

- Q1. (a)** Floating point arithmetic is known to be commutative. Let  $A = 1.0$ ,  $B = 1.5 \times 10^{38}$  and  $C = -1.5 \times 10^{38}$ . Prove that FP arithmetic is not associative using the given values.
- (b)** Perform addition of the floating point numbers 0x3FC00000 and 0x40500000 in the IEEE 754 single precision format. Give the sum in hexa form. Show all the steps clearly.
- (c)** What 8-bit pattern would represent the number  $3/256$ , in an 8-bit hypothetical floating-point format, including support for denormalized numbers and nonnumeric values where the first bit is for sign, next four bits for excess-7 exponent and 3 bits for the mantissa?
- (d)** Consider the 64-bit, IEEE 754 floating point format.  $b_{51}..b_0$  is the mantissa,  $b_{63}$  is the sign bit  $s$  and  $E'$  be the biased exponent. Derive an expression to compute the base-10 value of a normalized double precision number in terms of the given info.
- (e)** Perform the 32-bit division of the IEEE 754 Floating point numbers  $(X_1 / X_2)$ , where  $X_1 = 127.03125$  and  $X_2 = 16.9375$ . Show the steps clearly.

(2 x 5 = 10 marks)

- Q2. (a)** Translate the for loop control structure given below to its MIPS equivalent in the template given below.

| for loop control structure  | Equivalent MIPS code |
|---|----------------------|
| \$a0 = 0;   | li \$a0, ____        |
| for ( \$t0 = 10; \$t0 > 0; \$t0 = \$t0 - 1) do { \$a0 = \$a0 + \$t0 } | li \$t0, ____        |
|   | loop:                |
|   | add ____             |
|   | addi ____            |
|   | bgtz ____            |

- (b)** You are quite familiar with MIPS 32-bit ISA. If the instruction set of the MIPS processor is modified, the instruction format must also be changed. For each of the suggested changes, show the size of the bit fields of an R-type format instruction. Also, what is the total number of bits needed for each instruction?

| Suggested modification   | Size of bit fields and instruction format | Total no. of bits needed for each instruction |
|--|---|---|
| 8 registers  |   |   |
| 10-bit immediate constant  |   |   |
| 128 registers  |   |   |
| All arithmetic instructions can use base addressing mode with the last argument (Example: add \$a0, \$a2, 0[\$t1]) |   |   |

- (c)** This problem pertains to reading a byte-addressable memory in MIPS. Refer to the following byte-addressable scheme in MIPS and complete the table.

| Word Address | Data            |        |
|--------------|-----------------|--------|
| 0000000C     | 4 0 F 3 0 7 8 8 | Word 3 |
| 00000008     | 0 1 E E 2 8 4 2 | Word 2 |
| 00000004     | F 2 F 1 A C 0 7 | Word 1 |
| 00000000     | A B C D E F 7 8 | Word 0 |

width = 4 bytes

After the instruction lw \$s3, 4(\$0) is executed, value at the destination operand is \_\_\_\_\_

The instruction to store the value held in \$t7 into the 11<sup>th</sup> 32-bit memory location is \_\_\_\_\_

(d) Consider the big-endian and little-endian schemes in MIPS. Assume that \$t0 initially contains 0x23456789. After the following code is run, what are the respective outcomes? Answer as per the table given below.

sw \$t0, 0(\$0)

lb \$s0, 1(\$0)

| Scheme        | Outcome |
|---------------|---------|
| Big-endian    |         |
| Little endian |         |

(e) Consider the MIPS *div* instruction. The instruction format is **000000** | **R<sub>s</sub>** | **R<sub>t</sub>** | **00000** | **00000** | **011010**. Show in psedo code form, the 3 effects of execution of this instruction.

(2 x 5 = 10 marks)

3. (a) The 5 stages of a processor have the following latencies. Assume that when pipelining, each pipeline stage costs 20ps extra for the registers between the pipeline stages.

| Fetch | Decode | Execute | Memory | Writeback |
|-------|--------|---------|--------|-----------|
| 300ps | 400ps  | 350ps   | 500ps  | 100ps     |

(i) What is the cycle time for the processor, if the execution is non-pipelined? (ii) If the execution is pipelined, then what is the execution time?

(b) Assume the distribution of instructions that run on a processor is: 50% ALU; 25% BEQ; 15% LW and 10% SW. Assuming there are no stalls or hazards, what is the utilization of the data memory?

(c) We have a 'k' segment pipeline with clock cycle time as 'T<sub>p</sub>'. Let there be 'n' tasks to be completed in the pipelined processor. Compute the time taken ET<sub>pipeline</sub> to execute 'n' instructions in this pipelined processor.

(d) Using the data given in Part (c), compute the speedup (S) of a pipelined processor over a non-pipelined processor, when 'n' tasks are executed on the same processor. Given, the number of tasks 'n' is significantly larger than k, that is, n >> k.

(e) This problem pertains to memory alignment in ISAs. For the given address, please fill in the details in the table.

| Memory Address | Alignment (8 bits) | Alignment (16 bits) | Alignment (32 bits) | Alignment (64 bits) |
|----------------|--------------------|---------------------|---------------------|---------------------|
| 0x00000005     |                    |                     |                     |                     |

(2 x 5 = 10 marks)

4. The following table shows the template for implementation of Booth's algorithm for signed multiplication. Complete the implementation. Show all the steps clearly.

| Input: 0110, 0010 // Multiplicand M, Multiplier Q, Product in A |      |                 |      |                |               |
|---|------|-----------------|------|----------------|---------------|
| A   | Q    | Q <sub>-1</sub> | M    |                | step count sc |
| 0000  | 0010 | 0               | 0110 | Initial values |               |

(10 marks)

5. The following algorithm is meant to perform the non-restoring division of unsigned integers. Complete the algorithm.

Step 1: The registers are initialized with corresponding values (Q = Dividend, M = Divisor, A = 0, n = number of bits in dividend)  
 Step 2: Check the sign bit of the register \_\_\_\_  
 Step 3: If it is 1, then shift left the contents of \_\_\_\_ and perform A = A+M, otherwise shift left \_\_\_\_ and perform A = \_\_\_\_  
 Step 4: Again check the sign bit of register \_\_\_\_  
 Step 5: If sign bit is 1, Q<sub>0</sub> becomes \_\_\_\_ otherwise Q<sub>0</sub> becomes \_\_\_\_.  
 Step 6: Decrement the value of \_\_\_\_ by 1  
 Step 7: If n is not equal to zero, go to step \_\_\_\_ otherwise go to next step  
 Step 8: If sign bit of A is 1, then perform A = \_\_\_\_  
 Step 9: Register \_\_\_\_ contains quotient and \_\_\_\_ contains remainder

(12 marks)

6. (a) Let *lw* be the slowest instruction in a single-cycle implementation scenario with a latency of 5 (IF) + 3 (ID) + 5 (EX) + 5 (MEM) + 3 (WB) = 21 ns. There are N instructions and the CPI of a single-cycle implementation is 1. What is the running time?

(b) Now, if the ALU is made 25% faster, it takes 5/1.25 = 4 ns to do the ALU computation; the clock-cycle time is reduced to 20 ns. What is the new running time? What is the percentage of improvement?

(4 + 4 = 8 marks)