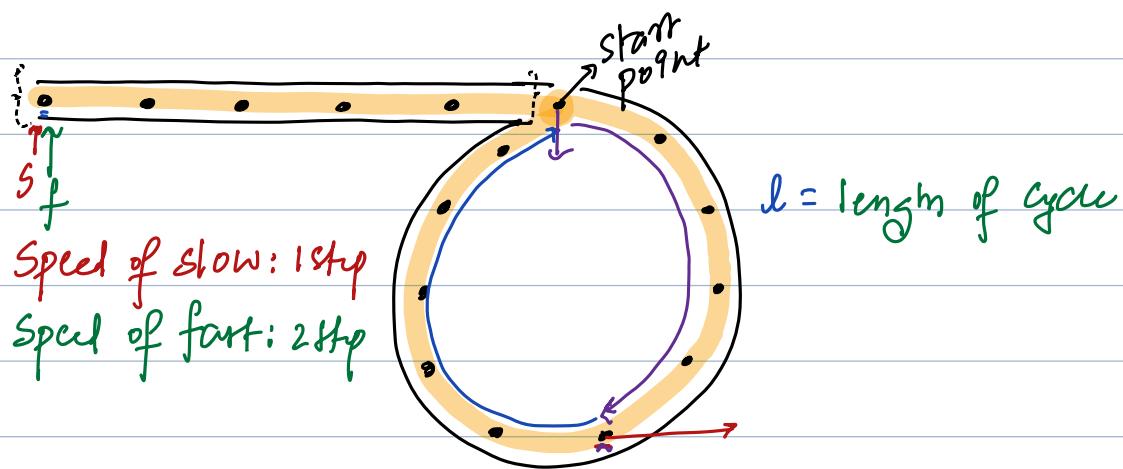
```
while (h1 != nullptr) {  
    if [ h1 == h2 ] { return h1; }  
    h1 = h1 → next;  
    h2 = h2 → next  
}
```

return nullptr

#Idcas:



#Cycle Detection: Proof.



#Note: Both s_g , f move same amount of time.

$ds: a + kl + d \quad // k: \text{No. of times } s \text{ points iterating in cycle}$

$df: a + ml + d \quad // m: \text{No. of times } f \text{ points iterating in cycle}$

$$df = 2^*ds;$$

$$a + ml + d = 2[a + kl + d]$$

$$\cancel{a} + \cancel{ml} + \cancel{d} = \cancel{2a} + \cancel{2kl} + \cancel{2d}$$

$$ml = a + 2kl + d$$

$$ml - 2kl - d = a$$

$$a = l[m - 2k] - dj$$

$$a = ly - dj$$