LalCheetah ICPC Team Notebook (2016-17)

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Combinatorial optimization

1.1 Sparse max-flow(aka Modified Ford Fulkerson)

```
// Adjacency list implementation of Dinic's blocking flow
     algorithm.
// This is very fast in practice, and only loses to push-
     relabel flow.
// Running time:
      O(|V|^2 |E|)
      - graph, constructed using AddEdge()
      - source and sink
     - maximum flow value
      - To obtain actual flow values, look at edges with
     capacity > 0
        (zero capacity edges are residual edges).
#include < cstdio >
#include<vector>
#include<queue>
using namespace std;
```

```
typedef long long LL;
struct Edge {
 int u, v;
 LL cap, flow;
 Edge(int u, int v, LL cap): u(u), v(v), cap(cap), flow(0)
struct Dinic {
 int N:
 vector<Edge> E;
 vector<vector<int>> q;
 vector<int> d, pt;
 Dinic(int N): N(N), E(0), g(N), d(N), pt(N) {}
  void AddEdge(int u, int v, LL cap) {
   if (u != v) {
     E.emplace_back(Edge(u, v, cap));
      g[u].emplace_back(E.size() - 1);
      E.emplace_back(Edge(v, u, 0));
     g[v].emplace_back(E.size() - 1);
 bool BFS(int S, int T) {
   queue<int> q({S});
    fill(d.begin(), d.end(), N + 1);
    d[S] = 0;
    while(!q.empty()) {
     int u = q.front(); q.pop();
     if (u == T) break;
     for (int k: g[u]) {
   Edge &e = E[k];
        if (e.flow < e.cap && d[e.v] > d[e.u] + 1) {
         d[e.v] = d[e.u] + 1;
          q.emplace(e.v);
    return d[T] != N + 1;
 LL DFS (int u, int T, LL flow = -1) {
   if (u == T || flow == 0) return flow;
    for (int &i = pt[u]; i < g[u].size(); ++i) {</pre>
     Edge &e = E[g[u][i]];
     Edge &oe = E[g[u][i]^1];
     if (d[e.v] == d[e.u] + 1) {
        LL amt = e.cap - e.flow;
        if (flow !=-1 && amt > flow) amt = flow;
        if (LL pushed = DFS(e.v, T, amt)) {
          e.flow += pushed;
          oe flow -= pushed;
          return pushed;
    return 0;
 LL MaxFlow(int S, int T) {
    LL total = 0;
    while (BFS(S, T)) {
      fill(pt.begin(), pt.end(), 0);
      while (LL flow = DFS(S, T))
       total += flow;
    return total;
};
// The following code solves SPOJ problem #4110: Fast
     Maximum Flow (FASTFLOW)
int main()
 int N, E;
scanf("%d%d", &N, &E);
 Dinic dinic(N);
  for (int i = 0; i < E; i++)
   int u, v;
   LL cap; scanf("%d%d%lld", &u, &v, &cap);
    dinic.AddEdge(u - 1, v - 1, cap);
```

```
dinic.AddEdge(v - 1, u - 1, cap);
 printf("%lld\n", dinic.MaxFlow(0, N - 1));
 return 0:
// END CUT
```

1.2 Global min-cut

```
// Adjacency matrix implementation of Stoer-Wagner min cut
     algorithm.
// Running time:
      0(|V|^3)
       - graph, constructed using AddEdge()
// OUTPUT:
       - (min cut value, nodes in half of min cut)
#include <cmath>
#include <vector>
#include <iostream>
using namespace std;
typedef vector<int> VI;
typedef vector<VI> VVI;
const int INF = 1000000000;
pair<int, VI> GetMinCut(VVI &weights) {
 int N = weights.size();
  VI used(N), cut, best_cut;
 int best_weight = -1;
  for (int phase = N-1; phase >= 0; phase--) {
   VI w = weights[0];
    VI added = used;
    int prev, last = 0;
    for (int i = 0; i < phase; i++) {</pre>
      prev = last;
last = -1;
      for (int j = 1; j < N; j++)</pre>
        if (!added[j] && (last == -1 || w[j] > w[last]))
      if (i == phase-1) {
        for (int j = 0; j < N; j++) weights[prev][j] +=</pre>
             weights[last][j];
        for (int j = 0; j < N; j++) weights[j][prev] =
    weights[prev][j];</pre>
        used[last] = true;
        cut.push back(last);
        if (best_weight == -1 || w[last] < best_weight) {</pre>
          best_weight = w[last];
        for (int j = 0; j < N; j++)
        w[j] += weights[last][j];
added[last] = true;
  return make_pair(best_weight, best_cut);
// The following code solves UVA problem #10989: Bomb,
    Divide and Conquer
int main() {
 int N;
 cin >> N;
  for (int i = 0; i < N; i++) {
   int n, m;
    cin >> n >> m;
    VVI weights(n, VI(n));
    for (int j = 0; j < m; j++) {
     int a, b, c;
      cin >> a >> b >> c;
      weights[a-1][b-1] = weights[b-1][a-1] = c_i
```

```
pair<int, VI> res = GetMinCut(weights);
  cout << "Case #" << i+1 << ": " << res.first << endl;
}
}// END CUT</pre>
```

