

✓ Path in graph: Let $G=(V,E)$ be a graph. Let v_i and v_j be two vertices in V_G . We say that there is a path between v_i and v_j , if there exists an ordered tuple of edges $(e_1, e_2, e_3, \dots, e_n) \in E_G$ such that first end-point of $e_1 = v_i$, the last end-point of $e_n = v_j$ and for all edges $e_1, e_2, \dots, e_k, \dots, e_{n-1}$, second end-point of $e_k =$ first-end point of e_{k+1} .

✓ Comment: If there exists a path between v_i to v_j then we denote it as follows

$$v_i \overset{P}{\rightsquigarrow} v_j$$

$$v_i \overset{P_1}{\rightsquigarrow} v_j \mid v_i \overset{P_2}{\rightsquigarrow} v_j \mid v_i \overset{P_3}{\rightsquigarrow} v_j \mid$$

$P_1 = (e_1, \dots)$ tuple of edges.

$P_2 = \emptyset$

P_3

$$v_i \overset{P}{\rightsquigarrow} v_j \mid P = \underline{(e_1, \dots, e_n)} \text{ where } e_1, \dots, e_n \in E_G, \\ v_i, v_j \in V_G.$$

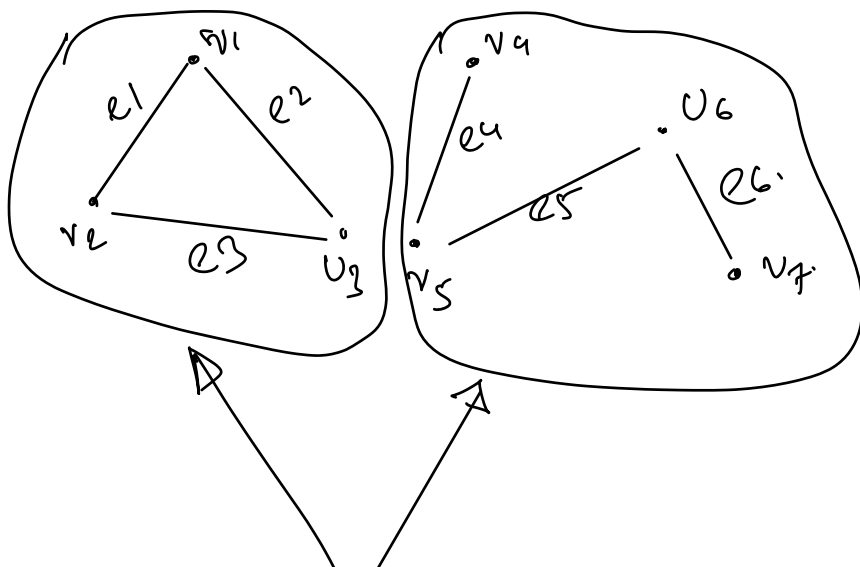
Cycle in graph: Let $G = (V, E)$ be a graph. Let $v_i \in V_G$.
 If there exists a path p between v_i to v_i which is
 NOT a self edge, then such path is known as a
 cycle.

$(\exists v_i \in V_G \ni \exists p \text{ } v_i \overset{p}{\rightsquigarrow} v_i) \Rightarrow G \text{ contains a cycle!}$
 (Annotations:
 - $\exists v_i \in V_G$: [Existential Quantifier] (there exists, belongs)
 - \ni : such that
 - $\overset{p}{\rightsquigarrow}$: path
 - \Rightarrow : implies)

Connected Graphs

Let $G = (V, E)$ be a graph. If there exists a path between
 any two distinct vertices v_i and $v_j \in V_G$, then G is a
 connected graph.

$G = (V, E)$ is connected $\equiv [\forall v_i, v_j \in V_G, v_i \neq v_j,$
 if $\exists p, v_i \overset{p}{\rightsquigarrow} v_j]$



$V_G = (v_1, \dots, v_7)$

$E_G = (e_1, \dots, e_6)$

$v_1 \dots v_7$
 e_1

Connected Components of graph or forest of graph.

$$\vdash \boxed{\forall x. p(x)} \equiv \exists x. \neg p(x)$$

$$\neg (\forall x. p(x)) \equiv \exists x. \neg p(x)$$

$G = (V, E)$ is connected $\equiv [\forall v_i, v_j \in V_G, v_i \neq v_j, \underline{\underline{\text{ADV}}}$
if $\exists p, v_i \overset{p}{\rightsquigarrow} v_j]$

