

Principles of Programming Languages - Homework 5

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1 Problem 1

(a)

$$\begin{aligned}
 \text{(i)} & \frac{\frac{\frac{\{x \rightarrow 3, y \rightarrow -2\} \vdash 3 \Downarrow 3}{\{x \rightarrow 3, y \rightarrow -2\} \vdash x \Downarrow 3} \text{EvalVal} \quad \frac{x \in \text{dom}(\{x \rightarrow 3, y \rightarrow -2\})}{\{x \rightarrow 3, y \rightarrow -2\} \vdash x \Downarrow 3} \text{EvalVar} \quad 9 = 3 * 3}{\{x \rightarrow 3, y \rightarrow -2\} \vdash x * 3 \Downarrow 9} \text{EvalTimes} \quad \frac{\{x \rightarrow 3, y \rightarrow -2\} \vdash 2 \Downarrow 2}{\{x \rightarrow 3, y \rightarrow -2\} \vdash 11 = 9 + 2} \text{EvalVal} \quad 11 = 9 + 2}{\{x \rightarrow 3, y \rightarrow -2\} \vdash 3 * x + 2 \Downarrow 11} \text{EvalPlus} \\
 \text{(ii)} & \frac{\frac{\frac{\{x \rightarrow 3, y \rightarrow -2\} \vdash 2 \Downarrow 2}{\{x \rightarrow 3, y \rightarrow -2\} \vdash y \Downarrow -2} \text{EvalVal} \quad \frac{y \in \text{dom}(\{x \rightarrow 3, y \rightarrow -2\})}{\{x \rightarrow 3, y \rightarrow -2\} \vdash y \Downarrow -2} \text{EvalVar} \quad 0 = 2 + -2}{\{x \rightarrow 3, y \rightarrow -2\} \vdash 2 + y \Downarrow 0} \text{EvalPlus} \quad \text{env}' = \text{env}[b \rightarrow 0] \quad \frac{\frac{b \in \text{dom}(\{x \rightarrow 3, y \rightarrow -2, b \rightarrow 0\})}{\text{env}' \vdash b \Downarrow 0} \text{EvalVar} \quad \text{toBool}(0) = \text{false}}{\{x \rightarrow 3, y \rightarrow -2, b \rightarrow 0\} \vdash y \Downarrow -2} \text{EvalIfElse} \quad \text{EvalIfThen}
 \end{aligned}$$

(b)

$$\text{(i)} \quad 3 + (1 \ \&\& \ 5) \xrightarrow{a} \underline{3 + 5} \xrightarrow{b} 8$$

a : SearchBop2, DoAndTrue

b : DoPlus

$$\begin{aligned}
 \text{(ii)} & \text{const } x = \underline{2 + 1}; x * 0 ? x : x + x \xrightarrow{a} \text{const } x = 3; \underline{x} * 0 ? x : x + x \xrightarrow{b} \text{const } x = 3; \underline{3 * 0} ? x : x + x \xrightarrow{c} \text{const } x = 3; \underline{0 ? x : x + x} \xrightarrow{d} \text{const } x = 3; \underline{x} + x \\
 & \xrightarrow{e} \text{const } x = 3; 3 + \underline{x} \xrightarrow{f} \text{const } x = 3; 3 + 3 \xrightarrow{g} \text{const } x = 3; \underline{3 + 3} \xrightarrow{h} 6
 \end{aligned}$$

a : SearchConstDecl1, DoPlus

b : SearchConstDecl2, DoVar

c : SearchConstDecl2, DoTimes

d : SearchConstDecl2, DoIfFalse

e : SearchConstDecl2, DoVar

f : SearchConstDecl2, DoVar

g : SearchConstDecl1, DoPlus

h : DoConstDecl

(c)

$$\text{SearchPlus1} \quad \frac{\text{env} \vdash e_2 \rightarrow e'_2}{\text{env} \vdash e_1 + e_2 \rightarrow e'_1 + e'_2}$$

$$\text{SearchPlus2} \quad \frac{\text{env} \vdash e_1 \rightarrow e'_1}{\text{env} \vdash e_1 + v_2 \rightarrow e'_1 + v_2}$$

$$\text{DoPlus} \quad \frac{v = \text{toNum}(v_1) + \text{toNum}(v_2)}{\text{env} \vdash v_1 + v_2 = v}$$

Order for right-to-left evaluation: SearchPlus1, SearchPlus2, DoPlus

(d)

Big-step SOS $\frac{env \vdash e_1 \rightarrow v_1 \quad env \vdash e_2 \rightarrow v_2 \quad v = v_2}{env \vdash e_1, e_2 \Downarrow v}$

Small-step SOS

SearchSeq1 $\frac{env \vdash e_1 \rightarrow e'_1}{env \vdash e_1, e_2 \rightarrow e'_1, e_2}$

SearchSeq2 $\frac{env \vdash e'_1 \rightarrow e'_2}{env \vdash v_1, e_2 \rightarrow v_1, e'_2}$

DoReturnValue $\frac{}{env \vdash v_1, v_2 = v}$

Order for right-to-left evaluation: SearchSeq1, SearchSeq2, DoReturnValue

(e)

(i) This program evaluates to 5. During evaluation, the EvalVar is applied three times as follows:

The first application is the using occurrence of y in the definition of the function f on line 3. In this case, y was bound to the variable 3 in the call to f on line 4.

The second application is the occurrence of x in the definition of the function g on line 2. x is bound to the value 2 in the constant declaration of 2 on line 1.

The third application is for the occurrence of y in the definition of the function g on line 2. This occurrence of y was bound to the value 3 in the call to g on line 3.

(ii) This program evaluates to 6. During evaluation, the EvalVar rule is applied four times as follows:

The first application is the using occurrence of y in the definition of function f on line 3. In this case, y was bound to the value 3 in the call to f on line 4.

The second application is the using occurrence of x in the definition of the function g on line 2. This occurrence was bound to the value 3 in the call to g on line 3.

The third application is the using occurrence of y in the definition of the function g on line 2. This occurrence was bound to the value that will be passed into the function by the function call in line 3. The function call to g returns a function, and that function is called on line 3 using the value 3. So the occurrence was bound to the value 3 in the call to the anonymous function on line 3.

The fourth application is the using occurrence of y in the definition of function f on line 3. This occurrence was bound to the value 3 in the call to the anonymous function that is returned by calling g on line 3.

2 Problem 2

(a)

$$e_1 = (3 * y) + 4$$

(b)

$$e_1 = (x * y) + 4$$

(c)

$$e_2 = \text{const } y = y; 3 + y$$

(d)

$$e_2 = \textit{const } y = 3; x + y$$

(e)

$$e_3 = \textit{const } x = (\textit{function}(z)(x(z))); x(y(2))$$

(f)

$$e_3 = \textit{const } x = (\textit{function}(z)(y(x(z)))); x(y)$$