2 Geometry

2.1 Convex hull

```
// Compute the 2D convex hull of a set of points using the
     monotone chain
// algorithm. Eliminate redundant points from the hull if
     REMOVE_REDUNDANT is
// #defined.
// Running time: O(n log n)
     INPUT: a vector of input points, unordered.
     OUTPUT: a vector of points in the convex hull,
     counterclockwise, starting
              with bottommost/leftmost point
#include <cstdio>
#include <cassert>
#include <vector>
#include <algorithm>
#include <cmath>
  / BEGIN CUT
#include <map>
// END CUT
using namespace std;
#define REMOVE REDUNDANT
typedef double T:
const T EPS = 1e-7:
struct PT {
  T x, y;
PT() {}
  PT(T x, T y) : x(x), y(y) {}
 bool operator<(const PT &rhs) const { return make_pair(y,
       x) < make_pair(rhs.y,rhs.x); }</pre>
 bool operator==(const PT &rhs) const { return make_pair(y
       ,x) == make_pair(rhs.y,rhs.x); }
T cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
T area2(PT a, PT b, PT c) { return cross(a,b) + cross(b,c)
     + cross(c,a); }
#ifdef REMOVE REDUNDANT
bool between (const PT &a, const PT &b, const PT &c) {
  return (fabs(area2(a,b,c)) < EPS && (a.x-b.x)*(c.x-b.x)
       <= 0 && (a.y-b.y) * (c.y-b.y) <= 0);
#endif
void ConvexHull(vector<PT> &pts) {
  sort(pts.begin(), pts.end());
  pts.erase(unique(pts.begin(), pts.end()), pts.end());
  vector<PT> up, dn;
  for (int i = 0; i < pts.size(); i++) {</pre>
    while (up.size() > 1 && area2(up[up.size()-2], up.back
         (), pts[i]) >= 0) up.pop_back();
    while (dn.size() > 1 && area2(dn[dn.size()-2], dn.back
         (), pts[i]) <= 0) dn.pop_back();
    up.push_back(pts[i]);
    dn.push_back(pts[i]);
  pts = dn;
  for (int i = (int) up.size() - 2; i >= 1; i--) pts.
       push_back(up[i]);
#ifdef REMOVE_REDUNDANT
 if (pts.size() <= 2) return;</pre>
  dn.clear();
dn.push_back(pts[0]);
  dn.push_back(pts[1]);
  for (int i = 2; i < pts.size(); i++) {</pre>
    if (between(dn[dn.size()-2], dn[dn.size()-1], pts[i]))
         dn.pop_back();
```

```
dn.push back(pts[i]);
  if (dn.size() >= 3 && between(dn.back(), dn[0], dn[1])) {
    dn[0] = dn.back();
    dn.pop_back();
  pts = dn:
#endif
// BEGIN CUT
// The following code solves SPOJ problem #26: Build the
     Fence (BSHEEP)
int main() {
 int t;
scanf("%d", &t);
  for (int caseno = 0; caseno < t; caseno++) {</pre>
    int n:
    scanf("%d", &n);
    vector<PT> v(n);
    for (int i = 0; i < n; i++) scanf("%lf%lf", &v[i].x, &v</pre>
    [i].y);
vector<PT> h(v);
    map<PT,int> index;
    for (int i = n-1; i \ge 0; i--) index[v[i]] = i+1;
    ConvexHull(h);
    double len = 0;
    for (int i = 0; i < h.size(); i++) {</pre>
      double dx = h[i].x - h[(i+1)%h.size()].x;
      double dy = h[i].y - h[(i+1)%h.size()].y;
      len += sqrt (dx*dx+dy*dy);
    if (caseno > 0) printf("\n");
    printf("%.2f\n", len);
    for (int i = 0; i < h.size(); i++) {</pre>
      if (i > 0) printf(" ");
      printf("%d", index[h[i]]);
    printf("\n");
// END CUT
```