$\neg (G = (V, E) \text{ is connected}) \equiv [G = (V, E)] \text{ is not connected}$

$\neg (\forall (v_i, v_j) \in V_G, \exists p (v_i \overset{p}{\rightsquigarrow} v_j)) \equiv$
 $\exists (v_i, v_j) \in V_G \ni (\exists p (v_i \overset{p}{\rightsquigarrow} v_j) \equiv F)$

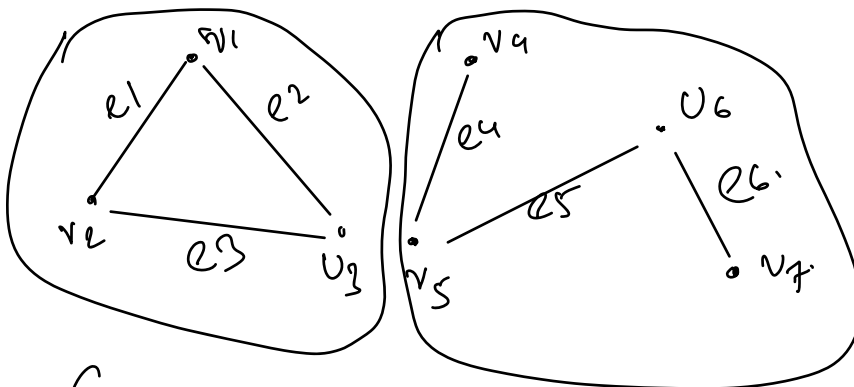
$\exists (v_i, v_j) \in V_G \ni v_i \not\overset{p}{\rightsquigarrow} v_j$

Let $G = (V, E)$ be a graph. Let V' be a non-empty subset of V . If $\forall v_i, v_j \in V', \exists p v_i \overset{p}{\rightsquigarrow} v_j$ then V' is a connected component.

Let $G = (V, E)$ be a graph. Let $V_S \subseteq V$ and $V_S \neq \emptyset$ and let $E_S \subseteq E$ and $E_S \neq \emptyset$ such that for any two vertices in V_S , there exists a path and all edges

belonging to any path of any two vertices in V_S belong to E_S then (V_S, E_S) is a connected component or a forest of G .

Let $G = (V, E)$ be a graph. Let $E_S \subseteq E$ and $E_S \neq \emptyset$ and V_S is set of all vertices such that each vertex in V_S is an end point of at least one edge in E_S . Then $(V_S, E_S) \subseteq (V, E)$



$$G = (\{v_1, v_2, v_3, v_4, v_5, v_6, v_7\}, \{e_1, e_2, e_3, e_4, e_5, e_6\})$$

$$G_1 = (\{v_1, v_2, v_3\}, \{e_1, e_2, e_3\}) \mid G_1 \text{ is a forest.}$$

$$G_2 = (\{v_4, v_5, v_6, v_7\}, \{e_4, e_5, e_6\}) \mid G_2 \text{ is a connected component}$$

$$G \mid G_1 \mid G_2. \quad G_1 \subseteq G \mid G_2 \subseteq G.$$

$$\{e_1, e_3, e_6\} \subseteq \{e_1, e_2, e_3, e_4, e_5, e_6\}$$

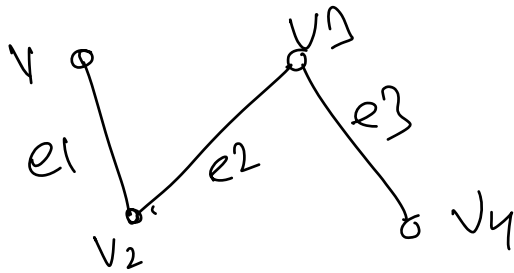
$$\{v_1, v_2, v_3, v_6, v_7\} \subseteq V$$

derive

$$(\{v_1, v_2, v_3, v_6, v_7\}, \{e_1, e_3, e_6\}) \subseteq G.$$

Construct a subgraph of G .

- 1) Select subset $E_S \neq \emptyset$, of E_G . ($E_S \subseteq E_G$).
- 2) Collect all end points of edges in E_S without repetition in set V_S .
- 3) $(V_S, E_S) \subseteq G$.



$$\{e_1, e_3\}$$

$$\{v_1, v_2, v_3, v_4\}$$