2.2 Miscellaneous geometry

```
// C++ routines for computational geometry.
#include <iostream>
#include <vector>
#include <cmath>
#include <cassert>
using namespace std;
double INF = 1e100;
double EPS = 1e-12;
struct PT {
  double x, y;
  PT(double x, double y) : x(x), y(y) {}
PT(const PT &p) : x(p.x), y(p.y) {}
PT operator + (const PT &p) const { return PT(x+p.x, y+p)
  PT operator - (const PT &p) const { return PT(x-p.x, y-p
        .v): }
  PT operator * (double c)
                                    const { return PT(x*c, y*c
          ); }
  PT operator / (double c)
                                    const { return PT(x/c, y/c
double dot(PT p, PT q)
                               { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q) { return dot(p-q,p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream & operator << (ostream &os, const PT &p) {
 os << "(" << p.x << "," << p.y << ")";
// rotate a point CCW or CW around the origin
PT RotateCCW90 (PT p)
                            return PT(-p.y,p.x);
PT RotateCW90 (PT p)
                           { return PT(p.y,-p.x); }
PT RotateCCW(PT p, double t) {
```

```
return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
  return a + (b-a) *dot(c-a, b-a) /dot(b-a, b-a);
// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
  double r = dot(b-a, b-a);
 if (fabs(r) < EPS) return a;</pre>
  r = dot(c-a, b-a)/r;
  if (r < 0) return a;</pre>
 if (r > 1) return b;
  return a + (b-a) *r;
// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
 return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
// compute distance between point (x,y,z) and plane ax+by+
double d)
  return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
// determine if lines from a to b and c to d are parallel
     or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
 return fabs(cross(b-a, c-d)) < EPS;
bool LinesCollinear (PT a, PT b, PT c, PT d) {
  return LinesParallel(a, b, c, d)
      && fabs(cross(a-b, a-c)) < EPS
      && fabs(cross(c-d, c-a)) < EPS;
// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect (PT a, PT b, PT c, PT d) {
 if (LinesCollinear(a, b, c, d)) {
    if (dist2(a, c) < EPS || dist2(a, d) < EPS ||</pre>
      dist2(b, c) < EPS || dist2(b, d) < EPS) return true;</pre>
    if (dot(c-a, c-b) > 0 \&\& dot(d-a, d-b) > 0 \&\& dot(c-b, d-b)
         d-b) > 0)
      return false;
    return true;
  if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return false;
  if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return false;
  return true;
// compute intersection of line passing through a and b
// with line passing through c and d, assuming that unique
// intersection exists; for segment intersection, check if
// seaments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
  b=b-a; d=c-d; c=c-a;
  assert(dot(b, b) > EPS && dot(d, d) > EPS);
  return a + b*cross(c, d)/cross(b, d);
// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
 b = (a+b)/2;
  c = (a+c)/2;
  return ComputeLineIntersection(b, b+RotateCW90(a-b), c, c
       +RotateCW90(a-c));
// determine if point is in a possibly non-convex polygon (
     by William
// Randolph Franklin); returns 1 for strictly interior
     points, 0 for
// strictly exterior points, and 0 or 1 for the remaining
    points.
// Note that it is possible to convert this into an *exact*
      test usina
// integer arithmetic by taking care of the division
```

```
appropriately
// (making sure to deal with signs properly) and then by
     writing exact
// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++){</pre>
    int j = (i+1)%p.size();
    if ((p[i].y <= q.y && q.y < p[j].y ||</pre>
      \begin{array}{l} p[j].y <= q.y & \& & q.y < p[i].y) & \& \\ q.x < p[i].x + & (p[j].x - p[i].x) & * & (q.y - p[i].y) / \end{array} \label{eq:pp}
            p[j].y - p[i].y))
  return c;
// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
  for (int i = 0; i < p.size(); i++)</pre>
    if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()],
      q), q) < EPS) return true;
    return false:
// compute intersection of line through points a and b with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c, double
     r) {
  vector<PT> ret;
  b = b-a;
  a = a-c;

double A = dot(b, b);

double B = dot(a, b);
  double C = dot(a, a) - r * r;
  double D = B*B - A*C;
  if (D < -EPS) return ret;</pre>
  ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
  if (D > EPS)
    ret.push_back(c+a+b*(-B-sqrt(D))/A);
  return ret;
// compute intersection of circle centered at a with radius
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r,
     double R) {
  double d = sqrt(dist2(a, b));
  if (d > r+R \mid | d+min(r, R) < max(r, R)) return ret;
double x = (d*d-R*R+r*r)/(2*d);
  double y = sqrt(r*r-x*x);
  PT v = (b-a)/d;
  ret push_back(a+v*x + RotateCCW90(v)*y);
  if (y > 0)
    ret.push_back(a+v*x - RotateCCW90(v)*y);
  return ret;
// This code computes the area or centroid of a (possibly
     nonconvex)
// polygon, assuming that the coordinates are listed in a
     clockwise or
// counterclockwise fashion. Note that the centroid is
     often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
  double area = 0;
  for(int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    area += p[i].x*p[j].y - p[j].x*p[i].y;
  return area / 2.0;
double ComputeArea(const vector<PT> &p) {
  return fabs(ComputeSignedArea(p));
PT ComputeCentroid(const vector<PT> &p) {
  PT c(0,0);
  double scale = 6.0 * ComputeSignedArea(p);
  for (int i = 0; i < p.size(); i++) {</pre>
    int j = (i+1) % p.size();
    c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
  return c / scale;
```

```
// tests whether or not a given polygon (in CW or CCW order
    ) is simple
bool IsSimple(const vector<PT> &p) {
 for (int i = 0; i < p.size(); i++) {</pre>
   for (int k = i+1; k < p.size(); k++) {</pre>
     int j = (i+1) % p.size();
     int \tilde{1} = (k+1) % p.size();
     if (i == 1 \mid | j == k) continue;
     if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
       return false:
 return true;
int main() {
  // expected: (-5,2)
 cerr << RotateCCW90(PT(2,5)) << endl;</pre>
  // expected: (5,-2)
  cerr << RotateCW90(PT(2,5)) << endl;</pre>
  // expected: (-5,2)
 cerr << RotateCCW(PT(2,5),M_PI/2) << endl;</pre>
  // expected: (5,2)
 cerr << ProjectPointLine(PT(-5,-2), PT(10,4), PT(3,7)) <</pre>
       endl;
  // expected: (5,2) (7.5,3) (2.5,1)
 << ProjectPointSegment(PT(-5,-2), PT(2.5,1), PT(3,7)
           ) << endl;
  // expected: 6.78903
 cerr << DistancePointPlane(4,-4,3,2,-2,5,-8) << endl;</pre>
  // expected: 1 0 1
 << LinesParallel(PT(1,1), PT(3,5), PT(2,0), PT(4,5))
      << LinesParallel(PT(1,1), PT(3,5), PT(5,9), PT(7,13)
  // expected: 0 0 1
 << LinesCollinear(PT(1,1), PT(3,5), PT(2,0), PT(4,5)
           ) << " "
      << LinesCollinear(PT(1,1), PT(3,5), PT(5,9), PT
           (7,13)) << endl;
  // expected: 1 1 1 0

.<</pre>
SegmentsIntersect(PT(0,0), PT(2,4), PT(4,3), PT
(0,5)) << " "
      << SegmentsIntersect (PT(0,0), PT(2,4), PT(2,-1), PT (-2,1)) << " "
       << SegmentsIntersect(PT(0,0), PT(2,4), PT(5,5), PT
           (1,7)) << endl;
  cerr << ComputeLineIntersection(PT(0,0), PT(2,4), PT(3,1)</pre>
       , PT(-1,3)) << endl;
  // expected: (1,1)
  cerr << ComputeCircleCenter(PT(-3,4), PT(6,1), PT(4,5))</pre>
      << endl;
  vector<PT> v:
 v.push_back(PT(0,0));
 v.push_back(PT(5,0));
v.push_back(PT(5,5));
 v.push back (PT(0,5));
  // expected: 1 1 1 0 0
 cerr << PointInPolygon(v, PT(2,2)) << " "</pre>
      << PointInPolygon(v, PT(2,0)) << " "
      << PointInPolygon(v, PT(0,2)) << " "
      << PointInPolygon(v, PT(5,2)) << " "
      << PointInPolygon(v, PT(2,5)) << endl;
 // expected: 0 1 1 1 1
```

```
cerr << PointOnPolygon(v, PT(2,2)) << " "</pre>
     << PointOnPolygon(v, PT(2,0)) << " "
     << PointOnPolygon(v, PT(0,2)) << " "
     << PointOnPolygon(v, PT(5,2)) << " "
     << PointOnPolygon(v, PT(2,5)) << endl;
              (5,4) (4,5)
             blank line
             (4,5) (5,4)
             blank line
             (4,5) (5,4)
vector<PT> u = CircleLineIntersection(PT(0,6), PT(2,6),
    PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " ";</pre>
     cerr << endl:
u = CircleLineIntersection(PT(0,9), PT(9,0), PT(1,1), 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << "</pre>
     cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(10,10), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " ";</pre>
    cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(8,8), 5, 5);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " ";</pre>
     cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 10,
     sqrt (2.0) /2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " ";</pre>
     cerr << endl;
u = CircleCircleIntersection(PT(1,1), PT(4.5,4.5), 5,
     sgrt (2.0)/2.0);
for (int i = 0; i < u.size(); i++) cerr << u[i] << " ";</pre>
     cerr << endl:
// area should be 5.0
 // centroid should be (1.1666666, 1.166666)
PT pa[] = { PT(0,0), PT(5,0), PT(1,1), PT(0,5) };
vector<PT> p(pa, pa+4);
PT c = ComputeCentroid(p);
cerr << "Area: " << ComputeArea(p) << endl;</pre>
cerr << "Centroid: " << c << endl;</pre>
return 0:
```

3 Numerical algorithms

3.1 Fast Fourier transform

```
#include <cassert>
#include <cstdio>
#include <cmath>
struct cpx
 cpx (double aa):a(aa),b(0){}
  cpx(double aa, double bb):a(aa),b(bb){}
 double a;
 double b;
 double modsq(void) const
    return a * a + b * b;
  cpx bar (void) const
    return cpx(a, -b);
cpx operator + (cpx a, cpx b)
 return cpx(a.a + b.a. a.b + b.b);
cpx operator *(cpx a, cpx b)
 return cpx(a.a * b.a - a.b * b.b, a.a * b.b + a.b * b.a);
cpx operator / (cpx a, cpx b)
  cpx r = a * b.bar();
 return cpx(r.a / b.modsq(), r.b / b.modsq());
```

```
cpx EXP (double theta)
 return cpx(cos(theta), sin(theta));
const double two_pi = 4 * acos(0);
// in:
           input array
// out:
           output array
// step:
           {SET TO 1} (used internally)
// size:
           length of the input/output {MUST BE A POWER OF
// dir.
           either plus or minus one (direction of the FFT)
// RESULT: out[k] = \sum_{j=0}^{size - 1} in[j] * exp(dir *
      2pi * i * j * k / size)
void FFT(cpx *in, cpx *out, int step, int size, int dir)
 if(size < 1) return;</pre>
  if(size == 1)
    out[0] = in[0];
    return;
  FFT(in, out, step * 2, size / 2, dir);
  FFT(in + step, out + size / 2, step * 2, size / 2, dir);
  for (int i = 0; i < size / 2; i++)
    cpx odd = out[i + size / 2];
    out[i] = even + EXP(dir * two_pi * i / size) * odd;
    out[i + size / 2] = even + EXP(dir * two pi * (i + size)
          / 2) / size) * odd;
// f[0...N-1] and g[0..N-1] are numbers
// Want to compute the convolution h, defined by
// h[n] = sum of f[k]g[n-k] (k = 0, ..., N-1).
// Here, the index is cyclic; f[-1] = f[N-1], f[-2] = f[N
// Let F[0...N-1] be FFT(f), and similarly, define G and H.
// The convolution theorem says H[n] = F[n]G[n] (element-
     wise product).
// To compute h[] in O(N log N) time, do the following:
   1. Compute F and G (pass dir = 1 as the argument).
    2. Get H by element-wise multiplying F and G.
    3. Get h by taking the inverse FFT (use dir = -1 as
     the argument)
       and *dividing by N*. DO NOT FORGET THIS SCALING
     FACTOR.
int main (void)
 printf("If rows come in identical pairs, then everything
       works.\n");
 cpx a[8] = \{0, 1, cpx(1,3), cpx(0,5), 1, 0, 2, 0\}; cpx b[8] = \{1, cpx(0,-2), cpx(0,1), 3, -1, -3, 1, -2\};
  cpx B[8];
 FFT(a, A, 1, 8, 1);
FFT(b, B, 1, 8, 1);
  for (int i = 0; i < 8; i++)
    printf("%7.21f%7.21f", A[i].a, A[i].b);
  printf("\n");
  for (int i = 0; i < 8; i++)
    cpx Ai(0,0);
    for (int j = 0; j < 8; j++)
     Ai = Ai + a[j] * EXP(j * i * two_pi / 8);
   printf("%7.21f%7.21f", Ai.a, Ai.b);
  printf("\n");
  cpx AB[8];
  for (int i = 0; i < 8; i++)
   AB[i] = A[i] * B[i];
  cpx aconvb[8];
  FFT (AB, aconvb, 1, 8, -1);
  for (int i = 0; i < 8; i++)
   aconvb[i] = aconvb[i] / 8;
```

```
for(int i = 0; i < 8; i++)
{
    printf("%7.21f%7.21f", aconvb[i].a, aconvb[i].b);
}
printf("\n");
for(int i = 0; i < 8; i++)
{
    cpx aconvbi(0,0);
    for(int j = 0; j < 8; j++)
    {
        aconvbi = aconvbi + a[j] * b[(8 + i - j) % 8];
    }
    printf("%7.21f%7.21f", aconvbi.a, aconvbi.b);
}
printf("\n");
return 0;</pre>
```

4 Graph algorithms

4.1 Fast Dijkstra's algorithm

```
// Implementation of Dijkstra's algorithm using adjacency
// and priority queue for efficiency.
// Running time: O(|E| log |V|)
#include <queue>
#include <cstdio>
using namespace std;
const int INF = 2000000000;
typedef pair<int, int> PII;
int main() {
        scanf("%d%d%d", &N, &s, &t);
        vector<vector<PII> > edges(N);
        for (int i = 0; i < N; i++) {
                int M;
scanf("%d", &M);
                for (int j = 0; j < M; j++) {
                        int vertex, dist;
                        scanf("%d%d", &vertex, &dist);
                        edges[i].push_back(make_pair(dist,
                             vertex)); // note order of
                             arguments here
        // use priority queue in which top element has the
             "smallest" priority
        priority_queue<PII, vector<PII>, greater<PII> > Q;
        vector<int> dist(N, INF), dad(N, -1);
       Q.push(make_pair(0, s));
        dist[s] = 0;
       while (!Q.empty()) {
                PII p = Q.top();
                Q.pop();
                int here = p.second;
                if (here == t) break;
                if (dist[here] != p.first) continue;
                for (vector<PII>::iterator it = edges[here
                     ].begin(); it != edges[here].end(); it
                        if (dist[here] + it->first < dist[</pre>
                             it->second]) {
                                dist[it->second] = dist[
                                     here] + it->first;
                                dad[it->second] = here;
                                Q.push (make_pair (dist[it->
                                     second], it->second));
       printf("%d\n", dist[t]);
```

4.2 Strongly connected components

```
#include<memory.h>
struct edge{int e, nxt;};
int V, E;
edge e[MAXE], er[MAXE];
int sp[MAXV], spr[MAXV];
int group_cnt, group_num[MAXV];
bool v[MAXV];
int stk[MAXV];
void fill_forward(int x)
 v[x]=true;
 for(i=sp[x];i;i=e[i].nxt) if(!v[e[i].e]) fill_forward(e[i
 stk[++stk[0]]=x;
void fill backward(int x)
 v[x]=false;
 group_num[x]=group_cnt;
 for(i=spr[x];i;i=er[i].nxt) if(v[er[i].e]) fill_backward(
      er[i].e);
void add_edge(int v1, int v2) //add edge v1->v2
  e [++E].e=v2; e [E].nxt=sp [v1]; sp [v1]=E;
 er[ E].e=v1; er[E].nxt=spr[v2]; spr[v2]=E;
void SCC()
 int i:
 stk[0]=0;
 memset(v, false, sizeof(v));
 for(i=1;i<=V;i++) if(!v[i]) fill_forward(i);</pre>
 group_cnt=0;
  for (i=stk[0];i>=1;i--) if (v[stk[i]]) {group_cnt++;
       fill_backward(stk[i]);}
```

4.3 Bellman Ford's algorithm

```
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool BellmanFord (const VVT &w, VT &dist, VI &prev, int
  int n = w.size();
 prev = VI(n, -1);
 dist = VT(n, 1000000000);
 dist[start] = 0;
 for (int k = 0; k < n; k++) {
    for (int i = 0; i < n; i++) {
     for (int j = 0; j < n; j++) {
       if (dist[j] > dist[i] + w[i][j]){
         if (k == n-1) return false;
          dist[j] = dist[i] + w[i][j];
         prev[j] = i;
 return true;
```

4.4 Minimum Spanning Tree: Kruskal

```
Uses Kruskal's Algorithm to calculate the weight of the
    minimum spanning
forest (union of minimum spanning trees of each connected
    component) of
a possibly disjoint graph, given in the form of a matrix of
      edge weights
(-1 \text{ if no edge exists}). Returns the weight of the minimum
forest (also calculates the actual edges - stored in T).
    Note: uses a
disjoint-set data structure with amortized (effectively)
     constant time per
union/find. Runs in O(E*log(E)) time.
#include "template.h"
typedef int T:
struct edge {
 int u, v;
  T d;
struct edgeCmp{
 int operator()(const edge& a, const edge& b) { return a.d
        > b.d; }
int find(vector <int>& C, int x) { return (C[x] == x)?x: C[
     x]=find(C, C[x]); }
T Kruskal(vii Alist[], int n) {
  T weight = 0;
  vector <int> C(n), R(n);
 for(int i=0; i<n; i++) { C[i] = i; R[i] = 0; }</pre>
  vector <edge> T:
 priority_queue <edge, vector <edge>, edgeCmp> E;
  rep(i, n)
    rep(j, Alist[i].size()) {
      e.u = i, e.v = Alist[i][j].F, e.d = Alist[i][j].S;
      E.push(e):
```

```
while (T.size() < n-1 \&\& !E.empty()) {
    edge cur = E.top(); E.pop();
    int uc = find(C, cur.u), vc = find(C, cur.v);
    if(uc != vc) {
      T.push_back(cur); weight += cur.d;
      if(R[uc] > R[vc])
C[vc] = uc;
      else if(R[vc] > R[uc])
        C[uc] = vc;
      else
        C[vc] = uc; R[uc]++;
 return weight;
int main() {
 int n;
 cin >> n;
 vii Alist[Lim];
 cout << Kruskal(Alist, n) << endl;</pre>
```

5 Data structures

5.1 Suffix array

```
// computing the length of the longest common prefix of any
// suffixes in O(log L) time.
// INPUT: string s
// OUTPUT: array suffix[] such that suffix[i] = index (
     from 0 to L-1)
           of substring s[i...L-1] in the list of sorted
     suffixes.
           That is, if we take the inverse of the
     permutation suffix[],
           we get the actual suffix array.
#include <vector>
#include <iostream>
#include <string>
using namespace std;
struct SuffixArray {
 const int L:
 string s;
 vector<vector<int> > P;
 vector<pair<int,int>,int> > M;
 SuffixArray(const string &s) : L(s.length()), s(s), P(1,
       vector<int>(L, 0)), M(L) {
    for (int i = 0; i < L; i++) P[0][i] = int(s[i]);
    for (int skip = 1, level = 1; skip < L; skip \star= 2,
         level++) {
     P.push_back(vector<int>(L, 0));
     for (int i = 0; i < L; i++)
       M[i] = make_pair(make_pair(P[level-1][i], i + skip
            < L ? P[level-1][i + skip] : -1000), i);
     sort(M.begin(), M.end());
     for (int i = 0; i < L; i++)
       P[level][M[i].second] = (i > 0 && M[i].first == M[i]
             -1].first) ? P[level][M[i-1].second] : i;
 vector<int> GetSuffixArray() { return P.back(); }
  // returns the length of the longest common prefix of s[i
       \dots L-1] and s[j\dots L-1]
  int LongestCommonPrefix(int i, int j) {
   int len = 0;
   if (i == j) return L - i;
```

// Suffix array construction in O(L log^2 L) time. Routine

```
for (int k = P.size() - 1; k \ge 0 && i < L && j < L; k
      if (P[k][i] == P[k][j]) {
       i += 1 << k;
        i += 1 << k;
        len += 1 << k;
   return len;
// The following code solves UVA problem 11512: GATTACA.
#define TESTING
#ifdef TESTING
int main() {
 int T;
 cin >> T;
 for (int caseno = 0; caseno < T; caseno++) {</pre>
   string s;
   cin >> s;
   SuffixArray array(s);
    vector<int> v = array.GetSuffixArray();
    int bestlen = -1, bestpos = -1, bestcount = 0;
    for (int i = 0; i < s.length(); i++) {</pre>
     int len = 0, count = 0;
      for (int j = i+1; j < s.length(); j++) {</pre>
       int 1 = array.LongestCommonPrefix(i, j);
        if (1 >= len) {
          if (1 > len) count = 2; else count++;
          len = 1;
      if (len > bestlen || len == bestlen && s.substr(
          bestpos, bestlen) > s.substr(i, len)) {
        bestlen = len:
        bestcount = count;
       bestpos = i;
    if (bestlen == 0) {
     cout << "No repetitions found!" << endl;</pre>
      cout << s.substr(bestpos, bestlen) << " " <<
           bestcount << endl;
#else
// END CUT
int main() {
 // bobocel is the O'th suffix
 // obocel is the 5'th suffix
     bocel is the 1'st suffix
       ocel is the 6'th suffix
        cel is the 2'nd suffix
         el is the 3'rd suffix
          l is the 4'th suffix
 SuffixArray suffix("bobocel");
  vector<int> v = suffix.GetSuffixArray();
  // Expected output: 0 5 1 6 2 3 4
 for (int i = 0; i < v.size(); i++) cout << v[i] << " ";</pre>
 cout << endl;
 cout << suffix.LongestCommonPrefix(0, 2) << endl;</pre>
// BEGIN CUT
#endif
// END CUT
```

5.2 Binary Indexed Tree

```
#include <iostream>
using namespace std;
#define LOGSZ 17
int tree[(1<<LOGSZ)+1];
int N = (1<<LOGSZ);</pre>
```

```
// add v to value at x
void set(int x, int v) {
  while (x \le N)
   tree[x] += v;
    x += (x \& -x);
// get cumulative sum up to and including x
int get(int x) {
 int res = 0;
 while(x) {
   res += tree[x];
   x = (x \& -x);
  return res;
// get largest value with cumulative sum less than or equal
// for smallest, pass x-1 and add 1 to result
int getind(int x) {
 int idx = 0, mask = N;
  while(mask && idx < N)</pre>
   int t = idx + mask;
    if(x >= tree[t]) {
     idx = t;
     x -= tree[t];
   mask >>= 1;
 return idx:
```

5.3 Union-find set(aka DSU)

5.4 Lowest common ancestor

```
const int max_nodes, log_max_nodes;
int num_nodes, log_num_nodes, root;
vector<int> children[max nodes];
                                        // children[i]
     contains the children of node i
int A[max nodes][log max nodes+1];
                                       // A[i][j] is the
     2^j-th ancestor of node i, or -1 if that ancestor does
int L[max_nodes];
                                       // L[i] is the
     distance between node i and the root
// floor of the binary logarithm of n
int lb(unsigned int n)
    if(n==0)
       return -1;
    int p = 0;
    if (n >= 1<<16) { n >>= 16; p += 16; }
   if (n >= 1<< 8) { n >>= 8; p += 8; }
    if (n >= 1 << 4) { n >>= 4; p += 4; }
    if (n >= 1<< 2) { n >>= 2; p += 2; }
    if (n >= 1<< 1) {
                               p += 1; }
```

```
return p;
void DFS(int i, int 1)
    L[i] = 1;
    for(int j = 0; j < children[i].size(); j++)</pre>
        DFS(children[i][j], l+1);
int LCA(int p, int q)
    // ensure node p is at least as deep as node q
    if(L[p] < L[q])
        swap(p, q);
    // "binary search" for the ancestor of node p situated
         on the same level as q
    for(int i = log_num_nodes; i >= 0; i--)
        if(L[p] - (1 << i) >= L[q])
           p = A[p][i];
    if(p == q)
        return p;
    // "binary search" for the LCA
    for(int i = log_num_nodes; i >= 0; i--)
        if(A[p][i] != -1 && A[p][i] != A[q][i])
            p = A[p][i];
            q = A[q][i];
    return A[p][0];
int main(int argc,char* argv[])
    // read num_nodes, the total number of nodes
    log_num_nodes=1b(num_nodes);
    for(int i = 0; i < num_nodes; i++)</pre>
        int p;
        // read p, the parent of node i or -1 if node i is
        A[i][0] = p;
        if(p != -1)
           children[p].push_back(i);
        else
            root = i;
    // precompute A using dynamic programming
    for(int j = 1; j <= log_num_nodes; j++)</pre>
        for (int i = 0; i < num nodes; i++)
            if(A[i][j-1] != -1)
               A[i][j] = A[A[i][j-1]][j-1];
                A[i][j] = -1;
    // precompute L
    DFS(root, 0);
    return 0;
```

5.5 Segment tree for range minima query

```
} else if (a < b) {</pre>
        int mid=(a+b)>>1, left=Node<<1;</pre>
        int right=left|1;
        buildTree(left, a, mid);
        buildTree(right, mid+1, b);
        Tree[Node] = min(Tree[left], Tree[right]);
void updateTree(int Node, ll value, int a, int b, int index
    if (a > index || b < index) {</pre>
    } else if (a == b) {
        Tree[Node] = value;
        Arr[index] = value;
    } else if (a <= index && b >= index) {
        int mid=(a+b)>>1, left=Node<<1;</pre>
        int right=left|1;
        updateTree(left, index, value, a, mid);
        updateTree(right, index, value, mid+1, b);
        Tree[Node] = min(Tree[left], Tree[right]);
11 queryTree(int Node, int start, int end, int a, int b) {
    int mid=(a+b)>>1, left=Node<<1;</pre>
    int right=left|1;
    11 Ans = Inf;
    if (start <= a && b <= end) {</pre>
        return Tree[Node];
    } else {
        if(mid >= start)
            Ans = queryTree(left, start, end, a, mid);
        if(mid < end)</pre>
            Ans = min(Ans, queryTree(right, start, end, mid
        return Ans;
```

5.6 Lazy Propogation for Range update and Query

```
#include "template.h"
 A lazy tree implementation of Range Updation & Range
11 Arr[Lim], Tree[4*Lim], lazy[4*Lim];
void build_tree(int Node, int a, int b) {
  // Do not forget to clear lazy Array before calling build
 if(a == b) {
  Tree[Node] = Arr[a];
  } else if (a < b) {</pre>
    int mid = (a+b)>>1, left=Node<<1, right=left|1;</pre>
    build_tree(left, a, mid); build_tree(right, mid+1, b);
    Tree[Node] = Tree[left]+Tree[right];
void Propogate(int Node, int a, int b) {
 int left=Node<<1, right=left|1;</pre>
  Tree [Node] += lazy [Node] * (b-a+1);
  if(a != b) {
    lazy[left]+=lazy[Node];
    lazy[right]+=lazy[Node];
  lazv[Node] = 0;
void update_tree (int Node, int start, int end, ll value,
    int a, int b) {
  int mid=(a+b)>>1, left=Node<<1, right=left|1;</pre>
 if(lazy[Node] != 0)
   Propogate (Node, a, b);
```

```
if(a > b || a > end || b < start) {</pre>
    if(start <= a && b <= end) {
      if (a != b) {
        lazy[left] += value;
        lazy[right] += value;
      Tree[Node] += value * (b - a + 1);
      update_tree(left, start, end, value, a, mid);
      update_tree(right, start, end, value, mid+1, b);
      Tree[Node] = Tree[left] + Tree[right];
11 query(int Node, int start, int end, int a, int b) {
  int mid=(a+b)>>1, left=Node<<1, right=left|1;</pre>
  if(lazy[Node] != 0)
    Propogate (Node, a, b);
  if (a > b || a > end || b < start) {</pre>
    return 0:
  } else {
    11 Sum1, Sum2;
    if (start <= a && b <= end) {</pre>
      return Tree[Node];
    } else {
      Sum1 = query(left, start, end, a, mid);
      Sum2 = query(right, start, end, mid + 1, b);
return Sum1+Sum2;
```

6 Miscellaneous

6.1 C++ template

```
#include <bits/stdc++.h>
using namespace std;
const long long Mod = 1e9 + 7;
const long long Inf = 1e18;
const long long Lim = 1e5 + 1e3;
const double eps = 1e-10;
typedef long long 11;
typedef vector <int> vi;
typedef vector <11> v1;
typedef pair <int, int> pii;
typedef pair <11, 11> pll;
typedef vector <pii> vii;
typedef vector <pll> vll;
#define F first
#define S second
#define uint unsigned int
#define mp make_pair
#define pb push back
#define pi 2*acos(0.0)
#define rep2(i,b,a) for(ll i = (ll)b, a = (ll)a; i \ge a;
    i --)
#define repl(i,a,b) for(ll i = (ll)a, _b = (ll)b; i <= _b;
#define rep(i,n) for(ll i = 0, _n = (ll)n; i < _n ; i++)
#define mem(a, val) memset(a, val, sizeof(a))
#define fast ios_base::sync_with_stdio(false),cin.tie(0),
    cout.tie(0);
```

6.2 C++ input/output

```
#include <iostream>
#include <iomanip>
using namespace std;
```

6.3 Sieve for Prime Numbers

```
#include "template.h"
    isPrime stores the largest prime number which divides
    vector primeNum contains all the prime numbers
vi primeNum;
int isPrime[Lim];
void pop_isPrime(int limit) {
    mem(isPrime, 0);
    repl(i, 2, limit)
        if (isPrime[i])
            continue;
        if (i <= (int) (sqrt(limit)+10))</pre>
            for (ll j = i * i; j <= limit; j += i)
                isPrime[j] = i;
        primeNum.pb(i);
        isPrime[i]=i:
int main() {
    fast;
    pop_isPrime(500);
    rep1(i, 1, 500)
        cout << i << ' ' << isPrime[i] << '\n';
```

6.4 Longest increasing subsequence

```
#define STRICTLY INCREASING
VI LongestIncreasingSubsequence(VI v) {
 VPII best:
 VI dad(v.size(), -1);
 for (int i = 0; i < v.size(); i++) {</pre>
#ifdef STRICTLY_INCREASNG
   PII item = make pair(v[i], 0);
    VPII::iterator it = lower_bound(best.begin(), best.end
        (), item);
    item.second = i:
#else
    PII item = make_pair(v[i], i);
    VPII::iterator it = upper_bound(best.begin(), best.end
         (), item);
#endif
   if (it == best.end()) {
      dad[i] = (best.size() == 0 ? -1 : best.back().second)
     best.push_back(item);
    } else {
      dad[i] = dad[it->second];
      *it = item;
 for (int i = best.back().second; i >= 0; i = dad[i])
   ret.push_back(v[i]);
  reverse(ret.begin(), ret.end());
 return ret;
```

6.5 Knuth-Morris-Pratt

```
Searches for the string w in the string s (of length k).
0-based index of the first match (k if no match is found).
    Algorithm
runs in O(k) time.
#include <iostream>
#include <string>
#include <vector>
using namespace std;
typedef vector<int> VI;
void buildTable(string& w, VI& t)
 t = VI(w.length());
 int i = 2, j = 0;
t[0] = -1; t[1] = 0;
  while(i < w.length())</pre>
    if(w[i-1] == w[j]) { t[i] = j+1; i++; j++; }
    else if(j > 0) j = t[j];
    else { t[i] = 0; i++; }
int KMP (string& s, string& w)
  int m = 0, i = 0;
  VI t:
  buildTable(w, t);
  while(m+i < s.length())</pre>
    if(w[i] == s[m+i])
      if(i == w.length()) return m;
    else
      m += i-t[i];
      if(i > 0) i = t[i];
```

```
return s.length();
int main()
  string a = (string) "The example above illustrates the
       general technique for assembling "+
    "the table with a minimum of fuss. The principle is that of the overall search: "+
    "most of the work was already done in getting to the
         current position, so very "+
    "little needs to be done in leaving it. The only minor
         complication is that the "+
    "logic which is correct late in the string erroneously
         gives non-proper "+
    "substrings at the beginning. This necessitates some
         initialization code.";
  string b = "table";
  int p = KMP(a, b); cout << p << ": " << a.substr(p, b.length()) << " " << b
       << endl;
```

6.6 Topological sort (C++)

```
// This function uses performs a non-recursive topological
    sort.
// Running time: O(|V|^2). If you use adjacency lists (
     vector<map<int> >),
                 the running time is reduced to O(|E|).
    INPUT: w[i][j] = 1 if i should come before j, 0
     otherwise
    OUTPUT: a permutation of 0,...,n-1 (stored in a
     vector)
              which represents an ordering of the nodes
     which
              is consistent with w
// If no ordering is possible, false is returned.
#include <iostream>
#include <queue>
#include <cmath>
#include <vector>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
bool TopologicalSort (const VVI &w, VI &order) {
 int n = w.size();
 VI parents (n);
 queue<int> q;
 order.clear();
  for (int i = 0; i < n; i++) {</pre>
   for (int j = 0; j < n; j++)
     if (w[j][i]) parents[i]++;
      if (parents[i] == 0) q.push (i);
  while (q.size() > 0){
    int i = q.front();
   q.pop();
    order.push_back (i);
    for (int j = 0; j < n; j++) if (w[i][j]) {</pre>
     parents[j]--;
      if (parents[j] == 0) q.push (j);
 return (order.size() == n);
```

6.7 Fast exponentiation

```
Uses powers of two to exponentiate numbers and matrices.
     Calculates
n^k in O(\log(k)) time when n is a number. If A is an n x n
     matrix.
calculates A^k in O(n^3*log(k)) time.
#include <iostream>
#include <vector>
using namespace std;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
T power(T x, int k) {
  T ret = 1;
  while(k) {
   if(k & 1) ret *= x;
    k \gg 1; x \star = x;
 return ret:
VVT multiply(VVT& A, VVT& B) {
 int n = A.size(), m = A[0].size(), k = B[0].size();
 VVT C(n, VT(k, 0));
 for (int i = 0; i < n; i++)
    for (int j = 0; j < k; j++)
     for(int 1 = 0; 1 < m; 1++)
       C[i][j] += A[i][1] * B[1][j];
VVT power(VVT& A, int k) {
 int n = A.size();
  VVT ret(n, VT(n)), B = A;
  for(int i = 0; i < n; i++) ret[i][i]=1;</pre>
  while(k) {
   if(k & 1) ret = multiply(ret, B);
   k >>= 1; B = multiply(B, B);
 return ret;
int main()
 /* Expected Output:
    2.37^48 = 9.72569e+17
     376 264 285 220 265
     550 376 529 285 484
     484 265 376 264 285
     285 220 265 156 264
     529 285 484 265 376 */
  double n = 2.37;
 int k = 48;
 cout << n << "^" << k << " = " << power(n, k) << endl;
 double At [5] [5]
    { 0, 0, 1, 0, 0 }, 
{ 1, 0, 0, 1, 0 },
     0, 0, 0, 0, 1 },
     1, 0, 0, 0, 0 },
  vector <vector <double> > A(5, vector <double>(5));
  for (int i = 0; i < 5; i++)
   for (int j = 0; j < 5; j++)
     A[i][j] = At[i][j];
 vector <vector <double> > Ap = power(A, k);
  cout << endl;
  for(int i = 0; i < 5; i++) {
    for (int j = 0; j < 5; j++)
     cout << Ap[i][j] << " ";
    cout << endl;
```

6.8 Longest common subsequence

```
Calculates the length of the longest common subsequence of
    two vectors.
Backtracks to find a single subsequence or all subsequences
     . Runs in
O(m*n) time except for finding all longest common
    subsequences, which
may be slow depending on how many there are.
#include <iostream>
#include <vector>
#include <set>
#include <algorithm>
using namespace std;
typedef int T;
typedef vector<T> VT;
typedef vector<VT> VVT;
typedef vector<int> VI;
typedef vector<VI> VVI;
void backtrack(VVI& dp, VT& res, VT& A, VT& B, int i, int j
 if(!i || !j) return;
 if(A[i-1] == B[j-1]) \{ res.push\_back(A[i-1]); backtrack(
      dp, res, A, B, i-1, j-1); }
   if(dp[i][j-1] >= dp[i-1][j]) backtrack(dp, res, A, B, i
         , i-1);
    else backtrack(dp, res, A, B, i-1, j);
void backtrackall(VVI& dp, set<VT>& res, VT& A, VT& B, int
    i, int j)
 if(!i || !j) { res.insert(VI()); return; }
 if(A[i-1] == B[j-1])
   backtrackall(dp, tempres, A, B, i-1, j-1);
    for(set<VT>::iterator it=tempres.begin(); it!=tempres.
         end(); it++)
      VT temp = *it;
     temp.push_back(A[i-1]);
      res.insert(temp);
 else
    if(dp[i][j-1] >= dp[i-1][j]) backtrackall(dp, res, A, B
         , i, j-1);
    if(dp[i][j-1] \le dp[i-1][j]) backtrackall(dp, res, A, B
        , i-1, j);
VT LCS(VT& A, VT& B)
 VVI dp;
 int n = A.size(), m = B.size();
 dp.resize(n+1);
 for(int i=0; i<=n; i++) dp[i].resize(m+1, 0);</pre>
 for (int i=1; i<=n; i++)</pre>
   for(int j=1; j<=m; j++)</pre>
      if(A[i-1] == B[j-1]) dp[i][j] = dp[i-1][j-1]+1;
      else dp[i][j] = max(dp[i-1][j], dp[i][j-1]);
 backtrack(dp, res, A, B, n, m);
```

```
reverse(res.begin(), res.end());
 return res;
set < VT > LCSall (VT& A, VT& B)
 VVI dp;
 int n = A.size(), m = B.size();
 dp.resize(n+1);
 for(int i=0; i<=n; i++) dp[i].resize(m+1, 0);</pre>
 for(int i=1; i<=n; i++)</pre>
   for(int j=1; j<=m; j++)</pre>
     if(A[i-1] == B[j-1]) dp[i][j] = dp[i-1][j-1]+1;
     else dp[i][j] = max(dp[i-1][j], dp[i][j-1]);
  set<VT> res;
 backtrackall(dp, res, A, B, n, m);
 return res;
int main()
 VI C = LCS(A, B);
 set <VI> D = LCSall(A, B);
 for(set<VI>::iterator it = D.begin(); it != D.end(); it
      ++)
   for (int i=0; i<(*it).size(); i++) cout << (*it)[i] << "</pre>
```

```
cout << endl;
}</pre>
```

6.9 Miller-Rabin Primality Test (C)

```
// Randomized Primality Test (Miller-Rabin):
// Error rate: 2^(-TRIAL)
// Almost constant time. srand is needed
#include <stdlib.h>
#define EPS 1e-7
typedef long long LL;
LL ModularMultiplication(LL a, LL b, LL m)
       LL ret=0, c=a;
       while(b)
                if (b&1) ret=(ret+c) %m;
               b>>=1; c=(c+c)%m;
        return ret;
LL ModularExponentiation(LL a, LL n, LL m)
       LL ret=1, c=a;
       while(n)
                if(n&1) ret=ModularMultiplication(ret, c, m
               n>>=1; c=ModularMultiplication(c, c, m);
```

```
return ret;
bool Witness(LL a, LL n)
        LL u=n-1:
  int t=0;
        while (!(u\&1))\{u>>=1; t++;\}
        LL x0=ModularExponentiation(a, u, n), x1;
        for (int i=1; i <=t; i++)</pre>
                x1=ModularMultiplication(x0, x0, n);
                if(x1==1 && x0!=1 && x0!=n-1) return true;
        if(x0!=1) return true;
        return false;
LL Random(LL n)
  LL ret=rand(); ret*=32768;
        ret+=rand(); ret*=32768;
        ret+=rand(); ret*=32768;
        ret+=rand();
  return ret%n;
bool IsPrimeFast(LL n, int TRIAL)
  while (TRIAL--)
    LL a=Random(n-2)+1;
    if(Witness(a, n)) return false;
  return true;
```